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Chapter 19. Mentor Graphics Precision Synthesis Support
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Revised: June 2012
Part Number: QII51010-12.0.0

Chapter 21. Analyzing Designs with Quartus II Netlist Viewers
Revised: November 2012
Part Number: QII51013-12.1.0
Section I. Design Flows

The Altera® Quartus® II design software provides a complete design environment that easily adapts to your specific design requirements. This handbook is arranged in chapters, sections, and volumes that correspond to the major stages in the overall design flow. For a general introduction to features and the standard design flow in the software, refer to the *Introduction to the Quartus II Software* manual.

This section is an introduction to design planning. It documents various specialized design flows in the following chapters:

- **Chapter 1, Managing Quartus II Projects**
  Describes how to manage all the elements in your Quartus II project. You can save multiple revisions of your project to experiment with settings that achieve your design goals. Quartus II projects also support team-based, distributed work flows and a scripting interface.

- **Chapter 2, Design Planning with the Quartus II Software**
  This chapter is an overview of various design planning considerations: device selection, early power estimation, I/O pin planning, and design planning. To help you improve design productivity, it provides recommendations and describes various tools available for Altera FPGAs.

- **Chapter 3, Quartus II Incremental Compilation for Hierarchical and Team-Based Design**
  This chapter documents Altera’s incremental design and compilation flow, which allows you to preserve the results and performance for unchanged logic in your design as you make changes elsewhere, reduces design iteration time by up to 70% so you achieve timing closure efficiently, and facilitates modular hierarchical and team-based design flows using top-down or bottom-up methodologies.

- **Chapter 4, Design Planning for Partial Reconfiguration**
  This chapter provides a high-level guide to the use of partial reconfiguration in the Quartus II software. Partial reconfiguration allows you to reconfigure a portion of the FPGA dynamically, while the remainder of the device continues to operate.

- **Chapter 5, Designing HardCopy Series Devices**
  With the Quartus II software, you can use an FPGA device as a prototype and seamlessly migrate your design to a HardCopy ASIC to reduce cost for volume production. This chapter describes the Quartus II support for HardCopy flows.
Chapter 6, Quartus II Design Separation Flow

This chapter describes rules and guidelines for creating a floorplan with the Design Separation flow. The Quartus II Design Separation flow provides the ability to design physically independent structures on a single device. This allows system designers to achieve a higher level of integration on a single FPGA, and alleviates increasingly strict Size Weight and Power (SWaP) requirements.
The Quartus® II software organizes and manages the elements of your design within a project. The project encapsulates information about your design hierarchy, libraries, constraints, and project settings. Click File > New Project Wizard to quickly create a new project and specify basic project settings.

When you open a project, a unified GUI displays integrated project information. The Project Navigator allows you to view and edit the elements of your project. The Messages window lists important information about project processing.

You can save multiple revisions of your project to experiment with settings that achieve your design goals. Quartus II projects support team-based, distributed work flows and a scripting interface.

Quick Start

To quickly create a project and specify basic settings, click File > New Project Wizard.

Figure 1–1. Quick Project Setup with New Project Wizard
Understanding Quartus II Projects

A single Quartus II Project File (.qpf) represents each project. The text-based .qpf references the Quartus II Settings File (.qsf), that lists all project files and stores project and entity settings. When you make project changes in the GUI, these text files automatically store the changes. The GUI provides access to all project settings and helps to manage all aspects of your project, including:

- Creating and viewing projects
- Managing logic design, EDA, IP core, and Qsys system files
- Specifying and optimizing project settings and constraints
- Archiving and migrating projects

Table 1–1. Quartus II Project At a Glance (Gray Files Optional)

<table>
<thead>
<tr>
<th>File Type</th>
<th>Stores</th>
<th>Click to Access</th>
<th>File Format(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Project file</td>
<td>Project and revision name</td>
<td>File&gt; New Project Wizard View&gt;Project Navigator Project&gt; Revisions</td>
<td>Quartus II Project File (.qpf)</td>
</tr>
<tr>
<td>Project settings</td>
<td>Files list, settings, device, synthesis directives, and pin placement</td>
<td>Assignments&gt; Settings Assignments&gt; Device Assignments&gt; Assignment Editor</td>
<td>Quartus II Settings File (.qsf)</td>
</tr>
<tr>
<td>Project database</td>
<td>Compilation results</td>
<td>Project&gt; Export Database Project&gt; Export Design Partition Project&gt; Clean Project</td>
<td>Quartus II Exported Partition (.qxp)</td>
</tr>
<tr>
<td>Timing constraints</td>
<td>Clock properties, exceptions, setup/hold time</td>
<td>Tools&gt; TimeQuest Timing Analyzer</td>
<td>Synopsys Design Constraints (.sdc)</td>
</tr>
<tr>
<td>Logic design files</td>
<td>RTL and other design logic source files</td>
<td>View&gt; Project Navigator File&gt; New</td>
<td>Verilog Design File (.v)</td>
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<tr>
<td></td>
<td></td>
<td></td>
<td>VHDL Design File (.vhd)</td>
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<td></td>
<td></td>
<td></td>
<td>Block Design File (.bdf)</td>
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<tr>
<td></td>
<td></td>
<td></td>
<td>EDA Tool Synthesis File (.vqm)</td>
</tr>
<tr>
<td>Programming files</td>
<td>Device programming options and information</td>
<td>Assignments&gt; Settings Tools&gt; Programmer</td>
<td>Chain Description File (.cdf)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>SRAM Object File (.sop)</td>
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<tr>
<td></td>
<td></td>
<td></td>
<td>Programmer Object File (.pot)</td>
</tr>
<tr>
<td>Project libraries</td>
<td>Project and global library information</td>
<td>Assignments&gt; Settings</td>
<td>quartus2.ini file (global)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>.qsf (project)</td>
</tr>
<tr>
<td>IP files</td>
<td>IP core logic, synthesis, and simulation information</td>
<td>View&gt; Project Navigator Tools&gt; Qsys Project&gt; Upgrade IP Components Tools&gt; MegaWizard Plug-In Manager</td>
<td>Verilog Design File (.v)</td>
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<td></td>
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<td></td>
<td>System Verilog File (.sv)</td>
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<td></td>
<td>VHDL Design File (.vhd)</td>
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<td>Quartus II IP File (.qip)</td>
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<td></td>
<td>Quartus II Simulation IP (.sip)</td>
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<td></td>
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<td>Various EDA simulation files</td>
</tr>
<tr>
<td>EDA tool files</td>
<td>Files generated for third-party EDA tools</td>
<td>Assignments&gt; Settings Tools&gt; Options&gt; EDA Tool Options</td>
<td>Verilog Output File (.vo)</td>
</tr>
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<td></td>
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<td></td>
<td>VHDL Output File (.vo)</td>
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<td></td>
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<td></td>
<td>Verilog Quartus Mapping (.vqm)</td>
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<td></td>
<td>Stamp model files</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>PartMiner XML-Format (.xml)</td>
</tr>
<tr>
<td>Archive files</td>
<td>Complete project as single compressed file</td>
<td>Project&gt; Archive Project</td>
<td>Quartus II Archive File (.qar)</td>
</tr>
</tbody>
</table>
Figure 1–2. Basic Project Directory (Gray Files and Directories Optional)

- `<Quartus II Project Directory>`
  - `<project_name>.qpf` - Quartus II Project file
  - `<revision_name>bsf` - represents design in schematics
  - `<revision_name>.qsf` - stores revision's project settings and constraints
  - `<revision_name>.assignments.default.qdf` - stores default project settings and constraints
  - `<revision_name>.sdc` - stores timing constraints in Synopsys Design Constraints format
  - `<logic_design_file>.v` or `<logic_design_file>.vhd` - RTL source code
  - `<Qsys_system_name>.qsys` - Qsys system file
  - `<logic_design_file>.vqm` - logic from EDA synthesis tool
  - `<instance name> - QII IP synthesis files`
  - `<instance name>_sim - QII IP simulation files`
  - `<Qsys_system_name> - Qsys system and IP files`
  - `simulation - EDA simulation files`
  - `symbols - EDA board-level symbol tool files`
  - `board - EDA board-level signal integrity tool files`
  - `timing - EDA board-level timing analysis tool files`

**Viewing Your Project**

View basic information about your project in the Project Navigator, Report panel, and Messages window.

**Viewing Basic Project Information**

View project elements in the Project Navigator (View > Utility Windows > Project Navigator). The Project Navigator displays key project information, including design files, IP components, and revisions of your project. Use the Project Navigator to:

- View and modify the design hierarchy (right-click > Set as Top-Level Entity)
- Set the project revision (right-click > Set Current Revision)
- View and update logic design files and constraint files (right-click > Open)
- Update IP component version information (right-click > Upgrade IP Component)

Figure 1–3. Project Navigator Hierarchy, Files, Revisions, and IP Components
Viewing Project Reports

The Report panel (Processing > Compilation Report) displays detailed reports after project processing, including the following:

- Analysis & Synthesis reports
- Fitter reports
- Timing analysis reports
- Power analysis reports
- Signal integrity reports

Refer to the List of Compilation Reports in Quartus II Help for a complete list.

Analyze the detailed project information in these reports to determine correct implementation. Right-click report data to locate and edit the source in project files.

Figure 1–4. Report Panel

Viewing Project Messages

The Messages window (View > Utility Windows > Messages) displays information, warning, and error messages about Quartus II processes. Right-click messages to locate the source or get message help.

- Processing tab—displays messages from the most recent process
- System tab—displays messages unrelated to design processing
- Search—locates specific messages

Messages are written to stdout when you use command-line executables.

For more information about the Messages window and message suppression, refer to About the Messages Window and About Message Suppression in Quartus II Help.
Suppressing Messages

Suppress display of unimportant messages so they do not obscure valid messages. Right-click messages and choose any of the following:

- **Suppress Message**—suppresses all messages matching exact text
- **Suppress Messages with Matching ID**—suppresses all messages matching the message ID number, ignoring variables
- **Suppress Messages with Matching Keyword**—suppresses all messages matching keyword or hierarchy path
- **Message Suppression Manager**—manages all message suppression rules

Figure 1–5. Messages Window

Figure 1–6. Message Suppression by Message ID Number
Message Suppression Guidelines

- You cannot suppress error or Altera® legal agreement messages.
- Suppressing a message also suppresses any submessages.
- Message suppression is revision-specific. Derivative revisions inherit any suppression.
- You cannot edit messages or suppression rules during compilation.

Managing Logic Design Files

The Quartus II software helps you create and manage the logic design files in your project. Logic design files contain the logic that implements your design. When you add a logic design file to the project, the Compiler automatically compiles that file as part of the project. The Compiler synthesizes your logic design files to generate programming files for your target device.

The Quartus II software includes full-featured schematic and text editors, as well as HDL templates to accelerate your design work. The Quartus II software supports VHDL Design Files (.vhd), Verilog HDL Design Files (.v), SystemVerilog (.sv) and schematic Block Design Files (.bdf). The Quartus II software also supports Verilog Quartus Mapping (.vqm) design files generated by other design entry and synthesis tools. In addition, you can combine your logic design files with Altera and third-party IP core design files, including combining components into a Qsys system (.qsys).

The New Project Wizard prompts you to identify logic design files. Add or remove project files by clicking Project > Add/Remove Files in Project. View the project’s logic design files in the Project Navigator.

Figure 1–7. Design and IP Files in Project Navigator

Right-click files in the Project Navigator to:

- Open and edit the file
- Remove File from Project
- Set as Top-Level Entity for the project revision
- Create a Symbol File for Current File for display in schematic editors
- Edit file Properties
Including Design Libraries

You can include design files libraries in your project. Specify libraries for a single project, or for all Quartus II projects. The .qsf stores project library information. The quartus2.ini file stores global library information. Refer to “Migrating Design Libraries” on page 1–21 for migration guidelines.

Specifying Design Libraries

To specify project libraries from the GUI:

1. Click Assignment > Settings.
2. Click Libraries and specify the Project Library name or Global Library name.

Alternatively, you can specify project libraries with SEARCH_PATH in the .qsf, and global libraries in the quartus2.ini file.

Refer to Recommended Design Practices and Recommended HDL Coding Styles in the Quartus II Handbook for more information about creating logic design files.

Managing Project Settings

The New Project Wizard helps you initially assign basic project settings. Optimizing project settings enables the Compiler to generate programming files that meet or exceed your specifications. The .qsf stores each revision’s project settings.

Click Assignments > Settings to access global project settings, including:

- Project files list
- Synthesis directives and constraints
- Logic options and compiler effort levels
- Placement constraints
- Timing constraint files
- Operating temperature limits and conditions
- File generation for other EDA tools
- Target device (click Assignments > Device)

The Quartus II Default Settings File (<revision name>_assignment_defaults.qdf) stores initial settings and constraints for each new project revision.
The Assignment Editor (Tools > Assignment Editor) provides a spreadsheet-like interface for assigning all instance-specific settings and constraints.

### Optimizing Project Settings

Optimize project settings to meet your design goals. The Quartus II Design Space Explorer iteratively compiles your project with various setting combinations to find the optimal setting for your goals. Alternatively, you can create a project revision or project copy to manually compare various project settings and design combinations.
Optimizing with Design Space Explorer

Use the Design Space Explorer (Tools > Launch Design Space Explorer) to find optimal project settings for resource, performance, or power optimization goals. Design Space Explorer (DSE) processes your design using various setting and constraint combinations, and reports the best settings for your design. DSE attempts multiple seeds to identify one meeting your requirements. DSE can run different compilations on multiple computers in parallel to streamline timing closure.

Figure 1–10. Design Space Explorer

Optimizing with Project Revisions

You can save multiple, named project revisions within your Quartus II project (Project > Revisions). Each revision captures a unique set of project settings and constraints, but does not capture any logic design file changes. Use revisions to experiment with different settings while preserving the original. You can compare revisions to determine the best combination, or optimize different revisions for various applications. Use revisions for the following:

- Create a unique revision to optimize a design for different criteria, such as by area in one revision and by \( f_{\text{MAX}} \) in another revision. When you create a new revision, the default Quartus II settings initially apply.

- Create a revision of a revision to experiment with settings and constraints. The child revision includes all the assignments and settings of the parent revision.

You create, delete, specify current, and compare revisions in the Revisions dialog box. Each time you create a new project revision, the Quartus II software creates a new .qsf using the revision name.

To compare each revision’s synthesis, fitting, and timing analysis results side-by-side, click Project > Revisions and then click Compare.
In addition to viewing the compilation results of each revision, you can also compare the assignments for each revision. This comparison reveals how different optimization options affect your design.

**Figure 1–11. Comparing Project Revisions**

![Comparing Project Revisions](image)

### Copying Your Project

Click **Project > Copy Project** to create a separate copy of your project, rather than just a revision within the same project. The project copy includes all design files, .qsf(s), and project revisions. Use this technique to optimize project copies for different applications. For example, optimize one project to interface with a 32-bit data bus, and optimize a project copy to interface with a 64-bit data bus.

### Managing Timing Constraints

Apply appropriate timing constraints to correctly optimize fitting and analyze timing for your design. The Fitter optimizes the placement of logic in the device to meet your specified timing and routing constraints.

Specify timing constraints in the TimeQuest Timing Analyzer (**Tools > TimeQuest Timing Analyzer**), or in an .sdc file. Specify constraints for clock characteristics, timing exceptions, and external signal setup and hold times before running analysis. TimeQuest reports the detailed information about the performance of your design compared with constraints in the Compilation Report panel.

Save the constraints you specify in the GUI in an industry-standard Synopsys Design Constraints File (.sdc). You can subsequently edit the text-based .sdc file directly.

For more information about TimeQuest analyzer and SDC constraints, refer to *Quartus II TimeQuest Timing Analyzer* in the *Quartus II Handbook*.
Managing System and IP Components

Virtually all complex FPGA designs include integrated IP cores. The Quartus II GUI helps you define, integrate, and update the IP files in your project. Use Altera’s optimized and verified IP in your project to shorten design cycles and maximize performance. The Quartus II software includes many basic and complex IP cores, and supports IP from other sources. You can combine IP with other design elements to quickly create a complete system using the Qsys system integration tool.

Integrating System and IP Files

You can easily customize and quickly integrate Qsys system and IP core files in your project. The Quartus II software implements your specified system or IP core parameters and generates files for synthesis and simulation in the Quartus II software and other EDA tools.

IP components are represented as design elements in your project. The Quartus II software includes the following IP and system integration tools:

Table 1–2. IP Integration Tools

<table>
<thead>
<tr>
<th>IP Integration Tool</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MegaWizard™ Plug-In Manager</td>
<td>Parameterize individual IP cores and generate HDL synthesis files, simulation models, and testbenches.</td>
</tr>
<tr>
<td>Qsys</td>
<td>Parameterize and connect all components in a system-level hardware design, automating integration of customized HDL components.</td>
</tr>
</tbody>
</table>
FIGURE 1–13. Qsys System Integration Tool and MegaWizard IP Core Editor

Updating Outdated IP Files

Some Altera IP components are version-specific with the Quartus II software. Click Project > Upgrade IP Components to easily upgrade outdated IP in the Project Navigator. Failure to upgrade outdated IP components can result in a mismatch between the outdated IP core variation and the current supporting libraries.

Altera verifies that the current version of the Quartus II software compiles the previous version of each IP core. The MegaCore IP Library Release Notes and Errata reports any verification exceptions. Altera does not verify compilation for IP cores older than the previous release.

FIGURE 1–14. Upgrading IP Components in Project Navigator
System and IP File Locations

When you generate an Altera IP core variation with the MegaWizard Plug-In Manager or Qsys, the Quartus II software generates files in the following locations.

Figure 1–15. System and IP Files Generated by MegaWizard Plug-In Manager and Qsys

Processing Encrypted IP Files

Projects may include encrypted Altera or third-party IP cores that prevent unlicensed viewing of source code. The Compiler processes encrypted IP files along with the rest of your project. The Quartus II software provides a black-box representation of Altera megafuntions and encrypted IP cores for synthesis in other EDA tools.

The Quartus II software also includes IEEE-encrypted Verilog HDL models for both Verilog HDL and VHDL simulation models for Altera IP cores. Use these files to simulate encrypted IP in other EDA tools. The Quartus II software does not provide IP core encryption or decryption functions.

IP File Search Path

If your project includes two IP core files of the same name, the search path precedence rules how similarly named files are resolved. The Quartus II software recognizes the following file naming precedence:

1. Project directory.
2. Project database directory.
3. Project libraries specified in Assignments > Settings > Libraries, or with the SEARCH_PATH assignment in the revision .qsf.
4. Global libraries specified in Assignments > Settings > Libraries, or with the SEARCH_PATH assignment in the quartus2.ini file.
5. Quartus II software libraries directory, such as <Quartus II Installation>\libraries.
Use the `SEARCH_PATH` assignment to define the project libraries. The Quartus II software supports multiple `SEARCH_PATH` assignments. Specify only one source directory for each `SEARCH_PATH` assignment.

For more information, refer to IP core user guides on the IP and Megafunctions Documentation section of the Altera website, and to Creating a System with Qsys in the Quartus II Handbook.

**Integrating Other EDA Tools**

You can integrate supported EDA design entry, synthesis, simulation, physical synthesis, and formal verification tools into the Quartus II design flow. The Quartus II software supports netlist files from other EDA design entry and synthesis tools. The Quartus II software optionally generates various files for use in other EDA tools.

The Quartus II software manages EDA tool files and provides the following integration capabilities:

- Automatically generate files for synthesis and simulation and automatically launch other EDA tools (Assignments > Settings > EDA Tool Settings > NativeLink settings).
- Compile all RTL and gate-level simulation model libraries for your device, simulator, and design language automatically (Tools > Launch Simulation Library Compiler).
- Include files (.edf, .vqm) generated by other EDA design entry or synthesis tools in your project as synthesized design files (Project > Add/Remove File from Project)
- Automatically generate optional files for board-level verification (Assignments > Settings > EDA Tool Settings).

Figure 1–16. EDA Tool Settings
The Quartus II software optionally generates the following files for other EDA tools:

**Figure 1–17. Quartus II Generated Files for Other EDA Tools**


**Managing Team-based Projects**

The Quartus II software supports multiple designers, design iterations, and platforms. You can use the following techniques to preserve and track project changes in a team-based environment. These techniques may also be helpful for individual designers.

- Preserving Compilation Results
- Archiving Projects
- Using External Revision Control
- Migrating Projects Across Operating Systems

**Preserving Compilation Results**

The Quartus II software maintains a database of compilation results for each project revision. The databases files store results of incremental or full compilation. Do not edit these files directly. However, you can use the database files in the following ways:
■ Preserve compilation results for migration to a new version of the Quartus II software. Export a post-synthesis or post-fit, version-compatible database (Project > Export Database), and then import it into a newer version of the Quartus II software (Project > Import Database), or into another project.

■ Optimize and lock down the compilation results for individual blocks. Export the post-synthesis or post-fit netlist as a Quartus II Exported Partition File (.qxp) (Project > Export Design Partition). You can then import the partition as a new project design file.

■ Purge the content of the project database (Project > Clean Project) to remove unwanted previous compilation results at any time.

Factors Affecting Compilation Results
Changes to any of the following factors can impact compilation results:

■ Project Files—project settings (.qsf), design files, and timing constraints (.sdc)

■ Hardware—CPU architecture, not including hard disk or memory size differences. Windows XP x32 results are not identical to Windows XP x64 results. Linux x86 results is not identical to Linux x86_64.

■ Quartus II Software Version—including build number and installed patches. Click Help > About to obtain this information.

■ Operating System—Windows or Linux operating system, excluding version updates. For example, Windows XP, Windows Vista, and Windows 7 results are identical. Similarly, Linux RHEL, CentOS 4, and CentOS 5 results are identical.

Refer to Quartus II Incremental Compilation for Hierarchical and Team-Based Design and Design Planning for Partial Reconfiguration in the Quartus II Handbook for more information about partitions, incremental compilation, and device reconfiguration.

Migrating Results Across Quartus II Software Versions
To preserve compilation results for migration to a later version of the Quartus II software, export a version-compatible database file, and then import it into the later version of the Quartus II software. A few device families do not support version-compatible database generation, as indicated by project messages.

Exporting and Importing the Results Database
To save the compilation results in a version-compatible format for migration to a later version of the Quartus II software, follow these steps:

1. Open the project for migration in the original version of the Quartus II software.
2. Generate the project database and netlist with one of the following:
   ■ Click Processing > Start > Start Analysis & Synthesis to generate a post-synthesis netlist.
   ■ Click Processing > Start Compilation to generate a post-fit netlist.
3. Click Project > Export Database and specify the Export directory.
4. In a later version of the Quartus II software, click New Project Wizard and create a new project with the same top-level design entity name as the migrated project.
5. Click **Project > Import Database** and select the `<project directory>/export_db/` exported database directory. The Quartus II software opens the compiled project and displays compilation results.

---

You can turn on **Assignments > Settings > Compilation Process Settings > Export version-compatible database** if you want to always export the database following compilation.

---

**Figure 1–18. Quartus II Version-Compatible Database Structure**

---

**Cleaning the Project Database**

To clean the project database and remove all prior compilation results, follow these steps:

1. Click **Project > Clean Project**.
2. Select **All revisions** to remove the databases for all revisions of the current project, or specify a **Revision name** to remove only that revision’s database.
3. Click **OK**. A message indicates when the database is clean.

---

**Archiving Projects**

You can save the elements of a project in a single, compressed Quartus II Archive File (.qar) by clicking **Project > Archive Project**. The .qar captures logic design, project, and settings files required to restore the project. Use this technique to share projects between designers, or to transfer your project to a new version of the Quartus II software, or to Altera support.

You can optionally add compilation results, Qsys system files, and third-party EDA tool files to the archive. If you restore the archive in a different version of the Quartus II software, you must include the original .qdf in the archive to preserve original compilation results.

---

**Manually Adding Files To Archives**

To manually add files to an archive:

1. Click **Project > Archive Project** and specify the archive file name.
2. Click Advanced.

3. Select the File set for archive or select Custom. Turn on File subsets for archive.

4. Click Add and select Qsys system or EDA tool files, as detailed in Figure 1–15 and Figure 1–17. Click OK.

5. Click Archive.

Archiving Compilation Results

You can include compilation results in a project archive to avoid recompilation and preserve original results in the restored project. To archive compilation results, export the post-synthesis or post-fit version compatible database and include this file in the archive.

1. Export the project database as described in “Exporting and Importing the Results Database”.

2. Click Project > Archive Project and specify the archive file name.

3. Click Advanced.

4. Under File subsets, turn on Version-compatible database files and click OK.

5. Click Archive.

To restore an archive containing a version-compatible database, follow these steps:

1. Click Project > Restore Archived Project.

2. Select the archive name and destination folder and click OK.

3. After restoring the archived project, click Project > Import Database and import the version-compatible database.

Archiving Projects for Altera Service Requests

When archiving projects for an Altera service request, include all of the following file types for proper debugging by Altera Support:

To quickly identify and include appropriate archive files for an Altera service request:

1. Click Project > Archive Project and specify the archive file name.

2. Click Advanced.

3. In File set, select Service Request to include files for Altera Support.

   ■ Project source and setting files (.v, .vhd, .vqm, .qsf, .sdc, .qip, .qpf, .cmp, .sip)
   ■ Automatically detected source files (various)
   ■ Programming output files (.jdi, .sof, .pof)
   ■ Report files (.rpt, .pin, .summary, .smsg)
   ■ Qsys system and IP files (.qsys, .qip)

4. Click OK, and then click Archive.
Managing Team-based Projects

Your project may involve different team members with distributed responsibilities, such as sub-module design, device and system integration, simulation, and timing closure. In such cases, it may be useful to track and protect file revisions in an external revision control system.

While Quartus II project revisions preserve various project setting and constraint combinations, external revision control systems can also track and merge RTL source code, simulation testbenches, and build scripts. External revision control supports design file version experimentation through branching and merging different versions of source code from multiple designers. Refer to your external revision control documentation for setup information.

**Files to Include In External Revision Control**

Include the following Quartus II project file types in external revision control systems:

- Logic design files (.v, .vhdl, .bdf, .edf, .vqm)
- Timing constraint files (.sdc)
- Quartus II project settings and constraints (.qdf, .qpf, .qsf)
- MegaWizard-generated IP files (.v, .sv, .vhdl, .qip, .sip)
- Qsys-generated files (.qsys, .qip, .sip)
- EDA tool files (.vo, .vho)

You can generate or modify these files manually if you use a scripted design flow. If you use an external source code control system, you can check-in project files anytime you modify assignments and settings in the Quartus II software. Refer to Figure 1–15 for a list of IP and Qsys generated files.
Migrating Projects Across Operating Systems

Consider the following cross-platform issues when moving your project from one operating system to another (for example, from Windows to Linux).

Migrating Design Files and Libraries

Consider the following file naming differences when migrating projects across operating systems:

- Use appropriate case for your platform in file path references.
- Use a character set common to both platforms.
- Do not change the forward-slash (/) and back-slash (\) path separators in the .qsf. The Quartus II software automatically changes all back-slash (\) path separators to forward-slashes (/) in the .qsf.
- Observe the target platform’s file name length limit.
- Use underscore instead of spaces in file and directory names.
- Change library absolute path references to relative paths in the .qsf.
- Ensure that any external project library exists in the new platform’s file system.
- Specify file and directory paths as relative to the project directory. For example, for a project titled foo_design, specify the source files as: top.v, foo_folder/foo1.v, foo_folder/foo2.v, and foo_folder/bar_folder/bar1.vhdl.
- Ensure that all the subdirectories are in the same hierarchical structure and relative path as in the original platform.

Figure 1–20. All Inclusive Project Directory Structure
Relative Paths

Express file paths using relative path notation (./). For example, in the directory structure shown in Figure 1–21, you can specify top.v as ./source/top.v and foo1.v as ../../source/foo_folder/foo1.v.

Figure 1–21. Quartus II Project Directory Separate from Design Files

Migrating Design Libraries

The following guidelines apply to library migration across computing platforms:

- The project directory takes precedence over the project libraries.
- For Linux, the Quartus II software creates the file in the altera.quartus directory under the <home> directory.
- All library files are relative to the libraries. For example, if you specify the user_lib1 directory as a project library and you want to add the /user_lib1/foo1.v file to the library, you can specify the foo1.v file in the .qsf as foo1.v. The Quartus II software includes files in specified libraries.
- If the directory is outside of the project directory, an absolute path is created by default. Change the absolute path to a relative path before migration.
- When copying projects that include libraries, you must either copy your project library files along with the project directory or ensure that your project library files exist in the target platform.
- On Windows, the Quartus II software searches for the quartus2.ini file in the following directories and order:
  a. USERPROFILE, for example, C:\Documents and Settings\<user name>
  b. Directory specified by the TMP environmental variable
  c. Directory specified by the TEMP environmental variable
  d. Root directory, for example, C:\
Scripting API

You can use command-line executables or scripts to execute project commands, rather than using the GUI. The following commands are available for scripting project management.

Scripting Project Settings

You can use a Tcl script to specify settings and constraints, rather than using the GUI. This can be helpful if you have many settings and wish to track them in a single file or spreadsheet for iterative comparison. The .qsf supports only a limited subset of Tcl commands. Therefore, pass settings and constraints using a Tcl script:

1. Create a text file with the extension .tcl that contains your assignments in Tcl format.
2. Source the Tcl script file by adding the following line to the .qsf:
   ```tcl
   set_global_assignment -name SOURCE_TCL_SCRIPT_FILE <file name>.
   ```

Project Revision Commands

Use the following commands for scripting project revisions.

Create Revision Command

-`-based_on` and `-set_current` are optional. `-copy_results` copies results from “based_on” revision.

**Example 1–1. Create Revisions from other Revision and Set as Current Revision**

```bash
create_revision <name> -based_on <name> -set_current
```

Set Current Revision Command

The `-force` option enables you to open the revision that you specify under revision name and overwrite the compilation database if the database version is incompatible.

**Example 1–2. Specify the Current Revision**

```bash
set_current_revision -force <revision name>
```

Get Project Revisions Command

**Example 1–3. Get List of Revisions in Open Project**

```bash
get_project_revisions <project_name>
```

Delete Revision Command

**Example 1–4. Delete a Revision**

```bash
delete_revision <revision name>
```
**Project Archive Commands**

Use the following commands for scripting project archives.

**Project Archive Command**

**Example 1–5. Overwrite the Project Archive with New Archive with Default Settings**

```
project_archive <name>.qar -overwrite
```

You can specify the following other options:

- `all_revisions`
- `include_libraries`
- `include_outputs`
- `use_file_set <file_set>`
- `version_compatible_database`

Version-compatible databases are not be available for some device families. If you require the database files to reproduce the compilation results in the same Quartus II software version, use the `-use_file_set full_db` option to archive the complete database.

**Example 1–6. Create a Project Archive**

```
quartus_sh --archive <name>
```

**Project Restore Commands**

Use the following commands for scripting project restore.

**Example 1–7. Restore a Project Archive**

```
project_restore <name>.qar -destination restored -overwrite
```

**Example 1–8. Restore a Project Archive**

```
quartus_sh --restore archive.qar
```

**Project Database Commands**

Use the following commands for scripting project database import and export.

**Export and Import Database Commands**

**Example 1–9. Import and Export Version-Compatible Databases**

```
export_database <directory>
import_database <directory>
```
Example 1–10 shows the Tcl commands from the flow package to import or export version-compatible databases. If you use the flow package, you must specify the database directory variable name. flow and database_manager packages contain commands to manage version-compatible databases.

**Example 1–10. Import and Export Version-Compatible Databases from flow Package**

```tcl
set_global_assignment -name VER_COMPATIBLE_DB_DIR <directory>
execute_flow -flow export_database
execute_flow -flow import_database
```

**Example 1–11. Generate Version-Compatible Databases After Every Compilation**

```tcl
set_global_assignment -name AUTO_EXPORT_VER_COMPATIBLE_DB ON
set_global_assignment -name VER_COMPATIBLE_DB_DIR <directory>
```

Example 1–12 shows the quartus_cdb and the quartus_sh executables to manage version-compatible databases:

**Example 1–12. quartus_cdb and quartus_sh Executable**

```bash
quartus_cdb <project> -c <revision> --export_database=<directory>
quartus_cdb <project> -c <revision> --import_database=<directory>
quartus_sh --flow export_database <project> -c <revision>
quartus_sh --flow import_database <project> -c <revision>
```

Example 1–13 archives the version-compatible database with the project for restoration in the same version of the Quartus II software:

**Example 1–13. Archive Compilation Database with Project**

```bash
quartus_sh --archive -use_file_set full_db [-revision <revisionname>] <project_name>
```

**Project Library Commands**

In Tcl, use commands in the :quartus::project package to specify project libraries. To specify project libraries, use the set_global_assignment command.

**Example 1–14. Specify Project Libraries with SEARCH_PATH Assignment**

```tcl
set_global_assignment -name SEARCH_PATH *./other_dir/library1
set_global_assignment -name SEARCH_PATH *./other_dir/library2
set_global_assignment -name SEARCH_PATH *./other_dir/library3
```

To report any project libraries specified for a project and any global libraries specified for the current installation of the Quartus II software, use the get_global_assignment and get_user_option Tcl commands.
Example 1–15 shows that the Tcl script outputs the user paths and global libraries for an open Quartus II project:


```
get_global_assignment -name SEARCH_PATH
get_user_option -name SEARCH_PATH
```

For more information about scripting, refer to *Tcl Scripting* or *Command-Line Scripting* in the *Quartus II Handbook*. For comprehensive scripting reference, refer to the *Quartus II Settings File Manual*.

### Document Revision History

Table 1–3 shows the revision history for this chapter.

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<td>May 2013</td>
<td>13.0.0</td>
<td>■ Overhaul for improved usability and updated information.</td>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Removed survey link.</td>
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<td>■ Moved “Hiding Messages” to Help.</td>
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<td>■ Removed references about the set_user_option command.</td>
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<td>■ Updated “Working with Messages” on page 4–17.</td>
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<td>■ Added a link to Help.</td>
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<td>■ Removed Figure 4–2 on page 4–7, Figure 4–11 on page 23, and Figure 4–12 on page.</td>
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<td>■ Updated “Specifying Libraries” on page 4–14 section.</td>
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<td>■ Changed “User Libraries” to “Libraries”.</td>
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<td>■ Added “Managing Projects in a Team-Based Design Environment” on page 4–22 and “File Association” on page 4–2.</td>
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<td>■ Updated Figure 4–1 on page 4–6, Figure 4–2 on page 4–8, Figure 4–6 on page 4–18, Figure 4–6 on page 4–19, and Figure 4–7 on page 4–21.</td>
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<td>■ Added “Quartus II Text Editor” on page 4–2, “Reducing Compilation Time” on page 4–32.</td>
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<td>■ Updated Figure 4–4 on page 4–9, Figure 4–7 on page 4–19.</td>
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Table 1–3. Document Revision History (Part 2 of 2)

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<td>- Updated Example 4–3, Example 4–4, Example 4–5, and Example 4–6.</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter discusses key FPGA design planning considerations, provides recommendations, and describes various tools available for you to improve your design productivity with Altera® FPGAs.

Because of the significant increase in FPGA device densities, designs are complex and can sometimes involve multiple designers. System architects must also resolve design issues when integrating design blocks. However, you can solve potential problems early in the design cycle by following the design planning considerations in this chapter.

This chapter contains the following sections:

- “Creating Design Specifications” on page 2–2
- “Selecting Intellectual Property” on page 2–2
- “Using Qsys and Standard Interfaces in System Design” on page 2–3
- “Selecting a Device” on page 2–3
- “Planning for Device Programming or Configuration” on page 2–5
- “Estimating Power” on page 2–5
- “Early Pin Planning and I/O Analysis” on page 2–6
- “Selecting Third-Party EDA Tools” on page 2–8
- “Planning for On-Chip Debugging Tools” on page 2–10
- “Design Practices and HDL Coding Styles” on page 2–11
- “Planning for Hierarchical and Team-Based Design” on page 2–13
- “Fast Synthesis and Early Timing Estimation” on page 2–15

This chapter provides only an introduction to various design planning features in the Quartus® II software. For more information about Quartus II features and methodologies, this chapter provides references to other appropriate chapters in the Quartus II Handbook.

Before reading the design planning guidelines discussed in this chapter, consider your design priorities. More device features, density, or performance requirements can increase system cost. Signal integrity and board issues can impact I/O pin locations. Power, timing performance, and area utilization all affect each other, and compilation time is affected when optimizing these priorities.
Creating Design Specifications

Before you create your design logic or complete your system design, create detailed design specifications that define the system, specify the I/O interfaces for the FPGA, identify the different clock domains, and include a block diagram of basic design functions.

In addition, creating a test plan helps you to design for verification and manufacturability. For example, you might need to validate interfaces incorporated in your design. To perform any built-in self-test functions to drive interfaces, you can use a UART interface with a Nios® II processor inside the FPGA device. For guidelines related to analyzing and debugging the device after it is in the system, refer to “Planning for On-Chip Debugging Tools” on page 2–10.

If more than one designer works on your design, you must consider a common design directory structure or source control system to make design integration easier. For more suggestions on team-based designs, refer to “Planning for Hierarchical and Team-Based Design” on page 2–13. Consider whether you want to standardize on an interface protocol for each design block. To improve reusability and ease of integration, refer to “Using Qsys and Standard Interfaces in System Design”.

Selecting Intellectual Property

Altera and its third-party intellectual property (IP) partners offer a large selection of standardized IP cores optimized for Altera devices. The IP you select often affects system design, especially if the FPGA interfaces with other devices in the system. Consider which I/O interfaces or other blocks in your system design are implemented using IP cores, and plan to incorporate these cores in your FPGA design.

The OpenCore Plus feature, which is available for many IP cores, allows you to program the FPGA to verify your design in the hardware before you purchase the IP license. The evaluation supports the following modes:

- Untethered—the design runs for a limited time.
- Tethered—the design requires an Altera serial JTAG cable connected between the JTAG port on your board and a host computer running the Quartus II Programmer for the duration of the hardware evaluation period.

For descriptions of available IP cores, refer to the Intellectual Property page of the Altera website.
Using Qsys and Standard Interfaces in System Design

You can use the Quartus II Qsys system integration tool to create your design with fast and easy system-level integration. With Qsys, you can specify system components in a GUI and generate the required interconnect logic automatically, along with adapters for clock crossing and width differences. Because system design tools change the design entry methodology, you must plan to start developing your design within the tool. Ensure all design blocks use appropriate standard interfaces from the beginning of the design cycle so that you do not need to make changes later.

Qsys components use Avalon® standard interfaces for the physical connection of components, and you can connect any logical device (either on-chip or off-chip) that has an Avalon interface. The Avalon Memory-Mapped interface allows a component to use an address mapped read or write protocol that enables flexible topologies for connecting master components to any slave components. The Avalon Streaming interface enables point-to-point connections between streaming components that send and receive data using a high-speed, unidirectional system interconnect between source and sink ports.

In addition to enabling the use of a system integration tool such as Qsys, using standard interfaces ensures compatibility between design blocks from different design teams or vendors. Standard interfaces simplify the interface logic to each design block and enable individual team members to test their individual design blocks against the specification for the interface protocol to ease system integration.

For more information about using Qsys to improve your productivity, refer to the System Design with Qsys section in volume 1 of the Quartus II Handbook.

Qsys replaces the SOPC Builder system integration tool for new designs. For more information about SOPC Builder, refer to the SOPC Builder User Guide.

Selecting a Device

The device you choose affects board specification and layout. This section provides guidelines in the device selection process.

Choose the device family that best suits your design requirements. Families differ in cost, performance, logic and memory density, I/O density, power utilization, and packaging. You must also consider feature requirements, such as I/O standards support, high-speed transceivers, global or regional clock networks, and the number of phase-locked loops (PLLs) available in the device.

You can use the Altera Product Selector available on the Altera website to help you choose your device. You can also review important features of each device family in the Selector Guides page of the Altera website. Each device family also has a device handbook, including a data sheet, which documents device features in detail. You can also see a summary of the resources for each device in the Device dialog box in the Quartus II software.

For a list of device selection guides, refer to Devices and Adapters in Quartus II Help.
Carefully study the device density requirements for your design. Devices with more logic resources and higher I/O counts can implement larger and more complex designs, but at a higher cost. Smaller devices use lower static power. Select a device larger than what your design requires if you want to add more logic later in the design cycle to upgrade or expand your design, and reserve logic and memory for on-chip debugging (refer to “Planning for On-Chip Debugging Tools” on page 2–10). Consider requirements for types of dedicated logic blocks, such as memory blocks of different sizes, or digital signal processing (DSP) blocks to implement certain arithmetic functions.

If you have older designs that target an Altera device, you can use their resources as an estimate for your design. Compile existing designs in the Quartus II software with the Auto device selected by the Fitter option in the Settings dialog box. Review the resource utilization to learn which device density fits your design. Consider coding style, device architecture, and the optimization options used in the Quartus II software, which can significantly affect the resource utilization and timing performance of your design.

To obtain resource utilization estimates for certain configurations of Altera’s IP, refer to the user guides for Altera megafUNCTIONS and IP MegaCores on the IP and MegafUNCTIONS literature page of the Altera website.

Device Migration Planning

Determine whether you want to migrate your design to another device density to allow flexibility when your design nears completion, or whether you want to migrate to a HardCopy® ASIC when your design reaches volume production. You may want to target a smaller (and less expensive) device and then move to a larger device if necessary to meet your design requirements. Other designers may prototype their design in a larger device to reduce optimization time and achieve timing closure more quickly, and then migrate to a smaller device after prototyping. Similarly, many designers compile and optimize their design for an FPGA device and then migrate to a HardCopy ASIC when the design is complete and ready for higher-volume production. If you want the flexibility to migrate your design, you must specify these migration options in the Quartus II software at the beginning of your design cycle.

For more information about specifying the target migration devices, refer to Specifying Devices for Device Migration in Quartus II Help.

Selecting a migration device impacts pin placement because some pins may serve different functions in different device densities or package sizes. If you make pin assignments in the Quartus II software, the Pin Migration View in the Pin Planner highlights pins that change function between your migration devices. (For more information, refer to “Early Pin Planning and I/O Analysis” on page 2–6.) Selecting a companion device might restrict logic utilization to ensure that your design is compatible with a selected HardCopy device. Adding migration or companion devices later in the design cycle is possible, but requires extra effort to check pin assignments, and might require design changes to fit into the new target device. Consider these issues early in the design cycle rather than at the end, when your design is near completion and ready for migration.
Additionally, if you plan to migrate your design to a HardCopy ASIC, review HardCopy guidelines early in the design cycle for any Quartus II settings that you must use or other restrictions that you must consider. You must use complete timing constraints if you want to migrate to a HardCopy ASIC because of the rigorous verification requirements for ASIC devices.

For more information about timing requirements and analysis for HardCopy designs, refer to the HardCopy Series Handbook, and the Quartus II Support for HardCopy Series Devices chapter in volume 1 of the Quartus II Handbook.

Planning for Device Programming or Configuration

Planning how to program or configure the device in your system allows system and board designers to determine what companion devices, if any, your system requires. Your board layout also depends on the type of programming or configuration method you plan to use for programmable devices. Many programming options require a JTAG interface to connect to the devices, so you might have to set up a JTAG chain on the board. Additionally, the Quartus II software uses the settings for the configuration scheme, configuration device, and configuration device voltage to enable the appropriate dual purpose pins as regular I/O pins after you complete configuration. The Quartus II software performs voltage compatibility checks of those pins during I/O assignment analysis and compilation of your design. You can use the Configuration tab of the Device and Pin Options dialog box to select your configuration scheme.

The device family handbooks describe the configuration options available for a device family. For more details about configuration options, refer to the Configuration Handbook. For information about programming CPLD devices, refer to your device data sheet or handbook.

Estimating Power

You can use the Quartus II power estimation and analysis tools to provide information to PCB board and system designers. Power consumption in FPGA devices depends on the design logic, which can make planning difficult. You can estimate power before you create any source code, or when you have a preliminary version of the design source code, and then perform the most accurate analysis with the PowerPlay Power Analyzer when you complete your design.

You must accurately estimate device power consumption to develop an appropriate power budget and to design the power supplies, voltage regulators, heat sink, and cooling system. Power estimation and analysis helps you satisfy two important planning requirements:

- **Thermal**—ensure that the cooling solution is sufficient to dissipate the heat generated by the device. The computed junction temperature must fall within normal device specifications.
- **Power supply**—ensure that the power supplies provide adequate current to support device operation.

The PowerPlay Early Power Estimator (EPE) spreadsheet allows you to estimate power utilization for your design.
You can manually enter data into the EPE spreadsheet, or use the Quartus II software to generate device resource information for your design.

To manually enter data into the EPE spreadsheet, enter the device resources, operating frequency, toggle rates, and other parameters for your design. If you do not have an existing design, estimate the number of device resources used in your design, and then enter the data into the EPE spreadsheet manually.

If you have an existing design or a partially completed design, you can use the Quartus II software to generate the PowerPlay Early Power Estimator File (.txt, .csv) to assist you in completing the PowerPlay EPE spreadsheet.

For more information about generating the PowerPlay EPE File, refer to *Performing an Early Power Estimate Using the PowerPlay Early Power Estimator* in Quartus II Help.

The PowerPlay EPE spreadsheet includes the Import Data macro that parses the information in the PowerPlay EPE File and transfers the information into the spreadsheet. If you do not want to use the macro, you can manually transfer the data into the EPE spreadsheet. For example, after importing the PowerPlay EPE File information into the PowerPlay EPE spreadsheet, you can add device resource information. If the existing Quartus II project represents only a portion of your full design, manually enter the additional device resources you use in the final design.

Estimating power consumption early in the design cycle allows planning of power budgets and avoids unexpected results when designing the PCB.

The PowerPlay EPE spreadsheets for each supported device family are available on the PowerPlay Early Power Estimator and Power Analyzer page of the Altera website.

When you complete your design, perform a complete power analysis to check the power consumption more accurately. The PowerPlay Power Analyzer tool in the Quartus II software provides an accurate estimation of power, ensuring that thermal and supply limitations are met.

For more information about power estimation and analysis, refer to the *PowerPlay Power Analysis* chapter in volume 3 of the *Quartus II Handbook*.

**Early Pin Planning and I/O Analysis**

This section describes early pin planning and I/O analysis features for different stages of the design flow.

In many design environments, FPGA designers want to plan the top-level FPGA I/O pins early to help board designers begin the PCB design and layout. The I/O capabilities and board layout guidelines of the FPGA device influence pin locations and other types of assignments. If the board design team specifies an FPGA pin-out, the pin locations must be verified in the FPGA placement and routing software to avoid board design changes.
You can create a preliminary pin-out for an Altera FPGA with the Quartus II Pin Planner before you develop the source code, based on standard I/O interfaces (such as memory and bus interfaces) and any other I/O requirements for your system. The Quartus II I/O Assignment Analysis checks that the pin locations and assignments are supported in the target FPGA architecture. You can then use I/O Assignment Analysis to validate I/O-related assignments that you create or modify throughout the design process. When you compile your design in the Quartus II software, I/O Assignment Analysis runs automatically in the Fitter to validate that the assignments meet all the device requirements and generates error messages.

Early in the design process, before creating the source code, the system architect has information about the standard I/O interfaces (such as memory and bus interfaces), the IP cores in your design, and any other I/O-related assignments defined by system requirements. You can use this information with the Early Pin Planning feature in the Pin Planner to specify details about the design I/O interfaces. You can then create a top-level design file that includes all I/O information.

The Pin Planner interfaces with the IP core parameter editor, which allows you to create or import custom megafunctions and IP cores that use I/O interfaces. You can configure how to connect the functions and cores to each other by specifying matching node names for selected ports. You can create other I/O-related assignments for these interfaces or other design I/O pins in the Pin Planner, as described in this section. The Pin Planner creates virtual pin assignments for internal nodes, so internal nodes are not assigned to device pins during compilation. After analysis and synthesis of the newly generated top-level wrapper file, use the generated netlist to perform I/O Analysis with the Start I/O Assignment Analysis command.

For more information about setting up the nodes in your design, refer to Set Up Top-Level Design File Window (Edit Menu) in Quartus II Help.

You can use the I/O analysis results to change pin assignments or IP parameters even before you create your design, and repeat the checking process until the I/O interface meets your design requirements and passes the pin checks in the Quartus II software. When you complete initial pin planning, you can create a revision based on the Quartus II-generated netlist. You can then use the generated netlist to develop the top-level design file for your design, or disregard the generated netlist and use the generated Quartus II Settings File (.qsf) with your design.

During this early pin planning, after you have generated a top-level design file, or when you have developed your design source code, you can assign pin locations and assignments with the Pin Planner.

With the Pin Planner, you can identify I/O banks, voltage reference (VREF) groups, and differential pin pairings to help you through the I/O planning process. If migration devices are selected (including HardCopy devices) as described in “Device Migration Planning” on page 2–4, the Pin Migration View highlights the pins that have changed functions in the migration device when compared to the currently selected device. Selecting the pins in the Device Migration view cross-probes to the rest of the Pin Planner, so that you can use device migration information when planning your pin assignments. You can also configure board trace models of selected pins for use in “board-aware” signal integrity reports generated with the Enable
Advanced I/O Timing option. This option ensures that you get accurate I/O timing analysis. You can use a Microsoft Excel spreadsheet to start the I/O planning process if you normally use a spreadsheet in your design flow, and you can export a Comma-Separated Value File (.csv) containing your I/O assignments for spreadsheet use when you assign all pins.

When you complete your pin planning, you can pass pin location information to PCB designers. The Pin Planner is tightly integrated with certain PCB design EDA tools, and can read pin location changes from these tools to check suggested changes. Your pin assignments must match between the Quartus II software and your schematic and board layout tools to ensure the FPGA works correctly on the board, especially if you must make changes to the pin-out. The system architect uses the Quartus II software to pass pin information to team members designing individual logic blocks, allowing them to achieve better timing closure when they compile their design.

Start FPGA planning before you complete the HDL for your design to improve the confidence in early board layouts, reduce the chance of error, and improve the overall time to market of the design. When you complete your design, use the Fitter reports for the final sign-off of pin assignments. After compilation, the Quartus II software generates the Pin-Out File (.pin), and you can use this file to verify that each pin is correctly connected in board schematics.

For more information about I/O assignment and analysis, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook. For more information about passing I/O information between the Quartus II software and third-party EDA tools, refer to the Mentor Graphics PCB Design Tools Support and Cadence PCB Design Tools Support chapters in the I/O and PCB Tools section in volume 2 of the Quartus II Handbook.

Simultaneous Switching Noise Analysis

Simultaneous switching noise (SSN) is a noise voltage inducted onto a victim I/O pin of a device due to the switching behavior of other aggressor I/O pins in the device. Altera provides tools for SSN analysis and estimation, including SSN characterization reports, an Early SSN Estimator (ESE) spreadsheet tool, and the SSN Analyzer in the Quartus II software. SSN often leads to the degradation of signal integrity by causing signal distortion, thereby reducing the noise margin of a system. You must address SSN with estimation early in your system design, to minimize later board design changes. When your design is complete, verify your board design by performing a complete SSN analysis of your FPGA in the Quartus II software.

For more information and device support for the ESE spreadsheet tool, refer to Altera’s Signal Integrity Center on the Altera website. For more information about the SSN Analyzer, refer to the Simultaneous Switching Noise (SSN) Analysis and Optimizations chapter in volume 2 of the Quartus II Handbook.

Selecting Third-Party EDA Tools

Your complete FPGA design flow may include third-party EDA tools in addition to the Quartus II software. Determine which tools you want to use with the Quartus II software to ensure that they are supported and set up properly, and that you are aware of their capabilities.
Synthesis Tool

The Quartus II software includes integrated synthesis that supports Verilog HDL, VHDL, Altera Hardware Description Language (AHDL), and schematic design entry. You can also use supported standard third-party EDA synthesis tools to synthesize your Verilog HDL or VHDL design, and then compile the resulting output netlist file in the Quartus II software. Different synthesis tools may give different results for each design. To determine the best tool for your application, you can experiment by synthesizing typical designs for your application and coding style. Perform placement and routing in the Quartus II software to get accurate timing analysis and logic utilization results.

Because tool vendors frequently add new features, fix tool issues, and enhance performance for Altera devices, you must use the most recent version of third-party synthesis tools. The Quartus II Software Release Notes lists the version of each synthesis tool that is supported by a given version of the Quartus II software.

The synthesis tool you choose may allow you to create a Quartus II project and pass constraints, such as the EDA tool setting, device selection, and timing requirements that you specified in your synthesis project. You can save time when setting up your Quartus II project for placement and routing.

To use incremental compilation, you must partition your design for synthesis and generate multiple output netlist files. For more information, refer to “Incremental Compilation with Design Partitions” on page 2–14.

For more information about synthesis tool flows, refer to Volume 1: Design and Synthesis of the Quartus II Handbook.

Simulation Tool

Altera provides the Mentor Graphics ModelSim®-Altera Starter Edition with the Quartus II software. You can also purchase the ModelSim-Altera Edition or a full license of the ModelSim software to support large designs and achieve faster simulation performance. The Quartus II software can generate both functional and timing netlist files for ModelSim and other third-party simulators.

Use the simulator version that your Quartus II software version supports for best results. You must also use the model libraries provided with your Quartus II software version. Libraries can change between versions, which might cause a mismatch with your simulation netlist.

For a list of the version of each simulation tool that is supported with a given version of the Quartus II software, refer to the Quartus II Software Release Notes.

For more information about simulation tool flows, refer to the appropriate chapter in the Simulation section in volume 3 of the Quartus II Handbook.

Formal Verification Tool

Consider whether the Quartus II software supports the formal verification tool that you want to use, and whether the flow impacts your design and compilation stages of your design.
For more information about formal verification flows and the supported tools, refer to Volume 3: Verification of the Quartus II Handbook.

Using a formal verification tool can impact performance results because performing formal verification requires turning off certain logic optimizations, such as register retiming, and forces you to preserve hierarchy blocks, which can restrict optimization. Formal verification treats memory blocks as black boxes. Therefore, you must keep memory in a separate hierarchy block so other logic does not get incorporated into the black box for verification. Other restrictions may limit your design, and you must consult Volume 3: Verification of the Quartus II Handbook for details. If formal verification is important to your design, plan for limitations and restrictions at the beginning of the design cycle rather than make changes later.

Planning for On-Chip Debugging Tools

In-system debugging tools offer different advantages and trade-offs. A particular debugging tool may work better for different systems and designers. You must evaluate on-chip debugging tools early in your design process, to ensure that your system board, Quartus II project, and design can support the appropriate tools. You can reduce debugging time and avoid making changes to accommodate your preferred debugging tools later.

For an overview of debugging tools that can help you decide which tools to use, refer to the System Debugging Tools Overview chapter in volume 3 of the Quartus II Handbook.

If you intend to use any of these tools, you may have to plan for the tools when developing your system board, Quartus II project, and design. Consider the following debugging requirements when you plan your design:

- **JTAG connections**—required to perform in-system debugging with JTAG tools. Plan your system and board with JTAG ports that are available for debugging.

- **Additional logic resources**—required to implement JTAG hub logic. If you set up the appropriate tool early in your design cycle, you can include these device resources in your early resource estimations to ensure that you do not overload the device with logic.

- **Reserve device memory**—required if your tool uses device memory to capture data during system operation. To ensure that you have enough memory resources to take advantage of this debugging technique, consider reserving device memory to use during debugging.

- **Reserve I/O pins**—required if you use the Logic Analyzer Interface (LAI) or SignalProbe tools, which require I/O pins for debugging. If you reserve I/O pins for debugging, you do not have to later change your design or board. The LAI can multiplex signals with design I/O pins if required. Ensure that your board supports a debugging mode, in which debugging signals do not affect system operation.

- **Instantiate a megafunction in your HDL code**—required if your debugging tool uses a Quartus II megafunction.
Design Practices and HDL Coding Styles

When you develop complex FPGA designs, design practices and coding styles have an enormous impact on the timing performance, logic utilization, and system reliability of your device.

Design Recommendations

Use synchronous design practices to consistently meet your design goals. Problems with asynchronous design techniques include reliance on propagation delays in a device, incomplete timing analysis, and possible glitches. In a synchronous design, a clock signal triggers all events. When you meet all register timing requirements, a synchronous design behaves in a predictable and reliable manner for all process, voltage, and temperature (PVT) conditions. You can easily target synchronous designs to different device families or speed grades.

Table 2–1 lists which factors are important for each debugging tool.

Table 2–1. Factors to Consider When Using Debugging Tools During Design Planning Stages

<table>
<thead>
<tr>
<th>Design Planning Factor</th>
<th>SignapTap II Logic Analyzer</th>
<th>System Console</th>
<th>In-System Memory Content Editor</th>
<th>Logic Analyzer Interface (LAI)</th>
<th>SignalProbe</th>
<th>In-System Sources and Probes</th>
<th>Virtual JTAG Megafunction</th>
</tr>
</thead>
<tbody>
<tr>
<td>JTAG connections</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>Additional logic resources</td>
<td>—</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓</td>
</tr>
<tr>
<td>Reserve device memory</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>Reserve I/O pins</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>Instantiate a megafunction in your HDL code</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓</td>
</tr>
</tbody>
</table>

For more information, refer to Design Debugging Using the SignalTap II Logic Analyzer chapter in volume 3 of the Quartus II Handbook.
Clock signals have a large effect on the timing accuracy, performance, and reliability of your design. Problems with clock signals can cause functional and timing problems in your design. Use dedicated clock pins and clock routing for best results, and if you have PLLs in your target device, use the PLLs for clock inversion, multiplication, and division. For clock multiplexing and gating, use the dedicated clock control block or PLL clock switchover feature instead of combinational logic, if these features are available in your device. If you must use internally-generated clock signals, register the output of any combinational logic used as a clock signal to reduce glitches.

The Design Assistant in the Quartus II software is a design-rule checking tool that enables you to verify design issues. The Design Assistant checks your design for adherence to Altera-recommended design guidelines. You can also use third-party lint tools to check your coding style.

For more information about running the Design Assistant, refer to About the Design Assistant in Quartus II Help.

Consider the architecture of the device you choose so that you can use specific features in your design. For example, the control signals should use the dedicated control signals in the device architecture. Sometimes, you might need to limit the number of different control signals used in your design to achieve the best results.

For more information about design recommendations and using the Design Assistant, refer to the Recommended Design Practices chapter in volume 1 of the Quartus II Handbook. You can also refer to industry papers for more information about multiple clock design. For a good analysis, refer to Synthesis and Scripting Techniques for Designing Multi-Asynchronous Clock Designs under Papers (www.sunburst-design.com).

**Recommended HDL Coding Styles**

HDL coding styles can have a significant effect on the quality of results for programmable logic designs. If you design memory and DSP functions, you must understand the target architecture of your device so you can use the dedicated logic block sizes and configurations. Follow the coding guidelines for inferring megafunctions and targeting dedicated device hardware, such as memory and DSP blocks.

For HDL coding examples and recommendations, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. For additional tool-specific guidelines, refer to the documentation of your synthesis tool.

**Managing Metastability**

Metastability problems can occur in digital design when a signal is transferred between circuitry in unrelated or asynchronous clock domains, because the designer cannot guarantee that the signal meets the setup and hold time requirements during the signal transfer. Designers commonly use a synchronization chain to minimize the occurrence of metastable events. Ensure that your design accounts for synchronization between any asynchronous clock domains. Consider using a synchronizer chain of more than two registers for high-frequency clocks and frequently-toggling data signals to reduce the chance of a metastability failure.
You can use the Quartus II software to analyze the average mean time between failures (MTBF) due to metastability when a design synchronizes asynchronous signals, and optimize your design to improve the metastability MTBF. The MTBF due to metastability is an estimate of the average time between instances when metastability could cause a design failure. A high MTBF (such as hundreds or thousands of years between metastability failures) indicates a more robust design. Determine an acceptable target MTBF given the context of your entire system and the fact that MTBF calculations are statistical estimates.

The Quartus II software can help you determine whether you have enough synchronization registers in your design to produce a high enough MTBF at your clock and data frequencies.

For information about metastability analysis, reporting, and optimization features in the Quartus II software, refer to the Managing Metastability with the Quartus II Software chapter in volume 1 of the Quartus II Handbook.

Planning for Hierarchical and Team-Based Design

To create a hierarchical design so that you can use compilation-time savings and performance preservation with the Quartus II software incremental compilation feature, plan for an incremental compilation flow from the beginning of your design cycle. The following subsections describe the flat compilation flow, in which the design hierarchy is flattened without design partitions, and then the incremental compilation flow that uses design partitions. Incremental compilation flows offer several advantages, but require more design planning to ensure effective results. The last subsections discuss planning an incremental compilation flow, planning design partitions, and optionally creating a design floorplan.

For information about using the incremental compilation flow methodology in the Quartus II software, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Flat Compilation Flow with No Design Partitions

In the flat compilation flow with no design partitions in the Quartus II software, the Quartus II software compiles the entire design in a “flat” netlist. Your source code can have hierarchy, but the Quartus II software flattens your design during compilation and synthesizes all the design source code and fits in the target device whenever the software recompile your design after any change in your design. By processing the entire design, the software performs all available logic and placement optimizations on the entire design to improve area and performance. You can use debugging tools in an incremental design flow, such as the SignalTap II Logic Analyzer, but you do not specify any design partitions to preserve design hierarchy during compilation.
The flat compilation flow is easy to use; you do not have to plan any design partitions. However, because the Quartus II software recompiles the entire design whenever you change your design, compilation times can be slow for large devices. Additionally, you may find that the results for one part of the design change when you change a different part of your design. You can turn on the Rapid Re compile option to instruct the software to preserve compatible placement and routing results when the design changes in subsequent compilations. This option can reduce your compilation time in a flat or partitioned design when you make small changes to your design.

**Incremental Compilation with Design Partitions**

In an incremental compilation flow, the system architect splits a large design into partitions. When hierarchical design partitions are well chosen and placed in the device floorplan, you can speed up your design compilation time while maintaining the quality of results.

Incremental compilation preserves the compilation results and performance of unchanged partitions in the design, greatly reducing design iteration time by focusing new compilations on changed design partitions only. Incremental compilation then merges new compilation results with the previous compilation results from unchanged design partitions. Additionally, you can target optimization techniques, such as physical synthesis, to specific design partitions while leaving other partitions unchanged. You can also use empty partitions to indicate that parts of your design are incomplete or missing, while you compile the rest of your design.

Third-party IP designers can also export logic blocks to be integrated into the top-level design. Team members can work on partitions independently, which can simplify the design process and reduce compilation time. With exported partitions, the system architect must provide guidance to designers or IP providers to ensure that each partition uses the appropriate device resources. Because the designs may be developed independently, each designer has no information about the overall design or how their partition connects with other partitions. This lack of information can lead to problems during system integration. The top-level project information, including pin locations, physical constraints, and timing requirements, must be communicated to the designers of lower-level partitions before they start their design.

The system architect plans design partitions at the top level and allows third-party designs to access the top-level project framework. By designing in a copy of the top-level project (or by checking out the project files in a source control environment), the designers of the lower-level block have full information about the entire project, which helps to ensure optimal results.

When you plan your design code and hierarchy, ensure that each design entity is created in a separate file so that the entities remain independent when you make source code changes in the file. If you use a third-party synthesis tool, create separate Verilog Quartus Mapping or EDIF netlists for each design partition in your synthesis tool. You may have to create separate projects in your synthesis tool, so that the tool synthesizes each partition separately and generates separate output netlist files. The netlists are then considered the source files for incremental compilation.

For more information about support for Quartus II incremental compilation, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter of the *Quartus II Handbook*. 
Planning Design Partitions and Floorplan Location Assignments

Partitioning a design for an FPGA requires planning to ensure optimal results when you integrate the partitions. Following Altera’s recommendations for creating design partitions should improve the overall quality of results. For example, registering partition I/O boundaries keeps critical timing paths inside one partition that can be optimized independently. When you specify the design partitions, you can use the Incremental Compilation Advisor to ensure that partitions meet Altera’s recommendations.

If you have timing-critical partitions that are changing through the design flow, or partitions exported from another Quartus II project, you can create design floorplan assignments to constrain the placement of the affected partitions. Good partition and floorplan design helps partitions meet top-level design requirements when integrated with the rest of your design, reducing time you spend integrating and verifying the timing of the top-level design.

For detailed guidelines about creating design partitions and organizing your source code, as well as information about when and how to create floorplan assignments, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

For more information about creating floorplan assignments in the Chip Planner, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

Fast Synthesis and Early Timing Estimation

You save time when you find design issues early in the design cycle rather than in the final timing closure stages. When the first version of the design source code is complete, you might want to perform a quick compilation to create a kind of silicon virtual prototype (SVP) that you can use to perform timing analysis.

If you synthesize with the Quartus II software, you can choose to perform a Fast synthesis, which reduces the compilation time, but may give reduced quality of results.

For more information about Fast synthesis, refer to Synthesis Effort logic option in Quartus II Help.

Regardless of your compilation flow, you can run an early timing estimate to perform a quick placement and routing, and a timing analysis of your design. The software chooses a device automatically if required, places any LogicLock regions to create a floorplan, finds a quick initial placement for all the design logic, and provides a useful estimate of the final design performance. If you have entered timing constraints, timing analysis reports on these constraints.

For more information about how to run an early timing estimate, refer to Running a Timing Analysis in Quartus II Help.
If you design individual design blocks or partitions separately, you can use the Fast synthesis and early timing estimate features as you develop your design. Any issues highlighted in the lower-level design blocks are communicated to the system architect. Resolving these issues might require allocating additional device resources to the individual partition, or changing the timing budget of the partition.

Expert designers can also use fast synthesis and early timing estimation to prototype the entire design. Incomplete partitions are marked as empty in an incremental compilation flow, while the rest of the design is compiled to get an early timing estimate and detect any problems with design integration.

**Conclusion**

Modern FPGAs support large, complex designs with fast timing performance. By planning several aspects of your design early, you can reduce time in later stages of the development cycle. Use features of the Quartus II software to quickly plan your design and achieve the best possible results. Following the guidelines presented in this chapter can improve productivity, which can reduce cost and development time.

**Document Revision History**

Table 2–2 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tr>
<td>November, 2012</td>
<td>12.1.0</td>
<td>Update for changes to early pin planning feature</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Editorial update.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>■ Added link to System Design with Qsys in “Creating Design Specifications” on page 1–2&lt;br&gt; ■ Updated “Simultaneous Switching Noise Analysis” on page 1–8&lt;br&gt; ■ Updated “Planning for On-Chip Debugging Tools” on page 1–10&lt;br&gt; ■ Removed information from “Planning Design Partitions and Floorplan Location Assignments” on page 1–15</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Changed to new document template&lt;br&gt; ■ Updated “System Design and Standard Interfaces” on page 1–3 to include information about the Qsys system integration tool&lt;br&gt; ■ Added link to the Altera Product Selector in “Device Selection” on page 1–3&lt;br&gt; ■ Converted information into new table (Table 1–1) in “Planning for On-Chip Debugging Options” on page 1–10&lt;br&gt; ■ Simplified description of incremental compilation usages in “Incremental Compilation with Design Partitions” on page 1–14&lt;br&gt; ■ Added information about the Rapid Recompile option in “Flat Compilation Flow with No Design Partitions” on page 1–14&lt;br&gt; ■ Removed details and linked to Quartus II Help in “Fast Synthesis and Early Timing Estimation” on page 1–16</td>
</tr>
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### Table 2–2. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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| July 2010  | 10.0.0  | Added new section “System Design” on page 1–3  
|            |         | Removed details about debugging tools from “Planning for On-Chip Debugging Options” on page 1–10 and referred to other handbook chapters for more information  
|            |         | Updated information on recommended design flows in “Incremental Compilation with Design Partitions” on page 1–14 and removed “Single-Project Versus Multiple-Project Incremental Flows” heading  
|            |         | Merged the “Planning Design Partitions” section with the “Creating a Design Floorplan” section. Changed heading title to “Planning Design Partitions and Floorplan Location Assignments” on page 1–15  
|            |         | Removed “Creating a Design Floorplan” section  
|            |         | Removed “Referenced Documents” section  
|            |         | Minor updates throughout chapter |
| November 2009 | 9.1.0 | Added details to “Creating Design Specifications” on page 1–2  
|            |         | Added details to “Intellectual Property Selection” on page 1–2  
|            |         | Updated information on “Device Selection” on page 1–3  
|            |         | Added reference to “Device Migration Planning” on page 1–4  
|            |         | Removed information from “Planning for Device Programming or Configuration” on page 1–4  
| March 2009  | 9.0.0   | No change to content |
| November 2008 | 8.1.0 | Changed to 8-1/2 x 11 page size. No change to content. |
| May 2008    | 8.0.0   | Organization changes  
|            |         | Added “Creating Design Specifications” section  
|            |         | Added reference to new details in the In-System Design Debugging section of volume 3  
|            |         | Added more details to the “Design Practices and HDL Coding Styles” section  
|            |         | Added references to the new Best Practices for Incremental Compilation and Floorplan Assignments chapter  
|            |         | Added reference to the Quartus II Language Templates |
This chapter provides information and design scenarios to help you partition your design to take advantage of the Quartus® II incremental compilation feature.

The ability to iterate rapidly through FPGA design and debugging stages is critical. The Quartus II software introduced the FPGA industry’s first true incremental design and compilation flow, with the following benefits:

- Preserves the results and performance for unchanged logic in your design as you make changes elsewhere.
- Reduces design iteration time by an average of 75% for small changes in large designs, so that you can perform more design iterations per day and achieve timing closure efficiently.
- Facilitates modular hierarchical and team-based design flows, as well as design reuse and intellectual property (IP) delivery.

Quartus II incremental compilation supports the Arria®, Stratix®, and Cyclone® series of devices, with limited support for HardCopy® ASICs (for details, refer to “Limitations for HardCopy Compilation and Migration Flows” on page 3–48).

This document contains the following sections:

- “Deciding Whether to Use an Incremental Compilation Flow” on page 3–2
- “Incremental Compilation Summary” on page 3–7
- “Common Design Scenarios Using Incremental Compilation” on page 3–10
- “Deciding Which Design Blocks Should Be Design Partitions” on page 3–14
- “Specifying the Level of Results Preservation for Subsequent Compilations” on page 3–20
- “Exporting Design Partitions from Separate Quartus II Projects” on page 3–26
- “Team-Based Design Optimization and Third-Party IP Delivery Scenarios” on page 3–35
- “Creating a Design Floorplan With LogicLock Regions” on page 3–44
- “Incremental Compilation Restrictions” on page 3–47
- “Scripting Support” on page 3–54
Deciding Whether to Use an Incremental Compilation Flow

The Quartus II incremental compilation feature enhances the standard Quartus II design flow by allowing you to preserve satisfactory compilation results and performance of unchanged blocks of your design.

This section outlines the flat compilation flow with no design partitions, the incremental flow when you divide the design into partitions, and the differences between the flat compilation and incremental compilation flows. This section also explains when a flat compilation flow is satisfactory, and highlights some of the reasons why you might want to create design partitions and use the incremental compilation flow. A discussion about incremental and team design flows in “Team-Based Design Flows and IP Delivery” on page 3–6 describes how it is beneficial to keep your design within one project, as well as when it might be necessary for other team members or IP providers to develop particular design blocks or partitions separately, and then later integrate their partitions into the top-level design.

Flat Compilation Flow with No Design Partitions

In the flat compilation flow with no design partitions, all the source code is processed and mapped during the Analysis and Synthesis stage, and placed and routed during the Fitter stage whenever the design is recompiled after a change in any part of the design. One reason for this behavior is to ensure optimal push-button quality of results. By processing the entire design, the Compiler can perform global optimizations to improve area and performance.

You can use a flat compilation flow for small designs, such as designs in CPLD devices or low-density FPGA devices, when the timing requirements are met easily with a single compilation. A flat design is satisfactory when compilation time and preserving results for timing closure are not concerns.

For more information on how to reduce compilation time when you use a flat compilation for your design, refer to the Reducing Compilation Time chapter in volume 2 of the Quartus II Handbook.

Incremental Capabilities Available When A Design Has No Partitions

The Quartus II software has incremental compilation features available even when you do not partition your design, including Smart Compilation, incremental debugging, and Rapid Recompile. These features work in either an incremental or flat compilation flow.

In any Quartus II compilation flow, you can use Smart Compilation to allow the Compiler to determine which compilation stages are required, based on the changes made to the design since the last smart compilation, and then skip any stages that are not required. For example, when Smart Compilation is turned on, the Compiler skips the Analysis and Synthesis stage if all the design source files are unchanged. When Smart Compilation is turned on, if you make any changes to the logic of a design, the Compiler does not skip any compilation stage. You can turn on Smart Compilation on the Compilation Process Settings page of the Setting dialog box.
The Quartus II software also includes a Rapid Recompile feature that instructs the Compiler to reuse the compatible compilation results if most of the design has not changed since the last compilation. This feature reduces compilation times for small and isolated design changes. You do not have control over which parts of the design are recompiled using this option; the Compiler determines which parts of the design must be recompiled. The Rapid Recompile feature preserves performance and can save compilation time by reducing the amount of changed logic that must be recompiled. You can turn on the **Rapid Recompile** option in the Quartus II software on the **Incremental Compilation** page in the **Settings** dialog box.

During the debugging stage of the design cycle, you can use incremental compilation to add the SignalTap® II Logic Analyzer incrementally to your design, even if the design does not have partitions. To preserve the compilation netlist for the entire design, instruct the software to reuse the compilation results for the automatically-created "Top" partition that contains the entire design. For more information, refer to “Debugging Incrementally With the SignalTap II Logic Analyzer” on page 3–13.

### Incremental Compilation Flow With Design Partitions

In the standard incremental compilation design flow, the top-level design is divided into design partitions, which can be compiled and optimized together in the top-level Quartus II project. You can preserve fitting results and performance for completed partitions while other parts of the design are changing, which reduces the compilation times for each design iteration.

Incremental compilation is recommended for large designs and high resource densities when preserving results is important to achieve timing closure. The incremental compilation feature also facilitates team-based design flows that allow designers to create and optimize design blocks independently, when necessary. Refer to “Team-Based Design Flows and IP Delivery” on page 3–6 for more information.

To take advantage of incremental compilation, start by splitting your design along any of its hierarchical boundaries into design blocks to be compiled incrementally, and set each block as a design partition. The Quartus II software synthesizes each individual hierarchical design partition separately, and then merges the partitions into a complete netlist for subsequent stages of the compilation flow. When recompiling your design, you can use source code, post-synthesis results, or post-fitting results to preserve satisfactory results for each partition. Refer to “Incremental Compilation Summary” on page 3–7 for more information.

In a team-based environment, part of your design may be incomplete, or it may have been developed by another designer or IP provider. In this scenario, you can add the completed partitions to the design incrementally. Alternatively, other designers or IP providers can develop and optimize partitions independently and the project lead can later integrate the partitions into the top-level design. Refer to “Team-Based Design Flows and IP Delivery” on page 3–6 for more information.
Table 3–1 shows a summary of the impact the Quartus II incremental compilation feature has on compilation results.

### Table 3–1. Impact Summary of Using Incremental Compilation

<table>
<thead>
<tr>
<th>Characteristic</th>
<th>Impact of Incremental Compilation with Design Partitions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Compilation Time Savings</td>
<td>Typically saves an average of 75% of compilation time for small design changes in large designs when post-fit netlists are preserved; there are savings in both Quartus II Integrated Synthesis and the Fitter. (1)</td>
</tr>
<tr>
<td>Performance Preservation</td>
<td>Excellent performance preservation when timing critical paths are contained within a partition, because you can preserve post-fitting information for unchanged partitions.</td>
</tr>
<tr>
<td>Node Name Preservation</td>
<td>Preserves post-fitting node names for unchanged partitions.</td>
</tr>
<tr>
<td>Area Changes</td>
<td>The area (logic resource utilization) might increase because cross-boundary optimizations are limited, and placement and register packing are restricted.</td>
</tr>
<tr>
<td>$f_{\text{MAX}}$ Changes</td>
<td>The design’s maximum frequency might be reduced because cross-boundary optimizations are limited. If the design is partitioned and the floorplan location assignments are created appropriately, there might be no negative impact on $f_{\text{MAX}}$.</td>
</tr>
</tbody>
</table>

**Note to Table 3–1:**

(1) Quartus II incremental compilation does not reduce processing time for the early "pre-fitter" operations, such as determining pin locations and clock routing, so the feature cannot reduce compilation time if runtime is dominated by those operations.

If you use the incremental compilation feature at any point in your design flow, it is easier to accommodate the guidelines for partitioning a design and creating a floorplan if you start planning for incremental compilation at the beginning of your design cycle.

For more information and recommendations on how to prepare your design to use the Quartus II incremental compilation feature, and how to avoid negative impact on your design results, refer to the *Best Practices for Incremental Compilation Partitions and Floorplan Assignments* chapter in volume 1 of the *Quartus II Handbook*. 

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(1) Quartus II incremental compilation does not reduce processing time for the early "pre-fitter" operations, such as determining pin locations and clock routing, so the feature cannot reduce compilation time if runtime is dominated by those operations.
Figure 3–1 shows a diagram of the Quartus II design flow using incremental compilation with design partitions.

**Figure 3–1. Quartus II Design Flow Using Incremental Compilation**

The diagram in Figure 3–1 shows a top-level partition and two lower-level partitions. If any part of the design changes, Analysis and Synthesis processes the changed partitions and keeps the existing netlists for the unchanged partitions. After completion of Analysis and Synthesis, there is one post-synthesis netlist for each partition.

**Note to Figure 3–1:**

(1) When you use EDIF or VQM netlists created by third-party EDA synthesis tools, Analysis and Synthesis creates the design database, but logic synthesis and technology mapping are performed only for black boxes.
The Partition Merge step creates a single, complete netlist that consists of post-synthesis netlists, post-fit netlists, and netlists exported from other Quartus II projects, depending on the netlist type that you specify for each partition.

The Fitter then processes the merged netlist, preserves the placement and routing of unchanged partitions, and refits only those partitions that have changed. The Fitter generates the complete netlist for use in future stages of the compilation flow, including timing analysis and programming file generation, which can take place in parallel if more than one processor is enabled for use in the Quartus II software. The Fitter also generates individual netlists for each partition so that the Partition Merge stage can use the post-fit netlist to preserve the placement and routing of a partition, if specified, for future compilations.

If you define partitions, but want to check your compilation results without partitions in a “what if” scenario, you can direct the Compiler to ignore all partitions assignments in your project and compile the design as a “flat” netlist. When you turn on the Ignore partitions assignments during compilation option on the Incremental Compilation page, the Quartus II software disables all design partition assignments in your project and runs a full compilation ignoring all partition boundaries and netlists. Turning off the Ignore partitions assignments during compilation option restores all partition assignments and netlists for subsequent compilations.

For more information on incremental compilation settings, refer to Incremental Compilation Page and Design Partition Properties Dialog Box in Quartus II Help.

**Team-Based Design Flows and IP Delivery**

The Quartus II software supports various design flows to enable team-based design and third-party IP delivery. A top-level design can include one or more partitions that are designed or optimized by different designers or IP providers, as well as partitions that will be developed as part of a standard incremental methodology.

In a team-based environment, part of your design may be incomplete because it is being developed elsewhere. The project lead or system architect can create empty placeholders in the top-level design for partitions that are not yet complete. Designers or IP providers can create and verify HDL code separately, and then the project lead later integrates the code into the single top-level Quartus II project. In this scenario, you can add the completed partitions to the design incrementally, however, the design flow allows all design optimization to occur in the top-level design for easiest design integration. Altera recommends using a single Quartus II project whenever possible because using multiple projects can add significant up-front and debugging time to the development cycle.

Alternatively, partition designers can design their partition in a copy of the top-level design or in a separate Quartus II project. Designers export their completed partition as either a post-synthesis netlist or optimized placed and routed netlist, or both, along with assignments such as LogicLock™ regions, as appropriate. The project lead then integrates each design block as a design partition into the top-level design. Altera recommends that designers export and reuse post-synthesis netlists, unless optimized post-fit results are required in the top-level design, to simplify design optimization.
Teams with a bottom-up design approach often want to optimize placement and routing of design partitions independently and may want to create separate Quartus II projects for each partition. However, optimizing design partitions in separate Quartus II projects, and then later integrating the results into a top-level design, can have the following potential drawbacks that require careful planning:

- Achieving timing closure for the full design may be more difficult if you compile partitions independently without information about other partitions in the design. This problem may be avoided by careful timing budgeting and special design rules, such as always registering the ports at the module boundaries.

- Resource budgeting and allocation may be required to avoid resource conflicts and overuse. Creating a floorplan with LogicLock regions is recommended when design partitions are developed independently in separate Quartus II projects.

- Maintaining consistency of assignments and timing constraints can be more difficult if there are separate Quartus II projects. The project lead must ensure that the top-level design and the separate projects are consistent in their assignments.

A unique challenge of team-based design and IP delivery for FPGAs is the fact that the partitions being developed independently must share a common set of resources. To minimize issues that might arise from sharing a common set of resources, you can design partitions within a single Quartus II project or a copy of the top-level design. A common project ensures that designers have a consistent view of the top-level project framework.

For timing-critical partitions being developed and optimized by another designer, it is important that each designer has complete information about the top-level design in order to maintain timing closure during integration, and to obtain the best results. When you want to integrate partitions from separate Quartus II projects, the project lead can perform most of the design planning, and then pass the top-level design constraints to the partition designers. Preferably, partition designers can obtain a copy of the top-level design by checking out the required files from a source control system. Alternatively, the project lead can provide a copy of the top-level project framework, or pass design information using Quartus II-generated design partition scripts. In the case that a third-party designer has no information about the top-level design, developers can export their partition from an independent project if required.

For more information about managing team-based design flows, refer to “Exporting Design Partitions from Separate Quartus II Projects” on page 3–26 and “Project Management—Making the Top-Level Design Available to Other Designers” on page 3–28.

Exporting partitions is not supported in HardCopy or FPGA companion device compilations when there is a migration device setting. For more information, refer to “Limitations for HardCopy Compilation and Migration Flows” on page 3–48.

**Incremental Compilation Summary**

This section provides a summary of the standard incremental compilation design flow and describes how to create design partitions.
Figure 3–2 illustrates the incremental compilation design flow when all partitions are contained in one top-level design.

**Steps for Incremental Compilation**

This section summarizes the steps in an incremental compilation flow; preparing a design to use the incremental compilation feature, and then preserving satisfactory results and performance in subsequent incremental compilations.

For an interactive introduction to implementing an incremental compilation design flow, refer to the *Getting Started Tutorial* on the Help menu in the Quartus II software. For step-by-step instructions on how to use the incremental compilation feature, refer to *Using the Incremental Compilation Design Flow* in Quartus II Help.

**Preparing a Design for Incremental Compilation**

To begin, elaborate your design, or run any compilation flow (such as a full compilation) that includes the elaboration step. Elaboration is the part of the synthesis process that identifies your design’s hierarchy.

Next, designate specific instances in the design hierarchy as design partitions, as described in “Creating Design Partitions” on page 3–9.

If required for your design flow, create a floorplan with LogicLock regions location assignments for timing-critical partitions that change with future compilations. Assigning a partition to a physical region on the device can help maintain quality of results and avoid conflicts in certain situations. For more information about LogicLock region assignments, refer to “Creating a Design Floorplan With LogicLock Regions” on page 3–44.
Compiling a Design Using Incremental Compilation

The first compilation after making partition assignments is a full compilation, and prepares the design for subsequent incremental compilations. In subsequent compilations of your design, you can preserve satisfactory compilation results and performance of unchanged partitions with the Netlist Type setting in the Design Partitions window. The Netlist Type setting determines which type of netlist or source file the Partition Merge stage uses in the next incremental compilation. You can choose the Source File, Post-Synthesis netlist, or Post-Fit netlist. For more information about the Netlist Type setting, refer to “Specifying the Level of Results Preservation for Subsequent Compilations” on page 3–20.

Creating Design Partitions

There are several ways to designate a design instance as a design partition. This section provides an overview of tools you can use to create partitions in the Quartus II software. For more information on selecting which design blocks to assign as partitions and how to analyze the quality of your partition assignments, refer to “Deciding Which Design Blocks Should Be Design Partitions” on page 3–14.

Creating Design Partitions in the Project Navigator

You can right-click an instance in the list under the Hierarchy tab in the Project Navigator and use the sub-menu to create and delete design partitions.

For more information about how to create design partitions in the Quartus II Project Navigator, refer to Creating Design Partitions in Quartus II Help.

Creating Design Partitions in the Design Partitions Window

The Design Partitions window, available from the Assignments menu, allows you to create, delete, and merge partitions, and is the main window for setting the netlist type to specify the level of results preservation for each partition on subsequent compilations. For information about how to set the netlist type and the available settings, refer to “Netlist Type for Design Partitions” on page 3–21.

The Design Partitions window also lists recommendations at the bottom of the window with links to the Incremental Compilation Advisor, where you can view additional recommendations about partitions. The Color column indicates the color of each partition as it appears in the Design Partition Planner and Chip Planner.

You can right-click a partition in the window to perform various common tasks, such as viewing property information about a partition, including the time and date of the compilation netlists and the partition statistics.

When you create a partition, the Quartus II software automatically generates a name based on the instance name and hierarchy path. You can edit the partition name in the Design Partitions Window so that you avoid referring to them by their hierarchy path, which can sometimes be long. This is especially useful when using command-line commands or assignments, or when you merge partitions to give the partition a meaningful name. Partition names can be from 1 to 1024 characters in length and must be unique. The name can consist of alphanumeric characters and the pipe ( | ), colon ( : ), and underscore ( _ ) characters.
For more information about how to create and manage design partitions in the Design Partitions window, refer to Creating Design Partitions in Quartus II Help.

**Creating Design Partitions With the Design Partition Planner**

The Design Partition Planner allows you to view design connectivity and hierarchy, and can assist you in creating effective design partitions that follow Altera’s guidelines.

The Design Partition Planner displays a visual representation of design connectivity and hierarchy, as well as partitions and entity relationships. You can explore the connectivity between entities in the design, evaluate existing partitions with respect to connectivity between entities, and try new partitioning schemes in “what if” scenarios.

When you extract design blocks from the top-level design and drag them into the Design Partition Planner, connection bundles are drawn between entities, showing the number of connections existing between pairs of entities. In the Design Partition Planner, you can then set extracted design blocks as design partitions.

The Design Partition Planner also has an Auto-Partition feature that creates partitions based on the size and connectivity of the hierarchical design blocks.

For more information about how to use the Design Partition Planner, refer to Using the Design Partition Planner in Quartus II Help and the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

**Creating Design Partitions With Tcl Scripting**

You can also create partitions with Tcl scripting commands. For more information about the command-line and scripting flow, refer to “Scripting Support” on page 3–54.

**Automatically-Generated Partitions**

The Compiler creates some partitions automatically as part of the compilation process, which appear in some post-compilation reports. For example, the sld_hub partition is created for tools that use JTAG hub connections, such as the SignalTap II Logic Analyzer. The hard_block partition is created to contain certain “hard” or dedicated logic blocks in the device that are implemented in a separate partition so that they can be shared throughout the design.

**Common Design Scenarios Using Incremental Compilation**

This section provides recommended applications of the incremental compilation flow after you have set up your design with partitions for incremental compilation as described in, “Steps for Incremental Compilation” on page 3–8.

This section contains the following design scenarios:

- “Reducing Compilation Time When Changing Source Files for One Partition” on page 3–11
- “Optimizing a Timing-Critical Partition” on page 3–11
Reducing Compilation Time When Changing Source Files for One Partition

Scenario background: You set up your design to include partitions for several of the major design blocks, and now you have just performed a lengthy compilation of the entire design. An error is found in the HDL source file for one partition and it is being fixed. Because the design is currently meeting timing requirements, and the fix is not expected to affect timing performance, it makes sense to compile only the affected partition and preserve the rest of the design.

Use the flow in this example to update the source file in one partition without having to recompile the other parts of the design. To reduce the compilation time, instruct the software to reuse the post-fit netlists for the unchanged partitions. This flow also preserves the performance of these blocks, which reduces additional timing closure efforts.

Perform the following steps to update a single source file:

1. Apply and save the fix to the HDL source file.
2. On the Assignments menu, open the Design Partitions window.
3. Change the netlist type of each partition, including the top-level entity, to Post-Fit to preserve as much as possible for the next compilation.

   The Quartus II software recompiles partitions by default when changes are detected in a source file. You can refer to the Partition Dependent Files table in the Analysis and Synthesis report to determine which partitions were recompiled. If you change an assignment but do not change the logic in a source file, you can set the netlist type to Source File for that partition to instruct the software to recompile the partition’s source design files and its assignments.

4. Click Start Compilation to incrementally compile the fixed HDL code. This compilation should take much less time than the initial full compilation.
5. Simulate the design to ensure that the error is fixed, and use the TimeQuest Timing Analyzer report to ensure that timing results have not degraded.

Optimizing a Timing-Critical Partition

Scenario background: You have just performed a lengthy full compilation of a design that consists of multiple partitions. The TimeQuest Timing Analyzer reports that the clock timing requirement is not met, and you have to optimize one particular partition. You want to try optimization techniques such as raising the Placement Effort Multiplier, enabling Physical Synthesis, and running the Design Space Explorer. Because these techniques all involve significant compilation time, you should apply them to only the partition in question.
Use the flow in this example to optimize the results of one partition when the other partitions in the design have already met their requirements. You can use this flow iteratively to lock down the performance of one partition, and then move on to optimization of another partition.

Perform the following steps to preserve the results for partitions that meet their timing requirements, and to recompile a timing-critical partition with new optimization settings:

1. Open the Design Partitions window.
2. For the partition in question, set the netlist type to Source File.

   If you change a setting that affects only the Fitter, you can save additional compilation time by setting the netlist type to Post-Synthesis to reuse the synthesis results and refit the partition.

3. For the remaining partitions (including the top-level entity), set the netlist type to Post-Fit.

   You can optionally set the Fitter Preservation Level on the Advanced tab in the Design Partitions Properties dialog box to Placement to allow for the most flexibility during routing.

4. Apply the desired optimization settings.
5. Click Start Compilation to perform incremental compilation on the design with the new settings. During this compilation, the Partition Merge stage automatically merges the critical partition’s new synthesis netlist with the post-fit netlists of the remaining partitions. The Fitter then refits only the required partition. Because the effort is reduced as compared to the initial full compilation, the compilation time is also reduced.

To use the Design Space Explorer, perform the following steps:

1. Repeat steps 1–3 of the previous procedure.
2. Save the project and run the Design Space Explorer.

**Adding Design Logic Incrementally or Working With an Incomplete Design**

Scenario background: You have one or more partitions that are known to be timing-critical in your full design. You want to focus on developing and optimizing this subset of the design first, before adding the rest of the design logic.

Use this flow to compile a timing-critical partition or partitions in isolation, optionally with extra optimizations turned on. After timing closure is achieved for the critical logic, you can preserve its content and placement and compile the remaining partitions with normal or reduced optimization levels. For example, you may want to compile an IP block that comes with instructions to perform optimization before you incorporate the rest of your custom logic.
To implement this design flow, perform the following steps:

1. Partition the design and create floorplan location assignments. For best results, ensure that the top-level design includes the entire project framework, even if some parts of the design are incomplete and are represented by an empty wrapper file.

2. For the partitions to be compiled first, in the Design Partitions window, set the netlist type to **Source File**.

3. For the remaining partitions, set the netlist type to **Empty**.

4. To compile with the desired optimizations turned on, click **Start Compilation**.

5. Check the Timing Analyzer reports to ensure that timing requirements are met. If so, proceed to step 6. Otherwise, repeat steps 4 and 5 until the requirements are met.

6. In the Design Partitions window, set the netlist type to **Post-Fit** for the first partitions. You can set the **Fitter Preservation Level** on the **Advanced** tab in the **Design Partitions Properties** dialog box to **Placement** to allow more flexibility during routing if exact placement and routing preservation is not required.

7. Change the netlist type from **Empty** to **Source File** for the remaining partitions, and ensure that the completed source files are added to the project.

8. Set the appropriate level of optimizations and compile the design. Changing the optimizations at this point does not affect any fitted partitions, because each partition has its netlist type set to **Post-Fit**.

9. Check the Timing Analyzer reports to ensure that timing requirements are met. If not, make design or option changes and repeat step 8 and step 9 until the requirements are met.

The flow in this example is similar to design flows in which a module is implemented separately and is later merged into the top-level, such as in the team-based design flow described in “Designing in a Team-Based Environment” on page 3–38. Generally, optimization in this flow works only if each critical path is contained within a single partition due to the effects described in “Deciding Which Design Blocks Should Be Design Partitions” on page 3–14. Ensure that if there are any partitions representing a design file that is missing from the project, you create a placeholder wrapper file to define the port interface. For more information, refer to “Empty Partitions” on page 3–28.

### Debugging Incrementally With the SignalTap II Logic Analyzer

Scenario background: Your design is not functioning as expected, and you want to debug the design using the SignalTap II Logic Analyzer. To maintain reduced compilation times and to ensure that you do not negatively affect the current version of your design, you want to preserve the synthesis and fitting results and add the SignalTap II Logic Analyzer to your design without recompiling the source code.

Use this flow to reduce compilation times when you add the logic analyzer to debug your design, or when you want to modify the configuration of the SignalTap II File without modifying your design logic or its placement.
It is not necessary to create design partitions in order to use the SignalTap II incremental compilation feature. The SignalTap II Logic Analyzer acts as its own separate design partition.

Perform the following steps to use the SignalTap II Logic Analyzer in an incremental compilation flow:

1. Open the Design Partitions window.
2. Set the netlist type to **Post-fit** for all partitions to preserve their placement.

   The netlist type for the top-level partition defaults to **Source File**, so be sure to change this “Top” partition in addition to any design partitions that you have created.

3. If you have not already compiled the design with the current set of partitions, perform a full compilation. If the design has already been compiled with the current set of partitions, the design is ready to add the SignalTap II Logic Analyzer.

4. Set up your SignalTap II File using the **post-fitting** filter in the **Node Finder** to add signals for logic analysis. This allows the Fitter to add the SignalTap II logic to the post-fit netlist without modifying the design results.

   To add signals from the pre-synthesis netlist, set the partition’s netlist type to **Source File** and use the **pre-synthesis** filter in the **Node Finder**. This allows the software to resynthesize the partition and to tap directly to the pre-synthesis node names that you choose. In this case, the partition is resynthesized and refit, so the placement is typically different from previous fitting results.

For more information about setting up the SignalTap II Logic Analyzer, refer to the **Design Debugging Using the SignalTap II Embedded Logic Analyzer** chapter in volume 3 of the **Quartus II Handbook**.

**Deciding Which Design Blocks Should Be Design Partitions**

The incremental compilation design flow requires more planning than flat compilations. For example, you might have to structure your source code or design hierarchy to ensure that logic is grouped correctly for optimization.

It is a common design practice to create modular or hierarchical designs in which you develop each design entity separately, and then instantiate them in a higher-level entity, forming a complete design. The Quartus II software does not automatically consider each design entity or instance to be a design partition for incremental compilation; instead, you must designate one or more design hierarchies below the top-level project as a design partition. Creating partitions might prevent the Compiler from performing optimizations across partition boundaries, as discussed in “Impact of Design Partitions on Design Optimization” on page 3–16. However, this allows for separate synthesis and placement for each partition, making incremental compilation possible.
Partitions must have the same boundaries as hierarchical blocks in the design because a partition cannot be a portion of the logic within a hierarchical entity. You can merge partitions that have the same immediate parent partition to create a single partition that includes more than one hierarchical entity in the design. When you declare a partition, every hierarchical instance within that partition becomes part of the same partition. You can create new partitions for hierarchical instances within an existing partition, in which case the instances within the new partition are no longer included in the higher-level partition, as described in the following example.

In Figure 3–3, a complete design is made up of instances A, B, C, D, E, F, and G. The shaded boxes in Representation i indicate design partitions in a “tree” representation of the hierarchy. In Representation ii, the lower-level instances are represented inside the higher-level instances, and the partitions are illustrated with different colored shading. The top-level partition, called “Top”, automatically contains the top-level entity in the design, and contains any logic not defined as part of another partition. The design file for the top level may be just a wrapper for the hierarchical instances below it, or it may contain its own logic. In this example, partition B contains the logic in instances B, D, and E. Entities F and G were first identified as separate partitions, and then merged together to create a partition F-G. The partition for the top-level entity A, called “Top”, includes the logic in one of its lower-level instances, C, because C was not defined as part of any other partition.

Figure 3–3. Partitions in a Hierarchical Design

You can create partition assignments to any design instance. The instance can be defined in HDL or schematic design, or come from a third-party synthesis tool as a VQM or EDIF netlist instance.
To take advantage of incremental compilation when source files change, create separate design files for each partition. If you define two different entities as separate partitions but they are in the same design file, you cannot maintain incremental compilation because the software would have to recompile both partitions if you changed either entity in the design file. Similarly, if two partitions rely on the same lower-level entity definition, changes in that lower-level affect both partitions.

The remainder of this section provides information to help you choose which design blocks you should assign as partitions.

**Impact of Design Partitions on Design Optimization**

The boundaries of your design partitions can impact the design’s quality of results. Creating partitions might prevent the Compiler from performing logic optimizations across partition boundaries, which allows the software to synthesize and place each partition separately in an incremental flow. Therefore, consider partitioning guidelines to help reduce the effect of partition boundaries.

Whenever possible, register all inputs and outputs of each partition. This helps avoid any delay penalty on signals that cross partition boundaries and keeps each register-to-register timing path within one partition for optimization. In addition, minimize the number of paths that cross partition boundaries. If there are timing-critical paths that cross partition boundaries, rework the partitions to avoid these inter-partition paths. Including as many of the timing-critical connections as possible inside a partition allows you to effectively apply optimizations to that partition to improve timing, while leaving the rest of the design unchanged.

Avoid constant partition inputs and outputs. You can also merge two or more partitions to allow cross-boundary optimizations for paths that cross between the partitions, as long as the partitions have the same parent partition. Merging related logic from different hierarchy blocks into one partition can be useful if you cannot change the design hierarchy to accommodate partition assignments.

The Design Partition Planner can help you create good assignments, as described in “Creating Design Partitions” on page 3–9. Refer to “Partition Statistics Reports” on page 3–19 for information about the number of I/O connections and how many are unregistered or driven by a constant value. For information on timing reports and additional design guidelines, refer to “Partition Timing Reports” on page 3–19 and “Incremental Compilation Advisor” on page 3–19.

If critical timing paths cross partition boundaries, you can perform timing budgeting and make timing assignments to constrain the logic in each partition so that the entire timing path meets its requirements. In addition, because each partition is optimized independently during synthesis, you may have to perform resource allocation to ensure that each partition uses an appropriate number of device resources. If design partitions are compiled in separate Quartus II projects, there may be conflicts related to global routing resources for clock signals when the design is integrated into the top-level design. You can use the Global Signal logic option to specify which clocks should use global or regional routing, use the ALTCLK_CTRL megafunction to instantiate a clock control block and connect it appropriately in both the partitions being developed in separate Quartus II projects, or find the compiler-generated clock control node in your design and make clock control location assignments in the Assignment Editor.
Turning On Supported Cross-boundary Optimizations

You can improve the optimizations performed between design partitions by turning on supported cross-boundary optimizations. These optimizations are turned on a per partition basis and you can select the optimizations as individual assignments. This allows the cross-boundary optimization feature to give you more control over the optimizations that work best for your design. You can turn on the cross-boundary optimizations for your design partitions on the Advanced tab of the Design Partition Properties dialog box. Once you change the optimization settings, the Quartus II software recompiles your partition from source automatically. Cross-boundary optimizations include the following: propagate constants, propagate inversions on partition inputs, merge inputs fed by a common source, merge electrically equivalent bidirectional pins, absorb internal paths, and remove logic connected to dangling outputs.

Cross-boundary optimizations are implemented top-down from the parent partition into the child partition, but not vice-versa. Also, cross-boundary optimizations cannot be enabled for partitions that allow multiple personas (partial reconfiguration partitions).

For more information about cross-boundary optimizations in the Quartus II software, refer to Design Partition Properties Dialog Box in Quartus II Help.

For more partitioning guidelines and specific recommendations for fixing common design issues, as well as information on resource allocation, global signal usage, and timing budgeting, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

Design Partition Assignments Compared to Physical Placement Assignments

Design partitions for incremental compilation are logical partitions, which is different from physical placement assignments in the device floorplan. A logical design partition does not refer to a physical area of the device and does not directly control the placement of instances. A logical design partition sets up a virtual boundary between design hierarchies so that each is compiled separately, preventing logical optimizations from occurring between them. When the software compiles the design source code, the logic in each partition can be placed anywhere in the device unless you make additional placement assignments.

If you preserve the compilation results using a Post-Fit netlist, it is not necessary for you to back-annotate or make any location assignments for specific logic nodes. You should not use the incremental compilation and logic placement back-annotation features in the same Quartus II project. The incremental compilation feature does not use placement “assignments” to preserve placement results; it simply reuses the netlist database that includes the placement information.
You can assign design partitions to physical regions in the device floorplan using LogicLock region assignments. In the Quartus II software, LogicLock regions are used to constrain blocks of a design to a particular region of the device. Altera recommends using LogicLock regions for timing-critical design blocks that will change in subsequent compilations, or to improve the quality of results and avoid placement conflicts in some cases. Creating floorplan location assignments for design partitions using LogicLock regions is discussed in “Creating a Design Floorplan With LogicLock Regions” on page 3–44.

For more information about when and why to create a design floorplan, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

**Using Partitions With Third-Party Synthesis Tools**

If you are using a third-party synthesis tool, set up your tool to create a separate VQM or EDIF netlist for each hierarchical partition. In the Quartus II software, assign the top-level entity from each netlist to be a design partition. The VQM or EDIF netlist file is treated as the source file for the partition in the Quartus II software.

**Synopsys Synplify Pro/Premier and Mentor Graphics Precision RTL Plus**

The Synplify Pro and Synplify Premier software include the MultiPoint synthesis feature to perform incremental synthesis for each design block assigned as a Compile Point in the user interface or a script. The Precision RTL Plus software includes an incremental synthesis feature that performs block-based synthesis based on Partition assignments in the source HDL code. These features provide automated block-based incremental synthesis flows and create different output netlist files for each block when set up for an Altera device.

Using incremental synthesis within your synthesis tool ensures that only those sections of a design that have been updated are resynthesized when the design is compiled, reducing synthesis run time and preserving the results for the unchanged blocks. You can change and resynthesize one section of a design without affecting other sections of the design.

For more information about these incremental synthesis flows, refer to your tool vendor’s documentation, or the Synopsys Synplify Support chapter or Mentor Graphics Precision Synthesis Support chapter in volume 1 of the Quartus II Handbook.

**Other Synthesis Tools**

You can also partition your design and create different netlist files manually with the basic Synplify software (non-Pro/Premier), the basic Precision RTL software (non-Plus), or any other supported synthesis tool by creating a separate project or implementation for each partition, including the top level. Set up each higher-level project to instantiate the lower-level VQM/EDIF netlists as black boxes. Synplify, Precision, and most synthesis tools automatically treat a design block as a black box if the logic definition is missing from the project. Each tool also includes options or attributes to specify that the design block should be treated as a black box, which you can use to avoid warnings about the missing logic.
Assessing Partition Quality

The Quartus II software provides various tools to assess the quality of your assigned design partitions. You can take advantage of these tools to assess your partition quality, and use the information to improve your design or assignments as required to achieve the best results.

Partition Statistics Reports

After compilation, you can view statistics about design partitions in the Partition Merge Partition Statistics report, and on the Statistics tab in the Design Partitions Properties dialog box.

The Partition Merge Partition Statistics report lists statistics about each partition. The statistics for each partition (each row in the table) include the number of logic cells it contains, as well as the number of input and output pins it contains, and how many are registered or unconnected. This report is useful when optimizing your design partitions, ensuring that the partitions meet the guidelines presented in the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

You can also view post-compilation statistics about the resource usage and port connections for a particular partition on the Statistics tab in the Design Partition Properties dialog box.

Partition Timing Reports

You can generate a Partition Timing Overview report and a Partition Timing Details report by clicking Report Partitions in the Tasks pane in the TimeQuest Timing Analyzer, or using the report_partitions Tcl command.

The Partition Timing Overview report shows the total number of failing paths for each partition and the worst-case slack for any path involving the partition.

The Partition Timing Details report shows the number of failing partition-to-partition paths and worst-case slack for partition-to-partition paths, to provide a more detailed breakdown of where the critical paths in the design are located with respect to design partitions.

Incremental Compilation Advisor

You can use the Incremental Compilation Advisor to check that your design follows Altera’s recommendations for creating design partitions and floorplan location assignments.

Recommendations are split into General Recommendations, Timing Recommendations, and Team-Based Design Recommendations that apply to design flows in which partitions are compiled independently in separate Quartus II projects before being integrated into the top-level design. Each recommendation provides an explanation, describes the effect of the recommendation, and provides the action required to make a suggested change. In some cases, there is a link to the appropriate Quartus II settings page where you can make a suggested change to assignments or
settings. For some items, if your design does not follow the recommendation, the Check Recommendations operation creates a table that lists any nodes or paths in your design that could be improved. The relevant timing-independent recommendations for the design are also listed in the Design Partitions window and the LogicLock Regions window.

To verify that your design follows the recommendations, go to the Timing Independent Recommendations page or the Timing Dependent Recommendations page, and then click Check Recommendations. For large designs, these operations can take a few minutes.

After you perform a check operation, symbols appear next to each recommendation to indicate whether the design or project setting follows the recommendations, or if some or all of the design or project settings do not follow the recommendations. Following these recommendations is not mandatory to use the incremental compilation feature. The recommendations are most important to ensure good results for timing-critical partitions.

For some items in the Advisor, if your design does not follow the recommendation, the Check Recommendations operation lists any parts of the design that could be improved. For example, if not all of the partition I/O ports follow the Register All Non-Global Ports recommendation, the advisor displays a list of unregistered ports with the partition name and the node name associated with the port.

When the advisor provides a list of nodes, you can right-click a node, and then click Locate to cross-probe to other Quartus II features, such as the RTL Viewer, Chip Planner, or the design source code in the text editor.

Opening a new TimeQuest report resets the Incremental Compilation Advisor results, so you must rerun the Check Recommendations process.

**Specifying the Level of Results Preservation for Subsequent Compilations**

As introduced in “Incremental Compilation Summary” on page 3–7 and “Common Design Scenarios Using Incremental Compilation” on page 3–10, the netlist type of each design partition allows you to specify the level of results preservation. The netlist type determines which type of netlist or source file the Partition Merge stage uses in the next incremental compilation.

When you choose to preserve a post-fit compilation netlist, the default level of Fitter preservation is the highest degree of placement and routing preservation supported by the device family. The advanced Fitter Preservation Level setting allows you to specify the amount of information that you want to preserve from the post-fit netlist file.
Netlist Type for Design Partitions

Before starting a new compilation, ensure that the appropriate netlist type is set for each partition to preserve the desired level of compilation results. Table 3–2 describes the settings for the netlist type, explains the behavior of the Quartus II software for each setting, and provides guidance on when to use each setting.

Table 3–2. Partition Netlist Type Settings (Part 1 of 2)

<table>
<thead>
<tr>
<th>Netlist Type</th>
<th>Quartus II Software Behavior for Partition During Compilation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Source File</td>
<td>Always compiles the partition using the associated design source file(s). (1) Use this netlist type to recompile a partition from the source code using new synthesis or Fitter settings.</td>
</tr>
</tbody>
</table>
| Post-Synthesis| Preserves post-synthesis results for the partition and reuses the post-synthesis netlist when the following conditions are true:  
  ■ A post-synthesis netlist is available from a previous synthesis.  
  ■ No change that initiates an automatic resynthesis has been made to the partition since the previous synthesis. (2) For details, refer to “What Changes Initiate the Automatic Resynthesis of a Partition?” on page 3–24.  
  Compiles the partition from the source files if resynthesis is initiated or if a post-synthesis netlist is not available. (1) Use this netlist type to preserve the synthesis results unless you make design changes, but allow the Fitter to refit the partition using any new Fitter settings. |
| Post-Fit       | Preserves post-fit results for the partition and reuses the post-fit netlist when the following conditions are true:  
  ■ A post-fit netlist is available from a previous fitting.  
  ■ No change that initiates an automatic resynthesis has been made to the partition since the previous fitting. (2) For details, refer to “What Changes Initiate the Automatic Resynthesis of a Partition?” on page 3–24.  
  When a post-fit netlist is not available, the software reuses the post-synthesis netlist if it is available, or otherwise compiles from the source files. Compiles the partition from the source files if resynthesis is initiated. (1) The Fitter Preservation Level specifies what level of information is preserved from the post-fit netlist. For details, refer to “Fitter Preservation Level for Design Partitions” on page 3–22.  
  Assignment changes, such as Fitter optimization settings, do not cause a partition set to Post-Fit to recompile. |
### Table 3–2. Partition Netlist Type Settings (Part 2 of 2)

<table>
<thead>
<tr>
<th>Netlist Type</th>
<th>Quartus II Software Behavior for Partition During Compilation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Empty</td>
<td>Uses an empty placeholder netlist for the partition. The partition’s port interface information is required during Analysis and Synthesis to connect the partition correctly to other logic and partitions in the design, and peripheral nodes in the source file including pins and PLLs are preserved to help connect the empty partition to the rest of the design and preserve timing of any lower-level non-empty partitions within empty partitions. If the source file is not available, you can create a wrapper file that defines the design block and specifies the input, output, and bidirectional ports. In Verilog HDL: a module declaration, and in VHDL: an entity and architecture declaration. You can use this netlist type to skip the compilation of a partition that is incomplete or missing from the top-level design. You can also set an empty partition if you want to compile only some partitions in the design, such as to optimize the placement of a timing-critical block such as an IP core before incorporating other design logic, or if the compilation time is large for one partition and you want to exclude it. If the project database includes a previously generated post-synthesis or post-fit netlist for an unchanged Empty partition, you can set the netlist type from Empty directly to Post-Synthesis or Post-Fit and the software reuses the previous netlist information without recompiling from the source files.</td>
</tr>
</tbody>
</table>

**Notes to Table 3–2:**

1. If you turn on the Rapid Recompile option, the Quartus II software may not recompile the entire partition from the source code as described in this table; it will reuse compatible results if there have been only small changes to the logic in the partition. Refer to “Incremental Capabilities Available When A Design Has No Partitions” on page 3–2 for more information.

2. You can turn on the Ignore changes in source files and strictly use the specified netlist, if available option on the Advanced tab in the Design Partitions Properties dialog box to specify whether the Compiler should ignore source file changes when deciding whether to recompile the partition.

### Fitter Preservation Level for Design Partitions

The default Fitter Preservation Level for partitions with a Post-Fit netlist type is the highest level of preservation available for the target device family and provides the most compilation time reduction.

You can change the advanced Fitter Preservation Level setting to provide more flexibility in the Fitter during placement and routing. You can set the Fitter Preservation Level on the Advanced tab in the Design Partitions Properties dialog box. Table 3–3 describes the Fitter Preservation Level settings.

### Table 3–3. Fitter Preservation Level Settings (Part 1 of 2)

<table>
<thead>
<tr>
<th>Fitter Preservation Level</th>
<th>Quartus II Behavior for Partition During Compilation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Placement and Routing</td>
<td>Preserves the design partition’s netlist atoms and their placement and routing. This setting reduces compilation times compared to Placement only, but provides less flexibility to the router to make changes if there are changes in other parts of the design. By default, the Fitter preserves the usage of high-speed programmable power tiles contained within the selected partition, for devices that support high-speed and low-power tiles. You can turn off the Preserve high-speed tiles when preserving placement and routing option on the Advanced tab in the Design Partitions Properties dialog box.</td>
</tr>
</tbody>
</table>
Table 3–3. Fitter Preservation Level Settings  (Part 2 of 2)

<table>
<thead>
<tr>
<th>Fitter Preservation Level</th>
<th>Quartus II Behavior for Partition During Compilation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Placement</td>
<td>Preserves the netlist atoms and their placement in the design partition. Reroutes the design partition and does not preserve high-speed power tile usage.</td>
</tr>
</tbody>
</table>
| Netlist Only             | Preserves the netlist atoms of the design partition, but replaces and reroutes the design partition. A post-fit netlist with the atoms preserved can be different than the Post-Synthesis netlist because it contains Fitter optimizations; for example, Physical Synthesis changes made during a previous Fitting. You can use this setting to:  
  ■ Preserve Fitter optimizations but allow the software to perform placement and routing again.  
  ■ Reapply certain Fitter optimizations that would otherwise be impossible when the placement is locked down.  
  ■ Resolve resource conflicts between two imported partitions. |

For more information about how to set the Netlist Type and Fitter Preservation Level settings in the Quartus II software, refer to Setting the Netlist Type and Fitter Preservation Level for Design Partitions in Quartus II Help.

Where Are the Netlist Databases Saved?

The incremental compilation database folder (\incremental\_db) includes all the netlist information from previous compilations. To avoid unnecessary recompilations, these database files must not be altered or deleted.

If you archive or reproduce the project in another location, you can use a Quartus II Archive File (.qar). Include the incremental compilation database files to preserve post-synthesis or post-fit compilation results. For more information, refer to “Using Incremental Compilation With Quartus II Archive Files” on page 3–48.

To manually create a project archive that preserves compilation results without keeping the incremental compilation database, you can keep all source and settings files, and create and save a Quartus II Settings File (.qxp) for each partition in the design that will be integrated into the top-level design. For more information about how to create a .qxp for a partition within your design, refer to “Exporting Design Partitions from Separate Quartus II Projects” on page 3–26.

Deleting Netlists

You can choose to abandon all levels of results preservation and remove all netlists that exist for a particular partition with the Delete Netlists command in the Design Partitions window. When you delete netlists for a partition, the partition is compiled using the associated design source file(s) in the next compilation. Resetting the netlist type for a partition to Source would have the same effect, though the netlists would not be permanently deleted and would be available for use in subsequent compilations. For an imported partition, the Delete Netlists command also optionally allows you to remove the imported .qxp.
What Changes Initiate the Automatic Resynthesis of a Partition?

A partition is synthesized from its source files if there is no post-synthesis netlist available from a previous synthesis, or if the netlist type is set to **Source File**. Additionally, certain changes to a partition initiate an automatic resynthesis of the partition when the netlist type is **Post-Synthesis** or **Post-Fit**. The software resynthesizes the partition in these cases to ensure that the design description matches the post-place-and-route programming files. If you do not want resynthesis to occur automatically, refer to “Forcing Use of the Compilation Netlist When a Partition has Changed” on page 3–26.

The following list explains the changes that initiate a partition’s automatic resynthesis when the netlist type is set to **Post-Synthesis** or **Post-Fit**:

- The device family setting has changed.
- Any dependent source design file has changed. For more information, refer to “Resynthesis Due to Source Code Changes” on page 3–25.
- The partition boundary was changed by an addition, removal, or change to the port boundaries of a partition (for example, a new partition has been defined for a lower-level instance within this partition).
- A dependent source file was compiled into a different library (so it has a different -library argument).
- A dependent source file was added or removed; that is, the partition depends on a different set of source files.
- The partition’s root instance has a different entity binding. In VHDL, an instance may be bound to a specific entity and architecture. If the target entity or architecture changes, it triggers resynthesis.
- The partition has different parameters on its root hierarchy or on an internal AHDL hierarchy (AHDL automatically inherits parameters from its parent hierarchies). This occurs if you modified the parameters on the hierarchy directly, or if you modified them indirectly by changing the parameters in a parent design hierarchy.
- You have moved the project and compiled database between a Windows and Linux system. Due to the differences in the way new line feeds are handled between the operating systems, the internal checksum algorithm may detect a design file change in this case.

The software reuses the post-synthesis results but re-fits the design if you change the device setting within the same device family. The software reuses the post-fitting netlist if you change only the device speed grade.

Synthesis and Fitter assignments, such as optimization settings, timing assignments, or Fitter location assignments including pin assignments, do not trigger automatic recompilation in the incremental compilation flow. To recompile a partition with new assignments, change the netlist type for that partition to one of the following:

- **Source File** to recompile with all new settings
- **Post-Synthesis** to recompile using existing synthesis results but new Fitter settings
Post-Fit with the Fitter Preservation Level set to Placement to rerun routing using existing placement results, but new routing settings (such as delay chain settings)

You can use the LogicLock Origin location assignment to change or fine-tune the previous Fitter results from a Post-Fit netlist. For details about how you can affect placement with LogicLock regions, refer to “Changing Partition Placement with LogicLock Changes” on page 3–46.

Resynthesis Due to Source Code Changes
The Quartus II software uses an internal checksum algorithm to determine whether the contents of a source file have changed. Source files are the design description files used to create the design, and include Memory Initialization Files (.mif) as well as .qxp from exported partitions. When design files in a partition have dependencies on other files, changing one file may initiate an automatic recompilation of another file. The Partition Dependent Files table in the Analysis and Synthesis report lists the design files that contribute to each design partition. You can use this table to determine which partitions are recompiled when a specific file is changed.

For example, if a design has file A.v that contains entity A, B.v that contains entity B, and C.v that contains entity C, then the Partition Dependent Files table for the partition containing entity A lists file A.v, the table for the partition containing entity B lists file B.v, and the table for the partition containing entity C lists file C.v. Any dependencies are transitive, so if file A.v depends on B.v, and B.v depends on C.v, the entities in file A.v depend on files B.v and C.v. In this case, files B.v and C.v are listed in the report table as dependent files for the partition containing entity A.

If you turn on the Rapid Recompile option, the Quartus II software may not recompile the entire partition from the source code as described in this section; it will reuse compatible results if there have been only small changes to the logic in the partition. Refer to “Incremental Capabilities Available When A Design Has No Partitions” on page 3–2 for more information.

If you define module parameters in a higher-level module, the Quartus II software checks the parameter values when determining which partitions require resynthesis. If you change a parameter in a higher-level module that affects a lower-level module, the lower-level module is resynthesized. Parameter dependencies are tracked separately from source file dependencies; therefore, parameter definitions are not listed in the Partition Dependent Files list.

If a design contains common files, such as an includes.v file that is referenced in each entity by the command ‘include includes.v’, all partitions are dependent on this file. A change to includes.v causes the entire design to be recompiled. The VHDL statement use work.all also typically results in unnecessary recompileations, because it makes all entities in the work library visible in the current entity, which results in the current entity being dependent on all other entities in the design.

To avoid this type of problem, ensure that files common to all entities, such as a common include file, contain only the set of information that is truly common to all entities. Remove use work.all statements in your VHDL file or replace them by including only the specific design units needed for each entity.
Forcing Use of the Compilation Netlist When a Partition has Changed

Forcing the use of a post-compilation netlist when the contents of a source file have changed is recommended only for advanced users who understand when a partition must be recompiled. You might use this assignment, for example, if you are making source code changes but do not want to recompile the partition until you finish debugging a different partition, or if you are adding simple comments to the source file but you know the design logic itself is not being changed and you want to keep the previous compilation results.

To force the Fitter to use a previously generated netlist even when there are changes to the source files, right-click the partition in the Design Partitions window and then click Design Partition Properties. On the Advanced tab, turn on the Ignore changes in source files and strictly use the specified netlist, if available option.

Turning on this option can result in the generation of a functionally incorrect netlist when source design files change, because source file updates will not be recompiled. Use caution when setting this option.

Exporting Design Partitions from Separate Quartus II Projects

Partitions that are developed by other designers or team members in the same company or third-party IP providers can be exported as design partitions to a Quartus II Exported Partition File (qxp), and then integrated into a top-level design. A qxp is a binary file that contains compilation results describing the exported design partition and includes a post-synthesis netlist, a post-fit netlist, or both, and a set of assignments, sometimes including LogicLock placement constraints. The qxp does not contain the source design files from the original Quartus II project.

To enable team-based development and third-party IP delivery, you can design and optimize partitions in separate copies of the top-level Quartus II project framework, or even in isolation. If the designers have access to the top-level project framework through a source control system, they can access project files as read-only and develop their partition within the source control system. If designers do not have access to a source control system, the project lead can provide the designer with a copy of the top-level project framework to use as they develop their partitions. The project lead also has the option to generate design partition scripts to manage resource and timing budgets in the top-level design when partitions are developed outside the top-level project framework.

The exported compilation results of completed partitions are given to the project lead, preferably using a source control system, who is then responsible for integrating them into the top-level design to obtain a fully functional design. This type of design flow is required only if partition designers want to optimize their placement and routing independently, and pass their design to the project lead to reuse placement and routing results. Otherwise, a project lead can integrate source HDL from several designers in a single Quartus II project, and use the standard incremental compilation flow described previously.
The diagram in Figure 3–4 illustrates the team-based incremental compilation design flow using a methodology in which partitions are compiled in separate Quartus II projects before being integrated into the top-level design. This flow can be used when partitions are developed by other designers or IP providers.

Figure 3–4. Summary of Team-Based Incremental Compilation Flow

You cannot export or import partitions that have been merged. For more information about merged partitions, refer to “Deciding Which Design Blocks Should Be Design Partitions” on page 3–14.

The topics in this section provide a description of the team-based design flow using exported partitions, describe how to generate a .qxp for a design partition, and explain how to integrate the .qxp into the top-level design:

There are some additional restrictions related to design flows using exported partitions, described in “Incremental Compilation Restrictions” on page 3–47.

Preparing the Top-Level Design

To prepare your design to incorporate exported partitions, first create the top-level project framework of the design to define the hierarchy for the subdesigns that will be implemented by other team members, designers, or IP providers.

In the top-level design, create project-wide settings, for example, device selection, global assignments for clocks and device I/O ports, and any global signal constraints to specify which signals can use global routing resources.
Next, create the appropriate design partition assignments and set the netlist type for each design partition that will be developed in a separate Quartus II project to **Empty**. Refer to “Empty Partitions” below for details. It may be necessary to constrain the location of partitions with LogicLock region assignments if they are timing-critical and are expected to change in future compilations, or if the designer or IP provider wants to place and route their design partition independently, to avoid location conflicts. For details, refer to “Creating a Design Floorplan With LogicLock Regions” on page 3–44.

Finally, provide the top-level project framework to the partition designers, preferably through a source control system. Refer to “Project Management—Making the Top-Level Design Available to Other Designers” on page 3–28 for more information.

**Empty Partitions**

You can use a design flow in which some partitions are set to an **Empty** netlist type to develop pieces of the design separately, and then integrate them into the top-level design at a later time. In a team-based design environment, you can set the netlist type to **Empty** for partitions in your design that will be developed by other designers or IP providers. The **Empty** setting directs the Compiler to skip the compilation of a partition and use an empty placeholder netlist for the partition.

When a netlist type is set to **Empty**, peripheral nodes including pins and PLLs are preserved and all other logic is removed. The peripheral nodes including pins help connect the empty partition to the design, and the PLLs help preserve timing of non-empty partitions within empty partitions.

When you set a design partition to **Empty**, a design file is required during Analysis and Synthesis to specify the port interface information so that it can connect the partition correctly to other logic and partitions in the design. If a partition is exported from another project, the .qxp contains this information. If there is no .qxp or design file to represent the design entity, you must create a wrapper file that defines the design block and specifies the input, output, and bidirectional ports. For example, in Verilog HDL, you should include a module declaration, and in VHDL, you should include an entity and architecture declaration.

**Project Management—Making the Top-Level Design Available to Other Designers**

In team-based incremental compilation flows, whenever possible, all designers or IP providers should work within the same top-level project framework. Using the same project framework among team members ensures that designers have the settings and constraints needed for their partition, and makes timing closure easier when integrating the partitions into the top-level design. If other designers do not have access to the top-level project framework, the Quartus II software provides tools for passing project information to partition designers.

**Distributing the Top-Level Quartus II Project**

There are several methods that the project lead can use to distribute the “skeleton” or top-level project framework to other partition designers or IP providers.
If partition designers have access to the top-level project framework, the project will already include all the settings and constraints needed for the design. This framework should include PLLs and other interface logic if this information is important to optimize partitions.

If designers are part of the same design environment, they can check out the required project files from the same source control system. This is the recommended way to share a set of project files.

Otherwise, the project lead can provide a copy of the top-level project framework so that each design develops their partition within the same project framework.

If a partition designer does not have access to the top-level project framework, the project lead can give the partition designer a Tcl script or other documentation to create the separate Quartus II project and all the assignments from the top-level design.

For details about project management scripts you can create with the Quartus II software, refer to “Optimizing the Placement for a Timing-Critical Partition” on page 3–57.

If the partition designers provide the project lead with a post-synthesis .qxp and fitting is performed in the top-level design, integrating the design partitions should be quite easy. If you plan to develop a partition in a separate Quartus II project and integrate the optimized post-fitting results into the top-level design, use the following guidelines to improve the integration process:

- Ensure that a LogicLock region constrains the partition placement and uses only the resources allocated by the project lead.
- Ensure that you know which clocks should be allocated to global routing resources so that there are no resource conflicts in the top-level design.
- Set the Global Signal assignment to On for the high fan-out signals that should be routed on global routing lines.
- To avoid other signals being placed on global routing lines, turn off Auto Global Clock and Auto Global Register Controls under More Settings on the Fitter page in the Settings dialog box. Alternatively, you can set the Global Signal assignment to Off for signals that should not be placed on global routing lines.

Placement for LABs depends on whether the inputs to the logic cells within the LAB use a global clock. You may encounter problems if signals do not use global lines in the partition, but use global routing in the top-level design.

- Use the Virtual Pin assignment to indicate pins of a partition that do not drive pins in the top-level design. This is critical when a partition has more output ports than the number of pins available in the target device. Using virtual pins also helps optimize cross-partition paths for a complete design by enabling you to provide more information about the partition ports, such as location and timing assignments.
When partitions are compiled independently without any information about each other, you might need to provide more information about the timing paths that may be affected by other partitions in the top-level design. You can apply location assignments for each pin to indicate the port location after incorporation in the top-level design. You can also apply timing assignments to the I/O ports of the partition to perform timing budgeting.

For more information about resource balancing and timing allocation between partitions, refer to the *Best Practices for Incremental Compilation Partitions and Floorplan Assignments* chapter in volume 1 of the *Quartus II Handbook.*

**Generating Design Partition Scripts**

If IP providers or designers on a team want to optimize their design blocks independently and do not have access to a shared project framework, the project lead must perform some or all of the following tasks to ensure successful integration of the design blocks:

- Determine which assignments should be propagated from the top-level design to the partitions. This requires detailed knowledge of which assignments are required to set up low-level designs.
- Communicate the top-level assignments to the partitions. This requires detailed knowledge of Tcl or other scripting languages to efficiently communicate project constraints.
- Determine appropriate timing and location assignments that help overcome the limitations of team-based design. This requires examination of the logic in the partitions to determine appropriate timing constraints.
- Perform final timing closure and resource conflict avoidance in the top-level design. Because the partitions have no information about each other, meeting constraints at the lower levels does not guarantee they are met when integrated at the top-level. It then becomes the project lead’s responsibility to resolve the issues, even though information about the partition implementation may not be available.

Design partition scripts automate the process of transferring the top-level project framework to partition designers in a flow where each design block is developed in separate Quartus II projects before being integrated into the top-level design. If the project lead cannot provide each designer with a copy of the top-level project framework, the Quartus II software provides an interface for managing resources and timing budgets in the top-level design. Design partition scripts make it easier for partition designers to implement the instructions from the project lead, and avoid conflicts between projects when integrating the partitions into the top-level design. This flow also helps to reduce the need to further optimize the designs after integration.

You can use options in the *Generate Design Partition Scripts* dialog box to choose which types of assignments you want to pass down and create in the partitions being developed in separate Quartus II projects.

For an example design scenario using design partition scripts, refer to “Enabling Designers on a Team to Optimize Independently” on page 3–39.
For step-by-step information on how to generate design partition scripts, and a description of each option that can be included in design partition scripts, refer to *Generating Design Partition Scripts for Project Management*, and *Generate Design Partition Scripts Dialog Box* in Quartus II Help.

## Exporting Partitions

When partition designers achieve the design requirements in their separate Quartus II projects, each designer can export their design as a partition so it can be integrated into the top-level design by the project lead. The Export Design Partition dialog box, available from the Project menu, allows designers to export a design partition to a Quartus II Exported Partition File (.qxp) with a post-synthesis netlist, a post-fit netlist, or both. The project lead then adds the .qxp to the top-level design to integrate the partition.

A designer developing a timing-critical partition or who wants to optimize their partition on their own would opt to export their completed partition with a post-fit netlist, allowing for the partition to more reliably meet timing requirements after integration. In this case, you must ensure that resources are allocated appropriately to avoid conflicts. If the placement and routing optimization can be performed in the top-level design, exporting a post-synthesis netlist allows the most flexibility in the top-level design and avoids potential placement or routing conflicts with other partitions.

When designing the partition logic to be exported into another project, you can add logic around the design block to be exported as a design partition. You can instantiate additional design components for the Quartus II project so that it matches the top-level design environment, especially in cases where you do not have access to the full top-level design project. For example, you can include a top-level PLL in the project, outside of the partition to be exported, so that you can optimize the design with information about the frequency multipliers, phase shifts, compensation delays, and any other PLL parameters. The software then captures timing and resource requirements more accurately while ensuring that the timing analysis in the partition is complete and accurate. You can export the partition for the top-level design without any auxiliary components that are instantiated outside the partition being exported.

If your design team uses makefiles and design partition scripts, the project lead can use the make command with the master_makefile command created by the scripts to export the partitions and create .qxp files. When a partition has been compiled and is ready to be integrated into the top-level design, you can export the partition with option on the Export Design Partition dialog box, available from the Project menu.

For more information about how to export a design partition, refer to *Using a Team-Based Incremental Compilation Design Flow* in the Quartus II Help.

### Viewing the Contents of a Quartus II Exported Partition File (.qxp)

The QXP report allows you to view a summary of the contents in a .qxp when you open the file in the Quartus II software. The .qxp is a binary file that contains compilation results so the file cannot be read in a text editor. The QXP report opens in the main Quartus II window and contains summary information including a list of the I/O ports, resource usage summary, and a list of the assignments used for the exported partition.
Integrating Partitions into the Top-Level Design

To integrate a partition developed in a separate Quartus II project into the top-level design, you can simply add the .qxp as a source file in your top-level design (just like a Verilog or VHDL source file). You can also use the Import Design Partition dialog box to import the partition, in certain situations, described in “Advanced Importing Options” on page 3–33.

The .qxp contains the design block exported from the partition and has the same name as the partition. When you instantiate the design block into a top-level design and include the .qxp as a source file, the software adds the exported netlist to the database for the top-level design. The .qxp port names are case sensitive if the original HDL of the partition was case sensitive.

When you use a .qxp as a source file in this way, you can choose whether you want the .qxp to be a partition in the top-level design. If you do not designate the .qxp instance as a partition, the software reuses just the post-synthesis compilation results from the .qxp, removes unconnected ports and unused logic just like a regular source file, and then performs placement and routing.

If you assigned the .qxp instance as a partition, you can set the netlist type in the Design Partitions Window to choose the level of results to preserve from the .qxp. To preserve the placement and routing results from the exported partition, set the netlist type to Post-Fit for the .qxp partition in the top-level design. If you assign the instance as a design partition, the partition boundary is preserved, as discussed in “Impact of Design Partitions on Design Optimization” on page 3–16.

Integrating Assignments from the .qxp

The Quartus II software filters assignments from .qxp files to include appropriate assignments in the top-level design. The assignments in the .qxp are treated like assignments made in an HDL source file, and are not listed in the Quartus II Settings File (.qsf) for the top-level design. Most assignments from the .qxp can be overridden by assignments in the top-level design.

The following subsections provide more details about specific assignment types:

Design Partition Assignments Within the Exported Partition

Design partition assignments defined within a separate Quartus II project are not added to the top-level design. All logic under the exported partition in the project hierarchy is treated as single instance in the .qxp.

Synopsys Design Constraint Files for the Quartus II TimeQuest Timing Analyzer

Timing assignments made for the Quartus II TimeQuest analyzer in a Synopsys Design Constraint File (.sdc) in the lower-level partition project are not added to the top-level design. Ensure that the top-level design includes all of the timing requirements for the entire project.

For recommendations about managing SDC constraints for the top-level design and independent lower-level partition projects, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.
Global Assignments
The project lead should make all global project-wide assignments in the top-level design. Global assignments from the exported partition’s project are not added to the top-level design. When it is possible for a particular constraint, the global assignment is converted to an instance-specific assignment for the exported design partition.

LogicLock Region Assignments
The project lead typically creates LogicLock region assignments in the top-level design for any lower-level partition designs where designer or IP providers plan to export post-fit information to be used in the top-level design, to help avoid placement conflicts between partitions. When you use the .qxp as a source file, LogicLock constraints from the exported partition are applied in the top-level design, but will not appear in your .qsf file or LogicLock Regions window for you to view or edit. The LogicLock region itself is not required to constrain the partition placement in the top-level design if the netlist type is set to Post-Fit, because the netlist contains all the placement information. For information on how to control LogicLock region assignments for exported partitions, refer to the “Advanced Importing Options” on page 3–33.

Integrating Encrypted IP Cores from .qxp Files
Proper license information is required to compile encrypted IP cores. If an IP core is exported as a .qxp from another Quartus II project, the top-level designer instantiating the .qxp must have the correct license. The software requires a full license to generate an unrestricted programming file. If you do not have a license, but the IP in the .qxp was compiled with OpenCore Plus hardware evaluation support, you can generate an evaluation programming file without a license. If the IP supports OpenCore simulation only, you can fully compile the design and generate a simulation netlist, but you cannot create programming files unless you have a full license.

Advanced Importing Options
You can use advanced options in the Import Design Partition dialog box to integrate a partition developed in a separate Quartus II project into the top-level design. The import process adds more control than using the .qxp as a source file, and is useful only in the following circumstances:

- **If you want LogicLock regions in your top-level design (.qsf)**—If you have regions in your partitions that are not also in the top-level design, the regions will be added to your .qsf during the import process.
- **If you want different settings or placement for different instantiations of the same entity**—You can control the setting import process with the advanced import options, and specify different settings for different instances of the same .qxp design block.

When you use the Import Design Partition dialog box to integrate a partition into the top-level design, the import process sets the partition’s netlist type to Imported in the Design Partitions window.
After you compile the entire design, if you make changes to the place-and-route results (such as movement of an imported LogicLock region), use the **Post-Fit** netlist type on subsequent compilations. To discard an imported netlist and recompile from source code, you can compile the partition with the netlist type set to **Source File** and be sure to include the relevant source code in the top-level design. Refer to “Netlist Type for Design Partitions” on page 3–21 for details. The import process sets the partition’s Fitter Preservation Level to the setting with the highest degree of preservation supported by the imported netlist. For example, if a post-fit netlist is imported with placement information, the Fitter Preservation Level is set to **Placement**, but you can change it to the **Netlist Only** value. For more information about preserving previous compilation results, refer to “Netlist Type for Design Partitions” on page 3–21 and “Fitter Preservation Level for Design Partitions” on page 3–22.

When you import a partition from a `.qxp`, the `.qxp` itself is not part of the top-level design because the netlists from the file have been imported into the project database. Therefore if a new version of a `.qxp` is exported, the top-level designer must perform another import of the `.qxp`.

When you import a partition into a top-level design with the **Import Design Partition** dialog box, the software imports relevant assignments from the partition into the top-level design, as described for the source file integration flow in “Integrating Assignments from the `.qxp`” on page 3–32. If required, you can change the way some assignments are imported, as described in the following subsections.

**Importing LogicLock Assignments**

LogicLock regions are set to a fixed size when imported. If you instantiate multiple instances of a subdesign in the top-level design, the imported LogicLock regions are set to a Floating location. Otherwise, they are set to a Fixed location. You can change the location of LogicLock regions after they are imported, or change them to a Floating location to allow the software to place each region but keep the relative locations of nodes within the region wherever possible. For details, refer to “Changing Partition Placement with LogicLock Changes” on page 3–46. To preserve changes made to a partition after compilation, use the **Post-Fit** netlist type.

The LogicLock Member State assignment is set to **Locked** to signify that it is a preserved region.

LogicLock back-annotation and node location data is not imported because the `.qxp` contains all of the relevant placement information. Altera strongly recommends that you do not add to or delete members from an imported LogicLock region.

**Advanced Import Settings**

The **Advanced Import Settings** dialog box allows you to disable assignment import and specify additional options that control how assignments and regions are integrated when importing a partition into a top-level design, including how to resolve assignment conflicts.

For descriptions of the advanced import options available, refer to Advanced Import Settings Dialog Box in Quartus II Help.
Team-Based Design Optimization and Third-Party IP Delivery Scenarios

This section includes the following design flows with step-by-step descriptions when you have partitions being developed in separate Quartus II projects, or by a third-party IP provider.

- “Using an Exported Partition to Send to a Design Without Including Source Files” on page 3–35
- “Creating Precompiled Design Blocks (or Hard-Wired Macros) for Reuse” on page 3–36
- “Designing in a Team-Based Environment” on page 3–38
- “Enabling Designers on a Team to Optimize Independently” on page 3–39
- “Performing Design Iterations With Lower-Level Partitions” on page 3–42

Using an Exported Partition to Send to a Design Without Including Source Files

Scenario background: A designer wants to produce a design block and needs to send out their design, but to preserve their IP, they prefer to send a synthesized netlist instead of providing the HDL source code to the recipient.

Use this flow to package a full design as a single source file to send to an end customer or another design location.

As the sender in this scenario perform the following steps to export a design block:

1. Provide the device family name to the recipient. If you send placement information with the synthesized netlist, also provide the exact device selection so they can set up their project to match.

2. Create a black box wrapper file that defines the port interface for the design block and provide it to the recipient for instantiating the block as an empty partition in the top-level design.

3. Create a Quartus II project for the design block, and complete the design.

4. Export the level of hierarchy into a single .qxp. Following a successful compilation of the project, you can generate a .qxp from the GUI, the command-line, or with Tcl commands, as described in the following:
   - If you are using the Quartus II GUI, use the Export Design Partition command.
   - If you are using command-line executables, run quartus_cdb with the --incremental_compilation_export option.
   - If you are using Tcl commands, use the following command:
     ```
     execute_flow -incremental_compilation_export
     ```

5. Select the option to include just the Post-synthesis netlist if you do not have to send placement information. If the recipient wants to reproduce your exact Fitter results, you can select the Post-fitting netlist option, and optionally enable Export routing.

6. Provide the .qxp to the recipient. Note that you do not have to send any of your design source code.
As the recipient in this example, first create a Quartus II project for your top-level design and ensure that your project targets the same device (or at least the same device family if the .qxp does not include placement information), as specified by the IP designer sending the design block. Instantiate the design block using the port information provided, and then incorporate the design block into a top-level design.

Add the .qxp from the IP designer as a source file in your Quartus II project to replace any empty wrapper file. If you want to use just the post-synthesis information, you can choose whether you want the file to be a partition in the top-level design. To use the post-fit information from the .qxp, assign the instance as a design partition and set the netlist type to Post-Fit. Refer to “Creating Design Partitions” on page 3–9 and “Netlist Type for Design Partitions” on page 3–21.

Creating Precompiled Design Blocks (or Hard-Wired Macros) for Reuse

Scenario background: An IP provider wants to produce and sell an IP core for a component to be used in higher-level systems. The IP provider wants to optimize the placement of their block for maximum performance in a specific Altera device and then deliver the placement information to their end customer. To preserve their IP, they also prefer to send a compiled netlist instead of providing the HDL source code to their customer.

Use this design flow to create a precompiled IP block (sometimes known as a hard-wired macro) that can be instantiated in a top-level design. This flow provides the ability to export a design block with post-synthesis or placement (and, optionally, routing) information and to import any number of copies of this pre-compiled block into another design.

The customer first specifies which Altera device is being used for this project and provides the design specifications.

As the IP provider in this example, perform the following steps to export a preplaced IP core (or hard macro):

1. Create a black box wrapper file that defines the port interface for the IP core and provide the file to the customer to instantiate as an empty partition in the top-level design.
2. Create a Quartus II project for the IP core.
3. Create a LogicLock region for the design hierarchy to be exported.

Using a LogicLock region for the IP core allows the customer to create an empty placeholder region to reserve space for the IP in the design floorplan and ensures that there are no conflicts with the top-level design logic. Reserved space also helps ensure the IP core does not affect the timing performance of other logic in the top-level design. Additionally, with a LogicLock region, you can preserve placement either absolutely or relative to the origin of the associated region. This is important when a .qxp is imported for multiple partition hierarchies in the same project, because in this case, the location of at least one instance in the top-level design does not match the location used by the IP provider.
4. If required, add any logic (such as PLLs or other logic defined in the customer’s top-level design) around the design hierarchy to be exported. If you do so, create a design partition for the design hierarchy that will exported as an IP core.

5. Optimize the design and close timing to meet the design specifications.

6. Export the level of hierarchy for the IP core into a single .qxp.

7. Provide the .qxp to the customer. Note that you do not have to send any of your design source code to the customer; the design netlist and placement and routing information is contained within the .qxp.

As the customer in this example, incorporate the IP core in your design by performing the following steps:

1. Create a Quartus II project for the top-level design that targets the same device and instantiate a copy or multiple copies of the IP core. Use a black box wrapper file to define the port interface of the IP core.

2. Perform Analysis and Elaboration to identify the design hierarchy.

3. Create a design partition for each instance of the IP core (refer to “Creating Design Partitions” on page 3–54) with the netlist type set to Empty (refer to “Netlist Type for Design Partitions” on page 3–21).

4. You can now continue work on your part of the design and accept the IP core from the IP provider when it is ready.

5. Include the .qxp from the IP provider in your project to replace the empty wrapper-file for the IP instance. Or, if you are importing multiple copies of the design block and want to import relative placement, follow these additional steps:
   a. Use the Import command to select each appropriate partition hierarchy. You can import a .qxp from the GUI, the command-line, or with Tcl commands:
      ■ If you are using the Quartus II GUI, use the Import Design Partition command.
      ■ If you are using command-line executables, run quartus_cdb with the --incremental_compilation_import option.
      ■ If you are using Tcl commands, use the following command:
        execute_flow -incremental_compilation_import.
   b. When you have multiple instances of the IP block, you can set the imported LogicLock regions to floating, or move them to a new location, and the software preserves the relative placement for each of the imported modules (relative to the origin of the LogicLock region). Routing information is preserved whenever possible. Refer to “Changing Partition Placement with LogicLock Changes” on page 3–46

   The Fitter ignores relative placement assignments if the LogicLock region’s location in the top-level design is not compatible with the locations exported in the .qxp.

6. You can control the level of results preservation with the Netlist Type setting. Refer to “Netlist Type for Design Partitions” on page 3–21.
If the IP provider did not define a LogicLock region in the exported partition, the software preserves absolute placement locations and this leads to placement conflicts if the partition is imported for more than one instance.

**Designing in a Team-Based Environment**

Scenario background: A project includes several lower-level design blocks that are developed separately by different designers and instantiated exactly once in the top-level design.

This scenario describes how to use incremental compilation in a team-based design environment where each designer has access to the top-level project framework, but wants to optimize their design in a separate Quartus II project before integrating their design block into the top-level design.

As the project lead in this scenario, perform the following steps to prepare the top-level design:

1. Create a new Quartus II project to ultimately contain the full implementation of the entire design and include a "skeleton" or framework of the design that defines the hierarchy for the subdesigns implemented by separate designers. The top-level design implements the top-level entity in the design and instantiates wrapper files that represent each subdesign by defining only the port interfaces but not the implementation.

2. Make project-wide settings. Select the device, make global assignments such as device I/O ports, define the top-level timing constraints, and make any global signal allocation constraints to specify which signals can use global routing resources.

3. Make design partition assignments for each subdesign and set the netlist type for each design partition to be imported to **Empty** in the Design Partitions window.

4. Create LogicLock regions to create a design floorplan for each of the partitions that will be developed separately. This floorplan should consider the connectivity between partitions and estimates of the size of each partition based on any initial implementation numbers and knowledge of the design specifications.

5. Provide the top-level project framework to partition designers using one of the following procedures:
   - Allow access to the full project for all designers through a source control system. Each designer can check out the projects files as read-only and work on their blocks independently. This design flow provides each designer with the most information about the full design, which helps avoid resource conflicts and makes design integration easy.
   - Provide a copy of the top-level Quartus II project framework for each designer. You can use the **Copy Project** command on the Project menu or create a project archive.

As the designer of a lower-level design block in this scenario, design and optimize your partition in your copy of the top-level design, and then follow these steps when you have achieved the desired compilation results:

1. On the Project menu, click **Export Design Partition**.
2. In the Export Design Partition dialog box, choose the netlist(s) to export. You can export a Post-synthesis netlist if placement or performance preservation is not required, to provide the most flexibility for the Fitter in the top-level design. Select Post-fit netlist to preserve the placement and performance of the lower-level design block, and turn on Export routing to include the routing information, if required. One .qxp can include both post-synthesis and post-fitting netlists.

3. Provide the .qxp to the project lead.

Finally, as the project lead in this scenario, perform these steps to integrate the .qxp files received from designers of each partition:

1. Add the .qxp as a source file in the Quartus II project, to replace any empty wrapper file for the previously Empty partition.

2. Change the netlist type for the partition from Empty to the required level of results preservation.

Enabling Designers on a Team to Optimize Independently

Scenario background: A project consists of several lower-level design blocks that are developed separately by different designers who do not have access to a shared top-level project framework. This scenario is similar to “Creating Precompiled Design Blocks (or Hard-Wired Macros) for Reuse” on page 3–36, but assumes that there are several design blocks being developed independently (instead of just one IP block), and the project lead can provide some information about the design to the individual designers. If the designers have shared access to the top-level design, use the previous scenario “Designing in a Team-Based Environment” on page 3–38.

This scenario describes how to use incremental compilation in a team-based design environment where designers or IP developers want to fully optimize the placement and routing of their design independently in a separate Quartus II project before sending the design to the project lead. This design flow requires more planning and careful resource allocation because design blocks are developed independently.

As the project lead in this scenario, perform the following steps to prepare the top-level design:

1. Create a new Quartus II project to ultimately contain the full implementation of the entire design and include a “skeleton” or framework of the design that defines the hierarchy for the subdesigns implemented by separate designers. The top-level design implements the top-level entity in the design and instantiates wrapper files that represent each subdesign by defining only the port interfaces but not the implementation.

2. Make project-wide settings. Select the device, make global assignments such as device I/O ports, define the top-level timing constraints, and make any global signal constraints to specify which signals can use global routing resources.

3. Make design partition assignments for each subdesign and set the netlist type for each design partition to be imported to Empty in the Design Partitions window.

4. Create LogicLock regions. This floorplan should consider the connectivity between partitions and estimates of the size of each partition based on any initial implementation numbers and knowledge of the design specifications.
5. Provide the constraints from the top-level design to partition designers using one of the following procedures:

- Use design partition scripts to pass constraints and generate separate Quartus II projects. On the Project menu, use the **Generate Design Partition Scripts** command, or run the script generator from a Tcl or command prompt. Make changes to the default script options as required for your project. Altera recommends that you pass all the default constraints, including LogicLock regions, for all partitions and virtual pin location assignments. If partitions have not already been created by the other designers, use the partition script to set up the projects so that you can easily take advantage of makefiles. Provide each partition designer with the Tcl file to create their project with the appropriate constraints. If you are using makefiles, provide the makefile for each partition.

- Use documentation or manually-created scripts to pass all constraints and assignments to each partition designer.

As the designer of a lower-level design block in this scenario, perform the appropriate set of steps to successfully export your design, whether the design team is using makefiles or exporting and importing the design manually.

If you are using makefiles with the design partition scripts, perform the following steps:

1. Use the **make** command and the makefile provided by the project lead to create a Quartus II project with all design constraints, and compile the project.

2. The information about which source file should be associated with which partition is not available to the software automatically, so you must specify this information in the makefile. You must specify the dependencies before the software rebuilds the project after the initial call to the makefile.

3. When you have achieved the desired compilation results and the design is ready to be imported into the top-level design, the project lead can use the **master_makefile** command to export this partition and create a `.qxp`, and then import it into the top-level design.

If you are not using makefiles, perform the following steps:

1. If you are using design partition scripts, source the Tcl script provided by the Project Lead to create a project with the required settings:

   - To source the Tcl script in the Quartus II software, on the Tools menu, click **Utility Windows** to open the Tcl console. Navigate to the script’s directory, and type the following command: `source <filename>`

   - To source the Tcl script at the system command prompt, type the following command: `quartus_cdb -t <filename>.tcl`
2. If you are not using design partition scripts, create a new Quartus II project for the subdesign, and then apply the following settings and constraints to ensure successful integration:
   - Make LogicLock region assignments and global assignments (including clock settings) as specified by the project lead.
   - Make Virtual Pin assignments for ports which represent connections to core logic instead of external device pins in the top-level design.
   - Make floorplan location assignments to the Virtual Pins so they are placed in their corresponding regions as determined by the top-level design. This provides the Fitter with more information about the timing constraints between modules. Alternatively, you can apply timing I/O constraints to the paths that connect to virtual pins.

3. Proceed to compile and optimize the design as needed.

4. When you have achieved the desired compilation results, on the Project menu, click Export Design Partition.

5. In the Export Design Partition dialog box, choose the netlist(s) to export. You can export a Post-synthesis netlist instead if placement or performance preservation is not required, to provide the most flexibility for the Fitter in the top-level design. Select Post-fit to preserve the placement and performance of the lower-level design block, and turn on Export routing to include the routing information, if required. One .qxp can include both post-synthesis and post-fitting netlists.

6. Provide the .qxp to the project lead.

Finally, as the project lead in this scenario, perform the appropriate set of steps to import the .qxp files received from designers of each partition.

If you are using makefiles with the design partition scripts, perform the following steps:

1. Use the master_makefile command to export each partition and create .qxp files, and then import them into the top-level design.

2. The software does not have all the information about which source files should be associated with which partition, so you must specify this information in the makefile. The software cannot rebuild the project if source files change unless you specify the dependencies.

If you are not using makefiles, perform the following steps:

1. Add the .qxp as a source file in the Quartus II project, to replace any empty wrapper file for the previously Empty partition.

2. Change the netlist type for the partition from Empty to the required level of results preservation.

**Resolving Assignment Conflicts During Integration**

When integrating lower-level design blocks, the project lead may notice some assignment conflicts. This can occur, for example, if the lower-level design block designers changed their LogicLock regions to account for additional logic or placement constraints, or if the designers applied I/O port timing constraints that differ from constraints added to the top-level design by the project lead. The project
lead can address these conflicts by explicitly importing the partitions into the top-level design, and using options in the Advanced Import Settings dialog box, as described in “Advanced Importing Options” on page 3–33. After the project lead obtains the .qxp for each lower-level design block from the other designers, use the Import Design Partition command on the Project menu and specify the partition in the top-level design that is represented by the lower-level design block .qxp. Repeat this import process for each partition in the design. After you have imported each partition once, you can select all the design partitions and use the Reimport using latest import files at previous locations option to import all the files from their previous locations at one time. To address assignment conflicts, the project lead can take one or both of the following actions:

■ Allow new assignments to be imported

■ Allow existing assignments to be replaced or updated

When LogicLock region assignment conflicts occur, the project lead may take one of the following actions:

■ Allow the imported region to replace the existing region

■ Allow the imported region to update the existing region

■ Skip assignment import for regions with conflicts

If the placement of different lower-level design blocks conflict, the project lead can also set the set the partition’s Fitter Preservation Level to Netlist Only, which allows the software to re-perform placement and routing with the imported netlist.

**Importing a Partition to be Instantiated Multiple Times**

In this variation of the design scenario, one of the lower-level design blocks is instantiated more than once in the top-level design. The designer of the lower-level design block may want to compile and optimize the entity once under a partition, and then import the results as multiple partitions in the top-level design.

If you import multiple instances of a lower-level design block into the top-level design, the imported LogicLock regions are automatically set to Floating status.

If you resolve conflicts manually, you can use the import options and manual LogicLock assignments to specify the placement of each instance in the top-level design.

**Performing Design Iterations With Lower-Level Partitions**

Scenario background: A project consists of several lower-level subdesigns that have been exported from separate Quartus II projects and imported into the top-level design. In this example, integration at the top level has failed because the timing requirements are not met. The timing requirements might have been met in each individual lower-level project, but critical inter-partition paths in the top-level design are causing timing requirements to fail.

After trying various optimizations in the top-level design, the project lead determines that the design cannot meet the timing requirements given the current partition placements that were imported. The project lead decides to pass additional information to the lower-level partitions to improve the placement.
Use this flow if you re-optimize partitions exported from separate Quartus II projects by incorporating additional constraints from the integrated top-level design.

The best way to provide top-level design information to designers of lower-level partitions is to provide the complete top-level project framework using the following steps:

1. For all partitions other than the one(s) being optimized by a designer(s) in a separate Quartus II project(s), set the netlist type to **Post-Fit**.
2. Make the top-level design directory available in a shared source control system, if possible. Otherwise, copy the entire top-level design project directory (including database files), or create a project archive including the post-compilation database.
3. Provide each partition designer with a checked-out version or copy of the top-level design.
4. The partition designers recompile their designs within the new project framework that includes the rest of the design’s placement and routing information as well top-level resource allocations and assignments, and optimize as needed.
5. When the results are satisfactory and the timing requirements are met, export the updated partition as a `.qxp`.

If this design flow is not possible, you can generate partition-specific scripts for individual designs to provide information about the top-level project framework with these steps:

1. In the top-level design, on the Project menu, click **Generate Design Partition Scripts**, or launch the script generator from Tcl or the command line.
2. If lower-level projects have already been created for each partition, you can turn off the **Create lower-level project if one does not exist** option.
3. Make additional changes to the default script options, as necessary. Altera recommends that you pass all the default constraints, including LogicLock regions, for all partitions and virtual pin location assignments. Altera also recommends that you add a maximum delay timing constraint for the virtual I/O connections in each partition.
4. The Quartus II software generates Tcl scripts for all partitions, but in this scenario, you would focus on the partitions that make up the cross-partition critical paths. The following assignments are important in the script:
   - Virtual pin assignments for module pins not connected to device I/O ports in the top-level design.
   - Location constraints for the virtual pins that reflect the initial top-level placement of the pin’s source or destination. These help make the lower-level placement “aware” of its surroundings in the top-level design, leading to a greater chance of timing closure during integration at the top level.
   - `INPUT_MAX_DELAY` and `OUTPUT_MAX_DELAY` timing constraints on the paths to and from the I/O pins of the partition. These constrain the pins to optimize the timing paths to and from the pins.
5. The partition designers source the file provided by the project lead.
   - To source the Tcl script from the Quartus II GUI, on the Tools menu, click **Utility Windows** and open the Tcl console. Navigate to the script’s directory, and type the following command: `source <filename>`
   - To source the Tcl script at the system command prompt, type the following command: `quartus_cdb -t <filename>.tcl`

6. The partition designers recompile their designs with the new project information or assignments and optimize as needed. When the results are satisfactory and the timing requirements are met, export the updated partition as a `.qxp`.

   The project lead obtains the updated `.qxp` files from the partition designers and adds them to the top-level design. When a new `.qxp` is added to the files list, the software will detect the change in the “source file” and use the new `.qxp` results during the next compilation. If the project uses the advanced import flow, the project lead must perform another import of the new `.qxp`.

   You can now analyze the design to determine whether the timing requirements have been achieved. Because the partitions were compiled with more information about connectivity at the top level, it is more likely that the inter-partition paths have improved placement which helps to meet the timing requirements.

**Creating a Design Floorplan With LogicLock Regions**

A floorplan represents the layout of the physical resources on the device. Creating a design floorplan, or floorplanning, describe the process of mapping the logical design hierarchy onto physical regions in the device floorplan. After you have partitioned the design, you can create floorplan location assignments for the design to improve the quality of results when using the incremental compilation design flow. Creating a design floorplan is not a requirement to use an incremental compilation flow, but it is recommended in certain cases. Floorplan location planning can be important for a design that uses incremental compilation for the following reasons:

   - To avoid resource conflicts between partitions, predominantly when partitions are imported from another Quartus II project

   Design floorplan assignments prevent the situation in which the Fitter must place a partition in an area of the device where most resources are already used by other partitions. A physical region assignment provides a reasonable region to re-place logic after a change, so the Fitter does not have to scatter logic throughout the available space in the device.

   Floorplan assignments are not required for non-critical partitions compiled as part of the top-level design. The logic for partitions that are not timing-critical (such as simple top-level glue logic) can be placed anywhere in the device on each recompilation, if that is best for your design.
The simplest way to create a floorplan for a partitioned design is to create one LogicLock region per partition (including the top-level partition). If you have a compilation result for a partitioned design with no LogicLock regions, you can use the Chip Planner with the Design Partition Planner to view the partition placement in the device floorplan. You can draw regions in the floorplan that match the general location and size of the logic in each partition. Or, initially, you can set each region with the default settings of *Auto* size and *Floating* location to allow the Quartus II software to determine the preliminary size and location for the regions. Then, after compilation, use the Fitter-determined size and origin location as a starting point for your design floorplan. Check the quality of results obtained for your floorplan location assignments and make changes to the regions as needed. Alternatively, you can perform synthesis, and then set the regions to the required size based on resource estimates. In this case, use your knowledge of the connections between partitions to place the regions in the floorplan.

Once you have created an initial floorplan, you can refine the region using tools in the Quartus II software. You can also use advanced techniques such as creating non-rectangular regions by merging LogicLock regions.

For more information about when creating a design floorplan can be important, as well as guidelines for creating the floorplan, refer to the *Best Practices for Incremental Compilation Partitions and Floorplan Assignments* chapter in volume 1 of the *Quartus II Handbook*.

You can use the Incremental Compilation Advisor to check that your LogicLock regions meet Altera’s guidelines, as described in “Incremental Compilation Advisor” on page 3–19.

**Creating and Manipulating LogicLock Regions**

Options in the *LogicLock Regions Properties* dialog box, available from the Assignments menu, allow you to enter specific sizing and location requirements for a region. You can also view and refine the size and location of LogicLock regions in the Quartus II Chip Planner. You can select a region in the graphical interface in the Chip Planner and use handles to move or resize the region.

Options in the *Layer Settings* panel in the Chip Planner allow you to create, delete, and modify tasks to determine which objects, including LogicLock regions and design partitions, to display in the Chip Planner.

For more information about creating and viewing LogicLock regions in the LogicLock Regions window and Chip Planner, refer to *Creating and Manipulating LogicLock Regions* in Quartus II Help.
Changing Partition Placement with LogicLock Changes

When a partition is assigned to a LogicLock region as part of a design floorplan, you can modify the placement of a post-fit partition by moving the LogicLock region. As described in “What Changes Initiate the Automatic Resynthesis of a Partition?” on page 3–24, most assignment changes do not initiate a recompilation of a partition if the netlist type specifies that Fitter results should be preserved. For example, changing a pin assignment does not initiate a recompilation; therefore, the design does not use the new pin assignment unless you change the netlist type to Post-Synthesis or Source File.

Similarly, if a partition’s placement is preserved, and the partition is assigned to a LogicLock region, the Fitter always reuses the corresponding LogicLock region size specified in the post-fit netlist. That is, changes to the LogicLock Size setting do not initiate refitting if a partition’s placement is preserved with the Post-Fit netlist type, or with .qxp that includes post-fit information.

However, you can use the LogicLock Origin location assignment to change or fine-tune the previous Fitter results. When you change the Origin setting for a region, the Fitter can move the region in the following manner, depending upon how the placement is preserved for that region’s members:

- When you set a new region Origin, the Fitter uses the new origin and replaces the logic, preserving the relative placement of the member logic.
- When you set the region Origin to Floating, the following conditions apply:
  - If the region’s member placement is preserved with an imported partition, the Fitter chooses a new Origin and re-places the logic, preserving the relative placement of the member logic within the region.
  - If the region’s member placement is preserved with a Post-Fit netlist type, the Fitter does not change the Origin location, and reuses the previous placement results.

Taking Advantage of the Early Timing Estimator

When creating a floorplan you can take advantage of the Early Timing Estimator to enable quick compilations of the design while creating assignments. The Early Timing Estimator feature provides a timing estimate for a design without having to run a full compilation. You can use the Chip Planner to view the “placement estimate” created by this feature, identify critical paths by locating from the timing analyzer reports, and, if necessary, add or modify floorplan constraints. You can then rerun the Early Timing Estimator to quickly assess the impact of any floorplan location assignments or logic changes, enabling rapid iterations on design variants to help you find the best solution. This faster placement has an impact on the quality of results. If getting the best quality of results is important in a given design iteration, perform a full compilation with the Fitter instead of using the Early Timing Estimate feature.
Incremental Compilation Restrictions

This section documents the following restrictions and limitations that you may encounter when using incremental compilation, including interactions with other Quartus II features:

- “When Timing Performance May Not Be Preserved Exactly” on page 3–47
- “When Placement and Routing May Not Be Preserved Exactly” on page 3–47
- “Using Incremental Compilation With Quartus II Archive Files” on page 3–48
- “Limitations for HardCopy Compilation and Migration Flows” on page 3–48
- “Formal Verification Support” on page 3–49
- “SignalProbe Pins and Engineering Change Orders” on page 3–49
- “SignalTap II Logic Analyzer in Exported Partitions” on page 3–49
- “External Logic Analyzer Interface in Exported Partitions” on page 3–50
- “Assignments Made in HDL Source Code in Exported Partitions” on page 3–50
- “Design Partition Script Limitations” on page 3–50
- “Restrictions on MegafUNCTION Partitions” on page 3–52
- “Register Packing and Partition Boundaries” on page 3–53
- “I/O Register Packing” on page 3–53

When Timing Performance May Not Be Preserved Exactly

Timing performance might change slightly in a partition with placement and routing preserved when other partitions are incorporated or re-placed and routed. Timing changes are due to changes in parasitic loading or crosstalk introduced by the other (changed) partitions. These timing changes are very small, typically less than 30 ps on a timing path. Additional fan-out on routing lines when partitions are added can also degrade timing performance.

To ensure that a partition continues to meet its timing requirements when other partitions change, a very small timing margin might be required. The Fitter automatically works to achieve such margin when compiling any design, so you do not need to take any action.

When Placement and Routing May Not Be Preserved Exactly

The Fitter may have to refit affected nodes if the two nodes are assigned to the same location, due to imported netlists or empty partitions set to re-use a previous post-fit netlist. There are two cases in which routing information cannot be preserved exactly. First, when multiple partitions are imported, there might be routing conflicts because two lower-level blocks could be using the same routing wire, even if the floorplan assignments of the lower-level blocks do not overlap. These routing conflicts are automatically resolved by the Quartus II Fitter re-routing on the affected nets. Second, if an imported LogicLock region is moved in the top-level design, the relative placement of the nodes is preserved but the routing cannot be preserved, because the routing connectivity is not perfectly uniform throughout a device.
Using Incremental Compilation With Quartus II Archive Files

The post-synthesis and post-fitting netlist information for each design partition is stored in the project database, the `incremental_db` directory. When you archive a project, the database information is not included in the archive unless you include the compilation database in the `.qar` file.

If you want to re-use post-synthesis or post-fitting results, include the database files in the `Archive Project` dialog box so compilation results are preserved. Click `Advanced`, and choose a file set that includes the compilation database, or turn on `Incremental compilation database files` to create a Custom file set.

When you include the database, the file size of the `.qar` archive file may be significantly larger than an archive without the database.

The netlist information for imported partitions is already saved in the corresponding `.qxp`. Imported `.qxp` files are automatically saved in a subdirectory called `imported_partitions`, so you do not need to archive the project database to keep the results for imported partitions. When you restore a project archive, the partition is automatically reimported from the `.qxp` in this directory if it is available.

For new device families with advanced support, a version-compatible database might not be available. In this case, the archive will not include the compilation database. If you require the database files to reproduce the compilation results in the same Quartus II version, you can use the following command-line option to archive a full database:

```
quartus_sh --archive -use_file_set full_db [-revision <revision name>] <project name>
```

Limitations for HardCopy Compilation and Migration Flows

Incremental compilation within a single Quartus II project is supported for the base family in HardCopy migration flows for both the FPGA first and HardCopy first flows. Design partition assignments are migrated to the companion device. However, you cannot make changes to the design after migration because the design would not match the compilation results for the base family. Therefore, you can perform incremental compilation on one device family, but cannot add new partitions or remove existing partitions after migration.

The Netlist Only preservation level is not supported for Post-fit netlists for FPGA or HardCopy ASIC device compilations when a migration device is specified (that is, for HardCopy ASIC device compilations with a FPGA migration device, or FPGA device compilations with a HardCopy ASIC migration device).

Exporting and importing partitions is not supported in HardCopy ASIC or FPGA device compilations when there is a migration device setting.

The Revision Compare feature requires that the HardCopy ASIC and FPGA netlists are the same. Therefore, all operations performed on one revision must also occur on the other revision. This is accomplished by logging all operations and replaying them on the other revision. Importing partitions does not support this requirement. You can often use `Empty` partitions to implement behavior similar to an exported partition flow, as long as you do not change any global assignments between compilations. All global assignments must be the same for all compiled partitions, so the assignments can be reproduced in the companion device after migration.
Formal Verification Support

You cannot use design partitions for incremental compilation if you are creating a netlist for a formal verification tool.

SignalProbe Pins and Engineering Change Orders

ECO and SignalProbe changes are performed only during ECO and SignalProbe compilations. Other compilation flows do not preserve these netlist changes.

When incremental compilation is turned on and your design contains one or more design partitions, partition boundaries are ignored while making ECO changes and SignalProbe signal settings. However, the presence of ECO and/or SignalProbe changes does not affect partition boundaries for incremental compilation. During subsequent compilations, ECO and SignalProbe changes are not preserved regardless of the Netlist Type or Fitter Preservation Level settings. To recover ECO changes and SignalProbe signals, you must use the Change Manager to re-apply the ECOs after compilation.

For partitions developed independently in separate Quartus II projects, the exported netlist includes all currently saved ECO changes and SignalProbe signals. If you make any ECO or SignalProbe changes that affect the interface to the lower-level partition, the software issues a warning message during the export process that this netlist does not work in the top-level design without modifying the top-level HDL code to reflect the lower-level change. After integrating the .qxp partition into the top-level design, the ECO changes will not appear in the Change Manager.

For more information about using the SignalProbe feature to debug your design, refer to the Quick Design Debugging Using SignalProbe chapter in volume 3 of the Quartus II Handbook. For more information about using the Chip Planner and the Resource Property Editor to make ECOs, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

SignalTap II Logic Analyzer in Exported Partitions

You can use the SignalTap II Embedded Logic Analyzer in any project that you can compile and program into an Altera device.

When incremental compilation is turned on, debugging logic is added to your design incrementally and you can tap post-fitting nodes and modify triggers and configuration without recompiling the full design. Use the appropriate filter in the Node Finder to find your node names. Use SignalTap II: post-fitting if the netlist type is Post-Fit to incrementally tap node names in the post-fit netlist database. Use SignalTap II: pre-synthesis if the netlist type is Source File to make connections to the source file (pre-synthesis) node names when you synthesize the partition from the source code.

If incremental compilation is turned off, the debugging logic is added to the design during Analysis and Elaboration, and you cannot tap post-fitting nodes or modify debug settings without fully compiling the design.
For design partitions that are being developed independently in separate Quartus II projects and contain the logic analyzer, when you export the partition, the Quartus II software automatically removes the SignalTap II logic analyzer and related SLD_HUB logic. You can tap any nodes in a Quartus II project, including nodes within .qxp partitions. Therefore, you can use the logic analyzer within the full top-level design to tap signals from the .qxp partition.

You can also instantiate the SignalTap II megafunction directly in your lower-level design (instead of using an .stp file) and export the entire design to the top level to include the logic analyzer in the top-level design.

For details about using the SignalTap II logic analyzer in an incremental design flow, refer to the Design Debugging Using the SignalTap II Embedded Logic Analyzer chapter in volume 3 of the Quartus II Handbook.

External Logic Analyzer Interface in Exported Partitions

You can use the Logic Analyzer Interface in any project that you can compile and program into an Altera device. You cannot export a partition that uses the Logic Analyzer Interface. You must disable the Logic Analyzer Interface feature and recompile the design before you export the design as a partition.

For more information about the Logic Analyzer Interface, refer to the In-System Debugging Using External Logic Analyzers chapter in volume 3 of the Quartus II Handbook.

Assignments Made in HDL Source Code in Exported Partitions

Assignments made with I/O primitives or the altera_attribute HDL synthesis attribute in lower-level partitions are passed to the top-level design, but do not appear in the top-level .qsf file or Assignment Editor. These assignments are considered part of the source netlist files. You can override assignments made in these source files by changing the value with an assignment in the top-level design.

Design Partition Script Limitations

The Quartus II software has some additional limitations related to the design partition scripts described in “Generating Design Partition Scripts” on page 3–30.

Warnings About Extra Clocks Due to Design Partition Scripts

The generated scripts include applicable clock information for all clock signals in the top-level design. Some of those clocks may not exist in the lower-level projects, so you may see warning messages related to clocks that do not exist in the project. You can ignore these warnings or edit your constraints so the messages are not generated.
Synopsys Design Constraint Files for the TimeQuest Timing Analyzer in Design Partition Scripts

After you have compiled a design using TimeQuest constraints, and the timing assignments option is turned on in the scripts, a separate Tcl script is generated to create an .sdc file for each lower-level project. This script includes only clock constraints and minimum and maximum delay settings for the TimeQuest Timing Analyzer.

PLL settings and timing exceptions are not passed to lower-level designs in the scripts. For suggestions on managing SDC constraints between top-level and lower-level projects, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

Wildcard Support in Design Partition Scripts

When applying constraints with wildcards, note that wildcards are not analyzed across hierarchical boundaries. For example, an assignment could be made to these nodes: Top|A:inst|B:inst|*, where A and B are lower-level partitions, and hierarchy B is a child of A, that is B is instantiated in hierarchy A. This assignment is applied to modules A, B, and all children instances of B. However, the assignment Top|A:inst|B:inst* is applied to hierarchy A, but is not applied to the B instances because the single level of hierarchy represented by B:inst* is not expanded into multiple levels of hierarchy. To avoid this issue, ensure that you apply the wildcard to the hierarchical boundary if it should represent multiple levels of hierarchy.

When using the wildcard to represent a level of hierarchy, only single wildcards are supported. This means assignments such as Top|A:inst*|B:inst* are not supported. The Quartus II software issues a warning in these cases.

Derived Clocks and PLLs in Design Partition Scripts

If a clock in the top level is not directly connected to a pin of a lower-level partition, the lower-level partition does not receive assignments and constraints from the top-level pin in the design partition scripts.

This issue is of particular importance for clock pins that require timing constraints and clock group settings. Problems can occur if your design uses logic or inversion to derive a new clock from a clock input pin. Make appropriate timing assignments in your lower-level Quartus II project to ensure that clocks are not unconstrained.

If the lower-level design uses the top-level project framework from the project lead, the design will have all the required information about the clock and PLL settings. Otherwise, if you use a PLL in your top-level design and connect it to lower-level partitions, the lower-level partitions do not have information about the multiplication or phase shift factors in the PLL. Make appropriate timing assignments in your lower-level Quartus II project to ensure that clocks are not unconstrained or constrained with the incorrect frequency. Alternatively, you can manually duplicate the top-level derived clock logic or PLL in the lower-level design file to ensure that you have the correct multiplication or phase-shift factors, compensation delays and other PLL parameters for complete and accurate timing analysis. Create a design partition for the rest of the lower-level design logic for export to the top level. When the lower-level design is complete, export only the partition that contains the relevant logic.
Pin Assignments for GXB and LVDS Blocks in Design Partition Scripts

Pin assignments for high-speed GXB transceivers and hard LVDS blocks are not written in the scripts. You must add the pin assignments for these hard IP blocks in the lower-level projects manually.

Virtual Pin Timing Assignments in Design Partition Scripts

Design partition scripts use `INPUT_MAX_DELAY` and `OUTPUT_MAX_DELAY` assignments to specify inter-partition delays associated with input and output pins, which would not otherwise be visible to the project. These assignments require that the software specify the clock domain for the assignment and set this clock domain to “*”.

This clock domain assignment means that there may be some paths constrained and reported by the timing analysis engine that are not required. To restrict which clock domains are included in these assignments, edit the generated scripts or change the assignments in your lower-level Quartus II project. In addition, because there is no known clock associated with the delay assignments, the software assumes the worst-case skew, which makes the paths seem more timing critical than they are in the top-level design. To make the paths appear less timing-critical, lower the delay values from the scripts. If required, enter negative numbers for input and output delay values.

Top-Level Ports that Feed Multiple Lower-Level Pins in Design Partition Scripts

When a single top-level I/O port drives multiple pins on a lower-level module, it unnecessarily restricts the quality of the synthesis and placement at the lower-level. This occurs because in the lower-level design, the software must maintain the hierarchical boundary and cannot use any information about pins being logically equivalent at the top level. In addition, because I/O constraints are passed from the top-level pin to each of the children, it is possible to have more pins in the lower level than at the top level. These pins use top-level I/O constraints and placement options that might make them impossible to place at the lower level. The software avoids this situation whenever possible, but it is best to avoid this design practice to avoid these potential problems. Restructure your design so that the single I/O port feeds the design partition boundary and the single connection is split into multiple signals within the lower-level partition.

Restrictions on Megafunction Partitions

The Quartus II software does not support partitions for megafunction instantiations. If you use the MegaWizard™ Plug-In Manager to customize a megafunction variation, the MegaWizard-generated wrapper file instantiates the megafunction. You can create a partition for the MegaWizard-generated megafunction custom variation wrapper file.

The Quartus II software does not support creating a partition for inferred megafucntions (that is, where the software infers a megafunction to implement logic in your design). If you have a module or entity for the logic that is inferred, you can create a partition for that hierarchy level in the design.
The Quartus II software does not support creating a partition for any Quartus II internal hierarchy that is dynamically generated during compilation to implement the contents of a megafunction.

**Register Packing and Partition Boundaries**

The Quartus II software performs register packing during compilation automatically. However, when incremental compilation is enabled, logic in different partitions cannot be packed together because partition boundaries might prevent cross-boundary optimization. This restriction applies to all types of register packing, including I/O cells, DSP blocks, sequential logic, and unrelated logic. Similarly, logic from two partitions cannot be packed into the same ALM.

**I/O Register Packing**

Cross-partition register packing of I/O registers is allowed in certain cases where your input and output pins exist in the top-level hierarchy (and the Top partition), but the corresponding I/O registers exist in other partitions.

The following specific circumstances are required for input pin cross-partition register packing:

- The input pin feeds exactly one register.
- The path between the input pin and register includes only input ports of partitions that have one fan-out each.

The following specific circumstances are required for output register cross-partition register packing:

- The register feeds exactly one output pin.
- The output pin is fed by only one signal.
- The path between the register and output pin includes only output ports of partitions that have one fan-out each.

Output pins with an output enable signal cannot be packed into the device I/O cell if the output enable logic is part of a different partition from the output register. To allow register packing for output pins with an output enable signal, structure your HDL code or design partition assignments so that the register and tri-state logic are defined in the same partition.

Bidirectional pins are handled in the same way as output pins with an output enable signal. If the registers that need to be packed are in the same partition as the tri-state logic, you can perform register packing.

The restrictions on tri-state logic exist because the I/O atom (device primitive) is created as part of the partition that contains tri-state logic. If an I/O register and its tri-state logic are contained in the same partition, the register can always be packed with tri-state logic into the I/O atom. The same cross-partition register packing restrictions also apply to I/O atoms for input and output pins. The I/O atom must feed the I/O pin directly with exactly one signal. The path between the I/O atom and the I/O pin must include only ports of partitions that have one fan-out each.
For more information and examples of cross-partition boundary I/O packing, refer to the
Best Practices for Incremental Compilation Partitions and Floorplan Assignments
chapter in volume 1 of the Quartus II Handbook.

Scripting Support

You can run procedures and make settings described in this chapter in a Tcl script or
at a command-line prompt. This section provides scripting examples that cover some
of the topics discussed in this chapter.

Tcl Scripting and Command-Line Examples

For information about the `::quartus:: incremental_compilation` Tcl package that
contains a set of functions for manipulating design partitions and settings related to
the incremental compilation feature, refer to `::quartus:: incremental_compilation` in
Quartus II Help.

For scripting support information, including design examples and training, refer to
the Quartus II Software Scripting Support page of the Altera website. For detailed Tcl
scripting and command-line information, including design examples, refer to the Tcl
Scripting and Command-Line Scripting chapters in volume 2 of the Quartus II Handbook.

Creating Design Partitions

To create a design partition to a specified hierarchy name, use the following
command:

```
create_partition [-h | -help] [-long_help] -contents <hierarchy name>
-partition <partition name>
```

<table>
<thead>
<tr>
<th>Argument</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-h</td>
<td>-help</td>
</tr>
<tr>
<td>-long_help</td>
<td>Long help with examples and possible return values</td>
</tr>
<tr>
<td>-contents &lt;hierarchy name&gt;</td>
<td>Partition contents (hierarchy assigned to Partition)</td>
</tr>
<tr>
<td>-partition &lt;partition name&gt;</td>
<td>Partition name</td>
</tr>
</tbody>
</table>

Enabling or Disabling Design Partition Assignments During Compilation

To direct the Quartus II Compiler to enable or disable design partition assignments
during compilation, use the following command:

```
set_global_assignment -name IGNORE_PARTITIONS <value>
```
Chapter 3: Quartus II Incremental Compilation for Hierarchical and Team-Based Design

Scripting Support

Setting the Netlist Type

To set the partition netlist type, use the following command:

```
set_global_assignment -name PARTITION_NETLIST_TYPE <value> \\
-section_id <partition name>
```

The `PARTITION_NETLIST_TYPE` command accepts the following values: `SOURCE`, `POST_SYNTH`, `POST_FIT`, and `EMPTY`. For descriptions for these values, refer to “Partition Netlist Type Settings” on page 3–21.

Setting the Fitter Preservation Level for a Post-fit or Imported Netlist

To set the Fitter Preservation Level for a post-fit or imported netlist, use the following command:

```
set_global_assignment -name PARTITION_FITTER_PRESERVATION_LEVEL \\
-<value> -section_id <partition name>
```

The `PARTITION_FITTER_PRESERVATION` command accepts the following values: `NETLIST_ONLY`, `PLACEMENT`, and `PLACEMENT_AND_ROUTING`. For descriptions for these values, refer to “Fitter Preservation Level Settings” on page 3–22.

Preserving High-Speed Optimization

To preserve high-speed optimization for tiles contained within the selected partition, use the following command:

```
set_global_assignment -name PARTITION_PRESERVE_HIGH_SPEED_TILES_ON
```

Specifying the Software Should Use the Specified Netlist and Ignore Source File Changes

To specify that the software should use the specified netlist and ignore source file changes, even if the source file has changed since the netlist was created, use the following command:

```
set_global_assignment -name PARTITION_IGNORE_SOURCE_FILE_CHANGES ON \\
-section_id "<partition name>".
```
Reducing Opening a Project, Creating Design Partitions, and Performing an Initial Compilation

Scenario background: You open a project called AB_project, set up two design partitions, entities A and B, and then perform an initial full compilation.

Example 3–1. AB_project

```bash
set project AB_project
load_package incremental_compilation
load_package flow
project_open $project

# Ensure that design partition assignments are not ignored
set_global_assignment -name IGNORE_PARTITIONS OFF

# Set up the partitions
create_partition -contents A -name "Partition_A"
create_partition -contents B -name "Partition_B"

# Set the netlist types to post-fit for subsequent
# compilations (all partitions are compiled during the
# initial compilation since there are no post-fit
# netlists)
set_partition -partition "Partition_A" -netlist_type POST_FIT
set_partition -partition "Partition_B" -netlist_type POST_FIT

# Run initial compilation:
export_assignments
execute_flow -full_compile

project_close
```
Optimizing the Placement for a Timing-Critical Partition

Scenario background: You have run the initial compilation shown in the example script under Example 3–1. You would like to apply Fitter optimizations, such as physical synthesis, only to partition A. No changes have been made to the HDL files. To ensure the previous compilation result for partition B is preserved, and to ensure that Fitter optimizations are applied to the post-synthesis netlist of partition A, set the netlist type of B to Post-Fit (which was already done in the initial compilation, but is repeated here for safety), and the netlist type of A to Post-Synthesis, as shown in the following example:

Example 3–2. AB_project (2)

```bash
set project AB_project

load_package flow
load_package incremental_compilation
load_package project
project_open $project

# Turn on Physical Synthesis Optimization
set_high_effort_fmax_optimization_assignments

# For A, set the netlist type to post-synthesis
set_partition -partition "Partition_A" -netlist_type POST_SYNTH

# For B, set the netlist type to post-fit
set_partition -partition "Partition_B" -netlist_type POST_FIT

# Also set Top to post-fit
set_partition -partition "Top" -netlist_type POST_FIT

# Run incremental compilation:
export_assignments
execute_flow -full_compile

project_close
```

Generating Design Partition Scripts

To generate design partition scripts, use the following script:

```bash
# load required package
load_package database_manager

# name and open the project
set project <project_path/project_name>
project_open $project

# generate the design partition scripts
generate_bottom_up_scripts <options>

#close project
project_close
```

The options map to the same as those in the Quartus II software in the Generate Design Partition Scripts dialog box. For detailed information about each option, refer to Generate Design Partition Scripts Dialog Box in Quartus II Help.
Exporting a Partition

To open a project and load the `::quartus::incremental_compilation` package before you use the Tcl commands to export a partition to a `.qxp` that contains both a post-synthesis and post-fit netlist, with routing, use the following script:

```tcl
# load required package
load_package incremental_compilation

# open project
project_open <project name>

# export partition to the .qxp and set preservation level
export_partition -partition <partition name> -qxp <.qxp file name> -<options>

# close project
project_close
```

Importing a Partition into the Top-Level Design

To import a `.qxp` into a top-level design, use the following script:

```tcl
# load required packages
load_package incremental_compilation
load_package project
load_package flow

# open project
project_open <project name>

# import partition
import_partition -partition <partition name> -qxp <.qxp file> -<options>

# close project
project_close
```

Makefiles

For an example of how to use incremental compilation with a `makefile` as part of the team-based incremental compilation design flow, refer to the `read_me.txt` file that accompanies the `incr_comp` example located in the `/qdesigns/incr_comp_makefile` subdirectory.

When using a team-based incremental compilation design flow, the Generate Design Partition Scripts dialog box can write makefiles that automatically export lower-level design partitions and import them into the top-level design whenever design files change. For more information about the Generate Design Partition Scripts dialog box, refer to Generate Design Partition Scripts Dialog Box in Quartus II Help.

Conclusion

With the Quartus II incremental compilation feature described in this chapter, you can preserve the results and performance of unchanged logic in your design as you make changes elsewhere. The various applications of incremental compilation enable you to improve your productivity while designing for high-density FPGAs.
Document Revision History

Table 3–6 shows the revision history for this document.

Table 3–6. Document Revision History (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>■ Updated “Tcl Scripting and Command-Line Examples”.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Changed to new document template.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Reorganized Tcl scripting section. Added description for new feature:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Ignore partitions assignments during compilation option.</td>
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<tr>
<td></td>
<td></td>
<td>■ Reorganized “Incremental Compilation Summary” section.</td>
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<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Removed the explanation of the “bottom-up design flow” where designers</td>
</tr>
<tr>
<td></td>
<td></td>
<td>work completely independently, and replaced with Altera’s recommendations</td>
</tr>
<tr>
<td></td>
<td></td>
<td>for team-based environments where partitions are developed in the same</td>
</tr>
<tr>
<td></td>
<td></td>
<td>top-level project framework, plus an explanation of the bottom-up process</td>
</tr>
<tr>
<td></td>
<td></td>
<td>for including independent partitions from third-party IP designers.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Expanded the Merge command explanation to explain how it now accommodates</td>
</tr>
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<td></td>
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<td>cross-partition boundary optimizations.</td>
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<td></td>
<td></td>
<td>■ Restructured Altera recommendations for when to use a floorplan.</td>
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<tr>
<td></td>
<td></td>
<td>■ Added “Viewing the Contents of a Quartus II Exported Partition File (.qxp)” section.</td>
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<tr>
<td></td>
<td></td>
<td>■ Reorganized chapter to make design flow scenarios more visible; integrated into various sections rather than at the end of the chapter.</td>
</tr>
<tr>
<td>October 2009</td>
<td>9.1.0</td>
<td>■ Redefined the bottom-up design flow as team-based and reorganized previous design flow examples to include steps on how to pass top-level design information to lower-level designers.</td>
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<tr>
<td></td>
<td></td>
<td>■ Moved SDC Constraints from Lower-Level Partitions section to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.</td>
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<tr>
<td></td>
<td></td>
<td>■ Reorganized the “Conclusion” section.</td>
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<tr>
<td></td>
<td></td>
<td>■ Removed HardCopy APEX and HardCopy Stratix Devices section.</td>
</tr>
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</table>
Table 3–6. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
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</table>
| March 2009  | 9.0.0   | ■ Split up netlist types table  
■ Moved “Team-Based Incremental Compilation Design Flow” into the “Including or Integrating partitions into the Top-Level Design” section.  
■ Added new section “Including or Integrating Partition into the Top-Level Design”.  
■ Removed “Exporting a Lower-Level Partition that Uses a JTAG Feature” restriction  
■ Other edits throughout chapter                                                                 |
| November 2008 | 8.1.0 | ■ Added new section “Importing SDC Constraints from Lower-Level Partitions” on page 2–44  
■ Removed the Incremental Synthesis Only option  
■ Removed section “OpenCore Plus Feature for MegaCore Functions in Bottom-Up Flows”  
■ Removed section “Compilation Time with Physical Synthesis Optimizations”  
■ Added information about using a .qxp as a source design file without importing  
■ Reorganized several sections  
■ Updated Figure 2–10                                                                 |

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive
The Partial Reconfiguration (PR) feature in the Quartus II software allows you to reconfigure a portion of the FPGA dynamically, while the remainder of the device continues to operate. The Quartus II software supports the PR feature for the Altera® Stratix® V device family. This chapter provides an overview of partial reconfiguration and provides design guidelines for using this feature with supported FPGA devices.

For assistance with support for partial reconfiguration with the Arria® V or Cyclone® V device families, file a service request at mySupport.

This chapter assumes a basic knowledge of Altera's FPGA design flow, incremental compilation and LogicLock™ region features available in the Quartus II software. It also assumes knowledge of the internal FPGA resources such as logic array blocks (LABs), memory logic array blocks (MLABs), memory types (RAM and ROM), DSP blocks, clock networks.

This chapter discusses the following topics:

- “Terminology” on page 4–1
- “An Example of a Partial Reconfiguration Design” on page 4–4
- “Partial Reconfiguration Design Flow” on page 4–7
- “Partial Reconfiguration with an External Host” on page 4–28
- “Partial Reconfiguration with an Internal Host” on page 4–30
- “Managing a Partial Reconfiguration Project” on page 4–31
- “Creating Programming Files for a Partial Reconfiguration Project” on page 4–33
- “Partial Reconfiguration Known Limitations” on page 4–38

**Terminology**

This section describes commonly used terminology in this chapter.

**project**: A Quartus II project contains the design files, settings, and constraints files required for the compilation of your design.

**revision**: In the Quartus II software, a revision is a set of assignments and settings for one version of your design. A Quartus II project can have several revisions, and each revision has its own set of assignments and settings. A revision helps you to organize several versions of your design into a single project.
incremental compilation: This is a feature of the Quartus II software that allows you to preserve results of previous compilations of unchanged parts of the design, while changing the implementation of the parts of your design that you have modified since your previous compilation of the project. The key benefits include timing preservation and compile time reduction by only compiling the logic that has changed.

partition: You can partition your design along logical hierarchical boundaries. Each design partition is independently synthesized and then merged into a complete netlist for further stages of compilation. With the Quartus II incremental compilation flow, you can preserve results of unchanged partitions at specific preservation levels. For example, you can set the preservation levels at post-synthesis or post-fit, for iterative compilations in which some part of the design is changed. A partition is only a logical partition of the design, and does not necessarily refer to a physical location on the device. However, you may associate a partition with a specific area of the FPGA by using a floorplan assignment.

For more information on design partitions, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in the Quartus II Handbook.

LogicLock region: A LogicLock region constrains the placement of logic in your design. You can associate a design partition with a LogicLock region to constrain the placement of the logic in the partition to a specific physical area of the FPGA.

For more information about LogicLock regions, refer to the Analyzing and Optimizing the Design Floorplan with the Chip Planner chapter in the Quartus II Handbook.

PR project: Any Quartus II design project that uses the PR feature.
PR region: A design partition with an associated contiguous LogicLock region in a PR project. A PR project can have one or more PR regions that can be partially reconfigured independently. A PR region may also be referred to as a PR partition.
static region: The region outside of all the PR regions in a PR project that cannot be reprogrammed with partial reconfiguration (unless you reprogram the entire FPGA). This region is called the static region, or fixed region.
persona: A PR region has multiple implementations. Each implementation is called a persona. PR regions can have multiple personas. In contrast, static regions have a single implementation or persona.
PR control block: Dedicated block in the FPGA that processes the PR requests, handshake protocols, and verifies the CRC.

Determining Resources for Partial Reconfiguration
You can use partial reconfiguration to configure only the resources such as LABs, embedded memory blocks, and DSP blocks in the FPGA core fabric that are controlled by configuration RAM (CRAM). The functions in the periphery, such as GPIOs or I/O Registers, are controlled by I/O configuration bits and therefore cannot be partially reconfigured. Clock multiplexors for $GCLK$ and $QCLK$ are also not partially reconfigurable because they are controlled by I/O periphery bits.
Figure 4–1 shows the types of resource blocks in a Stratix V device.

**Figure 4–1. Partially Reconfigurable Resources**

![Diagram of Partially Reconfigurable Resources]

Table 4–1 describes the reconfiguration type supported by each FPGA resource block, which are shown in Figure 4–1.

**Table 4–1. Reconfiguration Modes of the FPGA Resource Block**

<table>
<thead>
<tr>
<th>Hardware Resource Block</th>
<th>Reconfiguration Mode</th>
</tr>
</thead>
<tbody>
<tr>
<td>Logic Block</td>
<td>Partial Reconfiguration</td>
</tr>
<tr>
<td>Digital Signal Processing</td>
<td>Partial Reconfiguration</td>
</tr>
<tr>
<td>Memory Block</td>
<td>Partial Reconfiguration</td>
</tr>
<tr>
<td>Transceivers</td>
<td>Dynamic Reconfiguration ALTGX_Reconfig</td>
</tr>
<tr>
<td>PLL</td>
<td>Dynamic Reconfiguration ALTGX_Reconfig</td>
</tr>
<tr>
<td>Core Routing</td>
<td>Partial Reconfiguration</td>
</tr>
<tr>
<td>Clock Networks</td>
<td>Clock network sources cannot be changed, but a PLL driving a clock network can be dynamically reconfigured</td>
</tr>
<tr>
<td>I/O Blocks and Other Periphery</td>
<td>Not supported</td>
</tr>
</tbody>
</table>

The transceivers and PLLs in Altera FPGAs can be reconfigured using dynamic reconfiguration. For more information on dynamic reconfiguration, refer to the *Dynamic Reconfiguration in Stratix V Devices* chapter in the *Stratix V Handbook*. 
An Example of a Partial Reconfiguration Design

A PR design is divided into two parts. The static region where the design logic does not change, and one or more PR regions. Each PR region can have different design personas, that change with partial reconfiguration.

Figure 4–2 shows the top-level of a PR design, which includes a static region and two PR regions. PR Region A has three personas associated with it; A1, A2, and A3. PR Region B has two personas; B1 and B2. Each persona for the two PR regions can implement different application specific logic, and using partial reconfiguration, the persona for each PR region can be modified without interrupting the operation of the device in the static or other PR region.

When a region can access more than one persona, you must create control logic to swap between personas for a PR region.

Partial Reconfiguration Modes

When you implement a design on an Altera FPGA device, your design implementation is controlled by bits stored in CRAM inside the FPGA. The CRAM bits control individual LABs, MLABs, M20K memory blocks, DSP blocks, and routing multiplexers in a design. The CRAM bits are organized into a frame structure representing vertical areas that correspond to specific locations on the FPGA.

If you change a design and reconfigure the FPGA in a non-PR flow, the process reloads all the CRAM bits to a new functionality.

Configuration bitstreams used in a non-PR flow are different than those used in a PR flow. In addition to standard data and CRC check bits, configuration bitstreams for partial reconfiguration also include instructions that direct the PR control block to process the data for partial reconfiguration.

The configuration bitstream written into the CRAM is organized into configuration frames. If a LAB column passes through multiple PR regions, those regions share some programming frames.

You can use partial reconfiguration in the SCRUB mode or the AND/OR mode. The mode you select affects your PR flow in ways detailed later in this chapter.
SCRUB Mode

In the SCRUB mode, the unchanging CRAM bits from the static region are “scrubbed” back to their original values. They are neither erased nor reset. The static regions controlled by the CRAM bits from the same programming frame as the PR region continue to operate. All the CRAM bits corresponding to a PR region are overwritten with new data, regardless of what was previously contained in the region.

The SCRUB mode of partial reconfiguration involves re-writing all the bits in an entire LAB column of the CRAM, including bits controlling any PR regions above or below the region being reconfigured. As a result, it is not currently possible to correctly determine the bits associated with a PR region above or below the region being reconfigured, because those bits could have already been reconfigured and changed to an unknown value. This restriction does not apply to static bits above or below the PR region, since those bits never change and you can rewrite them with the same value as the current state of the configuration bit. You cannot use the SCRUB mode when two PR regions have a vertically overlapping column in the device.

The advantage of using the SCRUB mode is that the programming file size is much smaller than the AND/OR mode.

Figure 4–3 shows the floorplan of a FPGA using SCRUB mode, with two PR regions, whose columns do not overlap.

AND/OR Mode

The AND/OR mode refers to how the bits are rewritten. Partial reconfiguration with AND/OR uses a two-pass method. Simplistically, this can be compared to bits being ANDed with a MASK, and ORed with new values, allowing multiple PR regions to vertically overlap a single column. In the first pass, all the bits in the CRAM region for a column passing through a PR region are ANDed with 0’s while those outside the PR region are ANDed with 1’s. After the first pass, all the CRAM bits corresponding to
the PR region are reset without modifying the static region. In the second pass for each CRAM frame, new data is ORed with the current value of 0 inside the PR region, and in the static region, the bits are ORed with 0's so they remain unchanged. The programming file size of a PR region using the AND/OR mode could be twice the programming file size of the same PR region using SCRUB mode.

Figure 4–4 shows two PR regions that overlap the same device columns using the AND/OR mode.

**Figure 4–4. Partial Reconfiguration AND/OR Mode**

![Diagram of Partial Reconfiguration AND/OR Mode](image)

If you have overlapping PR regions in your design, you must use AND/OR mode to program all PR regions, including PR regions with no overlap. The Quartus II software will not permit the use of SCRUB mode when there are overlapping regions. If none of your regions overlap, you can use AND/OR, SCRUB, or a mixture of both.

### Programming File Sizes for a Partial Reconfiguration Project

The programming file size for a partial reconfiguration is proportional to the area of the PR region. A partial reconfiguration programming bitstream for AND/OR mode makes two passes on the PR region; the first pass clears all relevant bits, and the second pass sets the necessary bits. Due to this two-pass sequence, the size of a partial bitstream can be larger than a full FPGA programming bitstream depending on the size of the PR region.

When using the AND/OR mode for partial reconfiguration, the formula which describes the approximate file size within ten percent is:

\[
\text{PR bitstream size} = \frac{((\text{Size of region in the horizontal direction})}{(\text{full horizontal dimension of the part})} \times 2 \times (\text{size of full bitstream})
\]

The way the Fitter reserves routing for partial reconfiguration increases the effective size for small PR regions from a bitstream perspective. PR bitstream sizes in designs with a single small PR region will not match the file size computed by this equation.

The PR bitstream size is approximately half of the size computed above when using SCRUB mode.
Partial Reconfiguration Design Flow

The primary building block of partial reconfiguration is the revision. Your initial design is the base revision, where you define the boundaries of the static region and reconfigurable regions on the FPGA. From the base revision, you create multiple revisions, which contain the static region and describe the differences in the reconfigurable regions.

Two types of revisions are specific to partial reconfiguration: reconfigurable and aggregate. Both import the persona for the static region from the base revision. A reconfigurable revision generates personas for PR regions. An aggregate revision is used to combine personas from multiple reconfigurable revisions to create a complete design suitable for timing analysis.

The design flow for partial reconfiguration also utilizes the Quartus II incremental compilation flow. To take advantage of incremental compilation for partial reconfiguration, you must organize your design into logical and physical partitions for synthesis and fitting. For the PR flow, these partitions are treated as PR regions that must also have associated LogicLock assignments.

Revisions make use of personas, which are subsidiary archives describing the characteristics of both static and reconfigurable regions, that contain unique logic which implements a specific set of functions to reconfigure a PR region of the FPGA. Partial reconfiguration uses personas to pass this logic from one revision to another.

The flow chart in Figure 4–5 illustrates the steps involved in the PR flow.

Figure 4–5. Partial Reconfiguration Design Flow

Plan Your System for Partial Reconfiguration

Identify the Design Blocks Designated to be Partially Reconfigured

Code the Design Using HDL

Develop the Personas for the Partial Blocks

Simulate the Design Functionality

Is Functionality Verified?

yes

no

Functionality is Verified?

yes

no

Designate All Partial Block(s) as Design Partition(s) for the Use with Incremental Compilation

Assign All PR Partition(s) to LogicLock Regions

Create Revisions and Compile the Design for Each Revision

Debug the Timing Failure & Revise the Appropriate Step

Is Timing Met for Each Revision?

yes

no

Generate Configuration Files

Program the Device
The PR design flow requires more initial planning than a standard design flow. Planning requires setting up the design logic for partitioning, and determining placement assignments to create a floorplan. Well-planned partitions can help improve design area utilization and performance, and make timing closure easier. You should also decide whether your system requires partial reconfiguration to originate from the FPGA pins or internally, and which mode you are using; the AND/OR mode or the SCRUB mode, because this influences some of the planning steps described in this section.

You must structure your source code or design hierarchy to ensure that logic is grouped correctly for optimization. Implementing the correct logic grouping early in the design cycle is more efficient than restructuring the code later. The PR flow requires you to be more rigorous about following good design practices. The guidelines for creating partitions for incremental compilation also include creating partitions for partial reconfiguration.

Use the following best practice guidelines for designing in the PR flow, which are described in detail in this section:

- Determining resources for partial reconfiguration
- Partitioning the design for partial reconfiguration
- Creating incremental compilation partitions for partial reconfiguration
- Instantiating the PR controller in the design
- Creating wrapper logic for PR regions
- Creating freeze logic for PR regions
- Planning clocks and other global signals for the PR design
- Creating floorplan assignments for the PR design

### Partitioning the Design for Partial Reconfiguration

You must create design partitions for each PR region that you want to partially reconfigure. Optionally, you can also create partitions for the static parts of the design for timing preservation and/or for reducing compilation time.

There is no limit on the number of independent partitions or PR regions you can create in your design. You can designate any partition as a PR partition by enabling that feature in the LogicLock Regions window in the Quartus II software.

### Creating Incremental Compilation Partitions for Partial Reconfiguration

Use the following best practices guidelines when creating partitions for PR regions in your design:

- Register all partition boundaries; register all inputs and outputs of each partition when possible. This practice prevents any delay penalties on signals that cross partition boundaries and keeps each register-to-register timing path within one partition for optimization.
- Minimize the number of paths that cross partition boundaries.
Minimize the timing-critical paths passing in or out of PR regions. If there are
timing-critical paths that cross PR region boundaries, rework the PR regions to
avoid these paths.

The Quartus II software can optimize some types of paths between design
partitions for non-PR designs. However, for PR designs, such inter-partition paths
are strictly not optimized.

For more information about incremental compilation, refer to the Quartus II
Incremental Compilation for Hierarchical and Team-Based Design chapter in the Quartus II
Handbook.

**Instantiating the Partial Reconfiguration Controller in the Design**

You must instantiate the Stratix V PR control block and the Stratix V CRC block in
your design in order to use the PR feature in Stratix V devices. You may find that
adding the PR control block and CRC block at the top level of the design offers the
most convenience. For example, in a design named Core_Top, all the logic is
contained under the Core_Top module hierarchy. Create a wrapper (Chip_Top) at the
top-level of the hierarchy that instantiates this Core_Top module, the Stratix V PR
control block, and the Stratix V CRC check modules.

If you are performing partial reconfiguration from pins, then the required pins should
be on the I/O list for the top-level (Chip_Top) of the project, as shown in the code in
Example 4–1 on page 4–10. If you are performing partial reconfiguration from within
the core, you may choose another configuration scheme, such as Active Serial, to
transmit the reconfiguration data into the core, and then assemble it to 16-bit wide
data inside the FPGA within your logic. In such cases, the PR pins are not part of the
FPGA I/O.
The following code sample has the component declaration in VHDL, showing the ports of the Stratix V PR control block and the Stratix V CRC block. In Example 4–1, the PR function is performed from within the core (code located in Core_Top) and you must add additional ports to Core_Top to connect to both components.

**Example 4–1. Component Declaration of the PR Control Block and CRC Block in VHDL**

```vhdl
-- The Stratix V control block interface
component stratixv_prblock is
    port(
        corectl: in    STD_LOGIC ;
        prrequest: in STD_LOGIC ;
        data : in STD_LOGIC_VECTOR(15 downto 0);
        error : out STD_LOGIC ;
        ready : out STD_LOGIC ;
        done : out STD_LOGIC
    ) ;
end component ;

-- The Stratix V CRC block for diagnosing CRC errors
component stratixv_crcblock is
    port(
        shiftnld: in STD_LOGIC ;
        clk : in STD_LOGIC ;
        crcerror: out STD_LOGIC
    ) ;
end component ;
```

The following rules apply when connecting the PR control block to the rest of your design:

- The **corectl** signal must be set to ‘1’ (when using partial reconfiguration from core) or to ‘0’ (when using partial reconfiguration from pins).
- The **corectl** signal has to match the **Enable PR pins** option setting in the Device and Pin Options dialog box on the Setting page; if you have turned on **Enable PR pins**, then the **corectl** signal on the PR control block instantiation must be toggled to ‘0’.
- When performing partial reconfiguration from pins the Quartus II software automatically assigns the PR unassigned pins. If you so choose, you can make pin assignments to all the dedicated PR pins in **Pin Planner** or **Assignment Editor**.
- When performing partial reconfiguration from core, you can connect the **prblock** signals to either core logic or I/O pins, excluding the dedicated programming pin such as **DCLK**.

Verilog HDL does not require a component declaration. You can instantiate the PR control block as shown in Example 4–2 on page 4–11.
Example 4-2 shows how to instantiate PR control blocks inside your top-level project, Chip_Top, in Verilog HDL:

**Example 4-2. Instantiating the PR Control Block and CRC Block in Verilog HDL**

```verilog
module Chip_Top (  
//User I/O signals (excluding PR related signals)  
  ..  
  ..  
  //PR interface & configuration signals  
  pr_request,  
  pr_ready,  
  pr_done,  
  crc_error,  
  dclk,  
  pr_data,  
  init_done  
);  
//user I/O signal declaration  
  ..  
  ..  
  //PR interface and configuration signals declaration  
  input pr_request;  
  output pr_ready;  
  output pr_done;  
  output crc_error;  
  input dclk;  
  input [15:0] pr_data;  
  output init_done  
  // Following shows the connectivity within the Chip_Top module  
Core_Top : Core_Top  
port_map (  
  ..  
  ..  
);  
m_pr : stratixv_prblock  
port_map(  
  clk => dclk,  
  corectl=> '0', //1 - when using PR from inside  
                     //0 - for PR from pins; You must also enable  
                     // the appropriate option in Quartus II settings  
  prrequest=> pr_request,  
  data => pr_data,  
  error => pr_error,  
  ready => pr_ready,  
  done => pr_done  
);  
m_crc : stratixv_crcblock  
port_map(  
  shiftnld=> '1', //If you want to read the EMR register when  
                //error occurs, refer to AN539 for the  
                //connectivity forthis signal. If you only want  
                //to detect CRC errors, but plan to take no  
                //further action, you can tie the shiftnld  
                //signal to logical high.  
  clk=> dummy_clk,  
  crcerror=> crc_error  
);```
Example 4–3 shows how to instantiate PR control blocks inside your top-level project, Chip_Top, in VHDL.

**Example 4–3. Instantiating the PR Control Block in VHDL (Part 1 of 2)**

```vhdl
entity Chip_Top is
  --PR INTERFACE & CONFIGURATION SIGNALS
  port(
    pr_request: in STD_LOGIC;
    dclk : in STD_LOGIC;
    pr_data : in STD_LOGIC_VECTOR (15 downto 0);
    pr_ready: out STD_LOGIC;
    pr_done : out STD_LOGIC;
    crc_error:out STD_LOGIC;
    init_done:out STD_LOGIC;
    ..
    ..
  )
end Chip_top;

--declare all components to be instantiated

component Core_Top is
  port( ..
    ..
  );
end component ;

component stratixv_prblock is
  port( corectl: in STD_LOGIC; --1 - when using PR from inside
      --0 for PR from pins; You must also enable
      --the appropriate option in Quartus II settings
      clk: in STD_LOGIC;
      prrequest:in STD_LOGIC;
      data: in STD_LOGIC_VECTOR (15 downto 0);
      error: out STD_LOGIC;
      ready: out STD_LOGIC;
      done: out STD_LOGIC);
end component;

component stratixv_crcblock is
  port(shiftnld:in  STD_LOGIC;
      clk: in STD_LOGIC;
      crcerror:out STD_LOGIC);
end component;

architecture struct of Chip_Top is
  -- signal declaration
  ..
 ..
  signal pr_error: STD_LOGIC;
  signal crcblkclk: STD_LOGIC; -- this can be a dummy clock, but
  ..
  ..
end struct;
```
Example 4–3. Instantiating the PR Control Block in VHDL (Part 2 of 2)

```vhdl
-- component instantiation
begin
  m_pr : stratixv_prblock
  port map
    (corecktl =>'0',-- pr from pins
     clk => dclk, -- uses the dclk as pr clock
     prrequest=>pr_request,
     data => pr_data,
     error => pr_error,
     ready => pr_ready,
     done => pr_done);

  m_crc: stratixv_crcblock
  port map
    (shiftnld=>'1',
     clk => crcblkclk,
     (shiftnld=> '1', -- If you want to read the EMR register
     clk=>  dummy_clk, -- when an error occurs, refer to
               -- AN539 for the connectivity for this
               -- signal. If you only want to
               -- detect CRC errors, but intend
               -- no further action, you can tie the
               -- shiftnld signal to logical high.
     crcerror=>crc_error);

  n_Core_Top: Core_Top
  port map ( )
  -- preferably no other glue logic exists at this level
end struct;
```

For more information on port connectivity for reading the Error Message Register (EMR), refer to AN539: Test Methodology of Error Detection and Recovery using CRC in Altera FPGA Devices.

Creating Wrapper Logic for PR Regions

Each persona of a PR region must implement the same input and output boundary ports. These ports act as the boundary between static and reconfigurable logic. Implementing the same boundary ports ensures that all ports of a PR region remain stationary regardless of the underlying persona, so that the routing from the static logic does not change with different PR persona implementations.
The Quartus II software automatically instantiates a wire-LUT for each port of the PR region to lock down the same location for all instances of the PR persona, as shown in Figure 4–6.

**Figure 4–6. Wire-LUTs at PR Region Boundary**

If one persona of your PR region has a different number of ports than others, then you must create a wrapper so that the static region always communicates with this wrapper. In this wrapper, you can create dummy ports to ensure that all of the PR personas of a PR region have the same connection to the static region.

The sample code in Example 4–4 on page 4–15 and Example 4–5 on page 4–16 each create two personas; persona_1 and persona_2 are different functions of one PR region. Note that one persona has a few dummy ports.
Example 4–4 is sample partial reconfiguration wrapper logic in Verilog HDL.

**Example 4–4. Wrapper Logic for PR in Verilog HDL**

```verilog
module persona_1
  (  
    input reset, 
    input [2:0] a, 
    input [2:0] b, 
    input [2:0] c, 
    output [3:0] p, 
    output [7:0] q 
  ); 
  reg [3:0] p, q; 
  always@(a or b) begin 
    p = a + b ; 
  end 
  always@(a or b or c or p) begin 
    q = (p*a - b*c ) 
  end 
endmodule

module persona_2 
  (  
    input reset, 
    input [2:0] a, 
    input [2:0] b, 
    input [2:0] c, //never used in this persona 
    output [3:0] p, 
    output [7:0] q //never assigned in this persona 
  ); 
  reg [3:0] p, q; 
  always@(a or b) begin 
    p = a * b; 
    // note q is not assigned value in this persona 
  end 
endmodule
```
Creating Freeze Logic for PR Regions

When you use partial reconfiguration, you must freeze all non-global inputs of a PR region except global clocks. Locally routed signals are not considered global signals, and must also be frozen during partial reconfiguration. Freezing refers to driving a '1' on those PR region inputs. When you start a partial reconfiguration process, the chip is in user mode, with the device still running.
Freezing all non-global inputs for the PR region ensures there is no contention between current values that may result in unexpected behavior of the design after partial reconfiguration is complete. Global signals going into the PR region should not be frozen to high. The Quartus II software freezes the outputs from the PR region; therefore the logic outside of the PR region is not affected. Figure 4–7 shows how to freeze the inputs.

**Figure 4–7. Freezing at PR Region Boundary**

During partial reconfiguration, the static region logic should not depend on the outputs from PR regions to be at a specific logic level for the continued operation of the static region.

The easiest way to control the inputs to PR regions is by creating a wrapper around the PR region in RTL. In addition to freezing all inputs high, you can also drive the outputs from the PR block to a specific value, if required by your design. For example, if the output drives a signal that is active high, then your wrapper could freeze the output to GND.
Example 4–6 implements a freeze wrapper in Verilog HDL, on a module named pr_module.

**Example 4–6. Freeze Wrapper Verilog HDL Sample Code**

```verilog
module freeze_wrapper
(
    input reset,
    input freeze, //PR process active, generated by user logic
    input clk1, //global clock signal
    input clk2, // non-global clock signal
    input [3:0] control_mode, // synchronous to clk1
    input [3:0] framer_ctl, // synchronous to clk2
    output [15:0] data_out
);

reg [3:0]control_mode_sync, framer_ctl_sync;
wire clk2_to_use;

//instantiate pr_module
pr_module pr_module
(
    .reset (reset), //input
    .clk1 (clk1), //input, global clock
    .clk2 (clk2_to_use), // input, non-global clock
    .control_mode (control_mode_sync), //input
    .framer_ctl (framer_ctl_sync), //input
    .pr_module_out (data_out)// collection of outputs from pr_module
);

always@(posedge clk1) begin
    control_mode_sync <= freeze ? 4'hF: control_mode;
end

always@(posedge clk2) begin
    framer_ctl_sync <= freeze ? 4'hF: framer_ctl;
end

assign clk2_to_use = freeze ? 1'b1 : clk2;

endmodule
```
Example 4–7 implements a freeze wrapper in VHDL, on a module named pr_module.

**Example 4–7. Freeze Wrapper VHDL Sample Code**

```vhdl
entity freeze_wrapper is
  port( reset:in STD_LOGIC;
       freeze:in STD_LOGIC;
       clk1:in STD_LOGIC; --global signal
       clk2:in STD_LOGIC; --non-global signal
       control_mode: in STD_LOGIC_VECTOR (3 downto 0);
       framer_ctl:in STD_LOGIC_VECTOR (3 downto 0);
       data_out:out STD_LOGIC_VECTOR (15 downto 0));
end freeze_wrapper;

architecture behv of freeze_wrapper is
  component pr_module
    port(reset:in STD_LOGIC;
         clk1:in STD_LOGIC;
         clk2:in STD_LOGIC;
         control_mode:in STD_LOGIC_VECTOR (3 downto 0);
         framer_ctl:in STD_LOGIC_VECTOR (3 downto 0);
         pr_module_out:out STD_LOGIC_VECTOR (15 downto 0));
  end component
  signal control_mode_sync: in STD_LOGIC_VECTOR (3 downto 0);
  signal framer_ctl_sync  : in STD_LOGIC_VECTOR (3 downto 0);
  signal clk2_to_use : STD_LOGIC;
  signal data_out_temp : STD_LOGIC_VECTOR (15 downto 0);
  begin
    data_out(15 downto 0) <= data_out_temp(15 downto 0);
    m_pr_module: pr_module
      port map (reset => reset,
                 clk1 => clk1,
                 clk2 => clk2,
                 control_mode =>control_mode_sync,
                 framer_ctl=> framer_ctl_sync,
                 pr_module_out=> data_out_temp);
    -- freeze all inputs
    process(clk1) begin
      if clk1'event and clk1 = '1' then
        if freeze = '1' then
          control_mode_sync <= "1111";
        else control_mode_sync <= control_mode;
          end if;
        end if;
      end process;
    -- freeze the non-global clocks as well
    process(clk2, freeze) begin
      if clk2'event and clk2 = '1' then
        if freeze = '1' then
          framer_ctl_sync <= "1111";
        else framer_ctl_sync <= framer_ctl;
          end if;
        end if;
      end process;
  end behv;
```

Planning Clocks and Other Global Signals for a PR Design

For non-PR designs, the Quartus II software automatically promotes high fan-out signals onto available clocks or other forms of global signals using a process called global promotion during the pre-fitter stage of design compilation. For PR designs, however, automatic global promotion is disabled by default for PR regions, and you must assign the global clock resources necessary for PR partitions.

There are 16 global clock networks in a Stratix V device. However, only six unique clocks can drive a row clock region limiting you to a maximum of six global signals in each PR region. The Quartus II software must ensure that any global clock can feed every location in the PR region.

The limit of six global signals to a PR region includes the GCLK, QCLK and PCLKs used inside of the PR region. Make QSF assignments for global signals in your project's Quartus II Settings File (.qsf), based on the clocking requirements for your design. In designs with multiple clocks that are external to the PR region, it may be beneficial to align the PR region boundaries to be within the global clock boundary (such as QCLK or PCLK).

If your PR region requires more than six global signals, modify the region architecture to reduce the number of global signals within this to six or fewer. For example, you can split a PR region into multiple regions, each of which uses only a subset of the clock domains, so that each region does not use more than six.

Every instance of a PR region that uses the global signals (for example, PCLK, QCLK, GCLK, ACLR) must use a global signal for that input.

PR regions are allowed to contain output ports that are used outside of the PR region as global signals.

Global signals can only be used to route certain secondary signals into a PR region. Table 4–2 shows the restrictions for each block. Data signals and other secondary signals not listed in the table, such as synchronous clears and clock enables are not supported.

### Table 4–2. Supported Signal Types for Driving Clock Networks in a PR Region

<table>
<thead>
<tr>
<th>Block Types</th>
<th>Supported Signals for Global/Periphery/Quadrant Clock Networks</th>
</tr>
</thead>
<tbody>
<tr>
<td>LAB</td>
<td>Clock, ACLR</td>
</tr>
<tr>
<td>RAM</td>
<td>Clock, ACLR, Write Enable(WE), Read Enable(RE)</td>
</tr>
<tr>
<td>DSP</td>
<td>Clock, ACLR</td>
</tr>
</tbody>
</table>

If a global signal feeds both static and reconfigurable logic, the restrictions in Table 4–2 also apply to destinations in the static region. For example, the same global signal cannot be used as an SCLR in the static region and an ACLR in the PR region.
Creating Floorplan Assignments for PR designs

You must create a LogicLock region so the interface of the PR region with the static region is the same for any persona you implement. If different personas of a PR region have different area requirements, you must make a LogicLock region assignment that contains enough resources to fit the largest persona for the region. The static regions in your project do not necessarily require a floorplan, but depending on any other design requirement, you may choose to create a floorplan for a specific static region.

There is no minimum or maximum size for the LogicLock region assigned for a PR region. Because wire-LUTs are added on the periphery of a PR region by the Quartus II software, the LogicLock region for a PR region must be slightly larger than an equivalent non-PR region. Make sure the PR regions include only the resources that can be partially reconfigured; limit the LogicLock regions for PR regions to contain only LABs, DSPs, and RAM blocks. When multiple PR regions are present in a design, the shape and alignment of the region determines the PR mode (SCRUB or AND/OR) you can use.

You can use the default Auto size and Floating location LogicLock region properties to estimate the preliminary size and location for the PR region.

You can also define regions in the floorplan that match the general location and size of the logic in each partition. You may choose to create a LogicLock region assignment that is non-rectangular, depending on the design requirements, but disjoint LogicLock regions are not allowed for PR regions.

After compilation, use the Fitter-determined size and origin location as a starting point for your design floorplan. Check the quality of results obtained for your floorplan location assignments and make changes to the regions as needed.

Alternatively, you can perform Analysis and Synthesis, and then set the regions to the required size based on resource estimates. In this case, use your knowledge of the connections between partitions to place the regions in the floorplan.

For more information on making design partitions and using an incremental design flow, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in the Quartus II Handbook. For more design guidelines to ensure good quality of results, and suggestions on making design floorplan assignments with LogicLock regions, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in the Quartus II Handbook.

Implementation Details for Partial Reconfiguration

This section describes implementation details that help you create your PR design.

Partial Reconfiguration Pins

Partial reconfiguration can be performed through external pins or from inside the core of the FPGA. When using PR from pins, some of the I/O pins are dedicated for implementing partial reconfiguration functionality. If you perform partial reconfiguration from pins, then you must use the passive parallel with 16 data bits (FPPx16) configuration mode.

To enable partial reconfiguration from pins in the Quartus II software, perform the following steps:
1. From the Assignments menu, click Device, then click Device and Pin Options.

2. In the Device and Pin Options dialog box, select General in the Category list and turn on Enable PR pins from the Options list.

3. Click Configuration in the Category list and select Passive Parallel x16 from the Configuration scheme list.

4. Click OK, or continue to modify other settings in the Device and Pin Options dialog box.

5. Click OK.

Table 4–3 lists the dedicated pins available for use with partial reconfiguration:

<table>
<thead>
<tr>
<th>Pin Name</th>
<th>Pin Type</th>
<th>Pin Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PR_REQUEST</td>
<td>Input</td>
<td>Dedicated input when Enable PR pins is turned on; otherwise, available as user I/O. Logic high on pin indicates the PR host is requesting partial reconfiguration.</td>
</tr>
<tr>
<td>PR_READY</td>
<td>Output</td>
<td>Dedicated output when Enable PR pins is turned on; otherwise, available as user I/O. Logic high on this pin indicates the Stratix V control block is ready to begin partial reconfiguration.</td>
</tr>
<tr>
<td>PR_DONE</td>
<td>Output</td>
<td>Dedicated output when Enable PR pins is turned on; otherwise, available as user I/O. Logic high on this pin indicates that partial reconfiguration is complete.</td>
</tr>
<tr>
<td>PR_ERROR</td>
<td>Output</td>
<td>Dedicated output when Enable PR pins is turned on; otherwise, available as user I/O. Logic high on this pin indicates the device has encountered an error during partial reconfiguration.</td>
</tr>
<tr>
<td>DATA[15:0]</td>
<td>Input</td>
<td>Dedicated input when Enable PR pins is turned on; otherwise, available as user I/O. These pins provide connectivity for PR_DATA when Enable PR pins is turned on.</td>
</tr>
<tr>
<td>DCLK</td>
<td>Bidirectional</td>
<td>Dedicated input when Enable PR pins is turned on; otherwise, available as user I/O. PR_DATA is sent synchronous to this clock. This is a dedicated programming pin, and is not available as user I/O even if Enable PR pins is turned off.</td>
</tr>
</tbody>
</table>


You can enable open drain on PR pins from the Device and Pins Options dialog box in the Settings page of the Quartus II software.
Interfacing with the PR Control Block through a PR Host

The control signals described in Table 4–3 are used for communication between your PR control IP and the PR Control Block (CB) while executing partial reconfiguration. You can communicate with the PR control block through handshake signals, which can be accessed through external pins (known as using an external host), or within the FPGA core through internal control signals (known as using an internal host). The internal PR host can be user logic or a Nios® II processor.

Figure 4–8 shows how these blocks should be connected to the PR control block (CB). In your system, you will have either the External Host or the Internal Host, but not both.

The PR mode is independent of the full chip programming mode. For example, you can configure the full chip using a JTAG download cable, or other supported configuration modes. When configuring PR regions, you must use the FPPx16 interface to the PR control block whether you choose to partially reconfigure the chip from an external or internal host.

When using an external host, you must implement the control logic for managing system aspects of partial reconfiguration on an external device. By using an internal host, you can implement all of your logic necessary for partial reconfiguration in the FPGA, therefore external devices are not required to support partial reconfiguration. When using an internal host, you can use any interface to load the PR bitstream data to the FPGA, for example, from a serial or a parallel flash device, and then format the PR bitstream data to fit the FPPx16 interface on the PR Control Block.

To use the external host for your design, turn on the Enable PR Pins option in the Device and Pin Options dialog box in the Quartus II software when you compile your design. If this setting is turned off, then you must use an internal host. Also, you must tie the corectl port on the PR control block instance in the top-level of the design to the appropriate level for the selected mode.

PR Control Signals Interface

The Quartus II Programmer allows you to generate the different bit-streams necessary for full chip configuration and for partial reconfiguration. The programming bit-stream for partial reconfiguration contains the instructions (opcodes) as well as the configuration bits, necessary for reconfiguring each of the partial regions.
Figure 4–9 shows the handshaking control signals used for partial reconfiguration. When using an external host, the interface ports on the control block are mapped to FPGA pins. When using an internal host, these signals are within the core of the FPGA.

**Figure 4–9. Partial Reconfiguration Interface Signals**

- **PR_DATA**: The configuration bitstream is sent on \( \text{PR DATA}[15:0] \), synchronous to the \( \text{Clk} \).
- **PR_DONE**: Sent from CB to control logic indicating the PR process is complete.
- **PR_READY**: Sent from CB to control logic indicating the CB is ready to accept PR data from the control logic.
- **CRC_Error**: The CRC_Error generated from the device’s CRC block, is used to determine whether to partially reconfigure a region again, when encountering a CRC_Error.
- **PR_ERROR**: Sent from CB to control logic indicating an error during partial reconfiguration.
- **PR_REQUEST**: Sent from your control logic to CB indicating readiness to begin the PR process.
- **corectl**: Determines whether partial reconfiguration is performed internally or through pins.
Reconfiguring a PR Region

Figure 4–10 shows a system in which your PR Control logic is implemented inside the FPGA. However, this section is also applicable for partial reconfiguration with an external host. The PR control block (CB) represents the Stratix V PR controller inside the FPGA. PR1 and PR2 are two PR regions in a user design. In addition to the four control signals (PR_REQUEST, PR_READY, PR_DONE, PR_ERROR) and the data/clock signals interfacing with the PR control block, your PR Control IP should also send a control signal (PR_CONTROL) to each PR region. This signal implements the freezing and unfreezing of the PR Interface signals. This is necessary to avoid contention on the FPGA routing fabric.

![Figure 4–10. Example of a PR System with Two PR Regions](image)

After the FPGA device has been configured with a full chip configuration at least once, the INIT_DONE signal is released, and the signal is asserted high due to the external resistor on this pin. The INIT_DONE signal must be assigned to a pin to monitor it externally. When a full chip configuration is complete, and the device is in user mode, the following steps describe the PR sequence:

1. Begin a partial reconfiguration process from your PR Control logic, which initiates the PR process for one or more of the PR regions (asserting PR1_Control or PR2_Control in Figure 4–10). The wrapper HDL described earlier freezes (pulls high) all non-global inputs of the PR region before the PR process.
2. Send `PR_REQUEST` signal from your control logic to the PR Control Block (CB). If your design uses an external controller, monitor `INIT_DONE` to verify that the chip is in user mode before asserting the `PR_REQUEST` signal. The CB initializes itself to accept the PR data and clock stream. After that, the CB asserts a `PR_READY` signal to indicate it can accept PR data. Exactly four clock cycles must occur before sending the PR data to make sure the PR process progresses correctly. Data and clock signals are sent to the PR control block to partially reconfigure the PR region interface.

- If there are multiple PR personas for the PR region, your PR Control IP must determine the programming file data for partial reconfiguration.
- When there are multiple PR regions in the design, then the same PR control IP determines which regions require reconfiguration based on system requirements.
- At the end of the PR process, the PR control block asserts a `PR_DONE` signal and de-asserts the `PR_READY` signal.
- If you want to suspend sending data, you can implement logic to pause the clock at any point.

3. Your PR control logic must de-assert the `PR_REQUEST` signal within eight clock cycles after the `PR_DONE` signal goes high. If your logic does not de-assert the `PR_REQUEST` signal within eight clock cycles, a new PR cycle starts.

4. If your design includes additional PR regions, repeat steps 2 – 3 for each region. Otherwise, proceed to step 5.

5. Your PR Control logic de-asserts the `PR_CONTROL` signal(s) to the PR region. The freeze wrapper releases all input signals of the PR region, thus the PR region is ready for normal user operation.

6. You must perform a reset cycle to the PR region to bring all logic in the region to a known state. After partial reconfiguration is complete for a PR region, the states in which the logic in the region come up is unknown.

The PR event is now complete, and you can resume operation of the FPGA with the newly configured PR region.

At any time after the start of a partial reconfiguration cycle, the PR host can suspend sending the `PR_DATA`, but the host must suspend sending the `PR_CLK` at the same time. If the `PR_CLK` is suspended after a PR process, there must be at least 20 clock cycles after the `PR_DONE` or `PR_ERROR` signal is asserted to prevent incorrect behavior. Table 4–4 on page 4–27 contains other clock requirements for partial reconfiguration.

For an overview of different reset schemes in Altera devices, please refer to the Recommended Design Practices chapter in the Quartus II Handbook.

### Partial Reconfiguration Cycle Waveform

The PR host initiates the PR request, transfers the data to the FPGA device when it is ready, and monitors the PR process for any errors or until it is done.
A PR cycle is initiated by the host (internal or external) by asserting the PR_REQUEST signal high. When the FPGA device is ready to begin partial reconfiguration, it responds by asserting the PR_READY signal high. The PR host responds by sending configuration data on DATA [15:0]. The data is sent synchronous to PR_CLK. When the FPGA device receives all PR data successfully, it asserts the PR_DONE high, and de-asserts PR_READY to indicate the completion of the PR cycle.

If there is an error encountered during partial reconfiguration, the FPGA device asserts the PR_ERROR signal high and de-asserts the PR_READY signal low.

The PR host must continuously monitor the PR_DONE and PR_ERROR signals status. Whenever either of these two signals are asserted, the host must de-assert PR_REQUEST within eight PR_CLK cycles. As a response to PR_ERROR error, the host can optionally request another partial reconfiguration or perform a full FPGA configuration.

Note that the PR_CLK signal has a nominal maximum frequency of 62.5 MHz. To prevent incorrect behavior, the PR_CLK signal must be active a minimum of twenty clock cycles after PR_DONE or PR_ERROR signal is asserted high. Once PR_DONE is asserted, PR_REQUEST must be de-asserted within eight clock cycles. PR_DONE is de-asserted by the device within twenty PR_CLK cycles. The host can assert PR_REQUEST again after the 20 clocks after PR_DONE is de-asserted.

Table 4–4 lists interface signal timing requirements for partial reconfiguration.

<table>
<thead>
<tr>
<th>Timing Parameters</th>
<th>Value (clock cycles)</th>
</tr>
</thead>
<tbody>
<tr>
<td>PRREADY to first data</td>
<td>4 (exact)</td>
</tr>
<tr>
<td>PERROR to last clock</td>
<td>20 (minimum)</td>
</tr>
<tr>
<td>PDONE to last clock</td>
<td>20 (minimum)</td>
</tr>
<tr>
<td>DONETO_REQ_low</td>
<td>8 (maximum)</td>
</tr>
</tbody>
</table>

At any time during partial reconfiguration, to pause sending PR_DATA, the PR host can stop toggling PR_CLK. The clock can be stopped either high or low.

At any time during partial reconfiguration, the PR host can terminate the process by de-asserting the PR request. A partially completed PR process results in a PR error. You can have the PR host restart the PR process after a failed process by sending out a new PR request 20 cycles later.

In case you terminate a PR process before completion, and follow it up with a FPGA reset using the nConfig signal, you must keep the PR_CLK signal running through the FPGA reset cycle to avoid causing the partial reconfiguration to lock up.
Figure 4–11 shows the partial reconfiguration cycle waveform for the hand-shaking protocol.

**Figure 4–11. Partial Reconfiguration Timing Diagram**

During these steps, the PR control block might assert a **PR_ERROR** or a **CRC_ERROR** signal to indicate that there was an error during the partial reconfiguration process. Assertion of **PR_ERROR** indicates that the PR bitstream data was corrupt, and the assertion of CRC error indicates a CRAM CRC error either during or after completion of PR process. If the **PR_ERROR** or **CRC_ERROR** signals are asserted, you must plan whether to reconfigure the PR region or reconfigure the whole FPGA, or leave it unconfigured.

**Partial Reconfiguration with an External Host**

For partial reconfiguration using an external host, you must set the **MSEL [4:0]** pins for FPPx16 configuration scheme. You can use a microcontroller, another FPGA, or a CPLD such as a MAX V device, to implement the configuration and PR controller. In this setup, the Stratix V device configures in FPPx16 mode during power-up. Alternatively, you can use a JTAG interface to configure the Stratix V device.

At any time during user-mode, the external host can initiate partial reconfiguration and monitor the status using the external PR dedicated pins: **PR_REQUEST**, **PR_READY**, **PR_DONE**, and **PR_ERROR**. In this mode, the external host must respond appropriately to the hand-shaking signals for a successful partial reconfiguration. This includes acquiring the data from the flash memory and loading it into the Stratix V device on **DATA [15:0]**. The waveform for handshaking signals between the external host and the FPGA device is also shown in Figure 4–11.

---

Note to Figure 4–11

(1) For external mode the **PR_CLK** is the same as the configuration clock **DCLK**. The maximum specification for **DCLK** = 80MHz when using partial reconfiguration.
Figure 4–12 shows the connection setup for partial reconfiguration with an external host in the FPPx16 configuration scheme.

**Figure 4–12. Connecting to an External Host**
Using an External Host with Multiple Devices

Figure 4–13 shows an example of an external host controlling multiple Stratix V devices on a board. You must design the external host to accommodate the arbitration scheme that is required for your system, as well as the partial reconfiguration interface requirement for each device.

Partial Reconfiguration with an Internal Host

The PR internal host is a piece of soft logic implemented in the FPGA that you must design to accommodate the hand-shaking protocol with the PR control block. For example, PR programming bitstream(s) stored in an external flash device can be routed through the regular I/Os of the FPGA device, or received through the high speed transceiver channel (PCI Express, SRIO or Gigabit Ethernet), and can be stored in on-chip memory such as MLABs or M20K blocks, for processing by the internal logic. This data must be formatted into the 16 bit wide data so that it can be transmitted to the PR control block by the internal IP, because the PR control block can only accept PR data via its FPPx16 interface.
The PR dedicated pins (PR_REQUEST, PR_READY, PR_DONE, and PR_ERROR) can be used as regular I/Os when performing partial reconfiguration with an internal host. For the full FPGA configuration upon power-up, you can set the MSEL[4:0] pins to match the configuration scheme, for example, AS, PS, FPPx8, FPPx16, or FPPx32. Alternatively, you can use the JTAG interface to configure the FPGA device. At any time during user-mode, you can initiate partial reconfiguration through the FPGA core fabric using the PR internal host.

Figure 4–14 shows an example of the configuration setup when performing partial reconfiguration using the internal host. In this example, the programming bitstream for partial reconfiguration is received through the PCI Express link, and your logic converts the data to the FPPx16 mode.

**Figure 4–14. Connecting to an Internal Host**

---

**Managing a Partial Reconfiguration Project**

When compiling your PR project, you must create a base revision, and one or more reconfigurable revisions. The project revision you start out is termed the base revision.

**Creating Reconfigurable Revisions**

To create a reconfigurable revision, use the **Revisions** tab of the **Project Navigator** window in the Quartus II software. When you create a reconfigurable revision, the Quartus II software adds the required assignments to associate the reconfigurable revision with the base revision of the PR project. You can add the necessary files to each revision with the **Add/Remove Files** option in the **Project** option under the **Project** menu in the Quartus II software. With this step, you can associate the right implementation files for each revision of the PR project.
Compiling Reconfigurable Revisions

Altera recommends that you use the largest persona of the PR region for the base compilation so that the Quartus II software can automatically budget sufficient routing.

Here are the typical steps involved in a PR design flow.
1. Compile the base revision with the largest persona for each PR region.
2. Create reconfigurable revisions for other personas of the PR regions by right-clicking in the Revisions tab in the Project Navigator.
3. Compile your reconfigurable revisions.
4. Analyze timing on each reconfigurable revision to make sure the design performs correctly to specifications.
5. Create aggregate revisions as needed.
6. Create programming files.

For more information on compiling a partial reconfiguration project, refer to Performing Partial Reconfiguration in Quartus II Help.

Timing Closure for a Partial Reconfiguration Project

As with any other FPGA design project, simulate the functionality of various PR personas to make sure they perform to your system specifications. You must also make sure there are no timing violations in the implementation of any of the personas for every PR region in your design project. In the Quartus II software, this process is manual, and you must run multiple timing analyses, on the base, reconfigurable, and aggregate revisions. The different timing requirements for each PR persona can be met by using different SDC constraints for each of the personas.

The interface between the partial and static partitions remains identical for each reconfigurable and aggregate revision in the PR flow. If all the interface signals between the static and the PR regions are registered, and there are no timing violations within the static region as well as within the PR regions, the reconfigurable and aggregate revisions should not have any timing violations.

However, you should perform timing analysis on the reconfigurable and aggregate revisions, in case you have any unregistered signals on the interface between partial reconfiguration and static regions.

Bitstream Compression and Encryption for PR Designs

You can choose to independently compress and encrypt the base bitstream as well as the PR bitstream for your PR project using options available in the Quartus II software.

When you choose to compress the bitstreams, you can compress the base and PR programming bitstreams independently, based on your design requirements. However, if you want to encrypt only the base image, you can choose either to encrypt or not encrypt the PR images.
Creating Programming Files for a Partial Reconfiguration Project

You must generate PR bitstream(s) based on the designs and send them to the control block for partial reconfiguration. Compile the PR project, including the base revision and at least one reconfigurable revision before generating the PR bitstreams. The Quartus II Programmer generates PR bitstreams. This generated bitstream can be sent to the PR ports on the control block for partial reconfiguration.

Consider a partial reconfiguration design that has three revisions and one PR region, as shown in Figure 4–15; a base revision with persona a, one PR revision with persona b, and a second PR revision with persona c.

When these individual revisions are compiled in the Quartus II software, the assembler produces Masked SRAM Object Files (.msf) and the SRAM Object Files (.sof) for each revision, as shown in Figure 4–15. The .sof files are created as before (for non-PR designs). Additionally, .msf files are created specifically for partial reconfiguration, one for each revision. The pr_region.msf file is the one of interest for generating the PR bitstream. It contains the mask bits for the PR region. Similarly, the static.msf file has the mask bits for the static region. The .sof files have the information on how to configure the static region as well as the corresponding PR
The `pr_region.msf` file is used to mask out the static region so that the bitstream can be computed for the PR region. The default file name of the `pr_region.msf` corresponds to the location of the PR region unless the associated LogicLock region has a non-default name. In that case, the `.msf` file is named after the region.

Altera recommends naming all LogicLock regions to enhance documenting your design.

You can convert files in the Convert Programming Files window or run the `quartus_cpf -p` command to process the `pr_region.msf` files in conjunction with the `.sof` files and generate the Partial-Masked SRAM Object File (.pmsf) files as shown in Figure 4–16 by merging the `.msf` and the `.sof` files. The `.msf` file helps determine the PR region from each of the `.sof` files during the PR bitstream computation.

Once all the `.pmsf` files are created, process the PR bitstreams by running the `quartus_cpf -o` command to produce the raw binary `.rbf` files for reconfiguration. Figure 4–17 shows how three bitstreams can be created to partially reconfigure the region with persona a, persona b, or persona c as desired.

If one wishes to partially reconfigure the PR region with persona a, use the `a.rbf` bitstream file, and so on for the other personas.

In the Quartus II software, the Convert Programming Files window supports the generation of the required programming bitstreams. This is described “Generating PR Programming Files Using the CPF GUI” on page 4–35. When using the `quartus_cpf` from the command line, the following options for generating the programming files are read from an option text file, for example, `option.txt`.

- If you want to use SCRUB mode, before generating the bitstreams create an option text file, with the following line:
  
  `use_scrub=on`
■ If you have initialized M20K blocks in the PR region (ROM/Initialized RAM), then add the following line in the option text file, before generating the bitstreams:

```plaintext
write_block_memory_contents=on
```

■ If you want to compress the programming bitstream files, add the following line in the option text file. This option is available when converting base .sof to any supported programming file types, such as .rbf, .pof and JTAG Indirect Configuration File (.jic).

```plaintext
bitstream_compression=on
```

### Generating Required Programming Files

1. Generate .sof and .msf files (part of a full compilation of the base and PR revisions).

2. Generate a Partial-Masked SRAM Object File (.pmsf) using the following commands:

   ```plaintext
   quartus_cpf -p <pr_revision>.msf <pr_revision>.sof <new_filename>.pmsf
   ```

   For example:

   ```plaintext
   quartus_cpf -p x7y48.msf switchPRBS.sof x7y48_new.pmsf
   ```

3. Convert the .pmsf file for every PR region in your design to .rbf file format. The .rbf format is used to store the bitstream in an external flash memory. This command should be run in the same directory where the files are located:

   ```plaintext
   quartus_cpf -o scrub.txt -c <pr_revision>.pmsf <pr_revision>.rbf
   ```

   For example:

   ```plaintext
   quartus_cpf -o scrub.txt -c x7y48_new.pmsf x7y48.rbf
   ```

   When you do not have an option text file such as `scrub.txt`, the files generated would be for AND/OR mode of PR, rather than SCRUB mode.

### Generating PR Programming Files Using the CPF GUI

In the Quartus II software, the PR programming file generation flow is supported in the Convert Programming Files window, creating a .pmsf output file, from .msf and .sof input files as well as creating a .rbf output file, with a .pmsf input file.

Convert Programming Files also provides the option to enable the option bit for bitstream decompression during partial reconfiguration, when converting the base .sof (full design .sof) to any supported file type.

For additional details, please refer to the Quartus II Programmer chapter in the Quartus II Handbook.

### Generating a .pmsf File from a .msf and .sof Input File

Perform the following steps in the Quartus II software to generate the .pmsf file in the Convert Programming Files window.

1. Open the Convert Programming Files window.

2. Specify the programming file type as Partial-Masked SRAM Object File (.pmsf).

3. Specify the output file name.
4. Select input files to convert (only a single .msf and .sof file are allowed). Click Add.

5. Click Generate to generate the .pmsf file.

Generating a .rbf File from a .pmsf Input File

Perform the following steps in the Quartus II software to generate the partial reconfiguration .rbf file in the Convert Programming Files window.

1. From the File menu, click Convert Programming Files.

2. Specify the programming file type as Raw Binary File for Partial Reconfiguration (.rbf).

3. Specify the output file name.

4. Select input file to convert. Only a single .pmsf input file is allowed. Click Add.

5. Select the new .pmsf file and click Properties.

6. Turn the Compression, Enable SCRUB mode, Write memory contents, and Generate encrypted bitstream options on or off depending on the requirements of your design. Click Generate to generate the .rbf file for partial reconfiguration.

- **Compression**: Enables compression on the PR bitstream.

- **Enable SCRUB mode**: Default is based on AND/OR mode. This option is valid only when your design does not contain vertically overlapped PR masks. The .rbf file generation fails otherwise.

- **Write memory contents**: Turn this on when you have a .mif file that was used during compilation. Otherwise, turning this option on forces you to use double PR.

- **Generate encrypted bitstream**: If this option is enabled, you must specify the Encrypted Key Programming (.ekp) file, which generated when converting a base .sof to an encrypted bitstream. The same .ekp must be used to encrypt the PR bitstream.

When you turn on Compression, you must present each PR_DATA[15:0] word for exactly four clock cycles. Other timing requirements shown in Figure 4–11 on page 4–28 are listed in Table 4–5 for bitstream compression.

<table>
<thead>
<tr>
<th>Timing Parameters</th>
<th>Value (clock cycles)</th>
</tr>
</thead>
<tbody>
<tr>
<td>PR_READY to first data</td>
<td>4 (exact)</td>
</tr>
<tr>
<td>PR_ERROR to last clock</td>
<td>80 (minimum)</td>
</tr>
<tr>
<td>PR_DONE to last clock</td>
<td>80 (minimum)</td>
</tr>
<tr>
<td>DONE_to_REQ_low</td>
<td>8 (maximum)</td>
</tr>
</tbody>
</table>

Turn on the Write memory contents option only if you have M20K blocks in your PR design that need to be initialized. When you check this box, you must to perform double PR for regions with initialized M20K blocks. For more information, refer to “Using Double PR Cycle for Initializing M20K blocks” on page 4–45.
Option to Enable Bitstream Decompression during Partial Reconfiguration

In the Quartus II software, the Convert Programming Files window provides the option in the .sof file properties to enable bitstream decompression during partial reconfiguration. This option is available when converting base .sof to any supported programming file types, such as .rbf, .pof, and .jic.

In order to view this option, the base .sof must be targeted on Stratix V devices in the .sof File Properties. This option must be turned on if you turned on the Compression option during .pmsf to .rbf file generation.

Option to Enable Bitstream Decryption During Partial Reconfiguration

The Convert Programming Files window provides the option in the .sof file properties to enable bitstream decryption during partial reconfiguration. This option is available when converting base .sof to any supported programming file types, such as .rbf, .pof, and .jic.

The base .sof must have partial reconfiguration enabled and the base .sof generated from a design that has a PR Control Block instantiated, to view this option in the .sof File Properties. This option must be turned on if you wants to turn on the Generate encrypted bitstream option during .pmsf to .rbf file generation.

Using On-Chip Debug for PR designs

You cannot instantiate a SignalTap II block inside a PR region. If you must monitor signals within a PR region for debug purposes, bring those signals to the ports of the PR region. Figure 4–18 shows how you can then instantiate the SignalTap II block in the static region of the design and probe the signals you want to monitor.

Figure 4–18. Using SignalTap II with a PR Design

The Quartus II software does not support the Incremental SignalTap feature for PR designs. After you instantiate the SignalTap II block inside the static region, you must recompile your design. When you recompile your design, the static region may have a modified implementation and you must also recompile your PR revisions. If you modify an existing SignalTap II instance you must also recompile your entire design; base revision and reconfigurable revisions.

You can use other on-chip debug features in the Quartus II software, such as the In-System Sources and Probes or SignalProbe, to debug a PR design. As in the case of SignalTap, In-System Sources and Probes can only be instantiated within the static region of a PR design. If you have to probe any signal inside the PR region, you must bring those signals to the ports of the PR region in order to monitor them within the static region of the design.
Partial Reconfiguration Known Limitations

The restrictions in the following sections apply only if your design uses M20K blocks as RAMs or ROMs in your PR project. The restrictions derive from hardware limitations in specific Stratix V devices.

Memory Blocks Initialization Requirement for PR Designs

For a non-PR design, the power up value for the contents of a M20K RAM or a MLAB RAM are all set at zero. However, at the end of performing a partial reconfiguration, the contents of a M20K or MLAB memory block are unknown, and user must intentionally initialize the contents of all the memory to zero, if required by the functionality of the design, and not rely upon the power on values.

Usage of M20K RAM Blocks in PR Designs

When your PR design uses M20K RAM blocks in Stratix V devices, there are some restrictions which limit how you utilize the respective memory blocks as ROMs or as RAMs with initial content. Additionally, if your design must have initialized memory content either as a ROM or a RAM inside a PR region, you must follow the partial reconfiguration sequence described in “Implementing Memories with Initialized Content in PR Designs” on page 4–43 of this chapter to properly design the PR host logic.

Limitations with Stratix V ES Devices

For more information including a list of the Stratix V ES devices, refer to the Stratix V ES Errata Sheet and Guidelines.

Partial reconfiguration related errata items are discussed in this section.

If you implement a M20K block in your PR region as a ROM or a RAM with initialized content, then during the time the PR region is being reconfigured, the data read from the memory blocks in static regions in columns that cross the PR region are incorrect.
If the functionality of the static region depends on any data read out from M20K RAMs in the static region, the design will malfunction. Figure 4–19 shows this limitation.

Figure 4–19. Limitations for Using M20Ks in PR Regions

Use one of the following workarounds, which are applicable to both AND/OR and SCRUB modes of partial reconfiguration:

- Do not use ROMs or RAMs with initialized content inside PR regions.
- If this is not possible for your design, you can program the memory content for M20K blocks with a .mif using the suggested workarounds.
- Make sure your PR region extends vertically all the way through the device, in such a way that the M20K column lies entirely inside a PR region.
Figure 4–20 shows the LogicLock region extended as a rectangle reducing the area available for the static region. However, you can create non-rectangular LogicLock regions for optimally allocating the required resources for the partition. If area saving is a concern, extend the LogicLock region to include M20K columns entirely.

**Figure 4–20. Workaround for Using M20Ks in PR Regions**

Block all the M20K columns not inside a PR region, but are in columns above or below a PR region, using Reserved LogicLock Regions as shown in Figure 4–21. In this case, you may choose to under-utilize M20K resources, for gaining ROM functionality within the PR region.

**Figure 4–21. Alternative Workaround for Using M20Ks in PR Region**

**Limitations When Using Stratix V Production Devices**

Some Stratix V production devices have the same restriction and workaround described above for ES devices.
For more information, including a list of the Stratix V devices that have this limitation, refer to the *Stratix V Errata Sheet and Guidelines*.

**Using MLAB blocks in PR designs**

Stratix V devices include dual-purpose blocks called MLABs, which can be used to implement RAMs or LABs for user logic. This section describes the restrictions while using MLAB blocks (sometimes also referred to as LUT-RAM) in Stratix V devices for your PR designs.

If your design uses MLABS as LUT RAM, you must use all available MLAB bits within the region.

If your design does not use any MLAB blocks as RAMs, the following discussion does not apply. The restrictions listed below are the result of hardware limitations in specific devices.

**Limitations with Stratix V ES devices**

- LUT-RAMs without initialization in MLABs are supported in both AND/OR and SCRUB modes, but with the following restrictions.
  - MLAB blocks contain 640 bits of memory. The LUT RAMs in PR regions in your design must occupy all MLAB bits, you should not use partial MLABs.
  - You must include control logic in your design with which you can write to all MLAB locations used inside the PR region.
  - Using this control logic, write ‘1’ at each MLAB RAM bit location in the PR region before starting the PR process. This is to work around a false EDCRC error during partial reconfiguration.
  - When using the AND-OR mode, you must also specify a `.mif` that sets all MLAB RAM bits to ‘1’ immediately after PR is complete. When using SCRUB mode, you do not have to use a `.mif`.

- LUT-RAMs with initialized content in MLABs are supported only when using the SCRUB mode of partial reconfiguration, but with the following caveat.
  - They have the same restrictions as LUT-RAMs without initialization, as described in the previous bullet.
  - There are no restrictions to using MLABs in the static region of your PR design.

**Limitations with Stratix V Production Devices**

When using SCRUB mode:

- LUT-RAMs without initialized content, LUT-RAMs with initialized content, and LUT-ROMs can be implemented in MLABs within PR regions without any restriction.

When using AND/OR mode:

- LUT-RAMs with initialized content or LUT-ROMs cannot be implemented in a PR region.
LUT-RAMs without initialized content in MLABs inside PR regions are supported with the following restrictions.

- MLAB blocks contain 640 bits of memory. The LUT RAMs in PR regions in your design must occupy all MLAB bits, you should not use partial MLABs.
- You must include control logic in your design with which you can write to all MLAB locations used inside PR region.
  - Using this control logic, write '1' at each MLAB RAM bit location in the PR region before starting the PR process. This is to work around a false EDCRC error during partial reconfiguration.
- You must also specify a .mif that sets all MLAB RAM bits to 1 immediately after PR is complete.
- ROMs cannot be implemented in MLABs (LUT-ROMs), with either AND/OR or SCRUB modes of partial reconfiguration.
- There are no restrictions to using MLABs in the static region of your PR design.

Table 4–6 shows a summary of the LUT-RAM Restrictions.

Table 4–6. RAM Implementation Restrictions Summary

<table>
<thead>
<tr>
<th>PR Mode</th>
<th>Type of memory in PR region</th>
<th>Stratix V ES (1)</th>
<th>Stratix V Production (2)</th>
</tr>
</thead>
<tbody>
<tr>
<td>SCRUB mode</td>
<td>LUT RAM (no initial content)</td>
<td>Write 1’s to all locations before partial reconfiguration</td>
<td>OK</td>
</tr>
<tr>
<td></td>
<td>LUT ROM and LUT RAM with your initial content</td>
<td>No</td>
<td>OK</td>
</tr>
<tr>
<td>AND/OR mode</td>
<td>LUT RAM (no initial content)</td>
<td>While design is running: Write 1s to all locations before partial reconfiguration At compile time: Explicitly initialize all memory locations in each new persona to 1 via initialization file (.mif).</td>
<td>While design is running: Write 1s to all locations before partial reconfiguration At compile time: Explicitly initialize all memory locations in each new persona to 1 via initialization file (.mif).</td>
</tr>
<tr>
<td></td>
<td>LUT ROM and LUT RAM with your initial content</td>
<td>No</td>
<td>No</td>
</tr>
</tbody>
</table>

Note to table:
(1) Stratix V ES devices require the listed workarounds. For more information refer to the Stratix V ES Errata Sheet and Guidelines in the Stratix V handbook.
(2) Stratix V production devices require the listed workarounds. For more information refer to the Stratix V Errata Sheet and Guidelines in the Stratix V handbook.
Implementing Memories with Initialized Content in PR Designs

If your Stratix V PR design implements ROMs, RAMs with initialization, or ROMs within the PR regions, using either M20K blocks or LUT-RAMs, then you must follow the design guidelines in Table 4–7 for what is applicable in your case.

Table 4–7. Implementing Memory with Initialized Content in PR Designs (Part 1 of 2)

<table>
<thead>
<tr>
<th>Mode</th>
<th>Production Devices</th>
<th>ES Devices</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>AND/OR</td>
<td>SCRUB</td>
</tr>
<tr>
<td><strong>LUT-RAM without</strong></td>
<td></td>
<td></td>
</tr>
<tr>
<td><strong>initialization</strong></td>
<td><strong>Suggested</strong></td>
<td><strong>No special</strong></td>
</tr>
<tr>
<td><strong>Method</strong></td>
<td><strong>Method</strong></td>
<td><strong>method</strong></td>
</tr>
<tr>
<td></td>
<td><strong>While design is</strong></td>
<td><strong>Make sure</strong></td>
</tr>
<tr>
<td></td>
<td><strong>running:</strong> Write</td>
<td>no spurious</td>
</tr>
<tr>
<td></td>
<td>‘1’ to all</td>
<td>write on PR</td>
</tr>
<tr>
<td></td>
<td>locations before</td>
<td>entry (1)</td>
</tr>
<tr>
<td></td>
<td>partial reconfiguration.</td>
<td></td>
</tr>
<tr>
<td></td>
<td><strong>At compile</strong></td>
<td></td>
</tr>
<tr>
<td></td>
<td><strong>time:</strong> Explicitly</td>
<td></td>
</tr>
<tr>
<td></td>
<td>initialize all</td>
<td></td>
</tr>
<tr>
<td></td>
<td>memory locations in</td>
<td></td>
</tr>
<tr>
<td></td>
<td>each new persona to</td>
<td></td>
</tr>
<tr>
<td></td>
<td>‘1’ via initialization</td>
<td></td>
</tr>
<tr>
<td></td>
<td>file (.mif)</td>
<td>(1)</td>
</tr>
<tr>
<td></td>
<td>Make sure no</td>
<td></td>
</tr>
<tr>
<td></td>
<td>spurious write on</td>
<td></td>
</tr>
<tr>
<td></td>
<td>PR entry (1)</td>
<td></td>
</tr>
<tr>
<td></td>
<td><strong>Without</strong></td>
<td></td>
</tr>
<tr>
<td><strong>Suggested</strong></td>
<td><strong>Method</strong></td>
<td></td>
</tr>
<tr>
<td><strong>Method</strong></td>
<td><strong>CRC Error</strong></td>
<td><strong>N/A</strong></td>
</tr>
</tbody>
</table>

| **LUT-RAM with**           |                    |            |
| **initialization**         | **Not supported**  |            |            |
| **without**                | **Supported**      |            |            |
| **initialization**         | **Not supported**  |            |            |
|                            | **Incorrect**      |            |            |
|                            | **results**        |            |            |
|                            | **Make sure no**   |            |            |
|                            | **spurious write on** |            |
|                            | **PR exit**        |            |            |
|                            | **(1)**            |            |            |

| **M20K without**           | **Suggested**      |            |            |
| **initialization**         | **No special**     |            |            |
|                            | **Method**         |            |
|                            | **Without**        |**N/A**     |            |
| **Suggested**              | **Method**         |            |
| **Method**                 | **Without**        |            |            |
Partial Reconfiguration Known Limitations

To avoid spurious writes during PR entry and exit, implement the following clock enable circuit in the same PR region as the RAM. The circuit depends on an active-high clear signal from the static region. Before entering PR, freeze this signal in the same manner as all PR inputs. Your host control logic should de-assert the clear signal as the final step in the PR process.

**Table 4–7. Implementing Memory with Initialized Content in PR Designs (Part 2 of 2)**

<table>
<thead>
<tr>
<th>Mode</th>
<th>Production Devices</th>
<th>ES Devices</th>
</tr>
</thead>
<tbody>
<tr>
<td>M20K with initialization</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Suggested Method</td>
<td>Use double PR cycle (2) Make sure no spurious write on PR exit (1)</td>
<td>Use double PR cycle (2) Make sure no spurious write on PR exit (1)</td>
</tr>
<tr>
<td>Without Suggested Method</td>
<td>Incorrect results</td>
<td></td>
</tr>
</tbody>
</table>

**Note to table:**
(1) Use the circuit shown in Figure 4–22 to create clock enable logic to safely exit partial reconfiguration without spurious writes.
(2) Double partial reconfiguration is describe in “Using Double PR Cycle for Initializing M20K blocks” on page 4–45.

To avoid spurious writes during PR entry and exit, implement the following clock enable circuit in the same PR region as the RAM. The circuit depends on an active-high clear signal from the static region. Before entering PR, freeze this signal in the same manner as all PR inputs. Your host control logic should de-assert the clear signal as the final step in the PR process.

**Figure 4–22. M20K/LUTRAM**
Using Double PR Cycle for Initializing M20K blocks

When a PR region in your PR design contains an initialized M20K block and is reconfigured via AND/OR mode, your host logic must complete a double PR cycle, instead of a single PR cycle.

Figure 4–23. Next PR Request Assertion During Double PR Cycle

![Diagram of double PR cycle]

**Note to Figure 4–23**

(1) Error_to_LAST_CLOCK and DONE_to_LAST_CLOCK values are valid after the second PR cycle for double partial reconfiguration.

Figure 4–23 displays the second phase of a double PR cycle, where the host logic must issue another PR_REQUEST signal after exactly seven clock cycles after the PR_DONE signal is asserted. If the PR compression feature is enabled, the host logic must issue another PR_REQUEST signal exactly two clock cycles after PR_DONE is asserted. The FPGA responds with PR_READY signal to the second PR_REQUEST signal assertion.

The PR host must continue sending PR_DATA signal exactly four clock cycles after the PR_READY signal, just as in the first PR cycle. The data on PR_DATA pins can be don't care between the first PR_DONE signal and until four clock cycles after the PR_READY signal is asserted for the second PR cycle.

The host must continue sending a PR_DATA signal for the second PR cycle, until it receives the PR_DONE signal for the second request, similar to the first PR cycle. After the PR_DONE signal is asserted for the second time, the host should de-assert the PR_REQUEST signal and continue with other operations needed for region bring up, such as issuing a reset to bring the region to a known state.

Using Double PR with Compressed Programming Bitstream

You can use bitstream compression along with PR designs that also require memory initialization for M20K blocks. For a compressed bitstream requiring a double PR cycle, the PR host must stop sending the PR_DATA signal in the bitstream as soon as the first PR_DONE is asserted. The PR host must resume sending the PR_DATA signal immediately after the second PR_READY signal is asserted.
Document Revision History

Table 4–8 lists the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>Added support for encrypted bitstreams.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated support for double PR.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>
This chapter describes the flow for designing HardCopy® series devices in the Quartus® II software.

Altera® HardCopy ASICs are the lowest risk, lowest total cost ASICs. The HardCopy system development methodology offers fast time-to-market, low risk, and with the Quartus II software, you can design with one set of RTL code and one set of IP for both FPGA and ASIC implementations. This flow enables you to conduct true hardware/software co-design and completely prepare your system for production prior to ASIC design hand-off. Altera provides a turn-key process to convert your design to a HardCopy ASIC for production.

In this chapter, the term FPGA refers to a Stratix® II, Stratix III, or Stratix IV device, which is the prototype device for a HardCopy II, HardCopy III, or HardCopy IV device, respectively.

This chapter discusses the following topics:

- “HardCopy Development Flow” on page 5–2
- “HardCopy Utilities” on page 5–6
- “Selecting the Prototype and Companion Devices” on page 5–7
- “Applying Design Constraints” on page 5–10
- “Compiling the Design and Creating Companion Revisions” on page 5–15
- “Timing Closure and Verification” on page 5–21
- “Performing ECOs with Quartus II Engineering Change Management with the Chip Planner” on page 5–25
- “Preparing the Design for Handoff” on page 5–29

For more information about HardCopy series devices, refer to the respective HardCopy device handbook, which is available on the Literature page of the Altera website at www.altera.com.
HardCopy Series Design Benefits

Designing with HardCopy devices offers the following substantial benefits:

■ Seamless prototyping using an FPGA for at-speed system verification and system development, which reduces total project development time and cost

■ Dependable conversion from an FPGA prototype to a HardCopy ASIC expands product planning options

■ Unified design methodology for FPGA and HardCopy designs reduces the need for ASIC development software and two sets of intellectual property, which reduces project risk

■ System development methodology delivers lowest total cost

HardCopy Development Flow

In the HardCopy development flow, you design a FPGA and a HardCopy companion device together in one Quartus II project using one of the following design flows:

■ FPGA first flow—Design the FPGA first for in-system verification, and then create a HardCopy companion device second. Performing system verification early helps reduce overall total project development time. The FPGA first flow is the default flow and the rest of this chapter is based on this flow.

■ HardCopy first flow—Design the HardCopy device first, and then create the FPGA companion device second for in-system verification. This method more accurately predicts the maximum performance of the HardCopy device during development. If you optimize your design to maximize HardCopy performance, but cannot meet your performance requirements with the FPGA, you can still map your design with decreased performance requirements for in-system verification.

These two flows are illustrated at a high level in Figure 5–1.
Designing the FPGA First

The FPGA first flow begins with seamless FPGA prototyping and is similar to the traditional FPGA design flow, but requires you to perform additional tasks necessary to convert the design to the HardCopy companion device within the same project. The following steps provide an overview of the tasks necessary in the FPGA first flow:

1. Specify an FPGA device for prototyping and a HardCopy companion device. Refer to “Selecting the Prototype and Companion Devices” on page 5–7 for more information.

2. Apply design constraints. Refer to “Applying Design Constraints” on page 5–10 for more information.

3. Compile the FPGA design, and then create and compile the HardCopy companion revision. Refer to “Compiling the Design and Creating Companion Revisions” on page 5–15 for more information.


5. Generate the handoff files, reports, and archive, and arrange for its submission to the Altera HardCopy Design Center for back-end implementation. Refer to “Preparing the Design for Handoff” on page 5–29 for more information.
Figure 5–2 provides an overview of the FPGA first flow.

**Figure 5–2. Designing the FPGA First Flow**

HardCopy Device Development with the FPGA Device First Flow

- Prepare FPGA Design
- Select HardCopy Companion Device
- Review HardCopy Advisor
- Apply Design Constraints
- Compile FPGA Design
- Any Violations?
  - Yes:
    - Fix Violations
  - No:
    - Create or Overwrite HardCopy Companion Revision
    - Compile HardCopy Companion Revision
    - Any Violations?
      - Yes:
        - No
      - No:
        - Select a Larger HardCopy Companion Device
        - Fits in HardCopy Device?
          - Yes:
            - Compare FPGA and HardCopy Revisions
            - Any Violations?
              - Yes
              - No
            - No
          - No:
            - Compare FPGA and HardCopy Revisions
            - Any Violations?
              - Yes
              - No
            - No

Design Submission & Back End Implementation Phase

- Generate Handoff Report
- Archive Project for Handoff
Designing the HardCopy Device First

The HardCopy device first flow in the Quartus II software enables you to maximize performance in the HardCopy device and map the design to the FPGA prototype for in-system verification. This approach is preferred if you use the HardCopy device to achieve higher performance than the FPGA prototype, because you can see your potential maximum performance in the HardCopy device immediately during development, and you can create a slower performing FPGA prototype of the design for in-system verification. This design process is similar to the FPGA first flow development flow, but you begin the design with a different initial device family instead. The remaining tasks to complete your design for both the FPGA and HardCopy devices essentially follow the same process.

The complete HardCopy first flow is shown in Figure 5–3.

Figure 5–3. Designing the HardCopy First Flow
HardCopy Advisor

The HardCopy Advisor provides an interactive list of tasks to help you through the development of your FPGA prototype and HardCopy design. The following tasks highlight the checkpoints that the HardCopy Advisor reviews, which includes the major checkpoints in the design process, but not every step in the process.

1. Select an FPGA device.
2. Select a HardCopy companion device.
3. Set up the FPGA revision.
4. Confirm FPGA junction temperature range settings (HardCopy III and HardCopy IV devices only).
5. Compile and check the FPGA design.
6. Create or overwrite the HardCopy companion revision.
7. Confirm HardCopy junction temperature range settings (HardCopy III and HardCopy IV devices only).
8. Compile and check the HardCopy companion results.
9. Compare companion revisions.
11. Archive handoff files and send them to Altera.

For more information about the HardCopy Advisor in the Quartus II software, refer to About the HardCopy Advisor in Quartus II Help.

HardCopy Utilities

The HardCopy Utilities menu contains the main functions you use to develop your HardCopy design and FPGA prototype companion revision. To access this menu in the Quartus II software, on the Project menu, click HardCopy Utilities. Each HardCopy Utilities menu feature is summarized in Table 5–1.

Table 5–1. HardCopy Utilities Menu Options (Part 1 of 2)

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
<th>Applicable Design Revision</th>
<th>Restrictions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Create/Overwrite HardCopy Companion Revision</td>
<td>Creates a new companion revision or overwrites an existing companion revision for your FPGA and HardCopy design</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>• The <strong>Auto device selected by the Fitter</strong> option must be turned off</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>• An FPGA device and a HardCopy companion device must be set</td>
</tr>
<tr>
<td>Set Current HardCopy Companion Revision</td>
<td>Specifies the companion revision to associate with the current design revision</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>A companion revision must already exist</td>
</tr>
<tr>
<td>Compare HardCopy Companion Revisions</td>
<td>Compares the FPGA design revision with the HardCopy companion design revision and generates a report</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>Both revisions must be compiled</td>
</tr>
</tbody>
</table>
Selecting the Prototype and Companion Devices

For both HardCopy device and FPGA prototype planning, the first stage is to choose the device family, device density, speed grade, and package that best suits your design requirements. You can select the appropriate companion device based on the device that you select for prototyping.

You can use the HardCopy Device Resource Guide to help you select the appropriate companion device, as described in “HardCopy Device Resource Guide”.

For information about the features available in each device density, including logic, memory blocks, multipliers, and phase-locked loops (PLLs), as well as the various package offerings and I/O pin counts, refer to the respective HardCopy device handbook, which is available on the Literature page of the Altera website at www.altera.com.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
<th>Applicable Design Revision</th>
<th>Restrictions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Generate HardCopy Handoff Report</td>
<td>Generates a report containing important design information files and messages generated by the Quartus II Compiler</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>Both revisions must be compiled</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>The Compare HardCopy Companion Revisions command must be successfully run</td>
</tr>
<tr>
<td>Archive HardCopy Handoff Files</td>
<td>Generates a Quartus II Archive File (.qar) specifically for submitting the design to the Altera HardCopy Design Center</td>
<td>HardCopy companion revision</td>
<td>Both revisions must be compiled</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>The Compare HardCopy Companion Revisions command must be successfully run</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>The Generate HardCopy Handoff Report command must be successfully run</td>
</tr>
<tr>
<td>Start HardCopy Design Readiness Check</td>
<td>Generates a report with the design’s settings, I/O check, PLL, and RAM usage checks</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>None</td>
</tr>
<tr>
<td>HardCopy Advisor</td>
<td>Opens the HardCopy Advisor, which guides you through the process of creating a HardCopy project</td>
<td>FPGA prototype design and HardCopy companion revision</td>
<td>None</td>
</tr>
</tbody>
</table>

Table 5–1. HardCopy Utilities Menu Options (Part 2 of 2)
HardCopy Device Resource Guide

The HardCopy Device Resource Guide compares the resources required to successfully compile a design with the resources available in the various HardCopy devices. The report rates each HardCopy device and each device resource according to how well it fits the design. The Quartus II software generates the HardCopy Device Resource Guide for all designs successfully compiled for FPGA devices. You can find this guide in the Fitter folder of the Compilation report. Table 5–2 describes the color codes used in the guide.

Table 5–2. HardCopy Device Resource Guide Color Legend

<table>
<thead>
<tr>
<th>Color</th>
<th>Package Resource (1)</th>
<th>Device Resources</th>
</tr>
</thead>
<tbody>
<tr>
<td>Green (High)</td>
<td>The design can map to the HardCopy package and has been fitted with target device migration enabled in the HardCopy Companion Device dialog box.</td>
<td>The resource quantity is within the range of the HardCopy device and the design can likely map if all other resources also fit. You still must compile the HardCopy revision to ensure the design is able to route and close timing.</td>
</tr>
<tr>
<td>Orange (Medium)</td>
<td>The design can map to the HardCopy package; however, the design has not been fitted with the target device migration enabled in the HardCopy Companion Device dialog box.</td>
<td>The resource quantity is within the range of the HardCopy device; however, the resource is at risk of exceeding the range for the HardCopy package. Compile your design targeting the HardCopy device as soon as possible to check if the design fits and is able to route and migrate all other resources. You might have to select a larger device.</td>
</tr>
<tr>
<td>Red (None)</td>
<td>The design cannot map to the HardCopy package.</td>
<td>The resource quantity exceeds the range of the HardCopy device. The design cannot migrate to this HardCopy device.</td>
</tr>
</tbody>
</table>

Note to Table 5–2:

(1) The package resource is constrained by the FPGA for which the design was compiled. Only vertical migration devices within the same package are able to migrate to HardCopy devices.

Use this report to identify potential HardCopy device candidates for your design. The HardCopy and FPGA device packages must be compatible. A logic resource usage greater than 100% or a ratio greater than 1:1 in a category indicates that the design will probably not fit in that specific HardCopy device.

The HardCopy architecture consists of an array of fine-grained HCells to build logic equivalent to FPGA adaptive logic modules (ALMs) and digital signal processing (DSP) blocks. The DSP blocks in HardCopy devices match the functionality of the FPGA DSP blocks, though timing of these blocks is different than the FPGA DSP blocks because they are constructed of HCell macros.

Memory blocks in HardCopy devices and FPGAs are equivalent. Preliminary timing reports of the HardCopy device are available in the Quartus II software. Final timing results of the HardCopy device are provided by the Altera HardCopy Design Center after the HardCopy back-end implementation process is complete.

For more information about the HardCopy device resources, refer to Fitter Resources Reports in Quartus II Help.
For more information about the HardCopy device resources, refer to the respective HardCopy series device handbook, which is available on the Literature page of the Altera website at www.altera.com.

The report example in Figure 5–4 shows the resource comparisons for a design compiled for an EP4SE230F29C2 device. Based on the report, the HC4E25FF484 device in the 484-pin FineLine BGA package is an appropriate HardCopy device. The EP4SE230F29C2 device is rated green because the device is specified as a migration target in the example. If the HC4E25FF484 device is not specified as a migration target during the compilation, its package and migration compatibility is rated medium (orange). The migration compatibilities of the other HardCopy devices are rated none (red), because the package types are incompatible with the FPGA device.

**Figure 5–4. HardCopy Device Resource Guide with Target Migration Enabled**

<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1 Migration Compatibility</td>
<td>High</td>
<td>Medium</td>
<td>Medium</td>
<td>Medium</td>
<td>Medium</td>
</tr>
<tr>
<td>2 Primary Migration Constraint</td>
<td>Package</td>
<td>Package</td>
<td>Package</td>
<td>Package</td>
<td>Package</td>
</tr>
<tr>
<td>3 Package</td>
<td>FBGA - 780</td>
<td>FBGA - 780</td>
<td>FBGA - 484</td>
<td>FBGA - 484</td>
<td>FBGA - 780</td>
</tr>
</tbody>
</table>

**Selecting the Companion Device**

In the Quartus II software, you can select a HardCopy companion device to ensure compatibility between the FPGA design and the HardCopy device’s resources. To select your HardCopy companion device, on the Assignments menu, click **Device** and select a companion device from the **Companion device** list in the **Device** dialog box.

Selecting a HardCopy companion device for your FPGA prototype constrains the memory blocks, DSP blocks, and pin assignments, so that your design fits in the HardCopy device resources. Pin assignments are constrained in the FPGA design revision, so that the HardCopy device selected is pin-compatible. The Quartus II software also constrains the FPGA design revision so that identical device resources are targeted in both the FPGA and the HardCopy ASIC.

Although not all FPGA ALM configurations are available in HardCopy devices, no restriction is made during synthesis of the FPGA. Unsupported configurations are converted to multiple cells for the HardCopy device.

You can also specify your HardCopy companion device using the following tool command language (Tcl) command:

```
set_global_assignment -name DEVICE_TECHNOLOGY_MIGRATION_LIST <HardCopy Device Part Number>
```

For example, to select the HC4E25FF484 device as your HardCopy companion device for the EP4SE230F29C2 FPGA, use the following the Tcl command:

```
set_global_assignment -name DEVICE_TECHNOLOGY_MIGRATION_LIST HC4E25FF484C
```
Applying Design Constraints

The HardCopy development flow requires that you plan specific steps in addition to the standard FPGA design flow, because you are developing your design for implementation in two devices: a prototype of your design in an FPGA and a companion revision in a HardCopy device for production. Additional settings and constraints in the Quartus II software are required to make the FPGA design compatible with the HardCopy device, and in some cases, you must remove certain settings in the design. This section explains the additional design constraints necessary for your design to be successful in both FPGA and HardCopy devices.

Limit DSP and RAM to HardCopy Device Resources

To maintain compatibility between the FPGA and HardCopy devices, your design must use resources that are common to both families. You must turn on the Limit DSP & RAM to HardCopy device resources option in the Device dialog box before submitting the design to the Altera HardCopy Design Center for back-end implementation. Turning on this option ensures that your design does not use resources in the FPGA device that are not available in the selected HardCopy device or vice versa.

For more information about the Limit DSP & RAM to HardCopy device resources option in the Quartus II software, refer to Device Dialog Box in Quartus II Help.

Enabling Design Assistant to Run During Compilation

You must use the Design Assistant in the Quartus II software to check all HardCopy designs for design rule violations before submitting the designs to the Altera HardCopy Design Center. Additionally, you must fix all critical and high-level errors reported by the Quartus II Design Assistant.

Altera recommends turning on the Design Assistant to run automatically during each compilation so that you can review the violations to determine which errors you must fix or which you can waive, iteratively.

For more information about the Design Assistant and its rules in the Quartus II software, refer to About the Design Assistant in Quartus II Help.

I/O Assignment Settings

Due to the complex rules governing the use of I/O cells and their availability for specific pins and packages, Altera recommends that I/O assignments be completed using the Pin Planner and the Assignment Editor in the Quartus II software. These tools ensure that all of the rules regarding each pin and I/O cell are applied correctly. The Quartus II software can export a .Tcl script containing all I/O assignments.

For more information about I/O location and type assignments using the Quartus II Assignment Editor and Pin Planner tools, refer to the Constraining Designs chapter in volume 2 of the Quartus II Handbook.
To ensure that the HardCopy mapping is successful, you must make accurate I/O assignments that include pin locations, I/O standards, drive strengths, and capacitance loading for the design. Ensure that the I/O assignments are compatible with all selected devices. Altera recommends assigning I/O assignments for all I/O pins. Leaving unassigned I/O assignments may result in incompatible assignments.

When mapping between the FPGA device and a HardCopy device, the I/O pin location must be assigned to the available common groups, called modular I/O banks, for both devices. Because HardCopy devices have fewer I/O banks than FPGA devices, the Quartus II software limits the I/O banks to only those available in HardCopy devices.

For more information about I/O banks and pins in HardCopy series devices, refer to the respective HardCopy series device handbook, which is available on the Literature page of the Altera website at www.altera.com.

HardCopy III I/O buffers support only the 3.0-V I/O standard with a maximum supply voltage (VCCIO) of 3.0 V. Therefore, when specifying the I/O standard for the Stratix III FPGA device with the HardCopy III companion device already selected, you must choose an I/O standard with a VCCIO of 3.0 V or less. Selecting an I/O standard that requires a VCCIO of 3.3 V results in a compilation error.


HardCopy IV I/O buffers support 3.3 V I/O standards, which you can use as transmitters or receivers in your system. The 3.3 V I/O standard can be supported by using the bank VCCIO at 3.0 V. In this method, the clamp diode (on-chip or off-chip), when enabled, can sufficiently clamp overshoot voltage to within the DC and AC input voltage specification. The clamped voltage can be expressed as the sum of the VCCIO and the diode forward voltage.

For more information about HardCopy IV I/O buffers, refer to the DC and Switching Characteristics of HardCopy IV Devices chapter of the HardCopy IV Device Handbook.

You must constrain the I/O standards for the design specifically for your HardCopy device. If you do not assign an I/O standard to an I/O pin, the Quartus II software assigns the I/O standard to 2.5 V by default, which may not be compatible with your design. To check supported I/O standards and identify incompatible I/O settings on the assigned I/O pins, run I/O assignment analysis by pointing to Start on the Processing menu, and then clicking Start I/O Assignment Analysis. The Start I/O Assignment Analysis command verifies the I/O settings and assignments.

Altera recommends verifying the correct output drive strength for the design because the default value in the Quartus II software might not be appropriate for your application. Assigning the right output drive strength improves signal integrity while achieving timing requirements. In addition, the output capacitance loading for both the output and bidirectional pins must be set in the I/O assignment for a successful HardCopy compilation.
Quartus II Fitter Settings

To make the HardCopy device implementation more robust across process, temperature, and voltage variations, the Altera HardCopy Design Center requires that you turn on the Optimize multi-corner timing option and set the Timing-driven compilation option to Optimize hold timing for the Quartus II Fitter.

The Optimize multi-corner timing option directs the Fitter to optimize a design to meet timing requirements at both the fast-timing and the slow-timing process corners and operating conditions. Setting the Timing-driven compilation option to Optimize hold timing enables the Fitter to optimize hold time by adding delay to the appropriate paths. You can set these Fitter options in the Fitter Settings page of the Settings dialog box.

For more information about Fitter settings in the Quartus II software, refer to Fitter Settings Page (Settings Dialog Box) in Quartus II Help.

Physical Synthesis Optimization

The physical synthesis optimizations performed in the FPGA device are mapped to the HardCopy companion revision for placement and timing closure. When designing with a HardCopy device first, you can enable physical synthesis optimizations for the HardCopy device. These post-fit optimizations are then passed to the FPGA revision. The optimizations in the base revision are mapped to the companion device architecture and the post-fit netlists of both devices are generated and compared. Therefore, you must have the identical physical synthesis settings for both the HardCopy ASIC and FPGA revisions in order to avoid revision comparison failure.

The Effort level on the Physical Synthesis Optimizations page of the Settings dialog box for HardCopy III and HardCopy IV devices must be set to Fast because the performance gain achieved compared to the compilation time is very limited.

Not all Physical Synthesis Optimizations settings are available when you map your FPGA device to a HardCopy companion revision.

For more information about setting physical synthesis optimizations for the FPGA revision of the designs in the Quartus II software, refer to Setting up and Running the Fitter in Quartus II Help.

Timing Settings

The TimeQuest Timing Analyzer is a complete static timing analysis tool that you use as a sign-off tool for FPGAs and HardCopy ASICs. The TimeQuest analyzer guides the Fitter and analyzes timing results after compilation and is the required timing analysis tool for all Quartus II software designs.

For more information about the TimeQuest Timing Analyzer, refer to The Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook and About TimeQuest Timing Analysis in Quartus II Help.
**TimeQuest Timing Analyzer Settings**

The Altera HardCopy Design Center requires that all HardCopy handoff files include a TimeQuest analyzer timing report for design review. In the TimeQuest analyzer timing report, you must include both fast- and slow-corner timing analysis for setup, hold, and I/O paths by turning on the `Enable multicorner timing analysis during compilation` option. This option directs the TimeQuest analyzer to analyze the design and generate slack reports for the slow and fast corners.

You must also direct the TimeQuest analyzer to remove the common clock path pessimism during slack computation by turning on the `Enable common clock path pessimism removal` option.

You can turn on these TimeQuest analyzer options in the TimeQuest Timing Analyzer page of the Settings dialog box in the Quartus II software.

For more information about TimeQuest analyzer options in the Quartus II software, refer to `TimeQuest Timing Analyzer Page (Settings Dialog Box)` in Quartus II Help.

**Constraints for Clock Effect Characteristics**

The `create_clock` and `create_generated_clock` commands create ideal clocks, but do not account for board effects. To account for clock effect characteristics, you can use the `set_clock_latency` and `set_clock_uncertainty` commands.

For more information about how to use these commands, refer to the `The Quartus II TimeQuest Timing Analyzer` chapter in volume 3 of the `Quartus II Handbook`.

You can use the `derive_clock_uncertainty` command to automatically derive the clock uncertainties in your Synopsys Design Constraints File (.sdc). This command is useful when you are unsure of the clock uncertainties. The calculated clock uncertainty values are based on I/O buffer, static phase errors (SPE) and jitter in the PLLs, clock networks, and core noise.

The following syntax is for the `derive_clock_uncertainty` command:

```
```

For more information about the `derive_clock_uncertainty` command in the Quartus II software, refer to `derive_clock_uncertainty` in Quartus II Help.

When the `derive_clock_uncertainty` command is used, a `PLL_PLLSPE_INFO.txt` file is automatically generated in the project directory. This file lists the names of the PLLs, as well as their jitter and SPE values in the design. This text file can be used by the HCII_DTW_CU_Calculator.

Altera strongly recommends that you use the `derive_clock_uncertainty` command in the HardCopy revision. The Altera HardCopy Design Center does not accept designs that do not have clock uncertainty constraints applied by either using the `derive_clock_uncertainty` command or the HardCopy II Clock Uncertainty Calculator, and then using the `set_clock_uncertainty` command.

For more information about how to use the HardCopy II Clock Uncertainty Calculator, refer to the `HardCopy II Clock Uncertainty Calculator User Guide`. 
**LogicLock Regions**

LogicLock regions are flexible floorplan location constraints that help you place logic on the target device. You can use LogicLock regions in FPGA designs targeted to HardCopy devices, which are also passed onto the HardCopy companion revision.

When LogicLock regions are created in a HardCopy device, they start with width and height dimensions set to (1,1), and the origin coordinates for placement are at \( X_1 \ Y_1 \) in the lower left corner of the floorplan. You must adjust the size and location of the LogicLock regions in HardCopy devices before compiling the design.

Altera recommends that you do not use floating LogicLock regions for HardCopy devices because floating LogicLock regions may affect the design’s ability to meet timing closure. Additionally, you must manually size and place HardCopy device LogicLock regions in the floorplan; you cannot set the LogicLock regions to Auto.

For more information about using LogicLock regions, refer to the *Analyzing and Optimizing the Design Floorplan* chapter in volume 2 of the *Quartus II Handbook*.

**PowerPlay Power Analyzer**

You can perform initial power estimation and analysis of your HardCopy and FPGA devices using the PowerPlay Early Power Estimator. You can then use the PowerPlay Power Analyzer for a more accurate estimation of your device’s power consumption.

For more information about using the PowerPlay Power Analyzer, refer to the *PowerPlay Power Analysis* chapter in volume 3 of the *Quartus II Handbook*.

**Incremental Compilation**

The Quartus II software offers incremental compilation to preserve the compilation results for unchanged logic in your design. This feature dramatically reduces your design iteration time by focusing new compilations only on changed design partitions. New compilation results are then merged with the previous compilation results from unchanged design partitions.

Quartus II incremental compilation within a single Quartus II project is supported for the base family for both the FPGA first and HardCopy first flows. Exporting and importing partitions is not supported in HardCopy ASIC or FPGA device compilations when there is a migration device setting.

For more information about using Quartus II incremental compilation, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook* and *About Incremental Compilation* in Quartus II Help.

**External Memory Interfaces**

The HardCopy I/O structure is equivalent to the Stratix I/O structure, providing high-performance support for existing and emerging external memory standards such as DDR, DDR2, DDR3, QDRII, QDRII+, and RLDRAM II.
A self-calibrating soft IP core (UniPHY) optimized to take advantage of HardCopy device I/Os in conjunction with the Quartus II TimeQuest Timing Analyzer, provides the total solution for the highest reliable frequency of operation across process, voltage, and temperature (PVT).

**Compiling the Design and Creating Companion Revisions**

After you finish applying constraints to your prototype design, you can compile your design and review the messages generated by the Quartus II software during compilation to check for potential problems. If you do not have violations that you must fix, you can proceed to creating or overwriting a companion revision, as described in “Creating a Companion Revision” on page 5–15. If you have violations that you must fix, you must fix the violations, recompile the design, and recheck for violations before proceeding to creating a companion revision.

**Creating a Companion Revision**

The Quartus II software creates specific HardCopy design revisions of the project in conjunction with the primary project revisions. These parallel design revisions for HardCopy devices are called companion revisions. You can create multiple design revisions for both the FPGA and the HardCopy device. For example, if your initial FPGA revision is named `top` and the corresponding HardCopy revision is named `top_hc`, you could create another FPGA revision, named `top_fpga`, and the corresponding HardCopy revision would be named `top_fpga_hc`.

Although you can create multiple design revisions, Altera recommends maintaining only one FPGA revision after you create the HardCopy companion revision.

After you have successfully compiled your FPGA prototype, you can create and compile the HardCopy companion revision of your design. You can associate only one FPGA revision to one HardCopy companion revision. If you create more than one revision or companion revision, set the current companion for the revision you are working on.

For more information about creating or setting a companion revision in the Quartus II software, refer to *Migrating a Design to a HardCopy or FPGA Device* in Quartus II Help.
Compiling the HardCopy Companion Revision

The Quartus II software contains preliminary timing models for HardCopy devices, and you can gauge the degree of performance improvement you can achieve in the HardCopy device compared to the FPGA by compiling your HardCopy design with preliminary timing information in the Quartus II software. The timing constraints for the HardCopy companion revision can be the same as the FPGA design that you use to create the revision. Altera verifies that the HardCopy companion device timing requirements are met in the Altera HardCopy Design Center.

After you create your HardCopy companion revision from your compiled FPGA design, select the companion revision in the Quartus II software design revision pull-down list as shown in Figure 5–5 or from the Revisions list, and then compile the HardCopy companion revision. After you compile your design in the Quartus II software, you can perform a comparison check of the HardCopy companion revision to the FPGA prototype revision as described in “Comparing HardCopy and FPGA Companion Revisions” on page 5–22.

HardCopy Design Readiness Check

The HardCopy Design Readiness Check (HCDRC) is available as one of the processing steps in the default compilation of both the FPGA first and the HardCopy first flows. This feature checks issues that must be addressed prior to handing off the HardCopy design to the Altera HardCopy Design Center for the HardCopy back-end process. The HCDRC is different from the user-driven approach in the HardCopy Advisor, in which you must manually open the advisor to check for violations.

The checks performed in the HCDRC include settings (global, instance, and operating settings), I/O, PLL, RAM, and ALTGX checks.

Turning the HardCopy Design Readiness Check On and Off

You can turn the HCDRC on or off in the More Compilation Process Settings dialog box, or by using the following Quartus II Settings File (.qsf) assignments:

```
set_global_assignment -name \FLOW_HARDCOPY_DESIGN_READINESS_CHECK ON
set_global_assignment -name \FLOW_HARDCOPY_DESIGN_READINESS_CHECK OFF
```

Setting Check

The Setting Check report lists the results of the setting checks from the Handoff report. The Setting Check report contains the sections described below.

Summary

The Summary section displays the number of settings that do not meet recommendations. One of the following messages appears:
<number> global setting(s) do not meet recommendation. Please review the recommendation and do appropriate correction as it may affect the result of the migration to HardCopy.

or

<number> instance setting(s) do not meet recommendation. Please review the recommendation and do appropriate correction as it may affect the result of the migration to HardCopy.

**Global Setting**
The Global Setting section displays recommendations for global settings only. Global settings with values other than the recommended values are highlighted in red.

**Instance Setting**
The Instance Setting section displays recommendations for instances assignments only. Instance settings with values other than the recommended values are highlighted in red.

**Operating Setting**
The Operating Setting section displays checks related to the recommended operating settings for the FPGA and the HardCopy device.

This check is primarily applicable to Stratix III devices used as prototype FPGAs because HardCopy III devices only support 0.9 V core voltage, whereas Stratix III devices support both 1.1 V and 0.9 V core voltage.

The Setting Check reports also include checking for illegal assignments in the HardCopy design flow.

An example of illegal assignment checks is shown in **Example 5–1**.

**Example 5–1. Illegal Assignment Checks**

<table>
<thead>
<tr>
<th>Setting</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT</td>
<td>ON (1)</td>
</tr>
<tr>
<td>SIGNAL_PROBE_ENABLE</td>
<td>ON</td>
</tr>
<tr>
<td>SIGNAL_PROBE_SOURCE</td>
<td>ON</td>
</tr>
</tbody>
</table>

Notes to **Example 5–1**:

(1) Refer to the section “RAM Usage Check” on page 5–19.
(2) SignalProbe is not supported in HardCopy ASICS.

**I/O Check**
The I/O check ensures that you have assigned location assignments for the pins, I/O standards, current strength assignments, output pin load assignments, termination assignments, and also checks for unconnected pins.

The following message appears in the message panel during compilation when the HCDRC detects missing I/O standard assignments:

<number> pin(s) have no explicit I/O Standard assignments provided in the setting file and default values are being used. Please add a specific I/O Standard assignment for these pins.
**Input Pin Placement for Global and Regional Clock Check**

Due to the difference in the interconnect delays between the FPGA and HardCopy device, using non-primary clock inputs as clock inputs in a design can cause timing closure to be a problem when migrating the FPGA to the HardCopy device. The Input Pin Placement for Global and Regional Clock check informs you of potential problems before finalizing the pin location, so that clock inputs can be moved to the primary clock input.

This check lists all the pins that drive the global or regional clock, but are not placed in a dedicated clock pad. All pins are required to have manual location assignments. Pins that are missing location assignments are listed in the Missing Pin Location Assignment report.

The following message appears in the message panel during compilation and also appears in the I/O Check Summary:

<\number> pin(s) drives global or regional clock, but is not placed in a dedicated clock pin position. Clock insertion delay will be different between FPGA and HardCopy companion revisions because of differences in local routing interconnect delays.

**PLL Usage Check**

The PLL Usage Check report lists PLL usage requirements and violations checks.

**PLL Real-Time Reconfigurable Check**

This check highlights the PLLs without PLL reconfiguration. PLL reconfiguration enables fine tuning of the PLLs in the design after manufacturing. PLL elements without PLL reconfiguration are listed in a table.

The following message appears in the message panel during compilation and also appears in the Logic Check Summary:

<\number> PLL(s) don’t have real time reconfiguration. It is highly recommended that each PLL to have PLL reconfiguration for designs migrating to HardCopy.

**PLL Clock Outputs Driving Multiple Clock Network Types Check**

This check is derived from the Design Assistant rule check for HardCopy devices (Rule ID H102) and lists all PLL instances in the current design that have clock outputs driving multiple clock network types.

The following message appears in the message panel during compilation if the HCDRC detects this type of violation:

Found <\number> PLL(s) with clock outputs that drives multiple clock network types.

**PLL with No Compensation Mode Check**

This check lists all PLLs that are in No Compensation operating mode. This setting is not recommended for a design migrating to a HardCopy device because of differences in the clock networks and the clock delays between the FPGA and HardCopy device.

The following warning message appears during compilation when a PLL is in a No Compensation mode:
PLL(s) is operating in a "No compensation" mode.

PLL with Normal or Source Synchronous Mode Feeding Output Pin Check

When a PLL is directly feeding an output pin, it must be set to Zero Delay Buffer operating mode. If a PLL is set either in Normal Compensation mode or Source Synchronous mode, a warning message is issued during compilation.

The following warning message appears during the runtime of HC Ready:

PLL(s) is in normal or source synchronous mode that is not fully compensated because it feeds an output pin -- only PLLs in zero delay buffer mode can fully compensate output pins.

RAM Usage Check

HardCopy series devices do not support initialized RAM blocks upon power-up. However, you can use the ALTMEM_INIT megafunction to initialize the RAMs of a HardCopy series device in your design with the content of a ROM.

For more information about the RAM Initializer megafunction, refer to the RAM Initializer (ALTMEM_INIT) Megafunction User Guide.

In HardCopy series devices, RAM blocks power up uninitialized. During the RAM Usage check, the HCDRC checks for RAMs that are initialized using a Memory Initialization File (.mif). RAM with a .mif is listed in a table.

The following warning message appears during compilation when the HCDRC detects this type of error:

RAM(s) have Memory Initialization File (MIF). HardCopy devices do not allow initialized RAM. Please ensure that no RAM is initialized by a MIF file.

Initialized Memory Dependency Testing

The Assembler module of the Compiler optionally enables you to write an FPGA programming file with an initialized checkerboard pattern for memory contents of M4K memory blocks for the FPGA revision. Use this option only on a parallel copy of your compiled FPGA design that you want to test on your board. Using this option in a FPGA revision to migrate to the HardCopy revision creates irreconcilable revision differences between the FPGA and HardCopy designs because the HardCopy handoff design cannot physically have initialized memory content.

Table 5–3 lists the checkerboard patterns that you can include to a revision .qsf.

<table>
<thead>
<tr>
<th>Checkerboard Patterns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT ON</td>
<td>Sets bit stream of 0101 in RAM.</td>
</tr>
<tr>
<td>set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT 0101</td>
<td>Default checkerboard pattern.</td>
</tr>
<tr>
<td></td>
<td>Sets bit stream of 0101 in RAM.</td>
</tr>
<tr>
<td></td>
<td>Same with setting ON.</td>
</tr>
</tbody>
</table>
To create a programming file with an initialized checkerboard pattern, perform the following steps in the Quartus II software:

1. Compile your completed FPGA revision to use for prototype testing. You must eventually use this FPGA revision to create your HardCopy companion revision.
2. Create and compile the HardCopy companion revision.
3. Compare your HardCopy companion revision.
4. Generate and archive the HardCopy handoff files for your design.
5. Switch back to your FPGA revision, and on the Project menu, click **Revisions**, and then double click **<<new revision>>** in the **Revisions** table.
6. In the **Create Revision** dialog box, type a revision name in the **Revision name** box and turn on **Copy database** and **Set as current revision**. This step copies your FPGA revision and sets the new revision as the current open revision.

### Table 5–3. Checkerboard Patterns (Part 2 of 2)

<table>
<thead>
<tr>
<th>Checkerboard Patterns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT 1010</td>
<td>Sets bit stream of 1010 in RAM.</td>
</tr>
<tr>
<td>set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT 0000</td>
<td>Sets bit stream of 0000 in RAM.</td>
</tr>
<tr>
<td>set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT 1111</td>
<td>Sets bit stream of 1111 in RAM.</td>
</tr>
</tbody>
</table>
| set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT RANDOM | Pseudo-random pattern.
- Enables you to choose pseudo-random seed in the **quartus.ini** (for example, setting pseudo-random seed to 2:
  
  `asm_randomized_checkerboard_pattern_seed = 2`.
- If pseudo-random seed is not set in the **quartus.ini**, the pseudo-random pattern defaults to pseudo-random seed 1. As long as the pseudo-random seed is fixed, then random pattern sets in RAM is consistent for each run, ensuring the exact same pattern is reproducible. |
| set_global_assignment -name USE_CHECKERED_PATTERN_AS_UNINITIALIZED_RAM_CONTENT OFF | Turns off uninitialized RAM dependency check. |
7. On the Assignments menu, click **Settings**, and then click **Assembler** in the **Category** list. Turn on **Use checkered pattern as uninitialized RAM content** on the **Assembler** page, or add one of the checkerboard patterns listed in Table 5–3 to the revision .qsf.

8. Run the Assembler in the FPGA revision to generate a new programming file for your FPGA.

9. Test the new programming file in your prototype environment to determine if your design has a dependency for FPGA RAM contents initialized with zeros after power-up and configuration.

Because the checkerboard pattern is used for testing, the patterns written in the RAM blocks for the new programming file may not detect all cases of zero-initialized RAM content dependencies. Some designs may detect only one bit as zero (for example, the LSB of a memory word), so this method may not detect all cases. This checkerboard pattern test will detect a case when a full RAM word line is expected as zeros at startup.

**ALTGX Usage Check**

Beginning in the Quartus II software version 10.0, the ALTGX Usage check performs checks on ALTGX instance usage for designs targeting Stratix IV GX and HardCopy IV GX devices.

The HCDRC checks all the ALTGX instances that are initialized in the design for connectivity with the ALTGX_RECONFIG instance. Example 5–2 shows the warning message, with the respective instance HSSI_CMU atom name, for ALTGX instances that do not connect to an ALTGX_RECONFIG instance:

**Example 5–2. Warning Message**

ALTGX megafunctions do not have ALTGX_RECONFIG megafunctions connected. Altera recommends connecting ALTGX_RECONFIG megafunction to each ALTGX megafunction when migrating your designs to HardCopy devices.

**Timing Closure and Verification**

After compiling the project for the FPGA and HardCopy designs, verify that the design meets your timing requirements. Review the messages generated by the Quartus II software during compilation to check for potential problems. Also, verify the design functionality between the FPGA and HardCopy devices with the **HardCopy Companion Revision Comparison** command as described in “Comparing HardCopy and FPGA Companion Revisions” on page 5–22.

You can also use third-party formal verification software, Cadence Encounter Conformal verification software, to run formal verification, and then compare the companion revisions. Formal verification is described in more detail in “Formal Verification of FPGA and HardCopy Revisions” on page 5–22.
Timing Closure with the TimeQuest Timing Analyzer

The TimeQuest Timing Analyzer is the timing analysis tool for all HardCopy devices during the front-end design process in the Quartus II software. After you specify the initial timing constraints that describe the clock characteristics, timing exceptions, and signal transition arrival and required time in the .sdc, the TimeQuest analyzer analyzes the timing paths in the design, calculates the propagation delay along each path, checks for timing constraint violations, and reports timing results.

Comparing HardCopy and FPGA Companion Revisions

The Quartus II software uses the companion revisions in a single Quartus II project to maintain compatibility between the FPGA and HardCopy ASIC. This methodology enables you to design with one set of RTL code that is used in both the FPGA and HardCopy ASIC, guaranteeing functional equivalency.

When making changes to your design in a companion revision, use the Compare HardCopy Companion Revisions command to ensure that your design matches your HardCopy design functionality and compilation settings. You must perform this command after both the FPGA and HardCopy designs are compiled and before you hand off the design to the Altera HardCopy Design Center.

The Comparison Revision Summary in the Compilation report identifies where assignments were changed between revisions or if there is a change in the logic resource count due to different compilation settings.

For more information about comparing companion revisions in the Quartus II software, refer to Migrating a Design to a HardCopy or FPGA Device in Quartus II Help.

Formal Verification of FPGA and HardCopy Revisions

Third-party formal verification software, Cadence Encounter Conformal verification software, is used for several FPGA and HardCopy families. The formal verification flow for HardCopy ASIC designs is a two-step process. First, run formal verification on the FPGA netlist to ensure that the FPGA netlist matches the RTL. Second, use the Compare HardCopy Revisions command in the Quartus II software to ensure that the HardCopy implementation matches the FPGA.

Although this flow is enabled, performing formal verification is not necessary due to the one-to-one mapping of logic between the FPGA prototype and the HardCopy ASIC.

To use the Conformal software with the Quartus II software project for your FPGA design revision, you must automatically run the EDA Netlist Writer during compilation so it can generate the necessary netlist and command files required to run the Conformal software.
To automatically run the EDA Netlist Writer during the compilation of your FPGA revision, perform the following steps:

1. On the Assignments menu, click **Settings**.

2. In the **Category** list, under **EDA Tool Settings**, click **Formal Verification**, and then in the **Tool name** list, select **Conformal LEC**.

3. Compile your FPGA and HardCopy design revisions.

The Quartus II EDA Netlist Writer produces the netlist for the FPGA revision. You can compare your FPGA post-compilation netlist to your RTL source code using the scripts generated by the EDA Netlist Writer.

After compiling both the FPGA and HardCopy revisions, you can run the **Compare HardCopy Revisions** command, as described in “Comparing HardCopy and FPGA Companion Revisions” on page 5–22 to ensure that the HardCopy implementation matches the FPGA.

For more information about using the Cadence Encounter Conformal verification software, refer to the **Cadence Encounter Conformal Support** chapter in volume 3 of the **Quartus II Handbook**.

### HardCopy Floorplan View

The Quartus II software displays the floorplan and placement of your HardCopy companion revision. This floorplan shows the preliminary placement and connectivity of all I/O pins, PLLs, memory blocks, HCell macros, and DSP HCell macros. Congestion mapping of routing connections can be viewed using the **Layers Setting** dialog box, which is available from the View menu of the Chip Planner. Congestion mapping is useful in analyzing densely packed areas of your floorplan that can reduce the peak performance of your design. The Altera HardCopy Design Center verifies final HCell macro timing and placement to guarantee that timing closure is achieved. Figure 5–6 shows an example of the HC4E25FF484 device floorplan.
In this small example design, you can see the placement of a DSP block constructed of HCell macros, various logic HCell macros, and an M9K memory block. A close-up view of this region is shown in Figure 5–7.

Figure 5–6. HC4E25FF484 Device Floorplan

Figure 5–7. Close-Up View of Floorplan
Performing ECOs with Quartus II Engineering Change Management with the Chip Planner

The Altera HardCopy Design Center performs final placement and timing closure on your HardCopy design based on the timing constraints provided in the FPGA design.

For more information about the Altera HardCopy Design Center process, refer to the respective HardCopy series device handbook, which is available on the Literature page of the Altera website at www.altera.com.

Performing ECOs with Quartus II Engineering Change Management with the Chip Planner

During the last stage of the design cycle, the ability to implement a specific portion of the design without affecting the rest of its logic is critical. As described in “Incremental Compilation” on page 5–14, incremental compilation enables you to implement and manage certain partitions of the design and preserve the optimization results for the rest of the design. However, implementing changes may become difficult to manage because Engineering Change Orders (ECOs) are often implemented as last-minute changes to your design.

With the Altera Chip Planner, you can shorten the design cycle time significantly. When changes are made to your design as ECOs, you do not have to perform a full compilation in the Quartus II software. Instead, you can make changes directly to the post place-and-route netlist, generate a new programming file, test the revised design by performing a gate-level simulation and timing analysis, and then verify the change on the system. When the change has been verified on the FPGA, switch to the HardCopy revision, apply the same ECOs, run timing analysis and the Assembler, compare the revisions, and then run the HardCopy Netlist Writer for design submission.

There are three types of migration scenarios:

- One-to-one changes, which are changes that can be implemented on both FPGA and HardCopy architectures
- Changes that must be implemented differently on the two architectures to achieve the same result
- Changes that cannot be implemented on both architectures

The following sections describe the methods for migrating each type of changes.
Migrating One-to-One Changes

One-to-one changes are implemented using identical commands in both architectures, and include changes that affect only I/O cells or PLL cells. Some examples of one-to-one changes include creating, deleting, or moving pins, changing pin or PLL properties, or changing pin connectivity (provided the source and destination of the connectivity changes are I/Os or PLLs). These types of changes can be implemented identically on both architectures.

To duplicate the same engineering change order (ECO) in the Quartus II software, use the Change Manager and record all ECOs for the FPGA revision. Ensure that the same ECO operations occur on each revision for both the FPGA and HardCopy ASIC revisions to avoid a revision comparison failure.

If such changes are exported to Tcl, directly reapplying the generated Tcl script (with a minor text edit) on the companion revision implements the appropriate changes as described in the following steps:

1. In the FPGA revision, open the Change Manager.
2. On the View menu, point to Utility Windows and click Change Manager.
3. Perform the ECO in the Chip Planner or Resource Property Editor. You will see the ECO operations in the Change Manager.
4. Export the changes from the Change Manager to Tcl, by right-clicking the entry, pointing to Export, and then clicking Export All Changes As…
5. Save the .tcl script, which you will use in the HardCopy revision.

In the HardCopy revision, apply the .tcl script to the companion revision by following these steps:

1. Open the generated Tcl script and change the project_open <project> -revision <revision> line to refer to the appropriate companion revision.
2. Save the .tcl script.
3. Apply the Tcl script to the companion revision. On the Tools menu, click Tcl Scripts and in the Tcl Scripts dialog box, select ECO Tcl and click Run.

Migrating Changes that Must Be Implemented Differently

Some changes must be implemented differently on the two architectures, such as changes affecting the logic of the design. Some examples include LUTMASK changes, LC_COMB/HSADDER creation and deletion, connectivity changes not described in the previous section, and different PLL settings for the FPGA and the HardCopy revisions.

For more information about how to use different PLL settings for the FPGA and HardCopy devices, refer to AN 432: Using Different PLL Settings Between Stratix II and HardCopy II Devices.
Table 5–4 summarizes the suggested implementation of various changes that must be implemented differently on the FPGA and HardCopy architectures.

<table>
<thead>
<tr>
<th>Change Type</th>
<th>Suggested Implementation</th>
</tr>
</thead>
<tbody>
<tr>
<td>LUTMASK changes</td>
<td>Because a single FPGA atom can require multiple HardCopy atoms to implement, you may need to change multiple HardCopy atoms to implement the change, including adding or modifying connectivity.</td>
</tr>
<tr>
<td>Make/Delete LC_COMB</td>
<td>If you use an FPGA LC_COMB in extended mode (7-LUT) or are using a SHARE chain, you must create multiple atoms to implement the same logic functions in the HardCopy device. Additionally, the placement of the LC_COMB cell has no meaning in the companion revision because the underlying resources are different.</td>
</tr>
<tr>
<td>Make/Delete LC_FF</td>
<td>Basic creation and deletion is the same on both architectures; however, similar to LC_COMB creation and deletion, the location of an LC_FF in a HardCopy and FPGA revision do not translate.</td>
</tr>
<tr>
<td>Editing logic connectivity</td>
<td>Because a LCELL_COMB atom may be required to be broken up into several HardCopy LCELL_COMB atoms, the source or destination ports for connectivity changes might have to be analyzed to properly implement the change in the companion revision.</td>
</tr>
</tbody>
</table>

### Changes That Cannot Be Migrated

A small set of changes are incompatible and cannot be implemented in both architectures. The best example of this incompatibility occurs when moving logic in a design; because the logic fabric is different between the two architectures, locations in the FPGA and HardCopy device are not compatible with each other.

### Overall Migration Flow

This section outlines the migration flow and the suggested procedure for implementing changes in both FPGA and HardCopy revisions to ensure a successful revision comparison so the design can be submitted to the Altera HardCopy Design Center.

#### Preparing the Revisions

The general process for migrating changes from the FPGA to HardCopy device or vice versa is the same and is described below:

1. Compile the design on the initial device.
2. Migrate the design from the initial device to the target device in the companion revision.
3. Compile the companion revision.
4. Run the Compare HardCopy Companion Revisions command. Both revisions must pass the revision comparison.

If testing identifies problems requiring ECO changes, equivalent changes can be applied to both FPGA and HardCopy revisions, as described in the following section.
Applying ECO Changes

The general process for applying equivalent changes in companion revisions is described below:

1. To make changes in one revision using the Chip Planner, and then verify and export these changes, follow these steps:
   a. Make changes using a Chip Planner tool (Chip Planner, Resource Property Editor, or Change Manager).
   b. Perform a netlist check using the Check and Save All Netlist Changes command.
   c. Verify correctness using timing analysis, simulation, and prototyping (FPGA only). If more changes are required, repeat steps a and b.
   d. Export change records from the Change Manager to Tcl scripts, or Comma-Separated Value File (.csv) or Tab-Separated Value File (.txt) formats. This exported file makes the equivalent changes in the companion revision.

2. Open the companion revision in the Quartus II software.

3. Using the exported file, manually reapply the changes using a Chip Planner tool. As stated previously, some changes can be reapplied directly to the companion revision (either manually or by applying the Tcl commands), while others require some modifications.

4. Run the Compare HardCopy Revisions command. The revisions must match.

5. Verify the correctness of all changes, which may require running timing analysis.

6. Run the HardCopy Assembler command and the HardCopy Netlist Writer command for design submission along with handoff files.
   - Use the following Tcl command to run the HardCopy Assembler:

     ```tcl
     execute_module -tool asm -args "--read_settings_files=off --write_settings_files=off"
     ```

   - Use the following Tcl command to run the HardCopy Netlist Writer:

     ```tcl
     execute_module -tool cdb \
     -args "--generate_hardcopy_files"
     ```

   For more information about using the Chip Planner, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.
Preparing the Design for Handoff

To submit a design to the Altera HardCopy Design Center for design review and back-end implementation, you must generate a HardCopy Handoff report and archive the HardCopy project.

Generating a HardCopy Handoff Report

The Generate HardCopy Handoff Report command creates the HardCopy Handoff report, which provides important information about the design for the Altera HardCopy Design Center to review.

After you generate the HardCopy Handoff report, you can archive the design using the Archive HardCopy Handoff Files command, which is described in “Archiving HardCopy Handoff Files”.

For more information about the Generate HardCopy Handoff Report command in the Quartus II software, refer to Generate HardCopy Handoff Report Command (Project Menu) in Quartus II Help.

Archiving HardCopy Handoff Files

The last step in the HardCopy design flow is to archive the HardCopy project for submission to the Altera HardCopy Design Center for HardCopy back-end implementation. The Archive HardCopy Handoff command creates a unique .qar, which is different than the standard file the Quartus II project archive utility generates. The HardCopy archive file contains only the necessary data from the Quartus II project required to implement the design in the Altera HardCopy Design Center.

For more information about the Archive HardCopy Handoff Files command in the Quartus II software, refer to Archive HardCopy Handoff Files Dialog Box in Quartus II Help.
Document Revision History

Table 5–5 shows the revision history for this chapter.

Table 5–5. Document Revision History (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
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<tbody>
<tr>
<td></td>
<td></td>
<td>Added Table 5–3 on page 5–19.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated Figure 5–4 on page 5–9, Figure 5–6 on page 5–24, and Figure 5–7 on page 5–24.</td>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
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<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Reorganized the chapter</td>
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<tr>
<td></td>
<td></td>
<td>Removed the “Quartus II Features for HardCopy Planning” section</td>
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<td>Changed the “HardCopy Recommended Settings in the Quartus II software” section to “Applying Design Constraints”</td>
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<td></td>
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<td>Organized the “HardCopy Utilities” subsections by design flow</td>
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<td></td>
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<td>Added the “Selecting the Prototype and Companion Devices” section</td>
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<td>Added the “I/O Assignments” section</td>
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<td>Added the “External Memory Interfaces” section</td>
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<td></td>
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<td>Added the “Compiling the Design and Creating Companion Revisions” section</td>
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<td></td>
<td></td>
<td>Added the “Timing Closure and Verification” section</td>
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<td></td>
<td>Added the “Preparing the Design for Handoff” section</td>
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<tr>
<td></td>
<td></td>
<td>Linked applicable sections to Quartus II Help</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Edited the “Timing Settings” section to remove support for the Classic Timing Analyzer</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Changed to new document template</td>
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<tr>
<td></td>
<td></td>
<td>Editorial changes</td>
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<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Added the “ALTGX Usage Check” section</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated the “LogicLock Regions” section for updated companion revision support</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated the “Incremental Compilation” section for updated companion revision support</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Linked sections throughout the chapter to Quartus II Help</td>
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<tr>
<td></td>
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<td>Removed the “Referenced Documents” section</td>
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<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Removed HardCopy Stratix legacy support information</td>
</tr>
<tr>
<td></td>
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<td>Updated the “Physical Synthesis Optimization” section</td>
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<td>Updated the “Quartus II Software Features Supported for HardCopy Designs” section</td>
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<td></td>
<td>Updated the “Referenced Documents” section</td>
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<tr>
<td></td>
<td></td>
<td>Updated the tables and figures for HardCopy Series devices</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Updated the “RAM Usage Check” section</td>
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Table 5–5. Document Revision History (Part 2 of 2)

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<th>Changes</th>
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<td>November 2008</td>
<td>8.1.0</td>
<td>• Added HardCopy IV E support information</td>
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<tr>
<td></td>
<td></td>
<td>• Added notes for Initialized Memory Dependency testing</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Changed page size to 8.5” × 11”</td>
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<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>• Updated the “RAM Usage Check” section</td>
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<td></td>
<td>• Updated “Referenced Documents”</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
6. Quartus II Design Separation Flow

This chapter contains rules and guidelines for creating a floorplan with the design separation flow, and assumes familiarity with the Quartus® II incremental compilation flow and floorplanning with the LogicLock™ feature.

The basic principle of a secure and reliable system is that critical subsystems in the design have physical and functional independence. Systems with redundancy require physical independence to ensure fault isolation—that a failure or corruption of any single subsystem does not affect any other part of the system adversely. Furthermore, if errors occur, physical independence simplifies analysis by allowing developers to evaluate each subsystem separately.

Traditionally, systems that require redundancy implement critical IP structures using multiple devices. The Quartus II design separation flow, used in Cyclone® III LS devices, allows you to design physically independent structures on a single device. This functionality allows system designers to achieve a higher level of integration on a single FPGA, and alleviates increasingly strict Size Weight and Power (SWaP) requirements. Figure 6–1 shows this concept.

**Figure 6–1. Achieving Higher Level Integration on a Single Cyclone III LS Device**

The Quartus II design separation flow introduces the constraints necessary to create secured regions and floorplan a secured system. When implemented in Cyclone III LS devices, a secured region provides physical independence through controlled routing and a boundary of unused resources. Restricting routing resources and providing a physical guard band of unused logic array blocks (LABs) prevent faults or unintended signals originating in one secured region from adversely affecting other design blocks on the device.

The Quartus II design separation flow features require specific licensing in addition to licensing the Quartus II software. For more information, contact your local Altera sales representative or Altera distributor.
The Quartus II design separation flow incorporates additional LogicLock and floorplanning features into the incremental compilation flow. The following three chapters in the *Quartus II Handbook* serve as companion references to this chapter:

- **Quartus II Incremental Compilation for Hierarchical and Team-Based Design**—Describes the Quartus II incremental compilation flow
- **Best Practices for Incremental Compilation Partitions and Floorplan Assignments**—Contains guidelines for using the incremental compilation flow and creating a design floorplan
- **Analyzing and Optimizing the Design Floorplan with the Chip Planner**—Describes various attributes associated with LogicLock location constraints and introduces the Chip Planner for creating and modifying a floorplan

### Design Flow Overview

The design separation flow is based on the incremental compilation flow. You begin with an incremental compilation design flow and then apply design separation constraints to each design partition that you want to physically isolate from the rest of your design. This section provides an overview of the design separation flow steps.

Figure 6–2 shows a flow chart of the design separation flow. Red boxes in the flow chart highlight steps that are specific to the design separation flow, while the remaining boxes in the flow chart are common to both the design separation and incremental compilation flows. This section provides a brief description for each step in Figure 6–2 and serves as a quick-start guide for the design separation flow.
1. **Set up design hierarchy for secured partitioning**—Prepare your design for implementation of the design separation flow, by setting up your design hierarchy for secured partitioning along logical hierarchical boundaries. If necessary, create wrapper files to create logical boundaries in the design hierarchy to support the design entities that you must separate from the remainder of the design.

2. **Perform analysis and elaboration**—Run analysis and elaboration to identify the hierarchy in your design.

3. **Create design partitions for secured regions**—For each design entity that requires physical independence, create a logical design partition for each design entity. Partition logic using guidelines from the incremental compilation flow.
For more information, refer to “Creating Design Partitions for the Design Separation Flow” on page 6–5.

4. **Create a design floorplan with security attributes**—After creating design partitions, create LogicLock location assignments and a floorplan to secure all the entities in your design. Use the security attributes in the LogicLock Regions window to specify the security level of each LogicLock region. These attributes create fencing regions in your floorplan to isolate the secured LogicLock regions. For more information, refer to “Creating a Design Floorplan with Secured Regions” on page 6–6.

5. **Assign design partitions to secured regions**—Assign design partitions to secured LogicLock regions to separate them from each other and from all other hierarchy blocks. Refer to “Using Secured Regions” on page 6–9 for more information.

6. **Add I/O pins that directly interface with a secured region as a member of the secured region**—If a secured region interfaces with one or more I/O pins, make the I/O pins members of the secured region. If a secured region has I/O pins as members, that region must overlap the I/O pads. Refer to “Adding I/O Pins as Members of Secured Regions” on page 6–9 for more information.

7. **Create security routing interfaces to and from secured regions**—Create security routing interfaces by applying the security routing interface attribute to LogicLock regions.

You can use only routing resources in a security routing interface; you cannot place any logic. Each security routing interface must abut one or two secured regions. After you create an interface region for each signal or group of signals entering or exiting a secured region, assign the signals to the appropriate routing interfaces.

For signals routing between secured regions with different security attributes or between a secured region and an unsecured region, you must lower the security attribute for the signal exiting the stricter security region. For more information, refer to “Making Signal Security Assignments” on page 6–19.

8. **Assign I/O pins**—After creating secured regions and security routing interfaces, if the secured regions contain I/O pins as members, assign I/O pins to meet design separation flow requirements. For example, secured regions cannot share I/O banks. If a secured region contains I/O pins as members, the entire I/O bank is usable only by the secured region that sinks or sources the I/O pin. For more information, refer to “Assigning I/O Pins” on page 6–25.

9. **Make design changes, set the netlist type for each design partition, and compile the design**—After making the necessary I/O pin assignments, you complete the design separation flow-specific steps, and you can start the iterative process of making design changes, setting the netlist type for each design partition, and then compiling your design until you achieve a floorplan that meets your design requirements.

> Subsequent sections in this chapter describes the design separation flow-specific steps (step 1 and steps 4 through 8).
Creating Design Partitions for the Design Separation Flow

After setting up your design to support secured partitioning and running analysis and elaboration, you can create design partitions.

Each secured region floorplan assignment uses a single design partition in the incremental compilation flow to identify the functional elements that belong to a secured region. You can make design partition assignments along entity boundaries in the RTL design hierarchy.

Only a single partition may be used in a secure region. Plan your design entities such that logic that requires physical isolation from the rest of your design are contained in a single design entity. Additionally, you must create wrapper files where necessary to reorganize your hierarchy, so that a single entity or module in your RTL contains all your secured regions. The incremental compilation feature allows functional independence of each design partition because it disables netlist optimizations across partition boundaries.

Most of the rules, guidelines, and tools for creating design partitions used in the incremental compilation flow are applicable in the design separation flow. You can use the Incremental Compilation Advisor, the Design Partition Planner, and the Chip Planner features in the Quartus II software to help you create design partition assignments.

When creating design partitions, the following considerations are important:

- Register the inputs and outputs of a design partition to avoid cross-boundary logic optimizations and to maintain timing performance along the signal path.
- Minimize the number of I/O paths that cross partition boundaries to keep logic paths in a single partition for optimization. Minimizing the number of cross-boundary I/O paths makes partitions more independent for both logic and placement optimization.
- Avoid logic that requires cross-boundary logic optimizations.

For more information about guidelines for creating design partitions, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

When creating your design in the design separation flow, you must be aware of some restrictions and special considerations that differ from the incremental compilation flow. For more information about these considerations, refer to the “Merging PLL Resources” and “Avoiding Multiple Design Partitions With a Secured Region” sections.

Merging PLL Resources

In the Quartus II incremental compilation flow without design separation constraints, the Fitter can use the same PLL resource on the device when multiple design partitions instantiate a PLL with the same parameters. When you enable the design separation flow and a design contains one or more secured regions, the Quartus II software disables PLL merging across design partitions, which helps to maintain the physical separation between design partitions. The Quartus II software also disables PLL merging for the entire design, even if LogicLock regions in a Cyclone III LS
design contain no security attributes. For partitions that require shared PLL resources, you must instantiate the PLL outside of the design partitions.

**Avoiding Multiple Design Partitions With a Secured Region**

The Quartus II software does not allow multiple design partitions, including child partitions and multi-hierarchy partitions. Each secured region, which you designate after creating design partitions, must contain only a single design partition.

Child partitions are design partitions from a subentity of an existing design partition and would potentially create multiple design partitions in a secured region, so they are not allowed in the design separation flow.

You can create multi-hierarchy partitions by merging multiple design partitions from different branches of the hierarchy. Merge these partitions into a single netlist during elaboration to enable cross-boundary optimizations during synthesis and fitting, and result in a single incremental result for each multi-hierarchy partition. Multi-hierarchy partitions function similarly as single-hierarchy partitions, but must contain hierarchies from a common parent partition. The Quartus II software does not allow multi-hierarchy partitions in the design separation flow.

**Creating a Design Floorplan with Secured Regions**

After creating design partitions, you can create a design floorplan with secured regions with the Chip Planner and security attributes in the LogicLock Regions window.

The Quartus II software uses LogicLock location assignments to map logic in your design hierarchy to physical resources on the device. The Chip Planner provides a visual floorplan of the entire device and allows you to move and resize your LogicLock location constraints on the floorplan of the device. The design separation flow adds a security attribute constraint to each LogicLock region to further constrain routing to achieve physical isolation between LogicLock regions. Assign signals that require connectivity between two secured regions or between a secured region and unsecured logic to a special LogicLock region known as a security routing interface. A security routing interface is a controlled region that limits the routing of the contained signals to only the one or two LogicLock regions that this region abuts.

To create fault isolation between secured regions, the design separation flow selectively shuts off routing around the periphery of a secured region. Because signal connectivity at the boundary of the secured region is unused, any faults that occur in the secured region are prevented from adversely affecting neighboring regions. Fault isolation, when using the design separation flow, is possible because no physical connection exists to propagate the fault outside of the region.

Cyclone III LS devices use a MultiTrack interconnect architecture that consists of row and column interconnects that span fixed distances to achieve signal connectivity between LABs. In the horizontal direction, row interconnects use wire resources that span 1 LAB, 4 LABs, and 24 LABs. These row-routing resources are direct link interconnects, R4 interconnects, and R24 interconnects, respectively. In the vertical direction, routing resources span distances of 1 LAB, 4 LABs, and 16 LABs. These column routing resources are register chain interconnects, C4 interconnects, and C16 interconnects, respectively. In the design separation flow, the Quartus II software disables LogicLock region routing wires (C4, C16, R4, and R16) that cross outside the
border of a boundary. Each secured region uses an unused boundary (or a fence) of LABs to guard against the faults from wire resources spanning a length of one LAB (direct link and register chain routing resources) from affecting a neighboring region.

The rules and guidelines for floorplanning in the design separation flow are similar to those in a typical compilation flow. However, there are some special considerations for the relative placement of secured regions in your design floorplan. Because each secured region is a keep-out region for routing resources from other LogicLock regions, ensure that a routing path with valid communication interfaces exists between secured regions. Furthermore, the routing path (encapsulated in a security routing interface) should not follow a circuitous path and must be simple enough to meet your timing requirements.

The Fitter cannot generate a placement for LogicLock regions with security attributes. You must manually place LogicLock regions with security attributes; that is, the size attribute cannot be Auto, and the state attribute cannot be Floating for any LogicLock region in a secured design.

You can use a Fitter-generated floorplan, created without security attributes, as a starting point to create a final floorplan for the design separation flow.

To use a Fitter-generated floorplan as an initial floorplan, apply Reserved attributes to LogicLock regions that must be physically isolated from the rest of your design. A Fitter-generated floorplan with Reserved attributes generates non-overlapping LogicLock regions. You can modify the initial floorplan by adjusting the relative placement for each secured region, taking into account the connectivity requirements for each region.

For more information about using the Chip Planner settings and options, refer to the Analyzing and Optimizing the Design Floorplan with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

### Using Security Attributes

The Security Attributes column in the LogicLock Regions window and the Security tab in the LogicLock Regions Properties dialog box are available when you license your version of the Quartus II software specifically for the design separation feature. Setting the Security attribute applies a constraint to a LogicLock region, making the region either a secured region or a security routing interface, from where signals enter or exit a secured region.

The software populates the Signals list with the inputs and outputs of secured regions after analysis and synthesis. Columns in the Signals list describe the Security Level, the security routing interface the signal is assigned to, and whether the signal is an output or input to the region.
Figures 6–3 and 6–4 show the design separation flow security features in the LogicLock Regions window and the LogicLock Regions Properties dialog box.

Table 6–1 lists a summary of the Security Attributes available for the design separation flow.

Table 6–1. Security Attributes for LogicLock Regions (Part 1 of 2)

<table>
<thead>
<tr>
<th>Security Attribute</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Unsecured</td>
<td>Removes the constraint for physical isolation.</td>
</tr>
<tr>
<td>1</td>
<td>Creates a secured region. Physically isolates the LogicLock region by restricting routing resources from leaving the region. Creates a one-LAB width border of unused logic (LABs, DSP blocks, or embedded memory blocks) around the LogicLock region. Applying this attribute to a LogicLock region sets the global assignment LL_REGION_SECURITY_LEVEL 1 for the LogicLock region.</td>
</tr>
</tbody>
</table>
Creating a Design Floorplan with Secured Regions

When you apply a secured region attribute (1 or 2) to an existing LogicLock region, the LogicLock region must have a fixed size with a locked origin. Each secured region must have a minimum size of eight-LABs in both the horizontal and vertical dimensions. A region smaller than 8 × 8 LABs may be non-routable when using the design separation flow.

The Quartus II software does not allow child regions when creating a secured region because a secured region contains only a single partition.

Adding I/O Pins as Members of Secured Regions

A secured region must contain all required physical device resources to complete compilation. I/O pads that are members of a secured region must be contained in the boundaries of the secured region that sources or sinks it. That is, a secured region must overlap the I/O pads that are members of the region. If the logic in the secured region instantiates a PLL or a clock block, those physical device resources must also be overlapped by the region.

You can add I/O pins as members of a secured region using the LogicLock Region Properties dialog box.

Using Security Routing Interfaces

A LogicLock region with the security routing interface security attribute creates a routing channel for signals to and from a secured region. You may not place logic in a security routing interface. Each security routing interface can connect two secured regions, or a secured region with one or more unsecured regions. If you are connecting two secured regions, the Quartus II software places a fencing region around the interface region automatically. You can assign each signal entering or exiting a secured region to a security routing interface on the Security tab in the LogicLock Regions Properties dialog box.

For information about assigning signals to a security routing interface, refer to “Making Signal Security Assignments” on page 6–19.

For information about the number of signals that you can contain in a security routing interface, refer to “Routing Restrictions” on page 6–27.
Making Design Separation Flow Location Assignments in the Chip Planner

The Chip Planner allows you to modify the size and location of LogicLock regions visually. This section describes the attributes of LogicLock regions in the context of the design separation flow.

When you enable the design separation flow, the Chip Planner shades the fencing region around each secured region in gray and security routing interfaces in green. The Chip Planner highlights illegal placements that violate secured region boundaries in red at the location in which the violation occurs. Figure 6–5 shows the LogicLock regions with security attributes in the Chip Planner.

Figure 6–5. LogicLock Regions With Security Attributes

Notes to Figure 6–5:
(1) Floorplan Editing Mode task.
(2) Unused fence around a secured region.
(3) Security violation, created by a LogicLock region placement in a fencing region of a secured region.
(4) Security routing interface region connecting two secured regions.
(5) Security routing interface region connecting secured region and unsecured logic.
Understanding Fencing Regions

The Quartus II software automatically adds a fencing region, which is a border of unused logic (LABs, DSP blocks, or embedded memory blocks) when you apply security attribute 1 or security attribute 2 to a LogicLock region. You may not place any logic into a fencing region. The Fitter does not use any routing wires that exit the fence boundary of a secured region. Because you can use direct drive and carry chains at the edge of a secured region, the fencing region prevents signals driven on one length one wires (in the horizontal and vertical directions) from exiting the secured region.

The fencing region around a secured region uses one unit of unused logic horizontally and one unit of unused logic vertically for security attribute 1. A fencing region for LogicLock regions of security attribute 2 uses one unused logic block horizontally and two unused logic blocks vertically. The following regions require special fencing regions:

- Vertical I/O regions
- Areas around the configuration engine

I/O banks along the top and bottom of the chip use only vertical routing wires to and from the I/O Elements (IOEs). The heavy use of C4 wires from IOEs creates a four-LAB fence between the vertical I/O banks and a secured region. Secured regions requiring a connection to I/O in the top or bottom banks of the device optimally use resources if you add the I/O signals as members of the secured region and overlap the corresponding I/O pads in the floorplan. In Figure 6–6, Secured Region2 is five LABs away from the bottom of the device and Secured Region1 is four LABs away from the bottom of the device.

Figure 6–6. Vertical Fencing Near I/O Banks
A configuration engine is a hard IP block that manages the configuration of the device. Additionally, the configuration engine routes the control signals for the CRC detection circuit and the internal oscillator into the core logic on the device. In the design separation flow, the Quartus II software automatically adds a one-LAB fence around the configuration engine whenever a secured region occupies the same LAB column as the configuration engine. The configuration engine is a region notched out of the left side in the middle of the device.

All control signals to and from the configuration engine route from its right edge. If you use an instantiated WYSIWYG that uses any control signals to and from the configuration engine, the signals must either interface with unsecured logic or they must interface with a secured region through a security routing interface.

If your design routes signals to and from the configuration engine, then do not place a secured region that directly abuts the configuration engine signal interface (along the right side of the configuration engine) to avoid a Fitter error.

Figure 6–7 shows a configuration engine with a fencing region in the floorplan.

![Configuration Engine](image)

You can overlap fencing regions between two secured regions. That is, you can separate two adjacent secured regions by a one-row fence. The Chip Planner issues a security warning violation if you place a LogicLock region in the boundary of a secured region. The Chip Planner highlights security violations in red and the tooltip of a secured region indicates the locations of all security violations. You may receive an error if you try to compile a design with a security violation. Figure 6–8 shows two regions with overlapping fences and a security violation from an unsecured region.
Creating Non-Rectangular Regions

You can create non-rectangular regions by creating multiple rectangular regions and then merging them.

For more information about creating non-rectangular regions in the Chip Planner in the Quartus II software, refer to Creating and Manipulating LogicLock Regions in Quartus II Help.

Non-rectangular LogicLock regions in the design separation flow make circuitous routes more likely. As such, non-rectangular regions can have an adverse affect on performance when used with the design separation flow.

If a secured non-rectangular region contains a subregion that is less than 8 × 8 LABs, the chances of a non-routable situation occurring increases. Subregions that deterministically require the use of certain routing resources may not fit successfully if a violation of the secured region occurs. As a general guideline, each subregion should be 8 × 8 LABs or larger, to ensure that routing resources with a length of four LABs are readily available. In Figure 6–9, each subregion of Region 2 (labeled A, B, C, and D) are less than 8 × 8 LABs in dimension. These subregions can potentially cause a no-fit error. Depending on the placement and connectivity of LABs, certain routes may be difficult to achieve. For example, the Fitter would not be able to route a connection from LAB 1 to LAB 2 in region A directly. While another path may be possible, a series of hops that do not leave the LogicLock region may not be available and may not satisfy the timing requirements of the route.
Guidelines for the Relative Placement of Secured LogicLock Regions

Because each secured region is a keep-out region for placement and routing of any logic that is not a member of the secured region, you must be aware of the guidelines in this section as you lay out your floorplan. Placement that does not account for the connectivity requirements between LogicLock regions may cause poor performance or a non-routable design. The guidelines for floorplanning when using the design separation flow include:

- Create a complete floorplan, including location assignments for unsecured logic.
- Create a non-circuitous route between secured regions requiring a routing region. Routing regions between secured regions should be rectangular.
- Create security routing interfaces between secured regions that do not intersect with other routing regions; secured regions and their routing edges must fit on a single plane. A secured region must overlap any physical resources (such as I/Os, PLLs, and CLKCTRL) that the design partition in the secured region instantiates.
- Abut the secured region to the edge of the device whenever possible.

Creating a Complete Floorplan

You must allocate a region for all logic in your design. If you have a large secured region that divides the device into multiple disjointed regions, and you have unsecured logic that is not floorplanned, the design may not be routable.

If an unsecured partition does not contain any location assignments, the placement algorithms may make logic assignments on any unallocated space on the device. In the floorplan shown in Figure 6–10, the source and sink registers do not have a valid path through the device, because Secured Region 1 and Secured Region 2 occupy all routing channels.
If a complete floorplan is impossible for all partitions in your design, you can use empty LogicLock regions with the `Reserved` attribute to prevent the Fitter from placing any logic in a region that can potentially cause a no-fit error. For the example provided in Figure 6–10, you can place an empty region in the upper-left corner of your device to prevent the Quartus II software from placing any logic that has not been floorplanned there, as shown in Figure 6–11.
Ensuring Routability Between Regions

The Quartus II software cannot create auto-generated location constraints for any region with a security attribute. If you use a Fitter-generated placement as a starting point for a floorplan with security attributes, an optimal floorplan in a design without separation may not work in the same design. In a floorplan without secured regions, the Quartus II software restricts only the placement of logic. The Fitter may use all routing resources on your device, and may route through a LogicLock region to reach a destination. Secured regions reserve all routing resources in the LogicLock boundary to the design partition contained in the region.

Having a circuitous route between two regions degrades performance and may cause a non-routable design. Modify any regions that have signal connectivity and must route around a secured region to achieve a connection. Figure 6–12 shows a floorplan that does not contain disjointed parts. However, the source region must route around a secured region to connect to the sink region.
Ensuring Planarity

The Quartus II software automatically creates a fence around a security routing interface connecting two secured regions. Because no other routing resources may pass through a security routing interface connecting two secured regions, you must model all secured regions as nodes in a routing graph and all security routing interfaces as the edges, and all nodes and their edges must fit on a planar graph (that is, none of the edges can intersect). If you have five or more secured regions on the device, and each secured region contains signals that fan out to multiple secured regions, a planar floorplan may not be possible. Figure 6–13 shows a routing graph with five nodes. A complete graph with each pair of distinct vertices connected by an edge is impossible without having any of the edges cross. If the topology of your floorplan contains such a non-routable arrangement, you must rearrange your design hierarchy to collapse related design partitions into a single design partition.

If you can model your secured regions and security routing interfaces as a planar graph, but have a high degree of connectivity between the components, then you may have to rearrange the shape, size, or location of the secured regions to generate a routable floorplan. For instance, the hypothetical floorplan shown in Figure 6–14 does not have a valid routing path BD (between region B and region D). The modified floorplan in Figure 6–15 shows how you can achieve all the required connections on a planar surface.
You can use the Design Partition Planner for a visual representation of the connectivity between design partitions. This tool helps you determine if you can arrange the secured regions in your design on a planar floorplan. Figure 6–16 shows the Design Partition Planner.
Placing Physical Resources
You must contain all physical resources that the secured region requires inside the boundary of the secured region, including I/O pins that connect to the secured region and primitives that you instantiate in the secured region, such as PLLs and clock control blocks.

Making Signal Security Assignments
Each signal that enters or exits a secured region must have a security level attribute. You must also explicitly assign the signal to a security routing interface. The Quartus II software automatically assigns the security level for each signal a default value. The default value matches the secured region that is the source of the signal. Possible security levels of a signal include: Auto, Unsecured, 1, and 2. Auto sets the default security level for the signal.

A signal with a security attribute may connect to a region with an equivalent or higher security level. For example, a signal with a security level of Unsecured can drive logic in a region set to Unsecured, 1, or 2 and a signal with a security level of 1 can drive logic in a region set to 1 or 2. A signal originating from a secured region may not drive logic in a region with a lower security level. If you have a signal from a higher security level that must drive logic in a lower security level, you can direct the Fitter to honor the connection by explicitly lowering the security level of the signal.

At most, each security routing interface connects two regions. If a signal fans out to multiple regions, assign the signal to multiple security routing interface regions; one interface region per destination.

You can assign signals to security routing interfaces and the security level of signals with the Security tab in the LogicLock Region Properties dialog box, as shown in Figure 6-4.
To assign a signal to a security routing interface, follow these steps:

1. On the Security tab of the LogicLock Regions Properties dialog box, select a signal name in the Signals list, and then click Edit. The Edit Security Assignments for Signal dialog box appears, as shown in Figure 6–17.

   Alternatively, you can select multiple names in the Signal list by pressing the Ctrl key, clicking multiple names, and then clicking Edit.

The Quartus II software populates the Signals list with the names of signals entering and exiting the secured region after analysis and synthesis and after completing a successful partition merge.

### Figure 6–17. Edit Security Assignments for Signal Dialog Box

2. If necessary, lower the security level of the signal by specifying the Security level.

3. Select the security routing interface for signal or signals assignment. You can assign signals that fan-out or fan-in to multiple regions to multiple security routing interfaces.

### Understanding Signal Names

The list of signals entering and exiting a secured region consists of signal names from the post-map netlist. Signal outputs from a secured region derive from the output port name, as specified in the top-level RTL entity in the secured region. Signal inputs to a secured region derive from the name of the output port name that feeds the secured region. In the design separation flow, the Quartus II software preserves output port names through the compilation process. The output port name is also an alias for the logic or register that feed them.

The post-map region output signals listed in the signal list coincide with the signal name in the post-fit netlist. However, combinational signal names from unsecured or unpartitioned logic that feed a secured region may change through the compilation process. The Quartus II software optimizes many of the RTL signals during synthesis.
and placement and routing. Frequently, RTL signal names may not appear in the post-fit netlist after optimization. For example, the compilation process can add tildes (~) to nets that fan-out from a node, making it difficult to decipher which signal nets they actually represent. Use the post-compilation filter in the node finder to add additional signals to a security routing interface. When possible, use registered signals as inputs into a secured region, and register the output signals from a secured design partition.

**Working with Global Signals**

Global signals are low-skew routing lines that drive throughout the device. Global signals do not require an interface region to drive into a secured region. Cyclone III LS devices contain 20 global routing resources for use with high fan-out signals, such as clocks or control signals. A clock control block accesses each global signal. You can drive each clock control block directly by external clock pins, PLL outputs, or a signal generated from internal logic. You can locate a clock control block on the periphery boundary of your device.

For more information about the clock networks in Cyclone III LS devices, refer to the Clock Networks and PLLs in Cyclone III Device Family chapter in volume 1 of the Cyclone III Device Handbook.

In a compilation flow without security assignments, the Quartus II software automatically promotes signals with a high fan-out (such as clock pins and control signals) to use global clock resources. In the design separation flow, the Quartus II software disables automatic global promotion. You must manually promote signals with high fan-out requiring global routing resources to drive a clock control block.

The Quartus II software cannot promote signals onto a global routing resource through a global signal assignment from within a secured region. The Fitter only allows a clock promotion assignment to a signal if the signal is in an unsecured region. If you have a signal inside of a secured region that must use a global routing resource, you must first route the signal outside of the secured region before applying a global promotion assignment. You must assign the signal to a security routing interface and lower the security level of the signal.
To honor a global promotion assignment, a clock control block that is not overlapped by a secured region and a routing path to the clock control block must be available. There are five clock control blocks located on each side of the device, along the horizontal and vertical axes that run through the center of the device. Figure 6–18 shows the location of the clock control blocks and the PLLs for a 3CLS70 device in the Chip Planner floorplan.

Figure 6–18. PLL and Clock Control Block Location on a EPC3SL70 Device
You can manually instantiate PLLs and clock control blocks in the design partition of a secured region using the ALTPLL and ALTCLKCTRL megafunctions, respectively. Instantiation of the ALTCLKCTRL megafuntion in a secured partition forces the global promotion of the signal driving the clock control block. To generate a valid placement when you instantiate PLLs or a clock control block, the secured region containing the physical resource must overlap a free PLL, a free clock control block, or both.

You must be aware of certain restrictions when you instantiate a PLL in a secured region. Secured regions with a PLL fed by an external clock pin must contain the PLL and a valid clock pin that can drive the PLL. Each PLL has a set of dedicated clock control blocks that it can access, located to the right (clockwise) of the PLL in the device floorplan.

Because automatic promotion of signals onto a global resource is not allowed, you must not place a PLL and the clock control block that the PLL drives in the same secured region. If your design has a PLL inside of a secured region, you must assign the PLL output to a security routing interface and then lower the security level of the PLL output.

A secured region must not cover the clock control block associated with the PLL. There are two sets of dedicated clock pins that can drive a PLL input. The pads for the clock input pins are co-located with the clock control blocks. If you use the clock input pin that is co-located with the clock control block associated with the PLL, you cannot add the clock pin as a member of the secured region. Instead, you must either assign the clock pin to a security routing interface that connects with the secured region, or you can apply the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment to relax the fitter restriction on the clock input pin.

For more information about the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment, refer to “Assigning I/O Pins” on page 6–25.

Figure 6–19 shows examples of valid placement and invalid placement of secured regions that instantiate PLLs, before applying the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment.
Figure 6–19. Location of Valid and Invalid PLL, Clock Pin, and Clock Control Block Placement in a Cyclone III LS Device

Notes to Figure 6–19:
(1) Five clock control blocks are available on each side.
(2) You cannot use remote clocks to feed the PLLs.
(3) Dedicated clock paths can feed into this PLL. However, these are not fully-compensated paths.
(4) This secured region contains a PLL that an external clock pin feeds, whose outputs drive the clock control block through an unsecured region.
(5) This secured region contains a PLL whose output drives a clock control block in the same secured region. This placement is invalid.
Assigning I/O Pins

After ensuring that signals that enter or exit a secured region contain a security level attribute and after you have explicitly assign the signals to a security routing interface, you must also ensure that I/O pin assignments adhere to design separation flow guidelines. Consider the following three rules, in addition to the typical pin assignment rules, when assigning I/O pins with the design separation flow enabled:

- You must assign I/O pins that connect to a secured region as a member of that secured region or to a security routing interface region that abuts the secured region.
- You must ensure that secured regions with I/O pins as members do not share I/O banks with any other region.
- You must ensure that I/O pins associated with different secured regions or security levels do not use adjacent pins.

When I/O pins directly connect to the secured region, you may add I/O pins as members of a secured region. To add I/O pins as members of a secured region, in the LogicLock Regions Properties dialog box, on the General tab, click Add node. If an I/O pin is a member of a secured region, the I/O pad must be physically in the region, and the secured region must overlap the I/O resource.

If you do not add the I/O pin as a member of the secured region, you must assign the I/O signal to a security routing interface that abuts the secured region. This security routing interface must connect the secured region to the root region or another unsecured region. Explicitly lower the security level of any output signals from the secured region that connects to I/O pins.

**CAUTION**

I/O signals that route out to unsecured logic are no longer guaranteed to be physically isolated from other signals in your design.

Each I/O pin is adjacent to eight other pins: four along the horizontal and vertical axes, and four in the two diagonal axes, as shown in Figure 6–20.

**Figure 6–20. Pin Adjacency**

![Pin Adjacency Diagram](image)
Pins from different I/O banks may not share an adjacent I/O pin if one of the I/O banks contains pins that are members of a secured region. You must assign user I/O pins that are adjacent to a signal in a secured region, which belong to a different I/O bank than the secured signal, to GND in the Quartus II software. For example, in Figure 6–20, pin E4 is assigned a signal from a secured region, and I/O banks 1 and 7 belong to different LogicLock regions. Pins D4 and D5 are assigned to GND to ensure that no signal adjacencies exist between the I/O banks.

As a rule, you must assign all unused I/O pins to GND in the Quartus II software and to a ground plane on the PCB. By default, the Quartus II software assigns unused pins to GND. You can configure this option in the Unused Pins page of the Device and Pin Options dialog box.

If you must relax a particular I/O restriction for specific signals to meet your design requirements, you may use the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment. The Quartus II software uses the assignment to bypass normal I/O pin checks for a specific signal. For example, you can apply this assignment to a clock pin assigned to one of the dedicated clock inputs.

Apply the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment in the Quartus Settings File (.qsf) located in the project directory of the active design. Each project revision contains a single .qsf.

To disable the I/O signal rule check for the specified pin name in the .qsf, add the assignment line:

```
set_instance_assignment -name LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT ON -to <pin_name>
```

For more information about the pinouts and pin adjacencies for Cyclone III LS devices, refer to the Cyclone III Device Pin-Out tables. For more information and guidance about I/O assignments, refer to the Cyclone III Device Family Pin Connection Guidelines for Cyclone III LS devices and the I/O Management chapter in volume 2 of the Quartus II Handbook.

### Making Post Compilation Edits

Engineering Change Orders (ECOs) and the rapid recompile feature make incremental changes to routing in a post-fit netlist. ECOs are small changes made to the functionality of a design after the design has been fully compiled. A design is fully compiled when synthesis and placement and routing are completed.

The design separation flow supports any ECOs that do not affect routing, such as changing the LUT mask on an ALM. The design separation flow does not permit ECOs that affect routing or make incremental changes to the routing in a post-fit netlist.

For more information about Rapid Recompile option in the Quartus II software, refer to Incremental Compilation Page (Settings Dialog Box) in Quartus II Help.
Routing Restrictions

During the overall planning of your design, you must be aware of specific design separation flow routing restrictions, especially during the floorplanning stages. This section discusses these routing restrictions.

Column and row interconnect routing resources on Cyclone III LS devices are staggered, with a group of routing elements that starts at each LAB location. The LAB location in which the wire starts drives each routing element. The routing element can reach any LAB destination along the length of the routing element. Figure 6–21 shows a set of staggered R4 interconnects.

Figure 6–21. Staggered R4 Interconnects
The Fitter disables routing wires near the edge of a secured region, in which routing is confined in the region. Figure 6–22 shows the Chip Planner displaying used routing elements in a design with secured regions, using options in the Layer Settings dialog box and using the background color map I/O banks, with only the Global Routing and Used Resources options turned on.

Figure 6–22. Chip Planner View of Used Resources

Figure 6–22 shows that no routing resources reach outside of LogicLock region boundaries, except for global routing signals and signals through interface regions.

Long wires are often unusable in secured regions because their length extends beyond the border of the region. If a secured region abuts the device boundary, you can often attain an increase in routability, because you can use all the routing interconnects that start inside the region and drive toward the edge of your device.

I/O pads along the top and bottom of the device can only use column interconnects to drive into the device fabric. The shortest routing element from the I/O to core logic is a C4 routing wire. I/O pads on the left and right sides of the device can use both C4 and R4 routing elements to reach their LAB destinations. Because the Quartus II software restricts column I/Os from using C4 interconnects going into your device, the Quartus II software creates a four-LAB fence around secured regions when the boundary of the secured region is in four-LABs of the top and bottom I/O pads.
Secured regions should be sized at a minimum of 8 × 8 LABs. If a region is smaller than 8 × 8 LABs, a connection between two LABs that violates the secured region boundary may occur. For example, in Figure 6–23, any elements along the middle axis of the 7 × 7 LAB array cannot use any C4 or R4 routing elements, because a C4 routing element would reach outside the secured region.

**Figure 6–23. 7x7 LAB Array**

---

### Number of Signals in Routing Interfaces

In Cyclone III LS devices, every LAB location has 68 routing elements (R4) driving horizontally in each direction and 48 routing elements (C4) driving vertically in each direction. An individual LAB can directly drive 17 connections in the horizontal direction and 12 in the vertical direction. To guarantee routability, Altera recommends that you have a routing interface height of at least one LAB for every 17 signals routing either left or right, and a routing interface width of one LAB for every 12 signals routing either up or down.

Figure 6–24, Table 6–2, and Table 6–3 illustrate this concept. Figure 6–24 shows three secured regions with two security routing regions; one routing signals horizontally and the other routing signals vertically. Table 6–2 and Table 6–3 list the maximum and the recommended number of signals crossing each security region.
In Figure 6–24, $H_{AB}$ is both the smaller of the height of the region and the height of the routing interface. The minimum $W_{AB}$ is one LAB. $W_{BC}$ is both the smaller width of the region and the width of the routing interface. The minimum $H_{BC}$ is one LAB. Changing $W_{AB}$ or $H_{BC}$ does not affect the values in Table 6–3.

**Figure 6–24. Signals Crossing a Routing Interface**

![Routing Interface Diagram]

**Table 6–2. Maximum Number of Signals Assigned in an Interface Region**

<table>
<thead>
<tr>
<th>To</th>
<th>From</th>
<th>A</th>
<th>B</th>
<th>C</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>—</td>
<td>68 $\times$ $H_{AB}$</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>B</td>
<td>68 $\times$ $H_{AB}$</td>
<td>—</td>
<td>—</td>
<td>48 $\times$ $W_{BC}$</td>
</tr>
<tr>
<td>C</td>
<td>—</td>
<td>48 $\times$ $W_{BC}$</td>
<td>—</td>
<td>—</td>
</tr>
</tbody>
</table>

**Table 6–3. Recommended Number of Signals to Ensure Routability**

<table>
<thead>
<tr>
<th>To</th>
<th>From</th>
<th>A</th>
<th>B</th>
<th>C</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>—</td>
<td>17 $\times$ $H_{AB}$</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>B</td>
<td>17 $\times$ $H_{AB}$</td>
<td>—</td>
<td>—</td>
<td>12 $\times$ $W_{BC}$</td>
</tr>
<tr>
<td>C</td>
<td>—</td>
<td>12 $\times$ $W_{BC}$</td>
<td>—</td>
<td>—</td>
</tr>
</tbody>
</table>

As a general guideline, keep the security routing interface channel width between the two connecting secured regions as short as possible and the depth of the channel as wide as possible. The channel width is the number of LABs that a security routing interface abuts and the depth of the channel is the number of LABs a signal passes as it goes through the routing channel.
In Figure 6–25, an optimal security interface for routing AB would have a channel width equal to the height of secured region A (H\textsubscript{AB}) and a channel depth of one LAB (W\textsubscript{AB}). Having a wide channel with a short depth increases the number of routing resources available between two secured regions.

You can use the Routing Congestion task in the Chip Planner for a visual representation of the routing utilization between secured regions. The Routing Congestion task filters routing resources by type. Utilization of each routing resource type is highlighted on a color gradient over the range that you specify. This tool can help you adjust region sizes and security routing interface channel widths to help you achieve an optimal floorplan. Figure 6–25 shows a design with the Routing Congestion task in the Chip Planner and R24 routing utilization.

*Figure 6–25. Routing Congestion*
Application Example: Modifying a Fitter-Generated Floorplan for the Design Separation Flow

In this application example, the design contains five partitions that you must pack into secured regions. Figure 6–26 shows a block diagram of the design, the entities of the design, and the connectivity between the five secured partitions.

Figure 6–26. Connectivity between Five Secured Partitions

The following steps outline a recommended design flow for creating a floorplan for this design:

1. Create a LogicLock region for each partition that you must pack into a secured region.

2. Set each LogicLock region with the following settings:
   - **Size** set to *Auto*,
   - **State** set to *Floating*,
   - **Reserved** set to *On*, and
   - **Security Attributes** set to *Unsecured*.

   Running an initial placement with these settings generates non-overlapping LogicLock regions that can be used as an initial floorplan.

3. On the Processing menu, point to **Start** and click **Start Early Timing Estimate** to run an initial placement and routing. The initial placement and routing approximates the size of each region and the general placement of the LogicLock regions relative to other LogicLock regions to achieve timing closure. Figure 6–27 shows the floorplan that the early timing estimate generates.
4. In the LogicLock Regions window, select the LogicLock regions, right-click, and then click **Set Size and Origin to Previous Fitter Results**.

5. Use the Design Partition Planner to view the connectivity between the different regions. You can experiment with the relative placement of the blocks by dragging and dropping each design partition. The wire bundles between design partitions help you to determine a placement that has non-overlapping routing channels.

   You must also consider the connectivity to the I/O banks when arranging your floorplan. You can toggle the display of the connections between the partitions and the I/O banks in the Design Partition Planner to help you properly allocate I/O resources and to avoid conflicts between I/O connections and inter-partition signals. To display routing between partitions and the I/O banks, turn on **Display connections to I/O banks** in the **Bundle Configuration** dialog box.

6. Set each LogicLock region to the necessary security attribute.

7. In the Chip Planner, adjust the size and placement of each LogicLock region using the relative placement you created with the Design Partition Planner. Altera recommends the following considerations when modifying your floorplan:

   - The floorplan must be complete. If unsecured logic that is non-contiguous due to the placement of a secured region is present, use an empty reserved LogicLock region to prevent a non-routable placement.

   - Each secured region must be a minimum of 8 × 8 LABs.

   - Each region that has I/O pins added as members of the LogicLock region should overlap the I/O bank to which it is connected. You can use the I/O bank background color map to visualize the boundaries between the I/O banks (Figure 6–28).

   - A secured region must not cover all global resources that unsecured logic require (such as clock pins and PLLs).
8. Create security routing interfaces between each of the secured regions. Assign all signals entering or exiting a region to a security routing interface.

Figure 6–29 shows the final floorplan result for this application example.

Figure 6–29. Final Floorplan
Report Panels

After the Fitter successfully places and routes your design with secured regions, the Quartus II software generates security reports. Use the security reports to review the secured regions, their associated routing interfaces, all inputs and outputs from each secured region, and the I/O bank usage for each secured region. You can locate the security reports in the Fitter section of the Compilation reports.

Secured LogicLock Region Summary

This report provides a summary of all secured regions in your design. Table 6–4 describes each column in the Secured LogicLock Region Summary report.

Table 6–4. Secured LogicLock Region Report

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Secured LogicLock Region</td>
<td>Lists all secured LogicLock regions in the design.</td>
</tr>
<tr>
<td>Security Attribute</td>
<td>Lists the security attribute (unsecured, 1, 2, or security routing interface) of the LogicLock region.</td>
</tr>
<tr>
<td>Partition Assigned</td>
<td>Lists the design partition assigned to the secured region.</td>
</tr>
<tr>
<td>Number of Input Signals</td>
<td>Lists the number of inputs and fan-outs into a region. The input counts the number of unique drivers that feed a secured region. The fan-out counts the total number of unique destinations being fed by the input signals into the secured region. Figure 6–30 shows input signals and fan-outs to a region.</td>
</tr>
<tr>
<td>Number of Output Signals</td>
<td>Lists the number of outputs and fan-outs from a region. The output counts the number of unique drivers sourcing a signal from the secured region. The fan-out counts the total number of unique destinations fed by the output signal.</td>
</tr>
</tbody>
</table>

Figure 6–30. Input Signals and Fan-Outs to a Region

Security Routing Interfaces

This report summarizes the security routing interfaces. Table 6–5 describes each column in the Security Routing Interfaces report.
### Table 6–5. Security Routing Interface Report

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Name</td>
<td>Lists all security routing interfaces in the design.</td>
</tr>
<tr>
<td>Abutting Region A</td>
<td>First region that the security routing interface abuts (touches the border of the secured region).</td>
</tr>
<tr>
<td>Abutting Region B</td>
<td>Second region that the security routing interface abuts (touches the border of the secured region).</td>
</tr>
<tr>
<td>Number of Signals A to B</td>
<td>Lists the number of signal connections between region A and region B. The counts are shown as signals and fan-outs. Signals list the number of unique drivers from region A. Fan-out lists the number of unique destinations in region B that are fed by region A.</td>
</tr>
<tr>
<td>(Total Fanout in B)</td>
<td></td>
</tr>
<tr>
<td>Number of Signals B to A</td>
<td>Lists the number of signal connections between region B and region A. The counts are shown as signals and fan-outs. Signals list the number of unique drivers from region B. Fan-out lists the number of unique destinations in region A that are fed by region B.</td>
</tr>
<tr>
<td>(Total Fanout in A)</td>
<td></td>
</tr>
</tbody>
</table>
Secured LogicLock Region Inputs and Outputs

This set of reports provides a detailed list of every signal that enters or exits a secured region. There is one report per secured region.

Security I/O Bank Usage

This report displays the secured LogicLock region associated with each I/O bank, lists the number of pins within each region, and lists the number of pins in use. Table 6–6 describes each column in the Secured LogicLock Region Inputs and Outputs report.

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>I/O Bank</td>
<td>Lists all available I/O banks on the device.</td>
</tr>
<tr>
<td>Associated Region</td>
<td>An I/O bank becomes associated with a secured LogicLock region if any portion of the I/O bank is covered by the region. If no secured region covers an I/O bank, “Unsecured Logic” is displayed, and all pins of the I/O bank are available for unsecured use.</td>
</tr>
<tr>
<td>Pin Locations Used / Pin Locations Covered by Region</td>
<td>Displays the ratio of pins with a signal assignment in the I/O bank to the number of possible I/O pin assignments.</td>
</tr>
</tbody>
</table>

Quartus Settings File Syntax

This section contains the syntax description for each Quartus Settings File (.qsf) assignment in the design separation flow.

**LL_SECURITY_ROUTING_INTERFACE**

This command changes a LogicLock region assignment to a security routing interface.

Type: Boolean; (ON/OFF—Defaults to OFF)

Syntax:

```
set_global_assignment -name LL_SECURITY_ROUTING_INTERFACE <value> \ -section_id <section_identifier> LL_REGION_SECURITY_LEVEL
```

**LL_REGION_SECURITY_LEVEL**

This command identifies the security level of a LogicLock region.

Type: Enumeration—Defaults to UNSECURED

- 1
- 2
- UNSECURED

Syntax:

```
set_global_assignment -name LL_REGION_SECURITY_LEVEL <value> \ -section_id <section_identifier>
```
LL_MEMBER_OF_SECURITY_ROUTING_INTERFACE

This command assigns an I/O pin from a secured region to a security routing interface. `<value>` and `<section_id>` denote the name of the routing interface region. `<to>` specifies the name of the signal.

Type: String

Syntax:

```
set_instance_assignment -name \ LL_MEMBER_OF_SECURITY_ROUTING_INTERFACE
<value> -to <to> \n-section_id <section_id>
```

LL_SIGNAL_SECURITY_LEVEL

This command sets the security level of a signal. The default value is the security level of the region that generates the signal. This assignment may be used only to lower a security level.

Type: Enumeration

- UNSECURED
- 1
- 2

Syntax:

```
set_instance_assignment -name LL_SIGNAL_SECURITY_LEVEL <value> \n-to <to> -section_id <section_id>
```

Document Revision History

Table 6–7 lists the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>• Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>• Updated Table 6–1 on page 6–8.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Updated “Using Secured Regions” on page 6–9 and “Understanding Fencing Regions” on page 6–11.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• General editorial update.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Template update.</td>
</tr>
</tbody>
</table>
Table 6–7. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
</table>
| December 2010 | 10.1.0  | ■ Modified the former “Avoiding Child Partitions” section into the new “Avoiding Multiple Design Partitions With a Secured Region” on page 4–6 section and added information about multi-hierarchy partitions.  
■ Updated the “Using Secured Regions” on page 4–9 section.  
■ Updated the “Making Design Separation Flow Location Assignments in the Chip Planner” on page 4–10 section.  
■ Updated the “Creating a Complete Floorplan” on page 4–14.  
■ Updated the “Working with Global Signals” on page 4–21 and “Assigning I/O Pins” on page 6–25 sections with information about the LL_IGNORE_IO_PIN_SECURITY_CONSTRAINT assignment.  
■ Added the “Making Post Compilation Edits” on page 4–26.  
■ Updated the “Number of Signals in Routing Interfaces” on page 4–29.  
■ Added feature licensing information.  
■ Updated figures and overall editorial update.  
■ Template update. |
| July 2010   | 10.0.0  | Initial release. Content originated from AN 567: Quartus II Design Separation Flow.                                                                                                                  |
Section 2. System Design with Qsys

This section provides information about Qsys. Qsys is a powerful system integration tool which is included as part of the Quartus II software. Qsys automates the task of capturing of integrating customized HDL components, which may include IP cores, verification IP, and other design modules. You can use Qsys to integrate your own components with the components that Altera® or third-party developers provide. In some cases, you can implement an entire design using components from the Qsys component library.

This section includes the following chapters:

- **Chapter 7, Creating a System With Qsys**
  This chapter provides an overview of the Qsys system integration tool, including an introduction to hierarchical system design.

- **Chapter 8, Creating Qsys Components**
  This chapter introduces Qsys components and the Qsys component library. It also provides an overview of the Qsys component editor which you can use to define custom components.

- **Chapter 9, Qsys Interconnect**
  This chapter discusses the Qsys interconnect, a high-bandwidth structure for connecting components that use Avalon® interfaces.

- **Chapter 10, Optimizing Qsys System Performance**
  This chapter provides information on optimizing system performance with the Qsys system integration tool. Following the design practices recommended in this chapter can improve the maximum clock frequency, concurrency and throughput, logic utilization, or even power utilization of your system.

- **Chapter 11, Component Interface Tcl Reference**
  This chapter describes an alternative method for defining Qsys components by declaring their properties and behaviors in a Hardware Component Description File (_hw.tcl). It also provides a reference for the Tool Command Language (Tcl) commands that describe Qsys components.

- **Chapter 12, Qsys System Design Components**
  This chapter describes the structure of Qsys components, with an emphasis on using the Qsys Component Editor to create the Hardware Component Description File (_hw.tcl), which describe and package components that you can use in a Qsys system.
7. Creating a System With Qsys

Qsys is a system integration tool included as part of the Quartus® II software. Qsys captures system-level hardware designs at a high level of abstraction and automates the task of defining and integrating customized HDL components, which may include IP cores, verification IP, and other design modules. Qsys facilitates design reuse by packaging and making available your custom components and systems, and integrates your custom components with Altera® and third-party developer components.

Qsys automatically creates interconnect logic from the connectivity options you specify, eliminating the error-prone and time-consuming task of writing HDL to specify the system-level connections.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

Qsys provides the following advantages for system design:

- Automates the process of customizing and integrating components
- Supports 64-bit addressing
- Supports modular system design
- Supports visualization of systems
- Supports optimization of interconnect and pipelining within the system
- Provides full integration with the Quartus II software

For descriptions of unique or exceptional AXI and APB support in the Qsys software, refer to the Qsys Interconnect chapter in volume 1 of the Quartus II Handbook. For more information about Avalon, AXI, and APB interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website.

Qsys 64-Bit Addressing Support

Qsys interconnect supports up to 64-bit addressing for all Qsys interfaces and components, with a range of 0x0000 0000 0000 0000 to 0xFFFF FFFF FFFF FFFF, inclusive.

In Qsys, address parameters appear in the Base and End columns on the System Contents tab, on the Address Map tab, in the parameter editor, and in validation messages. The Qsys GUI displays as many digits as needed in order to display the top-most set bit, for example, 12 hex digits for a 48-bit address.
A Qsys system can have multiple 64-bit masters, with each master establishing its own address space. Slaves may be shared among masters and masters can map slaves in different ways; for example, one master may interact with slave 0 at base address 0000_0000_0000, and another master may see the same slave at base address c000_000_000.

64-bit are also supported for narrow-to-wide and wide-to-narrow transactions across Avalon and AXI interfaces, though bursts that exceed 32-bits are legal only within the Avalon interface. AXI3 allows bursts of 1 - 16 transfers. AXI4 allows burst lengths of 256.

Quartus II debug tools that provide access to the state of an addressable system via the Avalon-MM interconnect are also 64-bit compatible and process within a 64-bit address space, including a JTAG to Avalon master bridge.

### Ports and Bridges

You can configure address ports within memory-mapped interfaces to be 64-bits wide. When a component’s master port is not 64-bit capable, you can use the window bridge component (address span extender) to enable it to access a specific 32-bit segment of a 64-bit address map. The address span extender enables a 32-bit master to access a windowed portion of a larger memory map. The slave interface has an address port size corresponding to the address window.

For more information about the Address Span Extender feature, refer to the [Qsys System Design Components](#) chapter in volume 1 of the Quartus II Handbook.

### DMA Controllers

DMA controllers are limited to 32-bit addressing. As a workaround, you can use the window bridge component, as described in “Ports and Bridges” above.

### Qsys Interface Support

Qsys interconnect connects the following interface types:

- **Memory-Mapped**—Implements a partial crossbar interconnect structure (Avalon-MM, AXI, and APB) that provides concurrent paths between master and slaves. Interconnect consists of synchronous logic and routing resources inside the FPGA, and implementation is based on a network-on-chip architecture.

- **Streaming**—Connects Avalon Streaming (Avalon-ST) sources and sinks that stream unidirectional data, as well as high-bandwidth, low-latency components. Streaming creates datapaths for unidirectional traffic including multichannel streams, packets, and DSP data. The Avalon-ST interconnect is flexible and can be used to implement on-chip interfaces for industry standard telecommunications and data communications cores, such as Ethernet, Interlaken, and video. In all cases, you can define bus widths, packets, and error conditions.
Interrupts—Connects interrupt senders and the interrupt receivers of the component that serves them. In systems with interrupt request sender (IRQ) interfaces, Qsys interconnect includes several components to implement interrupt processing. Qsys processes individual, single-bit interrupt requests (IRQs). In the event that multiple senders assert their IRQs simultaneously, the receiver logic (typically under software control) determines which IRQ has highest priority, then responds appropriately.

Clocks—Connects clock sources with clock input interfaces.

Resets—Connects reset sources with reset input interfaces. If your system requires a particular positive-edge or negative-edge synchronized reset, Qsys inserts a reset controller to create the appropriate reset signal. If you design a system with multiple reset inputs, the Reset Controller ORs all reset inputs and generate a single reset output. A Reset Bridge allows you to use a reset signal in two or more subsystems of your Qsys system.

Conduits—Connects point-to-point conduit interfaces. Conduit interfaces are brought to the top level of the system as additional ports, and are always point-to-point connections. Exported signals are usually either application-specific signals or the interface signals. Application-specific signals are exported to the top level of the system by the conduit interface(s) defined in a component. These are I/O signals in a component’s HDL logic that are not part of any Avalon interfaces and connect to an external device, for example DDR SDRAM memory, or logic defined outside of the Qsys system.
**Understanding the Qsys Design Flow**

Figure 7–1 illustrates a bottom-up design flow example in Qsys.

**Figure 7–1. Complete Qsys Design Flow**

In the alternative top-down design flow, you begin by designing the Qsys system, and then define and instantiate custom Qsys components. The top-down design flow clarifies the system requirements earlier in the design process.

Designs targeting HardCopy series devices require specific design constraints. Consequently, if you are targeting a HardCopy series device, you must verify your design for the HardCopy companion device.
Follow these guidelines to verify your design for HardCopy devices:

1. In the Device dialog box in the Quartus II software, select both the FPGA and the appropriate HardCopy companion device.

2. In Step 8 of the design flow shown in Figure 7–1 compile for both the FPGA and HardCopy device.

3. After Step 10 of the design flow shown in Figure 7–1, if the FPGA passes all functional simulation and hardware verification tests, generate the HardCopy handoff archive file and send this archive file to the HardCopy Design Center for the backend flow implementation.

For more information about designing for HardCopy devices, refer to About Designing HardCopy Devices in Quartus II Help.

Searching for Component Files to Add to the Component Library

The component library includes the design elements that you use in your Qsys systems. Components can include Altera-provided IP cores, third-party IP cores, and custom IP cores that you provide. Previously created Qsys systems can also appear in the component library, and you can use these systems in other designs if they have exported interfaces.

Altera and third-party developers provide ready-to-use components, which are installed automatically with the Quartus II software and are available in the Qsys component library. The Qsys component library includes the following components:

- Microprocessors, such as the Nios® II processor
- DSP IP cores, such as the Reed Solomon II core
- Interface protocols, such as the IP Compiler for PCI Express
- Memory controllers, such as the RLDRAM II Controller with UniPHY
- Avalon® Streaming (Avalon-ST) components, such as the Avalon-ST Multiplexer IP core
- Qsys Interconnect components
- Verification IP (VIP) Bus Functional Models (BFMs)

You can set the IP Search Path option to specify custom and third-party components that you want to appear in the component library. Qsys searches for component files each time you open the tool, and locates and displays the list of available components in the component library.

Qsys searches the directories listed in the IP Search Path for the following component file types:

- Hardware Component Description Files (_hw.tcl) files. Each _hw.tcl file defines a single component.
- IP Index (.ipx) files. Each file indexes a collection of available components, or a reference to other directories to search. In general, .ipx files facilitate faster startup for Qsys and other tools because fewer directories need to be searched and analyzed.
Qsys searches some directories recursively and other directories only to a specific depth. When a directory is recursively searched, the search stops at any directory containing a _hw.tcl or .ipx file; subdirectories are not searched. In the following list of search locations, a recursive descent is annotated by **. The * signifies any file.

- PROJECT_DIR/*
- PROJECT_DIR/ip/**/*
- QUARTUS_INSTALLDIR/../ip/**/*

Complete the following steps to extend the default search path by specifying additional directories:

1. On the Tools menu, click Options.
2. In the Category list, click IP Search Path.
3. Click Add.
4. Browse to locate additional directories and click Open to add them to your search path.

You do not need to include the components specified in the IP Search Path as part of your Quartus II project.

**Adding Components to the Component Library**

Use one of the following methods to add components to the component library:

**Copy Components to the Install Directory**

The simplest method to add a new component to the Qsys Component Library is to copy your components into the default <install_dir>/ip directory provided by Altera. This approach is useful in the following situations:

- You want to associate your components with a specific release of the Quartus II software.
- You want to have the same components available across multiple projects.
Figure 7–2 illustrates this approach.

Figure 7–2. User Library Included In Subdirectory of <install_dir>/ip/

In Figure 7–2, the circled numbers identify a typical directory structure for the Quartus II software. For the directory structure above, Qsys performs the component discovery algorithm described below to locate .ipx and _hw.tcl files and initiate the component library:

1. Qsys recursively searches the <install_dir>/ip/ directory by default. The recursive search stops when Qsys finds an .ipx file.
2. As part of the recursive search, Qsys also looks in the user_components directory because this directory path appears as an IP Search Path in the Options dialog box. Qsys finds the component1 directory, which contains component1_hw.tcl. When Qsys finds that component, the recursive search ends, and no components in subdirectories of component1 are found.
3. Qsys then searches the component2 directory, because this directory path also appears as an IP Search Path, and discovers component2_hw.tcl. When Qsys finds component2_hw.tcl, the recursive search ends.

If you save your _hw.tcl file in the <install_dir>/ip/ directory, Qsys finds your _hw.tcl file and stops. Qsys does not conduct the component discovery algorithm just described.

Reference Components in an .ipx File

You can specify the search path in a user_components.ipx file under the <install_dir>/ip directory. This method allows you to store components in a location that is not linked to a specific Quartus II installation, and to add a location that is independent of the default search path. You can also save the .ipx file in any of the default search locations, for example, the Quartus II project directory, or the ip directory in the project directory.
The user_components.ipx file includes a single line of code redirecting Qsys to the location of each user library. Example 7–1 shows the redirection code.

Example 7–1. Redirect to User Library

```xml
<library>
    <path path="<user_lib_dir>/user_ip/**/**"/>
</library>
```

You can verify that components are available with the ip-catalog command. You can use the ip-make-ipx command to create an .ipx file for a directory tree, which can reduce the startup time for Qsys. The following sections describe these commands.

**ip-catalog**

This command displays the catalog of available components relative to the current project directory in either plain text or XML format.

**Usage**

```
```

**Options**

- **--project-dir=**<directory>—Optional. Components are found in locations relative to the project, if any. By default, the current directory, ‘.’ is used. To exclude any project directory, leave the value empty.
- **--name=**<value>—Optional. This argument provides a pattern to filter the names of the components found. To show all components, use a * or ‘’. By default, all components are shown. The argument is not case sensitive.
- **--verbose**—Optional. If set, reports the progress of the command.
- **--xml**—Optional. If set, generates the output in XML format instead of a line-and colon-delimited format.
- **--help**—Shows Help for the ip-catalog command.

**ip-make-ipx**

This command creates an ip-make-ipx (.ipx) file and is a convenient way to include a collection of components from an arbitrary directory in the Qsys search path. You can also edit the .ipx file to disable visibility of one or more components in the Qsys component library.

**Usage**

```
```
Options

- **--source-directory=<directory>**—Optional. Specifies the root directory(ies) that Qsys uses to find the component files. The default directory is ".". You can also provide a comma-separated list of directories.

- **--output=<file>**—Optional. Specifies the name of the file to generate. The default name is `/components.ipx`.

- **--relative-vars=<value>**—Optional. Causes the output file to include references relative to the specified variable or variables where possible. You can specify multiple variables as a comma-separated list.

- **--thorough-descent**—Optional. If set, a component or .ipx file in a directory does not prevent subdirectories from being searched.

- **--message-before=<value>**—Optional. A message to print to `stdout` when indexing begins.

- **--message-after=<value>**—Optional. A message to send to `stdout` when indexing completes.

- **--help**—Show Help for this command.

Understanding IPX File Syntax

An `.ipx` file is an XML file that describes the search path used to discover components that are available for a Qsys system. A `<path>` entry specifies a directory in which components may be found. A `<component>` entry specifies the path to a single component. Example 7–2 illustrates this format.

**Example 7–2. .ipx File Structure**

```xml
<library>
    <path path="...<user directory>" />
    <path path="...<user directory>" />
    ...
    <component ... file="...<user directory>" />
    ...
</library>
```

A `<path>` element contains a path attribute, which specifies the path to a directory, or the path to another .ipx file, and can use wildcards in its definition. An asterisk matches any file name. If you use an asterisk as a directory name, it matches any number of subdirectories.

When searching the specified path, the following three types of files are identified:

- `.ipx`—additional index files
- `_hw.tcl`—Qsys component definitions
- `_sw.tcl`—Nios II board support package (BSP) software component definitions
A `<component>` element contains several attributes to define a component. If you provide the required details for each component in an .ipx file, the startup time for Qsys is less than if Qsys must discover the files in a directory. Example 7–3 shows two `<component>` elements. Note that the paths for file names are specified relative to the .ipx file.

**Example 7–3. Component Elements**

```xml
<library>
  <component
    name="A Qsys Component"
    displayName="Qsys FIR Filter Component"
    version="2.1"
    file="/components/qsys_filters/fir_hw.tcl"
  />
  <component
    name="rgb2cmyk_component"
    displayName="RGB2CMYK Converter(Color Conversion Category!)
    version="0.9"
    file="/components/qsys_converters/color/rgb2cmyk_hw.tcl"
  />
</library>
```

**Integrating Third-Party Components**

You can use Qsys components created by third-party IP developers. Altera awards the Qsys Compliant label to IP cores that are fully supported in Qsys. These cores support Avalon AXI interfaces and may include timing and placement constraints, software drivers, simulation models, and reference designs.

To find supported third-party Qsys components, on the Intellectual Property & Reference Designs web page, type Qsys Certified in the Search box, select IP Core & Reference Designs, and then press Enter.

**Creating a Qsys System**

You can create a Qsys system in the Quartus II software by clicking Qsys System File in the New dialog box, or opening Qsys from the Tools menu. To open a previously created Qsys design, click Open on the File menu in the Quartus II software window, or the Qsys window.

For more information about the Qsys GUI, refer to About Qsys in Quartus II Help.

Qsys is more powerful if you design your custom components using standard interfaces. By using standard interfaces, your components inter-operate with the components in the Qsys component library. In addition, you can take advantage of bus functional models (BFMs), monitors, and other verification IP to verify your design.

**Adding System Contents**

The Component Library tab displays the components that you add to your system.
Adding Components

To add a component to your system, select the component in the Component Library, and then click Add. A parameter editor appears allowing you to configure the component instance.

You can type some or all of the component’s name in the Component Library search box to help locate a particular component type. For example, you can type memory to locate memory-mapped components, or axi to locate AXI interconnect components.

Working With Presets for Supported IP Components

When you add a component to your system, the Qsys Presets Editor opens for IP components whose parameters you are allowed to modify and lists presets that you can apply to your component, depending on the design protocol. When you apply a preset to a component, the parameters with specific required values for the protocol are automatically set for you.

You can search for text to filter the Presets list. For example, if you select the DDR3 SDRAM Controller with UniPHY component, and then type 1g micron 256, the Presets list shows only those presets that apply to the 1g micron 256 protocol. Presets whose parameter values match the current parameter settings are shown in bold.

Selecting a preset does not prevent you from changing any parameter to meet the requirements of your design. Clicking Update allows you to update parameter values for a custom preset. The Update Preset dialog box displays the default value, which you can edit, and the current value, which is static.

You can also create your own preset by clicking New. When you create a preset, you specify a name, description and the list of parameters whose values are set by the preset.

You can remove a preset from the Quartus II project directory by clicking Delete.

For more information about presets, refer to Presets Editor in Quartus II Help.

Connecting Components

When you add connections to a Qsys system, you connect the interfaces of the modules in the system. The individual signals in each interface are connected by the Qsys interconnect when the HDL for the system is generated. You connect interfaces of compatible types and opposite directions. For example, you can connect a memory-mapped master interface to a slave interface, and an Interrupt sender interface to an Interrupt receiver interface.
To view possible connections for an interface on the System Contents by hovering your pointer in the Connections column. In this view, open circles represent possible connections, and filled circles indicate connections that you have made. To make a connection, click the open circle at the intersection of the two interface names. Clicking a filled-in circle removes the connection. Figure 7–3 illustrates the connections display.

**Figure 7–3. Connections Column**

For more information about connecting components, refer to *Connecting Qsys Components* in Quartus II Help.

**Filtering Components**

You can use the Filters dialog box to filter the display of your system in the System Contents tab. You can filter the display of your system by interface type, instance name, or by using custom tags. For example, you can use filtering to view only instances that include memory-mapped interfaces, instances that are connected to a particular Nios II processor, or to temporarily hide clock and reset interfaces to simplify the display.

For more information about filtering components, refer to the *Filters Dialog Box* in Quartus II Help.

**Using the System Inspector**

The System Inspector tab displays the underlying model of your complete system, and provides comprehensive details about your system such as the following information:

- The connections between signals
- The names of signals included in exported interfaces
The internal connections of Qsys subsystems that are included as components

- In contrast, the **System Contents** tab displays only the exported interfaces of Qsys subsystems included as components.

The global parameter settings that you specified on the **Project Settings** tab

You can use the **System Inspector** tab to review and change component parameters and to review interface timing. For example, Figure 7–4 shows the timing for the Avalon-MM DMA write master for the PCI Express-to-Ethernet system illustrated in Figure 7–12 on page 7–30.

**Figure 7–4. Avalon-MM Write Master Timing Waveforms Available on the Project Settings Tab**

To display the timing for an interface, expand the component, and then click the interface name.

**Defining the Address Map**

The **Address Map** tab provides a table including all the memory-mapped slaves in your design and the address range that each connected memory-mapped master uses to address that slave. The table shows the slaves on the left and masters across the top, with the address span of the connection shown in each cell. A blank cell implies that there is no connection between that master and slave.

Follow these steps to change or create a connection between master and slave components:

1. In Qsys, click the **Address Map** tab.
2. Locate the table cell that represents the connection between the master and slave component pair.
3. Either type in a base address or update the current base address in the cell.
You can design a system where two masters access a slave at different addresses. If you use this feature, the **Base** and **End** address columns of the **System Contents** tab are labeled *mixed* rather than providing the address range.

### Specifying Clock Settings

The **Clock Settings** tab defines the **Name**, **Source**, and frequency (MHz) of each clock in your system. Clicking the **Add** button adds a new clock.

For more information, refer to the *Adding Components to a Qsys System* in Quartus II Help.

### Specifying Project Settings

The **Project Settings** tab allows you to view and change the properties of your Qsys system. Table 7–1 describes system-level parameters available on the **Project Settings** tab.

<table>
<thead>
<tr>
<th>Parameter Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Device Family</strong></td>
<td>Specifies the Altera device family.</td>
</tr>
<tr>
<td><strong>Device</strong></td>
<td>Specifies the target device for the selected device family.</td>
</tr>
<tr>
<td><strong>Clock Crossing Adapter Type</strong></td>
<td>Specifies the default implementation for automatically inserted clock crossing adapters. The following choices are available:</td>
</tr>
<tr>
<td></td>
<td>■ <strong>Handshake</strong>—This adapter uses a simple hand-shaking protocol to propagate transfer control signals and responses across the clock boundary. This methodology uses fewer hardware resources because each transfer is safely propagated to the target domain before the next transfer can begin. The <strong>Handshake</strong> adapter is appropriate for systems with low throughput requirements.</td>
</tr>
<tr>
<td></td>
<td>■ <strong>FIFO</strong>—This adapter uses dual-clock FIFOs for synchronization. The latency of the FIFO-based adapter is a couple of clock cycles more than the handshaking clock crossing component, but the FIFO-based adapter can sustain higher throughput because it supports multiple transactions at any given time. The FIFO-based clock crossers require more resources. The <strong>FIFO</strong> adapter is appropriate for memory-mapped transfers requiring high throughput across clock domains.</td>
</tr>
<tr>
<td></td>
<td>■ <strong>Auto</strong>—If you select <strong>Auto</strong>, Qsys specifies the <strong>FIFO</strong> adapter for bursting links, and the <strong>Handshake</strong> adapter for all other links.</td>
</tr>
<tr>
<td><strong>Limit interconnect pipeline stages to</strong></td>
<td>Specifies the maximum number of pipeline stages that Qsys may insert in each command and response path to increase the f(_{\text{MAX}}) at the expense of additional latency. You can specify between 0–4 pipeline stages, where 0 means that the interconnect has a combinational data path. Choosing 3 or 4 pipeline stages may significantly increase the logic utilization of the system. This setting is per Qsys system or subsystem, meaning that each subsystem can have a different setting. Note that the additional latency is for both the command and response directions for the two Qsys systems, even if you combine them into a single Quartus II project.</td>
</tr>
<tr>
<td><strong>Generation Id</strong></td>
<td>A unique integer value that is set to a timestamp just before Qsys system generation that Qsys uses to check for software compatibility.</td>
</tr>
</tbody>
</table>
Qsys generates a warning message if the selected device family and target device do not match the Quartus II software project settings. Also, when you open Qsys from within the Quartus II software, the device type in your Qsys project is replaced with the selected device in your open Quartus II software project.

**Defining Qsys Instance Parameters**

The **Instance Parameters** tab allows you to define parameters for a Qsys system. You can use instance parameters to modify a Qsys system when you use the system as a subcomponent in another Qsys system. The higher-level Qsys system can assign values to these instance parameters.

The **Instance Script** on the **Instance Parameters** tab defines how the specified values for the instance parameters should affect your Qsys design subcomponents. The instance script allows you to make queries about the instance parameters you define and set the values of the parameters for the subcomponents in your design.

When you click Preview Instance, Qsys creates a preview of the current Qsys system with the specified parameters and instance script, and shows the parameter editor for the instance. This allows you to see how an instance of this system appears when you use it in another system. The preview instance does not affect your saved system.

For more information, refer to Working with Instance Parameters in Qsys in Quartus II Help.

To use Instance Parameters, the components or subsystems in your Qsys system must have parameters that can be set when they are instantiated in a higher-level system. Many components in the Component Library have parameters that you can set when adding the component to your system. If you create your own IP components, you use the _hw.tcl file to specify which parameters can be set when the component is added to a system. If you create hierarchical Qsys systems, each Qsys system in the hierarchy can include instance parameters to pass parameter values through multiple levels of hierarchy.

For more information on creating your own components and specifying parameters, refer to the Component Interface Tcl Reference chapter in the Quartus II Handbook.

**Creating an Instance Script**

The first command in an instance script must specify the version of the Tcl commands to be used in the script. This command ensures the Tcl commands behave identically in future versions of the tool. Use the following Tcl command to specify the version of the Tcl commands, where <version> is a Quartus II software version number, such as 13.0:

```tcl
package require -exact qsys <version>
```

To use Tcl commands that work with instance parameters in the instance script, you must specify the commands within a Tcl procedure called a composition callback. In the instance script, you specify the name for the composition callback with the following command:

```tcl
set_module_property COMPOSITION_CALLBACK <name of callback procedure>
```
Specify the appropriate Tcl commands inside the Tcl procedure with the following syntax:

```tcl
proc <name of procedure defined in previous command> {} {
    #Tcl commands to query and set parameters go here
}
```

Use Tcl commands in the procedure to query the parameters of a Qsys system, or to set the values of the parameters of the subcomponents instantiated in the system. Table 7–2 describes the supported Tcl commands.

**Example 7–4. Simple Instance Script**

```tcl
# Request a specific version of the scripting API
package require -exact qsys 13.0

# Set the name of the procedure to manipulate parameters:
set_module_property COMPOSITION_CALLBACK compose

proc compose {} {
    #Get the pio_width parameter value from this Qsys system and pass the
    #value to the width parameter of the pio_0 instance
    set_instance_parameter_value pio_0 width [get_parameter_value \pio_width]
}
```

For more information about _hw.tcl syntax and manipulating parameters, refer to the Component Interface Tcl Reference chapter in the Quartus II Handbook.
For another example, refer to “Hierarchical System Using Instance Parameters Example” on page 7–33.

**Viewing the HDL Example**

The HDL Example tab provides the top-level HDL definition of your system in either Verilog HDL or VHDL, and also displays VHDL component declarations. You can copy and paste the example into a top-level HDL file that instantiates the Qsys system, if the system is not the top-level module in your Quartus II project.

**Creating Hierarchical Systems**

Qsys supports team-based and hierarchical system design. You can include any Qsys system that exports an interface as a component in another Qsys system. In a team-based design flow, you can have one or more systems in your design developed simultaneously by other team members, decreasing time-to-market for the complete design.

Figure 7–5 shows the top-level of a Qsys hierarchical design that implements a PCI Express™ to Ethernet bridge. This example combines separate PCI Express and Ethernet subsystems with Altera’s DDR3 SDRAM Controller with UniPHY IP core.

**Figure 7–5. Top-Level for a PCI Express to Ethernet Bridge**

Hierarchical system design in Qsys offers the following advantages:

- Enables team-based, modular design by dividing large designs into subsystems.
- Enables design reuse by allowing you to use any Qsys system as a component.
- Enables scalability by allowing you to instantiate multiple instances of a Qsys system.

For more information about hierarchical design, refer to “PCI Express Subsystem Example” on page 7–29.
Adding Systems to the Component Library

Any Qsys system that exports an interface is available for use in other Qsys systems. Figure 7–6 shows the component library, including the PCI Express and Ethernet subsystems as components in the component library for the PCI Express to Ethernet Bridge example system in Figure 7–15 on page 7–31. To include systems as components in other designs, you can add the system to the component library, or include the directory for the system in component search path for Qsys.

Creating a Component Based on a System

The Export System as hw.tcl Component command on the File menu allows you to save the system currently open in Qsys as an _hw.tcl file in the current working directory. The saved system appears in the System list under Project in the Qsys Component Library.

Because Qsys systems become components in the component library, be careful not to give your system a name that is already in use.

Creating Secure Systems (TrustZones)

TrustZone refers to the security extension of the ARM architecture, which includes the concept of secure and non-secure transactions, and a protocol for processing between the designations. TrustZone security support is a part of the Qsys interconnect.

In Qsys, AXI masters are treated as TrustZone-aware; all other memory-mapped interfaces are set to secure, non-secure or TrustZone-aware (only for AXI slaves with TrustZone support). The default value for non-AXI master interfaces is non-secure. Unless specified, all non-TrustZone-aware components are treated as non-secure, for example, Avalon master and slave components.

Qsys provides compilation-time TrustZone support for non-TrustZone-aware components, for cases such as when an Avalon master needs to communicate with a secure AXI slave. For example, the designer can specify whether the connection point is secure or non-secure at compilation time. You can specify secure address ranges on memory slaves, if a per-interface security setting is not sufficient.
For TrustZone-aware masters, the interconnect uses the master’s `axprot` signal to determine the security status of each transaction.

The table below summarizes secure and non-secure access between master, slave, and memory components in Qsys. Per-access refers to allowing a TrustZone-aware master to allow or disallow a particular access (or transactions).

<table>
<thead>
<tr>
<th>Table 7–3.</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Transaction Type</strong></td>
</tr>
<tr>
<td>TrustZone-aware slave/memory</td>
</tr>
<tr>
<td>Non-TrustZone-aware slave (secure)</td>
</tr>
<tr>
<td>Non-TrustZone-aware slave (non-secure)</td>
</tr>
<tr>
<td>Non-TrustZone-aware memory (secure region)</td>
</tr>
<tr>
<td>Non-TrustZone-aware memory (non-secure region)</td>
</tr>
</tbody>
</table>

If a master issues transactions that fall into the per-access or not allowed cells, as described in the table above, your design must contain a default slave. A transaction that violates security is rerouted to the default slave and subsequently terminated with an error. You can connect any slave as the default that is able to respond to the master that requires a default slave with errors. You can share the default slave between multiple masters. Altera recommends that you have one default slave for each domain. Altera also recommends that you use the `altera_axi_default_slave` component as the default slave because this component has the required TrustZone features.

In Qsys, you can achieve an optimized secure system by planning how you partition your design. For example, for masters and slaves under the same hierarchy, it is possible for a non-secure master to initiate continuous transactions resulting in unsuccessful transfer to a secure slave. In the case of a memory aliasing, you must carefully designate secure or non-secure address maps to maintain reliable data.

**Managing Secure Settings in Qsys**

To create a secure design, you must first add masters and slaves and the connections between them. Once you establish connections between the masters and slaves, you can then set the security options, as required, with options in the **Security** column.

On the **System Contents** tab, in the **Security** column, the following selections are available for master, slave, and memory components:

- **Non-secure**—Master issues only non-secure transactions. There is no security available for the slave.
- **Secure**—Master issues only secure transactions. For the slave, Qsys prevents non-secure transactions from reaching the slave, and routes them to the default slave for the master that issued the transaction.
- **Secure Ranges**—Slave only, the specified address ranges within the slave’s address space are secure; all others are not. The format is a comma-separated list of inclusiveLow:inclusiveHigh addresses, for example, `0x0:0xffff,0x2000:0x20ff`.
- **TrustZone-aware**—Master issues either secure or non-secure transactions at run-time. The slave accepts either secure or non-secure transactions at run-time.
After setting security options for the masters and slaves, you must identify those masters that require a default slave before generation. To designate a slave as the default slave, turn on Default Slave in the Systems Contents tab. A master can have only one default slave.

**Understanding Compilation-Time Security Configuration Options**

The following compile-time configurations are available when creating secure designs that have mixed secure and non-secure components:

- Masters that support TrustZone and are connected to slaves that are compile-time secure. This configuration requires a default slave.
- Slaves that support TrustZone and are connected to masters that have compile-time secure settings. This configuration does not require a default slave.
- Master connected to slaves with secure address ranges. This configuration requires a default slave.

**Generating Output Files From a Qsys System**

Qsys system generation creates the interconnect between components and generates files that you use to synthesize or simulate the design. You specify the files that you want to generate on the Generation tab. You can generate simulation models, simulation testbench files, as well as HDL files for Quartus II synthesis, or a Block Symbol File (.bsf) for schematic design.

For your simulation model and testbench system, you can select Verilog or VHDL for the top-level module language, which applies to the system’s top-level definition and child instance that support generation for the selected target language.

For synthesis, you can select the top-level module language as Verilog or VHDL, which applies to the system’s top-level definition. If the design contains a composed _hw.tcl component or .qsys sub-modules, the language selection also applies to the sub-modules. For non-composed _hw.tcl sub-modules, a Verilog synthesis file is generated.

The default target language for simulation, testbench system, and synthesis is Verilog. Qsys places the generated output files in a subdirectory of your project directory, along with an HTML report file. To change the default behavior, on the Generation tab, specify a new directory under Output Directory.
Figure 7–7 illustrates the directory structure for the output files.

**Figure 7–7. Qsys Generated Files Directory Structure**

![Diagram of Qsys Generated Files Directory Structure]

Table 7–4 describes the files that Qsys generates. Each time you generate your system, Qsys overwrites these files, therefore, you should not edit Qsys-generated output files. If you have constraints, such as board-level timing constraints, Altera recommends that you create a separate Synopsys Design Constraints File (.sdc) and include that file in your Quartus II project. If you need to change top-level I/O pin names or instance name, Altera recommends you create a top-level HDL file that instantiates the Qsys system, so that the Qsys-generated output is instantiated in your design without any changes to the Qsys output files.

<table>
<thead>
<tr>
<th>File Name or Directory Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&lt;qsys_design&gt;</code></td>
<td>The top-level project directory.</td>
</tr>
<tr>
<td><code>&lt;qsys_design&gt;.bsf</code></td>
<td>A Block Symbol File (.bsf) representation of the top-level Qsys system for use in Quartus II Block Diagram Files (.bdf).</td>
</tr>
</tbody>
</table>
| `<qsys_design>.html`        | A report for the system, which provides a system overview including the following information:  
- External connections for the system  
- A memory map showing the address of each slave with respect to each master to which it is connected  
- Parameter assignments for each component |
| `<qsys_design>.sopcinfo`    | Describes the components and connections in your system. This file is a complete system description and is used by downstream tools such as the Nios II tool chain. It also describes the parameterization of each component in the system; consequently, you can parse its contents to get requirements when developing software drivers for Qsys components. This file and the `system.h` file generated for the Nios II tool chain include address map information for each slave relative to each master that accesses the slave. Different masters may have a different address map to access a particular slave component. |
| `/synthesis`               | This directory includes the Qsys-generated output files that the Quartus II software uses to synthesize your design. |
| `<qsys_design>.v`           | An HDL file for the top-level Qsys system that instantiates each component in the system. |
CMSIS Support for Qsys Systems With An HPS Component

Qsys systems that contain a Hard Processor System (HPS) component generate a System View Description (.svd) file that lists peripherals connected to the ARM processor. The .svd (or CMSIS-SVD) file format is an XML schema specified as part of the Cortex Microcontroller Software Interface Standard (CMSIS) provided by ARM. The CMSIS-SVD file allows HPS System Debug tools (such as the DS-5 Debugger) to gain visibility into the register maps of peripherals connected to the HPS within a Qsys system.

Qsys supports the ability for IP component designers to specify register map information on their slave interfaces. This allows components with slave interfaces that are connected to an HPS component to include their internal register description in the generated .svd file. To specify their internal register map, the IP component designer must write and generate their own .svd file and attach it to the slave interface using the following command:

```
set_interface_property <slave interface> CMSIS_SVD_FILE <file path>
```
For information about the set_interface_property command and its properties, refer to the Component Interface Tcl Reference chapter in the Quartus II Handbook. For complete CMSIS specifications, refer to CMSIS - Cortex Microcontroller Software Interface Standard on the ARM website.

Using Qsys With the Quartus II Software

This section describes the Quartus® II software features that integrate with Qsys, including the following:

- Quartus II IP File
- Synopsys Design Constraint
- Quartus II Simulation IP File
- PLLs and Clocks

Quartus II Project Files

The Quartus II IP File (.qip) provides the Quartus II software with all required information about your Qsys system. Qsys creates the .qip during system generation and adds a reference to it in the Quartus II Settings File (.qsf). The information required to process most Qsys components is included in the system's single .qip file, though some more complex components provide their own .qip file, in which case the system's .qip file references the component’s .qip file.

You must add the Qsys-generated Quartus II IP File (.qip) to your Quartus II project before you compile a design that includes a Qsys system. The .qip file is stored in the synthesis directory after generation, and lists the files necessary for compilation, and includes references to the following information:

- HDL files used in the Qsys system
- TimeQuest Timing Analyzer Synopsys Design Constraint (.sdc) files
- Component definition files for archiving purposes

Qsys automatically generates an .sdc file for Qsys systems and components. In most cases, you use TimeQuest constraints to declare false paths for signals that cross clock domains within a component, so that the TimeQuest Timing Analyzer does not perform setup and hold analysis for them. You can add .sdc files for custom components, with the Add Files command on Files tab in the Component Editor.

To use Nativelink simulation integration with a Qsys system, you must add the Quartus II Simulation IP File (.sip) file to your Quartus II project. The .sip file lists the files necessary for simulation with Nativelink. The .sip file is stored in the synthesis directory after generation.

Add the generated .qip file, not the .qsys file, to your Quartus II project.

Refer to the Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook for further description of the TimeQuest Timing Analyzer.

For more information about adding files to your Quartus II project, refer to Managing Files in a Project in Quartus II Help.
Working With PLLs and Clocks

You must provide clock and timing constraints in Synopsys Design Constraint File (.sdc) format to direct Quartus II synthesis and fitting to optimize the design appropriately, and to set up the TimeQuest timing analyzer to check that the design meets timing performance requirements.

You must specify a base clock assignment for each clock input with the create_clock command, and then you can use the derive_pll_clocks command to define the PLL clock output frequencies and phase shifts for all PLLs in the Quartus II project.

The Qsys system shown in Figure 7–8 illustrates the .sdc commands required for the case of a single clock input signal called clk, and one PLL with a single output.

Figure 7–8. Single Clock Input Signal

For this system, use the following commands in your .sdc file for the TimeQuest Timing Analyzer:

create_clock -name master_clk -period 20 [get_ports {clk}]
derive_pll_clocks

These commands create the input clock and the derived clock output of the PLL. The TimeQuest Timing Analyzer analyzes and reports performance of the constrained clocks in the Clocks Summary report, as shown in Figure 7–9.

Figure 7–9. Clocks Summary Report

<table>
<thead>
<tr>
<th>Clock Name</th>
<th>Type</th>
<th>Period</th>
</tr>
</thead>
<tbody>
<tr>
<td>1 master_clk</td>
<td>Base</td>
<td>20.000</td>
</tr>
<tr>
<td>2 the_my_pll_componentauto_generatedp1 clk(0)</td>
<td>Generated</td>
<td>30.000</td>
</tr>
</tbody>
</table>

master_clk is defined by the create_clock command, and the my_pll clock is derived from the derive_pll_clocks command.
Simulating a Qsys System

The Qsys Generation tab provides the following options for simulating a Qsys system:

- Generate the Verilog or VHDL simulation model for your system to use in your own simulation environment.
- Generate a standard or simple testbench system with BFM or Mentor Verification IP (for AXI3/AXI4) components that drive the external interfaces of your system, and generate a Verilog or VHDL simulation model for the testbench system to use in your simulation tool.
- First generate a testbench system, and then modify the testbench system in Qsys before generating its simulation model.

In most cases, you should select only one of the simulation model options, that is generate a simulation model for the original system, or for the testbench system.

Table 7–5 summarizes the options on the Generation tab that correspond to the simulation flows described above.

<table>
<thead>
<tr>
<th>Simulation Setting</th>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Create simulation model</td>
<td>None</td>
<td>Creates simulation model files and simulation scripts. Use this option to include the simulation model in your own custom testbench or simulation environment. You can also use this option to generate models for a testbench system that you have modified.</td>
</tr>
<tr>
<td></td>
<td>Verilog, VHDL</td>
<td></td>
</tr>
<tr>
<td>Create testbench Qsys system</td>
<td>Standard, BFMs for standard Qsys Interconnect</td>
<td>Creates a testbench Qsys system with BFM components attached to exported Avalon and AXI3 interfaces. Includes any simulation partner modules specified by IP cores in the system. In Qsys 13.0, the testbench generator supports AXI interfaces and can connect AXI3/AXI4 interfaces to Mentor Graphics AXI3/AXI4 master/slave BFM. For more information, refer to the Mentor Verification IP (VIP) Altera Edition (AE) document. However, BFM supports only an address width of up to 32-bits.</td>
</tr>
<tr>
<td></td>
<td>Simple, BFMs for clocks and resets</td>
<td>Creates a testbench Qsys system with BFM components driving only clocks and reset interfaces. Includes any simulation partner modules specified by IP cores in the system.</td>
</tr>
<tr>
<td>Create testbench simulation model</td>
<td>None, Verilog, VHDL</td>
<td>Creates simulation model files and simulation scripts for the testbench Qsys system specified in the setting above. Use this option if you do not need to modify the Qsys-generated testbench before running the simulation.</td>
</tr>
<tr>
<td>Create HDL design files for synthesis</td>
<td>On/Off</td>
<td>Creates Verilog or VHDL design files.</td>
</tr>
<tr>
<td>Top-level module language for synthesis</td>
<td>Verilog, VHDL</td>
<td>Creates the top-level module in the system in the selected language.</td>
</tr>
<tr>
<td>Create block symbol files (.bsf)</td>
<td>On/Off</td>
<td>You can optionally create a (.bsf) file to use in schematic Block Diagram File (.bdf) designs.</td>
</tr>
<tr>
<td>Output Directory</td>
<td>&lt;directory name&gt;</td>
<td>Allows you to browse and locate an alternate directory than the project directory for each generation target.</td>
</tr>
</tbody>
</table>
For more information about using bus functional models (BFMs) and monitors to simulate Avalon standard interfaces, including tutorials demonstrating sample systems, refer to the *Avalon Verification IP Suite User Guide*. For AXI verification protocol information, refer to the *Mentor Verification IP (VIP) Altera Edition (AE)* document.

For more information about generating system synthesis or simulation models, and a standard Qsys testbench, refer to *Generating a System for Synthesis or Simulation* and *Generation Tab (Qsys)* in Quartus II Help.

### Testbench Design Flow

You can use the following design flows to create a testbench system of your Verilog or VHDL design.

**Generate the Testbench System and a Simulation model at the Same Time (Verilog only)**

1. Create a Qsys system.
2. Generate a Verilog testbench system and the simulation model for the testbench system on the Qsys **Generation** tab.
3. Create a custom test program for the BFMs.
4. Compile and load the Qsys design and testbench in your simulator, and then run the simulation.

**Generate the Testbench System (Verilog and VHDL)**

1. Create a Qsys system.
2. Generate a Verilog or VHDL testbench system on the Qsys **Generation** tab.
3. Open the testbench system in Qsys. Make changes, as needed, to the BFMs, such as changing the BFM instance names and BFM’s **VHDL ID** value. You can modify the **VHDL ID** value in the **Altera Avalon Interrupt Source** component.
4. If you modified a BFM, generate the simulation model for the testbench system on the Qsys **Generation** tab.
5. Create a custom test program for the BFMs.
6. Compile and load the Qsys design and testbench in your simulator, and then run the simulation.

### Adding Assertion Monitors

You can add monitors to Avalon-MM, AXI, and Avalon-ST interfaces in your system to verify protocol correctness and test coverage with a simulator that supports SystemVerilog assertions.

*Modelsim Altera Edition does not support SystemVerilog assertions. If you want to use assertion monitors, you will need to use an advanced simulator such as Mentor Questasim, Synopsys VCS, or Cadence Incisive.*
Figure 7–10 demonstrates the use of monitors with an Avalon-MM monitor between the previously connected pcie_compiler bar1_0_Prefetchable Avalon-MM master interface and the dma_0 control_port_slave Avalon-MM slave interface.

Similarly, you can insert an Avalon-ST monitor between Avalon-ST source and sink interfaces.

**Simulation Scripts**

Qsys generates simulation scripts to script the simulation environment set up for Mentor Graphics Modelsim and Questasim, Synopsys VCS and VCS MX, Cadence® Incisive® Enterprise Simulator (NCSIM), and the Aldec Riviera-PRO Simulator. You can use the scripts to compile the required device libraries and system design files in the correct order and elaborate or load the top-level design for simulation.

The simulation scripts provide the following variables that allow flexibility in your simulation environment:

- **TOP_LEVEL_NAME**—If the Qsys testbench system is not the top-level instance in your simulation environment because you instantiate the Qsys testbench within your own top-level simulation file, set the `TOP_LEVEL_NAME` variable to the top-level hierarchy name.
- **QSYS_SIMDIR**—If the simulation files generated by Qsys are not in the simulation working directory, use the `QSYS_SIMDIR` variable to specify the directory location of the Qsys simulation files.
- **QUARTUS_INSTALL_DIR**—Points to the device family library.

**Example 7–5** shows a simple top-level simulation HDL file for a testbench system pattern_generator_tb, which was generated for a Qsys system called pattern_generator. The `top.sv` file defines the top-level module that instantiates the pattern_generator_tb simulation model as well as a custom SystemVerilog test program with BFM transactions, called test_program.

**Example 7–5. Top-level Simulation HDL File**

```verilog
module top();
    pattern_generator_tb tb();
    test_program pgm();
endmodule
```
Refer to the following documents for simulation script examples:

- ModelSim-Altera software, Mentor Graphics ModelSim support
- Synopsys VCS and VCS MX support
- Cadence Incisive Enterprise Simulator (IES) support
- Aldec Active-HDL and Rivera-PRO support

**Simulating Software Running on a Nios II Processor**

To simulate the software in a system driven by a Nios II embedded processor, generate the simulation model for a simple Qsys testbench system by completing the following steps:

1. On the Generation tab, set Create testbench Qsys system to Simple, BFMs for clocks and resets.
2. Set Create testbench simulation model to Verilog or VHDL.
3. Click Generate.

Follow these steps to use the software build tools for simulation:

1. Open the Nios II Software Build Tools for Eclipse.
2. Set up an application project and board support package (BSP) for the `<qsys_system>.sopcinfo` file.
3. To optimize the BSP for simulation and disable hardware programming, right-click the BSP project and click Properties, and then click Nios II BSP Properties, and turn on ModelSim only, no hardware support.
4. To simulate, right-click the application project in Eclipse, point to Run as, and then click 4 Nios II ModelSim. The Run As Nios II ModelSim command sets up the ModelSim simulation environment, compiles and loads the Nios II software simulation.
5. To run the simulation in ModelSim, type `run -all` in the ModelSim transcript window.
6. If prompted, set ModelSim configuration settings and select the correct Qsys Testbench Simulation Package Descriptor (.spd) file, `<qsys_system>_tb.spd`. The .spd file is generated with the testbench simulation model for Nios II designs and specifies all the files required for the Nios II software simulation.

For more information about the Nios II SBT for Eclipse, refer to *Getting Started with the Graphical User Interface* in the Nios II Software Developer's Handbook. For more information about the Nios II SBT command-line options, refer to *Getting Started from the Command-Line in the Nios II Software Developer's Handbook*.

**System Examples**

This section includes a detailed system example that demonstrates design hierarchy and the use of pipeline bridges, and an example that shows the use of instance parameters to control the instantiation of subcomponents in a hierarchical system.
PCI Express Subsystem Example

Figure 7–11 shows the details of the PCI Express example subsystem, which is also illustrated at a high level in Figure 7–5 on page 7–17. In this example, an application running on the root complex processor programs the DMA controller. The DMA controller’s Avalon-MM read and write master interfaces initiate transfers to and from the DDR3 memory and to the PCI Express Avalon-MM TX data port. The system exports the DMA master interfaces through an Avalon-MM pipeline bridge. As Figure 7–11 illustrates, all three masters connect to a single slave interface. During system generation, Qsys automatically inserts arbitration logic to control access to this slave interface. By default, the arbiter provides equal access to all requesting masters; however, you can weight the arbitration by changing the number of arbitration shares for the requesting masters. The second pipeline bridge allows an external master, such as a host processor, to also issue transactions to the CSR interfaces.

For more information, refer to “Arbitration” in the Qsys Interconnect chapter in volume 1 of the Quartus II Handbook.

Figure 7–11. PCI Express Subsystem
Figure 7–12 shows the Qsys representation of the PCI Express subsystem.

**Figure 7–12. Qsys Representation of the PCI Express Subsystem**

### Ethernet Subsystem Example

Figure 7–13 expands the details of the Ethernet subsystem example from Figure 7–5. In this subsystem, the transmit (TX) DMA receives data from the DDR3 memory and writes it to the Altera Triple-Speed Ethernet IP core using an Avalon-ST source interface. The receive (RX) DMA accepts data from the Triple-Speed Ethernet IP core on its Avalon-ST sink interface and writes it to DDR3 memory.

The read and write masters of both Scatter-Gather DMA controllers and the Triple-Speed Ethernet IP core connect to the DDR3 memory through an Avalon-MM pipeline bridge. This Ethernet example subsystem exports all three control and status interfaces through an Avalon-MM pipeline bridge, which connects to a controller outside of the Qsys system.

**Figure 7–13. Scatter-Gather DMA-to-Ethernet Subsystem**
Figure 7–14 shows the Qsys representation of the Ethernet subsystem.

### Figure 7–14. Qsys Representation of the Ethernet Subsystem

<table>
<thead>
<tr>
<th>Use</th>
<th>Connections</th>
<th>Module</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>csr</td>
<td>csr</td>
<td>Scatter-Catcher DMA Controller</td>
</tr>
<tr>
<td></td>
<td>descriptor_read</td>
<td>descriptor_read</td>
<td>Avalon Memory Mapped Slave</td>
</tr>
<tr>
<td></td>
<td>descriptor_write</td>
<td>descriptor_write</td>
<td>Avalon Memory Mapped Master</td>
</tr>
<tr>
<td></td>
<td>csr_irq</td>
<td>csr_irq</td>
<td>Avalon Memory Mapped Master</td>
</tr>
<tr>
<td></td>
<td>write</td>
<td>write</td>
<td>Avalon Memory Mapped Master</td>
</tr>
<tr>
<td></td>
<td>m_read</td>
<td>m_read</td>
<td>Avalon Memory Mapped Master</td>
</tr>
<tr>
<td></td>
<td>n_read</td>
<td>n_read</td>
<td>Avalon Memory Mapped Master</td>
</tr>
<tr>
<td></td>
<td>triple_speed_ethernet_0</td>
<td>transmit</td>
<td>Triple-Speed Ethernet</td>
</tr>
<tr>
<td></td>
<td></td>
<td>receive</td>
<td>Avalon Streaming Sink</td>
</tr>
<tr>
<td></td>
<td></td>
<td>control_host</td>
<td>Avalon Streaming Source</td>
</tr>
<tr>
<td></td>
<td></td>
<td>conduct_connection</td>
<td>Avalon Memory Mapped Slave</td>
</tr>
<tr>
<td></td>
<td></td>
<td>cal_clk</td>
<td>Conduct Endpoint</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Avalon Streaming Source</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Avalon Memory Mapped Slave</td>
</tr>
</tbody>
</table>

### PCI Express to Ethernet Bridge Example

The PCI Express-to-Ethernet Bridge example in Figure 7–15 includes two clock domains and an Ethernet subsystem. The PCI Express and Ethernet subsystems run at 125 MHz. The DDR3 SDRAM controller runs at 200 MHz. Qsys automatically inserts clock crossing logic to synchronize the DDR3 SDRAM Controller with the PCI Express and Ethernet subsystems.

### Figure 7–15. PCI Express-to-Ethernet Bridge Example System

![Diagram of PCI Express to Ethernet Bridge Example System](image)
Figure 7–16 shows the Qsys representation of the PCI Express-to-Ethernet Bridge example.

### Pipeline Bridges

The PCI Express to Ethernet bridge example system uses several pipeline bridges. These bridges must be configured to accommodate the address range of all of connected components, including the components in the originating subsystem and the components in the next higher level of the system hierarchy. As the name suggests, the pipeline bridge inserts a pipeline stage between the connected components. Altera recommends registering signals at the subsystem interface level for the following reasons:

- Registering interface signals decreases the amount of combinational logic that must be completed in one cycle, making it easier to meet timing constraints.
- Registering interface signals raises the potential frequency, or f\(_{\text{MAX}}\), of your design at the expense of an additional cycle of latency, which might adversely affect system throughput.
- The Quartus II incremental compilation feature can achieve better f\(_{\text{MAX}}\) results if the subsystem boundary is registered.

For more information about optimizing a Qsys design for performance using bridges and other techniques, refer to *Optimizing System Performance for Qsys* in volume 1 of the *Quartus II Handbook*.

AXI bridge components are not available in the Quartus II software, but you can connect AXI interfaces with other bridge types. Connections between AXI and Avalon interfaces are made without requiring the use of explicitly instantiated bridges; the interconnect provides all necessary bridging logic. For more information about AXI support, refer to the *Qsys System Design Components* chapter in volume 1 of the *Quartus II Handbook*. 
Hierarchical System Using Instance Parameters Example

You can use an instance parameter to control the implementation of system components from a higher-level Qsys system. You define instance parameters on the Instance Parameters tab in Qsys.

In this example, a Qsys design called `my_system.qsys` has two instances of the same IP component, `My_IP`. `My_IP` is a Qsys component with a system identification parameter called `MY_SYSTEM_ID`. When `my_system.qsys` is instantiated within another higher-level Qsys system, the two `My_IP` subcomponents require different values for their `MY_SYSTEM_ID` parameters based on a value determined by the higher-level system. In this example, the value specified by the top-level system is designated `top_id` and in `my_system.qsys`, the component instance `comp0` requires `MY_SYSTEM_ID` set to `top_id + 1`, and instance `comp1` requires `MY_SYSTEM_ID` set to `top_id + 2`.

The following `_hw.tcl` code defines the `MY_SYSTEM_ID` system ID parameter in the IP component `My_IP`:

```tcl
add_parameter MY_SYSTEM_ID int 8
set_parameter_property MY_SYSTEM_ID DISPLAY_NAME \ MY_SYSTEM_ID_PARAM
set_parameter_property MY_SYSTEM_ID UNITS None
```

To satisfy the design requirements for this example, you define an instance parameter in `my_system.qsys` that is set by the higher-level system, and then define an instance script to specify how the values of the parameters of the `My_IP` components instantiated in `my_system.qsys` are affected by the value set on the instance parameter.

To do this, in Qsys, open the `my_system.qsys` Qsys system that instantiates the two instances of the `My_IP` components. On the Instance Parameters tab, create a parameter called `system_id`. For this example, you can set this parameter to be of type Integer and choose 0 as the default value.

Next, you provide a Tcl Instance Script that defines how the value of the `system_id` parameter should affect the parameters of `comp0` and `comp1` subcomponents in `my_system.qsys`.

The example script in Example 7–6 gets the value of the parameter `system_id` from the top-level system and saves it as `top_id`, and then increments the value by 1 and 2. The script then uses the new calculated values to set the `MY_SYSTEM_ID` parameter in the `My_IP` component for the instances `comp0` and `comp1`. The script uses informational messages to print the status of the parameter settings when the `my_system.qsys` system is added to the higher-level system.
You can click **Preview Instance** to see a parameters panel that allows you to modify the parameter value interactively and see the effect of the scripts in the message panel which can be useful for debugging the script. In this example, if you change the parameter value in the **Preview** screen, the component generates messages to report the top-level ID parameter value and the parameter values used for the two instances of the component.

For more information on creating a parameter on the **Instance Parameters** tab, refer to *Working with Instance Parameters in Qsys* in Quartus II Help.

### Using Qsys Command-Line with Utilities and Scripts

You can perform many of the functions available in the Qsys GUI from the command-line with the `qsys-generate`, `qsys-script`, `ip-generate`, and `ip-make-simscript` utilities. You run these command-line executables from the Quartus II installation directory, as follows:

```
<Quartus II installation directory>\quartus\sopc_builder\bin
```

You can use `qsys-generate`, `ip-generate`, and `ip-make-simscript` to generate Qsys output files outside of the Qsys GUI. You can use `qsys-script` to create and manipulate or manage a Qsys system with command-line scripting. The following subsections provide information about using Qsys from the command-line and with scripts. For command-line help listing options for these executables, type the following command:

```
<Quartus II installation directory>\quartus\sopc_builder\bin\<executable name> --help
```

Example 7–7 below shows an example using Qsys command-line scripting.

---

**Example 7–6. Using an Instance Script To Set Parameters On Subcomponents**

```bash
package require qsys 13.0
Set_module_property Composition_callback My_callback
proc My_callback { } {
    # Get The Value Of System_id Parameter From The Higher-level System
    Set Top_id [Get_parameter_value System_id]
    # Print Info Message
    Send_message Info "System_id Value Specified: $top_id"
    # Use Above Value To Set Parameter Values For The Subcomponents
    Set Child_id_0 [Expr {$top_id + 1} ]
    Set Child_id_1 [Expr {$top_id + 2} ]
    # Set The Parameter Values On The Subcomponent Instances
    Set_instance_parameter_value Comp0 My_system_id $child_id_0
    Set_instance_parameter_value Comp1 My_system_id $child_id_1
    # Print Info Messages
    Send_message Info "System_id Value Used In Comp0: $child_id_0"
    Send_message Info "System_id Value Used In Comp1: $child_id_1"
}
```

---

Example 7–7 below shows an example using Qsys command-line scripting.
Chapter 7: Creating a System With Qsys

Using Qsys Command-Line with Utilities and Scripts

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Example 7–7. Using an Instance Script To Display Component Names

```
qsys-script --script=my_script.tcl --system-file=fancy.qsys
my_script.tcl contains:
package require -exact qsys 13.0
# get all the instance names in the system and print them one by one
set instances [ get_instances ]
foreach instance $instances {
    send_message Info "$instance"
}
```

For more information about Qsys utilities and scripting, including examples, refer to the Altera Wiki Qsys Scripts page.

Generating Qsys Systems with the qsys-generate Utility

You can use the qsys-generate utility to generate RTL for your Qsys system, to compile in Quartus II, simulation models and scripts, and to create testbench systems for testing your Qsys system in a simulator using BFMs. Output from the qsys-generate command is the same as when generating using the Qsys GUI.

When possible, you should use qsys-generate instead of ip-generate and ip-make-simscript. The command-line options for qsys-generate are simpler, and the generation options and output directory structure always match those from the Qsys GUI generation.

The following is a list of options that you can use with the qsys-generate utility:

- **<1st arg file>**—Required. The name of the .qsys system file to generate.
- **--synthesis=<VERILOG | VHDL>**—Optional. Creates synthesis HDL files that Qsys uses to compile the system in a Quartus II project. You must specify the preferred generation language for the top-level RTL file for the generated Qsys system.
- **--block-symbol-file**—Optional. Creates a block symbol file (.bsf) for the system.
- **--simulation=<VERILOG | VHDL>**—Optional. Creates a simulation model for the system. The simulation model contains generated HDL files for the simulator, and may include simulation-only features. You must specify the preferred simulation language.
- **--testbench=<SIMPLE | STANDARD>**—Optional. Creates a testbench system. The testbench system instantiates the original system, adding bus functional models to drive the top-level interfaces. Once generated, the bus functional models interact with the system in the simulator.
- **--testbench-simulation=<VERILOG | VHDL>**—Optional. After creating the testbench system, also create a simulation model for the testbench system.
- **--output-directory=<value>**—Optional. Sets the output directory. Each generation target is created in a subdirectory of the output directory. If you do not specify the output directory, a subdirectory of the current working directory matching the name of the system is used.
- **--search-path=**<value>—Optional. If omitted, a standard default path is used. If provided, a comma-separated list of paths is searched. To include the standard path in your replacement, use "$", for example, "/extra/dir,$".

- **--jvm-max-heap-size=**<value>—Optional. The maximum memory size that Qsys uses for allocations when running this tool. The value is specified as <size><unit> where unit can be m (or M) for multiples of megabytes or g (or G) for multiples of gigabytes. The default value is 512m.

### Generating Qsys Systems with the ip-generate Utility

You use `ip-generate` to configure parameters and generate HDL and other output files for Qsys systems and IP cores.

When you generate a system in the Qsys GUI, the generation output messages include the command-lines that you can use to run generation with the same settings using the `ip-generate` utility. For example, when you generate synthesis files for a system called `test.qsys` in `c:/my_dir`, Qsys outputs a message such as the following specifying the command-lines for the `ip-generate` utility:

```
```

The following is a list of options that you can use with the `ip-generate` utility:

- **--project-directory=**<directory>—Optional. Components are found in the locations relative to the project, if any. By default, the current directory '.' is used. To exclude any project directory, use "".

- **--output-directory=**<directory>—Optional. This directory will contain the output file set(s). The directory is created if required. If omitted, the current directory is used.

- **--file-set=**<QUARTUS_SYNTH | SIM_VERILOG | SIM_VHDL>—Optional. Type of output to generate. QUARTUS_SYNTH produces HDL that is compiled by the Quartus II software integrated synthesis. SIM_VERILOG and SIM_VHDL produce simulation models in the respective languages.

- **--report-file=**<type><filename>—Optional. Partial or complete path for the generated report file, for example, `html:report.html`. A partial path is relative to the current directory. To assign multiple files, use this option multiple times. The following are common report types.
  - Block Symbol File (.bsf)
  - Hypertext Markup Language File (.html)
  - Quartus II IP File (.qip)
  - Quartus II Simulation IP File (.sip)
  - SOPC Information File (.sopcinfo)
  - Simulation Package Descriptor File (.spd)
Chapter 7: Creating a System With Qsys

Using Qsys Command-Line with Utilities and Scripts

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Volume 1: Design and Synthesis

- **--search-path=value**—Optional. If omitted, a standard default path is used. If provided, a comma-separated list of paths is searched. To include the standard installation directory path, use $, for example, /<directory path>/dir,$. You can also use any directory path, or a path to an .ipx file. Multiple directory references are separated with a comma.
- **--component-file=file**—Optional. A file from which to extract an IP component or system, for example, "my_system.qsys" or "my_component_hw.tcl".
- **--component-name=value**—Optional. The name of an IP component to instantiate, for example, "altera_avalon_uart". If a component file is specified, the component must be found within that file.
- **--component-parameter=value**—Optional. A single value assignment for a component parameter, for example, "--component-param=WIDTH=11". To assign multiple parameters, use this option multiple times.
- **--system-info=parameter=value**—Optional. A single value assignment, for example, "--system-info=DEVICE_FAMILY=Stratix IV". To assign multiple system parameters, use this option multiple times. Common parameters are: DEVICE_FAMILY, DEVICE, and DEVICE_SPEEDGRADE.
- **--language=value**—Optional. Supported values are Verilog and VHDL, with Verilog as the default value.
- **--jvm-max-heap-size=value**—Optional. The maximum memory size that Qsys uses for allocations when running this utility. The value is specified as <size><unit>, where unit can be m (or M) for multiples of megabytes or g (or G) for multiples of gigabytes. The default value is 512m.

Generating Qsys Simulation Scripts with the ip-make-simscript Utility

You can use the ip-make-simscript utility to create simulation scripts, and specify options related to simulation scripts that are not available in the Qsys GUI. qsys-generate, by default, generates these simulation scripts; this utility is only necessary if you use the ip-generate flow. For example, you can use the --compile-to-work option of the ip-make-simscript command if you want to use a single directory for your simulation files rather than the default compilation structure. To use this utility, you must specify an .spd file that lists the required simulation files.

The following is a list of options that you can use for the ip-make-simscript utility:
- **--spd=file**—Required. The .spd files describe the list of files to be compiled, and the memory models hierarchy.
- **--output-directory=directory**—Optional. Directory path specifying the location of output files. If not specified, defaults to the directory from which the ip-make-simscript is run.
- **--compile-to-work**—Optional—Compiles all design files to the default library, ../work.
- **--use-relative-paths**—Optional. Uses relative paths whenever possible.
**Creating and Managing a System with qsys-script**

You can use the `qsys-script` tool to create and manipulate a Qsys system with Tcl scripting commands.

You must provide a package version for the `qsys-script`. If you do not specify the `--package-version=<value> qsys-script` command, you must then provide a Tcl script and request the system scripting API directly with the `package require -exact qsys <version>` command.

The following is a list of options that you can use with the `qsys-script` utility:

- **--system-file=**<file>—Optional. Specifies the path to a .qsys system file. This system is loaded before running scripting commands.

- **--script=**<file>—Optional. A file containing Tcl scripting commands for creating or manipulating Qsys systems. If you specify both `--cmd` and `--script`, the `--cmd` commands are run before the script specified by `--script`.

- **--cmd=**<value>—Optional. A string that contains Tcl scripting commands to create or manipulate a Qsys system. If you specify both `--cmd` and `--script`, the `--cmd` commands are run before the script specified by `--script`.

- **--package-version=**<value>—Optional. Specifies which system scripting Tcl API version to use and determines the functionality and behavior of the Tcl commands. The Quartus II software supports the Tcl API scripting commands. If you do not specify the version on the command-line, your Tcl script must request the system scripting API directly with the `package require -exact qsys <version>` command.

- **--help**—Optional. Displays help for the `qsys-script` tool.

- **--search-path=**<value>—Optional. If omitted, a standard default path is used. If provided, a comma-separated list of paths is searched. To include the standard path in your replacement, use "$", for example, `/<directory path>/dir,$`. Multiple directory references are separated with a comma.

- **--jvm-max-heap-size=**<value>—Optional. The maximum memory size that is used by the `qsys-script` tool. You specify this value as `<size><unit>` where unit can be `m` or `M` for multiples of megabytes or `g` or `G` for multiples of gigabytes.

- **--naivelink-modes**—Optional. Generates files for Quartus II NativeLink RTL simulation.

- **--jvm-max-heap-size=**<value>—Optional. The maximum memory size that the `qsys-script` tool uses. You specify this value as `<size><unit>` when `unit` is `m` or `M` for multiples of megabytes, or `g` or `G` for multiples of gigabytes.
Qsys Scripting Command Reference

Table 7–6 summarizes the Qsys scripting commands and provides a reference to the full command description.

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_connection &lt;start&gt; [&lt;end&gt;]</td>
<td>page 7–41</td>
</tr>
<tr>
<td>add_instance &lt;name&gt; &lt;type&gt; [&lt;version&gt;]</td>
<td>page 7–41</td>
</tr>
<tr>
<td>add_interface &lt;name&gt; &lt;type&gt; &lt;direction&gt;</td>
<td>page 7–41</td>
</tr>
<tr>
<td>auto_assign_base_addresses &lt;instance&gt;</td>
<td>page 7–42</td>
</tr>
<tr>
<td>auto_assign_irqs &lt;instance&gt;</td>
<td>page 7–42</td>
</tr>
<tr>
<td>auto_connect &lt;element&gt;</td>
<td>page 7–42</td>
</tr>
<tr>
<td>create_system [&lt;name&gt;]</td>
<td>page 7–42</td>
</tr>
<tr>
<td>get_composed_connection_parameter_value &lt;instance&gt; &lt;childConnection&gt; &lt;parameter&gt;</td>
<td>page 7–43</td>
</tr>
<tr>
<td>get_composed_connection_parameters &lt;instance&gt; &lt;childConnection&gt;</td>
<td>page 7–43</td>
</tr>
<tr>
<td>get_composed_connections &lt;instance&gt;</td>
<td>page 7–43</td>
</tr>
<tr>
<td>get_composed_instance_assignment &lt;instance&gt; &lt;childInstance&gt; &lt;key&gt;</td>
<td>page 7–44</td>
</tr>
<tr>
<td>get_composed_instance_assignments &lt;instance&gt; &lt;childInstance&gt;</td>
<td>page 7–44</td>
</tr>
<tr>
<td>get_composed_instance_parameter_value &lt;instance&gt; &lt;childInstance&gt; &lt;parameter&gt;</td>
<td>page 7–44</td>
</tr>
<tr>
<td>get_composed_instance_parameters &lt;instance&gt; &lt;childInstance&gt;</td>
<td>page 7–45</td>
</tr>
<tr>
<td>get_composed_instances &lt;instance&gt;</td>
<td>page 7–45</td>
</tr>
<tr>
<td>get_connection_parameter_property &lt;connection&gt; &lt;parameter&gt; &lt;property&gt;</td>
<td>page 7–45</td>
</tr>
<tr>
<td>get_connection_parameter_value &lt;connection&gt; &lt;parameter&gt;</td>
<td>page 7–45</td>
</tr>
<tr>
<td>get_connection_parameters &lt;connection&gt;</td>
<td>page 7–46</td>
</tr>
<tr>
<td>get_connection_properties</td>
<td>page 7–46</td>
</tr>
<tr>
<td>get_connection_property &lt;connection&gt; &lt;property&gt;</td>
<td>page 7–46</td>
</tr>
<tr>
<td>get_connections [&lt;element&gt;]</td>
<td>page 7–46</td>
</tr>
<tr>
<td>get_instance_assignment &lt;instance&gt; &lt;key&gt;</td>
<td>page 7–47</td>
</tr>
<tr>
<td>get_instance_assignments &lt;instance&gt;</td>
<td>page 7–47</td>
</tr>
<tr>
<td>get_instance_interface_assignment &lt;instance&gt; &lt;interface&gt; &lt;key&gt;</td>
<td>page 7–47</td>
</tr>
<tr>
<td>get_instance_interface_assignments &lt;instance&gt; &lt;interface&gt;</td>
<td>page 7–47</td>
</tr>
<tr>
<td>get_instance_interface_parameter_property &lt;instance&gt; &lt;interface&gt; &lt;parameter&gt; &lt;property&gt;</td>
<td>page 7–48</td>
</tr>
<tr>
<td>get_instance_interface_parameter_value &lt;instance&gt; &lt;interface&gt; &lt;parameter&gt;</td>
<td>page 7–48</td>
</tr>
<tr>
<td>get_instance_interface_parameters &lt;instance&gt; &lt;interface&gt;</td>
<td>page 7–48</td>
</tr>
<tr>
<td>get_instance_interface_port_property &lt;instance&gt; &lt;interface&gt; &lt;port&gt; &lt;property&gt;</td>
<td>page 7–49</td>
</tr>
<tr>
<td>get_instance_interface_ports &lt;instance&gt; &lt;interface&gt;</td>
<td>page 7–49</td>
</tr>
<tr>
<td>get_instance_interface_properties</td>
<td>page 7–49</td>
</tr>
<tr>
<td>get_instance_interface_property &lt;instance&gt; &lt;interface&gt; &lt;property&gt;</td>
<td>page 7–49</td>
</tr>
<tr>
<td>get_instance_interfaces &lt;instance&gt;</td>
<td>page 7–50</td>
</tr>
<tr>
<td>get_instance_parameter_property &lt;instance&gt; &lt;parameter&gt; &lt;property&gt;</td>
<td>page 7–50</td>
</tr>
</tbody>
</table>
Table 7–6. Qsys System Scripting Command Reference (Part 2 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>get_instance_parameter_value &lt;instance&gt; &lt;parameter&gt;</code></td>
<td>page 7–50</td>
</tr>
<tr>
<td><code>get_instance_parameters &lt;instance&gt;</code></td>
<td>page 7–50</td>
</tr>
<tr>
<td><code>get_instance_port_property &lt;instance&gt; &lt;port&gt; &lt;property&gt;</code></td>
<td>page 7–51</td>
</tr>
<tr>
<td><code>get_instance_properties</code></td>
<td>page 7–51</td>
</tr>
<tr>
<td><code>get_instance_property &lt;instance&gt; &lt;property&gt;</code></td>
<td>page 7–51</td>
</tr>
<tr>
<td><code>get_instances</code></td>
<td>page 7–51</td>
</tr>
<tr>
<td><code>get_interface_port_property &lt;interface&gt; &lt;port&gt; &lt;property&gt;</code></td>
<td>page 7–52</td>
</tr>
<tr>
<td><code>get_interface_ports &lt;interface&gt;</code></td>
<td>page 7–52</td>
</tr>
<tr>
<td><code>get_interface_properties</code></td>
<td>page 7–52</td>
</tr>
<tr>
<td><code>get_interface_property &lt;interface&gt; &lt;property&gt;</code></td>
<td>page 7–52</td>
</tr>
<tr>
<td><code>get_interfaces</code></td>
<td>page 7–53</td>
</tr>
<tr>
<td><code>get_module_properties</code></td>
<td>page 7–53</td>
</tr>
<tr>
<td><code>get_module_property &lt;property&gt;</code></td>
<td>page 7–53</td>
</tr>
<tr>
<td><code>get_parameter_properties</code></td>
<td>page 7–53</td>
</tr>
<tr>
<td><code>get_port_properties</code></td>
<td>page 7–54</td>
</tr>
<tr>
<td><code>get_project_properties</code></td>
<td>page 7–54</td>
</tr>
<tr>
<td><code>get_project_property &lt;property&gt;</code></td>
<td>page 7–58</td>
</tr>
<tr>
<td><code>load_system &lt;file&gt;</code></td>
<td>page 7–54</td>
</tr>
<tr>
<td><code>lock_avalon_base_address &lt;instance.interface&gt;</code></td>
<td>page 7–55</td>
</tr>
<tr>
<td><code>preview_insert_avalon_streaming_adapters</code></td>
<td>page 7–55</td>
</tr>
<tr>
<td><code>remove_connection &lt;connection&gt;</code></td>
<td>page 7–55</td>
</tr>
<tr>
<td><code>remove_instance &lt;instance&gt;</code></td>
<td>page 7–55</td>
</tr>
<tr>
<td><code>remove_interface &lt;interface&gt;</code></td>
<td>page 7–56</td>
</tr>
<tr>
<td><code>save_system [&lt;file&gt;]</code></td>
<td>page 7–56</td>
</tr>
<tr>
<td><code>send_message &lt;level&gt; &lt;message&gt;</code></td>
<td>page 7–56</td>
</tr>
<tr>
<td><code>set_connection_parameter_value &lt;connection&gt; &lt;parameter&gt; &lt;value&gt;</code></td>
<td>page 7–57</td>
</tr>
<tr>
<td><code>set_instance_parameter_value &lt;instance&gt; &lt;parameter&gt; &lt;value&gt;</code></td>
<td>page 7–57</td>
</tr>
<tr>
<td><code>set_instance_property &lt;instance&gt; &lt;property&gt; &lt;value&gt;</code></td>
<td>page 7–57</td>
</tr>
<tr>
<td><code>set_interface_property &lt;interface&gt; &lt;property&gt; &lt;value&gt;</code></td>
<td>page 7–58</td>
</tr>
<tr>
<td><code>set_module_property &lt;property&gt; &lt;value&gt;</code></td>
<td>page 7–53</td>
</tr>
<tr>
<td><code>set_project_property &lt;property&gt; &lt;value&gt;</code></td>
<td>page 7–58</td>
</tr>
<tr>
<td><code>set_validation_property &lt;property&gt; &lt;value&gt;</code></td>
<td>page 7–59</td>
</tr>
<tr>
<td><code>unlock_avalon_base_address &lt;instance.interface&gt;</code></td>
<td>page 7–59</td>
</tr>
<tr>
<td><code>upgrade_sopc_system &lt;filename&gt;</code></td>
<td>page 7–59</td>
</tr>
<tr>
<td><code>validate_connection &lt;connection&gt;</code></td>
<td>page 7–60</td>
</tr>
<tr>
<td><code>validate_instance &lt;instance&gt;</code></td>
<td>page 7–60</td>
</tr>
<tr>
<td><code>validate_instance_interface &lt;instance&gt; &lt;interface&gt;</code></td>
<td>page 7–60</td>
</tr>
<tr>
<td><code>validate_system</code></td>
<td>page 7–60</td>
</tr>
</tbody>
</table>
Interface properties work differently for qsys scripting than with _hw.tcl scripting. In _hw.tcl, interfaces do not distinguish between properties and parameters; in qsys scripting, properties and parameters are unique.

add_connection
This command connects interfaces using an appropriate connection type. Interface names consist of a child instance name, followed by the name of an interface provided by that module, for example, mux0.out is the interface out on the instance named mux0.

<table>
<thead>
<tr>
<th><strong>add_connection</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

add_instance
This command adds an instance of a component, referred to as a child or child instance, to the system.

<table>
<thead>
<tr>
<th><strong>add_instance</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

add_interface
This command adds an interface to your system, which you can use to export an interface from within the system. You specify the exported interface with the command set_interface_property EXPORT_OF <instance.interface>.

<table>
<thead>
<tr>
<th><strong>add_interface</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
<tr>
<td></td>
</tr>
</tbody>
</table>
### auto_assign_base_addresses

This command assigns base addresses to memory-mapped interfaces on an instance in the system. Instance interfaces that are locked with `lock_avalon_base_address` command keep their addresses during address auto-assignment.

<table>
<thead>
<tr>
<th>auto_assign_base_addresses</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### auto_assign_irqs

This command assigns interrupt numbers to all connected interrupt senders on an instance in the system.

<table>
<thead>
<tr>
<th>auto_assign_irqs</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### auto_connect

This command creates connections from an instance or instance interface to matching interfaces in other instances in the system. For example, Avalon-MM slaves are connected to Avalon-MM masters.

<table>
<thead>
<tr>
<th>auto_connect</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
<tr>
<td></td>
</tr>
</tbody>
</table>

### create_system

This command replaces the current system in the system script with a new system with the specified name.

<table>
<thead>
<tr>
<th>create_system</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
get_composed_connection_parameter_value

This command returns the value of a parameter in a connection in the subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>get_composed_connection_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_composed_connection_parameters

This command returns a list of parameters on a connection in the subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>get_composed_connection_parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_composed_connections

This command returns a list of all connections in a subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>get_composed_connections</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
**get_composed_instance_assignment**

This command returns the value of an assignment on an instance of a subsystem, for an instance that models a subsystem.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_composed_instance_assignment &lt;instance&gt; &lt;childInstance&gt; &lt;key&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String The value of the assignment</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance The child instance containing a subsystem</td>
</tr>
<tr>
<td></td>
<td>childInstance The name of a child instance found in the subsystem</td>
</tr>
<tr>
<td></td>
<td>key The assignment key</td>
</tr>
</tbody>
</table>
| Example  | get_composed_instance_assignment subsystem_0 video_0 "embeddedsw.CMacro.colorSpace"

**get_composed_instance_assignments**

This command returns a list of assignments on an instance of a subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_composed_instance_assignments &lt;instance&gt; &lt;childInstance&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[] A list of assignment names</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance The child instance containing a subsystem</td>
</tr>
<tr>
<td></td>
<td>childInstance The name of a child instance found in the subsystem</td>
</tr>
<tr>
<td>Example</td>
<td>get_composed_instance_assignments subsystem_0 cpu</td>
</tr>
</tbody>
</table>

**get_composed_instance_parameter_value**

This command returns the value of a parameters on an instance in a subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_composed_instance_parameter_value &lt;instance&gt; &lt;childInstance&gt; &lt;parameter&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String The value of a parameter on an instance of a subsystem</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance The child instance containing a subsystem</td>
</tr>
<tr>
<td></td>
<td>childInstance The name of a child instance found in the subsystem</td>
</tr>
<tr>
<td></td>
<td>parameter The name of the parameter to query on an instance of a subsystem</td>
</tr>
<tr>
<td>Example</td>
<td>get_composed_instance_parameter_value subsystem_0 cpu DATA_WIDTH</td>
</tr>
</tbody>
</table>
**get_composed_instance_parameters**

This command returns a list of parameters on an instance of a subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_composed_instance_parameters &lt;instance&gt; &lt;childInstance&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[] A list of parameter names</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance The child instance containing a subsystem</td>
</tr>
<tr>
<td></td>
<td>childInstance The name of a child instance found in the subsystem</td>
</tr>
<tr>
<td>Example</td>
<td>get_composed_instance_parameters subsystem_0 cpu</td>
</tr>
</tbody>
</table>

**get_composed_instances**

This command returns a list of child instances in the subsystem, for an instance that contains a subsystem.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_composed_instances &lt;instance&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[] A list of instance names found in the subsystem</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance The child instance containing a subsystem</td>
</tr>
<tr>
<td>Example</td>
<td>get_composed_instances subsystem_0</td>
</tr>
</tbody>
</table>

**get_connection_parameter_property**

This command returns the value of a parameter property in a connection.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_connection_parameter_property &lt;connection&gt; &lt;parameter&gt; &lt;property&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>various The value of the parameter property</td>
</tr>
<tr>
<td>Arguments</td>
<td>connection The connection to query</td>
</tr>
<tr>
<td></td>
<td>parameter The name of the parameter</td>
</tr>
<tr>
<td></td>
<td>property The property of the connection</td>
</tr>
<tr>
<td>Example</td>
<td>get_connection_parameter_property cpu.data_master/dma0.csr baseAddress UNITS</td>
</tr>
</tbody>
</table>

**get_connection_parameter_value**

This command gets the value of a parameter on the connection. Parameters represent aspects of the connection that can be modified once the connection is created, such as the base address for an Avalon-MM connection.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_connection_parameter_value &lt;connection&gt; &lt;parameter&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>various The value of the parameter</td>
</tr>
<tr>
<td>Arguments</td>
<td>connection The connection to query</td>
</tr>
<tr>
<td></td>
<td>parameter The name of the parameter</td>
</tr>
<tr>
<td>Example</td>
<td>get_connection_parameter_value cpu.data_master/dma0.csr baseAddress</td>
</tr>
</tbody>
</table>
get_connection_parameters
This command returns a list of parameters found on a connection. The list of connection parameters is the same for all connections of the same type.

<table>
<thead>
<tr>
<th>get_connection_parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_connection_properties
This command returns a list of properties found on a connection. The list of connection properties is the same for all connections, regardless of type.

<table>
<thead>
<tr>
<th>get_connection_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_connection_property
This command returns the value of a connection property.

<table>
<thead>
<tr>
<th>get_connection_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_connections
This command returns a list of connections in the system if no element is specified. If a child instance is specified, for example cpu, all connections to any interface on the instance are returned. If an interface on a child instance is specified, for example cpu.instruction_master, only connections to that interface are returned.

<table>
<thead>
<tr>
<th>get_connections</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
### get_instance_assignment

This command returns the value of an assignment on a child instance.

<table>
<thead>
<tr>
<th>get_instance_assignment</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_assignments

This command returns a list of assignment keys for any assignments defined for the instance.

<table>
<thead>
<tr>
<th>get_instance_assignments</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_interface_assignment

This command returns the value of an assignment on an interface of a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_assignment</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_interface_assignments

This command returns the value of an assignment on an interface of a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_assignments</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
get_instance_interface_parameter_property
This command returns the property value on a parameter in an interface of a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_parameter_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_instance_interface_parameter_value
This command returns the value of a parameter of an interface in a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

get_instance_interface_parameters
This command returns a list of parameters for an interface in a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
### get_instance_interface_port_property

This command returns the property value of a port in the interface of a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_port_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_interface_ports

This command returns a list of ports in an interface of a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_ports</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_interface_properties

This command returns a list of properties that you can be query for an interface in a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_instance_interface_property

This command returns the property value for an interface in a child instance.

<table>
<thead>
<tr>
<th>get_instance_interface_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
get_instance_interfaces

This command returns a list of interfaces in a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_interfaces &lt;instance&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
</tr>
<tr>
<td></td>
<td>A list of interface names</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>The name of the child instance</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_interfaces uart_0</td>
</tr>
</tbody>
</table>

get_instance_parameter_property

This command returns the value of a property on a parameter in a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_parameter_property &lt;instance&gt; &lt;parameter&gt; &lt;property&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>various</td>
</tr>
<tr>
<td></td>
<td>The value of the parameter property</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>The name of the child instance</td>
</tr>
<tr>
<td></td>
<td>parameter</td>
</tr>
<tr>
<td></td>
<td>The name of the parameter in the instance</td>
</tr>
<tr>
<td></td>
<td>property</td>
</tr>
<tr>
<td></td>
<td>The name of the property of the parameter</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_parameter_property uart_0 baudRate ENABLED</td>
</tr>
</tbody>
</table>

get_instance_parameter_value

This command returns the value of a property in a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_parameter_value &lt;instance&gt; &lt;parameter&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>various</td>
</tr>
<tr>
<td></td>
<td>The value of the parameter</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>The name of the child instance</td>
</tr>
<tr>
<td></td>
<td>parameter</td>
</tr>
<tr>
<td></td>
<td>The name of the parameter in the instance</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_parameter_value uart_0 baudRate</td>
</tr>
</tbody>
</table>

get_instance_parameters

This command returns a list of parameters in a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_parameters &lt;instance&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
</tr>
<tr>
<td></td>
<td>A list of parameters in the instance</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>The name of the child instance</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_parameters uart_0</td>
</tr>
</tbody>
</table>
**get_instance_port_property**

This command returns the value of a property of a port contained by an interface in a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_port_property &lt;instance&gt; &lt;port&gt; &lt;property&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>various</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>port</td>
</tr>
<tr>
<td></td>
<td>property</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_port_property uart_0 tx WIDTH</td>
</tr>
</tbody>
</table>

**get_instance_properties**

This command returns a list of properties for a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_properties</td>
</tr>
</tbody>
</table>

**get_instance_property**

This command returns the value of a property for a child instance.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instance_property &lt;instance&gt; &lt;property&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>instance</td>
</tr>
<tr>
<td></td>
<td>property</td>
</tr>
<tr>
<td>Example</td>
<td>get_instance_property cpu ENABLED</td>
</tr>
</tbody>
</table>

**get_instances**

This command returns a list of the instance names for all child instances in the system.

<table>
<thead>
<tr>
<th>Usage</th>
<th>get_instances</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td>get_instances</td>
</tr>
</tbody>
</table>
**get_interface_port_property**

This command returns the value of a property of a port contained by one of the top-level exported interfaces.

<table>
<thead>
<tr>
<th>get_interface_port_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_interface_ports**

This command returns the names of all of the ports that have been added to an interface.

<table>
<thead>
<tr>
<th>get_interface_ports</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_interface_properties**

This command returns the names of all the available interface properties. The list of interface properties is the same for all interface types.

<table>
<thead>
<tr>
<th>get_interface_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_interface_property**

This command returns the value of a property from the specified interface.

<table>
<thead>
<tr>
<th>get_interface_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
### get_interfaces
This command returns a list of top-level interfaces in the system.

<table>
<thead>
<tr>
<th>get_interfaces</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_module_properties
This command returns the properties that you can manage for the top-level module.

<table>
<thead>
<tr>
<th>get_module_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_module_property
This command returns the value of a top-level system property.

<table>
<thead>
<tr>
<th>get_module_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

### get_parameter_properties
This command returns a list of properties that you can query on parameters. These properties can be queried on any parameter, such as parameters on instances, interfaces, instance interfaces, and connections.

<table>
<thead>
<tr>
<th>get_parameter_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
**get_port_properties**
This command returns a list of properties that you can query on ports.

<table>
<thead>
<tr>
<th>get_port_properties</th>
<th>Usage</th>
<th>get_port_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
<td>A list of port properties</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
<td></td>
</tr>
<tr>
<td>Example</td>
<td>get_port_properties</td>
<td></td>
</tr>
</tbody>
</table>

**get_project_properties**
This command returns a list of properties that you can query for the Quartus II project.

<table>
<thead>
<tr>
<th>get_project_properties</th>
<th>Usage</th>
<th>get_project_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String[]</td>
<td>A list of project properties</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
<td></td>
</tr>
<tr>
<td>Example</td>
<td>get_project_properties</td>
<td></td>
</tr>
</tbody>
</table>

**get_project_property**
This command returns the value of a Quartus II project property.

<table>
<thead>
<tr>
<th>get_project_property</th>
<th>Usage</th>
<th>get_project_property &lt;property&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>String</td>
<td>The value of the property</td>
</tr>
<tr>
<td>Arguments</td>
<td>property</td>
<td>The name of the project property; DEVICE_FAMILY</td>
</tr>
<tr>
<td>Example</td>
<td>get_project_property DEVICE_FAMILY</td>
<td></td>
</tr>
</tbody>
</table>

**load_system**
This command loads a Qsys system from a file, and uses the system as the current system for scripting commands.

<table>
<thead>
<tr>
<th>load_system</th>
<th>Usage</th>
<th>load_system &lt;file&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>None</td>
<td></td>
</tr>
<tr>
<td>Arguments</td>
<td>file</td>
<td>The path to a .qsys file</td>
</tr>
<tr>
<td>Example</td>
<td>load_system example.qsys</td>
<td></td>
</tr>
</tbody>
</table>
**lock_avalon_base_address**

This command prevents the memory-mapped base address from being changed for connections to an interface on an instance when the `auto_assign_base_addresses` or `auto_assign_system_base_addresses` commands are run.

<table>
<thead>
<tr>
<th>lock_avalon_base_address</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**preview_insert_avalon_streaming_adapters**

This command runs the adapter insertion for Avalon-ST connections, which adapt connections with mismatched configuration, such as mismatched data widths.

<table>
<thead>
<tr>
<th>preview_insert_avalon_streaming_adapters</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**remove_connection**

This command removes a connection from the system.

<table>
<thead>
<tr>
<th>remove_connection</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**remove_instance**

This command removes a child instance from the system.

<table>
<thead>
<tr>
<th>remove_instance</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
**remove_interface**

This command removes an exported top-level interface from the system.

```plaintext
remove_interface
Usage: remove_interface <interface>
Returns: None
Arguments:
- interface: The name of the exported top-level interface

Example: remove_interface clk_out
```

**save_system**

This command saves the current in-memory system to the named file. If the file is not specified, the system saves to the same file that was opened with the `load_system` command.

```plaintext
save_system
Usage: save_system [file]
Returns: None
Arguments:
- file (optional): If present, the path of the .qsys file to save

Example:
- save_system
- save_system example.qsys
```

**send_message**

This command sends a message to the user of the script. The message text is normally interpreted as HTML. You can use the `<b>` element to provide emphasis.

```plaintext
send_message
Usage: send_message <level> <message>
Returns: None
Arguments:
- level: The following message levels are supported:
  - ERROR—Provides an error message.
  - WARNING—Provides a warning message.
  - INFO—Provides an informational message.
  - PROGRESS—Provides a progress message.
  - DEBUG—Provides a debug message when debug mode is enabled.
- message: The text of the message.

Example: send_message ERROR "The system is down!"
```
set_connection_parameter_value

This command sets the parameter value for a connection.

<table>
<thead>
<tr>
<th>set_connection_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

set_instance_parameter_value

This command set the parameter value for a child instance. Derived parameters and SYSTEM_INFO parameters for the child instance can not be set with this command.

<table>
<thead>
<tr>
<th>set_instance_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

set_instance_property

This command sets the property value of a child instance. Most instance properties are read-only and can only be set by the instance itself. The primary use for this command is to update the ENABLED parameter, which includes or excludes a child instance when generating the system.

<table>
<thead>
<tr>
<th>set_instance_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
**set_interface_property**
This command sets the property value on an exported top-level interface. This command is used to set the EXPORT_OF property to specify which interface of a child instance is exported by the top-level interface.

**set_interface_property**
<table>
<thead>
<tr>
<th>Usage</th>
<th>set_interface_property &lt;interface&gt; &lt;property&gt; &lt;value&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td>interface</td>
</tr>
<tr>
<td></td>
<td>property</td>
</tr>
<tr>
<td></td>
<td>value</td>
</tr>
<tr>
<td>Example</td>
<td>set_interface_property clk_out EXPORT_OF clk.clk_out</td>
</tr>
</tbody>
</table>

**set_module_property**
This command sets the system property value, such as the name of the system using the NAME property.

**set_module_property**
<table>
<thead>
<tr>
<th>Usage</th>
<th>set_module_property &lt;property&gt; &lt;value&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td>property</td>
</tr>
<tr>
<td></td>
<td>value</td>
</tr>
<tr>
<td>Example</td>
<td>set_module_property NAME &quot;new_system_name&quot;</td>
</tr>
</tbody>
</table>

**set_project_property**
This command sets the project property value, such as the device family.

**set_project_property**
<table>
<thead>
<tr>
<th>Usage</th>
<th>set_project_property &lt;property&gt; &lt;value&gt;</th>
</tr>
</thead>
<tbody>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td>property</td>
</tr>
<tr>
<td></td>
<td>value</td>
</tr>
<tr>
<td>Example</td>
<td>set_project_property DEVICE_FAMILY &quot;Cyclone IV GX&quot;</td>
</tr>
</tbody>
</table>
set_validation_property

This command sets a property that affects how and when validation is run during system scripting. To disable system validation after each scripting command, set AUTOMATIC_VALIDATION to false.

| set_validation_property
| Usage | set_validation_property <property> <value>
| Returns | None
| Arguments | property | The name of the property
| value | The new property value.
| Example | set_validation_property AUTOMATIC_VALIDATION false

unlock_avalon_base_address

This command allows the memory-mapped base address to be changed for connections to an interface on an instance when the auto_assign_base_addresses or auto_assign_system_base_addresses commands are run.

| unlock_avalon_base_address
| Usage | unlock_avalon_base_address <instance.interface>
| Returns | None
| Arguments | instance.interface | The qualified name of the interface of an instance, in <instance>.<interface> format
| Example | unlock_avalon_base_address sram.s1

upgrade_sopc_system

This command loads the specified .sopc file, which then upgrades the file as a Qsys-compatible system. Some child instances and interconnect are replaced so that the system functions in Qsys. You must save the new Qsys-compatible system with the save_system command.

| upgrade_sopc_system
| Usage | upgrade_sopc_system <filename>
| Returns | None
| Arguments | filename | The path to the .sopc file being upgraded. The upgrade will move the .sopc file and related generation files to a backup directory.
| Example | upgrade_sopc_system old_system.sopc
validate_connection
This command validates the specified connection, and returns the during validation messages.

<table>
<thead>
<tr>
<th>validate_connection</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

validate_instance
This command validates the specified child instance, and returns the validation messages.

<table>
<thead>
<tr>
<th>validate_instance</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

validate_instance_interface
This command validates an interface on a child instance, and returns the validation messages.

<table>
<thead>
<tr>
<th>validate_instance_interface</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

validate_system
This command validates the system, and returns the validation messages.

<table>
<thead>
<tr>
<th>validate_system</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
## Document Revision History

Table 7–7 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>- Added AMBA APB support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added qsys-generate utility.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added VHDL BFM ID support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added “Creating Secure Systems (TrustZones)”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added “CMSIS Support for Qsys Systems With An HPS Component”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added VHDL language support options.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>- Added AMBA AXI4 support.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>- Added AMBA AXI3 support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added Preset Editor updates.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added command-line utilities, and scripts.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>- Added Synopsys VCS and VCS MX Simulation Shell Script.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added Cadence Incisive Enterprise (NCSIM) Simulation Shell Script.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added Using Instance Parameters and Example Hierarchical System Using Parameters.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>- Added simulation support in Verilog HDL and VHDL.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added testbench generation support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated simulation and file generation sections.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the *Quartus II Handbook*, refer to the [Quartus II Handbook Archive](#)

For more information about the benefits of using Qsys, refer to *Five Reasons to Switch from SOPC Builder to Qsys* on the Webcasts and Videos page of the Altera website.
This chapter describes the structure of Qsys components, with an emphasis on using the Qsys Component Editor to create the Hardware Component Description File (_hw.tcl), which describe and package components that you can use in a Qsys system.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

This chapter uses the Demo AXI Memory example available on the Qsys Design Examples page of the Altera® web site.

For Tcl command reference information, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

### Qsys Components

A Qsys component includes the following elements:

- Information about the component type, such as name, version, and author.
- HDL description of the component’s hardware, including SystemVerilog, Verilog HDL, or VHDL files
- Constraint files (.sdc or Quartus) that define the component for synthesis and simulation.
- A component’s interfaces, including I/O signals.
- The parameters that configure the operation of the component.

### Component Interfaces

Components can have any number of interfaces in any combination. For example, a component might provide both an Avalon-ST source port for high-throughput data, and a memory-mapped slave port for register configuration.
The following interfaces are available in Qsys:

- **Memory-Mapped**—For Avalon-MM, AXI masters and slaves, and APB masters and slaves that communicate using memory-mapped read and write commands.

- **Avalon Streaming (Avalon-ST)**—For point-to-point connections between Avalon-ST sources and sinks that stream data.

- **Interrupts**—For point-to-point connections between interrupt senders that generate interrupts, and interrupt receivers that service interrupts.

- **Clocks**—For point-to-point connections between clock sources and clock sinks.

- **Resets**—For point-to-point connections between reset sources and reset sinks.

- **Avalon Tri-State Conduit (Avalon-TC)**—For a tri-state conduit controller in the Qsys system connected to tri-state devices on the PCB.

- **Conduits**—For point-to-point connections between conduit interfaces. You can use the conduit interface type to define a custom collection of signals that do not fit into any other interface category.

### Component Structure and Types

Altera provides components automatically installed with the Quartus® II software. You can obtain a list of Qsys-compliant components provided by third-party IP developers on the Intellectual Property & Reference Designs web page by typing Qsys Certified in the Search box, and then selecting IP Core & Reference Designs. Components are also provided with Altera development kits, which are listed on the All Development Kits web page.

Every component is defined with a `<component_name>_hw.tcl` file, a text file written in the Tcl scripting language that describes the component to Qsys. When you design your own custom component, you can create the `_hw.tcl` file manually, or by using the Qsys Component Editor.

The Component Editor simplifies the process of creating `_hw.tcl` files by creating a file that you can edit outside of the Component Editor to add advanced procedures. When you edit a previously saved `_hw.tcl` file, Qsys automatically backs up the earlier version as `_hw.tcl~`.

You can move component files into a new directory, such as a network location, so that other users can use the component in their systems. The `_hw.tcl` file contains relative paths to the other files, so if you move an `_hw.tcl` file, you should also move all the HDL and other files associated with it.

### Static Components

You implement a static component using the same HDL files for all instances of the component. If the top-level HDL module is parameterized, instances may have unique behavior, depending on the parameter values.

### Generated Components

A generated component’s fileset callback allows an instance of the component to create unique HDL design files based on the instance’s parameter values.
For example, you can write a fileset callback to include a control and status interface based on the value of a parameter. The callback overcomes a limitation of HDL languages, which do not allow runtime parameters.

For information about defining fileset callback procedures, refer to “Controlling File Generation Dynamically with Parameters and a Fileset Callback” on page 8–27.

**Composed Components**

Composed components are subsystems constructed from instances of other components. You can use a composition callback to manage the subsystem in a composed component.

For information about creating a composed component, refer to “Creating a Composed Component or Subsystem” on page 8–28.

**Component File Organization**

A typical component uses the following directory structure. The names of the directories are not significant.

- `<component_directory>`
  - `<hdl>`—Contains the component HDL design files, for example `.v`, `.sv`, or `.vhd` files that contain the top-level module, along with any required constraint files.
  - `<component_name>_hw.tcl`—The component description file.
  - `<component_name>_sw.tcl`—The software driver configuration file. This file specifies the paths for the `.c` and `.h` files associated with the component, when required.
  - `<software>`—Contains software drivers or libraries related to the component. Altera recommends that the software directory be a subdirectory of the directory that contains the `_hw.tcl` file.

For information about writing a device driver or software package suitable for use with the Nios II processor, refer to the Hardware Abstraction Layer section of the Nios II Software Developer’s Handbook. The Nios II Software Build Tool Reference chapter of the Nios II Software Developer’s Handbook describes the commands you can use in the Tcl script.

**Component Versions**

Qsys systems support multiple versions of the same component within the same system; you can create and maintain multiple versions of the same component.

If you have multiple `_hw.tcl` files for components with the same NAME module properties and different VERSION module properties, both versions of the component are available.

If multiple versions of the component are available in the Qsys Component Library, you can add a specific version of a component by right-clicking the component, and then selecting Add version `<version_number>`.
Upgrading IP Components to the Latest Version

When you open a Qsys design, if Qsys detects IP components that require regeneration, the Upgrade IP Cores dialog box appears and allows you to upgrade outdated components. Components that you must upgrade in order to successfully compile your design appear in red. Status icons indicate whether a component is currently being regenerated, the component is encrypted, or that there is not enough information to determine the status of component. To upgrade a component, in the Upgrade IP Cores dialog box, select the component that you want to upgrade, and then click Upgrade. The Quartus II software maintains a list of all IP components associated with your design on the Components tab in the Project Navigator.

For more information about upgrading IP components, refer to Upgrade IP Components Dialog Box in Quartus II Help.

Life Cycle of a Component

When you define a component with the Qsys Component Editor, or a custom _hw.tcl file, you specify the information that Qsys requires to instantiate the component in a Qsys system and to generate the appropriate output files for synthesis and simulation. The following phases describe the process when working with components in Qsys:

- **Discovery**—During the discovery phase, Qsys reads the _hw.tcl file to identify information that appears in the Qsys Component Library, such as the component’s name, version, and documentation URLs. Each time you open Qsys, the tool searches for the following file types using the default search locations and entries in the IP Search Path:
  - _hw.tcl files—Each _hw.tcl file defines a single component.
  - IP Index (.ipx) files—Each .ipx file indexes a collection of available components, or a reference to other directories to search.

- **Static Component Definition**—During the static component definition phase, Qsys reads the _hw.tcl file to identify static parameter declarations, interface properties, interface signals, and HDL files that define the component. At this stage of the life cycle, the component interfaces might be only partially defined.

- **Parameterization**—During the parameterization phase, after an instance of the component is added to a Qsys system, the user of the component specifies parameters with the component’s parameter editor.

- **Validation**—During the validation phase, Qsys validates the values of each instance’s parameters against the allowed ranges specified for each parameter. You can use callback procedures that run during the validation phase to provide validation messages. For example, if there are dependencies between parameters where only certain combinations of values are supported, you can report errors for the unsupported values.
Elaboration—During the elaboration phase, Qsys queries the component for its interface information. Elaboration is triggered when an instance of a component is added to a system, when its parameters are changed, or when a system property changes. You can use callback procedures that run during the elaboration phase to dynamically control interfaces, signals, and HDL files based on the values of parameters. For example, interfaces defined with static declarations can be enabled or disabled during elaboration. When elaboration is complete, the component’s interfaces and design logic must be completely defined.

Composition—During the composition phase, a component can manipulate the instances in the component’s subsystem. The _hw.tcl file uses a callback procedure to provide parameterization and connectivity of subcomponents.

Generation—During the generation phase, Qsys generates synthesis or simulation files for each component in the system into the appropriate output directories, as well as any additional files that support associated tools.

Creating Qsys Components in the Component Editor

The Qsys Component Editor, accessed by clicking New Component in the Qsys Component Library, allows you to create and package a component for use in Qsys. When you use the Component Editor to define a component the Component Editor writes the information to the _hw.tcl file. The Component Editor allows you to perform the following tasks:

- Specify component’s identifying information, such as name, version, author, etc.
- Specify the SystemVerilog, Verilog HDL, or VHDL files, and constraint files that define the component for synthesis and simulation.
- Create an HDL template for a component by first defining its parameters, signals, and interfaces.
- Associate and define signals for a component’s interfaces.
- Set parameters on interfaces, which specify characteristics.
- Specify relationships between interfaces.
- Declare parameters that alter the component structure or functionality.

If the component is HDL-based, you must define the parameters and signals in the HDL file, and cannot add or remove them in the Component Editor. If you have not yet created the top-level HDL file, you declare the parameters and signals in the Component Editor, and they are then included in the HDL template file that Qsys creates.

In a Qsys system, the interfaces of a component are connected within the system, or exported as top-level signals from the system.

If you are creating the component using an existing HDL file, the order in which the tabs appear in the Component Editor reflects the recommended design flow for component development. You can use the Prev and Next buttons at the bottom of the Component Editor window to guide you through the tabs.
If the component is not based on an existing HDL file, enter the parameters, signals, and interfaces first, and then return to the Files tab to create the top-level HDL file template. When you click Finish, Qsys creates the component _hw.tcl file with the details provided on the Component Editor tabs.

After the component is saved, it is available in the Qsys Component Library.

If you require features in the component that are not supported by the Component Editor, such as callback procedures, you can use the Component Editor to create the _hw.tcl file, and then manually edit the file to complete the component definition. Subsequent sections of this chapter document the hw.tcl commands that are generated by the Component Editor, as well as some of the advanced features that you can add with your own hw.tcl commands.

For full syntax and details about all hw.tcl commands, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Saving a Component and Creating an _hw.tcl File

You save a component by clicking Finish in the Component Editor. The Component Editor saves the component to a file with the file name <component_name>_hw.tcl.

Altera recommends that you save _hw.tcl files and their associated files in an ip/<class-name> directory within your Quartus II project directory.

You can publish component information for use by software, such as a C compiler and a board support package (BSP) generator. For information on how to publish hardware component information for embedded software tools, refer to the Publishing Component Information to Embedded Software chapter in the Nios II Software Developer’s Handbook.

Editing a Component

In Qsys, you make changes to a component by right-clicking the component in the Component Library, and then clicking Edit. After making changes, click Finish to save the changes to the _hw.tcl file. You can open the _hw.tcl file in a text editor to view the hardware Tcl for the component. If you edit the _hw.tcl file to use advanced features, you cannot use the Component Editor to make further changes without over-writing the original file.

You cannot use the Component Editor to edit components installed with the Quartus II software, such as Altera-provided components. If you edit the HDL for a component and change the interface to the top-level module, you must edit the component to reflect the changes you made to the HDL.

For more information about the procedures for creating components in the Component Editor, refer to Creating Qsys Components in Quartus II Help.

Specifying Basic Component Information

The Component Type tab in the Component Editor allows you to specify the following information about the component:
■ **Name**—Specifies the name used in the `_hw.tcl` filename, as well as in the top-level module name when you create a synthesis wrapper file for a non HDL-based component.

■ **Display name**—Identifies the component in the parameter editor, which you use to configure and instance of the component, and also appears in the Component Library under Project and on the System Contents tab.

■ **Version**—Specifies the version number of the component.

■ **Group**—Represents the category of the component in the list of available components in the Component Library. You can select an existing group from the list, or define a new group by typing a name in the Group box. Separating entries in the Group box with a slash defines a subcategory. For example, if you type **Memories and Memory Controllers/On-Chip**, the component appears in the Component Library under the On-Chip group, which is a subcategory of the **Memories and Memory Controllers** group. If you save the Qsys design in the project directory, the component appears in the Component Library in the group you specified under Project. Alternatively, if you save the design in the Quartus II installation directory, the component appears in the specified group under Library.

■ **Description**—Allows you to describe the component.

■ **Created By**—Allows you to specify the author of the component.

■ **Icon**—Allows you to enter the relative path to an icon file (.gif, .jpg, or .png format) that represents the component and appears as the header in the parameter editor for the component. The default image is the Altera MegaCore function icon.

■ **Documentation**—Allows you to add links to documentation for the component, and appears when you right-click the component in the Component Library, and then select Details.

■ To specify an Internet file, begin your path with `http://`, for example: `http://mydomain.com/datasheets/my_memory_controller.html`.

■ To specify a file in the file system, begin your path with `file:///` for Linux, and `file:///` for Windows; for example (Windows): `file:///company_server/datasheets/my_memory_controller.pdf`.

The **Display name**, **Group**, **Description**, **Created By**, **Icon**, and **Documentation** entries are optional.
Figure 8–1 shows an example of the **Component Type** tab with the component information.

**Figure 8–1. Component Type Tab in the Component Editor**

When you use the Component Editor to create a component, it writes this basic component information in the _hw.tcl file. Example 8–1 shows the component hardware Tcl code related to the entries for the **Component Type** tab in Figure 8–1. The package require command specifies the ACDS version that Qsys uses to create the _hw.tcl file, and ensures compatibility with this version of the Qsys API in future ACDS releases.

The component defines its basic information with various module properties using the set_module_property command. For example, set_module_property NAME specifies the name of the component, while set_module_property VERSION allows you to specify the version of the component. When you apply a version to the _hw.tcl file, it allows the file to behave exactly the same way in future releases of the Quartus II software.

**Example 8–1. Hardware Tcl Commands Created from the Component Type Tab in the Component Editor**

```tcl
# request TCL package from ACDS 13.0
# package require -exact qsys 13.0
# module demo_axi_memory
# set_module_property DESCRIPTION "Demo AXI-3 memory with optional Avalon-ST port"
set_module_property NAME demo_axi_memory
set_module_property VERSION 1.0
set_module_property GROUP "My Components"
set_module_property AUTHOR Altera
set_module_property DISPLAY_NAME "Demo AXI Memory"
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the **Component Interface Tcl Reference** chapter in volume 1 of the *Quartus II Handbook*. 
Specifying Files for Synthesis and Simulation

The Files tab in the Component Editor allows you to specify files for synthesis and simulation. If you already have HDL code that describes the Qsys component that you want to create, you can specify the files on the Files tab, as described in “Specifying HDL Files for Synthesis” below.

If you have not yet created the HDL code that describes the component, but you have identified the signals and parameters that you want in the component, you can use the Files tab to create a top-level HDL template file. The Component Editor generates the appropriate hw.tcl commands to specify the files. You can also write your own hw.tcl file with the same commands, if you are not using the Component Editor.

A component uses filesets to specify the different sets of files that can be generated for an instance of the component. The supported fileset types include: QUARTUS_SYNTH, for synthesis and compilation in the Quartus II software, SIM_VERILOG, for Verilog HDL simulation, and SIM_VHDL, for VHDL simulation.

You add a fileset to a component with the add_fileset command. You then add files with the add_fileset_file command, which adds files to the most recently declared fileset. The add_fileset_property command allows you to add properties such as TOP_LEVEL, which specifies the top-level HDL module for the component.

You can populate a fileset with a a fixed list of files, add different files based on a parameter value, or even generate an HDL file with a custom HDL generator function outside of the _hw.tcl file.

Specifying HDL Files for Synthesis

You can add HDL files and other support files that should be included when this component is created to the list of Synthesis Files by clicking +, and then selecting the files in the Open dialog box.

A component must specify an HDL file as the top-level file, which contains the top-level module. The Synthesis Files list might also include supporting HDL files, such as timing constraints, or other files required to successfully synthesize and compile in the Quartus II software. The synthesis files for a component are copied to the generation output directory during Qsys system generation.
Figure 8–2 indicates the `demo_axi_memory.sv` file as the top-level file for the component in the Synthesis Files section on the Files tab.

Figure 8–2. Using HDL Files to Define a Component

Creating a New HDL File for Synthesis

If you do not already have an HDL implementation of the component, you can use the Component Editor to define the component, and then create a simple top-level synthesis file containing the signals and parameters for the component. You can then edit this HDL file to add the logic that directs the component's behavior.

To begin, you first specify the information about the component on the Parameters, Signals, and Interfaces tabs, as described in the upcoming sections of this chapter. Then, you return to the Files tab to create an HDL file by clicking Create Synthesis File from Signals. The Component Editor creates an HDL file from the specified parameters and signals.

Analyzing Synthesis Files

After the top-level HDL file is specified, click Analyze Synthesis Files to analyze the parameters and signals in the top-level, and then select the top-level module from the Top Level Module list. If there is a single module or entity in the HDL file, Qsys automatically populates the Top-level Module list.

Once analysis is complete and the top-level module is selected, the parameters and signals found in the top-level module are used as the parameters and signals for the component, and you can view them on the Parameters and Signals tabs. The Component Editor might report errors or warnings at this stage, because the signals and interfaces are not yet fully defined.

At this stage in the Component Editor flow, you cannot add or remove parameters or signals created from a specified HDL file without editing the HDL file itself.

Example 8–2 shows the component hardware Tcl code related to the entries for the Files Type tab in the Synthesis Files section in Figure 8–2.
The synthesis files are added to a fileset with the name `QUARTUS_SYNTH` and type `QUARTUS_SYNTH`. The top-level module is used to specify the `TOP_LEVEL` fileset property. Each synthesis file is individually added to the fileset. If the source files are saved in a different directory from the working directory where the `_hw.tcl` is located, you can use standard fixed or relative path notation to identify the file location for the `PATH` variable.

Example 8–2. Hardware Tcl Commands Created from the Files Tab - Synthesis Files Section of the Component Editor

```tcl
# file sets
#
add_fileset QUARTUS_SYNTH QUARTUS_SYNTH "" ""
set_fileset_property QUARTUS_SYNTH TOP_LEVEL demo_axi_memory
add_fileset_file demo_axi_memory.sv SYSTEM_VERILOG PATH demo_axi_memory.sv
add_fileset_file single_clk_ram.v VERILOG PATH single_clk_ram.v
```

For more information about `_hw.tcl` syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

**Naming HDL Signals for Automatic Interface and Type Recognition**

If you create the component's top-level HDL file before using the Component Editor, the Component Editor recognizes the interface and signal types based on the signal names in the source HDL file. This auto-recognition feature eliminates the task of manually assigning each interface and signal type in the Component Editor. To enable this auto-recognition feature, you must create signal names using the following naming convention:

```
<interface type prefix>_<interface name>_<signal type>
```

Specifying an interface name with `<interface name>` is optional if you have only one interface of each type in the component definition. For interfaces with only one signal, such as clock and reset inputs, the `<interface type prefix>` is also optional. When the Component Editor recognizes a valid prefix and signal type for a signal, it automatically assigns an interface and signal type to the signal based on the naming convention. If no interface name is specified for a signal, you can choose an interface name on the Interfaces tab in the Component Editor.
Table 8–1 lists valid values for an `<interface type prefix>`.

<table>
<thead>
<tr>
<th>Interface Prefix</th>
<th>Interface Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>asi</td>
<td>Avalon-ST sink (input)</td>
</tr>
<tr>
<td>aso</td>
<td>Avalon-ST source (output)</td>
</tr>
<tr>
<td>avm</td>
<td>Avalon-MM master</td>
</tr>
<tr>
<td>avs</td>
<td>Avalon-MM slave</td>
</tr>
<tr>
<td>axm</td>
<td>AXI master</td>
</tr>
<tr>
<td>axs</td>
<td>AXI slave</td>
</tr>
<tr>
<td>apm</td>
<td>APB master</td>
</tr>
<tr>
<td>aps</td>
<td>APB slave</td>
</tr>
<tr>
<td>coe</td>
<td>Conduit</td>
</tr>
<tr>
<td>csi</td>
<td>Clock Sink (input)</td>
</tr>
<tr>
<td>cso</td>
<td>Clock Source (output)</td>
</tr>
<tr>
<td>inr</td>
<td>Interrupt receiver</td>
</tr>
<tr>
<td>ins</td>
<td>Interrupt sender</td>
</tr>
<tr>
<td>ncm</td>
<td>Nios II custom instruction master</td>
</tr>
<tr>
<td>ncs</td>
<td>Nios II custom instruction slave</td>
</tr>
<tr>
<td>rsi</td>
<td>Reset sink (input)</td>
</tr>
<tr>
<td>rso</td>
<td>Reset source (output)</td>
</tr>
<tr>
<td>tcm</td>
<td>Avalon-TC master</td>
</tr>
<tr>
<td>tcs</td>
<td>Avalon-TC slave</td>
</tr>
</tbody>
</table>

Refer to the *Avalon Interface Specifications* or the *AMBA Protocol Specification* for the signal types available for each interface type.

### Specifying Files for Simulation

To support Qsys system generation for simulation, a component must specify the VHDL or Verilog simulation files. Simulation files are generated when a user adds the component to a Qsys system and chooses to generate Verilog or VHDL simulation files. In most cases, these files are the same as the synthesis files. If there are simulation-specific HDL files, you can use them in addition to, or in place of the synthesis files in the Component Editor. To use your synthesis files as your simulation files, on the Files tab, click **Copy From Synthesis Files** to copy the list of synthesis files to the **Verilog Simulation Files** or **VHDL Simulation Files** lists.

You specify the simulation files in a similar way as the synthesis files with the fileset commands. Example 8–3 shows `SIM_VERILOG` and `SIM_VHDL` filesets for Verilog and VHDL simulation output files. In this example, the same Verilog files are used for both Verilog and VHDL outputs, and there is one additional System Verilog file added.

This method works for designers of Verilog IP to support users who want to generate a VHDL top-level simulation file when they have a mixed-language simulation tool and license that can read the Verilog output for the component.
Figure 8–3 shows the files specified for simulation on the Files tab.

Example 8–3 shows the component hardware Tcl code related to the entries for the Files Type tab in the Simulation Files section.

```tcl
add_fileset SIM_VERILOG SIM_VERILOG "" ""
set_fileset_property SIM_VERILOG TOP_LEVEL demo_axi_memory
add_fileset_file single_clk_ram.v VERILOG PATH single_clk_ram.v
add_fileset_file verbosity_pkg.sv SYSTEM_VERILOG PATH verification_lib/verbosity_pkg.sv
add_fileset_file demo_axi_memory.sv SYSTEM_VERILOG PATH demo_axi_memory.sv
add_fileset SIM_VHDL SIM_VHDL "" ""
set_fileset_property SIM_VHDL TOP_LEVEL demo_axi_memory
set_fileset_property SIM_VHDL ENABLE_RELATIVE_INCLUDE_PATHS false
add_fileset_file demo_axi_memory.sv SYSTEM_VERILOG PATH demo_axi_memory.sv
add_fileset_file single_clk_ram.v VERILOG PATH single_clk_ram.v
add_fileset_file verbosity_pkg.sv SYSTEM_VERILOG PATH verification_lib/verbosity_pkg.sv
```

For more information about _hw.tcl_ syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

**Specifying Component Parameters**

Components can include parameterized HDL, which allows users of the component flexibility in meeting their system requirements. For example, a component might have a configurable memory size or data width, where one HDL implementation can be used in many different systems, each with unique parameters values.
The **Parameters** tab in the Component Editor allows you specify the parameters that are used to configure instances of the component in a Qsys system. You can specify various properties for each parameter that describe how the parameter is displayed and used. You can also specify a range of allowed values that are checked during the Validation phase. The **Parameters** table displays the HDL parameters that are declared in the top-level HDL module. If you have not yet created the top-level HDL file, the parameters that you create on the **Parameters** tab are included in the top-level synthesis file template created from the **Files** tab.

When the component includes HDL files, the parameters match those defined in the top-level module, and you cannot be add or remove them on the **Parameters** tab. To add or remove the parameters, edit your HDL source, and then re-analyze the file.

If you used the Component Editor to create a top-level template HDL file for synthesis, you can remove the newly-created file from the **Synthesis Files** list on the **Files** tab, make your parameter changes, and then re-analyze the top-level synthesis file.

You can use the **Parameters** table to specify the following information about each parameter:

- **Name**—Allows you to name the parameter.
- **Default Value**—Sets the default value used in new instances of the component.
- **Editable**—Specifies whether or not the user can edit the parameter value.
- **Type**—Defines the parameter type as string, integer, boolean, std_logic, logic vector, natural, or positive.
- **Group**—Allows you to group parameters in parameter editor.
- **Tooltip**—Allows you to add a description of the parameter that appears when the user of the component points to the parameter in the parameter editor.
On the **Parameters** tab, you can click **Preview the GUI** at any time to see how the declared parameters appear in the parameter editor. Figure 8–4 shows parameters with their default values defined, with checks in the **Editable** column indicating that users of this component are allowed to modify the parameter value. Editable parameters cannot contain computed expressions. You can group parameters under a common heading or section in the parameter editor with the **Group** column, and a tooltip helps users of the component understand the function of the parameter. Various parameter properties allow you to customize the component’s parameter editor, such as using radio buttons for parameter selections, or displaying an image.

**Figure 8–4. Parameters Tab in the Components Editor**

If a parameter \(<n>\) defines the width of a signal, the signal width must follow the format: \(<n-1>:0\).
Example 8–4 shows the component hardware Tcl code related to the entries for the Parameters tab in Figure 8–4. In this example, the first add_parameter command includes commonly-specified properties. The set_parameter_property command specifies each property individually. The Tooltip column on the Parameters tab maps to the DESCRIPTION property, and there is an additional unused UNITS property created in the code. The HDL_PARAMETER property specifies that the value of the parameter is specified in the HDL instance wrapper when creating instances of the component. The Group column in the Parameters tab maps to the display items section with the add_display_item commands.

**Example 8–4. Hardware Tcl Commands Created from the Parameters Tab in the Component Editor**

```tcl
# parameters
#
add_parameter AXI_ID_W INTEGER 4 "Width of ID fields"
set_parameter_property AXI_ID_W DEFAULT_VALUE 4
set_parameter_property AXI_ID_W DISPLAY_NAME AXI_ID_W
set_parameter_property AXI_ID_W TYPE INTEGER
set_parameter_property AXI_ID_W UNITS None
set_parameter_property AXI_ID_W DESCRIPTION "Width of ID fields"
add_parameter AXI_ADDRESS_W INTEGER 12
set_parameter_property AXI_ADDRESS_W DEFAULT_VALUE 12

add_parameter AXI_DATA_W INTEGER 32

# display items
#
add_display_item "AXI Port Widths" AXI_ID_W PARAMETER ""
```

If an AXI slave’s ID bit width is smaller than the formula, the AXI slave response might not reach AXI masters. The formula of an AXI slave ID bit width is calculated as follows:

\[
\text{maximum_master_id_width_in_the_interconnect} + \log_2(\text{number_of_masters_in_the_same_interconnect})
\]

For example, if an AXI slave connects to three AXI masters and the maximum AXI master ID length of the three masters is 5 bits, then the AXI slave ID is 7 bits, then:

\[
5 \text{ bits} + 2 \text{ bits (log}_2(3 \text{ masters})) = 7
\]

For more information about hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

**Allowed Ranges Parameter Property**

In a component’s hw.tcl file, you can specify valid ranges for parameters. In Qsys, validation checks each parameter value against the ALLOWED_RANGES property. If the values specified are outside of the allowed ranges, Qsys displays an error message. Specifying choices for the allowed values enables users of the component to choose the parameter value from a drop-down list or radio button in the parameter editor GUI instead of entering a value.
The \texttt{ALLOWED\_RANGES} property is a list of valid ranges, where each range is a single value, or a range of values defined by a start and end value. Table 8–2 shows examples of the \texttt{ALLOWED\_RANGES} property. For a parameter editor GUI example of the \texttt{ALLOWED\_RANGES} property, refer to "Parameter Editor Illustrating Parameter Declarations" on page 8–19. Refer to the “Declaring Parameters in hw.tcl, including Allowed Ranges and a Derived Parameter” on page 8–19 for an \texttt{hw.tcl} code example that uses the \texttt{ALLOWED\_RANGES} property.

### Table 8–2. ALLOWED\_RANGES Property

<table>
<thead>
<tr>
<th>ALLOWED_RANGES</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>{a b c}</td>
<td>a or b or c</td>
</tr>
<tr>
<td>{&quot;No Control&quot; &quot;Single Control&quot; &quot;Dual Controls&quot;}</td>
<td>Unique string values. Quotation marks are required if the strings include spaces</td>
</tr>
<tr>
<td>{1 2 4 8 16}</td>
<td>1, 2, 4, 8, or 16</td>
</tr>
<tr>
<td>{1:3}</td>
<td>1 through 3, inclusive</td>
</tr>
<tr>
<td>{1 2 3 7:10}</td>
<td>1, 2, 3, or 7 through 10 inclusive</td>
</tr>
</tbody>
</table>

#### Types of Parameters

Qsys uses the following parameter types: user parameters, system information parameters, and derived parameters.

##### User Parameters

User parameters are parameters that users of a component can control, and appear in the parameter editor for instances of the component. User parameters map directly to parameters in the component HDL, such as \texttt{AXI\_DATA\_W} in Example 8–5 on page 8–19, or control commands in an elaboration callback, such as \texttt{ENABLE\_STREAM\_OUTPUT}, also in Example 8–5.

##### System Information Parameters

A \texttt{SYSTEM\_INFO} parameter is a parameter whose value is set automatically by the Qsys system. When you define a \texttt{SYSTEM\_INFO} parameter, you provide an information type, and additional arguments. For example, you can configure a parameter to store the clock frequency driving a clock input for your component. To do this, define the parameter as \texttt{SYSTEM\_INFO} of type, \texttt{CLOCK\_RATE}:

```
set_parameter_property <param> SYSTEM\_INFO CLOCK\_RATE
```

You then set the name of the clock interface as the \texttt{SYSTEM\_INFO} argument:

```
set_parameter_property <param> SYSTEM\_INFO\_ARG <clfname>
```

##### Derived Parameters

Derived parameter values are calculated from other parameters during the Elaboration phase, and are specified in the \texttt{hw.tcl} file with the \texttt{DERIVED} property. For example, you can derive a clock period parameter from a data rate parameter. Derived parameters are sometimes used to perform operations that are difficult to perform in HDL, such as using logarithmic functions to determine the number of address bits that a component requires.
Parameterized Parameter Widths

Qsys allows a `std_logic_vector` parameter to have a width that is defined by another parameter, similar to derived parameters. The width can be a constant or the name of another parameter.

Declaring Parameters with Custom `hw.tcl` Commands

Example 8–5 illustrates a custom `hw.tcl` file, with more advanced parameter commands than those generated when you specify parameters in the Component Editor. Commands include the `ALLOWED_RANGES` property described in “Allowed Ranges Parameter Property” on page 8–16 to provide a range of values for the `AXI_ADDRESS_W` (Address Width) parameter, and a list of parameter values for the `AXI_DATA_W` (Data Width) parameter. Example 8–5 also shows the parameter `AXI_NUMBYTES` (Data width in bytes) parameter; that uses the `DERIVED` property, as described in “Derived Parameters” on page 8–17. In addition, these commands illustrate the use of the `GROUP` property, which groups some parameters under a heading in the parameter editor GUI. You use the `ENABLE_STREAM_OUTPUT_GROUP` (Include Avalon streaming source port) parameter to enable or disable the optional Avalon-ST interface in this design, and is displayed as a check box in the parameter editor GUI because the parameter is of type `BOOLEAN`. Refer to Figure 8–5 to see the parameter editor GUI resulting from these `hw.tcl` commands.

Example 8–5 illustrates parameter declaration statements and includes a parameter whose value is derived during the Elaboration phase based on another parameter, instead of being assigned to a specific value. `AXI_NUMBYTES` describes the number of bytes in a word of data. Qsys calculates the `AXI_NUMBYTES` parameter from the `DATA_WIDTH` parameter by dividing by 8. The `.hw.tcl` code defines the `AXI_NUMBYTES` parameter as a derived parameter, since its value is calculated in an elaboration callback procedure.

For information about creating an elaboration callback procedure, refer to “Controlling Interfaces Dynamically with an Elaboration Callback” on page 8–26.
In the parameter editor, the AXI_NUMBYTES parameter value is not editable, because its value is based on another parameter value.

**Example 8–5. Declaring Parameters in hw.tcl, including Allowed Ranges and a Derived Parameter**

```tcl
add_parameter AXI_ADDRESS_W INTEGER 12
set_parameter_property AXI_ADDRESS_W DISPLAY_NAME "AXI Slave Address Width"
set_parameter_property AXI_ADDRESS_W DESCRIPTION "Address width."
set_parameter_property AXI_ADDRESS_W UNITS bits
set_parameter_property AXI_ADDRESS_W ALLOWED_RANGES 4:16
set_parameter_property AXI_ADDRESS_W HDL_PARAMETER true
set_parameter_property AXI_ADDRESS_W GROUP "AXI Port Widths"

add_parameter AXI_DATA_W INTEGER 32
set_parameter_property AXI_DATA_W DISPLAY_NAME "Data Width"
set_parameter_property AXI_DATA_W DESCRIPTION "Width of data buses."
set_parameter_property AXI_DATA_W UNITS bits
set_parameter_property AXI_DATA_W ALLOWED_RANGES {8 16 32 64 128 256 512 1024}
set_parameter_property AXI_DATA_W HDL_PARAMETER true
set_parameter_property AXI_DATA_W GROUP "AXI Port Widths"

add_parameter AXI_NUMBYTES INTEGER 4
set_parameter_property AXI_NUMBYTES DERIVED true
set_parameter_property AXI_NUMBYTES DISPLAY_NAME "Data Width in bytes; Data Width/8"
set_parameter_property AXI_NUMBYTES DESCRIPTION "Number of bytes in one word"
set_parameter_property AXI_NUMBYTES UNITS bytes
set_parameter_property AXI_NUMBYTES HDL_PARAMETER true
set_parameter_property AXI_NUMBYTES GROUP "AXI Port Widths"

add_parameter ENABLE_STREAM_OUTPUT BOOLEAN true
set_parameter_property ENABLE_STREAM_OUTPUT DISPLAY_NAME "Include Avalon Streaming Source Port"
set_parameter_property ENABLE_STREAM_OUTPUT DESCRIPTION "Include optional Avalon-ST source (default), or hide the interface"
set_parameter_property ENABLE_STREAM_OUTPUT GROUP "Streaming Port Control"
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Figure 8–5 shows how the parameter editor GUI generated from Example 8–5.

**Figure 8–5. Parameter Editor Illustrating Parameter Declarations**
Validating Parameter Values with a Validation Callback

You can use a validation callback procedure to validate parameter values with more complex validation operations than the ALLOWED_RANGES property allows. You define a validation callback by setting the VALIDATION_CALLBACK module property to the name of the Tcl callback procedure that runs during the validation phase. In the validation callback procedure, the current parameter values is queried, and warnings or errors are reported about the component’s configuration.

In the Demo AXI Memory example in Example 8–6, if the optional Avalon streaming interface is enabled, then the control registers must be wide enough to hold an AXI RAM address, so the designer can add an error message to ensure that the user enters allowable parameter values.

Example 8–6. set_module_property VALIDATION_CALLBACK

```tcl
set_module_property VALIDATION_CALLBACK validate
proc validate {} {
  if {
    [get_parameter_value ENABLE_STREAM_OUTPUT ] &&
    ([get_parameter_value AXI_ADDRESS_W] >
     [get_parameter_value AV_DATA_W])
  }
  {send_message error "If the optional Avalon streaming port is enabled, the AXI Data Width must be equal to or greater than the Avalon control port Address Width"}
}
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Specifying Interface and Signal Types

The Signals tab in the Components Editor allows you to specify the interface and signal type of each signal in the component. When you add HDL files to the Synthesis Files table on the Files tab, and then click Analyze Synthesis Files, the signals on the top-level module appear on the Signals tab.

If you have not yet created your top-level HDL file, you can click Add Signal to specify each top-level signal in the component. For each signal that you add, you must provide the appropriate values in the Name, Interface, Signal Type, Width, and Direction columns. You can use the error and warning messages at the bottom of the window to guide your selections. You can edit the signal name by double-clicking the Name column, and then typing the new name.

After you have analyzed the component’s top-level HDL file on the Files tab, you cannot add or remove signals or change the signal names on the Signals tab. To change the signals, edit your HDL source, and then re-analyze the file.

If you used the Component Editor to create a top-level template HDL file for synthesis, you can remove the newly-created file from the Synthesis Files list on the Files tab, make your signal changes, and then re-analyze the top-level synthesis file.
The **Interface** column allows you assign a signal to an interface. Each signal must belong to an interface and be assigned a legal signal type for that interface. To create a new interface of a specific type, select `new <interface type>` from the list; this new interface then become available in the list for subsequent signal assignments. You can highlight all of the signals in an interface and then select an Interface from the list to apply the Interface name to each signal in the interface.

You edit the interface name on the **Interface** tab; you cannot edit the interface name on the **Signals** tab.

**Figure 8–6** shows the `altera_axi_slave` selection available for the `axs_awaddr` signal.

**Figure 8–6. Signals Tab in the Components Editor**
Example 8–7 shows the component hardware Tcl code related to the entries for the Signals tab in Figure 8–6.

Example 8–7. Hardware Tcl Commands Created from the Signals Tab

```tcl
add_interface_port altera_axi_slave axs_awid awid Input AXI_ID_W
add_interface_port altera_axi_slave axs_awaddr awaddr Input AXI_ADDRESS_W
add_interface_port altera_axi_slave axs_awslen awlen Input 4
add_interface_port altera_axi_slave axs_awsizew awsize Input 3
add_interface_port altera_axi_slave axs_awburst awburst Input 2
add_interface_port altera_axi_slave axs_awlock awlock Input 2
add_interface_port altera_axi_slave axs_awcache awcache Input 4
add_interface_port altera_axi_slave axs_awprot awprot Input 3
add_interface_port altera_axi_slave axs_awvalid awvalid Input 1
add_interface_port altera_axi_slave axs_awready awready Output 1
add_interface_port altera_axi_slave axs_wid wid Input AXI_ID_W
add_interface_port altera_axi_slave axs_wdata wdata Input AXI_DATA_W
add_interface_port altera_axi_slave axs_wstrb wstrb Input AXI_NUMBYTES
add_interface_port altera_axi_slave axs_wlast wlast Input 1
add_interface_port altera_axi_slave axs_wready wready Output 1
add_interface_port altera_axi_slave axs_bid bid Output AXI_ID_W
add_interface_port altera_axi_slave axs_bresp bresp Output 2
add_interface_port altera_axi_slave axs_bvalid bvalid Output 1
add_interface_port altera_axi_slave axs_bready bready Input 1
add_interface_port altera_axi_slave axs_arid arid Input AXI_ID_W
add_interface_port altera_axi_slave axs_araddr araddr Input AXI_ADDRESS_W
add_interface_port altera_axi_slave axs_arlen arlen Input 4
add_interface_port altera_axi_slave axs_arsize arsize Input 3
add_interface_port altera_axi_slave axs_arburst arburst Input 2
add_interface_port altera_axi_slave axs_arlock arlock Input 2
add_interface_port altera_axi_slave axs_arcache arcache Input 4
add_interface_port altera_axi_slave axs_arprot arprot Input 3
add_interface_port altera_axi_slave axs_arvalid arvalid Input 1
add_interface_port altera_axi_slave axs_arready arready Output 1
add_interface_port altera_axi_slave axs_rid rid Output AXI_ID_W
add_interface_port altera_axi_slave axs_rdata rdata Output AXI_DATA_W
add_interface_port altera_axi_slave axs_rlast rlast Output 1
add_interface_port altera_axi_slave axs_rvalid rvalid Output 1
add_interface_port altera_axi_slave axs_rready rready Input 1
add_interface_port altera_axi_slave axs_rresp rresp Output 2
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Adding Interfaces and Managing Interface Settings

The Interfaces tab in the Component Editor allows you to manage settings for each interface of the component. The interface name appears on the Signals tab, and in the Qsys System Contents tab when the component is added to a system.

You can configure the type and properties of each interface. Some interfaces display waveforms that illustrate the timing for the interface. If you update timing parameters, the waveforms automatically update.

You add additional interfaces by clicking Add Interface, and then you must specify the signals for the added interface on the Signals tab. You can remove interfaces that have no assigned signals by clicking Remove Interfaces With No Signals.
Figure 8–7 shows the **Avalon Streaming Source** interface, named *streaming*.

**Figure 8–7. Interfaces Tab in the Components Editor**
Example 8–8 shows the component hardware Tcl code related to the entries for the Interfaces tab in Figure 8–7. In this example, each interface is created with the `add_interface` command. You specify the properties of each interface with the `set_interface_property` command. The interface's signals are specified with the `add_interface_port` command.

Example 8–8. Hardware Tcl Commands Created from the Signals and Interfaces Tabs

```
# # connection point clock
#
add_interface clock clock end
set_interface_property clock clockRate 0
set_interface_property clock ENABLED true

add_interface_port clock clk clk Input 1

# # connection point reset
#
add_interface reset reset end
set_interface_property reset associatedClock clock
set_interface_property reset synchronousEdges DEASSERT
set_interface_property reset ENABLED true

add_interface_port reset reset_n reset_n Input 1

# # connection point streaming
#
add_interface streaming avalon_streaming start
set_interface_property streaming associatedClock clock
set_interface_property streaming associatedReset reset
set_interface_property streaming dataBitsPerSymbol 8
set_interface_property streaming errorDescriptor ""
set_interface_property streaming firstSymbolInHighOrderBits true
set_interface_property streaming maxChannel 0
set_interface_property streaming readyLatency 0
set_interface_property streaming ENABLED true

add_interface_port streaming aso_data data Output 8
add_interface_port streaming aso_valid valid Output 1
add_interface_port streaming aso_ready ready Input 1

# # connection point slave
#
add_interface slave axi end
set_interface_property slave associatedClock clock
set_interface_property slave associatedReset reset
set_interface_property slave readAcceptanceCapability 1
set_interface_property slave writeAcceptanceCapability 1
set_interface_property slave combinedAcceptanceCapability 1
set_interface_property slave readDataReorderingDepth 1
set_interface_property slave ENABLED true

add_interface_port slave axs_awid awid Input AXI_ID_W
...
add_interface_port slave axs_rresp rresp Output 2
```
Qsys refers to AXI component parameters to build AXI interconnect. If these parameter settings are incompatible with RTL component's behavior, Qsys interconnect and transactions might not work correctly. To prevent unexpected interconnect behavior, you must set the AXI component parameters described in Table 8–3 below.

Table 8–3. AXI Master and Slave Parameters

<table>
<thead>
<tr>
<th>AXI Master Parameters</th>
<th>AXI Slave Parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>readIssuingCapability</td>
<td>readAcceptanceCapability</td>
</tr>
<tr>
<td>writeIssuingCapability</td>
<td>writeAcceptanceCapability</td>
</tr>
<tr>
<td>combinedIssuingCapability</td>
<td>combinedAcceptanceCapability</td>
</tr>
<tr>
<td>readDataReorderingDepth</td>
<td></td>
</tr>
</tbody>
</table>

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Creating Custom .hw_tcl Interface Settings and Properties

Example 8–9 shows clock, reset, AXI slave, and Avalon streaming interfaces using variables for the interface names that make the file easier to read and update. The interface declaration statement includes the name, type, and direction of the interface, as well as the associated clock and reset interfaces. Also in Example 8–9, some of the AXI memory signals use parameters to specify their width.

Example 8–9. Interface Signals and Properties

```tcl
set CLOCK_INTERFACE  "clk"
add_interface $CLOCK_INTERFACE clock end
add_interface_port $CLOCK_INTERFACE clk clk Input 1

set RESET_INTERFACE  "reset"
add_interface $RESET_INTERFACE reset end
set_interface_property $RESET_INTERFACE associatedClock clk
set_interface_property $RESET_INTERFACE synchronousEdges DEASSERT
add_interface_port reset reset_n reset_n Input 1

set SLAVE_INTERFACE  "slave"
add_interface $SLAVE_INTERFACE axi end
set_interface_property $SLAVE_INTERFACE associatedClock "clk"
set_interface_property $SLAVE_INTERFACE associatedReset "reset"
set_interface_property $SLAVE_INTERFACE readAcceptanceCapability 1
... add_interface_port $SLAVE_INTERFACE axs_wdata wdata Input AXI_DATA_W
add_interface_port $SLAVE_INTERFACE axs_wstrb wstrb Input AXI_NUMBYTES
add_interface_port $SLAVE_INTERFACE axs_wlast wlast Input 1
... set STREAMING_INTERFACE "streaming"
add_interface $STREAMING_INTERFACE avalon_streaming start
set_interface_property $STREAMING_INTERFACE associatedClock "clk"
... add_interface_port $STREAMING_INTERFACE aso_data data Output 8
add_interface_port $STREAMING_INTERFACE aso_valid valid Output 1
add_interface_port $STREAMING_INTERFACE aso_ready ready Input 1
```
For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Controlling Interfaces Dynamically with an Elaboration Callback

You can allow user parameters to dynamically control your component's behavior with a an elaboration callback procedure during the elaboration phase. Using an elaboration callback allows you to change interface properties, remove interfaces, or add new interfaces as a function of parameter values. You define an elaboration callback by setting the module property ELABORATION_CALLBACK to the name of the Tcl callback procedure that runs during the elaboration phase. In the callback procedure, you can query the parameter values of the component instance, and then change the interfaces accordingly.

Example 8–10 shows an Avalon-ST source interface included in an instance of the component, based on the ENABLE_STREAM_OUTPUT parameter. The streaming interface was defined previously in the static portion of the HDL file, shown in Example 8–9 on page 8–25, and the ENABLE_STREAM_OUTPUT parameter was defined previously in Example 8–6 on page 8–20.

Example 8–10. Elaboration Callback Example

```tcl
set_module_property ELABORATION_CALLBACK         elaborate
proc elaborate {} {
    # Optionally disable the Avalon-ST data output
    if { [ get_parameter_value ENABLE_STREAM_OUTPUT ] == "false" } {
        set_port_property aso_data    termination true
        set_port_property aso_valid   termination true
        set_port_property aso_ready   termination true
        set_port_property aso_ready   termination_value 0
    }
    # Calculate the Data Bus Width in bytes
    set bytewidth_var [ expr [ get_parameter_value AXI_DATA_W] /8 ]
    set_parameter_value AXI_NUMBYTES $bytewidth_var
}
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.
Controlling File Generation Dynamically with Parameters and a Fileset Callback

You can use a fileset callback to control which files are created in the output directories during the generation phase based on parameter values, instead of providing a fixed list of files, as was described “Specifying Files for Synthesis and Simulation” on page 8–9. In a callback procedure, you can query the values of the parameters and use them to generate the appropriate files. To define a fileset callback, you specify a callback procedure name as an argument in the add_fileset command. You can use the same fileset callback procedure for all of the filesets, or create separate procedures for synthesis and simulation, or Verilog and VHDL.

Example 8–11 shows a complete fileset callback example using parameters to control filesets in two different ways. The RAM_VERSION parameter chooses between two different source files to control the implementation of a RAM block. For the top-level source file, a custom Tcl routine generates HDL that optionally includes control and status registers, depending on the value of the CSR_ENABLED parameter.

During the generation phase, Qsys creates a a top-level Qsys system HDL wrapper module to instantiate the component top-level module, and applies the component’s parameters, for any parameter whose parameter property HDL_PARAMETER is set to true.
Creating a Composed Component or Subsystem

A composed component is a subsystem containing instances of other components. Unlike an HDL-based component, a composed component’s HDL is created by generating HDL for the components in the subsystem, in addition to the Qsys interconnect to connect the subsystem instances.

You can add child instances in either the static part of the _hw.tcl file, or in a composition callback.

---

Example 8–11. Fileset Callback Example

```
#Create synthesis fileset with fileset_callback and set top level
add_fileset my_synthesis_fileset QUARTUS_SYNTH fileset_callback
set_fileset_property my_synthesis_fileset TOP_LEVEL demo_axi_memory

#Create Verilog simulation fileset with same fileset_callback and set top level
add_fileset my_verilog_sim_fileset SIM_VERILOG fileset_callback
set_fileset_property my_verilog_sim_fileset TOP_LEVEL demo_axi_memory
#Add extra file needed for simulation only
add_fileset_file verbosity_pkg.sv SYSTEM_VERILOG PATH verification_lib/verbosity_pkg.sv

#Create VHDL simulation fileset (with Verilog files; for mixed-language VHDL simulation)
add_fileset my_vhdl_sim_fileset SIM_VHDL fileset_callback
set_fileset_property my_vhdl_sim_fileset TOP_LEVEL demo_axi_memory
add_fileset_file verbosity_pkg.sv SYSTEM_VERILOG PATH verification_lib/verbosity_pkg.sv

#Define parameters required for fileset_callback
add_parameter RAM_VERSION INTEGER 1
set_parameter_property RAM_VERSION ALLOWED_RANGES {1 2}
set_parameter_property RAM_VERSION HDL_PARAMETER false
add_parameter CSR_ENABLED BOOLEAN enable
set_parameter_property CSR_ENABLED HDL_PARAMETER false

#Create Tcl callback procedure to add appropriate files to filesets based on parameters
proc fileset_callback { entityName } {
    send_message INFO "Generating top-level entity $entityName"
    set ram [get_parameter_value RAM_VERSION]
    set csr_enabled [get_parameter_value CSR_ENABLED]
    send_message INFO "Generating memory implementation based on RAM_VERSION $ram"
    if {$ram == 1} {
        add_fileset_file single_clk_ram1.v VERILOG PATH single_clk_ram1.v
    } else {
        add_fileset_file single_clk_ram2.v VERILOG PATH single_clk_ram2.v
    }
    send_message INFO "Generating top-level file for CSR_ENABLED $csr_enabled"
    generate_my_custom_hdl $csr_enabled demo_axi_memory_gen.sv
    add_fileset_file demo_axi_memory_gen.sv VERILOG PATH demo_axi_memory_gen.sv
}
```

For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.
■ **Static commands**—In the main static part of the _hw.tcl file, you can use composition commands such as add_instance, set_instance_parameter_value, and add_connection to create and parameterize subcomponent instances.

■ **Composition callback**—You can use a composition callback to modify the subsystem that is created in the static program, or to define the subsystem. With a composition callback, you can also instantiate and parameterize subcomponents as a function of the composed component’s parameter values. You define a composition callback by setting the COMPOSITION_CALLBACK module property to the name of the composition callback procedures.

A composition callback replaces the validation and elaboration phases. HDL for the subsystem is generated by generating all of the subcomponents and the top-level that combines them.

To connect instances of your component, you must define the component’s interfaces. Unlike an HDL-based component, a composed component does not directly specify the signals that are exported. Instead, interfaces of submodules are chosen as the external interface, and each internal interface’s ports are connected through the exported interface.

Exporting an interface means that you are making the interface visible from the outside of your component, instead of connecting it internally. You can set the EXPORT_OF property of the externally visible interface from the main program or the composition callback, to indicate that it is an exported view of the submodule’s interface.

Exporting an interface is different than defining an interface. An exported interface is an exact copy of the subcomponent’s interface, and you are not allowed to change properties on the exported interface. For example, if the internal interface is a 32-bit or 64-bit master without bursting, then the exported interface is the same. An interface on a subcomponent cannot be exported and also connected within the subsystem.

When you create an exported interface, the properties of the exported interface are copied from the subcomponent’s interface without modification. Ports are copied from the subcomponent’s interface with only one modification; the names of the exported ports on the composed component are chosen to ensure that they are unique.
Figure 8–8 shows a block diagram for the composed component in Example 8–12.

**Example 8–12** provides an example of a composed `_hw.tcl` file that instantiates two subcomponents. It connects the components, and also connects the clocks and resets. Note that clock and reset bridge components are required to allow both subcomponents to see common clock and reset inputs.

```tcl
package require -exact qsys 13.0
set_module_property name my_component

add_instance clk altera_clock_bridge
add_instance phy my_phy_microcore
add_interface clk clock end
add_instance reset altera_reset_bridge
set_interface_property clk EXPORT_OF clk.in_clk
add_instance regs my_regs_microcore
set_instance_property_value reset synchronous_edges deassert

add_interface reset reset end
set_interface_property reset EXPORT_OF reset.in_reset

add_interface pins conduit end
set_interface_property pins EXPORT_OF phy.pins

add_interface slave avalon slave
set_interface_property slave EXPORT_OF regs.slave

add_connection clk.out_clk reset.clk
add_connection clk.out_clk phy.clk
add_connection reset.out_reset phy.clk_reset
add_connection clk.out_clk regs.clk
add_connection reset.out_reset regs.reset
add_connection phy.output regs.input
add_connection regs.output phy.input
```

---

**Figure 8–8. Top-Level of a Composed Component**

![Block diagram of the composed component showing the connections between `my_component`, `altera_clock_bridge`, `my_phy_microcore`, `altera_reset_bridge`, `my_regs_microcore`, and `pins`. The diagram illustrates the flow of signals and the connections between the components.](image-url)
For more information about _hw.tcl syntax and advanced features that you can add to a component definition, refer to the Component Interface Tcl Reference chapter in volume 1 of the Quartus II Handbook.

Adding Component Instances to a Static or Generated Component

Static or generated components can create instances of other components. You can add child instances in either the static part of the _hw.tcl file, or in an elaboration callback. The following applies when:

- **Static commands**—In the main static part of the _hw.tcl file, you can use instance commands such as add_hdl_instance and set_instance_parameter_value to create and parameterize subcomponent instances.

- **Elaboration callback**—You can use an elaboration callback to modify the child components that you create in the static part of the _hw.tcl file with set_instance_parameter_value, or to define subcomponents. With an elaboration callback, you can also instantiate and parameterize subcomponents with add_hdl_instance as a function of the parent component’s parameter values.

Design Guidelines

When instantiating multiple nested components, you must create a unique variation name for each component with the add_hdl_instance command. Prefixing a variation name with the parent component name can prevent conflicts in a system. The variation name can be the same across multiple parent components if the generated parameterization of the nested component is exactly the same.

If you do not adhere to the above naming variation guidelines, Qsys validation-time errors occur, which are often difficult to debug.

A component is either generated or static. Generated components change their generation output (RTL) based on their parameterization. Static components always generate the same output, regardless of their parameterization.

If a component is generated, then any component that might instantiate it with multiple parameter sets must also be considered generated, since its RTL changes with its parameterization. This case has an effect that propagates up to the top level of a design.

Furthermore, components that instantiate static components must have only static children. Additionally, any design file that is static between all parameterizations of a component can only instantiate other static design files. Like the generated rule above, this propagates down the hierarchy.

In order to promote standard and predictable results when generating both generated and static components, Altera recommends the following best-practices:

- For two different parameterizations of a component, a component must never generate a file of the same name with different instantiations. The contents of a file of the same name must be identical for every parameterization of the component.
If a component generates a nested component, it must never instantiate two different parameterizations of the nested component using the same instance name. If the parent component's parameterization affects the parameters of the nested component, the parent component must use a unique instance name for each unique parameterization of the nested component.

Static components that generate differently based on parameterization have the potential to cause problems in the following cases:

- Different file names with the same entity names, results in same entity conflicts at compilation-time.
- Different contents with the same file name, results in overwriting other instances of the component, and results in either compile-time conflicts or unexpected behavior.

Generated components that generate files not based on the output name and that have different content, results in either compile-time conflicts, or unexpected behavior.

**Document Revision History**

Table 8–4 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>- Consolidated content from other Qsys chapters.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added “Upgrading IP Components to the Latest Version”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated for AMBA APB support.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>- Added AMBA AXI4 support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added the demo_axi_memory example with screen shots and example _hw.tcl code.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>- Added new tab structure for the Component Editor.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added AMBA AXI3 support.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>Template update.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>- Removed beta status.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added Avalon Tri-state Conduit (Avalon-TC) interface type.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added many interface templates for Nios custom instructions and Avalon-TC interfaces.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the [Quartus II Handbook Archive](#).
This chapter describes Qsys interconnect, which is a high-bandwidth structure for connecting components, and that allows you to connect IP cores to other IP cores with various interfaces.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

For more information about how you use Qsys to connect IP components and take advantage of the automatically-generated interconnect, refer to the Creating a System With Qsys chapter in volume 1 of the Quartus II Handbook.

For more information about the Avalon interfaces, refer to the Avalon Interface Specifications, and the Qsys System Design Components chapter in volume 1 of the Quartus II Handbook.

The following interconnect interfaces are discussed in this chapters:

- “Memory-Mapped Interfaces” on page 9–1
- “Streaming Interfaces” on page 9–21
- “Interrupt Interfaces” on page 9–28
- “Clock and Reset Interfaces” on page 9–33
- “Conduits” on page 9–35
- “Qsys Interconnect Functions” on page 9–35
- “AMBA AXI3 (version 1.0) Specification Support” on page 9–40
- “AMBA AXI4 (version 2.0) Specification Support” on page 9–43
- “AMBA APB (version 1.0) Specification Support” on page 9–45

Memory-Mapped Interfaces

This section describes the implementation and structure of the Qsys interconnect for memory-mapped interfaces. Content pertains to both Avalon and AXI memory-mapped interfaces, unless noted otherwise.
Memory-mapped transactions between masters and slaves are encapsulated in packets and transmitted on a network that carries the packets between masters and slaves. The command network transports read and write command packets from master interfaces to slave interfaces. The response network transports response packets from slave interfaces to master interfaces.

For each component interface, Qsys interconnect manages memory-mapped transfers, interacting with signals on the connected component. Master and slave interfaces can contain different signals and the interconnect processes any adaptation necessary between them. In the path between master and slaves, the Qsys interconnect might introduce registers for timing synchronization, finite state machines for event sequencing, or nothing at all, depending on the services required by the specific interfaces.

Qsys interconnect supports the following implementation scenarios:

- Any number of components with master and slave interfaces. The master-to-slave relationship can be one-to-one, one-to-many, many-to-one, or many-to-many.
- Master and slaves of different data widths.
- Master and slaves operating in different clock domains.
- Components with different interface properties and signals. Qsys adapts the component interfaces so that interfaces with the following differences can be connected:
  - Avalon and AXI interfaces that use active-high and active-low signalling

    AXI signals are active high, except for the reset signal.

  - Interfaces with different burst characteristics
  - Interfaces with different latencies
  - Interfaces with different data widths
  - Interfaces with different port signatures

    AXI3 to AXI3 interface connections declare a fixed set of signals with variable latency, so there is no need for adapting between active-high/low signalling, burst characteristics, different latencies, or port signatures. Some adaptation is necessary when going to or from Avalon interfaces.
Packet Format for Memory-Mapped Interfaces

The Qsys packet format supports Avalon, AXI, and APB transactions. For Avalon systems without AXI or APB interfaces, some fields are ignored or removed.
Table 9–1 describes the fields of the Qsys packet that encapsulate the memory-mapped master commands and memory-mapped slave responses.

**Table 9–1. Qsys Packet Format (Part 1 of 2)**

<table>
<thead>
<tr>
<th>Field</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Address</td>
<td>Specifies the byte address for the lowest byte in the current cycle. There are no restrictions on address alignment.</td>
</tr>
<tr>
<td>Size</td>
<td>Encodes the run-time size of the transaction.</td>
</tr>
<tr>
<td></td>
<td>In conjunction with address, this field describes the segment of the payload that contains valid data for a beat within the packet.</td>
</tr>
<tr>
<td>Address Sideband</td>
<td>Carries “address” sideband signals. The interconnect passes this field from master to slave. This field is valid for each beat in a packet, even though it is only produced and consumed by an address cycle. Up to 8-bit sideband signals are supported for both read and write address channels.</td>
</tr>
<tr>
<td>Cache</td>
<td>Carries the AXI cache signals.</td>
</tr>
<tr>
<td>Transaction (Exclusive)</td>
<td>Indicates whether the transaction has exclusive access.</td>
</tr>
<tr>
<td>Transaction (Posted)</td>
<td>Used to indicate non-posted writes (writes that require responses).</td>
</tr>
<tr>
<td>Data</td>
<td>For command packets, carries the data to be written. For read response packets, carries the data that has been read.</td>
</tr>
<tr>
<td>Byteenable</td>
<td>Specifies which symbols are valid. AXI can issue or accept any byteenable pattern. For compatibility with Avalon, Altera recommends that you use the following legal values for 32-bit data transactions between Avalon masters and slaves:</td>
</tr>
<tr>
<td></td>
<td>■ 1111—writes full 32 bits</td>
</tr>
<tr>
<td></td>
<td>■ 0011—writes lower 2 bytes</td>
</tr>
<tr>
<td></td>
<td>■ 1100—writes upper 2 bytes</td>
</tr>
<tr>
<td></td>
<td>■ 0001—writes byte 0 only</td>
</tr>
<tr>
<td></td>
<td>■ 0010—writes byte 1 only</td>
</tr>
<tr>
<td></td>
<td>■ 0100—writes byte 2 only</td>
</tr>
<tr>
<td></td>
<td>■ 1000—writes byte 3 only</td>
</tr>
<tr>
<td>Source_ID</td>
<td>The ID of the master or slave that initiated the command or response.</td>
</tr>
<tr>
<td>Destination_ID</td>
<td>The ID of the master or slave to which the command or response is directed.</td>
</tr>
<tr>
<td>Response</td>
<td>Carries the AXI response signals.</td>
</tr>
<tr>
<td>Thread ID</td>
<td>Carries the AXI transaction ID values.</td>
</tr>
<tr>
<td>Byte count</td>
<td>The number of bytes remaining in the transaction, including this beat. Number of bytes requested by the packet.</td>
</tr>
</tbody>
</table>
Table 9–2 lists the transaction type encodings.

<table>
<thead>
<tr>
<th>Bit</th>
<th>Name</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>PKT_TRANS_READ</td>
<td>When asserted, indicates a read transaction.</td>
</tr>
<tr>
<td>1</td>
<td>PKT_TRANS_COMPRESSED_READ</td>
<td>For read transactions, specifies whether or not the read command can be expressed in a single cycle, that is whether or not it has all byte enable asserted on every cycle.</td>
</tr>
<tr>
<td>2</td>
<td>PKT_TRANS_WRITE</td>
<td>When asserted, indicates a write transaction.</td>
</tr>
</tbody>
</table>
The fields of the Qsys packet format are variable length to minimize the resources used. However, if the majority of components in a design have a single data width, for example 32 bits, and a single component has a data width of 64 bits, Qsys inserts a width adapter to accommodate 64-bit transfers.

### Interconnect Domains

A group of connected memory-mapped masters and slaves is called an *interconnect domain*. The components in a single interconnect domain share the same packet format. The following two examples illustrate this point.

### Using Two Separate Domains

Figure 9–2 illustrates the use of two separate domains. The first domain includes two 64-bit masters connected to two 64-bit slaves. The second domain includes one 16-bit master connected to two 16-bit slaves. Because the interfaces in Domain 1 and Domain 2 do not share any connections, Qsys can optimize the packet format for the two separate domains. In this example, the first domain uses a 64-bit data width and the second domain uses 16-bit data.

### Using One Domain with Width Adaptation

Figure 9–3 illustrates a Qsys system that includes two 64-bit masters that access two 64-bit slaves. It also includes one 16-bit master, that accesses two 16-bit slaves and two 64-bit slaves. Because one of the masters connects to all of the slaves, Qsys creates a single domain with two packet formats: one with 64-bit data, and one with 16-bit data. A width adapter manages accesses between the 16-bit master and 64-bit slaves.

### Table 9–2. Transaction Types

<table>
<thead>
<tr>
<th>Bit</th>
<th>Name</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>PKT_TRANS_POSTED</td>
<td>When asserted, no response is required.</td>
</tr>
<tr>
<td>4</td>
<td>PKT_TRANS_LOCK</td>
<td>When asserted, indicates arbitration is locked. Applies to write packets.</td>
</tr>
</tbody>
</table>

3 PKT_TRANS_POSTED
When asserted, no response is required.

4 PKT_TRANS_LOCK
When asserted, indicates arbitration is locked. Applies to write packets.
When there are AXI interfaces in the domain, width adapters are automatically placed on the network boundaries, for example, they are treated as part of the network interface modules. As a result, the switches for a particular AXI domain have a common data width.

In Figure 9–3, the 16-bit Avalon master connects through a 1:4 adapter, then a 4:1 adapter to reach its 16-bit slaves.

**Figure 9–3. One Domain with 1:4 and 4:1 Width Adapters**

---

**Qsys Transformations**

Figure 9–4 provides a detailed view of the transformation that occurs when you generate a Qsys system with memory-mapped master and slave components. The memory-mapped master and slave components connect to network interface modules that encapsulate the transaction in Avalon-ST packets. The memory-mapped interfaces have no information about the encapsulation or the function of the layer transporting the packets and simply operate in accordance with memory-mapped protocol, using the read and write signals and transfers as defined in the Avalon or AXI interface specification.
Figure 9–4 shows a Qsys system with memory-mapped master and slave components.

**Figure 9–4. Qsys Transform from Memory-Mapped Interfaces to Avalon-ST**

The Qsys components that implement the blocks shaded grey in Figure 9–4 are described in “Master Network Interfaces” on page 9–8 and “Slave Network Interfaces” on page 9–11.

**Master Network Interfaces**

This section describes the following master network interfaces:

- “Avalon-MM Master Agent” on page 9–9
- “Avalon-MM Master Translator” on page 9–9
- “AXI Master Agent” on page 9–10
- “AXI Translator” on page 9–10
- “APB Master Agent” on page 9–10
- “APB Slave Agent” on page 9–10
- “Memory-Mapped Router” on page 9–11
- “Memory-Mapped Traffic Limiter” on page 9–11
Avalon network interfaces drive default values for the QoS and RUSER, WUSER, and RUSER packet fields in the master agent, and drop the packet fields in the slave agent. Figure 9–5 shows the Avalon-MM Master network interface.

Figure 9–5. Avalon-MM Master Network Interface

An AXI4 master supports INCR bursts up to 256 beats, QoS signals, and Data Sideband signals. Figure 9–6 shows the AXI Master Network Interface.

Figure 9–6. AXI Master Network Interface

Avalon-MM Master Agent
The Avalon-MM Master Agent translates Avalon-MM master transactions into Qsys command packets and translates the Qsys Avalon-MM slave response packets into Avalon-MM responses.

Avalon-MM Master Translator
The Avalon-MM Master Translator interfaces with an Avalon-MM master component and converts the Avalon-MM master interface to a simpler representation for use in Qsys. It performs the following functions:

- Translates active low signalling to active high signalling
- Inserts wait states to prevent an Avalon-MM master from reading invalid data
- Translates word and symbol addresses
Translates word and symbol burst counts
- Manages re-timing and re-sequencing bursts
- Removes unnecessary address bits

**AXI Master Agent**

An AXI Master Agent accepts AXI commands and produces Qsys command packets. It also accepts Qsys response packets and converts those into AXI responses. This component has separate packet channels for read commands, write commands, read responses, and write responses.

Avalon master agent drives the QoS and BUSER, WUSER, and RUSER packet fields with default values \texttt{AXQO} and \texttt{b0000}, respectively.

For signal descriptions, refer to “Qsys Packet Format” on page 9–4.

**AXI Translator**

AXI4 allows some signals to be omitted from interfaces. The translator bridges between these “incomplete” AXI4 interfaces and the “complete” AXI4 interface on the network interfaces. The AXI translator is inserted for both AXI4 master and slave, and performs the following functions:

- Matches ID widths between the master and slave in 1x1 systems.
- Drives default values as defined in the AMBA Protocol Specifications for missing signals.
- Performs lock transaction bit conversion when an AXI3 master connects to an AXI4 slave in 1x1 systems.

**APB Master Agent**

An APB master agent accepts APB commands and produces or generates Qsys command packets. It also converts Qsys response packets to APB responses.

**APB Slave Agent**

An APB slave agent issues resulting transaction to the APB interface. It also accepts creates Qsys response packets.

**APB Translator**

An APB peripheral does not require \texttt{pslverr} signals to support additional signals for the APB debug interface. The APB translator is inserted for both the master and slave and performs the following functions:

- Sets the response value default to \texttt{OKAY} if the APB slave does not have a \texttt{pslverr} signal.
- Turns on or off additional signals between the APB debug interface, which is used with HPS (Altera SoC’s Hard Processor System).
Memory-Mapped Router

The Memory-Mapped Router routes command packets from the master to the slave, and response packets from the slave to the master. For master command packets, the router uses the address to set the Destination_ID and Avalon-ST channel. For the slave response packet, the router uses the Destination_ID to set the Avalon-ST channel. The demultiplexers use the Avalon-ST channel to route the packet to the correct destination.

Memory-Mapped Traffic Limiter

The Memory-Mapped Traffic Limiter ensures the responses arrive in order. It prevents any command from being sent if the response could conflict with the response for a command that has already been issued. By guaranteeing in-order responses, the Traffic Limiter simplifies the response network.

Slave Network Interfaces

This section describes the following slave network interfaces:
- “Avalon-MM Slave Translator” on page 9–12
- “Avalon-MM Slave Agent” on page 9–13
- “AXI Slave Agent” on page 9–13

An AXI4 slave supports up to 256 beat INCR bursts, QoS signals, and data sideband signals. Figure 9–7 shows an Avalon slave network interface.

Figure 9–7. Avalon-MM Slave Network Interface
Figure 9–8 shows the AXI slave network interface.

Figure 9–8. AXI Slave Network Interface

Avalon-MM Slave Translator

The Avalon-MM Slave Translator interfaces to an Avalon-MM slave component as Figure 9–7 illustrates. It converts the Avalon-MM slave interface to a simplified representation that the Qsys network can use. An Avalon-MM Merlin Slave Translator performs the following functions:

- Drives the `begintransfer`, `beginbursttransfer`, and `byteenable` signals.
- Supports Avalon-MM slaves that operate using fixed timing and or slaves that use the `readdatavalid` signal to identify valid data.
- Translates the `read`, `write`, and `chipselect` signals into the representation that the Avalon-ST slave response network uses.
- Converts active low signals to active high signals.
- Translates word and symbol addresses and burstcounts.
- Handles burstcount timing and sequencing.
- Removes unnecessary address bits.

AXI Translator

AXI4 allows some signals to be omitted from interfaces. The translator bridges between these “incomplete” AXI4 interfaces and the “complete” AXI4 interface on the network interfaces. The AXI translator is inserted for both AXI4 master and slave, and performs the following functions:

- Matches ID widths between master and slave in 1x1 systems.
- Drives default values as defined in the *AMBA Protocol Specifications* for missing signals.
- Performs lock transaction bit conversion when an AXI3 master connects to an AXI4 slave in 1x1 systems.
Avalon-MM Slave Agent

The Avalon-MM Slave Agent accepts command packets and issues the resulting transactions to the Avalon interface. For pipelined slaves, an Avalon-ST FIFO stores information about pending transactions. The size of this FIFO is the maximum number of pending responses that you specify when creating the slave component.

The Avalon-MM Slave Agent also backpressures the Avalon-MM master command interface when the FIFO is full if the slave component includes the waitrequest signal.

AXI Slave Agent

An AXI Slave Agent works similar to a master agent in reverse. The AXI slave Agent accepts Qsys command packets to create AXI commands, and accepts AXI responses to create Qsys response packets. This component has separate packet channels for read commands, write commands, read responses, and write responses.

Arbitration

When multiple masters contend for access to a slave, Qsys automatically inserts arbitration logic which grants access in fairness-based, round-robin order.

In a fairness-based arbitration scheme, each master has an integer value of transfer shares with respect to a slave. One share represents permission to perform one transfer. The default arbitration scheme is equal share round-robin that grants equal, sequential access to all requesting masters. You can change the arbitration scheme to weighted round-robin by specifying a relative number of arbitration shares to the masters that access a particular slave. AXI slaves have separate arbitration for their independent read and write channels, and the Arbitration Shares setting affects both the read and write arbitration. To display arbitration settings, right-click an instance on the System Contents tab, and then click Show Arbitration Shares.

Figure 9–9 illustrates arbitration shares indicated in the Connections column.

Arbitration Timing

Figure 9–10 illustrates the arbitration timing. As this figure illustrates, a device can drive valid data in the granted cycle. Figure 9–10 shows the following sequence of events:

1. In cycle one, the arbiter grants a request. The granted device drives valid data in cycles one and two.
2. In cycle 4, the arbiter grants a request. The granted device drives valid data in cycles 4 and 5.

3. In cycle 6, the arbiter grants a request. The granted device drives valid data in cycles 6 and 7.

4. At the positive edge of cycle 3, a grant signal is not present, although there is a request signal asserted. Therefore, cycle 3 does not contain valid data.

**Arbitration Rules**

This section describes the rules by which the arbiter grants access to masters.

**Fairness-Based Shares**

In a fairness-based arbitration scheme, each master-to-slave connection provides a transfer share count. This count is a request for the arbiter to grant a specific number of transfers to this master before giving control to a different master. One share represents permission to perform one transfer.

For example, for two masters that continuously attempt to perform back-to-back transfers to a slave, the arbiter grants Master 1 access for three transfers, and Master 2 for four transfers. This cycle repeats indefinitely. Figure 9–11 shows each master’s transfer request output, wait request input (which is driven by the arbiter logic), and the current master with control of the slave.
If a master stops requesting transfers before it exhausts its shares, it forfeits all of its remaining shares, and the arbiter grants access to another requesting master, as shown in Figure 9–14. After completing one transfer, Master 2 stops requesting for one clock cycle. As a result, the arbiter grants access back to Master 1, which gets a replenished supply of shares.

**Figure 9–12. Arbitration of Two Masters with a Gap in Transfer Requests**

![Diagram](image)

**Round-Robin Scheduling**

When multiple masters contend for access to a slave, the arbiter grants shares in round-robin order. Round-robin scheduling drives a request interface according to space available and data available credit interfaces. At every slave transfer, only requesting masters are included in the arbitration.

**Burst Transfers**

For burst transfer arbitration, there is a one-to-one relationship between a single share and a transaction, so that arbitration shares are respected during burst transactions. For example, for an arbitration share of 3, a master is granted 3 burst transactions. Once a burst begins between a master-slave pair, arbiter logic does not allow any other master to access the slave until the burst completes.

**Arbitration Examples**

Figure 9–13 illustrates the timing for two Avalon-MM masters continuously accessing a single Avalon-MM slave to perform back-to-back transfers. Master 1 has three shares and Master 2 has four shares. The arbiter grants Master 1 access for three transfers, then Master 2 for four transfers. This cycle repeats indefinitely.

**Figure 9–13. Arbitration of Continuous Transfer Requests from Two Masters**

![Diagram](image)
If a master stops requesting transfers before it exhausts its shares, it forfeits all of its remaining shares, and the arbiter grants access to another requesting master as Figure 9–14 illustrates. After completing one transfer, Master 2 stops requesting for one clock cycle. As a result, the arbiter grants access back to Master 1, which gets three shares.

**Figure 9–14. Arbitration of Two Masters with a Gap in Transfer Requests**

Memory-Mapped Arbiter

The input to the Memory-Mapped Arbiter is the command packet for all masters requesting access to a particular slave. The arbiter outputs the channel number for the selected master. This channel number controls the output of a multiplexer that selects the slave device. Figure 9–15 illustrates this logic.
In Figure 9–15, four Avalon-MM masters connect to four Avalon-MM slaves. In each cycle, an arbiter positioned in front of each Avalon-MM slave selects among the requesting Avalon-MM masters.

**Figure 9–15. Arbitration Logic**

If you specified a Limit interconnect pipeline stages to parameter greater than zero on the Qsys Project Settings tab, the output of the Arbiter is registered. Registering this output reduces the amount of combinational logic between the master and interconnect, increasing the $f_{\text{MAX}}$ of the system.

**Width Adaptation**

Qsys width adaptation converts between Avalon memory-mapped master and slaves with different data and byte enable widths, and manages the run-time size requirements of AXI. Width adaptation for AXI to Avalon interfaces is also supported.
Memory-Mapped Width Adapter

The Memory-Mapped Width Adapter is used in the Avalon-ST domain and operates with information contained in the packet format illustrated Figure 9–2 on page 9–6. It accepts packets on its sink interface with one data width and produces output packets on its source interface with a different data width. The ratio of the narrow data width must be a power of two, such as 4:1, 8:1, and 16:1. The ratio of the wider data width to the narrower width must be a power of two, such as 4:1, 8:1, and 16:1. These output packets may have a different size if the input size exceeds the output data bus width, or if data packing is enabled.

This adapter assumes that the field ordering of the input and output packets is the same, with the only difference being the width of the data and accompanying byte enable fields.

When the width adapter converts from wide data to narrow data, the narrower data is transmitted over several beats. The first output beat contains the lowest addressed segment of the input data and byte enables. Figure 1-14 illustrates the timing for a 4:1 width adapter.

When the width adapter converts from narrow data to wide data, each input beat’s data and byte enables are copied to the appropriate segment of the wider output data and byte enables signals.

Figure 9–16. Width Adapter Timing for a 4:1 Adapter
AXI Wide to Narrow Adaptation

For all cases shown in Table 9–3, read data is repacked to match the original size. Responses are merged, with the following error precedence: DBERR, SLVERR, OKAY, and EXOKAY.

Table 9–3. Wide to Narrow Adaptation (Downsizing)

<table>
<thead>
<tr>
<th>Burst Type</th>
<th>Behavior</th>
</tr>
</thead>
<tbody>
<tr>
<td>Incrementing</td>
<td>If the transaction size is less than or equal to the output width, the burst is unmodified. Otherwise, it is converted to an incrementing burst with a larger length and size equal to the output width. If the resulting burst is unsuitable for the slave, the burst is converted to multiple sequential bursts of the largest allowable lengths. For example, for a 2:1 downsizing ratio, an INCR9 burst is converted into INCR16 + INCR2 bursts. This is true if the maximum burstcount a slave can accept is 16, which is the case for AXI3 slaves. Avalon slaves have a maximum burstcount of 64.</td>
</tr>
<tr>
<td>Wrapping</td>
<td>If the transaction size is less than or equal to the output width, the burst is unmodified. Otherwise, it is converted to a wrapping burst with a larger length, with a size equal to the output width. If the resulting burst is unsuitable for the slave, the burst is converted to multiple sequential bursts of the largest allowable lengths; respecting wrap boundaries. For example, for a 2:1 downsizing ratio, a WRAP16 burst is converted into two or three INCR bursts, depending on the address.</td>
</tr>
<tr>
<td>Fixed</td>
<td>If the transaction size is less than or equal to the output width, the burst is unmodified. Otherwise, it is converted into repeated sequential bursts over the same addresses. For example, for a 2:1 downsizing ratio, a FIXED single burst is converted into an INCR2 burst.</td>
</tr>
</tbody>
</table>

AXI Narrow to Wide Adaptation

Table 9–4 describes burst behavior for AXI narrow-to-wide adaptation.

Table 9–4. Narrow to Wide Adaptation (Upsizing)

<table>
<thead>
<tr>
<th>Burst Type</th>
<th>Behavior</th>
</tr>
</thead>
<tbody>
<tr>
<td>Incrementing</td>
<td>The burst (and its response) passes through unmodified. Data and write strobes are placed in the correct output segment.</td>
</tr>
<tr>
<td>Wrapping</td>
<td>The burst (and its response) passes through unmodified.</td>
</tr>
<tr>
<td>Fixed</td>
<td>The burst (and its response) passes through unmodified.</td>
</tr>
</tbody>
</table>

Burst Transfers

Avalon-MM and AXI burst transactions grant a master uninterrupted access to a slave for a specified number of transfers. The master specifies the number of transfers when it initiates the burst. Once a burst begins between a master and slave pair, arbiter logic is locked until the burst completes. For burst masters, the length of the burst is the number of cycles that the master has access to the slave, and the selected arbitration shares have no effect.
Memory-Mapped Burst Adapter

The Qsys interconnect uses the memory-mapped burst adapter to accommodate the burst capabilities of each interface in the system, including interfaces that do not support burst transfers. The maximum burst length for each interface is a property of the component interface and is independent of other interfaces in the system. Therefore, a particular master might be capable of initiating a burst longer than a slave’s maximum supported burst length. In this case, the burst adapter translates the large master burst into smaller bursts, or into individual slave transfers if the slave does not support bursting. Until the master completes the burst, the arbiter logic prevents other masters from accessing the target slave. For example, if a master initiates a burst of 16 transfers to a slave with maximum burst length of 8, the burst adapter initiates 2 bursts of length 8 to the slave.

AXI masters can issue burst types that Avalon cannot accept, for example, fixed bursts. In this case, the burst adapter converts the fixed burst into a sequence of transactions to the same address.

For AXI4 slaves, Qsys allows 256-beat INCR bursts, though you must ensure that 256-beat narrow-sized INCR bursts are shortened to 16-beat narrow-sized INCR bursts for AXI3 slaves.

Avalon-MM masters always issue addresses that are aligned to the size of the transfer. However, in some cases, when a narrow-to-wide width adaptation is used, the resulting address may be unaligned. In the case of unaligned addresses, the burst adapter issues the maximum possible sized bursts, with appropriate byte enables, to bring the burst-in-progress up to an aligned slave address. Then, it completes the burst on aligned addresses.

The burst adapter supports variable wrap or sequential burst types to accommodate the different properties of memory-mapped masters. Some bursting masters can issue more than one burst type.

Burst adaptation is available for Avalon to Avalon, Avalon to AXI, and AXI to Avalon connections, and AXI to AXI connections.

For AXI4 to AXI3 connections, Qsys follows an AXI4 256 burst length to AXI3 16 burst length.

Table 9–5 and Table 9–6 specify the behavior when converting between AXI and Avalon burst types.
Table 9–5 describes behavior when converting from AXI and Avalon burst types.

Table 9–5. Burst Adaptation: AXI to Avalon

<table>
<thead>
<tr>
<th>Burst Type</th>
<th>Behavior</th>
</tr>
</thead>
<tbody>
<tr>
<td>Incrementing</td>
<td><strong>Sequential Slave</strong>&lt;br&gt;Bursts that exceed ( \text{slave_max_burst_length} ) are converted to multiple sequential bursts of a length less than or equal to the ( \text{slave_max_burst_length} ). Otherwise, the burst is unconverted. For example, for an Avalon slave with a maximum burst length of 4, an \text{INCR7} burst is converted to \text{INCR4} + \text{INCR3}.</td>
</tr>
<tr>
<td>Wrapping</td>
<td><strong>Wrapping Slave</strong>&lt;br&gt;Bursts that exceed the ( \text{slave_max_burst_length} ) are converted to multiple sequential bursts of length less than or equal to the ( \text{slave_max_burst_length} ). Bursts that exceed the wrapping boundary are converted to multiple sequential bursts that respect the slave’s wrapping boundary.</td>
</tr>
<tr>
<td>Fixed</td>
<td><strong>Sequential Slave</strong>&lt;br&gt;A WRAP burst is converted to multiple sequential bursts. The sequential bursts are less than or equal to the ( \text{max_burst_length} ) and respect the transaction’s wrapping boundary.</td>
</tr>
<tr>
<td>Narrow</td>
<td><strong>Wrapping Slave</strong>&lt;br&gt;If the WRAP transaction’s boundary matches the slave’s boundary, then the burst passes through. Otherwise, the burst is converted to sequential bursts that respect both the transaction and slave wrap boundaries.</td>
</tr>
</tbody>
</table>

Table 9–6 describes behavior when converting from Avalon to AXI burst types.

Table 9–6. Burst Adaptation: Avalon to AXI

<table>
<thead>
<tr>
<th>Burst Type</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sequential</td>
<td>Bursts of length greater than 16 are converted to multiple \text{INCR} bursts of a length less than or equal to 16. Bursts of length less than or equal to 16 are not converted.</td>
</tr>
<tr>
<td>Wrapping</td>
<td>Only Avalon masters with ( \text{alwaysBurstMaxBurst} = \text{true} ) are supported. The WRAP burst is passed through if the length is less than or equal to 16. Otherwise, it is converted to two or more \text{INCR} bursts that respect the transaction’s wrap boundary.</td>
</tr>
</tbody>
</table>

For information about AXI-to-AXI adaptation, refer to “AXI Wide to Narrow Adaptation” on page 9–19.

**Streaming Interfaces**

High bandwidth components with streaming data typically use Avalon-ST interfaces for the high throughput datapath. These components can also use memory-mapped connection interfaces to provide an access point for control. In contrast to the memory-mapped interconnect, which you can use to create a wide variety of topologies, the Avalon-ST interconnect always creates a point-to-point connection between a single data source and data sink, as Figure 9–17 illustrates.
There are two connection pairs in this figure:

- The data source in the Rx Interface transfers data to the data sink in the FIFO.
- The data source in the FIFO transfers data to the Tx Interface data sink.

In Figure 9–17, the memory-mapped interface allows a processor to access the data source, FIFO, or data sink to provide system control.

**Figure 9–17. Use of the Memory-Mapped and Avalon-ST Interfaces**

If your source and sink interfaces have different formats, for example, a 32-bit source and an 8-bit sink, Qsys automatically inserts the necessary adapters, which are then visible in the System Contents tab. For more information, refer to “Avalon-ST Adapters” on page 9–24.

Figure 9–18 illustrates the simplest system example with an Avalon-ST connection between the source and sink. This source-sink pair includes only the data signal. The sink must be able to receive data as soon as the source interface comes out of reset.

**Figure 9–18. Interconnect for a Simple Avalon Streaming Source-Sink Pair**
Figure 9–19 illustrates a more extensive interface that includes signals indicating the start and end of packets, channel numbers, error conditions, and back pressure.

**Figure 9–19. Avalon Streaming Interface for Packet Data**

All data transfers using Avalon-ST interconnect occur synchronously to the rising edge of the associated clock interface. Throughput and frequency of a system depends on the components and how they are connected.

AXI streaming components are not available in Qsys, version 13.0. For details about the Avalon-ST interface protocol, refer to the *Avalon Interface Specification*.

The Qsys Component Library includes a number of Avalon-ST components that you can use to create datapaths, including datapaths whose input and output streams have different properties. Generated systems that include memory-mapped master and slave components may also use these Avalon-ST components because the generation process creates an interconnect whose structure resembles a network topology as “Qsys Transformations” on page 9–7 describes. The following sections introduce the Avalon-ST components.

**Avalon-ST Multiplexer**

The Avalon-ST Multiplexer accepts data on its Avalon-ST sink interface and multiplexes the data for transmission on its Avalon-ST source interface. You can parameterize the multiplexer to append channel information on the source to indicate which sink is driving the source data. The multiplexer includes internal arbitration logic which selects between inputs using a round-robin arbitration algorithm. Figure 9–20 illustrates the Avalon-ST multiplexer. Among the parameters that you can specify are the option to use packet scheduling, which guarantees that the multiplexer only changes inputs at the end of a packet.

**Figure 9–20. Avalon-ST Multiplexer**
**Avalon-ST Demultiplexer**

The Avalon-ST Demultiplexer accepts channelized data on its sink interface, and transmits the data on one of its source interfaces. The channel bits of the source stream indicate which port the drives the output data. Figure 9–21 illustrates the Multiplexer. Among the parameters that you can specify are the number of output ports and the width of the channel signal.

**Figure 9–21. Avalon-ST Demultiplexer**

---

**Avalon-ST Adapters**

Qsys automatically adds Avalon-ST adapters between two components during system generation when it detects mismatched interfaces. If you connect mismatched Avalon-ST sources and sinks, for example, a 32-bit source and an 8-bit sink, Qsys inserts the appropriate adapter type, as described below, to connect the mismatched interfaces. After generation, you can view the inserted adapters with the **Show System With Qsys Fabric Components** command on the System menu. Qsys reports the mismatched interfaces and inserted adapter with informational messages.

You can turn off the auto-inserted adapters feature by adding the `qsys_enable_avalon_streaming_transform=off` command to the `quartus.ini` file. When you turn off the auto-inserted adapters feature, if mismatched interfaces are detected during system generation, Qsys does not insert adapters and reports the mismatched interface with an error message.

The auto-inserted adapters feature does not work for video IP core connections.

Qsys includes the following adapter types:

- Data Format Adapter
- Timing Adapter
- Channel Adapter
- Error Adapter

The following sections provide an overview of these adapters.

**Data Format Adapter**

The data format adapter allows you to connect interfaces that have different values for the parameters defining the data signal. One of the most common uses of this adapter is to convert data streams of different widths.
The available data format adaptations are listed in Table 9–7.

<table>
<thead>
<tr>
<th>Condition</th>
<th>Description of Adapter Logic</th>
</tr>
</thead>
<tbody>
<tr>
<td>The source and sink’s bits per symbol are different.</td>
<td>The connection cannot be made.</td>
</tr>
<tr>
<td>The source and sink have a different number of symbols per beat.</td>
<td>The adapter converts the source’s width to the sink’s width.</td>
</tr>
<tr>
<td></td>
<td>If the adaptation is from a wider to a narrower interface, a beat of data at the input</td>
</tr>
<tr>
<td></td>
<td>corresponds to multiple beats of data at the output. If the input error signal is asserted</td>
</tr>
<tr>
<td></td>
<td>for a single beat, it is asserted on output for multiple beats.</td>
</tr>
<tr>
<td></td>
<td>If the adaptation is from a narrow to a wider interface, multiple input beats are required</td>
</tr>
<tr>
<td></td>
<td>to fill a single output beat, and the output error is the logical OR of the input error</td>
</tr>
</tbody>
</table>

Figure 9–22 shows a data format adapter that allows a connection between a 128-bit input data stream and three 32-bit output data streams.

Figure 9–22. Avalon Streaming Interconnect with Data Format Adapter
Timing Adapter

The timing adapter allows you to connect component interfaces that require a different number of cycles before driving or receiving data. This adapter inserts a FIFO buffer between the source and sink to buffer data or pipeline stages to delay the back pressure signals. You can also use the timing adapter to connect interfaces that support the ready signal, and those that do not. The timing adapter treats all signals other than the ready and valid signals as payload, and simply drives them from the source to the sink. Table 9–8 outlines the adaptations that the timing adapter provides.

<table>
<thead>
<tr>
<th>Condition</th>
<th>Adaptation</th>
</tr>
</thead>
<tbody>
<tr>
<td>The source has ready, but the sink does not.</td>
<td>In this case, the source can respond to backpressure, but the sink never needs to apply it. The ready input to the source interface is connected directly to logical 1.</td>
</tr>
<tr>
<td>The source does not have ready, but the sink does.</td>
<td>The sink may apply backpressure, but the source is unable to respond to it. There is no logic that the adapter can insert that prevents data loss when the source asserts valid but the sink is not ready. The adapter provides simulation time error messages and an error indication if data is ever lost. The user is presented with a warning, and the connection is allowed.</td>
</tr>
<tr>
<td>The source and sink both support backpressure, but the sink’s ready latency is greater than the source’s.</td>
<td>The source responds to ready assertion or deassertion faster than the sink requires it. A number of pipeline stages equal to the difference in ready latency are inserted in the ready path from the sink back to the source, causing the source and the sink to see the same cycles as ready cycles.</td>
</tr>
<tr>
<td>The source and sink both support backpressure, but the sink’s ready latency is less than the source’s.</td>
<td>The source cannot respond to ready assertion or deassertion in time to satisfy the sink. A buffer whose depth is equal to the difference in ready latency is inserted to compensate for the source’s inability to respond in time.</td>
</tr>
</tbody>
</table>

Channel Adapter

The channel adapter provides adaptations between interfaces that have different support for the channel signal, the maximum number of channels supported, or channel-related parameters. The adaptations for the channel adapter are described in Table 9–9.

<table>
<thead>
<tr>
<th>Condition</th>
<th>Description of Adapter Logic</th>
</tr>
</thead>
<tbody>
<tr>
<td>The source uses channels, but the sink does not.</td>
<td>You are given a warning at generation time. The adapter provides a simulation error and signals an error for data for any channel from the source other than 0.</td>
</tr>
<tr>
<td>The sink has channel, but the source does not.</td>
<td>You are given a warning, and the channel inputs to the sink are all tied to a logical 0.</td>
</tr>
</tbody>
</table>
The error adapter ensures that per-bit-error information provided by the source interface is correctly connected to the sink interface’s input error signal. Matching error conditions processed by the source and sink are connected. If the source has an error condition that is not supported by the sink, the signal is left unconnected; the adapter provides a simulation error message and an error indication if this error is ever asserted. If the sink has an error condition that is not supported by the source, the sink’s input is tied to zero.

Table 9–10 describes the available options for the error adapter on the Parameter Settings page of the configuration wizard.

### Table 9–9. Channel Adapter

<table>
<thead>
<tr>
<th>Condition</th>
<th>Description of Adapter Logic</th>
</tr>
</thead>
<tbody>
<tr>
<td>The source and sink both support channels, and the source’s maximum number of channels is less than the sink’s.</td>
<td>The source’s channel is connected to the sink’s channel unchanged. If the sink’s channel signal has more bits, the higher bits are tied to a logical 0.</td>
</tr>
<tr>
<td>The source and sink both support channels, but the source’s maximum number of channels is greater than the sink’s.</td>
<td>The source’s channel is connected to the sink’s channel unchanged. If the source’s channel signal has more bits, the higher bits are left unconnected. You are given a warning that channel information may be lost. An adapter provides a simulation error message and an error indication if the value of channel from the source is greater than the sink’s maximum number of channels. In addition, the valid signal to the sink is deasserted so that the sink never sees data for channels that are out of range.</td>
</tr>
</tbody>
</table>

### Error Adapter

The error adapter ensures that per-bit-error information provided by the source interface is correctly connected to the sink interface’s input error signal. Matching error conditions processed by the source and sink are connected. If the source has an error condition that is not supported by the sink, the signal is left unconnected; the adapter provides a simulation error message and an error indication if this error is ever asserted. If the sink has an error condition that is not supported by the source, the sink’s input is tied to zero.

Table 9–10 describes the available options for the error adapter on the Parameter Settings page of the configuration wizard.

### Table 9–10. Avalon-ST Error Adapter Parameters (Part 1 of 2)

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Input Interface Parameters</strong></td>
<td></td>
</tr>
<tr>
<td>Error Signal Width (bits)</td>
<td>Type the width of the error signal. Valid values are 0–31 bits. Type 0 if the error signal is not used.</td>
</tr>
<tr>
<td>Error Signal Description</td>
<td>Type the description for each of the error bits. Separate the description fields by commas. For a connection to be made, the description of the error bits in the source and sink must match. Refer to “Error Adapter” on page 9–27 for the adaptations that can be made when the bits do not match.</td>
</tr>
<tr>
<td><strong>Output Interface Parameters</strong></td>
<td></td>
</tr>
<tr>
<td>Error Signal Width (bits)</td>
<td>Type the width of the error signal. Valid values are 0–31 bits. Type 0 if you do not need to send error values.</td>
</tr>
<tr>
<td>Error Signal Description</td>
<td>Type the description for each of the error bits. Separate the description fields by commas. For a connection to be made, the description of the error bits in the source and sink must match. Refer to “Error Adapter” on page 9–27 for the adaptations that can be made when the bits do not match.</td>
</tr>
<tr>
<td><strong>Common to Input and Output Interfaces</strong></td>
<td></td>
</tr>
<tr>
<td>Support Backpressure with the ready signal</td>
<td>Turn on this option to add the backpressure functionality to the interface.</td>
</tr>
</tbody>
</table>
Interrupt Interfaces

Using individual requests, the interrupt logic can process up to 32 IRQ inputs connected to each interrupt receiver. With this logic, the interrupt sender connected to interrupt receiver_0 is the highest priority with sequential receivers being successively lower priority. You can redefine the priority of interrupt senders by instantiating the IRQ mapper component. For more information refer to the “IRQ Mapper” on page 9–31.

You can define the interrupt sender interface as asynchronous with no associated clock or reset interfaces; the properties associatedClock and associatedReset can thus be left undefined for Avalon masters and slaves. AXI masters and slaves are required to specify associatedClock and associatedReset.

For clock crossing adaption on interrupts, Qsys inserts a synchronizer, which is clocked with the interrupt end point interface clock when the corresponding starting point interrupt interface has no clock or a different clock (than the end point). Qsys inserts the adapter if there is any kind of mismatch between the start and end points.

Individual Requests IRQ Scheme

In the individual requests IRQ scheme, Qsys interconnect passes IRQs directly from the sender to the receiver, without making assumptions about IRQ priority. In the event that multiple senders assert their IRQs simultaneously, the receiver logic determines which IRQ has highest priority, and then responds appropriately.

Using individual requests, the interrupt controller can process up to 32 IRQ inputs. The interrupt controller generates a 32-bit signal irq[31:0] to the receiver, and maps slave IRQ signals to the bits of irq[31:0]. Any unassigned bits of irq[31:0] are disabled.
Figure 9–23 shows an example of the interrupt controller mapping the IRQs on four senders to irq[31:0] on a receiver.

**Figure 9–23. IRQ Mapping Using Software Priority**

![Figure 9–23](image)

**Priority-Encoded Interrupt Scheme**

In the priority-encoded interrupt scheme, in the event that multiple slaves assert their IRQs simultaneously, Qsys interconnect provides the interrupt receiver with a 1-bit interrupt signal, and the number of the highest priority active interrupt. An IRQ of lesser priority is undetectable until all IRQs of higher priority are serviced.

Using priority-encoded interrupts, the interrupt controller can process up to 64 slave IRQ signals. The interrupt controller generates a 1-bit irq signal to the receiver, signifying that one or more senders have generated an IRQ. The controller also generates a 6-bit irqnumber signal, which outputs the encoded value of the highest pending IRQ.
Figure 9–24 shows an example of IRQ mapping using the encoded interrupt scheme.

**Figure 9–24. IRQ Mapping Using Hardware Priority**

Assigning IRQS in Qsys

You assign IRQ connections on the **System Contents** tab of Qsys. After adding all components to the system, you connect interrupt senders and receivers. You can use the **IRQ** column to specify an IRQ number with respect to each receiver, or to specify a receiver’s IRQ as unconnected.

Qsys uses the following three components to implement interrupt handling:

- **IRQ Bridge**
- **IRQ Mapper**
- **IRQ Clock Crosser**

**IRQ Bridge**

The IRQ Bridge allows you to route interrupt wires between Qsys subsystems. These interrupts are routed to the IRQ receiver bridge in the CPU Subsystem.

> Nios II BSP tools do not fully support the IRQ Bridge. Interrupts connected via an IRQ Bridge will not appear in the generated `system.h` file.
In Figure 9–25, the Peripheral Subsystem has three interrupt senders that are exported to the top level of the subsystem.

**Figure 9–25. Qsys IRQ Bridge Application**

**IRQ Mapper**

Qsys inserts the IRQ Mapper automatically during generation. The IRQ Mapper converts individual interrupt wires to a bus, and then maps the appropriate IRQ priority number onto the bus.

By default, the interrupt sender connected to the `receiver0` interface of the IRQ mapper is the highest priority, and sequential receivers are successively lower priority. You can modify the interrupt priority of each IRQ wire by modifying the IRQ priority number in Qsys under the **IRQ** column. The modified priority is reflected in the **IRQ_MAP** parameter for the auto-inserted IRQ Mapper.
Figure 9–26 shows the IRQ column in Qsys and the default interrupt priority allocated for the CPU subsystem shown in Figure 9–25 above.

**Figure 9–26. Qsys IRQ Column**

<table>
<thead>
<tr>
<th>clk</th>
<th>Cloc...</th>
<th>DOutclk_0</th>
</tr>
</thead>
<tbody>
<tr>
<td>clk_reset</td>
<td>Reset...</td>
<td>DOut...</td>
</tr>
<tr>
<td>irq_bridge</td>
<td>IRQ0...</td>
<td>DOutclk_0</td>
</tr>
<tr>
<td>clk</td>
<td>Cloc...</td>
<td>DOutclk_0</td>
</tr>
<tr>
<td>receiver_irq</td>
<td>Inter...</td>
<td>irq... [clk]</td>
</tr>
<tr>
<td>clk_reset</td>
<td>Reset...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>sender0_irq</td>
<td>Inter...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>sender1_irq</td>
<td>Inter...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>sender2_irq</td>
<td>Inter...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>interrupt_sender_4</td>
<td>Inter...</td>
<td></td>
</tr>
<tr>
<td>clk</td>
<td>Cloc...</td>
<td>DOutclk_0</td>
</tr>
<tr>
<td>reset</td>
<td>Reset...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>p1</td>
<td></td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>irq</td>
<td>Inter...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>nios2_processor</td>
<td>NIOS...</td>
<td></td>
</tr>
<tr>
<td>clk</td>
<td>Cloc...</td>
<td>DOutclk_0</td>
</tr>
<tr>
<td>reset_n</td>
<td>Reset...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>data_master</td>
<td>Avai...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>instruction_master</td>
<td>Avai...</td>
<td>DOut[clk]</td>
</tr>
<tr>
<td>d_irq</td>
<td>Inter...</td>
<td>DOut[clk]</td>
</tr>
</tbody>
</table>

**IRQ Clock Crosser**

The IRQ Clock Crosser synchronizes interrupt senders and receivers that are in different clock domains. To use this component, connect the clocks for both the interrupt sender and receiver, and for both the interrupt sender and receiver interfaces. Qsys automatically inserts this component when it is required.
Clock and Reset Interfaces

Clock interfaces define the clocks used by a component. Components can have clock inputs, clock outputs, or both.

Reset interfaces provide both soft and hard reset functionality. Soft reset logic typically re-initializes registers and memories without powering down the device. Hard reset logic initializes the device after power-on. You can define separate reset sources for each clock domain, a single reset source for all clocks, or any combination in between.

Clock Interfaces

You can use the Clock Settings tab to define external clock sources, for example an oscillator on your board. The Clock Source parameters allows you to set the following options:

- **Clock frequency**—The frequency of the output clock from this clock source.
- **Clock frequency is known**—When turned on, the clock frequency is known. When turned off, the frequency is set from outside the system.

If turned off, system generation may fail because the components do not receive the necessary clock information. For best results, turn this option on before system generation.

- **Reset synchronous edges**
  - **None**—The reset is asserted and deasserted asynchronously. You can use this setting if you have internal synchronization circuitry that matches the reset required for the IP in the system.
  - **Both**—The reset is asserted and deasserted synchronously.
  - **Deassert**—The reset is deasserted synchronously and asserted asynchronously.

For more information about synchronous design practices, refer to Recommended Design Practices in volume 1 of the Quartus II Handbook.

Reset Interfaces

You can choose to create a single global reset domain by selecting Create Global Reset Network on the System menu. If your design requires more than one reset domain, you can implement your own reset logic and connectivity.

The Component Library includes a reset controller and a reset bridge to implement the reset functionality. You can also design your own reset logic.

If you design your own reset circuitry you must carefully consider situations which might result in system lockup. For example, if an Avalon-MM slave is reset in the middle of a transaction, the Avalon-MM master might wait forever.
Single Global Reset Signal Implemented by Qsys

If you select Create Global Reset Network on the System menu, the Qsys interconnect creates a global reset bus. All of the reset requests are ORed together, synchronized to each clock domain, and fed to the reset inputs. The duration of the reset signal is at least one clock period.

The Qsys interconnect inserts the system-wide reset under the following conditions:

- The global reset input to the Qsys system is asserted.
- Any component asserts its reset request signal.

Reset Controller

The Reset Controller has the following parameters that you can specify to customize its behavior:

- **Number of inputs**—Indicates the number of individual reset interfaces the controller ORs to create a signal reset output.

- **Output reset synchronous edges**—Specifies the level of synchronization. You can select one of the following options:
  - **None**—The reset is asserted and deasserted asynchronously. You can use this setting if you have designed internal synchronization circuitry that matches the reset style required for the IP in the system.
  - **Both**—The reset is asserted and deasserted synchronously.
  - **Deassert**—The reset is deasserted synchronously and asserted asynchronously.

- **Synchronization depth**—Specifies the number of register stages the synchronizer uses to eliminate the propagation of metastable events.

Qsys automatically inserts reset synchronizers under the following conditions:

- More than one reset source is connected to a reset sink
- There is a mismatch between the reset source’s synchronous edges and the reset sinks’ synchronous edges

Reset Bridge

The Reset Bridge allows you to use a reset signal in two or more subsystems of your Qsys system. You can connect one reset source to local components, and export one or more to other subsystems, as required. The Reset Bridge parameters are used to describe the incoming reset and include the following options:

- **Active low reset**—When turned on, reset is asserted low.

- **Synchronous edges**—Specifies the level of synchronization and includes the following options:
  - **None**—The reset is asserted and deasserted asynchronously. Use this setting if you have internal synchronization circuitry.
  - **Both**—The reset is asserted and deasserted synchronously.
  - **Deassert**—The reset is deasserted synchronously, and asserted asynchronously.
Number of reset outputs—The number of reset interfaces that are exported.

Qsys supports multiple reset sink connections to a single reset source interface. However, there are situations in composed systems where an internally generated reset must be exported from the composed system in addition to being used to connect internal components. In this situation, you must declare one reset output interface as an export, and use another reset output to connect internal components.

Conduits

You can use the conduit interface type for interfaces that do not fit any of the other interface types, and to group any arbitrary collection of signals. Like other interface types, you can export or connect conduit interfaces. The PCI Express-to-Ethernet Bridge Example System shown in Figure 7–15 on page 7–31 in the Creating a System With Qsys chapter is an example of the use of the conduit interface for export.

You can declare an associated clock interface for conduit interfaces in the same way as memory-mapped interfaces with the associatedClock.

To connect two conduit interfaces inside Qsys, the following conditions must be met:

- The interfaces must match exactly with the same signal roles and widths.
- The interfaces must be the opposite directions.
- Clocked conduit connections must have matching associatedClocks on each of their endpoint interfaces.

To connect a conduit output to more than one input conduit interface, you can create a custom component. The custom component could have one input that connects to two outputs, and you can use this component between other conduits that you want to connect.

For more information about the Avalon Conduit interface, refer to the Avalon Interface Specifications.

Qsys Interconnect Functions

Qsys supports the following underlying system interconnect functionality:

- “Address Decoding” on page 9–35
- “Datapath Multiplexing” on page 9–37
- “Wait State Insertion” on page 9–38
- “Interconnect Pipelining” on page 9–39

Address Decoding

Address decoding logic forwards appropriate addresses to each slave. Address decoding logic simplifies component design in the following ways:

- The interconnect selects a slave whenever it is being addressed by a master. Slave components do not need to decode the address to determine when they are selected.
- Slave addresses are properly aligned to the slave interface.
- Changing the system memory map does not involve manually editing HDL.

Separate address-decoding logic is generated for every master in a system. The address decoding logic processes the difference between the master address width ($<M>$) and the individual slave address widths ($<S>$ and $<T>$). The address decoding logic also maps only the necessary master address bits to access words in each slave’s address space.

Figure 9–27 shows a block diagram of the address-decoding logic for one master and two slaves.

**Figure 9–27. Block Diagram of Address Decoding Logic**
In Qsys, the base addresses are controlled by the Base setting of active components on the System Contents tab, as shown in Figure 9–28.

Figure 9–28. Base Settings in Qsys Address Decoding

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Export</th>
<th>Clk. Base</th>
<th>End</th>
<th>IRQ</th>
</tr>
</thead>
<tbody>
<tr>
<td>clk_0</td>
<td>Clock Source</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>cpu</td>
<td>Nios II Processor</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>reset_n</td>
<td>Reset Input</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>data_master</td>
<td>Avalon Memory Mapped Master</td>
<td>Double-click to export</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>instruction_master</td>
<td>Avalon Memory Mapped Master</td>
<td>Double-click to export</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>tag_debug_module</td>
<td>Avalon Memory Mapped Slave</td>
<td>Double-click to export</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ext_flash</td>
<td>Flash Memory Interface (CFI)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ext_ram</td>
<td>DT71V416 SRAM</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ext_ram_bus</td>
<td>Generic Tri-State Controller</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>button_pio</td>
<td>PIO (Parallel I/O)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>high_res_timer</td>
<td>Interval Timer</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Datapath Multiplexing

Datapath multiplexing logic drives the writedata signal from the granted master to the selected slave, and the readdata signal from the selected slave back to the requesting master. Qsys generates separate datapath multiplexing logic for every master in the system.
Figure 9–29 shows a block diagram of the datapath multiplexing logic for one master and two slaves.

**Figure 9–29. Block Diagram of Datapath Multiplexing Logic**

---

**Wait State Insertion**

Wait states extend the duration of a transfer by one or more cycles. Wait state insertion logic accommodates the timing needs of each slave, and causes the master to wait until the slave can proceed. Qsys interconnect inserts wait states into a transfer when the target slave cannot respond in a single clock cycle, as well as in cases when slave read_enable and write_enable signals have setup or hold time requirements.

Wait state insertion logic is a small finite-state machine that translates control signal sequencing between the slave side and the master side. Qsys interconnect can force a master to wait for the wait state needs of a slave. For example, arbitration logic in a multimaster system. Qsys generates wait state insertion logic based on the properties of all slaves in the system.

Figure 9–30 shows a block diagram of the wait state insertion logic between one master and one slave.

**Figure 9–30. Block Diagram of Wait State Insertion Logic**
Interconnect Pipelining

If you set the **Limit interconnect pipeline stages** parameter to a value greater than 0 on the **Project Settings** tab, Qsys automatically inserts Avalon-ST pipeline stages when you generate your design. The pipeline stages increase the \( f_{\text{MAX}} \) of your design by reducing the combinational logic depth. The cost is additional latency and logic.

The insertion of pipeline stages depends upon the existence of certain interconnect components. For example, in a single-slave system, no multiplexer exists; therefore multiplexer pipelining does not occur. In an extreme case, of a single-master to single-slave system, no pipelining occurs, regardless of the value of **Limit interconnect pipeline stages**.

**Figure 9–31** shows the placement of up to four potential pipeline stages inserted by Qsys before the input to the demultiplexer, at the output of the multiplexer, between the arbiter and the multiplexer, and at the outputs of the demultiplexer.

**Figure 9–31. Pipeline Placement in Arbitration Logic**

[Diagram showing the placement of pipeline stages in arbitration logic.]

- Red lines indicate a pipeline stage, masters 0-3.
- Blue lines indicate a pipeline stage, selected request.

---

**Logic included in the Avalon-ST Command Network**

- Command packet for master 0
- Command packet for master 1
- Command packet for master 2
- Command packet for master 3
- Arbiter for slave 0
- Arbiter for slave 1
- Arbiter for slave 2
- Arbiter for slave 3
- Selected request

---

= Pipeline stage, masters 0-3

= Pipeline stage, selected request
For more information about pipelined Avalon-MM Interfaces, refer to the Optimizing Qsys System Performance chapter in volume 1 of the Quartus II Handbook.

AMBA AXI3 (version 1.0) Specification Support

Qsys allows connections between AXI3 components, AXI4 components, and Avalon memory-mapped interface types. The sections that follow describe unique or exceptional AXI3 support in the Qsys software.

Read this section in conjunction with the AMBA Protocol Specifications for AXI3 on the ARM website.

AXI3 Channels

Read and Write Address Channels

All signals are allowed with the following limitations:

- Supports 64-bit addressing.
- ID width limited to 18-bits.
- HPS-FPGA master interface has a 12-bit ID.

Write Data, Write Response, and Read Data Channels

All signals are allowed with the following limitations:

- Data widths limited to a maximum of 1024-bits
- Limited to a fixed byte width of 8-bits

Low Power Channel

Low power extensions are not supported in Qsys, version 13.0.

Cache Support

AWCACHE and ARCACHE are passed to an AXI slave unmodified.

Bufferable

The interconnect treats AXI transactions as non-bufferable; all responses must come from the terminal slave. When connecting to Avalon-MM slaves, since they do not have write responses, the following exceptions applies:

- For Avalon-MM slaves, the write response are generated by the slave agent once the write transaction is accepted by the slave. The following limitation exists for an Avalon bridge:

- For an Avalon bridge, the response is generated before the write reaches the endpoint; users must be aware of this limitation and avoid multiple paths past the bridge to any endpoint slave, or only perform bufferable transactions to an Avalon bridge.
Cacheable (Modifiable)

The interconnect acknowledges the cacheable (modifiable) attribute of AXI transactions. It does not change the address, burst length, or burst size of non-modifiable transactions, with the following exceptions:

- A wide transaction to a narrow slave is treated as modifiable because the size needs to be reduced.
- AXI read and write transactions might be treated as modifiable when the destination is an Avalon slave. The AXI transaction might be split into multiple Avalon transactions if the slave is unable to accept the transaction, which might occur because of burst lengths, narrow sizes, or burst types.

All other bits, for example, read allocate or write allocate, are ignored because the interconnect does not perform caching.

By default, transactions issued by Avalon masters are treated as non-bufferable and non-cacheable, with the allocate bits tied low. Qsys provides compile-time options to control the cache behavior of Avalon transactions on a per-master basis.

Security Support

TrustZone refers to the security extension of the ARM architecture, which includes the concept of "secure" and "non-secure" transactions, and a protocol for processing between the designations. TrustZone security support is a part of the Qsys 13.0 interconnect.

The interconnect passes the `AWPROT` and `ARPROT` signals to the endpoint slave without modification. It does not use or modify the `PROT` bits.

For more information about secure systems and the TrustZone feature, refer to the Creating a System With Qsys chapter in volume 1 of the Quartus II Handbook.

Atomic Accesses

Exclusive accesses are supported for AXI slaves by passing the lock, transaction ID, and response signals from master to slave, with the limitation that slaves that do not reorder responses.

Avalon slaves do not support exclusive accesses, and always return `OKAY` as a response.

Locked accesses are also not supported.

Response Signaling

Full response signalling is supported. Avalon slaves always return `OKAY` as a response.

Ordering Model

Qsys interconnect provides responses in the same order as the commands were issued.
To prevent reordering, for slaves that accept reordering depths greater than 0, Qsys does not transfer the transaction ID from the master, but provides a constant transaction ID of 0. For slaves that do not reorder, Qsys allows the transaction ID to be transferred to the slave.

To avoid cyclic dependencies, Qsys supports a single outstanding slave scheme for both reads and writes. Changing the targeted slave before all responses have returned stalls the master, regardless of transaction ID.

**AXI and Avalon Ordering**

According to the *AMBA Protocol Specifications*, there is no ordering requirement between reads and writes. However, Avalon has an implicit ordering model that requires transactions from a master to the same slave to be in order. As a result, there is a potential read-after-write risk when Avalon masters transact to AXI slaves.

In response to this potential risk, Avalon interfaces provide a compile-time option to enforce strict order. When turned on, the Avalon interface waits for outstanding write responses before issuing reads.

**Data Buses**

Narrow bus transfers are supported.

AXI write strobes can have any pattern that is compatible with the address and size information. Altera recommends that transactions to Avalon slaves follow Avalon byteenable limitations for maximum compatibility.

Byte 0 is always bits [7:0] in the interconnect, following AXI’s and Avalon’s byte (address) invariance scheme.

**Unaligned Address Commands**

Unaligned address commands are commands with addresses that do not conform to the data width of a slave. Since Avalon-MM slaves accept only aligned addresses, Qsys modifies unaligned commands from AXI masters to the correct data width. Qsys must preserve commands issued by AXI masters when passing the commands to AXI slaves.

Unaligned transfers are aligned if downsizing occurs. For example, when downsizing to a bus width narrower than that required by the transaction size, AWSIZE or ARSIZE, the transaction must be modified.

**Avalon and AXI Transactions**

The following sections describe limitations and solutions for AXI and Avalon transactions.

**Transaction Cannot Cross 4KB Boundaries**

When an Avalon master issues a transaction to an AXI slave, the transaction cannot cross 4KB boundaries. Non-bursting Avalon masters already follow this boundary restriction.
Handling Read Side Effects

Read side effects can occur when more bytes than necessary are read by the slave, and the unwanted data that are read are later inaccessible on subsequent reads.

For write commands, the correct byteenable paths are asserted based on the size of the transactions. For read commands, narrow-sized bursts are broken up into multiple non-bursting commands, and each command with the correct byteenable paths asserted.

Qsys always assumes that the byteenable is asserted based on the size of the command, not the address of the command. For example, for a 32-bit AXI master that issues a read command with unaligned address starting at address 0x01, and a burstcount of 2 to a 32-bit Avalon slave, are treated as having a starting address of 0x00.

AMBA AXI4 (version 2.0) Specification Support

Qsys allows connections between AXI4 components, and AXI3 and Avalon memory-mapped interface types. The sections that follow describe unique or exceptional AXI4 support in the Qsys software.

Read this section in conjunction with the AMBA Protocol Specifications for AXI4 on the ARM website.

Burst Support

Qsys supports for INCR bursts up to 256 beats. Qsys converts long bursts to multiple bursts in a packet with each burst having a length less than or equal to MAX_BURST when going to AXI3 or Avalon slaves.

For narrow-sized transfers, bursts with Avalon slaves as destinations are shortened to multiple non-bursting transactions in order to transmit the correct address to the slaves, since Avalon slaves always perform full-sized datawidth transactions.

Bursts with AXI3 slaves as destinations are shortened to multiple bursts, with each burst length less than or equal to 16. Bursts with AXI4 slaves as destinations are not shortened.

QoS

Qsys routes 4-bit QoS signals (Quality of Service Signaling) on the read and write address channels directly from the master to the slave.

Transactions from AXI3 and Avalon masters have a default value of 4'b0000, which indicates that the transactions are not part of the QoS flow. QoS values are not used for slaves that do not support QoS.

For Qsys 13.0, there are no programmable QoS registers or compile-time QoS options for a master that overrides its real or default value.
Regions

For Qsys 13.0, there is no support for the optional regions feature. AXI4 slaves with AXREGION signals are allowed. AXREGION signals are driven with the default value of 0x0, and are limited to one entry in a master’s address map.

Write Response Dependency

Write response dependency as specified in the AMBA Protocol Specifications for AXI4 is not supported.

AWCACHE and ARCACHE

For AXI4, Qsys meets the requirement for modifiable and non-modifiable transactions. The modifiable bit refers to ARCACHE[1] and AWCACHE[1].

Width Adaptation and Data Packing in Qsys

Data packing applies only to systems where the data width of masters is less than the data width of slaves, and the following applies:

- Data packing is supported when masters and slaves are Avalon-MM.
- Data packing is not supported when any master or slave is an AXI3, AXI4, or APB component.

For example, for a read/write command with a 32-bit master connected to a 64-bit slave, and a transaction of 2 burstcounts, Qsys sends 2 separate read/write commands to access the 64-bit data width of the slave. Data packing is only supported if the system does not contain AXI3, AXI4, or APB masters or slaves.

Ordering Model

Out of order support is not implemented in Qsys, version 13.0. Qsys processes AXI slaves as device non-bufferable memory types.

The following describes the required behavior for the device non-bufferable memory type:

- Write response must be obtained from the final destination.
- Read data must be obtained from the final destination.
- Transaction characteristics must not be modified.
- Reads must not be pre-fetched. Writes must not be merged.
- Non-modifiable read and write transactions.

(AWCACHE[1] = 0 or ARCACHE[1] = 0) from the same ID to the same slave must remain ordered. The interconnect always provides responses in the same order as the commands issued. Slaves that support reordering provide a constant transaction ID to prevent reordering. AXI slaves that do not reorder are provided with transaction IDs, which allows exclusive accesses to be used for such slaves.
Read and Write Allocate
Read and write allocate does not apply to Qsys interconnect, which does not have caching features, and always receives responses from an endpoint.

Locked Transactions
Locked transactions are not supported for Qsys, version 13.0.

Memory Types
For AXI4, Qsys processes transactions as though the endpoint is a device memory type. For device memory types, using non-bufferable transactions to force previous bufferable transactions to finish is irrelevant, because Qsys interconnect always identifies transactions as being non-bufferable.

Mismatched Attributes
There are rules for how multiple masters issue cache values to a shared memory region. The interconnect meets requirements as long as cache signals are not modified.

Signals
Qsys supports up to 64-bits for the BUSER, WUSER and RUSER sideband signals. AXI4 allows some signals to be omitted from interfaces by aligning them with the default values as defined in the AMBA Protocol Specifications on the ARM® website.

AMBA APB (version 1.0) Specification Support
AMBA APB provides a low-cost interface that is optimized for minimal power consumption and reduced interface complexity. You can use AMBA APB to interface to peripherals which are low-bandwidth and do not require the high performance of a pipelined bus interface. Signal transitions are sampled at the rising edge of the clock to enable the integration of APB peripherals easily into any design flow.

Qsys allows connections between APB components, and Avalon, AXI3, and AXI4 Avalon memory-mapped interface types. The following sections describe unique or exceptional APB support in the Qsys software.

Read this section in conjunction with the AMBA Protocol Specifications for APB on the ARM website.

Bridges
With APB, you cannot use bridge components that use multiple PSELx in Qsys. As a workaround, you can group PSELx, and then send the packet to the slave directly.

Altera recommends as an alternative that you instantiate the APB bridge and all the APB slaves in Qsys. You should then connect the slave side of the bridge to any high speed interface and connect the master side of the bridge to the APB slaves. Qsys creates the interconnect on either side of the APB bridge and creates only one PSEL signal.
Alternatively, you can connect a bridge to the APB bus outside of Qsys. Use an Avalon/AXI bridge to export the Avalon/AXI master to the top-level, and then connect this Avalon/AXI interface to the slave side of the APB bridge. Alternatively, instantiate the APB bridge in Qsys and export APB master to the top-level, and from there connect to APB bus outside of Qsys.

**Burst Adaptation**

APB is a non-bursting interface. Therefore, for any AXI or Avalon master with bursting support, a burst adapter is inserted before the slave interface and the burst transaction is translated into a series of non-bursting transactions before reaching the APB slave.

**Width Adaption**

Qsys allows different data width connections with APB. When connecting a wider master to a narrower APB slave, the width adapter converts the wider transactions to a narrower transaction to fit the APB slave data width.

APB does not support Write Strobe. Therefore, when connecting a narrower transaction to a wider APB slave, the slave cannot determine which byte lane to write, so the data at the slave might be overwritten or corrupted.

**Error Response**

Error responses are returned to the master. Qsys performs error mapping if the master is an AXI3 or AXI4 master, for example, \(RRESP/BRESP=SLVERR\). For the case when the slave does not use \(SLVERR\) signal, an \(OKAY\) response is sent back to master by default.

**Document Revision History**

Table 9–11 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>- Added AMBA APB support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added auto-inserted Avalon-ST adapters feature.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added AMBA AXI4 support.</td>
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<td>November 2012</td>
<td>12.1.0</td>
<td>- Added AMBA AXI3 support.</td>
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<td>- Added Avalon-ST.</td>
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<td></td>
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<td>- Added Address Span Extender.</td>
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<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
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<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Removed beta status.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>
For previous versions of the *Quartus II Handbook*, refer to the *Quartus II Handbook Archive.*
This chapter provides information on optimizing system interconnect performance for designs generated by the Altera® Qsys system integration tool.

The foundation of any large system is the interconnect logic used to connect hardware blocks or components. Creating interconnect logic is prone to errors, is time consuming to write, and is difficult to modify when design requirements change. The Qsys system integration tool addresses these issues by providing an automatically generated and optimized interconnect designed to satisfy your system requirements.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

For more discussion about determining which interface standard you want to use to create your Qsys design, refer to the Creating a System With Qsys chapter in volume 1 of the Quartus II Handbook.

Following the design practices recommended in this chapter may improve the clock frequency, throughput, logic utilization, or power consumption of your Qsys design. When you design a Qsys system, use your knowledge of your design intent and goals to further optimize system performance beyond the automated optimization available in Qsys.

The following sections describe Qsys support for optimization of interconnect logic:

- “Designing with Avalon and AXI Interfaces” on page 10–1
- “Using Hierarchy in Systems” on page 10–3
- “Using Concurrency in Memory-Mapped Systems” on page 10–5
- “Insert Pipeline Stages to Increase System Frequency” on page 10–10
- “Using Avalon Bridges” on page 10–10
- “Increasing Transfer Throughput” on page 10–21
- “Reducing Logic Utilization” on page 10–28
- “Reducing Power Consumption” on page 10–34
- “Design Examples” on page 10–39

**Designing with Avalon and AXI Interfaces**

Qsys Avalon and AXI interconnect for memory-mapped interfaces is flexible, partial crossbar logic that connects master and slave interfaces.
Avalon Streaming (Avalon-ST) links connect point-to-point, unidirectional interfaces, and are typically used in data stream applications. Each a pair of components is connected without any requirement to arbitrate between the data source and sink.

Because Qsys supports multiplexed memory-mapped and streaming connections, you can implement systems that use multiplexed logic for control and streaming logic for data in a single design.

For more information about designing streaming and memory-mapped components, refer to the Creating Qsys Components chapter in volume 1 of the Quartus II Handbook.

### Designing Streaming Components

When you design streaming component interfaces, you must consider integration and communication for each component in the system. One common consideration is buffering data internally to accommodate latency between components. For example, if the component’s Avalon-ST output or source of streaming data is back-pressured because the ready signal is deasserted, then the component must back-pressure its input or sink interface to avoid overflow.

You can use a FIFO to back-pressure internally on the output side of the component, so that the input can accept more data even if the output is back-pressured. Then, you use the FIFO almost full flag to back-pressure the sink interface or input data when the FIFO has only enough space left to satisfy the internal latency. You drive the data valid signal of the output or source interface with the not empty flag of the FIFO when that data is available.

AXI streaming and bridge components are not available in the Quartus II software, version 12.1.

### Designing Memory-Mapped Components

When designing with memory-mapped components, “Example of Control and Status Registers (CSR) in a Slave Component” on page 10–3 is an example that you can use to implement any component that contains multiple registers mapped to memory locations. Components that implement read and write memory-mapped transactions require three main building blocks: an address decoder, a register file, and a read multiplexer.
Figure 10–1 shows how to implement a set of four output registers to support software read back from logic.

**Figure 10–1. Example of Control and Status Registers (CSR) in a Slave Component**

The decoder enables the appropriate 32-bit or 64-bit register for writes. For reads, the address bits drive the multiplexer selection bits. The `read` signal registers the data from the multiplexer, adding a pipeline stage so that the component can achieve a higher clock frequency. This component has write wait states and one read wait state. Alternatively, if you want high throughput, you might set both the read and write wait states to zero, and then specify a read latency of one, because the component also supports pipelined reads.

**Using Hierarchy in Systems**

You can use hierarchy to sub-divide a system into smaller subsystems that can be connected together in a top-level Qsys system. You can use hierarchy to simplify verification control of slaves connected to each master in a memory-mapped system.

Before you begin implementing subsystems in your design, you should plan the system hierarchical blocks at the top level, using the following guidelines:
■ **Plan shared resources**—For example, determine the best location for shared resources in the system hierarchy. For example, if two subsystems share resources, you should add the components that use those resources to a higher-level system for easy access.

■ **Plan shared address space between subsystems**—Planning the address space ensures you can set appropriate sizes for bridges between subsystems.

■ **Plan how much latency you might add to your system**—When you add a pipeline bridge between subsystems, you might add more latency to the overall system. You can reduce the added latency by parameterizing the pipeline bridge with zero cycles of latency.

Figure 10–2 shows an example of two Nios II processor subsystems with shared resources for message passing. Bridges in each subsystem export the Nios II data master to the top-level system that includes the mutex (mutual exclusion component) and shared memory component (which could be another on-chip RAM, or a controller for an off-chip RAM device).

**Figure 10–2. Message Passing Between Subsystems**

If a design contains one or more identical functional units, the functional unit can be defined as a subsystem and instantiated multiple times within a top-level system. You can also design systems that process multiple data channels by instantiating the same subsystem for each channel. This approach is easier to maintain than a larger, non hierarchical system. In addition, such systems are easier to scale because you can calculate the required resources as a simple multiple of the subsystem requirements.
Figure 10–3 shows a design with three subsystems, each processing a unique channel.

Using Concurrency in Memory-Mapped Systems

Qsys interconnect takes advantage of parallel hardware in FPGAs, which allows you to design concurrency into your system and process multiple transactions simultaneously. The following sections describe design choices that can increase concurrency in your system.

Create Multiple Masters

Implementing concurrency requires multiple masters in the system. Systems that include a processor contain at least two master interfaces because the processors include separate instruction and data masters. Master components can be categorized as follows:

- General purpose processors, such as Nios II processors
- DMA (direct memory access) engines
- Communication interfaces, such as PCI Express
Because Qsys generates an interconnect with slave-side arbitration, every master interface in your system can issue transfers concurrently. Masters in the system can issue transfers concurrently as long as they are not posting transfers to the same slave. Concurrency is limited by the number of master interfaces sharing any particular slave interface. If your design requires higher data throughput, you can increase the number of master and slave interfaces to increase the number of transfers that occur simultaneously. Refer to “Create Multiple Slave Interfaces” on page 10–8 for more information.

Figure 10–4 shows a system with three master interfaces. The lines are examples of connections that can be active simultaneously.

**Figure 10–4. Avalon Multiple Master Parallel Access**

In this Avalon example, the DMA engine operates with Avalon-MM read and write masters. However, an AXI DMA interface typically has only one master, because in the AXI standard the write and read channels on the master are independent and can process transactions simultaneously.
Figure 10–5 shows an AXI example where the DMA engine operates with a single master, because in AXI the write and read channels on the master are independent and can process transactions simultaneously. This example shows concurrency between the read and write channels, with the yellow lines representing concurrent data paths.

**Figure 10–5. AXI Multi Master Parallel Access**
Create Multiple Slave Interfaces

You can create multiple slave interfaces for a particular function to increase concurrency in your design. Figure 10–6 shows two channel processing systems. In the first, four hosts must arbitrate for the single slave interface of the channel processor. In the second, each host drives a dedicated slave interface, allowing all master interfaces to simultaneously access the slave interfaces of the component. Arbitration is not necessary when there is a single host and slave interface.

Figure 10–6. Single Interface Vs Multiple Interfaces
Use DMA Engines

In some systems, you can use DMA engines to increase throughput. You can use a DMA engine to transfer blocks of data between interfaces, which then frees the CPU from carrying out this routine task. A DMA engine transfers data between a programmed start and end address without intervention, and the data throughput is dictated by the components connected to the DMA. Factors that affect data throughput include data width and clock frequency. Figure 10–7 shows a system that can sustain more concurrent read and write operations by including more DMA engines, for the case that accesses to the read and write buffers in the top system can be split between two DMA engines, as shown in the Dual DMA Channels system at the bottom of the figure.

In this example, the DMA engine operates with Avalon-MM write and read masters. An AXI DMA typically has only one master, because in AXI the write and read channels on the master are independent and can process transactions simultaneously.

Figure 10–7. Single or Dual DMA Channels

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![Single DMA Channel Diagram](image1)

**Single DMA Channel**
Maximum of One Read & One Write Per Clock Cycle

**Dual DMA Channels**
Maximum of two Reads & Two Writes Per Clock Cycle
Insert Pipeline Stages to Increase System Frequency

Qsys provides the **Limit interconnect pipeline stages to** option on the **Project Settings** tab to automatically add pipeline stages to the Qsys interconnect when you generate your system. You can specify between 0 to 4 pipeline stages, where 0 means that the interconnect has a combinational data path. You can specify a unique interconnect pipeline stage value for each subsystem.

Adding pipeline stages might increase the $f_{\text{MAX}}$ of your design by reducing the combinational logic depth, at the cost of additional latency and logic utilization.

The insertion of pipeline stages requires certain interconnect components. For example, in a system with a single slave interface, there is no multiplexer; therefore multiplexer pipelining does not occur. When there is an Avalon or AXI single-master to single-slave system, no pipelining occurs, regardless of the **Limit interconnect pipeline stages to** parameter.

> For more information about the **Limit interconnect pipeline stages to** parameter, refer to the **Qsys Interconnect** chapter in volume 1 of the **Quartus II Handbook**.

Using Avalon Bridges

You can use bridges to increase system frequency, minimize generated Qsys logic, minimize adapter logic, and to structure system topology when you want to control where Qsys adds pipelining. You can also use bridges with arbiters when there is concurrency in the system.

> AXI bridges are not supported in the Quartus II software, version 12.1; however, you can use Avalon bridges between AXI interfaces, and between Avalon domains. Qsys automatically creates interconnect logic between the AXI and Avalon interfaces, so you do not have to explicitly instantiate bridges between these domains. For more discussion about the benefits and disadvantages of shared and separate domains, refer to the **Qsys Interconnect** chapter in volume 1 of the **Quartus II Handbook**.

An Avalon bridge has an Avalon-MM slave interface and an Avalon-MM master interface. You can have many components connected to the bridge slave interface, or many components connected to the bridge master interface, or a single component connected to a single bridge slave or master interface. You can configure the data width of the bridge, which can affect how Qsys generates bus sizing logic in the interconnect. Both interfaces support Avalon-MM pipelined transfers with variable latency, and can also support configurable burst lengths.

Transfers to the bridge slave interface are propagated to the master interface, which connects to components downstream from the bridge. When you need greater control over the interconnect pipelining, you can use bridges instead of using the **Limit Interconnect Pipeline Stages to** parameter.

Increasing System Frequency

In Qsys, you can introduce interconnect pipeline stages or pipeline bridges to increase clock frequency in your system. Bridges control the system interconnect topology and allow you to subdivide the interconnect, giving you more control over pipelining and clock crossing functionality.
**Insert Pipeline Bridges**

You can insert an Avalon-MM pipeline bridge to insert registers in the path between the bridges and its master and slaves. If a critical register-to-register delay occurs in the Qsys interconnect, a pipeline bridge can help reduce this delay and improve system $f_{MAX}$.

The Avalon-MM pipeline bridge component integrates into any Qsys system. The pipeline bridge options can increase logic utilization and read latency. The change in topology may also reduce concurrency if multiple masters arbitrate for the bridge.

You can use the Avalon-MM pipeline bridge to control topology without adding a pipeline stage. A pipeline bridge that does not add a pipeline stage is optimal in some latency-sensitive applications. For example, a CPU may benefit from minimal latency when accessing memory.

*Figure 10–8* shows the architecture of an Avalon-MM pipeline bridge.

*Figure 10–8. Avalon-MM Pipeline Bridge*
Implement Command Pipelining (Master-to-Slave)

When many masters share a slave device, use command pipelining to improve performance. The arbitration logic for the slave interface must multiplex the address, writedata, and burstcount signals. The multiplexer width increases proportionally with the number of masters connecting to a single slave interface. The increased multiplexer width might become a timing critical path in the system. If a single pipeline bridge does not provide enough pipelining, you can instantiate multiple instances of the bridge in a tree structure to increase the pipelining and further reduce the width of the multiplexer at the slave interface, as Figure 10–9 shows.

Figure 10–9. Tree of Bridges

Response Pipelining (Slave-to-Master)

A system can benefit from slave-to-master pipelining for masters that connect to many slaves that support read transfers. The interconnect inserts a multiplexer for every read data path back to the master. As the number of slaves supporting read transfers connecting to the master increases, so does the width of the read data multiplexer. As with master-to-slave pipelining, if the performance increase is insufficient with one bridge, you can use multiple bridges in a tree structure to improve $f_{\text{MAX}}$.

Use Clock Crossing Bridges

Transfers to the slave interface are propagated to the master interface. The clock crossing bridge contains a pair of clock crossing FIFOs, which isolate the master and slave interfaces in separate, asynchronous clock domains.
When you use a FIFO clock crossing bridge for the clock domain crossing, you add data buffering. Buffering allows pipelined read masters to post multiple reads to the bridge, even if the slaves downstream from the bridge do not support pipelined transfers.

**Separate Component Frequencies**

You can use a clock crossing bridge to place high and low frequency components in separate clock domains. If you limit the fast clock domain to the portion of your design that requires high performance, you might achieve a higher \( f_{\text{MAX}} \) for this portion of the design.

For example, the majority of processor peripherals included in embedded designs do not need to operate at high frequencies, therefore you do not need to use a high-frequency clock for these components. When you compile a design with the Quartus II software, compilation may take more time when the clock frequency requirements are difficult to meet because the Fitter needs more time to place registers to achieve the required \( f_{\text{MAX}} \). To reduce the amount of effort that the Fitter uses on low priority and low performance components, you can place these behind a clock crossing bridge operating at a lower frequency, allowing the Fitter to increase the effort placed on the higher priority and higher frequency data paths.

**Minimize Design Logic**

Bridges can reduce the interconnect logic by reducing the amount of arbitration and multiplexer logic that Qsys generates. This reduction occurs because bridges limit the number of concurrent transfers that can occur. The following sections discuss how you can use bridges to minimize the logic generated by Qsys.

**Avoid Speed Optimizations That Increase Logic**

Adding an additional pipeline stage with a pipeline bridge between masters and slaves reduces the amount of combinational logic between registers, which can increase system performance, as described in the section “Increasing System Frequency” on page 10–10.

If you can increase the \( f_{\text{MAX}} \) of your design logic, you may be able to turn off the Quartus II optimization settings, such as the Perform register duplication setting. Register duplication creates duplicate registers to be placed in two or more physical locations in the FPGA to reduce register-to-register delays. You might also want to choose Speed for the optimization method, which typically results in higher logic utilization due to logic duplication. By making use of the registers or FIFOs available in the Avalon-MM bridges, you can increase the design speed and avoid needless logic duplication or speed optimizations, thereby reducing the logic utilization of the design.

**Reduced Concurrency**

The amount of logic generated for the interconnect often increases as the system becomes larger because Qsys creates arbitration logic for every slave interface that is shared by multiple master interfaces. Qsys inserts multiplexer logic between master interfaces that connect to multiple slave interfaces if both support read data paths. Most embedded processor designs contain components that are either incapable of...
supporting high data throughput, or do not need to be accessed frequently. These components can contain Avalon-MM master or slave interfaces. Because the interconnect supports concurrent accesses, you might want to limit concurrency by inserting bridges into the datapath to limit the amount of arbitration and multiplexer logic generated.

For example, if your system contains three masters and three slave interfaces that are interconnected, Qsys generates three arbiters and three multiplexers for the read data path. If these masters do not require a significant amount of simultaneous throughput, you can reduce the resources that your design consumes by connecting the three masters to a pipeline bridge. The bridge masters the three slave interfaces, and reduces the interconnect into a bus structure. Qsys creates one arbitration block between the bridge and the three masters, and a single read data path multiplexer between the bridge and three slaves, and prevents concurrency; similar to that of a standard bus architecture. You should not use this method for high throughput data paths to ensure that you do not limit overall system performance.

Figure 10–10 shows the difference in architecture between systems with or without a pipeline bridge.

**Figure 10–10. Switch Interconnect to Bus**
Minimizing Adapter Logic

Qsys generates adapter logic for clock crossing, width adaptation, and burst support when there is a mismatch between the clock domains, widths, or bursting capabilities of the master and slave interface pairs. Qsys creates burst adapters when the maximum burst length of the master is greater than the master burst length of the slave. The adapter logic creates extra logic resources, which can be substantial when your system contains master interfaces connected to many components that do not share the same characteristics. By placing bridges in your design, you can reduce the amount of adapter logic that Qsys generates.

Effective Placement of Bridges

To determine the effective placement of a bridge, you should initially analyze each master in your system to determine if the connected slave devices support different bursting capabilities or operate in a different clock domain. The maximum burstcount of a component is visible as the burstcount signal in the HDL file of the component. The maximum burst length is 2 (width(burstcount -1)), so that if the burstcount width is four bits, the maximum burstcount is eight. If no burstcount signal is present, the component does not support bursting or has a burst length of 1.

To determine if the system requires a clock crossing adapter between the master and slave interfaces, check the clock column beside the master and slave interfaces in Qsys. If the clock is different for the master and slave interfaces, Qsys inserts a clock crossing adapter between them. To avoid creating multiple adapters, you can place the components containing slave interfaces behind a bridge so that only one adapter is created. By placing multiple components with the same burst or clock characteristics behind a bridge, you limit concurrency and the number of adapters.

You can use a bridge to separate AXI and Avalon domains to minimize burst adaptation logic. For example, if there are multiple Avalon slaves that are connected to an AXI master, you can consider inserting a bridge to access the adaptation logic once before the bridge, instead once per slave. This costs latency, though, and you would also lose concurrency between reads and writes.

Changing the Response Buffer Depth

When you use automatic clock-crossing adapters, Qsys determines the required depth of FIFO buffering based on the slave properties. If a slave has a high Maximum Pending Reads parameter, the resulting deep response buffer FIFO that Qsys inserts between the master and slave can consume a lot of device resources. To control the response FIFO depth, you can use a clock crossing bridge and manually adjust its FIFO depth to trade off throughput with smaller memory utilization. For example, if you have masters that cannot saturate the slave, you do not need response buffering, so that using a bridge reduces the FIFO memory depth and reduces the Maximum Pending Reads available from the slave.
Consequences of Using Bridges

Before you use pipeline or clock crossing bridges in your design, you should carefully consider their effects. Bridges can have any combination of the following consequences on your design, which could be positive or negative. You can benchmark your system before and after inserting bridges to determine their impact. The following sections discuss the possible consequences of adding bridges to your system.

Increased Latency

Adding a bridge to your design has an effect on the read latency between the master and the slave. Depending on the system requirements and the type of master and slave, this latency increase may or may not be acceptable in your design.

Acceptable Latency Increase

For a pipeline bridge, a cycle of latency is added for each pipeline option that is enabled. The buffering in the clock crossing bridge also adds latency. If you use a pipelined or burst master that posts many read transfers, the increase in latency does not impact performance significantly because the latency increase is very small compared to the length of the data transfer.

For example, if you use a pipelined read master such as a DMA controller to read data from a component with a fixed read latency of four clock cycles, but only perform a single word transfer, the overhead is three clock cycles out of the total four, assuming there is no additional pipeline latency in the Qsys interconnect. The read throughput is only 25%. Figure 10–11 shows this type of low-efficiency read transfer.

However, if 100 words of data are transferred without interruptions, the overhead is three cycles out of the total of 103 clock cycles, corresponding to a read efficiency of approximately 97% when there is no additional pipeline latency in the interconnect. Adding a pipeline bridge to this read path adds two extra clock cycles of latency. The transfer requires 105 cycles to complete, corresponding to an efficiency of approximately 94%. Although the efficiency decreased by 3%, adding the bridge
might increase the f\textsubscript{MAX} by 5\%, for example, and in that case, if the clock frequency can be increased, the overall throughput would improve. As the number of words transferred increases, the efficiency increases to nearly 100\%, whether or not a pipeline bridge is present. Figure 10–12 shows this type of high-efficiency read transfer.

**Unacceptable Latency Increase**

Processors are sensitive to high latency read times and typically fetch data for use in calculations that cannot proceed until the data arrives. Before adding a bridge to the data path of a processor instruction or data master, determine whether the clock frequency increase justifies the added latency. Figure 10–13 shows the performance of a Nios II processor and memory operating at 100 MHz. The Nios II processor instruction master has a cache memory with a read latency of four cycles, that is eight sequential words of data return for each read. At 100 MHz, the first read takes 40 ns to complete. Each successive word takes 10 ns so that eight reads complete in 110 ns.

![Figure 10–12. High Efficiency Read Transfer](image)

![Figure 10–13. Processor System: Eight Reads with Four Cycles Latency](image)
Adding a clock crossing bridge allows the memory to operate at 125 MHz in this example. However, this increase in frequency is negated by the increase in latency for the following reasons, as shown in Figure 10–14. If the clock crossing bridge adds six clock cycles of latency at 100 MHz, then the memory continues to operate with a read latency of four clock cycles; consequently, the first read from memory takes 100 ns, and each successive word takes 10 ns because reads arrive at the frequency of the processor, which is 100 MHz. In total, eight reads complete after 170 ns. Although the memory operates at a higher clock frequency, the frequency at which the master operates limits the throughput.

Figure 10–14. Processor System: Eight Reads with Ten Cycles Latency

**Limited Concurrency**

Placing an bridge between multiple master and slave interfaces limits the number of concurrent transfers your system can initiate. This limitation is the same as connecting multiple master interfaces to a single slave interface. The slave interface of the bridge is shared by all the masters and, as a result, Qsys creates arbitration logic. If the components placed behind a bridge are infrequently accessed, this concurrency limitation might be acceptable.

Bridges can have a negative impact on system performance if you use them inappropriately. For example, if multiple memories are used by several masters, you should not place the memory components behind a bridge. The bridge limits memory performance by preventing concurrent memory accesses. Placing multiple memory components behind a bridge can cause the separate slave interfaces to appear as one large memory to the masters accessing the bridge; all masters must access the same slave interface.
Figure 10–15 shows a memory subsystem with one bridge that acts as a single slave interface for the Avalon-MM Nios II and DMA masters, which results in a bottleneck architecture. The bridge acts as a bottleneck between the two masters and the memories. An AXI DMA typically has only one master, because in the AXI standard the write and read channels on the master are independent and can process transactions simultaneously.

**Figure 10–15. Inappropriate Use of a Bridge in a Hierarchical System**
If the $f_{\text{MAX}}$ of your memory interfaces is low and you want to use a pipeline bridge between subsystems, you can place each memory behind its own bridge, which increases the $f_{\text{MAX}}$ of the system without sacrificing concurrency, as Figure 10–16 shows.

**Figure 10–16. Efficient Memory Pipelining Without a Bottleneck in a Hierarchical System**

<table>
<thead>
<tr>
<th>Address Space Translation</th>
</tr>
</thead>
<tbody>
<tr>
<td>The slave interface of a pipeline or clock crossing bridge has a base address and address span. You can set the base address or allow Qsys to set it automatically. The address of the slave interface is the base offset address of all the components connected to the bridge. The address of components connected to the bridge is the sum of the base offset and the address of that component.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Address Shifting</th>
</tr>
</thead>
<tbody>
<tr>
<td>The master interface of the bridge drives only the address bits that represent the offset from the base address of the bridge slave interface. Any time a master accesses a slave through a bridge, both addresses must be added together, otherwise the transfer fails. The <strong>Address Map</strong> tab in Qsys displays the addresses of the slaves connected to each master and includes address translations caused by system bridges.</td>
</tr>
</tbody>
</table>
Figure 10–17 shows how address translation functions. In this example, the Nios II processor connects to a bridge located at base address 0x1000, a slave connects to the bridge master interface at an offset of 0x20, and the processor performs a write transfer to the fourth 32-bit or 64-bit word within the slave. Nios II drives the address 0x102C to interconnect, which is within the address range of the bridge. The bridge master interface drives 0x2C, which is within the address range of the slave, and the transfer completes.

Address Coherency

To simplify the system design, all masters should access slaves at the same location. In many systems, a processor passes buffer locations to other mastering components, such as a DMA controller. If the processor and DMA controller do not access the slave at the same location, Qsys must compensate for the differences.

In Figure 10–18, a Nios II processor and DMA controller access a slave interface located at address 0x20. The processor connects directly to the slave interface. The DMA controller connects to a pipeline bridge located at address 0x1000, which then connects to the slave interface. Because the DMA controller accesses the pipeline bridge first, it must drive 0x1020 to access the first location of the slave interface. Because the processor accesses the slave from a different location, you must maintain two base addresses for the slave device.
To avoid the requirement for two addresses, you can add an additional bridge to the system, set its base address to 0x1000, and then disable all the pipelining options in the second bridge so that the bridge has minimal impact on system timing and resource utilization. Because this second bridge has the same base address as the original bridge, the DMA controller connects to both the processor and DMA controller and accesses the slave interface with the same address range, as shown in Figure 10–19.

Figure 10–19. Address Translation Corrected With Bridge

Increasing Transfer Throughput

Increasing the transfer efficiency of the master and slave interfaces in your system increases the throughput of your design. Designs with strict cost or power requirements benefit from increasing the transfer efficiency because you can then use less expensive, lower frequency devices. Designs requiring high performance also benefit from increased transfer efficiency because increased efficiency improves the performance of frequency–limited hardware.

Throughput is the number of symbols (such as bytes) of data that can be transferred in a given clock cycle of time period. Read latency is the number of clock cycles between the address and data phase of a transaction. For example, a read latency of two means that the data is valid two cycles after the address is posted. If the master has to wait for one request to finish before the next begins, such as with a processor, then the read latency is very important to the overall throughput.

You can measure throughput and latency in simulation by observing the waveforms, or using the verification IP monitors. For more information, refer to the *Avalon Verification IP Suite User Guide* or the *Mentor Graphics AXI Verification IP Suite - Altera Edition* on the Altera website.
Using Pipelined Transfers

Pipelined transfers increase the read efficiency by allowing a master to post multiple reads before data from an earlier read returns.

Masters that support pipelined transfers post transfers continuously, relying on the readdatavalid signal to indicate valid data. Avalon-MM slaves support pipelined transfers by including the readdatavalid signal or operating with a fixed read latency.

AXI masters declare how many outstanding writes and reads it can issue with the writeIssuingCapability and readIssuingCapability parameters. In the same way, a slave can declare how many reads it can accept with the readAcceptanceCapability parameter.

AXI masters with a read issuing capability greater than one are pipelined in the same way as Avalon masters and the readdatavalid signal.

Using the Maximum Pending Reads Parameter

If you create a custom component with a slave interface supporting variable-latency reads, you must specify the Maximum Pending Reads parameter in the Component Editor. Qsys uses the Maximum Pending Reads parameter to generate the appropriate interconnect, and represents the maximum number of read transfers that your pipelined slave component can process. If the number of reads presented to the slave interface exceeds the Maximum Pending Reads parameter, then the slave interface must assert waitrequest.

Optimizing the value of the Maximum Pending Reads parameter requires a good understanding of the latencies of your custom components. This parameter should be based on the component’s highest read latency for the various logic paths inside the component. For example, if your pipelined component has two modes, one requiring two clock cycles and the other five, set the Maximum Pending Reads parameter to 5, which allows your component to pipeline five transfers, eliminating dead cycles after the initial five-cycle latency.

You can also determine the correct value for the Maximum Pending Reads parameter by monitoring the number of reads that are pending during system simulation or while running the hardware. To use this method, set the Maximum Pending Reads to a very high value and use a master that issues read requests on every clock. You can use a DMA for this task as long as the data is written to a location that does not frequently assert waitrequest. If you implement this method with the hardware, you can observe your component with a logic analyzer or built-in monitoring hardware.

Choosing the correct value for the Maximum Pending Reads parameter of your custom pipelined read component is important. If you underestimate the Maximum Pending Reads value, you might cause a master interface to stall with a waitrequest until the slave responds to an earlier read request and frees a FIFO position.

The Maximum Pending Reads parameter controls the depth of the response FIFO inserted into the interconnect for each master connected to the slave. This FIFO does not use significant hardware resources. Overestimating the Maximum Pending Reads parameter for your custom component results in a slight increase in hardware utilization. For these reasons, if you are not sure of the optimal value, you should overestimate this value.
If your system includes a bridge, you must set the **Maximum Pending Reads** parameter on the bridge as well. To allow maximum throughput, this value should be equal to or greater than the **Maximum Pending Reads** value for the connected slave that has the highest value. As described in “Changing the Response Buffer Depth” on page 10–15, you can limit the maximum pending reads of a slave and reduce the buffer depth by reducing the parameter value on the bridge if the high throughput is not required. If you do not know the **Maximum Pending Reads** value for all your slave components, you can monitor the number of reads that are pending during system simulation while running the hardware. To use this method, set the **Maximum Pending Reads** parameter to a high value and use a master that issues read requests on every clock, such as a DMA. Then, reduce the number of maximum pending reads of the bridge until the bridge reduces the performance of any masters accessing the bridge.

**Arbitration Shares and Bursts**

Arbitration shares provide control over the arbitration process. By default, the arbitration algorithm allocates evenly, with all masters receiving one share.

You can adjust the arbitration process to your system requirements by assigning a larger number of shares to the masters that need greater throughput. The larger the arbitration share, the more transfers are allocated to the master to access a slave. The masters get uninterrupted access to the slave for its number of shares, as long as the master is transacting (reading or writing).

If a master cannot post a transfer and other masters are waiting to gain access to a particular slave, the arbiter grants another master access. This mechanism prevents a master from wasting arbitration cycles if it cannot post back-to-back transfers. A bursting transaction contains multiple beats (or words) of data, starting from a single address. Bursts allow a master to maintain access to a slave for more than a single word transfer. If a bursting master posts a write transfer with a burst length of eight, it is guaranteed arbitration for eight write cycles.

You can assign arbitration shares to Avalon-MM bursting master and AXI masters (which are always considered a bursting master). Each share consists of one burst transaction (such as multi-cycle write), and allows a master to complete a number of bursts before arbitration switches to the next master.

For more information about arbitration shares and bursts, refer to the *Avalon Interface Specifications*, or the *AMBA Protocol Specification* on the ARM website.

**Differences Between Arbitration Shares and Bursts**

The following three key characteristics distinguish arbitration shares and bursts:

- Arbitration lock
- Sequential addressing
- Burst adapters
Arbitration Lock

When a master posts a burst transfer, the arbitration is locked for that master; consequently, the bursting master should be capable of sustaining transfers for the duration of the locked period. If, after the fourth write, the master deasserts the write (Avalon-MM write or AXI wvalid) signal for fifty cycles, all other masters continue to wait for access during this stalled period.

To avoid wasted bandwidth, your master designs should wait until a full burst transfer is ready before requesting access to a slave device. Alternatively, you can avoid wasted bandwidth by posting burstcounts equal to the amount of data that is ready. For example, if you create a custom bursting write master with a maximum burstcount of eight, but only three words of data are ready, you can simply present a burstcount of three. This strategy does not result in optimal use of the system bandwidth if the slave is capable of handling a larger burst; however, this strategy prevents stalling and allows access for other masters in the system.

Avalon-MM Sequential Addressing

An Avalon-MM burst transfer includes a base address and a burstcount. The burstcount represents the number of words of data to be transferred, starting from the base address and incrementing sequentially. Burst transfers are common for processors, DMAs, and buffer processing accelerators; however, sometimes when a master must access non-sequential addresses. Consequently, a bursting master must set the burstcount to the number of sequential addresses, and then reset the burstcount for the next location.

The arbitration share algorithm has no restrictions on addresses; therefore, your custom master can update the address it presents to the interconnect for every read or write transaction.

AXI has different burst types than the Avalon interface. For more information about AXI burst types, refer to the Qsys Interconnect chapter in volume 1 of the Quartus II Handbook, and the AMBA AXI Protocol Specification on the ARM website.

Burst Adapters

Qsys allows you to create systems that mix bursting and non-bursting master and slave interfaces. This design strategy allows you to connect bursting master and slave interfaces that support different maximum burst lengths, and Qsys generates burst adapters when appropriate.

Qsys inserts a burst adapter whenever a master interface burst length exceeds the burst length of the slave interface, or if the master issues a burst type that the slave cannot support. For example, if you connect an AXI master to an Avalon slave, a burst adapter is inserted.

Qsys assigns non-bursting masters and slave interfaces a burst length of one. The burst adapter divides long bursts into shorter bursts. As a result, the burst adapter adds logic to the address and burstcount paths between the master and slave interfaces.
Choosing Avalon-MM Interface Types

To avoid inefficient Avalon-MM transfers, custom master or slave interfaces must use the appropriate simple, pipelined, or burst interfaces. The three possible transfer types are described below.

Simple Avalon-MM Interfaces

Simple interface transfers do not support pipelining or bursting for reads or writes; consequently, their performance is limited. Simple interfaces are appropriate for transfers between masters and infrequently used slave interfaces. In Qsys, the PIO, UART, and Timer include slave interfaces that use simple transfers.

Pipelined Avalon-MM Interfaces

Pipelined read transfers allow a pipelined master interface to start multiple read transfers in succession without waiting for the prior transfers to complete. Pipelined transfers allow master-slave pairs to achieve higher throughput, even though the slave port might require one or more cycles of latency to return data for each transfer.

In many systems, read throughput becomes inadequate if simple reads are used and pipelined transfers can increase throughput. If you define a component with a fixed read latency, Qsys automatically provides the pipelining logic necessary to support pipelined reads. Altera recommends using fixed latency pipelining as the default design starting point for slave interfaces. If your slave interface has a variable latency response time, use the readdatavalid signal to indicate when valid data is available. The interconnect implements read response FIFO buffering to handle the maximum number of pending read requests.

To use components that support pipelined read transfers, and to use a pipelined system interconnect efficiently, your system must contain pipelined masters. Refer to the “Avalon Pipelined Read Master Example” on page 10–39 for an example of a pipelined read master. Altera recommends using pipelined masters as the default starting point for new master components. Use the readdatavalid signal for these master interfaces.

Because master and slaves often have mismatched pipeline latency, interconnect often contains logic to reconcile the differences. Many cases of pipeline latency are possible, as shown in Table 10–1.

Table 10–1. Various Cases of Pipeline Latency in a Master-Slave Pair (Part 1 of 2)

<table>
<thead>
<tr>
<th>Master</th>
<th>Slave</th>
<th>Pipeline Management Logic Structure</th>
</tr>
</thead>
<tbody>
<tr>
<td>No pipeline</td>
<td>No pipeline</td>
<td>The Qsys interconnect does not instantiate logic to handle pipeline latency.</td>
</tr>
<tr>
<td>No pipeline</td>
<td>Pipelined with fixed or variable latency</td>
<td>The Qsys interconnect forces the master to wait through any slave-side latency cycles. This master-slave pair gains no benefits from pipelining, because the master waits for each transfer to complete before beginning a new transfer. However, while the master is waiting, the slave can accept transfers from a different master.</td>
</tr>
<tr>
<td>Pipelined</td>
<td>No pipeline</td>
<td>The Qsys interconnect carries out the transfer as if neither master nor slave were pipelined, causing the master to wait until the slave returns data. An example of a non-pipeline slave is an asynchronous off-chip interface.</td>
</tr>
</tbody>
</table>
Burst Avalon-MM Interfaces

Burst transfers are commonly used for latent memories such as SDRAM and off-chip communication interfaces such as PCI Express. To use a burst-capable slave interface efficiently, you must connect to a bursting master. Components that require bursting to operate efficiently typically have an overhead penalty associated with short bursts or non-bursting transfers.

Altera recommends that you design a burst-capable slave interface if you know that your component requires sequential transfers to operate efficiently. Because SDRAM memories incur a penalty when switching banks or rows, performance improves when SDRAM memories are accessed sequentially with bursts.

Architectures that use the same signals to transfer address and data also benefit from bursting. Whenever an address is transferred over shared address and data signals, the throughput of the data transfer is reduced. Because the address phase adds overhead, using large bursts increases the throughput of the connection.

Avalon-MM Burst Master Example

Figure 10–20 shows the architecture of a bursting write master that receives data from a FIFO and writes the contents to memory. You can use this master as a starting point for your own bursting components, such as custom DMAs, hardware accelerators, or off-chip communication interfaces. In Figure 10–20, the master performs word accesses and writes to sequential memory locations.

### Table 10–1. Various Cases of Pipeline Latency in a Master-Slave Pair (Part 2 of 2)

<table>
<thead>
<tr>
<th>Master</th>
<th>Slave</th>
<th>Pipeline Management Logic Structure</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pipelined</td>
<td>Pipelined with fixed latency</td>
<td>The Qsys interconnect allows the master to capture data at the exact clock cycle when data from the slave is valid, to enable maximum throughput. An example of a fixed latency slave is an on-chip memory.</td>
</tr>
<tr>
<td>Pipelined</td>
<td>Pipelined with variable latency</td>
<td>The slave asserts a signal when its readdata is valid, and the master captures the data. The master-slave pair can achieve maximum throughput if the slave has variable latency. Examples of variable latency slaves include SDRAM and FIFO memories.</td>
</tr>
</tbody>
</table>
For more information about the example in Figure 10–20, refer to the write master design in the *Avalon Memory-Mapped Master Templates* on the Altera website.

**Figure 10–20. Avalon Bursting Write Master**

When *go* is asserted, the *start_address* and *transfer_length* are registered. On the next clock cycle, the control logic asserts *burst_begin*. The *burst_begin* signal synchronizes the internal control signals in addition to the *master_address* and *master_burstcount* presented to the interconnect. The timing of these two signals is important because during bursting write transfers *address*, *bytenable*, and *burstcount* must be held constant for the entire burst.

To avoid inefficient writes, the master only posts a burst when enough data has been buffered in the FIFO. To maximize the burst efficiency, the master should stall only when a slave asserts *waitrequest*. In this example, the FIFO’s *used* signal tracks the number of words of data that are stored in the FIFO and determines when enough data has been buffered.

The *address* register increments after every word transfer, and the *length* register decrements after every word transfer. The address remains constant throughout the burst. Because a transfer is not guaranteed to complete on burst boundaries, additional logic is necessary to recognize the completion of short bursts and complete the transfer.
Reducing Logic Utilization

This section describes how to minimize logic size of Qsys systems. Typically, there is a trade-off between logic utilization and performance. Information in this section applies to both Avalon and AXI interfaces.

Minimize Interconnect Logic

In Qsys, changes to the connections between master and slave reduce the amount of interconnect logic required in the system.

Create Dedicated Master and Slave Connections

You might be able to create a system so that a master interface connects to a single slave interface. This configuration eliminates address decoding, arbitration, and return data multiplexing, which simplifies the interconnect. Dedicated master-to-slave connections attain the same clock frequencies as Avalon-ST connections.

Typically, these one-to-one connections include an Avalon memory-mapped bridge or hardware accelerator. For example, if you insert a pipeline bridge between a slave and all other master interfaces, the logic between the bridge master and slave interface is reduced to wires. Figure 10–16 on page 10–19 shows this technique. If a hardware accelerator connects only to a dedicated memory, no system interconnect logic is generated between the master and slave pair.

Removing Unnecessary Connections

The number of connections between master and slave interfaces affects the $f_{\text{MAX}}$ of your system. Every master interface that you connect to a slave interface increases the width of the multiplexer width. As a multiplexer width increases, so does the logic depth and width that implements the multiplexer in the FPGA. To improve your system performance, connect masters and slaves only when necessary.

When you connect a master interface to many slave interfaces, the multiplexer for the read data signal grows. Avalon typically uses a $\text{readdata}$ signal, and AXI read data signals add a response status and last indicator to the read response channel using the commands $\text{rdata}$, $\text{rresp}$, and $\text{rlast}$. Use bridges to help control the depth of multiplexers, as shown in Figure 10–9.

Simplifying Address Decode Logic

If address code logic is in the critical path, you may be able to change the address map to simplify the decode logic. Experiment with different address maps, including a one-hot encoding, to see if results improve.

Minimize Arbitration Logic by Consolidating Multiple Interfaces Into One

As the number of components in your design increases, the amount of logic required to implement the interconnect also increases. The number of arbitration blocks increases for every slave interface that is shared by multiple master interfaces. The width of the read data multiplexer increases as the number of slave interfaces supporting read transfers increases on a per master interface basis. For these reasons, consider implementing multiple blocks of logic as a single interface to reduce interconnect logic utilization.
Logic Consolidation Trade-Offs

You should consider the following trade-offs before making modifications to your system or interfaces.

Refer to “Using Concurrency in Memory-Mapped Systems” on page 10–5 for additional discussion on concurrency trade-offs.

First, consider the impact on concurrency that results when you consolidate components. When your system has four master components and four slave interfaces, it can initiate four concurrent accesses. If you consolidate the four slave interfaces into a single interface, then the four masters must compete for access. Consequently, you should only combine low priority interfaces such as low speed parallel I/O devices if the combination does not impact the performance.

Second, determine whether consolidation introduces new decode and multiplexing logic for the slave interface that the interconnect previously included. If an interface contains multiple read and write address locations, the interface already contains the necessary decode and multiplexing logic. When you consolidate interfaces, you typically reuse the decoder and multiplexer blocks already present in one of the original interfaces; however, combining interfaces might simply move the decode and multiplexer logic, rather than eliminate duplication.

Finally, consider whether consolidating interfaces makes the design complicated. If so, Altera recommends that you do not consolidate interfaces.

System Example of Consolidating Interfaces

In this example, the Nios II/e core maintains communication between the Nios II/f core and external processors. The Nios II/f core supports a maximum burst size of eight. The external processor interface supports a maximum burst length of 64. The Nios II/e core does not support bursting. The only memory in the system is SDRAM with an Avalon maximum burst length of two.
Figure 10–21 shows a system with a mix of components with different burst capabilities. It includes a Nios II/e core, a Nios II/f core, and an external processor, which off-loads some processing tasks to the Nios II/f core.

Qsys automatically inserts burst adapters to compensate for burst length mismatches. The adapters reduce bursts to a single transfer, or the length of two transfers. For the external processor interface connecting to DDR SDRAM, a burst of 64 words is divided into 32 burst transfers, each with a burst length of two.

When you generate a system, Qsys inserts burst adapters based on maximum burstcount values; consequently, the interconnect logic includes burst adapters between masters and slave pairs that do not require bursting, if the master is capable of bursts. In Figure 10–21, Qsys inserts a burst adapter between the Nios II processors and the timer, system ID, and PIO peripherals. These components do not support bursting and the Nios II processor performs only single word read and write accesses to these components.
To reduce the number of adapters, you can add pipeline bridges, as Figure 10–22 shows. The pipeline bridge between the Nios II/f core and the peripherals that do not support bursts eliminates three burst adapters from Figure 10–21. A second pipeline bridge between the Nios II/f core and the DDR SDRAM, with its maximum burst size set to eight, eliminates another burst adapter.

**Figure 10–22. Mixed Bursting System with Bridges**

**Implementing Multiple Clock Domains**

You specify clock domains in Qsys on the System Contents tab. Clock sources can be driven by external input signals to Qsys, or by PLLs inside Qsys. Clock domains are differentiated based on the name of the clock. You may create multiple asynchronous clocks with the same frequency.
Clock Domain Crossing Logic

Qsys generates Clock Domain Crossing Logic (CDC) that hides the details of interfacing components operating in different clock domains. The system interconnect supports the memory-mapped protocol with each port independently, and therefore masters do not need to incorporate clock adapters in order to interface to slaves on a different domain. Qsys interconnect logic propagates transfers across clock domain boundaries automatically.

Clock-domain adapters provide the following benefits:

- Allow component interfaces to operate at different clock frequencies.
- Eliminates the need to design CDC hardware.
- Allows each memory-mapped port to operate in only one clock domain, which reduces design complexity of components.
- Enable masters to access any slave without communication with the slave clock domain.
- Allows you to focus performance optimization efforts on components that require fast clock speed.

A clock domain adapter consists of two finite state machines (FSM), one in each clock domain, that use a simple hand-shaking protocol to propagate transfer control signals (read_request, write_request, and the master waitrequest signals) across the clock boundary.

Figure 10–23 shows illustrates a clock domain adapter between one master and one slave.

Figure 10–23. Block Diagram of Clock Crossing Adapter

The synchronizer blocks in Figure 10–23 use multiple stages of flipflops to eliminate the propagation of metastable events on the control signals that enter the handshake FSMs. The CDC logic works with any clock ratio.
The typical sequence of events for a transfer across the CDC logic is described as follows:

1. Master asserts address, data, and control signals.
2. The master handshake FSM captures the control signals, and immediately forces the master to wait.
   - The FSM uses only the control signals, not address and data. For example, the master simply holds the address signal constant until the slave side has safely captured it.
3. Master handshake FSM initiates a transfer request to the slave handshake FSM.
4. The transfer request is synchronized to the slave clock domain.
5. The slave handshake FSM processes the request, performing the requested transfer with the slave.
6. When the slave transfer completes, the slave handshake FSM sends an acknowledge back to the master handshake FSM.
7. The acknowledge is synchronized back to the master clock domain.
8. The master handshake FSM completes the transaction by releasing the master from the wait condition.

Transfers proceed as normal on the slave and the master side, without a special protocol to handle crossing clock domains. From the perspective of a slave, there is nothing different about a transfer initiated by a master in a different clock domain. From the perspective of a master, a transfer across clock domains simply requires extra clock cycles. Similar to other transfer delay cases (for example, arbitration delay or wait states on the slave side), the Qsys forces the master to wait until the transfer terminates. As a result, pipeline master ports do not benefit from pipelining when performing transfers to a different clock domain.

Qsys automatically determines where to insert CDC logic based on the system contents and the connections between components, and places CDC logic to maintain the highest transfer rate for all components. Qsys evaluates the need for CDC logic for each master and slave pair independently, and generates CDC logic wherever necessary.

**Duration of Transfers Crossing Clock Domains**

CDC logic extends the duration of master transfers across clock domain boundaries. In the worst case which is for reads, each transfer is extended by five master clock cycles and five slave clock cycles. Assuming the default value of 2 for the Master domain synchronizer length and the Slave domain synchronizer length, the components of this delay are the following:

- Four additional master clock cycles, due to the master-side clock synchronizer
- Four additional slave clock cycles, due to the slave-side clock synchronizer
- One additional clock in each direction, due to potential metastable events as the control signals cross clock domains
Systems that require a higher performance clock should use the Avalon-MM clock crossing bridge instead of the automatically inserted CDC logic. The clock crossing bridge includes a buffering mechanism, so that multiple reads and writes can be pipelined. After paying the initial penalty for the first read or write, there is no additional latency penalty for pending reads and writes, increasing throughput by up to four times, at the expense of added logic resources.

For more information, refer to *Avalon Memory-Mapped Design Optimizations* in the *Embedded Design Handbook*.

### Reducing Power Consumption

This section describes various low power design changes that you can make to reduce the power consumption of the interconnect and your custom components.

Qsys does not support AXI standard low power extensions in the current version of the QII software.

### Use Multiple Clock Domains

When you use multiple clock domains, you should put non-critical logic in the slower clock domain. Qsys automatically reconciles data crossing over asynchronous clock domains by inserting clock crossing logic (handshake or FIFO).

You can use clock crossing in Qsys to reduce the clock frequency of the logic that does not require a high frequency clock, allowing you to reduce power consumption. You can use either handshaking clock crossing bridges or handshaking clock crossing adapters to separate clock domains.

### Clock Crossing Bridge

You can use the clock crossing bridge to connect master interfaces operating at a higher frequency to slave interfaces running a a lower frequency. Only connect low throughput or low priority components to a clock crossing bridge that operates at a reduced clock frequency. The following are examples of low throughput or low priority components:

- PIOs
- UARTs (JTAG or RS-232)
- System identification (SysID)
- Timers
- PLL (instantiated within Qsys)
- Serial peripheral interface (SPI)
- EPCS controller
- Tristate bridge and the components connected to the bridge

By reducing the clock frequency of the components connected to the bridge, you reduce the dynamic power consumption of your design. Dynamic power is a function of toggle rates and decreasing the clock frequency decreases the toggle rate.
Figure 10–24 shows a system where a bridge reduces power consumption.

Qsys automatically inserts clock crossing adapters between master and slave interfaces that operate at different clock frequencies. You can choose the type of clock crossing adapter in the Qsys Project Settings tab. There are three types of clock crossing adapter types available in Qsys, as described below. Adapters do not appear in the Qsys Connection column because you do not insert them.

**Clock Crossing Adapter Types**

Specifies the default implementation for automatically inserted clock crossing adapters. The following adapter types are available:

- **Handshake**—Uses a simple handshaking protocol to propagate transfer control signals and responses across the clock boundary. This adapter uses fewer hardware resources because each transfer is safely propagated to the target domain before the next transfer can begin. The Handshake adapter is appropriate for systems with low throughput requirements.
FIFO—Uses dual-clock FIFOs for synchronization. The latency of the FIFO adapter is approximately two clock cycles more than the handshake clock crossing component, but the FIFO-based adapter can sustain higher throughput because it can support multiple transactions simultaneously. The FIFO adapter requires more resources. The FIFO adapter is appropriate for memory-mapped transfers requiring high throughput across clock domains.

Auto—Qsys specifies the appropriate FIFO adapter for bursting links and the Handshake adapter for all other links.

Throughput

Because the clock crossing bridge uses FIFOs to implement the clock crossing logic, it buffers transfers and data. Clock crossing adapters are not pipelined, so that each transaction is blocking until the transaction completes. Blocking transactions may lower the throughput substantially; consequently, if you want to reduce power consumption without limiting the throughput significantly, you should use the clock crossing bridge or the FIFO clock crossing adapter. However, if the design simply requires single read transfers, a clock crossing adapter is preferable because the latency is lower.

Resource Utilization

The clock crossing bridge requires few logic resources besides on-chip memory. The number of on-chip memory blocks used is proportional to the address span, data width, buffering depth, and bursting capabilities of the bridge. The clock crossing adapter does not use on-chip memory and requires a moderate number of logic resources. The address span, data width, and the bursting capabilities of the clock crossing adapter determine the resource utilization of the device.

Throughput versus Memory Trade-Offs

When you decide to use a clock crossing bridge or clock crossing adapter, you must consider the effects of throughput and memory utilization in your design. If on-chip memory resources are limited, you may be forced to choose the clock crossing adapter. Using the clock crossing bridge to reduce the power of a single component may not justify using more resources. However, if you can place all your low priority components behind a single clock crossing bridge, you reduce power consumption in your design.

Minimizing Toggle Rates

Your design consumes power whenever logic transitions between on and off states. When the state is held constant between clock edges, no charging or discharging occurs. This section discusses the following three design techniques that you can use to reduce the toggle rates of your system:

- Registering component boundaries
- Using clock enable signals
- Inserting bridges
Registering Component Boundaries

Qsys interconnect is uniquely combinational when no adapters or bridges are present and there is no interconnect pipelining. When a slave interface is not selected by a master, various signals may toggle and propagate into the component. By registering the boundary of your component at the master or slave interface, you can minimize the toggling of the interconnect and your component. In addition, registering boundaries can improve operating frequency. When you register the signals at the interface level, you must ensure that the component continues to operate within the interface standard specification.

Avalon-MM waitrequest is a difficult signal to synchronize when you add registers to your component. waitrequest must be asserted during the same clock cycle that a master asserts read or write to, in order to prolong the transfer. A master interface may read the waitrequest signal too early and post more reads and writes prematurely.

There is no direct AXI equivalent for waitrequest and burstcount, though the AMBA Protocol Specification implies that ready (the equivalent of Avalon-MM waitrequest) cannot depend combinatorially on AXI valid. Therefore, Qsys typically buffers AXI component boundaries (at least for the ready signal).

For slave interfaces, the interconnect manages the begintransfer signal, which is asserted during the first clock cycle of any read or write transfer. If your waitrequest is one clock cycle late, you can logically OR your waitrequest and the begintransfer signals to form a new waitrequest signal that is properly synchronized, as shown in Figure 10–25.

Figure 10–25. Variable Latency

Alternatively, your component can assert waitrequest before it is selected, guaranteeing that the waitrequest is already asserted during the first clock cycle of a transfer.
Using Clock Enables

You can use clock enables to hold your logic in a steady state. You can use the write and read signals as clock enables for slave components. Even if you add registers to your component boundaries, your interface can potentially toggle without the use of clock enables.

You can also use the clock enable to disable combinational portions of your component. For example, you can use an active high clock enable to mask the inputs into your combinational logic to prevent it from toggling when the component is inactive. Before preventing inactive logic from toggling, you must determine if the masking causes your circuit to function differently. If masking causes a functional failure, it might be possible to use a register stage to hold the combinational logic constant between clock cycles.

Inserting Bridges

You can use bridges to reduce toggle rates, if you do not want to modify the component by using boundary registers or clock enables. A bridge acts as a repeater where transfers to the slave interface are repeated on the master interface. If the bridge is not accessed, the components connected to its master interface are also not accessed. The master interface of the bridge remains idle until a master accesses the bridge slave interface.

Bridges can also reduce the toggle rates of signals that are inputs to other master interfaces. These signals are typically readdata, readdatavalid, and waitrequest. Slave interfaces that support read accesses drive the readdata, readdatavalid, and waitrequest signals. A bridge inserts either a register or clock crossing FIFO between the slave interface and the master to reduce the toggle rate of the master input signals.

Disabling Logic

There are typically two types of low power modes: volatile and non-volatile. A volatile low power mode holds the component in a reset state. When the logic is reactivated, the previous operational state is lost. A non-volatile low power mode restores the previous operational state. This section discusses using either software-controlled or hardware-controlled sleep modes to disable a component in order to reduce power consumption.

Software-Controlled Sleep Mode

To design a component that supports software controlled sleep mode, create a single memory mapped location that enables and disables logic, by writing a zero or one. Use the register’s output as a clock enable or reset, depending on whether the component has non-volatile requirements. The slave interface must remain active during sleep mode so that the enable bit can be set when the component needs to be activated.

If multiple masters can access a component that supports sleep mode, you can use the mutex core available in Qsys to provide mutually exclusive accesses to your component. You can also build in the logic to re-enable the component on the very first access by any master in your system. If the component requires multiple clock cycles to re-activate, then it must assert wait request to prolong the transfer as it exits sleep mode.
For more information about the mutex core, refer to the *Mutex Core* chapter of the *Embedded Peripherals IP User Guide*.

**Hardware-Controlled Sleep Mode**

You can implement a timer in your component that automatically causes the component to enter a sleep mode based on a timeout value specified in clock cycles between read or write accesses. Each access resets the timer to the timeout value. Each cycle with no accesses decrements the timeout value by one. If the counter reaches zero, the hardware enters sleep mode until the next access. Figure 10–26 provides a schematic for this logic. If restoring the component to an active state takes a long time, use a long timeout value so that the component is not continuously entering and exiting sleep mode.

The slave interface must remain functional while the rest of the component is in sleep mode. When the component exits sleep mode, the component must assert the *waitrequest* signal until it is ready for read or write accesses.

For more information on reducing power utilization, refer to *Power Optimization* in the *Quartus II Handbook*.

**Design Examples**

The following examples illustrate the resolution of Qsys system design challenges.

**Avalon Pipelined Read Master Example**

For a high throughput system using the Avalon-MM standard, you can design a pipelined read master that allows your system to issue multiple read requests before data returns. Pipelined read masters hide the latency of read operations by posting reads as frequently as every clock cycle. You can use this type of master when the address logic is not dependent on the data returning.

**Design Requirements**

You must carefully design the logic for the control and data paths of pipelined read masters. The control logic must extend a read cycle whenever the *waitrequest* signal is asserted. This logic must also control the master address, byteenable, and read signals. To achieve maximum throughput, pipelined read masters should post reads continuously as long as *waitrequest* is deasserted. While *read* is asserted, the address presented to the interconnect is stored.
The data path logic includes the readdata and readdatavalid signals. If your master can accept data on every clock cycle, you can register the data with the readdatavalid as an enable bit. If your master cannot process a continuous stream of read data, it must buffer the data in a FIFO. The control logic must stop issuing reads when the FIFO reaches a predetermined fill level to prevent FIFO overflow.

Refer to the *Avalon Interface Specifications* to learn more about the signals that implement an Avalon pipelined read master.

**Expected Throughput Improvement**

The throughput improvement that you can achieve with a pipelined read master is typically directly proportional to the pipeline depth of the interconnect and the slave interface. For example, if the total latency is two cycles, you can double the throughput by inserting a pipelined read master, assuming the slave interface also supports pipeline transfers. If either the master or slave does not support pipelined read transfers, then the interconnect asserts waitrequest until the transfer completes. You can also gain throughput when there are some cycles of overhead before a read response.

The “Increased Latency” on page 10–15 describes an example in which both the master and slave interfaces support pipelined read transfers. In this example, data can flow on a continuous stream after the initial latency. Where reads are not pipelined, the throughput is reduced. When both the master and slave interfaces support pipelined read transfers, data flows in a continuous stream after the initial latency. Figure 10–27 illustrates reads that are not pipelined. The system uses three cycles of latency for each read, achieving an overall throughput of 25%. Figure 10–20 shows reads that are pipelined. After the three cycles of latency, the data flows continuously.

You can use a pipelined read master that stores data in a FIFO to implement a custom DMA, hardware accelerator, or off-chip communication interface. Figure 10–27 shows a pipeline read master that stores data in a FIFO. The master performs word accesses that are word-aligned and reads from sequential memory addresses. The transfer length is a multiple of the word size. In Figure 10–27, the master performs word accesses that are word-aligned and reads from sequential memory addresses. The transfer length is a multiple of the word size.
Figure 10–27 shows a pipeline read master that stores data in a FIFO.

When the `go` bit is asserted, the master registers the start_address and transfer_length signals. The master begins issuing reads continuously on the next clock until the length register reaches zero. In this example, the word size is four bytes so that the address always increments by four and the length decrements by four. The read signal remains asserted unless the FIFO fills to a predetermined level. The address register increments and the length register decrements if the length has not reached 0 and a read is posted.

The master posts a read transfer every time the read signal is asserted and the waitrequest is deasserted. The master issues reads until the entire buffer has been read or waitrequest is asserted. An optional tracking block monitors the done bit. When the length register reaches zero, some reads are outstanding. The tracking logic prevents assertion of done until last read completes. The tracking logic monitors the number of reads posted to the interconnect so that it does not exceed the space remaining in the readdata FIFO. This logic includes a counter that verifies the following conditions are met:

- If a read is posted and readdatavalid is deasserted, the counter increments.
- If a read is not posted and readdatavalid is asserted, the counter decrements.

When the length register and the tracking logic counter reach zero, all the reads have completed and the done bit is asserted. The done bit is important if a second master overwrites the memory locations that the pipelined read master accesses. This bit guarantees that the reads have completed before the original data is overwritten.
Multiplexer Examples

You can combine adapters with streaming components to create datapaths whose input and output streams have different properties. The following sections provide examples of datapaths in which the output stream is higher performance than the input stream. Figure 10–28 shows an output with double the throughput of each interface with a corresponding doubling of the clock frequency. Figure 10–29 doubles the data width. Figure 10–30 boosts the frequency of a stream by 10% by multiplexing input data from two sources.

Example to Double Clock Frequency

Figure 10–28 illustrates a datapath that uses the dual clock version of the on-chip FIFO memory and Avalon-ST channel multiplexer to merge the 100 MHz input from two streaming data sources into a single 200 MHz streaming output. This example increases throughput by increasing the frequency and combining inputs.

Example to Double Data Width and Maintain Frequency

Figure 10–29 illustrates a datapath that uses the data format adapter and Avalon-ST channel multiplexer to convert two, 8-bit inputs running at 100 MHz to a single 16-bit output at 100 MHz.
Example to Boost the Frequency

Figure 10–30 illustrates a datapath that uses the dual clock version of the on-chip FIFO memory to boost the frequency of input data from 100 MHz to 110 MHz by sampling two input streams at differential rates. In this example, the on-chip FIFO memory has an input clock frequency of 100 MHz and an output clock frequency of 110 MHz. The channel multiplexer runs at 110 MHz and samples one input stream 27.3 percent of the time and the second 72.7 percent of the time.

You do not need to know what the typical and maximum input channel utilizations are before this type of design. For example, if the first channel hits 50% utilization, the output stream exceeds 100% utilization.

Conclusion

Recommendations presented in this chapter may improve your system’s maximum clock frequency, concurrency and throughput, logic utilization, or even power utilization. When you design a Qsys system, use your knowledge of the design intent and goals to further optimize system performance beyond the automated optimization available within Qsys.

Document Revision History

Table 10–2 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>■ Added AMBA APB support.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Added AMBA AXI4 support.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Added AMBA AXI3 support.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>■ New document release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter contains descriptions of Tcl commands and indicates the Qsys phases during which it is available: in the main static body of the program (main), or during the elaboration, composition, and fileset callback phases, or any combination. Table 11–1 summarizes the commands and provides a reference to the full description.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

For more information about procedures for creating component _hw.tcl files in the Qsys Component Editor, and supported interface standards, refer to the Creating Qsys Components and the Qsys Interconnect chapters in volume 1 of the Quartus II Handbook.

If you are developing a component to work with the Nios II processor, refer to the Publishing Component Information to Embedded Software chapter, which describes how to publish hardware component information for embedded software tools, such as a C compiler and a Board Support Package (BSP) generator.

This section provides a reference for hardware Tcl commands, as follows:

- “Module Definition” on page 11–3
- “Parameters” on page 11–8
- “Display Items” on page 11–18
- “Interfaces and Ports” on page 11–22
- “Composition” on page 11–30
- “Fileset Generation” on page 11–37

Table 11–1. Command Summary

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_documentation_link &lt;title&gt; &lt;fileOrUrl&gt;</td>
<td>page 11–4</td>
</tr>
<tr>
<td>get_module_assignment &lt;moduleName&gt;</td>
<td>page 11–6</td>
</tr>
<tr>
<td>get_module_assignments</td>
<td>page 11–4</td>
</tr>
<tr>
<td>get_module_ports</td>
<td>page 11–5</td>
</tr>
<tr>
<td>get_module_properties</td>
<td>page 11–5</td>
</tr>
<tr>
<td>get_module_property &lt;propertyName&gt;</td>
<td>page 11–5</td>
</tr>
<tr>
<td>package &lt;require&gt; -exact qsys &lt;version&gt;</td>
<td>page 11–5</td>
</tr>
</tbody>
</table>
### Table 11–1. Command Summary *(1)* (Part 2 of 3)

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>send_message &lt;messageLevel&gt; &lt;messageText&gt;</td>
<td>page 11–6</td>
</tr>
<tr>
<td>set_module_assignment &lt;moduleName&gt; [value]</td>
<td>page 11–6</td>
</tr>
<tr>
<td>set_module_property &lt;propertyName&gt; &lt;propertyValue&gt;</td>
<td>page 11–7</td>
</tr>
</tbody>
</table>

#### Parameters

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_parameter &lt;parameterName&gt; &lt;parameterType&gt; [defaultValue description]</td>
<td>page 11–8</td>
</tr>
<tr>
<td>decode_address_map &lt;address_map_XML_string&gt;</td>
<td>page 11–9</td>
</tr>
<tr>
<td>get_parameters</td>
<td>page 11–10</td>
</tr>
<tr>
<td>get_parameter_properties</td>
<td></td>
</tr>
<tr>
<td>get_parameter_property &lt;parameterName&gt; &lt;propertyName&gt;</td>
<td>page 11–11</td>
</tr>
<tr>
<td>get_parameter_value &lt;parameterName&gt;</td>
<td>page 11–11</td>
</tr>
<tr>
<td>get_string &lt;identifier&gt;</td>
<td>page 11–11</td>
</tr>
<tr>
<td>load_strings &lt;fileName&gt;</td>
<td>page 11–12</td>
</tr>
<tr>
<td>set_parameter_property &lt;parameterName&gt; &lt;propertyName&gt; &lt;value&gt;</td>
<td>page 11–11</td>
</tr>
<tr>
<td>set_parameter_value &lt;parameterName&gt; &lt;value&gt;</td>
<td>page 11–11</td>
</tr>
</tbody>
</table>

#### Display Items

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_display_item &lt;groupName&gt; &lt;id&gt; &lt;type&gt; [additionalInfo]</td>
<td>page 11–18</td>
</tr>
<tr>
<td>get_display_items</td>
<td>page 11–20</td>
</tr>
<tr>
<td>get_display_item_properties</td>
<td>page 11–20</td>
</tr>
<tr>
<td>get_display_item_property &lt;itemName&gt; &lt;propertyName&gt;</td>
<td>page 11–20</td>
</tr>
<tr>
<td>set_display_item_property &lt;itemName&gt; &lt;propertyName&gt; &lt;value&gt;</td>
<td>page 11–21</td>
</tr>
</tbody>
</table>

#### Interfaces and Ports

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_interface &lt;interfaceName&gt; &lt;interfaceType&gt; &lt;direction&gt; [associatedClock]</td>
<td>page 11–22</td>
</tr>
<tr>
<td>add_interface_port &lt;interfaceName&gt; &lt;portName&gt; &lt;portRole&gt; [direction width_expr]</td>
<td>page 11–23</td>
</tr>
<tr>
<td>get_interfaces &lt;interfaceName&gt;</td>
<td>page 11–23</td>
</tr>
<tr>
<td>get_interface_assignment &lt;interfaceName&gt; &lt;name&gt;</td>
<td>page 11–24</td>
</tr>
<tr>
<td>get_interface_assignments</td>
<td>page 11–24</td>
</tr>
<tr>
<td>get_interface_ports [&lt;interfaceName&gt;]</td>
<td>page 11–24</td>
</tr>
<tr>
<td>get_interface_property &lt;interfaceName&gt; &lt;propertyName&gt;</td>
<td>page 11–25</td>
</tr>
<tr>
<td>get_port_properties</td>
<td>page 11–26</td>
</tr>
<tr>
<td>get_port_property &lt;portName&gt; &lt;propertyName&gt;</td>
<td>page 11–26</td>
</tr>
<tr>
<td>set_interface_assignment &lt;interfaceName&gt; &lt;name&gt; [value]</td>
<td>page 11–24</td>
</tr>
<tr>
<td>set_interface_property &lt;interfaceName&gt; &lt;propertyName&gt; &lt;value&gt;</td>
<td>page 11–27</td>
</tr>
<tr>
<td>set_port_property &lt;portName&gt; &lt;propertyName&gt; [value]</td>
<td>page 11–28</td>
</tr>
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</table>

#### Composition

<table>
<thead>
<tr>
<th>Command</th>
<th>Full Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>add_connection &lt;startInterface&gt; &lt;endInterface&gt; [type]</td>
<td>page 11–30</td>
</tr>
<tr>
<td>get_connections</td>
<td>page 11–31</td>
</tr>
<tr>
<td>get_connection_parameters &lt;instanceName&gt;</td>
<td>page 11–31</td>
</tr>
</tbody>
</table>
Module Definition

This section provides information about the commands that you use to define and query a module.
add_documentation_link

This command allows you to link to documentation for your component.

<table>
<thead>
<tr>
<th>add_documentation_link</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
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<tr>
<td></td>
</tr>
</tbody>
</table>

get_module_assignment

This command returns the value of an assignment. You can use the get_module_assignment and set_module_assignment and the get_interface_assignment and set_interface_assignment commands to provide information about hardware components to embedded software tools and applications.

<table>
<thead>
<tr>
<th>get_module_assignment</th>
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<tbody>
<tr>
<td>Callback availability</td>
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<tr>
<td>Usage</td>
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<tr>
<td>Returns</td>
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<tr>
<td>Arguments</td>
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<tr>
<td>Example</td>
</tr>
</tbody>
</table>

get_module_assignments

This command returns names of the module assignments.

<table>
<thead>
<tr>
<th>get_module_assignments</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
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<tr>
<td>Example</td>
</tr>
</tbody>
</table>
**get_module_ports**

This command returns a list of the names of all the ports that are currently defined.

<table>
<thead>
<tr>
<th>get_module_ports</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

**get_module_properties**

This command returns the names of all the available module properties as a list of strings. You can use the `get_module_property` and `set_module_property` commands to get and set values of individual properties. The value returned by this command is always the same for a particular version of Qsys.

<table>
<thead>
<tr>
<th>get_module_properties</th>
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</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

**get_module_property**

This command returns the value of a single module property.

<table>
<thead>
<tr>
<th>get_module_property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

**package**

The `package` command allows you to specify a particular version of the Qsys software to avoid software compatibility issues. You must use the `package` command at the beginning of your `_hw.tcl` file. When used, the component files behave as if they are interpreted by the version of the Qsys software that you specify.
This document describes the behavior of components which start with
package require -exact qsys 12.1 For earlier versions of the _hw.tcl commands,
refer to the documentation for that release. If you do not request a package, you get
_hw.tcl commands compatible with Quartus II version 9.0.

**package**

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main (before any other commands in the file)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>package require -exact qsys &lt;version&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td>version The version of Qsys that you require, such as 12.1.</td>
</tr>
<tr>
<td>Example</td>
<td>package require -exact qsys 12.1</td>
</tr>
</tbody>
</table>

**send_message**

This command sends a message to the user of the component. The message text is
normally interpreted as HTML. The `<b>` element can be used to provide emphasis. If
you do not want the message text to be interpreted as HTML, then pass a list such as
`{info text}` as the message level.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>send_message &lt;messageLevel&gt; &lt;messageText&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
</tbody>
</table>

**Arguments**

<table>
<thead>
<tr>
<th>messageLevel</th>
<th>The following message levels are supported:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Error</td>
<td>Provides an error message. The Qsys system cannot be generated with existing error messages.</td>
</tr>
<tr>
<td>Warning</td>
<td>Provides a warning message.</td>
</tr>
<tr>
<td>Info</td>
<td>Provides an informational message.</td>
</tr>
<tr>
<td>Progress</td>
<td>Reports progress during generation.</td>
</tr>
<tr>
<td>Debug</td>
<td>Provides messages when debug mode is enabled.</td>
</tr>
<tr>
<td>messageText</td>
<td>The text of the message.</td>
</tr>
</tbody>
</table>

**set_module_assignment**

This command sets the value of the specified assignment.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_module_assignment &lt;name&gt; [&lt;value&gt;]</td>
</tr>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
</tbody>
</table>
**set_module_assignment**

This command allows you to set the values for module properties.

<table>
<thead>
<tr>
<th>Arguments</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>The assignment whose value is being set.</td>
</tr>
<tr>
<td>value</td>
<td>The value of the assignment.</td>
</tr>
</tbody>
</table>

**Example**

```
set_module_assignment embeddedsw.CMacro.colorSpace CMYK
```

**set_module_property**

This command allows you to set the values for module properties.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_module_property &lt;propertyName&gt; &lt;propertyValue&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td>propertyName One of the properties listed in Table 11–2 on page 11–7. propertyValue The new value of the property.</td>
</tr>
</tbody>
</table>

**Example**

```
set_module_property VERSION 10.0
```

Table 11–2 lists the available module properties, their use, and the phases in which they can be set.

### Module Properties Table

**Table 11–2. Module Properties (Part 1 of 2)**

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Property Type</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AUTHOR</td>
<td>String</td>
<td>Main program</td>
<td>The module’s author.</td>
</tr>
<tr>
<td>COMPOSITION_CALLBACK</td>
<td>String</td>
<td>String</td>
<td>The name of the composition callback. If you define a composition callback, you cannot define the generation or elaboration callbacks.</td>
</tr>
<tr>
<td>DESCRIPTION</td>
<td>String</td>
<td>Main program</td>
<td>The description of the module, such as &quot;Example Qsys Module.&quot;</td>
</tr>
<tr>
<td>DISPLAY_NAME</td>
<td>String</td>
<td>Main program</td>
<td>The name to display when referencing the module, such as &quot;My Component.&quot;</td>
</tr>
<tr>
<td>EDITABLE</td>
<td>Boolean</td>
<td>Main program</td>
<td>Indicates whether you can edit the component in the Component Editor.</td>
</tr>
<tr>
<td>ELABORATION_CALLBACK</td>
<td>String</td>
<td>Main program</td>
<td>The name of the elaboration callback. When set, the component’s elaboration callback is called to validate and elaborate interfaces for instances of the component.</td>
</tr>
<tr>
<td>GROUP</td>
<td>String</td>
<td>Main program</td>
<td>The group in the Component Library that includes this component.</td>
</tr>
<tr>
<td>ICON_PATH</td>
<td>String</td>
<td>Main program</td>
<td>A path to an icon to display in the module’s parameter editor.</td>
</tr>
</tbody>
</table>
Parameters

Parameters allow users of your component to affect its operation in the same manner as Verilog HDL parameters or VHDL generics.

**add_parameter**

This command adds a parameter to your component. Most of the parameter types are found in the C programming language or HDL. However, the `string_list` and `integer_list` parameters that are used to create tables in GUIs require some explanation.

- When you add a parameter of type `string_list` or `integer_list`, the parameter is displayed in a single-column table that can add or remove items in the list.
- If you add multiple parameters of type `string_list` or `integer_list` to the same group, and add the group with a `TABLE` hint, the entire group is displayed as a multi-column table. Example 11–1 illustrates the use of the `integer_list` parameter types to create a multi-column table.

**Example 11–1. Creating Tables Using the string_list and integer_list Parameter Types**

```
add_parameter names STRING_LIST
add_parameter counts INTEGER_LIST

add_display_item "" myTable GROUP TABLE
add_display_item myTable names PARAMETER
add_display_item myTable counts PARAMETER
```
add_parameter

Callback availability  Main

Usage  add_parameter <parameterName> <parameterType> [defaultValue description]

Returns  String

Arguments

| parameterName | A name that you, the component author, choose for your parameter. |
| parameterType | The following types are supported: integer, natural, positive, boolean, float, long, std_logic, std_logic_vector, string, string_list, and integer_list. |
| defaultValue  | The value for this parameter if the parameter’s value is never explicitly set. |
| description   | Explains the use of the parameter. |

Example  add_parameter seed integer 17 "The seed to use for data generation."

decode_address_map

This utility function converts an XML–formatted address map into a list of Tcl lists. Each inner list is in the correct format for conversion to an array. The XML code describing each slave includes: its name, start address, and end address + 1.

Figure 11–1 shows a portion of a Qsys system with three slave devices.

Figure 11–1. Qsys System with Three Avalon-MM Slaves

Example 11–2 shows the XML code that describes the address map for the master that accesses these slaves. The format of the XML string provided may differ from that described here; it may have different white space between the elements and could include additional attributes or elements. Using the decode_address_map command to decode the XML representing an master’s address map is easier and ensures that your code will work with future versions of the XML address map.

Altera recommends that you use the code provided in the description of Example 11–2 to enumerate over the components within an address map, rather than writing your own parser.

Example 11–2. Address Map for Master

```xml
<address-map>
  <slave name='ext_ssram' start='0x01000000' end='0x01200000' />
  <slave name='sys_clk_timer' start='0x02120800' end='0x02120820' />
  <slave name='sysid' start='0x021208B8' end='0x021208C0' />
</address-map>
```
**decode_address_map**

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Elaboration, generation, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>decode_address_map &lt;address_map_XML_string&gt;</code></td>
</tr>
<tr>
<td>Returns</td>
<td>List of Tcl lists, each one suitable for passing to array set</td>
</tr>
<tr>
<td>Arguments</td>
<td><code>address_map_XML_string</code> An XML string describing the address map of a master.</td>
</tr>
</tbody>
</table>
| Example               | `set address_map_xml [get_parameter_value my_map_param]`  
                         `set address_map_dec [decode_address_map $address_map_xml]`  
                         `foreach i $address_map_dec {`  
                         `    array set info $i`  
                         `    send_message info "Connected to slave $info(name)"`  
                         `}` |

**get_parameters**

This command returns the names of all parameters that have been previously defined by `add_parameter` as a space-separated list.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>get_parameters</code></td>
</tr>
<tr>
<td>Returns</td>
<td>List of strings</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td><code>set parameter_summary [get_parameters]</code></td>
</tr>
</tbody>
</table>

**get_parameter_properties**

This command returns a list of all the available parameter properties as a list of strings. The `get_parameter_property` and `set_parameter_property` commands are used to get and set the values of these properties, respectively.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>get_parameter_properties</code></td>
</tr>
<tr>
<td>Returns</td>
<td>List of strings</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td><code>set property_summary [get_parameter_properties]</code></td>
</tr>
</tbody>
</table>
**get_parameter_property**

This command returns a single parameter property.

<table>
<thead>
<tr>
<th><strong>get_parameter_property</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td>Main, elaboration, fileset, and composition</td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><code>get_parameter_property &lt;parameterName&gt; &lt;propertyName&gt;</code></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td>string, boolean, or units, depending on property. Refer to Table 11–3 on page 11–14.</td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><code>parameterName</code></td>
</tr>
<tr>
<td><code>propertyName</code></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
<tr>
<td><code>get_parameter_property parameter1 GROUP</code></td>
</tr>
</tbody>
</table>

**get_parameter_value**

This command returns the current value of a parameter defined previously with the `add_parameter` command.

<table>
<thead>
<tr>
<th><strong>get_parameter_value</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td>Elaboration (1), fileset, and composition</td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><code>get_parameter_value &lt;parameterName&gt;</code></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td>String</td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><code>parameterName</code></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
<tr>
<td><code>set fifo_width [get_parameter_value fifo_width]</code></td>
</tr>
</tbody>
</table>

**Note:**

(1) If `AFFECTS_ELABORATION=false` for a given parameter, `get_parameter_value` is not available for that parameter from the elaboration callback. If `AFFECTS_GENERATION=false` then it is not available from the generation callback.

**get_string**

This command returns the value of an externalized string previously loaded by the `load_strings` command.

**Example 11–3. get_string**

```tcl
package require -exact qsys <version>
load_strings test.properties
set_module_property NAME test
set_module_property VERSION [get_string VERSION]
set_module_property DISPLAY_NAME [get_string DISPLAY_NAME]
add_parameter firepower INTEGER 0 ""
set_parameter_property firepower DISPLAY_NAME [get_string PARAM_DISPLAY_NAME]
set_parameter_property firepower TYPE INTEGER
set_parameter_property firepower DESCRIPTION [get_string PARAM_DESCRIPTION]
DISPLAY_NAME = Trogdor!
VERSION = 1.0
PARAM_DISPLAY_NAME = Firepower
PARAM_DESCRIPTION = The amount of force to use when breathing fire.
```
### get_string

<table>
<thead>
<tr>
<th>Availability</th>
<th>You can use this command in any phase.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_string &lt;identifier&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>string</td>
</tr>
<tr>
<td>Arguments</td>
<td>identifier</td>
</tr>
<tr>
<td>Example</td>
<td>get_string MY_STRING</td>
</tr>
</tbody>
</table>

#### load_strings

This command loads strings from an external `.properties` file. The format of the properties file is in the Java Properties File format.

**Example 11–4. load_strings**

```tcl
package require -exact qsys 12.1
load_strings test.properties
set_module_property NAME test
set_module_property VERSION [get_string VERSION]
set_module_property DISPLAY_NAME [get_string DISPLAY_NAME]
add_parameter firepower INTEGER 0 ""
set_parameter_property firepower DISPLAY_NAME [get_string PARAM_DISPLAY_NAME]
set_parameter_property firepower TYPE INTEGER
set_parameter_property firepower DESCRIPTION [get_string PARAM_DESCRIPTION]
DISPLAY_NAME = Trogdor!
VERSION = 1.0
PARAM_DISPLAY_NAME = Firepower
PARAM_DESCRIPTION = The amount of force to use when breathing fire.
```

<table>
<thead>
<tr>
<th>availability</th>
<th>You can use this command in any phase.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>load_strings &lt;path&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>none</td>
</tr>
<tr>
<td>Arguments</td>
<td>path</td>
</tr>
<tr>
<td>Example</td>
<td>load_strings_my_externalized_strings.properties</td>
</tr>
</tbody>
</table>
### set_parameter_property

This command sets a single parameter property.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_parameter_property &lt;parameterName&gt; &lt;propertyName&gt; &lt;value&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>string, boolean, or units depending on property</td>
</tr>
<tr>
<td>Arguments</td>
<td><strong>parameterName</strong> Specifies the parameter that is being set.</td>
</tr>
<tr>
<td></td>
<td><strong>propertyName</strong> Specifies the property of parameterName that is being set, refer to Table 11–3 on page 11–14 for a list of properties.</td>
</tr>
<tr>
<td></td>
<td><strong>value</strong> The new value for the property.</td>
</tr>
<tr>
<td>Example</td>
<td>set_parameter_property BAUD_RATE ALLOWED_RANGES {9600 19200 38400}</td>
</tr>
</tbody>
</table>

### set_parameter_value

This command sets a parameter value. The values of derived parameters can be set from the elaboration callback.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Elaboration and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_parameter_value &lt;parameterName&gt; &lt;value&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>None</td>
</tr>
<tr>
<td>Arguments</td>
<td><strong>parameterName</strong> Specifies the parameter that is being set.</td>
</tr>
<tr>
<td></td>
<td><strong>value</strong> Specifies the value of parameterName.</td>
</tr>
<tr>
<td>Example</td>
<td>set_parameter_value BAUD_RATE 19200</td>
</tr>
</tbody>
</table>
### Parameter Properties Table

Table 11–3 describes the properties available to describe the behaviors of each of the parameters you can specify, their use, and when they can be set.

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Type/Default</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AFFECTS_ELABORATION</td>
<td>Boolean, true</td>
<td>Main program</td>
<td>Set AFFECTS_ELABORATION to false for parameters that do not affect the external interface of the module. An example of a parameter that does not affect the external interface is isNonVolatileStorage. An example of a parameter that does affect the external interface is width. When the value of a parameter changes, if that parameter has set AFFECTS_ELABORATION=false, the elaboration phase (calling the callback or hardware analysis) is not repeated, improving performance. Because the default value of AFFECTS_ELABORATION is true, the provided HDL file is normally re-analyzed to determine the new port widths and configuration every time a parameter changes.</td>
</tr>
<tr>
<td>AFFECTS_GENERATION</td>
<td>Boolean, refer to description</td>
<td>Main program</td>
<td>The default value of AFFECTS_GENERATION is false if you provide a top-level HDL module; it is true if you provide a fileset callback. Set AFFECTS_GENERATION to false if the value of a parameter does not change the results of fileset generation.</td>
</tr>
<tr>
<td>AFFECTS_VALIDATION</td>
<td>Boolean, refer to description</td>
<td>Main program</td>
<td>The AFFECTS_VALIDATION property marks whether a parameter’s value is used to set derived parameters, and whether the value affects validation messages. When set to false, this may improve response time in the parameter editor UI when the value is changed. The default value is true.</td>
</tr>
<tr>
<td>ALLOWED_RANGES</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>Indicates the range or ranges that the parameter value can have. For integers, The ALLOWED_RANGES property is a list of ranges that the parameter can take on, where each range is a single value, or a range of values defined by a start and end value separated by a colon, such as 11:15. This property can also specify legal values and display strings for integers, such as {0:None 1:Monophonic 2:Stereo 4:Quadrophonic} meaning 0,1,2,4 are the legal values. You can also assign longer strings to be displayed in the parameter editor to string variables. For example, ALLOWED_RANGES = {&quot;dev1:Cyclone IV GX&quot; &quot;dev2:Stratix V GT&quot;} Refer to Table 11–1 on page 11–1.</td>
</tr>
<tr>
<td>DEFAULT_VALUE</td>
<td>String or Boolean</td>
<td>Main program</td>
<td>The default value.</td>
</tr>
<tr>
<td>DERIVED</td>
<td>Boolean, false</td>
<td>Elaboration callback</td>
<td>When true, indicates that this parameter’s value is computed by the component during elaboration or composition callback. The default value is false.</td>
</tr>
</tbody>
</table>
### Table 11–3. Parameter Properties (Part 2 of 3)

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Type/Default</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DESCRIPTION</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>DESCRIPTION—A tooltip description of the parameter that appears in the parameter editor.</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>LONG_DESCRIPTION—A description of the parameter that appears in a popup dialog box when you click Documentation in the parameter editor.</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>If LONG_DESCRIPTION is not set, then DESCRIPTION appears in the documentation window.</td>
</tr>
<tr>
<td>DISPLAY_NAME</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>The GUI label that identifies the parameter.</td>
</tr>
<tr>
<td>DISPLAY_UNITS</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>The GUI label that describes the units that the parameter’s value represents, such as “Megabytes”.</td>
</tr>
<tr>
<td>ENABLED</td>
<td>Boolean, true</td>
<td>Main program</td>
<td>When false, the parameter is disabled, meaning that it is displayed, but greyed out, indicating that it is not editable on the parameter editor.</td>
</tr>
<tr>
<td>HDL_PARAMETER</td>
<td>Boolean, false</td>
<td>Main program</td>
<td>When true, the parameter must be passed to the HDL component description. The default value is false.</td>
</tr>
<tr>
<td>NEW_INSTANCE_VALUE</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>This property allows you to change the default value of a parameter without affecting older components that have assigned a default value to this parameter using the defaultValue argument. The practical result is that older instances of the component will continue to use defaultValue for the parameter and newer instances can use the value assigned by NEW_INSTANCE_VALUE.</td>
</tr>
<tr>
<td>SYSTEM_INFO</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>Allows you to assign information about the instantiating system to a parameter. A SYSTEM_INFO property requires an argument specifying the type of information requested, &lt;info-type&gt;. &lt;info-type&gt; may also take an argument. The syntax of the Tcl command is:</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>set_parameter_property my_parameter SYSTEM_INFO &lt;info-type&gt; [&lt;arg&gt;] Refer to Table 11–4 for descriptions of the &lt;info_type&gt; argument. You can use the SYSTEM_INFO property to set the SYSTEM_INFO_TYPE and SYSTEM_INFO_ARG properties at the same time.</td>
</tr>
<tr>
<td>SYSTEM_INFO_ARG</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>Defines an argument to be passed to a particular SYSTEM_INFO function.</td>
</tr>
<tr>
<td>SYSTEM_INFO_TYPE</td>
<td>Various</td>
<td>Main program</td>
<td>Specifies one of the types of information listed in Table 11–4 on page 11–17.</td>
</tr>
<tr>
<td>TYPE</td>
<td>String, &quot;&quot;</td>
<td>Main program</td>
<td>Specifies one of the following types: INTEGER, NATURAL, POSITIVE, BOOLEAN, STD_LOGIC, STD_LOGIC_VECTOR, STRING, STRING_LIST, INTEGER_LIST, LONG, or FLOAT.</td>
</tr>
</tbody>
</table>
Table 11–3. Parameter Properties (Part 3 of 3)

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Type/Default</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
</table>
| UNITS         | String, ""   | Main program | Sets the units of the parameter. The following values are possible:  
|               |              |            | - NONE  
|               |              |            | - ADDRESS  
|               |              |            | - BITS  
|               |              |            | - BITSPERSECOND  
|               |              |            | - BYTES  
|               |              |            | - CYCLES  
|               |              |            | - GIGABYTES  
|               |              |            | - GIGABITSPERSECOND  
|               |              |            | - GIGAHERTZ  
|               |              |            | - HERTZ  
|               |              |            | - KILOBYTES  
|               |              |            | - KILOHERTZ  
|               |              |            | - KILOBITSPERSECOND  
|               |              |            | - MEGABYTES  
|               |              |            | - MEGABITSPERSECOND  
|               |              |            | - MEGAHERTZ  
|               |              |            | - MICROSECONDS  
|               |              |            | - MILISECONDS  
|               |              |            | - NANOSECONDS  
|               |              |            | - PERCENT  
|               |              |            | - PICOSECONDS  
|               |              |            | - SECONDS  
|               |              |            | For example, set parameter property frequency UNITS GIGAHERTZ |
| VISIBLE       | Boolean, true | Main program elaboration, callbacks | Indicates whether or not to display the parameter in the parameterization GUI. |
| WIDTH         | String, ""   | Main program | For a STD_LOGIC_VECTOR parameter, this indicates the width of the logic vector. |
## SYSTEM_INFO Properties Table

Table 11–4 describes the properties that you can use with the `SYSTEM_INFO` parameter property.

<table>
<thead>
<tr>
<th>SYSTEM_INFO Type</th>
<th>Type of Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ADDRESS_MAP</td>
<td>String</td>
<td>Assigns an XML-formatted string describing the address map to the parameter you specify.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {ADDRESS_MAP &lt;my_avalon-mm_master&gt;}</code></td>
</tr>
<tr>
<td>ADDRESS_WIDTH</td>
<td>Integer</td>
<td>Assigns an integer to the parameter you specify that is the number of bits an master must drive to address all of its slaves, using byte addresses.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {ADDRESS_WIDTH &lt;my_avalon-mm_master&gt;}</code></td>
</tr>
<tr>
<td>CLOCK_DOMAIN</td>
<td>Integer</td>
<td>Assigns an integer representing the clock domain to the parameter you specify. You can use this command to determine whether multiple interfaces in your module are on the same clock domain. The absolute value of the integer value is arbitrary, but if two interfaces are on the same clock domain, the <code>CLOCK_DOMAIN</code> value is guaranteed to be the same and greater than zero.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {CLOCK_DOMAIN &lt;my_clk&gt;}</code></td>
</tr>
<tr>
<td>CLOCK_RATE</td>
<td>Integer or String</td>
<td>Assigns a positive number, which is the clock frequency in Hz to the clock input interface you specify. Assigns 0 if the clock rate is not known.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {CLOCK_RATE &lt;my_clk&gt;}</code></td>
</tr>
<tr>
<td>CLOCK_RESET_INFO</td>
<td>String</td>
<td>Specifies the name of the module’s clock or reset sink interface. (Specifies the clock sink interface for designs that use a global reset.)</td>
</tr>
<tr>
<td>CUSTOM_INSTRUCTION_SLAVES</td>
<td>String</td>
<td>Provides custom instruction slave information, including the name, base address, address span, and clock cycle type.</td>
</tr>
<tr>
<td>DEVICE_FAMILY</td>
<td>String</td>
<td>Assigns the family name (not the specific device part number) of the currently selected device to the parameter you specify.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {DEVICE_FAMILY}</code></td>
</tr>
<tr>
<td>DEVICE_FEATURES</td>
<td>String</td>
<td>Creates a list of key/value pairs delineated by spaces indicating whether a particular device feature is available in the currently selected device family. The format of the list is suitable for passing to the Tcl <code>array set</code> command. This list is assigned to the parameter you specify. The following features are supported: M512_MEMORY, M4K_MEMORY, M9K_MEMORY, M144K_MEMORY, MRAM_MEMORY, MLAB_MEMORY, ESB, DSP, and EMUL.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {DEVICE_FEATURES}</code></td>
</tr>
<tr>
<td>INTERRUPTS_USED</td>
<td>Integer or string</td>
<td>Creates a mask indicating which bits of the interrupt receiver vector are connected to an interrupt sender. This mask is assigned to the parameter you specify. You can use this interrupt mask to optimize logic that handles interrupts.</td>
</tr>
<tr>
<td></td>
<td></td>
<td><code>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {INTERRUPTS_USED &lt;my_interrupt_receiver&gt;}</code></td>
</tr>
</tbody>
</table>
You specify your component GUI using the display commands.

**add_display_item**

You can use this command to specify the following aspects of component display:

- You can create logical groups for a component’s parameters. For example, you might want to create separate groups for the component’s timing, size, and simulation parameters. A component displays the groups and parameters in the order that you specify the display items for them in the _hw.tcl_ file.

<table>
<thead>
<tr>
<th>SYSTEM_INFO Type</th>
<th>Type of Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MAX_SLAVE_DATA_WIDTH</td>
<td>Integer</td>
<td>Assigns an integer to the parameter you specify that is the data width of the widest slave connected to the specified master.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {MAX_SLAVE_DATA_WIDTH &lt;my_avalon_mm_master&gt;}</td>
</tr>
<tr>
<td>RESET_DOMAIN</td>
<td>Integer</td>
<td>Assigns an integer representing the reset domain to the parameter you specify. You can use this command to determine whether multiple interfaces in your module are on the same reset domain. The absolute value of the integer value is arbitrary, but if two interfaces are on the same reset domain, the RESET_DOMAIN value is guaranteed to be the same and greater than zero.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_parameter_property &lt;my_parameter&gt; SYSTEM_INFO {RESET_DOMAIN &lt;my_reset&gt;}</td>
</tr>
<tr>
<td>TRISTATECONDUIT_MASTERS</td>
<td>String</td>
<td>Specifies the name or names of the module’s interfaces that are tri-state conduit slaves.</td>
</tr>
<tr>
<td>TRISTATECONDUIT_INFO</td>
<td>String</td>
<td>Returns an XML string containing information about the Avalon-TC masters connected to the specified Avalon-TC slave interface on a given component. The returned string may include all of the following information:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ The Avalon-TC slave interface name</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ The Avalon-TC master module and interface names</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ The Avalon-TC signal names, directions, and widths</td>
</tr>
<tr>
<td></td>
<td></td>
<td>The argument to SYSTEM_INFO_ARG is a regular expression that specifies the interface or interfaces of interest. The following example returns an XML string named TC_slave_info for TC slave interface named CFI_FLASH.uas:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>add_parameter TC_slave_info string &quot;&quot;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_parameter_property TC_slave_info SYSTEM_INFO_TYP TRISTATECONDUIT_INFO</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_parameter_property TC_slave_info SYSTEM_INFO_ARG &quot;uas&quot;</td>
</tr>
<tr>
<td></td>
<td></td>
<td>To retrieve information about all the slave interfaces on a module substitute &quot;***&quot; for the interface name, as the following example illustrates:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_parameter_property TC_slave_info SYSTEM_INFO_ARG &quot;***&quot;</td>
</tr>
<tr>
<td>UNIQUE_ID</td>
<td>String</td>
<td>A string guaranteed to be unique to this module.</td>
</tr>
</tbody>
</table>
You can create multicolumn tables to present a component’s parameters. Refer to Example 11–1 on page 11–8 for an example that illustrates multicolumn tables.

You can specify an image to provide a pictorial representation of a parameter or parameter group.

You can create a button by adding a display item of type action. The display item includes the name of the callback to run when the action is performed.

You create a display group by adding display items to it.

---

### add_display_item

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>add_display_item &lt;groupName&gt; &lt;id&gt; &lt;type&gt; [additionalInfo]</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
</tbody>
</table>

- **groupName**: Specifies the group to which a display item belongs.
- **id**: Specifies the parameter or icon to be displayed in a group. Each display item associated with a component must have a different ID.
- **type**: Specifies the category of the display item. The following types are defined:
  - **icon**: a .gif, .jpg, or .png file
  - **parameter**: a parameter in the instance
  - **text**: a block of text
  - **group**: a group. If the **groupName** is also defined, the new group is a child of the **groupName** group. If **groupName** is an empty string, the group is top-level.
  - **action**: an action defined by a callback procedure when you click the button labeled by **actionName**.
- **additionalInfo**: Provides extra information required for display items. The following examples illustrate how you use the **additionalInfo** argument for the various types:
  - add_display_item **groupName** **id** icon path-to-image-file
  - add_display_item **groupName** parameterName parameter (additionalInfo not required)
  - add_display_item **groupName** **id** text "your-text"
    The **your-text** argument is a block of text that is displayed in the GUI. Some simple HTML formatting is allowed, such as `<b>` and `<i>`, if the text starts with "html>".
  - add_display_item **parentGroupName** **childGroupName** group
    The **tab** is an optional parameter. If present, the group appears in separate tab in the GUI for the instance.
  - add_display_item **parentGroupName** **actionName** action buttonClickCallbackProc

| Examples | add_display_item timing read_latency parameter  
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>add_display_item sound speaker icon speaker.jpg</td>
</tr>
</tbody>
</table>
**get_display_items**
This command returns a list of all items to be displayed as part of the parameterization GUI.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_display_items</td>
</tr>
<tr>
<td>Returns</td>
<td>List of strings</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td>get_display_items</td>
</tr>
</tbody>
</table>

**get_display_item_properties**
This command returns a list of properties that can be set on display items.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_display_item_properties</td>
</tr>
<tr>
<td>Returns</td>
<td>List of strings</td>
</tr>
<tr>
<td>Arguments</td>
<td>None</td>
</tr>
<tr>
<td>Example</td>
<td>get_display_item_properties</td>
</tr>
</tbody>
</table>

**get_display_item_property**
This command returns the value of a property that can be set on a display item.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_display_item_property &lt;itemName&gt; &lt;propertyName&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>itemName</td>
</tr>
<tr>
<td></td>
<td>propertyName</td>
</tr>
<tr>
<td>Example</td>
<td>set my_label [get_display_item_property my_action DISPLAY_NAME]</td>
</tr>
</tbody>
</table>
**set_display_item_property**
This command sets the value of a property of a display item that is part of the parameterization GUI.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_display_item_property &lt;itemName&gt; &lt;propertyName&gt; &lt;value&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
</tbody>
</table>
| Arguments             | itemName: The name of the display item whose property value is being set.  
                      | propertyName: The property whose value is being set.  
                      | value: The value to set. |
| Example               | set_display_item_property my_action DISPLAY_NAME "Click Me"  
                      | set_display_item_property my_action DESCRIPTION "clicking this button runs the click_me_callback procedure in the hw.tcl file" |

### Display Item Properties Table

**Table 11-5. Display Items Properties** (Part 1 of 2)

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Type</th>
<th>Default</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DESCRIPTION</td>
<td>String</td>
<td></td>
<td>Main program</td>
<td>For an ACTION display item, updates the description /tooltip for the action button.</td>
</tr>
</tbody>
</table>
| DISPLAY_HINT     | String,"" |         | Main program     | Provides a hint about how to display a parameter. The following values are possible:  
                      |           |                     | - boolean—for integer parameters whose value can be 0 or 1. The parameter displays as an option that you can turn on or off.  
                      |           |                     | - radio—displays a parameter with a list of values as radio buttons instead of a drop-down list.  
                      |           |                     | - hexadecimal—for integer parameters, display and interpret the value as a hexadecimal number, for example:  
                      |           |                     | - fixed_size—if the parameter is displayed in a table, the fixed_size DISPLAY_HINT eliminates the add and remove buttons from tables. |
| ENABLED          | Boolean   | true    | Main program and elaboration callbacks | You can be use this property to enable or disable PARAMETER, GROUP, and ACTION display items. For a GROUP display item, this enables or disables all items in the group. |
| GROUP            | String,"" |         | Main             | Controls the grouping of parameters in GUI. |
| PATH             | String    |         | Main program     | For an ICON display item, updates path to the file for the displayed graphics, and allows the image to be switched dynamically. |
Interfaces and Ports

You can use the interface and port commands to define interfaces and ports and retrieve their properties.

add_interface

This command adds an interface to your module. As the component author, you choose the name of the interface. By default, interfaces are enabled. You can set the interface property `ENABLED` to `false` to disable a component interface. If an interface is disabled, it is hidden and its ports are automatically terminated to their default values. Active high signals are terminated to 0, and active low signals are terminated 1.

Table 11–5. Display Items Properties (Part 2 of 2)

<table>
<thead>
<tr>
<th>Property Name</th>
<th>Type</th>
<th>Default</th>
<th>Can Be Set</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>TEXT</td>
<td>string</td>
<td></td>
<td>Main program and elaboration callbacks</td>
<td>For a TEXT display item, updates the displayed text to some new text. You can use a TEXT display item as a replacement for derived parameters to display text to users; provides a cleaner UI and reduces clutter when authoring parameters.</td>
</tr>
<tr>
<td>VISIBLE</td>
<td>Boolean</td>
<td>true</td>
<td>Main program elaboration, callbacks</td>
<td>Makes PARAMETER, ICON, GROUP, TEXT, and ACTION display items visible or hidden. For a GROUP display item, shows or hides all items in the group.</td>
</tr>
</tbody>
</table>

**Interfaces and Ports**

You can use the interface and port commands to define interfaces and ports and retrieve their properties.

**add_interface**

This command adds an interface to your module. As the component author, you choose the name of the interface. By default, interfaces are enabled. You can set the interface property `ENABLED` to `false` to disable a component interface. If an interface is disabled, it is hidden and its ports are automatically terminated to their default values. Active high signals are terminated to 0, and active low signals are terminated 1.
add_interface (Part 2 of 2)

Example

```
add_interface mm_slave avalon slave
```

Notes:

(1) The terms master, source, and start are interchangeable. The terms slave, sink, and end are interchangeable.

---

### add_interface_port

This command adds a port to an interface on your module. The name must match the name of a signal on the top-level module in the HDL of your component. The port width and direction must be set by the end of the elaboration phase. The port width can be set with one of the following mechanisms:

- A variable-width expression can be set in the main program.
- A constant width set in the main program and updated in the elaboration callback.

### Arguments

<table>
<thead>
<tr>
<th><strong>Argument</strong></th>
<th><strong>Description</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><code>interfaceName</code></td>
<td>The name of the interface to which the port belongs.</td>
</tr>
<tr>
<td><code>portName</code></td>
<td>The name of the port that matches a signal name on the top-level module in the component’s HDL files.</td>
</tr>
<tr>
<td><code>portRole</code></td>
<td>The role of this port within the interface. Port roles are referred to as signal types in the Avalon Interface Specification. Refer to the Avalon Interface Specifications for the signal types available for each interface type.</td>
</tr>
<tr>
<td><code>direction</code></td>
<td>The direction can be input, output, or bidir.</td>
</tr>
<tr>
<td><code>width_expr</code></td>
<td>The port’s width expression. In simple cases, this is just the width of the port in bits.</td>
</tr>
</tbody>
</table>

### Example

```
add_interface_port mm_slave s0_rdata readdata output 32.
```
**get_interface_assignment**

This command returns the value of the specified name for the specified interface.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_interface_assignments &lt;interfaceName&gt; &lt;name&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>interfaceName</td>
</tr>
<tr>
<td></td>
<td>name</td>
</tr>
<tr>
<td>Example</td>
<td>get_interface_assignment s1 embeddedsw.configuration.isFlash</td>
</tr>
</tbody>
</table>

**get_interface_assignments**

This command returns the value of all interface assignments for the specified interface.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_interface_assignments &lt;interfaceName&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>interfaceName</td>
</tr>
<tr>
<td>Example</td>
<td>get_interface_assignments s1</td>
</tr>
</tbody>
</table>

**get_interface_ports**

This command returns the names of all of the ports that have been added to a given interface. If the interface name is omitted, all ports for all interfaces are returned.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and fileset</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_interface_ports {&lt;interfaceName&gt;}</td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>interfaceName</td>
</tr>
<tr>
<td>Example</td>
<td>get_interface_ports mm_slave</td>
</tr>
</tbody>
</table>
**get_interface_properties**

This command returns the names of all the available interface properties for the specified interface as a space separated list.

<table>
<thead>
<tr>
<th>get_interface_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

The properties available for each interface type are different. Refer to the *Avalon Interface Specifications* for more information about interface properties. The interface properties that are common to all interface types are listed below in Table 11–6.

**get_interface_property**

This command returns the value of a single interface property from the specified interface.

<table>
<thead>
<tr>
<th>get_interface_property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>
get_port_properties

This command returns a list of all available port properties.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_port_properties &lt;portName&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>String, boolean, or units, depending on property. Refer to Table 11–7 on page 11–29 for more information.</td>
</tr>
<tr>
<td>Arguments</td>
<td>portName</td>
</tr>
<tr>
<td>Example</td>
<td>get_port_properties mm_slave</td>
</tr>
</tbody>
</table>

get_port_property

This command returns the value of single port property for the specified port.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and fileset</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>get_port_property &lt;portName&gt; &lt;propertyName&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>Depends on the type of the property</td>
</tr>
<tr>
<td>Arguments</td>
<td>portName</td>
</tr>
<tr>
<td></td>
<td>propertyName</td>
</tr>
<tr>
<td>Example</td>
<td>get_port_property rdata WIDTH_VALUE</td>
</tr>
</tbody>
</table>
set_interface_assignment

This command sets the value of the specified assignment for the specified interface.

<table>
<thead>
<tr>
<th>set_interface_assignment</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

For more information about the use of the set_interface_assignment command, refer to the Publishing Component Information to Embedded Software chapter in the Nios II Software Developer’s Handbook.

set_interface_property

This command sets a single interface property for an interface.

<table>
<thead>
<tr>
<th>set_interface_property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

The properties available for each interface type are different. The ENABLED property applies to all interface types. Refer to the Avalon Interface Specifications for a description of other properties.
**set_port_property**
This command sets a single port property.

<table>
<thead>
<tr>
<th>Property</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EXPORT_OF</td>
<td>String</td>
<td>String for composed _hw1.tcl files, the EXPORT_OF property indicates which interface of a child instance is to be exported through this interface. Before using this command, you must have created the border interface using add_interface. The interface to be exported is of the form &lt;instanceName.interfaceName&gt;. Example: set_interface_property CSC_input EXPORT_OF my_colorSpaceConverter.input_port.</td>
</tr>
<tr>
<td>ENABLED</td>
<td>Boolean</td>
<td>Specifies whether or not interface is enabled.</td>
</tr>
<tr>
<td>SVD_ADDRESS_GROUP &lt;int&gt;</td>
<td>Integer</td>
<td>Integer &lt;int&gt; is the number of the address group. If two masters with the same SVD_ADDRESS_GROUP group number and one of the masters has a System View Description (.svd) file associated with it, then all the slaves connected to the second master are added to the merged .svd file of the first master. Refer to “set_interface_property” on page 11–27.</td>
</tr>
<tr>
<td>SVD_ADDRESS_OFFSET</td>
<td>Integer</td>
<td>Integer Slaves connected to the master of name &lt;interface_name&gt; are adjusted by the &lt;offset&gt; number. Refer to “set_interface_property” on page 11–27.</td>
</tr>
<tr>
<td>CMSIS_SVD_FILE &lt;path_to_svd_file&gt;</td>
<td>String</td>
<td>The path to the .svd file.</td>
</tr>
</tbody>
</table>

**set_port_property**

Callback availability: Main program and elaboration

**Usage**

`set_port_property <portName> <propertyName> [value]`

**Returns**

String, boolean, or units, depending on property. Refer to Table 11–7 on page 11–29.

**Arguments**

- **portName**: The name of the port.
- **propertyName**: One of the supported properties described in Table 11–7.
- **value**: The value to set.

**Example**

`set_port_property rdata WIDTH_EXPR 32`
## Port Properties Table

**Table 11-7. Port Properties**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>DIRECTION</strong></td>
<td>input, output, bidir</td>
<td>The direction of the port from the component’s perspective.</td>
</tr>
<tr>
<td><strong>TERMINATION</strong></td>
<td>boolean</td>
<td>When true, instead of connecting the port to the Qsys system, it is left</td>
</tr>
<tr>
<td></td>
<td></td>
<td>unconnected for output and bidir or set to a fixed value for input. Has no</td>
</tr>
<tr>
<td></td>
<td></td>
<td>effect for components that implement a generation callback instead of using</td>
</tr>
<tr>
<td></td>
<td></td>
<td>the default wrapper generation.</td>
</tr>
<tr>
<td><strong>TERMINATION_VALUE</strong></td>
<td>integer</td>
<td>The constant value to drive an input port.</td>
</tr>
<tr>
<td><strong>VHDL_TYPE</strong></td>
<td>std_logic, std_logic_vector, auto</td>
<td>indicates the type of a VHDL port. The default value, auto, selects std_logic if the width is fixed at 1, and std_logic_vector otherwise.</td>
</tr>
<tr>
<td><strong>WIDTH_VALUE</strong></td>
<td>integer</td>
<td>The width of the port in bits. Cannot be set directly. Any changes</td>
</tr>
<tr>
<td></td>
<td></td>
<td>must be set through the WIDTH_EXPR property.</td>
</tr>
<tr>
<td><strong>WIDTH_EXPR</strong></td>
<td>string</td>
<td>The width expression of a port. The width_value_expr property can be set</td>
</tr>
<tr>
<td></td>
<td></td>
<td>directly to an integer if desired.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>When get_port_property is used, width always returns the current integer</td>
</tr>
<tr>
<td></td>
<td></td>
<td>width of the port while width_expr always returns the unevaluated width</td>
</tr>
<tr>
<td></td>
<td></td>
<td>expression.</td>
</tr>
<tr>
<td><strong>DRIVEN_BY</strong></td>
<td>integer, input</td>
<td>Indicates that this output port is always driven to a constant value or by</td>
</tr>
<tr>
<td></td>
<td></td>
<td>an input port. If all outputs on a component have their driven_by property</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set to a valid value then the component’s HDL is generated automatically.</td>
</tr>
<tr>
<td><strong>ROLE</strong></td>
<td>string</td>
<td>Specifies an Avalon signal type such as waitrequest, readdata, or read.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>For a complete list of signal types, refer to the [Avalon Interface</td>
</tr>
<tr>
<td><strong>FRAGMENT_LIST</strong></td>
<td>string</td>
<td>Specifies a split or join from the component definition in the Qsys component library and the instantiation in a Qsys system.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>For example, if the library component has two output ports, data1[7:0] and</td>
</tr>
<tr>
<td></td>
<td></td>
<td>data2[7:0], which the instantiation joins to create a single output port,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>data1_2[7:0], the fragment_list defines data1_2[7:0].fragment_list =</td>
</tr>
<tr>
<td></td>
<td></td>
<td>{data1[7:0],data2[7:0]} using the following command:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_port_property data_1_2 FRAGMENT_LIST {</td>
</tr>
<tr>
<td></td>
<td></td>
<td>data1[7:0], data2[7:0]}.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>If the library component has a vectored signal, bus[3:0], that is split</td>
</tr>
<tr>
<td></td>
<td></td>
<td>into 4 separate ports in the instantiation, bus_0..bus_3, then the</td>
</tr>
<tr>
<td></td>
<td></td>
<td>following commands specify the separate ports:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_port_property bus_0 FRAGMENT_LIST { bus@0 }</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_port_property bus_1 FRAGMENT_LIST { bus@1 }</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_port_property bus_2 FRAGMENT_LIST { bus@2 }</td>
</tr>
<tr>
<td></td>
<td></td>
<td>set_port_property bus_3 FRAGMENT_LIST { bus@3 }</td>
</tr>
</tbody>
</table>
Composition

This section describes the commands that allow you to build a component by combining instances of other components. It also includes commands to query the child instances in the component.

```add_instance```

The `add_instance` command adds an instance of a component, referred to as a child or child module, to the component. You can use this command to create components that are composed of other components.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>add_instance &lt;instanceName&gt; &lt;type&gt; [&lt;version&gt;]</code></td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td><strong>instanceName</strong> Specifies a unique local name that you can use to manipulate the module. This name is used in the generated HDL to identify the module.</td>
</tr>
<tr>
<td></td>
<td><strong>type</strong> The type refers to a module available in the component library, for example <code>altera_avalon_uart</code>.</td>
</tr>
<tr>
<td></td>
<td><strong>version</strong> The required version of the specified module. If no version is specified, the latest version is used.</td>
</tr>
<tr>
<td>Example</td>
<td><code>add_instance my_uart altera_avalon_uart</code></td>
</tr>
</tbody>
</table>

```add_connection```

This command connects the named interfaces on child instances together using an appropriate connection type. Both interface names consist of a child instance name, followed by the name of an interface provided by that module. For example, `mux0.out` is the interface named `out` on the instance named `mux0`. The command returns the name of the newly added connection in `start.point/end.point` format. Be careful to connect the start to the end, and not the other way around.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>add_connection &lt;start.Interface&gt; [&lt;end.Interface&gt;] [kind] [name]</code></td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td><strong>start.interface</strong> The start interface to be connected, of the form, <code>&lt;instance_name&gt;.&lt;interface_name&gt;</code></td>
</tr>
<tr>
<td></td>
<td><strong>end.interface</strong> The end interface to be connected, <code>&lt;instance_name&gt;.&lt;interface_name&gt;</code></td>
</tr>
<tr>
<td></td>
<td><strong>kind</strong> Indicates the interface type. For a list of interface types refer to &quot;add_interface&quot; on page 11–22.</td>
</tr>
<tr>
<td></td>
<td><strong>name</strong> Specifies the name of the connection. If omitted, the name is of the form <code>start-module.start-interface/end-module.end-interface</code>.</td>
</tr>
<tr>
<td>Example</td>
<td><code>add_connection dma.read_master sdram.s1</code></td>
</tr>
</tbody>
</table>
get_connections
This command returns a list of connections. If no argument is specified, all connections in the component are returned. If a child instance is specified, all connections to interfaces on the instance are returned. If an interface on a child instance is specified, only connections to that interface are returned.

<table>
<thead>
<tr>
<th>get_connections</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

get_connection_parameters
This command gets the names of all parameters for the connection specified.

<table>
<thead>
<tr>
<th>get_connection_parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>

get_connection_parameter_value
This command gets the value of a parameter on the connection.

<table>
<thead>
<tr>
<th>get_connection_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Callback availability</td>
</tr>
<tr>
<td>Usage</td>
</tr>
<tr>
<td>Returns</td>
</tr>
<tr>
<td>Arguments</td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td>Example</td>
</tr>
</tbody>
</table>
**get_instances**
This command lists the instance names of all child modules in the component.

<table>
<thead>
<tr>
<th>get_instances</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_instance_interfaces**
This command returns the names of all of the interfaces of a child module. The interfaces can change when the parameterization of the module changes.

<table>
<thead>
<tr>
<th>get_instance_interfaces</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_instance_interface_ports**
This command returns a list of the names of the ports on the specified interface.

<table>
<thead>
<tr>
<th>get_instance_interface_ports</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**get_instance_interface_properties**
This command returns the names of all of the properties of the specified interface.

<table>
<thead>
<tr>
<th>get_instance_interface_properties</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
</tbody>
</table>
**get_instance_interface_properties**

Returns: String

Arguments:
- `instanceName` Specifies the instance name of the module.
- `interfaceName` Specifies an interface of instance.

Example:
```
get_instance_interface_properties my_ColorSpaceConverter
inputInterface
```

**get_instance_interface_property**

This command returns the value of a property associated with the specified module interface.

**get_instance_parameters**

This command gets the parameters for an existing instance where the return value is an array of key/value pairs. It omits parameters that are derived and those that have the `SYSTEM_INFO` parameter property set.

**get_instance_parameter_property**

This command returns the names of the specified instance parameter property.
### `get_instance_parameter_property`

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
<td><code>get_instance_parameter_property &lt;instanceName&gt; &lt;parameterName&gt; &lt;propertyName&gt;</code></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
<td>String, boolean, or units, depending on property. Refer to Table 11–3 on page 11–14.</td>
</tr>
</tbody>
</table>
| **Arguments**         | `
| `instanceName`        | Specifies the instance name of the module. |
| `parameterName`       | Specifies the parameter for which a property is being retrieved. |
| `propertyName`        | Specifies the property whose value is being retrieved. |
| **Example**           | `get_instance_parameter_property my_stereo separate_control DISPLAY_NAME` |
get_instance_parameter_value
This command returns the value of a parameter in a child instance. You cannot use this command to get the value of parameters whose values are derived or those that are defined using the `SYSTEM_INFO` parameter property.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, validation, and composition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>get_instance_parameter_value &lt;instanceName&gt; &lt;parameterName&gt;</code></td>
</tr>
<tr>
<td>Returns</td>
<td>String, boolean, or units, depending on property. Refer to Table 11-3 on page 11-14.</td>
</tr>
<tr>
<td>Arguments</td>
<td></td>
</tr>
<tr>
<td>instanceName</td>
<td>Specifies the name of the instance whose parameter is being retrieved.</td>
</tr>
<tr>
<td>parameterName</td>
<td>Specifies the parameter whose value is being retrieved.</td>
</tr>
<tr>
<td>Examples</td>
<td><code>get_instance_parameter_value pixel_converter input_DPI</code></td>
</tr>
<tr>
<td></td>
<td><code>proc get_instance_parameter_value {instance_name param_name} {...}</code></td>
</tr>
<tr>
<td>Notes</td>
<td>You can use this command with instances created by the <code>add_instance</code> command.</td>
</tr>
<tr>
<td>See Also</td>
<td>&quot;get_instance_parameters&quot;,  &quot;get_instances&quot;,  &quot;set_instance_parameter_value&quot;</td>
</tr>
</tbody>
</table>
**get_instance_port_property**

This command returns a information about the port property specified.

<table>
<thead>
<tr>
<th>get_instance_port_property</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>

**set_connection_parameter_value**

This command sets a property of the connection. The start and end are each interface names of the format <instance>.<interface>. Connection parameters depend on the type of connection, for memory-mapped they include base addresses and arbitration priorities.

<table>
<thead>
<tr>
<th>set_connection_parameter_value</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Callback availability</strong></td>
</tr>
<tr>
<td><strong>Usage</strong></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
</tr>
<tr>
<td><strong>Arguments</strong></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td></td>
</tr>
<tr>
<td><strong>Example</strong></td>
</tr>
</tbody>
</table>
### set_instance_parameter_value

This command sets the value of a parameter for a child instance. Derived parameters and `SYSTEM_INFO` parameters for the child instance can not be set using this command.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, composition</th>
</tr>
</thead>
</table>

**Usage**

```tcl
set_instance_parameter_value <instance> <parameter> <value>
```

**Returns**

None

**Arguments**

- **instanceName**
  
  Specifies the name of the child module.

- **parameterName**
  
  Specifies the parameter that is being set.

- **parameterValue**
  
  Specifies the value of the parameter that is being set.

**Examples**

```tcl
set_instance_parameter_value pixel_converter input_DPI 1200
proc set_instance_parameter_value {instance_name param_name value}{…}
```

**Notes**

You can use this command with instances created by the `add_instance` command.

**See Also**

"get_instance_parameter_value", "get_instances"

### Fileset Generation

This section covers the commands that create files to define the component and provide information to downstream tools.

**add_fileset**

This command adds a generation fileset for a particular target as specified by `<filesetKind>`. This target (SIM_VHDL, SIM_VERILOG, QUARTUS_SYNTH, or EXAMPLE_DESIGN) is called by Qsys when the specified generation target is requested. You may define multiple filesets for each kind of fileset. The specified callback procedure must have a single argument. The value of this argument is a generated name which must be used in the top-level module or entity declaration of your component. To override this generated name, you may set the fileset property `TOP_LEVEL`.

---

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Overriding the generated name is only possible if all parameterizations of a core yield identical HDL.

### add_fileset

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Usage</strong></td>
<td><code>add_fileset &lt;filesetName&gt; &lt;filesetKind&gt; &lt;callbackProcName&gt; [&lt;displayName&gt;]</code></td>
</tr>
<tr>
<td><strong>Returns</strong></td>
<td>String</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Arguments</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>filesetName</code></td>
<td>The name of the fileset.</td>
</tr>
<tr>
<td><code>filesetKind</code></td>
<td>Files support the following kinds:</td>
</tr>
<tr>
<td></td>
<td>- <code>SIM_VHDL</code></td>
</tr>
<tr>
<td></td>
<td>- <code>SIM_VERILOG</code></td>
</tr>
<tr>
<td></td>
<td>- <code>QUARTUS_SYNTH</code></td>
</tr>
<tr>
<td></td>
<td>- <code>EXAMPLE_DESIGN</code></td>
</tr>
<tr>
<td><code>callbackProcName</code></td>
<td>A string identifying the name of the callback procedure. If you add files in the global section, you can then specify a blank callback procedure.</td>
</tr>
<tr>
<td><code>displayName</code></td>
<td>A display string to identify the fileset.</td>
</tr>
</tbody>
</table>

**Example** `add_fileset PCIE_SYNTHESIS QUARTUS_SYNTH mySynthProc`
add_fileset_file

This command adds an output file for the generation directory. You can specify source file locations using either an absolute path or a path that is relative to the component’s _hw.tcl file.

**Callback availability** Fileset

**Usage**

```tcl
add_fileset_file <fileDestination> <fileKind> <fileSource> <contentsOrPath>
```

**Returns** String

<table>
<thead>
<tr>
<th>Arguments</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>fileDestination</td>
<td>Specifies the output file for the file after Qsys generation. Files support the following kinds:</td>
</tr>
<tr>
<td></td>
<td>■ VERILOG</td>
</tr>
<tr>
<td></td>
<td>■ SYSTEM_VERILOG</td>
</tr>
<tr>
<td></td>
<td>■ SYSTEM_VERILOG_INCLUDE</td>
</tr>
<tr>
<td></td>
<td>■ VHDL</td>
</tr>
<tr>
<td></td>
<td>■ SDC</td>
</tr>
<tr>
<td></td>
<td>■ MIF</td>
</tr>
<tr>
<td></td>
<td>■ HEX</td>
</tr>
<tr>
<td></td>
<td>■ DAT</td>
</tr>
<tr>
<td></td>
<td>■ OTHER</td>
</tr>
<tr>
<td>fileKind</td>
<td>The following sources are defined:</td>
</tr>
<tr>
<td></td>
<td>■ PATH—specifies the original source file to be copied to filePath.</td>
</tr>
<tr>
<td></td>
<td>■ TEXT—specifies an arbitrary text string for the contents of the file.</td>
</tr>
<tr>
<td>fileSource</td>
<td>When the fileSource is PATH, specifies the file to be copied to filePath. When the fileSource is TEXT, specifies the text string to be stored in the file.</td>
</tr>
<tr>
<td>contentsOrPath</td>
<td>When the fileSource is PATH, specifies the file to be copied to filePath. When the fileSource is TEXT, specifies the text string to be stored in the file.</td>
</tr>
</tbody>
</table>

**Example**

```tcl
add_fileset_file "./implementation/rx_pma.sv" SYSTEM_VERILOG PATH synth_rx_pma.sv
add_fileset_file gui.sv SYSTEM_VERILOG TEXT "Customize your IP core"
```

create_temp_file

This command creates a temporary file which can be manipulated inside the fileset callbacks of a _hw.tcl file. This temporary file can serve as a scratch pad or can be included in the generation output if it is included using the add_fileset_file command.

**Callback availability** Fileset

**Usage**

```tcl
create_temp_file <fileName>
```

**Returns** String

<table>
<thead>
<tr>
<th>Arguments</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>fileName</td>
<td>The name of the created file.</td>
</tr>
</tbody>
</table>

**Example**

```tcl
set filelocation 
add_fileset_file compute_frequency.v VERILOG PATH "$filelocation"
```
**set_fileset_property**

Allows a user to set the properties of a fileset.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Fileset Generation</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>set_fileset_property &lt;filename&gt; &lt;filesetProperty&gt; &lt;value&gt;</code></td>
</tr>
<tr>
<td>Returns</td>
<td>String</td>
</tr>
<tr>
<td>Arguments</td>
<td>filename The name of the fileset.</td>
</tr>
<tr>
<td></td>
<td>filesetProperty TOP_LEVEL The of the top-level HDL module produced by this fileset</td>
</tr>
<tr>
<td>Example</td>
<td><code>set_fileset_property mySynthFileset TOP_LEVEL simple_uart</code></td>
</tr>
</tbody>
</table>

**Miscellaneous**

**check_device_family_equivalence**

This command returns 1 if the device family is equivalent to one of the families in the list, and returns 0 if the device family is not equivalent to any families.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, and fileset</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>check_device_family_equivalence &lt;deviceName&gt; &lt;deviceList&gt;</code></td>
</tr>
<tr>
<td>Returns</td>
<td>1 or 0 Based on equivalence.</td>
</tr>
<tr>
<td>Arguments</td>
<td>deviceName The device name of device being checked.</td>
</tr>
<tr>
<td></td>
<td>deviceList The device string list to check against.</td>
</tr>
<tr>
<td>Example</td>
<td><code>check_device_family_equivalence &quot;CYLCONE III LS&quot; { &quot;stratixv&quot; &quot;Cyclone IV&quot; &quot;cycloneiiils&quot; }</code></td>
</tr>
</tbody>
</table>

**get_device_family_displayname**

This command returns the display name of a given device family.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration, fileset</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td><code>get_device_family_displayname &lt;deviceName&gt;</code></td>
</tr>
<tr>
<td>Returns</td>
<td>A user-suitable display name for a device family.</td>
</tr>
<tr>
<td>Arguments</td>
<td>A device family name (example <code>stratixv</code>, <code>arraii</code>).</td>
</tr>
<tr>
<td>Example</td>
<td><code>get_device_family_displayname cycloneiiils</code> (Returns: Cyclone III LS)</td>
</tr>
</tbody>
</table>
set_qip_strings

This command places strings in the Quartus II IP File (.qip) file. The .qip file contains paths to the files for an IP core. You add the .qip file to your Quartus II project in the under Files in the Settings dialog box. Successive calls to set_qip_strings are not additive, they replace the previously declared value.

<table>
<thead>
<tr>
<th>Callback availability</th>
<th>Main, elaboration</th>
</tr>
</thead>
<tbody>
<tr>
<td>Usage</td>
<td>set_qip_strings &lt;qip Entries&gt;</td>
</tr>
<tr>
<td>Returns</td>
<td>The Tcl list which was set</td>
</tr>
<tr>
<td>Arguments</td>
<td>qip Entries</td>
</tr>
<tr>
<td>Macros</td>
<td>%entityName%</td>
</tr>
<tr>
<td></td>
<td>%libraryName%</td>
</tr>
<tr>
<td>Example</td>
<td>set_qip_strings {&quot;QIP Entry 1&quot; &quot;QIP Entry 2&quot;}</td>
</tr>
</tbody>
</table>

Port Roles (Interface Signal Types)

Each interfaces defines a number of signal roles and their behavior. Many signal roles are optional, allowing component designers the flexibility to select only the signal roles necessary to implement the required functionality.

**AXI Interface Signal Types**

For complete AXI interface specifications, refer to the AMBA Protocol Specifications on the ARM® website

**AXI Master Interface Signal Types**

<table>
<thead>
<tr>
<th>Name</th>
<th>Direction</th>
<th>Width</th>
</tr>
</thead>
<tbody>
<tr>
<td>araddr</td>
<td>Output</td>
<td>1 to 64</td>
</tr>
<tr>
<td>arburst</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>arcache</td>
<td>Output</td>
<td>4</td>
</tr>
<tr>
<td>arid</td>
<td>Output</td>
<td>1 to 18</td>
</tr>
<tr>
<td>arlen</td>
<td>Output</td>
<td>4</td>
</tr>
<tr>
<td>arlock</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>arprot</td>
<td>Output</td>
<td>3</td>
</tr>
<tr>
<td>arready</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>arsize</td>
<td>Output</td>
<td>3</td>
</tr>
<tr>
<td>aruser</td>
<td>Output</td>
<td>1 to 64</td>
</tr>
<tr>
<td>arvalid</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>awaddr</td>
<td>Output</td>
<td>1 to 64</td>
</tr>
</tbody>
</table>
### Table 11–8. altera_axi_master (Part 2 of 2)

<table>
<thead>
<tr>
<th>Name</th>
<th>Direction</th>
<th>Width</th>
</tr>
</thead>
<tbody>
<tr>
<td>awburst</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>awcache</td>
<td>Output</td>
<td>4</td>
</tr>
<tr>
<td>awid</td>
<td>Output</td>
<td>1 to 18</td>
</tr>
<tr>
<td>awlen</td>
<td>Output</td>
<td>4</td>
</tr>
<tr>
<td>awlock</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>awprot</td>
<td>Output</td>
<td>3</td>
</tr>
<tr>
<td>awready</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>awsize</td>
<td>Output</td>
<td>3</td>
</tr>
<tr>
<td>awuser</td>
<td>Output</td>
<td>1 to 64</td>
</tr>
<tr>
<td>awvalid</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>bid</td>
<td>Input</td>
<td>1 to 18</td>
</tr>
<tr>
<td>bready</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>bresp</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>bvalid</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>rdata</td>
<td>Input</td>
<td>8, 16, 32, 64, 128, 256, 512, 1024</td>
</tr>
<tr>
<td>rid</td>
<td>Input</td>
<td>1 to 18</td>
</tr>
<tr>
<td>rlast</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>rready</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>rresp</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>rvalid</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>wdata</td>
<td>Output</td>
<td>8, 16, 32, 64, 128, 256, 512, 1024</td>
</tr>
<tr>
<td>wid</td>
<td>Output</td>
<td>1 to 18</td>
</tr>
<tr>
<td>wlast</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>wready</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>wstrb</td>
<td>Output</td>
<td>1, 2, 4, 8, 16, 32, 64, 128</td>
</tr>
<tr>
<td>wvalid</td>
<td>Output</td>
<td>1</td>
</tr>
</tbody>
</table>

### AXI Slave Interface Signals

### Table 11–9. altera_axi_slave (Part 1 of 2)

<table>
<thead>
<tr>
<th>Name</th>
<th>Direction</th>
<th>Width</th>
</tr>
</thead>
<tbody>
<tr>
<td>araddr</td>
<td>Input</td>
<td>1 to 64</td>
</tr>
<tr>
<td>arburst</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>arcache</td>
<td>Input</td>
<td>4</td>
</tr>
<tr>
<td>arid</td>
<td>Input</td>
<td>1 to 18</td>
</tr>
<tr>
<td>arlen</td>
<td>Input</td>
<td>4</td>
</tr>
<tr>
<td>arlock</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>arprot</td>
<td>Input</td>
<td>3</td>
</tr>
<tr>
<td>arready</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>arsize</td>
<td>Input</td>
<td>3</td>
</tr>
</tbody>
</table>
### Table 11–9. altera_axi_slave (Part 2 of 2)

<table>
<thead>
<tr>
<th>Name</th>
<th>Direction</th>
<th>Width</th>
</tr>
</thead>
<tbody>
<tr>
<td>aruser</td>
<td>Input</td>
<td>1 to 64</td>
</tr>
<tr>
<td>arvalid</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>awaddr</td>
<td>Input</td>
<td>1 to 64</td>
</tr>
<tr>
<td>awburst</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>awcache</td>
<td>Input</td>
<td>4</td>
</tr>
<tr>
<td>awid</td>
<td>Input</td>
<td>1 to 18</td>
</tr>
<tr>
<td>awlen</td>
<td>Input</td>
<td>4</td>
</tr>
<tr>
<td>awlock</td>
<td>Input</td>
<td>2</td>
</tr>
<tr>
<td>awprot</td>
<td>Input</td>
<td>3</td>
</tr>
<tr>
<td>awready</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>awsize</td>
<td>Input</td>
<td>3</td>
</tr>
<tr>
<td>awuser</td>
<td>Input</td>
<td>1 to 64</td>
</tr>
<tr>
<td>awvalid</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>bid</td>
<td>Output</td>
<td>1 to 18</td>
</tr>
<tr>
<td>bready</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>bresp</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>bvalid</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>rdata</td>
<td>Output</td>
<td>8, 16, 32, 64, 128, 256, 512, 1024</td>
</tr>
<tr>
<td>rid</td>
<td>Output</td>
<td>1 to 18</td>
</tr>
<tr>
<td>rlast</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>rready</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>rresp</td>
<td>Output</td>
<td>2</td>
</tr>
<tr>
<td>rvalid</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>wdata</td>
<td>Input</td>
<td>8, 16, 32, 64, 128, 256, 512, 1024</td>
</tr>
<tr>
<td>wid</td>
<td>Input</td>
<td>1 to 18</td>
</tr>
<tr>
<td>wlast</td>
<td>Input</td>
<td>1</td>
</tr>
<tr>
<td>wready</td>
<td>Output</td>
<td>1</td>
</tr>
<tr>
<td>wstrb</td>
<td>Input</td>
<td>1, 2, 4, 8, 16, 32, 64, 128</td>
</tr>
<tr>
<td>wvalid</td>
<td>Input</td>
<td>1</td>
</tr>
</tbody>
</table>

### APB Interface Signal Types

### Table 11–10. APB Interface Signal Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Width</th>
<th>Direction (APB master)</th>
<th>Direction (APB Slave)</th>
<th>Required</th>
</tr>
</thead>
<tbody>
<tr>
<td>paddr</td>
<td>[1:32]</td>
<td>Output</td>
<td>Input</td>
<td>Yes</td>
</tr>
<tr>
<td>psel</td>
<td>[1:16]</td>
<td>Output</td>
<td>Input</td>
<td>Yes</td>
</tr>
<tr>
<td>penable</td>
<td>1</td>
<td>Output</td>
<td>Input</td>
<td>Yes</td>
</tr>
<tr>
<td>pwrite</td>
<td>1</td>
<td>Output</td>
<td>Input</td>
<td>Yes</td>
</tr>
</tbody>
</table>
Table 11–10. APB Interface Signal Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Width</th>
<th>Direction (APB master)</th>
<th>Direction (APB Slave)</th>
<th>Required</th>
</tr>
</thead>
<tbody>
<tr>
<td>pwdata</td>
<td>[1:32]</td>
<td>Output</td>
<td>Input</td>
<td>Yes</td>
</tr>
<tr>
<td>prdata</td>
<td>[1:32]</td>
<td>Input</td>
<td>Output</td>
<td>Yes</td>
</tr>
<tr>
<td>psilerr</td>
<td>1</td>
<td>Input</td>
<td>Output</td>
<td>No</td>
</tr>
<tr>
<td>pready</td>
<td>1</td>
<td>Input</td>
<td>Output</td>
<td>Yes</td>
</tr>
<tr>
<td>paddr31</td>
<td>1</td>
<td>Output</td>
<td>Input</td>
<td>No</td>
</tr>
</tbody>
</table>

### Avalon Interface Signal Types

**Avalon Memory-Mapped Signals**

Table 11–11. Avalon-MM Signals (1) (Part 1 of 4)

<table>
<thead>
<tr>
<th>Signal role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Fundamental Signals</strong></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>address</td>
<td>1–32</td>
<td>Master → Slave</td>
<td>For masters, the address signal represents a byte address. The value of the address must be aligned to the data width. To write to specific bytes within a data word, the master must use the byteenable signal. For slaves, the interconnect translates the byte address into a word address in the slave’s address space so that each slave access is for a word of data from the perspective of the slave. For example, address 0 selects the first word of the slave and address 1 selects the second word of the slave.</td>
</tr>
<tr>
<td>begintransfer</td>
<td>1</td>
<td>Master → Slave</td>
<td>Asserted by the interconnect for the first cycle of each transfer regardless of waitrequest and other signals. If you do not include this signal in your Avalon-MM master interface, Qsys automatically generates this signal for you.</td>
</tr>
</tbody>
</table>
### Table 11–11. Avalon-MM Signals *(1)* (Part 2 of 4)

<table>
<thead>
<tr>
<th>Signal role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteenable</td>
<td>1, 2, 4, 8, 16, 32, 64, 128</td>
<td>Master → Slave</td>
<td>Enables specific byte lane(s) during transfers on ports of width greater than 8 bits. Each bit in byteenable corresponds to a byte in writedata and readdata. The master bit (&lt;n&gt;) of byteenable indicates whether byte (&lt;n&gt;) is being written to. During writes, byteenables specify which bytes are being written to; other bytes should be ignored by the slave. During reads, byteenables indicates which bytes the master is reading. Slaves that simply return readdata with no side effects are free to ignore byteenables during reads. If an interface does not have a byteenable signal, the transfer proceeds as if all byteenables are asserted. When more than one bit of the byteenable signal is asserted, all asserted lanes are adjacent. The number of adjacent lines must be a power of 2, and the specified bytes must be aligned on an address boundary for the size of the data. For example, the following values are legal for a 32-bit slave:</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>1111 writes full 32 bits</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>0011 writes lower 2 bytes</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>1100 writes upper 2 bytes</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>0001 writes byte 0 only</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>0010 writes byte 1 only</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>0100 writes byte 2 only</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>1000 writes byte 3 only</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Altera strongly recommends that you use the byteenable signal in components that are used in systems with different word sizes. Using the byteenable signal prevents unintended side effects in systems that include width adapters.</td>
</tr>
<tr>
<td>chipselect</td>
<td>1</td>
<td>Master → Slave</td>
<td>When present, a slave port ignores all Avalon-MM signals unless chipselect is asserted. chipselect must be used in combination with read or write. The chipselect signal is not necessary; Altera does not recommend using it.</td>
</tr>
<tr>
<td>chipselect_n</td>
<td>1</td>
<td>Slave</td>
<td></td>
</tr>
<tr>
<td>debugaccess</td>
<td>1</td>
<td>Master → Slave</td>
<td>When asserted, allows internal memories that are normally write-protected to be written. For example, on-chip ROM memories can only be written when debugaccess is asserted.</td>
</tr>
<tr>
<td>read</td>
<td>1</td>
<td>Master → Slave</td>
<td>Asserted to indicate a read transfer. If present, readdata is required.</td>
</tr>
<tr>
<td>read_n</td>
<td>1</td>
<td>Slave</td>
<td></td>
</tr>
<tr>
<td>readdata</td>
<td>8,16, 32, 64, 128, 256, 512, 1024</td>
<td>Slave → Master</td>
<td>The readdata driven from the slave to the master in response to a read transfer.</td>
</tr>
<tr>
<td>write</td>
<td>1</td>
<td>Master → Slave</td>
<td>Asserted to indicate a write transfer. If present, writedata is required.</td>
</tr>
<tr>
<td>write_n</td>
<td>1</td>
<td>Slave</td>
<td></td>
</tr>
<tr>
<td>writedata</td>
<td>8,16, 32, 64, 128, 256, 512, 1024</td>
<td>Master → Slave</td>
<td>Data for write transfers. The width must be the same as the width of readdata if both are present.</td>
</tr>
</tbody>
</table>
Port Roles (Interface Signal Types)

### Wait-State Signals

<table>
<thead>
<tr>
<th>Signal role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>lock</td>
<td>1</td>
<td>Master → Slave</td>
<td>lock ensures that once a master wins arbitration, it maintains access to the slave for multiple transactions. It is asserted coincident with the first read or write of a locked sequence of transactions, and is deasserted on the final transaction of a locked sequence of transactions. lock assertion does not guarantee that arbitration is won, but after the lock-asserting master has been granted, it retains grant until it is deasserted. A master equipped with lock cannot be a burst master. Arbitration priority values for lock-equipped masters are ignored. lock is particularly useful for read-modify-write operations, where master A reads 32-bit data that has multiple bit fields, changes one field, and writes the 32-bit data back. If lock is not used, a master B could perform a write between Master A's read and write and master A's write would overwrite master B's changes.</td>
</tr>
<tr>
<td>waitrequest</td>
<td>1</td>
<td>Slave → Master</td>
<td>Asserted by the slave when it is unable to respond to a read or write request. Forces the master to wait until the interconnect is ready to proceed with the transfer. At the start of all transfers, a master initiates the transfer and waits until waitrequest is deasserted. A master must make no assumption about the assertion state of waitrequest when the master is idle: waitrequest may be high or low, depending on system properties. When waitrequest is asserted, master control signals to the slave remain constant with the exception of begintransfer and beginbursttransfer. An Avalon-MM slave may assert waitrequest during idle cycles. An Avalon-MM master may initiate a transaction when waitrequest is asserted and wait for that signal to be deasserted. To avoid system lockup, a slave device should assert waitrequest when in reset.</td>
</tr>
<tr>
<td>readdatavalid</td>
<td>1</td>
<td>Slave → Master</td>
<td>Used for variable-latency, pipelined read transfers. Asserted by the slave to indicate that the readdata signal contains valid data in response to a previous read request. A slave with readdatavalid must assert this signal for one cycle for each read access it has received. There must be at least one cycle of latency between acceptance of the read and assertion of readdatavalid. Required if the master supports pipelined reads. Bursting masters with read functionality must include the readdatavalid signal.</td>
</tr>
</tbody>
</table>

### Pipeline Signals

<table>
<thead>
<tr>
<th>Signal role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
</table>

Table 11–11. Avalon-MM Signals (1) (Part 3 of 4)
Port Roles (Interface Signal Types)

Avalon-ST Interface Signals

Table 11–12. Avalon-ST Interface Signals (Part 1 of 2)

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>channel</td>
<td>1–128</td>
<td>Source → Sink</td>
<td>The channel number for data being transferred on the current cycle. If an interface supports the channel signal, it must also define the maxChannel parameter.</td>
</tr>
<tr>
<td>data</td>
<td>1–4096</td>
<td>Source → Sink</td>
<td>The data signal from the source to the sink, typically carries the bulk of the information being transferred. The contents and format of the data signal is further defined by parameters.</td>
</tr>
<tr>
<td>error</td>
<td>1–256</td>
<td>Source → Sink</td>
<td>A bit mask used to mark errors affecting the data being transferred in the current cycle. A single bit in error is used for each of the errors recognized by the component, as defined by the errorDescriptor property.</td>
</tr>
<tr>
<td>ready</td>
<td>1</td>
<td>Sink → Source</td>
<td>Asserted high to indicate that the sink can accept data. ready is asserted by the sink on cycle &lt;n&gt; to mark cycle &lt;n + readyLatency&gt; as a ready cycle, during which the source may assert valid and transfer data. Sources without a ready input cannot be backpressured, and sinks without a ready output never need to backpressure.</td>
</tr>
</tbody>
</table>
Table 11–12. Avalon-ST Interface Signals (Part 2 of 2)

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>valid</td>
<td>1</td>
<td>Source → Sink</td>
<td>Asserted by the source to qualify all other source to sink signals. On ready cycles where valid is asserted, the data bus and other source to sink signals are sampled by the sink, and on other cycles are ignored. Sources without a valid output implicitly provide valid data on every cycle that they are not being backpressured, and sinks without a valid input expect valid data on every cycle that they are not backpressuring.</td>
</tr>
</tbody>
</table>

Packet Transfer Signals

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>empty</td>
<td>1–8</td>
<td>Source → Sink</td>
<td>Indicates the number of symbols that are empty during cycles that contain the end of a packet. The empty signal is not used on interfaces where there is one symbol per beat. If endofpacket is not asserted, this signal is not interpreted.</td>
</tr>
<tr>
<td>endofpacket</td>
<td>1</td>
<td>Source → Sink</td>
<td>Asserted by the source to mark the end of a packet.</td>
</tr>
<tr>
<td>startofpacket</td>
<td>1</td>
<td>Source → Sink</td>
<td>Asserted by the source to mark the beginning of a packet.</td>
</tr>
</tbody>
</table>

Tri-state Slave Interface Signals

Table 11–13. Avalon-MM tri-state Slave Signals (1) (Part 1 of 2)

<table>
<thead>
<tr>
<th>Signal Type</th>
<th>Width</th>
<th>Direction</th>
<th>Req’d</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>address</td>
<td>1-32</td>
<td>In</td>
<td>No</td>
<td>Address lines to the slave port. Specifies a byte offset into the slave’s address space.</td>
</tr>
<tr>
<td>read</td>
<td>1</td>
<td>In</td>
<td>No</td>
<td>Read-request signal. Not required if the slave port never outputs data. If present, data must also be used.</td>
</tr>
<tr>
<td>read_n</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>write</td>
<td>1</td>
<td>In</td>
<td>No</td>
<td>Write-request signal. Not required if the slave port never receives data from a master. If present, data must also be present, and writebyteenable cannot be present.</td>
</tr>
<tr>
<td>write_n</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>chipselect</td>
<td>1</td>
<td>In</td>
<td>No</td>
<td>When present, the slave port ignores all Avalon-MM signals unless chipselect is asserted. chipselect is always present in combination with read or write.</td>
</tr>
<tr>
<td>chipselect_n</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>outputenable</td>
<td>1</td>
<td>In</td>
<td>Yes</td>
<td>Output-enable signal. When deasserted, a tri-state slave port must not drive its data lines otherwise data contention may occur.</td>
</tr>
<tr>
<td>outputenable_n</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>data</td>
<td>8,16, 32, 64, 128, 256, 512, 1024</td>
<td>Bidir</td>
<td>No</td>
<td>Bidirectional data. During write transfers, the FPGA drives the data lines. During read transfers the slave device drives the data lines, and the FPGA captures the data signals and provides them to the master.</td>
</tr>
</tbody>
</table>
Port Roles (Interface Signal Types)

Enables specific byte lane(s) during transfers. Each bit in `byteenable` corresponds to a byte lane in `data`. During writes, `byteenables` specify which bytes the master is writing to the slave. During reads, `byteenables` indicates which bytes the master is reading. Slaves that simply return `data` with no side effects are free to ignore `byteenables` during reads.

When more than one byte lane is asserted, all asserted lanes are guaranteed to be adjacent. The number of adjacent lines must be a power of 2, and the specified bytes must be aligned on an address boundary for the size of the data. The are legal values for a 32-bit slave:

- `1111` writes full 32 bits
- `0011` writes lower 2 bytes
- `1100` writes upper 2 bytes
- `0001` writes byte 0 only
- `0010` writes byte 1 only
- `0100` writes byte 2 only
- `1000` writes byte 3 only

Equivalent to the logical AND of the `byteenable` and `write` signals. When used, the `write` signal is not used.

 asserted for the first cycle of each transfer.

Table 11–13. Avalon-MM tri-state Slave Signals (1) (Part 2 of 2)

<table>
<thead>
<tr>
<th>Signal Type</th>
<th>Width</th>
<th>Direction</th>
<th>Req’ d</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>byteenable</code></td>
<td>2, 4, 8, 16, 32, 64, 128</td>
<td>In</td>
<td>No</td>
<td>Enables specific byte lane(s) during transfers. Each bit in <code>byteenable</code> corresponds to a byte lane in <code>data</code>. During writes, <code>byteenables</code> specify which bytes the master is writing to the slave. During reads, <code>byteenables</code> indicates which bytes the master is reading. Slaves that simply return <code>data</code> with no side effects are free to ignore <code>byteenables</code> during reads. When more than one byte lane is asserted, all asserted lanes are guaranteed to be adjacent. The number of adjacent lines must be a power of 2, and the specified bytes must be aligned on an address boundary for the size of the data. The are legal values for a 32-bit slave: <code>1111</code> writes full 32 bits <code>0011</code> writes lower 2 bytes <code>1100</code> writes upper 2 bytes <code>0001</code> writes byte 0 only <code>0010</code> writes byte 1 only <code>0100</code> writes byte 2 only <code>1000</code> writes byte 3 only</td>
</tr>
<tr>
<td><code>writebyteenable</code></td>
<td>2, 4, 8, 16, 32, 64, 128</td>
<td>In</td>
<td>No</td>
<td>Equivalent to the logical AND of the <code>byteenable</code> and <code>write</code> signals. When used, the <code>write</code> signal is not used.</td>
</tr>
<tr>
<td><code>begintransfer</code></td>
<td>1</td>
<td>In</td>
<td>No</td>
<td>Asserted for the first cycle of each transfer.</td>
</tr>
</tbody>
</table>

Note to Table 11–11:

(1) All Avalon signals are active high. Avalon signals that can also be asserted low list both versions in the Signal Type column.
**Tri-state Conduit Interface Signals**

Table 11–14. Tristate Conduit Interface Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Required</th>
<th>Description</th>
</tr>
</thead>
</table>
| request     | 1     | Master → Slave | Yes      | The meaning of request depends on the state of the grant signal, as the following rules dictate.  
1. When request is asserted and grant is deasserted, request is requesting access for the current cycle.  
2. When request is asserted and grant is asserted, request is requesting access for the next cycle; consequently, request should be deasserted on the final cycle of an access.  
Because request is deasserted in the last cycle of a bus access, it can be reassigned immediately following the final cycle of a transfer, making both rearbitration and continuous bus access possible if no other masters are requesting access.  
Once asserted, request must remain asserted until granted; consequently, the shortest bus access is 2 cycles. |
| grant       | 1     | Slave → Master | Yes      | When asserted, indicates that a tristate conduit master has been granted access to perform transactions. grant is asserted in response to the request signal and remains asserted until 1 cycle following the deassertion of request.  
The design of the Avalon-TC Interface does not allow a default Avalon-TC master to be granted when no masters are requesting. |
| <name>_in   | 1–1024 | Slave → Master | No       | The input signal of a logical tristate signal. |
| <name>_out  | 1–1024 | Master → Slave | No       | The output signal of a logical tristate signal. |
| <name>_outen| 1     | Master → Slave | No       | The output enable for a logical tristate signal. |

**Avalon Clock Sink Interface Signals**

Table 11–15. Clock Input Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Required</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>clk</td>
<td>1</td>
<td>Input</td>
<td>Yes</td>
<td>A clock signal. Provides synchronization for internal logic and for other interfaces.</td>
</tr>
</tbody>
</table>

**Avalon Clock Source Interface Signals**

Table 11–16. Clock Source Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Required</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>clk</td>
<td>1</td>
<td>Output</td>
<td>Yes</td>
<td>An output clock signal.</td>
</tr>
</tbody>
</table>
Avalon Conduit Interface Signals

Table 11–17. Conduit Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>export</td>
<td>(&lt;n&gt;)</td>
<td>In, out or bidirectional</td>
<td>A conduit interface consists of one or more signals of arbitrary width of direction input or output. Compatible conduit interfaces can be connected inside the Qsys system, exported to the next level of the hierarchical design, or to the top-level of the Qsys system.</td>
</tr>
</tbody>
</table>

Interrupt Sender Interface Signals

Table 11–18. Interrupt Sender Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Required</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>irq</td>
<td>1</td>
<td>Output</td>
<td>Yes</td>
<td>Interrupt Request. A slave asserts irq when it needs to be serviced.</td>
</tr>
<tr>
<td>irq_n</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Interrupt Receiver Interface Signals

Table 11–19. Interrupt Receiver Signal Roles

<table>
<thead>
<tr>
<th>Signal Role</th>
<th>Width</th>
<th>Direction</th>
<th>Required</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>irq</td>
<td>1–32</td>
<td>Input</td>
<td>Yes</td>
<td>irq is an (&lt;n&gt;)-bit vector, where each bit corresponds directly to one IRQ sender, with no inherent assumption of priority.</td>
</tr>
</tbody>
</table>

Document Revision History

Table 11–20 shows the revision history for this document.

Table 11–20. Document Revision History (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>■ Consolidated content from other Qsys chapters.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added AMBA APB support.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Added the demo_axi_memory example with screen shots and example_hw.tcl code.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Added AMBA AXI3 support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added the set_display_item_property, set_parameter_property parameter LONG_DESCRIPTION, and static filesets.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>■ Template update.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added commands set_qip_strings, get_qip_strings, get_device_family_displayname, check_device_family_equivalence</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated text to reflect changes in 11.1.</td>
</tr>
</tbody>
</table>
Table 11–20. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
</table>
| May 2011     | 11.0.0  | ■ Removed Beta status.  
■ Revised section describing HDL and composed component implementations.  
■ Separated reset and clock interfaces in examples.  
■ Added the TRISTATECONDUIT_INFO, GENERATION_ID, and UNIQUE_ID SYSTEM_INFO properties.  
■ Added the WIDTH and SYSTEM_INFO_ARG parameter properties.  
■ Removed the doc_type argument from the add_documentation_link command.  
■ Removed get_instance_parameter_properties command.  
(get_instance_parameter_property is available.)  
■ Added the add_fileset, add_fileset_file and create_temp_file commands.  
■ Updated Tcl examples to show separate clock and reset interfaces. |
| December 2010| 10.1.0  | Initial release.                                                                                                                      |

For more information about Tcl syntax, refer to the Tcl Developer Xchange website.

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter describes components and IP cores that you can use to design your Qsys systems. The Qsys interfaces define components appropriate for streaming high-speed data, reading and writing registers and memory, and controlling off-chip devices.

Qsys supports standard Avalon®, AMBA® AXI3™ (version 1.0), AMBA AXI4™ (version 2.0), and AMBA APB™ 3 (version 1.0) interfaces. For more information about Avalon and AMBA interfaces, refer to the Avalon Interface Specifications and the AMBA Protocol Specifications on the ARM® website. AXI4-Lite is not supported.

Bridges

Qsys provides bridge components to provide flexibility and control in your system implementation. You can use bridges to control the topology of a Qsys system. Bridges are not end points for data, but rather affect the way data is transported between components. By inserting bridges between masters and slaves, you control system topology, which in turn affects the interconnect that Qsys generates.

You can also use bridges to separate components in different clock domains and to isolate clock domain crossing logic.

This section describes the following bridge components:

- Clock Bridge
- Avalon-MM Clock Crossing Bridge
- Avalon-MM Pipeline Bridge
- Address Span Extender

AXI bridge components are not available in the Quartus II software, but you can connect AXI interfaces with other bridge types.
### Structure of a Bridge

A bridge has one slave interface and one master interface, as shown in Figure 12–1. In Qsys, one or more masters interfaces from other components connect to the bridge slave; then, the bridge master connects to one or more slave interfaces on other components. In Figure 12–1, all three masters have logical connections to all three slaves, although physically each master connects only to the bridge.

#### Figure 12–1. Example of Bridge in a Qsys System

Transfers initiated to the bridge’s slave propagate to the bridge master in the same order in which they are initiated on the bridge slave.

### Bridges Between Avalon and AXI Interfaces

Connections between AXI and Avalon interfaces are made without requiring the use of explicitly instantiated bridges; the interconnect provides all necessary bridging logic. However, this does not prevent the use of explicit bridges to separate the AXI and Avalon domains. Though, using an explicit Avalon-MM bridge to separate the AXI and Avalon domains reduces the amount of bridging logic in the interconnect, at the expense of concurrency.
Figure 12–2 shows an Avalon-MM pipeline bridge between Avalon-MM and AXI networks.

**Figure 12–2. Connecting Avalon and AXI Interfaces**

![Diagram of Avalon-MM and AXI Interfaces](image)

**Clock Bridge**

The Clock Bridge allows you to connect a clock source to multiple clock input interfaces. You can use this bridge to connect a clock source that’s outside the Qsys system through an exported interface to multiple clock input interfaces in the system. Figure 12–3 illustrates the use of a clock bridge.

**Figure 12–3. Clock Bridge**

![Diagram of Clock Bridge](image)

**Avalon-MM Clock Crossing Bridge**

The Avalon-MM Clock Crossing Bridge transfers Avalon-MM commands and responses between different clock domains. It uses asynchronous FIFOs to implement the clock crossing logic. The Clock Crossing Bridge has a number of parameters, including parameters to control the depth of the command and response FIFOs in both the master and slave clock domains. If the number of in-flight reads exceeds the depth of the response FIFO, the Clock Crossing Bridge stops sending reads. To maintain throughput for high-performance applications, increase the response FIFO depth from the default minimum depth, which is twice the maximum burst size.
You can use the Avalon-MM Clock Crossing Bridge to bridge between AXI Masters and Slaves of different clock domains. For more information about Handshake and FiFO clock crossing types, refer to the Creating a System With Qsys chapter in volume 1 of the Quartus II Handbook.

The Avalon-MM Clock Crossing Bridge core is implemented to work with the Qsys interconnect. If you migrate a Qsys design that was created with an older version of Qsys or SOPC Builder, and that includes an Avalon-MM Clock Crossing Bridge, Qsys automatically updates your design to the current version.

**Avalon-MM Pipeline Bridge**

The Avalon-MM Pipeline Bridge inserts a register stage in the Avalon-MM command and response paths. It accepts commands on its Avalon-MM slave port and propagates them to its Avalon-MM master port. It provides separate parameters to turn on pipelining in the command and response networks.

You can also use the Avalon-MM bridge to export a single Avalon-MM slave interface that can be used to control multiple Avalon-MM slave devices, and you can optionally turn off the pipelining feature of this bridge. In this configuration, the bridge transfers commands received on its Avalon-MM slave interface to its Avalon-MM master port. Figure 12–4 illustrates its use in a XAUI PHY transceiver IP core.

Because the Avalon-MM slave interface is exported to the pins of the device, having a single Avalon-MM slave port (rather than separate ports for each Avalon-MM slave device) reduces the pin count of the FPGA.
Address Span Extender

The Address Span Extender component creates a windowed bridge and allows memory-mapped master interfaces to access a larger address map than the width of their address signals allow. When connected to an address span extender, an address span restricted master can access a broader address range. The extender splits up the larger addressable space into separate windows so that the master with a smaller address span can access the appropriate part of the memory.

For example, if the fast variant of a processor can address only 2GB of address span, and you need that processor to access a broader span, then you can use the address span extender to access the broader span by providing a window with a smaller address span. The same issue occurs with SoC devices. For example, an HPS subsystem in SoC devices can address only 1GB of address span within the FPGA. You can use the address span extender in this case, as well.

When you implement the address span extender in Qsys for a master with limited addressing space, you must first decide how large of an address space you want a particular slave to occupy in a master’s address map.

This component allows you to define between 1 and 64 address windows, and accordingly, a given number of registers to hold the upper address bits for each window. In the component GUI, you must select the number of bits you want to access (Expanded Master Byte Address Width), the number of bits you want the master to see (Slave Word Address Width), and the number of sub-windows.

The upper bits of the slave address are used to pick which window is used. For example, if you specify 4 windows, then the top 2 bits of the slave address are used to specify window[0, 1, 2, 3]. Therefore having more windows does require the windows to be smaller, for example having 4 windows requires the windows themselves to be $1/4$ the size of the slave address space. The total windowed address space is still equal to the original slave address space, but the windows allow access to memory regions in a larger overall address space.

You can parametrize the address span extender with an initial fixed address value by entering an address for the Reset Default for Master Window option, and selecting True for the Disable Slave Control Port option, which allows the address span extender to function as a fixed, non-programmable component.

Each sub-window is equal in size and is stacked sequentially in the windowed slave interface’s address space. To control the fixed address bits of a particular sub-window, you can write to the sub-window’s register in the register control slave interface. Qsys structures the logic so that the Quartus II software can optimize away all unneeded bits.

Burst Properties

For read commands, when the Address Span Extender must propagate byteenables to prevent read side-effects, the Burstcount Width option should be set to 1. If Burstcount Width is set greater than 1, the read burst command is expressed in a single cycle, and assumes all byteenables are asserted on every cycle.
Tri-state Components

The tri-state interface type allows you to design Qsys subsystems that connect to tri-state devices on your PCB. The following three components implement the tri-state conduit functionality:

- **Generic Tri-state Controller**
- **Tri-state Conduit Pin Sharer**
- **Tri-state Conduit Bridge**

You can use these components to implement pin sharing, convert between unidirectional and bidirectional signals, and create tri-state controllers for devices whose interfaces can be described using the tri-state signal types.

For more information about the tri-state signal types, refer to the *Avalon Tri-state Conduit Interfaces* chapter in the *Avalon Interface Specifications*, and the *Avalon Tri-State Conduit Components Use Guide*.

Figure 12–5 illustrates the typical use of tri-state components, and includes two Generic Tri-state Conduit Controllers. The first is customized to control a flash memory. The second is customized to control an off-chip SSRAM. The Tri-state Conduit Pin Sharer multiplexes between these two controllers, and the Tri-state Conduit Bridge converts between an on-chip encoding of tri-state signals and true bidirectional signals.

Figure 12–5. Tri-state Conduit System to Control Off-Chip SRAM and Flash Devices
By default, the Tri-state Conduit Pin Sharer and Tri-State Conduit Bridge present byte addresses. Each address location in many memory devices contains more than one byte of data. In the example in Figure 12–5, the flash device operates on 16-bit words and must ignore the least-significant bit of the Avalon-MM address. The SSRAM memory operates on 32-bit words and must ignore the two, low-order memory bits. Because neither device requires a byte address, addr[0] is not routed on the PCB. Figure 12–6 shows addr[0] as unconnected.

Figure 12–6. Address Connections from Qsys System to PCB

In this example design, the flash device responds to address range 0 MBytes to 8 MBytes-1. The SSRAM responds to address range 8 MBytes to 10 MBytes-1. The PCB schematic for the PCB connects addr[20:2] to addr[18:0] of the SSRAM device because the SSRAM responds to 32-bit word address. The 8 MByte flash device accesses 16-bit words; consequently, the schematic does not connect addr[0]. The chipselect signals select between the two devices.

If you create a custom tri-state conduit master with word-aligned addresses, the Tri-state Conduit Pin Sharer does nothing to change or align the address signals.
Figure 12–7 illustrates this example system in Qsys.

**Generic Tri-state Controller**

The Generic Tri-state Controller provides a template for a controller that you can parameterize to reflect the behavior of an off-chip device. The Generic Tri-state Controller has many parameters that you can use to customize this component such as the following examples:

- The width of the address and data signals
- The read and write wait times
- The bus-turnaround time

In calculating delays, the Generic Tri-state Controller chooses the larger of the bus-turnaround time and read latency. Turnaround time is measured from the time that a command is accepted, not from the time that the previous read returned data.

- The data hold time

The Generic Tri-state Controller always includes the following interfaces:
■ Memory-mapped slave interface—This interface connects to an memory-mapped master, such as a processor.
■ Tri-state master interface usually connects to the tri-state conduit slave interface of the tri-state conduit pin sharer.
■ Clock sink—The component’s clock reference. This interface must be connected to a clock source.
■ Resets sink—This interface connects to a reset source interface.

**Tri-state Conduit Pin Sharer**

The Tri-state Conduit Pin Sharer multiplexes between the signals of the connected tri-state controllers. You connect all signals from the tri-state controllers to the Tri-state Conduit Pin Sharer and use the parameter editor to specify the signals that are shared. The parameter editor includes a Shared Signal Name column for you to type the shared signal name as Figure 12–8 illustrates.

Figure 12–8. Specifying Shared Signals Using the Tri-state Conduit Pin Sharer

If the widths of shared signals differ, the signals are aligned on their 0th bit and the higher-order pins are driven to 0 whenever the smaller signal has control of the bus. Unshared signals always propagate through the pin sharer. The tri-state conduit pin sharer uses the round-robin arbiter to select between tri-state conduit controllers.

All tri-state conduit components connected to a given pin sharer must be in the same clock domain.
Tri-state Conduit Bridge

The Tri-state Conduit Bridge instantiates bidirectional signals for each tri-state signal while passing all other signals straight through the component. The Tri-state Conduit Bridge registers all outgoing and incoming signals, which adds two cycles of latency for a read request. You must account for this additional pipelining when designing a custom controller. During reset, all outputs are placed in a high-impedance state; outputs are enabled in the first clock cycle after reset is deasserted. The Quartus II software labels these output signals bidirectional.

Test Pattern Generator and Checker Cores

The data generation and monitoring solution for Avalon-ST consists of two components: a test pattern generator core that generates data, and sends it out on an Avalon-ST data interface, and a test pattern checker core that receives the same data and verifies it. Optionally, the data can be formatted as packets, with accompanying start_of_packet and end_of_packet signals.

The test pattern generator inserts different error conditions, and the test pattern checker reports these error conditions to the control interface, each via an Avalon Memory-Mapped (Avalon-MM) slave.

Resource Utilization and Performance

Resource utilization and performance for the test pattern generator and checker cores depend on the data width, number of channels, and whether the streaming data uses the optional packet protocol.
Table 12–1 provides estimated resource utilization and performance values for the test pattern generator core.

<table>
<thead>
<tr>
<th>No. of Channels</th>
<th>Datawidth (No. of 8-bit Symbols Per Beat)</th>
<th>Packet Support</th>
<th>Stratix® II and Stratix II GX</th>
<th>Cyclone® II</th>
<th>Stratix</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>$f_{max}$ (MHz)</td>
<td>ALM Count</td>
<td>Memory (bits)</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>Yes</td>
<td>284</td>
<td>233</td>
<td>560</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>No</td>
<td>293</td>
<td>222</td>
<td>496</td>
</tr>
<tr>
<td>32</td>
<td>4</td>
<td>Yes</td>
<td>276</td>
<td>270</td>
<td>912</td>
</tr>
<tr>
<td>32</td>
<td>4</td>
<td>No</td>
<td>323</td>
<td>227</td>
<td>848</td>
</tr>
<tr>
<td>1</td>
<td>16</td>
<td>Yes</td>
<td>298</td>
<td>361</td>
<td>560</td>
</tr>
<tr>
<td>1</td>
<td>16</td>
<td>No</td>
<td>340</td>
<td>330</td>
<td>496</td>
</tr>
<tr>
<td>32</td>
<td>16</td>
<td>Yes</td>
<td>295</td>
<td>410</td>
<td>912</td>
</tr>
<tr>
<td>32</td>
<td>16</td>
<td>No</td>
<td>269</td>
<td>409</td>
<td>848</td>
</tr>
</tbody>
</table>

Table 12–2 provides estimated resource utilization and performance values for the test pattern checker core.

<table>
<thead>
<tr>
<th>No. of Channels</th>
<th>Datawidth (No. of 8-bit Symbols Per Beat)</th>
<th>Packet Support</th>
<th>Stratix® II and Stratix II GX</th>
<th>Cyclone® II</th>
<th>Stratix</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>$f_{max}$ (MHz)</td>
<td>ALM Count</td>
<td>Memory (bits)</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>Yes</td>
<td>270</td>
<td>271</td>
<td>96</td>
</tr>
<tr>
<td>1</td>
<td>4</td>
<td>No</td>
<td>371</td>
<td>187</td>
<td>32</td>
</tr>
<tr>
<td>32</td>
<td>4</td>
<td>Yes</td>
<td>185</td>
<td>396</td>
<td>3616</td>
</tr>
<tr>
<td>32</td>
<td>4</td>
<td>No</td>
<td>221</td>
<td>363</td>
<td>3520</td>
</tr>
<tr>
<td>1</td>
<td>16</td>
<td>Yes</td>
<td>253</td>
<td>462</td>
<td>96</td>
</tr>
<tr>
<td>1</td>
<td>16</td>
<td>No</td>
<td>277</td>
<td>306</td>
<td>32</td>
</tr>
<tr>
<td>32</td>
<td>16</td>
<td>Yes</td>
<td>182</td>
<td>582</td>
<td>3616</td>
</tr>
<tr>
<td>32</td>
<td>16</td>
<td>No</td>
<td>218</td>
<td>473</td>
<td>3520</td>
</tr>
</tbody>
</table>
Test Pattern Generator

The test pattern generator core accepts commands to generate data via an Avalon-MM command interface, and drives the generated data to an Avalon-ST data interface. You can parameterize most aspects of the Avalon-ST data interface, such as the number of error bits and data signal width, thus allowing you to test components with different interfaces.

Figure 12–9 shows a diagram of the test pattern generator.

![Test Pattern Generator Core Block Diagram](image)

The data pattern is calculated as: Symbol Value = Symbol Position in Packet XOR Data Error Mask. Data that is not organized in packets is a single stream with no beginning or end.

The test pattern generator has a throttle register that is set via the Avalon-MM control interface. The value of the throttle register is used in conjunction with a pseudo-random number generator to throttle the data generation rate.

Command Interface

The command interface is a 32-bit Avalon-MM write slave that accepts data generation commands. It is connected to a 16-element deep FIFO, thus allowing a master peripheral to drive a number of commands into the test pattern generator.

The command interface maps to the following registers: cmd_lo and cmd_hi. The command is pushed into the FIFO when the register cmd_lo (address 0) is addressed. When the FIFO is full, the command interface asserts the waitrequest signal. You can create errors by writing to the register cmd_hi (address 1). The errors are cleared when 0 is written to this register, or its respective fields. See page “Test Pattern Generator Command Registers” on page 12–17 for more information about the register fields.

Control and Status Interface

The control and status interface is a 32-bit Avalon-MM slave that allows you to enable or disable the data generation, as well as set the throttle. This interface also provides generation-time information, such as the number of channels and whether or not data packets are supported.
Output Interface

The output interface is an Avalon-ST interface that optionally supports data packets. You can configure the output interface to align with your system requirements. Depending on the incoming stream of commands, the output data may contain interleaved packet fragments for different channels. To keep track of the current symbol’s position within each packet, the test pattern generator maintains an internal state for each channel.

Functional Parameter

The functional parameter allows you to configure the test pattern generator as a whole system. The Throttle Seed is the starting value for the throttle control random number generator. Altera recommends a value that is unique to each instance of the test pattern generator and checker cores in a system.

Output Interface

You can configure the output interface of the test pattern generator with the following parameters:

- **Number of Channels**—The number of channels that the test pattern generator supports. Valid values are 1 to 256.
- **Data Bits Per Symbol**—The bits per symbol is related to the width of readdata and writedata signals, which must be a multiple of the bits per symbol.
  
  If you change only bits per symbol, and do not change the data width, you will get errors.

- **Data Symbols Per Beat**—The number of symbols (words) that are transferred per beat. Valid values are 1 to 256.
- **Include Packet Support**—Indicates whether or not packet transfers are supported. Packet support includes the startofpacket, endofpacket, and empty signals.
- **Error Signal Width (bits)**—The width of the error signal on the output interface. Valid values are 0 to 31. A value of 0 indicates that the error signal is not used.

Test Pattern Checker

The test pattern checker core accepts data via an Avalon-ST interface, verifies it against the same predetermined pattern used by the test pattern generator to produce the data, and reports any exceptions to the control interface. You can parameterize most aspects of the test pattern checker’s Avalon-ST interface such as the number of error bits and the data signal width, thus allowing you to test components with different interfaces.

The test pattern checker has a throttle register that is set via the Avalon-MM control interface. The value of the throttle register controls the rate at which data is accepted.
Figure 12–10 shows a diagram of the test pattern checker.

The test pattern checker detects exceptions and reports them to the control interface via a 32-element deep internal FIFO. Possible exceptions are data error, missing start-of-packet (SOP), missing end-of-packet (EOP), and signalled error.

As each exception occurs, an exception descriptor is pushed into the FIFO. If the same exception occurs more than once consecutively, only one exception descriptor is pushed into the FIFO. All exceptions are ignored when the FIFO is full. Exception descriptors are deleted from the FIFO after they are read by the control and status interface.

**Input Interface**

The input interface is an Avalon-ST interface that optionally supports data packets. You can configure the input interface to align with your system requirements.

Incoming data may contain interleaved packet fragments. To keep track of the current symbol’s position, the test pattern checker maintains an internal state for each channel.

**Control and Status Interface**

The control and status interface is a 32-bit Avalon-MM slave that allows you to enable or disable data acceptance, as well as set the throttle. This interface provides generation-time information, such as the number of channels and whether the test pattern checker supports data packets. The control and status interface also provides information on the exceptions detected by the test pattern checker. The interface obtains this information by reading from the exception FIFO.

**Functional Parameter**

The functional parameter allows you to configure the test pattern checker as a whole system. The **Throttle Seed** is the starting value for the throttle control random number generator. Altera recommends a unique value for each instance of the test pattern generator and checker cores in a system.

**Input Parameters**

You can configure the input interface of the test pattern checker using the following parameters:
- **Data Bits Per Symbol**—The bits per symbol is related to the width of `readdata` and `writedata` signals, which must be a multiple of the bits per symbol.
  
  If you change only bits per symbol, and do not change the data width, you will get errors.

- **Data Symbols Per Beat**—The number of symbols (words) that are transferred per beat. Valid values are 1 to 32.

- **Include Packet Support**—Indicates whether or not data packet transfers are supported. Packet support includes the `startofpacket`, `endofpacket`, and `empty` signals.

- **Number of Channels**—The number of channels that the test pattern checker supports. Valid values are 1 to 256.

- **Error Signal Width (bits)**—The width of the `error` signal on the input interface. Valid values are 0 to 31. A value of 0 indicates that the `error` signal is not used.

### Hardware Simulation Considerations

The test pattern generator and checker cores do not provide a simulation testbench for simulating a stand-alone instance of the component. However, you can use the standard SOPC Builder simulation flow to simulate the component design files inside an SOPC Builder system.

### Software Programming Model

This section describes the software programming model for the test pattern generator and checker cores.

#### HAL System Library Support

For Nios II processor users, Altera provides HAL system library drivers that allow you to initialize and access the test pattern generator and checker cores. Altera recommends you to use the provided drivers to access the cores instead of accessing the registers directly.

For Nios II IDE users, copy the provided drivers from the following installation folders to your software application directory:

- `<IP installation directory>/ip/sopc_builder_ip/altera_avalon_data_source/HAL`
- `<IP installation directory>/ip/sopc_builder_ip/altera_avalon_data_sink/HAL`

This instruction does not apply if you use the Nios II command-line tools.

### Software Files

The following software files define the low-level access to the hardware, and provide the routines for the HAL device drivers.

Do not modify the software files.
Software files provided with the test pattern generator core:
- data_source_regs.h—The header file that defines the test pattern generator’s register maps.
- data_source_util.h, data_source_util.c—The header and source code for the functions and variables required to integrate the driver into the HAL system library.

Software files provided with the test pattern checker core:
- data_sink_regs.h—The header file that defines the core’s register maps.
- data_sink_util.h, data_sink_util.c—The header and source code for the functions and variables required to integrate the driver into the HAL system library.

Register Maps
This section describes the register maps for the test pattern generator and checker cores.

Test Pattern Generator Control and Status Registers
Table 12–3 shows the offset for the test pattern generator control and status registers. Each register is 32 bits wide.

<table>
<thead>
<tr>
<th>Offset</th>
<th>Register Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>base + 0</td>
<td>status</td>
</tr>
<tr>
<td>base + 1</td>
<td>control</td>
</tr>
<tr>
<td>base + 2</td>
<td>fill</td>
</tr>
</tbody>
</table>

Table 12–4 describes the status register bits.

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[15:0]</td>
<td>ID</td>
<td>RO</td>
<td>A constant value of 0x64.</td>
</tr>
<tr>
<td>[23:16]</td>
<td>NUMCHANNELS</td>
<td>RO</td>
<td>The configured number of channels.</td>
</tr>
<tr>
<td>[30:24]</td>
<td>NUMSYMBOLS</td>
<td>RO</td>
<td>The configured number of symbols per beat.</td>
</tr>
<tr>
<td>[31]</td>
<td>SUPPORTPACKETS</td>
<td>RO</td>
<td>A value of 1 indicates data packet support.</td>
</tr>
</tbody>
</table>

Table 12–5 describes the control register bits.

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0]</td>
<td>ENABLE</td>
<td>RW</td>
<td>Setting this bit to 1 enables the test pattern generator core.</td>
</tr>
<tr>
<td>[7:1]</td>
<td></td>
<td></td>
<td>Reserved</td>
</tr>
</tbody>
</table>
Table 12–5. Control Field Descriptions (Part 2 of 2)

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[16:8]</td>
<td>THROTTLE</td>
<td>RW</td>
<td>Specifies the throttle value which can be between 0–256, inclusively. This value is used in conjunction with a pseudorandom number generator to throttle the data generation rate. Setting THROTTLE to 0 stops the test pattern generator core. Setting it to 256 causes the test pattern generator core to run at full throttle. Values between 0–256 result in a data rate proportional to the throttle value.</td>
</tr>
<tr>
<td>[17]</td>
<td>SOFT RESET</td>
<td>RW</td>
<td>When this bit is set to 1, all internal counters and statistics are reset. Write 0 to this bit to exit reset.</td>
</tr>
<tr>
<td>[31:18]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 12–6 describes the fill register bits.

Table 12–6. Fill Field Descriptions

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0]</td>
<td>BUSY</td>
<td>RO</td>
<td>A value of 1 indicates that data transmission is in progress, or that there is at least one command in the command queue.</td>
</tr>
<tr>
<td>[6:1]</td>
<td></td>
<td></td>
<td>Reserved</td>
</tr>
<tr>
<td>[15:7]</td>
<td>FILL</td>
<td>RO</td>
<td>The number of commands currently in the command FIFO.</td>
</tr>
<tr>
<td>[31:16]</td>
<td></td>
<td></td>
<td>Reserved</td>
</tr>
</tbody>
</table>

Test Pattern Generator Command Registers

Table 12–7 shows the offset for the command registers. Each register is 32 bits wide.

Table 12–7. Test Pattern Command Register Map

<table>
<thead>
<tr>
<th>Offset</th>
<th>Register Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>base + 0</td>
<td>cmd_lo</td>
</tr>
<tr>
<td>base + 1</td>
<td>cmd_hi</td>
</tr>
</tbody>
</table>

Table 12–8 describes the cmd_lo register bits. The command is pushed into the FIFO only when the cmd_lo register is addressed.

Table 12–8. cmd_lo Field Descriptions

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[15:0]</td>
<td>SIZE</td>
<td>RW</td>
<td>The segment size in symbols. Except for the last segment in a packet, the size of all segments must be a multiple of the configured number of symbols per beat. If this condition is not met, the test pattern generator core inserts additional symbols to the segment to ensure the condition is fulfilled.</td>
</tr>
<tr>
<td>[29:16]</td>
<td>CHANNEL</td>
<td>RW</td>
<td>The channel to send the segment on. If the channel signal is less than 14 bits wide, the low order bits of this register are used to drive the signal.</td>
</tr>
<tr>
<td>[30]</td>
<td>SOP</td>
<td>RW</td>
<td>Set this bit to 1 when sending the first segment in a packet. This bit is ignored when data packets are not supported.</td>
</tr>
<tr>
<td>[31]</td>
<td>EOP</td>
<td>RW</td>
<td>Set this bit to 1 when sending the last segment in a packet. This bit is ignored when data packets are not supported.</td>
</tr>
</tbody>
</table>
Table 12–9 describes the cmd_hi register bits.

Table 12–9. cmd_hi Field Descriptions

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[15:0]</td>
<td>SIGNALLED ERROR</td>
<td>RW</td>
<td>Specifies the value to drive the error signal. A non-zero value creates a signalled error.</td>
</tr>
<tr>
<td>[23:16]</td>
<td>DATA ERROR</td>
<td>RW</td>
<td>The output data is XORed with the contents of this register to create data errors. To stop creating data errors, set this register to 0.</td>
</tr>
<tr>
<td>[24]</td>
<td>SUPRESS SOP</td>
<td>RW</td>
<td>Set this bit to 1 to suppress the assertion of the startofpacket signal when the first segment in a packet is sent.</td>
</tr>
<tr>
<td>[25]</td>
<td>SUPRESS EOP</td>
<td>RW</td>
<td>Set this bit to 1 to suppress the assertion of the endofpacket signal when the last segment in a packet is sent.</td>
</tr>
</tbody>
</table>

Test Pattern Checker Control and Status Registers

Table 12–10 shows the offset for the control and status registers. Each register is 32 bits wide.

Table 12–10. Test Pattern Checker Control and Status Register Map

<table>
<thead>
<tr>
<th>Offset</th>
<th>Register Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>base + 0</td>
<td>status</td>
</tr>
<tr>
<td>base + 1</td>
<td>control</td>
</tr>
<tr>
<td>base + 2</td>
<td></td>
</tr>
<tr>
<td>base + 3</td>
<td>Reserved</td>
</tr>
<tr>
<td>base + 4</td>
<td></td>
</tr>
<tr>
<td>base + 5</td>
<td>exception_descriptor</td>
</tr>
<tr>
<td>base + 6</td>
<td>indirect_select</td>
</tr>
<tr>
<td>base + 7</td>
<td>indirect_count</td>
</tr>
</tbody>
</table>

Table 12–11 describes the status register bits.

Table 12–11. Status Field Descriptions

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[15:0]</td>
<td>ID</td>
<td>RO</td>
<td>Contains a constant value of 0x65.</td>
</tr>
<tr>
<td>[23:16]</td>
<td>NUMCHANNELS</td>
<td>RO</td>
<td>The configured number of channels.</td>
</tr>
<tr>
<td>[30:24]</td>
<td>NUMSYMBOLS</td>
<td>RO</td>
<td>The configured number of symbols per beat.</td>
</tr>
<tr>
<td>[31]</td>
<td>SUPPORTPACKETS</td>
<td>RO</td>
<td>A value of 1 indicates packet support.</td>
</tr>
</tbody>
</table>

Table 12–12 describes the control register bits.

Table 12–12. Control Field Descriptions (Part 1 of 2)

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0]</td>
<td>ENABLE</td>
<td>RW</td>
<td>Setting this bit to 1 enables the test pattern checker.</td>
</tr>
<tr>
<td>[7:1]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### Table 12–12. Control Field Descriptions (Part 2 of 2)

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[16:8]</td>
<td>THROTTLE</td>
<td>RW</td>
<td>Specifies the throttle value which can be between 0–256, inclusively. This</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>value is used in conjunction with a pseudorandom number generator to</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>throttle the data generation rate. Setting THROTTLE to 0 stops the test</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>pattern generator core. Setting it to 256 causes the test pattern generator</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>core to run at full throttle. Values between 0–256 result in a data rate</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>proportional to the throttle value.</td>
</tr>
<tr>
<td>[17]</td>
<td>SOFT_RESET</td>
<td>RW</td>
<td>When this bit is set to 1, all internal counters and statistics are reset.</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Write 0 to this bit to exit reset.</td>
</tr>
<tr>
<td>[31:18]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 12–14 describes the exception_descriptor register bits. If there is no exception, reading this register returns 0.

### Table 12–13. exception_descriptor Field Descriptions

<table>
<thead>
<tr>
<th>Bit(s)</th>
<th>Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0]</td>
<td>DATA ERROR</td>
<td>RO</td>
<td>A value of 1 indicates that an error is detected in the incoming data.</td>
</tr>
<tr>
<td>[7:3]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>[15:8]</td>
<td>SIGNALLED</td>
<td>RO</td>
<td>The value of the error signal.</td>
</tr>
<tr>
<td>[23:16]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>[31:24]</td>
<td>CHANNEL</td>
<td>RO</td>
<td>The channel on which the exception was detected.</td>
</tr>
</tbody>
</table>

Table 12–15 describes the indirect_select register bits.

### Table 12–15. indirect_select Field Descriptions

<table>
<thead>
<tr>
<th>Bit</th>
<th>Bits Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[7:0]</td>
<td>INDIRECT CHANNEL</td>
<td>RW</td>
<td>Specifies the channel number that applies to the INDIRECT PACKET COUNT,</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INDIRECT SYMBOL COUNT, and INDIRECT ERROR COUNT registers.</td>
</tr>
<tr>
<td>[15:8]</td>
<td>Reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>[31:16]</td>
<td>INDIRECT ERROR</td>
<td>RO</td>
<td>The number of data errors that occurred on the channel specified by</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INDIRECT CHANNEL.</td>
</tr>
</tbody>
</table>

Table 12–16 describes the indirect_count register bits.

### Table 12–16. indirect_count Field Descriptions

<table>
<thead>
<tr>
<th>Bit</th>
<th>Bits Name</th>
<th>Access</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[15:0]</td>
<td>INDIRECT PACKET COUNT</td>
<td>RO</td>
<td>The number of data packets received on the channel specified by INDIRECT CHANNEL.</td>
</tr>
<tr>
<td>[31:16]</td>
<td>INDIRECT SYMBOL COUNT</td>
<td>RO</td>
<td>The number of symbols received on the channel specified by INDIRECT CHANNEL.</td>
</tr>
</tbody>
</table>
Test Pattern Generator API

This section describes the application programming interface (API) for the test pattern generator.

API functions are currently not available from the interrupt service routine (ISR).

**data_source_reset()**

Prototype: 
```c
void data_source_reset(alt_u32 base);
```

Thread-safe: No.

Include: `<data_source_util.h>`

Parameters: base—The base address of the control and status slave.

Returns: void.

Description: This function resets the test pattern generator core including all internal counters and FIFOs. The control and status registers are not reset by this function.

**data_source_init()**

Prototype: 
```c
int data_source_init(alt_u32 base, alt_u32 command_base);
```

Thread-safe: No.

Include: `<data_source_util.h>`

Parameters: base—The base address of the control and status slave.
command_base—The base address of the command slave.

Returns: 
1—Initialization is successful.
0—Initialization is unsuccessful.

Description: This function performs the following operations to initialize the test pattern generator core:
- Resets and disables the test pattern generator core.
- Sets the maximum throttle.
- Clears all inserted errors.

**data_source_get_id()**

Prototype: 
```c
int data_source_get_id(alt_u32 base);
```

Thread-safe: Yes.

Include: `<data_source_util.h>`

Parameters: base—The base address of the control and status slave.

Returns: The test pattern generator core’s identifier.

Description: This function retrieves the test pattern generator core’s identifier.

**data_source_get_supports_packets()**

Prototype: 
```c
int data_source_get_supports_packets(alt_u32 base);
```

Thread-safe: Yes.

Include: `<data_source_util.h>`
**Parameters:**

base—The base address of the control and status slave.

**Returns:**

1—Data packets are supported.

0—Data packets are not supported.

**Description:**

This function checks if the test pattern generator core supports data packets.

### data_source_get_num_channels()

**Prototype:**

```c
int data_source_get_num_channels(alt_u32 base);
```

**Thread-safe:**

Yes.

**Include:**

`<data_source_util.h>`

**Parameters:**

base—The base address of the control and status slave.

**Returns:**

The number of channels supported.

**Description:**

This function retrieves the number of channels supported by the test pattern generator core.

### data_source_get_symbols_per_cycle()

**Prototype:**

```c
int data_source_get_symbols(alt_u32 base);
```

**Thread-safe:**

Yes.

**Include:**

`<data_source_util.h>`

**Parameters:**

base—The base address of the control and status slave.

**Returns:**

The number of symbols transferred in a beat.

**Description:**

This function retrieves the number of symbols transferred by the test pattern generator core in each beat.

### data_source_set_enable()

**Prototype:**

```c
void data_source_set_enable(alt_u32 base, alt_u32 value);
```

**Thread-safe:**

No.

**Include:**

`<data_source_util.h>`

**Parameters:**

base—The base address of the control and status slave.

value—The ENABLE bit is set to the value of this parameter.

**Returns:**

void.

**Description:**

This function enables or disables the test pattern generator core. When disabled, the test pattern generator core stops data transmission but continues to accept commands and stores them in the FIFO.

### data_source_get_enable()

**Prototype:**

```c
int data_source_get_enable(alt_u32 base);
```

**Thread-safe:**

Yes.

**Include:**

`<data_source_util.h>`

**Parameters:**

base—The base address of the control and status slave.

**Returns:**

The value of the ENABLE bit.

**Description:**

This function retrieves the value of the ENABLE bit.
**data_source_set_throttle()**

Prototype: 

```c
void data_source_set_throttle(alt_u32 base, alt_u32 value);
```

Thread-safe:  

No.

Include:  

```c
<data_source_util.h>
```

Parameters:  

- `base`—The base address of the control and status slave.
- `value`—The throttle value.

Returns:  

`void`.

Description:  

This function sets the throttle value, which can be between 0–256 inclusively. The throttle value, when divided by 256 yields the rate at which the test pattern generator sends data.

**data_source_get_throttle()**

Prototype:  

```c
int data_source_get_throttle(alt_u32 base);
```

Thread-safe:  

Yes.

Include:  

```c
<data_source_util.h>
```

Parameters:  

- `base`—The base address of the control and status slave.

Returns:  

The throttle value.

Description:  

This function retrieves the current throttle value.

**data_source_is_busy()**

Prototype:  

```c
int data_source_is_busy(alt_u32 base);
```

Thread-safe:  

Yes.

Include:  

```c
<data_source_util.h>
```

Parameters:  

- `base`—The base address of the control and status slave.

Returns:  

- 1—The test pattern generator core is busy.
- 0—The core is not busy.

Description:  

This function checks if the test pattern generator is busy. The test pattern generator core is busy when it is sending data or has data in the command FIFO to be sent.

**data_source_fill_level()**

Prototype:  

```c
int data_source_fill_level(alt_u32 base);
```

Thread-safe:  

Yes.

Include:  

```c
<data_source_util.h>
```

Parameters:  

- `base`—The base address of the control and status slave.

Returns:  

The number of commands in the command FIFO.

Description:  

This function retrieves the number of commands currently in the command FIFO.
data_source_send_data()

Prototype:

int data_source_send_data(alt_u32 cmd_base, alt_u16 channel, alt_u16 size,
alt_u32 flags, alt_u16 error, alt_u8 data_error_mask);

Thread-safe: No.
Include: <data_source_util.h>
Parameters:
    cmd_base—The base address of the command slave.
    channel—The channel to send the data on.
    size—The data size.
    flags—Specifies whether to send or suppress SOP and EOP signals. Valid values are
DATA_SOURCE_SEND_SOP, DATA_SOURCE_SEND_EOP, DATA_SOURCE_SEND_SUPPRESS_SOP and
DATA_SOURCE_SEND_SUPPRESS_EOP.
    error—The value asserted on the error signal on the output interface.
    data_error_mask—This parameter and the data are XORed together to produce erroneous data.

Returns: Always returns 1.
Description:
This function sends a data fragment to the specified channel.
If data packets are supported, user applications must ensure the following conditions are met:
SOP and EOP are used consistently in each channel.
Except for the last segment in a packet, the length of each segment is a multiple of the data width.
If data packets are not supported, user applications must ensure the following conditions are met:
No SOP and EOP indicators in the data.
The length of each segment in a packet is a multiple of the data width.

Test Pattern Checker API

This section describes the API for the test pattern checker core. The API functions are
currently not available from the ISR.

data_sink_reset()

Prototype:

void data_sink_reset(alt_u32 base);

Thread-safe: No.
Include: <data_sink_util.h>
Parameters:
    base—The base address of the control and status slave.

Returns: void.
Description:
This function resets the test pattern checker core including all internal counters.
data_sink_init()

Prototype: int data_source_init(alt_u32 base);
Thread-safe: No.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: 1—Initialization is successful.
0—Initialization is unsuccessful.
Description: This function performs the following operations to initialize the test pattern checker core:
\n n Resets and disables the test pattern checker core.
 n Sets the throttle to the maximum value.

data_sink_get_id()

Prototype: int data_sink_get_id(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: The test pattern checker core’s identifier.
Description: This function retrieves the test pattern checker core’s identifier.

data_sink_get_supports_packets()

Prototype: int data_sink_init(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: 1—Data packets are supported.
0—Data packets are not supported.
Description: This function checks if the test pattern checker core supports data packets.

data_sink_get_num_channels()

Prototype: int data_sink_get_num_channels(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: The number of channels supported.
Description: This function retrieves the number of channels supported by the test pattern checker core.
data_sink_get_symbols_per_cycle()

Prototype: int data_sink_get_symbols(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: The number of symbols received in a beat.
Description: This function retrieves the number of symbols received by the test pattern checker core in each beat.

data_sink_set_enable()

Prototype: void data_sink_set_enable(alt_u32 base, alt_u32 value);
Thread-safe: No.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
value—The ENABLE bit is set to the value of this parameter.
Returns: void.
Description: This function enables the test pattern checker core.

data_sink_get_enable()

Prototype: int data_sink_get_enable(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: The value of the ENABLE bit.
Description: This function retrieves the value of the ENABLE bit.

data_sink_set_throttle()

Prototype: void data_sink_set_throttle(alt_u32 base, alt_u32 value);
Thread-safe: No.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
value—The throttle value.
Returns: void.
Description: This function sets the throttle value, which can be between 0–256 inclusively. The throttle value, when divided by 256 yields the rate at which the test pattern checker receives data.
**data_sink_get_throttle()**

Prototype:  
\[
\text{int data_sink_get_throttle(alt_u32 base);} 
\]

Thread-safe: Yes.

Include: `<data_sink_util.h>`

Parameters:  
base—The base address of the control and status slave.

Returns: The throttle value.

Description: This function retrieves the throttle value.

**data_sink_get_packet_count()**

Prototype:  
\[
\text{int data_sink_get_packet_count(alt_u32 base, alt_u32 channel);} 
\]

Thread-safe: No.

Include: `<data_sink_util.h>`

Parameters:  
base—The base address of the control and status slave.
channel—Channel number.

Returns: The number of data packets received on the given channel.

Description: This function retrieves the number of data packets received on a given channel.

**data_sink_get_symbol_count()**

Prototype:  
\[
\text{int data_sink_get_symbol_count(alt_u32 base, alt_u32 channel);} 
\]

Thread-safe: No.

Include: `<data_sink_util.h>`

Parameters:  
base—The base address of the control and status slave.
channel—Channel number.

Returns: The number of symbols received on the given channel.

Description: This function retrieves the number of symbols received on a given channel.

**data_sink_get_error_count()**

Prototype:  
\[
\text{int data_sink_get_error_count(alt_u32 base, alt_u32 channel);} 
\]

Thread-safe: No.

Include: `<data_sink_util.h>`

Parameters:  
base—The base address of the control and status slave.
channel—Channel number.

Returns: The number of errors received on the given channel.

Description: This function retrieves the number of errors received on a given channel.
data_sink_get_exception()

Prototype: int data_sink_get_exception(alt_u32 base);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: base—The base address of the control and status slave.
Returns: The first exception descriptor in the exception FIFO.
0—No exception descriptor found in the exception FIFO.
Description: This function retrieves the first exception descriptor in the exception FIFO and pops it off the FIFO.

data_sink_exception_is_exception()

Prototype: int data_sink_exception_is_exception(int exception);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: exception—Exception descriptor
Returns: 1—Indicates an exception.
0—No exception.
Description: This function checks if a given exception descriptor describes a valid exception.

data_sink_exception_has_data_error()

Prototype: int data_sink_exception_has_data_error(int exception);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: exception—Exception descriptor.
Returns: 1—Data has errors.
0—No errors.
Description: This function checks if a given exception indicates erroneous data.

data_sink_exception_has_missing_sop()

Prototype: int data_sink_exception_has_missing_sop(int exception);
Thread-safe: Yes.
Include: <data_sink_util.h>
Parameters: exception—Exception descriptor.
Returns: 1—Missing SOP.
0—Other exception types.
Description: This function checks if a given exception descriptor indicates missing SOP.
data_sink_exception_has_missing_eop()

Prototype: ```c
int data_sink_exception_has_missing_eop(int exception);
```
Thread-safe: Yes.
Include: ```c
<data_sink_util.h>
```
Parameters: exception—Exception descriptor.
Returns: 1—Missing EOP.
0—Other exception types.
Description: This function checks if a given exception descriptor indicates missing EOP.

data_sink_exception_signalled_error()

Prototype: ```c
int data_sink_exception_signalled_error(int exception);
```
Thread-safe: Yes.
Include: ```c
<data_sink_util.h>
```
Parameters: exception—Exception descriptor.
Returns: The signalled error value.
Description: This function retrieves the value of the signalled error from the exception.

data_sink_exception_channel()

Prototype: ```c
int data_sink_exception_channel(int exception);
```
Thread-safe: Yes.
Include: ```c
<data_sink_util.h>
```
Parameters: exception—Exception descriptor.
Returns: The channel number on which the given exception occurred.
Description: This function retrieves the channel number on which a given exception occurred.

Splitter Core

The Avalon-ST Splitter Core allows you to replicate transactions from an Avalon-ST source interface to multiple Avalon-ST sink interfaces. This core support from 1 to 16 outputs.
Figure 12–11 shows a block diagram of the Avalon-ST Splitter Core.

**Backpressure**

The Avalon-ST Splitter core integrates with backpressure by AND-ing the `ready` signals from the output interfaces and sending the result to the input interface. As a result, if an output interface deasserts the `ready` signal, the input interface receives the deasserted `ready` signal, as well. This functionality ensures that backpressure on the output interfaces is propagated to the input interface.

When the **Qualify Valid Out** parameter is set to 1, the `Out_Valid` signals on all other output interfaces are gated when backpressure is applied from one output interface. In this case, when any output interface deasserts its `ready` signal, the `Out_Valid` signals on the other output interfaces are deasserted, as well.

When the **Qualify Valid Out** parameter is set to 0, the output interfaces have a non-gated `Out_Valid` signal when backpressure is applied. In this case, when an output interface deasserts its `ready` signal, the `Out_Valid` signals on the other output interfaces are not affected.

Because the logic is combinational, the core introduces no latency.

**Interfaces**

The Avalon-ST Splitter core supports streaming data, with optional packet, channel, and error signals. The core propagates backpressure from any output interface to the input interface.
Table 12–17 shows the properties of the Avalon-ST interfaces.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Backpressure</td>
<td>Ready latency = 0.</td>
</tr>
<tr>
<td>Data Width</td>
<td>Configurable.</td>
</tr>
<tr>
<td>Channel</td>
<td>Supported (optional).</td>
</tr>
<tr>
<td>Error</td>
<td>Supported (optional).</td>
</tr>
<tr>
<td>Packet</td>
<td>Supported (optional).</td>
</tr>
</tbody>
</table>

### Parameters

Table 12–18 describes the parameters that you can configure for the Splitter core.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Legal Values</th>
<th>Default Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number Of Outputs</td>
<td>1 to 16</td>
<td>2</td>
<td>The number of output interfaces. The value of 1 is supported for some cases of parameterized systems in which no duplicated output is required.</td>
</tr>
<tr>
<td>Qualify Valid Out</td>
<td>0 or 1</td>
<td>1</td>
<td>Determines whether the Out_Valid signal is gated or non-gated when backpressure is applied.</td>
</tr>
<tr>
<td>Data Width</td>
<td>1–512</td>
<td>8</td>
<td>The width of the data on the Avalon-ST data interfaces.</td>
</tr>
<tr>
<td>Bits Per Symbol</td>
<td>1–512</td>
<td>8</td>
<td>The number of bits per symbol for the input and output interfaces. For example, byte-oriented interfaces have 8-bit symbols.</td>
</tr>
<tr>
<td>Use Packets</td>
<td>0 or 1</td>
<td>0</td>
<td>Indicates whether or not data packet transfers are supported. Packet support includes the startofpacket, endofpacket, and empty signals.</td>
</tr>
<tr>
<td>Use Channel</td>
<td>0 or 1</td>
<td>0</td>
<td>The option to enable or disable the channel signal.</td>
</tr>
<tr>
<td>Channel Width</td>
<td>0–8</td>
<td>1</td>
<td>The width of the channel signal on the data interfaces. This parameter is disabled when Use Channel is set to 0.</td>
</tr>
<tr>
<td>Max Channels</td>
<td>0–255</td>
<td>1</td>
<td>The maximum number of channels that a data interface can support. This parameter is disabled when Use Channel is set to 0.</td>
</tr>
<tr>
<td>Use Error</td>
<td>0 or 1</td>
<td>0</td>
<td>The option to enable or disable the error signal.</td>
</tr>
<tr>
<td>Error Width</td>
<td>0–31</td>
<td>1</td>
<td>The width of the error signal on the output interfaces. A value of 0 indicates that the error signal is not used. This parameter is disabled when Use Error is set to 0.</td>
</tr>
</tbody>
</table>

### Delay Core

The Avalon-ST Delay Core provides a solution to delay Avalon-ST transactions by a constant number of clock cycles. This core supports up to 16 clock cycle delays.
Figure 12–12 shows a diagram of the Avalon-ST Delay Core.

The Delay core adds a delay between the input and output interfaces. The core accepts transactions presented on the input interface and reproduces them on the output interface \( N \) cycles later without changing the transaction.

The input interface delays the input signals by a constant \( N \) number of clock cycles to the corresponding output signals of the output interface. The **Number Of Delay Clocks** parameter defines the constant \( N \), which must be between 0 and 16. The change of the **In_Valid** signal is reflected on the **Out_Valid** signal exactly \( N \) cycles later.

**Reset**

The Avalon-ST Delay core has a reset signal that is synchronous to the clk signal. When the core asserts the reset signal, the output signals are held at 0. After the reset signal is deasserted, the output signals are held at 0 for \( N \) clock cycles. The delayed values of the input signals are then reflected at the output signals after \( N \) clock cycles.

**Interfaces**

The Delay core supports streaming data, with optional packet, channel, and error signals. This core does not support backpressure.

Table 12–19 shows the Delay core properties.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Backpressure</td>
<td>Not supported.</td>
</tr>
<tr>
<td>Data Width</td>
<td>Configurable.</td>
</tr>
<tr>
<td>Channel</td>
<td>Supported (optional).</td>
</tr>
<tr>
<td>Error</td>
<td>Supported (optional).</td>
</tr>
<tr>
<td>Packet</td>
<td>Supported (optional).</td>
</tr>
</tbody>
</table>
Parameters

Table 12–20 describes the parameters that you can configure for the Delay core.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Legal Values</th>
<th>Default Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number Of Delay Clocks</td>
<td>0 to 16</td>
<td>1</td>
<td>Specifies the delay the core introduces, in clock cycles. The value of 0 is supported for some cases of parameterized systems in which no delay is required.</td>
</tr>
<tr>
<td>Data Width</td>
<td>1–512</td>
<td>8</td>
<td>The width of the data on the Avalon-ST data interfaces.</td>
</tr>
<tr>
<td>Bits Per Symbol</td>
<td>1–512</td>
<td>8</td>
<td>The number of bits per symbol for the input and output interfaces. For example, byte-oriented interfaces have 8-bit symbols.</td>
</tr>
<tr>
<td>Use Packets</td>
<td>0 or 1</td>
<td>0</td>
<td>Indicates whether or not data packet transfers are supported. Packet support includes the startofpacket, endofpacket, and empty signals.</td>
</tr>
<tr>
<td>Use Channel</td>
<td>0 or 1</td>
<td>0</td>
<td>The option to enable or disable the channel signal.</td>
</tr>
<tr>
<td>Channel Width</td>
<td>0–8</td>
<td>1</td>
<td>The width of the channel signal on the data interfaces. This parameter is disabled when Use Channel is set to 0.</td>
</tr>
<tr>
<td>Max Channels</td>
<td>0–255</td>
<td>1</td>
<td>The maximum number of channels that a data interface can support. This parameter is disabled when Use Channel is set to 0.</td>
</tr>
<tr>
<td>Use Error</td>
<td>0 or 1</td>
<td>0</td>
<td>The option to enable or disable the error signal.</td>
</tr>
<tr>
<td>Error Width</td>
<td>0–31</td>
<td>1</td>
<td>The width of the error signal on the output interfaces. A value of 0 indicates that the error signal is not in use. This parameter is disabled when Use Error is set to 0.</td>
</tr>
</tbody>
</table>

Round Robin Scheduler

The Avalon-ST Round Robin Scheduler core controls the read operations from a multi-channel Avalon-ST component that buffers data by channels. It reads the almost-full threshold values from the multiple channels in the multi-channel component and issues the read request to the Avalon-ST source according to a round-robin scheduling algorithm.

In a multi-channel component, the component can store data either in the sequence that it comes in (FIFO), or in segments according to the channel. When data is stored in segments according to channels, a scheduler is needed to schedule the read operations.
Figure 12–13 shows the block diagram of the Avalon-ST Round Robin Scheduler.

Figure 12–13. Avalon-ST Round Robin Scheduler Block Diagram

Performance and Resource Utilization

This section lists the resource utilization and performance data for various Altera device families. The estimates are obtained by compiling the core using the Quartus II software.

Table 12–21 shows the resource utilization and performance data for a Stratix® II GX device (EP2SGX130GF1508I4).

Table 12–21. Resource Utilization and Performance Data for Stratix II GX Devices

<table>
<thead>
<tr>
<th>Number of Channels</th>
<th>ALUTs</th>
<th>Logic Registers</th>
<th>Memory M512/M4K/M-RAM</th>
<th>f&lt;sub&gt;MAX&lt;/sub&gt; (MHz)</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>7</td>
<td>7</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
<tr>
<td>12</td>
<td>25</td>
<td>17</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
<tr>
<td>24</td>
<td>62</td>
<td>30</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
</tbody>
</table>

Table 12–22 shows the resource utilization and performance data for a Stratix III device (EP3SL340F1760C3). The performance of Stratix IV devices is similar to Stratix III devices.

Table 12–22. Resource Utilization and Performance Data for Stratix III Devices

<table>
<thead>
<tr>
<th>Number of Channels</th>
<th>ALUTs</th>
<th>Logic Registers</th>
<th>Memory M9K/M144K/MLAB</th>
<th>f&lt;sub&gt;MAX&lt;/sub&gt; (MHz)</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>7</td>
<td>7</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
<tr>
<td>12</td>
<td>25</td>
<td>17</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
<tr>
<td>24</td>
<td>67</td>
<td>30</td>
<td>0/0/0</td>
<td>&gt; 125</td>
</tr>
</tbody>
</table>

Table 12–23 shows the resource utilization and performance data for a Cyclone® III device (EP3C120F780I7).

Table 12–23. Resource Utilization and Performance Data for Cyclone III Devices (Part 1 of 2)

<table>
<thead>
<tr>
<th>Number of Channels</th>
<th>Total Logic Elements</th>
<th>Total Registers</th>
<th>Memory M9K</th>
<th>f&lt;sub&gt;MAX&lt;/sub&gt; (MHz)</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>12</td>
<td>7</td>
<td>0</td>
<td>&gt; 125</td>
</tr>
</tbody>
</table>
Interfaces

The following interfaces are available in the Avalon-ST Round Robin Scheduler core:

- Almost-Full Status Interface
- Request Interface

Almost-Full Status Interface

The Almost-Full Status interface is an Avalon-ST sink interface that collects the almost-full status from the sink components for the channels in the sequence provided.

Table 12–24 describes the features available for the Almost-Full Status interface.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Backpressure</td>
<td>Not supported</td>
</tr>
<tr>
<td>Data Width</td>
<td>Data width = 1; Bits per symbol = 1</td>
</tr>
<tr>
<td>Channel</td>
<td>Maximum channel = 32; Channel width = 5</td>
</tr>
<tr>
<td>Error</td>
<td>Not supported</td>
</tr>
<tr>
<td>Packet</td>
<td>Not supported</td>
</tr>
</tbody>
</table>

Request Interface

The Request Interface is an Avalon-MM write master interface that requests data from a specific channel. The Avalon-ST Round Robin Scheduler cycles through the channels it supports and schedules data to be read.

Operations

If a particular channel is almost full, the Round Robin Scheduler will not schedule data to be read from that channel in the source component. The scheduler only requests 1 beat of data from a channel at each transaction. To request 1 beat of data from channel \( n \), the scheduler writes the value 1 to address \( (4 \times n) \). For example, if the scheduler is requesting data from channel 3, the scheduler writes 1 to address 0xC. At every clock cycle, the scheduler requests data from the next channel. Therefore, if the scheduler starts requesting from channel 1, at the next clock cycle, it requests from channel 2. The scheduler does not request data from a particular channel if the almost-full status for the channel is asserted. In this case, one clock cycle is used without a request transaction.

The Avalon-ST Round Robin Scheduler cannot determine if the requested component is able to service the request transaction. The component asserts `waitrequest` when it cannot accept new requests.
Table 12–25 shows the list of ports for the Round Robin Scheduler.

### Table 12–25. Ports for the Avalon-ST Round Robin Scheduler

<table>
<thead>
<tr>
<th>Signal</th>
<th>Direction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Clock and Reset</strong></td>
<td></td>
<td></td>
</tr>
<tr>
<td>clk</td>
<td>In</td>
<td>Clock reference.</td>
</tr>
<tr>
<td>reset_n</td>
<td>In</td>
<td>Asynchronous active low reset.</td>
</tr>
<tr>
<td><strong>Avalon-MM Request Interface</strong></td>
<td></td>
<td></td>
</tr>
<tr>
<td>request_address (log&lt;sub&gt;2&lt;/sub&gt;Max_Channels–1:0)</td>
<td>Out</td>
<td>The write address used to indicate which channel has the request.</td>
</tr>
<tr>
<td>request_write</td>
<td>Out</td>
<td>Write enable signal.</td>
</tr>
<tr>
<td>request_writedata</td>
<td>Out</td>
<td>The amount of data requested from the particular channel. This value is always fixed at 1.</td>
</tr>
<tr>
<td>request_waitrequest</td>
<td>In</td>
<td>Wait request signal, used to pause the scheduler when the slave cannot accept a new request.</td>
</tr>
<tr>
<td><strong>Avalon-ST Almost-Full Status Interface</strong></td>
<td></td>
<td></td>
</tr>
<tr>
<td>almost_full_valid</td>
<td>In</td>
<td>Indicates that almost_full_channel and almost_full_data are valid.</td>
</tr>
<tr>
<td>almost_full_channel (Channel_Width–1:0)</td>
<td>In</td>
<td>Indicates the channel for the current status indication.</td>
</tr>
<tr>
<td>almost_full_data (log&lt;sub&gt;2&lt;/sub&gt;Max_Channels–1:0)</td>
<td>In</td>
<td>A 1-bit signal that is asserted high to indicate that the channel indicated by almost_full_channel is almost full.</td>
</tr>
</tbody>
</table>

### Parameters

Table 12–26 describes the parameters that you can configure for the Round Robin Scheduler.

### Table 12–26. Parameters for Avalon-ST Round Robin Scheduler Component

<table>
<thead>
<tr>
<th>Parameters</th>
<th>Values</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Number of channels</td>
<td>2–32</td>
<td>Specifies the number of channels the Avalon-ST Round Robin Scheduler supports.</td>
</tr>
<tr>
<td>Use almost-full status</td>
<td>0–1</td>
<td>Specifies whether the almost-full interface is used. If the interface is not used, the core always requests data from the next channel at the next clock cycle.</td>
</tr>
</tbody>
</table>

### Packets to Transactions Converter

The Avalon Packets to Transactions Converter core receives streaming data from upstream components and initiates Avalon-MM transactions. The core then returns Avalon-MM transaction responses to the requesting components.

The SPI Slave to Avalon Master Bridge and JTAG to Avalon Master Bridge are examples of how Packets to Transactions Converter core is used. For more information, refer to the "Avalon Interface Specifications."
Figure 12–14 shows a diagram of the Avalon Packets to Transactions Converter.

![Figure 12–14. Avalon Packets to Transactions Converter Core](image)

**Interfaces**

Table 12–27 shows the properties of Avalon-ST interfaces.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Backpressure</td>
<td>Ready latency = 0.</td>
</tr>
<tr>
<td>Data Width</td>
<td>Data width = 8 bits; Bits per symbol = 8.</td>
</tr>
<tr>
<td>Channel</td>
<td>Not supported.</td>
</tr>
<tr>
<td>Error</td>
<td>Not used.</td>
</tr>
<tr>
<td>Packet</td>
<td>Supported.</td>
</tr>
</tbody>
</table>

The Avalon-MM master interface supports read and write transactions. The data width is set to 32 bits, and burst transactions are not supported.

**Operation**

The Packets to Transactions Converter core receives streams of packets on its Avalon-ST sink interface and initiates Avalon-MM transactions. Upon receiving transaction responses from Avalon-MM slaves, the core transforms the responses to packets and returns them to the requesting components via its Avalon-ST source interface. The core does not report Avalon-ST errors.
Data Packet Formats

The Packets to Transactions Converter core expects incoming data streams to be in the format shown in Table 12–28. A response packet is returned for every write transaction. The core also returns a response packet if a no transaction (0x7f) is received. An invalid transaction code is regarded as a no transaction. For read transactions, the core returns the data read.

Table 12–28. Data Packet Formats

<table>
<thead>
<tr>
<th>Byte</th>
<th>Field</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Transaction code</td>
<td>Type of transaction. See Table 12–27.</td>
</tr>
<tr>
<td>0</td>
<td>Reserved</td>
<td>Reserved for future use.</td>
</tr>
<tr>
<td>[3:2]</td>
<td>Size</td>
<td>Transaction size in bytes. For write transactions, the size indicates the</td>
</tr>
<tr>
<td></td>
<td></td>
<td>size of the data field. For read transactions, the size indicates the total</td>
</tr>
<tr>
<td></td>
<td></td>
<td>number of bytes to read.</td>
</tr>
<tr>
<td>[n:8]</td>
<td>Data</td>
<td>Transaction data; data to be written for write transactions.</td>
</tr>
</tbody>
</table>

Response Packet Format

<table>
<thead>
<tr>
<th>Byte</th>
<th>Field</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>Transaction code</td>
<td>The transaction code with the most significant bit inversed.</td>
</tr>
<tr>
<td>1</td>
<td>Reserved</td>
<td>Reserved for future use.</td>
</tr>
<tr>
<td>[4:2]</td>
<td>Size</td>
<td>Total number of bytes read/written successfully.</td>
</tr>
</tbody>
</table>

Supported Transactions

Table 12–29 lists the Avalon-MM transactions supported by the Packets to Transactions Converter core.

Table 12–29. Transaction Supported

<table>
<thead>
<tr>
<th>Transaction Code</th>
<th>Avalon-MM Transaction</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0x00</td>
<td>Write, non-incrementing address.</td>
<td>Writes data to the given address until the total number of bytes written</td>
</tr>
<tr>
<td></td>
<td></td>
<td>to the same word address equals to the value specified in the size field.</td>
</tr>
<tr>
<td>0x04</td>
<td>Write, incrementing address.</td>
<td>Writes transaction data starting at the given address.</td>
</tr>
<tr>
<td>0x10</td>
<td>Read, non-incrementing address.</td>
<td>Reads 32 bits of data from the given address until the total number of</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bytes read from the same address equals to the value specified in the size</td>
</tr>
<tr>
<td></td>
<td></td>
<td>field.</td>
</tr>
<tr>
<td>0x14</td>
<td>Read, incrementing address.</td>
<td>Reads the number of bytes specified in the size field starting from the</td>
</tr>
<tr>
<td></td>
<td></td>
<td>given address.</td>
</tr>
<tr>
<td>0x7f</td>
<td>No transaction.</td>
<td>No transaction is initiated. You can use this transaction type for testing</td>
</tr>
<tr>
<td></td>
<td></td>
<td>purposes. Although no transaction is initiated on the Avalon-MM interface,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>the core still returns a response packet for this transaction code.</td>
</tr>
</tbody>
</table>

The Packets to Transactions Converter core can process only a single transaction at a time. The ready signal on the core's Avalon-ST sink interface is asserted only when the current transaction is completely processed.
No internal buffer is implemented on the data paths. Data received on the Avalon-ST interface is forwarded directly to the Avalon-MM interface and vice-versa. Asserting the `waitrequest` signal on the Avalon-MM interface backpressures the Avalon-ST sink interface. In the opposite direction, if the Avalon-ST source interface is backpressured, the `read` signal on the Avalon-MM interface is not asserted until the backpressure is alleviated. Backpressuring the Avalon-ST source in the middle of a read could result in data loss. In this case, the core returns the data that is successfully received.

A transaction is considered complete when the core receives an EOP. For write transactions, the actual data size is expected to be the same as the value of the `size` property. Whether or not both values agree, the core always uses the EOP to determine the end of data.

**Malformed Packets**

The following are examples of malformed packets:

- **Consecutive start of packet (SOP)**—An SOP marks the beginning of a transaction. If an SOP is received in the middle of a transaction, the core drops the current transaction without returning a response packet for the transaction, and initiates a new transaction. This effectively precesses packets without an end of packet (EOP).

- **Unsupported transaction codes**—The core processes unsupported transactions as a no transaction.

**Streaming Channel Multiplexer and Demultiplexer Cores**

The Avalon-ST channel multiplexer core receives data from various input interfaces and multiplexes the data into a single output interface, using the optional `channel` signal to indicate the origin of the data. The Avalon-ST channel demultiplexer core receives data from a channelized input interface and drives that data to multiple output interfaces, where the output interface is selected by the input `channel` signal.

The multiplexer and demultiplexer cores can transfer data between interfaces on cores that support the unidirectional flow of data. The multiplexer and demultiplexer allow you to create multiplexed or demultiplexer datapaths without having to write custom HDL code. The multiplexer includes a Round-Robin Scheduler.
Resource Usage and Performance

Resource utilization for the multiplexer and demultiplexer cores depends on the number of input and output interfaces, the width of the datapath, and whether the streaming data uses the optional packet protocol. For the multiplexer, the parameterization of the scheduler also affects resource utilization. Table 12–30 provides estimated resource utilization for eleven different configurations of the multiplexer.

Table 12–30. Multiplexer Estimated Resource Usage and Performance

<table>
<thead>
<tr>
<th>No. of Inputs</th>
<th>Data Width</th>
<th>Scheduling Size (Cycles)</th>
<th>Stratix® II and Stratix II GX (Approximate LEs)</th>
<th>Cyclone® II</th>
<th>Stratix II GX</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td>(f_{\text{MAX}}) (MHz)</td>
<td>ALM Count</td>
<td>(f_{\text{MAX}}) (MHz)</td>
</tr>
<tr>
<td>----------------</td>
<td>-------------</td>
<td>--------------------------</td>
<td>-----------------------------------------------</td>
<td>-------------</td>
<td>---------------</td>
</tr>
<tr>
<td>2</td>
<td>Y</td>
<td>1</td>
<td>500</td>
<td>31</td>
<td>420</td>
</tr>
<tr>
<td>2</td>
<td>Y</td>
<td>2</td>
<td>500</td>
<td>36</td>
<td>417</td>
</tr>
<tr>
<td>2</td>
<td>Y</td>
<td>32</td>
<td>451</td>
<td>43</td>
<td>364</td>
</tr>
<tr>
<td>8</td>
<td>Y</td>
<td>2</td>
<td>401</td>
<td>150</td>
<td>257</td>
</tr>
<tr>
<td>8</td>
<td>Y</td>
<td>32</td>
<td>356</td>
<td>151</td>
<td>219</td>
</tr>
<tr>
<td>16</td>
<td>Y</td>
<td>2</td>
<td>262</td>
<td>333</td>
<td>174</td>
</tr>
<tr>
<td>16</td>
<td>Y</td>
<td>32</td>
<td>310</td>
<td>337</td>
<td>161</td>
</tr>
<tr>
<td>2</td>
<td>N</td>
<td>1</td>
<td>500</td>
<td>23</td>
<td>400</td>
</tr>
<tr>
<td>2</td>
<td>N</td>
<td>9</td>
<td>500</td>
<td>30</td>
<td>420</td>
</tr>
<tr>
<td>11</td>
<td>N</td>
<td>9</td>
<td>292</td>
<td>275</td>
<td>197</td>
</tr>
<tr>
<td>16</td>
<td>N</td>
<td>9</td>
<td>262</td>
<td>295</td>
<td>182</td>
</tr>
</tbody>
</table>

Table 12–31 provides estimated resource utilization for six different configurations of the demultiplexer. The core operating frequency varies with the device, the number of interfaces, and the size of the datapath.

Table 12–31. Demultiplexer Estimated Resource Usage

<table>
<thead>
<tr>
<th>No. of Inputs</th>
<th>Data Width (Symbols per Beat)</th>
<th>Stratix II (Approximate LEs)</th>
<th>Cyclone II</th>
<th>Stratix II GX (Approximate LEs)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td>(f_{\text{MAX}}) (MHz)</td>
<td>ALM Count</td>
<td>(f_{\text{MAX}}) (MHz)</td>
</tr>
<tr>
<td>2</td>
<td>1</td>
<td>500</td>
<td>53</td>
<td>400</td>
</tr>
<tr>
<td>15</td>
<td>1</td>
<td>349</td>
<td>171</td>
<td>235</td>
</tr>
<tr>
<td>16</td>
<td>1</td>
<td>363</td>
<td>171</td>
<td>233</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>500</td>
<td>85</td>
<td>392</td>
</tr>
<tr>
<td>15</td>
<td>2</td>
<td>352</td>
<td>247</td>
<td>213</td>
</tr>
<tr>
<td>16</td>
<td>2</td>
<td>328</td>
<td>280</td>
<td>218</td>
</tr>
</tbody>
</table>
Multiplexer

The Avalon-ST multiplexer takes data from a variety of input data interfaces, and multiplexes the data onto a single output interface. The multiplexer includes a round-robin scheduler that selects from the next input interface that has data. Each input interface has the same width as the output interface, so that the other input interfaces are backpressured when the multiplexer is carrying data from a different input interface.

The multiplexer includes an optional channel signal that allows each input interface to carry channelized data. When the channel signal is enabled on an input interfaces, the multiplexer adds $\log_2(\text{num\_input\_interfaces})$ bits to make the output channel signal. Then the output channel signal has all of the bits of the input channel, plus the bits required to indicate the originating cycle of data. These bits are appended to either the most or least significant bits of the output channel signal.

Figure 12–15 shows a diagram of an Avalon-ST Multiplexer.

The scheduler processes one input interface at a time, selecting it for transfer. Once an input interface has been selected, data from that input interface is sent until one of the following scenarios occurs:

- The specified number of cycles have elapsed.
- The input interface has no more data to send and valid is deasserted on a ready cycle.
- When packets are supported, endofpacket is asserted.

Input Interfaces

Each input interface is an Avalon-ST data interface that optionally supports packets. The input interfaces are identical; they have the same symbol and data widths, error widths, and channel widths.
Output Interface

The output interface carries the multiplexed data stream with data from the inputs. The symbol, data, and error widths are the same as the input interfaces. The width of the channel signal is the same as the input interfaces, with the addition of the bits needed to indicate the origin of the data.

You can configure the following parameters for the output interface:

- **Data Bits Per Symbol**—The bits per symbol is related to the width of readdata and writedata signals, which must be a multiple of the bits per symbol.
  
  If you change only bits per symbol, and do not change the data width, you will get errors.

- **Data Symbols Per Beat**—The number of symbols (words) that are transferred per beat (transfer). Valid values are 1 to 32.

- **Include Packet Support**—Indicates whether or not packet transfers are supported. Packet support includes the startofpacket, endofpacket, and empty signals.

- **Channel Signal Width (bits)**—The number of bits used for the channel signal for input interfaces. A value of 0 indicates that input interfaces do not have channels. A value of 4 indicates that up to 16 channels share the same input interface. The input channel can have a width between 0 to 31 bits. A value of 0 means that the optional channel signal is not used.

- **Error Signal Width (bits)**—The width of the error signal for input and output interfaces. A value of 0 means the error signal is not used.

Multiplexer Parameters

You can configure the following parameters for the multiplexer:

- **Number of Input Ports**—The number of input interfaces that the multiplexer supports. Valid values are 2 to 16.

- **Scheduling Size (Cycles)**—The number of cycles that are sent from a single channel before changing to the next channel.

- **Use Packet Scheduling**—When this parameter is turned on, the multiplexer only switches the selected input interface on packet boundaries. Therefore, packets on the output interface are not interleaved.

- **Use high bits to indicate source port**—When this parameter is turned on, the high bits of the output channel signal are used to indicate the origin of the input interface of the data. For example, if the input interfaces have 4-bit channel signals, and the multiplexer has 4 input interfaces, the output interface has a 6-bit channel signal. If this parameter is turned on, bits [5:4] of the output channel signal indicate origin of the input interface of the data, and bits [3:0] are the channel bits that were presented at the input interface.
Demultiplexer

That Avalon-ST demultiplexer takes data from a channelized input data interface and provides that data to multiple output interfaces, where the output interface selected for a particular transfer is specified by the input channel signal. The data is delivered to the output interfaces in the same order it is received at the input interface, regardless of the value of channel, packet, frame, or any other signal. Each of the output interfaces has the same width as the input interface; each output interface is idle when the demultiplexer is driving data to a different output interface. The demultiplexer uses \( \log_2(\text{num_output_interfaces}) \) bits of the channel signal to select the output for the data; the remainder of the channel bits are forwarded to the appropriate output interface unchanged. Figure 12–16 shows a diagram of an Avalon-ST Demultiplexer.

Figure 12–16. Demultiplexer

Input Interface

Each input interface is an Avalon-ST data interface that optionally supports packets. You can configure the following parameters for the input interface:

- **Data Bits Per Symbol**—The bits per symbol is related to the width of `readdata` and `writedata` signals, which must be a multiple of the bits per symbol.
  
  If you change only bits per symbol, and do not change the data width, you will get errors.

- **Data Symbols Per Beat**—The number of symbols (words) that are transferred per beat (transfer). Valid values are 1 to 32.

- **Include Packet Support**—Indicates whether or not data packet transfers are supported. Packet support includes the `startofpacket`, `endofpacket`, and `empty` signals.

- **Channel Signal Width (bits)**—The number of bits used for the channel signal for output interfaces. A value of 0 means that output interfaces do not use the optional channel signal.

- **Error Signal Width (bits)**—The width of the error signal for input and output interfaces. A value of 0 means the error signal is not unused.
Output Interfaces

Each output interface carries data from a subset of channels from the input interface. Each output interface is identical; all have the same symbol and data widths, error widths, and channel widths. The symbol, data, and error widths are the same as the input interface. The width of the channel signal is the same as the input interface, without the bits that were used to select the output interface.

Parameters

You can configure the following parameters for the demultiplexer:

- **Number of Output Ports**—The number of output interfaces that the multiplexer supports Valid values are 2 to 16.

- **High channel bits select output**—When this option is turned on, the high bits of the input channel signal are used by the demultiplexing function and the low order bits are passed to the output. When this option is turned off, the low order bits are used and the high order bits are passed through.

- **Figure 12–17** illustrates the significance of the location of signals; for example, there is one input interface and two output interfaces. If the low-order bits of the channel signal select the output interfaces, the even channels goes to channel 0, and the odd channels goes to channel 1. If the high-order bits of the channel signal select the output interface, channels 0 to 7 goes to channel 0 and channels 8 to 15 goes to channel 1.

![Figure 12–17. Select Bits for Demultiplexer](image)

Hardware Simulation Considerations

The multiplexer and demultiplexer components do not provide a simulation testbench for simulating a stand-alone instance of the component. However, you can use the standard Qsys simulation flow to simulate the component design files.

Software Programming Model

The multiplexer and demultiplexer components do not have any user-visible control or status registers. Therefore, Qsys cannot control or configure any aspect of the multiplexer or demultiplexer at run-time. The components cannot generate interrupts.
Streaming Pipeline Stage

The Avalon-ST pipeline stage receives data from an Avalon-ST source interface, and outputs the data to an Avalon-ST sink interface. If the ready signal is not pipelined, the pipeline stage becomes a simple register as shown in Figure 12–18 on page 12–45. If the ready signal is pipelined, the pipeline stage must also include a second "holding" register as shown in Figure 12–19 on page 12–45. In the absence of back pressure, the Avalon-ST pipeline stage source interface outputs data one cycle after receiving the data on its sink interface. If the pipeline stage receives back pressure on its source interface, it continues to assert its source interface's current data output. While the pipeline stage is receiving back pressure on its source interface and it receives new data on its sink interface, the pipeline stage will internally buffer the new data, and assert back pressure on its sink interface.

Once the back pressure is deasserted, the pipeline stage’s source interface is de-asserted and the pipeline stage will assert internally buffered data (if present). Additionally, the pipeline stage de-asserts back pressure on its sink interface.

**Figure 12–18. Pipeline Stage Simple Register**

![Pipeline Stage Simple Register](image1)

**Figure 12–19. Pipeline Stage Holding Register**

![Pipeline Stage Holding Register](image2)

Single-Clock and Dual-Clock FIFO Cores

The Avalon-ST Single-Clock and Avalon-ST Dual-Clock FIFO cores are FIFO buffers which operate with a common clock and independent clocks for input and output ports respectively.

**Figure 12–20** and **Figure 12–21** show diagrams of the Avalon-ST Single-Clock FIFO and Avalon-ST Dual-Clock FIFO cores.
Interfaces

This section describes the interfaces implemented in the FIFO cores.

**Avalon-ST Data Interface**

Each FIFO core has an Avalon-ST data sink and source interfaces. The data sink and source interfaces in the dual-clock FIFO core are driven by different clocks.

Table 12–32 shows the properties of the Avalon-ST interfaces.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Backpressure</td>
<td>Ready latency = 0.</td>
</tr>
<tr>
<td>Data Width</td>
<td>Configurable.</td>
</tr>
</tbody>
</table>
You can configure the single-clock FIFO core to include an optional Avalon-MM interface, and the dual-clock FIFO core to include an Avalon-MM interface in each clock domain. The Avalon-MM interface provides access to 32-bit registers, which allows you to retrieve the FIFO buffer fill level and configure the almost-empty and almost-full thresholds. In the single-clock FIFO core, you can also configure the packet and error handling modes. See Table 12–34 and Table 12–35 for the register descriptions.

### Avalon-ST Status Interface

The single-clock FIFO core has two optional Avalon-ST status source interfaces from which you can obtain the FIFO buffer almost-full and almost empty statuses.

### Operating Modes

The FIFO operating modes are as follows:

- **Default mode**—The core accepts incoming data on the in interface (Avalon-ST data sink) and forwards it to the out interface (Avalon-ST data source). The core asserts the valid signal on the Avalon-ST source interface to indicate that data is available at the interface.

- **Store and forward mode**—This mode applies only to the single-clock FIFO core. The core asserts the valid signal on the out interface only when a full packet of data is available at the interface. In this mode, you can also enable the drop-on-error feature by setting the drop_on_error register to 1. When this feature is enabled, the core drops all packets received with the in_error signal asserted.

- **Cut-through mode**—This mode applies only to the single-clock FIFO core. The core asserts the valid signal on the out interface to indicate that data is available for consumption when the number of entries specified in the cut_through_threshold register are available in the FIFO buffer.

To use the store and forward or cut-through mode, turn on the Use store and forward parameter to include the csr interface (Avalon-MM slave). Set the cut_through_threshold register to 0 to enable the store and forward mode, and then set the register to any value greater than 0 to enable the cut-through mode. The non-zero value specifies the minimum number of FIFO entries that must be available before the data is ready for consumption. Setting the register to 1 provides you with the default mode.

### Table 12–32. Properties of Avalon-ST Interfaces

<table>
<thead>
<tr>
<th>Feature</th>
<th>Property</th>
</tr>
</thead>
<tbody>
<tr>
<td>Channel</td>
<td>Supported, up to 255 channels.</td>
</tr>
<tr>
<td>Error</td>
<td>Configurable.</td>
</tr>
<tr>
<td>Packet</td>
<td>Configurable.</td>
</tr>
</tbody>
</table>
Fill Level

You can obtain the fill level of the FIFO buffer via the optional Avalon-MM control and status interface. Turn on the Use fill level parameter (Use sink fill level and Use source fill level in the dual-clock FIFO core) and read the fill_level register.

The dual-clock FIFO core has two fill levels, one in each clock domain. Due to the latency of the clock crossing logic, the fill levels reported in the input and output clock domains may be different at any given instance. In both cases, the fill level may report badly for the clock domain; that is, the fill level is reported high in the input clock domain, and low in the output clock domain.

The dual-clock FIFO has an output pipeline stage to improve f<sub>MAX</sub>. This output stage is accounted for when calculating the output fill level, but not when calculating the input fill level. Therefore, the best measure of the amount of data in the FIFO is given by the fill level in the output clock domain, while the fill level in the input clock domain represents the amount of space available in the FIFO (available space = FIFO depth – input fill level).

Thresholds

You can use almost-full and almost-empty thresholds as a mechanism to prevent FIFO overflow and underflow. This feature is available only in the single-clock FIFO core.

To use the thresholds, turn on the Use fill level, Use almost-full status, and Use almost-empty status parameters. You can access the almost_full_threshold and almost_empty_threshold registers via the csr interface and set the registers to an optimal value for your application. See Table 12–34 on page 12–49 for the register description.

You can obtain the almost-full and almost-empty statuses from almost_full and almost_empty interfaces (Avalon-ST status source). The core asserts the almost_full signal when the fill level is equal to or higher than the almost-full threshold. Likewise, the core asserts the almost_empty signal when the fill level is equal to or lower than the almost-empty threshold.

Parameters

Table 12–33 describes the parameters that you can configure for the Single-Clock and Dual-Clock FIFO cores.

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Legal Values</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Bits per symbol</td>
<td>1–32</td>
<td>These parameters determine the width of the FIFO.</td>
</tr>
<tr>
<td>Symbols per beat</td>
<td>1–32</td>
<td>FIFO width = Bits per symbol * Symbols per beat, where: Bits per symbol is the number of bits in a symbol, and Symbols per beat is the number of symbols transferred in a beat.</td>
</tr>
<tr>
<td>Error width</td>
<td>0–32</td>
<td>The width of the error signal.</td>
</tr>
<tr>
<td>FIFO depth</td>
<td>1–32</td>
<td>The FIFO depth. An output pipeline stage is added to the FIFO to increase performance, which increases the FIFO depth by one.</td>
</tr>
<tr>
<td>Use packets</td>
<td>—</td>
<td>Turn on this parameter to enable data packet support on the Avalon-ST data interfaces.</td>
</tr>
</tbody>
</table>
For more information on metastability in Altera devices, refer to AN 42: Metastability in Altera Devices. For more information on metastability analysis and synchronization register chains, refer to the Area and Timing Optimization chapter in volume 2 of the Quartus II Handbook.

Register Description

The csr interface in the Avalon-ST Single Clock FIFO core provides access to registers. Table 12–34 describes the registers.

Table 12–34. Register Description for Avalon-ST Single-Clock FIFO (Part 1 of 2)

<table>
<thead>
<tr>
<th>32-Bit Word Offset</th>
<th>Name</th>
<th>Access</th>
<th>Reset</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>fill_level</td>
<td>R</td>
<td>0</td>
<td>24-bit FIFO fill level. Bits 24 to 31 are unused.</td>
</tr>
<tr>
<td>1</td>
<td>Reserved</td>
<td></td>
<td></td>
<td>Reserved for future use.</td>
</tr>
<tr>
<td>2</td>
<td>almost_full_threshold</td>
<td>RW</td>
<td>FIFO depth–1</td>
<td>Set this register to a value that indicates the FIFO buffer is getting full.</td>
</tr>
<tr>
<td>3</td>
<td>almost_empty_threshold</td>
<td>RW</td>
<td>0</td>
<td>Set this register to a value that indicates the FIFO buffer is getting empty.</td>
</tr>
</tbody>
</table>
The `in_csr` and `out_csr` interfaces in the Avalon-ST Dual Clock FIFO core reports the FIFO fill level. Table 12–35 describes the fill level.

### Table 12–35. Register Description for Avalon-ST Dual-Clock FIFO

<table>
<thead>
<tr>
<th>32-Bit Word Offset</th>
<th>Name</th>
<th>Access</th>
<th>Reset</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td><code>cut_through_threshold</code></td>
<td>RW</td>
<td>0</td>
<td>0—Enables store and forward mode. &gt;0—Enables cut-through mode and specifies the minimum of entries in the FIFO buffer before the valid signal on the Avalon-ST source interface is asserted. Once the FIFO core starts sending the data to the downstream component, it continues to do so until the end of the packet. This register applies only when the Use store and forward parameter is turned on.</td>
</tr>
<tr>
<td>5</td>
<td><code>drop_on_error</code></td>
<td>RW</td>
<td>0</td>
<td>0—Disables drop-on error. 1—Enables drop-on error. This register applies only when the Use packet and Use store and forward parameters are turned on.</td>
</tr>
</tbody>
</table>

For more information about the Avalon interfaces, refer to the Avalon Interface Specifications or the Avalon Memory-Mapped Design Optimizations chapter in the Embedded Design Handbook.

### Document Revision History

Table 12–36 shows the revision history for this document.

### Table 12–36. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>• Added Streaming Pipeline Stage support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Added AMBA APB support.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>• Moved relevant content from Embedded IP User Guide.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
When designing for large and complex FPGAs, your design and coding styles can impact your quality of results significantly. Designs reflecting synchronous design practices behave predictably reliably, even when re-targeted to different device families or speed grades. Using recommended HDL coding styles ensures that synthesis tools can infer the optimal device hardware to implement your design. Following best practices when creating your design hierarchy and logic provides the most flexibility when partitioning the design for incremental compilation, and leads to the best results. If you create floorplan location assignments to control the placement of different design blocks (useful in team-based designs so each designer can target a different area of the device floorplan), following best practices is important to maintaining good design performance.

This section presents design and coding style recommendations in the following chapters:

- **Chapter 13, Recommended Design Practices**
  This chapter describes synchronous design practices, and provides guidelines for combinational logic structures, clocking schemes, and best practices for physical implementation and timing closure. It also explains how to check design rules using the Quartus® II Design Assistant. Finally, it discusses use of clock and register-control features in device architecture.

- **Chapter 14, Recommended HDL Coding Styles**
  This chapter discusses Altera megafunctions and provides specific Verilog HDL and VHDL coding examples to insure the Quartus II software infers Altera dedicated logic such as memory and DSP blocks. It also provides device-specific coding recommendations for registers and certain logic functions such as tri-state signals, multiplexers, and cyclic redundancy check (CRC) functions, and includes references to other Altera documentation for low-level logic design information.

- **Chapter 15, Managing Metastability with the Quartus II Software**
  This chapter describes ways you can use the Quartus II software to analyze the average mean time between failures (MTBF) due to metastability caused by synchronization of asynchronous signals, and optimize the design to improve the metastability MTBF.

- **Chapter 16, Best Practices for Incremental Compilation Partitions and Floorplan Assignments**
  This chapter provides a set of guidelines to help you set up and partition your design to take advantage of the compilation time savings, performance preservation, and hierarchical design features offered by Quartus II incremental compilation, and to help you create a design floorplan (using LogicLock™ regions) to support the flow when required.
Use this chapter when setting up your design hierarchy and determining the interfaces between logic blocks in your design, as well as if/when you create a design floorplan. You can also use this chapter to make changes to a design that was not originally set up to take advantage of incremental compilation, because it provides tips on changing a design to work better with an incremental design flow.
This chapter provides design recommendations for Altera® devices and describes the Quartus® II Design Assistant, which helps you check your design for violations of Altera’s design recommendations.

Current FPGA applications have reached the complexity and performance requirements of ASICS. In the development of complex system designs, good design practices have an enormous impact on the timing performance, logic utilization, and system reliability of a device. Well-coded designs behave in a predictable and reliable manner even when retargeted to different families or speed grades. Good design practices also aid in successful design migration between FPGA and HardCopy® or ASIC implementations for prototyping and production.

For optimal performance, reliability, and faster time-to-market when designing with Altera devices, you should adhere to the following guidelines:

- Understand the impact of synchronous design practices
- Follow recommended design techniques, including hierarchical design partitioning, and timing closure guidelines
- Take advantage of the architectural features in the targeted device

This chapter contains the following sections:

- “Synchronous FPGA Design Practices” on page 13–2
- “Design Guidelines” on page 13–4
- “Optimizing for Physical Implementation and Timing Closure” on page 13–12
- “Checking Design Violations” on page 13–16
- “Targeting Clock and Register-Control Architectural Features” on page 13–23
- “Targeting Embedded RAM Architectural Features” on page 13–35

For specific HDL coding examples and recommendations, including coding guidelines for targeting dedicated device hardware, such as memory and digital signal processing (DSP) blocks, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. For information about partitioning a hierarchical design for incremental compilation, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Synchronous FPGA Design Practices

The first step in good design methodology is to understand the implications of your design practices and techniques. This section outlines the benefits of optimal synchronous design practices and the hazards involved in other techniques. Good synchronous design practices can help you meet your design goals consistently. Problems with other design techniques can include reliance on propagation delays in a device, which can lead to race conditions, incomplete timing analysis, and possible glitches.

In a synchronous design, some clock signals trigger every event. As long as you ensure that all the timing requirements of the registers are met, a synchronous design behaves in a predictable and reliable manner for all process, voltage, and temperature (PVT) conditions. You can easily target synchronous designs to different device families or speed grades. In addition, synchronous design practices help ensure successful migration if you plan to migrate your design to a high-volume solution such as a HardCopy device or if you are prototyping an ASIC design.

Fundamentals of Synchronous Design

In a synchronous design, the clock signal controls the activities of all inputs and outputs. On every active edge of the clock (usually the rising edge), the data inputs of registers are sampled and transferred to outputs. Following an active clock edge, the outputs of combinational logic feeding the data inputs of registers change values. This change triggers a period of instability due to propagation delays through the logic as the signals go through several transitions and finally settle to new values. Changes that occur on data inputs of registers do not affect the values of their outputs until after the next active clock edge.

Because the internal circuitry of registers isolates data outputs from inputs, instability in the combinational logic does not affect the operation of the design as long as you meet the following timing requirements:

- Before an active clock edge, you must ensure that the data input has been stable for at least the setup time of the register.
- After an active clock edge, you must ensure that the data input remains stable for at least the hold time of the register.

When you specify all of your clock frequencies and other timing requirements, the Quartus II TimeQuest Timing Analyzer reports actual hardware requirements for the setup times ($t_{SU}$) and hold times ($t_{H}$) for every pin in your design. By meeting these external pin requirements and following synchronous design techniques, you ensure that you satisfy the setup and hold times for all registers in your device.

To meet setup and hold time requirements on all input pins, any inputs to combinational logic that feed a register should have a synchronous relationship with the clock of the register. If signals are asynchronous, you can register the signals at the inputs of the device to help prevent a violation of the required setup and hold times.
When you violate the setup or hold time of a register, you might oscillate the output, or set the output to an intermediate voltage level between the high and low levels called a metastable state. In this unstable state, small perturbations such as noise in power rails can cause the register to assume either the high or low voltage level, resulting in an unpredictable valid state. Various undesirable effects can occur, including increased propagation delays and incorrect output states. In some cases, the output can even oscillate between the two valid states for a relatively long period of time.

For information about timing requirements and analysis in the Quartus II software, refer to *About TimeQuest Timing Analysis* in Quartus II Help.

### Hazards of Asynchronous Design

In the past, designers have often used asynchronous techniques such as ripple counters or pulse generators in programmable logic device (PLD) designs, enabling them to take “short cuts” to save device resources. Asynchronous design techniques have inherent problems such as relying on propagation delays in a device, which can vary with temperature and voltage fluctuations, resulting in incomplete timing constraints and possible glitches and spikes.

Some asynchronous design structures rely on the relative propagation delays of signals to function correctly. In these cases, race conditions can arise where the order of signal changes can affect the output of the logic. PLD designs can have varying timing delays, depending on how the design is placed and routed in the device with each compilation. Therefore, it is almost impossible to determine the timing delay associated with a particular block of logic ahead of time. As devices become faster due to device process improvements, the delays in an asynchronous design may decrease, resulting in a design that does not function as expected. Specific examples are provided in “Design Guidelines” on page 13–4. Relying on a particular delay also makes asynchronous designs difficult to migrate to different architectures, devices, or speed grades.

The timing of asynchronous design structures is often difficult or impossible to model with timing assignments and constraints. If you do not have complete or accurate timing constraints, the timing-driven algorithms used by your synthesis and place-and-route tools may not be able to perform the best optimizations, and the reported results may not be complete.

Some asynchronous design structures can generate harmful glitches, which are pulses that are very short compared with clock periods. Most glitches are generated by combinational logic. When the inputs of combinational logic change, the outputs exhibit several glitches before they settle to their new values. These glitches can propagate through the combinational logic, leading to incorrect values on the outputs in asynchronous designs. In a synchronous design, glitches on the data inputs of registers are normal events that have no negative consequences because the data is not processed until the clock edge.
Design Guidelines

When designing with HDL code, you should understand how a synthesis tool interprets different HDL design techniques and what results to expect. Your design techniques can affect logic utilization and timing performance, as well as the design’s reliability. This section describes basic design techniques that ensure optimal synthesis results for designs targeted to Altera devices while avoiding several common causes of unreliability and instability. Altera recommends that you design your combinational logic carefully to avoid potential problems and pay attention to your clocking schemes so that you can maintain synchronous functionality and avoid timing problems.

Combinational Logic Structures

Combinational logic structures consist of logic functions that depend only on the current state of the inputs. In Altera FPGAs, these functions are implemented in the look-up tables (LUTs) with either logic elements (LEs) or adaptive logic modules (ALMs). For cases where combinational logic feeds registers, the register control signals can implement part of the logic function to save LUT resources. By following the recommendations in this section, you can improve the reliability of your combinational design.

Combinational Loops

Combinational loops are among the most common causes of instability and unreliability in digital designs. Combinational loops generally violate synchronous design principles by establishing a direct feedback loop that contains no registers. You should avoid combinational loops whenever possible. In a synchronous design, feedback loops should include registers. For example, a combinational loop occurs when the left-hand side of an arithmetic expression also appears on the right-hand side in HDL code. A combinational loop also occurs when you feed back the output of a register to an asynchronous pin of the same register through combinational logic, as shown in Figure 13–1.

Figure 13–1. Combinational Loop Through Asynchronous Control Pin

Use recovery and removal analysis to perform timing analysis on asynchronous ports, such as clear or reset in the Quartus II software.

If you are using the TimeQuest Timing Analyzer, refer to Specifying Timing Constraints and Exceptions (TimeQuest Timing Analyzer) in Quartus II Help for details about how TimeQuest analyzer performs recovery and removal analysis.
Combinational loops are inherently high-risk design structures for the following reasons:

- Combinational loop behavior generally depends on relative propagation delays through the logic involved in the loop. As discussed, propagation delays can change, which means the behavior of the loop is unpredictable.

- Combinational loops can cause endless computation loops in many design tools. Most tools break open combinational loops to process the design. The various tools used in the design flow may open a given loop in a different manner, processing it in a way that is inconsistent with the original design intent.

### Latches

A latch is a small circuit with combinational feedback that holds a value until a new value is assigned. You can implement latches with the Quartus II Text Editor or Block Editor. It is common for mistakes in HDL code to cause unintended latch inference; Quartus II Synthesis issues a warning message if this occurs.

Unlike other technologies, a latch in FPGA architecture is not significantly smaller than a register. The architecture is not optimized for latch implementation and latches generally have slower timing performance compared to equivalent registered circuitry.

Latches have a transparent mode in which data flows continuously from input to output. A positive latch is in transparent mode when the enable signal is high (low for negative latch). In transparent mode, glitches on the input can pass through to the output because of the direct path created. This presents significant complexity for timing analysis. Typical latch schemes use multiple enable phases to prevent long transparent paths from occurring. However, timing analysis cannot identify these safe applications.

The TimeQuest analyzer analyzes latches as synchronous elements clocked on the falling edge of the positive latch signal by default, and allows you to treat latches as having nontransparent start and end points. Be aware that even an instantaneous transition through transparent mode can lead to glitch propagation. The TimeQuest analyzer cannot perform cycle-borrowing analysis.

Due to various timing complexities, latches have limited support in formal verification tools. Therefore, you should not rely on formal verification for a design that includes latches.

Avoid using latches to ensure that you can completely analyze the timing performance and reliability of your design.

### Delay Chains

You require delay chains when you use two or more consecutive nodes with a single fan-in and a single fan-out to cause delay. Inverters are often chained together to add delay. Delay chains are sometimes used to resolve race conditions created by other asynchronous design practices.
Delays in PLD designs can change with each placement and routing cycle. Effects such as rise and fall time differences and on-chip variation mean that delay chains, especially those placed on clock paths, can cause significant problems in your design. Refer to “Hazards of Asynchronous Design” on page 13–3 for examples of the kinds of problems that delay chains can cause. Avoid using delay chains to prevent these kinds of problems.

In some ASIC designs, delays are used for buffering signals as they are routed around the device. This functionality is not required in FPGA devices because the routing structure provides buffers throughout the device.

**Pulse Generators and Multivibrators**

You can use delay chains to generate either one pulse (pulse generators) or a series of pulses (multivibrators). There are two common methods for pulse generation, as shown in Figure 13–2. These techniques are purely asynchronous and must be avoided.

Figure 13–2. Asynchronous Pulse Generators

In Figure 13–2, a trigger signal feeds both inputs of a 2-input AND gate, but the design adds inverters to create a delay chain to one of the inputs. The width of the pulse depends on the time differences between path that feeds the gate directly, and the path that goes through the delay chain. This is the same mechanism responsible for the generation of glitches in combinational logic following a change of input values. This technique artificially increases the width of the glitch.

As also shown in Figure 13–2, a register’s output drives the same register’s asynchronous reset signal through a delay chain. The register resets itself asynchronously after a certain delay.

The width of pulses generated in this way are difficult for synthesis and place-and-route to determine, set, or verify. The actual pulse width can only be determined after placement and routing, when routing and propagation delays are known. You cannot reliably create a specific pulse width when creating HDL code, and it cannot be set by EDA tools. The pulse may not be wide enough for the application under all PVT conditions. Also, the pulse width changes if you change to a different device. Additionally, verification is difficult because static timing analysis cannot verify the pulse width.
Multivibrators use a glitch generator to create pulses, together with a combinational loop that turns the circuit into an oscillator. This creates additional problems because of the number of pulses involved. Additionally, when the structures generate multiple pulses, they also create a new artificial clock in the design which must be analyzed by design tools.

When you must use a pulse generator, use synchronous techniques, as shown in Figure 13–3.

**Figure 13–3. Recommended Pulse-Generation Technique**

In Figure 13–3, the pulse width is always equal to the clock period. This pulse generator is predictable, can be verified with timing analysis, and is easily moved to other architectures, devices, or speed grades.

**Clocking Schemes**

Like combinational logic, clocking schemes have a large effect on the performance and reliability of a design. Avoid using internally generated clocks (other than PLLs) wherever possible because they can cause functional and timing problems in the design. Clocks generated with combinational logic can introduce glitches that create functional problems, and the delay inherent in combinational logic can lead to timing problems.

Specify all clock relationships in the Quartus II software to allow for the best timing-driven optimizations during fitting and to allow correct timing analysis. Use clock setting assignments on any derived or internal clocks to specify their relationship to the base clock.

Use global device-wide, low-skew dedicated routing for all internally-generated clocks, instead of routing clocks on regular routing lines. For more information, refer to “Clock Network Resources” on page 13–23.

Avoid data transfers between different clocks wherever possible. If you require a data transfer between different clocks, use FIFO circuitry. You can use the clock uncertainty features in the Quartus II software to compensate for the variable delays between clock domains. Consider setting a clock setup uncertainty and clock hold uncertainty value of 10% to 15% of the clock delay.

The following sections provide specific examples and recommendations for avoiding clocking scheme problems.
**Internally Generated Clocks**

If you use the output from combinational logic as a clock signal or as an asynchronous reset signal, you can expect to see glitches in your design. In a synchronous design, glitches on data inputs of registers are normal events that have no consequences. However, a glitch or a spike on the clock input (or an asynchronous input) to a register can have significant consequences. Narrow glitches can violate the register’s minimum pulse width requirements. Setup and hold requirements might also be violated if the data input of the register changes when a glitch reaches the clock input. Even if the design does not violate timing requirements, the register output can change value unexpectedly and cause functional hazards elsewhere in the design.

To avoid these problems, you should always register the output of combinational logic before you use it as a clock signal (Figure 13–4).

**Figure 13–4. Recommended Clock-Generation Technique**

Registering the output of combinational logic ensures that glitches generated by the combinational logic are blocked at the data input of the register.

**Divided Clocks**

Designs often require clocks that you create by dividing a master clock. Most Altera FPGAs provide dedicated phase-locked loop (PLL) circuitry for clock division. Using dedicated PLL circuitry can help you to avoid many of the problems that can be introduced by asynchronous clock division logic.

When you must use logic to divide a master clock, always use synchronous counters or state machines. Additionally, create your design so that registers always directly generate divided clock signals, as described in “Internally Generated Clocks”, and route the clock on global clock resources. To avoid glitches, do not decode the outputs of a counter or a state machine to generate clock signals.

**Ripple Counters**

To simplify verification, avoid ripple counters in your design. In the past, FPGA designers implemented ripple counters to divide clocks by a power of two because the counters are easy to design and may use fewer gates than their synchronous counterparts. Ripple counters use cascaded registers, in which the output pin of one register feeds the clock pin of the register in the next stage. This cascading can cause problems because the counter creates a ripple clock at each stage. These ripple clocks must be handled properly during timing analysis, which can be difficult and may require you to make complicated timing assignments in your synthesis and placement and routing tools.
You can often use ripple clock structures to make ripple counters out of the smallest amount of logic possible. However, in all Altera devices supported by the Quartus II software, using a ripple clock structure to reduce the amount of logic used for a counter is unnecessary because the device allows you to construct a counter using one logic element per counter bit. You should avoid using ripple counters completely.

**Multiplexed Clocks**

Use clock multiplexing to operate the same logic function with different clock sources. In these designs, multiplexing selects a clock source, as shown in Figure 13–5. For example, telecommunications applications that deal with multiple frequency standards often use multiplexed clocks.

**Figure 13–5. Multiplexing Logic and Clock Sources**

Adding multiplexing logic to the clock signal can create the problems addressed in the previous sections, but requirements for multiplexed clocks vary widely, depending on the application. Clock multiplexing is acceptable when the clock signal uses global clock routing resources and if the following criteria are met:

- The clock multiplexing logic does not change after initial configuration
- The design uses multiplexing logic to select a clock for testing purposes
- Registers are always reset when the clock switches
- A temporarily incorrect response following clock switching has no negative consequences

If the design switches clocks in real time with no reset signal, and your design cannot tolerate a temporarily incorrect response, you must use a synchronous design so that there are no timing violations on the registers, no glitches on clock signals, and no race conditions or other logical problems. By default, the Quartus II software optimizes and analyzes all possible paths through the multiplexer and between both internal clocks that may come from the multiplexer. This may lead to more restrictive analysis than required if the multiplexer is always selecting one particular clock. If you do not require the more complete analysis, you can assign the output of the multiplexer as a base clock in the Quartus II software, so that all register-to-register paths are analyzed using that clock.
Use dedicated hardware to perform clock multiplexing when it is available, instead of using multiplexing logic. For example, you can use the clock-switchover feature or clock control block available in certain Altera devices. These dedicated hardware blocks ensure that you use global low-skew routing lines and avoid any possible hold time problems on the device due to logic delay on the clock line.

For device-specific information about clocking structures, refer to the appropriate device data sheet or handbook on the Literature page of the Altera website.

**Gated Clocks**

Gated clocks turn a clock signal on and off using an enable signal that controls gating circuitry, as shown in Figure 13–6. When a clock is turned off, the corresponding clock domain is shut down and becomes functionally inactive.

![Gated Clock](image)

You can use gated clocks to reduce power consumption in some device architectures by effectively shutting down portions of a digital circuit when they are not in use. When a clock is gated, both the clock network and the registers driven by it stop toggling, thereby eliminating their contributions to power consumption. However, gated clocks are not part of a synchronous scheme and therefore can significantly increase the effort required for design implementation and verification. Gated clocks contribute to clock skew and make device migration difficult. These clocks are also sensitive to glitches, which can cause design failure.

Use dedicated hardware to perform clock gating rather than an AND or OR gate. For example, you can use the clock control block in newer Altera devices to shut down an entire clock network. Dedicated hardware blocks ensure that you use global routing with low skew, and avoid any possible hold time problems on the device due to logic delay on the clock line.

From a functional point of view, you can shut down a clock domain in a purely synchronous manner using a synchronous clock enable signal. However, when using a synchronous clock enable scheme, the clock network continues toggling. This practice does not reduce power consumption as much as gating the clock at the source does. In most cases, use a synchronous scheme such as those described in “Synchronous Clock Enables”. For improved power reduction when gating clocks with logic, refer to “Recommended Clock-Gating Methods” on page 13–11.
Synchronous Clock Enables

To turn off a clock domain in a synchronous manner, use a synchronous clock enable signal. FPGAs efficiently support clock enable signals because there is a dedicated clock enable signal available on all device registers. This scheme does not reduce power consumption as much as gating the clock at the source because the clock network keeps toggling, and performs the same function as a gated clock by disabling a set of registers. Insert a multiplexer in front of the data input of every register to either load new data, or copy the output of the register (Figure 13–7).

Figure 13–7. Synchronous Clock Enable

---

Recommended Clock-Gating Methods

Use gated clocks only when your target application requires power reduction and when gated clocks are able to provide the required reduction in your device architecture. If you must use clocks gated by logic, implement these clocks using the robust clock-gating technique shown in Figure 13–8 and ensure that the gated clock signal uses dedicated global clock routing.

You can gate a clock signal at the source of the clock network, at each register, or somewhere in between. Because the clock network contributes to switching power consumption, gate the clock at the source whenever possible, so that you can shut down the entire clock network instead of gating it further along the clock network at the registers.

Figure 13–8. Recommended Clock-Gating Technique

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In the technique shown in Figure 13–8, a register generates the enable signal to ensure that the signal is free of glitches and spikes. The register that generates the enable signal is triggered on the inactive edge of the clock to be gated. Use the falling edge when gating a clock that is active on the rising edge, as shown in Figure 13–8. Using this technique, only one input of the gate that turns the clock on and off changes at a time. This prevents glitches or spikes on the output. Use an AND gate to gate a clock that is active on the rising edge. For a clock that is active on the falling edge, use an OR gate to gate the clock and register the enable command with a positive edge-triggered register.
When using this technique, pay close attention to the duty cycle of the clock and the delay through the logic that generates the enable signal because you must generate the enable command in one-half the clock cycle. This situation might cause problems if the logic that generates the enable command is particularly complex, or if the duty cycle of the clock is severely unbalanced. However, careful management of the duty cycle and logic delay may be an acceptable solution when compared with problems created by other methods of gating clocks.

Ensure that you apply a clock setting to the gated clock in the TimeQuest analyzer. As shown in Figure 13–8 on page 13–11, apply a clock setting to the output of the AND gate. Otherwise, the timing analyzer might analyze the circuit using the clock path through the register as the longest clock path and the path that skips the register as the shortest clock path, resulting in artificial clock skew.

In certain cases, converting the gated clocks to clock enables may help reduce glitch and clock skew, and eventually produce a more accurate timing analysis. You can set the Quartus II software to automatically convert gated clocks to clock enables by turning on the **Auto Gated Clock Conversion** option. The conversion applies to two types of gated clocking schemes: single-gated clock and cascaded-gated clock. The TimeQuest analyzer supports this option for Arria® II, Arria II GX, Cyclone® II, Cyclone III, Cyclone IV, HardCopy series, Stratix® II, Stratix II GX, Stratix III, Stratix IV, and Stratix V devices.

For information about the settings and limitations of this option, refer to the “Auto Gated Clock Conversion” section of the *Quartus II Integrated Synthesis* chapter in volume 1 of the *Quartus II Handbook*.

### Optimizing for Physical Implementation and Timing Closure

This section provides design and timing closure techniques for high speed or complex core logic designs with challenging timing requirements. These techniques may also be helpful for low or medium speed designs. Best practices for high-speed designs include the following:

- **Planning Physical Implementation**
- **Planning FPGA Resources**
- **Optimizing for Timing Closure**

#### Planning Physical Implementation

When planning a design, consider the following elements of physical implementation:

- The number of unique clock domains and their relationships
- The amount of logic in each functional block
- The location and direction of data flow between blocks
- How data routes to the functional blocks between I/O interfaces
Interface-wide control or status signals may have competing or opposing constraints. For example, when a functional block’s control or status signals interface with physical channels from both sides of the device. In such cases you must provide enough pipeline register stages to allow these signals to traverse the width of the device. In addition, you can structure the hierarchy of the design into separate logic modules for each side of the device. The side modules can generate and use registered control signals per side. This simplifies floorplanning, particularly in designs with transceivers, by placing per-side logic near the transceivers.

When adding register stages to pipeline control signals, turn off the Auto Shift Register Replacement option (Assignments > Settings > Analysis & Synthesis Settings > More Settings) for these registers. By default, chains of registers can be converted to a RAM-based implementation based on performance and resource estimates. Since pipelining helps meet timing requirements over long distance, this assignment ensures that control signals are not converted.

**Planning FPGA Resources**

The requirements of your design affect the use of FPGA resources. Plan functional blocks with appropriate global, regional, and dual-regional network signals in mind. In general, after allocating the clocks in a design, use global networks for the highest fan-out control signals. When a global network signal distributes a high fan-out control signal, the global signal can drive logic anywhere in the device. Similarly, when using a regional network signal, the driven must be in one quadrant of the device, or half the device for a dual-regional network signal. Depending on data flow and physical locations of the data entry and exit between the I/Os and the device, restricting a functional block to a quadrant or half the device may not be practical for performance or resource requirements.

When floorplanning a design, consider the balance of different types of device resources, such as memory, logic, and DSP blocks in the main functional blocks. For example, if a design is memory intensive with a small amount of logic, it may be difficult to develop an effective floorplan. Logic that interfaces with the memory would have to spread across the chip to access the memory. In this case, it is important to use enough register stages in the data and control paths to allow signals to traverse the chip to access the physically disparate resources needed.

**Optimizing for Timing Closure**

You can make changes to your design and constraints that help you achieve timing closure. Whenever you change the project settings, you must balance any performance improvement of the setting against any potential increase in compilation time associated with the setting. You can view the performance gain versus runtime cost by reviewing the Fitter messages after design processing.

**Physical Synthesis Optimization**

You can use physical synthesis optimizations for combinational logic, register retiming, and register duplication techniques to optimize your design for timing closure. Click Assignments > Settings > Physical Synthesis Optimizations to turn on physical synthesis options.
Physical synthesis for combinational logic—When the **Perform physical synthesis for combinational logic** is turned on, the report panel identifies logic that physical synthesis can modify. You can use this information to modify the design so that the associated optimization can be turned off to save compile time.

Register duplication—This technique is most useful where registers have high fan-out, or where the fan-out is in physically distant areas of the device. Review the netlist optimizations report and consider manually duplicating registers automatically added by physical synthesis. You can also locate the original and duplicate registers in the Chip Planner. Compare their locations, and if the fan-out is improved, modify the code and turn off register duplication to save compile time.

Register retiming—This technique is particularly useful where some combinatorial paths between registers exceed the timing goal while other paths fall short. If a design is already heavily pipelined, register retiming is less likely to provide significant performance gains since there should not be significantly unbalanced levels of logic across pipeline stages.

**Timing Constraint Optimization**

The application of appropriate timing constraints is essential to timing closure. Use the following general guidelines in applying timing constraints:

- Apply multicycle constraints in your design wherever single-cycle timing analysis is not required.

- Apply False Path constraints to all asynchronous clock domain crossings or resets in the design. This technique prevents overconstraining and the Fitter focuses only on critical paths to reduce compile time. However, over constraining timing critical clock domains can sometimes provide better timing results and lower compile times than physical synthesis.

- Overconstrain rather than using physical synthesis when the slack improvement from physical synthesis is near zero. Overconstrain the frequency requirement on timing critical clock domains by using setup uncertainty.

- When evaluating the effect of constraint changes on performance and runtime, compile the design with at least three different seeds to determine the average performance and runtime effects. Different constraint combinations produce various results. Three samples or more establishes a performance trend. Modify your constraints based on performance improvement or decline.

- Leave settings at the default value whenever possible. Increasing performance constraints can increase the compile time significantly. While those increases may be necessary to close timing on a design, using the default settings whenever possible minimizes compile time.

**Optimizing Critical Timing Paths**

To close timing in high speed designs, review paths with the largest timing failures. Correcting a single, large timing failure can result in a very significant timing improvement. Review the register placement and routing paths by clicking **Tools > Chip Planner**. Large timing failures on high fan-out control signals can be caused by any of the following conditions:

- Sub-optimal use of global networks
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- Signals that traverse the chip on local routing without pipelining
- Failure to correct high fan-out by register duplication

For high-speed and high-bandwidth designs, optimize speed by reducing bus width and wire usage. To reduce wire use, move the data as little as possible. For example, if a block of logic functions on a few bits of a word, store inactive bits in a fifo or memory. Memory is cheaper and denser than registers and reduces wire usage.

Power Optimization

The total FPGA power consumption is comprised of I/O power, core static power, and core dynamic power. Knowledge of the relationship between these components is fundamental in calculating the overall total power consumption. You can use various optimization techniques and tools to minimize power consumption when applied during FPGA design implementation. The Quartus II software offers power-driven compilation features to fully optimize device power consumption. Power-driven compilation focuses on reducing your design’s total power consumption using power-driven synthesis and power-driven placement and routing.

For more information about power-driven compilation flow and low-power design guidelines, refer to the Power Optimization chapter in volume 2 of the Quartus II Handbook.

For more information about power optimization techniques available for Stratix III devices, refer to AN 437: Power Optimization in Stratix III FPGAs. For more information about power optimization techniques available for Stratix IV devices, refer to AN 514: Power Optimization in Stratix IV FPGAs. For more information about power optimization techniques available for Stratix V devices, refer to Reducing Power Consumption and Increasing Bandwidth on 28-nm FPGAs white paper.

Additionally, you can use the Quartus II PowerPlay suite of power analysis and optimization tools to help you during the design process by delivering fast and accurate estimations of power consumption. For more information about the Quartus II PowerPlay suite of power analysis and optimization tools, refer to About Power Estimation and Analysis in Quartus II Help.

Metastability

Metastability in PLD designs can be caused by the synchronization of asynchronous signals. You can use the Quartus II software to analyze the mean time between failures (MTBF) due to metastability, thus optimizing the design to improve the metastability MTBF. A high metastability MTBF indicates a more robust design.

For more information about how to ensure complete and accurate metastability analysis, refer to the Managing Metastability With the Quartus II Software chapter in volume 1 of the Quartus II Handbook.

For more information about viewing metastability reports, refer to Viewing Metastability Reports in Quartus II Help.
Incremental Compilation

The incremental compilation feature in the Quartus II software allows you to partition your design hierarchy, separately compile partitions, and reuse the results for unchanged partitions. Incremental compilation flows require more up-front planning than flat compilations, and generally require you to be more rigorous about following good design practices than flat compilations.

For more information about incremental compilation and floorplan assignments, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

For more information about incremental compilation, refer to About Incremental Compilation in Quartus II Help.

Checking Design Violations

To improve the reliability, timing performance, and logic utilization of your design, you should practice good design methodology and understand how to avoid design rule violations. The Quartus II software provides the Design Assistant tool that automatically checks for design rule violations and reports their location.

The Design Assistant is a design rule checking tool that allows you to check for design issues early in the design flow. The Design Assistant checks your design for adherence to Altera-recommended design guidelines. You can specify which rules you want the Design Assistant to apply to your design. This is useful if you know that your design violates particular rules that are not critical and you can allow these rule violations. The Design Assistant generates design violation reports with details about each violation based on the settings that you specified.

This section provides an introduction to the Quartus II design flow with the Design Assistant, message severity levels, and an explanation about how to set up the Design Assistant. The last parts of the section describe the design rules and the reports generated by the Design Assistant. The Design Assistant supports all Altera devices supported by the Quartus II software.

Quartus II Design Flow with the Design Assistant

You can run the Design Assistant after Analysis and Elaboration, Analysis and Synthesis, fitting, or a full compilation. If you set the Design Assistant to run automatically during compilation, the Design Assistant performs a post-fitting netlist analysis of your design. The default is to apply all of the rules to your project. If there are some rules that are unimportant to your design, you can turn off the rules that you do not want the Design Assistant to use.

For more information about running the Design Assistant, refer to About the Design Assistant in Quartus II Help.
Figure 13–9 shows the Quartus II software design flow with the Design Assistant.

Figure 13–9. Quartus II Design Flow with the Design Assistant

The Design Assistant analyzes your design netlist at different stages of the compilation flow and may yield different warnings or errors, even though the netlists are functionally the same. Your pre-synthesis, post-synthesis, and post-fitting netlists might be different due to optimizations performed by the Quartus II software. For example, a warning message in a pre-synthesis netlist may be removed after the netlist has been synthesized into a post-synthesis or post-fitting netlist.

The exact operation of the Design Assistant depends on when you run it:

■ When you run the Design Assistant after running a full compilation or fitting, the Design Assistant performs a post-fitting analysis on the design.
■ When you run the Design Assistant after performing Analysis and Synthesis, the Design Assistant performs post-synthesis analysis on the design.
■ When you start the Design Assistant after performing Analysis and Elaboration, the Design Assistant performs a pre-synthesis analysis on the design. You can also perform pre-synthesis analysis with the Design Assistant using the command-line. You can use the -rtl option with the quartus_drc executable, as shown in the following example:

\`\`quartus_drc <project_name> --rtl=on \`

For more information about Design Assistant settings, refer to About the Design Assistant and Design Assistant Page (Settings Dialog Box) in Quartus II Help.
Enabling and Disabling Design Assistant Rules

For more information about enabling or disabling Design Assistant rules on individual nodes by making an assignment in the Assignment Editor, in the Quartus II Settings File (.qsf), with the altera_attribute synthesis attribute in Verilog HDL or VHDL, or with a Tcl command, refer to Enabling Design Assistant Rules on Nodes, Entities, or Instances, or Disabling Design Assistant Rules on Nodes, Entities, or Instances in Quartus II Help.

Viewing Design Assistant Results

If your design violates a design rule, the Design Assistant generates warning messages and information messages about the violated rule. The Design Assistant displays these messages in the Messages window, in the Design Assistant Messages report, and in the Design Assistant report files. You can find the Design Assistant report files called <project_name>.drc.rpt in the <project_name> subdirectory of the project directory.

For information about the contents of the reports generated by the Design Assistant, refer to Design Assistant Reports in Quartus II Help.

Custom Rules

In addition to the existing design rules that the Design Assistant offers, you can also create your own rules and specify your own reporting format in a text file (with any file extension) with the XML format. You then specify the path to that file in the Design Assistant settings page and run the Design Assistant for violation checking.

Refer to the following location to locate the file that contains the default rules for the Design Assistant:

<Quartus II install path>\quartus\libraries\design-assistant\da_golden_rule.xml

For more information about how to set the file path to your custom rules, refer to Custom Rules Settings Dialog Box in Quartus II Help. For more information about the basics of writing custom rules, the Design Assistant settings, and coding examples on how to check for clock relationship and node relationship in a design, refer to Creating Custom Design Assistant Rules in Quartus II Help. To specify the rules that you want the Design Assistant to use when checking for violations, refer to Design Assistant Page (Settings Dialog Box) in Quartus II Help.

Custom Rules Coding Examples

The following examples of custom rules show how to check node relationships and clock relationships in a design.
Checking SR Latch Structures In a Design

Example 13–1 shows the XML codes for checking SR latch structures in a design.

Example 13–1. Detecting SR Latches in a Design

```xml
<DA_RULE ID="EX01" SEVERITY="CRITICAL" NAME="Checking Design for SR Latch"
DEFAULT_RUN="YES">
  <RULE_DEFINITION>
    <FORBID>
      <OR>
        <NODE NAME="NODE_1" TYPE="SRLATCH" />
        <HAS_NODE NODE_LIST="NODE_1" />
        <NODE NAME="NODE_1" TOTAL_FANIN="EQ2" />
        <NODE NAME="NODE_2" TOTAL_FANIN="EQ2" />
      </OR>
      <AND>
        <NODE_RELATIONSHIP FROM_NAME="NODE_1" FROM_TYPE="NAND" TO_NAME="NODE_2"
TO_TYPE="NAND" />
        <NODE_RELATIONSHIP FROM_NAME="NODE_2" FROM_TYPE="NAND" TO_NAME="NODE_1"
TO_TYPE="NAND" />
      </AND>
      <OR>
        <NODE_RELATIONSHIP FROM_NAME="NODE_1" FROM_TYPE="NAND" TO_NAME="NODE_2"
TO_TYPE="NAND" />
        <NODE_RELATIONSHIP FROM_NAME="NODE_2" FROM_TYPE="NAND" TO_NAME="NODE_1"
TO_TYPE="NAND" />
      </OR>
    </FORBID>
  </RULE_DEFINITION>
  <REPORTING_ROOT>
    <MESSAGE NAME="Rule %ARG1%: Found %ARG2% node(s) related to this rule.">
      <MESSAGE_ARGUMENT NAME="ARG1" TYPE="ATTRIBUTE" VALUE="ID" />
      <MESSAGE_ARGUMENT NAME="ARG2" TYPE="TOTAL_NODE" VALUE="NODE_1" />
    </MESSAGE>
  </REPORTING_ROOT>
</DA_RULE>
```

In Example 13–1, the possible SR latch structures are specified in the rule definition section. Codes defined in the <AND></AND> block are tied together, meaning that each statement in the block must be true for the block to be fulfilled (AND gate similarity). In the <OR></OR> block, as long as one statement in the block is true, the block is fulfilled (OR gate similarity). If no <AND></AND> or <OR></OR> blocks are specified, the default is <AND></AND>.

The <FORBID></FORBID> section contains the undesirable condition for the design, which in this case is the SR latch structures. If the condition is fulfilled, the Design Assistant highlights a rule violation.

The following examples are the undesired conditions from Example 13–1 with their equivalent block diagrams (Figure 13–10 and Figure 13–11):

```xml
<AND>
  <NODE_RELATIONSHIP FROM_NAME="NODE_1" FROM_TYPE="NAND" TO_NAME="NODE_2"
TO_TYPE="NAND" />
  <NODE_RELATIONSHIP FROM_NAME="NODE_2" FROM_TYPE="NAND" TO_NAME="NODE_1"
TO_TYPE="NAND" />
</AND>
```
Figure 13–10. Undesired Condition 1

<AND>
<NODE_RELATIONSHIP FROM_NAME="NODE_1" FROM_TYPE="NOR" TO_NAME="NODE_2" TO_TYPE="NOR"/>
<NODE_RELATIONSHIP FROM_NAME="NODE_2" FROM_TYPE="NOR" TO_NAME="NODE_1" TO_TYPE="NOR"/>
</AND>

Figure 13–11. Undesired Condition 2

<AND>
<NODE_RELATIONSHIP FROM_NAME="NODE_1" FROM_TYPE="NOR" TO_NAME="NODE_2" TO_TYPE="NOR"/>
<NODE_RELATIONSHIP FROM_NAME="NODE_2" FROM_TYPE="NOR" TO_NAME="NODE_1" TO_TYPE="NOR"/>
</AND>
Relating Nodes to a Clock Domain

Example 13–2 shows how to use the CLOCK RELATIONSHIP attribute to relate nodes to clock domains. This example checks for correct synchronization in data transfer between asynchronous clock domains. Synchronization is done with cascaded registers, also called synchronizers, at the receiving clock domain. The code in Example 13–2 checks for the synchronizer configuration based on the following guidelines:

- The cascading registers need to be triggered on the same clock edge
- There is no logic between the register output of the transmitting clock domain and the cascaded registers in the receiving asynchronous clock domain

Example 13–2. Detecting Incorrect Synchronizer Configuration

```xml
<DA_RULE ID="EX02" SEVERITY="HIGH" NAME="Data Transfer Not Synch Correctly" DEFAULT_RUN="YES">

<RULE_DEFINITION>
<DECLARE>
  <NODE NAME="NODE_1" TYPE="REG" />
  <NODE NAME="NODE_2" TYPE="REG" />
  <NODE NAME="NODE_3" TYPE="REG" />
</DECLARE>

<FORBID>
  <NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
    CLOCK_RELATIONSHIP="ASYN" />
  <NODE_RELATIONSHIP FROM_NAME="NODE_2" TO_NAME="NODE_3" TO_PORT="D_PORT"
    CLOCK_RELATIONSHIP="!ASYN" />
  <OR>
    <NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
      REQUIRED_THROUGH="YES" THROUGH_TYPE="COMB" CLOCK_RELATIONSHIP="ASYN" />
    <CLOCK_RELATIONSHIP NAME="SEQ_EDGE|ASYN" NODE_LIST="NODE_2, NODE_3" />
  </OR>
</FORBID>

<REPORTING_ROOT>
  <MESSAGE NAME="Rule %ARG1%: Found %ARG2% node(s) related to this rule.">
    <MESSAGE_ARGUMENT NAME="ARG1" TYPE="ATTRIBUTE" VALUE="ID" />
    <MESSAGE_ARGUMENT NAME="ARG2" TYPE="TOTAL_NODE" VALUE="NODE_1" />
  </MESSAGE>
  <MESSAGE NAME="Source node(s): %ARG3%, Destination node(s): %ARG4%">
    <MESSAGE_ARGUMENT NAME="ARG3" TYPE="NODE" VALUE="NODE_1" />
    <MESSAGE_ARGUMENT NAME="ARG4" TYPE="NODE" VALUE="NODE_2" />
  </MESSAGE>
</REPORTING_ROOT>
</DA_RULE>
```

The codes differentiate the clock domains. ASYN means asynchronous, and !ASYN means non-asynchronous. This notation is useful for describing nodes that are in different clock domains. The following lines from Example 13–2 state that NODE_2 and NODE_3 are in the same clock domain, but NODE_1 is not.

```xml
<NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
  CLOCK_RELATIONSHIP="ASYN" />

<NODE_RELATIONSHIP FROM_NAME="NODE_2" TO_NAME="NODE_3" TO_PORT="D_PORT"
  CLOCK_RELATIONSHIP="!ASYN" />
```
The next line of code states that `NODE_2` and `NODE_3` have a clock relationship of either sequential edge or asynchronous.

```xml
<CLOCK_RELATIONSHIP NAME="SEQ_EDGE|ASYN" NODE_LIST="NODE_2, NODE_3" />
```

The `<FORBID>` section contains the undesirable condition for the design, which in this case is the undesired configuration of the synchronizer. If the condition is fulfilled, the Design Assistant highlights a rule violation.

The following examples are the undesired conditions from Example 13–2 with their equivalent block diagrams (Figure 13–12 and Figure 13–13):

**Example 13–3.**

```xml
<NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
                      CLOCK_RELATIONSHIP="ASYN" />

<NODE_RELATIONSHIP FROM_NAME="NODE_2" TO_NAME="NODE_3" TO_PORT="D_PORT"
                      CLOCK_RELATIONSHIP="!ASYN" />

<NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
                      REQUIRED_THROUGH="YES" THROUGH_TYPE="COMB"
                      CLOCK_RELATIONSHIP="ASYN" />
```

**Figure 13–12. Undesired Condition 3**

![Figure 13–12. Undesired Condition 3](image)

**Example 13–4.**

```xml
<NODE_RELATIONSHIP FROM_NAME="NODE_1" TO_NAME="NODE_2" TO_PORT="D_PORT"
                      CLOCK_RELATIONSHIP="ASYN" />

<NODE_RELATIONSHIP FROM_NAME="NODE_2" TO_NAME="NODE_3" TO_PORT="D_PORT"
                      CLOCK_RELATIONSHIP="!ASYN" />

<CLOCK_RELATIONSHIP NAME="SEQ_EDGE|ASYN" NODE_LIST="NODE_2, NODE_3" />
```

**Figure 13–13. Undesired Condition 4**

![Figure 13–13. Undesired Condition 4](image)
Targeting Clock and Register-Control Architectural Features

In addition to following general design guidelines, you must code your design with the device architecture in mind. FPGAs provide device-wide clocks and register control signals that can improve performance.

Clock Network Resources

Altera FPGAs provide device-wide global clock routing resources and dedicated inputs. Use the FPGA’s low-skew, high fan-out dedicated routing where available. By assigning a clock input to one of these dedicated clock pins or with a Quartus II logic option to assign global routing, you can take advantage of the dedicated routing available for clock signals.

In an ASIC design, you should balance the clock delay as it is distributed across the device. Because Altera FPGAs provide device-wide global clock routing resources and dedicated inputs, there is no need to manually balance delays on the clock network.

You should limit the number of clocks in your design to the number of dedicated global clock resources available in your FPGA. Clocks feeding multiple locations that do not use global routing may exhibit clock skew across the device that could lead to timing problems. In addition, when you use combinational logic to generate an internal clock, it adds delays on the clock path. In some cases, delay on a clock line can result in a clock skew greater than the data path length between two registers. If the clock skew is greater than the data delay, you violate the timing parameters of the register (such as hold time requirements) and the design does not function correctly.

FPGAs offer a number of low-skew global routing resources to distribute high fan-out signals to help with the implementation of large designs with many clock domains. Many large FPGA devices provide dedicated global clock networks, regional clock networks, and dedicated fast regional clock networks. These clocks are organized into a hierarchical clock structure that allows many clocks in each device region with low skew and delay. There are typically several dedicated clock pins to drive either global or regional clock networks, and both PLL outputs and internal clocks can drive various clock networks.

To reduce clock skew in a given clock domain and ensure that hold times are met in that clock domain, assign each clock signal to one of the global high fan-out, low-skew clock networks in the FPGA device. The Quartus II software automatically uses global routing for high fan-out control signals, PLL outputs, and signals feeding the global clock pins on the device. You can make explicit Global Signal logic option settings by turning on the Global Signal option setting. Use this option when it is necessary to force the software to use the global routing for particular signals.

To take full advantage of these routing resources, the sources of clock signals in a design (input clock pins or internally-generated clocks) need to drive only the clock input ports of registers. In older Altera device families, if a clock signal feeds the data ports of a register, the signal may not be able to use dedicated routing, which can lead to decreased performance and clock skew problems. In general, allowing clock signals to drive the data ports of registers is not considered synchronous design and can complicate timing analysis.
**Reset Resources**

ASIC designs may use local resets to avoid long routing delays. Take advantage of the device-wide asynchronous reset pin available on most FPGAs to eliminate these problems. This reset signal provides low-skew routing across the device.

The following are three types of resets used in synchronous circuits:

- Synchronous Reset
- Asynchronous Reset
- Synchronized Asynchronous Reset—preferred when designing an FPGA circuit

**Synchronous Reset**

The synchronous reset ensures that the circuit is fully synchronous. You can easily time the circuit with the Quartus II TimeQuest analyzer. Because clocks that are synchronous to each other launch and latch the reset signal, the data arrival and data required times are easily determined for proper slack analysis. The synchronous reset is easier to use with cycle-based simulators.

There are two methods by which a reset signal can reach a register; either by being gated in with the data input, as shown in Figure 13–14, or by using an LAB-wide control signal (\(\text{sync1r}\)), as shown in Figure 13–15. If you use the first method, you risk adding an additional gate delay to the circuit to accommodate the reset signal, which causes increased data arrival times and negatively impacts setup slack. The second method relies on dedicated routing in the LAB to each register, but this is slower than an asynchronous reset to the same register.

**Figure 13–14. Synchronous Reset**

![Figure 13–14. Synchronous Reset](image)
Consider two types of synchronous resets when you examine the timing analysis of synchronous resets—externally synchronized resets and internally synchronized resets. Externally synchronized resets are synchronized to the clock domain outside the FPGA, and are not very common. A power-on asynchronous reset is dual-rank synchronized externally to the system clock and then brought into the FPGA. Inside the FPGA, gate this reset with the data input to the registers to implement a synchronous reset.
Figure 13–16 shows the schematic for an externally synchronized reset.

Example 13–5 shows the Verilog equivalent of the schematic. When you use synchronous resets, the reset signal is not put in the sensitivity list.

Example 13–5. Verilog Code for Externally Synchronized Reset

```verilog
module sync_reset_ext (  
    input clock,  
    input reset_n,  
    input data_a,  
    input data_b,  
    output out_a,  
    output out_b  
);  
    reg reg1, reg2  
    assign out_a = reg1;  
    assign out_b = reg2;  
    always @ (posedge clock)  
    begin  
        if (!reset_n)  
            begin  
                reg1 <= 1'bo;  
                reg2 <= 1'b0;  
            end  
        else  
            begin  
                reg1 <= data_a;  
                reg2 <= data_b;  
            end  
    end  
endmodule // sync_reset_ext
```
Example 13–6 shows the constraints for the externally synchronous reset. Because the external reset is synchronous, you only need to constrain the reset_n signal as a normal input signal with set_input_delay constraint for -max and -min.

**Example 13–6. SDC Constraints for Externally Synchronous Reset**

```plaintext
# Input clock - 100 MHz
create_clock [get_ports {clock}] \
  -name {clock} \
  -period 10.0 \
  -waveform {0.0 5.0}

# Input constraints on low-active reset 
# and data
set_input_delay 7.0 \  
  -max \  
  -clock [get_clocks {clock}] \
  [get_ports {reset_n data_a data_b}]  
set_input_delay 1.0 \  
  -min \  
  -clock [get_clocks {clock}] \
  [get_ports {reset_n data_a data_b}]
```

More often, resets coming into the device are asynchronous, and must be synchronized internally before being sent to the registers. Figure 13–17 shows an internally synchronized reset.

**Figure 13–17. Internally Synchronized Reset**
Example 13–7 shows the Verilog equivalent of the schematic. Only the clock edge is in the sensitivity list for a synchronous reset.

Example 13–7. Verilog Code for Internally Synchronous Reset

```verilog
module sync_reset ( 
    input clock, 
    input reset_n, 
    input data_a, 
    input data_b, 
    output out_a, 
    output out_b 
); 
reg reg1, reg2 
reg reg3, reg4 
assign out_a = reg1; 
assign out_b = reg2; 
assign rst_n = reg4; 
always @ (posedge clock) 
begin 
    if (!rst_n) 
        begin 
            reg1 <= 1’bo; 
            reg2 <= 1’b0; 
        end 
    else 
        begin 
            reg1 <= data_a; 
            reg2 <= data_b; 
        end 
end 
always @ (posedge clock) 
begin 
    reg3 <= reset_n; 
    reg4 <= reg3; 
end 
endmodule // sync_reset
```
The SDC constraints are similar to the external synchronous reset, except that the input reset cannot be constrained because it is asynchronous and should be cut with a `set_false_path` statement (as shown in Example 13–8) to avoid these being considered as unconstrained paths.

**Example 13–8. SDC Constraints for Internally Synchronized Reset**

```asciidoc
# Input clock - 100 MHz
create_clock [get_ports {clock}] 
   -name {clock} 
   -period 10.0 
   -waveform {0.0 5.0}

# Input constraints on data
set_input_delay 7.0 
   -max 
   -clock [get_clocks {clock}] 
   [get_ports {data_a data_b}]
set_input_delay 1.0 
   -min 
   -clock [get_clocks {clock}] 
   [get_ports {data_a data_b}]

# Cut the asynchronous reset input
set_false_path 
   -from [get_ports {reset_n}] 
   -to [all_registers]
```

An issue with synchronous resets is their behavior with respect to short pulses (less than a period) on the asynchronous input to the synchronizer flipflops. This can be a disadvantage because the asynchronous reset requires a pulse width of at least one period wide to guarantee that it is captured by the first flipflop. However, this can also be viewed as an advantage in that this circuit increases noise immunity. Spurious pulses on the asynchronous input have a lower chance of being captured by the first flipflop, so the pulses do not trigger a synchronous reset. In some cases, you might want to increase the noise immunity further and reject any asynchronous input reset that is less than n periods wide to debounce an asynchronous input reset.
Figure 13–18 shows the necessary modifications that you should make to the internally synchronized reset.

**Figure 13–18. Internally Synchronized Reset with Pulse Extender**

Many designs have more than one clock signal. In these cases, use a separate reset synchronization circuit for each clock domain in the design. When you create synchronizers for PLL output clocks, these clock domains are not reset until you lock the PLL and the PLL output clocks are stable. If you use the reset to the PLL, this reset does not have to be synchronous with the input clock of the PLL. You can use an asynchronous reset for this. Using a reset to the PLL further delays the assertion of a synchronous reset to the PLL output clock domains when using internally synchronized resets.

**Asynchronous Reset**

Asynchronous resets are the most common form of reset in circuit designs, as well as the easiest to implement. Typically, you can insert the asynchronous reset into the device, turn on the global buffer, and connect to the asynchronous reset pin of every register in the device. This method is only advantageous under certain circumstances—you do not need to always reset the register. Unlike the synchronous reset, the asynchronous reset is not inserted in the data path, and does not negatively impact the data arrival times between registers. Reset takes effect immediately, and as soon as the registers receive the reset pulse, the registers are reset. The asynchronous reset is not dependent on the clock.

However, when the reset is deasserted and does not pass the recovery (μtSU) or removal (μtH) time check (the TimeQuest analyzer recovery and removal analysis checks both times), the edge is said to have fallen into the metastability zone. Additional time is required to determine the correct state, and the delay can cause the setup time to fail to register downstream, leading to system failure. To avoid this, add a few follower registers after the register with the asynchronous reset and use the output of these registers in the design. Use the follower registers to synchronize the
data to the clock to remove the metastability issues. You should place these registers close to each other in the device to keep the routing delays to a minimum, which decreases data arrival times and increases MTBF. Ensure that these follower registers themselves are not reset, but are initialized over a period of several clock cycles by “flushing out” their current or initial state.

Figure 13–19 shows a schematic example of this circuit.

Example 13–9 shows the equivalent Verilog code. The active edge of the reset is now in the sensitivity list for the procedural block, which infers a clock enable on the follower registers with the inverse of the reset signal tied to the clock enable. The follower registers should be in a separate procedural block as shown using non-blocking assignments.

Example 13–9. Verilog Code of Asynchronous Reset with Follower Registers

```verilog
module async_reset (input clock, input reset_n, input data_a, output out_a);

reg reg1, reg2, reg3;
assign out_a = reg3;

always @ (posedge clock, negedge reset_n)
begin
  if (!reset_n)
    reg1 <= 1'b0;
  else
    reg1 <= data_a;
end

always @ (posedge clock)
begin
  reg2 <= reg1;
  reg3 <= reg2;
end
endmodule // async_reset
```
You can easily constrain an asynchronous reset. By definition, asynchronous resets have a non-deterministic relationship to the clock domains of the registers they are resetting. Therefore, static timing analysis of these resets is not possible and you can use the set_false_path command to exclude the path from timing analysis (as shown in Example 13–10). Because the relationship of the reset to the clock at the register is not known, you cannot run recovery and removal analysis in the TimeQuest analyzer for this path. Attempting to do so even without the false path statement results in no paths reported for recovery and removal.

Example 13–10. SDC Constraints for Asynchronous Reset

```verbatim
# Input clock - 100 MHz
create_clock [get_ports {clock}] \
   -name {clock} \n   -period 10.0 \n   -waveform {0.0 5.0}

# Input constraints on data
set_input_delay 7.0 \n   -max \n   -clock [get_clocks {clock}]\n   [get_ports {data_a}]
set_input_delay 1.0 \n   -min \n   -clock [get_clocks {clock}] \n   [get_ports {data_a}]

# Cut the asynchronous reset input
set_false_path \n   -from [get_ports {reset_n}] \n   -to [all_registers]
```
Figure 13–20 shows a method for implementing the synchronized asynchronous reset. You should use synchronizer registers in a similar manner as synchronous resets. However, the asynchronous reset input is gated directly to the CLRN pin of the synchronizer registers and immediately asserts the resulting reset. When the reset is deasserted, logic “1” is clocked through the synchronizers to synchronously deassert the resulting reset.
Example 13–11 shows the equivalent Verilog code. Use the active edge of the reset in the sensitivity list for the blocks in Figure 13–20.

**Example 13–11. Verilog Code for Synchronized Asynchronous Reset**

```verilog
module sync_async_reset (  
    input clock,  
    input reset_n,  
    input data_a,  
    input data_b,  
    output out_a,  
    output out_b  
);

reg reg1, reg2;
reg reg3, reg4;
assign out_a = reg1;
assign out_b = reg2;
assign rst_n = reg4;
always @ (posedge clock, negedge reset_n)
begin
    if (!reset_n)
        begin
            reg3 <= 1'b0;
            reg4 <= 1'b0;
        end
    else
        begin
            reg3 <= 1'b1;
            reg4 <= reg3;
        end
end
always @ (posedge clock, negedge rst_n)
begin
    if (!rst_n)
        begin
            reg1 <= 1'b0;
            reg2 <= 1'b0;
        end
    else
        begin
            reg1 <= data_a;
            reg2 <= data_b;
        end
end
endmodule // sync_async_reset
```

To minimize the metastability effect between the two synchronization registers, and to increase the MTBF, the registers should be located as close as possible in the device to minimize routing delay. If possible, locate the registers in the same logic array block (LAB). The input reset signal (reset_n) must be excluded with a `set_false_path` command, so the reset that comes from the synchronization register (rst_n) can be timed in the TimeQuest analyzer with recovery and removal Analysis.
The instantaneous assertion of synchronized asynchronous resets is susceptible to noise and runt pulses. If possible, you should debounce the asynchronous reset and filter the reset before it enters the device. The circuit in Figure 13–20 on page 13–33 ensures that the synchronized asynchronous reset is at least one full clock period in length. To extend this time to n clock periods, you must increase the number of synchronizer registers to n + 1. You must connect the asynchronous input reset (reset_n) to the CLRN pin of all the synchronizer registers to maintain the asynchronous assertion of the synchronized asynchronous reset.

For more information about specifying the minimum routing delay, refer to the Best Practices for the Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.

Register Control Signals

Avoid using an asynchronous load signal if the design target device architecture does not include registers with dedicated circuitry for asynchronous loads. Also, avoid using both asynchronous clear and preset if the architecture provides only one of these control signals. Stratix III devices, for example, directly support an asynchronous clear function, but not a preset or load function. When the target device does not directly support the signals, the synthesis or placement and routing software must use combinational logic to implement the same functionality. In addition, if you use signals in a priority other than the inherent priority in the device architecture, combinational logic may be required to implement the necessary control signals. Combinational logic is less efficient and can cause glitches and other problems; it is best to avoid these implementations.

For Verilog HDL and VHDL examples of registers with various control signals, and information about the inherent priority order of register control signals in Altera device architecture, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

Targeting Embedded RAM Architectural Features

Altera’s dedicated memory architecture offers many advanced features that you can target easily with the MegaWizard™ Plug-In Manager or with the recommended HDL coding styles that infer the appropriate RAM megafunction (ALTSYNCRAM or ALTDPRAM). Use synchronous memory blocks for your design, so that the blocks can be mapped directly into the device dedicated memory blocks. You can use single-port, dual-port, or three-port RAM with a single- or dual-clocking method. You should not infer the asynchronous memory logic as a memory block or place the asynchronous memory logic in the dedicated memory block, but implement the asynchronous memory logic in regular logic cells.

Altera memory blocks have different read-during-write behaviors, depending on the targeted device family, memory mode, and block type. Read-during-write behavior refers to read and write from the same memory address in the same clock cycle; for example, you read from the same address to which you write in the same clock cycle.
You should check how you specify the memory in your HDL code when you use read-during-write behavior. The HDL code that describes the read returns either the old data stored at the memory location, or the new data being written to the memory location.

In some cases, when the device architecture cannot implement the memory behavior described in your HDL code, the memory block is not mapped to the dedicated RAM blocks, or the memory block is implemented using extra logic in addition to the dedicated RAM block. Implement the read-during-write behavior using single-port RAM in Arria GX devices and the Cyclone and Stratix series of devices to avoid this extra logic implementation.

For Verilog HDL and VHDL examples and guidelines for inferring RAM functions that match the dedicated memory architecture in Altera devices, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

In many synthesis tools, you can specify that the read-during-write behavior is not important to your design; if, for example, you never read and write from the same address in the same clock cycle. For Quartus II integrated synthesis, add the synthesis attribute `ramstyle="no_rw_check"` to allow the software to choose the read-during-write behavior of a RAM, rather than using the read-during-write behavior specified in your HDL code. Using this type of attribute prevents the synthesis tool from using extra logic to implement the memory block and, in some cases, can allow memory inference when it would otherwise be impossible.

For details about using the `ramstyle` attribute, refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook. For information about the synthesis attributes in other synthesis tools, refer to your synthesis tool documentation, or to the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

## Conclusion

Following the design practices described in this chapter can help you to consistently meet your design goals. Asynchronous design techniques may result in incomplete timing analysis, may cause glitches on data signals, and may rely on propagation delays in a device leading to race conditions and unpredictable results. Taking advantage of the architectural features in your FPGA device can also improve the quality of your results.

### Document Revision History

Table 13–1 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>▪ Added “Optimizing for Physical Implementation and Timing Closure” section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Removed PrimeTime support.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
</tr>
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</table>
### Table 13–1. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Added information to “Reset Resources” on page 13–24.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Removed documentation of obsolete rules.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>No change to content.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8-1/2 x 11 page size</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Updated Figure 5–9 on page 5–13; added custom rules file to the flow</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added notes to Figure 5–9 on page 5–13</td>
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<tr>
<td></td>
<td></td>
<td>Added new section: “Custom Rules Report” on page 5–34</td>
</tr>
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<td></td>
<td></td>
<td>Added new section: “Custom Rules” on page 5–34</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added new section: “Targeting Embedded RAM Architectural Features” on</td>
</tr>
<tr>
<td></td>
<td></td>
<td>page 5–38</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Minor editorial updates throughout the chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added hyperlinks to referenced documents throughout the chapter</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides Hardware Description Language (HDL) coding style recommendations to ensure optimal synthesis results when targeting Altera® devices. HDL coding styles can have a significant effect on the quality of results that you achieve for programmable logic designs. Synthesis tools optimize HDL code for both logic utilization and performance; however, synthesis tools have no information about the purpose or intent of the design. The best optimizations often require conscious interaction by you, the designer.

This chapter includes the following sections:

- “Using the Quartus II Templates” on page 14–2
- “Using Altera Megafunctions” on page 14–3
- “Instantiating Altera Megafunctions in HDL Code” on page 14–3
- “Inferring Multiplier and DSP Functions from HDL Code” on page 14–5
- “Inferring Memory Functions from HDL Code” on page 14–13
- “Coding Guidelines for Registers and Latches” on page 14–44
- “General Coding Guidelines” on page 14–54
- “Designing with Low-Level Primitives” on page 14–74

For additional guidelines about structuring your design, refer to the Design Recommendations for Altera Devices and the Quartus II Design Assistant chapter in volume 1 of the Quartus II Handbook. For additional handcrafted techniques you can use to optimize design blocks for the adaptive logic modules (ALMs) in many Altera devices, including a collection of circuit building blocks and related discussions, refer to the Advanced Synthesis Cookbook.

The Altera website also provides design examples for other types of functions and to target specific applications. For more information about design examples, refer to the Design Examples page and the Reference Designs page on the Altera website.

For style recommendations, options, or HDL attributes specific to your synthesis tool (including Quartus® II integrated synthesis and other EDA tools), refer to the tool vendor’s documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.
Using the Quartus II Templates

Many of the Verilog HDL and VHDL examples in this chapter correspond with examples in the Full Designs section of the Quartus II Templates. You can easily insert examples into your HDL source code with the Quartus II Text Editor, in the Quartus II software, or with your preferred text editor.

Inserting a Template with the Quartus II Text Editor

To use a template with the Quartus II software, follow these steps:

1. On the File menu, click New.
2. In the New dialog box, select the type of design file corresponding to the type of HDL you want to use, SystemVerilog HDL File, VHDL File, or Verilog HDL File.
3. Click the Insert Template button on the text editor menu, or, right-click in the blank Verilog or VHDL file, then click Insert Template.
4. In the Insert Template dialog box shown in Figure 14–1, expand the section corresponding to the appropriate HDL, then expand the Full Designs section.
5. Expand the design category, for example RAMs and ROMs.
6. Select a design. The HDL appears in the Preview pane.
7. Click Insert to paste the HDL design to the blank Verilog or VHDL file you created in step 2.
8. Click Close to close the Insert Template dialog box.

Figure 14–1. Insert Template Dialog Box
Using Altera Megafunctions

Altera provides parameterizable megafunctions that are optimized for Altera device architectures. Using megafunctions instead of coding your own logic saves valuable design time. Additionally, the Altera-provided megafunctions may offer more efficient logic synthesis and device implementation. You can scale the megafunction’s size and specify various options by setting parameters. Megafunctions include the library of parameterized modules (LPM) and Altera device-specific megafunctions.

To use megafunctions in your HDL code, you can instantiate them as described in “Instantiating Altera Megafunctions in HDL Code” on page 14–3.

Sometimes it is preferable to make your code independent of device family or vendor. In this case, you might not want to instantiate megafunctions directly. For some types of logic functions, such as memories and DSP functions, you can infer device-specific dedicated architecture blocks instead of instantiating a megafunction. Synthesis tools, including Quartus II integrated synthesis, recognize certain types of HDL code and automatically infer the appropriate megafunction or map directly to device atoms. Synthesis tools infer megafunctions to take advantage of logic that is optimized for Altera devices or to target dedicated architectural blocks.

In cases where you prefer to use generic HDL code instead of instantiating a specific function, follow the guidelines and coding examples in “Inferring Multiplier and DSP Functions from HDL Code” on page 14–5 and “Inferring Memory Functions from HDL Code” on page 14–13 to ensure your HDL code infers the appropriate function.

You can infer or instantiate megafunctions to target some Altera device-specific architecture features such as memory and DSP blocks. You must instantiate megafunctions to target certain other device and high-speed features, such as LVDS drivers, phase-locked loops (PLLs), transceivers, and double-data rate input/output (DDIO) circuitry.

Instantiating Altera Megafunctions in HDL Code

The following sections describe how to use megafunctions by instantiating them in your HDL code with the following methods:

- “Instantiating Megafunctions Using the MegaWizard Plug-In Manager”—You can use the MegaWizard™ Plug-In Manager to parameterize the function and create a wrapper file.
- “Creating a Netlist File for Other Synthesis Tools”—You can optionally create a netlist file instead of a wrapper file.
- “Instantiating Megafunctions Using the Port and Parameter Definition”—You can instantiate the function directly in your HDL code.
Instantiating Megafunctions Using the MegaWizard Plug-In Manager

Use the MegaWizard Plug-In Manager as described in this section to create megafunctions in the Quartus II software that you can instantiate in your HDL code. The MegaWizard Plug-In Manager provides a GUI to customize and parameterize megafunctions, and ensures that you set all megafunction parameters properly. When you finish setting parameters, you can specify which files you want generated. Depending on which language you choose, the MegaWizard Plug-In Manager instantiates the megafunction with the correct parameters and generates a megafunction variation file (wrapper file) in Verilog HDL (.v), VHDL (.vhd), or AHDL (.tdf), along with other supporting files.

The MegaWizard Plug-In Manager provides options to create the files listed in Table 14–1.

Table 14–1.  MegaWizard Plug-In Manager Generated Files

<table>
<thead>
<tr>
<th>File</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&lt;output file&gt;.v</code></td>
<td>Verilog HDL Variation Wrapper File—Megafunction wrapper file for instantiation in a Verilog HDL, VHDL, or AHDL design respectively. The MegaWizard Plug-In Manager generates a .v, .vhd, or .tdf file, depending on the language you select for the output file on the megafunction selection page of the wizard.</td>
</tr>
<tr>
<td><code>&lt;output file&gt;.inc</code></td>
<td>ADHL Include File—Used in AHDL Text Design Files (.tdf).</td>
</tr>
<tr>
<td><code>&lt;output file&gt;.cmp</code></td>
<td>Component Declaration File—Used in VHDL design files.</td>
</tr>
<tr>
<td><code>&lt;output file&gt;.bsf</code></td>
<td>Block Symbol File—Used in Quartus II schematic Block Design Files (.bdf).</td>
</tr>
<tr>
<td><code>&lt;output file&gt;_inst.v</code></td>
<td>HDL Instantiation Template for the language of the variation file—Sample instantiation of the Verilog HDL module, VHDL entity, or AHDL subdesign.</td>
</tr>
<tr>
<td><code>&lt;output file&gt;_bb.v</code></td>
<td>Black box Verilog HDL Module Declaration—Hollow-body module declaration that can be used in Verilog HDL designs to specify port directions when instantiating the megafunction as a black box in third-party synthesis tools.</td>
</tr>
<tr>
<td><code>&lt;output file&gt;_syn.v</code></td>
<td>Synthesis timing and resource estimation netlist—Additional synthesis netlist file created if you enable the option to generate a synthesis timing and resource estimation netlist. For more information, refer to “Creating a Netlist File for Other Synthesis Tools” for details.</td>
</tr>
</tbody>
</table>

Creating a Netlist File for Other Synthesis Tools

When you use certain megafunctions with other EDA synthesis tools (that is, tools other than Quartus II integrated synthesis), you can optionally create a netlist for timing and resource estimation instead of a wrapper file.

The netlist file is a representation of the customized logic used in the Quartus II software. The file provides the connectivity of architectural elements in the megafunction but may not represent true functionality. This information enables certain other EDA synthesis tools to better report timing and resource estimates. In addition, synthesis tools can use the timing information to focus timing-driven optimizations and improve the quality of results.

To generate the netlist, turn on Generate netlist under Timing and resource estimation on the EDA page of the MegaWizard Plug-In Manager. The netlist file is called `<output file>_syn.v`. If you use this netlist for synthesis, you must include the megafunction wrapper file, either `<output file>.v` or `<output file>.vhd`, for placement and routing in the project created with the Quartus II software.
Because your synthesis tool may call the Quartus II software in the background to generate this netlist, turning on the Generate Netlist option might be optional.

For information about support for timing and resource estimation netlists in your synthesis tool, refer to the tool vendor’s documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

**Instantiating Megafucntions Using the Port and Parameter Definition**

You can instantiate the megafucntion directly in your Verilog HDL, VHDL, or AHDL code by calling the megafucntion and setting its parameters as you would any other module, component, or subdesign.

For a list of the megafucntion ports and parameters, refer to the specific megafucntion in the Quartus II Help. You can also refer to the IP and Megafucntion page on the Altera website.

Altera recommends that you use the MegaWizard Plug-In Manager for complex megafucntions such as PLLs, transceivers, and LVDS drivers. For details about using the MegaWizard Plug-In Manager, refer to “Instantiating Megafucntions Using the MegaWizard Plug-In Manager” on page 14–4.

**Inferring Multiplier and DSP Functions from HDL Code**

The following sections describe how to infer multiplier and DSP functions from generic HDL code, and, if applicable, how to target the dedicated DSP block architecture in Altera devices:

- “Inferring Multipliers from HDL Code”
- “Inferring Multiply-Accumulators and Multiply-Adders from HDL Code” on page 14–8

For synthesis tool features and options, refer to your synthesis tool documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

For more design examples involving advanced multiply functions and complex DSP functions, refer to the DSP Design Examples page on the Altera website.

**Inferring Multipliers from HDL Code**

To infer multiplier functions, synthesis tools look for multipliers and convert them to LPM_MULT or ALTMULT_ADD megafucntions, or may map them directly to device atoms. For devices with DSP blocks, the software can implement the function in a DSP block instead of logic, depending on device utilization. The Quartus II Fitter can also place input and output registers in DSP blocks (that is, perform register packing) to improve performance and area utilization.

For more information about the DSP block and supported functions, refer to the appropriate Altera device family handbook and the Altera DSP Solutions Center website.
Example 14–1 and Example 14–2 show Verilog HDL code examples, and Example 14–3 and Example 14–4 show VHDL code examples, for unsigned and signed multipliers that synthesis tools can infer as a megafunction or DSP block atoms. Each example fits into one DSP block element. In addition, when register packing occurs, no extra logic cells for registers are required.

The signed declaration in Verilog HDL is a feature of the Verilog 2001 Standard.

**Example 14–1. Verilog HDL Unsigned Multiplier**

```verilog
module unsigned_mult (out, a, b);
    output [15:0] out;
    input  [7:0] a;
    input  [7:0] b;
    assign out = a * b;
endmodule
```

**Example 14–2. Verilog HDL Signed Multiplier with Input and Output Registers (Pipelining = 2)**

```verilog
module signed_mult (out, clk, a, b);
    output [15:0] out;
    input clk;
    input signed [7:0] a;
    input signed [7:0] b;

    reg signed [7:0] a_reg;
    reg signed [7:0] b_reg;
    reg signed [15:0] out;
    wire signed [15:0] mult_out;

    assign mult_out = a_reg * b_reg;

    always @(posedge clk)
    begin
        a_reg <= a;
        b_reg <= b;
        out <= mult_out;
    end
endmodule
```
Example 14–3. VHDL Unsigned Multiplier with Input and Output Registers (Pipelining = 2)

LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.numeric_std.all;

ENTITY unsigned_mult IS
  PORT (  
a: IN UNSIGNED (7 DOWNTO 0);
b: IN UNSIGNED (7 DOWNTO 0);
clk: IN STD_LOGIC;
aclr: IN STD_LOGIC;
result: OUT UNSIGNED (15 DOWNTO 0)
  );
END unsigned_mult;

ARCHITECTURE rtl OF unsigned_mult IS
  SIGNAL a_reg, b_reg: UNSIGNED (7 DOWNTO 0);
BEGIN
  PROCESS (clk, aclr)
  BEGIN
    IF (aclr = '1') THEN
      a_reg <= (OTHERS => '0');
b_reg <= (OTHERS => '0');
    result <= (OTHERS => '0');
    ELSIF (clk'event AND clk = '1') THEN
      a_reg <= a;
b_reg <= b;
    result <= a_reg * b_reg;
    END IF;
  END PROCESS;
END rtl;

Example 14–4. VHDL Signed Multiplier

LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.numeric_std.all;

ENTITY signed_mult IS
  PORT (  
a: IN SIGNED (7 DOWNTO 0);
b: IN SIGNED (7 DOWNTO 0);
result: OUT SIGNED (15 DOWNTO 0)
  );
END signed_mult;

ARCHITECTURE rtl OF signed_mult IS
BEGIN
  result <= a * b;
END rtl;
Inferring Multiply-Accumulators and Multiply-Adders from HDL Code

Synthesis tools detect multiply-accumulators or multiply-adders and convert them to ALTMULT_ACCUM or ALTMULT_ADD megafunctions, respectively, or may map them directly to device atoms. The Quartus II software then places these functions in DSP blocks during placement and routing.

Synthesis tools infer multiply-accumulator and multiply-adder functions only if the Altera device family has dedicated DSP blocks that support these functions.

A simple multiply-accumulator consists of a multiplier feeding an addition operator. The addition operator feeds a set of registers that then feeds the second input to the addition operator. A simple multiply-adder consists of two to four multipliers feeding one or two levels of addition, subtraction, or addition/subtraction operators. Addition is always the second-level operator, if it is used. In addition to the multiply-accumulator and multiply-adder, the Quartus II Fitter also places input and output registers into the DSP blocks to pack registers and improve performance and area utilization.

Some device families offer additional advanced multiply-add and accumulate functions, such as complex multiplication, input shift register, or larger multiplications.

For details about advanced DSP block features, refer to the appropriate device handbook. For more design examples of DSP functions and inferring advanced features in the multiply-add and multiply-accumulate circuitry, refer to the DSP Design Examples page and AN639: Inferring Stratix V DSP Blocks for FIR Filtering Applications on Altera’s website.

The Verilog HDL and VHDL code samples in Example 14–5 through Example 14–8 on pages 14–9 through 14–12 infer multiply-accumulators and multiply-adders with input, output, and pipeline registers, as well as an optional asynchronous clear signal. Using the three sets of registers provides the best performance through the function, with a latency of three. You can remove the registers in your design to reduce the latency.
Example 14–5. Verilog HDL Unsigned Multiply-Accumulator

module unsig_altmult_accum (dataout, dataa, datab, clk, aclr, clken);
    input [7:0] dataa, datab;
    input clk, aclr, clken;
    output reg[16:0] dataout;

    reg [7:0] dataa_reg, datab_reg;
    reg [15:0] multa_reg;
    wire [15:0] multa;
    wire [16:0] adder_out;
    assign multa = dataa_reg * datab_reg;
    assign adder_out = multa_reg + dataout;

    always @(posedge clk or posedge aclr)
    begin
        if (aclr)
            begin
                dataa_reg <= 8'b0;
                datab_reg <= 8'b0;
                multa_reg <= 16'b0;
                dataout <= 17'b0;
            end
        else if (clken)
            begin
                dataa_reg <= dataa;
                datab_reg <= datab;
                multa_reg <= multa;
                dataout <= adder_out;
            end
    end
endmodule
Example 14-6. Verilog HDL Signed Multiply-Adder

module sig_altmult_add (dataa, datab, datac, datad, clock, aclr, result);
  input signed [15:0] dataa, datab, datac, datad;
  input clock, aclr;
  output reg signed [32:0] result;
  reg signed [15:0] dataa_reg, datab_reg, datac_reg, datad_reg;
  reg signed [31:0] mult0_result, mult1_result;

  always @ (posedge clock or posedge aclr) begin
    if (aclr) begin
      dataa_reg <= 16'b0;
      datab_reg <= 16'b0;
      datac_reg <= 16'b0;
      datad_reg <= 16'b0;
      mult0_result <= 32'b0;
      mult1_result <= 32'b0;
      result <= 33'b0;
    end
    else begin
      dataa_reg <= dataa;
      datab_reg <= datab;
      datac_reg <= datac;
      datad_reg <= datad;
      mult0_result <= dataa_reg * datab_reg;
      mult1_result <= datac_reg * datad_reg;
      result <= mult0_result + mult1_result;
    end
  end
endmodule
Example 14–7. VHDL Signed Multiply-Accumulator

LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.numeric_std.all;

ENTITY sig_altmult_accum IS
  PORT (    a: IN SIGNED(7 DOWNTO 0);
            b: IN SIGNED (7 DOWNTO 0);
            clk: IN STD_LOGIC;
            aclr: IN STD_LOGIC;
            accum_out: OUT SIGNED (15 DOWNTO 0)
  ) ;
END sig_altmult_accum;

ARCHITECTURE rtl OF sig_altmult_accum IS
  SIGNAL a_reg, b_reg: SIGNED (7 DOWNTO 0);
  SIGNAL pdt_reg: SIGNED (15 DOWNTO 0);
  SIGNAL adder_out: SIGNED (15 DOWNTO 0);
BEGIN
  PROCESS (clk, aclr)
  BEGIN
    IF (aclr = '1') then
      a_reg <= (others => '0');
      b_reg <= (others => '0');
      pdt_reg <= (others => '0');
      adder_out <= (others => '0');
    ELSIF (clk'event and clk = '1') THEN
      a_reg <= (a);
      b_reg <= (b);
      pdt_reg <= a_reg * b_reg;
      adder_out <= adder_out + pdt_reg;
    END IF;
  END PROCESS;
  accum_out <= adder_out;
END rtl;
Example 14–8. VHDL Unsigned Multiply-Adder

LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.numeric_std.all;

ENTITY unsignedmult_add IS
  PORT (    
a: IN UNSIGNED (7 DOWNTO 0);
b: IN UNSIGNED (7 DOWNTO 0);
c: IN UNSIGNED (7 DOWNTO 0);
d: IN UNSIGNED (7 DOWNTO 0);
clk: IN STD_LOGIC;
aclr: IN STD_LOGIC;
result: OUT UNSIGNED (15 DOWNTO 0)   
  );
END unsignedmult_add;

ARCHITECTURE rtl OF unsignedmult_add IS
  SIGNAL a_reg, b_reg, c_reg, d_reg: UNSIGNED (7 DOWNTO 0);
  SIGNAL pdt_reg, pdt2_reg: UNSIGNED (15 DOWNTO 0);
  SIGNAL result_reg: UNSIGNED (15 DOWNTO 0);
BEGIN
  PROCESS (clk, aclr)
  BEGIN
    IF (aclr = '1') THEN
      a_reg <= (OTHERS => '0');
b_reg <= (OTHERS => '0');
c_reg <= (OTHERS => '0');
d_reg <= (OTHERS => '0');
pdt_reg <= (OTHERS => '0');
pdt2_reg <= (OTHERS => '0');
      END IF;
ELSIF (clk'event AND clk = '1') THEN
      a_reg <= a;
b_reg <= b;
c_reg <= c;
d_reg <= d;
pdt_reg <= a_reg * b_reg;
pdt2_reg <= c_reg * d_reg;
      result_reg <= pdt_reg + pdt2_reg;
      END IF;
  END PROCESS;
  result <= result_reg;
END rtl;
Inferring Memory Functions from HDL Code

The following sections describe how to infer memory functions from generic HDL code and, if applicable, to target the dedicated memory architecture in Altera devices:

- “Inferring RAM functions from HDL Code” on page 14–14
- “Inferring ROM Functions from HDL Code” on page 14–36
- “Shift Registers—Inferring the ALTSHIFT_TAPS Megafunction from HDL Code” on page 14–40

For synthesis tool features and options, refer to your synthesis tool documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

Altera’s dedicated memory architecture offers a number of advanced features that can be easily targeted using the MegaWizard Plug-In Manager, as described in “Instantiating Altera Megafunctions in HDL Code” on page 14–3. The coding recommendations in the following sections provide portable examples of generic HDL code that infer the appropriate megafunction. However, if you want to use some of the advanced memory features in Altera devices, consider using the megafunction directly so that you can control the ports and parameters easily.

You can also use the Quartus II templates provided in the Quartus II software as a starting point. For more information, refer to “Inserting a Template with the Quartus II Text Editor” on page 14–2. Table 14–2 lists the full designs for RAMs and ROMs available in the Quartus II templates.

Most of these designs can also be found on the Design Examples page on the Altera website.

### Table 14–2. RAM and ROM Full Designs from the Quartus II Templates (Part 1 of 2)

<table>
<thead>
<tr>
<th>Language</th>
<th>Full Design Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>VHDL</td>
<td>Single-Port RAM</td>
</tr>
<tr>
<td></td>
<td>Single-Port RAM with Initial Contents</td>
</tr>
<tr>
<td></td>
<td>Simple Dual-Port RAM (single clock)</td>
</tr>
<tr>
<td></td>
<td>Simple Dual-Port RAM (dual clock)</td>
</tr>
<tr>
<td></td>
<td>True Dual-Port RAM (single clock)</td>
</tr>
<tr>
<td></td>
<td>True Dual-Port RAM (dual clock)</td>
</tr>
<tr>
<td></td>
<td>Mixed-Width RAM</td>
</tr>
<tr>
<td></td>
<td>Mixed-Width True Dual-Port RAM</td>
</tr>
<tr>
<td></td>
<td>Byte-Enabled Simple Dual-Port RAM</td>
</tr>
<tr>
<td></td>
<td>Byte-Enabled True Dual-Port RAM</td>
</tr>
<tr>
<td></td>
<td>Single-Port ROM</td>
</tr>
<tr>
<td></td>
<td>Dual-Port ROM</td>
</tr>
</tbody>
</table>
Inferring Memory Functions from HDL Code

To infer RAM functions, synthesis tools detect sets of registers and logic that can be replaced with the ALTSYNCRAM or ALTDPRAM megafunctions for device families that have dedicated RAM blocks, or may map them directly to device memory atoms. Tools typically consider all signals and variables that have a multi-dimensional array type and then create a RAM block, if applicable, based on the way the signals, variables, or both are assigned, referenced, or both in the HDL source description.

Standard synthesis tools recognize single-port and simple dual-port (one read port and one write port) RAM blocks. Some tools (such as the Quartus II software) also recognize true dual-port (two read ports and two write ports) RAM blocks that map to the memory blocks in certain Altera devices.

Some tools (such as the Quartus II software) also infer memory blocks for array variables and signals that are referenced (read/written) by two indices, to recognize mixed-width and byte-enabled RAMs for certain coding styles.

If your design contains a RAM block that your synthesis tool does not recognize and infer, the design might require a large amount of system memory that can potentially cause compilation problems.

When you use a formal verification flow, Altera recommends that you create RAM blocks in separate entities or modules that contain only the RAM logic. In certain formal verification flows, for example, when using Quartus II integrated synthesis, the entity or module containing the inferred RAM is put into a black box automatically because formal verification tools do not support RAM blocks. The Quartus II software issues a warning message when this situation occurs. If the entity or module contains any additional logic outside the RAM block, this logic cannot be verified because it also must be treated as a black box for formal verification.

The following sections present several guidelines for inferring RAM functions that match the dedicated memory architecture in Altera devices, and then provide recommended HDL code for different types of memory logic.

### Table 14–2. RAM and ROM Full Designs from the Quartus II Templates (Part 2 of 2)

<table>
<thead>
<tr>
<th>Language</th>
<th>Full Design Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Verilog HDL</td>
<td>Single-Port RAM&lt;br&gt;Single-Port RAM with Initial Contents&lt;br&gt;Simple Dual-Port RAM (single clock)&lt;br&gt;Simple Dual-Port RAM (dual clock)&lt;br&gt;True Dual-Port RAM (single clock)&lt;br&gt;True Dual-Port RAM (dual clock)&lt;br&gt;Single-Port ROM&lt;br&gt;Dual-Port ROM</td>
</tr>
<tr>
<td>System Verilog</td>
<td>Mixed-Width Port RAM&lt;br&gt;Mixed-Width True Dual-Port RAM&lt;br&gt;Mixed-Width True Dual-Port RAM (new data on same port read during write)&lt;br&gt;Byte-Enabled Simple Dual Port RAM&lt;br&gt;Byte-Enabled True Dual-Port RAM</td>
</tr>
</tbody>
</table>
Use Synchronous Memory Blocks

Altera recommends using synchronous memory blocks for Altera designs. Because memory blocks in the newest devices from Altera are synchronous, RAM designs that are targeted towards architectures that contain these dedicated memory blocks must be synchronous to be mapped directly into the device architecture. For these devices, asynchronous memory logic is implemented in regular logic cells.

Synchronous memory offers several advantages over asynchronous memory, including higher frequencies and thus higher memory bandwidth, increased reliability, and less standby power. In many designs with asynchronous memory, the memory interfaces with synchronous logic so that the conversion to synchronous memory design is straightforward. To convert asynchronous memory you can move registers from the data path into the memory block.

Synchronous memories are supported in all Altera device families. A memory block is considered synchronous if it uses one of the following read behaviors:

- Memory read occurs in a Verilog `always` block with a clock signal or a VHDL clocked process. The recommended coding style for synchronous memories is to create your design with a registered read output.

- Memory read occurs outside a clocked block, but there is a synchronous read address (that is, the address used in the read statement is registered). This type of logic is not always inferred as a memory block, or may require external bypass logic, depending on the target device architecture.

The synchronous memory structures in Altera devices can differ from the structures in other vendors’ devices. For best results, match your design to the target device architecture.

Later sections provide coding recommendations for various memory types. All of these examples are synchronous to ensure that they can be directly mapped into the dedicated memory architecture available in Altera FPGAs.

For additional information about the dedicated memory blocks in your specific device, refer to the appropriate Altera device family data sheet on the Altera website at www.altera.com.

Avoid Unsupported Reset and Control Conditions

To ensure that your HDL code can be implemented in the target device architecture, avoid unsupported reset conditions or other control logic that does not exist in the device architecture.

The RAM contents of Altera memory blocks cannot be cleared with a reset signal during device operation. If your HDL code describes a RAM with a reset signal for the RAM contents, the logic is implemented in regular logic cells instead of a memory block. Altera recommends against putting RAM read or write operations in an `always` block or `process` block with a reset signal. If you want to specify memory contents, initialize the memory as described in “Specifying Initial Memory Contents at Power-Up” on page 14–33 or write the data to the RAM during device operation.
**Example 14–9** shows an example of undesirable code where there is a reset signal that clears part of the RAM contents. Avoid this coding style because it is not supported in Altera memories.

**Example 14–9. Verilog RAM with Reset Signal that Clears RAM Contents: Not Supported in Device Architecture**

```verilog
module clear_ram
(
    input clock, reset, we,
    input [7:0] data_in,
    input [4:0] address,
    output reg [7:0] data_out
);

reg [7:0] mem [0:31];
integer i;

always @ (posedge clock or posedge reset)
begin
    if (reset == 1'b1)
        mem[address] <= 0;
    else if (we == 1'b1)
        mem[address] <= data_in;

    data_out <= mem[address];
end
endmodule
```
Example 14–10 shows an example of undesirable code where the reset signal affects the RAM, although the effect may not be intended. Avoid this coding style because it is not supported in Altera memories.

Example 14–10. Verilog RAM with Reset Signal that Affects RAM: Not Supported in Device Architecture

```verilog
module bad_reset
    (input clock,
     input reset,
     input we,
     input [7:0] data_in,
     input [4:0] address,
     output reg [7:0] data_out,
     input d,
     output reg q
    );

    reg [7:0] mem [0:31];
    integer i;

    always @(posedge clock or posedge reset)
    begin
        if (reset == 1'b1)
            q <= 0;
        else
            begin
                if (we == 1'b1)
                    mem[address] <= data_in;
                data_out <= mem[address];
                q <= d;
            end
    end
endmodule
```

In addition to reset signals, other control logic can prevent memory logic from being inferred as a memory block. For example, you cannot use a clock enable on the read address registers in Stratix® devices because doing so affects the output latch of the RAM, and therefore the synthesized result in the device RAM architecture would not match the HDL description. You can use the address stall feature as a read address clock enable in Stratix II, Cyclone® II, Arria® GX, and other newer devices to avoid this limitation. Check the documentation for your device architecture to ensure that your code matches the hardware available in the device.

Check Read-During-Write Behavior

It is important to check the read-during-write behavior of the memory block described in your HDL design as compared to the behavior in your target device architecture. Your HDL source code specifies the memory behavior when you read and write from the same memory address in the same clock cycle. The code specifies that the read returns either the old data at the address, or the new data being written to the address. This behavior is referred to as the read-during-write behavior of the memory block. Altera memory blocks have different read-during-write behavior depending on the target device family, memory mode, and block type.
Synthesis tools map an HDL design into the target device architecture, with the goal of maintaining the functionality described in your source code. Therefore, if your source code specifies unsupported read-during-write behavior for the device RAM blocks, the software must implement the logic outside the RAM hardware in regular logic cells.

One common problem occurs when there is a continuous read in the HDL code, as in the following examples. You should avoid using these coding styles:

```verilog
//Verilog HDL concurrent signal assignment
assign q = ram[raddr_reg];
```

```vhdl
-- VHDL concurrent signal assignment
q <= ram(raddr_reg);
```

When a write operation occurs, this type of HDL implies that the read should immediately reflect the new data at the address, independent of the read clock. However, that is not the behavior of synchronous memory blocks. In the device architecture, the new data is not available until the next edge of the read clock. Therefore, if the synthesis tool mapped the logic directly to a synchronous memory block, the device functionality and gate-level simulation results would not match the HDL description or functional simulation results. If the write clock and read clock are the same, the synthesis tool can infer memory blocks and add extra bypass logic so that the device behavior matches the HDL behavior. If the write and read clocks are different, the synthesis tool cannot reliably add bypass logic, so the logic is implemented in regular logic cells instead of dedicated RAM blocks. The examples in the following sections discuss some of these differences for read-during-write conditions.

In addition, the MLAB feature in certain device logic array blocks (LABs) does not easily support old data or new data behavior for a read-during-write in the dedicated device architecture. Implementing the extra logic to support this behavior significantly reduces timing performance through the memory.

For best performance in MLAB memories, your design should not depend on the read data during a write operation.

In many synthesis tools, you can specify that the read-during-write behavior is not important to your design; for example, if you never read from the same address to which you write in the same clock cycle. For Quartus II integrated synthesis, add the synthesis attribute `ramstyle` set to "no_rw_check" to allow the software to choose the read-during-write behavior of a RAM, rather than use the behavior specified by your HDL code. In some cases, this attribute prevents the synthesis tool from using extra logic to implement the memory block, or can allow memory inference when it would otherwise be impossible.

Synchronous RAM blocks require a synchronous read, so Quartus II integrated synthesis packs either data output registers or read address registers into the RAM block. When the read address registers are packed into the RAM block, the read address signals connected to the RAM block contain the next value of the read address signals indexing the HDL variable, which impacts which clock cycle the read and the write occur, and changes the read-during-write conditions. Therefore, bypass logic may still be added to the design to preserve the read-during-write behavior, even if the "no_rw_check" attribute is set.
For more information about attribute syntax, the no_rw_check attribute value, or specific options for your synthesis tool, refer to your synthesis tool documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

The next section describes how you control the logic implementation in the Altera device, and the following sections provide coding recommendations for various memory types. Each example describes the read-during-write behavior and addresses the support for the memory type in Altera devices.

**Controlling Inference and Implementation in Device RAM Blocks**

Tools usually do not infer small RAM blocks because small RAM blocks typically can be implemented more efficiently using the registers in regular logic. If you are using Quartus II integrated synthesis, you can direct the software to infer RAM blocks for all sizes with the Allow Any RAM Size for Recognition option in the More Analysis & Synthesis Settings dialog box.

Some synthesis tools provide options to control the implementation of inferred RAM blocks for Altera devices with synchronous memory blocks. For example, Quartus II integrated synthesis provides the ramstyle synthesis attribute to specify the type of memory block or to specify the use of regular logic instead of a dedicated memory block. Quartus II integrated synthesis does not map inferred memory into MLABs unless the HDL code specifies the appropriate ramstyle attribute, although the Fitter may map some memories to MLABs.

For details about using the ramstyle attribute, refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook. For information about synthesis attributes in other synthesis tools, refer to the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

If you want to control the implementation after the RAM function is inferred during synthesis, you can set the ram_block_type parameter of the ALTSYNCRAM megafunction. In the Assignment Editor, select Parameters in the Categories list. You can use the Node Finder or drag the appropriate instance from the Project Navigator window to enter the RAM hierarchical instance name. Type ram_block_type as the Parameter Name and type one of the following memory types supported by your target device family in the Value field: "M-RAM", "M512", "M4K", "M9K", "M10K", "M20K", "M144K", or "MLAB".

You can also specify the maximum depth of memory blocks used to infer RAM or ROM in your design. Apply the max_depth synthesis attribute to the declaration of a variable that represents a RAM or ROM in your design file. For example:

```vhdl
-- Limit the depth of the memory blocks implement "ram" to 512
-- This forces the software to use two M512 blocks instead of one M4K block to implement this RAM
(* max_depth = 512 *) reg [7:0] ram[0:1023];
```

**Single-Clock Synchronous RAM with Old Data Read-During-Write Behavior**

The code examples in this section show Verilog HDL and VHDL code that infers simple dual-port, single-clock synchronous RAM. Single-port RAM blocks use a similar coding style.
The read-during-write behavior in these examples is to read the old data at the memory address. Refer to “Check Read-During-Write Behavior” on page 14–17 for details. Altera recommends that you use the Old Data Read-During-Write coding style for most RAM blocks as long as your design does not require the RAM location’s new value when you perform a simultaneous read and write to that RAM location. For best performance in MLAB memories, use the appropriate attribute so that your design does not depend on the read data during a write operation.

If you require that the read-during-write results in new data, refer to “Single-Clock Synchronous RAM with New Data Read-During-Write Behavior” on page 14–21.

The simple dual-port RAM code samples in Example 14–11 and Example 14–12 map directly into Altera synchronous memory.

Single-port versions of memory blocks (that is, using the same read address and write address signals) can allow better RAM utilization than dual-port memory blocks, depending on the device family.

Example 14–11. Verilog HDL Single-Clock Simple Dual-Port Synchronous RAM with Old Data Read-During-Write Behavior

```verilog
module single_clk_ram(
    output reg [7:0] q,
    input [7:0] d,
    input [6:0] write_address, read_address,
    input we, clk
);
    reg [7:0] mem [127:0];
    always @ (posedge clk) begin
        if (we)
            mem[write_address] <= d;
        q <= mem[read_address]; // q doesn't get d in this clock cycle
    end
endmodule
```
Single-Clock Synchronous RAM with New Data Read-During-Write Behavior

The examples in this section describe RAM blocks in which a simultaneous read and write to the same location reads the new value that is currently being written to that RAM location.

To implement this behavior in the target device, synthesis software adds bypass logic around the RAM block. This bypass logic increases the area utilization of the design and decreases the performance if the RAM block is part of the design’s critical path. For more information, refer to “Check Read-During-Write Behavior” on page 14–17 for details. If this behavior is not required for your design, use the examples from “Single-Clock Synchronous RAM with Old Data Read-During-Write Behavior” on page 14–19.

The simple dual-port RAM in Example 14–13 and Example 14–14 require the software to create bypass logic around the RAM block.

Single-port versions of the Verilog memory block (that is, using the same read address and write address signals) do not require any logic cells to create bypass logic in the Arria, Stratix, and Cyclone series of devices, because the device memory supports new data read-during-write behavior when in single-port mode (same clock, same read address, and same write address).
Example 14–13 is similar to Example 14–11, but Example 14–13 uses a blocking assignment for the write so that the data is assigned immediately.

An alternative way to create a single-clock RAM is to use an assign statement to read the address of \( \text{mem} \) to create the output \( q \), as shown in the following coding style example. By itself, the code describes new data read-during-write behavior. However, if the RAM output feeds a register in another hierarchy, a read-during-write results in the old data. Synthesis tools may not infer a RAM block if the tool cannot determine which behavior is described, such as when the memory feeds a hard hierarchical partition boundary. For this reason, avoid using this alternate type of coding style:

```
reg [7:0] mem [127:0];
reg [6:0] read_address_reg;

always @(posedge clk) begin
  if (we)
    mem[write_address] <= d;
  q = mem[read_address]; // q does get d in this clock cycle if // we is high
end

assign q = mem[read_address_reg];
```

```
module single_clock_wr_ram(
  output reg [7:0] q,
  input [7:0] d,
  input [6:0] write_address, read_address,
  input we, clk
);

  reg [7:0] mem [127:0];

  always @(posedge clk) begin
    if (we)
      mem[write_address] = d;
    q = mem[read_address]; // q does get d in this clock cycle if // we is high
  end

endmodule
```

Example 14–13 is similar to Example 14–11, but Example 14–13 uses a blocking assignment for the write so that the data is assigned immediately.

An alternative way to create a single-clock RAM is to use an assign statement to read the address of \( \text{mem} \) to create the output \( q \), as shown in the following coding style example. By itself, the code describes new data read-during-write behavior. However, if the RAM output feeds a register in another hierarchy, a read-during-write results in the old data. Synthesis tools may not infer a RAM block if the tool cannot determine which behavior is described, such as when the memory feeds a hard hierarchical partition boundary. For this reason, avoid using this alternate type of coding style:

```
reg [7:0] mem [127:0];
reg [6:0] read_address_reg;

always @(posedge clk) begin
  if (we)
    mem[write_address] <= d;
  read_address_reg <= read_address;
end

assign q = mem[read_address_reg];
```
The VHDL sample in Example 14–14 uses a concurrent signal assignment to read from the RAM. By itself, this example describes new data read-during-write behavior. However, if the RAM output feeds a register in another hierarchy, a read-during-write results in the old data. Synthesis tools may not infer a RAM block if the tool cannot determine which behavior is described, such as when the memory feeds a hard hierarchical partition boundary.

**Example 14–14. VHDL Single-Clock Simple Dual-Port Synchronous RAM with New Data Read-During-Write Behavior**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;

ENTITY single_clock_rw_ram IS
  PORT (
    clock: IN STD_LOGIC;
    data: IN STD_LOGIC_VECTOR (2 DOWNTO 0);
    write_address: IN INTEGER RANGE 0 to 31;
    read_address: IN INTEGER RANGE 0 to 31;
    we: IN STD_LOGIC;
    q: OUT STD_LOGIC_VECTOR (2 DOWNTO 0)
  );
END single_clock_rw_ram;

ARCHITECTURE rtl OF single_clock_rw_ram IS
  TYPE MEM IS ARRAY(0 TO 31) OF STD_LOGIC_VECTOR(2 DOWNTO 0);
  SIGNAL ram_block: MEM;
  SIGNAL read_address_reg: INTEGER RANGE 0 to 31;
BEGIN
  PROCESS (clock)
  BEGIN
    IF (clock'event AND clock = '1') THEN
      IF (we = '1') THEN
        ram_block(write_address) <= data;
      END IF;
      read_address_reg <= read_address;
    END IF;
  END PROCESS;
  q <= ram_block(read_address_reg);
END rtl;
```

For Quartus II integrated synthesis, if you do not require the read-through-write capability, add the synthesis attribute `ramstyle="no_rw_check"` to allow the software to choose the read-during-write behavior of a RAM, rather than using the behavior specified by your HDL code. As discussed in “Check Read-During-Write Behavior” on page 14–17, this attribute may prevent generation of extra bypass logic but it is not always possible to eliminate the requirement for bypass logic.

**Simple Dual-Port, Dual-Clock Synchronous RAM**

In dual clock designs, synthesis tools cannot accurately infer the read-during-write behavior because it depends on the timing of the two clocks within the target device. Therefore, the read-during-write behavior of the synthesized design is undefined and may differ from your original HDL code. When Quartus II integrated synthesis infers this type of RAM, it issues a warning because of the undefined read-during-write behavior. Refer to “Check Read-During-Write Behavior” on page 14–17 for details.
The code samples in Example 14–15 and Example 14–16 show Verilog HDL and VHDL code that infers dual-clock synchronous RAM. The exact behavior depends on the relationship between the clocks.

**Example 14–15. Verilog HDL Simple Dual-Port, Dual-Clock Synchronous RAM**

```verilog
module dual_clock_ram(
    output reg [7:0] q,
    input [7:0] d,
    input [6:0] write_address, read_address,
    input we, clk1, clk2
);
    reg [6:0] read_address_reg;
    reg [7:0] mem [127:0];
    
    always @(posedge clk1)
    begin
        if (we)
            mem[write_address] <= d;
    end
    
    always @(posedge clk2) begin
        q <= mem[read_address_reg];
        read_address_reg <= read_address;
    end
endmodule
```

**Example 14–16. VHDL Simple Dual-Port, Dual-Clock Synchronous RAM**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
ENTITY dual_clock_ram IS
    PORT (
        clock1, clock2: IN STD_LOGIC;
        data: IN STD_LOGIC_VECTOR (3 DOWNTO 0);
        write_address: IN INTEGER RANGE 0 to 31;
        read_address: IN INTEGER RANGE 0 to 31;
        we: IN STD_LOGIC;
        q: OUT STD_LOGIC_VECTOR (3 DOWNTO 0)
    );
END dual_clock_ram;
ARCHITECTURE rtl OF dual_clock_ram IS
    TYPE MEM IS ARRAY(0 TO 31) OF STD_LOGIC_VECTOR(3 DOWNTO 0);
    SIGNAL ram_block: MEM;
    SIGNAL read_address_reg : INTEGER RANGE 0 to 31;
BEGIN
    PROCESS (clock1)
    BEGIN
        IF (clock1'event AND clock1 = '1') THEN
            IF (we = '1') THEN
                ram_block(write_address) <= data;
            END IF;
        END IF;
    END PROCESS;
    PROCESS (clock2)
    BEGIN
        IF (clock2'event AND clock2 = '1') THEN
            q <= ram_block(read_address_reg);
            read_address_reg <= read_address;
        END IF;
    END PROCESS;
END rtl;
```
**True Dual-Port Synchronous RAM**

The code examples in this section show Verilog HDL and VHDL code that infers true dual-port synchronous RAM. Different synthesis tools may differ in their support for these types of memories. This section describes the inference rules for Quartus II integrated synthesis. This type of RAM inference is supported for the Arria GX, Stratix, and Cyclone series of devices.

Altera synchronous memory blocks have two independent address ports, allowing for operations on two unique addresses simultaneously. A read operation and a write operation can share the same port if they share the same address. The Quartus II software infers true dual-port RAMs in Verilog HDL and VHDL with any combination of independent read or write operations in the same clock cycle, with at most two unique port addresses, performing two reads and one write, two writes and one read, or two writes and two reads in one clock cycle with one or two unique addresses.

In the synchronous RAM block architecture, there is no priority between the two ports. Therefore, if you write to the same location on both ports at the same time, the result is indeterminate in the device architecture. You must ensure your HDL code does not imply priority for writes to the memory block, if you want the design to be implemented in a dedicated hardware memory block. For example, if both ports are defined in the same process block, the code is synthesized and simulated sequentially so that there is a priority between the two ports. If your code does imply a priority, the logic cannot be implemented in the device RAM blocks and is implemented in regular logic cells.

You must also consider the read-during-write behavior of the RAM block to ensure that it can be mapped directly to the device RAM architecture. Refer to “Check Read-During-Write Behavior” on page 14–17 for details.

When a read and write operation occurs on the same port for the same address, the read operation may behave as follows:

- **Read new data**—This mode matches the behavior of synchronous memory blocks.
- **Read old data**—This mode is supported only in device families that support M144, M9k, and MLAB memory blocks.

When a read and write operation occurs on different ports for the same address (also known as mixed port), the read operation may behave as follows:

- **Read new data**—Quartus II integrated synthesis supports this mode by creating bypass logic around the synchronous memory block.
- **Read old data**—Synchronous memory blocks support this behavior.
- **Read don’t care**—This behavior is supported on different ports in simple dual-port mode by synchronous memory blocks for all device families except Arria, Arria GX, Cyclone, Cyclone II, HardCopy, HardCopy II, MAX, Stratix, and Stratix II device families.

The Verilog HDL single-clock code sample in Example 14–17 maps directly into Altera synchronous memory. When a read and write operation occurs on the same port for the same address, the new data being written to the memory is read. When a read and write operation occurs on different ports for the same address, the old data in the memory is read. Simultaneous writes to the same location on both ports results in indeterminate behavior.
A dual-clock version of this design describes the same behavior, but the memory in the target device will have undefined mixed port read-during-write behavior because it depends on the relationship between the clocks.

**Example 14–17. Verilog HDL True Dual-Port RAM with Single Clock**

```verilog
module true_dual_port_ram_single_clock
(
    input [(DATA_WIDTH-1):0] data_a, data_b,
    input [(ADDR_WIDTH-1):0] addr_a, addr_b,
    input we_a, we_b, clk,
    output reg [(DATA_WIDTH-1):0] q_a, q_b
);

    parameter DATA_WIDTH = 8;
    parameter ADDR_WIDTH = 6;

    // Declare the RAM variable
    reg [DATA_WIDTH-1:0] ram[2**ADDR_WIDTH-1:0];

always @ (posedge clk)
begin // Port A
    if (we_a)
        begin
            ram[addr_a] <= data_a;
            q_a <= data_a;
        end
    else
        q_a <= ram[addr_a];
end
always @ (posedge clk)
begin // Port b
    if (we_b)
        begin
            ram[addr_b] <= data_b;
            q_b <= data_b;
        end
    else
        q_b <= ram[addr_b];
end
endmodule
```

If you use the following Verilog HDL read statements instead of the `if-else` statements in Example 14–17, the HDL code specifies that the read results in old data when a read operation and write operation occurs at the same time for the same address on the same port or mixed ports. This mode is supported only in device families that support M144, M9k, and MLAB memory blocks.

```verilog
always @ (posedge clk)
begin // Port A
    if (we_a)
        ram[addr_a] <= data_a;
    q_a <= ram[addr_a];
end
always @ (posedge clk)
begin // Port B
```
The VHDL single-clock code sample in Example 14–18 maps directly into Altera synchronous memory. When a read and write operation occurs on the same port for the same address, the new data being written to the memory is read. When a read and write operation occurs on different ports for the same address, the old data in the memory is read. Simultaneous write operations to the same location on both ports results in indeterminate behavior.

A dual-clock version of this design describes the same behavior, but the memory in the target device will have undefined mixed port read-during-write behavior because it depends on the relationship between the clocks.

Example 14–18. VHDL True Dual-Port RAM with Single Clock (Part 1 of 2)

```vhdl
library ieee;
use ieee.std_logic_1164.all;

entity true_dual_port_ram_single_clock is
  generic (
    DATA_WIDTH : natural := 8;
    ADDR_WIDTH : natural := 6
  );
  port (    clk: in std_logic;
    addr_a: in natural range 0 to 2**ADDR_WIDTH - 1;
    addr_b: in natural range 0 to 2**ADDR_WIDTH - 1;
    data_a: in std_logic_vector((DATA_WIDTH-1) downto 0);
    data_b: in std_logic_vector((DATA_WIDTH-1) downto 0);
    we_a: in std_logic := '1';
    we_b: in std_logic := '1';
    q_a: out std_logic_vector((DATA_WIDTH -1) downto 0);
    q_b: out std_logic_vector((DATA_WIDTH -1) downto 0)
  );
end true_dual_port_ram_single_clock;

architecture rtl of true_dual_port_ram_single_clock is
  -- Build a 2-D array type for the RAM
  subtype word_t is std_logic_vector((DATA_WIDTH-1) downto 0);
  type memory_t is array((2**ADDR_WIDTH - 1) downto 0) of word_t;
  -- Declare the RAM signal.
  shared variable ram : memory_t;
```
Mixed-Width Dual-Port RAM

The RAM code examples in Example 14–20 through Example 14–23 show SystemVerilog and VHDL code that infers RAM with data ports with different widths. This type of logic is not supported in Verilog-1995 or Verilog-2001 because of the requirement for a multi-dimensional array to model the different read width, write width, or both. Different synthesis tools may differ in their support for these memories. This section describes the inference rules for Quartus II integrated synthesis.

The first dimension of the multi-dimensional packed array represents the ratio of the wider port to the narrower port, and the second dimension represents the narrower port width. The read and write port widths must specify a read or write ratio supported by the memory blocks in the target device, or the synthesis tool does not infer a RAM.

Example 14–19. VHDL True Dual-Port RAM with Single Clock (Part 2 of 2)

```
begin
  process(clk)
  begin
    if(rising_edge(clk)) then -- Port A
      if(we_a = '1') then
        ram(addr_a) <= data_a;
        -- Read-during-write on the same port returns NEW data
        q_a <= data_a;
      else
        -- Read-during-write on the mixed port returns OLD data
        q_a <= ram(addr_a);
      end if;
    end if;
  end process;

  process(clk)
  begin
    if(rising_edge(clk)) then -- Port B
      if(we_b = '1') then
        ram(addr_b) <= data_b;
        -- Read-during-write on the same port returns NEW data
        q_b <= data_b;
      else
        -- Read-during-write on the mixed port returns OLD data
        q_b <= ram(addr_b);
      end if;
    end if;
  end process;
end rtl;
```
Refer to the Quartus II Templates for parameterized examples that you can use for supported combinations of read and write widths, and true dual port RAM examples with two read ports and two write ports for mixed-width writes and reads.

**Example 14–20. SystemVerilog Mixed-Width RAM with Read Width Smaller than Write Width**

```verilog
module mixed_width_ram    // 256x32 write and 1024x8 read
(
    input [7:0] waddr,
    input [31:0] wdata,
    input we, clk,
    input [9:0] raddr,
    output [7:0] q
);
logic [3:0][7:0] ram[0:255];
always_ff@(posedge clk)
begin
    if(we) ram[waddr] <= wdata;
    q <= ram[raddr / 4][raddr % 4];
end
endmodule : mixed_width_ram
```

**Example 14–21. SystemVerilog Mixed-Width RAM with Read Width Larger than Write Width**

```verilog
module mixed_width_ram     // 1024x8 write and 256x32 read
(
    input [9:0] waddr,
    input [31:0] wdata,
    input we, clk,
    input [7:0] raddr,
    output [9:0] q
);
logic [3:0][7:0] ram[0:255];
always_ff@(posedge clk)
begin
    if(we) ram[waddr / 4][waddr % 4] <= wdata;
    q <= ram[raddr];
end
endmodule : mixed_width_ram
```
Example 14–22. VHDL Mixed-Width RAM with Read Width Smaller than Write Width

library ieee;
use ieee.std_logic_1164.all;

package ram_types is
  type word_t is array (0 to 3) of std_logic_vector(7 downto 0);
  type ram_t is array (0 to 255) of word_t;
end ram_types;

library ieee;
use ieee.std_logic_1164.all;
library work;
use work.ram_types.all;

entity mixed_width_ram is
  port (    we, clk : in  std_logic;
            waddr  : in  integer range 0 to 255;
            wdata  : in  word_t;
            raddr  : in  integer range 0 to 1023;
            q      : out std_logic_vector(7 downto 0));
end mixed_width_ram;

architecture rtl of mixed_width_ram is
  signal ram : ram_t;
begin  -- rtl
  process(clk, we)
  begin
    if(rising_edge(clk)) then
      if(we = '1') then
        ram(waddr) <= wdata;
      end if;
      q <= ram(raddr / 4 )(raddr mod 4);
    end if;
  end process;
end rtl;
**Example 14–23. VHDL Mixed-Width RAM with Read Width Larger than Write Width**

```vhdl
library ieee;
use ieee.std_logic_1164.all;

package ram_types is
  type word_t is array (0 to 3) of std_logic_vector(7 downto 0);
  type ram_t is array (0 to 255) of word_t;
end ram_types;

library ieee;
use ieee.std_logic_1164.all;
library work;
use work.ram_types.all;

entity mixed_width_ram is
  port (  
    we, clk : in  std_logic;
    waddr   : in  integer range 0 to 1023;
    wdata   : in  std_logic_vector(7 downto 0);
    raddr   : in  integer range 0 to 255;
    q       : out word_t);
end mixed_width_ram;

architecture rtl of mixed_width_ram is
  signal ram : ram_t;
begin  -- rtl
  process(clk, we)
  begin
    if(rising_edge(clk)) then
      if(we = '1') then
        ram(waddr / 4)(waddr mod 4) <= wdata;
      end if;
      q <= ram(raddr);
    end if;
  end process;
end rtl;
```

**RAM with Byte-Enable Signals**

The RAM code examples in Example 14–24 and Example 14–25 show SystemVerilog and VHDL code that infers RAM with controls for writing single bytes into the memory word, or byte-enable signals. Byte enables are modeled by creating write expressions with two indices and writing part of a RAM "word." With these implementations, you can also write more than one byte at once by enabling the appropriate byte enables.

This type of logic is not supported in Verilog-1995 or Verilog-2001 because of the requirement for a multidimensional array. Different synthesis tools may differ in their support for these memories. This section describes the inference rules for Quartus II integrated synthesis.
Refer to the Quartus II Templates for parameterized examples that you can use for different address widths, and true dual port RAM examples with two read ports and two write ports.

Example 14–24. SystemVerilog Simple Dual-Port Synchronous RAM with Byte Enable

```
module byte_enabled_simple_dual_port_ram
  (
    input we, clk,
    input [5:0] waddr, raddr, // address width = 6
    input [3:0] be,     // 4 bytes per word
    input [31:0] wdata,  // byte width = 8, 4 bytes per word
    output reg [31:0] q  // byte width = 8, 4 bytes per word
  );
  // use a multi-dimensional packed array
  //to model individual bytes within the word
  logic [3:0][7:0] ram[0:63]; // # words = 1 << address width
  always_ff@(posedge clk)
    begin
      if(we) begin
        if(be[0]) ram[waddr][0] <= wdata[7:0];
        if(be[1]) ram[waddr][1] <= wdata[15:8];
        if(be[2]) ram[waddr][2] <= wdata[23:16];
        if(be[3]) ram[waddr][3] <= wdata[31:24];
      end
      q <= ram[raddr];
    end
endmodule
```
Specifying Initial Memory Contents at Power-Up

Your synthesis tool may offer various ways to specify the initial contents of an inferred memory.

Certain device memory types do not support initialized memory, such as the M-RAM blocks in Stratix and Stratix II devices.
There are slight power-up and initialization differences between dedicated RAM blocks and the MLAB memory due to the continuous read of the MLAB. Altera dedicated RAM block outputs always power-up to zero and are set to the initial value on the first read. For example, if address 0 is pre-initialized to FF, the RAM block powers up with the output at 0. A subsequent read after power-up from address 0 outputs the pre-initialized value of FF. Therefore, if a RAM is powered up and an enable (read enable or clock enable) is held low, the power-up output of 0 is maintained until the first valid read cycle. The MLAB is implemented using registers that power-up to 0, but are initialized to their initial value immediately at power-up or reset. Therefore, the initial value is seen, regardless of the enable status. The Quartus II software maps inferred memory to MLABs when the HDL code specifies an appropriate ramstyle attribute.

Quartus II integrated synthesis supports the ram_init_file synthesis attribute that allows you to specify a Memory Initialization File (.mif) for an inferred RAM block.

For information about the ram_init_file attribute, refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook. For information about synthesis attributes in other synthesis tools, refer to the tool vendor’s documentation.

In Verilog HDL, you can use an initial block to initialize the contents of an inferred memory. Quartus II integrated synthesis automatically converts the initial block into a .mif file for the inferred RAM. Example 14–26 shows Verilog HDL code that infers a simple dual-port RAM block and corresponding .mif file.

Example 14–26. Verilog HDL RAM with Initialized Contents

```verilog
module ram_with_init(
    output reg [7:0] q,
    input [7:0] d,
    input [4:0] write_address, read_address,
    input we, clk
);
    reg [7:0] mem [0:31];
    integer i;

    initial begin
        for (i = 0; i < 32; i = i + 1)
            mem[i] = i[7:0];
    end

    always @ (posedge clk) begin
        if (we)
            mem[write_address] <= d;
        q <= mem[read_address];
    end
endmodule
```

Quartus II integrated synthesis and other synthesis tools also support the $readmemb and $readmemh commands so that RAM initialization and ROM initialization work identically in synthesis and simulation. Example 14–27 shows an initial block that initializes an inferred RAM block using the $readmemb command.
Refer to the *Verilog Language Reference Manual (LRM)* 1364-2001 Section 17.2.8 or the example in the Templates for the Quartus II software for details about the format of the *ram.txt* file.

**Example 14–27. Verilog HDL RAM Initialized with the readmemb Command**

```verilog
reg [7:0] ram[0:15];
initial
begin
  $readmemb("ram.txt", ram);
end
```

In VHDL, you can initialize the contents of an inferred memory by specifying a default value for the corresponding signal. Quartus II integrated synthesis automatically converts the default value into a *mif* file for the inferred RAM. **Example 14–28** shows VHDL code that infers a simple dual-port RAM block and corresponding *mif* file.

**Example 14–28. VHDL RAM with Initialized Contents**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
use ieee.numeric_std.all;

ENTITY ram_with_init IS
  PORT(
    clock: IN STD_LOGIC;
    data: IN UNSIGNED (7 DOWNTO 0);
    write_address: IN integer RANGE 0 to 31;
    read_address: IN integer RANGE 0 to 31;
    we: IN std_logic;
    q: OUT UNSIGNED (7 DOWNTO 0));
END;

ARCHITECTURE rtl OF ram_with_init IS
  TYPE MEM IS ARRAY(31 DOWNTO 0) OF unsigned(7 DOWNTO 0);
  FUNCTION initialize_ram
    return MEM is
    variable result : MEM;
    BEGIN
        FOR i IN 31 DOWNTO 0 LOOP
            result(i) := to_unsigned(natural(i), natural'(8));
        END LOOP;
        RETURN result;
    END initialize_ram;

    SIGNAL ram_block : MEM := initialize_ram;
    BEGIN
        PROCESS (clock)
        BEGIN
            IF (clock'event AND clock = '1') THEN
                IF (we = '1') THEN
                    ram_block(write_address) <= data;
                END IF;
                q <= ram_block(read_address);
            END IF;
        END PROCESS;
    END rtl;
```
Inferring ROM Functions from HDL Code

To infer ROM functions, synthesis tools detect sets of registers and logic that can be replaced with the ALTSYNCRAM or LPM_ROM megafunctions, depending on the target device family, and only for device families that have dedicated memory blocks.

ROMs are inferred when a `CASE` statement exists in which a value is set to a constant for every choice in the `CASE` statement. Because small ROMs typically achieve the best performance when they are implemented using the registers in regular logic, each ROM function must meet a minimum size requirement to be inferred and placed into memory.

If you use Quartus II integrated synthesis, you can direct the software to infer ROM blocks for all sizes with the **Allow Any ROM Size for Recognition** option in the **More Analysis & Synthesis Settings** dialog box.

Some synthesis tools provide options to control the implementation of inferred ROM blocks for Altera devices with synchronous memory blocks. For example, Quartus II integrated synthesis provides the `romstyle` synthesis attribute to specify the type of memory block or to specify the use of regular logic instead of a dedicated memory block.

For details about using the `romstyle` attribute, refer to the *Quartus II Integrated Synthesis* chapter in volume 1 of the *Quartus II Handbook*. For information about synthesis attributes in other synthesis tools, refer to the appropriate chapter in the *Synthesis* section in volume 1 of the *Quartus II Handbook*.

Because formal verification tools do not support ROM megafunctions, Quartus II integrated synthesis does not infer ROM megafunctions when a formal verification tool is selected. When you are using a formal verification flow, Altera recommends that you instantiate ROM megafunction blocks in separate entities or modules that contain only the ROM logic, because you may need to treat the entity or module as a black box during formal verification.

The Verilog HDL and VHDL code samples in Example 14–29 through Example 14–32 on pages 14–37 through 14–39 infer synchronous ROM blocks. Depending on the device family’s dedicated RAM architecture, the ROM logic may have to be synchronous; refer to the device family handbook for details.

For device architectures with synchronous RAM blocks, such as the Arria series, Cyclone series, or Stratix series devices and newer device families, either the address or the output must be registered for synthesis software to infer a ROM block. When your design uses output registers, the synthesis software implements registers from the input registers of the RAM block without affecting the functionality of the ROM. If you register the address, the power-up state of the inferred ROM can be different from the HDL design. In this scenario, the synthesis software issues a warning. The Quartus II Help explains the condition under which the functionality changes when you use Quartus II integrated synthesis.
The ROM code examples in Example 14–29 through Example 14–32 on pages 14–37 through 14–39 map directly to the Altera memory architecture.

**Example 14–29. Verilog HDL Synchronous ROM**

```verilog
module sync_rom (clock, address, data_out);
    input clock;
    input [7:0] address;
    output [5:0] data_out;

    reg [5:0] data_out;

    always @ (posedge clock)
    begin
        case (address)
            8'b00000000: data_out = 6'b101111;
            8'b00000001: data_out = 6'b110110;
            ...
            8'b11111110: data_out = 6'b000001;
            8'b11111111: data_out = 6'b101010;
        endcase
    end
endmodule
```

**Example 14–30. VHDL Synchronous ROM**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;

ENTITY sync_rom IS
    PORT (
        clock: IN STD_LOGIC;
        address: IN STD_LOGIC_VECTOR(7 downto 0);
        data_out: OUT STD_LOGIC_VECTOR(5 downto 0)
    );
END sync_rom;

ARCHITECTURE rtl OF sync_rom IS
BEGIN
    PROCESS (clock)
    BEGIN
        IF rising_edge (clock) THEN
            CASE address IS
                WHEN "00000000" => data_out <= "101111";
                WHEN "00000001" => data_out <= "110110";
                ...
                WHEN "11111110" => data_out <= "000001";
                WHEN "11111111" => data_out <= "101010";
                WHEN OTHERS => data_out <= "101111";
            END CASE;
        END IF;
    END PROCESS;
END rtl;
```
Example 14–31. Verilog HDL Dual-Port Synchronous ROM Using readmemb

```verilog
module dual_port_rom (  
    input [(addr_width-1):0] addr_a, addr_b,  
    input clk,  
    output reg [(data_width-1):0] q_a, q_b  
);  
    parameter data_width = 8;  
    parameter addr_width = 8;  
    reg [data_width-1:0] rom[2**addr_width-1:0];  
    initial // Read the memory contents in the file  
    //dual_port_rom_init.txt.  
    begin  
        $readmemb("dual_port_rom_init.txt", rom);  
    end  
    always @ (posedge clk)  
    begin  
        q_a <= rom[addr_a];  
        q_b <= rom[addr_b];  
    end  
endmodule
```
Example 14–32. VHDL Dual-Port Synchronous ROM Using Initialization Function

library ieee;
use ieee.std_logic_1164.all;
use ieee.numeric_std.all;

entity dual_port_rom is
  generic (
    DATA_WIDTH : natural := 8;
    ADDR_WIDTH : natural := 8
  );
  port (
    clk : in std_logic;
    addr_a: in natural range 0 to 2**ADDR_WIDTH - 1;
    addr_b: in natural range 0 to 2**ADDR_WIDTH - 1;
    q_a : out std_logic_vector((DATA_WIDTH -1) downto 0);
    q_b : out std_logic_vector((DATA_WIDTH -1) downto 0)
  );
end entity;

architecture rtl of dual_port_rom is
  -- Build a 2-D array type for the ROM
  subtype word_t is std_logic_vector((DATA_WIDTH-1) downto 0);
  type memory_t is array(addr_a'high downto 0) of word_t;

  function init_rom
    return memory_t is
    variable tmp : memory_t := (others => (others => '0'));
    begin
      for addr_pos in 0 to 2**ADDR_WIDTH - 1 loop
        -- Initialize each address with the address itself
        tmp(addr_pos) := std_logic_vector(to_unsigned(addr_pos,
          DATA_WIDTH));
      end loop;
      return tmp;
    end init_rom;

  -- Declare the ROM signal and specify a default initialization value.
  signal rom : memory_t := init_rom;
begin
  process(clk)
  begin
    if (rising_edge(clk)) then
      q_a <= rom(addr_a);
      q_b <= rom(addr_b);
    end if;
  end process;
end rtl;
Shift Registers—Inferring the ALTSHIFT_TAPS Megafunction from HDL Code

To infer shift registers, synthesis tools detect a group of shift registers of the same length and convert them to an ALTSHIFT_TAPS megafunction. To be detected, all the shift registers must have the following characteristics:

- Use the same clock and clock enable
- Do not have any other secondary signals
- Have equally spaced taps that are at least three registers apart

When you use a formal verification flow, Altera recommends that you create shift register blocks in separate entities or modules containing only the shift register logic, because you might have to treat the entity or module as a black box during formal verification.

Because formal verification tools do not support shift register megafunctions, Quartus II integrated synthesis does not infer the ALTSHIFT_TAPS megafunction when a formal verification tool is selected. You can select EDA tools for use with your design on the EDA Tool Settings page of the Settings dialog box in the Quartus II software.

For more information about the ALTSHIFT_TAPS megafunction, refer to the ALTSHIFT_TAPS Megafunction User Guide.

Synthesis software recognizes shift registers only for device families that have dedicated RAM blocks, and the software uses certain guidelines to determine the best implementation.

Quartus II integrated synthesis uses the following guidelines which are common in other EDA tools. The Quartus II software determines whether to infer the ALTSHIFT_TAPS megafunction based on the width of the registered bus (W), the length between each tap (L), and the number of taps (N). If the Auto Shift Register Recognition setting is set to Auto, Quartus II integrated synthesis uses the Optimization Technique setting, logic and RAM utilization information about the design, and timing information from Timing-Driven Synthesis to determine which shift registers are implemented in RAM blocks for logic.

- If the registered bus width is one \( (W = 1) \), the software infers ALTSHIFT_TAPS if the number of taps times the length between each tap is greater than or equal to 64 \( (N \times L \geq 64) \).
- If the registered bus width is greater than one \( (W > 1) \), the software infers ALTSHIFT_TAPS if the registered bus width times the number of taps times the length between each tap is greater than or equal to 32 \( (W \times N \times L \geq 32) \).

If the length between each tap \( (L) \) is not a power of two, the software uses more logic to decode the read and write counters. This situation occurs because for different sizes of shift registers, external decode logic that uses logic elements (LEs) or ALMs is required to implement the function. This decode logic eliminates the performance and utilization advantages of implementing shift registers in memory.
The registers that the software maps to the ALTSHIFT_TAPS megafunction and places in RAM are not available in a Verilog HDL or VHDL output file for simulation tools because their node names do not exist after synthesis.

If your design uses a shift enable signal to infer a shift register, the shift register will not be implemented into MLAB memory, but can use only dedicated RAM blocks.

**Simple Shift Register**

The code samples in Example 14–33 and Example 14–34 show a simple, single-bit wide, 64-bit long shift register. The synthesis software implements the register \( W = 1 \) and \( M = 64 \) in an ALTSHIFT_TAPS megafunction for supported devices and maps it to RAM in supported devices, which may be placed in dedicated RAM blocks or MLAB memory. If the length of the register is less than 64 bits, the software implements the shift register in logic.

**Example 14–33. Verilog HDL Single-Bit Wide, 64-Bit Long Shift Register**

```verilog
code
module shift_1x64 (clk, shift, sr_in, sr_out);
    input clk, shift;
    input sr_in;
    output sr_out;
    reg [63:0] sr;
    always @(posedge clk)
    begin
        if (shift == 1'b1)
            begin
                sr[63:1] <= sr[62:0];
                sr[0] <= sr_in;
            end
        assign sr_out = sr[63];
    end
endmodule
code
```
Example 14–34. VHDL Single-Bit Wide, 64-Bit Long Shift Register

LIBRARY IEEE;
USE IEEE.STD_LOGIC_1164.all;
ENTITY shift_1x64 IS
    PORT ( 
        clk: IN STD_LOGIC;
        shift: IN STD_LOGIC;
        sr_in: IN STD_LOGIC;
        sr_out: OUT STD_LOGIC 
    );
END shift_1x64;

ARCHITECTURE arch OF shift_1x64 IS
    TYPE sr_length IS ARRAY (63 DOWNTO 0) OF STD_LOGIC;
    SIGNAL sr: sr_length;
BEGIN
    PROCESS (clk)
    BEGIN
        IF (clk'EVENT and clk = '1') THEN
            IF (shift = '1') THEN
                sr(63 DOWNTO 1) <= sr(62 DOWNTO 0);
                sr(0) <= sr_in;
            END IF;
        END IF;
    END PROCESS;
    sr_out <= sr(63);
END arch;
Shift Register with Evenly Spaced Taps

The code samples in Example 14–35 and Example 14–36 show a Verilog HDL and VHDL 8-bit wide, 64-bit long shift register ($W > 1$ and $M = 64$) with evenly spaced taps at 15, 31, and 47. The synthesis software implements this function in a single ALTSHIFT_TAPS megafunction and maps it to RAM in supported devices, which is allowed placement in dedicated RAM blocks or MLAB memory.

Example 14–35. Verilog HDL 8-Bit Wide, 64-Bit Long Shift Register with Evenly Spaced Taps

```verilog
module shift_8x64_taps (clk, shift, sr_in, sr_out, sr_tap_one, sr_tap_two, sr_tap_three);
    input clk, shift;
    input [7:0] sr_in;
    output [7:0] sr_tap_one, sr_tap_two, sr_tap_three, sr_out;

    reg [7:0] sr [63:0];
    integer n;

    always @ (posedge clk)
    begin
        if (shift == 1'b1)
            begin
                for (n = 63; n>0; n = n-1)
                    begin
                        sr[n] <= sr[n-1];
                    end
                sr[0] <= sr_in;
            end
        end

    assign sr_tap_one = sr[15];
    assign sr_tap_two = sr[31];
    assign sr_tap_three = sr[47];
    assign sr_out = sr[63];
endmodule
```
This section provides device-specific coding recommendations for Altera registers and latches. Understanding the architecture of the target Altera device helps ensure that your code produces the expected results and achieves the optimal quality of results.

This section provides guidelines in the following areas:

- “Register Power-Up Values in Altera Devices”
- “Secondary Register Control Signals Such as Clear and Clock Enable” on page 14-46
- “Latches” on page 14-50

Register Power-Up Values in Altera Devices

Registers in the device core always power up to a low (0) logic level on all Altera devices. If your design specifies a power-up level other than 0, synthesis tools can implement logic that causes registers to behave as if they were powering up to a high (1) logic level.

Example 14-36. VHDL 8-Bit Wide, 64-Bit Long Shift Register with Evenly Spaced Taps

LIBRARY IEEE;
USE IEEE.STD_LOGIC_1164.all;
ENTITY shift_8x64_taps IS
PORT (  
  clk: IN STD_LOGIC;
  shift: IN STD_LOGIC;
  sr_in: IN STD_LOGIC_VECTOR(7 DOWNTO 0);
  sr_tap_one: OUT STD_LOGIC_VECTOR(7 DOWNTO 0);
  sr_tap_two : OUT STD_LOGIC_VECTOR(7 DOWNTO 0);
  sr_tap_three: OUT STD_LOGIC_VECTOR(7 DOWNTO 0);
  sr_out: OUT STD_LOGIC_VECTOR(7 DOWNTO 0)
);
END shift_8x64_taps;

ARCHITECTURE arch OF shift_8x64_taps IS
  SUBTYPE sr_width IS STD_LOGIC_VECTOR(7 DOWNTO 0);
  TYPE sr_length IS ARRAY (63 DOWNTO 0) OF sr_width;
  SIGNAL sr: sr_length;
BEGIN
  PROCESS (clk)
  BEGIN
    IF (clk'EVENT and clk = '1') THEN
      IF (shift = '1') THEN
        sr(63 DOWNTO 1) <= sr(62 DOWNTO 0);
        sr(0) <= sr_in;
      END IF;
    END IF;
  END PROCESS;
  sr_tap_one <= sr(15);
  sr_tap_two <= sr(31);
  sr_tap_three <= sr(47);
  sr_out <= sr(63);
END arch;
If your design uses a preset signal on a device that does not support presets in the register architecture, your synthesis tool may convert the preset signal to a clear signal, which requires synthesis to perform an optimization referred to as NOT gate push-back. NOT gate push-back adds an inverter to the input and the output of the register so that the reset and power-up conditions will appear to be high, and the device operates as expected. In this case, your synthesis tool may issue a message informing you about the power-up condition. The register itself powers up low, but the register output is inverted, so the signal that arrives at all destinations is high.

Due to these effects, if you specify a non-zero reset value, you may cause your synthesis tool to use the asynchronous clear (aclr) signals available on the registers to implement the high bits with NOT gate push-back. In that case, the registers look as though they power up to the specified reset value.

When an asynchronous load (aload) signal is available in the device registers, your synthesis tools can implement a reset of 1 or 0 value by using an asynchronous load of 1 or 0. When the synthesis tool uses a load signal, it is not performing NOT gate push-back, so the registers power up to a 0 logic level.

For additional details, refer to the appropriate device family handbook or the appropriate handbook on the Altera website.

Designers typically use an explicit reset signal for the design, which forces all registers into their appropriate values after reset. Altera recommends this practice to reset the device after power-up to restore the proper state.

You can make your design more stable and avoid potential glitches by synchronizing external or combinational logic of the device architecture before you drive the asynchronous control ports of registers.

For additional information about good synchronous design practices, refer to the Design Recommendations for Altera Devices and the Quartus II Design Assistant chapter in volume 1 of the Quartus II Handbook.

**Specifying a Power-Up Value**

If you want to force a particular power-up condition for your design, you can use the synthesis options available in your synthesis tool. With Quartus II integrated synthesis, you can apply the **Power-Up Level** logic option. You can also apply the option with an altera_attribute assignment in your source code. Using this option forces synthesis to perform NOT gate push-back because synthesis tools cannot actually change the power-up states of core registers.

You can apply the Quartus II integrated synthesis **Power-Up Level** logic option to a specific register or to a design entity, module, or subdesign. If you do so, every register in that block receives the value. Registers power up to 0 by default; therefore, you can use this assignment to force all registers to power up to 1 using NOT gate push-back.

Setting the **Power-Up Level** to a logic level of high for a large design entity could degrade the quality of results due to the number of inverters that are required. In some situations, issues are caused by enable signal inference or secondary control logic inference. It may also be more difficult to migrate such a design to an ASIC or a HardCopy® device.
You can simulate the power-up behavior in a functional simulation if you use initialization.

The **Power-Up Level** option and the `altera_attribute` assignment are described in the *Quartus II Integrated Synthesis* chapter in volume 1 of the *Quartus II Handbook*.

Some synthesis tools can also read the default or initial values for registered signals and implement this behavior in the device. For example, Quartus II integrated synthesis converts default values for registered signals into **Power-Up Level** settings. When the Quartus II software reads the default values, the synthesized behavior matches the power-up state of the HDL code during a functional simulation.

For example, the code samples in Example 14–37 and Example 14–38 both infer a register for `q` and set its power-up level to high.

**Example 14–37. Verilog Register with High Power-Up Value**

```verilog
reg q = 1'b1; //q has a default value of '1'

always @(posedge clk)
begin
    q <= d;
end
```

**Example 14–38. VHDL Register with High Power-Up Level**

```vhdl
SIGNAL q : STD_LOGIC := '1'; -- q has a default value of '1'

PROCESS (clk, reset)
BEGIN
    IF (rising_edge(clk)) THEN
        q <= d;
    END IF;
END PROCESS;
```

There may also be undeclared default power-up conditions based on signal type. If you declare a VHDL register signal as an integer, Quartus II synthesis attempts to use the left end of the integer range as the power-up value. For the default signed integer type, the default power-up value is the highest magnitude negative integer (100…001). For an unsigned integer type, the default power-up value is 0.

If the target device architecture does not support two asynchronous control signals, such as aclr and aload, you cannot set a different power-up state and reset state. If the NOT gate push-back algorithm creates logic to set a register to 1, that register will power-up high. If you set a different power-up condition through a synthesis assignment or initial value, the power-up level is ignored during synthesis.

**Secondary Register Control Signals Such as Clear and Clock Enable**

The registers in Altera FPGAs provide a number of secondary control signals (such as clear and enable signals) that you can use to implement control logic for each register without using extra logic cells. Device families vary in their support for secondary signals, so consult the device family data sheet to verify which signals are available in your target device.
To make the most efficient use of the signals in the device, your HDL code should match the device architecture as closely as possible. The control signals have a certain priority due to the nature of the architecture, so your HDL code should follow that priority where possible.

Your synthesis tool can emulate any control signals using regular logic, so achieving functionally correct results is always possible. However, if your design requirements are flexible in terms of which control signals are used and in what priority, match your design to the target device architecture to achieve the most efficient results. If the priority of the signals in your design is not the same as that of the target architecture, extra logic may be required to implement the control signals. This extra logic uses additional device resources and can cause additional delays for the control signals.

In addition, there are certain cases where using logic other than the dedicated control logic in the device architecture can have a larger impact. For example, the clock enable signal has priority over the synchronous reset or clear signal in the device architecture. The clock enable turns off the clock line in the LAB, and the clear signal is synchronous. Therefore, in the device architecture, the synchronous clear takes effect only when a clock edge occurs.

If you code a register with a synchronous clear signal that has priority over the clock enable signal, the software must emulate the clock enable functionality using data inputs to the registers. Because the signal does not use the clock enable port of a register, you cannot apply a Clock Enable Multicycle constraint. In this case, following the priority of signals available in the device is clearly the best choice for the priority of these control signals, and using a different priority causes unexpected results with an assignment to the clock enable signal.

The priority order for secondary control signals in Altera devices differs from the order for other vendors’ devices. If your design requirements are flexible regarding priority, verify that the secondary control signals meet design performance requirements when migrating designs between FPGA vendors and try to match your target device architecture to achieve the best results.

The signal order is the same for all Altera device families, although, as noted previously, not all device families provide every signal. The following priority order is observed:

1. Asynchronous Clear, aclr—highest priority
2. Asynchronous Load, aload
3. Enable, ena
4. Synchronous Clear, sclr
5. Synchronous Load, sload
6. Data In, data—lowest priority

The following examples provide Verilog HDL and VHDL code that creates a register with the aclr, aload, and ena control signals.

MAX® 3000 and MAX 7000 devices include a dedicated preset signal, which has second priority after aclr, and is not included in the following examples.
The Verilog HDL example (Example 14–39) does not have adata on the sensitivity list, but the VHDL example (Example 14–40) does. This is a limitation of the Verilog HDL language—there is no way to describe an asynchronous load signal (in which q toggles if adata toggles while aload is high). All synthesis tools should infer an aload signal from this construct despite this limitation. When they perform such inference, you may see information or warning messages from the synthesis tool.

Example 14–39. Verilog HDL D-Type Flipflop (Register) with ena, aclr, and preload Control Signals

```verilog
module dff_control(clk, aclr, preload, ena, data, adata, q);
  input clk, aclr, preload, ena, data, adata;
  output q;
  reg q;
  always @ (posedge clk or posedge aclr or posedge preload)
    begin
      if (aclr)
        q <= 1'b0;
      else if (preload)
        q <= adata;
      else if (ena)
        q <= data;
    end
endmodule
```

Example 14–40. VHDL D-Type Flipflop (Register) with ena, aclr, and preload Control Signals (Part 1 of 2)

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;

ENTITY dff_control IS
  PORT (  
    clk: IN STD_LOGIC;
    aclr: IN STD_LOGIC;
    preload: IN STD_LOGIC;
    data: IN STD_LOGIC;
    ena: IN STD_LOGIC;
    data: IN STD_LOGIC;
    q: OUT STD_LOGIC  
  );
END dff_control;
```
Creating many registers with different sload and sclr signals can make packing the registers into LABs difficult for the Quartus II Fitter because the sclr and sload signals are LAB-wide signals. In addition, using the LAB-wide sload signal prevents the Fitter from packing registers using the quick feedback path in the device architecture, which means that some registers cannot be packed with other logic.

Synthesis tools typically restrict use of sload and sclr signals to cases in which there are enough registers with common signals to allow good LAB packing. Using the look-up table (LUT) to implement the signals is always more flexible if it is available. Because different device families offer different numbers of control signals, inference of these signals is also device-specific. For example, because Stratix II devices have more flexibility than Stratix devices with respect to secondary control signals, synthesis tools might infer more sload and sclr signals for Stratix II devices.

If you use these additional control signals, use them in the priority order that matches the device architecture. To achieve the most efficient results, ensure the sclr signal has a higher priority than the sload signal in the same way that aclr has higher priority thanaload in the previous examples. Remember that the register signals are not inferred unless the design meets the conditions described previously. However, if your HDL described the desired behavior, the software always implements logic with the correct functionality.

In Verilog HDL, the following code for sload and sclr could replace the if (ena) q <= data; statements in the Verilog HDL in Example 14–39 (after adding the control signals to the module declaration).

**Example 14–41. Verilog HDL sload and sclr Control Signals**

```verilog
if (ena) begin
    if (sclr)
        q <= 1'b0;
    else if (sload)
        q <= sdata;
    else
        q <= data;
end
```
In VHDL, the following code for `sload` and `sclr` could replace the IF (ena = '1')
THEN q <= data; END IF; statements in the VHDL in Example 14–40 on page 14–48
(after adding the control signals to the entity declaration).

**Example 14–42. VHDL sload and sclr Control Signals**

```
IF (ena = '1') THEN
    IF (sclr = '1') THEN
        q <= '0';
    ELSIF (sload = '1') THEN
        q <= sdata;
    ELSE
        q <= data;
    END IF;
END IF;
```

### Latches

A latch is a small combinational loop that holds the value of a signal until a new value
is assigned.

> Altera recommends that you design without the use of latches whenever possible.

For additional information about the issues involved in designing with latches and
combinational loops, refer to the *Design Recommendations for Altera Devices and the
Quartus II Design Assistant* chapter in volume 1 of the *Quartus II Handbook*.

Latches can be inferred from HDL code when you did not intend to use a latch, as
described in “Unintentional Latch Generation”. If you do intend to infer a latch, it is
important to infer it correctly to guarantee correct device operation as detailed in
“Inferring Latches Correctly” on page 14–51.

### Unintentional Latch Generation

When you are designing combinational logic, certain coding styles can create an
unintentional latch. For example, when `CASE` or `IF` statements do not cover all possible
input conditions, latches may be required to hold the output if a new output value is
not assigned. Check your synthesis tool messages for references to inferred latches. If
your code unintentionally creates a latch, make code changes to remove the latch.

A latch is required if a signal is assigned a value outside of a clock edge (for example,
with an asynchronous reset), but is not assigned a value in an edge-triggered design
block. An unintentional latch may be generated if your HDL code assigns a value to a
signal in an edge-triggered design block, but that logic is removed during synthesis.
For example, when a `CASE` or `IF` statement tests the value of a condition with a
parameter or generic that evaluates to FALSE, any logic or signal assignment in that
statement is not required and is optimized away during synthesis. This optimization
may result in a latch being generated for the signal.

> Latches have limited support in formal verification tools. Therefore, ensure that you
do not infer latches unintentionally.
The full_case attribute can be used in Verilog HDL designs to treat unspecified cases as don’t care values (X). However, using the full_case attribute can cause simulation mismatches because this attribute is a synthesis-only attribute, so simulation tools still treat the unspecified cases as latches.

For more information about using attributes in your synthesis tool, refer to the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook. The Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook provides an example explaining possible simulation mismatches.

Omitting the final else or when others clause in an if or case statement can also generate a latch. Don’t care (X) assignments on the default conditions are useful in preventing latch generation. For the best logic optimization, assign the default case or final else value to don’t care (X) instead of a logic value.

The VHDL code sample in Example 14–43 prevents unintentional latches. Without the final else clause, this code creates unintentional latches to cover the remaining combinations of the sel inputs. When you are targeting a Stratix device with this code, omitting the final else condition can cause the synthesis software to use up to six LEs, instead of the three it uses with the else statement. Additionally, assigning the final else clause to 1 instead of X can result in slightly more LEs, because the synthesis software cannot perform as much optimization when you specify a constant value compared to a don’t care value.

Example 14–43. VHDL Code Preventing Unintentional Latch Creation

LIBRARY ieee;
USE IEEE.std_logic_1164.all;

ENTITY nolatch IS
    PORT (a,b,c: IN STD_LOGIC;
          sel: IN STD_LOGIC_VECTOR (4 DOWNTO 0);
          oput: OUT STD_LOGIC);
END nolatch;

ARCHITECTURE rtl OF nolatch IS
BEGIN
    PROCESS (a,b,c,sel) BEGIN
        if sel = "00000" THEN
            oput <= a;
        ELSIF sel = "00001" THEN
            oput <= b;
        ELSIF sel = "00010" THEN
            oput <= c;
        ELSE
            oput <= 'X'; -- Prevents latch inference
        END if;
    END PROCESS;
END rtl;

Inferring Latches Correctly

Synthesis tools can infer a latch that does not exhibit the glitch and timing hazard problems typically associated with combinational loops.
Any use of latches generates warnings and is flagged if the design is migrated to a HardCopy ASIC. In addition, timing analysis does not completely model latch timing in some cases. Do not use latches unless required by your design, and you fully understand the impact of using the latches.

When using Quartus II integrated synthesis, latches that are inferred by the software are reported in the **User-Specified and Inferred Latches** section of the Compilation Report. This report indicates whether the latch is considered safe and free of timing hazards.

If a latch or combinational loop in your design is not listed in the **User-Specified and Inferred Latches** section, it means that it was not inferred as a safe latch by the software and is not considered glitch-free.

All combinational loops listed in the **Analysis & Synthesis Logic Cells Representing Combinational Loops** table in the Compilation Report are at risk of timing hazards. These entries indicate possible problems with your design that you should investigate. However, it is possible to have a correct design that includes combinational loops. For example, it is possible that the combinational loop cannot be sensitized. This can occur in cases where there is an electrical path in the hardware, but either the designer knows that the circuit never encounters data that causes that path to be activated, or the surrounding logic is set up in a mutually exclusive manner that prevents that path from ever being sensitized, independent of the data input.

For macrocell-based devices, such as MAX 7000 and MAX 3000, all data (D-type) latches and set-reset (S-R) latches listed in the **Analysis & Synthesis User-Specified and Inferred Latches** table have an implementation free of timing hazards, such as glitches. The implementation includes both a cover term to ensure there is no glitching and a single macrocell in the feedback loop.

For 4-input LUT-based devices, such as Stratix devices, the Cyclone series, and MAX II devices, all latches in the **User-Specified and Inferred Latches** table with a single LUT in the feedback loop are free of timing hazards when a single input changes. Because of the hardware behavior of the LUT, the output does not glitch when a single input toggles between two values that are supposed to produce the same output value, such as a D-type input toggling when the enable input is inactive or a set input toggling when a reset input with higher priority is active. This hardware behavior of the LUT means that no cover term is required for a loop around a single LUT. The Quartus II software uses a single LUT in the feedback loop whenever possible. A latch that has data, enable, set, and reset inputs in addition to the output fed back to the input cannot be implemented in a single 4-input LUT. If the Quartus II software cannot implement the latch with a single-LUT loop because there are too many inputs, the **User-Specified and Inferred Latches** table indicates that the latch is not free of timing hazards.

For 6-input LUT-based devices, the software can implement all latch inputs with a single adaptive look-up table (ALUT) in the combinational loop. Therefore, all latches in the **User-Specified and Inferred Latches** table are free of timing hazards when a single input changes.

If a latch is listed as a safe latch, other optimizations performed by the Quartus II software, such as physical synthesis netlist optimizations in the Fitter, maintain the hazard-free performance.
To ensure hazard-free behavior, only one control input can change at a time. Changing two inputs simultaneously, such as deasserting set and reset at the same time, or changing data and enable at the same time, can produce incorrect behavior in any latch.

Quartus II integrated synthesis infers latches from `always` blocks in Verilog HDL and `process` statements in VHDL, but not from continuous assignments in Verilog HDL or concurrent signal assignments in VHDL. These rules are the same as for register inference. The software infers registers or flipflops only from `always` blocks and `process` statements.

The Verilog HDL code sample in Example 14–44 infers a S-R latch correctly in the Quartus II software.

**Example 14–44. Verilog HDL Set-Reset Latch**

```verilog
module simple_latch (  
  input SetTerm,  
  input ResetTerm,  
  output reg LatchOut 
);   
always @ (SetTerm or ResetTerm) begin  
  if (SetTerm)    
    LatchOut = 1'b1  
  else if (ResetTerm)  
    LatchOut = 1'b0  
end
endmodule
```

The VHDL code sample in Example 14–45 infers a D-type latch correctly in the Quartus II software.

**Example 14–45. VHDL Data Type Latch**

```vhdl
LIBRARY IEEE;
USE IEEE.std_logic_1164.all;

ENTITY simple_latch IS  
  PORT (    
    enable, data : IN STD_LOGIC;    
    q : OUT STD_LOGIC  
  );  
END simple_latch;

ARCHITECTURE rtl OF simple_latch IS  
BEGIN  
  latch : PROCESS (enable, data)  
  BEGIN  
    IF (enable = '1') THEN  
      q <= data;  
    END IF;  
  END PROCESS latch;  
END rtl;
```

The following example shows a Verilog HDL continuous assignment that does not infer a latch in the Quartus II software:

```verilog
assign latch_out = (~en & latch_out) | (en & data);
```
The behavior of the assignment is similar to a latch, but it may not function correctly as a latch, and its timing is not analyzed as a latch.

Quartus II integrated synthesis also creates safe latches when possible for instantiations of the LPM_LATCH megafuction. You can use this megafuction to create a latch with any combination of data, enable, set, and reset inputs. The same limitations apply for creating safe latches as for inferring latches from HDL code.

Inferring the Altera LPM_LATCH function in another synthesis tool ensures that the implementation is also recognized as a latch in the Quartus II software. If a third-party synthesis tool implements a latch using the LPM_LATCH megafuction, the Quartus II integrated synthesis lists the latch in the User-Specified and Inferred Latches table in the same way as it lists latches created in HDL source code. The coding style necessary to produce an LPM_LATCH implementation may depend on your synthesis tool. Some third-party synthesis tools list the number of LPM_LATCH functions that are inferred.

For LUT-based families, the Fitter uses global routing for control signals, including signals that Analysis and Synthesis identifies as latch enables. In some cases the global insertion delay may decrease the timing performance. If necessary, you can turn off the Quartus II Global Signal logic option to manually prevent the use of global signals. Global latch enables are listed in the Global & Other Fast Signals table in the Compilation Report.

General Coding Guidelines

This section helps you understand how synthesis tools map various types of HDL code into the target Altera device. Following Altera recommended coding styles, and in some cases designing logic structures to match the appropriate device architecture, can provide significant improvements in the design’s quality of results.

This section provides coding guidelines for the following logic structures:

- “Tri-State Signals”. This section explains how to create tri-state signals for bidirectional I/O pins.
- “Clock Multiplexing” on page 14–56. This section provides recommendations for multiplexing clock signals.
- “Adder Trees” on page 14–59. This section explains the different coding styles that lead to optimal results for devices with 4-input LUTs and 6-input ALUTs.
- “State Machines” on page 14–61. This section helps ensure the best results when you use state machines.
- “Multiplexers” on page 14–68. This section explains how multiplexers can be synthesized, addresses common problems, and provides guidelines to achieve optimal resource utilization.
- “Cyclic Redundancy Check Functions” on page 14–71. This section provides guidelines for getting good results when designing Cyclic Redundancy Check (CRC) functions.
- “Comparators” on page 14–73. This section explains different comparator implementations and provides suggestions for controlling the implementation.
“Counters” on page 14–74. This section provides guidelines to ensure your counter design targets the device architecture optimally.

**Tri-State Signals**

When you target Altera devices, you should use tri-state signals only when they are attached to top-level bidirectional or output pins. Avoid lower-level bidirectional pins, and avoid using the Z logic value unless it is driving an output or bidirectional pin.

Synthesis tools implement designs with internal tri-state signals correctly in Altera devices using multiplexer logic, but Altera does not recommend this coding practice.

In hierarchical block-based or incremental design flows, a hierarchical boundary cannot contain any bidirectional ports, unless the lower-level bidirectional port is connected directly through the hierarchy to a top-level output pin without connecting to any other design logic. If you use boundary tri-states in a lower-level block, synthesis software must push the tri-states through the hierarchy to the top level to make use of the tri-state drivers on output pins of Altera devices. Because pushing tri-states requires optimizing through hierarchies, lower-level tri-states are restricted with block-based design methodologies.

The code in Example 14–46 and Example 14–47 show Verilog HDL and VHDL code that creates tri-state bidirectional signals.

**Example 14–46. Verilog HDL Tri-State Signal**

```verilog
module tristate (myinput, myenable, mybidir);
    input myinput, myenable;
    inout mybidir;
    assign mybidir = (myenable ? myinput : 1'bZ);
endmodule
```

**Example 14–47. VHDL Tri-State Signal**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_arith.all;

ENTITY tristate IS
    PORT (mybidir : INOUT STD_LOGIC;
        myinput : IN STD_LOGIC;
        myenable : IN STD_LOGIC);
END tristate;

ARCHITECTURE rtl OF tristate IS
BEGIN
    mybidir <= 'Z' WHEN (myenable = '0') ELSE myinput;
END rtl;
```
# Clock Multiplexing

Clock multiplexing is sometimes used to operate the same logic function with different clock sources. This type of logic can introduce glitches that create functional problems, and the delay inherent in the combinational logic can lead to timing problems. Clock multiplexers trigger warnings from a wide range of design rule check and timing analysis tools.

Altera recommends using dedicated hardware to perform clock multiplexing when it is available, instead of using multiplexing logic. For example, you can use the Clock Switchover feature or the Clock Control Block available in certain Altera devices. These dedicated hardware blocks avoid glitches, ensure that you use global low-skew routing lines, and avoid any possible hold time problems on the device due to logic delay on the clock line. Many Altera devices also support dynamic PLL reconfiguration, which is the safest and most robust method of changing clock rates during device operation.


If you implement a clock multiplexer in logic cells because the design has too many clocks to use the clock control block, or if dynamic reconfiguration is too complex for your design, it is important to consider simultaneous toggling inputs and ensure glitch-free transitions.

Figure 14–2 shows a simple representation of a clock multiplexer (mux) in a device with 6-input LUTs.

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**Figure 14–2. Simple Clock Multiplexer in a 6-Input LUT**

The data sheet for your target device describes how LUT outputs may glitch during a simultaneous toggle of input signals, independent of the LUT function. Although, in practice, the 4:1 MUX function does not generate detectable glitches during simultaneous data input toggles, it is possible to construct cell implementations that do exhibit significant glitches, so this simple clock mux structure is not recommended. An additional problem with this implementation is that the output behaves erratically during a change in the clk_select signals. This behavior could create timing violations on all registers fed by the system clock and result in possible metastability.
A more sophisticated clock select structure can eliminate the simultaneous toggle and switching problems, as in Figure 14–3.

**Figure 14–3. Glitch-Free Clock Multiplexer Structure**

This structure can be generalized for any number of clock channels. Example 14–48 contains a parameterized version in Verilog HDL. The design enforces that no clock activates until all others have been inactive for at least a few cycles, and that activation occurs while the clock is low. The design applies a *synthesis_keep* directive to the AND gates on the right side of the figure, which ensures there are no simultaneous toggles on the input of the *clk_out* OR gate.
Switching from clock $A$ to clock $B$ requires that clock $A$ continue to operate for at least a few cycles. If the old clock stops immediately, the design sticks. The select signals are implemented as a “one-hot” control in this example, but you can use other encoding if you prefer. The input side logic is asynchronous and is not critical. This design can tolerate extreme glitching during the switch process.

Example 14–48. Verilog HDL Clock Multiplexing Design to Avoid Glitches

```verilog
module clock_mux (clk, clk_select, clk_out);

parameter num_clocks = 4;

input [num_clocks-1:0] clk;
input [num_clocks-1:0] clk_select; // one hot
output clk_out;

genvar i;

reg [num_clocks-1:0] ena_r0;
reg [num_clocks-1:0] ena_r1;
reg [num_clocks-1:0] ena_r2;
wire [num_clocks-1:0] qualified_sel;

// A look-up-table (LUT) can glitch when multiple inputs change simultaneously. Use the keep attribute to insert a hard logic cell buffer and prevent the unrelated clocks from appearing on the same LUT.
wire [num_clocks-1:0] gated_clks /* synthesis keep */;

initial begin
  ena_r0 = 0;
  ena_r1 = 0;
  ena_r2 = 0;
end

generate
  for (i=0; i<num_clocks; i=i+1)
  begin : lp0
    wire [num_clocks-1:0] tmp_mask;
    assign tmp_mask = {num_clocks{1'b1}} ^ (1 << i);
    assign qualified_sel[i] = clk_select[i] & (~|(ena_r2 & tmp_mask));
    
    always @(posedge clk[i]) begin
      ena_r0[i] <= qualified_sel[i] & (~|ena_r2 & tmp_mask));
    end
    always @(negedge clk[i]) begin
      ena_r2[i] <= ena_r1[i];
    end
    assign gated_clks[i] = clk[i] & ena_r2[i];
  end
endgenerate

// These will not exhibit simultaneous toggle by construction
assign clk_out = |gated_clks;
endmodule
```
Adder Trees

Structuring adder trees appropriately to match your targeted Altera device architecture can result in significant performance and density improvements. A good example of an application using a large adder tree is a finite impulse response (FIR) correlator. Using a pipelined binary or ternary adder tree appropriately can greatly improve the quality of your results.

This section explains why coding recommendations are different for Altera 4-input LUT devices and 6-input LUT devices.

Architectures with 4-Input LUTs in Logic Elements

Architectures such as Stratix devices and the Cyclone series of devices contain 4-input LUTs as the standard combinational structure in the LE.

If your design can tolerate pipelining, the fastest way to add three numbers $A$, $B$, and $C$ in devices that use 4-input lookup tables is to add $A + B$, register the output, and then add the registered output to $C$. Adding $A + B$ takes one level of logic (one bit is added in one LE), so this runs at full clock speed. This can be extended to as many numbers as desired.

Example 14–49 shows five numbers $A$, $B$, $C$, $D$, and $E$ are added. Adding five numbers in devices that use 4-input lookup tables requires four adders and three levels of registers for a total of 64 LEs (for 16-bit numbers).

**Example 14–49. Verilog HDL Pipelined Binary Tree**

```verilog
module binary_adder_tree (a, b, c, d, e, clk, out);
    parameter width = 16;
    input [width-1:0] a, b, c, d, e;
    input clk;
    output [width-1:0] out;
    wire [width-1:0] sum1, sum2, sum3, sum4;
    reg [width-1:0] sumreg1, sumreg2, sumreg3, sumreg4;
    // Registers
    always @(posedge CLK)
        begin
            sumreg1 <= sum1;
            sumreg2 <= sum2;
            sumreg3 <= sum3;
            sumreg4 <= sum4;
        end
    // 2-bit additions
    assign sum1 = A + B;
    assign sum2 = C + D;
    assign sum3 = sumreg1 + sumreg2;
    assign sum4 = sumreg3 + E;
    assign out = sumreg4;
endmodule
```
Architectures with 6-Input LUTs in Adaptive Logic Modules

High-performance Altera device families use a 6-input LUT in their basic logic structure, so these devices benefit from a different coding style from the previous example presented for 4-input LUTs. Specifically, in these devices, ALMs can simultaneously add three bits. Therefore, the tree in Example 14–49 must be two levels deep and contain just two add-by-three inputs instead of four add-by-two inputs.

Although the code in the previous example compiles successfully for 6-input LUT devices, the code is inefficient and does not take advantage of the 6-input adaptive ALUT. By restructuring the tree as a ternary tree, the design becomes much more efficient, significantly improving density utilization. Therefore, when you are targeting with ALUTs and ALMs, large pipelined binary adder trees designed for 4-input LUT architectures should be rewritten to take advantage of the advanced device architecture.

Example 14–50 uses just 32 ALUTs in a Stratix II device—more than a 4:1 advantage over the number of LUTs in the prior example implemented in a Stratix device.

You cannot pack a LAB full when using this type of coding style because of the number of LAB inputs. However, in a typical design, the Quartus II Fitter can pack other logic into each LAB to take advantage of the unused ALMs.

Example 14–50. Verilog HDL Pipelined Ternary Tree

```verilog
module ternary_adder_tree (a, b, c, d, e, clk, out);
    parameter width = 16;
    input [width-1:0] a, b, c, d, e;
    input clk;
    output [width-1:0] out;
    wire [width-1:0] sum1, sum2;
    reg [width-1:0] sumreg1, sumreg2;
    // registers
    always @ (posedge clk)
        begin
            sumreg1 <= sum1;
            sumreg2 <= sum2;
        end
    // 3-bit additions
    assign sum1 = a + b + c;
    assign sum2 = sumreg1 + d + e;
    assign out = sumreg2;
endmodule
```

These examples show pipelined adders, but partitioning your addition operations can help you achieve better results in nonpipelined adders as well. If your design is not pipelined, a ternary tree provides much better performance than a binary tree. For example, depending on your synthesis tool, the HDL code

\[ \text{sum} = (A + B + C) + (D + E) \]

is more likely to create the optimal implementation of a 3-input adder for \(A + B + C\) followed by a 3-input adder for \(\text{sum} + D + E\) than the code without the parentheses. If you do not add the parentheses, the synthesis tool may partition the addition in a way that is not optimal for the architecture.
State Machines

Synthesis tools can recognize and encode Verilog HDL and VHDL state machines during synthesis. This section presents guidelines to ensure the best results when you use state machines. Ensuring that your synthesis tool recognizes a piece of code as a state machine allows the tool to recode the state variables to improve the quality of results, and allows the tool to use the known properties of state machines to optimize other parts of the design. When synthesis recognizes a state machine, it is often able to improve the design area and performance.

To achieve the best results on average, synthesis tools often use one-hot encoding for FPGA devices and minimal-bit encoding for CPLD devices, although the choice of implementation can vary for different state machines and different devices. Refer to your synthesis tool documentation for specific ways to control the manner in which state machines are encoded.

For information about state machine encoding in Quartus II integrated synthesis, refer to the State Machine Processing section in the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook.

To ensure proper recognition and inference of state machines and to improve the quality of results, Altera recommends that you observe the following guidelines, which apply to both Verilog HDL and VHDL:

- Assign default values to outputs derived from the state machine so that synthesis does not generate unwanted latches.
- Separate the state machine logic from all arithmetic functions and data paths, including assigning output values.
- If your design contains an operation that is used by more than one state, define the operation outside the state machine and cause the output logic of the state machine to use this value.
- Use a simple asynchronous or synchronous reset to ensure a defined power-up state. If your state machine design contains more elaborate reset logic, such as both an asynchronous reset and an asynchronous load, the Quartus II software generates regular logic rather than inferring a state machine.

If a state machine enters an illegal state due to a problem with the device, the design likely ceases to function correctly until the next reset of the state machine. Synthesis tools do not provide for this situation by default. The same issue applies to any other registers if there is some kind of fault in the system. A default or when others clause does not affect this operation, assuming that your design never deliberately enters this state. Synthesis tools remove any logic generated by a default state if it is not reachable by normal state machine operation.

Many synthesis tools (including Quartus II integrated synthesis) have an option to implement a safe state machine. The software inserts extra logic to detect an illegal state and force the state machine’s transition to the reset state. It is commonly used when the state machine can enter an illegal state. The most common cause of this situation is a state machine that has control inputs that come from another clock domain, such as the control logic for a dual-clock FIFO.
This option protects only state machines by forcing them into the reset state. All other registers in the design are not protected this way. If the design has asynchronous inputs, Altera recommends using a synchronization register chain instead of relying on the safe state machine option.

For additional information about tool-specific options for implementing state machines, refer to the tool vendor’s documentation or the appropriate chapter in the Synthesis section in volume 1 of the Quartus II Handbook.

The following two sections, “Verilog HDL State Machines” and “VHDL State Machines” on page 14–66, describe additional language-specific guidelines and coding examples.

**Verilog HDL State Machines**

To ensure proper recognition and inference of Verilog HDL state machines, observe the following additional Verilog HDL guidelines. Some of these guidelines may be specific to Quartus II integrated synthesis. Refer to your synthesis tool documentation for specific coding recommendations.

If the state machine is not recognized and inferred by the synthesis software (such as Quartus II integrated synthesis), the state machine is implemented as regular logic gates and registers, and the state machine is not listed as a state machine in the Analysis & Synthesis section of the Quartus II Compilation Report. In this case, the software does not perform any of the optimizations that are specific to state machines.

- If you are using the SystemVerilog standard, use enumerated types to describe state machines. For more information, refer too “SystemVerilog State Machine Coding Example” on page 14–65.
- Represent the states in a state machine with the parameter data types in Verilog-1995 and Verilog-2001, and use the parameters to make state assignments. For more information, refer too “Verilog-2001 State Machine Coding Example” on page 14–63. This parameter implementation makes the state machine easier to read and reduces the risk of errors during coding.

  ALTERA recommends against the direct use of integer values for state variables, such as \texttt{next\_state <= 0}. However, using an integer does not prevent inference in the Quartus II software.

- No state machine is inferred in the Quartus II software if the state transition logic uses arithmetic similar to that in the following example:

  ```verilog
  case (state)
    0: begin
      if (ena) next\_state <= state + 2;
      else next\_state <= state + 1;
    end
    1: begin
      ... begin
    endcase
  ```

- No state machine is inferred in the Quartus II software if the state variable is an output.

- No state machine is inferred in the Quartus II software for signed variables.
Verilog-2001 State Machine Coding Example

The following module verilog fsm is an example of a typical Verilog HDL state machine implementation (Example 14–51).

This state machine has five states. The asynchronous reset sets the variable state to state_0. The sum of in_1 and in_2 is an output of the state machine in state_1 and state_2. The difference (in_1 – in_2) is also used in state_1 and state_2. The temporary variables tmp_out_0 and tmp_out_1 store the sum and the difference of in_1 and in_2. Using these temporary variables in the various states of the state machine ensures proper resource sharing between the mutually exclusive states.
Example 14–51. Verilog-2001 State Machine

module verilog_fsm (clk, reset, in_1, in_2, out);
  input clk, reset;
  input [3:0] in_1, in_2;
  output [4:0] out;
  parameter state_0 = 3'b000;
  parameter state_1 = 3'b001;
  parameter state_2 = 3'b010;
  parameter state_3 = 3'b011;
  parameter state_4 = 3'b100;

  reg [4:0] tmp_out_0, tmp_out_1, tmp_out_2;
  reg [2:0] state, next_state;

  always @(posedge clk or posedge reset)
  begin
    if (reset)
      state <= state_0;
    else
      state <= next_state;
  end

  always @(*)
  begin
    tmp_out_0 = in_1 + in_2;
    tmp_out_1 = in_1 - in_2;
    case (state)
      state_0: begin
        tmp_out_2 = in_1 + 5'b00001;
        next_state = state_1;
      end
      state_1: begin
        if (in_1 < in_2) begin
          next_state = state_2;
          tmp_out_2 = tmp_out_0;
        end
        else begin
          next_state = state_3;
          tmp_out_2 = tmp_out_1;
        end
      end
      state_2: begin
        tmp_out_2 = tmp_out_0 - 5'b00001;
        next_state = state_3;
      end
      state_3: begin
        tmp_out_2 = tmp_out_1 + 5'b00001;
        next_state = state_0;
      end
      state_4: begin
        tmp_out_2 = in_2 + 5'b00001;
        next_state = state_0;
      end
      default: begin
        tmp_out_2 = 5'b00000;
        next_state = state_0;
      end
    endcase
  end

  assign out = tmp_out_2;
endmodule
An equivalent implementation of this state machine can be achieved by using `define instead of the parameter data type, as follows:

```vhdl`
define state_0 3'b000
define state_1 3'b001
define state_2 3'b010
define state_3 3'b011
define state_4 3'b100```

In this case, the `state` and `next_state` assignments are assigned a `state_x` instead of a `state_x`, for example:

```vhdl```
next_state <= `state_3;
```

Although the `define` construct is supported, Altera strongly recommends the use of the parameter data type because doing so preserves the state names throughout synthesis.

### SystemVerilog State Machine Coding Example

The module `enum_fsm` in Example 14–52 is an example of a SystemVerilog state machine implementation that uses enumerated types. Altera recommends using this coding style to describe state machines in SystemVerilog.

In Quartus II integrated synthesis, the enumerated type that defines the states for the state machine must be of an unsigned integer type as in Example 14–52. If you do not specify the enumerated type as `int unsigned`, a signed `int` type is used by default. In this case, the Quartus II integrated synthesis synthesizes the design, but does not infer or optimize the logic as a state machine.
VHDL State Machines

To ensure proper recognition and inference of VHDL state machines, represent the states in a state machine with enumerated types and use the corresponding types to make state assignments. This implementation makes the state machine easier to read and reduces the risk of errors during coding. If the state is not represented by an enumerated type, synthesis software (such as Quartus II integrated synthesis) does not recognize the state machine. Instead, the state machine is implemented as regular logic gates and registers and the state machine is not listed as a state machine in the Analysis & Synthesis section of the Quartus II Compilation Report. In this case, the software does not perform any of the optimizations that are specific to state machines.

VHDL State Machine Coding Example

The following entity, vhd1_fsm, is an example of a typical VHDL state machine implementation (Example 14–53).

This state machine has five states. The asynchronous reset sets the variable state to state_0. The sum of in1 and in2 is an output of the state machine in state_1 and state_2. The difference (in1 - in2) is also used in state_1 and state_2. The temporary variables tmp_out_0 and tmp_out_1 store the sum and the difference of in1 and in2. Using these temporary variables in the various states of the state machine ensures proper resource sharing between the mutually exclusive states.
**Example 14–53. VHDL State Machine**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.numeric_std.all;
ENTITY vhdl_fsm IS
  PORT(
    clk: IN STD_LOGIC;
    reset: IN STD_LOGIC;
    in1: IN UNSIGNED(4 downto 0);
    in2: IN UNSIGNED(4 downto 0);
    out_1: OUT UNSIGNED(4 downto 0)
  );
END vhdl_fsm;
ARCHITECTURE rtl OF vhdl_fsm IS
  TYPE Tstate IS (state_0, state_1, state_2, state_3, state_4);
  SIGNAL state: Tstate;
  SIGNAL next_state: Tstate;
BEGIN
  PROCESS(clk, reset)
  BEGIN
    IF reset = '1' THEN
      state <= state_0;
    ELSIF rising_edge(clk) THEN
      state <= next_state;
    END IF;
  END PROCESS;
  PROCESS (state, in1, in2)
  VARIABLE tmp_out_0: UNSIGNED (4 downto 0);
  VARIABLE tmp_out_1: UNSIGNED (4 downto 0);
  BEGIN
    tmp_out_0 := in1 + in2;
    tmp_out_1 := in1 - in2;
    CASE state IS
      WHEN state_0 =>
        out_1 <= in1;
        next_state <= state_1;
      WHEN state_1 =>
        IF (in1 < in2) then
          next_state <= state_2;
          out_1 <= tmp_out_0;
        ELSE
          next_state <= state_3;
          out_1 <= tmp_out_1;
        END IF;
      WHEN state_2 =>
        IF (in1 < "0100") then
          out_1 <= tmp_out_0;
        ELSE
          out_1 <= tmp_out_1;
        END IF;
        next_state <= state_3;
      WHEN state_3 =>
        out_1 <= "11111";
        next_state <= state_4;
      WHEN state_4 =>
        out_1 <= in2;
        next_state <= state_0;
      WHEN OTHERS =>
        out_1 <= "00000";
        next_state <= state_0;
    END CASE;
  END PROCESS;
END rtl;
```
Multiplexers

Multiplexers form a large portion of the logic utilization in many FPGA designs. By optimizing your multiplexer logic, you ensure the most efficient implementation in your Altera device. This section addresses common problems and provides design guidelines to achieve optimal resource utilization for multiplexer designs. The section also describes various types of multiplexers, and how they are implemented.

For more information, refer to the *Advanced Synthesis Cookbook*.

Quartus II Software Option for Multiplexer Restructuring

Quartus II integrated synthesis provides the Restructure Multiplexers logic option that extracts and optimizes buses of multiplexers during synthesis. The default setting Auto for this option uses the optimization when it is most likely to benefit the optimization targets for your design. You can turn the option on or off specifically to have more control over its use.

For details, refer to the Restructure Multiplexers section in the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook.

Even with this Quartus II-specific option turned on, it is beneficial to understand how your coding style can be interpreted by your synthesis tool, and avoid the situations that can cause problems in your design.

Multiplexer Types

This section addresses how multiplexers are created from various types of HDL code. CASE statements, IF statements, and state machines are all common sources of multiplexer logic in designs. These HDL structures create different types of multiplexers, including binary multiplexers, selector multiplexers, and priority multiplexers. Understanding how multiplexers are created from HDL code, and how they might be implemented during synthesis, is the first step toward optimizing multiplexer structures for best results.

Binary Multiplexers

Binary multiplexers select inputs based on binary-encoded selection bits. Example 14–54 shows Verilog HDL code for two ways to describe a simple 4:1 binary multiplexer.

Example 14–54. Verilog HDL Binary-Encoded Multiplexers

```verilog
case (sel)
    2'b00: z = a;
    2'b01: z = b;
    2'b10: z = c;
    2'b11: z = d;
endcase
```
Stratix series devices starting with the Stratix II device family feature 6-input look up tables (LUTs) which are perfectly suited for 4:1 multiplexer building blocks (4 data and 2 select inputs). The extended input mode facilitates implementing 8:1 blocks, and the fractured mode handles residual 2:1 multiplexer pairs. For device families using 4-input LUTs, such as the Cyclone series and Stratix devices, the 4:1 binary multiplexer is efficiently implemented by using two 4-input LUTs. Larger binary multiplexers are decomposed by the synthesis tool into 4:1 multiplexer blocks, possibly with a residual 2:1 multiplexer at the head.

Selector Multiplexers

Selector multiplexers have a separate select line for each data input. The select lines for the multiplexer are one-hot encoded. **Example 14–55** shows a simple Verilog HDL code example describing a one-hot selector multiplexer.

**Example 14–55. Verilog HDL One-Hot-Encoded Case Statement**

```verilog
case (sel)
  4'b0001: z = a;
  4'b0010: z = b;
  4'b0100: z = c;
  4'b1000: z = d;
  default: z = 1'bx;
endcase
```

Selector multiplexers are commonly built as a tree of AND and OR gates. An N-input selector multiplexer of this structure is slightly less efficient in implementation than a binary multiplexer. However, in many cases the select signal is the output of a decoder, in which case Quartus II Synthesis will try to combine the selector and decoder into a binary multiplexer.

Priority Multiplexers

In priority multiplexers, the select logic implies a priority. The options to select the correct item must be checked in a specific order based on signal priority. These structures commonly are created from IF, ELSE, WHEN, SELECT, and ?: statements in VHDL or Verilog HDL. The example VHDL code in **Example 14–56** probably results in the schematic implementation illustrated in Figure 14–4.

**Example 14–56. VHDL IF Statement Implying Priority**

```vhdl
IF cond1 THEN z <= a;
ELSIF cond2 THEN z <= b;
ELSIF cond3 THEN z <= c;
ELSE z <= d;
END IF;
```
The multiplexers in Figure 14–4 form a chain, evaluating each condition or select bit sequentially.

Figure 14–4. Priority Multiplexer Implementation of an IF Statement

Depending on the number of multiplexers in the chain, the timing delay through this chain can become large, especially for device families with 4-input LUTs.

To improve the timing delay through the multiplexer, avoid priority multiplexers if priority is not required. If the order of the choices is not important to the design, use a CASE statement to implement a binary or selector multiplexer instead of a priority multiplexer. If delay through the structure is important in a multiplexed design requiring priority, consider recoding the design to reduce the number of logic levels to minimize delay, especially along your critical paths.

Implicit Defaults in If Statements

The IF statements in Verilog HDL and VHDL can be a convenient way to specify conditions that do not easily lend themselves to a CASE-type approach. However, using IF statements can result in complicated multiplexer trees that are not easy for synthesis tools to optimize. In particular, every IF statement has an implicit ELSE condition, even when it is not specified. These implicit defaults can cause additional complexity in a multiplexed design.

There are several ways you can simplify multiplexed logic and remove unneeded defaults. The optimal method may be to recode the design so the logic takes the structure of a 4:1 CASE statement. Alternatively, if priority is important, you can restructure the code to reduce default cases and flatten the multiplexer. Examine whether the default "ELSE IF" conditions are don’t care cases. You may be able to create a default ELSE statement to make the behavior explicit. Avoid unnecessary default conditions in your multiplexer logic to reduce the complexity and logic utilization required to implement your design.

Default or Others Case Assignment

To fully specify the cases in a CASE statement, include a default (Verilog HDL) or OTHERS (VHDL) assignment. This assignment is especially important in one-hot encoding schemes where many combinations of the select lines are unused. Specifying a case for the unused select line combinations gives the synthesis tool information about how to synthesize these cases, and is required by the Verilog HDL and VHDL language specifications.
Some designs do not require that the outcome in the unused cases be considered, often because designers assume these cases will not occur. For these types of designs, you can specify any value for the default or OTHERS assignment. However, be aware that the assignment value you choose can have a large effect on the logic utilization required to implement the design due to the different ways synthesis tools treat different values for the assignment, and how the synthesis tools use different speed and area optimizations.

To obtain best results, explicitly define invalid CASE selections with a separate default or OTHERS statement instead of combining the invalid cases with one of the defined cases.

If the value in the invalid cases is not important, specify those cases explicitly by assigning the X (don’t care) logic value instead of choosing another value. This assignment allows your synthesis tool to perform the best area optimizations.

### Cyclic Redundancy Check Functions

CRC computations are used heavily by communications protocols and storage devices to detect any corruption of data. These functions are highly effective; there is a very low probability that corrupted data can pass a 32-bit CRC check.

CRC functions typically use wide XOR gates to compare the data. The way synthesis tools flatten and factor these XOR gates to implement the logic in FPGA LUTs can greatly impact the area and performance results for the design. XOR gates have a cancellation property that creates an exceptionally large number of reasonable factoring combinations, so synthesis tools cannot always choose the best result by default.

The 6-input ALUT has a significant advantage over 4-input LUTs for these designs. When properly synthesized, CRC processing designs can run at high speeds in devices with 6-input ALUTs.

The following guidelines help you improve the quality of results for CRC designs in Altera devices.

#### If Performance is Important, Optimize for Speed

Synthesis tools flatten XOR gates to minimize area and depth of levels of logic. Synthesis tools such as Quartus II integrated synthesis target area optimization by default for these logic structures. Therefore, for more focus on depth reduction, set the synthesis optimization technique to speed.

Flattening for depth sometimes causes a significant increase in area.

#### Use Separate CRC Blocks Instead of Cascaded Stages

Some designers optimize their CRC designs to use cascaded stages (for example, four stages of 8 bits). In such designs, intermediate calculations are used as required (such as the calculations after 8, 24, or 32 bits) depending on the data width. This design is not optimal in FPGA devices. The XOR cancellations that can be performed in CRC designs mean that the function does not require all the intermediate calculations to
determine the final result. Therefore, forcing the use of intermediate calculations increases the area required to implement the function, as well as increasing the logic depth because of the cascading. It is typically better to create full separate CRC blocks for each data width that you require in the design, and then multiplex them together to choose the appropriate mode at a given time.

**Use Separate CRC Blocks Instead of Allowing Blocks to Merge**

Synthesis tools often attempt to optimize CRC designs by sharing resources and extracting duplicates in two different CRC blocks because of the factoring options in the XOR logic. As addressed previously, the CRC logic allows significant reductions, but this works best when each CRC function is optimized separately. Check for duplicate extraction behavior if you have different CRC functions that are driven by common data signals or that feed the same destination signals.

If you are having problems with the quality of results and you see that two CRC functions are sharing logic, ensure that the blocks are synthesized independently using one of the following methods:

- Define each CRC block as a separate design partition in an incremental compilation design flow.

  \[\rightarrow\] For details, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*.

- Synthesize each CRC block as a separate project in your third-party synthesis tool and then write a separate Verilog Quartus Mapping (.vqm) or EDIF netlist file for each.

**Take Advantage of Latency if Available**

If your design can use more than one cycle to implement the CRC functionality, adding registers and retiming the design can help reduce area, improve performance, and reduce power utilization. If your synthesis tool offers a retiming feature (such as the Quartus II software **Perform gate-level register retiming** option), you can insert an extra bank of registers at the input and allow the retiming feature to move the registers for better results. You can also build the CRC unit half as wide and alternate between halves of the data in each clock cycle.

**Save Power by Disabling CRC Blocks When Not in Use**

CRC designs are heavy consumers of dynamic power because the logic toggles whenever there is a change in the design. To save power, use clock enables to disable the CRC function for every clock cycle that the logic is not required. Some designs don’t check the CRC results for a few clock cycles while other logic is performed. It is valuable to disable the CRC function even for this short amount of time.
Use the Device Synchronous Load (sload) Signal to Initialize

The data in many CRC designs must be initialized to 1’s before operation. If your target device supports the use of the sload signal, you should use it to set all the registers in your design to 1’s before operation. To enable use of the sload signal, follow the coding guidelines presented in “Secondary Register Control Signals Such as Clear and Clock Enable” on page 14–46. You can check the register equations in the Chip Planner to ensure that the signal was used as expected.

If you must force a register implementation using an sload signal, you can use low-level device primitives as described in the Designing with Low-Level Primitives User Guide.

Comparators

Synthesis software, including Quartus II integrated synthesis, uses device and context-specific implementation rules for comparators (<, >, or ==) and selects the best one for your design. This section provides some information about the different types of implementations available and provides suggestions on how you can code your design to encourage a specific implementation.

The == comparator is implemented in general logic cells. The < comparison can be implemented using the carry chain or general logic cells. In devices with 6-input ALUTs, the carry chain is capable of comparing up to three bits per cell. In devices with 4-input LUTs, the capacity is one bit of comparison per cell, which is similar to an add/subtract chain. The carry chain implementation tends to be faster than the general logic on standalone benchmark test cases, but can result in lower performance when it is part of a larger design due to the increased restriction on the Fitter. The area requirement is similar for most input patterns. The synthesis software selects an appropriate implementation based on the input pattern.

If you are using Quartus II integrated synthesis, you can guide the synthesis by using specific coding styles. To select a carry chain implementation explicitly, rephrase your comparison in terms of addition. As a simple example, the following coding style allows the synthesis tool to select the implementation, which is most likely using general logic cells in modern device families:

```vhdl
wire [6:0] a,b;
wire alb = a<b;
```

In the following coding style, the synthesis tool uses a carry chain (except for a few cases, such as when the chain is very short or the signals a and b minimize to the same signal):

```vhdl
wire [6:0] a,b;
wire [7:0] tmp = a - b;
wire alb = tmp[7]
```

This second coding style uses the top bit of the tmp signal, which is 1 in twos complement logic if a is less than b, because the subtraction a – b results in a negative number.

If you have any information about the range of the input, you have “don’t care” values that you can use to optimize the design. Because this information is not available to the synthesis tool, you can often reduce the device area required to implement the comparator with specific hand implementation of the logic.
You can also check whether a bus value is within a constant range with a small amount of logic area by using the logic structure in Figure 14–5. This type of logic occurs frequently in address decoders.

**Figure 14–5. Example Logic Structure for Using Comparators to Check a Bus Value Range**

![Figure 14–5](image)

### Counters

Implementing counters in HDL code is easy; they are implemented with an adder followed by registers. Remember that the register control signals, such as enable (ena), synchronous clear (sclr), and synchronous load (sload), are available. For the best area utilization, ensure that the up/down control or controls are expressed in terms of one addition instead of two separate addition operators.

If you use the following coding style, your synthesis tool may implement two separate carry chains for addition (if it doesn’t detect the issue and optimize the logic):

```vhdl
out <= count_up ? out + 1 : out - 1;
```

The following coding style requires only one adder along with some other logic:

```vhdl
out <= out + (count_up ? 1 : -1);
```

In this case, the coding style better matches the device hardware because there is only one carry chain adder, and the –1 constant logic is implemented in the LUT in front of the adder without adding extra area utilization.

### Designing with Low-Level Primitives

Low-level HDL design is the practice of using low-level primitives and assignments to dictate a particular hardware implementation for a piece of logic. Low-level primitives are small architectural building blocks that assist you in creating your design. With the Quartus II software, you can use low-level HDL design techniques to force a specific hardware implementation that can help you achieve better resource utilization or faster timing results.

Using low-level primitives is an advanced technique to help with specific design challenges, and is optional in the Altera design flow. For many designs, synthesizing generic HDL source code and Altera megafunctions gives you the best results.

Low-level primitives allow you to use the following types of coding techniques:

- Instantiate the logic cell or LCELL primitive to prevent Quartus II integrated synthesis from performing optimizations across a logic cell.
Create carry and cascade chains using `CARRY`, `CARRY_SUM`, and `CASCADE` primitives

- Instantiate registers with specific control signals using `DFF` primitives
- Specify the creation of LUT functions by identifying the LUT boundaries
- Use I/O buffers to specify I/O standards, current strengths, and other I/O assignments
- Use I/O buffers to specify differential pin names in your HDL code, instead of using the automatically-generated negative pin name for each pair

For details about and examples of using these types of assignments, refer to the *Designing with Low-Level Primitives User Guide*.

**Conclusion**

Because coding style and megafuntion implementation can have such a large effect on your design performance, it is important to match the coding style to the device architecture from the very beginning of the design process. To improve design performance and area utilization, take advantage of advanced device features, such as memory and DSP blocks, as well as the logic architecture of the targeted Altera device by following the coding recommendations presented in this chapter.

For additional optimization recommendations, refer to the *Area and Timing Optimization* chapter in volume 2 of the *Quartus II Handbook*.

**Document Revision History**

Table 14–3 shows the revision history for this document.

**Table 14–3. Document Revision History (Part 1 of 2)**

<table>
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<th>Version</th>
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<td>June 2012</td>
<td>12.0.0</td>
<td>- Revised section on inserting Altera templates.</td>
</tr>
<tr>
<td></td>
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<td>- Code update for Example 11-51.</td>
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<tr>
<td></td>
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<td>- Minor corrections and updates.</td>
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<td>- Updated support for no_rv_check for inferring RAM blocks</td>
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<tr>
<td></td>
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<td>- Added support for byte-enable</td>
</tr>
<tr>
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<td>9.1.0</td>
<td>- Updated support for Controlling Inference and Implementation in Device RAM Blocks</td>
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<td>- Updated support for Shift Registers</td>
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### Table 14–3. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
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<td>- Added information to “RAM Functions—Inferring ALTSYNCRAM and ALTDPRAM Megafuntions from HDL Code” on page 6–13</td>
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<td></td>
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<td>- Added information to “Avoid Unsupported Reset and Control Conditions” on page 6–14</td>
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<td>- Added information to “Check Read-During-Write Behavior” on page 6–16</td>
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<td>- Added two new examples to “ROM Functions—Inferring ALTSYNCRAM and LPM_ROM Megafuntions from HDL Code” on page 6–28: Example 6–24 and Example 6–25</td>
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<td>- Added new section: “Clock Multiplexing” on page 6–46</td>
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<td>- Added hyperlinks to references within the chapter</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter describes the industry-leading analysis, reporting, and optimization features that can help you manage metastability in Altera® devices. You can use the Quartus® II software to analyze the average mean time between failures (MTBF) due to metastability caused by synchronization of asynchronous signals, and optimize the design to improve the metastability MTBF. This chapter explains how to take advantage of these features in the Quartus II software, and provides guidelines to help you reduce the chance of metastability effects caused by signal synchronization.

Introduction

All registers in digital devices, such as FPGAs, have defined signal-timing requirements that allow each register to correctly capture data at its input ports and produce an output signal. To ensure reliable operation, the input to a register must be stable for a minimum amount of time before the clock edge (register setup time or tSU) and a minimum amount of time after the clock edge (register hold time or tH). The register output is available after a specified clock-to-output delay (tCO).

If the data violates the setup or hold time requirements, the output of the register might go into a metastable state. In a metastable state, the voltage at the register output hovers at a value between the high and low states, which means the output transition to a defined high or low state is delayed beyond the specified tCO. Different destination registers might capture different values for the metastable signal, which can cause the system to fail.

In synchronous systems, the input signals must always meet the register timing requirements, so that metastability does not occur. Metastability problems commonly occur when a signal is transferred between circuitry in unrelated or asynchronous clock domains, because the signal can arrive at any time relative to the destination clock.

The MTBF due to metastability is an estimate of the average time between instances when metastability could cause a design failure. A high MTBF (such as hundreds or thousands of years between metastability failures) indicates a more robust design. You should determine an acceptable target MTBF in the context of your entire system and taking in account that MTBF calculations are statistical estimates.

The metastability MTBF for a specific signal transfer, or all the transfers in a design, can be calculated using information about the design and the device characteristics. Improving the metastability MTBF for your design reduces the chance that signal transfers could cause metastability problems in your device.
For more information about metastability due to signal synchronization, its effects in FPGAs, and how MTBF is calculated, refer to the *Understanding Metastability in FPGAs* white paper on the Altera website. Your overall device MTBF is also affected by other FPGA failure mechanisms that you cannot control with your design. For information about Altera device reliability, refer to the *Reliability Report* on the Altera website.

The Quartus II software provides analysis, optimization, and reporting features to help manage metastability in Altera designs. These metastability features are supported only for designs constrained with the Quartus II Timing Analyzer. Both typical and worst-case MBTF values are generated for select device families.

For information about device and version support for the metastability features in the Quartus II software, refer to the Quartus II Help.

This chapter contains the following topics:
- “Metastability Analysis in the Quartus II Software”
- “Metastability and MTBF Reporting” on page 15–5
- “MTBF Optimization” on page 15–8
- “Reducing Metastability Effects” on page 15–9
- “Scripting Support” on page 15–11

**Metastability Analysis in the Quartus II Software**

When a signal transfers between circuitry in unrelated or asynchronous clock domains, the first register in the new clock domain acts as a synchronization register. To minimize the failures due to metastability in asynchronous signal transfers, circuit designers typically use a sequence of registers (a synchronization register chain or synchronizer) in the destination clock domain to resynchronize the signal to the new clock domain and allow additional time for a potentially metastable signal to resolve to a known value. Designers commonly use two registers to synchronize a new signal, but a standard of three registers provides better metastability protection.

The timing analyzer can analyze and report the MTBF for each identified synchronizer that meets its timing requirements, and can generate an estimate of the overall design MTBF. The software uses this information to optimize the design MTBF, and you can use this information to determine whether your design requires longer synchronizer chains.

This section contains the following subsections:
- “Synchronization Register Chains”
- “Identifying Synchronizers for Metastability Analysis” on page 15–4
- “How Timing Constraints Affect Synchronizer Identification and Metastability Analysis” on page 15–4

For information about the reports generated by the timing analyzer, refer to “Metastability and MTBF Reporting” on page 15–5. For more information about optimizing the MTBF, refer to “MTBF Optimization” on page 15–8.
Synchronization Register Chains

A synchronization register chain, or synchronizer, is defined as a sequence of registers that meets the following requirements:

- The registers in the chain are all clocked by the same clock or phase-related clocks.
- The first register in the chain is driven asynchronously or from an unrelated clock domain.
- Each register fans out to only one register, except the last register in the chain.

The length of the synchronization register chain is the number of registers in the synchronizing clock domain that meet the above requirements. Figure 15–1 shows a sample two-register synchronization chain.

The path between synchronization registers can contain combinational logic as long as all registers of the synchronization register chain are in the same clock domain. Figure 15–2 shows an example of a synchronization register chain that includes logic between the registers.

The Quartus II software uses the design timing constraints to determine which connections are asynchronous signal transfers, as described in “How Timing Constraints Affect Synchronizer Identification and Metastability Analysis” on page 15–4.
The timing slack available in the register-to-register paths of the synchronizer allows a metastable signal to settle, and is referred to as the available settling time. The available settling time in the MTBF calculation for a synchronizer is the sum of the output timing slacks for each register in the chain. Adding available settling time with additional synchronization registers improves the metastability MTBF.

### Identifying Synchronizers for Metastability Analysis

The first step in enabling metastability MTBF analysis and optimization in the Quartus II software is to identify which registers are part of a synchronization register chain. You can apply synchronizer identification settings globally to automatically list possible synchronizers with the Synchronizer identification option on the Timing Analyzer page in the Settings dialog box.

Synchronization chains are already identified within most Altera intellectual property (IP) cores.

For more information about how to enable metastability MTBF analysis and optimization in the Quartus II software, and more detailed descriptions of the synchronizer identification settings, refer to Identifying Synchronizers for Metastability Analysis in Quartus II Help.

### How Timing Constraints Affect Synchronizer Identification and Metastability Analysis

The timing analyzer can analyze metastability MTBF only if a synchronization chain meets its timing requirements. The metastability failure rate depends on the timing slack available in the synchronizer’s register-to-register connections, because that slack is the available settling time for a potential metastable signal. Therefore, you must ensure that your design is correctly constrained with the real application frequency requirements to get an accurate MTBF report.

In addition, the Auto and Forced If Asynchronous synchronizer identification options use timing constraints to automatically detect the synchronizer chains in the design. These options check for signal transfers between circuitry in unrelated or asynchronous clock domains, so clock domains must be related correctly with timing constraints.

The timing analyzer views input ports as asynchronous signals unless they are associated correctly with a clock domain. If an input port fans out to registers that are not acting as synchronization registers, apply a set_input_delay constraint to the input port; otherwise, the input register might be reported as a synchronization register. Constraining a synchronous input port with a set_max_delay constraint for a setup (tSU) requirement does not prevent synchronizer identification because the constraint does not associate the input port with a clock domain.

Instead, use the following command to specify an input setup requirement associated with a clock:

```
set_input_delay -max -clock <clock name> <latch – launch – tsu requirement> <input port name>
```
 Registers that are at the end of false paths are also considered synchronization registers because false paths are not timing-analyzed. Because there are no timing requirements for these paths, the signal may change at any point, which may violate the $t_{SU}$ and $t_H$ of the register. Therefore, these registers are identified as synchronization registers. If these registers are not used for synchronization, you can turn off synchronizer identification and analysis. To do so, set **Synchronizer Identification** to **Off** for the first synchronization register in these register chains.

### Metastability and MTBF Reporting

The Quartus II software reports the metastability analysis results in the Compilation Report and Timing Analyzer reports as described in “Metastability Reports”. The MTBF calculation uses timing and structural information about the design, silicon characteristics, and operating conditions, along with the data toggle rate described in “Synchronizer Data Toggle Rate in MTBF Calculation” on page 15–7.

If you change the **Synchronizer Identification** settings, you can generate new metastability reports by rerunning the timing analyzer. However, you should rerun the Fitter first so that the registers identified with the new setting can be optimized for metastability MTBF. For information about metastability optimization, refer to “MTBF Optimization” on page 15–8.

For more information about how metastability MTBF is calculated, refer to the *Understanding Metastability in FPGAs* white paper.

### Metastability Reports

Metastability reports provide summaries of the metastability analysis results. In addition to the MTBF Summary and Synchronizer Summary reports, the Timing Analyzer tool reports additional statistics in a report for each synchronizer chain.

For more information about how to access metastability reports in the Quartus II software, refer to *Viewing Metastability Reports* in Quartus II Help.

If the design uses only the **Auto Synchronizer Identification** setting, the reports list likely synchronizers but do not report MTBF. To obtain an MTBF for each register chain, force identification of synchronization registers as described in “Identifying Synchronizers for Metastability Analysis” on page 15–4.

If the synchronizer chain does not meet its timing requirements, the reports list identified synchronizers but do not report MTBF. To obtain MTBF calculations, ensure that the design is properly constrained and that the synchronizer meets its timing requirements, as described in “How Timing Constraints Affect Synchronizer Identification and Metastability Analysis” on page 15–4.

### MTBF Summary Report

The MTBF Summary reports an estimate of the overall robustness of cross-clock domain and asynchronous transfers in the design. This estimate uses the MTBF results of all synchronization chains in the design to calculate an MTBF for the entire design.
The MTBF Summary Report reports the **Typical MTBF of Design** and the **Worst-Case MTBF of Design** for supported fully-characterized devices. The typical MTBF result assumes typical conditions, defined as nominal silicon characteristics for the selected device speed grade, as well as nominal operating conditions. The worst case MTBF result uses the worst case silicon characteristics for the selected device speed grade.

When you analyze multiple timing corners in the timing analyzer, the MTBF calculation may vary because of changes in the operating conditions, and the timing slack or available metastability settling time. Altera recommends running multi-corner timing analysis to ensure that you analyze the worst MTBF results, because the worst timing corner for MTBF does not necessarily match the worst corner for timing performance.

For more information about turning on multi-corner timing analysis in the Quartus II software, refer to the Timing Analyzer page in Quartus II Help.

The MTBF Summary report also lists the **Number of Synchronizer Chains Found** and the length of the **Shortest Synchronizer Chain**, which can help you identify whether the report is based on accurate information. If the number of synchronizer chains found is different from what you expect, or if the length of the shortest synchronizer chain is less than you expect, you might have to add or change **Synchronizer Identification** settings for the design. The report also provides the **Worst Case Available Settling Time**, defined as the available settling time for the synchronizer with the worst MTBF.

You can use the reported **Fraction of Chains for which MTBFs Could Not be Calculated** to determine whether a high proportion of chains are missing in the metastability analysis. A fraction of 1, for example, means that MTBF could not be calculated for any chains in the design. MTBF is not calculated if you have not identified the chain with the appropriate **Synchronizer identification** option, or if paths are not timing-analyzed and therefore have no valid slack for metastability analysis. You might have to correct your timing constraints to enable complete analysis of the applicable register chains.

Finally, the MTBF Summary report specifies how an increase of 100ps in available settling time increases the MTBF values. If your MTBF is not satisfactory, this metric can help you determine how much extra slack would be required in your synchronizer chain to allow you to reach the desired design MTBF.

**Synchronizer Summary Report**

The **Synchronizer Summary** lists the synchronization register chains detected in the design depending on the Synchronizer Identification setting. The **Source Node** is the register or input port that is the source of the asynchronous transfer. The **Synchronization Node** is the first register of the synchronization chain. The **Source Clock** is the clock domain of the source node, and the **Synchronization Clock** is the clock domain of the synchronizer chain.

This summary reports the calculated **Worst-Case MTBF**, if available, and the **Typical MTBF**, for each appropriately identified synchronization register chain that meets its timing requirement. To see more detail about each synchronizer, refer to the statistics report described in the following section.
Synchronizer Chain Statistics Report in the Timing Analyzer

The timing analyzer provides an additional report for each synchronizer chain. The Chain Summary tab matches the Synchronizer Summary information described in the previous section, while the Statistics tab adds more details, including whether the Method of Synchronizer Identification was User Specified (with the Forced if Asynchronous or Forced settings for the Synchronizer Identification setting), or Automatic (with the Auto setting). The Number of Synchronization Registers in Chain report provides information about the parameters that affect the MTBF calculation, including the Available Settling Time for the chain and the Data Toggle Rate Used in MTBF Calculation.

For information about the toggle rate, see “Synchronizer Data Toggle Rate in MTBF Calculation” on page 15–7.

The following information is also included to help you locate the chain in your design:

- Source Clock and Asynchronous Source node of the signal.
- Synchronization Clock in the destination clock domain.
- Node names of the Synchronization Registers in the chain.

Synchronizer Data Toggle Rate in MTBF Calculation

The MTBF calculations assume the data being synchronized is switching at a toggle rate of 12.5% of the source clock frequency. That is, the arriving data is assumed to switch once every eight source clock cycles. If multiple clocks apply, the highest frequency is used. If no source clocks can be determined, the data rate is taken as 12.5% of the synchronization clock frequency.

If you know an approximate rate at which the data changes, specify it with the Synchronizer Toggle Rate assignment in the Assignment Editor. You can also apply this assignment to an entity or the entire design. Set the data toggle rate, in number of transitions per second, on the first register of a synchronization chain. The timing analyzer takes the specified rate into account when computing the MTBF of that particular register chain. If a data signal never toggles and does not affect the reliability of the design, you can set the Synchronizer Toggle Rate to 0 for the synchronization chain so the MTBF is not reported. To apply the assignment with Tcl, use the following command:

```tcl
set_instance_assignment -name SYNCHRONIZER_TOGGLE_RATE <toggle rate in transitions/second> -to <register name>
```

There are two other assignments associated with toggle rates, which are not used for metastability MTBF calculations. The I/O Maximum Toggle Rate is only used for pins, and specifies the worst-case toggle rates used for signal integrity purposes. The Power Toggle Rate assignment is used to specify the expected time-averaged toggle rate, and is used by the PowerPlay Power Analyzer to estimate time-averaged power consumption.
MTBF Optimization

In addition to reporting synchronization register chains and MTBF values found in the design, the Quartus II software can also protect these registers from optimizations that might negatively impact MTBF and can optimize the register placement and routing if the MTBF is too low. Synchronization register chains must first be explicitly identified as synchronizers, as described in “Identifying Synchronizers for Metastability Analysis” on page 15–4. Altera recommends that you set Synchronizer Identification to Forced If Asynchronous for all registers that are part of a synchronizer chain.

Optimization algorithms, such as register duplication and logic retiming in physical synthesis, are not performed on identified synchronization registers. The Fitter protects the number of synchronization registers specified by the Synchronizer Register Chain Length option which is described in the next section.

In addition, the Fitter optimizes identified synchronizers for improved MTBF by placing and routing the registers to increase their output setup slack values. Adding slack in the synchronizer chain increases the available settling time for a potentially metastable signal, which improves the chance that the signal resolves to a known value, and exponentially increases the design MTBF. The Fitter optimizes the number of synchronization registers specified by the Synchronizer Register Chain Length option.

Metastability optimization is on by default. To view or change the option, on the Assignments menu, click Settings. Under Fitter Settings, click More Settings. From the More Settings dialog box, you can turn on or off the Optimize Design for Metastability option. To turn the optimization on or off with Tcl, use the following command:

```
set_global_assignment -name OPTIMIZE_FOR_METASTABILITY <ON|OFF>
```

Synchronization Register Chain Length

The Synchronization Register Chain Length option specifies how many registers should be protected from optimizations that might reduce MTBF for each register chain, and controls how many registers should be optimized to increase MTBF with the Optimize Design for Metastability option. For example, if the Synchronization Register Chain Length option is set to 2, optimizations such as register duplication or logic retiming are prevented from being performed on the first two registers in all identified synchronization chains. The first two registers are also optimized to improve MTBF when the Optimize Design for Metastability option is turned on.

The default setting for the Synchronization Register Chain Length option is 2. The first register of a synchronization chain is always protected from operations that might reduce MTBF, but you should set the protection length to protect more of the synchronizer chain. Altera recommends that you set this option to the maximum length of synchronization chains you have in your design so that all synchronization registers are preserved and optimized for MTBF.

To change the global Synchronization Register Chain Length option, on the Assignments menu, click Settings. Under Analysis & Synthesis Settings, click More Settings. From the More Settings dialog box, you can set the Synchronization Register Chain Length.
You can also set the Synchronization Register Chain Length on a node or an entity in the Assignment Editor. You can set this value on the first register in a synchronization chain to specify how many registers to protect and optimize in this chain. This individual setting is useful if you want to protect and optimize extra registers that you have created in a specific synchronization chain that has low MTBF, or optimize less registers for MTBF in a specific chain where the maximum frequency or timing performance is not being met. To make the global setting with Tcl, use the following command:

```tcl
set_global_assignment -name SYNCHRONIZATION_REGISTER_CHAIN_LENGTH <number of registers>
```

To apply the assignment to a design instance or the first register in a specific chain with Tcl, use the following command:

```tcl
set_instance_assignment -name SYNCHRONIZATION_REGISTER_CHAIN_LENGTH <number of registers> -to <register or instance name>
```

## Reducing Metastability Effects

You can check your design’s metastability MTBF in the Metastability Summary report described in “Metastability Reports” on page 15–5, and determine an acceptable target MTBF given the context of your entire system and the fact that MTBF calculations are statistical estimates. A high metastability MTBF (such as hundreds or thousands of years between metastability failures) indicates a more robust design.

This section provides guidelines to ensure complete and accurate metastability analysis, and some suggestions to follow if the Quartus II metastability reports calculate an unacceptable MTBF value. The Timing Optimization Advisor (available from the Tools menu) gives similar suggestions in the Metastability Optimization section.

### Apply Complete System-Centric Timing Constraints for the Timing Analyzer

To enable the Quartus II metastability features, make sure that the timing analyzer is used for timing analysis.

Ensure that the design is fully timing constrained and that it meets its timing requirements. If the synchronization chain does not meet its timing requirements, MTBF cannot be calculated. If the clock domain constraints are set up incorrectly, the signal transfers between circuitry in unrelated or asynchronous clock domains might be identified incorrectly.

Use industry-standard system-centric I/O timing constraints instead of using FPGA-centric timing constraints. As described in “How Timing Constraints Affect Synchronizer Identification and Metastability Analysis” on page 15–4, you should use `set_input_delay` constraints in place of `set_max_delay` constraints to associate each input port with a clock domain to help eliminate false positives during synchronization register identification.

### Force the Identification of Synchronization Registers

Use the guidelines in “Identifying Synchronizers for Metastability Analysis” on page 15–4 to ensure the software reports and optimizes the appropriate register chains.
In summary, identify synchronization registers with the Synchronizer Identification set to Forced If Asynchronous in the Assignment Editor. If there are any registers that the software detects as synchronous but you want to be analyzed for metastability, apply the Forced setting to the first synchronizing register. Set Synchronizer Identification to Off for registers that are not synchronizers for asynchronous signals or unrelated clock domains.

To help you find the synchronizers in your design, you can set the global Synchronizer Identification setting on the Timing Analyzer page of the Settings dialog box to Auto to generate a list of all the possible synchronization chains in your design.

Set the Synchronizer Data Toggle Rate

The MTBF calculations assume the data being synchronized is switching at a toggle rate of 12.5% of the source clock frequency. To obtain a more accurate MTBF for a specific chain or all chains in your design, set the Synchronizer Toggle Rate as described in “Synchronizer Data Toggle Rate in MTBF Calculation” on page 15–7.

Optimize Metastability During Fitting

Ensure that the Optimize Design for Metastability setting described in “MTBF Optimization” on page 15–8 is turned on.

Increase the Length of Synchronizers to Protect and Optimize

Increase the Synchronizer Chain Length parameter to the maximum length of synchronization chains in your design, as described in “Synchronization Register Chain Length” on page 15–8. If you have synchronization chains longer than 2 identified in your design, you can protect the entire synchronization chain from operations that might reduce MTBF and allow metastability optimizations to improve the MTBF.

Set Fitter Effort to Standard Fit instead of Auto Fit

If your design MTBF is too low after following the previous guidelines in this section, you can try increasing the Fitter effort to perform more metastability optimization. The default Auto Fit setting reduces the Fitter’s effort after meeting the design’s timing and routing requirements to reduce compilation time. This effort reduction can result in less metastability optimization if the timing requirements are easy to meet. If Auto Fit reduces the Fitter’s effort during your design compilation, setting the Fitter effort to Standard Fit might improve the design’s MTBF results. In the Settings dialog box, on the Fitter Settings page, set Fitter effort to Standard Fit.

Increase the Number of Stages Used in Synchronizers, If Possible

Designers commonly use two registers in a synchronization chain to minimize the occurrence of metastable events, and a standard of three registers provides better metastability protection. However, synchronization chains with two or even three registers may not be enough to produce a high enough MTBF when the design runs at high clock and data frequencies.
If a synchronization chain is reported to have a low MTBF, consider adding an additional register stage to your synchronization chain. This additional stage increases the settling time of the synchronization chain, allowing more opportunity for the signal to resolve to a known state during a metastable event. Additional settling time increases the MTBF of the chain and improves the robustness of your design. However, adding a synchronization stage introduces an additional stage of latency on the signal.

If you use the Altera FIFO megafunction with separate read and write clocks to cross clock domains, increase the metastability protection (and latency) for better MTBF. In the MegaWizard™ Plug-In Manager for the DCFIFO function, choose the **Best metastability protection, best fmax, unsynchronized clocks** option to add three or more synchronization stages. You can increase the number of stages to more than three using the **How many sync stages?** setting.

**Select a Faster Speed Grade Device, if Possible**

The design MTBF depends on process parameters of the device used. Faster devices are less susceptible to metastability issues. If the design MTBF falls significantly below the target MTBF, switching to a faster speed grade can improve the MTBF substantially.

**Scripting Support**

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II Command-Line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the **Tcl Scripting** chapter in volume 2 of the **Quartus II Handbook**. For more information about settings and constraints in the Quartus II software, refer to the **Quartus II Settings File Reference Manual**. For more information about command-line scripting, refer to the **Command-Line Scripting** chapter in volume 2 of the **Quartus II Handbook** and **About Quartus II Scripting** in **Quartus II Help**.

**Identifying Synchronizers for Metastability Analysis**

To apply the global Synchronizer Identification assignment described in “Identifying Synchronizers for Metastability Analysis” on page 15–4, use the following command:

```
set_global_assignment -name SYNCHRONIZER_IDENTIFICATION <OFF|AUTO|"FORCED IF ASYNCHRONOUS">
```

To apply the **Synchronizer Identification** assignment to a specific register or instance, use the following command:

```
set_instance_assignment -name SYNCHRONIZER_IDENTIFICATION <AUTO|"FORCED IF ASYNCHRONOUS"|FORCED|OFF> -to <register or instance name>
```
Synchronizer Data Toggle Rate in MTBF Calculation

To specify a toggle rate for MTBF calculations as described on page “Synchronizer Data Toggle Rate in MTBF Calculation” on page 15–7, use the following command:

```
set_instance_assignment -name SYNCHRONIZER_TOGGLE_RATE <toggle rate in transitions/second> -to <register name>
```

report_metastability and Tcl Command

If you use a command-line or scripting flow, you can generate the metastability analysis reports described in “Metastability Reports” on page 15–5 outside of the Quartus II and user interfaces. Table 15–1 describes the options for the `report_metastability` and Tcl command.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-append</td>
<td>If output is sent to a file, this option appends the result to that file. Otherwise, the file is overwritten.</td>
</tr>
<tr>
<td>-file &lt;name&gt;</td>
<td>Sends the results to an ASCII or HTML file. The extension specified in the file name determines the file type—either *.txt or *.html.</td>
</tr>
<tr>
<td>-panel_name &lt;name&gt;</td>
<td>Sends the results to the panel and specifies the name of the new panel.</td>
</tr>
<tr>
<td>-stdout</td>
<td>Indicates the report be sent to the standard output, via messages. This option is required only if you have selected another output format, such as a file, and would also like to receive messages.</td>
</tr>
</tbody>
</table>

MTBF Optimization

To ensure that metastability optimization described on page “MTBF Optimization” on page 15–8 is turned on (or to turn it off), use the following command:

```
set_global_assignment -name OPTIMIZE_FOR_METASTABILITY <ON|OFF>
```

Synchronization Register Chain Length

To globally set the number of registers in a synchronization chain to be protected and optimized as described on page “Synchronization Register Chain Length” on page 15–8, use the following command:

```
set_global_assignment -name SYNCHRONIZATION_REGISTER_CHAIN_LENGTH <number of registers>
```

To apply the assignment to a design instance or the first register in a specific chain, use the following command:

```
set_instance_assignment -name SYNCHRONIZATION_REGISTER_CHAIN_LENGTH <number of registers> -to <register or instance name>
```
Conclusion

Altera’s Quartus II software provides industry-leading analysis and optimization features to help you manage metastability in your FPGA designs. Set up your Quartus II project with the appropriate constraints and settings to enable the software to analyze, report, and optimize the design MTBF. Take advantage of these features in the Quartus II software and follow the guidelines in this chapter to make your design more robust with respect to metastability.

Document Revision History

Table 15–2 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Changed to new document template.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Technical edit.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Clarified description of synchronizer identification settings.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Minor changes to text and figures throughout document.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides guidelines to help you partition your design to take advantage of Quartus® II incremental compilation, and to help you create a design floorplan using LogicLock™ regions when they are recommended to support the compilation flow.

The Quartus II incremental compilation feature allows you to partition a design, compile partitions separately, and reuse results for unchanged partitions. Incremental compilation provides the following benefits:

- Reduces compilation times by an average of 75% for large design changes
- Preserves performance for unchanged design blocks
- Provides repeatable results and reduces the number of compilations
- Enables team-based design flows

For more information about the incremental compilation feature and application examples, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook. For feature support, refer to About Incremental Compilation in Quartus II Help.

This document contains the following sections:

- “Overview: Incremental Compilation” on page 16–2
- “Design Flows Using Incremental Compilation” on page 16–3
- “Why Plan Partitions and Floorplan Assignments?” on page 16–6
- “General Partitioning Guidelines” on page 16–8
- “Design Partition Guidelines” on page 16–10
- “Checking Partition Quality” on page 16–30
- “Including SDC Constraints from Lower-Level Partitions for Third-Party IP Delivery” on page 16–36
- “Introduction to Design Floorplans” on page 16–40
- “Design Floorplan Placement Guidelines” on page 16–43
- “Checking Floorplan Quality” on page 16–49
- “Recommended Design Flows and Application Examples” on page 16–51
Overview: Incremental Compilation

Quartus II incremental compilation is an optional compilation flow that enhances the default Quartus II compilation. If you do not partition your design for incremental compilation, your design is compiled using the default “flat” compilation flow. This section provides an overview of the incremental compilation flow and highlights best practices.

To prepare your design for incremental compilation, you first determine which logical hierarchy boundaries should be defined as separate partitions in your design, and ensure your design hierarchy and source code is set up to support this partitioning. You can then create design partition assignments in the Quartus II software to specify which hierarchy blocks are compiled independently as partitions (including empty partitions for missing or incomplete logic blocks).

During compilation, Quartus II Analysis & Synthesis and the Fitter create separate netlists for each partition. Netlists are internal post-synthesis and post-fit database representations of your design.

In subsequent compilations, you can select which netlist to preserve for each partition. You can either reuse the synthesis or fitting netlist, or instruct the Quartus II software to resynthesize the source files. You can also use compilation results exported from another Quartus II project.

When you make changes to your design, the Quartus II software recompiles only the designated partitions and merges the new compilation results with existing netlists for other partitions, according to the degree of results preservation you set with the netlist for each design partition.

In some cases, as described in “Introduction to Design Floorplans” on page 16–40, Altera recommends that you create a design floorplan with placement assignments to constrain parts of the design to specific regions of the device.

For step-by-step information about using incremental compilation to recompile only changed parts of your design, refer to Using the Incremental Compilation Design Flow in Quartus II Help.

You must use incremental compilation in conjunction with the partial reconfiguration (PR) feature for Stratix® V device families. Partial reconfiguration allows you to reconfigure a portion of the FPGA dynamically, while the remainder of the device continues to operate as intended.

For more information about partial reconfiguration, refer to the Design Planning for Partial Reconfiguration chapter in volume 1 of the Quartus II Handbook.
Recommendations for the Netlist Type

For subsequent compilations, you specify which post-compilation netlist you want to use with the netlist type for each partition.

Use the following general guidelines to set the netlist type for each partition:

- **Source File**—Use this setting to resynthesize the source code (with any new assignments, and replace any previous synthesis or Fitter results).
  - If you modify the design source, the software automatically resynthesizes the partitions with the appropriate netlist type, which makes the **Source File** setting optional in this case.
  - Most assignments do not trigger an automatic recompilation, so you must set the netlist type to **Source File** to compile the source files with new assignments or constraints that affect synthesis.

- **Post-Synthesis** (default)—Use this setting to re-fit the design (with any new Fitter assignments), but preserve the synthesis results when the source files have not changed. If it is difficult to meet the required timing performance, you can use this setting to allow the Fitter the most flexibility in placement and routing. This setting does not reduce compilation time as much as the **Post-Fit** setting or preserve timing performance from the previous compilation.

- **Post-Fit**—Use this setting to preserve Fitter and performance results when the source files have not changed. This setting reduces compilation time the most, and preserves timing performance from the previous compilation.

- **Post-Fit with Fitter Preservation Level set to Placement**—Use the Advanced Fitter Preservation Level setting on the Advanced tab in the Design Partition Properties dialog box to allow more flexibility and find the best routing for all partitions given their placement.

The Quartus II software Rapid Recompile feature instructs the Compiler to reuse the compatible compilation results if most of the design has not changed since the last compilation. This feature reduces compilation time and preserves performance when there are small and isolated design changes within a partition, and works with all netlist type settings. With this feature, you do not have control over which parts of the design are recompiled; the Compiler determines which parts of the design must be recompiled. You can turn on the Rapid Recompile option in the Quartus II software on the Incremental Compilation page of the Settings dialog box.

Design Flows Using Incremental Compilation

The Quartus II incremental compilation feature supports various design flows. Your design flow affects the impact design partitions have on design optimization and the amount of design planning required to obtain optimal results.
In the standard incremental compilation flow, the top-level design is divided into partitions, which can be compiled and optimized together in one Quartus II project. If another team member or IP provider is developing source code for the top-level design, they can functionally verify their partition independently, and then simply provide the partition’s source code to the project lead for integration into the top-level design. If the project lead wants to compile the top-level design when source code is not yet complete for a partition, they can create an empty placeholder for the partition until the code is ready and added to the top-level design.

Compiling all design partitions in a single Quartus II project ensures that all design logic is compiled with a consistent set of assignments, and allows the software to perform global placement and routing optimizations. Compiling all design logic together is beneficial for FPGA design flows because all parts of the design must use the same shared set of device resources. Therefore, it is often easier to ensure good quality of results when partitions are developed within a single top-level Quartus II project.

In the team-based incremental compilation flow, you can design and optimize partitions by accessing the top-level project from a shared source control system or creating copies of the top-level Quartus II project framework. As development continues, designers export their partition so that the post-synthesis netlist or post-fitting results can be integrated into the top-level design.

If required for third-party IP delivery, or in cases where designers cannot access a shared or copied top-level project framework, you can create and compile a design partition logic in isolation and export a partition that is included in the top-level project. If this type of design flow is necessary, planning and rigorous design guidelines might be required to ensure that designers have a consistent view of project assignments and resource allocations. Therefore, developing partitions in completely separate Quartus II projects can be more challenging than having all source code within one project or developing design partitions within the same top-level project framework.

You can also combine design flows and use exported partitions only when it is necessary to support your design environment. For example, if the top-level design includes one or more design blocks that will be optimized by remote designers or IP providers, you can integrate those blocks into the reserved partitions in top-level design when the code is complete, but also have other partitions that will be developed within the top-level design.

If any partitions are developed independently, the project lead must ensure that top-level constraints (such as timing constraints, any relevant floorplan or pin assignments, and optimization settings) are consistent with those used by all designers working independently.

**Project Management in Team-Based Design Flows**

If possible, each team member should work within the same top-level project framework. Using the same project framework amongst team members ensures that designers have the settings and constraints needed for their partition and allows designers to analyze how their design block interacts with other partitions in the top-level design.
In a team-based environment where designers have access to the project through source control software, each designer can use project files as read-only and develop their partition within the source control system. As designers check in their completed partitions, other team members can see how their partitions interact. If designers do not have access to a source control system, the project lead can provide each designer with a copy of the top-level project framework to use as they develop their partitions. In both cases, each designer exports their completed design as a partition, and then the project lead integrates the partition into the top-level design. The project lead can choose to use only the post-synthesis netlist and rerun placement and routing, or to use the post-fitting results to preserve the placement and routing results from the other designer’s projects. Using post-synthesis partitions gives the Fitter the most flexibility and is likely to achieve a good result for all partitions, but if one partition has difficulty meeting timing, the designer can choose to preserve their successful fitting results.

Alternatively, designers can use their own Quartus II project for their independent design block. You might use this design flow if a designer, such as a third-party IP provider, does not have access to the entire top-level project framework. In this case, each designer must create their own project with all the relevant assignments and constraints. As mentioned at the beginning of this section, this type of design flow requires more planning and rigorous design guidelines. If the project lead plans to incorporate the post-fitting compilation results for the partition, this design flow requires especially careful planning to avoid resource conflicts.

The project lead also has the option to generate design partition scripts to manage resource and timing budgets in the top-level design when partitions are developed outside the top-level project framework. Scripts make it easier for designers of independent Quartus II projects to follow instructions from the project lead. The Quartus II design partition scripts feature creates Tcl scripts or .tcl files and makefiles that an independent designer can run to set up an independent Quartus II project.

For more information about how to generate design partition scripts, refer to Generating Design Partition Scripts for Project Management in Quartus II Help.

If designers create Quartus II assignments or timing constraints for their partitions, they must ensure that the constraints are integrated into the top-level design. If partition designers use the same top-level project framework (and design hierarchy), the constraints or Synopsys Design Constraints File (.sdc) can be easily copied or included in the top-level design. If partition designers use a separate Quartus II project with a different design hierarchy, they must ensure that constraints are applied to the appropriate level of hierarchy in the top-level design, and design the .sdc for easy delivery to the project lead, as described in “Including SDC Constraints from Lower-Level Partitions for Third-Party IP Delivery” on page 16–36.

You cannot use an exported partition if you want to migrate to a HardCopy ASIC device. The Revision Compare feature requires that the HardCopy and FPGA netlists are the same, and all operations performed on one revision must also occur on the other revision. Exporting partitions does not support this requirement.

For more information about the different types of incremental design flows and example applications, as well as documented restrictions and limitations, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Why Plan Partitions and Floorplan Assignments?

Incremental design flows typically require more planning than flat compilations, and require you to be more rigorous about following good design practices. For example, you might need to structure your source code or design hierarchy to ensure that logic is grouped correctly for optimization. It is easier to implement the correct logic grouping early in the design cycle than to restructure the code later.

Planning involves setting up the design logic for partitioning and may also involve planning placement assignments to create a floorplan. Not all design flows require floorplan assignments. For more information, refer to “Introduction to Design Floorplans” on page 16–40. If you decide to add floorplan assignments later, when the design is close to completion, well-planned partitions make floorplan creation easier. Poor partition or floorplan assignments can worsen design area utilization and performance and make timing closure more difficult.

As FPGA devices get larger and more complex, following good design practices become more important for all design flows. Adhering to recommended synchronous design practices makes designs more robust and easier to debug. Using an incremental compilation flow adds additional steps and requirements to your project, but can provide significant benefits in design productivity by preserving the performance of critical blocks and reducing compilation time.

Partition Boundaries and Optimization

The logical hierarchical boundaries between partitions are treated as hard boundaries for logic optimization (except for some limited cross-boundary optimizations) to allow the software to size and place each partition independently. Figure 16–1 shows the effects of partition boundaries during logic optimization.

Figure 16–1. Effects of Partition Boundaries During Logic Optimization

You can use the Merge command in the Design Partitions window to combine hierarchical partitions into a single partition, as long as they share the same immediate parent partition. Merging partitions allows additional optimizations for partition I/O ports that connect between or feed more than one of the merged hierarchical design blocks.
When partitions are placed together, the Fitter can perform placement optimizations on the design as a whole to optimize the placement of cross-boundary paths. However, the Fitter can never perform logic optimizations such as physical synthesis across the partition boundary. If partitions are fit separately in different projects, or if some partitions use previous post-fitting results, the Fitter does not place and route the entire cross-boundary path at the same time and cannot fully optimize placement across the partition boundaries. Good design partitions can be placed independently because cross-partition paths are not the critical timing paths in the design.

There are possible timing performance utilization effects due to partitioning and creating a floorplan. Not all designs encounter these issues, but you should consider these effects if a flat version of your design is very close to meeting its timing requirements, or is close to using all the device resources, before adding partition or floorplan assignments:

- Partitions can increase resource utilization due to cross-boundary optimization limitations if the design does not follow partitioning guidelines. For more information, refer to “Design Partition Guidelines” on page 16-10. Floorplan assignments can also increase resource utilization because regions can lead to unused logic. If your device is full with the flat version of your design, you can focus on creating partitions and floorplan assignments for timing-critical or often-changing blocks to benefit most from incremental compilation. For more information, refer to “Checking Floorplan Quality” on page 16-49.

- Partitions and floorplan assignments might increase routing utilization compared to a flat design. If long compilation times are due to routing congestion, you might not be able to use the incremental flow to reduce compilation time. Review the Fitter messages to check how much time is spent during routing optimizations to determine the percentage of routing utilization. When routing is difficult, you can use incremental compilation to lock the routing for routing-critical blocks only (with other partitions empty), and then compile the rest of the design after the critical blocks meet their requirements.

- Partitions can reduce timing performance in some cases because of the optimization and resource effects described above, causing longer logic delays. Floorplan assignments restrict logic placement, which can make it more difficult for the Fitter to meet timing requirements. Use the guidelines in this chapter to reduce any effect on your design performance.

**Turning On Supported Cross-boundary Optimizations**

You can improve the optimizations performed between design partitions by turning on the cross-boundary optimizations feature. You can select the optimizations as individual assignments for each partition. This allows the cross-boundary optimization feature to give you more control over the optimizations that work best for your design.

You can turn on the cross-boundary optimizations for your design partitions on the Advanced tab of the Design Partition Properties dialog box. Once you change the optimization settings, the Quartus II software recompiles your partition from source automatically. Cross-boundary optimizations include the following: propagate constants, propagate inversions on partition inputs, merge inputs fed by a common source, merge electrically equivalent bidirectional pins, absorb internal paths, and remove logic connected to dangling outputs.
Cross-boundary optimizations are implemented top-down from the parent partition into the child partition, but not vice-versa. The cross-boundary optimization feature cannot be used with partitions with multiple personas (partial reconfiguration partitions).

For more information about cross-boundary optimizations, refer to Design Partition Properties Dialog Box in Quartus II Help.

Although more partitions allow for a greater reduction in compilation time, consider limiting the number of partitions to prevent degradation in the quality of results. Creating good design partitions and good floorplan location assignments helps to improve the design resource utilization and timing performance results for cross-partition paths.

General Partitioning Guidelines

The first step in planning your design partitions is to organize your source code so that it supports good partition assignments. Although you can assign any hierarchical block of your design as a design partition or merge hierarchical blocks into the same partition, following the design guidelines presented in this section ensures better results.

Plan Design Hierarchy and Design Files

You begin the partitioning process by planning the design hierarchy. When you assign a hierarchical instance as a design partition, the partition includes the assigned instance and entities instantiated below that are not defined as separate partitions. You can use the Merge command in the Design Partitions window to combine hierarchical partitions into a single partition, as long as they have the same immediate parent partition.

When planning your design hierarchy, keep logic in the “leaves” of the hierarchy instead of having logic at the top-level of the design so that you can isolate partitions if required.

Create entities that can form partitions of approximately equal size. For example, do not instantiate small entities at the same hierarchy level, because it is more difficult to group them to form reasonably-sized partitions.

Create each entity in an independent file. The Quartus II software uses a file checksum to detect changes, and automatically recompiles a partition if its source file changes and its netlist type is set to either post-synthesis or post-fit. If the design entities for two partitions are defined in the same file, changes to the logic in one partition initiates recompilation for both partitions.

Design dependencies also affect which partitions are compiled when a source file changes. If two partitions rely on the same lower-level entity definition, changes in that lower-level entity affect both partitions. Commands such as VHDL use and Verilog HDL include create dependencies between files, so that changes to one file can trigger recompilations in all dependent files. Avoid these types of file dependencies if possible. The Partition Dependent Files report for each partition in the Analysis & Synthesis section of the Compilation report lists which files contribute to each partition.
Using Partitions with Third-Party Synthesis Tools

Incremental compilation works well with third-party synthesis tools in addition to Quartus II Integrated Synthesis. If you use a third-party synthesis tool, set up your tool to create a separate Verilog Quartus Mapping File (.vqm) or EDIF Input File (.edf) netlist for each hierarchical partition. In the Quartus II software, designate the top-level entity from each netlist as a design partition. The .vqm or .edf netlist file is treated as the source file for the partition in the Quartus II software.

For more information about incremental synthesis in third-party tools, refer to the documentation provided by your tool vendor, or the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Partition Design by Functionality and Block Size

Initially, you should partition your design along functional boundaries. In a top-level system block diagram, each block is often a natural design partition. Typically, each block of a system is relatively independent and has more signal interaction internally than interaction between blocks, which helps reduce optimizations between partition boundaries. Keeping functional blocks together means that synthesis and fitting can optimize related logic as a whole, which can lead to improved optimization.

Consider how many partitions you want to maintain in your design to determine the size of each partition. Your compilation time reduction goal is also a factor, because compiling small partitions is typically faster than compiling large partitions.

There is no minimum size for partitions; however, having too many partitions can reduce the quality of results by limiting optimization. Ensure that the design partitions are not too small. As a general guideline, each partition should contain more than approximately 2,000 logic elements (LEs) or adaptive logic modules (ALMs). If your design is incomplete when you partition the design, use previous designs to help estimate the size of each block.

Partition Design by Clock Domain and Timing Criticality

Consider which clock in your design feeds the logic in each partition. If possible, keep clock domains within one partition. When a clock signal is isolated to one partition, it reduces dependence on other partitions for timing optimization. Isolating a clock domain to one partition also allows better use of regional clock routing networks if the partition logic is constrained to one region of the design. Additionally, limiting the number of clocks within each partition simplifies the timing requirements for each partition during optimization. Use an appropriate subsystem to implement the required logic for any clock domain transfers (such as a synchronization circuit, dual-port RAM, or FIFO). You can include this logic inside the partition at one side of the transfer.

For more information about clock domains and their affect on partition design, refer to the Analyzing and Optimizing the Design Floorplan with the Chip Planner chapter in volume 2 of the Quartus II Handbook.
Try to isolate timing-critical logic from logic that you expect to easily meet timing requirements. Doing so allows you to preserve the satisfactory results for non-critical partitions and focus optimization iterations on only the timing-critical portions of the design to minimize compilation time.

**Consider What Is Changing**

When assigning partitions, you should consider what is changing in the design. Is there intellectual property (IP) or reused logic for which the source code will not change during future design iterations? If so, define the logic in its own partition so that you can compile one time and immediately preserve the results and not have to compile that part of the design again. Is logic being tuned or optimized, or are specifications changing for part of the design? If so, define changing logic in its own partition so that you can recompile only the changing part while the rest of the design remains unchanged.

As a general rule, create partitions to isolate logic that will change from logic that will not change. Partitioning a design in this way maximizes the preservation of unchanged logic and minimizes compilation time.

**Design Partition Guidelines**

Follow the partitioning guidelines presented in this section when you create or modify the HDL code for each design block that you might want to assign as a design partition. You do not need to follow all the recommendations exactly to achieve a good quality of results with the incremental compilation flow, but adhering to as many as possible maximizes your chances for success.

This section includes examples of the types of optimizations that are prevented by partition boundaries, and describes how you can structure or modify your partitions to avoid these limitations.

**Register Partition Inputs and Outputs**

Use registers at partition input and output connections that are potentially timing-critical. Registers minimize the delays on inter-partition paths and prevent the need for cross-boundary optimizations.
If every partition boundary has a register as shown in Figure 16–2, every register-to-register timing path between partitions includes only routing delay. Therefore, the timing paths between partitions are likely not timing-critical, and the Fitter can generally place each partition independently from other partitions. This advantage makes it easier to create floorplan location assignments for each separate partition, and is especially important for flows in which partitions are placed independently in separate Quartus II projects. Additionally, the partition boundary does not affect combinational logic optimization because each register-to-register logic path is contained within a single partition.

**Figure 16–2. Registering Partition I/O**

If a design cannot include both input and output registers for each partition due to latency or resource utilization concerns, choose to register one end of each connection. If you register every partition output, for example, the combinational logic that occurs in each cross-partition path is included in one partition so that it can be optimized together.

It is a good synchronous design practice to include registers for every output of a design block. Registered outputs ensure that the input timing performance for each design block is controlled exclusively within the destination logic block. For more information about I/O ports and registers for each partition, refer to “Partition Statistics Report” on page 16–34, and “Incremental Compilation Advisor” on page 16–49.

**Minimize Cross-Partition-Boundary I/O**

Minimize the number of I/O paths that cross between partition boundaries to keep logic paths within a single partition for optimization. Doing so makes partitions more independent for both logic and placement optimization.

This guideline is most important for timing-critical and high-speed connections between partitions, especially in cases where the input and output of each partition is not registered. Slow connections that are not timing-critical are acceptable because they should not impact the overall timing performance of the design. If there are timing-critical paths between partitions, rework or merge the partitions to avoid these inter-partition paths.
When dividing your design into partitions, consider the types of functions at the partition boundaries. Figure 16–3 shows an expansive function with more outputs than inputs in the left diagram, which makes a poor partition boundary, and, on the right side, a better place to assign the partition boundary that minimizes cross-partition I/Os. Adding registers to one or both sides of the cross-partition path in this example would further improve partition quality.

Another way to minimize connections between partitions is to avoid using combinational “glue logic” between partitions. You can often move the logic to the partition at one end of the connection to keep more logic paths within one partition. For example, the bottom diagram in Figure 16–4 includes a new level of hierarchy C defined as a partition instead of block B. Clearly, there are fewer I/O connections between partitions A and C than between partitions A and B.

Figure 16–3. Minimizing I/O Between Partitions by Moving the Partition Boundary

Figure 16–4. Minimizing I/O between Partitions by Modifying Glue Logic
For more information about the number of I/O ports, as well as the number of inter-partition connections for each partition, refer to “Partition Statistics Report” on page 16–34. For more information about the number of intra-partition (within a partition) and inter-partition (between partitions) timing edges, refer to “Incremental Compilation Advisor” on page 16–49.

**Examine the Need for Logic Optimization Across Partitions**

As discussed in “Partition Boundaries and Optimization” on page 16–6, partition boundaries prevent logic optimizations across partitions (except for some limited cross-boundary optimizations).

In some cases, especially if part of the design is complete or comes from another designer, the designer might not have followed these guidelines when the source code was created. These guidelines are not mandatory to implement an incremental compilation flow, but can improve the quality of results. If assigning a partition affects resource utilization or timing performance of a design block as compared to the flat design, it might be due to one of the issues described in this section. Many of the examples suggest simple changes to your partition definitions or hierarchy to move the partition boundary to improve your results.

The following guidelines ensure that your design does not require logic optimization across partition boundaries:

- “Keep Logic in the Same Partition for Optimization and Merging” on page 16–13
- “Keep Constants in the Same Partition as Logic” on page 16–15
- “Avoid Signals That Drive Multiple Partition I/O or Connect I/O Together” on page 16–16
- “Invert Clocks in Destination Partitions” on page 16–17
- “Connect I/O Pin Directly to I/O Register for Packing Across Partition Boundaries” on page 16–17
- “Do Not Use Internal Tri-States” on page 16–21
- “Include All Tri-State and Enable Logic in the Same Partition” on page 16–22
- “Include Bidirectional I/O Registers in the Same Partition (For Older Device Families)” on page 16–23
- “Summary of Guidelines Related to Logic Optimization Across Partitions” on page 16–23

**Keep Logic in the Same Partition for Optimization and Merging**

If your design logic requires logic optimization or merging to obtain optimal results, ensure that all the logic is part of the same partition because only limited cross-boundary optimizations are permitted.
If a combinational logic path is split across two partitions, the logic cannot be optimized or merged into one logic cell in the device. This effect can result in an extra logic cell in the path, increasing the logic delay. As a very simple example, consider two inverters on the same signal in two different partitions, A and B, as shown in the left diagram of Figure 16–5. To maintain correct incremental functionality, these two inverters cannot be removed from the design during optimization because they occur in different design partitions. The Quartus II software cannot use information about other partitions when it compiles each partition, because each partition is allowed to change independently from the other.

On the right side of the figure, partitions A and B are merged to group the logic in blocks A and B into one partition. If the two blocks A and B are not under the same immediate parent partition, you can create a wrapper file to define a new level of hierarchy that contains both blocks, and set this new hierarchy block as the partition. With the logic contained in one partition, the software can optimize the logic and remove the two inverters (shown in gray), which reduces the delay for that logic path. Removing two inverters is not a significant reduction in resource utilization because inversion logic is readily available in Altera device architecture. However, this example is a simple demonstration of the types of logic optimization that are prevented by partition boundaries.

In a flat design, the Fitter can also merge logical instantiations into the same physical device resource. With incremental compilation, logic defined in different partitions cannot be merged to use the same physical device resource.

For example, the Fitter can merge two single-port RAMs from a design into one dedicated RAM block in the device. If the two RAMs are defined in different partitions, the Fitter cannot merge them into one dedicated device RAM block.

This limitation is a only a concern if merging is required to fit the design in the target device. Therefore, you are more likely to encounter this issue during troubleshooting rather than during planning, if your design uses more logic than is available in the device.
Merging PLLs and Transceivers (GXB)

Multiple instances of the ALTPLL megafunction can use the same PLL resource on the device. Similarly, GXB transceiver instances can share high-speed serial interface (HSSI) resources in the same quad as other instances. The Fitter can merge multiple instantiations of these blocks into the same device resource, even if it requires optimization across partitions. Therefore, there are no restrictions for PLLs and high-speed transceiver blocks when setting up partitions.

Keep Constants in the Same Partition as Logic

Because the Quartus II software cannot fully optimize across a partition boundary, constants are not propagated across partition boundaries, except from parent partition to child partition. A signal that is constant (1/VCC or 0/GND) in one partition cannot affect another partition.

For example, the left diagram of Figure 16–6 shows part of a design in which partition A defines some signals as constants (and assumes that the other input connections come from elsewhere in the design and are not shown in the figure). Constants such as these could appear due to parameter or generic settings or configurations with parameters, setting a bus to a specific set of values, or could result from optimizations that occur within a group of logic. Because the blocks are independent, the software cannot optimize the logic in block B based on the information from block A. The right side of Figure 16–6 shows a merged partition that groups the logic in blocks A and B. If the two blocks A and B are not under the same immediate parent partition, you can create a wrapper file to define a new level of hierarchy that contains both blocks, and set this new hierarchical block as the partition.

Within the single merged partition, the Quartus II software can use the constants to optimize and remove much of the logic in block B (shown in gray), as shown in Figure 16–6.

Figure 16–6. Keeping Constants in the Same Partition as the Logic They Feed

For more information about how many input ports are fed by GND or VCC, refer to “Partition Statistics Report” on page 16–34. For more information about port connections, refer to “Incremental Compilation Advisor” on page 16–49.
Avoid Signals That Drive Multiple Partition I/O or Connect I/O Together

Do not use the same signal to drive multiple ports of a single partition or directly connect two ports of a partition. If the same signal drives multiple ports of a partition, or if two ports of a partition are directly connected, those ports are logically equivalent. However, the software has limited information about connections made in another partition (including the top-level partition), the compilation cannot take advantage of the equivalence. This restriction usually produces sub-optimal results.

If your design has these types of connections, redefine the partition boundaries to remove the affected ports. If one signal from a higher-level partition feeds two input ports of the same partition, feed the one signal into the partition, and then make the two connections within the partition. If an output port drives an input port of the same partition, the connection can be made internally without going through any I/O ports. If an input port drives an output port directly, the connection can likely be implemented without the ports in the lower-level partition by connecting the signals in a higher-level partition.

Figure 16–7 shows an example of one signal driving more than one port. The left diagram shows a design where a single clock signal is used to drive both the read and write clocks of a RAM block. Because the RAM block is compiled as a separate partition A, the RAM block is implemented as though there are two unique clocks. If you know that the port connectivity will not change (that is, the ports will always be driven by the same signal in the top-level partition), redefine the port interface so that there is only a single port that can drive both connections inside the partition. You can create a wrapper file to define a partition that has fewer ports, as shown in the diagram on the right side. With the single clock fed into the partition, the RAM can be optimized into a single-clock RAM instead of a dual-clock RAM. Single-clock RAM can provide better performance in the device architecture. Additionally, partition A might use two global routing lines for the two copies of the clock signal. Partition B can use one global line that fans out to all destinations. Using just the single port connection prevents overuse of global routing resources.

For more information about partition ports that have the same driving signal and ports that are directly connected together, refer to “Incremental Compilation Advisor” on page 16–49.
Invert Clocks in Destination Partitions

For best results, clock inversion should be performed in the destination logic array block (LAB) because each LAB contains clock inversion circuitry in the device architecture. In a flat compilation, the Quartus II software can optimize a clock inversion to propagate it to the destination LABs regardless of where the inversion takes place in the design hierarchy. However, clock inversion cannot propagate through a partition boundary (except from a parent partition to a child partition) to take advantage of the inversion architecture in the destination LABs.

With partition boundaries as shown in the left diagram of Figure 16–8, the Quartus II software uses logic to invert the signal in the partition that defines the inversion (the top-level partition in this example), and then routes the signal on a global clock resource to its destinations (in partitions A and B). The inverted clock acts as a gated clock with high skew. A better solution is to invert the clock signal in the destination partitions as shown on the right side of the diagram. In this case, the correct logic and routing resources can be used, and the signal does not behave like a gated clock.

Figure 16–8 shows the clock signal inversion in the destination partitions.

![Figure 16–8. Inverting Clock Signal in Destination Partitions](image)

Notice that this diagram also shows another example of a single pin feeding two ports of a partition boundary. In the left diagram, partition B does not have the information that the clock and inverted clock come from the same source. In the right diagram, partition B has more information to help optimize the design because the clock is connected as one port of the partition.

Connect I/O Pin Directly to I/O Register for Packing Across Partition Boundaries

The Quartus II software allows cross-partition register packing of I/O registers in certain cases where your input and output pins are defined in the top-level hierarchy (and the top-level partition), but the corresponding I/O registers are defined in other partitions.

Input pin cross-partition register packing requires the following specific circumstances:

- The input pin feeds exactly one register.
The path between the input pin and register includes only input ports of partitions that have one fan-out each.

Output pin cross-partition register packing requires the following specific circumstances:
- The register feeds exactly one output pin.
- The output pin is fed by only one signal.
- The path between the register and output pin includes only output ports of partitions that have one fan-out each.

The following examples of I/O register packing illustrate this point using Block Design File (.bdf) schematics to describe the design logic.

Example 1—Output Register in Partition Feeding Multiple Output Pins
In this example, a subdesign contains a single register, as shown in Figure 16–9.

If the top-level design instantiates the subdesign with a single fan-out directly feeding an output pin, and designates the subdesign as a separate design partition, the Quartus II software can perform cross-partition register packing because the single partition port feeds the output pin directly.

In Example 1, the top-level design instantiates the subdesign in Figure 16–9 as an output register with more than one fan-out signal, as shown in Figure 16–10.

In this case, the Quartus II software does not perform output register packing. If there is a Fast Output Register assignment on pin `out`, the software issues a warning that the Fitter cannot pack the node to an I/O pin because the node and the I/O cell are connected across a design partition boundary.
This type of cross-partition register packing is not allowed because it requires modification to the interface of the subdesign partition. To perform incremental compilation, the Quartus II software must preserve the interface of design partitions.

To allow the Quartus II software to pack the register in the subdesign from Figure 16–9 with the output pin `out` in Figure 16–10, restructure your HDL code so that output registers directly connect to output pins by making one of the following changes:

- Place the register in the same partition as the output pin. The simplest method is to move the register from the subdesign partition into the partition containing the output pin. Doing so guarantees that the Fitter can optimize the two nodes without violating partition boundaries.

- Duplicate the register in your subdesign HDL as shown in Figure 16–11 so that each register feeds only one pin, and then connect the extra output pin to the new port in the top-level design as shown in Figure 16–12. Doing so converts the cross-partition register packing into the simplest case where each register has a single fan-out.

Figure 16–11. Modified Subdesign from Figure 16–9 with Two Output Registers and Two Output Ports

Figure 16–12. Modified Top-Level Design from Figure 16–10 Connecting Two Output Ports to Output Pins
Example 2—Input Register in Partition Fed by an Inverted Input Pin or Output Register in Partition Feeding an Inverted Output Pin

In this example, a subdesign designated as a separate partition contains a register, as shown in Figure 16–9. The top-level design in Figure 16–13 instantiates the subdesign as an input register with the input pin inverted. The top-level design in Figure 16–14 instantiates the subdesign as an output register with the signal inverted before feeding an output pin.

**Figure 16–13. Top-Level Design Instantiating the Subdesign in Figure 16–9 as an Input Register with an Inverted Input Pin**

In these cases, the Quartus II software does not perform register packing. If there is a *Fast Input Register* assignment on pin \( \text{in} \), as shown in Figure 16–13, or a *Fast Output Register* assignment on pin \( \text{out} \), as shown in Figure 16–14, the Quartus II software issues a warning that the Fitter cannot pack the node to an I/O pin because the node and I/O cell are connected across a design partition boundary.

This type of register packing is not allowed because it requires moving logic across a design partition boundary to place into a single I/O device atom. To perform register packing, either the register must be moved out of the subdesign partition, or the inverter must be moved into the subdesign partition to be implemented in the register.

To allow the Quartus II software to pack the register in the subdesign from Figure 16–9 with the input pin \( \text{in} \), as shown in Figure 16–13 or the output pin \( \text{out} \), as shown in Figure 16–14, restructure your HDL code to place the register in the same partition as the inverter by making one of the following changes:

- Move the register from the subdesign partition into the top-level partition containing the pin. Doing so ensures that the Fitter can optimize the I/O register and inverter without violating partition boundaries.
Move the inverter from the top-level block into the subdesign, and then connect the subdesign directly to a pin in the top-level design. Doing so allows the Fitter to optimize the inverter into the register implementation, so that the register is directly connected to a pin, which enables register packing.

**Do Not Use Internal Tri-States**

Internal tri-state signals are not recommended for FPGAs because the device architecture does not include internal tri-state logic. If designs use internal tri-states in a flat design, the tri-state logic is usually converted to OR gates or multiplexing logic. If tri-state logic occurs on a hierarchical partition boundary, the Quartus II software cannot convert the logic to combinational gates because the partition could be connected to a top-level device I/O through another partition.

Figure 16–15 and Figure 16–16 show a design with partitions that are not supported for incremental compilation due to the internal tri-state output logic on the partition boundaries. Instead of using internal tri-state logic for partition outputs, implement the correct logic to select between the two signals. Doing so is good practice even when there are no partitions, because such logic explicitly defines the behavior for the internal signals instead of relying on the Quartus II software to convert the tri-state signals into logic.

**Figure 16–15. Unsupported Internal Tri-State Signals**

Design results in Quartus II error message:
The software cannot synthesize this design and maintain incremental functionality.

**Figure 16–16. Merged Partition Allows Synthesis to Convert Internal Tri-State Logic to Combinational Logic**

Merged partition allows synthesis to convert tri-state logic into combinational logic.
Do not use tri-state signals or bidirectional ports on hierarchical partition boundaries, unless the port is connected directly to a top-level I/O pin on the device. If you must use internal tri-state logic, ensure that all the control and destination logic is contained in the same partition, in which case the Quartus II software can convert the internal tri-state signals into combinational logic as in a flat design. In this example, you can also merge all three partitions into one partition, as shown in Figure 16–16, to allow the Quartus II software to treat the logic as internal tri-state and perform the same type of optimization as a flat design. If possible, you should avoid using internal tri-state logic in any Altera FPGA design to ensure that you get the desired implementation when the design is compiled for the target device architecture.

**Include All Tri-State and Enable Logic in the Same Partition**

When multiple output signals use tri-state logic to drive a device output pin, the Quartus II software merges the logic into one tri-state output pin. The Quartus II software cannot merge tri-state outputs into one output pin if any of the tri-state logic occurs on a partition boundary. Similarly, output pins with an output enable signal cannot be packed into the device I/O cell if the output enable logic is part of a different partition from the output register. To allow register packing for output pins with an output enable signal, structure your HDL code or design partition assignments so that the register and enable logic are defined in the same partition.

Figure 16–17 shows a design with tri-state output signals that feed a device bidirectional I/O pin (assuming that the input connection feeds elsewhere in the design and is not shown in the figure). In the left diagram below, the tri-state output signals appear as the outputs of two separate partitions. In this case, the Quartus II software cannot implement the specified logic and maintain incremental functionality. In the right diagram, partitions A and B are merged to group the logic from the two blocks. With this single partition, the Quartus II software can merge the two tri-state output signals and implement them in the tri-state logic available in the device I/O element.

**Figure 16–17. Including All Tri-State Output Logic in the Same Partition**

![Diagram showing tri-state signals and enable logic in partitions](image-url)
**Include Bidirectional I/O Registers in the Same Partition (For Older Device Families)**

For a bidirectional partition port that feeds a bidirectional I/O pin at the top level, all logic that forms the bidirectional I/O cell must reside in the same partition in the Stratix II, Stratix, Cyclone® II, and Cyclone device families (this restriction does not apply to newer devices). Additionally, as discussed in the previous two recommendations, the I/O logic must feed the I/O pin without any intervening logic.

In Figure 16–18, all the I/O logic must be defined inside the same partition for the Quartus II software to implement all three registers in the I/O element along with the tri-state logic in the affected devices. The logic connected to the registers can occur in the same partition or any other partition; only the I/O registers must be grouped with the tri-state logic definition. The bidirectional I/O port of the partition must be directly connected to the bidirectional device pin at the top level. The signal can go through several partition boundaries if necessary, as long as the connection path contains no logic.

**Figure 16–18. Including All Bidirectional I/O Registers in the Same Partition (for Older Devices)**

Bidirectional logic is within one partition, and I/O logic directly feeds I/O pin

**Summary of Guidelines Related to Logic Optimization Across Partitions**

To ensure that your design does not require logic optimization across partitions, follow the guidelines in this section:

- Include logic in the same partition for optimization and merging
- Include constants in the same partition as logic
- Avoid signals that drive multiple partition I/O or connect I/O together
- Invert clocks in destination partitions
- Connect I/O directly to I/O register for packing across partition boundaries
- Do not use internal tri-states
- Include all tri-state and enable logic in the same partition
- Include bidirectional I/O registers in the same partition (in older device families)
Remember that these guidelines are not mandatory when implementing an incremental compilation flow, but can improve the quality of results. When creating source design code, follow these guidelines and organize your HDL code to support good partition boundaries. For designs that are complete, assess whether assigning a partition affects the resource utilization or timing performance of a design block as compared to the flat design. Make the appropriate changes to your design or hierarchy, or merge partitions as required, to improve your results.

**Consider a Cascaded Reset Structure**

Designs typically have a global asynchronous reset signal where a top-level signal feeds all partitions. To minimize skew for the high fan-out signal, the global reset signal is typically placed onto a global routing resource.

In some cases, having one global reset signal can lead to recovery and removal time problems. This issue is not specific to incremental flows; it could be applicable in any large high-speed design. In an incremental flow, the global reset signal creates a timing dependency between the top-level partition and lower-level partitions.

For incremental compilation, it is helpful to minimize the impact of global structures. To isolate each partition, consider adding reset synchronizers. Using cascaded reset structures, the intent is to reduce the inter-partition fan-out of the reset signal, thereby minimizing the effect of the global signal. Reducing the fan-out of the global reset signal also provides more flexibility in routing the cascaded signals, and might help recovery and removal times in some cases.

This recommendation can help in large designs, regardless of whether you are using incremental compilation. However, if one global signal can feed all the logic in its domain and meet recovery and removal times, this recommendation may not be applicable for your design. Minimizing global structures is more relevant for high-performance designs where meeting timing on the reset logic can be challenging. Isolating each partition and allowing more flexibility in global routing structures is an additional advantage in incremental flows.

If you add additional reset synchronizers to your design, latency is also added to the reset path, so ensure that this is acceptable in your design. Additionally, parts of the design may come out of the reset state in different clock cycles. You can balance the latency or add hand-shaking logic between partitions, if necessary, to accommodate these differences.

The signal is first synchronized on the chip following good synchronous design practices, meaning that the design asynchronously resets, but synchronously releases from reset to avoid any race conditions or metastability problems. Then, to minimize the impact of global structures, the circuit employs a divide-and-conquer approach for the reset structure. By implementing a cascaded reset structure, the reset paths for each partition are independent. This structure reduces the effect of inter-partition dependency because the inter-partition reset signals can now be treated as false paths for timing analysis. In some cases, the reset signal of the partition can be placed on local lines to reduce the delay added by routing to a global routing line. In other cases, the signal can be routed on a regional or quadrant clock signal.
Figure 16–19 shows a cascaded reset structure.

**Figure 16–19. Cascaded Reset Structure**

This circuit design can help you achieve timing closure and partition independence for your global reset signal. Evaluate the circuit and consider how it works for your design.

For more information and design recommendations for reset structures, refer to the *Recommended Design Practices* chapter in volume 1 of the *Quartus II Handbook*.

**Design Partition Guidelines for Third-Party IP Delivery**

This section includes additional design guidelines that can improve incremental compilation flows where exported partitions are developed independently. These guidelines are not always required, but are usually recommended if the design includes partitions compiled in a separate Quartus II project, such as when delivering intellectual property (IP). A unique challenge of IP delivery for FPGAs is the fact that the partitions developed independently must share a common set of resources. To minimize issues that might arise from sharing a common set of resources, you can design partitions within a single Quartus II project, or a copy of the top-level design. A common project ensures that designers have a consistent view of the top-level design framework, as described in “Project Management in Team-Based Design Flows” on page 16–4.

Alternatively, an IP designer can export just the post-synthesis results to be integrated in the top-level design when the post-fitting results from the IP project are not required. Using a post-synthesis netlist provides more flexibility to the Quartus II Fitter, so that less resource allocation is required. If a common project is not possible, especially when the project lead plans to integrate the IP’s post-fitting results, it is important that the project lead and IP designer clearly communicate their requirements.
Allocate Logic Resources

In an incremental compilation design flow in which designers, such as third-party IP providers, optimize partitions and then export them to a top-level design, the Quartus II software places and routes each partition separately. In some cases, partitions can use conflicting resources when combined at the top level. Allocation of logic resources requires that you decide on a set of logic resources (including I/O, LAB logic blocks, RAM and DSP blocks) that the IP block will "own". This process can be interactive; the project lead and the IP designer might work together to determine what resources are required for the IP block and are available in the top-level design.

You can constrain logic utilization for the IP core using design floorplan location assignments, as described in “Introduction to Design Floorplans” on page 16–40. The design should specify I/O pin locations with pin assignments.

You can also specify limits for Quartus II synthesis to allocate and balance resources. This procedure can also help if device resources are overused in the individual partitions during synthesis.

In the standard synthesis flow, the Quartus II software can perform automated resource balancing for DSP blocks or RAM blocks and convert some of the logic into regular logic cells to prevent overuse.

You can use the Quartus II synthesis options to control inference of megafunctions that use the DSP, or RAM blocks. You can also use the MegaWizard™ Plug-In Manager to customize your RAM or DSP megafunctions to use regular logic instead of the dedicated hardware blocks.

For more information about resource balancing DSP and RAM blocks when using Quartus II synthesis, refer to the “Megafunction Inference Control” section in the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook. For tips about resource balancing and reducing resource utilization, refer to the appropriate “Resource Utilization Optimization Techniques” section in the Area and Timing Optimization chapter in volume 2 of the Quartus II Handbook.

For information about how to set global logic options for partitions, refer to More Analysis & Synthesis Settings Dialog Box in Quartus II Help.

Allocate Global Routing Signals and Clock Networks if Required

In most cases, you do not have to allocate global routing signals because the Quartus II software finds the best solution for the global signals. However, if your design is complex and has multiple clocks, especially for a partition developed by a third-party IP designer, you may have to allocate global routing resources between various partitions.

Global routing signals can cause conflicts when independent partitions are integrated into a top-level design. The Quartus II software automatically promotes high fan-out signals to use global routing resources available in the device. Third-party partitions can use the same global routing resources, thus causing conflicts in the top-level design. Additionally, LAB placement depends on whether the inputs to the logic cells within the LAB use a global clock signal. Problems can occur if a design does not use a global signal in a lower-level partition, but does use a global signal in the top-level design.
If the exported IP core is small, you can reduce the potential for problems by using constraints to promote clock and high fan-out signals to regional routing signals that cover only part of the device, instead of global routing signals. In this case, the Quartus II software is likely to find a routing solution in the top-level design because there are many regional routing signals available on most Altera devices, and designs do not typically overuse regional resources.

To ensure that an IP block can utilize a regional clock signal, view the resource coverage of regional clocks in the Chip Planner, and then align LogicLock regions that constrain partition placement with available global clock routing resources. For example, if the LogicLock region for a particular partition is limited to one device quadrant, that partition’s clock can use a regional clock routing type that covers only one device quadrant. When all partition logic is available, the project lead can compile the entire design at the top level with floorplan assignments to allow the use of regional clocks that span only a part of the device.

If global resources are heavily used in the overall design, or the IP designer requires global clocks for their partition, you can set up constraints to avoid signal overuse at the top-level by assigning the appropriate type of global signals or setting a maximum number of clock signals for the partition.

You can use the Global Signal assignment to force or prevent the use of a global routing line, making the assignment to a clock source node or signal. You can also assign certain types of global clock resources in some device families, such as regional clocks. For example, if you have an IP core, such as a memory interface that specifies the use of a dual regional clock, you can constrain the IP to part of the device covered by a regional clock and change the Global Signal assignment to use a regional clock. This type of assignment can reduce clocking congestion and conflicts.

Alternatively, partition designers can specify the number of clocks allowed in the project using the maximum clocks allowed options in the More Fitter Settings dialog box. Specify Maximum number of clocks of any type allowed, or use the Maximum number of global clocks allowed, Maximum number of regional clocks allowed, and Maximum number of periphery clocks allowed options to restrict the number of clock resources of a particular type in your design.

If you require more control when planning a design with integrated partitions, you can assign a specific signal to use a particular clock network in Stratix II and newer device families by assigning the clock control block instance called CLKCTRL. You can make a point-to-point assignment from a clock source node to a destination node, or a single-point assignment to a clock source node with the Global Clock CLKCTRL Location logic option. Set the assignment value to the name of the clock control block: CLKCTRL_G<global network number> for a global routing network, or CLKCTRL_R<regional network number> for a dedicated regional routing network in the device.

If you want to disable the automatic global promotion performed in the Fitter to prevent other signals from being placed on global (or regional) routing networks, turn off the Auto Global Clock and Auto Global Register Control Signals options in the More Fitter Settings dialog box.

For information about how to disable automatic global promotion, refer to More Fitter Settings Dialog Box in Quartus II Help.
If you are using design partition scripts for independent partitions, the Quartus II software can automatically write the commands to pass global constraints and turn off automatic options.

For more information about how to generate design partition scripts, refer to Generating Design Partition Scripts for Project Management in Quartus II Help.

For more information about how clock networks affect partition design, refer to the Analyzing and Optimizing the Design Floorplan with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

Alternatively, to avoid problems when integrating partitions into the top-level design, you can direct the Fitter to discard the placement and routing of the partition netlist by using the post-synthesis netlist, which forces the Fitter to reassign all the global signals for the partition when compiling the top-level design.

### Assign Virtual Pins

Virtual pins map lower-level design I/Os to internal cells. If you are developing an IP block in an independent Quartus II project, use virtual pins when the number of I/Os on a partition exceeds the device I/O count, and to increase the timing accuracy of cross-partition paths.

You can create a virtual pin assignment in the Assignment Editor for partition I/Os that will become internal nodes in the top-level design. Leave the clock pins mapped to I/O pins to ensure proper routing.

You can specify locations for the virtual pins that correspond to the placement of other partitions, and also make timing assignments to the virtual pins to define a timing budget, as described in the following section. Virtual pins are created automatically from the top-level design if you use design partition scripts. The scripts place the virtual pins to correspond with the placement of the other partitions in the top-level design.

For more information about how to generate design partition scripts, refer to Generating Design Partition Scripts for Project Management in Quartus II Help.

Tri-state outputs cannot be assigned as virtual pins because internal tri-state signals are not supported in Altera devices. Connect the signal in the design with regular logic, or allow the software to implement the signal as an external device I/O pin.

### Perform Timing Budgeting if Required

If you optimize partitions independently and integrate them to the top-level design, or compile with empty partitions, any unregistered paths that cross between partitions are not optimized as entire paths. In these cases, the Quartus II software has no information about the placement of the logic that connects to the I/O ports. If the logic in one partition is placed far away from logic in another partition, the routing
delay between the logic can lead to problems in meeting timing requirements. You can reduce this effect by ensuring that input and output ports of the partitions are registered whenever possible. Additionally, using the same top-level project framework helps to avoid this problem by providing the software with full information about other design partitions in the top-level design.

To ensure that the software correctly optimizes the input and output logic in any independent partitions, you might be required to perform some manual timing budgeting. For each unregistered timing path that crosses between partitions, make timing assignments on the corresponding I/O path in each partition to constrain both ends of the path to the budgeted timing delay. Assigning a timing budget for each part of the connection ensures that the software optimizes the paths appropriately.

When performing manual timing budgeting in a partition for I/O ports that become internal partition connections in a top-level design, you can assign location and timing constraints to the virtual pin that represents each connection to further improve the quality of the timing budget. Refer to “Assign Virtual Pins” on page 16–28 for a description of virtual pins.

If you use design partition scripts, the Quartus II software can write I/O timing budget constraints automatically for virtual pins.

For more information about how to generate design partition scripts, refer to Generating Design Partition Scripts for Project Management in Quartus II Help.

**Drive Clocks Directly**

When partitions are exported from another Quartus II project, you should drive partition clock inputs directly with device clock input pins.

Connecting the clock signal directly avoids any timing analysis difficulties with gated clocks. Clock gating is never recommended for FPGA designs because of potential glitches and clock skew. Clock gating can be especially problematic with exported partitions because the partitions have no information about gating that takes place at the top-level design or in another partition. If a gated clock is required in a partition, perform the gating within that partition, as described for clock inversion in “Invert Clocks in Destination Partitions” on page 16–17.

Direct connections to input clock pins also allows design partition scripts to send constraints from the top-level device pin to lower-level partitions.

**Recreate PLLs for Lower-Level Partitions if Required**

If you connect a PLL in your top-level design to partitions designed in separate Quartus II projects by third-party IP designers, the IP partitions do not have information about the multiplication, phase shift, or compensation delays for the PLL in the top-level design. To accommodate the PLL timing, you can make appropriate timing assignments in the projects created by IP designers to ensure that clocks are not left unconstrained or constrained with an incorrect frequency. Alternatively, you can duplicate the top-level PLL (or other derived clock logic) in the design file for the project created by the IP designer to ensure that you have the correct PLL parameters and clock delays for a complete and accurate timing analysis.
If the project lead creates a copy of the top-level project framework that includes all the settings and constraints needed for the design, this framework should include PLLs and other interface logic if this information is important to optimize partitions.

If you use a separate Quartus II project for an independent design block (such as when a designer or third-party IP provider does not have access to the entire design framework), include a copy of the top-level PLL in the lower-level partition as shown in Figure 16–20.

In either case, the IP partition in the separate Quartus II project should contain just the partition logic that will be exported to the top-level design, while the full project includes more information about the top-level design. When the partition is complete, you can export just the partition without exporting the auxiliary PLL components to the top-level design. When you export a partition, the Quartus II software exports any hierarchy under the specified partition into the Quartus II Exported Partition File (.qxp), but does not include logic defined outside the partition (the PLL in this example).

**Figure 16–20. Recreating a Top-Level PLL in a Lower-Level Partition**

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### Checking Partition Quality

This section provides an overview of tools you can use to create and analyze partitions in the Quartus II software. Take advantage of these tools to assess your partition quality, and use the information to improve your design or assignments as required to achieve the best results.

#### Incremental Compilation Advisor

You can use the Incremental Compilation Advisor to ensure that your design follows Altera’s recommendations for creating design partitions and implementing the incremental compilation design flow methodology. Each recommendation in the Incremental Compilation Advisor provides an explanation, describes the effect of the recommendation, and provides the action required to make the suggested change.

For more information about the Incremental Compilation Advisor, refer to [Incremental Compilation Advisor Command](#) and [Example of Using the Incremental Compilation Advisor to Identify Non-Global Ports That Are Not Registered](#) in Quartus II Help, and the [Quartus II Incremental Compilation for Hierarchical and Team-Based Design](#) chapter in volume 1 of the [Quartus II Handbook](#).
Design Partition Planner

The Design Partition Planner allows you to view design connectivity and hierarchy, and can assist you in creating effective design partitions that follow the guidelines in this chapter. You can also use the Design Partition Planner to optimize design performance by isolating and resolving failing paths on a partition-by-partition basis.

To view a design and create design partitions in the Design Partition Planner, you must first compile the design, or perform Analysis & Synthesis. In the Design Partition Planner, the design appears as a single top-level design block, with lower-level instances displayed as color-specific boxes.

In the Design Partition Planner, you can show connectivity between blocks and extract instances from the top-level design block. When you extract entities, connection bundles are drawn between entities, showing the number of connections existing between pairs of entities. When you have extracted a design block that you want to set as a design partition, right-click that design block, and then click Create Design Partition.

The Design Partition Planner also has an auto-partition feature that creates partitions based on the size and connectivity of the hierarchical design blocks. You can right-click the design block you want to partition (such as the top-level design hierarchy), and then click Auto-Partition Children. You can then analyze and adjust the partition assignments as required.
Figure 16–21 shows the Design Partition Planner after making a design partition assignment to one instance and dragging another instance away from the top-level block within the same partition (two design blocks in the pale blue shaded box). The figure shows the connections between each partition and information about the size of each design instance.

You can switch between connectivity display mode and hierarchical display mode, or temporarily to a view-only hierarchy display. You can also remove the connection lines between partitions and I/O banks by turning off Display connections to I/O banks, or use the settings on the Connection Counting tab in the Bundle Configuration dialog box to adjust how the connections are counted in the bundles.

To optimize design performance, confine failing paths within individual design partitions so that there are no failing paths passing between partitions, as discussed in earlier sections. In the top-level entity, child entities that contain failing paths are marked by a small red dot in the upper right corner of the entity box.

To view the critical timing paths from a timing analyzer report, first perform a timing analysis on your design, and then in the Design Partition Planner, click Show Timing Data on the View menu.

For more information about the Design Partition Planner, refer to About the Design Partition Planner and Using the Design Partition Planner in Quartus II Help.
Viewing Design Partition Planner and Floorplan Side-by-Side

You can use the Design Partition Planner together with the Chip Planner to analyze natural placement groupings. This information can help you decide whether the design blocks should be grouped together in one partition, or whether they will make good partitions in the next compilation. It can also help determine whether the logic can easily be constrained by a LogicLock region. If logic naturally groups together when compiled without placement constraints, you can probably assign a reasonably sized LogicLock region to constrain the placement for subsequent compilations. You can experiment by extracting different design blocks in the Design Partition Planner and viewing the placement results of those design blocks from the previous compilation.

To view the Design Partition Planner and Chip Planner side-by-side, open the Design Partition Planner, and then open the Chip Planner and select the Design Partition Planner task. The Design Partition Planner task displays the physical locations of design entities with the same colors as in the Design Partition Planner.

In the Design Partition Planner, you can extract instances of interest from their parents by dragging and dropping, or with the Extract from Parent command. Evaluate the physical locations of instances in the Chip Planner and the connectivity between instances displayed in the Design Partition Planner. An entity is generally not suitable to be set as a separate design partition or constrained in a LogicLock region if the Chip Planner shows it physically dispersed over a noncontiguous area of the device after compilation. Use the Design Partition Planner to analyze the design connections. Child instances that are unsuitable to be set as separate design partitions or placed in LogicLock regions can be returned to their parent by dragging and dropping, or with the Collapse to Parent command.
Figure 16–22 shows a design displayed in the Design Partition Planner and the Chip Planner with different colors for the top-level design and the three major design instances.

**Figure 16–22. Design Partition Planner and Chip Planner**

For more information about the Design Partition Planner, refer to *About the Design Partition Planner* and *Using the Design Partition Planner* in Quartus II Help.

### Partition Statistics Report

You can view statistics about design partitions in the Partition Merge Partition Statistics report and the Statistics tab of the Design Partitions Properties dialog box. These reports are useful when optimizing your design partitions, or when compiling the completed top-level design in a team-based compilation flow to ensure that partitions meet the guidelines discussed in this chapter.

The Partition Merge Partition Statistics report in the Partition Merge section of the Compilation report lists statistics about each partition. The statistics for each partition (each row in the table) include the number of logic cells, as well as the number of input and output pins and how many are registered. This report also lists how many ports are unconnected, or driven by a constant $V_{CC}$ or GND. You can use this information to assess whether you have followed the guidelines for partition boundaries.
You can also view statistics about the resource and port connections for a particular partition on the Statistics tab of the Design Partition Properties dialog box. The Show All Partitions button allows you to view all the partitions in the same report. The Partition Merge Partition Statistics report also shows statistics for the Internal Congestion: Total Connections and Registered Connections. This information represents how many signals are connected within the partition. It then lists the inter-partition connections for each partition, which helps you to see how partitions are connected to each other.

For more information about the Partition Merge Reports, refer to Partition Merge Reports in Quartus II Help.

Report Partition Timing in the TimeQuest Timing Analyzer

The Report Partitions diagnostic report and the report_partitions SDC command in the TimeQuest analyzer produce a Partition Timing Overview and Partition Timing Details table, which lists the partitions, the number of failing paths, and the worst case timing slack within each partition.

You can use these reports to analyze the location of the critical timing paths in the design in relation to partitions. If a certain partition contains many failing paths, or failing inter-partition paths, you might be able to change your partitioning scheme and improve timing performance.

For more information about the TimeQuest report_timing command and reports, see the Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.

Check if Partition Assignments Impact the Quality of Results

You can ensure that you limit negative effect on the quality of results by following an iterative methodology during the partitioning process. In any incremental compilation flow where you can compile the source code for every partition during the partition planning phase, Altera recommends the following iterative flow:

1. Start with a complete design that is not partitioned and has no location or LogicLock region assignments.

   After Analysis & Synthesis and Partition Merge, perform a placement and timing analysis estimate with the Start Early Timing Estimate command. To run a full compilation instead, use the Start Compilation command.

2. Record the quality of results from the Compilation report (timing slack or f_{MAX}, area and any other relevant results).

3. Create design partitions following the guidelines described in this chapter.

4. Perform another early timing estimate or a full compilation.

5. Record the quality of results from the Compilation report. If the quality of results is significantly worse than those obtained in the previous compilation, repeat step 3 through step 5 to change your partition assignments and use a different partitioning scheme.
6. Even if the quality of results is acceptable, you can repeat step 3 through step 5 by further dividing a large partition into several smaller partitions, which can improve compilation time in subsequent incremental compilations. You can repeat these steps until you achieve a good trade-off point (that is, all critical paths are localized within partitions, the quality of results is not negatively affected, and the size of each partition is reasonable).

You can also remove or disable partition assignments defined in the top-level design at any time during the design flow to compile the design as one flat compilation and get all possible design optimizations to assess the results. To disable the partitions without deleting the assignments, use the **Ignore partition assignments during compilation** option on the **Incremental Compilation** page of the **Settings** dialog box in the Quartus II software. This option disables all design partition assignments in your project and runs a full compilation, ignoring all partition boundaries and netlists. This option can be useful if you are using partitions to reduce compilation time as you develop various parts of the design, but can run a long compilation near the end of the design cycle to ensure the design meets its timing requirements.

**Including SDC Constraints from Lower-Level Partitions for Third-Party IP Delivery**

When exported partitions are compiled in a separate Quartus II project, such as when a third-party designer is delivering IP, the project lead must transfer the top-level project framework information and constraints to the partitions, so that each designer has a consistent view of the constraints that apply to the entire design. If the independent partition designers make any changes or add any constraints, they might have to transfer new constraints back to the project lead, so that these constraints are included in final timing sign-off of the entire design. Many assignments from the partition are carried with the partition into the top-level design; however, SDC format constraints for the TimeQuest analyzer are not copied into the top-level design automatically.

Passing additional timing constraints from a partition to the top-level design must be managed carefully. This section provides recommendations for managing the timing constraints in a third-party IP delivery flow. You can design within a single Quartus II project or a copy of the top-level design to simplify constraint management.

To ensure that there are no conflicts between the project lead’s top-level constraints and those added by the third-party IP designer, use two .sdc files for each separate Quartus II project: an .sdc created by the project lead that includes project-wide constraints, and an .sdc created by the IP designer that includes partition-specific constraints.
This section uses the example design shown in Figure 16–23 to illustrate these recommendations. The top-level design instantiates a lower-level design block called `module_A` that is set as a design partition and developed by an IP designer in a separate Quartus II project.

**Figure 16–23. Example Design to Illustrate SDC Constraints**

In this top-level design, there is a single clock setting called `clk` associated with the FPGA input called `top_level_clk`. The top-level `.sdc` contains the following constraint for the clock:

```
create_clock -name {clk} -period 3.000 -waveform { 0.000 1.500 } [get_ports {TOP_LEVEL_CLK}]
```

**Creating an .sdc File with Project-Wide Constraints**

The `.sdc` with project-wide constraints for the separate Quartus II project should contain all constraints that are not completely localized to the partition. The `.sdc` should be maintained by the project lead. The project lead must ensure that these timing constraints are delivered to the individual partition owners and that they are syntactically correct for each of the separate Quartus II projects. This communication can be challenging when the design is in flux and hierarchies change. The project lead can use design partition scripts to automatically pass some of these constraints to the separate Quartus II projects.

For more information about how to generate design partition scripts, refer to *Generating Design Partition Scripts for Project Management* in Quartus II Help.

The `.sdc` with project-wide constraints is used in the partition, but is not exported back to the top-level design. The partition designer should not modify this file. If changes are necessary, they should be communicated to the project lead, who can then update the SDC constraints and distribute new files to all partition designers as required.

The `.sdc` should include clock creation and clock constraints for any clock used by more than one partition. These constraints are particularly important when working with complex clocking structures, such as the following:

- Cascaded clock multiplexers
- Cascaded PLLs
- Multiple independent clocks on the same clock pin
- Redundant clocking structures required for secure applications
- Virtual clocks and generated clocks that are consistently used for source synchronous interfaces
- Clock uncertainties

Additionally, the .sdc with project-wide constraints should contain all project-wide timing exception assignments, such as the following:

- Multicycle assignments, `set_multicycle_path`
- False path assignments, `set_false_path`
- Maximum delay assignments, `set_max_delay`
- Minimum delay assignments, `set_min_delay`

The project-wide .sdc can also contain any `set_input_delay` or `set_output_delay` constraints that are used for ports in separate Quartus II projects, because these represent delays external to a given partition. If the partition designer wants to set these constraints within the separate Quartus II projects, the team must ensure that the I/O port names are identical in all projects so that the assignments can be integrated successfully without changes.

Similarly, a constraint on a path that crosses a partition boundary should be in the project-wide .sdc, because it is not completely localized in a separate Quartus II project.

**Example Step 1: Project Lead Produces .sdc with Project-Wide Constraints for Lower-Level Partitions**

The device input `top_level_clk` in Figure 16–23 drives the `input_clk` port of `module_A`. To make sure the clock constraint is passed correctly to the partition, the project lead creates an .sdc with project-wide constraints for `module_A` that contains the following command:

```
create_clock -name {clk} -period 3.000 -waveform { 0.000 1.500 }
[get_ports {INPUT_CLK}]
```

The designer of `module_A` includes this .sdc as part of the separate Quartus II project.

**Creating an .sdc with Partition-Specific Constraints**

The .sdc with partition-specific constraints should contain all constraints that affect only the partition. For example, a `set_false_path` or `set_multicycle_path` constraint for a path entirely within the partition should be in the partition-specific .sdc. These constraints are required for correct compilation of the partition, but need not be present in any other separate Quartus II projects.

The partition-specific .sdc should be maintained by the partition designer; they must add any constraints required to properly compile and analyze their partition.
The partition-specific .sdc is used in the separate Quartus II project and must be exported back to the project lead for the top-level design. The project lead must use the partition-specific constraints to properly constrain the placement, routing, or both, if the partition logic is fit at the top level, and to ensure that final timing sign-off is accurate. Use the following guidelines in the partition-specific .sdc to simplify these export and integration steps:

- Create a hierarchy variable for the partition (such as module_A_hierarchy) and set it to an empty string because the partition is the top-level instance in the separate Quartus II project. The project lead modifies this variable for the top-level hierarchy, reducing the effort of translating constraints on lower-level design hierarchies into constraints that apply in the top-level hierarchy. Use the following Tcl command first to check if the variable is already defined in the project, so that the top-level design does not use this empty hierarchy path: if {![info exists module_A_hierarchy]}

- Use the hierarchy variable in the partition-specific .sdc as a prefix for assignments in the project. For example, instead of naming a particular instance of a register reg:inst, use ${module_A_hierarchy}reg:inst. Also, use the hierarchy variable as a prefix to any wildcard characters (such as “*”).

- Pay attention to the location of the assignments to I/O ports of the partition. In most cases, these assignments should be specified in the .sdc with project-wide constraints, because the partition interface depends on the top-level design. If you want to set I/O constraints within the partition, the team must ensure that the I/O port names are identical in all projects so that the assignments can be integrated successfully without changes.

- Use caution with the derive_clocks and derive_pll_clocks commands. In most cases, the .sdc with project-wide constraints should call these commands. Because these commands impact the entire design, integrating them unexpectedly into the top-level design might cause problems.

If the design team follows these recommendations, the project lead should be able to include the .sdc with the partition-specific constraints provided by the partition designer directly in the top-level design.

Example Step 2: Partition Designer Creates .sdc with Partition-Specific Constraints

The partition designer compiles the design with the .sdc with project-wide constraints and might want to add some additional constraints. In this example, the designer realizes that he or she must specify a false path between the register called reg_in_1 and all destinations in this design block with the wildcard character (such as “*”). This constraint applies entirely within the partition and must be exported to the top-level design, so it qualifies for inclusion in the .sdc with partition-specific constraints. The designer first defines the module_A_hierarchy variable and uses it when writing the constraint as follows:

if {![info exists module_A_hierarchy]}
   {
      set module_A_hierarchy ""
   }

set_false_path -from [get_registers ${module_A_hierarchy}reg_in_1] -to [get_registers ${module_A_hierarchy}]*
Consolidating the .sdc in the Top-Level Design

When the partition designers complete their designs, they export the results to the project lead. The project lead receives the exported .qxp files and a copy of the .sdc with partition-specific constraints.

To set up the top-level .sdc constraint file to accept the .sdc files from the separate Quartus II projects, the top-level .sdc should define the hierarchy variables specified in the partition .sdc files. List the variable for each partition and set it to the hierarchy path, up to and including the instantiation of the partition in the top-level design, including the final hierarchy character “|”.

To ensure that the .sdc files are used in the correct order, the project lead can use the Tcl Source command to load each .sdc.

Example Step 3: Project Lead Performs Final Timing Analysis and Sign-off

With these commands, the top-level .sdc file looks like the following example:

```
create_clock -name {clk} -period 3.000 -waveform \( 0.000 1.500 \)
[get_ports {TOP_LEVEL_CLK}]
# Include the lower-level SDC file
set module_A_hierarchy "module_A:inst|" # Note the final '|' character
source <partition-specific constraint file such as
..\module_A\module_A_constraints>.sdc
```

When the project lead performs top-level timing analysis, the false path assignment from the lower-level module_A project expands to the following:

```
set_false_path -from module_A:inst|reg_in_1 -to module_A:inst|*
```

Adding the hierarchy path as a prefix to the SDC command makes the constraint legal in the top-level design, and ensures that the wildcard does not affect any nodes outside the partition that it was intended to target.

By following the guidelines in this section, constraint propagation between the separate Quartus II projects can be managed effectively.

Introduction to Design Floorplans

A floorplan represents the layout of the physical resources on the device. Creating a design floorplan, or floorplanning, describes the process of mapping the logical design hierarchy onto physical regions in the device.

In the Quartus II software, LogicLock regions can be used to constrain blocks of a design to a particular region of the device. LogicLock regions represent an area on the device with a user-defined or Fitter-defined size and location in the device layout.

For more information about design floorplans and LogicLock regions, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

The Difference between Logical Partitions and Physical Regions

Design partitions are logical entities based on the design hierarchy. LogicLock regions are physical placement assignments that constrain logic to a particular region on the device.
A common misconception is that logic from a design partition is always grouped together on the device when you use incremental compilation. Actually, logic from a partition can be placed anywhere in the device if it is not constrained to a LogicLock region, although the Fitter can pack related logic together to improve timing performance. A logical design partition does not refer to any physical area on the device and does not directly control where instances are placed on the device.

If you want to control the placement of logic from a design partition and isolate it to a particular part of the device, you can assign the logical design partition to a physical region in the device floorplan with a LogicLock region assignment. Altera recommends creating a design floorplan by assigning design partitions to LogicLock regions to improve the quality of results and avoid placement conflicts in some situations for incremental compilation. For more information, refer to “Why Create a Floorplan?” on page 16–41.

Another misconception is that LogicLock assignments are used to preserve placement results for incremental compilation. Actually, LogicLock regions only constrain logic to a physical region on the device. Incremental compilation does not use LogicLock assignments or any location assignments to preserve the placement results; it simply reuses the results stored in the database netlist from a previous compilation.

**Why Create a Floorplan?**

Creating a design floorplan is usually required if you want to preserve placement for partitions that will be exported, to avoid resource conflicts between partitions in the top-level design. Floorplan location planning can be important for a design that uses incremental compilation, for the following reasons:

- To avoid resource conflicts between partitions, predominantly when integrating partitions exported from another Quartus II project.
- To ensure good quality of results when recompiling individual timing-critical partitions.

Location assignments for each partition ensure that there are no placement conflicts between partitions. If there are no LogicLock region assignments, or if LogicLock regions are set to auto-size or floating location, no device resources are specifically allocated for the logic associated with the region. If you do not clearly define resource allocation, logic placement can conflict when you integrate the partitions in the top-level design if you reuse the placement information from the exported netlist.

Creating a floorplan is also recommended for timing-critical partitions that have little timing margin to maintain good quality of results when the design changes.

Floorplan assignments are not required for non-critical partitions compiled in the same Quartus II project. The logic for partitions that are not timing-critical can be placed anywhere in the device on each recompilation if that is best for your design.

Design floorplan assignments prevent the situation in which the Fitter must place a partition in an area of the device where most resources are used by other partitions. A LogicLock region provides a reasonable region to re-place logic after a change, so the Fitter does not have to scatter logic throughout the available space in the device.
Figure 16–24 illustrates the problems that may be associated with refitting designs that do not have floorplan location assignments. The left floorplan shows the initial placement of a four-partition design (P1-P4) without any floorplan location assignments. The right floorplan shows the device if a change occurs to P3. After removing the logic for the changed partition, the Fitter must re-place and reroute the new logic for P3 in the scattered white space shown in Figure 16–24. The placement of the post-fit netlists for other partitions forces the Fitter to implement P3 with the device resources that have not been used.

The Fitter has a more difficult task because of more difficult physical constraints, and as a result, compilation time often increases. The Fitter might not be able to find any legal placement for the logic in partition P3, even if it could in the initial compilation. Additionally, if the Fitter can find a legal placement, the quality of results often decreases in these cases, sometimes dramatically, because the new partition is now scattered throughout the device.

Figure 16–25 shows the initial placement of a four-partition design with floorplan location assignments. Each partition is assigned to a LogicLock region. The second part of the figure shows the device after partition P3 is removed. This placement presents a much more reasonable task to the Fitter and yields better results.

Altera recommends that you create a LogicLock floorplan assignment for timing-critical blocks with little timing margin that will be recompiled as you make changes to the design.
When to Create a Floorplan

It is important that you plan early to incorporate partitions into the design, and ensure that each partition follows partitioning guidelines. You can create floorplan assignments at different stages of the design flow, early or late in the flow. These guidelines help ensure better results as you begin creating floorplan location assignments.

Early Floorplan

An early floorplan is created before the design stage. You can plan an early floorplan at the top level of a design to allocate each partition a portion of the device resources. Doing so allows the designer for each block to create the logic for their design partition without conflicting with other logic. Each partition can be optimized in a separate Quartus II project if required, and the design can still be easily integrated in the top-level design. Even within one Quartus II project, each partition can be locked down with a post-fit netlist, and you can be sure there is space in the device floorplan for other partitions.

When you have compiled your complete design, or after you have integrated the first versions of partitions developed in separate Quartus II projects, you can use the design information and Quartus II features to tune and improve the floorplan, as described in the following section.

Late Floorplan

A late floorplan is created or modified after the design is created, when the code is close to complete and the design structure is likely to remain stable. Creating a late floorplan is typically necessary only if you are starting to use incremental compilation late in the design flow, or need to reserve space for a logic block that becomes timing-critical but still has HDL changes to be integrated. When the design is complete, you can take advantage of the Quartus II analysis features to check the floorplan quality. To adjust the floorplan, you can perform iterative compilations as required and assess the results of different assignments.

It may not be possible to create a good-quality late floorplan if you do not create partitions in the early stages of the design.

Design Floorplan Placement Guidelines

The following guidelines are key to creating a good design floorplan:

- Capture correct resources in each region.
- Use good region placement to maintain design performance compared to flat compilation.

A common misconception is that creating a floorplan enhances timing performance, as compared to a flat compilation with no location assignments. The Fitter does not usually require guidance to get optimal results for a full design.
Floorplan assignments can help maintain good performance when designs change incrementally, as described in “Why Create a Floorplan?” on page 16–41. However, poor placement assignments in an incremental compilation can often adversely affect performance results, as compared to a flat compilation, because the assignments limit the options for the Fitter. Investing time to find good region placement is required to match the performance of a full flat compilation.

Use the following general procedure to create a floorplan:

1. Divide the design into partitions.
2. Assign the partitions to LogicLock regions.
3. Compile the design.
4. Analyze the results.
5. Modify the placement and size of regions, as required.

You might have to perform these steps several times to find the best combination of design partitions and LogicLock regions that meet the resource and timing goals of the design.

For more information about performing these steps, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Assigning Partitions to LogicLock Regions

Before compiling a design with new LogicLock assignments, ensure that the partition netlist type is set to Post-Synthesis or Source File, so that the Fitter does not reuse previous placement results.

In most cases, you should include logic from one partition in each LogicLock region. This organization helps to prevent resource conflicts when partitions are exported and can lead to better performance preservation when locking down parts of a design in a single project.

The Quartus II software is flexible and allows exceptions to this rule. For example, you can place more than one partition in the same LogicLock region if the partitions are tightly connected, but you do not want to merge the partitions into one larger partition. For best results, ensure that you recompile all partitions in the LogicLock region every time the logic in one partition changes. Additionally, if a partition contains multiple lower-level entities, you can place those entities in different areas of the device with multiple LogicLock regions, even if they are defined in the same partition.

You can use the Reserved LogicLock option to ensure that you avoid conflicts with other logic that is not locked into a LogicLock region. This option prevents other logic from being placed in the region, and is useful if you have empty partitions at any point during your design flow, so that you can reserve space in the floorplan. Do not make reserved regions too large to prevent unused area because no other logic can be placed in a region with the Reserved LogicLock option.

For more information about LogicLock region properties, refer to the LogicLock Region Properties Dialog Box in Quartus II Help.
How to Size and Place Regions

In an early floorplan, assign physical locations based on design specifications. Use information about the connections between partitions, the partition size, and the type of device resources required.

In a late floorplan, when the design is complete, you can use locations or regions chosen by the Fitter as a guideline. If you have compiled the full design, you can view the location of the partition logic in the Chip Planner. Refer to “Checking Partition Quality” on page 16–30 for information about viewing placement results for each partition in the device floorplan. You can use the natural grouping of each unconstrained partition as a starting point for a LogicLock region constraint. View the placement for each partition that requires a floorplan constraint, and create a new LogicLock region by drawing a box around the area on the floorplan, and then assigning the partition to the region to constrain the partition placement.

For a step-by-step procedure to create a LogicLock region, refer to Creating and Manipulating LogicLock Regions in Quartus II Help.

Instead of creating regions based on the previous compilation results, you can start with the Fitter results for a default auto size and floating origin location for each new region when the design logic is complete. After compilation, lock the size and origin location. Instead of a full compilation, you can use the Start Early Timing Estimate command to perform a fast placement.

Alternatively, if the design logic is complete with auto-sized or floating location regions, you can specify the size based on the synthesis results and use the locations chosen by the Fitter with the Set to Estimated Size command. Like the previous option, start with floating origin location. After compilation, lock the origin location. Again, instead of a full compilation, you can use the Start Early Timing Estimate command to perform a fast placement. You can also enable the Fast Synthesis Effort setting to reduce synthesis time.

After a compilation or early timing estimate, save the Fitter size and origin location of the Fitter with the Set Size and Origin to Previous Fitter Results command.

It is important that you use the Fitter-chosen locations only as a starting point to give the regions a good fixed size and location. Ensure that all LogicLock regions in the design have a fixed size and have their origin locked to a specific location on the device. On average, regions with fixed size and location yield better timing performance than auto-sized regions.

Modifying Region Size and Origin

After saving the Fitter results from an initial compilation for a late floorplan, modify the regions using your knowledge of the design to set a specific size and location. If you have a good understanding of how the design fits together, you can often improve upon the regions placed in the initial compilation. In an early floorplan, when the design has not yet been created, you can use the guidelines in this section to set the size and origin, even though there is no initial Fitter placement.
The easiest way to move and resize regions is to drag the region location and borders
in the Chip Planner. Make sure that you select the **User-Defined** region in the
floorplan (as opposed to the **Fitter-Placed** region from the last compilation) so that
you can change the region.

Generally, you can keep the Fitter-determined relative placement of the regions, but
make adjustments if required to meet timing performance. If you find that the early
timing estimate did not result in good relative placements, try performing a full
compilation so that the Fitter can optimize for a full placement and routing.

If two LogicLock regions have several connections between them, ensure they are
placed near each other to improve timing performance. By placing connected regions
near each other, the Fitter has more opportunity to optimize inter-region paths when
both partitions are recompiled. Reducing the criticality of inter-region paths also
allows the Fitter more flexibility when placing other logic in each region.

If resource utilization is low in the overall device, enlarge the regions. Doing so
usually improves the final results because it gives the Fitter more freedom to place
additional or modified logic added to the partition during subsequent incremental
compilations. It also allows room for optimizations such as pipelining and physical
synthesis logic duplication.

Try to have each region evenly full, with the same “fullness” that the complete design
would have without LogicLock regions; Altera recommends approximately 75% full.

Allow more area for regions that are densely populated, because overly congested
regions can lead to poor results. Allow more empty space for timing-critical partitions
to improve results. However, do not make regions too large for their logic. Regions
that are too large can result in wasted resources and also lead to suboptimal results.

Ideally, almost the entire device should be covered by LogicLock regions if all
partitions are assigned to regions.

Regions should not overlap in the device floorplan. If two partitions are allocated on
an overlapping portion of the chip, each may independently claim common resources
in this region. This leads to resource conflicts when integrating results into a top-level
design. In a single project, overlapping regions give more difficult constraints to the
Fitter and can lead to reduced quality of results.

You can create hierarchical LogicLock regions to ensure that the logic in a child
partition is physically placed inside the LogicLock region for its parent partition. This
can be useful when the parent partition does not contain registers at the boundary
with the lower-level child partition and has a lot of signal connectivity. To create a
hierarchical relationship between regions in the LogicLock Regions window, drag and
drop the child region to the parent region.

**I/O Connections**

Consider I/O timing when placing regions. Using I/O registers can minimize I/O
timing problems, and using boundary registers on partitions can minimize problems
connecting regions or partitions. However, I/O timing might still be a concern. It is
most important for flows where each partition is compiled independently, because the
Fitter can optimize the placement for paths between partitions if the partitions are
compiled at the same time.
Place regions close to the appropriate I/O, if necessary. For example, DDR memory interfaces have very strict placement rules to meet timing requirements. Incorporate any specific placement requirements into your floorplan as required. You should create LogicLock regions for internal logic only, and provide pin location assignments for external device I/O pins (instead of including the I/O cells in a LogicLock region to control placement).

**LogicLock Resource Exclusions**

You can exclude certain resource types from a LogicLock region to manage the ratio of logic to dedicated DSP and RAM resources in the region.

If your design contains memory or Digital Signal Processing (DSP) elements, you may want to exclude these elements from the LogicLock region. LogicLock resource exceptions prevent certain types of elements from being assigned to a region. Therefore, those elements are not required to be placed inside the region boundaries. The option does not prevent them from being placed inside the region boundaries unless the **Reserved** property of the region is turned on.

Resource exceptions are useful in cases where it is difficult to place rectangular regions for design blocks that contain memory and DSP elements, due to their placement in columns throughout the device floorplan. Exclude RAMs, DSPs, or logic cells to give the Fitter more flexibility with region sizing and placement. Excluding RAM or DSP elements can help to resolve no-fit errors that are caused by regions spanning too many resources, especially for designs that are memory-intensive, DSP-intensive, or both. Figure 16–26 shows an example of a design with an odd-shaped region to accommodate DSP blocks for a region that does not contain very much logic. The right side of the figure shows the result after excluding DSP blocks from the region. The region can be placed more easily without wasting logic resources.

**Figure 16–26. LogicLock Resource Exclusion Example**

![Diagram showing LogicLock Resource Exclusion Example](image-url)
To view any resource exceptions, right-click in the LogicLock Regions window, and then click LogicLock Regions Properties. In the LogicLock Regions Properties dialog box, select the design element (module or entity) in the Members box, and then click Edit. In the Edit Node dialog box, to set up a resource exception, click the Edit button next to the Excluded element types box, and then turn on the design element types to be excluded from the region. You can choose to exclude combinational logic or registers from logic cells, or any of the sizes of TriMatrix memory blocks, or DSP blocks.

If the excluded logic is in its own lower-level design entity (even if it is within the same design partition), you can assign the entity to a separate LogicLock region to constrain its placement in the device.

You can also use this feature with the LogicLock Reserved property to reserve specific resources for logic that will be added to the design.

Creating Floorplan Location Assignments With Tcl Commands—Excluding or Filtering Certain Device Elements (Such as RAM or DSP Blocks)

To assign a code block to a LogicLock region, with exclusions, use the following command:

```
set_logiclock_contents -region <LogicLock region name> -to <block> -exceptions "<keyword>:<keyword>"
```

- `<LogicLock region name>`—The name of the LogicLock region to which the code block is assigned.
- `<block>`—A code block in a Quartus II project hierarchy, which can also be a design partition.
- `<keyword>`—The list of exceptions made during assignment. For example, if DSP was in the keyword list, the named block of code would be assigned to the LogicLock region, except for any DSP block within the code block. You can include the following exceptions in the `set_logiclock_contents` command:

Keywords variables:
- `REGISTER`—Any registers in the logic cells.
- `COMBINATIONAL`—Any combinational elements in the logic cells.
- `SMALL_MEM`—Small TriMatrix memory blocks (M512 or MLAB).
- `MEDIUM_MEM`—Medium TriMatrix memory blocks (M4K or M9K).
- `LARGE_MEM`—Large TriMatrix memory blocks (M-RAM or M144K).
- `DSP`—Any DSP blocks.
- `VIRTUAL_PIN`—Any virtual pins.

Resource filtering uses the optional Tcl argument `-exclude_resources` in the `set_logiclock_contents` function. If left unspecified, no resource filter is created. In the `.qsf`, resource filtering uses an extra LogicLock membership assignment called LL_MEMBER_RESOURCE_EXCLUDE. For example, the following line in the `.qsf` is used to specify a resource filter for the alu:alu_unit entity assigned to the ALU region.

```
set_instance_assignment -name LL_MEMBER_RESOURCE_EXCLUDE "DSP:SMALL_MEM" -to "alu:alu_unit" -section_id ALU
```
Creating Non-Rectangular Regions

To constrain placement to non-rectangular or non-contiguous areas of the device, you can connect multiple rectangular regions together using the Merge command.

For devices that do not support the Merge command (Arria™GX, Cyclone, Cyclone II, HardCopy II, MAX™ II, Stratix, Stratix II, Stratix II GX, and Stratix GX devices), you can limit entity placement to a sub-area of a LogicLock region to create non-rectangular constraints. In these devices, construct a LogicLock hierarchy by creating child regions inside of parent regions, and then use the Reserved option to control which logic can be placed inside these child regions. Setting the Reserved option for the region prevents the Fitter from placing nodes that are not assigned to the region inside the boundary of the region.

For more information and examples of creating non-rectangular regions, refer to Creating and Manipulating LogicLock Regions in Quartus II Help.

Checking Floorplan Quality

This section provides an overview of tools that you can use as you create a floorplan in the Quartus II software. You can use these tools to assess your floorplan quality and use the information to improve your design or assignments as required to achieve the best results.

Incremental Compilation Advisor

You can use the Incremental Compilation Advisor to check that your design follows the recommendations for creating floorplan location assignments that are presented in this chapter. For more information, refer to “Incremental Compilation Advisor” on page 16–30.

LogicLock Region Resource Estimates

You can view resource estimates for a LogicLock region to determine the region’s resource coverage, and use this estimate before compilation to check region size. Using this estimate helps to ensure adequate resources when you are sizing or moving regions.

For information about how to view the properties of a LogicLock region, refer to LogicLock Region Properties Dialog Box in Quartus II Help.

LogicLock Region Properties Statistics Report

LogicLock region statistics are similar to design partition properties, but also include resource usage details after compilation.

The statistics report the number of resources used and the total resources covered by the region, and also list the number of I/O connections and how many I/Os are registered (good), as well as the number of internal connections and the number of inter-region connections (bad).
Locate the Quartus II TimeQuest Timing Analyzer Path in the Chip Planner

In the TimeQuest analyzer user interface, you can locate a specific path in the Chip Planner to view its placement and perform a report timing operation (for example, report timing for all paths with less than 0 ns slack).

Inter-Region Connection Bundles

The Chip Planner can display bundles of connections between LogicLock regions, with filtering options that allow you to choose the relevant data for display. These bundles can help you to visualize how many connections there are between each LogicLock region to improve floorplan assignments or to change partition assignments, if required.

Routing Utilization

The Chip Planner includes a feature to display a color map of routing congestion. This display helps identify areas of the chip that are too tightly packed.

In the Chip Planner, red LAB blocks indicate higher routing congestion. You can position the mouse pointer over a LAB to display a tooltip that reports the logic and routing utilization information.

Ensure Floorplan Assignments Do Not Significantly Impact Quality of Results

The end results of design partitioning and floorplan creation differ from design to design. However, it is important to evaluate your results to ensure that your scheme is successful. Compare your before and after results, and consider using another scheme if any of the following guidelines are not met:

- You should see only minor degradation in $f_{\text{MAX}}$ after the design is partitioned and floorplan location assignments are created. There is some performance cost associated with setting up a design for incremental compilation; approximately 3% is typical.
- The area increase should be no more than 5% after the design is partitioned and floorplan location assignments are created.
- The time spent in the routing stage should not significantly increase.
The amount of compilation time spent in the routing stage is reported in the Messages window with an Info message that indicates the elapsed time for Fitter routing operations. If you notice a dramatic increase in routing time, the floorplan location assignments may be creating substantial routing congestion. In this case, decrease the number of LogicLock regions, which typically reduces the compilation time in subsequent incremental compilations and may also improve design performance.

**Recommended Design Flows and Application Examples**

This section provides design flows for partitioning and creating a design floorplan during common timing closure and team-based design scenarios. Each flow describes the situation in which it should be used, and provides a step-by-step description of the commands required to implement the flow.

**Create a Floorplan for Major Design Blocks**

Use this incremental compilation flow for designs when you want to assign a floorplan location for each major block in your design. A full floorplan ensures that partitions do not interact as they are changed and recompiled—each partition has its own area of the device floorplan.

To create a floorplan for major design blocks, follow this general methodology:

1. In the Design Partitions window, ensure that all partitions have their netlist type set to **Source File** or **Post-Synthesis**. If the netlist type is set to **Post-Fit**, floorplan location assignments are not used when recompiling the design.
2. Create a LogicLock region for each partition (including the top-level entity, which is set as a partition by default).
3. Run a full compilation of your design to view the initial Fitter-chosen placement of the LogicLock regions as a guideline.
4. In the Chip Planner, view the placement results of each partition and LogicLock region on the device.
5. If required, modify the size and location of the LogicLock regions in the Chip Planner. For example, enlarge the regions to fill up the device and allow for future logic changes.

You can also, if needed, create a new LogicLock region by drawing a box around an area on the floorplan.

6. Run an early timing estimate with the **Start Early Timing Estimate** command to estimate the timing performance of your design with the modified or new LogicLock regions.
7. Repeat steps 5 and 6 until you are satisfied with the quality of results for your design floorplan. Once you are satisfied with your results, run a full compilation of your design.
Create a Floorplan Assignment for One Design Block with Difficult Timing

Use this flow when you have one timing-critical design block that requires more optimization than the rest of your design. You can take advantage of incremental compilation to reduce your compilation time without creating a full design floorplan.

In this scenario, you do not want to create floorplan assignments for the entire design. Instead, you can create a region to constrain the location of your critical design block, and allow the rest of the logic to be placed anywhere on the device. To create a region for critical design block, follow these steps:

1. Divide up your design into partitions. Consider the guidelines in “Design Partition Guidelines” on page 16–10 to determine partition boundaries. Ensure that you isolate the timing-critical logic in a separate partition.

2. Define a LogicLock region for the timing-critical partition. Ensure that you capture the correct amount of device resources in the region. Turn on the Reserved property to prevent any other logic from being placed in the region.
   - If the design block is not complete, reserve space in the design floorplan based on your knowledge of the design specifications, connectivity between design blocks, and estimates of the size of the partition based on any initial implementation numbers.
   - If the critical design block has initial source code ready, compile the design to place the LogicLock region. Save the Fitter-determined size and origin, and then enlarge the region to provide more flexibility and allow for future design changes.

As the rest of the design is completed, and the device fills up, the timing-critical region reserves an area of the floorplan. When you make changes to the design block, the logic will be re-placed in the same part of the device, which helps ensure good quality of results.

Create a Floorplan as the Project Lead in a Team-Based Flow

Use this approach when you have several designs that will be implemented in separate Quartus II projects by different designers, or third-party IP designers who want to optimize their designs independently and pass the results to the project lead.

As the project lead in this scenario, follow these steps to prepare the top-level design for a successful team-based design methodology with early floorplan planning:

1. Create a new Quartus II project that will ultimately contain the full implementation of the entire design.

2. Create a “skeleton” or framework of the design that defines the hierarchy for the subdesigns that will be implemented by separate designers. Consider the partitioning guidelines in this chapter when determining the design hierarchy.

3. Make project-wide settings. Select the device, make global assignments for clocks and device I/O ports, and make any global signal constraints to specify which signals can use global routing resources.

4. Make design partition assignments for each major subdesign. Set the netlist type for each partition that will be implemented in a separate Quartus II project and later exported and integrated with the top-level design set to Empty.
5. Create LogicLock regions for each partition to create a design floorplan. This floorplan should consider the connectivity between partitions and estimates of the size of each partition based on any initial implementation numbers and knowledge of the design specifications. Use the guidelines described in this chapter to choose a size and location for each LogicLock region.

6. Provide the constraints from the top-level design to partition designers using one of the following procedures:

   a. Create a copy of the top-level Quartus II project framework by checking out the appropriate files from a source control system, using the Copy Project command, or creating a project archive. Provide each partition designer with the copy of the project.

   b. Provide the constraints with documentation or scripts.

   To use design partition scripts to pass constraints and generate separate Quartus II projects, refer to *Generating Design Partition Scripts for Project Management* in Quartus II Help.

**Conclusion**

Incremental compilation can significantly improve your design productivity, especially for large, complex designs. To take advantage of the feature, it is worth spending time to create quality partition and floorplan assignments. Follow the guidelines to set up your design hierarchy and source code for incremental compilation.

Floorplan location assignments are required when design blocks are developed independently and are recommended for timing-critical partitions that are expected to change. Follow the guidelines to create and modify LogicLock regions to create good placement assignments for your design partitions.

Remember that you do not have to follow all the guidelines exactly to implement an incremental compilation design flow, but following the guidelines can maximize your chances of success.

**Document Revision History**

Table 16–1 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
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<tr>
<td>November 2011</td>
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<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>- Changed to new document template.</td>
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<td></td>
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<td>- Moved 'Creating Floorplan Location Assignments With Tcl Commands—Excluding or Filtering Certain Device Elements (Such as RAM or DSP Blocks)' from the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.</td>
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<td>- Consolidated Design Partition Planner and Incremental Compilation Advisor information between the Quartus II Incremental Compilation for Hierarchical and Team-Based Design and Best Practices for Incremental Compilation Partitions and Floorplan Assignments handbook chapters.</td>
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<td>- Removed the explanation of the “bottom-up design flow” where designers work completely independently, and replaced with Altera’s recommendations for team-based environments where partitions are developed in the same top-level project framework, plus an explanation of the bottom-up process for including independent partitions from third-party IP designers.</td>
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<td>- Expanded the Merge command explanation to explain how it now accommodates cross-partition boundary optimizations.</td>
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<td>- Added 'Including SDC Constraints from Lower-Level Partitions for Third-Party IP Delivery’ from the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.</td>
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<tr>
<td></td>
<td></td>
<td>- Reorganized the “Recommended Design Flows and Application Examples” section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Removed HardCopy APEX and HardCopy Stratix Devices section.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>- Added I/O register packing examples from Incremental Compilation for Hierarchical and Team-Based Designs chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Moved &quot;Incremental Compilation Advisor&quot; section</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added &quot;Viewing Design Partition Planner and Floorplan Side-by-Side&quot; section</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated Figure 15-22</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Chapter 8 was previously Chapter 7 in software release 8.1.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>- Changed to 8-1/2 x 11 page size. No change to content.</td>
</tr>
<tr>
<td>May 2007</td>
<td>8.0.0</td>
<td>- Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the [Quartus II Handbook Archive](#).
As programmable logic devices become more complex and require increased performance, advanced design synthesis has become an important part of the design flow. In the Quartus® II software you can use the integrated Analysis and Synthesis module of the Compiler to synthesize your design files and create the project database for future stages of the compilation flow. You can also use other EDA synthesis tools to first synthesize your designs, and then generate EDIF netlist files or Verilog Quartus Mapping Files (.vqm) that you can use with the Quartus II software. The Quartus II netlist viewers allow you to visually analyze the design netlist at different stages of synthesis and compilation. This section explains the options that are available for each of these flows and how they are supported in the Quartus II software.

This section includes the following chapters:

- **Chapter 17, Quartus II Integrated Synthesis**
  This chapter documents the integrated synthesis design flow and language support in the Quartus II software. It explains how you can improve synthesis results with Quartus II synthesis options and optimization techniques, and how you can control the inference of architecture-specific megafunctions. This chapter also explains some of the node-naming conventions used during synthesis to help you better understand your synthesized design and the messages issued during synthesis to improve your HDL code. Scripting techniques for applying all the options and settings described are also provided.

- **Chapter 18, Synopsys Synplify Support**
  This chapter documents support for the Synopsys Synplify software in the Quartus II software, as well as key design flows, methodologies, and techniques for achieving good results in Altera® devices with the Synplify software.

- **Chapter 19, Mentor Graphics Precision Synthesis Support**
  This chapter documents support for the Mentor Graphics® Precision Synthesis software in the Quartus II software, as well as key design flows, methodologies, and techniques for achieving good results in Altera® devices with the Precision Synthesis software.

- **Chapter 20, Mentor Graphics LeonardoSpectrum Support**
  This chapter documents key design methodologies and techniques for Altera devices using the Mentor Graphics LeonardoSpectrum™ software and Quartus II design flow. The LeonardoSpectrum software is a mature synthesis tool supporting legacy devices and many current devices. Altera recommends using the advanced Precision Synthesis software for new designs in new device families.
Chapter 21, Analyzing Designs with Quartus II Netlist Viewers

This chapter shows how to use the Quartus II netlist viewers to analyze your design at various stages of the design cycle. It also provides an introduction to the Quartus II design flow using netlist viewers, an overview of each viewer, and an explanation of the user interface.
This chapter describes the Integrated Synthesis design flow and provides scripting techniques for applying all the options and settings described in this chapter.

As programmable logic designs become more complex and require increased performance, advanced synthesis becomes an important part of a design flow. The Altera® Quartus® II software includes advanced Integrated Synthesis that fully supports VHDL, Verilog HDL, and Altera-specific design entry languages, and provides options to control the synthesis process. With this synthesis support, the Quartus II software provides a complete, easy-to-use solution.

This chapter contains the following sections:

- “Design Flow” on page 17–1
- “Language Support” on page 17–4
- “Incremental Compilation” on page 17–21
- “Quartus II Synthesis Options” on page 17–23
- “Analyzing Synthesis Results” on page 17–73
- “Analyzing and Controlling Synthesis Messages” on page 17–74
- “Node-Naming Conventions in Quartus II Integrated Synthesis” on page 17–78
- “Scripting Support” on page 17–84

For examples of Verilog HDL and VHDL code synthesized for specific logic functions, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. For more information about coding with primitives that describe specific low-level functions in Altera devices, refer to the Designing With Low-Level Primitives User Guide.

Design Flow

The Quartus II Analysis & Synthesis stage of the compilation flow runs Integrated Synthesis, which fully supports Verilog HDL, VHDL, and Altera-specific languages, and major features of the SystemVerilog language. For more information, refer to “Language Support” on page 17–4.

In the synthesis stage of the compilation flow, the Quartus II software performs logic synthesis to optimize design logic and performs technology mapping to implement the design logic in device resources such as logic elements (LEs) or adaptive logic modules (ALMs), and other dedicated logic blocks. The synthesis stage generates a single project database that integrates all your design files in a project (including any netlists from third-party synthesis tools).
You can use Analysis & Synthesis to perform the following compilation processes:

- **Analyze Current File**—parses your current design source file to check for syntax errors. This command does not report many semantic errors that require further design synthesis. To perform this analysis, on the Processing menu, click **Analyze Current File**.

- **Analysis & Elaboration**—checks your design for syntax and semantic errors and performs elaboration to identify your design hierarchy. To perform Analysis & Elaboration, on the Processing menu, point to **Start**, and then click **Start Analysis & Elaboration**.

- **Hierarchy Elaboration**—parses HDL designs and generates a skeleton of hierarchies. Hierarchy Elaboration is similar to the Analysis & Elaboration flow, but without any elaborated logic, thus making it much faster to generate.

  For more information about the Hierarchy Elaboration flow, refer to **Start Hierarchy Elaboration Command (Processing Menu)** in Quartus II Help.

- **Analysis & Synthesis**—performs complete Analysis & Synthesis on a design, including technology mapping. To perform Analysis & Synthesis, on the Processing menu, point to **Start**, and then click **Start Analysis & Synthesis**.

The Quartus II Integrated Synthesis design and compilation flow consists of the following steps:

1. Create a project in the Quartus II software and specify the general project information, including the top-level design entity name.

2. Create design files in the Quartus II software or with a text editor.

3. On the Project menu, click **Add/Remove Files in Project** and add all design files to your Quartus II project using the Files page of the Settings dialog box.

4. Specify Compiler settings that control the compilation and optimization of your design during synthesis and fitting. For synthesis settings, refer to “Quartus II Synthesis Options” on page 17–23.

5. Add timing constraints to specify the timing requirements.

   To partition your design to reduce compilation time, refer to “Incremental Compilation” on page 17–21.

6. Compile your design. To synthesize your design, on the Processing menu, point to **Start**, and then click **Start Analysis & Synthesis**. To run a complete compilation flow including placement, routing, creation of a programming file, and timing analysis, click **Start Compilation** on the Processing menu.

7. After obtaining synthesis and placement and routing results that meet your requirements, program or configure your Altera device.

Integrated Synthesis generates netlists that enable you to perform functional simulation or gate-level timing simulation, timing analysis, and formal verification.
Figure 17–1 shows the basic design flow using Quartus II Integrated Synthesis.

**Figure 17–1. Quartus II Design Flow Using Quartus II Integrated Synthesis**

**Notes to Figure 17–1:**
1. AHDL stands for the Altera Hardware Description Language.
2. BDF stands for the Altera schematic Block Design File (.bdf).
3. The Quartus II Exported Partition File (.qxp) is a precompiled netlist that you can use as a design source file. For more information about using .qxp as a design source file, refer to "Quartus II Exported Partition File as Source" on page 17–22.

For an overall summary of features in the Quartus II software, refer to the *Introduction to the Quartus II Software* manual.

For more information about Quartus II projects and the compilation flow, refer to *Managing Files in a Project* and *About Compilation Flows* in Quartus II Help.
Language Support

This section describes Quartus II Integrated Synthesis support for HDL, schematic design entry, graphical state machine entry, and how to specify the Verilog HDL or VHDL language version in your design. This section also describes language features such as Verilog HDL macros, initial constructs and memory system tasks, and VHDL libraries. “Design Libraries” on page 17–12 describes how to compile and reference design units in custom libraries, and “Using Parameters/Generics” on page 17–16 describes how to use parameters or generics and pass them between languages.

To ensure that the Quartus II software reads all associated project files, add each file to your Quartus II project by clicking Add/Remove Files in Project on the Project menu. You can add design files to your project. You can mix all supported languages and netlists generated by third-party synthesis tools in a single Quartus II project.

You can also use the available templates in the Quartus II Text Editor for various Verilog and VHDL features. For more information, refer to Insert Template Dialog Box in Quartus II Help.

Verilog HDL Support

The Quartus II Compiler’s Analysis & Synthesis module supports the following Verilog HDL standards:

- Verilog-2001 (IEEE Standard 1364-2001)
- SystemVerilog-2005 (IEEE Standard 1800-2005) (the Compiler does not support all constructs)

The Verilog HDL code samples provided in this document follow the Verilog-2001 standard unless otherwise specified. The Quartus II Compiler uses the Verilog-2001 standard by default for files that have the extension .v, and the SystemVerilog standard for files that have the extension .sv.

If you use scripts to add design files, you can use the -HDL_VERSION command to specify the HDL version for each design file. For more information, refer to “Adding an HDL File to a Project and Setting the HDL Version” on page 17–85.

The Quartus II software support for Verilog HDL is case sensitive in accordance with the Verilog HDL standard. The Quartus II software supports the compiler directive `define, in accordance with the Verilog HDL standard.

The Quartus II software supports the include compiler directive to include files with absolute paths (with either “/” or “\” as the separator), or relative paths. When searching for a relative path, the Quartus II software initially searches relative to the project directory. If the Quartus II software cannot find the file, the software then searches relative to all user libraries and then relative to the directory location of the current file.

For more information about specifying synthesis directives, refer to “Synthesis Directives” on page 17–27.

For more information about Verilog HDL, refer to About Verilog HDL in Quartus II Help.
For more information about Quartus II Verilog HDL support, refer to *Quartus II Verilog HDL Support* in Quartus II Help.

For more information about specifying a default Verilog HDL version for all files, refer to *Specifying Verilog Input Settings* in Quartus II Help.

For more information about controlling the Verilog HDL version that compiles your design in a design file with the `VERILOG_INPUT_VERSION` synthesis directive, refer to *verilog_input_version Synthesis Directive* in Quartus II Help.

For more information about Verilog HDL synthesis attributes and directives, refer to *Verilog HDL Synthesis Attributes and Directives* in Quartus II Help.

**Verilog HDL Configuration**

Verilog HDL configuration is a set of rules that specify the source code for particular instances.

Verilog HDL configuration allows you to perform the following tasks:

- Specify a library search order for resolving cell instances (as does a library mapping file)
- Specify overrides to the logical library search order for specified instances
- Specify overrides to the logical library search order for all instances of specified cells

For more information about these tasks, refer to Table 17–1.

**Configuration Syntax**

A configuration contains the following statements:

**Example 17–1. Verilog HDL Configuration Statement**

```vhdl
config config_identifier;
design [library_identifier.]cell_identifier;
config_rule_statement;
endconfig
```

Where:

- `config`—the keyword that begins the configuration.
- `config_identifier`—the name you enter for the configuration.
- `design`—the keyword that starts a design statement for specifying the top of the design.
- `[library_identifier.]cell_identifier`—specifies the top-level module (or top-level modules) in the design and the logical library for this module (modules).
- `config_rule_statement`—one or more of the following clauses: default, instance, or cell. For more information, refer to Table 17–1.
- `endconfig`—the keyword that ends a configuration.
Table 17–1 lists the type of clauses for the `config_rule_statement` keyword:

<table>
<thead>
<tr>
<th>Clause Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>default</td>
<td>Specifies the logical libraries to search to resolve a default cell instance. A default cell instance is an instance in the design that is not specified in a subsequent instance or cell clause in the configuration. You specify these libraries with the <code>liblist</code> keyword. The following is an example of a default clause: <code>default liblist lib1 lib2;</code> Also specifies resolving default instances in the logical libraries (<code>lib1</code> and <code>lib2</code>). Because libraries are inherited, some simulators (for example, VCS) also search the default (or current) library as well after the searching the logical libraries (<code>lib1</code> and <code>lib2</code>).</td>
</tr>
<tr>
<td>instance</td>
<td>Specifies a specific instance. The specified instance clause depends on the use of the following keywords: <code>liblist</code>—specifies the logical libraries to search to resolve the instance. <code>use</code>—specifies that the instance is an instance of the specified cell in the specified logical library. The following are examples of <code>instance</code> clauses: <code>instance top.dev1 liblist lib1 lib2;</code> This instance clause specifies to resolve <code>instance top.dev1</code> with the cells assigned to logical libraries <code>lib1</code> and <code>lib2</code>; <code>instance top.dev1.gm1 use lib2.gizmult;</code> This instance clause specifies that <code>top.dev1.gm1</code> is an instance of the cell named <code>gizmult</code> in logical library <code>lib2</code>.</td>
</tr>
<tr>
<td>cell</td>
<td>A <code>cell</code> clause is similar to an <code>instance</code> clause, except that the <code>cell</code> clause specifies all instances of a cell definition instead of specifying a particular instance. What it specifies depends on the use of the <code>liblist</code> or <code>use</code> keywords: <code>liblist</code>—specifies the logical libraries to search to resolve all instances of the cell. <code>use</code>—the specified cell's definition is in the specified library.</td>
</tr>
</tbody>
</table>

### Hierarchical Configurations

A design can have more than one configuration. For example, you can define a configuration that specifies the source code you use in particular instances in a sub hierarchy, then define a configuration for a higher level of the design.

Suppose, for example, a sub hierarchy of a design is an eight-bit adder and the RTL Verilog code describes the adder in a logical library named `rtllib` and the gate-level code describes the adder in a logical library named `gatelib`. If you want to use the gate-level code for the 0 (zero) bit of the adder and the RTL level code for the other seven bits, the configuration might appear as shown in Example 17–2:

**Example 17–2.**

```vhdl
config cfg1;
design aLib.eight_adder;
default liblist rtllib;
instance adder.fulladd0 liblist gatelib;
endconfig
```
If you are instantiating this eight-bit adder eight times to create a 64-bit adder, use configuration \texttt{cfg1} for the first instance of the eight-bit adder, but not in any other instance. A configuration that would perform this function is shown in Example 17–3:

\begin{verbatim}
Example 17–3.

\texttt{config \texttt{cfg2};
design bLib.64_adder;
default liblist bLib;
instance top.64add0 use work.cfg1:config;
endconfig}
\end{verbatim}

The name of the unbound module may be different than the name of the cell that is bounded to the instance.

\textbf{Suffix :config}

To distinguish between a module by the same name, use the optional extension :config to refer to configuration names. For example, you can always refer to a \texttt{cfg2} configuration as \texttt{cfg2:config} (even if the \texttt{cfg2} module does not exist).

\section*{SystemVerilog Support}

The Quartus II software supports the SystemVerilog constructs.

\begin{itemize}
\item Designs written to support the Verilog-2001 standard might not compile with the SystemVerilog setting because the SystemVerilog standard has several new reserved keywords.
\item For more information about the supported SystemVerilog constructs and the supported Verilog-2001 features, refer to \textit{Quartus II Support for SystemVerilog} and \textit{Quartus II Support for Verilog 2001} in Quartus II Help.
\end{itemize}

\index{initial}

\section*{Initial Constructs and Memory System Tasks}

The Quartus II software infers power-up conditions from Verilog HDL \texttt{initial} constructs. The Quartus II software also creates power-up settings for variables, including RAM blocks. If the Quartus II software encounters nonsynthesizable constructs in an \texttt{initial} block, it generates an error. To avoid such errors, enclose nonsynthesizable constructs (such as those intended only for simulation) in \texttt{translate\_off} and \texttt{translate\_on} synthesis directives, as described in \textit{“Translate Off and On / Synthesis Off and On”} on page 17–64. Synthesis of initial constructs enables the power-up state of the synthesized design to match the power-up state of the original HDL code in simulation. For more information, refer to \textit{“Power-Up Level”} on page 17–40.

\begin{itemize}
\item Initial blocks do not infer power-up conditions in some third-party EDA synthesis tools. If you convert between synthesis tools, you must set your power-up conditions correctly.
\end{itemize}
Quartus II Integrated Synthesis supports the $readmemb and $readmemh system tasks to initialize memories. Example 17–4 shows an initial construct that initializes an inferred RAM with $readmemb.

Example 17–4. Verilog HDL Code: Initializing RAM with the readmemb Command

```verilog
reg [7:0] ram[0:15];
initial
begin
  $readmemb("ram.txt", ram);
end
```

When creating a text file to use for memory initialization, specify the address using the format @<location> on a new line, and then specify the memory word such as 110101 or abcde on the next line. Example 17–5 shows a portion of a Memory Initialization File (.mif) for the RAM in Example 17–4.

Example 17–5. Text File Format: Initializing RAM with the readmemb Command

```plaintext
@0
00000000
@1
00000001
@2
00000010
... 
@e
00001110
@f
00001111
```

Verilog HDL Macros

The Quartus II software fully supports Verilog HDL macros, which you can define with the `define compiler directive in your source code. You can also define macros in the Quartus II software or on the command line.

Setting a Verilog HDL Macro Default Value in the Quartus II Software

To specify a macro in the Quartus II software, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, expand Analysis & Synthesis Settings and select Verilog HDL Input.
3. Under Verilog HDL macro, type the macro name in the Name box and the value in the Setting box.
4. Click Add.
Setting a Verilog HDL Macro Default Value on the Command Line

To set a default value for a Verilog HDL macro on the command line, use the `--verilog_macro` option, as shown in Example 17–6.

**Example 17–6. Command Syntax for Specifying a Verilog HDL Macro**

```
quartus_map <Design name> --verilog_macro="<Macro name>=<Macro setting>"
```

The command in Example 17–7 has the same effect as specifying `define a 2` in the Verilog HDL source code.

**Example 17–7. Specifying a Verilog HDL Macro a = 2**

```
quartus_map my_design --verilog_macro="a=2"
```

To specify multiple macros, you can repeat the option more than once, as in Example 17–8.

**Example 17–8. Specifying Verilog HDL Macros a = 2 and b = 3**

```
quartus_map my_design --verilog_macro="a=2" --verilog_macro="b=3"
```

VHDL Support

The Quartus II Compiler’s Analysis & Synthesis module supports the following VHDL standards:

The Quartus II Compiler uses the VHDL 1993 standard by default for files that have the extension `.vhdl` or `.vhd`.

The VHDL code samples provided in this chapter follow the VHDL 1993 standard.

To specify a default VHDL version for all files, follow these steps:
1. On the Assignments menu, click Settings.
2. In the Category list, expand Analysis & Synthesis Settings and select VHDL Input.
3. On the VHDL Input page, under VHDL version, select the appropriate version, and then click OK.

To override the default VHDL version for each VHDL design file, follow these steps:
1. On the Project menu, click Add/Remove Files in Project.
2. On the Files page, select the appropriate file in the list, and then click Properties.
3. In the HDL version list, select VHDL_2008, VHDL_1993, or VHDL_1987, and then click OK.
You can also specify the VHDL version that compiles your design for each design file with the `VHDL_INPUT_VERSION` synthesis directive, as shown in Example 17–9. This directive overrides the default HDL version and any HDL version specified in the File Properties dialog box.


```vhd```
```
--synthesis VHDL_INPUT_VERSION <language version>
```  
```

```vhd```
```
/* synthesis VHDL_INPUT_VERSION <language version> */
```  

The variable `<language version>` requires one of the following values:

- `VHDL_1987`
- `VHDL_1993`
- `VHDL_2008`

When the Quartus II software reads a `VHDL_INPUT_VERSION` synthesis directive, it changes the current language version as specified until after the file or until it reaches the next `VHDL_INPUT_VERSION` directive.

ℹ️ You cannot change the language version in a VHDL design unit.

For more information about specifying synthesis directives, refer to “Synthesis Directives” on page 17–27.

If you use scripts to add design files, you can use the `-HDL_VERSION` command to specify the HDL version for each design file. For more information, refer to “Adding an HDL File to a Project and Setting the HDL Version” on page 17–85.

The Quartus II software reads default values for registered signals defined in the VHDL code and converts the default values into power-up level settings. This enables the power-up state of the synthesized design to match, as closely as possible, the power-up state of the original HDL code in simulation. For more information, refer to “Power-Up Level” on page 17–40.

**VHDL-2008 Support**


⚠️ For more information, refer to *Quartus II Support for VHDL 2008* in Quartus II Help.
VHDL Standard Libraries and Packages

The Quartus II software includes the standard IEEE libraries and several vendor-specific VHDL libraries. For information about organizing your own design units into custom libraries, refer to “Design Libraries” on page 17–12.

The IEEE library includes the standard VHDL packages std_logic_1164, numeric_std, numeric_bit, and math_real. The STD library is part of the VHDL language standard and includes the packages standard (included in every project by default) and textio. For compatibility with older designs, the Quartus II software also supports the following vendor-specific packages and libraries:

- Synopsys packages such as std_logic_arith and std_logic_unsigned in the IEEE library
- Mentor Graphics® packages such as std_logic_arith in the ARITHMETIC library
- Altera primitive packages altera_primitives_components (for primitives such as GLOBAL and DFFE) and maxplus2 (for legacy support of MAX+PLUS® II primitives) in the ALTERA library
- Altera megafuction packages altera_mf_components and stratixgx_mf_components in the ALTERA_MF library (for Altera-specific megafuctions including LCELL), and lpm_components in the LPM library for library of parameterized modules (LPM) functions.

Altera recommends that you import component declarations for Altera primitives such as GLOBAL and DFFE from the altera_primitives_components package and not the altera_mf_components package.

VHDL wait Constructs

The Quartus II software supports one VHDL wait until statement per process block. However, the Quartus II software does not support other VHDL wait constructs, such as wait for and wait on statements, or processes with multiple wait statements.

Example 17–11 is a VHDL code example of a supported wait until construct.

```vhdl
architecture dff_arch of ls_dff is
begin
output: process begin
wait until (CLK'event and CLK='1');
Q <= D;
Qbar <= not D;
end process output;
end dff_arch;
```
AHDL Support

The Quartus II Compiler’s Analysis & Synthesis module fully supports the Altera Hardware Description Language (AHDL).

AHDL designs use Text Design Files (.tdf). You can import AHDL Include Files (.inc) into a .tdf with an AHDL include statement. Altera provides .inc files for all megafunctions shipped with the Quartus II software.

⚠️ The AHDL language does not support the synthesis directives or attributes in this chapter.

❓ For more information about AHDL, refer to About AHDL in the Quartus II Help.

Schematic Design Entry Support

The Quartus II Compiler’s Analysis & Synthesis module fully supports .bdf for schematic design entry.

⚠️ Schematic entry methods do not support the synthesis directives or attributes in this chapter.

❓ For information about creating and editing schematic designs, refer to About Schematic Design Entry in Quartus II Help.

State Machine Editor

The Quartus II software supports graphical state machine entry. To create a new finite state machine (FSM) design, on the File menu, click New. In the New dialog box, expand the Design Files list, and then select State Machine File.

❓ For more information about the State Machine Editor, refer to About the State Machine Editor in Quartus II Help.

Design Libraries

By default, the Quartus II software compiles all design files into the work library. If you do not specify a design library, if a file refers to a library that does not exist, or if the referenced library does not contain a referenced design unit, the Quartus II software searches the work library. This behavior allows the Quartus II software to compile most designs with minimal setup, but you have the option of creating separate custom design libraries.

To compile your design files into specific libraries (for example, when you have two or more functionally different design entities that share the same name), you can specify a destination library for each design file in various ways, as described in the following subsections:

- “Specifying a Destination Library Name in the Settings Dialog Box” on page 17–13
- “Specifying a Destination Library Name in the Quartus II Settings File or with Tcl” on page 17–13
When the Quartus II Compiler analyzes the file, it stores the analyzed design units in the destination library of the file.

A design can contain two or more entities with the same name if the Quartus II software compiles the entities into separate libraries.

When compiling a design instance, the Quartus II software initially searches for the entity in the library associated with the instance (which is the work library if you do not specify any library). If the Quartus II software could not locate the entity definition, the software searches for a unique entity definition in all design libraries. If the Quartus II software finds more than one entity with the same name, the software generates an error. If your design uses multiple entities with the same name, you must compile the entities into separate libraries.

In VHDL, you can associate an instance with an entity in several ways, as described in “Mapping a VHDL Instance to an Entity in a Specific Library” on page 17–14. In Verilog HDL, BDF schematic entry, AHDL, VQM and EDIF netlists, you can use different libraries for each of the entities that have the same name, and compile the instantiation into the same library as the appropriate entity.

Specifying a Destination Library Name in the Settings Dialog Box

To specify a library name for one of your design files, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, select Files.
3. Select the file in the File Name list.
4. Click Properties.
5. In the File Properties dialog box, select the type of design file from the Type list.
6. Type the library name in the Library field.
7. Click OK.

Specifying a Destination Library Name in the Quartus II Settings File or with Tcl

You can specify the library name with the -library option to the <language type>_FILE assignment in the Quartus II Settings File (.qsf) or with Tcl commands.

For example, the following assignments specify that the Quartus II software analyzes the my_file.vhd and stores its contents (design units) in the VHDL library my_lib, and then analyzes the Verilog HDL file my_header_file.h and stores its contents in a library called another_lib. Refer to Example 17–12.

Example 17–12. Specifying a Destination Library Name

```
set_global_assignment -name VHDL_FILE my_file.vhd -library my_lib
set_global_assignment -name VERILOG_FILE my_header_file.h -library another_lib
```

For more information about Tcl scripting, refer to “Scripting Support” on page 17–84.
Specifying a Destination Library Name in a VHDL File

You can use the `library` synthesis directive to specify a library name in your VHDL source file. This directive takes the name of the destination library as a single string argument. Specify the `library` directive in a VHDL comment before the context clause for a primary design unit (that is, a package declaration, an entity declaration, or a configuration), with one of the supported keywords for synthesis directives, that is, `altera`, `synthesis`, `pragma`, `synopsys`, or `exemplar`.

For more information about specifying synthesis directives, refer to “Synthesis Directives” on page 17–27.

The `library` directive overrides the default library destination `work`, the library setting specified for the current file in the Settings dialog box, any existing `.qsf` setting, any setting made through the Tcl interface, or any prior `library` directive in the current file. The directive remains effective until the end of the file or the next `library` synthesis directive.

Example 17–13 uses the `library` synthesis directive to create a library called `my_lib` that contains the design unit `my_entity`.


```vhdl
-- synthesis library my_lib
library ieee;
use ieee.std_logic_1164.all;
entity my_entity(...) end entity my_entity;
```

You can specify a single destination library for all your design units in a given source file by specifying the library name in the Settings dialog box, editing the `.qsf`, or using the Tcl interface. To organize your design units in a single file into different libraries rather than just a single library, you can use the `library` directive to change the destination VHDL library in a source file.

The Quartus II software generates an error if you use the `library` directive in a design unit.

**Mapping a VHDL Instance to an Entity in a Specific Library**

The VHDL language provides several ways to map or bind an instance to an entity in a specific library, as described in the following subsections.
Direct Entity Instantiation

In the direct entity instantiation method, the instantiation refers to an entity in a specific library, as shown in Example 17–14.

Example 17–14. VHDL Code: Direct Entity Instantiation

```vhdl
entity entity1 is
port(...);
end entity entity1;

architecture arch of entity1 is
begin
  inst: entity lib1.foo
  port map(...);
end architecture arch;
```

Component Instantiation—Explicit Binding Instantiation

You can bind a component to an entity in several mechanisms. In an explicit binding indication, you bind a component instance to a specific entity, as shown in Example 17–15.

Example 17–15. VHDL Code: Binding Instantiation

```vhdl
entity entity1 is
port(...);
end entity entity1;

package components is
  component entity1 is
    port map (...);
  end component entity1;
end package components;

entity top_entity is
port(...);
end entity top_entity;

use lib1.components.all;
architecture arch of top_entity is
  -- Explicitly bind instance I1 to entity1 from lib1
  for I1: entity1 use entity lib1.entity1
  port map(...);
  end for;
begin
  I1: entity1 port map(...);
end architecture arch;
```
Component Instantiation—Default Binding

If you do not provide an explicit binding indication, the Quartus II software binds a component instance to the nearest visible entity with the same name. If no such entity is visible in the current scope, the Quartus II software binds the instance to the entity in the library in which you declare the component. For example, if you declare the component in a package in the MY_LIB library, an instance of the component binds to the entity in the MY_LIB library. The code examples in Example 17–16 and Example 17–17 show this instantiation method:

Example 17–16. VHDL Code: Default Binding to the Entity in the Same Library as the Component Declaration

```
use mylib.pkg.foo; -- import component declaration from package "pkg" in
-- library "mylib"
architecture rtl of top
...
begin
  -- This instance will be bound to entity "foo" in library "mylib"
  inst: foo
  port map(...);
end architecture rtl;
```

Example 17–17. VHDL Code: Default Binding to the Directly Visible Entity

```
use mylib.foo; -- make entity "foo" in library "mylib" directly visible
architecture rtl of top
component foo is
  generic (...)
  port (...);
end component;
begin
  -- This instance will be bound to entity "foo" in library "mylib"
  inst: foo
  port map(...);
end architecture rtl;
```

Using Parameters/Generics

This section describes how the Quartus II software supports parameters (known as generics in VHDL) and how you can pass these parameters between design languages.

You can enter default parameter values for your design in the Default Parameters page under the Analysis & Synthesis Settings page in the Settings dialog box. Default parameters enable you to add, change, and delete global parameters for the current assignment. In AHDL, the Quartus II software inherits parameters, so any default parameters apply to all AHDL instances in your design. You can also specify parameters for instantiated modules in a .bdf. To specify parameters in a .bdf instance, double-click the parameter value box for the instance symbol, or right-click the symbol and click Properties, and then click the Parameters tab. For more information about the GUI-based entry methods, the interpretation of parameter values, and format recommendations, refer to “Setting Default Parameter Values and BDF Instance Parameter Values” on page 17–17.
You can specify parameters for instantiated modules in your design source files with the provided syntax for your chosen language. Some designs instantiate entities in a different language; for example, they might instantiate a VHDL entity from a Verilog HDL design file. You can pass parameters or generics between VHDL, Verilog HDL, AHDL, and BDF schematic entry, and from EDIF or VQM to any of these languages. You do not require an additional procedure to pass parameters from one language to another. However, sometimes you must specify the type of parameter you are passing. In those cases, you must follow certain guidelines to ensure that the Quartus II software correctly interprets the parameter value. For more information about parameter type rules, refer to “Passing Parameters Between Two Design Languages” on page 17–19.

Setting Default Parameter Values and BDF Instance Parameter Values

Default parameter values and BDF instance parameter values do not have an explicitly declared type. Usually, the Quartus II software can correctly infer the type from the value without ambiguity. For example, the Quartus II software interprets “ABC” as a string, 123 as an integer, and 15.4 as a floating-point value. In other cases, such as when the instantiated subdesign language is VHDL, the Quartus II software uses the type of the parameter, generic, or both in the instantiated entity to determine how to interpret the value, so that the Quartus II software interprets a value of 123 as a string if the VHDL parameter is of a type string. In addition, you can set the parameter value in a format that is legal in the language of the instantiated entity. For example, to pass an unsized bit literal value from .bdf to Verilog HDL, you can use '1 as the parameter value, and to pass a 4-bit binary vector from .bdf to Verilog HDL, you can use 4'b1111 as the parameter value.

In a few cases, the Quartus II software cannot infer the correct type of parameter value. To avoid ambiguity, specify the parameter value in a type-encoded format in which the first or first and second characters of the parameter indicate the type of the parameter, and the rest of the string indicates the value in a quoted sub-string. For example, to pass a binary string 1001 from .bdf to Verilog HDL, you cannot use the value 1001, because the Quartus II software interprets it as a decimal value. You also cannot use the string "1001" because the Quartus II software interprets it as an ASCII string. You must use the type-encoded string B"1001" for the Quartus II software to correctly interpret the parameter value.
Table 17–2 lists valid parameter strings and how the Quartus II software interprets the parameter strings. Use the type-encoded format only when necessary to resolve ambiguity.

<table>
<thead>
<tr>
<th>Parameter String</th>
<th>Quartus II Parameter Type, Format, and Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>S&quot;abc&quot;, s&quot;abc&quot;</td>
<td>String value abc</td>
</tr>
<tr>
<td>&quot;abc123&quot;, &quot;123abc&quot;</td>
<td>String value abc123 or 123abc</td>
</tr>
<tr>
<td>F&quot;12.3&quot;, f&quot;12.3&quot;</td>
<td>Floating point number 12.3</td>
</tr>
<tr>
<td>-5.4</td>
<td>Floating point -5.4</td>
</tr>
<tr>
<td>D&quot;123&quot;, d&quot;123&quot;</td>
<td>Decimal number 123</td>
</tr>
<tr>
<td>123, -123</td>
<td>Decimal number 123, -123</td>
</tr>
<tr>
<td>X&quot;ff&quot;, H&quot;ff&quot;</td>
<td>Hexadecimal value FF</td>
</tr>
<tr>
<td>Q&quot;77&quot;, O&quot;77&quot;</td>
<td>Octal value 77</td>
</tr>
<tr>
<td>B&quot;1010&quot;, b&quot;1010&quot;</td>
<td>Unsigned binary value 1010</td>
</tr>
<tr>
<td>SB&quot;1010&quot;, sb&quot;1010&quot;</td>
<td>Signed binary value 1010</td>
</tr>
<tr>
<td>R&quot;1&quot;, R&quot;0&quot;, R&quot;X&quot;, R&quot;Z&quot;, X&quot;1&quot;, R&quot;0&quot;, R&quot;X&quot;, R&quot;Z&quot;</td>
<td>Unsize bit literal</td>
</tr>
<tr>
<td>E&quot;apple&quot;, e&quot;apple&quot;</td>
<td>Enumeration type, value name is apple</td>
</tr>
<tr>
<td>P&quot;1 unit&quot;</td>
<td>Physical literal, the value is (1, unit)</td>
</tr>
<tr>
<td>A(...), a(...)</td>
<td>Array type or record type. The string (...) determines the array type or record type content</td>
</tr>
</tbody>
</table>

You can select the parameter type for global parameters or global constants with the pull-down list in the Parameter tab of the Symbol Properties dialog box. If you do not specify the parameter type, the Quartus II software interprets the parameter value and defines the parameter type. You must specify parameter type with the pull-down list to avoid ambiguity.

If you open a .bdf in the Quartus II software, the software automatically updates the parameter types of old symbol blocks by interpreting the parameter value based on the language-independent format. If the Quartus II software does not recognize the parameter value type, the software sets the parameter type as untyped.

The Quartus II software supports the following parameter types:

- Unsigned Integer
- Signed Integer
- Unsigned Binary
- Signed Binary
- Octal
- Hexadecimal
- Float
- Enum
- String
Passing Parameters Between Two Design Languages

When passing a parameter between two different languages, a design block that is higher in the design hierarchy instantiates a lower-level subdesign block and provides parameter information. The subdesign language (the design entity that you instantiate) must correctly interpret the parameter. Based on the information provided by the higher-level design and the value format, and sometimes by the parameter type of the subdesign entity, the Quartus II software interprets the type and value of the passed parameter.

When passing a parameter whose value is an enumerated type value or literal from a language that does not support enumerated types to one that does (for example, from Verilog HDL to VHDL), you must ensure that the enumeration literal is in the correct spelling in the language of the higher-level design block (block that is higher in the hierarchy). The Quartus II software passes the parameter value as a string literal, and the language of the lower-level design correctly converts the string literal into the correct enumeration literal.

If the language of the lower-level entity is SystemVerilog, you must ensure that the enum value is in the correct case. In SystemVerilog, two enumeration literals differ in more than just case. For example, `enum {item, ITEM}` is not a good choice of item names because these names can create confusion and is more difficult to pass parameters from case-insensitive HDLs, such as VHDL.

Arrays have different support in different design languages. For details about the array parameter format, refer to the Parameter section in the Analysis & Synthesis Report of a design that contains array parameters or generics.

The following code shows examples of passing parameters from one design entry language to a subdesign written in another language. Example 17–18 shows a VHDL subdesign that you instantiate in a top-level Verilog HDL design in Example 17–19. Example 17–20 shows a Verilog HDL subdesign that you instantiate in a top-level VHDL design in Example 17–21.

Example 17–18. VHDL Parameterized Subdesign Entity

```vhdl
type fruit is (apple, orange, grape);
entity vhdl_sub is
  generic (
    name : string := "default",
    width : integer := 8,
    number_string : string := "123",
    f : fruit := apple,
    binary_vector : std_logic_vector(3 downto 0) := "0101",
    signed_vector : signed (3 downto 0) := "1111"
  );
end vhdl_sub;
```
Example 17–19. Verilog HDL Top-Level Design Instantiating and Passing Parameters to VHDL
Entity from Example 17–18

```
vhdl_sub inst (...);
defparam inst.name = "lower";
defparam inst.width = 3;
defparam inst.num_string = "321";
defparam inst.f = "grape"; // Must exactly match enum value
defparam inst.binary_vector = 4'b1010;
defparam inst.signed_vector = 4'sb1010;
```

Example 17–20. Verilog HDL Parameterized Subdesign Module

```
module veri_sub (...)
parameter name = "default";
parameter width = 8;
parameter number_string = "123";
parameter binary_vector = 4'b0101;
parameter signed_vector = 4'sb1111;
```

Example 17–21. VHDL Top-Level Design Instantiating and Passing Parameters to the Verilog HDL
Module from Example 17–20

```
inst:veri_sub
generic map (name => "lower",
width => 3,
number_string => "321",
binary_vector = "1010",
signed_vector = "1010")
```

To use an HDL subdesign such as the one shown in Example 17–20 in a top-level .bdf
design, you must generate a symbol for the HDL file, as shown in Figure 17–2. Open
the HDL file in the Quartus II software, and then, on the File menu, point to
Create/Update, and then click Create Symbol Files for Current File.

To specify parameters on a .bdf instance, double-click the parameter value box for the
instance symbol, or right-click the symbol and click Properties, and then click the
Parameters tab. Right-click the symbol and click Update Design File from Selected
Block to pass the updated parameter to the HDL file.

Figure 17–2. BDF Top-Level Design Instantiating and Passing Parameters to the Verilog HDL
Module from Example 17–20

```
<table>
<thead>
<tr>
<th>Parameter</th>
<th>Value</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>&quot;lower&quot;</td>
<td>String</td>
</tr>
<tr>
<td>width</td>
<td>3</td>
<td>Signed Integer</td>
</tr>
<tr>
<td>number_string</td>
<td>&quot;321&quot;</td>
<td>String</td>
</tr>
<tr>
<td>binary_vector</td>
<td>4'b0101</td>
<td>Signed Binary</td>
</tr>
<tr>
<td>signed_vector</td>
<td>4'sb1111</td>
<td>Signed Binary</td>
</tr>
</tbody>
</table>
```
Incremental Compilation

Incremental compilation manages a design hierarchy for incremental design by allowing you to divide your design into multiple partitions. Incremental compilation ensures that the Quartus II software resynthesizes only the updated partitions of your design during compilation, to reduce the compilation time and the runtime memory usage. The feature maintains node names during synthesis for all registered and combinational nodes in unchanged partitions. You can perform incremental synthesis by setting the netlist type for all design partitions to `Post-Synthesis`.

You can also preserve the placement and routing information for unchanged partitions. This feature allows you to preserve performance of unchanged blocks in your design and reduces the time required for placement and routing, which significantly reduces your design compilation time.

For more information about incremental compilation, refer to `About Incremental Compilation` in Quartus II Help.

For more information about incremental compilation, refer to `Quartus II Incremental Compilation for Hierarchical and Team-Based Design` and `Best Practices for Incremental Compilation Partitions and Floorplan Assignments` chapters in volume 1 of the `Quartus II Handbook`.

Partitions for Preserving Hierarchical Boundaries

A design partition represents a portion of your design that you want to synthesize and fit incrementally.

If you want to preserve the `Optimization Technique` and `Restructure Multiplexers` logic options in any entity, you must create new partitions for the entity instead of using the `Preserve Hierarchical Boundary` logic option. If you have settings applied to specific existing design hierarchies, particularly those created in the Quartus II software versions before 9.0, you must create a design partition for the design hierarchy so that synthesis can optimize the design instance independently and preserve the hierarchical boundaries.

The `Preserve Hierarchical Boundary` logic option is available only in Quartus II software versions 8.1 and earlier. Altera recommends using design partitions if you want to preserve hierarchical boundaries through the synthesis and fitting process, because incremental compilation maintains the hierarchical boundaries of design partitions.

Parallel Synthesis

The `Parallel Synthesis` logic option reduces compilation time for synthesis. The option enables the Quartus II software to use multiple processors to synthesize multiple partitions in parallel.

This option is available when you perform the following tasks:

- Specifying the maximum number of processors allowed under `Parallel Compilation` options in the `Compilation Process Settings` page of the `Settings` dialog box.
■ Enabling the incremental compilation feature.
■ Using two or more partitions in your design.
■ Turning on the Parallel Synthesis option.

By default, the Quartus II software enables the Parallel Synthesis option. To disable parallel synthesis, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, click Analysis & Synthesis Settings, and then click More Settings to select Parallel Synthesis.

You can also set the Parallel Synthesis option with the following Tcl command, as shown in Example 17–22:

```
Example 17–22. Setting the Parallel Synthesis Option with Tcl Command

set_global_assignment -name parallel_synthesis off
```

If you use the command line, you can differentiate among the interleaved messages by turning on the Show partition that generated the message option in the Messages page. This option shows the partition ID in parenthesis for each message.

You can view all the interleaved messages from different partitions in the Messages window. The Partition column in the Messages window displays the partition ID of the partition referred to in the message. After compilation, you can sort the messages by partition.

For more information about displaying the Partition column, refer to About the Messages Window in Quartus II Help.

Quartus II Exported Partition File as Source

You can use a .qxp as a source file in incremental compilation. The .qxp contains the precompiled design netlist exported as a partition from another Quartus II project, and fully defines the entity. Project team members or intellectual property (IP) providers can use a .qxp to send their design to the project lead, instead of sending the original HDL source code. The .qxp preserves the compilation results and instance-specific assignments. Not all global assignments can function in a different Quartus II project. You can override the assignments for the entity in the .qxp by applying assignments in the top-level design.

For more information about .qxp, refer to Quartus II Exported Partition File (.qxp) in Quartus II Help.

For more information about exporting design partitions and using .qxp files, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Quartus II Synthesis Options

The Quartus II software offers several options to help you control the synthesis process and achieve optimal results for your design. “Setting Synthesis Options” on page 17–25 describes the Analysis & Synthesis Settings page of the Settings dialog box, in which you can set the most common global settings and options, and defines the following types of synthesis options: Quartus II logic options, synthesis attributes, and synthesis directives.

When you apply a Quartus II Synthesis option globally or to an entity, the option affects all lower-level entities in the hierarchy path, including entities instantiated with Altera and third-party IP.

The following subsections describe the following common synthesis options in the Quartus II software, and provide HDL examples on how to use each option:

- Major Optimization Settings:
  - “Optimization Technique” on page 17–28
  - “Auto Gated Clock Conversion” on page 17–29
  - “PowerPlay Power Optimization” on page 17–31
  - “Restructure Multiplexers” on page 17–34

- Settings Related to Timing Constraints:
  - “Timing-Driven Synthesis” on page 17–30
  - “Optimization Technique” on page 17–28
  - “Auto Gated Clock Conversion” on page 17–29
  - “SDC Constraint Protection” on page 17–31

- State Machine Settings and Enumerated Types:
  - “State Machine Processing” on page 17–34
  - “Manually Specifying State Assignments Using the syn_encoding Attribute” on page 17–36
  - “Manually Specifying Enumerated Types Using the enum_encoding Attribute” on page 17–37
  - “Safe State Machine” on page 17–38

- Register Power-Up Settings:
  - “Power-Up Level” on page 17–40
  - “Power-Up Don’t Care” on page 17–41
Controlling, Preserving, Removing, and Duplicating Logic and Registers:
- “Limiting Resource Usage in Partitions” on page 17–32
- “Remove Duplicate Registers” on page 17–41
- “Preserve Registers” on page 17–42
- “Disable Register Merging/Don’t Merge Register” on page 17–43
- “Noprune Synthesis Attribute/Preserve Fan-out Free Register Node” on page 17–43
- “Keep Combinational Node/Implement as Output of Logic Cell” on page 17–44
- “Disabling Synthesis Netlist Optimizations with dont_retime Attribute” on page 17–45
- “Disabling Synthesis Netlist Optimizations with dont_replicate Attribute” on page 17–46
- “Maximum Fan-Out” on page 17–47
- “Controlling Clock Enable Signals with Auto Clock Enable Replacement and direct_enable” on page 17–48
- “Auto Gated Clock Conversion” on page 17–29
- “Partitions for Preserving Hierarchical Boundaries” on page 17–21

Megafunction Inference Options:
- “Inferring Multiplier, DSP, and Memory Functions from HDL Code” on page 17–49
- “RAM Style and ROM Style—for Inferred Memory” on page 17–53
- “Turning Off the Add Pass-Through Logic to Inferred RAMs no_rw_check Attribute” on page 17–56
- “RAM Initialization File—for Inferred Memory” on page 17–59
- “Multiplier Style—for Inferred Multipliers” on page 17–60

Controlling Synthesis with Other Synthesis Directives:
- “Full Case Attribute” on page 17–62
- “Parallel Case” on page 17–63
- “Translate Off and On / Synthesis Off and On” on page 17–64
- “Ignore translate_off and synthesis_off Directives” on page 17–65
- “Read Comments as HDL” on page 17–66

Specifying I/O-Related Assignments:
- “Use I/O Flipflops” on page 17–67
- “Specifying Pin Locations with chip_pin” on page 17–68

Setting Quartus II Logic Options in Your HDL Source Code:
- “Using altera_attribute to Set Quartus II Logic Options” on page 17–70
Setting Synthesis Options

You can set synthesis options in the Settings dialog box, or with logic options in the Quartus II software, or you can use synthesis attributes and directives in your HDL source code.

The Analysis & Synthesis Settings page of the Settings dialog box allows you to set global synthesis options that apply to the entire project. You can also use a corresponding Tcl command.

You can set some of the advanced synthesis settings in the Physical Synthesis Optimizations page under Compilation Process Settings.

For more information about Physical Synthesis options, refer to the Netlist Optimizations and Physical Synthesis chapter in volume 2 of the Quartus II Handbook.

Quartus II Logic Options

The Quartus II logic options control many aspects of the synthesis and placement and routing process. To set logic options in the Quartus II software, on the Assignments menu, click Assignment Editor. You can also use a corresponding Tcl command to set global assignments. The Quartus II logic options enable you to set instance or node-specific assignments without editing the source HDL code.

For more information about using the Assignment Editor, refer to the About the Assignment Editor in Quartus II Help.

Synthesis Attributes

The Quartus II software supports synthesis attributes for Verilog HDL and VHDL, also commonly called pragmas. These attributes are not standard Verilog HDL or VHDL commands. Synthesis tools use attributes to control the synthesis process. The Quartus II software applies the attributes in the HDL source code, and attributes always apply to a specific design element. Some synthesis attributes are also available as Quartus II logic options via the Quartus II software or scripting. Each attribute description in this chapter indicates a corresponding setting or a logic option that you can set in the Quartus II software. You can specify only some attributes with HDL synthesis attributes.

Attributes specified in your HDL code are not visible in the Assignment Editor or in the .qsf. Assignments or settings made with the Quartus II software, the .qsf, or the Tcl interface take precedence over assignments or settings made with synthesis attributes in your HDL code. The Quartus II software generates warning messages if the software finds invalid attributes, but does not generate an error or stop the compilation. This behavior is necessary because attributes are specific to various design tools, and attributes not recognized in the Quartus II software might be for a different EDA tool. The Quartus II software lists the attributes specified in your HDL code in the Source assignments table of the Analysis & Synthesis report.
The Verilog-2001, SystemVerilog, and VHDL language definitions provide specific syntax for specifying attributes, but in Verilog-1995, you must embed attribute assignments in comments. You can enter attributes in your code using the syntax in Example 17–23 through Example 17–29, in which `<attribute>`, `<attribute type>`, `<value>`, `<object>`, and `<object type>` are variables, and the entry in brackets is optional. The examples in this chapter demonstrate each syntax form.

Verilog HDL is case sensitive; therefore, synthesis attributes in Verilog HDL files are also case sensitive.


```verbatim
// synthesis <attribute> [ = <value> ]
or
/* synthesis <attribute> [ = <value> ] */
```

You must use Verilog-1995 comment-embedded attributes as a suffix to the declaration of an item and must appear before a semicolon, when a semicolon is necessary (refer to Example 17–23).

You cannot use the open one-line comment in Verilog HDL when a semicolon is necessary after the line, because it is not clear to which HDL element that the attribute applies. For example, you cannot make an attribute assignment such as `reg r; // synthesis <attribute>` because the Quartus II software could read the attribute as part of the next line.

To apply multiple attributes to the same instance in Verilog-1995, separate the attributes with spaces, as shown in Example 17–24:

Example 17–24. Applying Multiple Attributes to the Same Instance in Verilog-1996

```verbatim
//synthesis <attribute1> [ = <value> ] <attribute2> [ = <value> ]
```

For example, to set the maxfan attribute to 16 (for details, refer to “Maximum Fan-Out” on page 17–47) and set the preserve attribute (for details, refer to “Preserve Registers” on page 17–42) on a register called my_reg, use the following syntax as shown in Example 17–25:

Example 17–25. Setting maxfan and preserve Attribute on a Register

```verbatim
reg my_reg /* synthesis maxfan = 16 preserve */;
```

In addition to the synthesis keyword shown above, the Quartus II software supports the pragma, synopsys, and exemplar keywords for compatibility with other synthesis tools. The software also supports the altera keyword, which allows you to add synthesis attributes that the Quartus II Integrated Synthesis feature recognizes and not by other tools that recognize the same synthesis attribute.

Because formal verification tools do not recognize the exemplar, pragma, and altera keywords, avoid using these attribute keywords when using formal verification.
You must use Verilog-2001 attributes as a prefix to a declaration, module item, statement, or port connection, and as a suffix to an operator or a Verilog HDL function name in an expression (refer to Example 17–26).


\[ (* <attribute> [ = <value> ] *) \]

Formal verification does not support the Verilog-2001 attribute syntax because the tools do not recognize the syntax.

To apply multiple attributes to the same instance in Verilog-2001 or SystemVerilog, separate the attributes with commas, as shown in Example 17–27:

**Example 17–27. Applying Multiple Attributes**

\[ (* <attribute1> [ = <value1> ], <attribute2> [ = <value2> ] *) \]

For example, to set the `maxfan` attribute to 16 (refer to “Maximum Fan-Out” on page 17–47 for details) and set the `preserve` attribute (refer to “Preserve Registers” on page 17–42 for details) on a register called `my_reg`, use the following syntax as shown in Example 17–28:

**Example 17–28. Setting Attribute**

\[ (* maxfan = 16, preserve *) reg my_reg; \]

VHDL attributes, as shown in Example 17–29, declare and apply the attribute type to the object you specify.

**Example 17–29. Synthesis Attributes in VHDL**

\[
\text{attribute <attribute> : <attribute type> ;}
\text{attribute <attribute> of <object> : <object type> is <value>;}
\]

The Quartus II software defines and applies each attribute separately to a given node. For VHDL designs, the software declares all supported synthesis attributes in the `altera_syn_attributes` package in the Altera library. You can call this library from your VHDL code to declare the synthesis attributes, as shown in Example 17–30:

**Example 17–30.**

\[
\text{LIBRARY altera;}
\text{USE altera.altera_syn_attributes.all;}
\]

**Synthesis Directives**

The Quartus II software supports synthesis directives, also commonly called compiler directives or pragmas. You can include synthesis directives in Verilog HDL or VHDL code as comments. These directives are not standard Verilog HDL or VHDL commands. Synthesis tools use directives to control the synthesis process. Directives do not apply to a specific design node, but change the behavior of the synthesis tool from the point in which they occur in the HDL source code. Other tools, such as simulators, ignore these directives and treat them as comments.
You can enter synthesis directives in your code using the syntax in Example 17–31, Example 17–32, and Example 17–33, in which <directive> and <value> are variables, and the entry in brackets are optional. For synthesis directives, no equal sign before the value is necessary; this is different than the Verilog syntax for synthesis attributes. The examples in this chapter demonstrate each syntax form.

Verilog HDL is case sensitive; therefore, all synthesis directives are also case sensitive.

**Example 17–31. Specifying Synthesis Directives with Verilog HDL**

```
// synthesis <directive> [ <value> ]
```

or

```
/* synthesis <directive> [ <value> ] */
```

**Example 17–32. Specifying Synthesis Directives with VHDL**

```
-- synthesis <directive> [ <value> ]
```

**Example 17–33. Specifying Synthesis Directives with VHDL-2008**

```
/* synthesis <directive> [<value>] */
```

In addition to the *synthesis* keyword shown above, the software supports the *pragma*, *synopsys*, and *exemplar* keywords in Verilog HDL and VHDL for compatibility with other synthesis tools. The Quartus II software also supports the keyword *altera*, which allows you to add synthesis directives that only Quartus II Integrated Synthesis feature recognizes, and not by other tools that recognize the same synthesis directives.

Because formal verification tools ignore the *exemplar*, *pragma*, and *altera* keywords, Altera recommends that you avoid using these directive keywords when you use formal verification to prevent mismatches with the Quartus II results.

**Optimization Technique**

The **Optimization Technique** logic option specifies the goal for logic optimization during compilation; that is, whether to attempt to achieve maximum speed performance or minimum area usage, or a balance between the two.

For more information about the **Optimization Technique** logic option, refer to *Optimization Technique logic option* in Quartus II Help.
Auto Gated Clock Conversion

Clock gating is a common optimization technique in ASIC designs to minimize power consumption. You can use the Auto Gated Clock Conversion logic option to optimize your prototype ASIC designs by converting gated clocks into clock enables when you use FPGAs in your ASIC prototyping. The automatic conversion of gated clocks to clock enables is more efficient than manually modifying source code. The Auto Gated Clock Conversion logic option automatically converts qualified gated clocks (base clocks as defined in the Synopsys Design Constraints [SDC]) to clock enables. To use Auto Gated Clock Conversion, you must select the option from the More Analysis & Synthesis Settings dialog box, in the Analysis & Synthesis Settings page.

The gated clock conversion occurs when all these conditions are met:

- Only one base clock drives a gated-clock
- For one set of gating input values, the value output of the gated clock remains constant and does not change as the base clock changes
- For one value of the base clock, changes in the gating inputs do not change the value output for the gated clock

The option supports combinational gates in clock gating network.

Figure 17–3 shows example of gated clock conversions.

Figure 17–3. Example Gated Clock Conversion
This option does not support registers in RAM, DSP blocks, or I/O related WYSIWYG primitives. Because the gated-clock conversion cannot trace the base clock from the gated clock, the gated clock conversion does not support multiple design partitions from incremental compilation in which the gated clock and base clock are not in the same hierarchical partition. A gated clock tree, instead of every gated clock, is the basis of each conversion. Therefore, if you cannot convert a gated clock from a root gated clock of a multiple cascaded gated clock, the conversion of the entire gated clock tree fails.

The Info tab in the Messages window lists all the converted gated clocks. You can view a list of converted and nonconverted gated clocks from the Compilation Report under the Optimization Results of the Analysis & Synthesis Report. The Gated Clock Conversion Details table lists the reasons for nonconverted gated clocks.

For more information about Auto Gated Clock Conversion logic option and a list of supported devices, refer to Auto Gated Clock Conversion logic option in Quartus II Help.

Timing-Driven Synthesis

The Timing-Driven Synthesis logic option specifies whether Analysis & Synthesis should use the SDC timing constraints of your design to better optimize the circuit. When you turn on this option, Analysis & Synthesis runs timing analysis to obtain timing information about the netlist, and then considers the SDC timing constraints to focus on critical portions of your design when optimizing for performance, while optimizing noncritical portions for area. When you turn on this option, Analysis & Synthesis also protects SDC constraints by not merging duplicate registers that have incompatible timing constraints. For more information, refer to “SDC Constraint Protection” on page 17–31.

When you turn on the Timing-Driven Synthesis logic option, Analysis & Synthesis increases performance by improving logic depth on critical portions of your design, and improving area on noncritical portions of your design. The increased performance affects the amount of area used, specifically adaptive look-up tables (ALUTs) and registers in your design. Depending on how much of your design is timing critical, overall area can increase or decrease when you turn on the Timing-Driven Synthesis logic option. Runtime and peak memory use increases slightly if you turn on the Timing-Driven Synthesis logic option.

When you turn on the Timing-Driven Synthesis logic option, the Optimization Technique logic option has the following effect. With Optimization Technique Speed, Timing-Driven Synthesis optimizes timing-critical portions of your design for performance at the cost of increasing area (logic and register utilization). With an Optimization Technique of Balanced, Timing-Driven Synthesis also optimizes the timing-critical portions of your design for performance, but the option allows only limited area increase. With Optimization Technique Area, Timing-Driven Synthesis optimizes your design only for area. Timing-Driven Synthesis prevents registers with incompatible timing constraints from merging for any Optimization Technique setting. If your design contains multiple partitions, you can select Timing-Driven Synthesis unique options for each partition. If you use a .qxp as a source file, or if your design uses partitions developed in separate Quartus II projects, the software cannot properly compute timing of paths that cross the partition boundaries.
Even with the Optimization Technique logic option set to Speed, the Timing-Driven Synthesis option still considers the resource usage in your design when increasing area to improve timing. For example, the Timing-Driven Synthesis option checks if a device has enough registers before deciding to implement the shift registers in logic cells instead of RAM for better timing performance.

When using incremental compilation, Integrated Synthesis allows each partition to use up all the registers in a device. You can use the Maximum Number of LABs settings to specify the number of LABs that every partition can use. If your design has only one partition, you can also use the Maximum Number of LABs settings to limit the number of resources that your design can use. This limitation is useful when you add more logic to your design.

To turn on or turn off the Timing-Driven Synthesis logic option, follow these steps:

1. On the Assignment menu, click Settings.
2. In the Category list, select Analysis & Synthesis Settings. In the Analysis & Synthesis Settings page, turn on or turn off Timing-Driven Synthesis.

Altera recommends that you select a specific device for timing-driven synthesis to have the most accurate timing information. When you select auto device, timing-driven synthesis uses the smallest device for the selected family to obtain timing information.

For more information about Timing-Driven Synthesis logic option and a list of supported devices, refer to Timing-Driven Synthesis logic option in Quartus II Help.

**SDC Constraint Protection**

The SDC Constraint Protection option specifies whether Analysis & Synthesis should protect registers from merging when they have incompatible timing constraints. For example, when you turn on this option, the software does not merge two registers that are duplicates of each other but have different multicycle constraints on them. When you turn on the Timing-Driven Synthesis option, the software detects registers with incompatible constraints, and you do not need to turn on SDC Constraint Protection. To use the SDC constraint protection option, you must turn on the option in the More Analysis & Synthesis Settings dialog box in the Analysis & Synthesis Settings page.

**PowerPlay Power Optimization**

The PowerPlay Power Optimization logic option controls the power-driven compilation setting of Analysis & Synthesis and determines how aggressively Analysis & Synthesis optimizes your design for power.

For more information about the available settings for the PowerPlay power optimization logic option and a list of supported devices, refer to PowerPlay Power Optimization logic option in Quartus II Help.
For more information about optimizing your design for power utilization, refer to the Power Optimization chapter in volume 2 of the Quartus II Handbook. For information about analyzing your power results, refer to the PowerPlay Power Analysis chapter in volume 3 of the Quartus II Handbook.

Limiting Resource Usage in Partitions

Resource balancing is important when performing Analysis & Synthesis. During resource balancing, Quartus II Integrated Synthesis considers the amount of used and available DSP and RAM blocks in the device, and tries to balance these resources to prevent no-fit errors.

For DSP blocks, resource balancing converts the remaining DSP blocks to equivalent logic if there are more DSP blocks in your design that the software can place in the device. For RAM blocks, resource balancing converts RAM blocks to different types of RAM blocks if there are not enough blocks of a certain type available in the device; however, Quartus II Integrated Synthesis does not convert RAM blocks to logic.

The RAM balancing feature does not support Stratix V devices because Stratix V has only M20K memory blocks.

By default, Quartus II Integrated Synthesis considers the information in the targeted device to identify the number of available DSP or RAM blocks. However, in incremental compilation, each partition considers the information in the device independently and consequently assumes that the partition has all the DSP and RAM blocks in the device available for use, resulting in over allocation of DSP or RAM blocks in your design, which means that the total number of DSP or RAM blocks used by all the partitions is greater than the number of DSP or RAM blocks available in the device, leading to a no-fit error during the fitting process.

The following sections describe the methods to prevent a no-fit error during the fitting process:

- “Creating LogicLock Regions” on page 17–32
- “Using Assignments to Limit the Number of RAM and DSP Blocks” on page 17–33

Creating LogicLock Regions

The floorplan-aware synthesis feature allows you to use LogicLock regions to define resource allocation for DSP blocks and RAM blocks. For example, if you assign a certain partition to a certain LogicLock region, resource balancing takes into account that all the DSP and RAM blocks in that partition need to fit in this LogicLock region. Resource balancing then balances the DSP and RAM blocks accordingly.

Because floorplan-aware balancing step considers only one partition at a time, it does not know that nodes from another partition may be using the same resources. When using this feature, Altera recommends that you do not manually assign nodes from different partitions to the same LogicLock region.

If you do not want the software to consider the LogicLock floorplan constraints when performing DSP and RAM balancing, you can turn off the floorplan-aware synthesis feature. You can turn off the Use LogicLock Constraints During Resource Balancing option in the More Analysis & Synthesis Settings dialog box in the Analysis & Synthesis Settings page.
For more information about using LogicLock regions to create a floorplan for incremental compilation, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Using Assignments to Limit the Number of RAM and DSP Blocks

For DSP and RAM block balancing, you can use assignments to limit the maximum number of blocks that the balancer allows. You can set these assignments globally or on individual partitions. For DSP block balancing, the Maximum DSP Block Usage logic option allows you to specify the maximum number of DSP blocks that the DSP block balancer assumes are available for the current partition. For RAM blocks, the floorplan-aware logic option allows you to specify maximum resources for different RAM types, such as Maximum Number of M4K/M9K/M20K/M10K Memory Blocks, Maximum Number of M512 Memory Blocks, Maximum Number of M-RAM/M144K Memory Blocks, or Maximum Number of LABs.

The partition-specific assignment overrides the global assignment, if any. However, each partition that does not have a partition-specific assignment uses the value set by the global assignment, or the value derived from the device size if no global assignment exists. This action can also lead to over allocation. Therefore, Altera recommends that you always set the assignment on each partition individually.

To select the Maximum Number <block type> Memory Blocks option or the Maximum DSP Block Usage option globally, follow these steps:

1. On the Assignment menu, click Settings.
2. Under Category, click Analysis & Synthesis Settings.
3. In the Analysis & Synthesis Settings dialog box, click More Settings.
4. In the Name list, select the required option and set the necessary value.

You can use the Assignment Editor to set this assignment on a partition by selecting the assignment, and setting it on the root entity of a partition. You can set any positive integer as the value of this assignment. If you set this assignment on a name other than a partition root, Analysis & Synthesis gives an error.

For more information about the logic options, including a list of supported device families, refer to Maximum DSP Block Usage logic option, Maximum Number of M4K/M9K/M20K/M10K Memory Blocks logic option, Maximum Number of M512 Memory Blocks logic option, Maximum Number of M-RAM/144K Memory Blocks logic option, and Maximum Number of LABs logic option in Quartus II Help.
Restructure Multiplexers

The Restructure Multiplexers logic option restructures multiplexers to create more efficient use of area, allowing you to implement multiplexers with a reduced number of LEs or ALMs.

When multiplexers from one part of your design feed multiplexers in another part of your design, trees of multiplexers form. Multiplexers may arise in different parts of your design through Verilog HDL or VHDL constructs such as the “if,” “case,” or “?:” statements. Multiplexer buses occur most often as a result of multiplexing together arrays in Verilog HDL, or STD_LOGIC_VECTOR signals in VHDL. The Restructure Multiplexers logic option identifies buses of multiplexer trees that have a similar structure. This logic option optimizes the structure of each multiplexer bus for the target device to reduce the overall amount of logic in your design.

Results of the multiplexer optimizations are design dependent, but area reductions as high as 20% are possible. The option can negatively affect your design’s fMAX.

For more information about optimizing for multiplexers, refer to the “Multiplexers” section of the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

For more information about the Multiplexer Restructuring Statistics report table for each bus of multiplexers, refer to Analysis & Synthesis Optimization Results Reports in Quartus II Help.

For more information about the Restructure Multiplexers logic option, including the settings and a list of supported device families, refer to Restructure Multiplexers logic option in Quartus II Help.

Synthesis Effort

The Synthesis Effort logic option specifies the overall synthesis effort level in the Quartus II software.

For more information about Synthesis Effort logic option, including a list of supported device families, refer to Synthesis Effort logic option in Quartus II Help.

Synthesis Seed

The Synthesis Seed option specifies the seed that Synthesis uses to randomly run synthesis in a slightly different way. You can use this seed when your design is close to meeting requirements, to get a slightly different result. The seeds that produce the best result for a design might change if your design changes.

To set the Synthesis Seed option from the Quartus II software, on the Analysis & Synthesis Settings page, click More Settings. The default value is 1. You can specify a positive integer value.

State Machine Processing

The State Machine Processing logic option specifies the processing style to synthesize a state machine.
Chapter 17: Quartus II Integrated Synthesis

Quartus II Synthesis Options

The default state machine encoding, **Auto**, uses one-hot encoding for FPGA devices and minimal-bits encoding for CPLDs. These settings achieve the best results on average, but another encoding style might be more appropriate for your design, so this option allows you to control the state machine encoding.

For guidelines on how to correctly infer and encode your state machine, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

For one-hot encoding, the Quartus II software does not guarantee that each state has one bit set to one and all other bits set to zero. Quartus II Integrated Synthesis creates one-hot register encoding with standard one-hot encoding and then inverts the first bit. This results in an initial state with all zero values, and the remaining states have two 1 values. Quartus II Integrated Synthesis encodes the initial state with all zeros for the state machine power-up because all device registers power up to a low value. This encoding has the same properties as true one-hot encoding: the software recognizes each state by the value of one bit. For example, in a one-hot-encoded state machine with five states, including an initial or reset state, the software uses the register encoding shown in Example 17–34:

**Example 17–34. Register Encoding**

<table>
<thead>
<tr>
<th>State</th>
<th>Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>State 0</td>
<td>0 0 0 0 0</td>
</tr>
<tr>
<td>State 1</td>
<td>0 0 0 1 1</td>
</tr>
<tr>
<td>State 2</td>
<td>0 0 1 0 1</td>
</tr>
<tr>
<td>State 3</td>
<td>0 1 0 0 1</td>
</tr>
<tr>
<td>State 4</td>
<td>1 0 0 0 1</td>
</tr>
</tbody>
</table>

If you set the **State Machine Processing** logic option to **User-Encoded** in a Verilog HDL design, the software starts with the original design values for the state constants. For example, a Verilog HDL design can contain a declaration such as shown in Example 17–35:

**Example 17–35.**

```verilog
class S0 = 4'b1010, S1 = 4'b0101, ...
```

If the software infers the states S0, S1, ... the software uses the encoding 4'b1010, 4'b0101, ... If necessary, the software inverts bits in a user-encoded state machine to ensure that all bits of the reset state of the state machine are zero.

You can view the state machine encoding from the Compilation Report under the State Machines of the Analysis & Synthesis Report. The State Machine Viewer displays only a graphical representation of the state machines as interpreted from your design.

For more information about the State Machine Viewer, refer to the State Machine Viewer section of the *Analyzing Designs with Quartus II Netlist Viewers* chapter in volume 1 of the *Quartus II Handbook*. 
To assign your own state encoding with the **User-Encoded** setting of the **State Machine Processing** option in a VHDL design, you must apply specific binary encoding to the elements of an enumerated type because enumeration literals have no numeric values in VHDL. Use the `syn_encoding` synthesis attribute to apply your encoding values. For more information, refer to “Manually Specifying State Assignments Using the `syn_encoding` Attribute”.

For information about the **State Machine Processing** logic option, including the settings and supported devices, refer to *State Machine Processing logic option* in Quartus II Help.

### Manually Specifying State Assignments Using the `syn_encoding` Attribute

The Quartus II software infers state machines from enumerated types and automatically assigns state encoding based on “State Machine Processing” on page 17–34. With this logic option, you can choose the value **User-Encoded** to use the encoding from your HDL code. However, in standard VHDL code, you cannot specify user encoding in the state machine description because enumeration literals have no numeric values in VHDL.

To assign your own state encoding for the **User-Encoded State Machine Processing** setting, use the `syn_encoding` synthesis attribute to apply specific binary encodings to the elements of an enumerated type or to specify an encoding style. The Quartus II software can implement Enumeration Types with different encoding styles, as shown in Table 17–3.

<table>
<thead>
<tr>
<th>Attribute Value</th>
<th>Enumeration Types</th>
</tr>
</thead>
<tbody>
<tr>
<td>&quot;default&quot;</td>
<td>Use an encoding based on the number of enumeration literals in the Enumeration Type. If the number of literals is less than five, use the &quot;sequential&quot; encoding. If the number of literals is more than five, but fewer than 50, use a &quot;one-hot&quot; encoding. Otherwise, use a &quot;gray&quot; encoding.</td>
</tr>
<tr>
<td>&quot;sequential&quot;</td>
<td>Use a binary encoding in which the first enumeration literal in the Enumeration Type has encoding 0 and the second 1.</td>
</tr>
<tr>
<td>&quot;gray&quot;</td>
<td>Use an encoding in which the encodings for adjacent enumeration literals differ by exactly one bit. An $N$-bit gray code can represent $2^N$ values.</td>
</tr>
<tr>
<td>&quot;johnson&quot;</td>
<td>Use an encoding similar to a gray code. An $N$-bit Johnson code can represent at most $2^N$ states, but requires less logic than a gray encoding.</td>
</tr>
<tr>
<td>&quot;one-hot&quot;</td>
<td>The default encoding style requiring $N$ bits, in which $N$ is the number of enumeration literals in the Enumeration Type.</td>
</tr>
<tr>
<td>&quot;compact&quot;</td>
<td>Use an encoding with the fewest bits.</td>
</tr>
<tr>
<td>&quot;user&quot;</td>
<td>Encode each state using its value in the Verilog source. By changing the values of your state constants, you can change the encoding of your state machine.</td>
</tr>
</tbody>
</table>

The `syn_encoding` attribute must follow the enumeration type definition, but precede its use.
Manually Specifying Enumerated Types Using the enum_encoding Attribute

By default, the Quartus II software one-hot encodes all enumerated types you defined. With the \texttt{enum_encoding} attribute, you can specify the logic encoding for an enumerated type and override the default one-hot encoding to improve the logic efficiency.

If an enumerated type represents the states of a state machine, using the \texttt{enum_encoding} attribute to specify a manual state encoding prevents the Compiler from recognizing state machines based on the enumerated type. Instead, the Compiler processes these state machines as regular logic with the encoding specified by the attribute, and the Report window for your project does not list these states machines as state machines. If you want to control the encoding for a recognized state machine, use the \textbf{State Machine Processing} logic option and the \texttt{syn_encoding synthesis} attribute.

To use the \texttt{enum_encoding} attribute in a VHDL design file, associate the attribute with the enumeration type whose encoding you want to control. The \texttt{enum_encoding} attribute must follow the enumeration type definition, but precede its use. In addition, the attribute value should be a string literal that specifies either an arbitrary user encoding or an encoding style of "default", "sequential", "gray", "johnson", or "one-hot".

An arbitrary user encoding consists of a space-delimited list of encodings. The list must contain as many encodings as the number of enumeration literals in your enumeration type. In addition, the encodings should have the same length, and each encoding must consist solely of values from the \texttt{std_ulogic} type declared by the \texttt{std_logic_1164} package in the IEEE library. In Example 17–36, the \texttt{enum_encoding} attribute specifies an arbitrary user encoding for the enumeration type \texttt{fruit}.

\textbf{Example 17–36. Specifying an Arbitrary User Encoding for Enumerated Type}

\begin{verbatim}
  type fruit is (apple, orange, pear, mango);
  attribute enum_encoding : string;
  attribute enum_encoding of fruit : type is "11 01 10 00";
\end{verbatim}

Example 17–37 shows the encoded enumeration literals:

\textbf{Example 17–37. Encoded Enumeration Literals}

\begin{verbatim}
  apple    = "11"
  orange   = "01"
  pear     = "10"
  mango    = "00"
\end{verbatim}
Altera recommends that you specify an encoding style, rather than a manual user encoding, especially when the enumeration type has a large number of enumeration literals. The Quartus II software can implement Enumeration Types with the different encoding styles, as shown in Table 17–4.

### Table 17–4. *enum_encoding* Attribute Values

<table>
<thead>
<tr>
<th>Attribute Value</th>
<th>Enumeration Types</th>
</tr>
</thead>
<tbody>
<tr>
<td>&quot;default&quot;</td>
<td>Use an encoding based on the number of enumeration literals in the enumeration type. If the number of literals are fewer than five, use the &quot;sequential&quot; encoding. If the number of literals are more than five, but fewer than 50 literals, use a &quot;one-hot&quot; encoding. Otherwise, use a &quot;gray&quot; encoding.</td>
</tr>
<tr>
<td>&quot;sequential&quot;</td>
<td>Use a binary encoding in which the first enumeration literal in the enumeration type has encoding 0 and the second 1.</td>
</tr>
<tr>
<td>&quot;gray&quot;</td>
<td>Use an encoding in which the encodings for adjacent enumeration literals differ by exactly one bit. An $N$-bit gray code can represent $2^N$ values.</td>
</tr>
<tr>
<td>&quot;johnson&quot;</td>
<td>Use an encoding similar to a gray code. An $N$-bit Johnson code can represent at most $2^N$ states, but requires less logic than a gray encoding.</td>
</tr>
<tr>
<td>&quot;one-hot&quot;</td>
<td>The default encoding style requiring $N$ bits, in which $N$ is the number of enumeration literals in the enumeration type.</td>
</tr>
</tbody>
</table>

In Example 17–36, the *enum_encoding* attribute manually specified a gray encoding for the enumeration type *fruit*. You can concisely write this example by specifying the "gray" encoding style instead of a manual encoding, as shown in Example 17–38.

**Example 17–38. Specifying the “gray” Encoding Style or Enumeration Type**

```plaintext
type fruit is (apple, orange, pear, mango);
attribute enum_encoding : string;
attribute enum_encoding of fruit : type is "gray";
```

### Safe State Machine

The **Safe State Machine** logic option and corresponding *syn_encoding* attribute value *safe* specify that the software must insert extra logic to detect an illegal state, and force the transition of the state machine to the reset state.

A finite state machine can enter an illegal state—meaning the state registers contain a value that does not correspond to any defined state. By default, the behavior of the state machine that enters an illegal state is undefined. However, you can set the *syn_encoding* attribute to *safe* or use the **Safe State Machine** logic option if you want the state machine to recover deterministically from an illegal state. The software inserts extra logic to detect an illegal state, and forces the transition of the state machine to the reset state. You can use this logic option when the state machine enters an illegal state. The most common cause of an illegal state is a state machine that has control inputs that come from another clock domain, such as the control logic for a clock-crossing FIFO, because the state machine must have inputs from another clock domain. This option protects only state machines (and not other registers) by forcing them into the reset state. You can use this option if your design has asynchronous inputs. However, Altera recommends using a synchronization register chain instead of relying on the safe state machine option.
The safe state machine value does not use any user-defined default logic from your HDL code that corresponds to unreachable states. Verilog HDL and VHDL enable you to specify a behavior for all states in the state machine explicitly, including unreachable states. However, synthesis tools detect if state machine logic is unreachable and minimize or remove the logic. Synthesis tools also remove any flag signals or logic that indicate such an illegal state. If the software implements the state machine as safe, the recovery logic added by Quartus II Integrated Synthesis forces its transition from an illegal state to the reset state.

You can set the Safe State Machine logic option globally, or on individual state machines. To set this logic option, on the Analysis & Synthesis Settings page, select More Settings. In the Existing option settings list, select Safe State Machine, and turn on this option in the Setting list.

You can set the syn_encoding safe attribute on a state machine in HDL, as shown in Example 17–39 through Example 17–41.


```verilog
reg [2:0] my_fsm /* synthesis syn_encoding = "safe" */;
```


```verilog
(* syn_encoding = "safe" *) reg [2:0] my_fsm;
```

**Example 17–41. VHDL Code: a Safe State Machine Attribute**

```vhdl
ATTRIBUTE syn_encoding OF my_fsm : TYPE IS "safe";
```

If you specify an encoding style (refer to “Manually Specifying State Assignments Using the syn_encoding Attribute” on page 17–36), separate the encoding style value in the quotation marks with the safe value with a comma, as follows: "safe, one-hot" or "safe, gray".

Safe state machine implementation can result in a noticeable area increase for your design. Therefore, Altera recommends that you set this option only on the critical state machines in your design in which the safe mode is necessary, such as a state machine that uses inputs from asynchronous clock domains. You may not need to use this option if you correctly synchronize inputs coming from other clock domains.

If you create the safe state machine assignment on an instance that the software fails to recognize as a state machine, or an entity that contains a state machine, the software takes no action. You must restructure the code, so that the software recognizes and infers the instance as a state machine.

For more information about the Safe State Machine logic option, refer to Safe State Machine logic option in Quartus II Help.

For guidelines to ensure that the software correctly infers your state machine, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
Power-Up Level

This logic option causes a register (flipflop) to power up with the specified logic level, either high (1) or low (0). The registers in the core hardware power up to 0 in all Altera devices. For the register to power up with a logic level high, the Compiler performs an optimization referred to as NOT-gate push back on the register. NOT-gate push back adds an inverter to the input and the output of the register, so that the reset and power-up conditions appear to be high and the device operates as expected. The register itself still powers up to 0, but the register output inverts so the signal arriving at all destinations is 1.

The Power-Up Level option supports wildcard characters, and you can apply this option to any register, registered logic cell WYSIWYG primitive, or to a design entity containing registers, if you want to set the power level for all registers in your design entity. If you assign this option to a registered logic cell WYSIWYG primitive, such as an atom primitive from a third-party synthesis tool, you must turn on the Perform WYSIWYG Primitive Resynthesis logic option for the option to take effect. You can also apply the option to a pin with the logic configurations described in the following list:

- If you turn on this option for an input pin, the option transfers to the register that the pin drives, if all these conditions are present:
  - No logic, other than inversion, between the pin and the register.
  - The input pin drives the data input of the register.
  - The input pin does not fan-out to any other logic.

- If you turn on this option for an output or bidirectional pin, the option transfers to the register that feeds the pin, if all these conditions are present:
  - No logic, other than inversion, between the register and the pin.
  - The register does not fan out to any other logic.

For more information about the Power-Up Level logic option, including information on the supported device families, refer to Power-Up Level logic option in Quartus II Help.

Inferred Power-Up Levels

Quartus II Integrated Synthesis reads default values for registered signals defined in Verilog HDL and VHDL code, and converts the default values into Power-Up Level settings. The software also synthesizes variables with assigned values in Verilog HDL initial blocks into power-up conditions. Synthesis of these default and initial constructs allows synthesized behavior of your design to match, as closely as possible, the power-up state of the HDL code during a functional simulation.
The following register declarations all set a power-up level of $V_{CC}$ or a logic value “1”, as shown in Example 17–42:

**Example 17–42.**

```vhdl
signal q : std_logic = '1'; -- power-up to VCC
reg q = 1'b1; // power-up to VCC
reg q;
initial begin q = 1'b1; end // power-up to VCC
```

For more information about NOT-gate push back, the power-up states for Altera devices, and how set and reset control signals affect the power-up level, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

**Power-Up Don’t Care**

This logic option allows the Compiler to optimize registers in your design that do not have a defined power-up condition.

For example, your design might have a register with its D input tied to $V_{CC}$, and with no clear signal or other secondary signals. If you turn on this option, the Compiler can choose for the register to power up to $V_{CC}$. Therefore, the output of the register is always $V_{CC}$. The Compiler can remove the register and connect its output to $V_{CC}$. If you turn this option off or if you set a **Power-Up Level** assignment of **Low** for this register, the register transitions from GND to $V_{CC}$ when your design starts up on the first clock signal. Thus, the register is at $V_{CC}$ and you cannot remove the register.

Similarly, if the register has a clear signal, the Compiler cannot remove the register because after asserting the clear signal, the register transitions again to GND and back to $V_{CC}$.

If the Compiler performs a **Power-Up Don’t Care** optimization that allows it to remove a register, it issues a message to indicate that it is doing so.

This project-wide option does not apply to registers that have the **Power-Up Level** logic option set to either **High** or **Low**.

For more information about **Power-Up Don’t Care** logic option and a list of supported devices, refer to **Power-Up Don’t Care logic option** in Quartus II Help.

**Remove Duplicate Registers**

The **Remove Duplicate Registers** logic option removes registers that are identical to other registers.

For more information about **Remove Duplicate Registers** logic option and the supported devices, refer to **Remove Duplicate Registers logic option** in Quartus II Help.
Preserve Registers

This attribute and logic option directs the Compiler not to minimize or remove a specified register during synthesis optimizations or register netlist optimizations. Optimizations can eliminate redundant registers and registers with constant drivers; this option prevents the software from reducing a register to a constant or merging with a duplicate register. This option can preserve a register so you can observe the register during simulation or with the SignalTap® II Logic Analyzer. Additionally, this option can preserve registers if you create a preliminary version of your design in which you have not specified the secondary signals. You can also use the attribute to preserve a duplicate of an I/O register so that you can place one copy of the I/O register in an I/O cell and the second in the core.

This option cannot preserve registers that have no fan-out. To prevent the removal of registers with no fan-out, refer to “Noprune Synthesis Attribute/Preserve Fan-out Free Register Node” on page 17–43.

The Preserve Registers logic option prevents the software from inferring a register as a state machine.

You can set the Preserve Registers logic option in the Quartus II software, or you can set the preserve attribute in your HDL code, as shown in Example 17–43 through Example 17–45. In these examples, the Quartus II software preserves the my_reg register.

In addition to preserve, the Quartus II software supports the syn_preserve attribute name for compatibility with other synthesis tools.

Example 17–43. Verilog HDL Code: syn_preserve Attribute

```
reg my_reg /* synthesis syn_preserve = 1 */;
```

Example 17–44. Verilog-2001 Code: syn_preserve Attribute

```
(* syn_preserve = 1 *) reg my_reg;
```

The = 1 after the preserve in Example 17–43 and Example 17–44 is optional, because the assignment uses a default value of 1 when you specify the assignment.

Example 17–45. VHDL Code: preserve Attribute

```
signal my_reg : stdlogic;
attribute preserve : boolean;
attribute preserve of my_reg : signal is true;
```

For more information about the Preserve Registers logic option and the supported devices, refer to Preserve Registers logic option in Quartus II Help.
Disable Register Merging/Don’t Merge Register

This logic option and attribute prevents the specified register from merging with other registers and prevents other registers from merging with the specified register. When applied to a design entity, it applies to all registers in the entity.

You can set the Disable Register Merging logic option in the Quartus II software, or you can set the dont_merge attribute in your HDL code, as shown in Example 17–46 through Example 17–48. In these examples, the logic option or the attribute prevents the my_reg register from merging.

Example 17–46. Verilog HDL Code: dont_merge Attribute

```vhdl
reg my_reg /* synthesis dont_merge */;
```

Example 17–47. Verilog-2001 and SystemVerilog Code: dont_merge Attribute

```vhdl
(* dont_merge *) reg my_reg;
```

Example 17–48. VHDL Code: dont_merge Attribute

```vhdl
signal my_reg : stdlogic;
attribute dont_merge : boolean;
attribute dont_merge of my_reg : signal is true;
```

Noprune Synthesis Attribute/Preserve Fan-out Free Register Node

This synthesis attribute and corresponding logic option direct the Compiler to preserve a fan-out-free register through the entire compilation flow. This option is different from the Preserve Registers option, which prevents the Quartus II software from reducing a register to a constant or merging with a duplicate register. Standard synthesis optimizations remove nodes that do not directly or indirectly feed a top-level output pin. This option can retain a register so you can observe the register in the Simulator or the SignalTap II Logic Analyzer. Additionally, this option can retain registers if you create a preliminary version of your design in which you have not specified the fan-out logic of the register.

You can set the Preserve Fan-out Free Register Node logic option in the Quartus II software, or you can set the noprune attribute in your HDL code, as shown in Example 17–49 though Example 17–51. In these examples, the logic option or the attribute preserves the my_reg register.

⚠️ You must use the noprune attribute instead of the logic option if the register has no immediate fan-out in its module or entity. If you do not use the synthesis attribute, the software removes (or “prunes”) registers with no fan-out during Analysis & Elaboration before the logic synthesis stage applies any logic options. If the register has no fan-out in the full design, but has fan-out in its module or entity, you can use the logic option to retain the register through compilation.
The software supports the attribute name `syn_noprune` for compatibility with other synthesis tools.

**Example 17–49. Verilog HDL Code: syn_noprune Attribute**

```verilog
reg my_reg /* synthesis syn_noprune */;
```


```verilog
(* noprune *) reg my_reg;
```

**Example 17–51. VHDL Code: noprune Attribute**

```vhdl
signal my_reg : stdlogic;
attribute noprune : boolean;
attribute noprune of my_reg : signal is true;
```

For more information about **Preserve Fan-out Free Register Node** logic option and a list of supported devices, refer to **Preserve Fan-out Free Register logic option** in Quartus II Help.

### Keep Combinational Node/Implement as Output of Logic Cell

This synthesis attribute and corresponding logic option direct the Compiler to keep a wire or combinational node through logic synthesis minimizations and netlist optimizations. A wire that has a `keep` attribute or a node that has the **Implement as Output of Logic Cell** logic option applied becomes the output of a logic cell in the final synthesis netlist, and the name of the logic cell remains the same as the name of the wire or node. You can use this directive to make combinational nodes visible to the SignalTap II Logic Analyzer.

The option cannot keep nodes that have no fan-out. You cannot maintain node names for wires with tri-state drivers, or if the signal feeds a top-level pin of the same name (the software changes the node name to a name such as `<net name>-buf0`).

You can use the **Ignore LCELL Buffers** logic option to direct Analysis & Synthesis to ignore logic cell buffers that the **Implement as Output of Logic Cell** logic option or the **LCELL** primitive created. If you apply this logic option to an entity, it affects all lower-level entities in the hierarchy path.

To avoid unintended design optimizations, ensure that any entity instantiated with Altera or third-party IP that relies on logic cell buffers for correct behavior does not inherit the **Ignore LCELL Buffers** logic option. For example, if an IP core uses logic cell buffers to manage high fan-out signals and inherits the **Ignore LCELL Buffers** logic option, the target device may no longer function properly.

You can turn off the **Ignore LCELL Buffers** logic option for a specific entity to override any assignments inherited from higher-level entities in the hierarchy path if logic cell buffers created by the **Implement as Output of Logic Cell** logic option or the **LCELL** primitive are required for correct behavior.
You can set the **Implement as Output of Logic Cell** logic option in the Quartus II software, or you can set the `keep` attribute in your HDL code, as shown in *Example 17–52* through *Example 17–54*. In these examples, the Compiler maintains the node name `my_wire`.

In addition to `keep`, the Quartus II software supports the `syn_keep` attribute name for compatibility with other synthesis tools.

**Example 17–52. Verilog HDL Code: keep Attribute**

```verilog
wire my_wire /* synthesis keep = 1 */;
```

**Example 17–53. Verilog-2001 Code: keep Attribute**

```verilog
(* keep = 1 *) wire my_wire;
```

**Example 17–54. VHDL Code: syn_keep Attribute**

```vhdl
signal my_wire: bit;
attribute syn_keep: boolean;
attribute syn_keep of my_wire: signal is true;
```

For more information about the **Implement as Output of Logic Cell** logic option and the supported devices, refer to *Implement as Output of Logic Cell logic option* in Quartus II Help.

### Disabling Synthesis Netlist Optimizations with `dont_retime` Attribute

This attribute disables synthesis retiming optimizations on the register you specify. When applied to a design entity, it applies to all registers in the entity.

You can turn off retiming optimizations with this option and prevent node name changes, so that the Compiler can correctly use your timing constraints for the register.
You can set the **Netlist Optimizations** logic option to **Never Allow** in the Quartus II software to disable retiming along with other synthesis netlist optimizations, or you can set the `dont_retime` attribute in your HDL code, as shown in Example 17–55 through Example 17–57. In these examples, the code prevents `my_reg` register from being retimed.

**Example 17–55. Verilog HDL Code: dont_retime Attribute**

```verilog
reg my_reg /* synthesis dont_retime */;
```


```verilog
(* dont_retime *) reg my_reg;
```

**Example 17–57. VHDL Code: dont_retime Attribute**

```vhdl
signal my_reg : std_logic;
attribute dont_retime : boolean;
attribute dont_retime of my_reg : signal is true;
```

For compatibility with third-party synthesis tools, Quartus II Integrated Synthesis also supports the attribute `syn_allow_retiming`. To disable retiming, set `syn_allow_retiming` to 0 (Verilog HDL) or false (VHDL). This attribute does not have any effect when you set the attribute to 1 or true.

**Disabling Synthesis Netlist Optimizations with dont_replicate Attribute**

This attribute disables synthesis replication optimizations on the register you specify. When applied to a design entity, it applies to all registers in the entity.

You can turn off register replication (or duplication) optimizations with this option, so that the Compiler uses your timing constraints for the register.

You can set the **Netlist Optimizations** logic option to **Never Allow** in the Quartus II software to disable replication along with other synthesis netlist optimizations, or you can set the `dont_replicate` attribute in your HDL code, as shown in Example 17–58 through Example 17–60. In these examples, the code prevents the replication of the `my_reg` register.

**Example 17–58. Verilog HDL Code: dont_replicate Attribute**

```verilog
reg my_reg /* synthesis dont_replicate */;
```


```verilog
(* dont_replicate *) reg my_reg;
```

**Example 17–60. VHDL Code: dont_replicate Attribute**

```vhdl
signal my_reg : std_logic;
attribute dont_replicate : boolean;
attribute dont_replicate of my_reg : signal is true;
```
For compatibility with third-party synthesis tools, Quartus II Integrated Synthesis also supports the attribute `syn_replicate`. To disable replication, set `syn_replicate` to 0 (Verilog HDL) or false (VHDL). This attribute does not have any effect when you set the attribute to 1 or true.

**Maximum Fan-Out**

This Maximum Fan-Out attribute and logic option direct the Compiler to control the number of destinations that a node feeds. The Compiler duplicates a node and splits its fan-out until the individual fan-out of each copy falls below the maximum fan-out restriction. You can apply this option to a register or a logic cell buffer, or to a design entity that contains these elements. You can use this option to reduce the load of critical signals, which can improve performance. You can use the option to instruct the Compiler to duplicate a register that feeds nodes in different locations on the target device. Duplicating the register can enable the Fitter to place these new registers closer to their destination logic to minimize routing delay.

To turn off the option for a given node if you set the option at a higher level of the design hierarchy, in the Netlist Optimizations logic option, select Never Allow. If not disabled by the Netlist Optimizations option, the Compiler acknowledges the maximum fan-out constraint as long as the following conditions are met:

- The node is not part of a cascade, carry, or register cascade chain.
- The node does not feed itself.
- The node feeds other logic cells, DSP blocks, RAM blocks, and pins through data, address, clock enable, and other ports, but not through any asynchronous control ports (such as asynchronous clear).

The Compiler does not create duplicate nodes in these cases, because there is no clear way to duplicate the node, or to avoid the small differences in timing which could produce functional differences in the implementation (in the third condition above in which asynchronous control signals are involved). If you cannot apply the constraint because you do not meet one of these conditions, the Compiler issues a message to indicate that the Compiler ignores the maximum fan-out assignment. To instruct the Compiler not to check node destinations for possible problems such as the third condition, you can set the Netlist Optimizations logic option to Always Allow for a given node.

If you have enabled any of the Quartus II netlist optimizations that affect registers, add the preserve attribute to any registers to which you have set a maxfan attribute. The preserve attribute ensures that the netlist optimization algorithms, such as register retiming, do not affect the registers.

For details about netlist optimizations, refer to the Netlist Optimizations and Physical Synthesis chapter in volume 2 of the Quartus II Handbook.

You can set the Maximum Fan-Out logic option in the Quartus II software. This option supports wildcard characters. You can also set the maxfan attribute in your HDL code, as shown in Example 17–61 through Example 17–63. In these examples, the Compiler duplicates the clk_gen register, so its fan-out is not greater than 50.
In addition to \texttt{maxfan}, the Quartus II software supports the \texttt{syn\_maxfan} attribute for compatibility with other synthesis tools.

**Example 17–61. Verilog HDL Code: \texttt{syn\_maxfan} Attribute**

\begin{verbatim}
reg clk_gen /* synthesis syn\_maxfan = 50 */;
\end{verbatim}

**Example 17–62. Verilog-2001 Code: \texttt{maxfan} Attribute**

\begin{verbatim}
(* maxfan = 50 *) reg clk_gen;
\end{verbatim}

**Example 17–63. VHDL Code: \texttt{maxfan} Attribute**

\begin{verbatim}
signal clk_gen : stdlogic;
attribute maxfan : signal ;
attribute maxfan of clk_gen : signal is 50;
\end{verbatim}

For more information about the \texttt{Maximum Fan-Out} logic option and the supported devices, refer to \textit{Maximum Fan-Out logic option} in Quartus II Help.

**Controlling Clock Enable Signals with Auto Clock Enable Replacement and \texttt{direct\_enable}**

The \texttt{Auto Clock Enable Replacement} logic option allows the software to find logic that feeds a register and move the logic to the register’s clock enable input port. To solve fitting or performance issues with designs that have many clock enables, you can turn off this option for individual registers or design entities. Turning the option off prevents the software from using the register’s clock enable port. The software implements the clock enable functionality using multiplexers in logic cells.

If the software does not move the specific logic to a clock enable input with the \texttt{Auto Clock Enable Replacement} logic option, you can instruct the software to use a direct clock enable signal. The attribute ensures that the signal drives the clock enable port, and the software does not optimize or combine the signal with other logic.

**Example 17–64 through Example 17–66** show how to set this attribute to ensure that the attribute preserves the signal and uses the signal as a clock enable.
In addition to `direct_enable`, the Quartus II software supports the `syn_direct_enable` attribute name for compatibility with other synthesis tools.

**Example 17–64. Verilog HDL Code: direct_enable attribute**

```verilog
wire my_enable /* synthesis direct_enable = 1 */ ;
```


```verilog
(* syn_direct_enable *) wire my_enable;
```

**Example 17–66. VHDL Code: direct_enable attribute**

```vhdl
attribute direct_enable: boolean;
attribute direct_enable of my_enable: signal is true;
```

For more information about the Auto Clock Enable Replacement logic option and the supported devices, refer to Auto Clock Enable Replacement logic option in Quartus II Help.

### Inferring Multiplier, DSP, and Memory Functions from HDL Code

The Quartus II Compiler automatically recognizes multipliers, multiply-accumulators, multiply-adders, or memory functions described in HDL code, and either converts the HDL code into respective megafuntion or maps them directly to device atoms or memory atoms. If the software converts the HDL code into a megafuntion, the software uses the Altera megafuntion code when you compile your design, even when you do not specifically instantiate the megafuntion. The software infers megafuntions to take advantage of logic that you optimize for Altera devices. The area and performance of such logic can be better than the results from inferring generic logic from the same HDL code.

Additionally, you must use megafuntions to access certain architecture-specific features, such as RAM, DSP blocks, and shift registers that provide improved performance compared with basic logic cells.

For details about coding style recommendations when targeting megafuntions in Altera devices, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

The Quartus II software provides options to control the inference of certain types of megafuntions, as described in the following subsections:

- “Multiply-Accumulators and Multiply-Adders”
- “Shift Registers” on page 17–50
- “RAM and ROM” on page 17–51
- “Resource Aware RAM, ROM, and Shift-Register Inference” on page 17–51
- “Auto RAM to Logic Cell Conversion” on page 17–52
Multiply-Accumulators and Multiply-Adders

Use the **Auto DSP Block Replacement** logic option to control DSP block inference for multiply-accumulations and multiply-adders. To disable inference, turn off this option for the entire project on the **Analysis & Synthesis Settings** page of the **Settings** dialog box, or turn off the option for a specific block with the Assignment Editor. By default, the software enables this logic option for Stratix V devices.

For more information about the **Auto DSP Block Replacement** logic option and the supported devices, refer to [Auto DSP Block Replacement logic option](#) in Quartus II Help.

Shift Registers

Use the **Auto Shift Register Replacement** logic option to control shift register inference. This option has three settings: **Off**, **Auto** and **Always**. **Auto** is the default setting in which Quartus II Integrated Synthesis decides which shift registers to replace or leave in registers. Placing shift registers in memory saves logic area, but can have a negative effect on \( f_{\text{max}} \). Quartus II Integrated Synthesis uses the optimization technique setting, logic and RAM utilization of your design, and timing information from **Timing-Driven Synthesis** to determine which shift registers are located in memory and which are located in registers. To disable inference, turn off this option for the entire project on the **Analysis & Synthesis Settings** page of the **Settings** dialog box by clicking **More Settings** and setting the option to **Off**. You can also disable the option for a specific block with the Assignment Editor. Even if you set the logic option to **On** or **Auto**, the software might not infer small shift registers because small shift registers do not benefit from implementation in dedicated memory. However, you can use the **Allow Any Shift Register Size for Recognition** logic option to instruct synthesis to infer a shift register even when its size is too small.

You can use the **Allow Shift Register Merging across Hierarchies** option to prevent the Compiler from merging shift registers in different hierarchies into one larger shift register. The option has three settings: **On**, **Off**, and **Auto**. The **Auto** setting is the default setting, and the Compiler decides whether or not to merge shift registers across hierarchies. When you turn on this option, the Compiler allows all shift registers to merge across hierarchies, and when you turn off this option, the Compiler does not allow any shift registers to merge across hierarchies. You can set this option globally or on entities or individual nodes.

The registers that the software maps to the ALTSHIFT_TAPS megafunction and places in RAM are not available in the Simulator because their node names do not exist after synthesis.

The Compiler turns off the **Auto Shift Register Replacement** logic option when you select a formal verification tool on the **EDA Tool Settings** page. If you do not select a formal verification tool, the Compiler issues a warning and the compilation report lists shift registers that the logic option might infer. To enable a megafunction for the shift register in the formal verification flow, you can either instantiate a shift register explicitly with the MegaWizard™ Plug-In Manager or make the shift register into a black box in a separate entity or module.

For more information about the **Auto Shift Register Replacement** logic option and the supported devices, refer to [Auto Shift Register Replacement logic option](#) in Quartus II Help.
RAM and ROM

Use the **Auto RAM Replacement** and **Auto ROM Replacement** logic options to control RAM and ROM inference, respectively. To disable inference, turn off the appropriate option for the entire project on the **Analysis & Synthesis Settings** page of the **Settings** dialog box by clicking **More Settings** and setting the option to **Off**. You can also disable the option for a specific block with the Assignment Editor.

Although the software implements inferred shift registers in RAM blocks, you cannot turn off the **Auto RAM Replacement** option to disable shift register replacement. Use the **Auto Shift Register Replacement** option (refer to “Shift Registers” on page 17–50).

The software might not infer very small RAM or ROM blocks because you can implement very small memory blocks with the registers in the logic. However, you can use the **Allow Any RAM Size for Recognition** and **Allow Any ROM Size for Recognition** logic options to instruct synthesis to infer a memory block even when its size is too small.

The software turns off the **Auto ROM Replacement** logic option when you select a formal verification tool in the **EDA Tool Settings** page. If you do not select a formal verification tool, the software issues a warning and a report panel provides a list of ROMs that the logic option might infer. To enable a megafunction for the shift register in the formal verification flow, you can either instantiate a ROM explicitly using the MegaWizard Plug-In Manager or create a black box for the ROM in a separate entity or in a separate module.

Although formal verification tools do not support inferred RAM blocks, due to the importance of inferring RAM in many designs, the software turns on the **Auto RAM Replacement** logic option when you select a formal verification tool in the **EDA Tool Settings** page. The software automatically performs black box instance for any module or entity that contains an inferred RAM block. The software issues a warning and lists the black box created in the compilation report. This black box allows formal verification tools to proceed; however, the formal verification tool cannot verify the entire module or entire entity that contains the RAM. Altera recommends that you explicitly instantiate RAM blocks in separate modules or in separate entities so that the formal verification tool can verify as much logic as possible.

For more information about the **Auto RAM Replacement** and **Auto ROM Replacement** logic options and their supported devices, refer to **Auto RAM Replacement logic option** and **Auto ROM Replacement logic option** in Quartus II Help.

Resource Aware RAM, ROM, and Shift-Register Inference

The Quartus II Integrated Synthesis considers resource usage when inferring RAM, ROM, and shift registers. During RAM, ROM, and shift register inferencing, synthesis looks at the number of memories available in the current device and does not infer more memory than is available to avoid a no-fit error. Synthesis tries to select the memories that are not inferred in a way that aims at the smallest increase in logic and registers.
Resource aware RAM, ROM and shift register inference is controlled by the Resource Aware Inference for Block RAM option. You can disable this option for the entire project in the More Analysis & Synthesis Settings dialog box, or per partition in the Assignment Editor.

When you select the Auto setting, resource aware RAM, ROM, and shift register inference use the resource counts from the largest device.

For designs with multiple partitions, Quartus II Integrated Synthesis considers one partition at a time. Therefore, for each partition, it assumes that all RAM blocks are available to that partition. If this causes a no-fit error, you can limit the number of RAM blocks available per partition with the Maximum Number of M512 Memory Blocks, Maximum Number of M4K/M9K/M20K/M10K Memory Blocks, Maximum Number of M-RAM/M144K Memory Blocks and Maximum Number of LABs settings in the Assignment Editor. The balancer also uses these options. For more information, refer to “Limiting Resource Usage in Partitions” on page 17–32.

**Auto RAM to Logic Cell Conversion**

The Auto RAM to Logic Cell Conversion logic option allows Quartus II Integrated Synthesis to convert small RAM blocks to logic cells if the logic cell implementation gives better quality of results. The software converts only single-port or simple-dual port RAMs with no initialization files to logic cells. You can set this option globally or apply it to individual RAM nodes. You can enable this option by turning on the appropriate option for the entire project in the More Analysis & Synthesis Settings dialog box.

For Arria GX and Stratix family of devices, the software uses the following rules to determine the placement of a RAM, either in logic cells or a dedicated RAM block:

- If the number of words is less than 16, use a RAM block if the total number of bits is greater than or equal to 64.
- If the number of words is greater than or equal to 16, use a RAM block if the total number of bits is greater than or equal to 32.
- Otherwise, implement the RAM in logic cells.

For the Cyclone family of devices, the software uses the following rules:

- If the number of words is greater than or equal to 64, use a RAM block.
- If the number of words is greater than or equal to 16 and less than 64, use a RAM block if the total number of bits is greater than or equal to 128.
- Otherwise, implement the RAM in logic cells.

For more information about the Auto RAM to Logic Cell Conversion logic options and the supported devices, refer to Auto RAM to Logic Cell Conversion logic option in Quartus II Help.
**RAM Style and ROM Style—for Inferred Memory**

These attributes specify the implementation for an inferred RAM or ROM block. You can specify the type of TriMatrix embedded memory block, or specify the use of standard logic cells (LEs or ALMs). The Quartus II software supports the attributes only for device families with TriMatrix embedded memory blocks.

The `ramstyle` and `romstyle` attributes take a single string value. The `M512`, `M4K`, `M-RAM`, `MLAB`, `M9K`, `M144K`, `M20K`, and `M10K` values (as applicable for the target device family) indicate the type of memory block to use for the inferred RAM or ROM. If you set the attribute to a block type that does not exist in the target device family, the software generates a warning and ignores the assignment. The `logic` value indicates that the Quartus II software implements the RAM or ROM in regular logic rather than dedicated memory blocks. You can set the attribute on a module or entity, in which case it specifies the default implementation style for all inferred memory blocks in the immediate hierarchy. You can also set the attribute on a specific signal (VHDL) or variable (Verilog HDL) declaration, in which case it specifies the preferred implementation style for that specific memory, overriding the default implementation style.

If you specify a `logic` value, the memory appears as a RAM or ROM block in the RTL Viewer, but Integrated Synthesis converts the memory to regular logic during synthesis.

In addition to `ramstyle` and `romstyle`, the Quartus II software supports the `syn_ramstyle` attribute name for compatibility with other synthesis tools.

Example 17–67 through Example 17–69 specify that you must implement all memory in the module or the `my_memory_blocks` entity with a specific type of block.

**Example 17–67. Verilog-1995 Code: Applying a romstyle Attribute to a Module Declaration**

```vhdl
module my_memory_blocks (...) /* synthesis romstyle = "M4K" */;
```

**Example 17–68. Verilog-2001 and SystemVerilog Code: Applying a ramstyle Attribute to a Module Declaration**

```vhdl
(* ramstyle = "M512" *) module my_memory_blocks (...);
```

**Example 17–69. VHDL Code: Applying a romstyle Attribute to an Architecture**

```vhdl
architecture rtl of my_ my_memory_blocks is
attribute romstyle : string;
attribute romstyle of rtl : architecture is "M-RAM";
begin
```
Example 17–70 through Example 17–72 specify that you must implement the inferred my_ram or my_rom memory with regular logic instead of a TriMatrix memory block.

**Example 17–70. Verilog-1995 Code: Applying a syn_ramstyle Attribute to a Variable Declaration**

```verilog
reg [0:7] my_ram[0:63] /* synthesis syn_ramstyle = "logic" */;
```

**Example 17–71. Verilog-2001 and SystemVerilog Code: Applying a romstyle Attribute to a Variable Declaration**

```verilog
(* romstyle = "logic" *) reg [0:7] my_rom[0:63];
```

**Example 17–72. VHDL Code: Applying a ramstyle Attribute to a Signal Declaration**

```vhdl
type memory_t is array (0 to 63) of std_logic_vector (0 to 7);
signal my_ram : memory_t;
attribute ramstyle : string;
attribute ramstyle of my_ram : signal is "logic";
```

You can control the depth of an inferred memory block and optimize its usage with the `max_depth` attribute. You can also optimize the usage of the memory block with this attribute. Example 17–73 through Example 17–75 specify the depth of the inferred memory `mem` using the `max_depth` synthesis attribute.

**Example 17–73. Verilog-1995 Code: Applying a max_depth Attribute to a Variable Declaration**

```verilog
reg [7:0] mem [127:0] /* synthesis max_depth = 2048 */
```

**Example 17–74. Verilog-2001 and SystemVerilog Code: Applying a max_depth Attribute to a Variable Declaration**

```verilog
(* max_depth = 2048*) reg [7:0] mem [127:0];
```

**Example 17–75. VHDL Code: Applying a max_depth Attribute to a Variable Declaration**

```vhdl
type ram_block is array (0 to 31) of std_logic_vector (2 downto 0);
signal mem : ram_block;
attribute max_depth : natural;
attribute max_depth OF mem : signal is 2048;
```

The syntax for setting these attributes in HDL is the same as the syntax for other synthesis attributes, as shown in “Synthesis Attributes” on page 17–25.
RAM Style Attribute—For Shift Registers Inference

The RAM style attribute for shift register allows you to use the RAM style attribute for shift registers, just as you use them for RAM or ROMs. The Quartus II Synthesis uses the RAM style attribute during shift register inference. If synthesis infers the shift register to RAM, it will be sent to the requested RAM block type. Shift registers are merged only if the RAM style attributes are compatible. If the RAM style is set to logic, a shift register does not get inferred to RAM.

Example 17–76 shows the code to set the RAM style attribute for shift registers in Verilog.

**Example 17–76. Verilog Code: Setting the RAM Style Attribute for Shift Registers**

(* ramstyle = "mlab" *) reg [N-1:0] sr;

Example 17–77 shows the code to set the RAM style attribute for shift registers in VHDL.

**Example 17–77. VHDL Code: Setting the RAM Style Attribute for Shift Registers**

attribute ramstyle : string;attribute ramstyle of sr : signal is "M20K";

You can also assign the RAM style attribute for shift registers globally, which will affect all shift registers.
Turning Off the Add Pass-Through Logic to Inferred RAMs no_rw_check Attribute

Setting the no_rw_check value for the ramstyle attribute, or turning off the corresponding global Add Pass-Through Logic to Inferred RAMs logic option indicates that your design does not depend on the behavior of the inferred RAM when there are reads and writes to the same address in the same clock cycle. If you specify the attribute or turn off the logic option, the Quartus II software can choose a read-during-write behavior instead of using the read-during-write behavior of your HDL source code.

Sometimes, you must map an inferred RAM into regular logic cells because the inferred RAM has a read-during-write behavior that the TriMatrix memory blocks in your target device do not support. In other cases, the Quartus II software must insert extra logic to mimic read-during-write behavior of the HDL source to increase the area of your design and potentially reduce its performance. In some of these cases, you can use the attribute to specify that the software can implement the RAM directly in a TriMatrix memory block without using logic. You can also use the attribute to prevent a warning message for dual-clock RAMs in the case that the inferred behavior in the device does not exactly match the read-during-write conditions described in the HDL code.

For more information about recommended styles for inferring RAM and some of the issues involved with different read-during-write conditions, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
To set the Add Pass-Through Logic to Inferred RAMs logic option with the Quartus II software, click More Settings on the Analysis & Synthesis Settings page of the Settings dialog box. Example 17–78 and Example 17–79 use two addresses and normally require extra logic after the RAM to ensure that the read-during-write conditions in the device match the HDL code. If your design does not require a defined read-during-write condition, the extra logic is not necessary. With the no_rw_check attribute, Quartus II Integrated Synthesis does not generate the extra logic.

Example 17–78. Verilog HDL Inferred RAM Using no_rw_check Attribute

```verilog
module ram_infer (q, wa, ra, d, we, clk);
    output [7:0] q;
    input [7:0] d;
    input [6:0] wa;
    input [6:0] ra;
    input we, clk;
    reg [6:0] read_add;
    (* ramstyle = "no_rw_check" *) reg [7:0] mem [127:0];
    always @(posedge clk) begin
        if (we)
            mem[wa] <= d;
        read_add <= ra;
    end
    assign q = mem[read_add];
endmodule
```

Example 17–79. VHDL Inferred RAM Using no_rw_check Attribute

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.ALL;

ENTITY ram IS
    PORT (
        clock: IN STD_LOGIC;
        data: IN STD_LOGIC_VECTOR (2 DOWNTO 0);
        write_address: IN INTEGER RANGE 0 to 31;
        read_address: IN INTEGER RANGE 0 to 31;
        we: IN STD_LOGIC;
        q: OUT STD_LOGIC_VECTOR (2 DOWNTO 0));
END ram;

ARCHITECTURE rtl OF ram IS
    TYPE MEM IS ARRAY(0 TO 31) OF STD_LOGIC_VECTOR(2 DOWNTO 0);
    SIGNAL ram_block: MEM;
    ATTRIBUTE ramstyle : string;
    ATTRIBUTE ramstyle of ram_block : signal is "no_rw_check";
    SIGNAL read_address_reg: INTEGER RANGE 0 to 31;
BEGIN
    PROCESS (clock)
    BEGIN
        IF (clock'event AND clock = '1') THEN
            IF (we = '1') THEN
                ram_block(write_address) <= data;
            END IF;
            read_address_reg <= read_address;
        END IF;
    END PROCESS;
    q <= ram_block(read_address_reg);
END rtl;
```
You can use a ramstyle attribute with the MLAB value, so that the Quartus II software can infer a small RAM block and place it in an MLAB.

You can use this attribute in cases in which some asynchronous RAM blocks might be coded with read-during-write behavior that does not match the Stratix III, Stratix IV, and Stratix V architectures. Thus, the device behavior would not exactly match the behavior that the code describes. If the difference in behavior is acceptable in your design, use the ramstyle attribute with the no_rw_check value to specify that the software should not check the read-during-write behavior when inferring the RAM. When you set this attribute, Quartus II Integrated Synthesis allows the behavior of the output to differ when the asynchronous read occurs on an address that had a write on the most recent clock edge. That is, the functional HDL simulation results do not match the hardware behavior if you write to an address that is being read. To include these attributes, set the value of the ramstyle attribute to `MLAB, no_rw_check`.

Example 17–80 and Example 17–81 show the method of setting two values to the ramstyle attribute with a small asynchronous RAM block, with the ramstyle synthesis attribute set, so that the software can implement the memory in the MLAB memory block and so that the read-during-write behavior is not important. Without the attribute, this design requires 512 registers and 240 ALUTs. With the attribute, the design requires eight memory ALUTs and only 15 registers.

**Example 17–80. Verilog HDL Inferred RAM Using no_rw_check and MLAB Attributes**

```verilog
module async_ram (  
   input   [5:0] addr,  
   input   [7:0] data_in,  
   input         clk,  
   input         write,  
   output  [7:0] data_out );  

(* ramstyle = "MLAB, no_rw_check" *) reg [7:0] mem[0:63];  

assign data_out = mem[addr];  

always @(posedge clk)  
begin  
   if (write)  
      mem[addr] = data_in;  
end  
endmodule
```
Example 17–81. VHDL Inferred RAM Using no_rw_check and MLAB Attributes

LIBRARY ieee;
USE ieee.std_logic_1164.ALL;
ENTITY ram IS
  PORT (  
    clock: IN STD_LOGIC;
    data: IN STD_LOGIC_VECTOR (2 DOWNTO 0);
    write_address: IN INTEGER RANGE 0 to 31;
    read_address: IN INTEGER RANGE 0 to 31;
    we: IN STD_LOGIC;
    q: OUT STD_LOGIC_VECTOR (2 DOWNTO 0));
END ram;

ARCHITECTURE rtl OF ram IS  
  TYPE MEM IS ARRAY(0 TO 31) OF STD_LOGIC_VECTOR(2 DOWNTO 0);
  SIGNAL ram_block: MEM;
 ATTRIBUTE ramstyle : string;
 ATTRIBUTE ramstyle of ram_block : signal is "MLAB , no_rw_check";
  SIGNAL read_address_reg: INTEGER RANGE 0 to 31;
BEGIN  
  PROCESS (clock)
  BEGIN  
    IF (clock’event AND clock = '1') THEN  
      IF (we = '1') THEN  
        ram_block(write_address) <= data;
      END IF;
      read_address_reg <= read_address;
    END IF;
  END PROCESS;
  q <= ram_block(read_address_reg);
END rtl;

For more information about the Add Pass-Through Logic to Inferred RAMs logic option and the supported devices, refer to Add Pass-Through Logic to Inferred RAMs logic option in Quartus II Help.

RAM Initialization File—for Inferred Memory

The ram_init_file attribute specifies the initial contents of an inferred memory with a .mif. The attribute takes a string value containing the name of the RAM initialization file.

Example 17–82. Verilog-1995 Code: Applying a ram_init_file Attribute

reg [7:0] mem[0:255] /* synthesis ram_init_file = "my_init_file.mif" */;


(* ram_init_file = "my_init_file.mif" *) reg [7:0] mem[0:255];
In VHDL, you can also initialize the contents of an inferred memory by specifying a
default value for the corresponding signal. In Verilog HDL, you can use an initial
block to specify the memory contents. Quartus II Integrated Synthesis automatically
converts the default value into a .mif for the inferred RAM.

Multiplier Style—for Inferred Multipliers

The multstyle attribute specifies the implementation style for multiplication
operations (*) in your HDL source code. You can use this attribute to specify whether
you prefer the Compiler to implement a multiplication operation in general logic or
dedicated hardware, if available in the target device.

The multstyle attribute takes a string value of "logic" or "dsp", indicating a
preferred implementation in logic or in dedicated hardware, respectively. In Verilog
HDL, apply the attribute to a module declaration, a variable declaration, or a specific
binary expression that contains the * operator. In VHDL, apply the synthesis attribute
to a signal, variable, entity, or architecture.

Specifying a multstyle of "dsp" does not guarantee that the Quartus II software can
implement a multiplication in dedicated DSP hardware. The final implementation
depends on several conditions, including the availability of dedicated hardware in the
target device, the size of the operands, and whether or not one or both operands are
constant.

In addition to multstyle, the Quartus II software supports the syn_multstyle
attribute name for compatibility with other synthesis tools.

When applied to a Verilog HDL module declaration, the attribute specifies the default
implementation style for all instances of the * operator in the module. For example, in
the following code examples, the multstyle attribute directs the Quartus II software
to implement all multiplications inside module my_module in the dedicated
multiplication hardware.

Example 17–85. Verilog-1995 Code: Applying a multstyle Attribute to a Module Declaration

module my_module (...) /* synthesis multstyle = "dsp" */;

Example 17–86. Verilog-2001 Code: Applying a multstyle Attribute to a Module Declaration

(* multstyle = "dsp" *) module my_module(...);
When applied to a Verilog HDL variable declaration, the attribute specifies the implementation style for a multiplication operator, which has a result directly assigned to the variable. The attribute overrides the multstyle attribute with the enclosing module, if present. In Example 17–87 and Example 17–88, the multstyle attribute applied to variable result directs the Quartus II software to implement $a \times b$ in logic rather than the dedicated hardware.

**Example 17–87. Verilog-2001 Code: Applying a multstyle Attribute to a Variable Declaration**

```verilog
wire [8:0] a, b;
(* multstyle = "logic" *) wire [17:0] result;
assign result = a * b;  //Multiplication must be
//directly assigned to result
```

**Example 17–88. Verilog-1995 Code: Applying a multstyle Attribute to a Variable Declaration**

```verilog
wire [8:0] a, b;
wire [17:0] result /* synthesis multstyle = "logic" */;
assign result = a * b;  //Multiplication must be
//directly assigned to result
```

When applied directly to a binary expression that contains the $\times$ operator, the attribute specifies the implementation style for that specific operator alone and overrides any multstyle attribute with the target variable or enclosing module. In Example 17–89, the multstyle attribute indicates that you must implement $a \times b$ in the dedicated hardware.

**Example 17–89. Verilog-2001 Code: Applying a multstyle Attribute to a Binary Expression**

```verilog
wire [8:0] a, b;
wire [17:0] result;
assign result = a * (* multstyle = "dsp" *) b;
```

You cannot use Verilog-1995 attribute syntax to apply the multstyle attribute to a binary expression.

When applied to a VHDL entity or architecture, the attribute specifies the default implementation style for all instances of the $\times$ operator in the entity or architecture. In Example 17–90, the multstyle attribute directs the Quartus II software to use dedicated hardware, if possible, for all multiplications inside architecture rtl of entity my_entity.

**Example 17–90. VHDL Code: Applying a multstyle Attribute to an Architecture**

```vhdl
architecture rtl of my_entity is
    attribute multstyle : string;
    attribute multstyle of rtl : architecture is "dsp";
begin
```
When applied to a VHDL signal or variable, the attribute specifies the implementation style for all instances of the * operator, which has a result directly assigned to the signal or variable. The attribute overrides the multstyle attribute with the enclosing entity or architecture, if present. In Example 17–91, the multstyle attribute associated with signal result directs the Quartus II software to implement a * b in logic rather than the dedicated hardware.

**Example 17–91. VHDL Code: Applying a multstyle Attribute to a Signal or Variable**

```vhdl
signal a, b : unsigned(8 downto 0);
signal result : unsigned(17 downto 0);
attribute multstyle : string;
attribute multstyle of result : signal is "logic";
result <= a * b;
```

**Full Case Attribute**

A Verilog HDL case statement is full when its case items cover all possible binary values of the case expression or when a default case statement is present. A full_case attribute attached to a case statement header that is not full forces synthesis to treat the unspecified states as a don’t care value. VHDL case statements must be full, so the attribute does not apply to VHDL.

Using this attribute on a case statement that is not full allows you to avoid the latch inference problems discussed in the Recommended Design Practices chapter in volume 1 of the Quartus II Handbook.

Latches have limited support in formal verification tools. Do not infer latches unintentionally, for example, through an incomplete case statement when using formal verification. Formal verification tools support the full_case synthesis attribute (with limited support for attribute syntax, as described in “Synthesis Attributes” on page 17–25).

Using the full_case attribute might cause a simulation mismatch between the Verilog HDL functional and the post-Quartus II simulation because unknown case statement cases can still function as latches during functional simulation. For example, a simulation mismatch can occur with the code in Example 17–92 when sel is 2'b11 because a functional HDL simulation output behaves as a latch and the Quartus II simulation output behaves as a don’t care value.

Altera recommends making the case statement “full” in your regular HDL code, instead of using the full_case attribute.
The case statement in Example 17–92 is not full because you do not specify some \texttt{sel} binary values. Because you use the \texttt{full_case} attribute, synthesis treats the output as “don’t care” when the \texttt{sel} input is \texttt{2'b11}.

**Example 17–92. Verilog HDL Code: a \texttt{full_case} Attribute**

```verilog
module full_case (a, sel, y);
    input [3:0] a;
    input [1:0] sel;
    output y;
    reg y;
    always @ (a or sel)
        case (sel) // synthesis full_case
            2'b00: y=a[0];
            2'b01: y=a[1];
            2'b10: y=a[2];
        endcase
endmodule
```

Verilog-2001 syntax also accepts the statements in Example 17–93 in the case header instead of the comment form as shown in Example 17–92.

**Example 17–93. Verilog-2001 Syntax for the \texttt{full_case} Attribute**

```verilog
(* full_case *) case (sel)
```

### Parallel Case

The \texttt{parallel_case} attribute indicates that you must consider a Verilog HDL case statement as parallel; that is, you can match only one case item at a time. Case items in Verilog HDL case statements might overlap. To resolve multiple matching case items, the Verilog HDL language defines a priority among case items in which the case statement always executes the first case item that matches the case expression value. By default, the Quartus II software implements the extra logic necessary to satisfy this priority relationship.

Attaching a \texttt{parallel_case} attribute to a case statement header allows the Quartus II software to consider its case items as inherently parallel; that is, at most one case item matches the case expression value. Parallel case items simplify the generated logic.

In VHDL, the individual choices in a case statement might not overlap, so they are always parallel and this attribute does not apply.

Altera recommends that you use this attribute only when the case statement is truly parallel. If you use the attribute in any other situation, the generated logic does not match the functional simulation behavior of the Verilog HDL.

Altera recommends that you avoid using the \texttt{parallel_case} attribute, because you may mismatch the Verilog HDL functional and the post-Quartus II simulation.

If you specify SystemVerilog-2005 as the supported Verilog HDL version for your design, you can use the SystemVerilog keyword \texttt{unique} to achieve the same result as the \texttt{parallel_case} directive without causing simulation mismatches.
Example 17–94 shows a casez statement with overlapping case items. In functional HDL simulation, the software prioritizes the three case items by the bits in sel. For example, sel[2] takes priority over sel[1], which takes priority over sel[0]. However, the synthesized design can simulate differently because the parallel_case attribute eliminates this priority. If more than one bit of sel is high, more than one output (a, b, or c) is high as well, a situation that cannot occur in functional HDL simulation.

Example 17–94. Verilog HDL Code: a parallel_case Attribute

```verilog
module parallel_case (sel, a, b, c);
    input [2:0] sel;
    output a, b, c;
    reg a, b, c;
    always @ (sel)
        begin
            {a, b, c} = 3'b0;
            casez (sel) // synthesis parallel_case
                3'b1???: a = 1'b1;
                3'b?1?: b = 1'b1;
                3'b??1: c = 1'b1;
            endcase
        end
    endmodule
```

Verilog-2001 syntax also accepts the statements as shown in Example 17–95 in the case (or casez) header instead of the comment form, as shown in Example 17–94.

Example 17–95. Verilog-2001 Syntax

```verilog
(* parallel_case *) casez (sel)
```

Translate Off and On / Synthesis Off and On

The translate_off and translate_on synthesis directives indicate whether the Quartus II software or a third-party synthesis tool should compile a portion of HDL code that is not relevant for synthesis. The translate_off directive marks the beginning of code that the synthesis tool should ignore; the translate_on directive indicates that synthesis should resume. You can also use the synthesis_on and synthesis_off directives as a synonym for translate on and off.
You can use these directives to indicate a portion of code for simulation only. The synthesis tool reads synthesis-specific directives and processes them during synthesis; however, third-party simulation tools read the directives as comments and ignore them. Example 17–96, Example 17–97, and Example 17–98 show these directives.

Example 17–96. Verilog HDL Code: Translate Off and On

```verilog
// synthesis translate_off
parameter tpd = 2;    // Delay for simulation
#tpd;
// synthesis translate_on
```

Example 17–97. VHDL Code: Translate Off and On

```vhdl
-- synthesis translate_off
use std.textio.all;
-- synthesis translate_on
```

Example 17–98. VHDL 2008 Code: Translate Off and On

```vhdl
/* synthesis translate_off */
use std.textio.all;
/* synthesis translate_on */
```

If you want to ignore only a portion of code in Quartus II Integrated Synthesis, you can use the Altera-specific attribute keyword `altera` for example, use the `// altera translate_off` and `// altera translate_on` directives to direct Quartus II Integrated Synthesis to ignore a portion of code that you intend only for other synthesis tools.

**Ignore translate_off and synthesis_off Directives**

The **Ignore translate_off and synthesis_off Directives** logic option directs Quartus II Integrated Synthesis to ignore the `translate_off` and `synthesis_off` directives described in the previous section. Turning on this logic option allows you to compile code that you want the third-party synthesis tools to ignore; for example, megafunction declarations that the other tools treat as black boxes but the Quartus II software can compile. To set the **Ignore translate_off and synthesis_off Directives** logic option, click **More Settings** on the **Analysis & Synthesis Settings** page of the **Settings** dialog box.

For more information about the **Ignore translate_off and synthesis_off Directives** logic option and the supported devices, refer to **Ignore translate_off and synthesis_off Directives logic option** in Quartus II Help.
Read Comments as HDL

The read_comments_as_HDL synthesis directive indicates that the Quartus II software should compile a portion of HDL code that you commented out. This directive allows you to comment out portions of HDL source code that are not relevant for simulation, while instructing the Quartus II software to read and synthesize that same source code. Setting the read_comments_as_HDL directive to on indicates the beginning of commented code that the synthesis tool should read; setting the read_comments_as_HDL directive to off indicates the end of the code.

You can use this directive with translate_off and translate_on to create one HDL source file that includes a megafunction instantiation for synthesis and a behavioral description for simulation.

Formal verification tools do not support the read_comments_as_HDL directive because the tools do not recognize the directive.

In Example 17–99, Example 17–100, and Example 17–101, the Compiler synthesizes the commented code enclosed by read_comments_as_HDL because the directive is visible to the Quartus II Compiler. VHDL 2008 allows block comments, which comments are also supported for synthesis directives.

Because synthesis directives are case sensitive in Verilog HDL, you must match the case of the directive, as shown in the following examples.

Example 17–99. Verilog HDL Code: Read Comments as HDL

```verilog
// synthesis read_comments_as_HDL on
// my_rom lpm_rom (.address (address),
// .data (data));
// synthesis read_comments_as_HDL off
```

Example 17–100. VHDL Code: Read Comments as HDL

```vhdl
-- synthesis read_comments_as_HDL on
-- my_rom : entity lpm_rom
-- port map (
-- address => address,
-- data => data, );
-- synthesis read_comments_as_HDL off
```

Example 17–101. VHDL 2008 Code: Read Block Comments as HDL

```vhdl
/* synthesis read_comments_as_HDL on */
/* my_rom : entity lpm_rom
port map (   
address => address,  
data => data, ); */ /*
synthesis read_comments_as_HDL off */
```
Use I/O Flipflops

The `useioff` attribute directs the Quartus II software to implement input, output, and output enable flipflops (or registers) in I/O cells that have fast, direct connections to an I/O pin, when possible. To improve I/O performance by minimizing setup, clock-to-output, and clock-to-output enable times, you can apply the `useioff` synthesis attribute. The **Fast Input Register**, **Fast Output Register**, and **Fast Output Enable Register** logic options support this synthesis attribute. You can also set this synthesis attribute in the Assignment Editor.

The `useioff` synthesis attribute takes a boolean value. You can apply the value only to the port declarations of a top-level Verilog HDL module or VHDL entity (it is ignored if applied elsewhere). Setting the value to 1 (Verilog HDL) or `TRUE` (VHDL) instructs the Quartus II software to pack registers into I/O cells. Setting the value to 0 (Verilog HDL) or `FALSE` (VHDL) prevents register packing into I/O cells.

In Example 17–102 and Example 17–103, the `useioff` synthesis attribute directs the Quartus II software to implement the `a_reg`, `b_reg`, and `o_reg` registers in the I/O cells corresponding to the `a`, `b`, and `o` ports, respectively.

**Example 17–102. Verilog HDL Code: the useioff Attribute**

```verilog
module top_level(clk, a, b, o);
    input clk;
    input [1:0] a, b /* synthesis useioff = 1 */;
    output [2:0] o /* synthesis useioff = 1 */;
    reg [1:0] a_reg, b_reg;
    reg [2:0] o_reg;
    always @(posedge clk)
        begin
            a_reg <= a;
            b_reg <= b;
            o_reg <= a_reg + b_reg;
        end
    assign o = o_reg;
endmodule
```
Example 17–103 and Example 17–104 show that the Verilog-2001 syntax also accepts the type of statements instead of the comment form in Example 17–102.

**Example 17–103. Verilog-2001 Code: the useioff Attribute**

```
(* useioff = 1 *) input [1:0] a, b;
(* useioff = 1 *) output [2:0] o;
```

**Example 17–104. VHDL Code: the useioff Attribute**

```
library ieee;
use ieee.std_logic_1164.all;
use ieee.numeric_std.all;
entity useioff_example is
  port (clk : in  std_logic;
        a, b : in  unsigned(1 downto 0);
        o    : out unsigned(1 downto 0));
attribute useioff : boolean;
attribute useioff of a : signal is true;
attribute useioff of b : signal is true;
attribute useioff of o : signal is true;
end useioff_example;
architecture rtl of useioff_example is
  signal o_reg, a_reg, b_reg : unsigned(1 downto 0);
begin
  process(clk)
  begin
    if (clk = '1' AND clk'event) then
      a_reg <= a;
      b_reg <= b;
      o_reg <= a_reg + b_reg;
    end if;
  end process;
  o <= o_reg;
end rtl;
```

**Specifying Pin Locations with chip_pin**

The `chip_pin` attribute allows you to assign pin locations in your HDL source. You can use the attribute only on the ports of the top-level entity or module in your design. You can assign pins only to single-bit or one-dimensional bus ports in your design.

For single-bit ports, the value of the `chip_pin` attribute is the name of the pin on the target device, as specified by the pin table of the device.

In addition to the `chip_pin` attribute, the Quartus II software supports the `altera_chip_pin_lc` attribute name for compatibility with other synthesis tools. When using this attribute in other synthesis tools, some older device families require an “@” symbol in front of each pin assignment. In the Quartus II software, the “@” is optional.
Example 17–105 through Example 17–107 show different ways of assigning my_pin1 to Pin C1 and my_pin2 to Pin 4 on a different target device.


```verbatim
input my_pin1 /* synthesis chip_pin = "C1" */;
input my_pin2 /* synthesis altera_chip_pin_lc = "@4" */;
```


```verbatim
(* chip_pin = "C1" *) input my_pin1;
(* altera_chip_pin_lc = "@4" *) input my_pin2;
```

**Example 17–107. VHDL Code: Applying Chip Pin to a Single Pin**

```vhdl
entity my_entity is
  port(my_pin1: in std_logic; my_pin2: in std_logic;...);
end my_entity;
attribute chip_pin : string;
attribute altera_chip_pin_lc : string;
attribute chip_pin of my_pin1 : signal is "C1";
attribute altera_chip_pin_lc of my_pin2 : signal is "@4";
```

For bus I/O ports, the value of the chip pin attribute is a comma-delimited list of pin assignments. The order in which you declare the range of the port determines the mapping of assignments to individual bits in the port. To leave a bit unassigned, leave its corresponding pin assignment blank.


```verbatim
input [2:0]  my_pin /* synthesis chip_pin = "4, 5, 6" */;
```

**Example 17–109** reverses the order of the signals in the bus, assigning my_pin[0] to Pin 4 and my_pin[2] to Pin 6 but leaves my_pin[1] unassigned.


```verbatim
input [0:2]  my_pin /* synthesis chip_pin = "4, ,6" */;
```

**Example 17–110** assigns my_pin[2] to Pin 4 and my_pin[0] to Pin 6, but leaves my_pin[1] unassigned.

**Example 17–110. VHDL Code: Applying Chip Pin to Part of a Bus of Pins**

```vhdl
entity my_entity is
  port(my_pin: in std_logic_vector(2 downto 0);...);
end my_entity;
attribute chip_pin of my_pin: signal is "4, ,6";
```
Example 17–111 shows a VHDL example on how to assign pin location and I/O standard.

**Example 17–111. VHDL Code: Assigning Pin Location and I/O Standard**

```
attribute altera_chip_pin_lc: string;
attribute  altera_attribute: string;
attribute altera_chip_pin_lc of clk: signal is "B13";
attribute altera_attribute of clk:signal is "-name IO_STANDARD "3.3-V LVCMOS"";
```

Example 17–112 shows a Verilog-2001 example on how to assign pin location and I/O standard.


```
(* altera_attribute = "-name IO_STANDARD \"3.3-V LVCMOS\"" *)
(* chip_pin = "L5" *)
input clk;

(* altera_attribute = "-name IO_STANDARD LVDS" *)
(* chip_pin = "L4" *)
input sel;
output [3:0] data_o, input [3:0] data_i);
```

**Using altera_attribute to Set Quartus II Logic Options**

The `altera_attribute` attribute allows you to apply Quartus II logic options and assignments to an object in your HDL source code. You can set this attribute on an entity, architecture, instance, register, RAM block, or I/O pin. You cannot set it on an arbitrary combinational node such as a net. With `altera_attribute`, you can control synthesis options from your HDL source even when the options lack a specific HDL synthesis attribute (such as many of the logic options presented earlier in this chapter). You can also use this attribute to pass entity-level settings and assignments to phases of the Compiler flow that follow Analysis & Synthesis, such as Fitting.

Assignments or settings made through the Quartus II software, the `.qsf`, or the Tcl interface take precedence over assignments or settings made with the `altera_attribute` synthesis attribute in your HDL code.

The syntax for setting this attribute in HDL is the same as the syntax for other synthesis attributes, as shown in “Synthesis Attributes” on page 17–25.

The attribute value is a single string containing a list of `.qsf` variable assignments separated by semicolons, as shown in Example 17–113.

**Example 17–113. Variable Assignments Separated by Semicolons**

```
-name <variable_1> <value_1> ; -name <variable_2> <value_2>[;...]
```

If the Quartus II option or assignment includes a target, source, and section tag, you must use the syntax in Example 17–114 for each `.qsf` variable assignment:

**Example 17–114. Syntax for Each .qsf Variable Assignment**

```
-name <variable> <value>
-from <source> -to <target> -section_id <section>
```
Example 17–115 shows the syntax for the full attribute value, including the optional target, source, and section tags for two different .qsf assignments.

Example 17–115. Syntax for Full Attribute Value

```
" -name <variable_1> <value_1> [-from <source_1>] [-to <target_1>] [-section_id <section_1>]; -name <variable_2> <value_2> [-from <source_2>] [-to <target_2>] [-section_id <section_2>]
```

If the assigned value of a variable is a string of text, you must use escaped quotes around the value in Verilog HDL (as shown in Example 17–116), or double-quotes in VHDL (as shown in Example 17–117):

Example 17–116. Assigned Value of a Variable in Verilog HDL (With Nonexistent Variable and Value Terms)

```
"VARIABLE_NAME "STRING_VALUE"
```

Example 17–117. Assigned Value of a Variable in VHDL (With Nonexistent Variable and Value Terms)

```
"VARIABLE_NAME ""STRING_VALUE"
```

To find the .qsf variable name or value corresponding to a specific Quartus II option or assignment, you can set the option setting or assignment in the Quartus II software, and then make the changes in the .qsf. You can also refer to the Quartus II Settings File Manual, which documents all variable names.

Example 17–118 through Example 17–120 use altera_attribute to set the power-up level of an inferred register.

For inferred instances, you cannot apply the attribute to the instance directly. Therefore, you must apply the attribute to one of the output nets of the instance. The Quartus II software automatically moves the attribute to the inferred instance.

Example 17–118. Verilog-1995 Code: Applying altera_attribute to an Instance

```
reg my_reg /* synthesis altera_attribute = "-name POWER_UP_LEVEL HIGH" */;
```

Example 17–119. Verilog-2001 Code: Applying altera_attribute to an Instance

```
(* altera_attribute = "-name POWER_UP_LEVEL HIGH" *) reg my_reg;
```

Example 17–120. VHDL Code: Applying altera_attribute to an Instance

```
signal my_reg : std_logic;
attribute altera_attribute : string;
attribute altera_attribute of my_reg: signal is "-name POWER_UP_LEVEL HIGH";
```
Example 17–121 through Example 17–123 use the altera_attribute to disable the Auto Shift Register Replacement synthesis option for an entity. To apply the Altera Attribute to a VHDL entity, you must set the attribute on its architecture rather than on the entity itself.

**Example 17–121. Verilog-1995 Code: Applying altera_attribute to an Entity**

```verilog
module my_entity(...) /* synthesis altera_attribute = "-name AUTO_SHIFT_REGISTER_RECOGNITION OFF" */;
```

**Example 17–122. Verilog-2001 Code: Applying altera_attribute to an Entity**

```verilog
(* altera_attribute = "-name AUTO_SHIFT_REGISTER_RECOGNITION OFF" *)
module my_entity(...) ;
```

**Example 17–123. VHDL Code: Applying altera_attribute to an Entity**

```vhdl
entity my_entity is
  -- Declare generics and ports
end my_entity;
architecture rtl of my_entity is
attribute altera_attribute : string;
  -- Attribute set on architecture, not entity
attribute altera_attribute of rtl: architecture is "-name AUTO_SHIFT_REGISTER_RECOGNITION OFF";
beg
  -- The architecture body
end rtl;
```

You can also use altera_attribute for more complex assignments that have more than one instance. In Example 17–125 through Example 17–127, the altera_attribute cuts all timing paths from reg1 to reg2, equivalent to this Tcl or .qsf command, as shown in Example 17–124:

**Example 17–124.**

```tcl
set_instance_assignment -name CUT ON -from reg1 -to reg2
```

**Example 17–125. Verilog-1995 Code: Applying altera_attribute with the -to Option**

```verilog
reg reg2;
reg regl /* synthesis altera_attribute = "-name CUT ON -to reg2" */;
```

**Example 17–126. Verilog-2001 and SystemVerilog Code: Applying altera_attribute with the -to Option**

```verilog
reg reg2;
(* altera_attribute = "-name CUT ON -to reg2" *) reg regl;
```

**Example 17–127. VHDL Code: Applying altera_attribute with the -to Option**

```vhdl
signal regl, reg2 : std_logic;
attribute altera_attribute : string;
attribute altera_attribute of regl : signal is "-name CUT ON -to reg2";
```
You can specify either the -to option or the -from option in a single
altera_attribute; Integrated Synthesis automatically sets the remaining option to
the target of the altera_attribute. You can also specify wildcards for either option.
For example, if you specify “*” for the -to option instead of reg2 in these examples,
the Quartus II software cuts all timing paths from reg1 to every other register in this
design entity.

You can use the altera_attribute only for entity-level settings, and the assignments
(including wildcards) apply only to the current entity.

Analyzing Synthesis Results

After performing synthesis, you can check your synthesis results in the Analysis &
Synthesis section of the Compilation Report and the Project Navigator.

Analysis & Synthesis Section of the Compilation Report

The Compilation Report, which provides a summary of results for the project, appears
after a successful compilation. After Analysis & Synthesis, the Summary section of the
Compilation Report provides a summary of utilization based on synthesis data,
before Fitter optimizations have occurred. The Analysis & Synthesis section lists
synthesis-specific information.

Analysis & Synthesis includes various report sections, including a list of the source
files read for the project, the resource utilization by entity after synthesis, and
information about state machines, latches, optimization results, and parameter
settings.

For more information about each report section, refer to the Analysis & Synthesis
Summary Reports in Quartus II Help.

Project Navigator

The Hierarchy tab of the Project Navigator provides a view of the project hierarchy
and a summary of resource and device information about the current project. After
Analysis & Synthesis, before the Fitter begins, the Project Navigator provides a
summary of utilization based on synthesis data, before Fitter optimizations have
occurred.

If an entity in the Hierarchy tab contains parameter settings, a tooltip displays the
settings when you hold the pointer over the entity.

Upgrade IP Components Dialog Box

In the Quartus II software version 12.1 SP1 and later, the Upgrade IP Components
dialog box allows you to upgrade all outdated IP in your project after you move to a
newer version of the Quartus II software.

For more information about the Upgrade IP Components dialog box, refer to Upgrade
IP Components dialog box in Quartus II Help.
Analyzing and Controlling Synthesis Messages

This section provides information about the generated messages during synthesis, and how you can control which messages appear during compilation.

Quartus II Messages

The messages that appear during Analysis & Synthesis describe many of the optimizations during the synthesis stage, and provide information about how the software interprets your design. Altera recommends checking the messages to analyze Critical Warnings and Warnings, because these messages can relate to important design problems. Read the Info messages to get more information about how the software processes your design.

The software groups the messages by following types: Info, Warning, Critical Warning, and Error.

For more information about the Messages window and message suppression, refer to About the Messages Window and About Message Suppression in Quartus II Help.

For more information about the Messages, refer to Managing Quartus II Projects chapter in volume 2 of the Quartus II Handbook.

You can specify the type of Analysis & Synthesis messages that you want to view by selecting the Analysis & Synthesis Message Level option. You can specify the display level by performing the following steps:

1. On the Assignments menu, click Settings.
2. In the Category list, click Analysis & Synthesis Settings.
3. Click More Settings. Select the level for the Analysis & Synthesis Message Level option.

VHDL and Verilog HDL Messages

The Quartus II software issues a variety of messages when it is analyzing and elaborating the Verilog HDL and VHDL files in your design. These HDL messages are a subset of all Quartus II messages that help you identify potential problems early in the design process.

HDL messages fall into the following categories:

- **Info message**—lists a property of your design.
- **Warning message**—indicates a potential problem in your design. Potential problems come from a variety of sources, including typos, inappropriate design practices, or the functional limitations of your target device. Though HDL warning messages do not always identify actual problems, Altera recommends investigating code that generates an HDL warning. Otherwise, the synthesized behavior of your design might not match your original intent or its simulated behavior.
- **Error message**—indicates an actual problem with your design. Your HDL code can be invalid due to a syntax or semantic error, or it might not be synthesizable as written.
In Example 17–128, the sensitivity list contains multiple copies of the variable \( i \). While the Verilog HDL language does not prohibit duplicate entries in a sensitivity list, it is clear that this design has a typing error: Variable \( j \) should be listed on the sensitivity list to avoid a possible simulation or synthesis mismatch.

**Example 17–128. Generating an HDL Warning Message**

```verilog
//dup.v
module dup(input i, input j, output reg o);
always @ (i or i)
  o = i & j;
endmodule
```

When processing the HDL code, the Quartus II software generates the warning message shown in Example 17–129:

**Example 17–129.**

Warning: (10276) Verilog HDL sensitivity list warning at dup.v(2): sensitivity list contains multiple entries for "i".

In Verilog HDL, variable names are case sensitive, so the variables \( \text{my\_reg} \) and \( \text{MY\_REG} \) in Example 17–130 are two different variables. However, declaring variables that have names in different cases is confusing, especially if you use VHDL, in which variables are not case sensitive.

**Example 17–130. Generating HDL Info Messages**

```verilog
// namecase.v
module namecase (input i, output o);
  reg my_reg;
  reg MY_REG;
  assign o = i;
endmodule
```

When processing the HDL code, the Quartus II software generates the following informational message, as shown in Example 17–131:

**Example 17–131.**

Info: (10281) Verilog HDL information at namecase.v(3): variable name "MY\_REG" and variable name "my\_reg" should not differ only in case.

In addition, the Quartus II software generates additional HDL info messages to inform you that this small design does not use neither \( \text{my\_reg} \) nor \( \text{MY\_REG} \), as shown in Example 17–132:

**Example 17–132.**

Info: (10035) Verilog HDL or VHDL information at namecase.v(3): object "my\_reg" declared but not used
Info: (10035) Verilog HDL or VHDL information at namecase.v(4): object "MY\_REG" declared but not used
The Quartus II software allows you to control how many HDL messages you can view during the Analysis & Elaboration of your design files. You can set the HDL Message Level to enable or disable groups of HDL messages, or you can enable or disable specific messages, as described in the following sections.

For more information about synthesis directives and their syntax, refer to “Synthesis Directives” on page 17–27.

Setting the HDL Message Level

The HDL Message Level specifies the types of messages that the Quartus II software displays when it is analyzing and elaborating your design files. Table 17–5 describes the HDL message levels.

Table 17–5. HDL Info Message Level

<table>
<thead>
<tr>
<th>Level</th>
<th>Purpose</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Level1</td>
<td>High-severity messages only</td>
<td>If you want to view only the HDL messages that identify likely problems with your design, select Level1. When you select Level1, the Quartus II software issues a message only if there is an actual problem with your design.</td>
</tr>
<tr>
<td>Level2</td>
<td>High-severity and medium-severity messages</td>
<td>If you want to view additional HDL messages that identify possible problems with your design, select Level2. Level2 is the default setting.</td>
</tr>
<tr>
<td>Level3</td>
<td>All messages, including low-severity messages</td>
<td>If you want to view all HDL info and warning messages, select Level3. This level includes extra “LINT” messages that suggest changes to improve the style of your HDL code.</td>
</tr>
</tbody>
</table>

You must address all issues reported at the Level1 setting. The default HDL message level is Level2.

To set the HDL Message Level in the Quartus II software, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, click Analysis & Synthesis Settings.
3. Set the necessary message level from the pull-down menu in the HDL Message Level list, and then click OK.

You can override this default setting in a source file with the message_level synthesis directive, which takes the values level1, level2, and level3, as shown in Example 17–133 and Example 17–134.

Example 17–133. Verilog HDL Examples of message_level Directive

```verilog
// altera message_level level1
or
/* altera message_level level3 */
```

Example 17–134. VHDL Code: message_level Directive

```vhdl
-- altera message_level level2
```
A `message_level` synthesis directive remains effective until the end of a file or until the next `message_level` directive. In VHDL, you can use the `message_level` synthesis directive to set the HDL Message Level for entities and architectures, but not for other design units. An HDL Message Level for an entity applies to its architectures, unless overridden by another `message_level` directive. In Verilog HDL, you can use the `message_level` directive to set the HDL Message Level for a module.

### Enabling or Disabling Specific HDL Messages by Module/Entity

Message ID is in parentheses at the beginning of the message. Use the Message ID to enable or disable a specific HDL info or warning message. Enabling or disabling a specific message overrides its HDL Message Level. This method is different from the message suppression in the Messages window because you can disable messages for a specific module or a specific entity. This method applies only to the HDL messages, and if you disable a message with this method, the Quartus II software lists the message as a suppressed message.

To disable specific HDL messages in the Quartus II software, follow these steps:

1. On the Assignments menu, click **Settings**.
2. In the **Category** list, expand **Analysis & Synthesis Settings** and select **Advanced**.
3. In the **Advanced Message Settings** dialog box, add the Message IDs you want to enable or disable.

To enable or disable specific HDL messages in your HDL, use the `message_on` and `message_off` synthesis directives. These directives require a space-separated list of Message IDs. You can enable or disable messages with these synthesis directives immediately before Verilog HDL modules, VHDL entities, or VHDL architectures. You cannot enable or disable a message during an HDL construct.

A message enabled or disabled via a `message_on` or `message_off` synthesis directive overrides its HDL Message Level or any `message_level` synthesis directive. The message remains disabled until the end of the source file or until you use another `message_on` or `message_off` directive to change the status of the message.

**Example 17–135. Verilog HDL message_off Directive for Message with ID 10000**

```plaintext
// altera message_off 10000
or
/* altera message_off 10000 */
```

**Example 17–136. VHDL message_off Directive for Message with ID 10000**

```plaintext
-- altera message_off 10000
```
Node-Naming Conventions in Quartus II Integrated Synthesis

This section provides an overview of the conventions that the Quartus II software uses during synthesis to name the nodes created from your HDL design. This information is useful for finding logic node names during verification and debugging of a design. This section focuses on the conventions for Verilog HDL and VHDL code, but discusses AHDL and .bdf files when appropriate.

Whenever possible, Quartus II Integrated Synthesis uses wire or signal names from your source code to name nodes such as LEs or ALMs. Some nodes, such as registers, have predictable names that do not change when a design is resynthesized, although certain optimizations can affect register names. The names of other nodes, particularly LEs or ALMs that contain only combinational logic, can change due to logic optimizations that the software performs.

This section describes the following topics:

- “Hierarchical Node-Naming Conventions”
- “Node-Naming Conventions for Registers (DFF or D Flipflop Atoms)”
- “Register Changes During Synthesis” on page 17–80
- “Preserving Register Names” on page 17–82
- “Node-Naming Conventions for Combinational Logic Cells” on page 17–82
- “Preserving Combinational Logic Names” on page 17–83

Hierarchical Node-Naming Conventions

To make each name in your design unique, the Quartus II software adds the hierarchy path to the beginning of each name. The “|” separator indicates a level of hierarchy. For each instance in the hierarchy, the software adds the entity name and the instance name of that entity, with the “;” separator between each entity name and its instance name. For example, if a design defines entity A with the name my_A_inst, the hierarchy path of that entity would be A:my_A_inst. You can obtain the full name of any node by starting with the hierarchical instance path, followed by a “|”, and ending with the node name inside that entity. Example 17–137 shows you the convention:

Example 17–137.

<entity 0>:<instance_name 0> | <entity 1>:<instance_name 1> | ... | <instance_name n> | <node_name>

For example, if entity A contains a register (DFF atom) called my_dff, its full hierarchy name would be A:my_A_inst | my_dff.

To instruct the Compiler to generate node names that do not contain entity names, on the Compilation Process Settings page of the Settings dialog box, click More Settings, and then turn off Display entity name for node name. With this option turned off, the node names use the convention in shown in Example 17–138:

Example 17–138.

<instance_name 0> | <instance_name 1> | ... | <instance_name n> | <node_name>
Node-Naming Conventions for Registers (DFF or D Flipflop Atoms)

In Verilog HDL and VHDL, inferred registers use the names of the `reg` or `signal` connected to the output.

Example 17–139 is an example of a register in Verilog HDL that creates a DFF primitive called `my_dff_out`.

**Example 17–139. Verilog HDL Register**

```verilog
wire dff_in, my_dff_out, clk;
always @(posedge clk)
my_dff_out <= dff_in;
```

Similarly, Example 17–140 is an example of a register in VHDL that creates a DFF primitive called `my_dff_out`.

**Example 17–140. VHDL Register**

```vhdl
signal dff_in, my_dff_out, clk;
process (clk)
begin
  if (rising_edge(clk)) then
    my_dff_out <= dff_in;
  end if;
end process;
```

AHDL designs explicitly declare DFF registers rather than infer, so the software uses the user-declared name for the register.

For schematic designs using a `.bdf`, your design names all elements when you instantiate the elements in your design, so the software uses the name you defined for the register or DFF.

In the special case that a wire or signal (such as `my_dff_out` in the preceding examples) is also an output pin of your top-level design, the Quartus II software cannot use that name for the register (for example, cannot use `my_dff_out`) because the software requires that all logic and I/O cells have unique names. Here, Quartus II Integrated Synthesis appends `-reg0` to the register name.

For example, the Verilog HDL code in Example 17–141 generates a register called `q-reg0`.

**Example 17–141. Verilog HDL Register Feeding Output Pin**

```verilog
module my_dff (input clk, input d, output q);
always @(posedge clk)
q <= d;
endmodule
```

This situation occurs only for registers driving top-level pins. If a register drives a port of a lower level of the hierarchy, the software removes the port during hierarchy flattening and the register retains its original name, in this case, `q`. 
Register Changes During Synthesis

On some occasions, you might not find registers that you expect to view in the synthesis netlist. Logic optimization might remove registers and synthesis optimizations might change the names of the registers. Common optimizations include inference of a state machine, counter, adder-subtractor, or shift register from registers and surrounding logic. Other common register changes occur when the software packs these registers into dedicated hardware on the FPGA, such as a DSP block or a RAM block.

This section describes the following factors that can affect register names:

- “Synthesis and Fitting Optimizations”
- “State Machines”
- “Inferred Adder-Subtractors, Shift Registers, Memory, and DSP Functions” on page 17–81
- “Packed Input and Output Registers of RAM and DSP Blocks” on page 17–82
- “Preserving Register Names” on page 17–82
- “Preserving Combinational Logic Names” on page 17–83

Synthesis and Fitting Optimizations

Logic optimization during synthesis might remove registers if you do not connect the registers to inputs or outputs in your design, or if you can simplify the logic due to constant signal values. Synthesis optimizations might change register names, such as when the software merges duplicate registers to reduce resource utilization.

NOT-gate push back optimizations can affect registers that use preset signals. This type of optimization can impact your timing assignments when the software uses registers as clock dividers. If this situation occurs in your design, change the clock settings to work on the new register name.

Synthesis netlist optimizations often change node names because the software can combine or duplicate registers to optimize your design.

For more information about the type of optimizations performed by synthesis netlist optimizations, refer to the Netlist Optimizations and Physical Synthesis chapter in volume 2 of the Quartus II Handbook.

The Quartus II Compilation Report provides a list of registers that synthesis optimizations remove, and a brief reason for the removal. To generate the Quartus II Compilation Report, follow these steps:

1. In the Analysis & Synthesis folder, open Optimization Results.
2. Open Register Statistics, and then click the Registers Removed During Synthesis report.
3. Click Removed Registers Triggering Further Register Optimizations.

The second report contains a list of registers that causes synthesis optimizations to remove other registers from your design. The report provides a brief reason for the removal, and a list of registers that synthesis optimizations remove due to the removal of the initial register.
Quartus II Integrated Synthesis creates synonyms for registers duplicated with the **Maximum Fan-Out** option (or `maxfan` attribute). Therefore, timing assignments applied to nodes that are duplicated with this option are applied to the new nodes as well.

The Quartus II Fitter can also change node names after synthesis (for example, when the Fitter uses register packing to pack a register into an I/O element, or when physical synthesis modifies logic). The Fitter creates synonyms for duplicated registers so timing analysis can use the existing node name when applying assignments.

You can instruct the Quartus II software to preserve certain nodes throughout compilation so you can use them for verification or making assignments. For more information, refer to “Preserving Register Names” on page 17–82.

**State Machines**

If your HDL code infers a state machine, the software maps the registers that represent the states into a new set of registers that implement the state machine. Most commonly, the software converts the state machine into a one-hot form in which one register represents each state. In this case, for Verilog HDL or VHDL designs, the registers take the name of the state register and the states.

For example, consider a Verilog HDL state machine in which the states are `parameter state0 = 1, state1 = 2, state2 = 3`, and in which the software declares the state machine register as `reg [1:0] my_fsm`. In this example, the three one-hot state registers are `my_fsm.state0, my_fsm.state1, and my_fsm.state2`.

An AHDL design explicitly specifies state machines with a state machine name. Your design names state machine registers with synthesized names based on the state machine name, but not the state names. For example, if a `my_fsm` state machine has four state bits, the software might synthesize these state bits with names such as `my_fsm~12, my_fsm~13, my_fsm~14, and my_fsm~15`.

**Inferred Adder-Subtractors, Shift Registers, Memory, and DSP Functions**

The Quartus II software infers megafunctions from Verilog HDL and VHDL code for logic that forms adder-subtractors, shift registers, RAM, ROM, and arithmetic functions that are placed in DSP blocks.

For information about inferring megafunctions, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

Because adder-subtractors are part of a megafunction instead of generic logic, the combinational logic exists in the design with different names. For shift registers, memory, and DSP functions, the software implements the registers and logic inside the dedicated RAM or DSP blocks in the device. Thus, the registers are not visible as separate LEs or ALMs.
Packed Input and Output Registers of RAM and DSP Blocks

The software packs registers into the input registers and output registers of RAM and DSP blocks, so that they are not visible as separate registers in LEs or ALMs.

For information about packing registers into RAM and DSP megafunctions, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

Preserving Register Names

Altera recommends that you preserve certain register names for verification or debugging, or to ensure that you applied timing assignments correctly. Quartus II Integrated Synthesis preserves certain nodes automatically if the software uses the nodes in a timing constraint.

Use the `preserve` attribute to instruct the Compiler not to minimize or remove a specified register during synthesis optimizations or register netlist optimizations. For details, refer to “Preserve Registers” on page 17–42.

Use the `noprune` attribute to preserve a fan-out-free register through the entire compilation flow. For details, refer to “Noprune Synthesis Attribute/Preserve Fan-out Free Register Node” on page 17–43.

Use the synthesis attribute `syn_dont_merge` to ensure that the Compiler does not merge registers with other registers. For more information, refer to “Disable Register Merging/Don’t Merge Register” on page 17–43.

Node-Naming Conventions for Combinational Logic Cells

Whenever possible for Verilog HDL, VHDL, and AHDL code, the Quartus II software uses wire names that are the targets of assignments, but can change the node names due to synthesis optimizations.

For example, consider the Verilog HDL code in Example 17–142. Quartus II Integrated Synthesis uses the names `c`, `d`, `e`, and `f` for the generated combinational logic cells.

### Example 17–142. Naming Nodes for Combinational Logic Cells in Verilog HDL

```verilog
wire c;
reg d, e, f;

assign c = a | b;
always @(a or b)
d = a & b;
always @(a or b) begin : my_label
e = a ^ b;
end

always @(a or b)
f = ~(a | b);
```

For schematic designs using a `.bdf`, your design names all elements when you instantiate the elements in your design and the software uses the name you defined when possible.
If logic cells, such as those created in Example 17–142, are packed with registers in device architectures such as the Stratix and Cyclone device families, those names might not appear in the netlist after fitting. In other devices, such as newer families in the Stratix and Cyclone series device families, the register and combinational nodes are kept separate throughout the compilation, so these names are more often maintained through fitting.

When logic optimizations occur during synthesis, it is not always possible to retain the initial names as described. Sometimes, synthesized names are used, which are the wire names with a tilde (~) and a number appended. For example, if a complex expression is assigned to wire \( w \) and that expression generates several logic cells, those cells can have names such as \( w, w-1 \), and \( w-2 \). Sometimes the original wire name \( w \) is removed, and an arbitrary name such as \( rt1-123 \) is created. Quartus II Integrated Synthesis attempts to retain user names whenever possible. Any node name ending with \( \sim <\text{number}>\) is a name created during synthesis, which can change if the design is changed and re-synthesized. Knowing these naming conventions helps you understand your post-synthesis results, helping you to debug your design or create assignments.

During synthesis, the software maintains combinational clock logic by not changing nodes that might be clocks. The software also maintains or protects multiplexers in clock trees, so that the TimeQuest analyzer has information about which paths are unate, to allow complete and correct analysis of combinational clocks. Multiplexers often occur in clock trees when the software selects between different clocks. To help with the analysis of clock trees, the software ensures that each multiplexer encountered in a clock tree is broken into 2:1 multiplexers, and each of those 2:1 multiplexers is mapped into one lookup table (independent of the device family). This optimization might result in a slight increase in area, and for some designs a decrease in timing performance. You can turn off this multiplexer protection with the option Clock MUX Protection under More Settings on the Analysis & Synthesis Settings page of the Settings dialog box.

For more information about Clock MUX Protection logic option and a list of supported devices, refer to Clock MUX Protection logic option in Quartus II Help.

**Preserving Combinational Logic Names**

You can preserve certain combinational logic node names for verification or debugging, or to ensure that timing assignments are applied correctly.

Use the keep attribute to keep a wire name or combinational node name through logic synthesis minimizations and netlist optimizations. For details, refer to “Keep Combinational Node/Implement as Output of Logic Cell” on page 17–44.

For any internal node in your design clock network, use keep to protect the name so that you can apply correct clock settings. Also, set the attribute for combinational logic involved in cut and -through assignments.

Setting the keep attribute for combinational logic can increase the area utilization and increase the delay of the final mapped logic because the attribute requires the insertion of extra combinational logic. Use the attribute only when necessary.
Scripting Support

You can run procedures and make settings in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II Command-Line and Tcl API Help browser.

To run the Help browser, type the command at the command prompt shown in Example 17–143:

**Example 17–143.**

```bash
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the *Quartus II Handbook*. For more information about all settings and constraints in the Quartus II software, refer to the *Quartus II Settings File Manual*. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the *Quartus II Handbook*.

You can specify many of the options described in this section either on an instance, at the global level, or both.

To make a global assignment, use the Tcl command shown in Example 17–144:

**Example 17–144.**

```tcl
set_global_assignment -name <QSF Variable Name> <Value>
```

To make an instance assignment, use the Tcl command shown in Example 17–145:

**Example 17–145.**

```tcl
set_instance_assignment -name <QSF Variable Name> <Value> -to <Instance Name>
```

To set the *Synthesis Effort* option at the command line, use the *--effort* option with the `quartus_map` executable, as shown in Example 17–146.

**Example 17–146. Command Syntax for Specifying Synthesis Effort Option**

```bash
quartus_map <Design name> --effort= "auto | fast"
```

The early timing estimate feature gives you preliminary timing estimates before running a full compilation, which results in a quicker iteration time; therefore, you can save significant compilation time to get a good estimation of the final timing of your design.
If you want to run fast synthesis with the Fitter Early Timing Estimate option, use the command shown in Example 17–147. This command runs the full flow with timing analysis:


```
quartus_sh --flow early_timing_estimate_with_synthesis <Design name>
```

For more information, refer to “Synthesis Effort” on page 17–34.

### Adding an HDL File to a Project and Setting the HDL Version

To add an HDL or schematic entry design file to your project, use the Tcl assignments shown in Example 17–148:

Example 17–148.

```
set_global_assignment -name VERILOG_FILE <file name>.<v|sv>
set_global_assignment -name SYSTEMVERILOG_FILE <file name>.sv
set_global_assignment -name VHDL_FILE <file name>.<vhdl|vhd|vhdl>
set_global_assignment -name AHDL_FILE <file name>.tdf
set_global_assignment -name BDF_FILE <file name>.bdf
```

You can use any file extension for design files, as long as you specify the correct language when adding the design file. For example, you can use .h for Verilog HDL header files.

To specify the Verilog HDL or VHDL version, use the option shown in Example 17–149, at the end of the VERILOG_FILE or VHDL_FILE command:

Example 17–149.

```
-HDL_VERSION <language version>
```

The variable `<language version>` takes one of the following values:

- VERILOG_1995
- VERILOG_2001
- SYSTEMVERILOG_2005
- VHDL_1987
- VHDL_1993
- VHDL_2008

For example, to add a Verilog HDL file called `my_file.v` written in Verilog-1995, use the command shown in Example 17–150:

Example 17–150.

```
set_global_assignment -name VERILOG_FILE my_file.v -HDL_VERSION \ VERILOG_1995
```
In Example 17–151, the syn_encoding attribute associates a binary encoding with the states in the enumerated type count_state. In this example, the states are encoded with the following values: zero = "11", one = "01", two = "10", three = "00".

**Example 17–151. Specifying User-Encoded States with the syn_encoding Attribute in VHDL**

```vhdl
ARCHITECTURE rtl OF my_fsm IS
  TYPE count_state IS (zero, one, two, three);
  ATTRIBUTE syn_encoding : STRING;
  ATTRIBUTE syn_encoding OF count_state : TYPE IS "11 01 10 00";
  SIGNAL present_state, next_state : count_state;
BEGIN
```

You can also use the syn_encoding attribute in Verilog HDL to direct the synthesis tool to use the encoding from your HDL code, instead of using the **State Machine Processing** option.

The syn_encoding value "user" instructs the Quartus II software to encode each state with its corresponding value from the Verilog HDL source code. By changing the values of your state constants, you can change the encoding of your state machine.

In Example 17–152, the states are encoded as follows:

```verilog
(* syn_encoding = "user" *) reg [1:0] state;
parameter init = 0, last = 3, next = 1, later = 2;
always @ (state) begin
  case (state)
    init:
      out = 2'b01;
    next:
      out = 2'b10;
    later:
      out = 2'b11;
    last:
      out = 2'b00;
  endcase
end
```

Without the syn_encoding attribute, the Quartus II software encodes the state machine based on the current value of the **State Machine Processing** logic option.

If you also specify a safe state machine (as described in “Safe State Machine” on page 17–38), separate the encoding style value in the quotation marks from the safe value with a comma, as follows: “safe, one-hot” or “safe, gray”.

For more information, refer to “Manually Specifying State Assignments Using the syn_encoding Attribute” on page 17–36.
Assigning a Pin

To assign a signal to a pin or device location, use the Tcl command shown in Example 17–153:

**Example 17–153.**

```tcl
set_location_assignment -to <signal name> <location>
```

Valid locations are pin location names. Some device families also support edge and I/O bank locations. Edge locations are `EDGE_BOTTOM`, `EDGE_LEFT`, `EDGE_TOP`, and `EDGE_RIGHT`. I/O bank locations include `IOBANK_1` to `IOBANK_n`, where `n` is the number of I/O banks in a device.

Creating Design Partitions for Incremental Compilation

To create a partition, use the command shown in Example 17–154:

**Example 17–154.**

```tcl
set_instance_assignment -name PARTITION_HIERARCHY \ 
<file name> -to <destination> -section_id <partition name>
```

The `<file name>` variable is the name used for internally generated netlist files during incremental compilation. If you create the partition in the Quartus II software, netlist files are named automatically by the Quartus II software based on the instance name. If you use Tcl to create your partitions, you must assign a custom file name that is unique across all partitions. For the top-level partition, the specified file name is ignored, and you can use any dummy value. To ensure the names are safe and platform independent, file names should be unique, regardless of case. For example, if a partition uses the file name `my_file`, no other partition can use the file name `MY_FILE`. To make file naming simple, Altera recommends that you base each file name on the corresponding instance name for the partition.

The `<destination>` is the short hierarchy path of the entity. A short hierarchy path is the full hierarchy path without the top-level name, for example: "ram:ram_unit|altsyncram:altsyncram_component" (with quotation marks). For the top-level partition, you can use the pipe (`|`) symbol to represent the top-level entity.

For more information about hierarchical naming conventions, refer to “Node-Naming Conventions in Quartus II Integrated Synthesis” on page 17–78.

The `<partition name>` is the partition name you designate, which should be unique and less than 1024 characters long. The name may only consist of alphanumeric characters, as well as pipe ( `|` ), colon ( `:` ), and underscore ( `_` ) characters. Altera recommends enclosing the name in double quotation marks (" ").
Quartus II Synthesis Options

For more information about the .qsf variable names and applicable values for the settings discussed in this chapter, refer to Logic options in Quartus II Help.

Conclusion

The Quartus II Integrated Synthesis supports Verilog HDL, SystemVerilog, VHDL, and Altera-specific languages, making the synthesis feature an easy-to-use, standalone solution for Altera designs. You can use the synthesis options in the software or in your HDL code to better control the way your design is synthesized, helping you improve your synthesis results. Use Quartus II reports and messages to analyze your compilation results.

Document Revision History

Table 17–6 lists the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>■ Added “Verilog HDL Configuration” on page 17–5.</td>
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<tr>
<td></td>
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<td>■ Added “RAM Style Attribute—For Shift Registers Inference” on page 17–55.</td>
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<tr>
<td></td>
<td></td>
<td>■ Added “Upgrade IP Components Dialog Box” on page 17–73.</td>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Updated “Design Flow” on page 17–1.</td>
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<tr>
<td></td>
<td></td>
<td>■ Added a link to Quartus II Help in “SystemVerilog Support” on page 14–5.</td>
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<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Updated “Verilog HDL Support” on page 13–4 to include Verilog-2001 support.</td>
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<tr>
<td></td>
<td></td>
<td>■ Updated “VHDL-2008 Support” on page 13–9 to include the condition operator (explicit and implicit) support.</td>
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<td></td>
<td></td>
<td>■ Rewrote “Limiting Resource Usage in Partitions” on page 13–32.</td>
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<td></td>
<td></td>
<td>■ Added “Creating LogicLock Regions” on page 13–32 and “Using Assignments to Limit the Number of RAM and DSP Blocks” on page 13–33.</td>
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<td></td>
<td></td>
<td>■ Updated “Turning Off the Add Pass-Through Logic to Inferred RAMs no_rw_check Attribute” on page 13–55.</td>
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<td>■ Updated “Auto Gated Clock Conversion” on page 13–28.</td>
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<td></td>
<td>■ Added links to Quartus II Help.</td>
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<td>July 2010</td>
<td>10.0.0</td>
<td>■ Removed Referenced Documents section.</td>
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<td></td>
<td></td>
<td>■ Added “Synthesis Seed” on page 9–36 section.</td>
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<td>■ Updated the following sections:</td>
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<td>■ “SystemVerilog Support” on page 9–5</td>
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<td></td>
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<td>■ “VHDL-2008 Support” on page 9–10</td>
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<td></td>
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<td>■ “Using Parameters/Generics” on page 9–16</td>
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<td>■ “Parallel Synthesis” on page 9–21</td>
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<td>■ “Limiting Resource Usage in Partitions” on page 9–32</td>
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<td>■ “Synthesis Effort” on page 9–35</td>
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<td>■ “Synthesis Attributes” on page 9–25</td>
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<td>■ “Synthesis Directives” on page 9–27</td>
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<td>■ “Auto Gated Clock Conversion” on page 9–29</td>
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<td>■ “State Machine Processing” on page 9–36</td>
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<td></td>
<td>■ “Multiply-Accumulators and Multiply-Adders” on page 9–50</td>
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<td>■ “Resource Aware RAM, ROM, and Shift-Register Inference” on page 9–52</td>
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<td>■ “RAM Style and ROM Style—for Inferred Memory” on page 9–53</td>
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<td>■ “Turning Off the Add Pass-Through Logic to Inferred RAMs no_rw_check Attribute” on page 9–55</td>
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<td></td>
<td></td>
<td>■ “Using altera_attribute to Set Quartus II Logic Options” on page 9–68</td>
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<tr>
<td></td>
<td></td>
<td>■ “Adding an HDL File to a Project and Setting the HDL Version” on page 9–83</td>
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<td>■ “Creating Design Partitions for Incremental Compilation” on page 9–85</td>
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<td></td>
<td>■ “Inferring Multiplier, DSP, and Memory Functions from HDL Code” on page 9–50</td>
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<td></td>
<td>■ Updated Table 9–9 on page 9–86.</td>
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<tr>
<td>December 2009</td>
<td>9.1.1</td>
<td>■ Added information clarifying inheritance of Synthesis settings by lower-level entities, including Altera and third-party IP</td>
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<td></td>
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<td>■ Updated “Keep Combinational Node/Implement as Output of Logic Cell” on page 9–46</td>
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Table 17–6. Document Revision History (Part 3 of 3)

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<td>■ Updated the following sections:</td>
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<td>■ “Initial Constructs and Memory System Tasks” on page 9–7</td>
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<td>■ “VHDL Support” on page 9–9</td>
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<td>■ “Parallel Synthesis” on page 9–21</td>
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<td></td>
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<td>■ “Synthesis Directives” on page 9–27</td>
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<td>■ “Timing-Driven Synthesis” on page 9–31</td>
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<td>■ “Safe State Machines” on page 9–40</td>
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<td>■ “RAM Style and ROM Style—for Inferred Memory” on page 9–53</td>
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<td>■ “Translate Off and On / Synthesis Off and On” on page 9–62</td>
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<td>■ “Read Comments as HDL” on page 9–63</td>
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<td>■ “Adding an HDL File to a Project and Setting the HDL Version” on page 9–81</td>
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<td></td>
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<td>■ Removed “Remove Redundant Logic Cells” section</td>
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<td></td>
<td>■ Added “Resource Aware RAM, ROM, and Shift-Register Inference” section</td>
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<td>■ Updated Table 9–9 on page 9–83</td>
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<td>March 2009</td>
<td>9.0.0</td>
<td>■ Updated Table 9–9.</td>
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<td>■ “Partitions for Preserving Hierarchical Boundaries” on page 9–20</td>
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<td>■ “Analysis &amp; Synthesis Settings Page of the Settings Dialog Box” on page 9–24</td>
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<td>■ “Timing-Driven Synthesis” on page 9–30</td>
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<td></td>
<td>■ Added “Parallel Synthesis” on page 9–21</td>
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<tr>
<td></td>
<td></td>
<td>■ Chapter 9 was previously Chapter 8 in software version 8.1</td>
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For previous versions of the Quartus II Handbook, refer to the ‘Quartus II Handbook Archive.’
This chapter documents support for the Synopsys Synplify software in the Quartus® II software, as well as key design flows, methodologies, and techniques for achieving optimal results in Altera® devices.

This chapter includes the following topics:

- General design flow with the Synplify and Quartus II software
- Synplify software optimization strategies, including timing-driven compilation settings, optimization options, and Altera-specific attributes
- Exporting designs and constraints to the Quartus II software using NativeLink integration
- Guidelines for Altera megafuctions and library of parameterized module (LPM) functions, instantiating them with the MegaWizard™ Plug-In Manager, and tips for inferring them from hardware description language (HDL) code
- Incremental compilation and block-based design, including the MultiPoint flow in the Synplify Pro and Synplify Premier software

The content in this chapter applies to the Synplify, Synplify Pro, and Synplify Premier software unless otherwise specified. This chapter includes the following sections:

- “Altera Device Family Support”
- “Design Flow” on page 18–2
- “Synplify Optimization Strategies” on page 18–6
- “Exporting Designs to the Quartus II Software Using NativeLink Integration” on page 18–13
- “Guidelines for Altera Megafuctions and Architecture-Specific Features” on page 18–16
- “Incremental Compilation and Block-Based Design” on page 18–29

This chapter assumes that you have set up, licensed, and are familiar with the Synplify software.

Altera Device Family Support

The Synplify software supports active devices available in the current version of the Quartus II software. Support for newly released device families may require an overlay. Contact Synopsys at www.synopsys.com for more information.
The Synplify software also supports the FLEX 8000 and MAX 9000 legacy devices that are supported only in the Altera MAX+PLUS® II software, as well as ACEX® 1K, APEX™ II, APEX 20K, APEX 20KC, APEX 20KE, FLEX® 10K, and FLEX 6000 legacy devices that are supported by the Quartus II software version 9.0 and earlier.

Design Flow

The following steps describe a basic Quartus II software design flow using the Synplify software:

1. Create Verilog HDL or VHDL design files.
2. Set up a project in the Synplify software and add the HDL design files for synthesis.
3. Select a target device and add timing constraints and compiler directives in the Synplify software to help optimize the design during synthesis.
4. Synthesize the project in the Synplify software.
5. Create a Quartus II project and import the following files generated by the Synplify software into the Quartus II software. Use the following files for placement and routing, and for performance evaluation:
   - The technology-specific Verilog Quartus Mapping File (.vqm) netlist or EDIF Input File (.edf) netlist for legacy devices also supported by the MAX+PLUS II software
   - The Synopsys Constraints Format (.scf) file for TimeQuest Timing Analyzer constraints

   If your design uses the Classic Timing Analyzer for timing analysis in the Quartus II software versions 10.0 and earlier, the Synplify software generates timing constraints in the Tcl Constraints File (.tcl). If you are using the Quartus II software versions 10.1 and later, you must use the TimeQuest Timing Analyzer for timing analysis.

   - The .tcl to set up your Quartus II project and pass constraints

   Alternatively, you can run the Quartus II software from within the Synplify software. For more information, refer to “Running the Quartus II Software from within the Synplify Software” on page 18–14.

6. After obtaining place-and-route results that meet your requirements, configure or program the Altera device.
Figure 18–1 shows the recommended design flow using the Synplify and the Quartus II software.

**Figure 18–1. Recommended Design Flow**

The Synplify software supports VHDL, Verilog HDL, and SystemVerilog source files. However, only the Synplify Pro and Premier software support mixed synthesis, allowing a combination of VHDL and Verilog HDL or SystemVerilog format source files.

Specify timing constraints and attributes for a design in a SCOPE Design Constraints File (.sdc) with the SCOPE window in the Synplify software using the standard Synopsys Design Constraint (SDC) format, or directly in the HDL source file. You can also define compiler directives in the HDL source file. Many of these constraints are forward-annotated for use by the Quartus II software. See Table 18–1 on page 18–4 for a list of the files generated by Synplify.
The HDL Analyst that is included in the Synplify software is a graphical tool for generating schematic views of the technology-independent RTL view netlist (.srs) and technology-view netlist (.srm) files. You can use the Synplify HDL Analyst to analyze and debug your design visually. The HDL Analyst supports cross-probing between the RTL and Technology views, the HDL source code, the Finite State Machine (FSM) viewer, and between the technology view and the timing report file in the Quartus II software.

A separate license file is required to enable the HDL Analyst in the Synplify software. The Synplify Pro and Premier software include the HDL Analyst.

After synthesis is complete, import the .vqm or .edf netlist to the Quartus II software for place-and-route. Use the .tcl file generated by the Synplify software to forward-annotate your project constraints including device selection, called the generated .scf file to forward-annotate TimeQuest Timing Analyzer timing constraints, and optionally to set up your project in the Quartus II software.

If area and timing requirements are satisfied, use the files generated by the Quartus II software to program or configure the Altera device. As shown in Figure 18–1, if your area or timing requirements are not met, you can change the constraints in the Synplify software or the Quartus II software and rerun synthesis. Altera recommends that you provide timing constraints in the Synplify software and any placement constraints in the Quartus II software. Repeat the process until area and timing requirements are met.

You can perform simulation and formal verification at various stages in the design process. You can perform final timing analysis after placement and routing is complete.

For more information about how the Synplify software supports formal verification, refer to Section V. Formal Verification in volume 3 of the Quartus II Handbook.

You can also use other options and techniques in the Quartus II software to meet area and timing requirements, such as WYSIWYG Primitive Resynthesis, which can perform optimizations on your .vqm netlist within the Quartus II software.

For information about netlist optimizations, refer to the Netlist Optimizations and Physical Synthesis chapter in volume 3 of the Quartus II Handbook.

In some cases, you might be required to modify the source code if the area and timing requirements cannot be met using options in the Synplify and Quartus II software.

During synthesis, the Synplify software produces several intermediate and output files, which are listed and described in Table 18–1.

### Table 18–1. Synplify Intermediate and Output Files (Part 1 of 2)

<table>
<thead>
<tr>
<th>File Extensions</th>
<th>File Description</th>
</tr>
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<tbody>
<tr>
<td>.srs</td>
<td>Technology-independent RTL netlist file that can be read only by the Synplify software.</td>
</tr>
<tr>
<td>.srm</td>
<td>Technology view netlist file.</td>
</tr>
</tbody>
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Table 18–1. Synplify Intermediate and Output Files (Part 2 of 2)

<table>
<thead>
<tr>
<th>File Extensions</th>
<th>File Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>.vqm/.edf</td>
<td>Technology-specific netlist in .vqm or .edf file format. A .vqm file is created for all Altera device families supported by the Quartus II software. An .edf file is created for devices supported by the MAX+PLUS II software.</td>
</tr>
<tr>
<td>.tcl</td>
<td>Forward-annotated constraints file containing constraints and assignments. A .tcl file for the Quartus II software is created for all devices. The .tcl file contains the appropriate Tcl commands to create and set up a Quartus II project and pass placement constraints.</td>
</tr>
<tr>
<td>.acf</td>
<td>Assignment and Configurations file for backward compatibility with the MAX+PLUS II software. For devices supported by the MAX+PLUS II software, the MAX+PLUS II assignments are imported from the MAX+PLUS II .acf file.</td>
</tr>
<tr>
<td>.scf</td>
<td>Synopsys Constraint Format file containing timing constraints for the TimeQuest Timing Analyzer.</td>
</tr>
</tbody>
</table>

Note to Table 18–1:
(1) This report file includes performance estimates that are often based on pre-place-and-route information. Use the fMax reported by the Quartus II software after place-and-route—it is the only reliable source of timing information. This report file includes post-synthesis device resource utilization statistics that might inaccurately predict resource usage after place-and-route. The Synplify software does not account for black box functions nor for logic usage reduction achieved through register packing performed by the Quartus II software. Register packing combines a single register and look-up table (LUT) into a single logic cell, reducing logic cell utilization below the Synplify software estimate. Use the device utilization reported by the Quartus II software after place-and-route.

Specifying the Output Netlist File Name and Result Format

The Result Format list specifies an .edf or .vqm netlist, depending on your device family. The software creates an .edf output netlist file only for devices supported by the MAX+PLUS II software. For current Altera devices, the software generates a .vqm-formatted netlist.

To specify the output netlist directory location, name, and format for the Synplify software, perform the following steps:
1. In the Synplify software, on the Project menu, click Implementation Options.
2. Click the Implementation Results tab.
3. In the Results Directory box, type your output netlist file directory location.
4. In the Result File Name box, type your output netlist file name.
5. In the Results Format list, specify either .vqm for devices using the Quartus II software or .edf for devices using the MAX+PLUS II software.

Specifying the Quartus II Software Version

You can specify your version of the Quartus II software in the Implementation Results tab of the Implementation Options dialog box. This option ensures that the netlist is compatible with the software version and supports the newest features. Altera recommends using the latest version of the Quartus II software whenever possible. If your Quartus II software version is newer than the versions available in the Quartus Version list, check if there is a newer version of the Synplify software available that supports the current Quartus II software version. Otherwise, select the latest version in the list for the best compatibility.

The Quartus Version list is available only after selecting an Altera device.
To specify the Quartus II software version used in the Synplify software, perform the following steps:

1. In the Synplify software, on the Project menu, click **Implementation Options**.
2. Click the **Implementation Results** tab.
3. Specify your version of the Quartus II software in the **Quartus Version** list.

Alternatively, type the following command at the command line:

```
set_option -quartus_version <version number>
```

## Synplify Optimization Strategies

Combining Synplify software constraints with VHDL and Verilog HDL coding techniques and Quartus II software options can help you obtain the results that you require. This section provides an overview of some of the techniques you can use to help improve the quality of your results.

For additional design and optimization techniques, refer to the *Design Recommendations for Altera Devices and the Quartus II Design Assistant* chapter in volume 1 and the *Area and Timing Optimization* chapter in volume 2 of the *Quartus II Handbook*.

For more information about applying the attributes discussed in this section, refer to the *Altera Constraints, Attributes, and Options* chapter of the *Synopsys FPGA Synthesis Reference Manual*.

## Using Synplify Premier to Optimize Your Design

Compared to other Synplify products, the Synplify Premier software offers additional physical synthesis optimizations. After typical logic synthesis, the Synplify Premier software places and routes the design and attempts to restructure the netlist based on the physical location of the logic in the Altera device. The Synplify Premier software forward-annotates the design netlist to the Quartus II software to perform the final placement and routing. In the default flow, the Synplify Premier software also forward-annotates placement information for the critical path(s) in the design, which can improve the compilation time in the Quartus II software.

The physical location annotation file is called `<design name>_plc.tcl`. If you open the Quartus II software from the Synplify Premier software user interface, the Quartus II software automatically uses this file for the placement information.

The Physical Analyst allows you to examine the placed netlist from the Synplify Premier software, which is similar to the HDL Analyst for a logical netlist. You can use this display to analyze and diagnose potential problems.
Using Implementations in Synplify Pro or Premier

To create different synthesis results without overwriting the existing results, in the Synplify Pro or Premier software, on the Project menu, click **New Implementation**. For each implementation, specify the target device, synthesis options, and constraint files. Each implementation generates its own subdirectory that contains all the resulting files, including `.vqm`, `.edf`, `.scf`, and `.tcl` files, from a compilation of the particular implementation. You can then compare the results of the different implementations to find the optimal set of synthesis options and constraints for a design.

Timing-Driven Synthesis Settings

The Synplify software supports timing-driven synthesis with user-assigned timing constraints to optimize the performance of the design. This section explains important timing constraints in the Synplify software.

The Quartus II NativeLink feature allows timing constraints that are applied in the Synplify software to be forward-annotated for the Quartus II software with an `.scf` file for timing-driven place and route. Refer to “Passing TimeQuest SDC Timing Constraints to the Quartus II Software” on page 18–15 for more details about how constraints such as clock frequencies, false paths, and multicycle paths are forward-annotated.

The Synplify Synthesis Report File (`.srr`) contains timing reports of estimated place-and-route delays. The Quartus II software can perform further optimizations on a post-synthesis netlist from third-party synthesis tools. In addition, designs might contain black boxes or intellectual property (IP) functions that have not been optimized by the third-party synthesis software. Actual timing results are obtained only after the design has been fully placed and routed in the Quartus II software. For these reasons, the Quartus II post place-and-route timing reports provide a more accurate representation of the design. Use the statistics in these reports to evaluate design performance.

Clock Frequencies

For single-clock designs, you can specify a global frequency when using the push-button flow. While this flow is simple and provides good results, it often does not meet the performance requirements for more advanced designs. You can use timing constraints, compiler directives, and other attributes to help optimize the performance of a design. You can enter these attributes and directives directly in the HDL code. Alternatively, you can enter attributes (not directives) into an `.sdc` file with the SCOPE window in the Synplify software.

Use the SCOPE window to set global frequency requirements for the entire design and individual clock settings. Use the **Clocks** tab in the SCOPE window to specify frequency (or period), rise times, fall times, duty cycle, and other settings. Assigning individual clock settings, rather than over-constraining the global frequency, helps the Quartus II software and the Synplify software achieve the fastest clock frequency for the overall design. The `define_clock` attribute assigns clock constraints.
Multiple Clock Domains

The Synplify software can perform timing analysis on unrelated clock domains. Each clock group is a different clock domain and is treated as unrelated to the clocks in all other clock groups. All clocks in a single clock group are assumed to be related, and the Synplify software automatically calculates the relationship between the clocks. You can assign clocks to a new clock group or put related clocks in the same clock group with the Clocks tab in the SCOPE window, or with the define_clock attribute.

Input and Output Delays

Specify the input and output delays for the ports of a design in the Input/Output tab of the SCOPE window, or with the define_input_delay and define_output_delay attributes. The Synplify software does not allow you to assign the tCO and tSU values directly to inputs and outputs. However, a tCO value can be inferred by setting an external output delay; a tSU value can be inferred by setting an external input delay.

Equation 18–1 illustrates the relationship between tCO and the output delay:

**Equation 18–1.**

\[ t_{CO} = \text{clock period} - \text{external output delay} \]

Equation 18–2 illustrates the relationship between tSU and the input delay:

**Equation 18–2.**

\[ t_{SU} = \text{clock period} - \text{external input delay} \]

When the syn_forward_io_constraints attribute is set to 1, the Synplify software passes the external input and output delays to the Quartus II software using NativeLink integration. The Quartus II software then uses the external delays to calculate the maximum system frequency.

Multicycle Paths

A multicycle path is a path that requires more than one clock cycle to propagate. Specify any multicycle paths in the design in the Multi-Cycle Paths tab of the SCOPE window, or with the define_multicycle_path attribute. You should specify which paths are multicycle to prevent the Quartus II and the Synplify compilers from working excessively on a non-critical path. Not specifying these paths can also result in an inaccurate critical path reported during timing analysis.

False Paths

False paths are paths that should be ignored during timing analysis, or be assigned low (or no) priority during optimization. Some examples of false paths include slow asynchronous resets, and test logic that has been added to the design. Set these paths in the False Paths tab of the SCOPE window, or use the define_false_path attribute.
FSM Compiler

If the FSM Compiler is turned on, the compiler automatically detects state machines in a design, which are then extracted and optimized. The FSM Compiler analyzes state machines and implements sequential, gray, or one-hot encoding, based on the number of states. The compiler also performs unused-state analysis, optimization of unreachable states, and minimization of transition logic. Implementation is based on the number of states, regardless of the coding style in the HDL code.

If the FSM Compiler is turned off, the compiler does not optimize logic as state machines. The state machines are implemented as coded in the HDL code. Thus, if the coding style for a state machine is sequential, the implementation is also sequential.

Use the `syn_state_machine` compiler directive to specify or prevent a state machine from being extracted and optimized. To override the default encoding of the FSM Compiler, use the `syn_encoding` directive.

The values for the `syn_encoding` directive are described in Table 18–2.

Table 18–2. *syn_encoding* Directive Values

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sequential</td>
<td>Generates state machines with the fewest possible flipflops. Sequential, also called binary, state machines are useful for area-critical designs when timing is not the primary concern.</td>
</tr>
<tr>
<td>Gray</td>
<td>Generates state machines where only one flipflop changes during each transition. Gray-encoded state machines tend to be free of glitches.</td>
</tr>
<tr>
<td>One-hot</td>
<td>Generates state machines containing one flipflop for each state. One-hot state machines typically provide the best performance and shortest clock-to-output delays. However, one-hot implementations are usually larger than sequential implementations.</td>
</tr>
<tr>
<td>Safe</td>
<td>Generates extra control logic to force the state machine to the reset state if an invalid state is reached. You can use the safe value in conjunction with any of the other three values, which results in the state machine being implemented with the requested encoding scheme and the generation of the reset logic.</td>
</tr>
</tbody>
</table>

Example 18–1 shows sample VHDL code for applying the `syn_encoding` directive.

**Example 18–1. Sample VHDL Code for syn_encoding**

```vhdl
SIGNAL current_state : STD_LOGIC_VECTOR(7 DOWNTO 0);
ATTRIBUTE syn_encoding : STRING;
ATTRIBUTE syn_encoding OF current_state : SIGNAL IS "sequential";
```

By default, the state machine logic is optimized for speed and area, which may be potentially undesirable for critical systems. The safe value generates extra control logic to force the state machine to the reset state if an invalid state is reached.

**FSM Explorer in Synplify Pro and Premier**

The Synplify Pro and Premier software use the FSM Explorer to explore different encoding styles for a state machine automatically, and then implement the best encoding based on the overall design constraints. The FSM Explorer uses the FSM Compiler to identify and extract state machines from a design. However, unlike the FSM Compiler, which chooses the encoding style based on the number of states, the FSM Explorer attempts several different encoding styles before choosing a specific one. The trade-off is that the compilation requires more time to analyze the state machine, but finds an optimal encoding scheme for the state machine.
Optimization Attributes and Options

The following sections describe more attributes and options that you can modify in the Synplify software to improve your design performance.

Retiming in Synplify Pro and Premier

The Synplify Pro and Premier software can retime a design, which can improve the timing performance of sequential circuits by moving registers (register balancing) across combinational elements. Be aware that retimed registers incur name changes.

To retime your design, turn on Retiming in the Device tab in the Implementation Options section, or use the syn_allow_retiming attribute.

Maximum Fan-Out

When your design has critical path nets with high fan-out, use the syn_maxfan attribute to control the fan-out of the net. Setting this attribute for a specific net results in the replication of the driver of the net to reduce overall fan-out. The syn_maxfan attribute takes an integer value and applies it to inputs or registers. The syn_maxfan attribute cannot be used to duplicate control signals. The minimum allowed value of the attribute is 4. Using this attribute might result in increased logic resource utilization, thus straining routing resources, which can lead to long compilation times and difficult fitting.

If you must duplicate an output register or an output enable register, you can create a register for each output pin by using the syn_useioff attribute. Refer to “Register Packing” on page 18–10.

Preserving Nets

During synthesis, the compiler maintains ports, registers, and instantiated components. However, some nets cannot be maintained to create an optimized circuit. Applying the syn_keep directive overrides the optimization of the compiler and preserves the net during synthesis. The syn_keep directive is a Boolean data type value and can be applied to wires (Verilog HDL) and signals (VHDL). Setting the value to true preserves the net through synthesis.

Register Packing

Altera devices allow register packing into I/O cells. Altera recommends allowing the Quartus II software to make the I/O register assignments. However, you can control register packing with the syn_useioff attribute. The syn_useioff attribute is a Boolean data type value that can be applied to ports or entire modules. Setting the value to 1 instructs the compiler to pack the register into an I/O cell. Setting the value to 0 prevents register packing in both the Synplify and Quartus II software.

Resource Sharing

The Synplify software uses resource sharing techniques during synthesis, by default, to reduce area. Turning off the Resource Sharing option on the Options tab of the Implementation Options dialog box improves performance results for some designs. You can also turn off the option for a specific module with the syn_sharing attribute. If you turn off this option, be sure to check the results to verify improvement in timing performance. If there is no improvement, turn on Resource Sharing.
Preserving Hierarchy

The Synplify software performs cross-boundary optimization by default, which causes the design to flatten to allow optimization. You can use the syn_hier attribute to override the default compiler settings. The syn_hier attribute applies a string value to modules, architectures, or both. Setting the value to hard maintains the boundaries of a module, architecture, or both, but allows constant propagation. Setting the value to locked prevents all cross-boundary optimizations. Use the locked setting with the partition setting to create separate design blocks and multiple output netlists for incremental compilation, as described in “Using MultiPoint Synthesis with Incremental Compilation” on page 18–31.

By default, the Synplify software generates a hierarchical .vqm file. To flatten the file, set the syn_netlist_hierarchy attribute to 0.

Register Input and Output Delays

Two advanced options, define_reg_input_delay and define_reg_output_delay, can speed up paths feeding a register, or coming from a register, by a specific number of nanoseconds. The Synplify software attempts to meet the global clock frequency goals for a design as well as the individual clock frequency goals (set with the define_clock attribute). You can use these attributes to add a delay to paths feeding into or out of registers to further constrain critical paths. You can slow down a path that is too highly optimized by setting this attributes to a negative number.

The define_reg_input_delay and define_reg_output_delay options are useful to close timing if your design does not meet timing goals, because the routing delay after placement and routing exceeds the delay predicted by the Synplify software. Rerun synthesis using these options, specifying the actual routing delay (from place-and-route results) so that the tool can meet the required clock frequency. Synopsys recommends that for best results, do not make these assignments too aggressively. For example, you can increase the routing delay value, but do not also use the full routing delay from the last compilation.

In the SCOPE constraint window, the registers panel contains the following options:

- **Register**—Specifies the name of the register. If you have initialized a compiled design, select the name from the list.
- **Type**—Specifies whether the delay is an input or output delay.
- **Route**—Shrinks the effective period for the constrained registers by the specified value without affecting the clock period that is forward-annotated to the Quartus II software.

Use the following Tcl command syntax to specify an input or output register delay in nanoseconds.

**Example 18–2. Specifying an Input or Output Register Delay Using Tcl Command Syntax**

```tcl
define_reg_input_delay [<register>] -route <delay in ns>
define_reg_output_delay [<register>] -route <delay in ns>
```
**syn_direct_enable**

This attribute controls the assignment of a clock-enable net to the dedicated enable pin of a register. With this attribute, you can direct the Synplify mapper to use a particular net as the only clock enable when the design has multiple clock enable candidates.

To use this attribute as a compiler directive to infer registers with clock enables, enter the `syn_direct_enable` directive in your source code, instead of the SCOPE spreadsheet.

The `syn_direct_enable` data type is Boolean. A value of `1` or `true` enables net assignment to the clock-enable pin. The following is the syntax for Verilog HDL:

```verilog
object /* synthesis syn_direct_enable = 1 */;
```

**I/O Standard**

For certain Altera devices, specify the I/O standard type for an I/O pad in the design with the **I/O Standard** panel in the Synplify SCOPE window.

**Example 18–3** shows the Synplify SDC syntax for the `define_io_standard` constraint, in which the `delay_type` must be either `input_delay` or `output_delay`.

---

**Example 18–3. Synplify SDC Syntax for the define_io_standard Constraint**

```tcl
define_io_standard [-disable|-enable] {<ObjectName>} -delay_type [input_delay|output_delay] <columnTclName>{<value>} [<columnTclName>{<value}>] ...
```

---

For details about supported I/O standards, refer to the *Altera I/O Standards* section in the *Synopsys FPGA Synthesis Reference Manual*.

### Altera-Specific Attributes

You can use the attributes described in this section with specific Altera device features, which are forward-annotated to the Quartus II project, and are used during place-and-route.

#### altera_chip_pin_lc

Use the `altera_chip_pin_lc` attribute to make pin assignments. This attribute applies a string value to inputs and outputs. Use the attribute only on the ports of the top-level entity in the design. Do not use this attribute to assign pin locations from entities at lower levels of the design hierarchy.

The `altera_chip_pin_lc` attribute is not supported for any MAX series device.

In the SCOPE window, set the value of the `altera_chip_pin_lc` attribute to a pin number or a list of pin numbers.
Example 18–4 shows VHDL code for making location assignments for supported Altera devices. Pin location assignments for these devices are written to the output .tcl file.

```
ENTITY sample (data_in : IN STD_LOGIC_VECTOR (3 DOWNTO 0);
data_out : OUT STD_LOGIC_VECTOR (3 DOWNTO 0));
ATTRIBUTE altera_chip_pin_lc : STRING;
ATTRIBUTE altera_chip_pin_lc OF data_out : SIGNAL IS "14, 5, 16, 15";
```

The `data_out` signal is a 4-bit signal; `data_out[3]` is assigned to pin 14 and `data_out[0]` is assigned to pin 15.

**altera_io_powerup**

Use the `altera_io_powerup` attribute to define the power-up value of an I/O register that has no set or reset. This attribute applies a string value (`high` | `low`) to ports with I/O registers. By default, the power-up value of the I/O register is set to `low`.

**altera_io_opendrain**

Use the `altera_io_opendrain` attribute to specify open-drain mode I/O ports. This attribute applies a boolean data type value to outputs or bidirectional ports for devices that support open-drain mode.

### Exporting Designs to the Quartus II Software Using NativeLink Integration

The NativeLink feature in the Quartus II software facilitates the seamless transfer of information between the Quartus II software and EDA tools, and allows you to run other EDA design entry or synthesis, simulation, and timing analysis tools automatically from within the Quartus II software. After a design is synthesized in the Synplify software, a .vqm or .edf netlist file, an .scf file for TimeQuest Timing Analyzer timing constraints, and .tcl files are used to import the design into the Quartus II software for place-and-route. You can run the Quartus II software from within the Synplify software or as a stand-alone application. After you import the design into the Quartus II software, you can specify different options to further optimize the design.

When you are using NativeLink integration, the path to your project must not contain empty spaces. The Synplify software uses Tcl scripts to communicate with the Quartus II software, and the Tcl language does not accept arguments with empty spaces in the path.

Use NativeLink integration to integrate the Synplify software and Quartus II software with a single GUI for both synthesis and place-and-route operations. NativeLink integration allows you to run the Quartus II software from within the Synplify software GUI, or to run the Synplify software from within the Quartus II software GUI.

This section explains the different NativeLink flows and provides details about how constraints are passed to the Quartus II software, and describes the following topics:
Running the Quartus II Software from within the Synplify Software

To run the Quartus II software from within the Synplify software, you must set the QUARTUS_ROOTDIR environment variable to the Quartus II software installation directory located in the \altera\quartus\<version number>. You must set this environment variable to use the Synplify and Quartus II software together. Synplify also uses this variable to open the Quartus II software in the background and obtain detailed information for Altera megafunctions used in the design.

In the Windows operating system, set the environment variable using the Control Panel, System options. In the Advanced tab, click Environment Variables, and create a QUARTUS_ROOTDIR system variable.

In the Linux operating system, create an environment variable QUARTUS_ROOTDIR that points to the <home directory>/altera <version number> location.

You can create new place and route implementations with the New P&R button in the Synplify software GUI. Under each implementation, the Synplify Pro software creates a place-and-route implementation called pr_<number> Altera Place and Route. To run the Quartus II software in command-line mode after each synthesis run, use the text box to turn on the place-and-route implementation. The results of the place-and-route are written to a log file in the pr_<number> directory under the current implementation directory.

You can also use the commands in the Quartus II menu to run the Quartus II software at any time following a successful completion of synthesis. In the Synplify software, on the Options menu, click Quartus II and then choose one of the following commands:

- **Launch Quartus**—Opens the Quartus II software GUI and creates a Quartus II project with the synthesized output file, forward-annotated timing constraints, and pin assignments. Use this command to configure options for the project and to execute any Quartus II commands.
- **Run Background Compile**—Runs the Quartus II software in command-line mode with the project settings from the synthesis run. The results of the place-and-route are written to a log file.

The <project_name>_cons.tcl file is used to set up the Quartus II project and directs the <project_name>.tcl file to pass constraints from the Synplify software to the Quartus II software. By default, the <project_name>.tcl file contains device, timing, and location assignments. The <project_name>.tcl file contains the command to use the Synplify-generated .scf constraints file with the TimeQuest Timing Analyzer.
Using the Quartus II Software to Run the Synplify Software

You can set up the Quartus II software to run the Synplify software for synthesis with NativeLink integration. This feature allows you to use the Synplify software to quickly synthesize a design as part of a standard compilation in the Quartus II software. When you use this feature, the Synplify software does not use any timing constraints or assignments, such as incremental compilation partitions, that you have set in the Quartus II software.

For best results, Synopsys recommends that you set constraints in the Synplify software and use a Tcl script to pass these constraints to the Quartus II software, instead of opening the Synplify software from within the Quartus II software.

To set up the Synplify software in the Quartus II software, on the Tools menu, click Options. In the Options dialog box, click EDA Tool Options and specify the path of the Synplify or Synplify Pro software under Location of Executable.

For more information about using NativeLink integration with the Synplify software in the Quartus II software, refer to About Using the Synplify Software with the Quartus II Software in Quartus II Help.

Running the Synplify software with NativeLink integration is supported on both floating network and node-locked fixed PC licenses. Both types of licenses support batch mode compilation.

Running the Quartus II Software Manually With the Synplify-Generated Tcl Script

You can also run the Quartus II software with a Synplify-generated Tcl script. To run the Tcl scrip to set up your project assignments, perform the following steps:

1. Ensure the .vqm/.edf, .scf, and .tcl files are located in the same directory.
2. In the Quartus II software, on the View menu, point to Utility Windows and click Tcl Console. The Quartus II Tcl Console opens.
3. At the Tcl Console command prompt, type the following:
   ```
   source <path>/<project name>_cons.tcl
   ```

Passing TimeQuest SDC Timing Constraints to the Quartus II Software

The TimeQuest Timing Analyzer is a powerful ASIC-style timing analysis tool that validates the timing performance of all logic in your design using an industry standard constraints format, Synopsys Design Constraints (SDC). This section explains how timing constraints set in the Synplify software are passed to the Quartus II software for use with the TimeQuest Timing Analyzer.

The Synplify-generated .tcl file contains constraints for the Quartus II software, such as the device specification and any location constraints. Timing constraints are forward-annotated in the Synopsys Constraints Format (.scf) file.

For additional information about the TimeQuest Timing Analyzer, refer to the Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.
Synopsys recommends that you modify constraints using the SCOPE constraint editor window, rather than using the generated .sdc, .scf, or .tcl file.

The following list of Synplify constraints are converted to the equivalent Quartus II SDC commands and are forward-annotated to the Quartus II software in the .scf file:

- define_clock
- define_input_delay
- define_output_delay
- define_multicycle_path
- define_false_path

All Synplify constraints described in the following sections are mapped to SDC commands for the TimeQuest Timing Analyzer.

For syntax and arguments for these commands, refer to the applicable subsection or refer to Synplify Help. For a list of corresponding commands in the Quartus II software, refer to the Quartus II Help.

**Individual Clocks and Frequencies**

Specify clock frequencies for individual clocks in the Synplify software with the define_clock command. This command is passed to the Quartus II software with the create_clock command.

**Input and Output Delay**

Specify input delay and output delay constraints in the Synplify software with the define_input_delay and define_output_delay commands, respectively. These commands are passed to the Quartus II software with the set_input_delay and set_output_delay commands.

**Multicycle Path**

Specify a multicycle path constraint in the Synplify software with the define_multicycle_path command. This command is passed to the Quartus II software with the set_multicycle_path command.

**False Path**

Specify a false path constraint in the Synplify software with the define_false_path command. This command is passed to the Quartus II software with the set_false_path command.

**Guidelines for Altera Megafunctions and Architecture-Specific Features**

Altera provides parameterizable megafunctions, including LPMs, device-specific Altera megafunctions, IP available as Altera MegaCore® functions, and IP available through the Altera Megafunction Partners Program (AMPPSM). You can use megafunctions and IP functions by instantiating them in your HDL code, or by inferring certain megafunctions from generic HDL code.
You can instantiate a megafunction in your HDL code with the MegaWizard Plug-In Manager to parameterize the function, or instantiate the function using the port and parameter definition. The MegaWizard Plug-In Manager provides a graphical interface within the Quartus II software for customizing and parameterizing any available megafunction for the design. For more information about the MegaWizard Plug-In Manager flow with the Synplify software, refer to “Instantiating Altera Megafunctions With the MegaWizard Plug-In Manager” on page 18–17 and “Instantiating Intellectual Property With the MegaWizard Plug-In Manager and IP Toolbench” on page 18–19.

For more information about specific Altera megafunctions, refer to the Quartus II Help. For more information about IP functions, refer to the appropriate IP documentation.

The Synplify software also automatically recognizes certain types of HDL code, and infers the appropriate megafunction when a megafunction provides optimal results. The Synplify software provides options to control inference of certain types of megafunctions, as described in “Inferring Altera Megafunctions from HDL Code” on page 18–22.

For more information about instantiating versus inferring megafunctions, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. This chapter also provides details about using the MegaWizard Plug-In Manager in the Quartus II software and explains the files generated by the wizard, as well as coding style recommendations and HDL examples for inferring megafunctions in Altera devices.

**Instantiating Altera Megafunctions With the MegaWizard Plug-In Manager**

This section describes how to instantiate Altera megafunctions with the MegaWizard Plug-In Manager.

When you use the MegaWizard Plug-In Manager to set up and parameterize a megafunction, the MegaWizard Plug-In Manager creates a VHDL or Verilog HDL wrapper file `<output file>.v` that instantiates the megafunction.

The Synplify software uses the Quartus II timing and resource estimation netlist feature to report more accurate resource utilization and timing performance estimates, and leverages timing-driven optimization, instead of treating the megafunction as a “black box.” Including the MegaWizard-generated megafunction variation wrapper file in your Synplify project, gives the Synplify software complete information about the megafunction.

There is an option in the MegaWizard Plug-In Manager to generate a netlist for resource and timing estimation. This option is not recommended for the Synplify software because the software automatically generates this information in the background without a separate netlist. If you do create a separate netlist `<output file>_syn.v` and use that file in your synthesis project, you must also include the `<output file>.v` file in your Quartus II project.
Verify that the correct Quartus II version is specified in the Synplify software before compiling the MegaWizard-generated file to ensure that the software uses the correct library definitions for the megafuncon. The **Quartus Version** setting must match the version of the Quartus II software used to generate the customized megafuncon in the MegaWizard Plug-In Manager.

For details about how to set the Quartus II version in the Synplify software, refer to “Specifying the Quartus II Software Version” on page 18–5.

In addition, ensure that the **QUARTUS_ROOTDIR** environment variable specifies the installation directory location of the correct Quartus II version. The Synplify software uses this information to launch the Quartus II software in the background. The environment variable setting must match the version of the Quartus II software used to generate the customized megafuncon in the MegaWizard Plug-In Manager. Refer to “Using the Quartus II Software to Run the Synplify Software” on page 18–15 for more details.

### Instantiating Megafuncsms With MegaWizard Plug-In Manager-Generated Verilog HDL Files

If you turn on the `<output file>_inst.v` option on the last page of the MegaWizard interface, the MegaWizard Plug-In Manager generates a Verilog HDL instantiation template file for use in your Synplify design. The instantiation template file, `<output file>_inst.v`, helps to instantiate the megafuncon variation wrapper file, `<output file>.v`, in your top-level design. Include the megafuncon variation wrapper file `<output file>.v` in your Synplify project. The Synplify software includes the megafuncon information in the output .vqm netlist file. You do not need to include the MegaWizard-generated megafuncon variation wrapper file in your Quartus II project.

### Instantiating Megafuncsms with MegaWizard Plug-In Manager-Generated VHDL Files

If you turn on the `<output file>.cmp` and `<output file>_inst.vhd` options on the last page of the MegaWizard interface, the MegaWizard Plug-In Manager generates a VHDL component declaration file and a VHDL instantiation template file for use in your Synplify design. These files can help you instantiate the megafuncon variation wrapper file, `<output file>.vhd`, in your top-level design. Include the `<output file>.vhd` in your Synplify project. The Synplify software includes the megafuncon information in the output .vqm netlist file. You do not need to include the MegaWizard-generated megafuncon variation wrapper file in your Quartus II project.

### Changing Synplify’s Default Behavior for Instantiated Altera Megafuncsms

By default, the Synplify software automatically opens the Quartus II software in the background to generate a resource and timing estimation netlist for megafuncsms, as described in the previous sections.

You may want to change this behavior to reduce run times in the Synplify software, because generating the netlist files can take several minutes for large designs, or if the Synplify software cannot access your Quartus II software installation to generate the files. Changing this behavior might speed up the compilation time in the Synplify software, but the Quality of Results (QoR) might be reduced.
The Synplify software directs the Quartus II software to generate information in two ways:

- Some megafunctions provide a “clear box” model—the Synplify software fully synthesizes this model and includes the device architecture-specific primitives in the output .vqm netlist file.
- Other megafunctions provide a “grey box” model—the Synplify software reads the resource information, but the netlist does not contain all the logic functionality.

For these functions, the Synplify software uses the logic information for resource and timing estimation and optimization, and then instantiates the megafunction in the output .vqm netlist file so the Quartus II software can implement the appropriate device primitives. By default, the Synplify software uses the clear box model when available, and otherwise uses the grey box model. To change this behavior, perform the following steps:

1. In the Synplify software, click Implementation Options.
2. On the Device tab, specify one of the following values for the Altera Models option:
   - On—uses the clearbox model when available and the grey box model when the clearbox model is unavailable
   - clearbox_only—enables the clear box model, but not the grey box model
   - Off—turns off the feature entirely

### Instantiating Intellectual Property With the MegaWizard Plug-In Manager and IP Toolbench

Many Altera IP functions include a resource and timing estimation netlist that the Synplify software uses to report more accurate resource utilization and timing performance estimates, and leverage timing-driven optimization rather than a black box function.

To create this netlist file, perform the following steps:

1. Select the IP function in the MegaWizard Plug-In Manager.
2. Click Next to open the IP Toolbench.
3. Click Set Up Simulation, which sets up all the EDA options.
4. Turn on the Generate netlist option to generate a netlist for resource and timing estimation and click OK.
5. Click Generate to generate the netlist file.

The Quartus II software generates a file <output file>_syn.v. This netlist contains the grey box information for resource and timing estimation, but does not contain the actual implementation. Include this netlist file in your Synplify project. Next, include the megafuction variation wrapper file <output file>.v | vhd in the Quartus II project along with your Synplify .vqm output netlist.

If your IP function does not include a resource and timing estimation netlist, the Synplify software must treat the IP function as a black box. In this case, refer to the following subsections for details about creating black boxes.
For information about including Quartus II-specific files in your Synplify project so they are automatically passed to the Quartus II software along with the output .vqm file, refer to “Including Files for Quartus II Placement and Routing Only” on page 18–22.

**Instantiating Black Box IP Functions With Generated Verilog HDL Files**

Use the `syn_black_box` compiler directive to declare a module as a black box. The top-level design files must contain the IP port-mapping and a hollow-body module declaration. Apply the `syn_black_box` directive to the module declaration in the top-level file or a separate file included in the project so that the Synplify software recognizes the module is a black box. The software compiles successfully without this directive, but reports an additional warning message. Using this directive allows you to add other directives, as discussed in “Other Synplify Software Attributes for Creating Black Boxes” on page 18–21.

Example 18–5 shows a sample top-level file that instantiates `my_verilogIP.v`, which is a simplified customized variation generated by the MegaWizard Plug-In Manager and the IP Toolbench.

**Example 18–5. Sample Top-Level Verilog HDL Code with Black Box Instantiation of IP**

```verilog
module top (clk, count);
    input clk;
    output[7:0] count;
    my_verilogIP verilogIP_inst (.clock (clk), .q (count));
endmodule
// Module declaration
// The following attribute is added to create a
// black box for this module.
module my_verilogIP (clock, q) /* synthesis syn_black_box */;
    input clock;
    output[7:0] q;
endmodule
```

**Instantiating Black Box IP Functions With Generated VHDL Files**

Use the `syn_black_box` compiler directive to declare a component as a black box. The top-level design files must contain the megafuction variation component declaration and port-mapping. Apply the `syn_black_box` directive to the component declaration in the top-level file. The software compiles successfully without this directive, but reports an additional warning message. Using this directive allows you to add other directives, such as the ones in the “Other Synplify Software Attributes for Creating Black Boxes” section.
Example 18–6 shows a sample top-level file that instantiates my_vhdlIP.vhd, which is a simplified customized variation generated by the MegaWizard Plug-In Manager and the IP Toolbench.

Example 18–6. Sample Top-Level VHDL Code with Black Box Instantiation of IP

LIBRARY ieee;
USE ieee.std_logic_1164.all;

ENTITY top IS
  PORT (
    clk: IN STD_LOGIC ;
    count: OUT STD_LOGIC_VECTOR (7 DOWNTO 0)
  );
END top;

ARCHITECTURE rtl OF top IS
COMPONENT my_vhdlIP
  PORT (  
    clock: IN STD_LOGIC ;
    q: OUT STD_LOGIC_VECTOR (7 DOWNTO 0)
  );
end COMPONENT;
attribute syn_black_box : boolean;
attribute syn_black_box of my_vhdlIP: component is true;
BEGIN
  vhdlip_inst : my_vhdlIP PORT MAP (  
    clock => clk,
    q => count
  );
END rtl;

Other Synplify Software Attributes for Creating Black Boxes

Instantiating a function as a black box does not provide visibility into the function module for the synthesis tool. Thus, it does not take full advantage of the synthesis tool’s timing-driven optimization. For better timing optimization, especially if the black box does not have registered inputs and outputs, add timing models to black boxes by adding the syn_tpd, syn_tsu, and syn_tco attributes. Refer to Example 18–7 for a Verilog HDL example.

Example 18–7. Adding Timing Models to Black Boxes in Verilog HDL

module ram32x4(z,d,addr,we,clk);
/* synthesis syn_black_box syn_tco1="clk->z[3:0]=4.0"
syn_tpd1="addr[3:0]->z[3:0]=8.0"
syn_tsu1="addr[3:0]->clk=2.0"
syn_tsu2="we->clk=3.0" */
output[3:0]z;
input[3:0]d;
input[3:0]addr;
input we
input clk
endmodule

The following additional attributes are supported by the Synplify software to communicate details about the characteristics of the black box module within the HDL code:

- syn_resources—Specifies the resources used in a particular black box
black_box_pad_pin—Prevents mapping to I/O cells
black_box_tri_pin—Indicates a tri-stated signal

For more information about applying these attributes, refer to the Altera Constraints, Attributes, and Options chapter of the Synopsys FPGA Synthesis Reference Manual.

Including Files for Quartus II Placement and Routing Only

In the Synplify software, you can add files to your project that are used only during placement and routing in the Quartus II software. This can be useful if you have grey or black boxes for Synplify synthesis that require the full design files to be compiled in the Quartus II software.

To include the files for Quartus II place-and-route only, perform the following steps:

1. Add the files to the Synplify project as source files.
2. Right-click the file, and on the shortcut menu, click File options.
3. Turn on Use for Place and Route Only. You can also set the option in a script using the -job_owner par option.

For example, the commands in Example 18–8 define files for a Synplify project that includes a top-level design file, a grey box netlist file, an IP wrapper file, and an encrypted IP file. With these files, the Synplify software writes an empty instantiation of “core” in the .vqm file and uses the grey box netlist for resource and timing estimation. The files core.v and core_enc8b10b.v are not compiled by the Synplify software, but are copied into the place-and-route directory. The Quartus II software compiles these files to implement the “core” IP block.

Example 18–8. Commands to Define Files for a Synplify Project

```bash
add_file -verilog -job_owner par "core_enc8b10b.v"
add_file -verilog -job_owner par "core.v"
add_file -verilog "core_gb.v"
add_file -verilog "top.v"
```

Inferring Altera Megafuntions from HDL Code

The Synplify software uses Behavior Extraction Synthesis Technology (BEST) algorithms to infer high-level structures such as RAMs, ROMs, operators, FSMs, and DSP multiplication operations. Then, the Synplify software keeps the structures abstract for as long as possible in the synthesis process. This allows the use of technology-specific resources to implement these structures by inferring the appropriate Altera megafuntion when a megafuntion provides optimal results. The following sections outline some of the Synplify-specific details when inferring Altera megafuntions. The Synplify software provides options to control inference of certain types of megafuntions, which is also described in the following sections.

For coding style recommendations and examples for inferring megafuntions in Altera devices, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
Inferring Multipliers

Figure 18–2 shows the HDL Analyst view of an unsigned 8 × 8 multiplier with two pipeline stages after synthesis in the Synplify software. This multiplier is converted into an ALTMULT_ADD or ALTMULT_ACCUM megafunction. For devices with DSP blocks, the software might implement the function in a DSP block instead of regular logic, depending on device utilization. For some devices, the software maps directly to DSP block device primitives instead of instantiating a megafunction in the .vqm file.

Resource Balancing

While mapping multipliers to DSP blocks, the Synplify software performs resource balancing for optimum performance.

Altera devices have a fixed number of DSP blocks, which includes a fixed number of embedded multipliers. If the design uses more multipliers than are available, the Synplify software automatically maps the extra multipliers to logic elements (LEs), or adaptive logic modules (ALMs).

If a design uses more multipliers than are available in the DSP blocks, the Synplify software maps the multipliers in the critical paths to DSP blocks. Next, any wide multipliers, which might or might not be in the critical paths, are mapped to DSP blocks. Smaller multipliers and multipliers that are not in the critical paths might then be implemented in the logic (LEs or ALMs). This ensures that the design fits successfully in the device.

Controlling the DSP Block Inference

You can implement multipliers in DSP blocks or in logic in Altera devices that contain DSP blocks. You can control this implementation through attribute settings in the Synplify software.

Signal Level Attribute

You can control the implementation of individual multipliers by using the syn_multstyle attribute as shown in the following Verilog HDL code:

```
<signal_name> /* synthesis syn_multstyle = "logic" */;
```

where `<signal_name>` is the name of the signal.

The syn_multstyle attribute applies to wires only; it cannot be applied to registers.
Table 18–3 describes the signal level attribute values that control the implementation of the multipliers in the DSP blocks or LEs in the Synplify software.

<table>
<thead>
<tr>
<th>Attribute Name</th>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>syn_multstyle</td>
<td>lpm_mult</td>
<td>LPM function inferred and multipliers implemented in DSP blocks</td>
</tr>
<tr>
<td></td>
<td>logic</td>
<td>LPM function not inferred and multipliers implemented LEs by the Synplify software</td>
</tr>
<tr>
<td></td>
<td>block_mult</td>
<td>DSP megafunction is inferred and multipliers are mapped directly to DSP block device primitives (for supported devices)</td>
</tr>
</tbody>
</table>

Example 18–9 and Example 18–10 show simple Verilog HDL and VHDL code using the `syn_multstyle` attribute.

**Example 18–9. Signal Attributes for Controlling DSP Block Inference in Verilog HDL Code**

```verilog
module mult(a,b,c,r,en);
    input [7:0] a,b;
    output [15:0] r;
    input [15:0] c;
    input en;
    wire [15:0] temp /* synthesis syn_multstyle="logic" */;

    assign temp = a*b;
    assign r = en ? temp : c;
endmodule
```

**Example 18–10. Signal Attributes for Controlling DSP Block Inference in VHDL Code**

```vhdl
library ieee;
use ieee.std_logic_1164.all;
use ieee.std_logic_unsigned.all;

entity onereg is port (
    r  : out std_logic_vector(15 downto 0);
    en : in std_logic;
    a  : in std_logic_vector(7 downto 0);
    b  : in std_logic_vector(7 downto 0);
    c  : in std_logic_vector(15 downto 0)
);
end onereg;

architecture beh of onereg is
    signal temp : std_logic_vector(15 downto 0);
    attribute syn_multstyle : string;
    attribute syn_multstyle of temp : signal is "logic";

begin
    temp <= a * b;
    r <= temp when en='1' else c;
end beh;
```
Inferring RAM

When a RAM block is inferred from an HDL design, the Synplify software uses an Altera megafunction to target the device memory architecture. For some devices, the Synplify software maps directly to memory block device primitives instead of instantiating a megafunction in the .vqm file.

Follow these guidelines for the Synplify software to successfully infer RAM in a design:

- The address line must be at least two bits wide.
- Resets on the memory are not supported. Refer to the device family documentation for information about whether read and write ports must be synchronous.
- Some Verilog HDL statements with blocking assignments might not be mapped to RAM blocks, so avoid blocking statements when modeling RAMs in Verilog HDL.

For some device families, the syn_ramstyle attribute specifies the implementation to use for an inferred RAM. You can apply the syn_ramstyle attribute globally, to a module, or to a RAM instance, to specify registers or block_ram values. To turn off RAM inference, set the attribute value to registers.

When inferring RAM for some Altera device families, the Synplify software generates additional bypass logic. This logic is generated to resolve a half-cycle read/write behavior difference between the RTL and post-synthesis simulations. The RTL simulation shows the memory being updated on the positive edge of the clock; the post-synthesis simulation shows the memory being updated on the negative edge of the clock. To eliminate bypass logic, the output of the RAM must be registered. By adding this register, the output of the RAM is seen after a full clock cycle, by which time the update has occurred, thus eliminating the need for bypass logic.

For devices with TriMatrix memory blocks, disable the creation of glue logic by setting the syn_ramstyle value to no_rw_check. Set syn_ramstyle to no_rw_check to disable the creation of glue logic in dual-port mode.
Example 18–11 shows sample VHDL code for inferring dual-port RAM.

**Example 18–11. VHDL Code for Inferred Dual-Port RAM**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_signed.all;

ENTITY dualport_ram IS
  PORT ( data_out: OUT STD_LOGIC_VECTOR (7 DOWNTO 0);
         data_in: IN STD_LOGIC_VECTOR (7 DOWNTO 0);
         wr_addr, rd_addr: IN STD_LOGIC_VECTOR (6 DOWNTO 0);
         we: IN STD_LOGIC;
         clk: IN STD_LOGIC);
END dualport_ram;

ARCHITECTURE ram_infer OF dualport_ram IS
  TYPE Mem_Type IS ARRAY (127 DOWNTO 0) OF STD_LOGIC_VECTOR (7 DOWNTO 0);
  SIGNAL mem: Mem_Type;
  SIGNAL addr_reg: STD_LOGIC_VECTOR (6 DOWNTO 0);
  BEGIN
    data_out <= mem (CONV_INTEGER(rd_addr));
    PROCESS (clk, we, data_in) BEGIN
      IF (clk='1' AND clk'EVENT) THEN
        IF (we='1') THEN
          mem(CONV_INTEGER(wr_addr)) <= data_in;
        END IF;
      END IF;
    END PROCESS;
  END ram_infer;
```
Example 18–12 shows an example of the VHDL code preventing bypass logic for inferring dual-port RAM. The extra latency behavior stems from the inferring methodology and is not required when instantiating a megafunction.

**Example 18–12. VHDL Code for Inferred Dual-Port RAM Preventing Bypass Logic**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_signed.all;

ENTITY dualport_ram IS
PORT ( data_out: OUT STD_LOGIC_VECTOR (7 DOWNTO 0);
data_in : IN STD_LOGIC_VECTOR (7 DOWNTO 0);
wr_addr, rd_addr : IN STD_LOGIC_VECTOR (6 DOWNTO 0);
we : IN STD_LOGIC;
clk : IN STD_LOGIC);
END dualport_ram;

ARCHITECTURE ram_infer OF dualport_ram IS
TYPE Mem_Type IS ARRAY (127 DOWNTO 0) OF STD_LOGIC_VECTOR (7 DOWNTO 0);
SIGNAL mem : Mem_Type;
SIGNAL addr_reg : STD_LOGIC_VECTOR (6 DOWNTO 0);
SIGNAL tmp_out : STD_LOGIC_VECTOR(7 DOWNTO 0); --output register
BEGIN
  tmp_out <= mem (CONV_INTEGER(rd_addr));
  PROCESS (clk, we, data_in) BEGIN
    IF (clk='1' AND clk'EVENT) THEN
      IF (we='1') THEN
        mem(CONV_INTEGER(wr_addr)) <= data_in;
      END IF;
      data_out <= tmp_out; --registers output preventing
END IF;
  END PROCESS;
END ram_infer;
```

**RAM Initialization**

Use the Verilog HDL $readmemb or $readmemh system tasks in your HDL code to initialize RAM memories. The Synplify compiler forward-annotates the initialization values in the .srs (technology-independent RTL netlist) file and the mapper generates a corresponding hexadecimal memory initialization (.hex) file. One .hex file is created for each of the altsyncram megafunctions that are inferred in the design. The .hex file is associated with the altsyncram instance in the .vqm file using the init_file attribute.
Example 18–13 and Example 18–14 illustrate how RAM memories can be initialized through HDL code, and how the corresponding .hex file is generated using Verilog HDL.

**Example 18–13. Using $readmemb System Task to Initialize an Inferred RAM in Verilog HDL Code**

```
initial
begin
  $readmemb("mem.ini", mem);
end

always @(posedge clk)
begin
  raddr_reg <= raddr;
  if(we)
    mem[waddr] <= data;
end
```

**Example 18–14. Sample of .vqm Instance Containing Memory Initialization File from Example 18–13**

```
altsyncram mem_hex( .wren_a(we), .wren_b(GND),...);
defparam mem_hex.lpm_type = "altsyncram";
defparam mem_hex.operation_mode = "Dual_Port";
...
defparam mem_hex.init_file = "mem_hex.hex";
```

**Inferring ROM**

When a ROM block is inferred from an HDL design, the Synplify software uses an Altera megafunction to target the device memory architecture. For some devices, the Synplify software maps directly to memory block device atoms instead of instantiating a megafunction in the .vqm file. Follow these guidelines for the Synplify software to successfully infer ROM in a design:

- The address line must be at least two bits wide.
- The ROM must be at least half full.
- A CASE or IF statement must make 16 or more assignments using constant values of the same width.

**Inferring Shift Registers**

The Synplify software infers shift registers for sequential shift components so that they can be placed in dedicated memory blocks in supported device architectures using the ALTSHIFT_TAPS megafunction.

If necessary, set the implementation style with the syn_srlstyle attribute. If you do not want the components automatically mapped to shift registers, set the value to registers. You can set the value globally, or on individual modules or registers.

For some designs, turning off shift register inference improves the design performance.
Incremental Compilation and Block-Based Design

As designs become more complex and designers work in teams, a block-based incremental design flow is often an effective design approach. In an incremental compilation flow, you can make changes to part of the design while maintaining the placement and performance of unchanged parts of the design. Design iterations are made dramatically faster by focusing new compilations on particular design partitions and merging results with previous compilation results of other partitions. You can perform optimization on individual subblocks and then preserve the results before you integrate the blocks into a final design and optimize it at the top-level.

MultiPoint synthesis, which is available for certain device technologies in the Synplify Pro and Premier software, provides an automated block-based incremental synthesis flow. The MultiPoint feature manages a design hierarchy to let you design incrementally and synthesize designs that take too long for synthesis of the entire project. MultiPoint synthesis allows different netlist files to be created for different sections of a design hierarchy and supports the Quartus II incremental compilation methodology. This feature also ensures that only those sections of a design that have been updated are resynthesized when the design is compiled, reducing synthesis run time and preserving the results for the unchanged blocks. You can change and resynthesize one section of a design without affecting other sections.

You can also partition your design and create different netlist files manually with the Synplify software by creating a separate project for the logic in each partition of the design. Creating different netlist files for each partition of the design also means that each partition can be independent of the others.

Hierarchical design methodologies can improve the efficiency of your design process, providing better design reuse opportunities and fewer integration problems when working in a team environment. When you use these incremental synthesis methodologies, you can take advantage of incremental compilation in the Quartus II software. You can perform placement and routing on only the changed partitions of the design, which reduces place-and-route time and preserves your fitting results. Follow the guidelines in this section to help you optimize results with these methodologies.

The following steps describe the general incremental compilation flow when using these features of the Quartus II software:

1. Create Verilog HDL or VHDL design files.
2. Determine which hierarchical blocks you want to treat as separate partitions in your design.
3. Set up your design using the MultiPoint synthesis feature or separate projects so that a separate netlist file is created for each design partition.
4. If using separate projects, disable I/O pad insertion in the implementations for lower-level partitions.
5. Compile and map each partition in the Synplify software, making constraints as you would in a non-incremental design flow.
6. Import the .vqm netlist and .tcl file for each partition into the Quartus II software and set up the Quartus II project(s) for incremental compilation.
7. Compile your design in the Quartus II software and preserve the compilation results with the post-fit netlist in incremental compilation.

8. When you make design or synthesis optimization changes to part of your design, resynthesize only the partition you modified to generate a new netlist and .tcl file. Do not regenerate netlist files for the unmodified partitions.

9. Import the new netlist and .tcl file into the Quartus II software and recompile the design in the Quartus II software with incremental compilation.

For more information about creating partitions and using the incremental compilation in the Quartus II software, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*.

### Creating a Design with Separate Netlist Files for Incremental Compilation

The first stage of a hierarchical or incremental design flow is to ensure that different parts of your design do not affect each other. Ensure that you have separate netlists for each partition in your design so you can take advantage of incremental compilation in the Quartus II software. If the entire design is in one netlist file, changes in one partition might affect other partitions because of possible node name changes when you resynthesize the design.

To ensure proper functionality of the synthesis flow, create separate netlist files only for modules and entities. In addition, each module or entity requires its own design file. If two different modules are in the same design file, but are defined as being part of different partitions, incremental compilation cannot be maintained since both partitions must be recompiled when one module is changed.

Altera recommends that you register all inputs and outputs of each partition. This makes logic synchronous, and avoids any delay penalty on signals that cross partition boundaries.

If you use boundary tri-states in a lower-level block, the Synplify software pushes, or bubbles, the tri-states through the hierarchy to the top-level to use the tri-state drivers on output pins of Altera devices. Because bubbling tri-states requires optimizing through hierarchies, lower-level tri-states are not supported with a block-based compilation methodology. Use tri-state drivers only at the external output pins of the device and in the top-level block in the hierarchy.

For more detailed recommendations about designing your hierarchy and creating partitions, refer to the *Best Practices for Incremental Compilation Partitions and Floorplan Assignments* chapter in volume 1 of the *Quartus II Handbook*.

You can generate multiple .vqm netlist files with the MultiPoint synthesis flow in the Synplify Pro and Premier software, or by manually creating separate Synplify projects and creating a black box for each block that you want to designate as a separate design partition.

In the MultiPoint synthesis flow in the Synplify Pro and Premier software, you create multiple .vqm netlist files from one easy-to-manage, top-level synthesis project. By using the manual black box method, you have multiple synthesis projects, which might be required for certain team-based or bottom-up designs where a single top-level project is not desired.
After you have created multiple .vqm files using one of these two methods, you must create the appropriate Quartus II projects to place-and-route the design.

**Using MultiPoint Synthesis with Incremental Compilation**

This section describes how to generate multiple .vqm files using the Synplify Pro and Premier software MultiPoint synthesis flow. You must first set up your constraint file and Synplify options, then apply the appropriate Compile Point settings to write multiple .vqm files and create design partition assignments for incremental compilation.

**Set Compile Points and Create Constraint Files**

The MultiPoint flow lets you segment a design into smaller synthesis units, called Compile Points. The synthesis software treats each Compile Point as a partition for incremental mapping, which allows you to isolate and work on each Compile Point module as independent segments of the larger design without impacting other design modules. A design can have any number of Compile Points, and Compile Points can be nested. The top-level module is always treated as a Compile Point.

Compile Points are optimized in isolation from their parent, which can be another Compile Point or a top-level design. Each block created with a Compile Point is unaffected by critical paths or constraints on its parent or other blocks. A Compile Point is independent, with its own individual constraints. During synthesis, any Compile Points that have not yet been synthesized are synthesized before the top level. Nested Compile Points are synthesized before the parent Compile Points in which they are contained. When you apply the appropriate setting for the Compile Point, a separate netlist is created for that Compile Point, isolating that logic from any other logic in the design.

**Figure 18–3** shows an example of a design hierarchy that is split into multiple partitions. The top-level block of each partition can be synthesized as a separate Compile Point.

**Figure 18–3. Partitions in a Hierarchical Design**

![Diagram of hierarchical design with partitions A, B, C, D, E, and F.](image-url)
In this case, modules A, B, and F are Compile Points. The top-level Compile Point consists of the top-level block in the design (that is, block A in this example), including the logic that is not defined under another Compile Point. In this example, the design for top-level Compile Point A also includes the logic in one of its subblocks, C. Because block F is defined as its own Compile Point, it is not treated as part of the top-level Compile Point A. Another separate Compile Point B contains the logic in blocks B, D, and E. One netlist is created for the top-level module A and submodule C, another netlist is created for B and its submodules D and E, while a third netlist is created for F.

Apply Compile Points to the module, or to the architecture in the Synplify Pro SCOPE spreadsheet, or to the .sdc file. You cannot set a Compile Point in the Verilog HDL or VHDL source code. You can set the constraints manually using Tcl or by editing the .sdc file, or you can use one of two methods in the GUI, as described in the following subsections.

**Defining Compile Points With .tcl or .sdc Files**

To set Compile Points with a .tcl or .sdc file, use the `define_compile_point` command, as shown in Example 18–15.

**Example 18–15. The define_compile_point Command**

```tcl
define_compile_point [-disable] {<objname>} -type {locked, partition}
```

In Example 18–15, `<objname>` represents any module in the design. The Compile Point type `{locked, partition}` indicates that the Compile Point represents a partition for the Quartus II incremental compilation flow.

Each Compile Point has a set of constraint files that begin with the `define_current_design` command to set up the SCOPE environment, as follows:

```tcl
define_current_design {<my_module>}
```

**Defining Compile Points in the Top-Level SCOPE Window**

The following method requires that you create separate constraint files for the top-level and lower-level Compile Points:

1. In the top-level SCOPE window, select the Compile Points tab.
2. Select the modules that you want to define as Compile Points and set Type to locked, partition.
3. Manually create a constraint file for each module to set constraints for each Compile Point.

**Defining Compile Points by Creating a New SCOPE File**

When you use the following method, the lower-level constraint file is created automatically:

1. On the File menu, click New and select Constraint File.
2. On the Select File Type tab of the Create a New SCOPE File dialog box, select Compile Point.
3. Select the module you want to designate as a Compile Point and click OK. The software automatically sets the Compile Points in the top-level constraint file and creates a lower-level constraint file for each Compile Point.

**Additional Considerations for Compile Points**

To ensure that changes to a Compile Point do not affect the top-level parent module, turn off the *Update Compile Point Timing Data* option in the *Implementation Options* dialog box. If this option is turned on, updates to a child module can impact the top-level module.

You can apply the `syn_allowed_resources` attribute to any Compile Point view to restrict the number of resources for a particular module.

When using Compile Points with incremental compilation, be aware of the following restrictions:

- To use Compile Points effectively, you must provide timing constraints (timing budgeting) for each Compile Point; the more accurate the constraints, the better your results are. Constraints are not automatically budgeted, so manual time budgeting is essential. Altera recommends that you register all inputs and outputs of each partition. This avoids any logic delay penalty on signals that cross-partition boundaries.

- When using the Synplify attribute `syn_useioff` to pack registers in the I/O Elements (IOEs) of Altera devices, these registers must be in the top-level module. Otherwise, you must direct the Quartus II software to perform I/O register packing instead of the `syn_useioff` attribute. You can use the *Fast Input Register* or *Fast Output Register* options, or set I/O timing constraints and turn on *Optimize I/O cell register placement for timing* on the *Fitter Settings* page of the *Settings* dialog box in the Quartus II software.

- There is no incremental synthesis support for top-level logic; any logic in the top-level is resynthesized during every compilation in the Synplify software.

For more information about using Compile Points and setting Synplify attributes and constraints for both top-level and lower-level Compile Points, refer to the *Synopsys FPGA Synthesis User Guide* and the *Synopsys FPGA Synthesis Reference Manual* in the Synplify software.

**Creating a Quartus II Project for Compile Points and Multiple .vqm Files**

During compilation, the Synplify Pro and Premier software creates a `<top-level project>.tcl` file that provides the Quartus II software with the appropriate constraints and design partition assignments, creating a partition for each `.vqm` file along with the information to set up a Quartus II project. For details about using this Tcl script to set up your Quartus II project and pass your constraints, refer to “Running the Quartus II Software Manually With the Synplify-Generated Tcl Script” on page 18–15.

Depending on your design methodology, you can create one Quartus II project for all netlists or a separate Quartus II project for each netlist. In the standard incremental compilation design flow, you create design partition assignments and optional LogicLock™ floorplan location assignments for each partition in the design within a single Quartus II project. This methodology allows for the best quality of results and performance preservation during incremental changes to your design.
You might require a bottom-up design flow if each partition must be optimized separately, such as for third-party IP delivery. If you use this flow, Altera recommends you create a design floorplan to avoid placement conflicts between each partition. To follow this design flow in the Quartus II software, create separate Quartus II projects, export each design partition and incorporate it into a top-level design using the incremental compilation features to maintain placement results.

The following sections describe how to create the Quartus II projects for these two design flows.

**Creating a Single Quartus II Project for a Standard Incremental Compilation Flow**

Use the `<top-level project>.tcl` file that contains the Synplify assignments for all partitions within the project. This method allows you to import all the partitions into one Quartus II project and optimize all modules within the project at once, while taking advantage of the performance preservation and compilation-time reduction that incremental compilation offers. Figure 18–4 shows a visual representation of the design flow for the example design in Figure 18–3 on page 18–31.

**Figure 18–4. Design Flow Using Multiple .vqm Files with One Quartus II Project**
Creating Multiple Quartus II Projects for a Bottom-Up Incremental Compilation Flow

Use the `<lower-level compile point>`.tcl files that contain the Synplify assignments for each Compile Point. Generate multiple Quartus II projects, one for each partition and netlist in the design. The designers in the project can optimize their own partitions separately within the Quartus II software and export the results for their own partitions. Figure 18–5 shows a visual representation of the design flow for the example design in Figure 18–3 on page 18–31. You can export the optimized subdesigns and then import them into one top-level Quartus II project using incremental compilation to complete the design.

Figure 18–5. Design Flow Using Multiple .vqm Files with Multiple Quartus II Projects

Creating Multiple .vqm Files for a Incremental Compilation Flow With Separate Synplify Projects

This section describes how to manually generate multiple .vqm files for a incremental compilation flow with black boxes and separate Synplify projects for each design partition. This manual flow is supported by versions of the Synplify software without the MultiPoint Synthesis feature.

Manually Creating Multiple .vqm Files With Black Boxes

To create multiple .vqm files manually in the Synplify software, create a separate project for each lower-level module and top-level design that you want to maintain as a separate .vqm file for an incremental compilation partition. Implement black box instantiations of lower-level partitions in your top-level project.

When synthesizing the projects for the lower-level modules, perform the following steps:

1. In the Implementation Options dialog box, turn on Disable I/O Insertion for the target technology.
2. Read the HDL files for the modules.

Modules might include black box instantiations of lower-level modules that are also maintained as separate .vqm files.
3. Add constraints with the SCOPE constraint window.
4. Enter the clock frequency to ensure that the sub-design is correctly optimized.
5. In the Attributes tab, set syn_netlist_hierarchy to 0.

When synthesizing the top-level design project, perform the following steps:
1. In the Implementation Options dialog box, turn off Disable I/O Insertion for the target technology.
2. Read the HDL files for top-level designs.
3. Create black boxes using lower-level modules in the top-level design.
4. Add constraints with the SCOPE constraint window.
5. Enter the clock frequency to ensure that the design is correctly optimized.
6. In the Attributes tab, set syn_netlist_hierarchy to 0.

The following sections describe an example of black box implementation to create separate .vqm files. Figure 18–3 on page 18–31 shows an example of a design hierarchy that is split into multiple partitions.

The partition top contains the top-level block in the design (block A) and the logic that is not defined as part of another partition. In this example, the partition for top-level block A also includes the logic in one of its sub-blocks, block C. Because block F is contained in its own partition, it is not treated as part of the top-level partition A. Another separate partition, partition B, contains the logic in blocks B, D, and E. In a team-based design, engineers can work independently on the logic in different partitions. One netlist is created for the top-level module A and its submodule C, another netlist is created for module B and its submodules D and E, while a third netlist is created for module F.

To create multiple .vqm files for this design, follow these steps:
2. Generate a .vqm file for module F. Use F.v/.vhd as the source files.
3. Generate a top-level .vqm file for module A. Use A.v/.vhd and C.v/.vhd as the source files. Ensure that you use black box modules B and F, which were optimized separately in the previous steps.

Creating Black Boxes in Verilog HDL

Any design block that is not defined in the project, or included in the list of files to be read for a project, is treated as a black box by the software. Use the syn_black_box attribute to indicate that you intend to create a black box for the module. In Verilog HDL, you must provide an empty module declaration for a module that is treated as a black box.
Example 18–16 shows an example of the A.v top-level file. Follow the same procedure for lower-level files that also contain a black box for any module beneath the current level hierarchy.

**Example 18–16. Verilog HDL Black Box for Top-Level File A.v**

```verilog
class A (data_in, clk, e, ld, data_out);
    input data_in, clk, e, ld;
    output [15:0] data_out;

    wire [15:0] cnt_out;
    B U1 (.data_in (data_in), .clk(clk), .ld (ld), .data_out(cnt_out));
    F U2 (.d(cnt_out), .clk(clk), .e(e), .q(data_out));

    // Any other code in A.v goes here.
endmodule

// Empty Module Declarations of Sub-Blocks B and F follow here.
// These module declarations (including ports) are required for black boxes.
module B (data_in, clk, ld, data_out) /* synthesis syn_black_box */ ;
    input data_in, clk, ld;
    output [15:0] data_out;
endmodule

module F (d, clk, e, q) /* synthesis syn_black_box */ ;
    input [15:0] d;
    input clk, e;
    output [15:0] q;
endmodule
```

**Creating Black Boxes in VHDL**

Any design block that is not defined in the project, or included in the list of files to be read for a project, is treated as a black box by the software. Use the `syn_black_box` attribute to indicate that you intend to treat the component as a black box. In VHDL, you must have a component declaration for the black box.

Although VHDL is not case-sensitive, a .vqm (a subset of Verilog HDL) file is case-sensitive. Entity names and their port declarations are forwarded to the .vqm file. Black box names and port declarations are also passed to the .vqm file. To prevent case-based mismatches, use the same capitalization for black box and entity declarations in VHDL designs.
Example 18–17 shows an example of the A.vhd top-level file. Follow this same procedure for any lower-level files that contain a black box for any block beneath the current level of hierarchy.

Example 18–17. VHDL Black Box for Top-Level File A.vhd

LIBRARY ieee;
USE ieee.std_logic_1164.all;
LIBRARY synplify;
USE synplify.attributes.all;

ENTITY A IS
PORT ( data_in : IN INTEGER RANGE 0 TO 15;
       clk, e, ld : IN STD_LOGIC;
       data_out : OUT INTEGER RANGE 0 TO 15 );
END A;

ARCHITECTURE a_arch OF A IS

COMPONENT B PORT(
       data_in : IN INTEGER RANGE 0 TO 15;
       clk, ld : IN STD_LOGIC;
       d_out : OUT INTEGER RANGE 0 TO 15 );
END COMPONENT;

COMPONENT F PORT(
       d : IN INTEGER RANGE 0 TO 15;
       clk, e: IN STD_LOGIC;
       q : OUT INTEGER RANGE 0 TO 15 );
END COMPONENT;

attribute syn_black_box of B: component is true;
attribute syn_black_box of F: component is true;

-- Other component declarations in A.vhd go here
signal cnt_out : INTEGER RANGE 0 TO 15;
BEGIN

U1 : B
PORT MAP (
       data_in => data_in,
       clk => clk,
       ld => ld,
       d_out => cnt_out );

U2 : F
PORT MAP (
       d => cnt_out,
       clk => clk,
       e => e,
       q => data_out );

-- Any other code in A.vhd goes here
END a_arch;

After you complete the steps described in this section, you have a netlist file for each partition of the design. These files are ready for use with the incremental compilation flow in the Quartus II software.
Creating a Quartus II Project for Multiple .vqm Files

The Synplify software creates a .tcl file for each .vqm file that provides the Quartus II software with the appropriate constraints and information to set up a project. For details about using the Tcl script generated by the Synplify software to set up your Quartus II project and pass your constraints, refer to “Running the Quartus II Software Manually With the Synplify-Generated Tcl Script” on page 18–15.

Depending on your design methodology, you can create one Quartus II project for all netlists or a separate Quartus II project for each netlist. In the standard incremental compilation design flow, you create design partition assignments and optional LogicLock floorplan location assignments for each partition in the design within a single Quartus II project. This methodology allows for the best quality of results and performance preservation during incremental changes to your design. You might require a bottom-up design flow where each partition must be optimized separately, such as for third-party IP delivery.

To perform follow this design flow in the Quartus II software, create separate Quartus II projects, export each design partition and incorporate it into a top-level design using the incremental compilation features to maintain the results.

The following sections describe how to create the Quartus II projects for these two design flows.

Creating a Single Quartus II Project for a Standard Incremental Compilation Flow

Use the <top-level project>.tcl file that contains the Synplify assignments for the top-level design. This method allows you to import all of the partitions into one Quartus II project and optimize all modules within the project at once, taking advantage of the performance preservation and compilation time reduction offered by incremental compilation. Figure 18–6 shows a visual representation of the design flow for the example design in Figure 18–3 on page 18–31.

All of the constraints from the top-level project are passed to the Quartus II software in the top-level .tcl file, but constraints made in the lower-level projects within the Synplify software are not forward-annotated. Enter these constraints manually in your Quartus II project.

Figure 18–6. Design Flow Using Multiple .vqm Files with One Quartus II Project
Creating Multiple Quartus II Projects for a Bottom-Up Incremental Compilation Flow

Use the .tcl file that is created for each .vqm file by the Synplify software for each Synplify project. This method generates multiple Quartus II projects, one for each block in the design. The designers in the project can optimize their own blocks separately within the Quartus II software and export the placement of their own blocks. Figure 18–7 shows a visual representation of the design flow for the example in Figure 18–3 on page 18–31.

Designers should create a LogicLock region to create a design floorplan for each block to avoid conflicts between partitions. The top-level designer then imports all the blocks and assignments into the top-level project. This method allows each block in the design to be optimized separately and then imported into one top-level project.

Performing Incremental Compilation in the Quartus II Software

In a standard design flow using Multipoint Synthesis, the Synplify software uses the Quartus II top-level .tcl file to ensure that the two tools databases stay synchronized. The Tcl file creates, changes, or deletes partition assignments in the Quartus II software for Compile Points that you create, change, or delete in the Synplify software. However, if you create, change, or delete a partition in the Quartus II software, the Synplify software does not change your Compile Point settings. Make any corresponding change in your Synplify project to ensure that you create the correct .vqm files.

If you use the NativeLink integration feature described in “Using the Quartus II Software to Run the Synplify Software” on page 18–15, the Synplify software does not use any information about design partition assignments that you have set in the Quartus II software.

If you create netlist files with multiple Synplify projects, or if you do not use the Synplify Pro or Premier-generated .tcl files to update constraints in your Quartus II project, you must ensure that your Synplify .vqm netlists align with your Quartus II partition settings.
After you have set up your Quartus II project with .vqm netlist files as separate design partitions, set the appropriate Quartus II options to preserve your compilation results. On the Assignments menu, click Design Partitions Window. Change the Netlist Type to Post-Fit to preserve the previous compilation’s post-fit placement results. If you do not make these settings, the Quartus II software does not reuse the placement or routing results from the previous compilation.

You can take advantage of incremental compilation with your Synplify design to reduce compilation time in the Quartus II software and preserve the results for unchanged design blocks.

For more information about using Quartus II incremental compilation, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Conclusion

Taking advantage of the Synopsys Synplify and Altera Quartus II design flows allow you to control how your design files are prepared for the Quartus II place-and-route process, as well as improve performance and optimize a design for use with Altera devices.

Document Revision History

Table 18–4. Document Revision History  (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Changed to new document template.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed Classic Timing Analyzer support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed the “altera_implement_in_esb or altera_implement_in_eab” section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Edited the “Creating a Quartus II Project for Compile Points and Multiple .vqm Files” on page 14–33 section for changes with the incremental compilation flow.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Edited the “Creating a Quartus II Project for Multiple .vqm Files” on page 14–39 section for changes with the incremental compilation flow.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Editorial changes.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Minor updates for the Quartus II software version 10.0 release.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Minor updates for the Quartus II software version 9.1 release.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Added new section “Exporting Designs to the Quartus II Software Using NativeLink Integration” on page 14–14.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor updates for the Quartus II software version 9.0 release.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Chapter 10 was previously Chapter 9 in software version 8.1.</td>
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### Table 18–4. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
</table>
| November 2008 | 8.1.0   | ■ Changed to 8-1/2 x 11 page size  
  ■ Changed the chapter title from “Synplicity Synplify & Synplify Pro Support” to “Synopsys Synplify Support”  
  ■ Replaced references to Synplicity with references to Synopsys  
  ■ Added information about Synplify Premier  
  ■ Updated supported device list  
  ■ Added SystemVerilog information to Figure 14–1 |
| May 2008    | 8.0.0   | ■ Updated supported device list  
  ■ Updated constraint annotation information for the TimeQuest Timing Analyzer  
  ■ Updated RAM and MAC constraint limitations  
  ■ Revised Table 9–1  
  ■ Added new section “Changing Synplify’s Default Behavior for Instantiated Altera Megafunctions”  
  ■ Added new section “Instantiating Intellectual Property Using the MegaWizard Plug-In Manager and IP Toolbench”  
  ■ Added new section “Including Files for Quartus II Placement and Routing Only”  
  ■ Added new section “Additional Considerations for Compile Points”  
  ■ Removed section “Apply the LogicLock Attributes”  
  ■ Modified Figure 9–4, 9–43, 9–47, and 9–48  
  ■ Added new section “Performing Incremental Compilation in the Quartus II Software”  
  ■ Numerous text changes and additions throughout the chapter  
  ■ Renamed several sections  
  ■ Updated “Referenced Documents” section |

For previous versions of the Quartus II Handbook, refer to the [Quartus II Handbook Archive](#).
This chapter documents support for the Mentor Graphics® Precision RTL Synthesis and Precision RTL Plus Synthesis software in the Quartus® II software design flow, as well as key design methodologies and techniques for improving your results for Altera® devices.

The topics discussed in this chapter include:

- “Altera Device Family Support”
- “Design Flow” on page 19–2
- “Creating and Compiling a Project in the Precision Synthesis Software” on page 19–5
- “Mapping the Precision Synthesis Design” on page 19–5
- “Synthesizing the Design and Evaluating the Results” on page 19–9
- “Exporting Designs to the Quartus II Software Using NativeLink Integration” on page 19–10
- “Guidelines for Altera Megafuncions and Architecture-Specific Features” on page 19–15
- “Incremental Compilation and Block-Based Design” on page 19–24

This chapter assumes that you have set up, licensed, and installed the Precision Synthesis software and the Quartus II software. You must set up, license, and install the Precision RTL Plus Synthesis software if you want to use the incremental synthesis feature for incremental compilation and block-based design.

To obtain and license the Precision Synthesis software, refer to the Mentor Graphics website at www.mentor.com. To install and run the Precision Synthesis software and to set up your work environment, refer to the Precision Synthesis Installation Guide in the Precision Manuals Bookcase. To access the Manuals Bookcase in the Precision Synthesis software, click Help and select Open Manuals Bookcase.

Altera Device Family Support

The Precision Synthesis software supports active devices available in the current version of the Quartus II software. Support for newly released device families may require an overlay. Contact Mentor Graphics for more information.
The Precision Synthesis software also supports the FLEX 8000 and MAX 9000 legacy devices that are supported only in the Altera MAX+PLUS® II software, as well as ACEX® 1K, APEX™ II, APEX 20K, APEX 20KC, APEX 20KE, FLEX® 10K, and FLEX 6000 legacy devices that are supported by the Quartus II software version 9.0 and earlier.

**Design Flow**

The following steps describe a basic Quartus II design flow using the Precision Synthesis software:

1. Create Verilog HDL or VHDL design files.
2. Create a project in the Precision Synthesis software that contains the HDL files for your design, select your target device, and set global constraints. Refer to “Creating and Compiling a Project in the Precision Synthesis Software” on page 19–5 for details.
3. Compile the project in the Precision Synthesis software.
4. Add specific timing constraints, optimization attributes, and compiler directives to optimize the design during synthesis.

   For best results, Mentor Graphics recommends specifying constraints that are as close as possible to actual operating requirements. Properly setting clock and I/O constraints, assigning clock domains, and indicating false and multicycle paths guide the synthesis algorithms more accurately toward a suitable solution in the shortest synthesis time.

5. Synthesize the project in the Precision Synthesis software. With the design analysis and cross-probing capabilities of the Precision Synthesis software, you can identify and improve circuit area and performance issues using prelayout timing estimates.
6. Create a Quartus II project and import the following files generated by the Precision Synthesis software into the Quartus II project:
   - The technology-specific EDIF (.edf) netlist or Verilog Quartus Mapping File (.vqm) netlist
   - Synopsys Design Constraints File (.sdc) for TimeQuest Timing Analyzer constraints
If your design uses the Classic Timing Analyzer for timing analysis in the Quartus II software versions 10.0 and earlier, the Precision Synthesis software generates timing constraints in the Tcl Constraints File (.tcl). If you are using the Quartus II software versions 10.1 and later, you must use the TimeQuest Timing Analyzer for timing analysis.

- Tcl Script Files (.tcl) to set up your Quartus II project and pass constraints

You can run the Quartus II software from within the Precision Synthesis software, or run the Precision Synthesis software using the Quartus II software. Refer to “Running the Quartus II Software from within the Precision Synthesis Software” on page 19–10 and “Using the Quartus II Software to Run the Precision Synthesis Software” on page 19–12 for more information.

7. After obtaining place-and-route results that meet your requirements, configure or program the Altera device.

Figure 19–1 shows the Quartus II design flow using the Precision Synthesis software as described in these steps, which are further described in detail in this chapter.

Figure 19–1. Design Flow Using the Precision Synthesis Software and Quartus II Software
If your area or timing requirements are not met, you can change the constraints and resynthesize the design in the Precision Synthesis software, or you can change the constraints to optimize the design during place-and-route in the Quartus II software. Repeat the process until the area and timing requirements are met.

You can use other options and techniques in the Quartus II software to meet area and timing requirements. For example, the **WYSIWYG Primitive Resynthesis** option can perform optimizations on your EDIF netlist in the Quartus II software.

For more information about netlist optimizations, refer to the **Netlist Optimizations and Physical Synthesis** chapter in volume 2 of the *Quartus II Handbook*. For more recommendations about how to optimize your design, refer to the **Area and Timing Optimization** chapter in volume 2 of the *Quartus II Handbook*.

While simulation and analysis can be performed at various points in the design process, final timing analysis should be performed after placement and routing is complete.

During synthesis, the Precision Synthesis software produces several intermediate and output files, which are described in Table 19–1.

**Table 19–1. Precision Synthesis Software Intermediate and Output Files**

<table>
<thead>
<tr>
<th>File Extension</th>
<th>File Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>.psp</td>
<td>Precision Synthesis Project File.</td>
</tr>
<tr>
<td>.rep (1)</td>
<td>Synthesis Area and Timing Report File. Technology-specific netlist in .vqm or .edf file format. By default, the Precision Synthesis software creates .vqm files for Arria series, Cyclone series, and Stratix series devices, and creates .edf files for ACEX, APEX, FLEX, and MAX series devices. The Precision Synthesis software can create .edf files for all Altera devices supported by the Quartus II software, but defaults to creating .vqm files when the device is supported.</td>
</tr>
<tr>
<td>.vqm/.edf (2)</td>
<td>Forward-annotated Tcl assignments and constraints file. The <code>&lt;project name&gt;.tcl</code> file is generated for all devices. The .tcl file acts as the Quartus II Project Configuration file and is used to make basic project and placement assignments, and to create and compile a Quartus II project.</td>
</tr>
<tr>
<td>.acf</td>
<td>Assignment and Configurations file for backward compatibility with the MAX+PLUS II software. For devices supported by the MAX+PLUS II software, the MAX+PLUS II assignments are imported from the MAX+PLUS II .acf file.</td>
</tr>
<tr>
<td>.sdc</td>
<td>Quartus II timing constraints file in Synopsys Design Constraints format. This file is generated automatically if the device uses the TimeQuest Timing Analyzer by default in the Quartus II software, and has the naming convention <code>&lt;project name&gt;_pnr_constraints.sdc</code>. For more information about generating a TimeQuest constraint file, refer to “Exporting Designs to the Quartus II Software Using NativeLink Integration” on page 19–10.</td>
</tr>
</tbody>
</table>

Notes to Table 19–1:

1. The timing report file includes performance estimates that are based on pre-place-and-route information. Use the f\text{max} reported by the Quartus II software after place-and-route for accurate post-place-and-route timing information. The area report file includes post-synthesis device resource utilization statistics that can differ from the resource usage after place-and-route due to black boxes or further optimizations performed during placement and routing. Use the device utilization reported by the Quartus II software after place-and-route for final resource utilization results. See “Synthesizing the Design and Evaluating the Results” on page 19–9 for details.

2. The Precision Synthesis software-generated VQM file is supported by the Quartus II software version 10.1 and later.
Creating and Compiling a Project in the Precision Synthesis Software

After creating your design files, create a project in the Precision Synthesis software that contains the basic settings for compiling the design.

To create a project, follow these steps:

1. In the Precision Synthesis software, click New Project in the Design Bar on the left side of the GUI.
2. Specify the Project Name and the Project Folder. The implementation name of the design corresponds to this project name.
3. Add input files to the project by clicking Add Input Files in the Design Bar. The Precision Synthesis software automatically detects the top-level module/entity of the design and uses it to name the current implementation directory, logs, reports, and netlist files.
4. In the Design Bar, click Setup Design.
5. To specify a target device family, expand Altera and select the target device and speed grade.
6. If you want, you can set a global design frequency and/or default input and output delays. This constrains all clock paths and I/O pins in your design. Modify the settings for individual paths or pins that do not require such a setting.
7. On the Design Center tab, right-click the Output Files folder and click Output Options.
8. To generate additional HDL netlists for post-synthesis simulation, select the desired output format. The Precision Synthesis software generates a separate file for each selected type of file: EDIF and Verilog HDL or VHDL.
9. To compile the design into a technology-independent implementation, in the Design Bar, click Compile.

Mapping the Precision Synthesis Design

In the next steps, you set constraints and map the design to technology-specific cells. The Precision Synthesis software maps the design by default to the fastest possible implementation that meets your timing constraints. To accomplish this, you must specify timing requirements for the automatically determined clock sources. With this information, the Precision Synthesis software performs static timing analysis to determine the location of the critical timing paths. The Precision Synthesis software achieves the best results for your design when you set as many realistic constraints as possible. Be sure to set constraints for timing, mapping, false paths, multicycle paths, and other factors that control the structure of the implemented design.

Mentor Graphics recommends creating an .sdc file and adding this file to the Constraint Files section of the Project Files list. You can create this file with a text editor, by issuing command-line constraint parameters, or by directing the Precision Synthesis software to generate the file automatically the first time you synthesize your design. To create a constraint file with the user interface, set constraints on design objects (such as clocks, design blocks, or pins) in the Design Hierarchy browser. By
default, the Precision Synthesis software saves all timing constraints and attributes in
two files: *precision rtl.sdc* and *precision tech.sdc*. The *precision rtl.sdc* file contains
constraints set on the RTL-level database (post- compilation) and the
*precision tech.sdc* file contains constraints set on the gate-level database
(post- synthesis) located in the current implementation directory.

You can also enter constraints at the command line. After adding constraints at the
command line, update the .sdc file with the `update constraint file` command. You can
add constraints that change infrequently directly to the HDL source files with HDL
attributes or pragmas.

The Precision .sdc file contains all the constraints for the Precision Synthesis project.
For the Quartus II software, placement constraints are written in a .tcl file and timing
constraints for the TimeQuest Timing Analyzer are written in the Quartus II .sdc file.

For details about the syntax of Synopsys Design Constraint commands, refer to the
more details and examples of attributes, refer to the Attributes chapter in the Precision
Synthesis Reference Manual. To access these manuals in the Precision Synthesis
software, click Help and select Open Manuals Bookcase.

### Setting Timing Constraints

The Precision Synthesis software uses timing constraints, based on the industry-
standard .sdc file format, to deliver optimal results. Missing timing constraints can
result in incomplete timing analysis and might prevent timing errors from being
detected. The Precision Synthesis software provides constraint analysis prior to
synthesis to ensure that designs are fully and accurately constrained. The
*<project name>_pnr_constraints.sdc* file, which contains timing constraints in SDC
format, is generated in the Quartus II software.

Because the .sdc file format requires that timing constraints be set relative to defined
clocks, you must specify your clock constraints before applying any other timing
constraints.

You also can use multicycle path and false path assignments to relax requirements or
exclude nodes from timing requirements, which can improve area utilization and
allow the software optimizations to focus on the most critical parts of the design.

For details about the syntax of Synopsys Design Constraint commands, refer to the
access these manuals in the Precision Synthesis software, click Help and select Open
Manuals Bookcase.

### Setting Mapping Constraints

Mapping constraints affect how your design is mapped into the target Altera device.
You can set mapping constraints in the user interface, in HDL code, or with the
`set_attribute` command in the constraint file.
Assigning Pin Numbers and I/O Settings

The Precision Synthesis software supports assigning device pin numbers, I/O standards, drive strengths, and slew-rate settings to top-level ports of the design. You can set these timing constraints with the `set_attribute` command, the GUI, or by specifying synthesis attributes in your HDL code. These constraints are forward-annotated in the `<project name>.tcl` file that is read by the Quartus II software during place-and-route and do not affect synthesis.

You can use the `set_attribute` command in the Precision Synthesis software .sdc file format to specify pin number constraints, I/O standards, drive strengths, and slow slew-rate settings. Table 19–2 describes the format to use for entries in the Precision Synthesis software constraint file.

You can also specify these options in the GUI. To specify a pin number or other I/O setting in the Precision Synthesis GUI, follow these steps:

1. After compiling the design, expand Ports in the Design Hierarchy Browser.
2. Under Ports, expand Inputs or Outputs.
3. Right-click the desired pin name and select Set Input Constraints under Inputs or Set Output Constraints under Outputs.
4. Type the desired pin number on the Altera device in the Pin Number box in the Port Constraints dialog box.
5. Select the I/O standard from the IO_STANDARD list.
6. For output pins, you can also select a drive strength setting and slew rate setting using the DRIVE and SLOW SLEW lists.

You also can use synthesis attributes or pragmas in your HDL code to make these assignments. Example 19–1 and Example 19–2 show code samples that make a pin assignment in your HDL code.

### Table 19–2. Constraint File Settings

<table>
<thead>
<tr>
<th>Constraint</th>
<th>Entry Format for Precision Constraint File</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pin number</td>
<td><code>set_attribute -name PIN_NUMBER -value &quot;&lt;pin number&gt;&quot; -port &lt;port name&gt;</code></td>
</tr>
<tr>
<td>I/O standard</td>
<td><code>set_attribute -name IOSTANDARD -value '&lt;I/O Standard&gt;' -port &lt;port name&gt;</code></td>
</tr>
<tr>
<td>Drive strength</td>
<td><code>set_attribute -name DRIVE -value '&lt;drive strength in mA&gt;' -port &lt;port name&gt;</code></td>
</tr>
<tr>
<td>Slew rate</td>
<td>`set_attribute -name SLEW -value &quot;TRUE</td>
</tr>
</tbody>
</table>

You can also specify these options in the GUI. To specify a pin number or other I/O setting in the Precision Synthesis GUI, follow these steps:

1. After compiling the design, expand Ports in the Design Hierarchy Browser.
2. Under Ports, expand Inputs or Outputs.

   You also can assign I/O settings by right-clicking the pin in the Schematic Viewer.

3. Right-click the desired pin name and select Set Input Constraints under Inputs or Set Output Constraints under Outputs.
4. Type the desired pin number on the Altera device in the Pin Number box in the Port Constraints dialog box.
5. Select the I/O standard from the IO_STANDARD list.
6. For output pins, you can also select a drive strength setting and slew rate setting using the DRIVE and SLOW SLEW lists.

You also can use synthesis attributes or pragmas in your HDL code to make these assignments. Example 19–1 and Example 19–2 show code samples that make a pin assignment in your HDL code.

### Example 19–1. Verilog HDL Pin Assignment

```verilog
//pragma attribute clk pin_number P10;
```

### Example 19–2. VHDL Pin Assignment

```vhdl
attribute pin_number : string attribute pin_number of clk : signal is "P10";
```
You can use the same syntax to assign the I/O standard using the `IOSTANDARD` attribute, drive strength using the attribute `DRIVE`, and slew rate using the `SLEW` attribute.

For more details about attributes and how to set these attributes in your HDL code, refer to the *Precision Synthesis Reference Manual*. To access this manual, in the Precision Synthesis software, click *Help* and select *Open Manuals Bookcase*.

### Assigning I/O Registers

The Precision Synthesis software performs timing-driven I/O register mapping by default. You can force a register to the device’s IO element (IOE) using the Complex I/O constraint. This option does not apply if you turn off I/O pad insertion. Refer to “Disabling I/O Pad Insertion” on page 19–8 for more information.

To force an I/O register into the device’s IOE using the GUI, follow these steps:

1. After compiling the design, expand *Ports* in the Design Hierarchy browser.
2. Under *Ports*, expand *Inputs* or *Outputs*.
3. Under *Inputs* or *Outputs*, right-click the desired pin name, point to *Map Input Register to IO* or *Map Output Register to IO*, for input or output respectively, and click *True*.

You also can make the assignment by right-clicking on the pin in the Schematic Viewer.

For the Stratix series, Cyclone series, and the MAX II device families, the Precision Synthesis software can move an internal register to an I/O register without any restrictions on design hierarchy.

For more mature devices, the Precision Synthesis software can move an internal register to an I/O register only when the register exists in the top-level of the hierarchy. If the register is buried in the hierarchy, you must flatten the hierarchy so that the buried registers are moved to the top-level of the design.

### Disabling I/O Pad Insertion

The Precision Synthesis software assigns I/O pad atoms (device primitives used to represent the I/O pins and I/O registers) to all ports in the top-level of a design by default. In certain situations, you might not want the software to add I/O pads to all I/O pins in the design. The Quartus II software can compile a design without I/O pads; however, including I/O pads provides the Precision Synthesis software with more information about the top-level pins in the design.

#### Preventing the Precision Synthesis Software from Adding I/O Pads

If you are compiling a subdesign as a separate project, I/O pins cannot be primary inputs or outputs of the device; therefore, the I/O pins should not have an I/O pad associated with them. To prevent the Precision Synthesis software from adding I/O pads, perform the following steps:

1. On the Tools menu, click *Set Options*. The *Options* dialog box appears.
2. On the Optimization page, turn off Add IO Pads.

3. Click Apply.

These steps add the following command to the project file:

setup_design -addio=false

**Preventing the Precision Synthesis Software from Adding an I/O Pad on an Individual Pin**

To prevent I/O pad insertion on an individual pin when you are using a black box, such as DDR or a phase-locked loop (PLL), at the external ports of the design, perform the following steps:

1. After compiling the design, in the Design Hierarchy browser, expand Ports.

2. Under Ports, expand Inputs or Outputs.

3. Under Inputs or Outputs, right-click the desired pin name and click Set Input Constraints.

4. In the Port Constraints dialog box for the selected pin name, turn off Insert Pad.

You also can make this assignment by right-clicking the pin in the Schematic Viewer or by attaching the nopad attribute to the port in the HDL source code.

**Controlling Fan-Out on Data Nets**

Fan-out is defined as the number of nodes driven by an instance or top-level port. High fan-out nets can cause significant delays that result in an unroutable net. On a critical path, high fan-out nets can cause longer delays in a single net segment that result in the timing constraints not being met. To prevent this behavior, each device family has a global fan-out value set in the Precision Synthesis software library. In addition, the Quartus II software automatically routes high fan-out signals on global routing lines in the Altera device whenever possible.

To eliminate routability and timing issues associated with high fan-out nets, the Precision Synthesis software also allows you to override the library default value on a global or individual net basis. You can override the library value by setting a max_fanout attribute on the net.

**Synthesizing the Design and Evaluating the Results**

To synthesize the design for the target device, click Synthesize in the Precision Synthesis Design Bar. During synthesis, the Precision Synthesis software optimizes the compiled design, and then writes out netlists and reports to the implementation subdirectory of your working directory after the implementation is saved, using the following naming convention:

<project name>_impl_<number>

After synthesis is complete, you can evaluate the results for area and timing. The Precision RTL Synthesis User’s Manual on the Mentor Graphics website describes different results that can be evaluated in the software.
There are several schematic viewers available in the Precision Synthesis software: RTL schematic, Technology-mapped schematic, and Critical Path schematic. These analysis tools allow you to quickly and easily isolate the source of timing or area issues, and to make additional constraint or code changes to optimize the design.

**Obtaining Accurate Logic Utilization and Timing Analysis Reports**

Historically, designers have relied on post-synthesis logic utilization and timing reports to determine the amount of logic their design requires, the size of the device required, and how fast the design runs. However, today’s FPGA devices provide a wide variety of advanced features in addition to basic registers and look-up tables (LUTs). The Quartus II software has advanced algorithms to take advantage of these features, as well as optimization techniques to increase performance and reduce the amount of logic required for a given design. In addition, designs can contain black boxes and functions that take advantage of specific device features. Because of these advances, synthesis tool reports provide post-synthesis area and timing estimates, but you should use the place-and-route software to obtain final logic utilization and timing reports.

**Exporting Designs to the Quartus II Software Using NativeLink Integration**

The NativeLink feature in the Quartus II software facilitates the seamless transfer of information between the Quartus II software and EDA tools, which allows you to run other EDA design entry/synthesis, simulation, and timing analysis tools automatically from within the Quartus II software.

After a design is synthesized in the Precision Synthesis software, the technology-mapped design is written to the current implementation directory as an EDIF netlist file, along with a Quartus II Project Configuration File and a place-and-route constraints file. You can use the Project Configuration script, `<project name>.tcl`, to create and compile a Quartus II project for your EDIF or VQM netlist. This script makes basic project assignments, such as assigning the target device specified in the Precision Synthesis software. If you select an Arria GX, Stratix III, Cyclone III, or newer device, the constraints are written in SDC format to the `<project name>_pnr_constraints.sdc` file by default, which is used by the Fitter and the TimeQuest Timing Analyzer in the Quartus II software.

Use the following Precision Synthesis software command before compilation to generate the `<project name>_pnr_constraints.sdc`:

```
setup_design -timequest_sdc
```

With this command, the file is generated after the synthesis.

**Running the Quartus II Software from within the Precision Synthesis Software**

The Precision Synthesis software also has a built-in place-and-route environment that allows you to run the Quartus II Fitter and view the results in the Precision Synthesis GUI. This feature is useful when performing an initial compilation of your design to view post-place-and-route timing and device utilization results, but not all the advanced Quartus II options that control the compilation process are available.
After you specify an Altera device as the target, set the options for the Quartus II software. On the Tools menu, click **Set Options**. On the **Integrated Place and Route** page, under **Quartus II Modular**, specify the path to the Quartus II executables in the **Path to Quartus II installation tree** box.

To automate the place-and-route process, click **Run Quartus II** in the **Quartus II Modular** window of the Precision Synthesis toolbar. The Quartus II software uses the current implementation directory as the Quartus II project directory and runs a full compilation in the background (that is, the user interface does not appear).

Two primary Precision Synthesis software commands control the place-and-route process. Use the **setup_place_and_route** command to set the place-and-route options. Start the process with the **place_and_route** command.

Precision Synthesis software uses individual Quartus II executables, such as analysis and synthesis (**quartus_map**), Fitter (**quartus_fit**), and the TimeQuest Timing Analyzer (**quartus_sta**) for improved runtime and memory utilization during place and route. This flow is referred to as the **Quartus II Modular** flow option in the Precision Synthesis software. By default, the Precision Synthesis software generates a Quartus II Project Configuration File (**.tcl** file) for current device families. Timing constraints that you set during synthesis are exported to the Quartus II place-and-route constraints file `<project name>_<pnr_constraints.sdc`.

After you compile the design in the Quartus II software from within the Precision Synthesis software, you can invoke the Quartus II GUI manually and then open the project using the generated Quartus II project file. You can view reports, run analysis tools, specify options, and run the various processing flows available in the Quartus II software.

For more information about running the Quartus II software from within the Precision Synthesis software, refer to the **Altera Quartus II Integration** chapter in the **Precision Synthesis Reference Manual**. To access this manual in the Precision Synthesis software, click **Help** and select **Open Manuals Bookcase**.

### Running the Quartus II Software Manually Using the Precision Synthesis-Generated Tcl Script

You can run the Quartus II software using a Tcl script generated by the Precision Synthesis software. To run the Tcl script generated by the Precision Synthesis software to set up your project and start a full compilation, perform the following steps:

1. Ensure the **.edf** or **.vqm** file, **.tcl** files, and **.sdc** file are located in the same directory. The files should be located in the implementation directory by default.
2. In the Quartus II software, on the **View** menu, point to Utility Windows and click **Tcl Console**.
3. At the Tcl Console command prompt, type the command:
   ```tcl
   source <path>/<project name>.tcl
   ```
4. On the **File** menu, click **Open Project**. Browse to the project name and click **Open**.
5. Compile the project in the Quartus II software.
Using the Quartus II Software to Run the Precision Synthesis Software

With NativeLink integration, you can set up the Quartus II software to run the Precision Synthesis software. This feature allows you to use the Precision Synthesis software to synthesize a design as part of a standard compilation. When you use this feature, the Precision Synthesis software does not use any timing constraints or assignments, such as incremental compilation partitions, that you have set in the Quartus II software.

For detailed information about using NativeLink integration with the Precision Synthesis software, refer to Using the NativeLink Feature with Other EDA Tools in the Quartus II Help.

Passing Constraints to the Quartus II Software

The place-and-route constraints script forward-annotates timing constraints that you made in the Precision Synthesis software. This integration allows you to enter these constraints once in the Precision Synthesis software, and then pass them automatically to the Quartus II software.

Refer to the introductory text in the section “Exporting Designs to the Quartus II Software Using NativeLink Integration” on page 19–10 for information on how to ensure the Precision Synthesis software targets the TimeQuest Timing Analyzer.

The following constraints are translated by the Precision Synthesis software and are applicable to the TimeQuest Timing Analyzer:

- `create_clock`
- `set_input_delay`
- `set_output_delay`
- `set_max_delay`
- `set_min_delay`
- `set_false_path`
- `set_multicycle_path`

**Example 19–3. Specifying a Clock using `create_clock`**

```bash
create_clock -name <clock_name> -period <period in ns> -waveform {<edge_list>} -domain \ <ClockDomain> <pin>
```
The period is specified in units of nanoseconds (ns). If no clock domain is specified, the clock belongs to a default clock domain main. All clocks in the same clock domain are treated as synchronous (related) clocks. If no <clock_name> is provided, the default name virtual_default is used. The <edge_list> sets the rise and fall edges of the clock signal over an entire clock period. The first value in the list is a rising transition, typically the first rising transition after time zero. The waveform can contain any even number of alternating edges, and the edges listed should alternate between rising and falling. The position of any edge can be equal to or greater than zero but must be equal to or less than the clock period.

If -waveform <edge_list> is not specified and -period <period in ns> is specified, the default waveform has a rising edge of 0.0 and a falling edge of <period_value>/2.

The Precision Synthesis software maps the clock constraint to the TimeQuest create_clock setting in the Quartus II software.

The Quartus II software supports only clock waveforms with two edges in a clock cycle. If the Precision Synthesis software finds a multi-edge clock, it issues an error message when you synthesize your design in the Precision Synthesis software.

set_input_delay

This port-specific input delay constraint is specified in the Precision Synthesis software, as shown in Example 19–4.

Example 19–4. Specifying set_input_delay

```
set_input_delay {<delay_value> <port_pin_list>} -clock <clock_name> -rise -fall -add_delay
```

This constraint is mapped to the set_input_delay setting in the Quartus II software.

When the reference clock <clock_name> is not specified, all clocks are assumed to be the reference clocks for this assignment. The input pin name for the assignment can be an input pin name of a time group. The software can use the clock_fall option to specify delay relative to the falling edge of the clock.

Although the Precision Synthesis software allows you to set input delays on pins inside the design, these constraints are not sent to the Quartus II software, and a message is displayed.

set_output_delay

This port-specific output delay constraint is specified in the Precision Synthesis software, as shown in Example 19–5.

Example 19–5. Using the set_output_delay Constraint

```
set_output_delay {<delay_value> <port_pin_list>} -clock <clock_name> -rise -fall -add_delay
```

This constraint is mapped to the set_output_delay setting in the Quartus II software.

When the reference clock <clock_name> is not specified, all clocks are assumed to be the reference clocks for this assignment. The output pin name for the assignment can be an output pin name of a time group.
Although the Precision Synthesis software allows you to set output delays on pins inside the design, these constraints are not sent to the Quartus II software.

**set_max_delay and set_min_delay**

The maximum delay for a point-to-point timing path constraint is specified in the Precision Synthesis software, as shown in Example 19–6. The minimum delay for a point-to-point timing path constraint is shown in Example 19–7.

**Example 19–6. Using the set_max_delay Constraint**

```
set_max_delay -from {<from_node_list>} -to {<to_node_list>} <delay_value>
```

**Example 19–7. Using the set_min_delay Constraint**

```
set_min_delay -from {<from_node_list>} -to {<to_node_list>} <delay_value>
```

The set_max_delay and set_min_delay commands specify that the maximum and minimum respectively, required delay for any start point in <from_node_list> to any endpoint in <to_node_list> must be less than or greater than <delay_value>. Typically, you use these commands to override the default setup constraint for any path with a specific maximum or minimum time value for the path.

The node lists can contain a collection of clocks, registers, ports, pins, or cells. The -from and -to parameters specify the source (start point) and the destination (endpoint) of the timing path, respectively. The source list (<from_node_list>) cannot include output ports, and the destination list (<to_node_list>) cannot include input ports. If you include more than one node on a list, you must enclose the nodes in quotes or in braces ({{}}).

If you specify a clock in the source list, you must specify a clock in the destination list. Applying set_max_delay or set_min_delay setting between clocks applies the exception from all registers or ports driven by the source clock to all registers or ports driven by the destination clock. Applying exceptions between clocks is more efficient than applying them for specific node-to-node, or node-to-clock paths. If you want to specify pin names in the list, the source must be a clock pin and the destination must be any non-clock input pin to a register. Assignments from clock pins, or to and from cells, apply to all registers in the cell or for those driven by the clock pin.

**set_false_path**

The false path constraint is specified in the Precision Synthesis software, as shown in Example 19–8.

**Example 19–8. Using the set_false_path Constraint**

```
set_false_path -to <to_node_list> -from <from_node_list> -reset_path
```

The node lists can be a list of clocks, ports, instances, and pins. Multiple elements in the list can be represented using wildcards such as * and ?.

In a place-and-route Tcl constraints file, this false path setting in the Precision Synthesis software is mapped to a set_false_path setting. The Quartus II software supports setup, hold, rise, or fall options for this assignment.
The node lists for this assignment represent top-level ports and/or nets connected to instances (end points of timing assignments).

Any false path setting in the Precision Synthesis software can be mapped to a setting in the Quartus II software with a through path specification.

**set_multicycle_path**

This multicycle path constraint is specified in the Precision Synthesis software, as shown in Example 19–9.

**Example 19–9. Using the set_multicycle_path Constraint**

```
set_multicycle_path <multiplier_value> [-start] [-end] -to <to_node_list> -from <from_node_list> \\
-reset_path
```

The node list can contain clocks, ports, instances, and pins. Multiple elements in the list can be represented using wildcards such as * and ?. Paths without multicycle path definitions are identical to paths with multipliers of 1. To add one additional cycle to the datapath, use a multiplier value of 2. The option start indicates that source clock cycles should be considered for the multiplier. The option end indicates that destination clock cycles should be considered for the multiplier. The default is to reference the end clock.

In the place-and-route Tcl constraints file, the multicycle path setting in the Precision Synthesis software is mapped to a set_multicycle_path setting. The Quartus II software supports the rise or fall options on this assignment.

The node lists represent top-level ports and/or nets connected to instances (end points of timing assignments). The node lists can contain wildcards (such as *); the Quartus II software automatically expands all wildcards.

Any multicycle path setting in Precision Synthesis software can be mapped to a setting in the Quartus II software with a through specification.

**Guidelines for Altera Megafunctions and Architecture-Specific Features**

Altera provides parameterizable megafUNCTIONS, including the LPMs, device-specific Altera megafUNCTIONS, IP available as Altera MegaCore functions, and IP available through the Altera MegafUNCTION Partners Program (AMPPSM). You can use megafUNCTIONS and IP functions by instantiating them in your HDL code or by inferring certain megafUNCTIONS from generic HDL code.

If you want to instantiate a megafUNCTION such as a PLL in your HDL code, you can instantiate and parameterize the function using the port and parameter definitions, or you can customize a function with the MegaWizard Plug-In Manager. Altera recommends using the MegaWizard Plug-In Manager, which provides a graphical interface within the Quartus II software for customizing and parameterizing any available megafUNCTION for the design. “Instantiating Altera MegafUNCTIONS Using the MegaWizard Plug-In Manager” and “Instantiating Intellectual Property With the MegaWizard Plug-In Manager and IP Toolbench” on page 19–17 describe the MegaWizard Plug-In Manager flow with the Precision Synthesis software.
For more information about specific Altera megafunctions and IP functions, refer to the IP and Megafunctions page of the Altera website.

The Precision Synthesis software automatically recognizes certain types of HDL code and infers the appropriate function. The Precision Synthesis software provides options to control inference of certain types of megafunctions, as described in “Inferring Altera Megafunctions from HDL Code” on page 19–19.

For a detailed discussion about instantiating functions versus inferring functions to target Altera architecture-specific features, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. This chapter also provides details on using the MegaWizard Plug-In Manager in the Quartus II software and explains the files generated by the wizard, as well as coding style recommendations and HDL examples for inferring functions in Altera devices.

**Instantiating Altera Megafaunctions Using the MegaWizard Plug-In Manager**

This section describes how to instantiate Altera megafunctions with the MegaWizard Plug-In Manager, and how to generate the files that are included in the Precision Synthesis project for synthesis.

You can run the stand-alone version of the MegaWizard Plug-In Manager by typing the following command at a command prompt:

```
qmegawiz
```

**Instantiating Megafunctions With MegaWizard Plug-In Manager-Generated Verilog HDL Files**

The MegaWizard Plug-In Manager generates a Verilog HDL instantiation template file `<output file>_inst.v` and a hollow-body black box module declaration `<output file>_bb.v` for use in your Precision Synthesis design. Incorporate the instantiation template file, `<output file>_inst.v`, into your top-level design to instantiate the megafunction wrapper file, `<output file>.v`.

Include the hollow-body black box module declaration `<output file>_bb.v` in your Precision Synthesis project to describe the port connections of the black box. Adding the megafunction wrapper file `<output file>.v` in your Precision Synthesis project is optional, but you must add it to your Quartus II project along with the Precision Synthesis-generated EDIF or VQM netlist.

Alternatively, you can include the megafunction wrapper file `<output file>.v` in your Precision Synthesis project and then right-click the file in the input file list, and select Properties. In the Input file properties dialog box, turn on Exclude file from Compile Phase and click OK. When this option is turned on, the Precision Synthesis software excludes the file from compilation and copies the file to the appropriate directory for use by the Quartus II software during place-and-route.
Instantiating Megafunctions With MegaWizard Plug-In Manager-Generated VHDL Files

The MegaWizard Plug-In Manager generates a VHDL component declaration file <output file>_cmp and a VHDL instantiation template file <output file>_inst.vhd for use in your Precision Synthesis design. Incorporate the component declaration and instantiation template into your top-level design to instantiate the megafunction wrapper file, <output file>.vhd.

Adding the megafunction wrapper file <output file>.vhd in your Precision Synthesis project is optional, but you must add the file to your Quartus II project along with the Precision Synthesis-generated EDIF or VQM netlist.

Alternatively, you can include the megafunction wrapper file <output file>.vhd in your Precision Synthesis project and then right-click the file in the input file list, and select Properties. In the Input file properties dialog box, turn on Exclude file from Compile Phase and click OK. When this option is turned on, the Precision Synthesis software excludes the file from compilation and copies the file to the appropriate directory for use by the Quartus II software during place-and-route.

Instantiating Intellectual Property With the MegaWizard Plug-In Manager and IP Toolbench

Many Altera IP functions include a resource and timing estimation netlist that the Precision Synthesis software can use to synthesize and optimize logic around the IP efficiently. As a result, the Precision Synthesis software provides better timing correlation, area estimates, and Quality of Results (QoR) than a black box approach.

To create this netlist file, perform the following steps:

1. Select the IP function in the MegaWizard Plug-In Manager.
2. Click Next to open the IP Toolbench.
3. Click Set Up Simulation, which sets up all the EDA options.
4. Turn on the Generate netlist option to generate a netlist for resource and timing estimation and click OK.
5. Click Generate to generate the netlist file.

The Quartus II software generates a file <output file>_syn.v. This netlist contains the “grey box” information for resource and timing estimation, but does not contain the actual implementation. Include this netlist file into your Precision Synthesis project as an input file. Then include the megafunction wrapper file <output file>.vhd in the Quartus II project along with your EDIF or VQM output netlist.

The generated “grey box” netlist file, <output file>_syn.v, is always in Verilog HDL format, even if you select VHDL as the output file format.

There is currently no grey box support for SOPC Builder systems in the MegaWizard Plug-In Manager. For information about creating a grey box netlist file from the command line, search Altera’s Knowledge Database. Alternatively, you can use a black box approach as described in "Instantiating Black Box IP Functions With Generated Verilog HDL Files".
**Instantiating Black Box IP Functions With Generated Verilog HDL Files**

You can use the `syn_black_box` or `black_box` compiler directives to declare a module as a black box. The top-level design files must contain the IP port mapping and a hollow-body module declaration. You can apply the directive to the module declaration in the top-level file or a separate file included in the project so that the Precision Synthesis software recognizes the module is a black box.

The `syn_black_box` and `black_box` directives are supported only on module or entity definitions.

Example 19–10 shows a sample top-level file that instantiates `my_verilogIP.v`, which is a simplified customized variation generated by the MegaWizard Plug-In Manager and IP Toolbench.

**Example 19–10.  Top-Level Verilog HDL Code with Black Box Instantiation of IP**

```verilog
top (clk, count);
  input clk;
  output[7:0] count;

  my_verilogIP verilogIP_inst (.clock (clk), .q (count));
endmodule

// Module declaration
// The following attribute is added to create a
// black box for this module.
module my_verilogIP (clock, q) /* synthesis syn_black_box */;
  input clock;
  output[7:0] q;
endmodule
```

**Instantiating Black Box IP Functions With Generated VHDL Files**

You can use the `syn_black_box` or `black_box` compiler directives to declare a component as a black box. The top-level design files must contain the megafuction variation component declaration and port mapping. Apply the directive to the component declaration in the top-level file.

The `syn_black_box` and `black_box` directives are supported only on module or entity definitions.
Example 19–11 shows a sample top-level file that instantiates my_vhdlIP.vhd, which is a simplified customized variation generated by the MegaWizard Plug-In Manager and IP Toolbench.

**Example 19–11. Top-Level VHDL Code with Black Box Instantiation of IP**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
ENTITY top IS
  PORT (clk: IN STD_LOGIC ;
        count: OUT STD_LOGIC_VECTOR (7 DOWNTO 0));
END top;
ARCHITECTURE rtl OF top IS
  COMPONENT my_vhdlIP
    PORT (clock: IN STD_LOGIC ;
        q: OUT STD_LOGIC_VECTOR (7 DOWNTO 0))
  end COMPONENT;
  attribute syn_black_box : boolean;
  attribute syn_black_box of my_vhdlIP: component is true;
BEGIN
  vhdlIP_inst : my_vhdlIP PORT MAP (clock => clk,
                                  q => count);
END rtl;
```

**Inferring Altera Megafuncions from HDL Code**

The Precision Synthesis software automatically recognizes certain types of HDL code and maps arithmetical and relational operators, and memory (RAM and ROM), to efficient technology-specific implementations. This functionality allows technology-specific resources to implement these structures by inferring the appropriate Altera function to provide optimal results. In some cases, the Precision Synthesis software has options that you can use to disable or control inference.

For coding style recommendations and examples for inferring technology-specific architecture in Altera devices, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook, and the Precision Synthesis Style Guide in the Precision Manuals Bookcase. To access these manuals, in the Precision Synthesis software, click Help and select Open Manuals Bookcase.

**Multipliers**

The Precision Synthesis software detects multipliers in HDL code and maps them directly to device atoms to implement the multiplier in the appropriate type of logic. The Precision Synthesis software also allows you to control the device resources that are used to implement individual multipliers, as described in the following section.
Controlling DSP Block Inference for Multipliers

By default, the Precision Synthesis software uses DSP blocks available in Stratix series devices to implement multipliers. The default setting is AUTO, which allows the Precision Synthesis software to map to logic look-up tables (LUTs) or DSP blocks, depending on the size of the multiplier. You can use the Precision Synthesis GUI or HDL attributes to direct mapping to only logic elements or to only DSP blocks.

The options for multiplier mapping in the Precision Synthesis software are described in Table 19–3.

Table 19–3. Options for dedicated_mult Parameter to Control Multiplier Implementation in Precision Synthesis

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ON</td>
<td>Use only DSP blocks to implement multipliers, regardless of the size of the multiplier.</td>
</tr>
<tr>
<td>OFF</td>
<td>Use only logic (LUTs) to implement multipliers, regardless of the size of the multiplier.</td>
</tr>
<tr>
<td>AUTO</td>
<td>Use logic (LUTs) or DSP blocks to implement multipliers, depending on the size of the multipliers.</td>
</tr>
</tbody>
</table>

Setting the Use Dedicated Multiplier Option

To set the Use Dedicated Multiplier option in the Precision Synthesis GUI, compile the design, and then in the Design Hierarchy browser, right-click the operator for the desired multiplier and click Use Dedicated Multiplier.

Setting the dedicated_mult Attribute

To control the implementation of a multiplier in your HDL code, use the dedicated_mult attribute with the appropriate value from Table 19–3, as shown in Example 19–12 and Example 19–13.

Example 19–12. Setting the dedicated_mult Attribute in Verilog HDL

```verilog
//synthesis attribute <signal name> dedicated_mult <value>
```

Example 19–13. Setting the dedicated_mult Attribute in VHDL

```vhdl
ATTRIBUTE dedicated_mult: STRING;
ATTRIBUTE dedicated_mult OF <signal name>: SIGNAL IS <value>;
```

The dedicated_mult attribute can be applied to signals and wires; it does not work when applied to a register. This attribute can be applied only to simple multiplier code, such as `a = b * c`.

Some signals for which the dedicated_mult attribute is set can be removed during synthesis by the Precision Synthesis software for design optimization. In such cases, if you want to force the implementation, you should preserve the signal by setting the preserve_signal attribute to TRUE, as shown in Example 19–14 and Example 19–15.

Example 19–14. Setting the preserve_signal Attribute in Verilog HDL

```verilog
//synthesis attribute <signal name> preserve_signal TRUE
```
Example 19–15. Setting the preserve_signal Attribute in VHDL

ATTRIBUTE preserve_signal: BOOLEAN;
ATTRIBUTE preserve_signal OF <signal name> SIGNAL IS TRUE;

Example 19–16 and Example 19–17 are examples, in Verilog HDL and VHDL, of using the dedicated_mult attribute to implement the given multiplier in regular logic in the Quartus II software.

Example 19–16. Verilog HDL Multiplier Implemented in Logic

module unsigned_mult (result, a, b);
   output [15:0] result;
   input [7:0] a;
   input [7:0] b;
   assign result = a * b;
   //synthesis attribute result dedicated_mult OFF
endmodule

Example 19–17. VHDL Multiplier Implemented in Logic

LIBRARY ieee;
USE ieee.std_logic_1164.ALL;
USE ieee.std_logic_arith.ALL;
USE ieee.std_logic_unsigned.ALL;

ENTITY unsigned_mult IS
   PORT(
      a: IN std_logic_vector (7 DOWNTO 0);
      b: IN std_logic_vector (7 DOWNTO 0);
      result: OUT std_logic_vector (15 DOWNTO 0));
   ATTRIBUTE dedicated_mult: STRING;
END unsigned_mult;

ARCHITECTURE rtl OF unsigned_mult IS
   SIGNAL a_int, b_int: UNSIGNED (7 downto 0);
   SIGNAL pdt_int:  UNSIGNED (15 downto 0);
   ATTRIBUTE dedicated_mult OF pdt_int: SIGNAL IS "OFF";
BEGIN
   a_int <= UNSIGNED (a);
   b_int <= UNSIGNED (b);
   pdt_int <= a_int * b_int;
   result <= std_logic_vector(pdt_int);
END rtl;

Multiplier-Accumulators and Multiplier-Adders

The Precision Synthesis software also allows you to control the device resources used to implement multiply-accumulators or multiply-adders in your project or in a particular module.

The Precision Synthesis software detects multiply-accumulators or multiply-adders in HDL code and infers an ALTMULT_ACCUM or ALTMULT_ADD megafuinction so that the logic can be placed in DSP blocks, or the software maps these functions directly to device atoms to implement the multiplier in the appropriate type of logic.
The Precision Synthesis software supports inference for these functions only if the target device family has dedicated DSP blocks. Refer to “Controlling DSP Block Inference” for more information.

For more information about DSP blocks in Altera devices, refer to the appropriate Altera device family handbook and device-specific documentation. For details about which functions a given DSP block can implement, refer to the DSP Solutions Center on the Altera website at www.altera.com.

For more information about inferring multiply-accumulator and multiply-adder megafunctions in HDL code, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook, and the Precision Synthesis Style Guide in the Precision Synthesis Manuals Bookcase.

Controlling DSP Block Inference

By default, the Precision Synthesis software infers the ALTMULT_ADD or ALTMULT_ACCUM megafunction appropriately in your design. These megafonctions allow the Quartus II software to select either logic or DSP blocks, depending on the device utilization and the size of the function.

You can use the extract_mac attribute to prevent inference of an ALTMULT_ADD or ALTMULT_ACCUM megafunction in a certain module or entity. The options for this attribute are described in Table 19–4.

Table 19–4. Options for extract_mac Attribute Controlling DSP Implementation

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>TRUE</td>
<td>The ALTMULT_ADD or ALTMULT_ACCUM megafunction is inferred.</td>
</tr>
<tr>
<td>FALSE</td>
<td>The ALTMULT_ADD or ALTMULT_ACCUM megafunction is not inferred.</td>
</tr>
</tbody>
</table>

To control inference, use the extract_mac attribute with the appropriate value from Table 19–4 in your HDL code, as shown in Example 19–18 and Example 19–19.

Example 19–18. Setting the extract_mac Attribute in Verilog HDL

```verilog
//synthesis attribute <module name> extract_mac <value>
```

Example 19–19. Setting the extract_mac Attribute in VHDL

```vhdl
ATTRIBUTE extract_mac: BOOLEAN;
ATTRIBUTE extract_mac OF <entity name>: ENTITY IS <value>;
```

To control the implementation of the multiplier portion of a multiply-accumulator or multiply-adder, you must use the dedicated_mult attribute.

Example 19–20 and Example 19–21 on page 19–23 use the extract_mac, dedicated_mult, and preserve_signal attributes (in Verilog HDL and VHDL) to implement the given DSP function in logic in the Quartus II software.
Example 19–20. Using extract_mac, dedicated_mult and preserve_signal in Verilog HDL

```verilog
module unsig_altmult_accum1 (dataout, dataa, datab, clk, aclr, clken);
    input [7:0] dataa, datab;
    input clk, aclr, clken;
    output [31:0] dataout;
    reg [31:0] dataout;
    wire [15:0] multa;
    wire [31:0] adder_out;
    assign multa = dataa * datab;
    //synthesis attribute multa preserve_signal TRUE
    //synthesis attribute multa dedicated_mult OFF
    assign adder_out = multa + dataout;
    always @ (posedge clk or posedge aclr)
    begin
        if (aclr)
            dataout <= 0;
        else if (clken)
            dataout <= adder_out;
    end
    //synthesis attribute unsig_altmult_accum1 extract_mac FALSE
endmodule
```

Example 19–21. Using extract_mac, dedicated_mult, and preserve_signal in VHDL

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_arith.all;
USE ieee.std_logic_signed.all;
ENTITY signedmult_add IS
    PORT(
        a, b, c, d: IN STD_LOGIC_VECTOR (7 DOWNTO 0);
        result: OUT STD_LOGIC_VECTOR (15 DOWNTO 0));
    ATTRIBUTE preserve_signal: BOOLEAN;
    ATTRIBUTE dedicated_mult: STRING;
    ATTRIBUTE extract_mac: BOOLEAN;
    ATTRIBUTE extract_mac OF signedmult_add: ENTITY IS FALSE;
END signedmult_add;
ARCHITECTURE rtl OF signedmult_add IS
    SIGNAL a_int, b_int, c_int, d_int : signed (7 DOWNTO 0);
    SIGNAL pdt_int, pdt2_int : signed (15 DOWNTO 0);
    SIGNAL result_int: signed (15 DOWNTO 0);
    ATTRIBUTE preserve_signal OF pdt_int: SIGNAL IS TRUE;
    ATTRIBUTE dedicated_mult OF pdt_int: SIGNAL IS "OFF";
    ATTRIBUTE preserve_signal OF pdt2_int: SIGNAL IS TRUE;
    ATTRIBUTE dedicated_mult OF pdt2_int: SIGNAL IS "OFF";
    BEGIN
        a_int <= signed (a);
        b_int <= signed (b);
        c_int <= signed (c);
        d_int <= signed (d);
        pdt_int <= a_int * b_int;
        pdt2_int <= c_int * d_int;
        result_int <= pdt_int + pdt2_int;
        result <= STD_LOGIC_VECTOR(result_int);
    END rtl;
```
RAM and ROM

The Precision Synthesis software detects memory structures in HDL code and converts them to an operator that infers an ALTSYNCRAM or LPM_RAM_DP megafunction, depending on the device family. The software then places these functions in memory blocks.

The software supports inference for these functions only if the target device family has dedicated memory blocks.

For more information about inferring RAM and ROM megafunctions in HDL code, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook, and the Precision Synthesis Style Guide in the Precision Synthesis Manuals Bookcase. To access these manuals, in the Precision Synthesis software, click Help and select Open Manuals Bookcase.

Incremental Compilation and Block-Based Design

As designs become more complex and designers work in teams, a block-based incremental design flow is often an effective design approach. In an incremental compilation flow, you can make changes to one part of the design while maintaining the placement and performance of unchanged parts of the design. Design iterations can be made dramatically faster by focusing new compilations on particular design partitions and merging results with the results of previous compilations of other partitions. You can perform optimization on individual blocks and then integrate them into a final design and optimize the design at the top-level.

The first step in an incremental design flow is to make sure that different parts of your design do not affect each other. You must ensure that you have separate netlists for each partition in your design. If the whole design is in one netlist file, changes in one partition affect other partitions because of possible node name changes when you resynthesize the design.

You can create different implementations for each partition in your Precision Synthesis project, which allows you to switch between partitions without leaving the current project file. You can also create a separate project for each partition if you require separate projects for a team-based design flow. Alternatively, you can use the incremental synthesis capability in the Precision RTL Plus software.

For more information about creating partitions and using incremental compilation in the Quartus II software, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Creating a Design with Precision RTL Plus Incremental Synthesis

The Precision RTL Plus incremental synthesis flow for Quartus II incremental compilation uses a partition-based approach to achieve faster design cycle time.

Using the incremental synthesis feature, you can create different netlist files for different partitions of a design hierarchy within one partition implementation, which makes each partition independent of the others in an incremental compilation flow. Only the portions of a design that have been updated must be recompiled during design iterations. You can make changes and resynthesize one partition in a design to create a new netlist without affecting the synthesis results or fitting of other partitions.
The following steps show a general flow for partition-based incremental synthesis with Quartus II incremental compilation.

1. Create Verilog HDL or VHDL design files.

2. Determine which hierarchical blocks you want to treat as separate partitions in your design, and designate the partitions with the `incr_partition` attribute. For the syntax to create partitions, refer to “Creating Partitions with the `incr_partition` Attribute” on page 19–25.

3. Create a project in the Precision RTL Plus Synthesis software and add the HDL design files to the project.

4. Enable incremental synthesis in the Precision RTL Plus Synthesis software using one of these methods:
   - Run the following command in the Transcript Window:
     ```
     setup_design -enable_incr_synth
     ```

5. Run the basic Precision Synthesis flow of compilation, synthesis, and place-and-route on your design. In subsequent runs, the Precision RTL Plus Synthesis software processes only the parts of the design that have changed, resulting in a shorter iteration than the initial run. The performance of the unchanged partitions is preserved.

   The Precision RTL Plus Synthesis software sets the netlist types of the unchanged partitions to Post-Fit and the changed partitions to Post-Synthesis. You can change the netlist type during timing closure in the Quartus II software to obtain the best QoR.

6. Import the EDIF or VQM netlist for each partition and the top-level .tcl file into the Quartus II software, and set up the Quartus II project to use incremental compilation.

7. Compile your Quartus II project.

8. If you want, you can change the Quartus II incremental compilation netlist type for a partition with the Design Partitions Window. You can change the Netlist Type to one of the following options:
   - To preserve the previous post-fit placement results, change the Netlist Type of the partition to Post-Fit.
   - To preserve the previous routing results, set the Fitter Preservation Level of the partition to Placement and Routing.

   **Creating Partitions with the `incr_partition` Attribute**

   Partitions are set using the HDL `incr_partition` attribute. The Precision Synthesis software creates or deletes partitions by reading this attribute during compilation iterations. The attribute can be attached to either the design unit definition or an instance. Example 19–22 and Example 19–23 show how to use the attribute to create partitions.

   To delete partitions, you can remove the attribute or set the attribute value to false.
The Precision Synthesis software ignores partitions set in a black box.

**Example 19–22. Using incr_partition Attribute to Create a Partition in Verilog HDL**

Design unit partition:

```verilog
module my_block(
    input clk;
    output reg [31:0] data_out) /* synthesis incr_partition */ ;
```

Instance partition:

```verilog
my_block my_block_inst(.clk(clk), .data_out(data_out));
// synthesis attribute my_block_inst incr_partition true
```

**Example 19–23. Using incr_partition Attribute to Create a Partition in VHDL**

Design unit partition:

```vhdl
entity my_block is
    port(
        clk : in std_logic;
        data_out : out std_logic_vector(31 downto 0)
    );
    attribute incr_partition : boolean;
    attribute incr_partition of my_block : entity is true;
end entity my_block;
```

Instance partition:

```vhdl
component my_block is
    port(
        clk : in std_logic;
        data_out : out std_logic_vector(31 downto 0)
    );
end component;
```

```vhdl
attribute incr_partition : boolean;
attribute incr_partition of my_block_inst : label is true;
my_block_inst my_block
    port map(clk, data_out);
```

Creating Multiple Mapped Netlist Files With Separate Precision Projects or Implementations

This section describes how to manually generate multiple netlist files, which can be VQM or EDIF files, for incremental compilation using black boxes and separate Precision projects or implementations for each design partition. This manual flow is supported in versions of the Precision software that do not include the incremental synthesis feature. You might also use this feature if you perform synthesis in a team-based environment without a top-level synthesis project that includes all of the lower-level design blocks.
In the Precision Synthesis software, create a separate implementation, or a separate project, for each lower-level module and for the top-level design that you want to maintain as a separate netlist file. Implement black box instantiations of lower-level modules in your top-level implementation or project.

For more information about managing implementations and projects, refer to the Precision RTL Synthesis User’s Manual. To access this manual, in the Precision Synthesis software, click Help and select Open Manuals Bookcase.

When synthesizing the implementations for lower-level modules, perform these steps in the Precision Synthesis software:

1. On the Tools menu, turn off Add IO Pads on the Optimization page under Set Options.

   You must turn off the Add IO Pads option while synthesizing the lower-level modules individually. Enable the Add IO Pads option only while synthesizing the top-level module.

2. Read the HDL files for the modules.

   Modules can include black box instantiations of lower-level modules that are also maintained as separate netlist files.

3. Add constraints for all partitions in the design.

When synthesizing the top-level design implementation, perform these steps:

1. Read the HDL files for top-level designs.


3. Create black boxes for lower-level modules in the top-level design.

4. Add constraints.

In a standard Quartus II incremental compilation flow, Precision Synthesis software constraints made on lower-level modules are not passed to the Quartus II software. Ensure that appropriate constraints are made in the top-level Precision Synthesis project, or in the Quartus II project.
Creating Black Boxes to Create EDIF Netlists

This section describes how to create black boxes to create separate EDIF netlists. Figure 19–2 shows an example of a design hierarchy separated into various partitions.

Figure 19–2. Partitions in a Hierarchical Design

In Figure 19–2, the top-level partition contains the top-level block in the design (block A) and the logic that is not defined as part of another partition. In this example, the partition for top-level block A also includes the logic in the sub-block C. Because block F is contained in its own partition, it is not treated as part of the top-level partition A. Another separate partition, B, contains the logic in blocks B, D, and E. In a team-based design, different engineers may work on the logic in different partitions. One netlist is created for the top-level module A and its submodule C, another netlist is created for module B and its submodules D and E, while a third netlist is created for module F. To create multiple EDIF netlist files for this design, follow these steps:

2. Generate an .edf file for module F. Use F.v/.vhd as the source file.
3. Generate a top-level .edf file for module A. Use A.v/.vhd and C.v/.vhd as the source files. Ensure that you create black boxes for modules B and F, which were optimized separately in the previous steps.

The goal is to individually synthesize and generate an .edf netlist file for each lower-level module and then instantiate these modules as black boxes in the top-level file. You can then synthesize the top-level file to generate the .edf netlist file for the top-level design. Finally, both the lower-level and top-level .edf netlist files are provided to your Quartus II project.

When you make design or synthesis optimization changes to part of your design, resynthesize only the changed partition to generate the new .edf netlist file. Do not resynthesize the implementations or projects for the unchanged partitions.

Creating Black Boxes in Verilog HDL

Any design block that is not defined in the project or included in the list of files to be read for a project is treated as a black box by the software. In Verilog HDL, you must provide an empty module declaration for any module that is treated as a black box.
A black box for the top-level file `A.v` is shown in the following example. Provide an empty module declaration for any lower-level files, which also contain a black box for any module beneath the current level of hierarchy.

**Example 19–24. Verilog HDL Black Box for Top-Level File A.v**

```verilog
module A (data_in, clk, e, ld, data_out);
    input data_in, clk, e, ld;
    output [15:0] data_out;
    wire [15:0] cnt_out;
    B U1 (.data_in (data_in), .clk(clk), .ld (ld), .data_out(cnt_out));
    F U2 (.d(cnt_out), .clk(clk), .e(e), .q(data_out));
    // Any other code in A.v goes here.
endmodule

// Empty Module Declarations of Sub-Blocks B and F follow here.
// These module declarations (including ports) are required for black
// boxes.
module B (data_in, clk, ld, data_out);
    input data_in, clk, ld;
    output [15:0] data_out;
endmodule

module F (d, clk, e, q);
    input [15:0] d;
    input clk, e;
    output [15:0] q;
endmodule
```

**Creating Black Boxes in VHDL**

Any design block that is not defined in the project or included in the list of files to be read for a project is treated as a black box by the software. In VHDL, you must provide a component declaration for the black box.

A black box for the top-level file `A.vhd` is shown in Example 19–25. Provide a component declaration for any lower-level files that also contain a black box or for any block beneath the current level of hierarchy.
After you complete the steps outlined in this section, you have different EDIF netlist files for each partition of the design. These files are ready for use with incremental compilation in the Quartus II software.

### Creating Quartus II Projects for Multiple EDIF Files

The Precision Synthesis software creates a .tcl file for each implementation, and provides the Quartus II software with the appropriate constraints and information to set up a project. When using incremental synthesis, the Precision RTL Plus Synthesis software creates only a single .tcl file, `<project name>_incr_partitions.tcl`, to pass the partition information to the Quartus II software. For details about using this Tcl script to set up your Quartus II project and to pass your top-level constraints, refer to “Running the Quartus II Software Manually Using the Precision Synthesis-Generated Tcl Script” on page 19–11.
Depending on your design methodology, you can create one Quartus II project for all EDIF netlists, or a separate Quartus II project for each EDIF netlist. In the standard incremental compilation design flow, you create design partition assignments for each partition in the design within a single Quartus II project. This methodology provides the best QoR and performance preservation during incremental changes to your design. You might require a bottom-up design flow if each partition must be optimized separately, such as for third-party IP delivery.

To follow this design flow in the Quartus II software, create separate Quartus II projects and export each design partition and incorporate it into a top-level design using the incremental compilation features to maintain placement results.

The following sections describe how to create the Quartus II projects for these two design flows.

Creating a Single Quartus II Project for a Standard Incremental Compilation Flow

Use the `<top-level project>.tcl` file generated for the top-level partition to create your Quartus II project and import all the netlists into this one Quartus II project for an incremental compilation flow. You can optimize all partitions within the single Quartus II project and take advantage of the performance preservation and compilation time reduction that incremental compilation provides. Figure 19–3 shows the design flow for the example design in Figure 19–2 on page 19–28.

All the constraints from the top-level implementation are passed to the Quartus II software in the top-level .tcl file, but any constraints made only in the lower-level implementations within the Precision Synthesis software are not forward-annotated. Enter these constraints manually in your Quartus II project.

**Figure 19–3. Design Flow Using Multiple EDIF Files with One Quartus II Project**
Creating Multiple Quartus II Projects for a Bottom-Up Flow

Use the .tcl files generated by the Precision Synthesis software for each Precision Synthesis software implementation or project to generate multiple Quartus II projects, one for each partition in the design. Each designer in the project can optimize their block separately in the Quartus II software and export the placement of their blocks using incremental compilation. Designers should create a LogicLock region to provide a floorplan location assignment for each block; the top-level designer should then import all the blocks and assignments into the top-level project. Figure 19–4 shows the design flow for the example design in Figure 19–2 on page 19–28.

Hierarchy and Design Considerations

To ensure the proper functioning of the synthesis flow, you can create separate partitions only for modules, entities, or existing netlist files. In addition, each module or entity must have its own design file. If two different modules are in the same design file, but are defined as being part of different partitions, incremental synthesis cannot be maintained because both regions must be recompiled when you change one of the modules.

Altera recommends that you register all inputs and outputs of each partition. This makes logic synchronous and avoids any delay penalty on signals that cross partition boundaries.

If you use boundary tri-states in a lower-level block, the Precision Synthesis software pushes the tri-states through the hierarchy to the top-level to make use of the tri-state drivers on output pins of Altera devices. Because pushing tri-states requires optimizing through hierarchies, lower-level tri-states are not supported with a block-based compilation methodology. You should use tri-state drivers only at the external output pins of the device and in the top-level block in the hierarchy.

For more tips on design partitioning, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.
Conclusion

The Mentor Graphics Precision Synthesis software and Quartus II design flow allow you to control how to prepare your design files for the Quartus II place-and-route process, which allows you to improve performance and optimizes your design for use with Altera devices. Several of the methodologies outlined in this chapter can help you optimize your design to achieve performance goals and decrease design time.

Document Revision History

Table 19–5 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>■ Template update.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial changes.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Changed to new document template.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed Classic Timing Analyzer support.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added support for .vqm netlist files.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Edited the “Creating Quartus II Projects for Multiple EDIF Files” on page 15–30 section for changes with the incremental compilation flow.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Editorial changes.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Minor updates for the Quartus II software version 10.0 release</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Minor updates for the Quartus II software version 9.1 release</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Updated list of supported devices for the Quartus II software version 9.0 release</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Chapter 11 was previously Chapter 10 in software version 8.1</td>
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<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>■ Changed to 8-1/2 x 11 page size</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Title changed to Mentor Graphics Precision Synthesis Support</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated list of supported devices</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about the Precision RTL Plus incremental synthesis flow</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Figure 10-1 to include SystemVerilog</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Guidelines for Altera Megafunctios and Architecture-Specific Features” on page 10–19</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Incremental Compilation and Block-Based Design” on page 10–28</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added section “Creating Partitions with the incr_partition Attribute” on page 10–29</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Removed Mercury from the list of supported devices</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed Precision version to 2007a update 3</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added note for Stratix IV support</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Renamed “Creating a Project and Compiling the Design” section to “Creating and Compiling a Project in the Precision RTL Synthesis Software”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about constraints in the Tcl file</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated document based on the Quartus II software version 8.0</td>
</tr>
</tbody>
</table>
For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter documents key design methodologies and techniques for Altera® devices using the LeonardoSpectrum and Quartus II design flow.

This chapter includes the following sections:

- “Altera Device Family Support”
- “Design Flow” on page 20–2
- “LeonardoSpectrum Optimization Strategies” on page 20–4
- “Timing Analysis with the LeonardoSpectrum Software” on page 20–6
- “Exporting Designs Using NativeLink Integration” on page 20–7
- “Guidelines for Altera Megafunctions and LPM Functions” on page 20–8
- “Block-Based Design with the Quartus II Software” on page 20–16

Altera recommends using the advanced Mentor Graphics Precision RTL Synthesis software for new designs in new device families. For more information about Precision RTL Synthesis, refer to the Mentor Graphics Precision Synthesis Support chapter in volume 1 of the Quartus II Handbook.

This chapter assumes that you have set up, licensed, and are familiar with the LeonardoSpectrum software.


### Altera Device Family Support

The LeonardoSpectrum software is a mature synthesis tool supporting legacy devices and many current devices; however, newly-released devices may not be supported.

Contact Mentor Graphics for more information about support for newly-released devices.

The LeonardoSpectrum software also supports the FLEX 8000 and MAX 9000 legacy devices that are supported only in the Altera MAX+PLUS® II software, as well as ACEX® 1K, APEX™ II, APEX 20K, APEX 20KC, APEX 20KE, FLEX® 10K, and FLEX 6000 legacy devices that are supported by the Quartus II software version 9.0 and earlier.
Design Flow

The following steps describe a basic Quartus II software design flow using the LeonardoSpectrum software:

1. Create Verilog HDL or VHDL design files.
2. Import the Verilog HDL or VHDL design files into the LeonardoSpectrum software for synthesis.
3. Select a target device and add timing constraints and compiler directives to help optimize the design during synthesis.
4. Synthesize the project in the LeonardoSpectrum software.
5. Create a Quartus II project and import the technology-specific EDIF Input File (.edf) netlist and the Tcl Script File (.tcl) generated by the LeonardoSpectrum software into the Quartus II software for placement and routing and performance evaluation.
6. After obtaining place-and-route results that meet your requirements, configure or program the Altera device.

Figure 20–1 on page 20–3 shows the recommended design flow using the LeonardoSpectrum and Quartus II software.

If your area and timing requirements are met, use the programming files generated by the Quartus II software to program or configure the Altera device. If the area or timing requirements are not met, change the constraints in the LeonardoSpectrum software and rerun synthesis. Repeat the process until the area and timing requirements are met. You can also use other Quartus II software options and techniques to meet the area and timing requirements.
The LeonardoSpectrum software supports both VHDL and Verilog HDL source files. With the appropriate license, the software also supports mixed synthesis, allowing a combination of VHDL and Verilog HDL source files. During synthesis, the LeonardoSpectrum software produces several intermediate and output files, which are listed and described in Table 20–1.

Table 20–1. LeonardoSpectrum Intermediate and Output Files

<table>
<thead>
<tr>
<th>File Extension(s)</th>
<th>File Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>.xdb</td>
<td>Technology-independent register transfer level (RTL) netlist file that can only be read by the LeonardoSpectrum software.</td>
</tr>
<tr>
<td>.edf</td>
<td>Technology-specific output netlist file in electronic design interchange format (EDIF).</td>
</tr>
<tr>
<td>.tcl</td>
<td>Forward-annotated constraints file containing constraints and assignments. A .tcl file for the Quartus II software is created for all devices. The .tcl file contains the appropriate Tcl commands to create and set up a Quartus II project and pass placement constraints.</td>
</tr>
<tr>
<td>.acf</td>
<td>Assignment and Configurations file for backward compatibility with the MAX+PLUS II software. For devices supported by the MAX+PLUS II software, the MAX+PLUS II assignments are imported from the MAX+PLUS II .acf file.</td>
</tr>
</tbody>
</table>
Altera recommends that you do not use project directory names that include spaces. Some file operations in the LeonardoSpectrum software do not work correctly if the path name contains spaces.

Specify timing constraints and compiler directives for the design in the LeonardoSpectrum software, or in a constraint file (.ctr). These constraints are forward-annotated in the .tcl file for use with the Classic Timing Analyzer in the Quartus II software version 10.0 and earlier.

The LeonardoSpectrum software does not generate a Synopsys Design Constraint (SDC) format file for the TimeQuest Timing Analyzer. If you are using a current version of the Quartus II software, you must convert your timing constraints to SDC format. Altera recommends using the advanced Precision RTL Synthesis software for new designs in new device families, instead of using the LeonardoSpectrum software.

The LeonardoInsight™ Schematic Viewer is an add-on graphical tool for schematic views of the technology-independent RTL netlist (.xdb) and the technology-specific gate-level results. You can use the Schematic Viewer to visually analyze and debug the design. It also supports cross-probing between the RTL and gate-level schematics, the design browser, and the source code in the HDLInventor™ text editor.

LeonardoSpectrum Optimization Strategies

You can configure most general settings in the Quick Setup tab in the LeonardoSpectrum user interface. FlowTabs provide additional options, some including multiple PowerTabs at the bottom of the screen, with more options. Advanced optimization options in the LeonardoSpectrum software include timing-driven synthesis, encoding style, resource sharing, and mapping I/O registers.

Timing-Driven Synthesis

The LeonardoSpectrum software supports timing-driven synthesis through user-assigned timing constraints that optimize the performance of the design. Constraints, such as clock frequency, can be specified globally or for individual clock signals. The following sections describe how to set the various types of timing constraints in the LeonardoSpectrum software.

The timing constraints described in “Global PowerTab” are set in the Constraints FlowTab. In this FlowTab, there are PowerTabs at the bottom, such as Global and Clock, for setting various constraints.

Global PowerTab

The Global PowerTab is the default PowerTab in the Constraints FlowTab where you can specify the global clock frequency. The Clock Frequency setting on the Quick Setup tab is equivalent to the Registers to Registers delay setting. You can also specify the Input Ports to Registers, Registers to Output Ports, and Inputs to Outputs delays that correspond to global tSU, tCO, and tPD requirements, respectively, in the Quartus II software. The timing diagram on the Global PowerTab reflects the settings you have made.
Clock PowerTab

You can set various constraints for each clock in your design. First, select the clock name in the Clock(s) window. The clock names appear after the design is read from the Input FlowTab. Configure settings for that particular clock and click Apply. You can also specify the Duty Cycle, which is set to 5% by default. The timing diagram shows these settings.

If a clock has an Offset from the main clock, which is considered to be time “0”, this constraint corresponds to the OFFSET_FROM_BASE_CLOCK setting in the Quartus II software.

You can specify the pin number for the clock input pin in the Pin Location box. This pin number is passed to the Quartus II software for place-and-route, but does not affect synthesis in the LeonardoSpectrum software.

Input and Output PowerTabs

Configure settings for individual input or output pins in the Input and Output tabs. First, select a name in the Input Ports or Output Ports window. The names appear after the design is read from the Input FlowTab. Then, make the following settings for that pin as described below.

The Arrival Time indicates that the input signal arrives a specified time after the rising clock edge (time “0”). This setting constrains the path from the pin to the first register by including the arrival time in the total delay, and corresponds to the EXTERNAL_INPUT_DELAY assignment in the Quartus II software.

The Required Time indicates the maximum delay after time “0” that the output signal should arrive at the output pin. This setting directly constrains the register to output delay, and corresponds to the EXTERNAL_OUTPUT_DELAY assignment in the Quartus II software.

Specify the pin number for the I/O pin in the Pin Location box. This pin number is passed to the Quartus II software for place-and-route, but does not affect synthesis in the LeonardoSpectrum software.

Other Constraints

The following sections describe other constraints that can be set with the LeonardoSpectrum user interface:

- “Encoding Style”
- “Resource Sharing” on page 20–6
- “Mapping I/O Registers” on page 20–6

Encoding Style

The LeonardoSpectrum software encodes state machines during the synthesis process. To improve performance when coding state machines, separate state machine logic from all arithmetic functions and data paths. When encoded, a design cannot be re-encoded later in the optimization process. You must follow a particular VHDL or Verilog HDL coding style for the LeonardoSpectrum software to identify the state machine.
Table 20–2 describes the state machine encoding styles supported by the LeonardoSpectrum software.

Table 20–2. State Machine Encoding Styles in the LeonardoSpectrum Software

<table>
<thead>
<tr>
<th>Style</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Binary</td>
<td>Generates state machines with the fewest possible flipflops. Binary state machines are useful for area-critical designs when timing is not the primary concern.</td>
</tr>
<tr>
<td>Gray</td>
<td>Generates state machines where only one flipflop changes during each transition. Gray-encoded state machines tend to be free of glitches.</td>
</tr>
<tr>
<td>One-hot</td>
<td>Generates state machines containing one flipflop for each state. One-hot state machines provide the best performance and shortest clock-to-output delays. However, one-hot implementations are usually larger than binary implementations.</td>
</tr>
<tr>
<td>Random</td>
<td>Generates state machines using random state machine encoding. Use random state machine encoding only when no other implementation achieves the desired results.</td>
</tr>
<tr>
<td>Auto (Default)</td>
<td>Implements binary or one-hot encoding, depending on the size of enumerated types in the state machine.</td>
</tr>
</tbody>
</table>

The **Encoding Style** is created in the **Input FlowTab**. The setting instructs the software to use a particular state machine encoding style for all state machines. The default **Auto** encoding style implements binary or one-hot encoding, depending on the size of enumerated types in the state machine.

To ensure proper recognition and improve performance when coding state machines, refer to the **Recommended HDL Coding Styles** chapter in volume 1 of the **Quartus II Handbook** for design guidelines.

**Resource Sharing**

You should turn on **Resource Sharing** in the **Input FlowTab** to allow optimization, which reduces device resources.

**Mapping I/O Registers**

The **Map I/O Registers** option in the **Technology FlowTab**, is available for Altera FPGAs containing I/O cells (IOC) or I/O elements (IOE). If this option is turned on, input or output registers are moved into the device’s I/O cells for faster setup or clock-to-output times.

**Timing Analysis with the LeonardoSpectrum Software**

The LeonardoSpectrum software reports successful synthesis with an information message in the Transcript or Information window. Estimated device usage and timing results are reported in the Device Utilization section of this window. **Figure 20–2 on page 20–7** shows an example of a LeonardoSpectrum compilation report.
Chapter 20: Mentor Graphics LeonardoSpectrum Support

Exporting Designs Using NativeLink Integration

June 2012 Altera Corporation Quartus II Handbook Version 13.0

Volume 1: Design and Synthesis

The LeonardoSpectrum software estimates the timing results based on timing models. The software does not have the design’s place-and-route information in the Quartus II software, so it cannot report accurate routing delays. Additionally, if the design includes any black boxed Altera-specific functions, the LeonardoSpectrum software does not report timing information for these functions.

Exporting Designs Using NativeLink Integration

You can use NativeLink® integration to integrate the LeonardoSpectrum and the Quartus II software with a single GUI for both the synthesis and place-and-route operations. You can run the Quartus II software from within the LeonardoSpectrum software GUI with NativeLink integration or you can run the LeonardoSpectrum software from within the Quartus II software GUI for device families supported in the Quartus II software.

Generating Netlist Files

The LeonardoSpectrum software generates an .edf netlist file readable as an input file in the Quartus II software for place-and-route. Select the .edf file option name in the Output FlowTab. The .edf file is also generated if the Auto option is turned on in the Output FlowTab.

Including Design Files for Black Boxed Modules

Be sure to include the MegaWizard™ Plug-In Manager-generated custom megafuction variation design file in the Quartus II project directory, or add it to the list of project files for place-and-route, if the design has black boxed megafuctions.

Figure 20–2. LeonardoSpectrum Compilation Report

<table>
<thead>
<tr>
<th>Resource</th>
<th>Used</th>
<th>Avail</th>
<th>Utilization</th>
</tr>
</thead>
<tbody>
<tr>
<td>IOs</td>
<td>22</td>
<td>136</td>
<td>16.18%</td>
</tr>
<tr>
<td>LFs</td>
<td>114</td>
<td>8320</td>
<td>1.37%</td>
</tr>
<tr>
<td>Memory Bits</td>
<td>0</td>
<td>106496</td>
<td>0.00%</td>
</tr>
</tbody>
</table>

Clock Frequency Report

- clk : 52.2 MHz
- clk2: 149.5 MHz

Critical Path Report
Passing Constraints with Scripts

The LeonardoSpectrum software can create a Tcl script (.tcl) file called <project name>.tcl. This file contains commands to create a Quartus II project along with constraints and other assignments. To output a .tcl file, turn on Write Vendor Constraint Files in the Output FlowTab.

To create and compile a Quartus II project using the .tcl file generated from the LeonardoSpectrum software, perform the following steps in the Quartus II software:

1. Place the .edf and .tcl files in the same directory.
2. On the View menu, point to Utility Windows, and click Tcl Console to open the Quartus II Tcl Console.
3. At the Tcl prompt, type source <path>/<project name>.tcl.
4. On the File menu, click Open Project to open the new project.
5. On the Processing menu, click Start Compilation.

The LeonardoSpectrum software does not generate a Synopsys Design Constraint (SDC) format file for the TimeQuest Timing Analyzer. If you are using a current version of the Quartus II software, you must convert your timing constraints to SDC format. Altera recommends using the advanced Precision RTL Synthesis software for new designs in new device families, instead of using the LeonardoSpectrum software.

Integration with the Quartus II Software

You can launch the Quartus II software from within the LeonardoSpectrum software with the options in the Place And Route section of the Quick Setup tab. Turn on the Run Integrated Place and Route option to start the compilation in the Quartus II software to show the fitting and performance results. You can also run place-and-route in the Quartus II software by turning on Run Quartus on the Physical FlowTab, and then clicking Run PR.

To use integrated place-and-route, on the Options menu, point to Place and Route Path and click Tools. Specify the location of the Quartus II software executable file (<Quartus II software installation directory>/bin).

Guidelines for Altera Megafunctions and LPM Functions

Altera provides parameterizable megafunctions ranging from simple arithmetic units, such as adders and counters, to advanced phase-locked loop (PLL) blocks, multipliers, and memory structures. These functions are performance-optimized for Altera devices. Megafunctions include the library of parameterized modules (LPM), device-specific megafunctions such as PLLs, LVDS, and digital signal processing (DSP) blocks, intellectual property (IP) available as Altera MegaCore® functions, and IP available through the Altera Megafuntion Partners Program (AMPP™).

Some IP cores require synthesis in the LeonardoSpectrum software. Refer to the user guide for the specific IP.
You can infer or instantiate megafunctions in the LeonardoSpectrum software. The LeonardoSpectrum software supports inferring some Altera megafunctions, such as multipliers, DSP functions, and RAM and ROM blocks. The LeonardoSpectrum software supports all Altera megafunctions through instantiation.

**Instantiating Altera Megafonctions**

There are two methods of instantiating Altera megafonctions in the LeonardoSpectrum software. The more common method, which maintains target technology awareness, is to set up and parameterize a megafonction variation in the MegaWizard Plug-In Manager in the Quartus II software. The megafonction wizard creates a wrapper file that instantiates the megafonction. The less common method is to directly instantiate the megafonction in the Verilog HDL or VHDL code. The advantage of using the megafonction wizard method over the instantiation method is that the megafonction wizard properly sets all the parameters. Also you do not need the library support, which is required in the instantiation method. This is referred to as black box methodology.

Altera recommends using the MegaWizard Plug-In Manager to ensure that the ports and parameters are set correctly.

When directly instantiating megafonctions, refer to the respective megafonction user guide on the Altera IP and Megafonction page.

**Inferring Altera Memory Elements**

The LeonardoSpectrum software can infer memory blocks from Verilog HDL or VHDL code. When the LeonardoSpectrum software detects a RAM or ROM from the style of the RTL code at a technology-independent level, the software maps the element to a generic module in the RTL database. During the technology-mapping phase of synthesis, the LeonardoSpectrum software maps the generic module to the most optimal primitive memory cells, or Altera megafonction, for the target Altera technology.

For more information about inferring RAM and ROM megafonctions, including examples of VHDL and Verilog HDL code, see the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

**Inferring RAM**

The LeonardoSpectrum software supports RAM inference for various device families. The following are the restrictions for the LeonardoSpectrum software to successfully infer RAM in a design:

- The write process must be synchronous.
- The read process can be asynchronous or synchronous, depending on the target Altera architecture.
- Resets on the memory are not supported.

The Stratix series and Cyclone series devices support the RAM primitive altsyncram with a minimum RAM size of two bits, and a minimum RAM address width of one bit.

To disable RAM inference, set the extract_ram and infer_ram variables to false.
To enter the value `false` when performing synthesis in the user interface with the **Advanced FlowTabs**, on the Tools menu, click **Variable Editor**, or add the `set extract_ram false` and `set infer_ram false` commands to your synthesis script.

**Inferring ROM**

You can implement ROM behavior in HDL source code with CASE statements or specify the ROM as a table. The LeonardoSpectrum software infers both synchronous and asynchronous ROM, depending on the target Altera device. For example, memory for Stratix series devices must be synchronous to be inferred.

To disable ROM inference, set the `extract_rom` variable to `false`. To enter the value `false` when performing synthesis in the user interface with the **Advanced FlowTabs**, on the Tools menu, click **Variable Editor**, or add the `set extract_rom false` commands to your synthesis script.

**Inferring Multipliers and DSP Functions**

Some Altera devices include dedicated DSP blocks optimized for DSP applications. The following Altera megafunctions are used with DSP block modes:

- **LPM_MULT**
- **ALTMULT_ACCUM**
- **ALTMULT_ADD**

You can instantiate these megafunctions in the design or direct the LeonardoSpectrum software to infer the appropriate megafunction by recognizing a multiplier, multiplier-accumulator (MAC), or multiplier-adder in the design. The Quartus II software maps the functions to the DSP blocks in the device during place-and-route.

For more information about inferring multipliers and DSP functions, including examples of VHDL and Verilog HDL code, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

**Simple Multipliers**

The LPM_MULT megafunction implements the DSP block in the simple multiplier mode. The following functionality is supported in this mode:

- The DSP block includes registers for the input and output stages, and an intermediate pipeline stage.
- Signed and unsigned arithmetic is supported.

**Multiplier Accumulators**

The ALTMULT_ACCUM megafunction implements the DSP block in the multiply-accumulator mode. The following functionality is supported in this mode:

- The DSP block includes registers for the input and output stages, and an intermediate pipeline stage.
- The output registers are required for the accumulator.
- The input and pipeline registers are optional.
- Signed and unsigned arithmetic is supported.
If the design requires input registers to be used as shift registers, use the black box method to instantiate the ALTMULT_ACCUM megafunction.

**Multiplier Adders**

The LeonardoSpectrum software can infer multiplier adders and map them to either the two-multiplier adder mode or the four-multiplier adder mode of the DSP blocks. The LeonardoSpectrum software maps the HDL code to the correct ALTMULT_ADD function.

The following functionality is supported in these modes:

- The DSP block includes registers for the input and output stages, and an intermediate pipeline stage.
- Signed and unsigned arithmetic is supported, but support for the Verilog HDL signed construct is limited.

### Controlling DSP Block Inference

Device features, such as dedicated DSP blocks, multipliers, multiply-accumulators, and multiply-adders can be implemented in DSP blocks or in logic. You can control this implementation through attribute settings in the LeonardoSpectrum software.

As described in Table 20–3, attribute settings in the LeonardoSpectrum software control the implementation of the multipliers in DSP blocks or logic at the signal block (or module), and project level.

**Table 20–3. Attribute Settings for DSP Blocks in the LeonardoSpectrum Software**

<table>
<thead>
<tr>
<th>Level</th>
<th>Attribute Name</th>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Global</td>
<td>extract_mac (2)</td>
<td>TRUE</td>
<td>All multipliers in the project are mapped to DSP blocks.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>FALSE</td>
<td>All multipliers in the project are mapped to logic.</td>
</tr>
<tr>
<td>Module</td>
<td>extract_mac (3)</td>
<td>TRUE</td>
<td>Multipliers inside the specified module are mapped to DSP blocks.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>FALSE</td>
<td>Multipliers inside the specified module are mapped to logic.</td>
</tr>
<tr>
<td>Signal</td>
<td>dedicated_mult</td>
<td>ON</td>
<td>LPM is inferred and multipliers are implemented in DSP block.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>OFF</td>
<td>LPM is inferred, but multipliers are implemented in logic by the Quartus II software.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>LCELL</td>
<td>LPM is not inferred, and multipliers are implemented in logic by the LeonardoSpectrum software.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>AUTO</td>
<td>LPM is inferred, but the Quartus II software automatically maps the multipliers to either logic or DSP blocks based on the Quartus II software place-and-route.</td>
</tr>
</tbody>
</table>

**Notes to Table 20–3:**

1. The `extract_mac` attribute takes precedence over the `dedicated_mult` attribute.
2. For devices with DSP blocks, the `extract_mac` attribute is set to “true” by default for the entire project.
3. For devices with DSP blocks, the `extract_mac` attribute is set to “true” by default for all modules.

**Global Attribute**

You can set the `extract_mac` global attribute to control the implementation of multipliers in DSP blocks for the entire project. You can set this attribute with the following script command:

```
set extract_mac <value>
```
Module Level Attributes

You can control the implementation of multipliers inside a module or component by setting the `extract_mac` attribute in the Verilog HDL source code. Setting this attribute for a module affects only the multipliers inside that module. Use the following command:

```verbatim
//synthesis attribute <module name> extract_mac <value>
```

The Verilog HDL and VHDL code samples in Example 20–1 and Example 20–2 show how to use the `extract_mac` attribute.

**Example 20–1. Using Module Level Attributes in Verilog HDL Code**

```verbatim
module mult_add (dataa, datab, datac, datad, result);
//synthesis attribute mult_add extract_mac FALSE
// Port Declaration
input [15:0] dataa;
input [15:0] datab;
input [15:0] datac;
input [15:0] datad;
output [32:0] result;

// Wire Declaration
wire [31:0] mult0_result;
wire [31:0] mult1_result;

// Implementation
// Each of these can go into one of the 4 mults in a
// DSP block
assign mult0_result = dataa * `signed datab;
//synthesis attribute mult0_result preserve_signal TRUE
assign mult1_result = datac * datad;

// This adder can go into the one-level adder in a DSP
// block
assign result = (mult0_result + mult1_result);
endmodule
```
Example 20–2. Using Module Level Attributes in VHDL Code

```vhdl
library ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_arith.all;

entity mult_acc is
  generic (size : integer := 4);
  port (a: in std_logic_vector (size-1 downto 0);
    b: in std_logic_vector (size-1 downto 0);
    clk : in std_logic;
    accum_out: inout std_logic_vector (2*size downto 0)
  );
  attribute extract_mac : boolean;
  attribute extract_mac of mult_acc : entity is FALSE;
end mult_acc;

architecture synthesis of mult_acc is
  signal a_int, b_int : signed (size-1 downto 0);
  signal pdt_int : signed (2*size-1 downto 0);
  signal adder_out : signed (2*size downto 0);
  begin
    a_int <= signed (a);
    b_int <= signed (b);
    pdt_int <= a_int * b_int;
    adder_out <= pdt_int + signed(accum_out);
    process (clk)
      begin
        if (clk'event and clk = '1') then
          accum_out <= std_logic_vector (adder_out);
        end if;
      end process;
  end synthesis;
```

### Signal Level Attributes

You can control the implementation of individual LPM_MULT multipliers with the `dedicated_mult` attribute, as shown in the following example:

```
//synthesis attribute <signal_name> dedicated_mult <value>
```

The `dedicated_mult` attribute is only applicable to signals or wires; it is not applicable to registers.

Table 20–4 describes the supported values for the `dedicated_mult` attribute.

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ON</td>
<td>LPM is inferred and multipliers are implemented in the DSP block.</td>
</tr>
<tr>
<td>OFF</td>
<td>LPM is inferred and multipliers are synthesized, implemented in logic, and optimized by the Quartus II software.</td>
</tr>
<tr>
<td>LCELL</td>
<td>LPM is not inferred and multipliers are synthesized, implemented in logic, and optimized by the LeonardoSpectrum software.</td>
</tr>
</tbody>
</table>
Some signals for which the `dedicated_mult` attribute is set may be removed during synthesis by the LeonardoSpectrum software due to design optimization. In such cases, if you want to force the implementation, you can preserve the signal from being removed during synthesis by setting the `preserve_signal` attribute to `true`.

The `extract_mac` attribute must be set to `false` for the module or project level when using the `dedicated_mult` attribute.

Example 20–3 and Example 20–4 are samples of Verilog HDL and VHDL codes, respectively, using the `dedicated_mult` attribute.

### Example 20–3. Signal Attributes for Controlling DSP Block Inference in Verilog HDL Code

```verilog
module mult (AX, AY, BX, BY, m, n, o, p);
  input [7:0] AX, AY, BX, BY;
  output [15:0] m, n, o, p;
  wire [15:0] m_i = AX * AY; // synthesis attribute m_i dedicated_mult ON
                    // synthesis attribute m_i preserve_signal TRUE
                    // Note that the preserve_signal attribute prevents
                    // signal m_i from being removed during synthesis
  wire [15:0] n_i = BX * BY; // synthesis attribute n_i dedicated_mult OFF
  wire [15:0] o_i = AX * BY; // synthesis attribute o_i dedicated_mult AUTO
  wire [15:0] p_i = BX * AY; // synthesis attribute p_i dedicated_mult LCELL
  // since n_i, o_i, p_i signals are not preserved,
  // they may be removed during synthesis based on the design
  assign m = m_i;
  assign n = n_i;
  assign o = o_i;
  assign p = p_i;
endmodule
```

### Table 20–4. Values for the `dedicated_mult` Attribute (Part 2 of 2)

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AUTO</td>
<td>LPM is inferred, but the Quartus II software maps the multipliers automatically to either the DSP block or logic based on resource availability.</td>
</tr>
</tbody>
</table>

Note to Table 20–4:

1. Although both `dedicated_mult=OFF` and `dedicated_mult=LCELLS` result in logic implementations, the optimized results in these two cases may differ.
Guidelines for Using DSP Blocks

In addition to the guidelines mentioned earlier in this section, use the following guidelines when designing with DSP blocks in the LeonardoSpectrum software:

- To access all the control signals for the DSP block, such as sign A, sign B, and dynamic addnsub, use the black box technique.

- While performing signed operations, ensure that the specified data width of the output port matches the data width of the expected result. Otherwise, the sign bit might be lost or data might be incorrect because the sign is not extended. For example, if the data widths of input A and B are width_a and width_b, respectively, the maximum data width of the result can be (width_a + width_b +2) for the four-multipliers adder mode. Thus, the data width of the output port should be less than or equal to (width_a + width_b +2).

- While using the accumulator, the data width of the output port should be equal to or greater than (width_a + width_b). The maximum width of the accumulator can be (width_a + width_b + 16). Accumulators wider than this are implemented in logic.

Example 20–4. Signal Attributes for Controlling DSP Block Inference in VHDL Code

```vhdl
library ieee;
USE ieee.std_logic_1164.all;
USE ieee.std_logic_arith.all;
USE ieee.std_logic_unsigned.all;
USE ieee.std_logic_signed.all;
ENTITY mult is
PORT( AX,AY,BX,BY: IN std_logic_vector (17 DOWNTO 0);
m,n,o,p: OUT std_logic_vector (35 DOWNTO 0));
attribute dedicated_mult: string;
attribute preserve_signal : boolean
END mult;
ARCHITECTURE struct of mult is
signal m_i, n_i, o_i, p_i : unsigned (35 downto 0);
attribute dedicated_mult of m_i:signal is "ON";
attribute dedicated_mult of n_i:signal is "OFF";
attribute dedicated_mult of o_i:signal is "AUTO";
attribute dedicated_mult of p_i:signal is "LCELL";
begnin
m_i <= unsigned (AX) * unsigned (AY);
n_i <= unsigned (BX) * unsigned (BY);
o_i <= unsigned (AX) * unsigned (BY);
p_i <= unsigned (BX) * unsigned (AY);
m <= std_logic_vector(m_i);
n <= std_logic_vector(n_i);
o <= std_logic_vector(o_i);
p <= std_logic_vector(p_i);
end struct;
```
If the design uses more multipliers than are available in a particular device, the Quartus II software may issue a no fit error. In such cases, use the attribute settings in the LeonardoSpectrum software to control the mapping of multipliers in your design to DSP blocks or logic.

**Block-Based Design with the Quartus II Software**

The incremental compilation design flow with LogicLock™ constraints enables users to design, optimize, and lock down a design one section at a time. You can independently create and implement each logic module into a hierarchical or team-based design. With this method, you can preserve the performance of each module during system integration and have more control over the placement of your design. To maximize the benefits of the incremental compilation in the Quartus II software, you can partition a new design into a hierarchy of netlist files during synthesis in the LeonardoSpectrum software.

You can create different netlist files with the LeonardoSpectrum software for different sections of a design hierarchy. Having different netlist files means that each section is independent of the others. When synthesizing the entire project, only portions of a design that have been updated must be resynthesized when you compile the design. You can make changes, optimize, and resynthesize your section of a design without affecting other sections.

For more information about incremental compilation, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*. For more information about the LogicLock feature, refer to the *Analyzing and Optimizing the Design Floorplan* chapter in volume 2 of the *Quartus II Handbook*.

**Hierarchy and Design Considerations**

You must plan your design’s structure and partitioning carefully to use incremental compilation and LogicLock features effectively. Optimal hierarchical design practices include partitioning the blocks at functional boundaries, registering the boundaries of each block, minimizing the I/O between each block, separating timing-critical blocks, and keeping the critical path within one hierarchical block.

For more recommendations for hierarchical design partitioning, refer to the *Design Recommendations for Altera Devices and the Quartus II Design Assistant* chapter in volume 1 of the *Quartus II Handbook*.

To ensure the proper functioning of the synthesis tool, you can apply the LogicLock option in the LeonardoSpectrum software only to modules, entities, or netlist files. In addition, each module or entity should have its own design file. It is difficult to maintain incremental synthesis if two different modules are in the same design file, but are defined as being part of different regions, because both regions must be recompiled when you change one of the modules or entities.
If you use boundary tri-states in a lower-level block, the LeonardoSpectrum software pushes, or bubbles, the tri-states through the hierarchy to the top-level to take advantage of the tri-state drivers on the output pins of the Altera device. Because bubbling tri-states require optimizing through hierarchies, lower-level tri-states are not supported with a block-level design methodology. You should use tri-state drivers only at the external output pins of the device and in the top-level block in the hierarchy.

If the hierarchy is flattened during synthesis, logic is optimized across boundaries, preventing you from making LogicLock assignments to the flattened blocks. Altera recommends preserving the hierarchy when compiling the design. In the Optimize command of your script, use the Hierarchy Preserve command, or in the user interface, specify Preserve in the Hierarchy section of the Optimize FlowTab.

If you are compiling your design with a script, you can use an alternative method for preventing optimization across boundaries. In this case, use the Auto hierarchy setting and set the auto_dissolve attribute to false on the instances or views that you want to preserve (that is, the modules with LogicLock assignments) using the following syntax:

```
set_attribute -name auto_dissolve -value false \
   .work.<block1>.INTERFACE
```

This alternative method flattens your design according to the auto_dissolve limits, but does not optimize across boundaries where you apply the attribute.

For more details about LeonardoSpectrum attributes and hierarchy levels, refer to the LeonardoSpectrum documentation in the Help menu.

Creating a Design with Multiple .edf Files

The first stage of a hierarchical design flow is to generate multiple .edf files, which allows you to take advantage of the incremental compilation flow in the Quartus II software. If the whole design is in one .edf file, changes in one block affect other blocks because of possible node name changes. You can generate multiple .edf files either by using the LogicLock option in the LeonardoSpectrum software, or by manually using a black box methodology on each block that you want to include in a LogicLock region.

After you create multiple .edf files with one of these methods, you must create the appropriate Quartus II project(s) to place-and-route the design.

Generating Multiple .edf Files Using the LogicLock Option

This section describes how to generate multiple .edf files using the LogicLock option in the LeonardoSpectrum software.

When synthesizing a top-level design that includes LogicLock regions, perform the following general steps:

1. Read in the Verilog HDL or VHDL source files.
2. Add LogicLock constraints.
3. Optimize and write output netlist files, or click Run Flow.
To set the correct constraints and compile the design, perform the following steps in the LeonardoSpectrum software:

1. On the Tools menu, switch to the **Advanced** FlowTab instead of the **Quick Setup** tab.
2. Set the target technology and speed grade for the device on the **Technology FlowTab**.
3. Open the input source files on the **Input FlowTab**.
4. Click **Read** on the **Input FlowTab** to read the source files, but not begin optimization.
5. Click the **Module PowerTab** located at the bottom of the **Constraints FlowTab**.
6. Select a module to be placed in a LogicLock region in the **Modules** section.
7. Turn on **LogicLock**.
8. Type the desired LogicLock region name under **LogicLock**.
9. Click **Apply**.
10. Repeat steps 6-9 for any other modules that you want to place in LogicLock regions.

In some cases, you are prompted to save your LogicLock and other non-global constraints in a Constraints File (.ctr) when you click anywhere off the **Constraints FlowTab**. The default name is `<project name>.ctr`. This file is added to your **Input** file list, and must be manually included later if you recreate the project.

The command written into the LeonardoSpectrum Information or Transcript Window is the Tcl command that is written into the .ctr file. The format of the “path” for the module specified in the command should be work.<module>.INTERFACE. To ensure that you do not see an optimized version of the module, do not perform a **Run Flow** on the **Quick Setup** tab prior to setting LogicLock constraints. Always use the **Read** command, as described in step 4.

11. Make any other settings as required on the **Constraints FlowTab**.
12. Select **Preserve** in the **Hierarchy** section of the **Optimize FlowTab** to ensure that the hierarchy names are not flattened during optimization.
13. Make any other settings as required on the **Optimize FlowTab**.
14. Run your synthesis flow with each FlowTab, or click **Run Flow**.

Synthesis creates an .edf file for each module that has a LogicLock assignment in the **Constraints FlowTab**. You can now use these files with the incremental compilation flow in the Quartus II software.

You might occasionally see multiple .edf files and LogicLock commands for the same module. An “unfolded” version of a module is created when you instantiate a module more than once and the boundary conditions of the instances are different. For example, if you apply a constant to one instance of the block, it might be optimized to eliminate unneeded logic. In this case, the LeonardoSpectrum software must create a
separate module for each instantiation (unfolding). If this unfolding occurs, you see more than one .edf file, and each .edf file has a LogicLock assignment to the same LogicLock region. When you import the .edf files to the Quartus II software, the .edf files created from the module are placed in different LogicLock regions. Any optimizations performed in the Quartus II software using the LogicLock methodology must be performed separately for each .edf file.

Creating a Quartus II Project for Multiple .edf Files Including LogicLock Regions

The LeonardoSpectrum software creates .tcl files that provide the Quartus II software with the appropriate LogicLock assignments, creating a region for each .edf file along with the information to set up a Quartus II project.

The .tcl file contains the commands shown in Example 20–5 for each LogicLock region. This example is for module taps where the name taps_region is typed as the LogicLock region name in the Constraints FlowTab in the LeonardoSpectrum software.

Example 20–5. Tcl File for Module Taps with taps_region as LogicLock Region Name

```tcl
project add_assignment {taps} {taps_region} {} {}
{LL_AUTO_SIZE} {ON}
project add_assignment {taps} {taps_region} {} {}
{LL_STATE} {FLOATING}
project add_assignment {taps} {taps_region} {} {}
{LL_MEMBER_OF} {taps_region}
```

These commands create a LogicLock region with auto-size and floating-origin properties. This flexible LogicLock region allows the Quartus II Compiler to select the size and location of the region.

For more information about Tcl commands, refer to the Tcl Scripting chapter in Volume 2 of the Quartus II Handbook.

You can use the following methods to import the .edf file and corresponding .tcl file into the Quartus II software:

- Use the .tcl file that is created for each .edf file by the LeonardoSpectrum software. This method allows you to generate multiple Quartus II projects, one for each block in the design. Each designer in the project can optimize their block separately in the Quartus II software and preserve their results. Altera recommends this method for bottom-up incremental and hierarchical design methodologies because it allows each block in the design to be treated separately. Each block can be brought into one top-level project with the import function.

or

- Use the <top-level project>.tcl file that contains the assignments for all blocks in the project. This method allows the top-level designer to import all the blocks into one Quartus II project. You can optimize all modules in the project at once in a top-down design flow. If additional optimization is required for individual blocks, each designer can use their .edf file to create a separate project at that time. You must then add new assignments to the top-level project using the import function.
In both methods, use the following steps to create the Quartus II project, import the appropriate LogicLock assignments, and compile the design:

1. Place the .edf and .tcl files in the same directory.
2. On the View menu, point to Utility Windows and click Tcl Console to open the Quartus II Tcl Console.
3. At the Tcl prompt, type `source <path>/<project name>.tcl`.
4. To open the newly completed project, on the File menu, click Open Project. Browse to and select the project name, and click Open.

For more information about importing a design using incremental compilation, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook. For more information about importing LogicLock assignments, see the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

**Generating Multiple .edf Files Using Black Boxes**

This section describes how to manually generate multiple .edf files using the black box technique. The manual flow, which was supported in older versions of the LeonardoSpectrum software, is discussed here because some designers want more control over the project for each submodule.

To create multiple .edf files in the LeonardoSpectrum software, create a separate project for each module and top-level design that you want to maintain as a separate .edf file. Implement black box instantiations of lower-level modules in your top-level project.

When synthesizing the projects for the lower-level modules and the top-level design, use the following general guidelines:

For lower-level modules:
- Turn off Map IO Registers for the target technology on the Technology FlowTab.
- Read the HDL files for the modules. Modules may include black box instantiations of lower-level modules that are also maintained as separate .edf files.
- Add constraints.
- Turn off Add I/O Pads on the Optimize FlowTab.

For the top-level design:
- Turn on Map IO Registers if you want to implement input and/or output registers in the IOEs for the target technology on the Technology FlowTab.
- Read the HDL files for the top-level design.
- Black box lower-level modules in the top-level design.
- Add constraints (clock settings should be made at this time).

The following sections describe examples of black box modules in a block-based and team-based design flow.
In Figure 20–3, the top-level module A is assigned to one engineer (designer 1), while two engineers work on the lower levels of the design. Designer 2 works on module B and its submodules D and E, while designer 3 works on module C and its submodule F.

**Figure 20–3. Block-Based and Team-Based Design Example**

One netlist is created for the top-level module A, another netlist is created for module B and its submodules D and E, and another netlist is created for module C and its submodule F. To create multiple .edf files, perform the following steps:

1. Generate an .edf file for module C. Use C.v and E.v as the source files.
2. Generate an .edf file for module B. Use B.v, D.v, and E.v as the source files.
3. Generate a top-level .edf file A.v for module A. Ensure that your black box modules B and C were optimized separately in steps 1 and 2.

**Black Box Methodology in Verilog HDL**

Any design block that is not defined in the project, or included in the list of files to be read for a project, is treated as a black box by the software. In Verilog HDL, you must also provide an empty module declaration for the module that you plan to treat as a black box.
Example 20–6 shows an example of the A.v top-level file. If any of your lower-level files also contain a black-boxed lower-level file in the next level of hierarchy, follow the same procedure.

**Example 20–6. Verilog HDL Top-Level File Black Boxing Example**

```verilog
module A (data_in, clk, e, ld, data_out);
    input data_in, clk, e, ld;
    output [15:0] data_out;

    reg [15:0] cnt_out;
    reg [15:0] reg_a_out;

    B U1 ( .data_in (data_in), .clk (clk), .e(e), .ld (ld),
           .data_out (cnt_out) );

    C U2 ( .d (cnt_out), .clk (clk), .e (e), .q (reg_out) );
    // Any other code in A.v goes here.
endmodule

// Empty Module Declarations of Sub-Blocks B and C follow here.
// These module declarations (including ports) are required for blackboxing.

module B (data_in, e, ld, data_out);
    input data_in, clk, e, ld;
    output [15:0] data_out;
endmodule

module C (d, clk, e, q);
    input d, clk, e;
    output [15:0] q;
endmodule
```

Previous versions of the LeonardoSpectrum software require the attribute statement `//exemplar attribute U1 NOOPT TRUE`, which instructs the software to treat the instance U1 as a black box. This attribute is no longer required, although it is still supported in the software.

**Black Boxing in VHDL**

Any design block that is not defined in the project, or included in the list of files to be read for a project, is treated as a black box by the LeonardoSpectrum software. In VHDL, a component declaration is required for the black box.
Example 20–7 shows an example of the A.vhd top-level file. If any of your lower-level files also contain a black-boxed lower-level file in the next level of hierarchy, follow the same procedure.

**Example 20–7. VHDL Top-Level File Black Boxing Example**

```vhdl
LIBRARY ieee;
USE ieee.std_logic_1164.all;
ENTITY A IS
PORT ( data_in : IN INTEGER RANGE 0 TO 15;
    clk : IN STD_LOGIC;
    e : IN STD_LOGIC;
    ld : IN STD_LOGIC;
    data_out : OUT INTEGER RANGE 0 TO 15
);
END A;

ARCHITECTURE a_arch OF A IS
COMPONENT B PORT(
    data_in : IN INTEGER RANGE 0 TO 15;
    clk : IN STD_LOGIC;
    e : IN STD_LOGIC;
    ld : IN STD_LOGIC;
    data_out : OUT INTEGER RANGE 0 TO 15
);
END COMPONENT;

COMPONENT C PORT(
    d : IN INTEGER RANGE 0 TO 15;
    clk : IN STD_LOGIC;
    e : IN STD_LOGIC;
    q : OUT INTEGER RANGE 0 TO 15
);
END COMPONENT;

-- Other component declarations in A.vhd go here

signal cnt_out : INTEGER RANGE 0 TO 15;
signal reg_a_out : INTEGER RANGE 0 TO 15;
BEGIN
CNT : C
PORT MAP (  
data_in => data_in,
    clk => clk,
    e => e,
    ld => ld,
    data_out => cnt_out
);

REG_A : D
PORT MAP (  
d => cnt_out,
    clk => clk,
    e => e,
    q => reg_a_out
);

-- Any other code in A.vhd goes here
END a_arch;
```
Previous versions of the LeonardoSpectrum software require the attribute statement
\texttt{noopt of C: component is TRUE}, which instructs the software to treat the component
C as a black box. This attribute is no longer required, although it is still supported in
the software.

After you have completed the steps outlined in this section, you have a different .edf
netlist file for each block of code. You can now use these files for the incremental
compilation flow in the Quartus II software.

Creating a Quartus II Project for Multiple .edf Files

The LeonardoSpectrum software creates a .tcl file for each .edf file, which provides
the Quartus II software with the information to set up a project.

There are two different methods for importing each .edf file and corresponding .tcl
file into the Quartus II software:

1. Use the .tcl file that is created for each .edf file by the LeonardoSpectrum software.
   This method generates multiple Quartus II projects, one for each block in the
design. Each designer in the project can optimize their block separately in the
Quartus II software and preserve their results. Designers should create a
LogicLock region for each block; the top-level designer should then import all the
blocks and assignments into the top-level project. Altera recommends this method
for bottom-up incremental and hierarchical design methodology because it allows
each block in the design to be treated separately; each block can be imported into
one top-level project.

2. Use the \texttt{<top-level project>.tcl} file that contains the information to set up the
top-level project. This method allows the top-level designer to create LogicLock
regions for each block and bring all the blocks into one Quartus II project.
Designers can optimize all modules in the project at once in a top-down design
flow. If additional optimization is required for individual blocks, each designer
creates a separate Quartus II project with each .edf file. New assignments must
then be added to the top-level project manually or through the import function.

For more information about importing designs using incremental compilation, refer
to the \textit{Quartus II Incremental Compilation for Hierarchical and Team-Based Design}
chapter in volume 1 of the \textit{Quartus II Handbook}. For more information about importing
LogicLock regions, refer to the \textit{Analyzing and Optimizing the Design Floorplan}
chapter in volume 2 of the \textit{Quartus II Handbook}.

With both methods, use the following steps to create the Quartus II project and
compile the design:

1. Place the .edf and .tcl files in the same directory.
2. On the View menu, point to \textbf{Utility Windows} and click \textbf{Tcl Console}. The
   Quartus II Tcl Console appears.
3. At a Tcl prompt, type \texttt{source <path>/<project name>.tcl}.
4. On the File menu, click \textbf{Open Project}. In the New Project window, browse to and
   select the project name. Click \textbf{Open}.
5. To create LogicLock assignments, on the Assignments menu, click **LogicLock Regions Window**. Use the LogicLock Regions window to create LogicLock regions.

6. On the Processing menu, click **Start Compilation**.

**Incremental Synthesis Flow**

If you make changes to one or more submodules, you can manually create new projects in the LeonardoSpectrum software to generate a new .edf file when there are changes to the source files. Alternatively, you can use incremental synthesis to generate a new netlist for the changed submodule(s). To perform incremental synthesis in the LeonardoSpectrum software, use the script described in this section to reoptimize and generate a new .edf file for only the affected modules using the LeonardoSpectrum top-level project. This method applies only when you are using the **LogicLock** option in the LeonardoSpectrum software.

**Modifications Required for the LogicLock_Incremental.tcl Script File**

There are three sets of entries in the file that must be modified before beginning incremental synthesis. The variables in the .tcl file are surrounded by angle brackets (< >).

1. Add the list of source files that are included in the project. You can enter the full path to the file or just the file name if the files are located in the working directory.

2. Indicate which modules in the design have changed. These modules are the .edf files that are regenerated by the LeonardoSpectrum software and contain a LogicLock assignment in the original compilation.

   Determine the LeonardoSpectrum software path for each module by looking at the .ctr file that contains the LogicLock assignments from the original project. Each LogicLock assignment is applied to a particular module in the design.

3. Enter the target device family using the appropriate device keyword. The device keyword is displayed on the Transcript or Information window when you select a target Technology and click **Load Library** or **Apply** on the **Technology** FlowTab.
Example 20–8 shows the LogicLock_Incremental.tcl file for the incremental synthesis flow. You must modify the .tcl file before you can use it for your project.

**Example 20–8. LogicLock_Interface.tcl Script File for Incremental Synthesis**

```
# You must indicate which modules have changed (based on the source files that have changed) and provide the complete path to each module
# You must also specify the list of design files and the target Altera technology being used

# Read the design source files.
read <list of design files separated by spaces (such as block1.v block2.v)>

# Get the list of modified modules in bottom-up "depth first search" order
# where the lower-level blocks are listed first (these should be modules that had LogicLock assignments and separate EDIF netlist files in the first pass and had their source code modified)
set list_of_modified_modules {.work.<block2>.INTERFACE .work.<block1>.INTERFACE}

foreach module $list_of_modified_modules {
    set err_rc [regexp {(.*)\.(.*)\.(.*)} $module unused lib module_name arch]
    present_design $module
    # Run optimization, preserving hierarchy. You must specify a technology.
    optimize -ta <technology> -hierarchy preserve
    # Ensure that the lower-level module is not optimized again when
    # optimizing higher-level modules.
    dont_touch $module
}

foreach module $list_of_modified_modules {
    set err_rc [regexp {(.*)\.(.*)\.(.*)} $module unused lib module_name arch]
    present_design $module
    undont_touch $module
    auto_write $module_name.edf
    # Ensure that the lower-level module is not written out in the EDIF file
    # of the higher-level module.
    noopt $module
}
```

**Running the Tcl Script File in LeonardoSpectrum**

After you modify the Tcl script, as described in “Modifications Required for the LogicLock_Incremental.tcl Script File” on page 20–25, you can compile your design using the script.

You can run the script in batch mode at the command-line prompt using the following command:

```
spectrum -file <Tcl_file>
```

To run the script from the GUI, on the File menu, click **Run Script**, then browse to your .tcl file and click **Open**.
The LogicLock incremental design flow uses module-based design to help you preserve performance of modules and have control over placement. By tagging the modules that require separate .edf files, you can make multiple .edf files for use with the Quartus II software from a single LeonardoSpectrum software project.

**Conclusion**

Taking advantage of the Mentor Graphics LeonardoSpectrum software and the Quartus II design flow allows you to control how your design files are prepared for the Quartus II place-and-route process, as well as to improve performance and optimize a design for use with Altera devices. The methodologies outlined in this chapter can help optimize a design to achieve performance goals and save design time with the LeonardoSpectrum software. For the best results with new designs in new device families, Altera recommends migrating to the advanced Mentor Graphics Precision RTL Synthesis software.

**Document Revision History**

Table 20–5 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
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<td>Changed to new document template.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed support for the Classic Timing Analyzer.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Editorial changes.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Updated supported devices.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed Referenced Documents section.</td>
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<tr>
<td></td>
<td></td>
<td>Minor updates for the version 10.0 release.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Minor updates for the Quartus II software version 9.1 release.</td>
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<tr>
<td></td>
<td></td>
<td>Removed Table 12–3, Inferring RAM Summary.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>No change to content.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Chapter 12 was previously Chapter 11 in software release 8.1.</td>
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<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8-1/2” x 11” page size.</td>
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<tr>
<td></td>
<td></td>
<td>Updated Table 12–3.</td>
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<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Updated date and part number and added hypertext links.</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter describes how you can use the Quartus® II Netlist Viewers to analyze and debug your designs.

As FPGA designs grow in size and complexity, the ability to analyze, debug, optimize, and constrain your design is critical. With today’s advanced designs, several design engineers are involved in coding and synthesizing different design blocks, making it difficult to analyze and debug the design. The Quartus II RTL Viewer, State Machine Viewer, and Technology Map Viewer provide powerful ways to view your initial and fully mapped synthesis results during the debugging, optimization, and constraint entry processes.

This chapter contains the following sections:

- “When to Use the Netlist Viewers: Analyzing Design Problems”
- “Quartus II Design Flow with the Netlist Viewers” on page 21–3
- “RTL Viewer Overview” on page 21–4
- “State Machine Viewer Overview” on page 21–5
- “Technology Map Viewer Overview” on page 21–5
- “Introduction to the User Interface” on page 21–6
- “Navigating the Schematic View” on page 21–22
- “Filtering in the Schematic View” on page 21–30
- “Probing to a Source Design File and Other Quartus II Windows” on page 21–36
- “Probing to the Netlist Viewers from Other Quartus II Windows” on page 21–37
- “Viewing a Timing Path” on page 21–38
- “Other Features in the Schematic Viewer” on page 21–39
When to Use the Netlist Viewers: Analyzing Design Problems

You can use the netlist viewers to analyze and debug your design. This section provides simple examples of how to use the RTL Viewer, State Machine Viewer, and Technology Map Viewer to analyze problems encountered in the design process.

The following sections contain information about how the netlist viewers display your design:

- “Quartus II Design Flow with the Netlist Viewers” on page 21–3
- “RTL Viewer Overview” on page 21–4
- “State Machine Viewer Overview” on page 21–5
- “Technology Map Viewer Overview” on page 21–5

Using the RTL Viewer is a good way to view your initial synthesis results to determine whether you have created the necessary logic, and that the logic and connections have been interpreted correctly by the software. You can use the RTL Viewer and State Machine Viewer to check your design visually before simulation or other verification processes. Catching design errors at this early stage of the design process can save you valuable time.

If you see unexpected behavior during verification, use the RTL Viewer to trace through the netlist and ensure that the connections and logic in your design are as expected. You can also view state machine transitions and transition equations with the State Machine Viewer. Viewing your design helps you find and analyze the source of design problems. If your design looks correct in the RTL Viewer, you know to focus your analysis on later stages of the design process and investigate potential timing violations or issues in the verification flow itself.

You can use the Technology Map Viewer to look at the results at the end of Analysis and Synthesis. If you have compiled your design through the Fitter stage, you can view your post-mapping netlist in the Technology Map Viewer (Post-Mapping) and your post-fitting netlist in the Technology Map Viewer. If you perform only Analysis and Synthesis, both the netlist viewers display the same post-mapping netlist.

In addition, you can use the RTL Viewer or Technology Map Viewer to locate the source of a particular signal, which can help you debug your design. Use the navigation techniques described in this chapter to search easily through your design. You can trace back from a point of interest to find the source of the signal and ensure the connections are as expected.

The Technology Map Viewer can help you locate post-synthesis nodes in your netlist and make assignments when optimizing your design. This functionality is useful when making a multicycle clock timing assignment between two registers in your design. Start at an I/O port and trace forward or backward through the design and through levels of hierarchy to find nodes of interest, or locate a specific register by visually inspecting the schematic.

You can use the RTL Viewer, State Machine Viewer, and Technology Map Viewer in many other ways throughout the design, debug, and optimization stages. This chapter shows you how to use the various features of the netlist viewers to increase your productivity when analyzing a design.
Quartus II Design Flow with the Netlist Viewers

When you first open one of the netlist viewers after compiling the design, a preprocessor stage runs automatically before the netlist viewer opens. If you close the netlist viewer and open it again later without recompiling the design, the netlist viewer opens immediately without performing the preprocessing stage.

Figure 21–1 shows how the netlist viewers fit into the basic Quartus II design flow.

Before the netlist viewer can run the preprocessor stage, you must compile your design:

- To open the RTL Viewer or State Machine Viewer, first perform Analysis and Elaboration.
- To open the Technology Map Viewer (Post-Fitting) or the Technology Map Viewer (Post-Mapping), first perform Analysis and Synthesis.

The netlist viewers display the results of the last successful compilation. Therefore, if you make a design change that causes an error during Analysis and Elaboration, you cannot view the netlist for the new design files, but you can still see the results from the last successfully compiled version of the design files. If you receive an error during compilation and you have not yet successfully run the appropriate compilation stage for your project, the netlist viewer cannot be displayed; in this case, the Quartus II software issues an error message when you try to open the netlist viewer.

If the netlist viewer is open when you start a new compilation, the netlist viewer closes automatically. You must open the netlist viewer again to view the new design netlist after compilation completes successfully.
RTL Viewer Overview

The Quartus II RTL Viewer allows you to view a register transfer level (RTL) graphical representation of your Quartus II integrated synthesis results or your third-party netlist file in the Quartus II software.

You can view results after Analysis and Elaboration when your design uses any supported Quartus II design entry method, including Verilog HDL Design Files (.v), SystemVerilog Design Files (.sv), VHDL Design Files (.vhd), AHDL Text Design Files (.tdf), schematic Block Design Files (.bdf), or schematic Graphic Design Files (.gdf) imported from the MAX+PLUS® II software. You can also view the hierarchy of atom primitives (such as device logic cells and I/O ports) when your design uses a synthesis tool to generate a Verilog Quartus Mapping File (.vqm) or Electronic Design Interchange Format (.edf) file. For a flow diagram, refer to Figure 21–1.

The Quartus II RTL Viewer displays a schematic view of the design netlist after Analysis and Elaboration or netlist extraction is performed by the Quartus II software, but before technology mapping and any synthesis or fitter optimizations. This view is not the final design structure because optimizations have not yet occurred. This view most closely represents your original source design. If you synthesized your design with the Quartus II integrated synthesis, this view shows how the Quartus II software interpreted your design files. If you use a third-party synthesis tool, this view shows the netlist written by your synthesis tool.

When displaying your design, the RTL Viewer optimizes the netlist to maximize readability in the following ways:

- Logic with no fan-out (its outputs are unconnected) and logic with no fan-in (its inputs are unconnected) are removed from the display.
- Default connections such as \( V_{CC} \) and GND are not shown.
- Pins, nets, wires, module ports, and certain logic are grouped into buses where appropriate.
- Constant bus connections are grouped.
- Values are displayed in hexadecimal format.
- NOT gates are converted to bubble inversion symbols in the schematic.
- Chains of equivalent combinational gates are merged into a single gate. For example, a 2-input AND gate feeding a 2-input AND gate is converted to a single 3-input AND gate.
- State machine logic is converted into a state diagram, state transition table, and state encoding table, which are displayed in the State Machine Viewer.

To run the RTL Viewer for a Quartus II project, first analyze the design to generate an RTL netlist. To analyze the design and generate an RTL netlist, on the Processing menu, point to Start and click Start Analysis & Elaboration. You can also perform a full compilation on any process that includes the initial Analysis and Elaboration stage of the Quartus II compilation flow.

To run the RTL Viewer, on the Tools menu, point to Netlist Viewers and click RTL Viewer.
You can set the RTL Viewer preprocessing to run during a full compilation, which allows you to open the RTL Viewer after Analysis and Synthesis has completed, but while the Fitter is still running. In this case, you do not have to wait for the Fitter to finish before viewing the schematic. This technique is useful for a large design that requires a substantial amount of time in the place-and-route stage.

To set the RTL Viewer preprocessing to run during compilation, on the Assignments menu, click Settings. In the Category list, select Compilation Process Settings and turn on Run RTL Viewer preprocessing during compilation. By default, this option is turned off.

**State Machine Viewer Overview**

The State Machine Viewer presents a high-level view of finite state machines in your design. The State Machine Viewer provides a graphical representation of the states and their related transitions, as well as a state transition table that displays the condition equation for each of the state transitions, and encoding information for each state.

To run the State Machine Viewer, on the Tools menu, point to Netlist Viewers and click State Machine Viewer. To open the State Machine Viewer for a particular state machine, double-click the state machine instance in the RTL Viewer or right-click the state machine instance and click Hierarchy Down.

**Technology Map Viewer Overview**

The Quartus II Technology Map Viewer provides a technology-specific, graphical representation of your design after Analysis and Synthesis or after the Fitter has mapped your design into the target device. The Technology Map Viewer shows the hierarchy of atom primitives (such as device logic cells and I/O ports) in your design. For supported families, you can also view internal registers and look-up tables (LUTs) inside logic cells (LCELLs) and registers in I/O atom primitives. For more information, refer to “Viewing Contents of Atom Primitives” on page 21–23.

Where possible, the port names of each hierarchy are maintained throughout synthesis; however, port names might change or be removed from the design. For example, if a port is unconnected or driven by GND or VCC, it is removed during synthesis. When a port name changes, the port is assigned a related user logic name in the design or a generic port name such as IN1 or OUT1.

You can view your Quartus II technology-mapped results after synthesis, fitting, or timing analysis. To run the Technology Map Viewer for a Quartus II project, on the Processing menu, point to Start and click Start Analysis & Synthesis to synthesize and map the design to the target technology. At this stage, the Technology Map Viewer shows the same post-mapping netlist as the Technology Map Viewer (Post-Mapping). You can also perform a full compilation, or any process that includes the synthesis stage in the compilation flow.
If you have completed the Fitter stage, the Technology Map Viewer shows the changes made to your netlist by the Fitter, such as physical synthesis optimizations, while the Technology Map Viewer (Post-Mapping) shows the post-mapping netlist. If you have completed the Timing Analysis stage, you can locate timing paths from the Timing Analyzer report in the Technology Map Viewer (for more information, refer to “Viewing a Timing Path” on page 21–38). For a flow diagram, refer to Figure 21–1 on page 21–3.

To run the Technology Map Viewer, on the Tools menu, point to Netlist Viewers and click Technology Map Viewer, or select Technology Map Viewer from the Applications toolbar.

To run the Technology Map Viewer (Post-Mapping), on the Tools menu, point to Netlist Viewers and click Technology Map Viewer (Post-Mapping).

**Introduction to the User Interface**

The RTL Viewer and Technology Map Viewer each consist of three main parts:

- The schematic view—displays a graphical representation of the internal structure of your design.
- The Netlist Navigator pane—displays a representation of the project hierarchy.
- The Find pane—allows you to find and locate specific design elements in the schematic view.

Figure 21–2 shows the RTL Viewer and indicates these three parts, along with other elements of the user interface. The netlist viewers also contain a toolbar that provides tools to use in the schematic view.

You can have only one RTL Viewer, one Technology Map Viewer, one Technology Map Viewer (Post-Mapping), and one State Machine Viewer window open at the same time, although each window can show multiple pages. For example, you cannot have two RTL Viewer windows open at the same time.
**Chapter 21: Analyzing Designs with Quartus II Netlist Viewers**

**Introduction to the User Interface**

Figure 21–2 shows the schematic view and the Netlist Navigator pane of the RTL Viewer.

**Schematic View**

The schematic view is shown on the right side of the RTL Viewer and Technology Map Viewer. The schematic view contains a schematic representing the design logic in the netlist. This view is the main screen for viewing your gate-level netlist in the RTL Viewer and your technology-mapped netlist in the Technology Map Viewer.

**Global Net Routing**

In the Quartus II software version 12.1 onwards, you can use the Global Net Routing feature to trace signal path in the schematic. To enable the feature, in the RTL Viewer, in the Tools menu, click **Options**. In the **Options** dialog box, turn on the **Enable global net routing** option.
Displaying Schematics in a Single Page

In the Quartus II software version 12.1 onwards, the RTL Viewer and Technology Map Viewer display schematic with less than 500 nodes in a single page view by default. You can view a schematic with more than 500 nodes in multiple pages. You can highlight a net and use the hand tool to trace the signal in a single page.

Displaying Schematics in Multiple Tabbed View

The RTL Viewer and Technology Map Viewer support multiple tabbed views. With multiple tabbed view, schematics can be displayed in different tabs. Each tabbed view has independent mode actions. Selection is synchronous among all tabbed views. The maximum number of tabs allowed is eight. To create an additional tab, right click in the schematic, and then click Duplicate View.

Schematic Symbols

The symbols for nodes in the schematic represent elements of your design netlist. These elements include input and output ports, registers, logic gates, Altera® primitives, high-level operators, and hierarchical instances.

Figure 21–3 shows an example of an RTL Viewer schematic for a 3-bit synchronous loadable counter. Example 21–1 shows the Verilog HDL code that produced this schematic. This example includes multiplexers and a group of registers (Table 21–1) in a bus along with an ADDER operator (Table 21–3 on page 21–13) inferred by the counting function in the HDL code.

The schematic in Figure 21–3 displays wire connections between nodes with a thin black line and bus connections with a thick black line.

Figure 21–3. Example Schematic Diagram in the RTL Viewer
**Example 21–1. Code Sample for Counter Schematic Shown in Figure 21–3**

```verilog
module counter (input [2:0] data, input clk, input load, output [2:0] result);
    reg [2:0] result_reg;
    always @(posedge clk)
        if (load)
            result_reg <= data;
        else
            result_reg <= result_reg + 1;
    assign result = result_reg;
endmodule
```

Figure 21–4 shows a portion of the corresponding Technology Map Viewer schematic with a compiled design that targets a Stratix® device. In this schematic, you can see the LCELL (logic cell) device-specific primitives that represent the counter function, labeled with their post-synthesis node names. The `REGOUT` port represents the output of the register in the LCELL; the `COMBOUT` port represents the output of the combinational logic in the LUT of the LCELL. The hexadecimal number in parentheses below each LCELL primitive represents the LUT mask, which is a hexadecimal representation of the logic function of the LCELL.

**Figure 21–4. Example Schematic Diagram in the Technology Map Viewer**

Table 21–1 lists and describes the primitives and basic symbols that you can display in the schematic view of the RTL Viewer and Technology Map Viewer. Table 21–3 on page 21–13 lists and describes the additional higher-level operator symbols in the RTL Viewer schematic view.
The logic gates and operator primitives appear only in the RTL Viewer. Logic in the Technology Map Viewer is represented by atom primitives, such as registers and LCELLs.

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>I/O Ports</strong></td>
<td>An input, output, or bidirectional port in the current level of hierarchy. A device input, output, or bidirectional pin when viewing the top-level hierarchy. The symbol can also represent a bus. Only one wire is shown connected to the bidirectional symbol, representing the input and output paths. Input symbols appear on the left-most side of the schematic. Output and bidirectional symbols appear on the right-most side of the schematic.</td>
</tr>
<tr>
<td><strong>I/O Connectors</strong></td>
<td>An input or output connector, representing a net that comes from another page of the same hierarchy (refer to “Partitioning the Schematic into Pages” on page 21–28). To go to the page that contains the source or the destination, right-click on the net and choose the page from the menu (refer to “Following Nets Across Schematic Pages” on page 21–29).</td>
</tr>
<tr>
<td><strong>Hierarchy Port Connector</strong></td>
<td>A connector representing a port relationship between two different hierarchies. A connector indicates that a path passes through a port connector in a different level of hierarchy.</td>
</tr>
<tr>
<td><strong>OR, AND, XOR Gates</strong></td>
<td>An OR, AND, or XOR gate primitive (the number of ports can vary). A small circle (bubble symbol) on an input or output port indicates the port is inverted.</td>
</tr>
<tr>
<td><strong>MULTIPLEXER</strong></td>
<td>A multiplexer primitive with a selector port that selects between port 0 and port 1. A multiplexer with more than two inputs is displayed as an operator (refer to “Operator Symbols in the RTL Viewer Schematic View” on page 21–13).</td>
</tr>
<tr>
<td><strong>BUFFER</strong></td>
<td>A buffer primitive. The figure shows the tri-state buffer, with an inverted output enable port. Other buffers without an enable port include LCELL, SOFT, CARRY, and GLOBAL. The NOT gate and EXP expander buffers use this symbol without an enable port and with an inverted output port.</td>
</tr>
</tbody>
</table>
| **CARRY_SUM** | A CARRY_SUM buffer primitive with the following ports:  
  - **SI** – SUM IN  
  - **SO** – SUM OUT  
  - **CI** – CARRY IN  
  - **CO** – CARRY OUT |
### Table 21–1. Symbols in the Schematic View (Part 2 of 3)

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
</table>
| **LATCH** | A latch primitive with the following ports:  
  - D – data input  
  - ENA – latch enable input  
  - Q – data output  
  - PRE – preset  
  - CLR – clear |
| **DFFE/DFFEAS** | A DFFE (data flipflop with clock enable) primitive, with the same ports as a latch and a clock trigger. The other flipflop primitives are similar:  
  - DFFE (data flipflop with enable and asynchronous load) primitive with additional ALOAD asynchronous load and ADATA data signals  
  - DFFEAS (data flipflop with enable and synchronous and asynchronous load), which has ASDATA as the secondary data port |
| **Atom Primitive** | An atom primitive. The symbol displays the atom name, the port names, and the atom type. The blue shading indicates an atom primitive for which you can view the internal details. For more information, refer to “Viewing Contents of Atom Primitives” on page 21–23. |
| **Other Primitive** | Any primitive that does not fall into the previous categories. Primitives are low-level nodes that cannot be expanded to any lower hierarchy. The symbol displays the port names, the primitive or operator type, and its name.  
  The figure shows an LCELL WYSIWYG primitive, with DATAA to DATAD and COMBOUT port connections. This type of LCELL primitive is found in the Technology Map Viewer for technology-specific atom primitives when the contents of the atom primitive cannot be viewed. The RTL Viewer contains similar primitives if the source design is a VQM or EDIF netlist. |
| **Instance** | An instance in the design that does not correspond to a primitive or operator (a user-defined hierarchy block). The symbol displays the port name and the instance name.  
  For more information about opening the schematic for the lower-level hierarchy, refer to “Traversing and Viewing the Design Hierarchy” on page 21–23. |
| **Encrypted Instance** | A user-defined encrypted instance in the design. The symbol displays the instance name. You cannot open the schematic for the lower-level hierarchy, because the source design is encrypted. |
| **State Machine Instance** | A finite state machine instance in the design. For more information, refer to “State Machine Viewer” on page 21–17. |
Table 21–1. Symbols in the Schematic View (Part 3 of 3)

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>RAM</td>
<td>A synchronous memory instance with registered inputs and optionally registered outputs. The symbol shows the device family and the type of memory block. This figure shows a true dual-port memory block in a Stratix M-RAM block.</td>
</tr>
</tbody>
</table>

Logic Cloud

A combinational logic cloud in the design. For more information, refer to “Grouping Combinational Logic into Logic Clouds” on page 21–25.

Constant

A constant signal value that is highlighted in gray and displayed in hexadecimal format by default throughout the schematic. To change the format, refer to “Changing the Constant Signal Value Formatting” on page 21–26.

Table 21–2 lists and describes the symbol open only in the State Machine Viewer.

Table 21–2. Symbol Available Only in the State Machine Viewer

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>State Node</td>
<td>The node representing a state in a finite state machine. State transitions are indicated with arcs between state nodes. The double circle border indicates the state connects to logic outside the state machine, and a single circle border indicates the state node does not feed outside logic.</td>
</tr>
</tbody>
</table>
Table 21–3 lists and describes the additional higher level operator symbols in the RTL Viewer schematic view.

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
</table>
| ![Adder](image) | An adder operator:  
\[
\text{OUT} = A + B
\]  |
| ![Multiplier](image) | A multiplier operator:  
\[
\text{OUT} = A \times B
\]  |
| ![Divider](image) | A divider operator:  
\[
\text{OUT} = A / B
\]  |
| ![Equals](image) | Equals |
| ![Left Shift](image) | A left shift operator:  
\[
\text{OUT} = (A \ll \text{COUNT})
\]  |
| ![Right Shift](image) | A right shift operator:  
\[
\text{OUT} = (A \gg \text{COUNT})
\]  |
| ![Modulo](image) | A modulo operator:  
\[
\text{OUT} = (A \% B)
\]  |
| ![Less Than](image) | A less than comparator:  
\[
\text{OUT} = (A < B : A > B)
\]  |
Selecting an Item in the Schematic View

To select an item in the schematic view, ensure that the Selection Tool is enabled in the netlist viewer toolbar (this tool is enabled by default). Click an item in the schematic view to highlight it in red.

Select multiple items by pressing the Shift or Ctrl key while selecting with your mouse. You can also select all nodes in a region by selecting a rectangular box area with your mouse cursor when the Selection Tool is enabled. To select nodes in a box, left-click at one corner of the area you want to select and drag the mouse to the diagonally opposite corner. By default, this highlights and selects all nodes in the selected area (instances, primitives, and pins), but not the nets. The Viewer Options dialog box provides an option to select nets. To include nets, right-click in the schematic and click Viewer Options. Under Net Selection, turn on the Select entire net when segment is selected option.

If you enable the Enable auto hierarchy expansion option, items selected in the schematic view are automatically selected in the Netlist Navigator pane (for more information about how to enable the auto hierarchy expansion option, refer to “Netlist Navigator Pane” on page 21–15). The folder then expands automatically if it is required to show the selected entry; however, the folder does not collapse automatically when you are not using or you have deselected the entries.

When you select a hierarchy box, node, or port in the schematic view, the item is highlighted in red but none of the connecting nets are highlighted. When you select a net (wire or bus) in the schematic view, all connected nets are highlighted in red. The selected nets are highlighted across all hierarchy levels and pages. Net selection can be useful when navigating a netlist because you see the net highlighted when you traverse between hierarchy levels or pages.

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
</table>
| ![A multiplexer](image1) | A multiplexer:  
OUT = DATA [SEL]  
The data range size is 2^sel range size |
| ![A selector](image2) | A selector:  
A multiplexer with one-hot select input and more than two input signals |
| ![A binary number decoder](image3) | A binary number decoder:  
OUT = (binary_number (IN) == x)  
for x = 0 to x = 2^(n+1) - 1 |
In some cases, when you select a net that connects to nets in other levels of the hierarchy, these connected nets are also highlighted in the current hierarchy. If you prefer that these nets not be highlighted, use the Viewer Options dialog box option to highlight a net only if the net is in the current hierarchy. Right-click in the schematic and click Viewer Options. In the Net Selection section, turn on the Limit selections to current hierarchy option.

**Moving and Panning in the Schematic View**

When the schematic view page is larger than the portion currently displayed, you can use the scroll bars at the bottom and right side of the schematic view to see other areas of the page.

You can also use the Hand Tool to “grab” the schematic page and drag it in any direction. Enable the Hand Tool with the toolbar button. Click and drag to move around the schematic view without using the scroll bars.

In addition to the scroll bars and Hand Tool, you can use the middle-mouse or wheel button to move and pan in the schematic view. Click the middle-mouse or wheel button once to enable the feature. Move the mouse or scroll the wheel to move around the schematic view. Click the middle-mouse or wheel button again to turn the feature off.

**Netlist Navigator Pane**

The Netlist Navigator pane displays the entire netlist in a tree format based on the hierarchical levels of the design. In each level, similar elements are grouped into subcategories. You can use the Netlist Navigator pane to traverse through the design hierarchy to view the logic schematic for each level. You can also select an element in the netlist navigator to highlight in the schematic view.

Nodes inside atom primitives are not listed in the Netlist Navigator pane.

For each module in the design hierarchy, the Netlist Navigator pane displays the applicable elements listed in Table 21–4.

<table>
<thead>
<tr>
<th>Elements</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Instances</td>
<td>Modules or instances in the design that can be expanded to lower hierarchy levels.</td>
</tr>
<tr>
<td>State Machines</td>
<td>State machine instances in the design that can be viewed in the State Machine Viewer.</td>
</tr>
<tr>
<td>Primitives</td>
<td>Low-level nodes that cannot be expanded to any lower hierarchy level. These primitives include:</td>
</tr>
<tr>
<td></td>
<td>■ Registers and gates that you can view in the RTL Viewer when using Quartus II integrated synthesis.</td>
</tr>
<tr>
<td></td>
<td>■ Logic cell atoms in the Technology Map Viewer or in the RTL Viewer when using a VQM or EDIF from third-party synthesis software.</td>
</tr>
<tr>
<td></td>
<td>In the Technology Map Viewer, you can view the internal implementation of certain atom primitives, but you cannot traverse into a lower-level of hierarchy.</td>
</tr>
</tbody>
</table>
Properties

The Properties pane is a dockable pane that displays only node type components such as pins, primitives, and instances. This pane consists of two viewers; the connectivity view and table view. The table view lists the fan-in, fan-out, parameters, and ports tables. The connectivity view displays the selected block connectivity information listed in the table view, such as fan-in or fan-out connection, schematic, truth table, and Karnaugh map.

To open the Properties Pane, follow these steps:

1. In the Quartus II software, on the Tools menu, point to Netlist Viewers and click RTL Viewer.
2. In the RTL Viewer window, on the View Toolbar, click the Property icon.

Table 21–4. Netlist Navigator Pane Elements (Part 2 of 2)

<table>
<thead>
<tr>
<th>Elements</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pins</td>
<td>The I/O ports in the current level of hierarchy.</td>
</tr>
<tr>
<td></td>
<td>■ Pins are device I/O pins when viewing the top hierarchy level and are I/O ports of the design when viewing the lower-levels.</td>
</tr>
<tr>
<td></td>
<td>■ When a pin represents a bus or an array of pins, expand the pin entry in the list view to see individual pin names.</td>
</tr>
<tr>
<td>Nets</td>
<td>Nets or wires connecting the nodes. When a net represents a bus or array of nets, expand the net entry in the tree to see individual net names.</td>
</tr>
<tr>
<td>Logic Clouds</td>
<td>A group of related combinational logics of a particular source. You can automatically or manually group combinational logics or ungroup logic clouds in your design.</td>
</tr>
</tbody>
</table>
Figure 21–5 shows the Properties pane.

**State Machine Viewer**

The State Machine Viewer displays a graphical representation of the state machines in your design. You can open the State Machine Viewer in any of the following ways:

- On the Tools menu, point to Netlist Viewers and click **State Machine Viewer**.
- Double-click a state machine instance in the RTL Viewer
- Right-click a state machine instance in the RTL Viewer and click **Hierarchy Down**.
- Select a state machine instance in the RTL Viewer, and on the Project menu, point to **Hierarchy** and click **Down**.
Figure 21–6 shows an example of the State Machine Viewer for a simple state machine.

**Figure 21–6. The State Machine Viewer**

State Diagram View

The state diagram view appears at the top of the State Machine Viewer. It contains a diagram of the states and state transitions.

The nodes that represent each state are arranged horizontally in the state diagram view with the initial state (the state node that receives the reset signal) in the left-most position. Nodes that connect to logic outside of the state machine instance are represented by a double circle. The state transition is represented by an arc with an arrow pointing in the direction of the transition.

When you select a node in the state diagram view, and turn on the Highlight Fan-in or Highlight Fan-out command from the View menu or the State Machine Viewer toolbar, the respective fan-in or fan-out transitions from the node are highlighted in red.

An encrypted block with a state machine displays encoding information in the state encoding table, but does not display a state transition diagram or table.
State Transition Table
The state transition table on the Transitions tab at the bottom of the State Machine Viewer displays the condition equation for each state transition. Each row in the table represents a transition (each arc in the state diagram view). The table has the following columns:
- **Source State**—the name of the source state for the transition
- **Destination State**—the name of the destination state for the transition
- **Condition**—the condition equation that causes the transition from source state to destination state

To see all of the transitions to and from each state name, click the appropriate column heading to sort on that column.

The text in each column is left-aligned by default; to change the alignment and to make it easier to see the relevant part of the text, right-click the column and click *Align Right*. To revert to left alignment, click *Align Left*.

Click in any cell in the table to select it. To select all cells, right-click in the cell and click Select All; or, on the Edit menu, click Select All. To copy selected cells to the clipboard, right-click the cells and click Copy Table; or, on the Edit menu, point to Copy and click Copy Table. You can paste the table into any text editor as tab-separated columns.

State Encoding Table
The state encoding table on the Encoding tab at the bottom of the State Machine Viewer displays encoding information for each state transition.

To view state encoding information in the State Machine Viewer, you must synthesize your design with the Start Analysis & Synthesis command. If you have only elaborated your design with the Start Analysis & Elaboration command, the encoding information is not displayed.

Selecting an Item in the State Machine Viewer
You can select and highlight each state node and transition in the State Machine Viewer. To select a state transition, click the arc that represents the transition.

When you select a node or transition arc in the state diagram view, the matching state node or equation conditions in the state transition table are highlighted; conversely, when you select a state node or equation condition in the state transition table, the corresponding state node or transition arc is highlighted in the state diagram view.

Switching Between State Machines
A design may contain multiple state machines. To choose which state machine to view, use the State Machine selection box located at the top of the State Machine Viewer. Click in the drop-down box and select the necessary state machine.

Options
The Options dialog box allows you to customize your settings, such as display settings, colors, fonts, tracing, customize view, and shortcut commands. To open the Options dialog box, in the RTL Viewer window, on the Tools menu, click Options.
Netlist Viewers

If you want to customize the display settings for your preferred viewing, you can direct the RTL Viewer and Technology Map Viewer to adjust the settings in the schematic view. To adjust the display settings, on the Tools menu, click Options. In the Netlist Viewers category, under Display Settings, you can select the options to customize your display.

Display Settings

You can divide the schematic representation of a large design into multiple pages. Dividing the display of the schematic into multiple pages does not affect your design, and only controls the number of elements per page. Under Display Settings, the Nodes per page option allows you to specify the number of nodes displayed per page. The default value is 500 nodes; however, you can view from one to 1,000 nodes per page. The Ports per page option allows you to specify the number of ports (or pins) displayed per page. The default value is 1,000 ports (or pins); the range is 1 to 2,000 ports (or pins). The netlist viewers partition your design into a new page if either the number of nodes or the number of ports exceeds the limit you specified. Occasionally, the number of ports displayed on the page might exceed the limit you specified, depending on the configuration of nodes on the page. If you turned on the Display boundary around hierarchy levels option and the total number of nodes or ports in the hierarchy exceeds the value you specified for the Nodes per page or Ports per page options, the netlist viewers display the boundary as a hierarchy port connector (refer to Table 21–1 on page 21–10).

To display net names in your schematic, turn on the Show Net Name option. If you turn on this option, the schematic view refreshes automatically to display the net names. To show node names in the schematic view, turn on the Show node name option. To change the value formatting, select the necessary format in the Constant signal format list.

To view highlighting around the design element in range of the mouse pointer, turn on the Enable rollover option. To more easily navigate through your design hierarchy, you can direct the netlist viewers to automatically expand the hierarchy list in Netlist Navigator pane and highlight the design element you selected in the schematic view. This automatic expansion and selection of corresponding design elements is useful when you have a complex schematic that spans multiple display pages in the netlist viewers. To use the automatic expansion and selection feature, turn on the Enable auto hierarchy expansion option.

To trace signal path in the schematic using the Global Net Routing feature, turn on the Enable global net routing option.

Radial Menu Settings

To enable the radial menu, turn on the Enable radial menu option. The radial menu is an octagonal menu with eight commands from which you can choose. This menu provides a quick way to perform any of the commands with a single click whenever you are in the schematic view.

To open the radial menu, right-click and hold anywhere in the schematic view and wait for the menu to appear. By default, the menu appears after 0.2 seconds. The radial menu appears with the mouse pointer always at the center point. The center point of the menu is a non-trigger boundary in which no command is started.
To run the necessary command, hold down the right mouse button, drag the mouse onto the command, and then press the left mouse button. If you decide not to trigger any command after the radial menu appears, press the Esc key or drag the pointer back into the center point and release the mouse button.

To change the delay time before the radial menu appears, select the necessary interval time in the drop-down list for Delay showing radial menu for. The default delay is 0.2 seconds.

**Automatic Texts Hiding**

This option allows you to automatically hide texts such as node names, port names, and net names, depending on your preferred zoom level. By default, when your schematic is zoomed 40% smaller, the text is hidden.

**Colors**

This option allows you to determine your preferred colors to represent the elements in the schematics.

**Fonts**

This option allows you to determine the type of text, font size, color and style for the node and net names in the schematics.

**Tracing**

If you want to filter information from the schematic view of your design to isolate specific design elements for further inspection, or if you want to expand specific design elements, you can direct the RTL Viewer and the Technology Map Viewer to adjust the elements the viewers show in the schematic view. To adjust the filtering and expansion settings in the RTL Viewer and Technology Map Viewer, on the Tools menu, click **Options**. Under **Tracing**, you can select the options to control filtering and expansion settings.

For all filtering commands, the netlist viewers stop tracing through the netlist when they reach one of the following objects:

- A pin
- A specified number of filtering levels, counting from the selected node or port
- A register

To specify the number of filtering levels, set the **Number of filtering levels** option to specify the number of levels to expand. You can specify a value from one to 100.

To enable the **Stop filtering at register** option, turn on the **Stop filtering at register** option. You can filter across hierarchies when you turn on the **Filter across hierarchy** option.

By default, the filtered schematic shows all possible connections between the nodes shown in the schematic. To remove the connections that are not directly part of the path that was traced to generate a filtered netlist, turn off the **Show all connections between nodes** option.
To set the amount of logic you want to expand, set the **Number of expansion levels** option to specify the number of levels to expand. You can specify a range from one to 100 levels. You can also set the **Stop expanding at register** option to specify whether netlist expansion should stop when a register is reached.

**Customize View**

If you want to customize the schematic display for better viewing and to speed up your debugging process, you can direct the RTL Viewer and the Technology Map Viewer to remove fan-out free nodes, show simplify logic, group or ungroup related nodes, and group combinational logic into a logic cloud. To adjust the options that control the schematic display in the RTL Viewer and the Technology Map Viewer, on the Tools menu, click **Options**. Under **Customize View**, you can select the options to customize your view. These options are also available in the **Customize View** tab of the **RTL/Technology Map Viewer Options** dialog box. To open the dialog box, right-click in the schematic and click **Viewer Options**.

When you change settings, the list of previously viewed pages is cleared. The settings are revision-specific, so different revisions can have different settings.

To remove fan-out free registers from your schematic display, turn on the **Remove registers without fan-out** option. To remove all single-input nodes and merge a chain of equivalent combinational gates that have direct connections (without inversion in between) into a single multiple-input gate, turn on the **Show simplified logic** option.

To group all related nodes into a single node, turn on the **Group all related nodes** option. You can manually group or ungroup any nodes by right-clicking the selected nodes in the schematic and selecting **Group Related Nodes** or **Ungroup Selected Nodes**. To group combinational logic into logic clouds, turn on the **Group combinational logic into logic cloud** option.

Customize Logic and Customize Group options are available for the RTL Viewer, whereas only the Customize Group option is available for the Technology Map Viewer.

**Shortcut Commands**

You can choose eight commands to appear on the radial menu from a list of 18 available commands. To customize the command list on the menu, first open the RTL Viewer or the Technology Map Viewer. On the Tools menu, select **Options**. On the **Shortcut Commands** category list, drag and drop the icon under **Shortcut buttons** into any region under **Shortcut commands popup**. You can click the icon under **Shortcut buttons** to see its description.

**Navigating the Schematic View**

This section describes methods to navigate through the pages and hierarchy levels in the schematic view of the RTL Viewer and the Technology Map Viewer.
Traversing and Viewing the Design Hierarchy

You can open different hierarchy levels in the schematic view from the Netlist Navigator pane (refer to “Netlist Navigator Pane” on page 21–15), or the Hierarchy Up and Hierarchy Down commands (Shortcut menu) in the schematic view.

Use the Hierarchy Down command to go down in an instance’s hierarchy, and open a lower-level schematic showing the internal logic of the instance. Use the Hierarchy Up command to go up in hierarchy or collapse a lower-level hierarchy, and open the parent higher level hierarchy. When the Selection Tool is selected, the appropriate option is available when your mouse pointer is located over an area of the schematic view that has a corresponding lower or higher level hierarchy.

The mouse pointer changes as it moves over different areas of the schematic to indicate whether you can move up, down, or both up and down in the hierarchy (Figure 21–7). To open the next hierarchy level, right-click in that area of the schematic and click Hierarchy Down or Hierarchy Up, as appropriate, or double-click in that area of the schematic.

![Figure 21–7. Mouse Pointers Indicate How to Traverse Hierarchy](image)

Flattening the Design Hierarchy

You can flatten the design hierarchy to view the design without hierarchical boundaries. To flatten the hierarchy from the current level and all lower-level hierarchies of the current design hierarchy, right-click in the schematic and click Flatten Netlist. To flatten the entire design, choose Flatten Netlist from the top-level schematic of the design.

Viewing the Contents of a Design Hierarchy in the Current Schematic

You can use the Display Content and Hide Content (Shortcut menu) commands to show or hide a lower hierarchy level for a specific instance in the schematic for the current hierarchy level.

To display the lower hierarchy netlist of an instance on the same schematic as the remaining logic in the currently viewed netlist, right-click the selected instance and click Display Content.

To hide all of the lower hierarchy logic of a hierarchy box into a closed instance, right-click the selected instance and click Hide Content.

Viewing Contents of Atom Primitives

In the Technology Map Viewer, you can view the contents of certain device atom primitives to see their underlying implementation details. For logic cell (LCELL) atoms in Arria® GX, Cyclone® series, MAX® II, and Stratix series of devices, you can view LUTs, registers, and logic gates. For I/O atoms in the Arria GX, Cyclone series, HardCopy® IV, and Stratix series of devices, you can view registers and logic gates.
In addition, you can view the implementation of RAM and DSP blocks in certain devices in the RTL Viewer or Technology Map Viewer. You can view the implementation of RAM blocks in the Arria GX, Cyclone series, and Stratix series of devices. You can view the implementation of DSP blocks only in Arria GX and Stratix series of devices.

If you can view the contents of an atom instance, the internal contents are shown in blue in the schematic view (Figure 21–8).

**Figure 21–8. Instance That Can Be Expanded to View Internal Contents**

![Instance That Can Be Expanded to View Internal Contents](image)

To view the contents of one or more atom primitive instances, select the necessary atom instances. Right-click a selected instance and click **Display Content**. You can also double-click the necessary atom instance to view the contents. Figure 21–9 shows an expanded version of the instance in Figure 21–8.

**Figure 21–9. Internal Contents of the Atom Instance in Figure 21–8.**

![Internal Contents of the Atom Instance](image)

To hide the contents (and revert to the compact format), select and right-click the atom instance or instances, and click **Hide Content**.

**Note:** In the schematic view, the internal details in an atom instance cannot be selected as individual nodes. Any mouse action on any of the internal details is treated as a mouse action on the atom instance.

### Viewing the Properties of Instances and Primitives

You can view the properties of an instance or primitive using the **Properties** dialog box. To view the properties of an instance or primitive in the RTL Viewer or Technology Map Viewer, right-click the node and click **Properties**.
The Properties dialog box contains the following information about the selected node:

- The parameter values of an instance.
- The active level of the port (for example, active high or active low). An active low port is denoted with an exclamation mark “!”.
- The port’s constant value (for example, $V_{CC}$ or GND). Table 21–5 describes the possible value of a port.

### Table 21–5. Possible Port Values

<table>
<thead>
<tr>
<th>Value</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>$V_{CC}$</td>
<td>The port is not connected and has $V_{CC}$ value (tied to $V_{CC}$)</td>
</tr>
<tr>
<td>GND</td>
<td>The port is not connected and has GND value (tied to GND)</td>
</tr>
<tr>
<td>--</td>
<td>The port is connected and has value (other than $V_{CC}$ or GND)</td>
</tr>
<tr>
<td>Unconnected</td>
<td>The port is not connected and has no value (hanging)</td>
</tr>
</tbody>
</table>

In the LUT of a logic cell (LCELL), the Properties dialog box contains the following additional information:

- The schematic of the LCELL
- The Truth Table representation of the LCELL
- The Karnaugh map representation of the LCELL

### Viewing LUT Representations in the Technology Map Viewer

You can view different representations of a LUT by right-clicking the selected LUT and clicking Properties. This feature is supported for the Arria GX, Cyclone series, MAX II, and Stratix series of devices only. You can view the LUT representations in the following three tabs in the Properties dialog box:

- The Schematic tab—the equivalent gate representations of the LUT.
- The Truth Table tab—the truth table representations.
- The Karnaugh Map tab—the Karnaugh map representations of the LUT. The Karnaugh map supports up to 6 input LUTs.

For more information about the Ports tab, refer to “Viewing the Properties of Instances and Primitives” on page 21–24.

### Grouping Combinational Logic into Logic Clouds

The following sections describe how to group combinational logic into logic clouds.

For the definition of a logic cloud, refer to Table 21–1 on page 21–10.

### Logic Clouds in the RTL Viewer

You can automatically group all combinational logic nodes in your design into logic clouds. For more information about how to group combinational logic clouds, refer to “Customize View” on page 21–22.
Logic Clouds in the Technology Map Viewer

In the Technology Map Viewer, the **Group combinational logic into logic clouds** option is supported for Cyclone II, HardCopy, and Stratix II devices only. For more information about how to group combinational logic clouds, refer to “Customize View” on page 21–22.

Grouping and Ungrouping Logic Clouds

To group logic nodes into a logic cloud manually, right-click the selected node or input port and click **Group source logic into logic cloud**. To ungroup a logic cloud manually, right-click the selected logic cloud and click **Ungroup source logic from logic cloud**. You can also ungroup a logic cloud manually by double-clicking the selected logic cloud. These options are not available if the nodes cannot be grouped.

Changing the Constant Signal Value Formatting

The constant signal value is highlighted in gray in the schematic view. By default, the value is displayed in hexadecimal format, but you can also choose binary or decimal format. For more information about changing the value formatting, refer to “Netlist Viewers” on page 21–20.

Changing the format affects all constant signal values throughout the schematic. For descriptions of constant signal values in the schematic, refer to Table 21–3 on page 21–13.

Zooming and Magnification

You can control the magnification of your schematic on the View menu, with the Zoom Tool in the toolbar, or the Ctrl key and mouse wheel button, as described in this section.

By default, the netlist viewer displays most pages sized to fit in the window. If the schematic page is very large, the schematic is displayed at the minimum zoom level, and the view is centered on the first node. Click **Zoom In** to view the image at a larger size, and click **Zoom Out** to view the image (when the entire image is not displayed) at a smaller size. The **Zoom** command allows you to specify a magnification percentage (100% is considered the normal size for the schematic symbols).

The **Fit Selection in Window** command zooms in on the selected nodes in a schematic to fit in the window. Use the Selection Tool to select one or more nodes (instances, primitives, pins, and nets), then click **Fit Selection in Window** to enlarge the area covered by the selection. This feature is helpful when you want to see a particular element in a large schematic. After you select a node, you can easily zoom in to view the particular node. You can temporarily enlarge a portion of the schematic with the magnifying glass tool in the toolbar.

You can also use the Zoom Tool on the netlist viewer toolbar to control magnification in the schematic view. When you select the Zoom Tool in the toolbar, clicking in the schematic zooms in and centers the view on the location you clicked. Right-click in the schematic to zoom out and center the view on the location you clicked. When you select the Zoom Tool, you can also zoom into a certain portion of the schematic by selecting a rectangular box area with your mouse cursor. The schematic is enlarged to show the selected area.
Alternatively, you can specify the magnification percentage by right-clicking the necessary area and dragging the mouse to the right to zoom in or to the left to zoom out with the Zoom Tool. You will see a red line with the zoom percentage above it. The zoom percentage is proportional to the length of the green line. Release the mouse button at the necessary zoom percentage.

By default, the netlist viewers maintain the zoom level when filtering on the schematic (refer to “Filtering in the Schematic View” on page 21–30). To change the behavior so that the zoom level is always reset to “Fit in Window,” on the Tools menu, click Options. In the Category list, select Netlist Viewers, and turn off Maintain zoom level.

**Schematic Debugging and Tracing Using the Bird’s Eye View**

Viewing the entire schematic can be useful when debugging and tracing through a large netlist. The Quartus II software allows you to quickly navigate to a specific section of the schematic using the Bird’s Eye View feature, which is available in the RTL Viewer and Technology Map Viewer.

The Bird’s Eye View shows the current area of interest. Select the necessary area by clicking and dragging the indicator or right-clicking to form a rectangular box around the necessary area. You can also click and drag the rectangular box to move around the schematic. To open the Bird’s Eye View, on the View menu, click **Bird’s Eye View**, or click on the **Bird’s Eye View** icon in the toolbar.

**Figure 21–10** shows the Bird’s Eye View window in the schematic diagram.

**Figure 21–10. Bird’s Eye View Window**
Partitioning the Schematic into Pages

For large design hierarchies, the RTL Viewer and Technology Map Viewer partition your netlist into multiple pages in the schematic view. For more information about controlling how much of the design is visible on each page, refer to “Netlist Viewers” on page 21–20.

When a hierarchy level is partitioned into multiple pages, the title bar for the schematic window indicates which page is displayed and how many total pages exist for this level of hierarchy (shown in the format: Page <current page number> of <total number of pages>), as shown in Figure 21–11.

When you change the number of nodes or ports per page, the change applies only to new pages that are shown or opened in the netlist viewer. To refresh the current page so that it displays the changed number of nodes or ports, click the Refresh button on the toolbar.

Moving Between Schematic Pages

To move to another schematic page, on the View menu, click Previous Page or Next Page, or click the Previous Page icon or the Next Page icon on the netlist viewer toolbar.

To go to a particular page of the schematic, on the Edit menu, click Go To, or right-click in the schematic view and click Go To. In the Page list, select the necessary page number. You can also go to a particular page by selecting the necessary page number from the pull-down list on the top right of the netlist viewer.

Moving Back and Forward Through Schematic Pages

To return to the previous view after changing the page view, click Back on the View menu, or click the Back icon on the netlist viewer toolbar. To go to the next view, click Forward on the View menu, or click the Forward icon on the netlist viewer toolbar.

You can go forward only if you have not made any changes to the view since going back. Use the Back and Forward commands to switch between page views. These commands do not undo an action, such as selecting a node.
Following Nets Across Schematic Pages

Input and output connectors indicate nodes that connect across pages of the same hierarchy. Right-click a connector to display a menu of commands that trace the net through the pages of the hierarchy.

After you right-click to follow a connector port, the netlist viewer opens a new page, which centers the view on the particular source or destination net using the same zoom factor as the previous page. To trace a specific net to the new page of the hierarchy, Altera recommends that you first select the necessary net, which highlights it in red, before you right-click to traverse pages.

**Input Connectors**

Figure 21–12 shows an example of the menu that appears when you right-click an input connector. The *From* command opens the page containing the source of the signal. The *Related* commands, if applicable, open the specified page containing another connection fed by the same source.

**Figure 21–12. Input Connector Shortcut Menu**

![Input Connector Shortcut Menu](image)

**Output Connectors**

Figure 21–13 shows an example of the menu that appears when you right-click an output connector. The *To* command opens the specified page that contains a destination of the signal.

**Figure 21–13. Output Connector Shortcut Menu**

![Output Connector Shortcut Menu](image)
Go to Net Driver

To locate the source of a particular net in the schematic view, right-click the net, point to Go to Net Driver and click Current page, Current hierarchy, or Across hierarchies. Table 21–6 lists the Go to Net Driver commands.

<table>
<thead>
<tr>
<th>Command</th>
<th>Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>Current page</td>
<td>Locates the source or driver on the current page of the schematic only.</td>
</tr>
<tr>
<td>Current hierarchy</td>
<td>Locates the source in the current level of hierarchy, even if the source is located on another page of the netlist schematic.</td>
</tr>
<tr>
<td>Across hierarchies</td>
<td>Locates the source across hierarchies until the software reaches the source at the top hierarchy level.</td>
</tr>
</tbody>
</table>

The schematic view opens the correct page of the schematic, if required, and adjusts the centering of the page so that you can see the net source. The schematic shows the default page for the net driver. The view is unfiltered, so no filtering results are kept.

Filtering in the Schematic View

Filtering allows you to filter out nodes and nets in your netlist to view only the logic elements of interest to you.

You can filter your netlist by selecting hierarchy boxes, nodes, ports of a node, nets, or states in a state machine that are part of the path you want to see. The following filter commands are available:

- **Sources**—displays the sources of the selection
- **Destinations**—displays the destinations of the selection
- **Sources & Destinations**—displays the sources and destinations of the selection
- **Selected Nodes and Nets**—displays only the selected nodes and nets with the connections between them
- **Between Selected Nodes**—displays nodes and connections in the path between the selected nodes
- **Bus Index**—displays the sources or destinations for one or more indices of an output or input bus port

To filter your netlist, select a hierarchy box, node, port, net, or state node, right-click in the window, point to Filter and click the appropriate filter command. The netlist viewer generates a new page showing the netlist that remains after filtering.

When filtering in a state diagram in the State Machine Viewer, sources and destinations refer to the previous and next transition states or paths between transition states in the state diagram. The transition table and encoding table also reflect the filtering.

You can go back to the netlist page before it was filtered using the Back command, as described in “Moving Back and Forward Through Schematic Pages” on page 21–28.
When viewing a filtered netlist, clicking an item in the Netlist Navigator pane causes the schematic view to display an unfiltered view of the appropriate hierarchy level. You cannot use the Netlist Navigator pane to select items or navigate in a filtered netlist.

**Filter Sources Command**

To filter out all but the source of the selected item, right-click the item, point to Filter, and click Sources. The selected object type determines what is displayed, as listed in Table 21–7 and shown in Figure 21–14.

<table>
<thead>
<tr>
<th>Selected Object</th>
<th>Result Shown in Filtered Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>Node or hierarchy box</td>
<td>Shows all the sources of the node’s input ports. For an example, refer to Figure 21–14.</td>
</tr>
<tr>
<td>Net</td>
<td>Shows the sources that feed the net.</td>
</tr>
<tr>
<td>Input port of a node</td>
<td>Shows only the input source nodes that feed this port.</td>
</tr>
<tr>
<td>Output port of a node</td>
<td>Shows only the selected node.</td>
</tr>
<tr>
<td>State node in a state machine</td>
<td>Shows the states that feed the selected state (previous transition states).</td>
</tr>
</tbody>
</table>

**Filter Destinations Command**

To filter out all but the destinations of the selected node or port as outlined in Table 21–8 and shown in Figure 21–14, right-click the node or port, point to Filter, and click Destinations.

<table>
<thead>
<tr>
<th>Selected Object</th>
<th>Result Shown in Filtered Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>Node or hierarchy box</td>
<td>Shows all the destinations of the node’s output ports. For an example, refer to Figure 21–14.</td>
</tr>
<tr>
<td>Net</td>
<td>Shows the destinations fed by the net.</td>
</tr>
<tr>
<td>Input port of a node</td>
<td>Shows only the selected node.</td>
</tr>
<tr>
<td>Output port of a node</td>
<td>Shows only the fan-out destination nodes fed by this port.</td>
</tr>
<tr>
<td>State node in a state machine</td>
<td>Shows the states that are fed by the selected states (next transition states).</td>
</tr>
</tbody>
</table>
Filter Sources and Destinations Command

The Sources & Destinations command is a combination of the Sources and Destinations filtering commands, in which the filtered page shows the sources and the destinations of the selected item. To select this option, right-click the necessary object, point to Filter, and click Sources & Destinations. Refer to the example in Figure 21–14.

Figure 21–14. Sources, Destinations, and Sources and Destinations Filtering for inst4

Filter Between Selected Nodes Command

To show the nodes in the path between two or more selected nodes or hierarchy boxes, right-click the necessary object, point to Filter, and click Between Selected Nodes. For this option, selecting a port of a node is the same as selecting the node.

Filter Selected Nodes and Nets Command

To create a filtered page that shows only the selected nodes, nets, or both, and, if applicable, the connections between the selected nodes, nets, or both, right-click the necessary object, point to Filter, and click Selected Nodes & Nets.
Figure 21–15 shows a schematic with several nodes selected.

Figure 21–15. Using Selected Nodes and Nets to Select Nodes

Figure 21–16 shows the schematic after filtering. If you select a net, the filtered page shows the immediate sources and destinations of the selected net.

Figure 21–16. Selected Nodes and Nets Filtering on Figure 21–15 Schematic
Filter Bus Index Command

To show the path related to a specific index of a bus input or output port in the RTL Viewer, right-click the port, point to Filter, and click Bus Index. The Select Bus Index dialog box allows you to select the indices of interest.

Filter Command Processing

The options to control filtering are available in the Tracing tab of the RTL/Technology Map Viewer Options dialog box. Right-click in the schematic view and click Viewer Options to open the dialog box. For more information about filter command processing, refer to “Tracing” on page 21–21.

Filtering Across Hierarchies

The filtering commands display nodes in all hierarchies by default. When the filtered path passes through levels of hierarchy on the same schematic page, green hierarchy boxes group the logic and show the hierarchy boundaries. A green rectangular symbol that appears on the border represents the port relationship between two different hierarchies (Figure 21–17 and Figure 21–18).

For more information about how to control filtering, refer to “Tracing” on page 21–21.

For more information about how to disable the box hierarchy display, refer to “Netlist Viewers” on page 21–20.

Netlists of the same hierarchy displayed over more than one page are not grouped with a box. Filtering and expanding on a blue atom primitive does not trace the underlying netlist, even when Filter across hierarchy is enabled.

Figure 21–17 and Figure 21–18 show examples of filtering across hierarchical boundaries. Figure 21–17 shows a smaller example of an input port of the taps instance after the Sources filter is applied, in which the input port of the lower-level hierarchical block connects directly to an input pin of the design. The name of the instance appears in the green border and as a tooltip when you move your mouse pointer over the instance.

Figure 21–17. Filtering Across Hierarchical Boundaries
Figure 21–18 shows a larger example of an input port of an instance after the **Sources** filter is applied, in which the source comes from input pins that are fed through another level of hierarchy.

**Expanding a Filtered Netlist**

After a netlist is filtered, some ports might not have connections displayed because their connections are not part of the main path through the netlist. Two expansion features, immediate expansion and the **Expand** command, allow you to add the fan-in or fan-out signals of these ports to the schematic display of a filtered netlist.

You can immediately expand any port whose connections are not displayed. When you double-click that port in the filtered schematic, one level of logic is expanded.

To expand more than one level of logic, right-click the port and click the **Expand** command. This command expands logic from the selected port by the amount specified in **Viewer Options**. To set these options, right-click in the schematic view and click **Viewer Options**. In the **Expansion** section, set the **Number of expansion levels** option to specify the number of levels to expand (the default value is 3 and the range is 1 to 15 levels). You can also set the **Stop expanding at register** option (which is turned on by default) to specify whether netlist expansion should stop when a register is reached. These options are also available under **Tracing** in the **Options** dialog box (refer to “Tracing” on page 21–21).

You can select multiple nodes to expand when you use the **Expand** command. If you select ports that are located on multiple schematic pages, only the ports on the currently viewed page appear in the expanded schematic.

In the State Machine Viewer, the **Expand** command has the following three options:

- **Sources**—Displays the states that feed the selected states (previous transition states)
- **Destinations**—Displays the states that are fed by the selected states (next transition states)
- **Sources & Destinations**—Displays the previous and next transition states

The state transition table and state encoding table also reflect the changes to the filter.
The expansion feature works across hierarchical boundaries if the filtered page containing the port you want to expand was generated with the Filter across hierarchy option turned on (for details about this option, refer to “Filtering in the Schematic View” on page 21–30). When viewing timing paths in the Technology Map Viewer, the Expand command always works across hierarchical boundaries because filtering across hierarchy is always turned on for these schematics (for details about these schematics, refer to “Viewing a Timing Path” on page 21–38).

Reducing a Filtered Netlist

In some cases, removing logic from a filtered schematic or state diagram makes the schematic view easier to read and minimizes distracting logic in the schematic that you do not need to view.

To reduce elements in the filtered schematic or state diagram view, right-click the node or nodes you want to remove and click Reduce.

Probing to a Source Design File and Other Quartus II Windows

The RTL Viewer, Technology Map Viewer, and State Machine Viewer allow you to cross-probe to the source design file and to various other windows in the Quartus II software. You can select one or more hierarchy boxes, nodes, nets, state nodes, or state transition arcs that interest you in the netlist viewer and locate the corresponding items in another applicable Quartus II software window. You can then view and make changes or assignments in the appropriate editor or floorplan.

To locate an item from the netlist viewer in another window, right-click the items of interest in the schematic or state diagram, point to Locate, and click the appropriate command. The following commands are available:

- Locate in Assignment Editor
- Locate in Pin Planner
- Locate in Chip Planner
- Locate in Resource Property Editor
- Locate in Technology Map Viewer
- Locate in RTL Viewer
- Locate in Design File

The options available for locating an item depend on the type of node and whether it exists after placement and routing. If a command is enabled in the menu, it is available for the selected node. You can use the Locate in Assignment Editor command for all nodes, but assignments might be ignored during placement and routing if they are applied to nodes that do not exist after synthesis.

The netlist viewer automatically opens another window for the appropriate editor or floorplan and highlights the selected node or net in the newly opened window. You can switch back to the netlist viewer by selecting it in the Window menu or by closing, minimizing, or moving the new window.

When probing to a logic cloud in the RTL Viewer, a message box appears, prompting you to ungroup the logic cloud or allow it to remain grouped.
Moving Selected Nodes to Other Quartus II Windows

You can drag selected nodes from the netlist viewers to the Text Editor, Block Editor, Pin Planner, SignalTap® II Embedded Logic Analyzer, and Waveform Editor windows in the Quartus II software. Whenever you see the drag-and-drop pointer on the selected node in the netlist viewers, it means that the node can be dragged to other child windows in the Quartus II software.

To tap a node from the schematic in the Technology Map Viewer to an open SignalTap II Embedded Logic Analyzer window or to a new SignalTap II file (.stp), right-click the selected node in the schematic diagram or in the netlist navigator, and then click Add Node to SignalTap II Logic Analyzer. If the node cannot be tapped, the option is unavailable.

Probing to the Netlist Viewers from Other Quartus II Windows

You can cross-probe to the RTL Viewer and Technology Map Viewer from other windows in the Quartus II software. You can select one or more nodes or nets in another window and locate them in one of the netlist viewers.

You can locate nodes between the RTL Viewer, State Machine Viewer, and Technology Map Viewer, and you can locate nodes in the RTL Viewer and Technology Map Viewer from the following Quartus II software windows:

- Project Navigator
- Timing Closure Floorplan
- Chip Planner
- Resource Property Editor
- Node Finder
- Assignment Editor
- Messages Window
- Compilation Report
- TimeQuest Timing Analyzer (supports the Technology Map Viewer only)

To locate elements in the netlist viewer from another Quartus II window, select the node or nodes in the appropriate window; for example, select an entity in the Entity list on the Hierarchy tab in the Project Navigator, or select nodes in the Timing Closure Floorplan, or select node names in the From or To column in the Assignment Editor. Next, right-click the selected object, point to Locate, and click Locate in RTL Viewer or Locate in Technology Map Viewer. After you click this command, the netlist viewer opens, or is brought to the foreground if the netlist viewer is open.

The first time the window opens after a compilation, the preprocessor stage runs before the netlist viewer opens.

The netlist viewer shows the selected nodes and, if applicable, the connections between the nodes. The display is similar to what you see if you right-click the object, point to Filter, and click Selected Nodes & Nets using Filter Across Hierarchy. If the nodes cannot be found in the netlist viewer, a message box displays the message: Can’t find requested location.
Viewing a Timing Path

To see a visual representation of a timing path, cross-probe from a report panel in the TimeQuest analyzer.

To take advantage of this feature, you must complete a full compilation of your design, including the timing analyzer stage. To see the timing results for your design, on the Processing menu, click **Compilation Report**. On the left side of the Compilation Report, select **Timing Analyzer** or **TimeQuest Timing Analyzer**. When you select a detailed report, the timing information is listed in a table format on the right side of the Compilation Report; each row of the table represents a timing path in the design. You can also view timing paths in TimeQuest analyzer report panels. To view a particular timing path in the Technology Map Viewer or RTL Viewer, right-click the appropriate row in the table, point to **Locate**, and click **Locate in Technology Map Viewer** or **Locate in RTL Viewer**.

To locate a path, on the **Tasks** pane, in the **Custom Reports** folder, double-click **Report Timing**. In the **Report Timing** dialog box, make necessary settings, and then click the **Report Timing** button. After the TimeQuest analyzer generates the report, right-click on the node in the table and select **Locate Path**. In the Technology Map Viewer, the schematic page displays the nodes along the timing path with a summary of the total delay.

When you locate the timing path from the TimeQuest analyzer to the Technology Map Viewer, the interconnect and cell delay associated with each node is displayed on top of the schematic symbols. The total slack of the selected timing path is displayed in the Page Title section of the schematic. If the nodes are grouped in a logic cloud, the delay information displayed with the logic cloud is the total sum delay of the grouped nodes. The delay information for each node in the logic cloud is displayed in a tooltip. Move the mouse pointer over the logic cloud to see the tooltip. For more information about tooltips, refer to “Tooltips” on page 21–39.

Figure 21–19 shows a portion of a timing path represented in the Technology Map Viewer.

![Figure 21–19. Timing Path Schematic in the Technology Map Viewer](image)

In the RTL Viewer, the schematic page displays the nodes in the paths between the source and destination registers with a summary of the total delay.

The RTL Viewer netlist is based on an initial stage of synthesis, so the post-fitting nodes might not exist in the RTL Viewer netlist. Therefore, the internal delay numbers are not displayed in the RTL Viewer as they are in the Technology Map Viewer, and the timing path might not be displayed exactly as it appears in the timing analysis report. If multiple paths exist between the source and destination registers, the RTL
Viewer might display more than just the timing path. There are also some cases in which the path cannot be displayed, such as paths through state machines, encrypted intellectual property (IP), or registers that are created during the fitting process. In cases where the timing path displayed in the RTL Viewer might not be the correct path, the compiler issues messages.

**Other Features in the Schematic Viewer**

This section describes other features in the schematic view that enhance usability and help you analyze your design.

**Tooltips**

A tooltip is displayed whenever the mouse pointer is held over an element in the schematic. The tooltip contains useful information about a node, net, logic cloud, input port, or output port. Table 21–9 lists the information contained in the tooltip for each type of node.

The tooltip information for an instance (the first row in Table 21–9) includes a list of the primitives found in that level of hierarchy and the number of each primitive contained in the current instance. The number includes all hierarchical blocks below the current instance in the hierarchy. This information lets you estimate the size and complexity of a hierarchical block without navigating into the block.

The tooltip information for atom primitives in the Technology Map Viewer (the second row in Table 21–9) shows the equation for the design atom. The equations are an expanded version of the equations you can view in the Equations window in the Timing Closure Floorplan. Advanced users can use these equations to analyze the design implementation in detail.

For more information about understanding equations, refer to the Quartus II Help.

To copy tooltips into the clipboard for use in other applications, right-click the necessary node or netlist and click **Copy Tooltip**.

To turn off tooltips or change the duration of time that a tooltip appears in the view, in the main Quartus II window, select **Options**. In the **Tooltip Settings** pane, turn on the **Enable tooltips** option.

The **Show names in tooltip for** option specifies the number of seconds to display the names of assigned nodes and pins in a tooltip when the pointer is over the assigned nodes and pins. Selecting **Unlimited** displays the tooltip as long as the pointer remains over the node or pin. Selecting 0 turns off tooltips. The default value is 5 seconds.
The **Delay showing tooltip for** option specifies the number of seconds you must hold the mouse pointer over assigned nodes and pins before the tooltip displays the names of the assigned nodes and pins. Selecting 0 displays the tooltip immediately when the pointer is over an assigned node or pin. Selecting **Unlimited** prevents the display of tooltips. The default value is 1 second.

<table>
<thead>
<tr>
<th>Table 21-9. Tooltip Information (Part 1 of 2)</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Tooltip Format</strong></td>
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<tr>
<td>Instance</td>
</tr>
<tr>
<td>Atom Primitive</td>
</tr>
<tr>
<td>Primitive</td>
</tr>
<tr>
<td>Pin</td>
</tr>
<tr>
<td>Connector</td>
</tr>
<tr>
<td>Net</td>
</tr>
<tr>
<td>Output Port</td>
</tr>
</tbody>
</table>
Finding Design Elements in the Netlist Viewers

You can narrow the range of the search process by setting the following options in the Find pane:

- Click Browse in the Find pane to specify the hierarchy level of the search. In the Select Hierarchy Level dialog box, select the particular instance you want to search.
- Turn on the Include subentities option to include child hierarchies of the parent instance during the search.
- Click Options to open the Find Options dialog box. Turn on Instances, Nodes, Pins, or any combination of the three to further refine the parameters of the search.

When you click the List button, a progress bar appears below the Find box.

All results that match the criteria you set are listed in a table. When you double-click an item in the table, the related node is highlighted in red in the schematic view.
Exporting and Copying a Schematic Image

You can export the schematic view of the RTL Viewer or Technology Map Viewer into various image formats. This allows you to include the schematic in project documentation or share it with other project members. The currently supported formats are JPEG File Interchange Format (.jpg), Portable Network Graphics (.png), Graphics Interchange Format (.gif), and Windows Bitmap (.bmp). To export the schematic view, on the File menu, click Export. In the Export dialog box, type a file name and location and select the necessary file type. The default file name is based on the current instance name; the default file type is .jpg. However, for pages that use filtering, expanding, or reducing operations, the default name is Filter<number of export operation>.jpg.

Nodes grouped as logic clouds are not shown in the exported or copied schematic image; the logic clouds are shown instead.

You can copy the entire image or a portion of the image. To copy the entire image, right-click on the schematic, point to Copy, and then click Full Image. To copy a portion of the image, right-click on the schematic, point to Copy, and then click Partial Image. The cursor changes to a “+” sign to indicate that you can draw a box shape. Drag the mouse pointer around the portion of the schematic you want to copy. When you release the mouse button, the partial image is copied to the clipboard.

Occasionally, due to the design size and objects selected, an image is too large to copy to the clipboard. In this case, the Quartus II software displays an error message.

To export or copy a schematic that is too large to copy in one piece, split the design into multiple pages to export or to copy smaller portions of the design. For more information about controlling how much of your design is shown on each schematic page, refer to “Partitioning the Schematic into Pages” on page 21–28. As an alternative, use the Partial Image feature to copy a portion of the image.

Printing

To print your schematic page, on the File menu, click Print. You can print each schematic page onto one page, or you can print selected parts of your schematic onto one page with the Selection option. To control how much of your design is shown on each schematic page, refer to “Partitioning the Schematic into Pages” on page 21–28.

You cannot print the Netlist Navigator pane in the RTL Viewer and Technology Map Viewer and the table view of the State Machine Viewer. You can use the State Machine Viewer Copy command to copy the table to a text editor and print from the text editor.

Conclusion

The Quartus II RTL Viewer, State Machine Viewer, and Technology Map Viewer allow you to explore and analyze your initial synthesis netlist, post-synthesis netlist, or post-fitting and physical synthesis netlist. The netlist viewers provide a number of features in the Netlist Navigator pane and schematic view to help you quickly trace through your netlist and find specific hierarchies or nodes of interest. These capabilities can help you debug, optimize, and constrain your design more efficiently to increase your productivity.
# Document Revision History

Table 21–10 shows the revision history for this chapter.

**Table 21–10. Document Revision History**

<table>
<thead>
<tr>
<th>Date</th>
<th>Document Version</th>
<th>Changes</th>
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</thead>
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<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Added the following sections:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Properties” on page 21–16</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Global Net Routing” on page 21–7</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Displaying Schematics in a Single Page” on page 21–8</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Displaying Schematics in Multiple Tabbed View” on page 21–8</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Colors” on page 21–21</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Fonts” on page 21–21</td>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Changed to new document template.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Updated screenshots</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated chapter for the Quartus II software version 10.0, including major user interface changes</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Updated devices</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor text edits</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Chapter 13 was formerly Chapter 12 in version 8.1.0</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Figure 13–2, Figure 13–3, Figure 13–4, Figure 13–14, and Figure 13–30</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “Enable or Disable the Auto Hierarchy List” on page 13–15</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Find Command” on page 13–44</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed page size to 8.5” × 11”</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Added Arria GX support</td>
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<td></td>
<td></td>
<td>■ Updated operator symbols</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated information about the radial menu feature</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated zooming feature</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated information about probing from schematic to SignalTap II Analyzer</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated constant signal information</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added .png and .gif to the list of supported image file formats</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated several figures and tables</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Renamed several sections</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed section “Customizing the Radial Menu”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Moved section “Grouping Combinational Logic into Logic Clouds”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated document content based on the Quartus II software version 8.0</td>
</tr>
</tbody>
</table>
For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides additional information about the document and Altera.

**About this Handbook**

This handbook provides comprehensive information about the current release of the Altera® Quartus® II design software.

**How to Contact Altera**

To locate the most up-to-date information about Altera products, refer to the following table.

<table>
<thead>
<tr>
<th>Contact (1)</th>
<th>Contact Method</th>
<th>Address</th>
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<tbody>
<tr>
<td>Technical support</td>
<td>Website</td>
<td><a href="http://www.altera.com/support">www.altera.com/support</a></td>
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<td>Technical training</td>
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<td>Product literature</td>
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<tr>
<td>Nontechnical support (general)</td>
<td>Email</td>
<td><a href="http://www.altera.com/literature">www.altera.com/literature</a></td>
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<tr>
<td>(software licensing)</td>
<td>Email</td>
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<td></td>
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<td><a href="mailto:authorization@altera.com">authorization@altera.com</a></td>
</tr>
</tbody>
</table>

*Note to Table:*
(1) You can also contact your local Altera sales office or sales representative.

**Third-Party Software Product Information**

Third-party software products described in this handbook are not Altera products, are licensed by Altera from third parties, and are subject to change without notice. Updates to these third-party software products may not be concurrent with Quartus II software releases. Altera has assumed responsibility for the selection of such third-party software products and its use in the Quartus II version 12.0 software release. To the extent that the software products described in this handbook are derived from third-party software, no third party warrants the software, assumes any liability regarding use of the software, or undertakes to furnish you any support or information relating to the software. EXCEPT AS EXPRESSLY SET FORTH IN THE APPLICABLE ALTERA PROGRAM LICENSE SUBSCRIPTION AGREEMENT UNDER WHICH THIS SOFTWARE WAS PROVIDED TO YOU, ALTERA AND THIRD-PARTY LICENSORS DISCLAIM ALL WARRANTIES WITH RESPECT TO THE USE OF SUCH THIRD-PARTY SOFTWARE CODE OR DOCUMENTATION IN THE SOFTWARE, INCLUDING, WITHOUT LIMITATION, ANY WARRANTIES OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE, TITLE, AND NONINFRINGEMENT. For more information, including the latest available version of specific third-party software products, refer to the documentation for the software in question.
# Typographic Conventions

The following table shows the typographic conventions this document uses.

<table>
<thead>
<tr>
<th>Visual Cue</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Bold Type with Initial Capital Letters</strong></td>
<td>Indicate command names, dialog box titles, dialog box options, and other GUI labels. For example, <strong>Save As</strong> dialog box. For GUI elements, capitalization matches the GUI.</td>
</tr>
<tr>
<td><strong>bold type</strong></td>
<td>Indicates directory names, project names, disk drive names, file names, file name extensions, software utility names, and GUI labels. For example, \qdesigns directory, D: drive, and chiptrip.gdf file.</td>
</tr>
<tr>
<td><strong>Italic Type with Initial Capital Letters</strong></td>
<td>Indicate document titles. For example, <em>Stratix IV Design Guidelines</em>.</td>
</tr>
<tr>
<td><strong>italic type</strong></td>
<td>Indicates variables. For example, $n + 1$. Variable names are enclosed in angle brackets (&lt;&gt;). For example, &lt;file name&gt; and &lt;project name&gt;.pof file.</td>
</tr>
<tr>
<td>Initial Capital Letters</td>
<td>Indicate keyboard keys and menu names. For example, the Delete key and the Options menu.</td>
</tr>
<tr>
<td>“Subheading Title”</td>
<td>Quotation marks indicate references to sections in a document and titles of Quartus II Help topics. For example, “Typographic Conventions.”</td>
</tr>
<tr>
<td><strong>Courier type</strong></td>
<td>Indicates signal, port, register, bit, block, and primitive names. For example, data1, tdi, and input. The suffix n denotes an active-low signal. For example, resetn. Indicates command line commands and anything that must be typed exactly as it appears. For example, c:\qdesigns\tutorial\chiptrip.gdf. Also indicates sections of an actual file, such as a Report File, references to parts of files (for example, the AHDL keyword SUBDESIGN), and logic function names (for example, TRI).</td>
</tr>
<tr>
<td>←</td>
<td>An angled arrow instructs you to press the Enter key.</td>
</tr>
<tr>
<td>1., 2., 3., and a., b., c., and so on</td>
<td>Numbered steps indicate a list of items when the sequence of the items is important, such as the steps listed in a procedure.</td>
</tr>
<tr>
<td>■ ■ ■</td>
<td>Bullets indicate a list of items when the sequence of the items is not important.</td>
</tr>
<tr>
<td></td>
<td>The hand points to information that requires special attention.</td>
</tr>
<tr>
<td></td>
<td>The question mark directs you to a software help system with related information.</td>
</tr>
<tr>
<td></td>
<td>The feet direct you to another document or website with related information.</td>
</tr>
<tr>
<td></td>
<td>The multimedia icon directs you to a related multimedia presentation.</td>
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<tr>
<td></td>
<td>A caution calls attention to a condition or possible situation that can damage or destroy the product or your work.</td>
</tr>
<tr>
<td></td>
<td>A warning calls attention to a condition or possible situation that can cause you injury.</td>
</tr>
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<td>The envelope links to the Email Subscription Management Center page of the Altera website, where you can sign up to receive update notifications for Altera documents.</td>
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<td>The feedback icon allows you to submit feedback to Altera about the document. Methods for collecting feedback vary as appropriate for each document.</td>
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<td></td>
<td>The social media icons allow you to inform others about Altera documents. Methods for submitting information vary as appropriate for each medium.</td>
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</table>
Quartus II Handbook Version 13.0

Volume 2: Design Implementation and Optimization
The Quartus II Handbook Volume 2: Design Implementation and Optimization was revised on the following dates.

Chapter 1. Constraining Designs
Revised: November 2012
Part Number: QII52001-12.1.0

Chapter 2. Command-Line Scripting
Revised: June 2012
Part Number: QII52002-12.0.0

Chapter 3. Tcl Scripting
Revised: June 2012
Part Number: QII52003-12.0.0

Chapter 4. I/O Management
Revised: May 2013
Part Number: QII52013-13.0.0

Chapter 5. Simultaneous Switching Noise (SSN) Analysis and Optimizations
Revised: June 2012
Part Number: QII52018-12.0.0

Chapter 6. Signal Integrity Analysis with Third-Party Tools
Revised: June 2012
Part Number: QII53020-12.0.0

Chapter 7. Mentor Graphics PCB Design Tools Support
Revised: June 2012
Part Number: QII52015-12.0.0

Chapter 8. Cadence PCB Design Tools Support
Revised: June 2012
Part Number: QII52014-12.0.0

Chapter 9. Reviewing Printed Circuit Board Schematics with the Quartus II Software
Revised: November 2012
Part Number: QII52019-12.1.0

Chapter 10. Design Optimization Overview
Revised: May 2013
Part Number: QII52021-13.0.0

Chapter 11. Reducing Compilation Time
Revised: May 2013
Part Number: QII52022-13.0.0

Chapter 12. Timing Closure and Optimization
Revised: May 2013
Chapter 13. Power Optimization
Revised: May 2013
Part Number: QII52016-13.0.0

Chapter 14. Area Optimization
Revised: May 2013
Part Number: QII52023-13.0.0

Chapter 15. Analyzing and Optimizing the Design Floorplan with the Chip Planner
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Part Number: QII52006-13.0.0

Chapter 16. Netlist Optimizations and Physical Synthesis
Revised: June 2012
Part Number: QII52007-12.0.0

Chapter 17. Engineering Change Management with the Chip Planner
Revised: June 2012
Part Number: QII52017-12.0.0
As a result of the increasing complexity of today’s FPGA designs and the demand for higher performance, designers must make a large number of complex timing and logic constraints to meet their performance requirements. After you create a project and design, you can use the Quartus® II software Assignment Editor and other GUI features to specify your initial design constraints, such as pin assignments, device options, logic options, and timing constraints.

This section describes how to constrain designs, how to take advantage of Quartus II modular executables, how to develop and run Tcl scripts to perform a wide range of functions, and how to manage the Quartus II projects.

This section includes the following chapters:

- **Chapter 1, Constraining Designs**
  This chapter discusses the ways to constrain designs in the Quartus II software, including the tools available in the Quartus II software GUI, as well as Tcl scripting flows.

- **Chapter 2, Command-Line Scripting**
  This chapter discusses Quartus II command-line executables, which provide command-line control over each step of the design flow. Each executable includes options to control commonly used software settings. Each executable also provides detailed, built-in help describing its function, available options, and settings.

- **Chapter 3, Tcl Scripting**
  This chapter discusses developing and running Tcl scripts in the Quartus II software to allow you to perform a wide range of functions, such as compiling a design or automating common tasks. This chapter includes sample Tcl scripts for automating the Quartus II software. You can modify these example scripts for use with your own designs.
This chapter discusses the various tools and methods for constraining and re-constraining Quartus II designs in different design flows, both with the Quartus II GUI and with Tcl to facilitate a scripted flow.

Constraints, sometimes known as assignments or logic options, control the way the Quartus II software implements a design for an FPGA. Constraints are also central in the way that the TimeQuest Timing Analyzer and the PowerPlay Power Analyzer inform synthesis, placement, and routing. There are several types of constraints:

- Global design constraints and software settings, such as device family selection, package type, and pin count.
- Entity-level constraints, such as logic options and placement assignments.
- Instance-level constraints.
- Pin assignments and I/O constraints.

User-created constraints are contained in one of two files: the Quartus II Settings File (.qsf) or, in the case of timing constraints, the Synopsys Design Constraints file (.sdc). Constraints and assignments made with the Device dialog box, Settings dialog box, Assignment Editor, Chip Planner, and Pin Planner are contained in the Quartus II Settings File. The .qsf file contains project-wide and instance-level assignments for the current revision of the project in Tcl syntax. You can create separate revisions of your project with different settings, and there is a separate .qsf file for each revision.

The TimeQuest Timing Analyzer uses industry-standard Synopsys Design Constraints, also using Tcl syntax, that are contained in Synopsys Design Constraints (.sdc) files. The TimeQuest Timing Analyzer GUI is a tool for making timing constraints and viewing the results of subsequent analysis.

There are several ways to constrain a design, each potentially more appropriate than the others, depending on your tool chain and design flow. You can constrain designs for compilation and analysis in the Quartus II software using the GUI, as well as using Tcl syntax and scripting. By combining the Tcl syntax of the .qsf files and the .sdc files with procedural Tcl, you can automate iteration over several different settings, changing constraints and recompiling.
Constraining Designs with the Quartus II GUI

In the Quartus II GUI, the New Project Wizard, Device dialog box, and Settings dialog box allow you to make global constraints and software settings. The Assignment Editor and Pin Planner are spreadsheet-style interfaces for constraining your design at the instance or entity level. The Assignment Editor and Pin Planner make constraint types and values available based on global design characteristics such as the targeted device. These tools help you verify that your constraints are valid before compilation by allowing you to pick only from valid values for each constraint.

The TimeQuest Timing Analyzer GUI allows you to make timing constraints in SDC format and view the effects of those constraints on the timing in your design. Before running the TimeQuest timing analyzer, you must specify initial timing constraints that describe the clock characteristics, timing exceptions, and external signal arrival and required times. The Quartus II Fitter optimizes the placement of logic in the device to meet your specified constraints.

For more information about timing constraints and the TimeQuest Timing Analyzer, refer to About TimeQuest Timing Analysis in Quartus II Help.

Global Constraints

Global constraints affect the entire Quartus II project and all of the applicable logic in the design. Many of these constraints are simply project settings, such as the targeted device selected for the design. Synthesis optimizations and global timing and power analysis settings can also be applied with globally. Global constraints are often made when running the New Project Wizard, or in the Device dialog box or the Settings dialog box, early project development.

The following are the most common types of global constraints:

- Target device specification
- Top-level entity of your design, and the names of the design files included in the project
- Operating temperature limits and conditions
- Physical synthesis optimizations
- Analysis and synthesis options and optimization techniques
- Verilog HDL and VHDL language versions used in your project
- Fitter effort and timing driven compilation settings
- .sdc files for the TimeQuest timing analyzer to use during analysis as part of a full compilation flow

Settings that direct compilation and analysis flows in the Quartus II software are also stored in the Quartus II Settings File for your project, including the following global software settings:

- Early Timing Estimate mode
- Settings for EDA tool integration such as third-party synthesis tools, simulation tools, timing analysis tools, and formal verification tools.
Chapter 1: Constraining Designs

Constraining Designs with the Quartus II GUI

November 2012 Altera Corporation Quartus II Handbook Version 13.0

Volume 2: Design Implementation and Optimization

Settings and settings file specifications for the Quartus II Assembler, SignalTap II Logic Analyzer, PowerPlay power analyzer, and SSN Analyzer.

Global constraints and software settings stored in the Quartus II settings file are specific to each revision of your design, allowing you to control the operation of the software differently for different revisions. For example, different revisions can specify different operating temperatures and different devices, so that you can compare results.

Only the valid assignments made in the Assignment Editor are saved in the Quartus II Settings File, which is located in the project directory. When you make a design constraint, the new assignment is placed on a new line at the end of the file.

When you create or update a constraint in the GUI, the Quartus II software displays the equivalent Tcl command in the System tab of the Messages window. You can use the displayed messages as references when making assignments using Tcl commands.

For more information about specifying initial global constraints and software settings, refer to Setting up and Running a Compilation in Quartus II Help.

For more information about how the Quartus II software uses Quartus II Settings Files, refer to the Managing Quartus II Projects chapter in volume 2 of the Quartus II Handbook.

Node, Entity, and Instance-Level Constraints

Node, entity, and instance-level constraints constrain a particular segment of the design hierarchy, as opposed to the entire design. In the Quartus II software GUI, most instance-level constraints are made with the Assignment Editor, Pin Planner, and Chip Planner. Both the Assignment Editor and Pin Planner aid you in correctly constraining your design, both passively, through device-and-assignment-determined pick lists, and actively, through live I/O checking.

You can assign logic functions to physical resources on the device, using location assignments with the Assignment Editor or the Chip Planner. Node, entity, and instance-level constraints take precedence over any global constraints that affect the same sections of the design hierarchy. You can edit and view all node and entity-level constraints you created in the Assignment Editor, or you can filter the assignments by choosing to view assignments only for specific locations, such as DSP blocks.

The Pin Planner helps you visualize, plan, and assign device I/O pins to ensure compatibility with your PCB layout. The Pin Planner provides a graphical view of the I/O resources in the target device package. You can quickly locate various I/O pins and assign them design elements or other properties. The Quartus II software uses these assignments to place and route your design during device programming. The Pin Planner also helps with early pin planning by allowing you to plan and assign IP interface or user nodes not yet defined in the design.

The Pin Planner Task window provides one-click access to common pin planning tasks. After clicking a pin planning task, you view and highlight the results in the Report window by selecting or deselecting I/O types. You can quickly identify I/O banks, VREF groups, edges, and differential pin pairings to assist you in the pin planning process. You can verify the legality of new and existing pin assignments with the live I/O check feature and view the results in the Live I/O Check Status window.
The Chip Planner allows you to view the device from a variety of different perspectives, and you can make precise assignments to specific floorplan locations. With the Chip Planner, you can adjust existing assignments to device resources, such as pins, logic cells, and LABs using drag and drop features and a graphical interface. You can also view equations and routing information, and demote assignments by dragging and dropping assignments to various regions in the Regions window.

For more information about the Assignment Editor, refer to About the Assignment Editor in Quartus II Help. For more information about the Chip Planner, refer to About the Chip Planner in Quartus II Help. For more information about the Pin Planner, refer to Assigning Device I/O Pins in Pin Planner in Quartus II Help.

Probing Between Components of the Quartus II GUI

The Assignment Editor, Chip Planner, and Pin Planner let you locate nodes and instances in the source files for your design in other Quartus II viewers. You can select a cell in the Assignment Editor spreadsheet and locate the corresponding item in another applicable Quartus II software window, such as the Chip Planner. To locate an item from the Assignment Editor in another window, right-click the item of interest in the spreadsheet, point to Locate, and click the appropriate command.

You can also locate nodes in the Assignment Editor and other constraint tools from other windows within the Quartus II software. First, select the node or nodes in the appropriate window. For example, select an entity in the Entity list in the Hierarchy tab in the Project Navigator, or select nodes in the Chip Planner. Next, right-click the selected object, point to Locate, and click Locate in Assignment Editor. The Assignment Editor opens, or it is brought to the foreground if it is already open.

For more information about the Assignment Editor, refer to About the Assignment Editor in Quartus II Help. For more information about the Chip Planner, refer to About the Chip Planner in Quartus II Help. For more information about the Pin Planner, refer to Assigning Device I/O Pins in Pin Planner in Quartus II Help.

SDC and the TimeQuest Timing Analyzer

You can make individual timing constraints for individual entities, nodes, and pins with the Constraints menu of the TimeQuest Timing Analyzer. The TimeQuest Timing Analyzer GUI provides easy access to timing constraints, and reporting, without requiring knowledge of SDC syntax. As you specify commands and options in the GUI, the corresponding SDC or Tcl command appears in the Console. This lets you know exactly what constraint you have added to your Synopsys Design Constraints file, and also enables you to learn SDC syntax for use in scripted flows. The GUI also provides enhanced graphical reporting features.

Individual timing assignments override project-wide requirements. You can also assign timing exceptions to nodes and paths to avoid reporting of incorrect or irrelevant timing violations. The TimeQuest timing analyzer supports point-to-point timing constraints, wildcards to identify specific nodes when making constraints, and assignment groups to make individual constraints to groups of nodes.

For more information about timing constraints and the TimeQuest Timing Analyzer, refer to About TimeQuest Timing Analysis in Quartus II Help.
Constraining Designs with Tcl

Because .sdc files and .qsf files are both in Tcl syntax, you can modify these files to be part of a scripted constraint and compilation flow. With Quartus II Tcl packages, Tcl scripts can open projects, make the assignments procedurally that would otherwise be specified in a .qsf file, compile a design, and compare compilation results against known goals and benchmarks for the design. Such a script can further automate the iterative process by modifying design constraints and recompiling the design.

For more information about controlling the Quartus II software with Tcl, refer to About Quartus II Tcl Scripting in Quartus II Help.

Quartus II Settings Files and Tcl

QSF files use Tcl syntax, but, unmodified, are not executable scripts. However, you can embed QSF constraints in a scripted iterative compilation flow, where the script that automates compilation and custom results reporting also contains the design constraints. Example 1–1 shows an example QSF file with boilerplate comments removed.
Example 1–1 shows the way that the `set_global_assignment` Quartus II Tcl command makes all global constraints and software settings, with `set_location_assignment` constraining each I/O node in the design to a physical pin on the device.

However, after you initially create the Quartus II Settings File for your design, you can export the contents to a procedural, executable Tcl (.tcl) file. You can then use that generated script to restore certain settings after experimenting with other constraints. You can also use the generated Tcl script to archive your assignments instead of archiving the Quartus II Settings file itself.
To export your constraints as an executable Tcl script, on the Project menu, click Generate Tcl File for Project. Example 1–2 shows the constraints in Example 1–1 converted to an executable Tcl script.

Example 1–2. Generated Tcl Script for a Quartus II Project (Part 1 of 2)

```tcl
# Quartus II: Generate Tcl File for Project
# File: chiptrip.tcl
# Generated on: Tue Jun 08 13:08:48 2010

# Load Quartus II Tcl Project package
package require ::quartus::project

set need_to_close_project 0
set make_assignments 1

# Check that the right project is open
if {{[is_project_open]}} {
  if {[string compare $quartus(project) "chiptrip"]} {
    puts "Project chiptrip is not open"
    set make_assignments 0
  }
} else {
  # Only open if not already open
  if {{[project_exists chiptrip]}} {
    project_open -revision chiptrip chiptrip
  } else {
    project_new -revision chiptrip chiptrip
  }
  set need_to_close_project 1
}

# Make assignments
if {$make_assignments} {
  set_global_assignment -name FAMILY "Cyclone II"
  set_global_assignment -name DEVICE EP2C35F672C6
  set_global_assignment -name TOP_LEVEL_ENTITY chiptrip
  set_global_assignment -name ORIGINAL_QUARTUS_VERSION 10.0
  set_global_assignment -name PROJECT_CREATION_TIME_DATE "11:45:02 JUNE 08, 2010"
  set_global_assignment -name LAST_QUARTUS_VERSION 10.0
  set_global_assignment -name MIN_CORE_JUNCTION_TEMP 0
  set_global_assignment -name MAX_CORE_JUNCTION_TEMP 85
  set_instance_assignment -name PARTITION_HIERARCHY root_partition -to | -section_id Top
  set_global_assignment -name PARTITION_NETLIST_TYPE SOURCE -section_id Top
  set_global_assignment -name PARTITION_FITTER_PRESERVATION_LEVEL PLACEMENT_AND_ROUTING \
  -section_id Top
```
After setting initial values for variables to control constraint creation and whether or not the project needs to be closed at the end of the script, the generated script checks to see if a project is open. If a project is open but it is not the correct project, in this case, chiptrip, the script prints Project chiptrip is not open to the console and does nothing else.

If no project is open, the script determines if chiptrip exists in the current directory. If the project exists, the script opens the project. If the project does not exist, the script creates a new project and opens the project.

The script then creates the constraints. After creating the constraints, the script writes the constraints to the Quartus II Settings File and then closes the project.
Timing Analysis with Synopsys Design Constraints and Tcl

Timing constraints used in analysis by the Quartus II TimeQuest Timing Analyzer are stored in .sdc files. Because they use Tcl syntax, the constraints in .sdc files can be incorporated into other scripts for iterative timing analysis. Example 1–3 shows a basic .sdc file for the chiptrip project.

Example 1–3. Initial .sdc file for the chiptrip Project

```tcl
# ------------------------------------------
set_time_unit ns
set_decimal_places 3
# ------------------------------------------
# create_clock -period 10.0 -waveform { 0 5.0 } clk2 -name clk2
create_clock -period 4.0 -waveform { 0 2.0 } clk1 -name clk1
# clk1 -> dir*: INPUT_MAX_DELAY = 1 ns
set_input_delay -max 1ns -clock clk1 [get_ports dir*]
# clk2 -> time*: OUTPUT_MAX_DELAY = -2 ns
set_output_delay -max -2ns -clock clk2 [get_ports time*]
```

Similar to the constraints in the Quartus II Settings File, you can make the SDC constraints in Example 1–3 part of an executable timing analysis script, as shown in example Example 1–4.

Example 1–4. Tcl Script Making Basic Timing Constraints and Performing Multi-Corner Timing Analysis

```tcl
project_open chiptrip
create_timing_netlist

# Create Constraints
# create_clock -period 10.0 -waveform { 0 5.0 } clk2 -name clk2
create_clock -period 4.0 -waveform { 0 2.0 } clk1 -name clk1
# clk1 -> dir*: INPUT_MAX_DELAY = 1 ns
set_input_delay -max 1ns -clock clk1 [get_ports dir*]
# clk2 -> time*: OUTPUT_MAX_DELAY = -2 ns
set_output_delay -max -2ns -clock clk2 [get_ports time*]

# Perform timing analysis for several different sets of operating conditions
foreach_in_collection oc [get_available_operating_conditions] {
    set_operating_conditions $oc
    update_timing_netlist
    report_timing -setup -npaths 1
    report_timing -hold -npaths 1
    report_timing -recovery -npaths 1
    report_timing -removal -npaths 1
    report_min_pulse_width -nworst 1
}

delete_timing_netlist
project_close
```
The script in Example 1–4 opens the project, creates a timing netlist, then constrains the two clocks in the design and applies input and output delay constraints. The clock settings and delay constraints are identical to those in the .sdc file shown in Example 1–3. The next section of the script updates the timing netlist for the constraints and performs multi-corner timing analysis on the design.

**A Fully Iterative Scripted Flow**

You can use the `::quartus::flow` Tcl package and other packages in the Quartus II Tcl API to add flow control to modify constraints and recompile your design in an automated flow. You can combine your timing constraints with the other constraints for your design, and embed them in an executable Tcl script that also iteratively compiles your design as different constraints are applied.

Each time such a modified generated script is run, it can modify the .qsf file and .sdc file for your project based on the results of iterative compilations, effectively replacing these files for the purposes of archiving and version control using industry-standard source control methods and practices.

This type of scripted flow can include automated compilation of a design, modification of design constraints, and recompilation of the design, based on how you foresee results and pre-determine next-step constraint changes in response to those results.

For more information about the Quartus II Tcl API, refer to *API Functions for Tcl* in Quartus II Help. For more information about controlling the Quartus II software with Tcl scripts, refer to *About Quartus II Tcl Scripting* in Quartus II Help.

**Document Revision History**

Table 1–1 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Update Pin Planner description for task and report windows.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Rewrote chapter to more broadly cover all design constraint methods.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed procedural steps and user interface details, and replaced with</td>
</tr>
<tr>
<td></td>
<td></td>
<td>links to Quartus II Help.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Added two notes.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor text edits.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Revised and reorganized the entire chapter.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added section “Probing to Source Design Files and Other Quartus II</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Windows” on page 1–2.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added description of node type icons (Table 1–3).</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added explanation of wildcard characters.</td>
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Table 1–1. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8½” × 11” page size. No change to content.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Updated Quartus II software 8.0 revision and date.</td>
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</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
FPGA design software that easily integrates into your design flow saves time and improves productivity. The Altera® Quartus® II software provides you with a command-line executable for each step of the FPGA design flow to make the design process customizable and flexible.

The benefits provided by command-line executables include:

- Command-line control over each step of the design flow
- Easy integration with scripted design flows including makefiles
- Reduced memory requirements
- Improved performance

The command-line executables are also completely interchangeable with the Quartus II GUI, allowing you to use the exact combination of tools that you prefer.

This chapter describes how to take advantage of Quartus II command-line executables, and provides several examples of scripts that automate different segments of the FPGA design flow. This chapter includes the following topics:

- “Benefits of Command-Line Executables”
- “Introductory Example” on page 2–2
- “Compilation with quartus_sh --flow” on page 2–7
- “The MegaWizard Plug-In Manager” on page 2–11
- “Command-Line Scripting Examples” on page 2–17

**Benefits of Command-Line Executables**

The Quartus II command-line executables provide control over each step of the design flow. Each executable includes options to control commonly used software settings. Each executable also provides detailed, built-in help describing its function, available options, and settings.

Command-line executables allow for easy integration with scripted design flows. You can easily create scripts with a series of commands. These scripts can be batch-processed, allowing for integration with distributed computing in server farms. You can also integrate the Quartus II command-line executables in makefile-based design flows. These features enhance the ease of integration between the Quartus II software and other EDA synthesis, simulation, and verification software.
Command-line executables add flexibility without sacrificing the ease-of-use of the Quartus II GUI. You can use the Quartus II GUI and command-line executables at different stages in the design flow. For example, you might use the Quartus II GUI to edit the floorplan for the design, use the command-line executables to perform place-and-route, and return to the Quartus II GUI to perform debugging with the Chip Editor.

Command-line executables reduce the amount of memory required during each step in the design flow. Because each executable targets only one step in the design flow, the executables themselves are relatively compact, both in file size and the amount of memory used during processing. This memory usage reduction improves performance, and is particularly beneficial in design environments where heavy usage of computing resources results in reduced memory availability.

For a complete list of the Quartus II command-line executables, refer to Using the Quartus II Executables in Shell Scripts in Quartus II Help.

### Introductory Example

The following introduction to command-line executables demonstrates how to create a project, fit the design, and generate programming files.

The tutorial design included with the Quartus II software is used to demonstrate this functionality. If installed, the tutorial design is found in the `<Quartus II directory>/qdesigns/fir_filter` directory.

Before making changes, copy the tutorial directory and type the four commands shown in Example 2–1 at a command prompt in the new project directory.

![Example 2–1. Introductory Example](quartus_map filtref --source=filtref.bdf --family="Cyclone III"

The `<Quartus II directory>/quartus/bin` directory must be in your `PATH` environment variable.

### Example 2–1. Introductory Example

```
quartus_map filtref --source=filtref.bdf --family="Cyclone III"
quartus_fit filtref --part=EP3C10F256C8 --pack_register=minimize_area
quartus_asm filtref
quartus_sta filtref
```

The `quartus_map filtref --source=filtref.bdf --family="Cyclone III"` command creates a new Quartus II project called `filtref` with `filtref.bdf` as the top-level file. It targets the Cyclone® III device family and performs logic synthesis and technology mapping on the design files.

The `quartus_fit filtref --part=EP3C10F256C8 --pack_register=minimize_area` command performs fitting on the `filtref` project. This command specifies an EP3C10F256C8 device, and the `--pack_register=minimize_area` option causes the Fitter to pack sequential and combinational functions into single logic cells to reduce device resource usage.

The `quartus_asm filtref` command creates programming files for the `filtref` project.

The `quartus_sta filtref` command performs basic timing analysis on the `filtref` project using the Quartus II TimeQuest Timing Analyzer, reporting worst-case setup slack, worst-case hold slack, and other measurements.
The TimeQuest Timing Analyzer employs Synopsys Design Constraints to fully analyze the timing of your design. For more information about using all of the features of the `quartus_sta` executable, refer to the *TimeQuest Timing Analyzer Quick Start Tutorial*.

You can put the four commands from Example 2–1 into a batch file or script file, and run them. For example, you can create a simple UNIX shell script called `compile.sh`, which includes the code shown in Example 2–2.

**Example 2–2. UNIX Shell Script: compile.sh**

```bash
#!/bin/sh
PROJECT=filtref
TOP_LEVEL_FILE=filtref.bdf
FAMILY="Cyclone III"
PART=EP3C10F256C8
PACKING_OPTION=minimize_area
quartus_map $PROJECT --source=$TOP_LEVEL_FILE --family=$FAMILY
quartus_fit $PROJECT --part=$PART --pack_register=$PACKING_OPTION
quartus_asm $PROJECT
quartus_sta $PROJECT
```

Edit the script as necessary and compile your project.

**Command-Line Scripting Help**

Help for command-line executables is available through different methods. You can access help built in to the executables with command-line options. You can use the Quartus II Command-Line and Tcl API Help browser for an easy graphical view of the help information.

To use the Quartus II Command-Line and Tcl API Help browser, type the following command:

`quartus_sh --qhelp`

This command starts the Quartus II Command-Line and Tcl API Help browser, a viewer for information about the Quartus II Command-Line executables and Tcl API (Figure 2–1).
Use the -h option with any of the Quartus II Command-Line executables to get a description and list of supported options. Use the --help=<option name> option for detailed information about each option.

**Figure 2–1. Quartus II Command-Line and Tcl API Help Browser**

**Project Settings with Command-Line Options**

Command-line options are provided for many common global project settings and for performing common tasks. You can use either of two methods to make assignments to an individual entity. If the project exists, open the project in the Quartus II GUI, change the assignment, and close the project. The changed assignment is updated in the.qsf. Any command-line executables that are run after this update use the updated assignment. For more information refer to “Option Precedence” on page 2–5. You can also make assignments using the Quartus II Tcl scripting API. If you want to completely script the creation of a Quartus II project, choose this method.

For more information about the Quartus II Tcl scripting API, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook. For more information about Quartus II project settings and assignments, refer to the QSF Reference Manual.
### Option Precedence

If you use command-line executables, you must be aware of the precedence of various project assignments and how to control the precedence. Assignments for a particular project exist in the Quartus II Settings File (.qsf) for the project. Before the .qsf is updated after assignment changes, the updated assignments are reflected in compiler database files that hold intermediate compilation results.

All command-line options override any conflicting assignments found in the .qsf or the compiler database files. There are two command-line options to specify whether the .qsf or compiler database files take precedence for any assignments not specified as command-line options.

Any assignment not specified as a command-line option or found in the .qsf or compiler database file is set to its default value.

The file precedence command-line options are `--read_settings_files` and `--write_settings_files`.

By default, the `--read_settings_files` and `--write_settings_files` options are turned on. Turning on the `--read_settings_files` option causes a command-line executable to read assignments from the .qsf instead of from the compiler database files. Turning on the `--write_settings_files` option causes a command-line executable to update the .qsf to reflect any specified options, as happens when you close a project in the Quartus II GUI.

If you use command-line executables, be aware of the precedence of various project assignments and how to control the precedence. Assignments for a particular project can exist in three places:

- The .qsf for the project
- The result of the last compilation, in the /db directory, which reflects the assignments that existed when the project was compiled
- Command-line options

Table 2–1 lists the precedence for reading assignments depending on the value of the `--read_settings_files` option.

<table>
<thead>
<tr>
<th>Option Specified</th>
<th>Precedence for Reading Assignments</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>--read_settings_files = on</code></td>
<td>1. Command-line options</td>
</tr>
<tr>
<td>(Default)</td>
<td>2. The .qsf for the project</td>
</tr>
<tr>
<td></td>
<td>3. Project database (db directory, if it exists)</td>
</tr>
<tr>
<td></td>
<td>4. Quartus II software defaults</td>
</tr>
<tr>
<td><code>--read_settings_files = off</code></td>
<td>1. Command-line options</td>
</tr>
<tr>
<td></td>
<td>2. Project database (db directory, if it exists)</td>
</tr>
<tr>
<td></td>
<td>3. Quartus II software defaults</td>
</tr>
</tbody>
</table>
Table 2–2 lists the locations to which assignments are written, depending on the value of the --write_settings_files command-line option.

### Table 2–2. Location for Writing Assignments

<table>
<thead>
<tr>
<th>Option Specified</th>
<th>Location for Writing Assignments</th>
</tr>
</thead>
<tbody>
<tr>
<td>--write_settings_files = on (Default)</td>
<td>.qsf and compiler database</td>
</tr>
<tr>
<td>--write_settings_files = off</td>
<td>Compiler database</td>
</tr>
</tbody>
</table>

Example 2–3 assumes that a project named `fir_filter` exists, and that the analysis and synthesis step has been performed (using the `quartus_map` executable).

#### Example 2–3. Write Settings Files

```
quartus_fit fir_filter --pack_register=off  
quartus_sta fir_filter  
mv fir_filter_sta.rpt fir_filter_1_sta.rpt  
quartus_fit fir_filter --pack_register=minimize_area --write_settings_files=off  
quartus_sta fir_filter  
mv fir_filter_sta.rpt fir_filter_2_sta.rpt  
```

The first command, `quartus_fit fir_filter --pack_register=off`, runs the `quartus_fit` executable with no aggressive attempts to reduce device resource usage.

The second command, `quartus_sta fir_filter`, performs basic timing analysis for the results of the previous fit.

The third command uses the UNIX `mv` command to copy the report file output from `quartus_sta` to a file with a new name, so that the results are not overwritten by subsequent timing analysis.

The fourth command runs `quartus_fit` a second time, and directs it to attempt to pack logic into registers to reduce device resource usage. With the --write_settings_files=off option, the command-line executable does not update the .qsf to reflect the changed register packing setting. Instead, only the compiler database files reflect the changed setting. If the --write_settings_files=off option is not specified, the command-line executable updates the .qsf to reflect the register packing setting.

The fifth command reruns timing analysis, and the sixth command renames the report file, so that it is not overwritten by subsequent timing analysis.

Use the options --read_settings_files=off and --write_settings_files=off (where appropriate) to optimize the way that the Quartus II software reads and updates settings files. In Example 2–4, the `quartus_asm` executable does not read or write settings files because doing so would not change any settings for the project.

#### Example 2–4. Avoiding Unnecessary Reading and Writing

```
quartus_map filtref --source=filtref --part=EP3C10F256C8  
quartus_fit filtref --pack_register=off --read_settings_files=off  
quartus_asm filtref --read_settings_files=off --write_settings_files=off  
```
Compilation with `quartus_sh --flow`

Figure 2–2 shows a typical Quartus II FPGA design flow using command-line executables.

**Figure 2–2. Typical Design Flow**

Use the `quartus_sh` executable with the `--flow` option to perform a complete compilation flow with a single command. The `--flow` option supports the smart recompile feature and efficiently sets command-line arguments for each executable in the flow.

The following example runs compilation, timing analysis, and programming file generation with a single command:

```shell
quartus_sh --flow compile filtre
```

For information about specialized flows, type `quartus_sh --help=flow` at a command prompt.
Text-Based Report Files

Each command-line executable creates a text report file when it is run. These files report success or failure, and contain information about the processing performed by the executable.

Report file names contain the revision name and the short-form name of the executable that generated the report file, in the format <revision>.<executable>.rpt. For example, using the quartus_fit executable to place and route a project with the revision name design_top generates a report file named design_top.fit.rpt. Similarly, using the quartus_sta executable to perform timing analysis on a project with the revision name fir_filter generates a report file named fir_filter.sta.rpt.

As an alternative to parsing text-based report files, you can use the `::quartus::report` Tcl package. For more information about this package, refer to `::quartus::report` in Quartus II Help.

Using Command-Line Executables In Scripts

You can use command-line executables in scripts that control other software in addition to the Quartus II software. For example, if your design flow uses third-party synthesis or simulation software, and if you can run the other software at a command prompt, you can include those commands with Quartus II executables in a single script.

The Quartus II command-line executables include options for common global project settings and operations, but you must use a Tcl script or the Quartus II GUI to set up a new project and apply individual constraints, such as pin location assignments and timing requirements. Command-line executables are very useful for working with existing projects, for making common global settings, and for performing common operations. For more flexibility in a flow, use a Tcl script, which makes it easier to pass data between different stages of the design flow and have more control during the flow.

For more information about Tcl scripts, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook, or About Quartus II Tcl Scripting in Quartus II Help.
For example, a UNIX shell script could run other synthesis software, then place-and-route the design in the Quartus II software, then generate output netlists for other simulation software. Example 2–5 shows a script that synthesizes a design with the Synopsys Synplify software, simulates the design using the Mentor Graphics ModelSim® software, and then compiles the design targeting a Cyclone III device.

Example 2–5. Script for End-to-End Flow

```bash
#!/bin/sh
# Run synthesis first.
# This example assumes you use Synplify software
#synplify -batch synthesize.tcl

# If your Quartus II project exists already, you can just
# recompile the design.
# You can also use the script described in a later example to
# create a new project from scratch
#quartus_sh --flow compile myproject

# Use the quartus_sta executable to do fast and slow-model
# timing analysis
quartus_sta myproject --model=slow
quartus_sta myproject --model=fast

# Use the quartus_eda executable to write out a gate-level
# Verilog simulation netlist for ModelSim
quartus_eda my_project --simulation --tool=modelsim --format=verilog

# Perform the simulation with the ModelSim software
vlib cycloneiii_ver
vlog -work cycloneiii_ver /opt/quartusii/eda/sim_lib/cycloneiii_atoms.v
vlib work
vlog -work work my_project.vo
vsim -L cycloneiii_ver -t 1ps work.my_project
```

Makefile Implementation

You can use the Quartus II command-line executables in conjunction with the make utility to automatically update files when other files they depend on change. The file dependencies and commands used to update files are specified in a text file called a makefile.

To facilitate easier development of efficient makefiles, the following “smart action” scripting command is provided with the Quartus II software:

`quartus_sh --determine_smart_action`

Because assignments for a Quartus II project are stored in the .qsf, including it in every rule results in unnecessary processing steps. For example, updating a setting related to programming file generation, which requires re-running only `quartus_asm`, modifies the .qsf, requiring a complete recompilation if the .qsf is included in every rule.

The smart action command determines the earliest command-line executable in the compilation flow that must be run based on the current .qsf, and generates a change file corresponding to that executable. For example, if `quartus_map` must be re-run, the smart action command creates or updates a file named `map.chg`. Thus, rather than including the .qsf in each makefile rule, include only the appropriate change file.
Example 2–6 uses change files and the smart action command. You can copy and modify it for your own use. A copy of this example is included in the help for the makefile option, which is available by typing:

```
quartus_sh --help=makefiles
```

### Example 2–6. Sample Makefile (Part 1 of 2)

```bash
# Project Configuration:
# Specify the name of the design (project), the Quartus II Settings
# File (.qsf), and the list of source files used.
PROJECT = chiptrip
SOURCE_FILES = auto_max.v chiptrip.v speed_ch.v tick_cnt.v time_cnt.v
ASSIGNMENT_FILES = chiptrip.qpf chiptrip.qsf

# Main Targets
# all: build everything
# clean: remove output files and database
all: smart.log $(PROJECT).asm.rpt $(PROJECT).sta.rpt

map: smart.log $(PROJECT).map.rpt
fit: smart.log $(PROJECT).fit.rpt
asm: smart.log $(PROJECT).asm.rpt
sta: smart.log $(PROJECT).sta.rpt
smart: smart.log

# Executable Configuration
MAP_ARGS = --family=Stratix
FIT_ARGS = --part=EP1S20F484C6
ASM_ARGS =
STA_ARGS =

# Target implementations
STAMP = echo done >

$(PROJECT).map.rpt: map.chg $(SOURCE_FILES)
quartus_map $(MAP_ARGS) $(PROJECT)
$(STAMP) fit.chg
```
The MegaWizard Plug-In Manager

The MegaWizard™ Plug-In Manager provides a GUI-based flow to configure megafunction and IP variation files. However, you can use command-line options for the qmegawiz executable to modify, update, or create variation files without using the GUI. This capability is useful in a fully scripted design flow, or in cases where you want to generate variation files without using the wizard GUI flow.

The MegaWizard Plug-In Manager has three functions:

■ Providing an interface for you to select the output file or files
■ Running a specific MegaWizard Plug-In
■ Creating output files (such as variation files, symbol files, and simulation netlist files)

Each MegaWizard Plug-In provides a user interface for configuring the variation, and performs validation and error checking of your selected ports and parameters. When you create or update a variation with the GUI, the parameters and values are entered through the GUI provided by the Plug-In. When you create a Plug-In variation with the command line, you provide the parameters and values as command-line options.

Example 2–7 shows how to create a new variation file at a system command prompt.

Example 2–7. MegaWizard Plug-In Manager Command-Line Executable

qmegawiz [options] [module=<module name>|wizard=<wizard name>] [<param>=<value> ...] [<port>=<used|unused> ...] [OPTIONAL_FILES=<optional files>] <variation file name>

When you use qmegawiz to update an existing variation file, the module or wizard name is not required.
If a megafunction changes between software versions, the variation files must be regenerated. To do this, use `qmegawiz -silent <variation file name>`. For example, if your design contains a variation file called `myrom.v`, type the following command:

```
qmegawiz -silent myrom.v
```

For more information on updating megafunction variation files as part of a scripted flow, refer to “Regenerating Megafuctions After Updating the Quartus II Software” on page 2–23.

Table 2–3 describes the supported options.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>-silent</code></td>
<td>Run the MegaWizard Plug-In Manager in command-line mode, without displaying the GUI.</td>
</tr>
<tr>
<td><code>-f:&lt;param file&gt;</code></td>
<td>A file that contains all options for the <code>qmegawiz</code> command. Refer to “Parameter File” on page 2–16.</td>
</tr>
</tbody>
</table>

For information about specifying the module name or wizard name, refer to “Module and Wizard Names” on page 2–13.

For information about specifying ports and parameters, refer to “Ports and Parameters” on page 2–14.

For information about generating optional files, refer to “Optional Files” on page 2–15.

For information about specifying the variation file name, refer to “Variation File Name” on page 2–17.

### Command-Line Support

Only the MegaWizard Plug-Ins listed in Table 2–4 support creation and update in command-line mode. For Plug-Ins not listed in the table, you must use the MegaWizard Plug-In Manager GUI for creation and updates.

<table>
<thead>
<tr>
<th>MegaWizard Plug-In</th>
<th>Wizard Name</th>
<th>Module Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>alt2gxb</td>
<td>ALT2GXB</td>
<td>alt2gxb</td>
</tr>
<tr>
<td>alt4gxb</td>
<td>ALTGX</td>
<td>alt4gxb</td>
</tr>
<tr>
<td>altasmii_parallel</td>
<td>ALTASMI_PARALLEL</td>
<td>altasmii_parallel</td>
</tr>
<tr>
<td>alttclkctrl</td>
<td>ALTTCLKCTRL</td>
<td>alttclkctrl</td>
</tr>
<tr>
<td>altddio_bidir</td>
<td>ALTDDIO_BIDIR</td>
<td>altddio_bidir</td>
</tr>
<tr>
<td>altddio_in</td>
<td>ALTDDIO_IN</td>
<td>altddio_in</td>
</tr>
<tr>
<td>altddio_out</td>
<td>ALTDDIO_OUT</td>
<td>altddio_out</td>
</tr>
<tr>
<td>altecc_decoder</td>
<td>ALTECC</td>
<td>altecc_decoder</td>
</tr>
<tr>
<td>altecc_encoder</td>
<td>ALTECC</td>
<td>altecc_encoder</td>
</tr>
<tr>
<td>altfp_abs</td>
<td>ALTFP_ABS</td>
<td>altfp_abs</td>
</tr>
</tbody>
</table>

Table 2–4. MegaWizard Plug-Ins with Command Line Support  (Part 1 of 2)
Table 2–4. MegaWizard Plug-Ins with Command Line Support (Part 2 of 2)

<table>
<thead>
<tr>
<th>MegaWizard Plug-In</th>
<th>Wizard Name</th>
<th>Module Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>altfp_add_sub</td>
<td>ALTFP_ADD_SUB</td>
<td>altfp_add_sub</td>
</tr>
<tr>
<td>altfp_atan</td>
<td>ALTFP_ATAN</td>
<td>altfp_atan</td>
</tr>
<tr>
<td>altfp_compare</td>
<td>ALTFP_COMPARE</td>
<td>altfp_compare</td>
</tr>
<tr>
<td>altfp_convert</td>
<td>ALTFP_CONVERT</td>
<td>altfp_convert</td>
</tr>
<tr>
<td>altfp_div</td>
<td>ALTFP_DIV</td>
<td>altfp_div</td>
</tr>
<tr>
<td>altfp_exp</td>
<td>ALTFP_EXP</td>
<td>altfp_exp</td>
</tr>
<tr>
<td>altfp_inv_sqrt</td>
<td>ALTFP_INV_SQRT</td>
<td>altfp_inv_sqrt</td>
</tr>
<tr>
<td>altfp_inv</td>
<td>ALTFP_INV</td>
<td>altfp_inv</td>
</tr>
<tr>
<td>altfp_log</td>
<td>ALTFP_LOG</td>
<td>altfp_log</td>
</tr>
<tr>
<td>altfp_matrix_inv</td>
<td>ALTFP_MATRIX_INV</td>
<td>altfp_matrix_inv</td>
</tr>
<tr>
<td>altfp_matrix_mult</td>
<td>ALTFP_MATRIXMULT</td>
<td>altfp_matrix_mult</td>
</tr>
<tr>
<td>altfp_mult</td>
<td>ALTFP_MULT</td>
<td>altfp_mult</td>
</tr>
<tr>
<td>altfp_sincos</td>
<td>ALTFP_SINCOS</td>
<td>altfp_sincos</td>
</tr>
<tr>
<td>altfp_sqrt</td>
<td>ALTFP_SQRT</td>
<td>altfp_sqrt</td>
</tr>
<tr>
<td>altiobuf_bidir</td>
<td>ALTIIOBUF</td>
<td>altiobuf_bidir</td>
</tr>
<tr>
<td>altiobuf_in</td>
<td>ALTIIOBUF</td>
<td>altiobuf_in</td>
</tr>
<tr>
<td>altiobuf_out</td>
<td>ALTIIOBUF</td>
<td>altiobuf_out</td>
</tr>
<tr>
<td>altlvds_rx</td>
<td>ALTLVDS</td>
<td>altlvds_rx</td>
</tr>
<tr>
<td>alt.lvds_tx</td>
<td>ALTLVDS</td>
<td>alt.lvds_tx</td>
</tr>
<tr>
<td>altmult_accum</td>
<td>ALTMULT_ACCUM(MAC)</td>
<td>altmult_accum</td>
</tr>
<tr>
<td>altmult_complex</td>
<td>ALTMULT_COMPLEX</td>
<td>altmult_complex</td>
</tr>
<tr>
<td>altotp</td>
<td>ALTOTP</td>
<td>altotp</td>
</tr>
<tr>
<td>altpll_reconfig</td>
<td>ALTPLL_RECONFIG</td>
<td>altpll_reconfig</td>
</tr>
<tr>
<td>altpll</td>
<td>ALTPLL</td>
<td>altpll</td>
</tr>
<tr>
<td>altremote_update</td>
<td>ALTREMOTE_UPDATE</td>
<td>altremote_update</td>
</tr>
<tr>
<td>altshift_taps</td>
<td>ALTSHIFT_TAPS</td>
<td>altshift_taps</td>
</tr>
<tr>
<td>altsyncram</td>
<td>RAM:2-PORT</td>
<td>altsyncram</td>
</tr>
<tr>
<td></td>
<td>RAM:1-PORT</td>
<td></td>
</tr>
<tr>
<td></td>
<td>ROM:1-PORT</td>
<td></td>
</tr>
<tr>
<td>alttemp_sense</td>
<td>ALTTEMP_SENSE</td>
<td>alttemp_sense</td>
</tr>
<tr>
<td>alt_c3gx</td>
<td>ALT_C3GXB</td>
<td>alt_c3gx</td>
</tr>
<tr>
<td>dcfifo</td>
<td>FIFO</td>
<td>dcfifo</td>
</tr>
<tr>
<td>scfifo</td>
<td>FIFO</td>
<td>scfifo</td>
</tr>
</tbody>
</table>

**Module and Wizard Names**

You must specify the wizard or module name, shown in Table 2–4, as a command-line option when you create a variation file. Use the option `module=<module_name>` to specify the module, or use the option `wizard=<wizard_name>` to specify the wizard. If there are spaces in the wizard or module name, enclose the name in double quotes, for example:
wizard="RAM: 2-PORT"

When there is a one-to-one mapping between the MegaWizard Plug-In, the wizard name, and the module name, you can use either the wizard option or the module option.

When there are multiple wizard names that correspond to one module name, use the wizard option to specify one wizard. For example, use the wizard option if you create a RAM, because one module is common to three wizards.

When there are multiple module names that correspond to one wizard name, use the module option to specify one module. For example, use the module option if you create a FIFO because one wizard is common to both modules.

If you edit or update an existing variation file, the wizard or module option is not necessary, because information about the wizard or module is already in the variation file.

**Ports and Parameters**

Ports and parameters for each MegaWizard Plug-In are described in Quartus II Help, and in the Megafuntion User Guides on the Altera website. Use these references to determine appropriate values for each port and parameter required for a particular variation configuration. Refer to “Strategies to Determine Port and Parameter Values” for more information. You do not have to specify every port and parameter supported by a Plug-In. The MegaWizard Plug-In Manager uses default values for any port or parameter you do not specify.

Specify ports as used or unused, for example:

```plaintext
<port>=used
<port>=unused
```

You can specify port names in any order. Grouping does not matter. Separate port configuration options from each other with spaces.

Specify a value for a parameter with the equal sign, for example:

```plaintext
<parameter>=<value>
```

You can specify parameters in any order. Grouping does not matter. Separate parameter configuration options from each other with spaces. You can specify port names and parameter names in upper or lower case; case does not matter.

All MegaWizard Plug-Ins allow you to specify the target device family with the `INTENDED_DEVICE_FAMILY` parameter, as shown in the following example:

```plaintext
gmegawiz wizard=<wizard> INTENDED_DEVICE_FAMILY="Cyclone III" <file>
```

You must specify enough ports and parameters to create a legal configuration of the Plug-In. When you use the GUI flow, each MegaWizard Plug-In performs validation and error checking for the particular ports and parameters you choose. When you use command-line options to specify ports and parameters, you must ensure that the ports and parameters you use are complete for your particular configuration.

For example, when you use a RAM Plug-In to configure a RAM to be 32 words deep, the Plug-In automatically configures an address port that is five bits wide. If you use the command-line flow to configure a RAM that is 32 words deep, you must use one option to specify the depth of the RAM, then calculate the width of the address port and specify that width with another option.
Invalid Configurations

If the combination of default and specified ports and parameters is not complete to create a legal configuration of the Plug-In, \texttt{qmegawiz} generates an error message that indicates what is missing and what values are supported. If the combination of default and specified ports and parameters results in an illegal configuration of the Plug-In, \texttt{qmegawiz} generates an error message that indicates what is illegal, and displays the legal values.

Strategies to Determine Port and Parameter Values

For simple Plug-In variations, it is often easy to determine appropriate port and parameter values with the information in Quartus II Help and other megafunction documentation. For example, determining that a 32-word-deep RAM requires an address port that is five bits wide is straightforward. For complex Plug-In variations, an option in the GUI might affect multiple port and parameter settings, so it can be difficult to determine a complete set of ports and parameters. In this case, use the GUI to generate a variation file that includes the ports and parameters for your desired configuration. Open the variation file in a text editor and use the port and parameter values in the variation file as command-line options.

Optional Files

In addition to the variation file, the MegaWizard Plug-In Manager can generate other files, such as instantiation templates, simulation netlists, and symbols for graphic design entry. Use the \texttt{OPTIONAL_FILES} parameter to control whether the MegaWizard Plug-In Manager generates optional files. Table 2–5 lists valid arguments for the \texttt{OPTIONAL_FILES} parameter.

Table 2–5. Arguments for the \texttt{OPTIONAL_FILES} Parameter

<table>
<thead>
<tr>
<th>Argument</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>INST</td>
<td>Controls the generation of the &lt;variation&gt;_inst.v file.</td>
</tr>
<tr>
<td>INC</td>
<td>Controls the generation of the &lt;variation&gt;.inc file.</td>
</tr>
<tr>
<td>CMP</td>
<td>Controls the generation of the &lt;variation&gt;.cmp file.</td>
</tr>
<tr>
<td>BSF</td>
<td>Controls the generation of the &lt;variation&gt;.bsf file.</td>
</tr>
<tr>
<td>BB</td>
<td>Controls the generation of the &lt;variation&gt;_bb.v file.</td>
</tr>
<tr>
<td>SIM_NETLIST</td>
<td>Controls the generation of the simulation netlist file, wherever there is wizard support.</td>
</tr>
<tr>
<td>SYNTH_NETLIST</td>
<td>Controls the generation of the synthesis netlist file, wherever there is wizard support.</td>
</tr>
<tr>
<td>ALL</td>
<td>Generates all applicable optional files.</td>
</tr>
<tr>
<td>NONE</td>
<td>Disables the generation of all optional files.</td>
</tr>
</tbody>
</table>

Specify a single optional file, for example:

\texttt{OPTIONALFILES=<argument>}

Specify multiple optional files separated by a vertical bar character, for example:

\texttt{OPTIONALFILES=<argument 1> | ... | <argument n>
If you prefix an argument with a dash (for example, -BB), it is excluded from the generated optional files. If any of the optional files exist when you run qmegawiz and they are excluded in the OPTIONAL_FILES parameter (with the NONE argument, or prefixed with a dash), they are deleted.

You can combine the ALL argument with other excluded arguments to generate “all files except <excluded files>.” You can combine the NONE argument with other included arguments to generate “no files except <files>.

When you combine multiple arguments, they are processed from left to right, and arguments evaluated later have precedence over arguments evaluated earlier. Therefore, use the ALL or NONE arguments first in a series of multiple arguments. When ALL is the first argument, all optional files are generated before exclusions are processed (deleted). When NONE is the first argument, none of the optional files are generated (in other words, any that exist are deleted), then any files you subsequently specify are generated.

Table 2–6 shows examples for the OPTIONAL_FILES parameter and describes the result of each example.

<table>
<thead>
<tr>
<th>Example Values for OPTIONAL_FILES</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>BB</td>
<td>The optional file &lt;variation&gt;_bb.v is generated, and no optional files are deleted</td>
</tr>
<tr>
<td>BB</td>
<td>INST</td>
</tr>
<tr>
<td>NONE</td>
<td>No optional files are generated, and any existing optional files are deleted.</td>
</tr>
<tr>
<td>NONE</td>
<td>INC</td>
</tr>
<tr>
<td>ALL</td>
<td>INST</td>
</tr>
<tr>
<td>-BB</td>
<td>The optional file &lt;variation&gt;_bb.v is deleted if it exists.</td>
</tr>
<tr>
<td>-BB</td>
<td>INST</td>
</tr>
</tbody>
</table>

The qmegawiz command accepts the ALL argument combined with other included file arguments, for example, ALL|BB, but that combination is equivalent to ALL because first all optional files are generated, and then the file <variation>_bb.v is generated a second time. Additionally, the software accepts the NONE argument combined with other excluded file arguments, for example, NONE|-BB, but that combination is equivalent to NONE because no optional files are generated, any that exist are deleted, and then the file <variation>_bb.v is deleted if it exists.

### Parameter File

You can put all parameter values and port values in a file, and pass the file name as an argument to qmegawiz with the `-f:<parameter file>` option. For example, the following command specifies a parameter file named `rom_params.txt`:

```
qmegawiz -silent module=altsyncram -f:rom_params.txt myrom.v
```

The `rom_params.txt` parameter file can include options similar to the following:
Working Directory

You can change the working directory that qmegawiz uses when it generates files. By default, the working directory is the current directory when you execute the qmegawiz command. Use the `-p` option to specify a different working directory, for example:

```
-p: <working directory>
```

You can specify the working directory with an absolute or relative path. Specify an alternative working directory any time you do not want files generated in the current directory. The alternative working directory can be useful if you generate multiple variations in a batch script, and keep generated files for the different Plug-In variations in separate directories.

If you use the `-f` option and the `-p` option together, the MegaWizard Plug-In Manager sources the parameter file in a directory specified with the `-p` option, or in a directory relative to that directory. For example, if you specify `C:\project\work\params.txt` with the `-f` option, the MegaWizard Plug-In Manager attempts to source the file `params.txt` in `C:\project\work\work`.

Variation File Name

The language used for a variation file depends on the file extension of the variation file name. The MegaWizard Plug-In Manager creates HDL output files in a language based on the file name extension. Therefore, you must always specify a complete file name, including file extension, as the last argument to the qmegawiz command. Table 2–7 shows the file extension that corresponds to supported HDL types.

<table>
<thead>
<tr>
<th>Variation File HDL Type</th>
<th>Required File Extension</th>
</tr>
</thead>
<tbody>
<tr>
<td>Verilog HDL</td>
<td><code>.v</code></td>
</tr>
<tr>
<td>VHDL</td>
<td><code>.vhd</code></td>
</tr>
<tr>
<td>AHDL</td>
<td><code>.tdf</code></td>
</tr>
</tbody>
</table>

Command-Line Scripting Examples

This section presents various examples of command-line executable use.

Create a Project and Apply Constraints

The command-line executables include options for common global project settings and commands. To apply constraints such as pin locations and timing assignments, run a Tcl script with the constraints in it. You can write a Tcl constraint file yourself, or generate one for an existing project. From the Project menu, click Generate Tcl File for Project.
Example 2–8 creates a project with a Tcl script and applies project constraints using the tutorial design files in the `<Quartus II installation directory>/qdesigns/fir_filter/` directory.

**Example 2–8. Tcl Script to Create Project and Apply Constraints**

```tcl
project_new filtref -overwrite
# Assign family, device, and top-level file
set_global_assignment -name FAMILY Cyclone
set_global_assignment -name DEVICE EP1C12F256C6
set_global_assignment -name BDF_FILE filtref.bdf
# Assign pins
set_location_assignment -to clk Pin_28
set_location_assignment -to clkx2 Pin_29
set_location_assignment -to d[0] Pin_139
set_location_assignment -to d[1] Pin_140
# Other assignments could follow
project_close
```

Save the script in a file called `setup_proj.tcl` and type the commands illustrated in Example 2–9 at a command prompt to create the design, apply constraints, compile the design, and perform fast-corner and slow-corner timing analysis. Timing analysis results are saved in two files, `filtref_sta_1.rpt` and `filtref_sta_2.rpt`.

**Example 2–9. Script to Create and Compile a Project**

```bash
quartus_sh -t setup_proj.tcl
quartus_map filtref
quartus_fit filtref
quartus_asm filtref
quartus_sta filtref --model=fast --export_settings=off
mv filtref_sta.rpt filtref_sta_1.rpt
quartus_sta filtref --export_settings=off
mv filtref_sta.rpt filtref_sta_2.rpt
```

Type the following commands to create the design, apply constraints, and compile the design, without performing timing analysis:

```bash
quartus_sh -t setup_proj.tcl
quartus_sh --flow compile filtref
```

The `quartus_sh --flow compile` command performs a full compilation, and is equivalent to clicking the **Start Compilation** button in the toolbar.

**Check Design File Syntax**

The UNIX shell script example shown in Example 2–10 assumes that the Quartus II software `fir_filter` tutorial project exists in the current directory. You can find the `fir_filter` project in the `<Quartus II directory>/qdesigns/fir_filter` directory unless the Quartus II software tutorial files are not installed.

The `--analyze_file` option causes the `quartus_map` executable to perform a syntax check on each file. The script checks the exit code of the `quartus_map` executable to determine whether there is an error during the syntax check. Files with syntax errors are added to the `FILES_WITH_ERRORS` variable, and when all files are checked, the script prints a message indicating syntax errors.
When options are not specified, the executable uses the project database values. If not specified in the project database, the executable uses the Quartus II software default values. For example, the `fir_filter` project is set to target the Cyclone device family, so it is not necessary to specify the `--family` option.

**Example 2–10. Shell Script to Check Design File Syntax**

```bash
#!/bin/sh
FILES_WITH_ERRORS=""
# Iterate over each file with a .bdf or .v extension
for filename in `ls *.bdf *.v`
do
  # Perform a syntax check on the specified file
  quartus_map fir_filter --analyze_file=$filename
  # If the exit code is non-zero, the file has a syntax error
  if [ $? -ne 0 ]
    then
      FILES_WITH_ERRORS="$FILES_WITH_ERRORS $filename"
    fi
  done
if [ -z "$FILES_WITH_ERRORS" ]
then
  echo "All files passed the syntax check"
  exit 0
else
  echo "There were syntax errors in the following file(s)"
  echo $FILES_WITH_ERRORS
  exit 1
fi
```

**Create a Project and Synthesize a Netlist Using Netlist Optimizations**

This example creates a new Quartus II project with a file `top.edf` as the top-level entity. The `--enable_register_retiming=on` and `--enable_wysiwyg_resynthesis=on` options cause `quartus_map` to optimize the design using gate-level register retiming and technology remapping.

For more information about register retiming, WYSIWYG primitive resynthesis, and other netlist optimization options, refer to Quartus II Help.

The `--part` option causes `quartus_map` to target an EP3C10F256C8 device. To create the project and synthesize it using the netlist optimizations described above, type the command shown in **Example 2–11** at a command prompt.

**Example 2–11. Creating a Project and Synthesizing a Netlist Using Netlist Optimizations**

```bash
quartus_map top --source=top.edf --enable_register_retiming=on
--enable_wysiwyg_resynthesis=on --part=EP3C10F256C8
```
Archive and Restore Projects

You can archive or restore a Quartus II Archive File (.qar) with a single command. This makes it easy to take snapshots of projects when you use batch files or shell scripts for compilation and project management. Use the --archive or --restore options for quartus_sh as appropriate. Type the command shown in Example 2–12 at a command prompt to archive your project.

Example 2–12. Archiving a Project

```
quartus_sh --archive <project name>
```

The archive file is automatically named <project name>.qar. If you want to use a different name, type the command with the -output option as shown in example Example 2–13.

Example 2–13. Archiving a Project

```
quartus_sh --archive <project name> -output <filename>
```

To restore a project archive, type the command shown in Example 2–14 at a command prompt.

Example 2–14. Restoring a Project Archive

```
quartus_sh --restore <archive name>
```

The command restores the project archive to the current directory and overwrites existing files.

For more information about archiving and restoring projects, refer to the Managing Quartus II Projects chapter in volume 2 of the Quartus II Handbook.

Perform I/O Assignment Analysis

You can perform I/O assignment analysis with a single command. I/O assignment analysis checks pin assignments to ensure they do not violate board layout guidelines. I/O assignment analysis does not require a complete place and route, so it can quickly verify that your pin assignments are correct. The command shown in Example 2–15 performs I/O assignment analysis for the specified project and revision.

Example 2–15. Performing I/O Assignment Analysis

```
quartus_fit --check_ios <project name> --rev=<revision name>
```

Update Memory Contents Without Recompiling

You can use two commands to update the contents of memory blocks in your design without recompiling. Use the quartus_cdb executable with the --update_mif option to update memory contents from .mif or .hexout files. Then, rerun the assembler with the quartus_asm executable to regenerate the .sof, .pof, and any other programming files.
Example 2–16 shows these two commands.

Example 2–16. Commands to Update Memory Contents Without Recompiling

```bash
quartus_cdb --update_mif <project name> [--rev=<revision name>]
quartus_asm <project name> [--rev=<revision name>]
```

Example 2–17 shows the commands for a DOS batch file for this example. With a DOS batch file, you can specify the project name and the revision name once for both commands. To create the DOS batch file, paste the following lines into a file called `update_memory.bat`.

Example 2–17. Batch file to Update Memory Contents Without Recompiling

```bash
quartus_cdb --update_mif %1 --rev=%2
quartus_asm %1 --rev=%2
```

To run the batch file, type the following command at a command prompt:

```bash
update_memory.bat <project name> <revision name>
```

Create a Compressed Configuration File

You can create a compressed configuration file in two ways. The first way is to run `quartus_cpf` with an option file that turns on compression.

To create an option file that turns on compression, type the following command at a command prompt:

```bash
quartus_cpf -w <filename>.opt
```

This interactive command guides you through some questions, then creates an option file based on your answers. Use `--option` to cause `quartus_cpf` to use the option file. For example, the following command creates a compressed `.pof` that targets an EPCS64 device:

```bash
quartus_cpf --convert --option=<filename>.opt --device=EPCS64 <file>.sof <file>.pof
```

Alternatively, you can use the Convert Programming Files utility in the Quartus II software GUI to create a Conversion Setup File (.cof). Configure any options you want, including compression, then save the conversion setup. Use the following command to run the conversion setup you specified:

```bash
quartus_cpf --convert <file>.cof
```

Fit a Design as Quickly as Possible

This example assumes that a project called `top` exists in the current directory, and that the name of the top-level entity is `top`. The `--effort=fast` option causes the `quartus_fit` to use the fast fit algorithm to increase compilation speed, possibly at the expense of reduced $f_{MAX}$ performance. The `--one_fit_attempt=on` option restricts the Fitter to only one fitting attempt for the design.
To attempt to fit the project called top as quickly as possible, type the command shown in Example 2–18 at a command prompt.

**Example 2–18. Fitting a Project Quickly**

```
quartus_fit top --effort=fast --one_fit_attempt=on
```

### Fit a Design Using Multiple Seeds

This shell script example assumes that the Quartus II software tutorial project called fir_filter exists in the current directory (defined in the file fir_filter.qpf). If the tutorial files are installed on your system, this project exists in the `<Quartus II directory>/qdesigns/quartus_version_number/fir_filter` directory. Because the top-level entity in the project does not have the same name as the project, you must specify the revision name for the top-level entity with the --rev option. The --seed option specifies the seeds to use for fitting.

A seed is a parameter that affects the random initial placement of the Quartus II Fitter. Varying the seed can result in better performance for some designs.

After each fitting attempt, the script creates new directories for the results of each fitting attempt and copies the complete project to the new directory so that the results are available for viewing and debugging after the script has completed.

**Example 2–19** is designed for use on UNIX systems using `sh` (the shell).

---

**Example 2–19. Shell Script to Fit a Design Using Multiple Seeds**

```bash
#!/bin/sh
ERROR_SEEDS=""
quartus_map fir_filter --rev=filtref
# Iterate over a number of seeds
for seed in 1 2 3 4 5
do
echo "Starting fit with seed=$seed"
# Perform a fitting attempt with the specified seed
quartus_fit fir_filter --seed=$seed --rev=filtref
# If the exit-code is non-zero, the fitting attempt was
# successful, so copy the project to a new directory
if [ $? -eq 0 ]
    then
        mkdir ../fir_filter-seed_$seed
        mkdir ../fir_filter-seed_$seed/db
        cp * ../fir_filter-seed_$seed
        cp db/* ../fir_filter-seed_$seed/db
    else
        ERROR_SEEDS="$ERROR_SEEDS $seed"
fi
done
if [ -z "$ERROR_SEEDS" ]
then
echo "Seed sweeping was successful"
exit 0
else
echo "There were errors with the following seed(s)"
echo $ERROR_SEEDS
exit 1
fi```

---
Use the Design Space Explorer (DSE) included with the Quartus II software script (by typing `quartus_sh --dse` at a command prompt) to improve design performance by performing automated seed sweeping.

For more information about the DSE, type `quartus_sh --help=dse` at a command prompt, or refer to Design Space Explorer in Quartus II Help.

**Regenerating Megafunctions After Updating the Quartus II Software**

Some megafunction variations may require regeneration when you update your installation of the Quartus II software. Read the release notes for the Quartus II software and any new documentation for the IP functions used in your design to determine if regeneration is necessary.

If regeneration is necessary, you can use a Tcl script to run the `qmegawiz` executable to update each function, allowing you to avoid regenerating each function in the Megawizard Plug-In Manager GUI.

Wizard-generated files are identified in the Source Files Used report panel (contained in `<project name>.map.rpt`) in the File Type column as “Auto-Found Wizard-Generated File”. In a Tcl script, use the commands in the `::quartus::report` package from the Quartus II Tcl API to recover the list of files. Use the `qexec` command in a loop to run `qmegawiz` for each variation file:

```
qexec "qmegawiz -silent <variation file name>"
```

For example, if your script determines that your design contains a variation file called `myrom.v`, in one iteration of the loop in your script, a combination of strings and variables passed to the `qexec` command would be equivalent to the following command:

```
qexec "qmegawiz -silent myrom.v"
```

If your design flow incorporates parameter files, those can be included in the `qmegawiz` call in the same way you would include them from a command prompt:

```
qexec "qmegawiz -silent -f:<parameter file>.txt <variation file name>"
```

For more information about the `::quartus::report` Tcl package, refer to `::quartus::report` in Quartus II Help.

For more information about the Quartus II Tcl scripting API, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook.

**The QFlow Script**

A Tcl/Tk-based graphical interface called QFlow is included with the command-line executables. You can use the QFlow interface to open projects, launch some of the command-line executables, view report files, and make some global project assignments. The QFlow interface can run the following command-line executables:

- `quartus_map` (Analysis and Synthesis)
- `quartus_fit` (Fitter)
- `quartus_sta` (TimeQuest timing analyzer)
■ quartus_asm (Assembler)
■ quartus_eda (EDA Netlist Writer)

To view floorplans or perform other GUI-intensive tasks, launch the Quartus II software.

Start QFlow by typing the following command at a command prompt:

\texttt{quartus_sh -g}

The QFlow script is located in the \texttt{<Quartus II directory>/common/tcl/apps/qflow/} directory.

### Document Revision History

Table 2–8 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Corrected quartus_qpf example usage.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated examples.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Template update.</td>
</tr>
</tbody>
</table>
|               |         | Added section on using a script to regenerate megafun
|               |         | ction variations.                                         |
|               |         | Removed references to the Classic Timing Analyzer (quartus_tan).|
|               |         | Removed QFlow illustration.                               |
| July 2010     | 10.0.0  | Updated script examples to use quartus_sta instead of quartus_tan, and other minor updates throughout document. |
| November 2009 | 9.1.0   | Updated Table 2–1 to add quartus_jli and quartus_jbcc executables and descriptions, and other minor updates throughout document. |
| March 2009    | 9.0.0   | No change to content.                                     |
Table 2–8. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Added the following sections:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “The MegaWizard Plug-In Manager” on page 2–11</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Command-Line Support” on page 2–12</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Module and Wizard Names” on page 2–13</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Ports and Parameters” on page 2–14</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Invalid Configurations” on page 2–15</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Strategies to Determine Port and Parameter Values” on page 2–15</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Optional Files” on page 2–15</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Parameter File” on page 2–16</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Working Directory” on page 2–17</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Variation File Name” on page 2–17</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Create a Compressed Configuration File” on page 2–21</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Option Precedence” on page 2–5 to clarify how to control precedence</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Corrected Example 2–5 on page 2–8</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed Example 2–1, Example 2–2, Example 2–4, and Example 2–7 to use the EP1C12F256C6 device</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated entire chapter using 8½” × 11” chapter template</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Updated “Referenced Documents” on page 2–20.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated references in document.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
Introduction

Developing and running Tcl scripts to control the Altera® Quartus® II software allows you to perform a wide range of functions, such as compiling a design or writing procedures to automate common tasks.

You can use Tcl scripts to manage a Quartus II project, make assignments, define design constraints, make device assignments, compile your design, perform timing analysis, and access reports. Tcl scripts also facilitate project or assignment migration. For example, when designing in different projects with the same prototype or development board, you can automate reassignment of pin locations in each new project. The Quartus II software can also generate a Tcl script based on all the current assignments in the project, which aids in switching assignments to another project.

The Quartus II software Tcl commands follow the EDA industry Tcl application programming interface (API) standards for command-line options. This simplifies learning and using Tcl commands. If you encounter an error with a command argument, the Tcl interpreter includes help information showing correct usage.

This chapter includes sample Tcl scripts for automating the Quartus II software. You can modify these example scripts for use with your own designs. You can find more Tcl scripts in the Design Examples section of the Support area on the Altera website.

This chapter includes the following topics:

- “Quartus II Tcl Packages” on page 3–2
- “Quartus II Tcl API Help” on page 3–3
- “Command-Line Options: -t, -s, and --tcl_eval” on page 3–5
- “End-to-End Design Flows” on page 3–7
- “Creating Projects and Making Assignments” on page 3–7
- “Compiling Designs” on page 3–8
- “Reporting” on page 3–9
- “Timing Analysis” on page 3–10
- “Automating Script Execution” on page 3–10
- “Other Scripting Features” on page 3–13
- “The Quartus II Tcl Shell in Interactive Mode” on page 3–17
Tool Command Language

Tcl (pronounced “tickle”) stands for Tool Command Language, a popular scripting language that is similar to many shell scripting and high-level programming languages. It provides support for control structures, variables, network socket access, and APIs. Tcl is the EDA industry-standard scripting language used by Synopsys, Mentor Graphics®, and Altera software. It allows you to create custom commands and works seamlessly across most development platforms. For a list of recommended literature on Tcl, refer to “External References” on page 3–25.

You can create your own procedures by writing scripts containing basic Tcl commands and Quartus II API functions. You can then automate your design flow, run the Quartus II software in batch mode, or execute the individual Tcl commands interactively in the Quartus II Tcl interactive shell.

If you are unfamiliar with Tcl scripting, or are a Tcl beginner, refer to “Tcl Scripting Basics” on page 3–18 for an introduction to Tcl scripting.

The Quartus II software supports Tcl/Tk version 8.5, supplied by the Tcl DeveloperXchange at tcl.activestate.com.

Quartus II Tcl Packages

The Quartus II Tcl commands are grouped in packages by function. Table 3–1 describes each Tcl package.

Table 3–1. Tcl Packages (Part 1 of 2)

<table>
<thead>
<tr>
<th>Package Name</th>
<th>Package Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>backannotate</td>
<td>Back annotate assignments</td>
</tr>
<tr>
<td>chip_planner</td>
<td>Identify and modify resource usage and routing with the Chip Editor</td>
</tr>
<tr>
<td>database_manager</td>
<td>Manage version-compatible database files</td>
</tr>
<tr>
<td>device</td>
<td>Get device and family information from the device database</td>
</tr>
<tr>
<td>flow</td>
<td>Compile a project, run command-line executables and other common flows</td>
</tr>
<tr>
<td>incremental compilation</td>
<td>Manipulate design partitions and LogicLock regions, and settings related to incremental compilation</td>
</tr>
<tr>
<td>insystem_memory_edit</td>
<td>Read and edit memory contents in Altera devices</td>
</tr>
<tr>
<td>insystem_source_probe</td>
<td>interact with the In-System Sources and Probes tool in an Altera device</td>
</tr>
<tr>
<td>jtag</td>
<td>Control the JTAG chain</td>
</tr>
<tr>
<td>logic_analyzer_interface</td>
<td>Query and modify the logic analyzer interface output pin state</td>
</tr>
<tr>
<td>misc</td>
<td>Perform miscellaneous tasks such as enabling natural bus naming, package loading, and message posting</td>
</tr>
<tr>
<td>project</td>
<td>Create and manage projects and revisions, make any project assignments including timing assignments</td>
</tr>
<tr>
<td>rapid_recompile</td>
<td>Manipulate Quartus II Rapid Recompile features</td>
</tr>
<tr>
<td>report</td>
<td>Get information from report tables, create custom reports</td>
</tr>
</tbody>
</table>
By default, only the minimum number of packages is loaded automatically with each Quartus II executable. This keeps the memory requirement for each executable as low as possible. Because the minimum number of packages is automatically loaded, you must load other packages before you can run commands in those packages.

Because different packages are available in different executables, you must run your scripts with executables that include the packages you use in the scripts. For example, if you use commands in the `sdc_ext` package, you must use the `quartus_sta` executable to run the script because the `quartus_sta` executable is the only one with support for the `sdc_ext` package.

The following command prints lists of the packages loaded or available to load for an executable, to the console:

```
<executable name> --tcl_eval help
```

For example, type the following command to list the packages loaded or available to load by the `quartus_fit` executable:

```
quartus_fit --tcl_eval help
```

### Loading Packages

To load a Quartus II Tcl package, use the `load_package` command as follows:

```
load_package [-version <version number>] <package name>
```

This command is similar to the `package require` Tcl command (described in Table 3–2 on page 3–4), but you can easily alternate between different versions of a Quartus II Tcl package with the `load_package` command because of the `-version` option.

---

**Quartus II Tcl API Help**

Access the Quartus II Tcl API Help reference by typing the following command at a system command prompt:

```
quartus_sh --qhelp
```

This command runs the Quartus II Command-Line and Tcl API help browser, which documents all commands and options in the Quartus II Tcl API.

---

**Table 3–1. Tcl Packages (Part 2 of 2)**

<table>
<thead>
<tr>
<th>Package Name</th>
<th>Package Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rtl</td>
<td>Traversing and querying the RTL netlist of your design</td>
</tr>
<tr>
<td>sdc</td>
<td>Specifies constraints and exceptions to the TimeQuest Timing Analyzer</td>
</tr>
<tr>
<td>sdc_ext</td>
<td>Altera-specific SDC commands</td>
</tr>
<tr>
<td>simulator</td>
<td>Configure and perform simulations</td>
</tr>
<tr>
<td>sta</td>
<td>Contains the set of Tcl functions for obtaining advanced information from the Quartus II TimeQuest Timing Analyzer</td>
</tr>
<tr>
<td>stp</td>
<td>Run the SignalTap® II Logic Analyzer</td>
</tr>
</tbody>
</table>
Quartus II Tcl help allows easy access to information about the Quartus II Tcl commands. To access the help information, type `help` at a Tcl prompt, as shown in Example 3–1.

**Example 3–1. Help Output**

```tcl
help
```

---

```
Available Quartus II Tcl Packages:
```

<table>
<thead>
<tr>
<th>Loaded</th>
<th>Not Loaded</th>
</tr>
</thead>
<tbody>
<tr>
<td>::quartus::misc</td>
<td>::quartus::device</td>
</tr>
<tr>
<td>::quartus::old_api</td>
<td>::quartus::backannotate</td>
</tr>
<tr>
<td>::quartus::project</td>
<td>::quartus::flow</td>
</tr>
<tr>
<td>::quartus::timing_assignment</td>
<td>::quartus::logiclock</td>
</tr>
<tr>
<td>::quartus::timing_report</td>
<td>::quartus::report</td>
</tr>
</tbody>
</table>

* Type "help -tcl"
to get an overview on Quartus II Tcl usages.

**Table 3–2. Help Options Available in the Quartus II Tcl Environment (Part 1 of 2)**

<table>
<thead>
<tr>
<th>Help Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>help</td>
<td>To view a list of available Quartus II Tcl packages, loaded and not loaded.</td>
</tr>
<tr>
<td>help -tcl</td>
<td>To view a list of commands used to load Tcl packages and access command-line help.</td>
</tr>
</tbody>
</table>
| help -pkg <package_name>      | To view help for a specified Quartus II package that includes the list of available Tcl commands. For convenience, you can omit the `::quartus::` package prefix, and type `help -pkg <package_name>`.
|     [-version <version number>] | If you do not specify the `-version` option, help for the currently loaded package is displayed by default. If the package for which you want help is not loaded, help for the latest version of the package is displayed by default. Examples: `help -pkg ::quartus::project` `help -pkg project` `help -pkg project -version 1.0` |
| <command_name> -h or <command_name> -help | To view short help for a Quartus II Tcl command for which the package is loaded. Examples: `project_open -h` `project_open -help` |
The Tcl API help is also available in Quartus II online help. Search for the command or package name to find details about that command or package.

### Command-Line Options: -t, -s, and --tcl_eval

Table 3–3 lists three command-line options you can use with executables that support Tcl.

#### Table 3–3. Command-Line Options Supporting Tcl Scripting (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command-Line Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>--script= &lt;script file&gt; [ &lt;script args&gt; ]</td>
<td>Run the specified Tcl script with optional arguments.</td>
</tr>
<tr>
<td>-t &lt;script file&gt; [ &lt;script args&gt; ]</td>
<td>Run the specified Tcl script with optional arguments. The -t option is the short form of the --script option.</td>
</tr>
<tr>
<td>--shell</td>
<td>Open the executable in the interactive Tcl shell mode.</td>
</tr>
</tbody>
</table>
Run a Tcl Script

Running an executable with the -t option runs the specified Tcl script. You can also specify arguments to the script. Access the arguments through the argv variable, or use a package such as cmdline, which supports arguments of the following form:

-<argument name> <argument value>

The cmdline package is included in the <Quartus II directory>/common/tcl/packages directory.

For example, to run a script called myscript.tcl with one argument, Stratix, type the following command at a system command prompt:

quartus_sh -t myscript.tcl Stratix

Refer to “Accessing Command-Line Arguments” on page 3–15 for more information.

Interactive Shell Mode

Running an executable with the -s option starts an interactive Tcl shell. For example, to open the Quartus II TimeQuest Timing Analyzer executable in interactive shell mode, type the following command:

quartus_sta -s

Commands you type in the Tcl shell are interpreted when you click Enter. You can run a Tcl script in the interactive shell with the following command:

source <script name>

If a command is not recognized by the shell, it is assumed to be an external command and executed with the exec command.

Evaluate as Tcl

Running an executable with the --tcl_eval option causes the executable to immediately evaluate the remaining command-line arguments as Tcl commands. This can be useful if you want to run simple Tcl commands from other scripting languages.

For example, the following command runs the Tcl command that prints out the commands available in the project package.

quartus_sh --tcl_eval help -pkg project
The Quartus II Tcl Console Window

You can run Tcl commands directly in the Quartus II Tcl Console window. On the View menu, click Utility Windows. By default, the Tcl Console window is docked in the bottom-right corner of the Quartus II GUI. All Tcl commands typed in the Tcl Console are interpreted by the Quartus II Tcl shell.

Some shell commands such as cd, ls, and others can be run in the Tcl Console window, with the Tcl exec command. However, for best results, run shell commands and Quartus II executables from a system command prompt outside of the Quartus II software GUI.

Tcl messages appear in the System tab (Messages window). Errors and messages written to stdout and stderr also are shown in the Quartus II Tcl Console window.

End-to-End Design Flows

You can use Tcl scripts to control all aspects of the design flow, including controlling other software, when the other software also includes a scripting interface.

Typically, EDA tools include their own script interpreters that extend core language functionality with tool-specific commands. For example, the Quartus II Tcl interpreter supports all core Tcl commands, and adds numerous commands specific to the Quartus II software. You can include commands in one Tcl script to run another script, which allows you to combine or chain together scripts to control different tools. Because scripts for different tools must be executed with different Tcl interpreters, it is difficult to pass information between the scripts unless one script writes information into a file and another script reads it.

Within the Quartus II software, you can perform many different operations in a design flow (such as synthesis, fitting, and timing analysis) from a single script, making it easy to maintain global state information and pass data between the operations. However, there are some limitations on the operations you can perform in a single script due to the various packages supported by each executable.

There are no limitations on running flows from any executable. Flows include operations found in the Start section of the Processing menu in the Quartus II GUI, and are also documented as options for the execute_flow Tcl command. If you can make settings in the Quartus II software and run a flow to get your desired result, you can make the same settings and run the same flow in a Tcl script.

Creating Projects and Making Assignments

You can easily create a script that makes all the assignments for an existing project, and then use the script at any time to restore your project settings to a known state. From the Project menu, click Generate Tcl File for Project to automatically generate a .tcl file with all of your assignments. You can source this file to recreate your project, and you can edit the file to add other commands, such as compiling the design. The file is a good starting point to learn about project management commands and assignment commands.
Refer to “Interactive Shell Mode” on page 3-6 for information about sourcing a script. Scripting information for all Quartus II project settings and assignments is located in the QSF Reference Manual. Refer to the Constraining Designs chapter in volume 2 of the Quartus II Handbook for more information on making assignments.

Example 3–2 shows how to create a project, make assignments, and compile the project. It uses the fir_filter tutorial design files in the qdesigns installation directory. Run this script in the fir_filter directory, with the quartus_sh executable.

Example 3–2. Create and Compile a Project

load_package flow

# Create the project and overwrite any settings # files that exist
project_new fir_filter -revision filtref -overwrite
# Set the device, the name of the top-level BDF, # and the name of the top level entity
set_global_assignment -name FAMILY Cyclone
set_global_assignment -name DEVICE EP1C6F256C6
set_global_assignment -name BDF_FILE filtref.bdf
set_global_assignment -name TOP_LEVEL_ENTITY filtref
# Add other pin assignments here
set_location_assignment -to clk Pin_G1
# compile the project
execute_flow -compile
project_close

The assignments created or modified while a project is open are not committed to the Quartus II Settings File (.qsf) unless you explicitly call export_assignments or project_close (unless -dont_export_assignments is specified). In some cases, such as when running execute_flow, the Quartus II software automatically commits the changes.


Compiling Designs

You can run the Quartus II command-line executables from Tcl scripts. Use the included flow package to run various Quartus II compilation flows, or run each executable directly.

The flow Package

The flow package includes two commands for running Quartus II command-line executables, either individually or together in standard compilation sequence. The execute_module command allows you to run an individual Quartus II command-line executable. The execute_flow command allows you to run some or all of the executables in commonly-used combinations. Use the flow package instead of system calls to run Quartus II executables from scripts or from the Quartus II Tcl Console.
Compile All Revisions

You can use a simple Tcl script to compile all revisions in your project. Save the script shown in Example 3–3 in a file called `compile_revisions.tcl` and type the following to run it:

```
quartus_sh -t compile_revisions.tcl <project name>
```

**Example 3–3. Compile All Revisions**

```tcl
load_package flow
project_open [lindex $quartus(args) 0]
set original_revision [get_current_revision]
foreach revision [get_project_revisions] {
    set_current_revision $revision
    execute flow -compile
}
set_current_revision $original_revision
project_close
```

Reporting

It is sometimes necessary to extract information from the Compilation Report to evaluate results. The Quartus II Tcl API provides easy access to report data so you do not have to write scripts to parse the text report files.

If you know the exact cell or cells you want to access, use the `get_report_panel_data` command and specify the row and column names (or x and y coordinates) and the name of the appropriate report panel. You can often search for data in a report panel. To do this, use a loop that reads the report one row at a time with the `get_report_panel_row` command.

Column headings in report panels are in row 0. If you use a loop that reads the report one row at a time, you can start with row 1 to skip the row with column headings. The `get_number_of_rows` command returns the number of rows in the report panel, including the column heading row. Because the number of rows includes the column heading row, continue your loop as long as the loop index is less than the number of rows.

Report panels are hierarchically arranged and each level of hierarchy is denoted by the string “||” in the panel name. For example, the name of the Fitter Settings report panel is `Fitter||Fitter Settings` because it is in the Fitter folder. Panels at the highest hierarchy level do not use the “||” string. For example, the Flow Settings report panel is named `Flow Settings`.

The code in Example 3–4 prints a list of all report panel names in your project. You can run this code with any executable that includes support for the report package.

**Example 3–4. Print All Report Panel Names**

```tcl
load_package report
project_open myproject
load_report
set panel_names [get_report_panel_names]
foreach panel_name $panel_names {
    post_message "$panel_name"
}
```
Viewing Report Data in Excel

The Microsoft Excel software is sometimes used to view or manipulate timing analysis results. You can create a Comma Separated Value (.csv) file from any Quartus II report to open with Excel. Example 3–5 shows a simple way to create a .csv file with data from the Fitter panel in a report. You could modify the script to use command-line arguments to pass in the name of the project, report panel, and output file to use. You can run this script example with any executable that supports the report package.

Example 3–5. Create .csv Files from Reports

```tcl
load_package report
project_open my-project

load_report

# This is the name of the report panel to save as a CSV file
set panel_name "Fitter||Fitter Settings"
set csv_file "output.csv"

set fh [open $csv_file w]
set num_rows [get_number_of_rows -name $panel_name]

# Go through all the rows in the report file, including the
# row with headings, and write out the comma-separated data
for { set i 0 } { $i < $num_rows } { incr i } {
    set row_data [get_report_panel_row -name $panel_name
        -row $i]
    puts $fh [join $row_data "," ]
}

close $fh
unload_report
```

Timing Analysis

The Quartus II TimeQuest Timing Analyzer includes support for industry-standard SDC commands in the sdc package. The Quartus II software also includes comprehensive Tcl APIs and SDC extensions for the TimeQuest Timing Analyzer in the sta, and sdc_ext packages.

Refer to the Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook for detailed information about how to perform timing analysis with the Quartus II TimeQuest Timing Analyzer.

Automating Script Execution

You can configure scripts to run automatically at various points during compilation. Use this capability to automatically run scripts that perform custom reporting, make specific assignments, and perform many other tasks.

The following three global assignments control when a script is run automatically:

- `PRE_FLOW_SCRIPT_FILE` — before a flow starts
- `POST_MODULE_SCRIPT_FILE` — after a module finishes
POST_FLOW_SCRIPT_FILE — after a flow finishes

A module is another term for a Quartus II executable that performs one step in a flow. For example, two modules are Analysis and Synthesis (quartus_map), and timing analysis (quartus_sta).

A flow is a series of modules that the Quartus II software runs with predefined options. For example, compiling a design is a flow that typically consists of the following steps (performed by the indicated module):

1. Analysis and synthesis (quartus_map)
2. Fitter (quartus_fit)
3. Assembler (quartus_asm)
4. Timing Analyzer (quartus_sta)

Other flows are described in the help for the execute_flow Tcl command. In addition, many commands in the Processing menu of the Quartus II GUI correspond to this design flow.

To make an assignment automatically run a script, add an assignment with the following form to the .qsf for your project:

```tcl
set_global_assignment -name <assignment name> <executable>:<script name>
```

The Quartus II software runs the scripts as shown in Example 3–6.

**Example 3–6.**

```tcl
<executable> -t <script name> <flow or module name> <project name> <revision name>
```

The first argument passed in the argv variable (or quartus(args) variable) is the name of the flow or module being executed, depending on the assignment you use. The second argument is the name of the project and the third argument is the name of the revision.

When you use the POST_MODULE_SCRIPT_FILE assignment, the specified script is automatically run after every executable in a flow. You can use a string comparison with the module name (the first argument passed in to the script) to isolate script processing to certain modules.

**Execution Example**

Example 3–7 illustrates how automatic script execution works in a complete flow, assuming you have a project called top with a current revision called rev_1, and you have the following assignments in the .qsf for your project:

**Example 3–7.**

```tcl
set_global_assignment -name PRE_FLOW_SCRIPT_FILE quartus_sh:first.tcl
set_global_assignment -name POST_MODULE_SCRIPT_FILE quartus_sh:next.tcl
set_global_assignment -name POST_FLOW_SCRIPT_FILE quartus_sh:last.tcl
```

When you compile your project, the PRE_FLOW_SCRIPT_FILE assignment causes the following command to be run before compilation begins:
Next, the Quartus II software starts compilation with analysis and synthesis, performed by the `quartus_map` executable. After the analysis and synthesis finishes, the `POST_MODULE_SCRIPT_FILE` assignment causes the following command to run:

```
quartus_sh -t next.tcl quartus_map top rev_1
```

Then, the Quartus II software continues compilation with the Fitter, performed by the `quartus_fit` executable. After the Fitter finishes, the `POST_MODULE_SCRIPT_FILE` assignment runs the following command:

```
quartus_sh -t next.tcl quartus_fit top rev_1
```

Corresponding commands are run after the other stages of the compilation. When the compilation is over, the `POST_FLOW_SCRIPT_FILE` assignment runs the following command:

```
quartus_sh -t last.tcl compile top rev_1
```

### Controlling Processing

The `POST_MODULE_SCRIPT_FILE` assignment causes a script to run after every module. Because the same script is run after every module, you might have to include some conditional statements that restrict processing in your script to certain modules.

For example, if you want a script to run only after timing analysis, use a conditional test like the one shown in Example 3–8. It checks the flow or module name passed as the first argument to the script and executes code when the module is `quartus_sta`.

**Example 3–8. Restrict Processing to a Single Module**

```tcl
set module [lindex $quartus(args) 0]
if [string match "quartus_sta" $module] {
    # Include commands here that are run after timing analysis
    # Use the post-message command to display messages
    post_message "Running after timing analysis"
}
```

### Displaying Messages

Because of the way the Quartus II software runs the scripts automatically, you must use the `post_message` command to display messages, instead of the `puts` command. This requirement applies only to scripts that are run by the three assignments listed in “Automating Script Execution” on page 3–10.

Refer to “The `post_message` Command” on page 3–14 for more information about this command.
Other Scripting Features

The Quartus II Tcl API includes other general-purpose commands and features described in this section.

Natural Bus Naming

The Quartus II software supports natural bus naming. Natural bus naming allows you to use square brackets to specify bus indexes in HDL without including escape characters to prevent Tcl from interpreting the square brackets as containing commands. For example, one signal in a bus named address can be identified as address[0] instead of address\[0\]. You can take advantage of natural bus naming when making assignments, as in Example 3–9.

Example 3–9. Natural Bus Naming

```
set_location_assignment -to address[10] Pin_M20
```

The Quartus II software defaults to natural bus naming. You can turn off natural bus naming with the disable_natural_bus_naming command. For more information about natural bus naming, type the following at a Quartus II Tcl prompt:

```
enable_natural_bus_naming -h
```

Short Option Names

You can use short versions of command options, as long as they are unambiguous. For example, the project_open command supports two options: -current_revision and -revision. You can use any of the following abbreviations of the -revision option: -r, -re, -rev, -revis, and -revis. You can use an option as short as -r because in the case of the project_open command no other option starts with the letter r. However, the report_timing command includes the options -recovery and -removal. You cannot use -r or -re to shorten either of those options, because the abbreviation would not be unique to only one option.

Collection Commands

Some Quartus II Tcl functions return very large sets of data that would be inefficient as Tcl lists. These data structures are referred to as collections. The Quartus II Tcl API uses a collection ID to access the collection. There are two Quartus II Tcl commands for working with collections, foreach_in_collection and get_collection_size. Use the set command to assign a collection ID to a variable.

For information about which Quartus II Tcl commands return collection IDs, refer to foreach_in_collection in Quartus II Help.
The `foreach_in_collection` Command

The `foreach_in_collection` command is similar to the `foreach` Tcl command. Use it to iterate through all elements in a collection. Example 3–10 prints all instance assignments in an open project.

Example 3–10. Collection Commands

```tcl
set all_instance_assignments [get_all_instance_assignments -name *]
foreach_in_collection asgn $all_instance_assignments {
    # Information about each assignment is
    # returned in a list. For information
    # about the list elements, refer to Help
    # for the get-all-instance-assignments command.
    set to [lindex $asgn 2]
    set name [lindex $asgn 3]
    set value [lindex $asgn 4]
    puts "Assignment to $to: $name = $value"
}
```

The `get_collection_size` Command

Use the `get_collection_size` command to get the number of elements in a collection. Example 3–11 prints the number of global assignments in an open project.

Example 3–11. `get_collection_size` Command

```tcl
set all_global_assignments [get_all_global_assignments -name *]
set num_global_assignments [get_collection_size $all_global_assignments]
puts "There are $num_global_assignments global assignments in your project"
```

The `post_message` Command

To print messages that are formatted like Quartus II software messages, use the `post_message` command. Messages printed by the `post_message` command appear in the `System` tab of the Messages window in the Quartus II GUI, and are written to standard output when scripts are run. Arguments for the `post_message` command include an optional message type and a required message string.

The message type can be one of the following:

- info (default)
- extra_info
- warning
- critical_warning
- error

If you do not specify a type, Quartus II software defaults to `info`.

With the Quartus II software in Windows, you can color code messages displayed at the system command prompt with the `post_message` command. Add the following line to your `quartus2.ini` file:

```
DISPLAY_COMMAND_LINE_MESSAGES_IN_COLOR = on
```
Example 3–12 shows how to use the `post_message` command.

**Example 3–12. post_message command**

```tcl
post_message -type warning "Design has gated clocks"
```

### Accessing Command-Line Arguments

Many Tcl scripts are designed to accept command-line arguments, such as the name of a project or revision. The global variable `quartus(args)` is a list of the arguments typed on the command-line following the name of the Tcl script. Example 3–13 shows code that prints all of the arguments in the `quartus(args)` variable.


```tcl
set i 0
foreach arg $quartus(args) {
    puts "The value at index $i is $arg"
    incr i
}
```

If you copy the script in the previous example to a file named `print_args.tcl`, it displays the following output when you type the command shown in Example 3–14 at a command prompt.

**Example 3–14. Passing Command-Line Arguments to Scripts**

```tcl
quartus_sh -t print_args.tcl my_project 100MHz
```

The value at index 0 is `my_project`
The value at index 1 is `100MHz`

### The cmdline Package

You can use the `cmdline` package included with the Quartus II software for more robust and self-documenting command-line argument passing. The `cmdline` package supports command-line arguments with the form `-<option> <value>`.

**Example 3–15 uses the cmdline package.**

**Example 3–15. cmdline Package**

```tcl
package require cmdline
variable ::argv0 $::quartus(args)
set options {
    { "project.arg" "" "Project name" }
    { "frequency.arg" "" "Frequency" }
}
set usage "You need to specify options and values"
array set optshash [::cmdline::getoptions ::argv $options $usage]
puts "The project name is $optshash(project)"
puts "The frequency is $optshash(frequency)"
```
If you save those commands in a Tcl script called `print_cmd_args.tcl` you see the following output when you type the command shown in Example 3–16 at a command prompt.

### Example 3–16. Passing Command-Line Arguments for Scripts

```
quartus_sh -t print_cmd_args.tcl -project my_project -frequency 100MHz r
```

The project name is `my_project`
The frequency is `100MHz`

Virtually all Quartus II Tcl scripts must open a project. Example 3–17 opens a project, and you can optionally specify a revision name. The example checks whether the specified project exists. If it does, the example opens the current revision, or the revision you specify.

### Example 3–17. Full-Featured Method to Open Projects

```tcl
package require cmdline
variable ::argv0 $::quartus(args)
set options { 
  { "project.arg" "Project Name" } 
  { "revision.arg" "Revision Name" } 
}
array set optshash [::cmdline::getoptions ::argv0 $options]

# Ensure the project exists before trying to open it
if { ![project_exists $optshash(project)]} {
  if { ![string equal "" $optshash(revision)]} {
    # There is no revision name specified, so default
    # to the current revision
    project_open $optshash(project) -current_revision
  } else {
    # There is a revision name specified, so open the
    # project with that revision
    project_open $optshash(project) -revision $optshash(revision)
  } else {
    puts "Project $optshash(project) does not exist"
    exit 1
  }
}
# The rest of your script goes here
```

If you do not require this flexibility or error checking, you can use just the `project_open` command, as shown in Example 3–18.

### Example 3–18. Simple Method to Open Projects

```tcl
set proj_name [lindex $argv 0]
project_open $proj_name
```
The quartus() Array

The scripts in the preceding examples parsed command line arguments found in quartus(args). The global quartus() Tcl array includes other information about your project and the current Quartus II executable that might be useful to your scripts. For information on the other elements of the quartus() array, type the following command at a Tcl prompt:

```
help -quartus
```

The Quartus II Tcl Shell in Interactive Mode

This section presents how to make project assignments and then compile the finite impulse response (FIR) filter tutorial project with the quartus_sh interactive shell. This example assumes that you already have the fir_filter tutorial design files in a project directory.

To begin, type the following at the system command prompt to run the interactive Tcl shell:

```
quartus_sh -s
```

Create a new project called fir_filter, with a revision called filtref by typing the following command at a Tcl prompt:

```
project_new -revision filtref fir_filter
```

If the project file and project name are the same, the Quartus II software gives the revision the same name as the project.

Because the revision named filtref matches the top-level file, all design files are automatically picked up from the hierarchy tree.

Next, set a global assignment for the device with the following command:

```
set_global_assignment -name family Cyclone
```

To learn more about assignment names that you can use with the -name option, refer to Quartus II Help.

For assignment values that contain spaces, enclose the value in quotation marks.

To quickly compile a design, use the ::quartus::flow package, which properly exports the new project assignments and compiles the design with the proper sequence of the command-line executables. First, load the package:

```
load_package flow
```

It returns the following:

```
1.0
```

To perform a full compilation of the FIR filter design, use the execute_flow command with the -compile option:

```
execute_flow -compile
```
This command compiles the FIR filter tutorial project, exporting the project assignments and running `quartus_map`, `quartus_fit`, `quartus_asm`, and `quartus_sta`. This sequence of events is the same as selecting **Start Compilation** from the Processing menu in the Quartus II GUI.

When you are finished with a project, close it with the `project_close` command as shown in **Example 3–19**.

**Example 3–19.**

```
project_close
```

To exit the interactive Tcl shell, type `exit` at a Tcl prompt.

---

**The tclsh Shell**

On the UNIX and Linux operating systems, the tclsh shell included with the Quartus II software is initialized with a minimal `PATH` environment variable. As a result, system commands might not be available within the tclsh shell because certain directories are not in the `PATH` environment variable. To include other directories in the path searched by the tclsh shell, set the `QUARTUS_INIT_PATH` environment variable before running the tclsh shell. Directories in the `QUARTUS_INIT_PATH` environment variable are searched by the tclsh shell when you execute a system command.

---

**Tcl Scripting Basics**

The core Tcl commands support variables, control structures, and procedures. Additionally, there are commands for accessing the file system and network sockets, and running other programs. You can create platform-independent graphical interfaces with the Tk widget set.

Tcl commands are executed immediately as they are typed in an interactive Tcl shell. You can also create scripts (including the examples in this chapter) in files and run them with the Quartus II executables or with the tclsh shell.

---

**Hello World Example**

The following shows the basic “Hello world” example in Tcl:

```
puts "Hello world"
```

Use double quotation marks to group the words `hello` and `world` as one argument. Double quotation marks allow substitutions to occur in the group. Substitutions can be simple variable substitutions, or the result of running a nested command, described in “**Substitutions**” on page 3–19. Use curly braces `{}` for grouping when you want to prevent substitutions.
Variables

Assign a value to a variable with the `set` command. You do not have to declare a variable before using it. Tcl variable names are case-sensitive. Example 3–20 assigns the value 1 to the variable named `a`.

**Example 3–20. Assigning Variables**

```
set a 1
```

To access the contents of a variable, use a dollar sign ("\$") before the variable name. Example 3–21 prints "Hello world" in a different way.

**Example 3–21. Accessing Variables**

```
set a Hello
set b world
puts "$a $b"
```

Substitutions

Tcl performs three types of substitution:

- Variable value substitution
- Nested command substitution
- Backslash substitution

**Variable Value Substitution**

Variable value substitution, as shown in Example 3–21, refers to accessing the value stored in a variable with a dollar sign ("\$") before the variable name.

**Nested Command Substitution**

Nested command substitution refers to how the Tcl interpreter evaluates Tcl code in square brackets. The Tcl interpreter evaluates nested commands, starting with the innermost nested command, and commands nested at the same level from left to right. Each nested command result is substituted in the outer command. Example 3–22 sets `a` to the length of the string `foo`.

**Example 3–22. Command Substitution**

```
set a [string length foo]
```
Backslash Substitution

Backslash substitution allows you to quote reserved characters in Tcl, such as dollar signs (“$”) and braces (“{ }”). You can also specify other special ASCII characters like tabs and new lines with backslash substitutions. The backslash character is the Tcl line continuation character, used when a Tcl command wraps to more than one line. Example 3–23 shows how to use the backslash character for line continuation.

Example 3–23. Backslash Substitution

```
set this_is_a_long_variable_name [string length "Hello \ world."]
```

Arithmetic

Use the `expr` command to perform arithmetic calculations. Use curly braces (“{ }”) to group the arguments of this command for greater efficiency and numeric precision. Example 3–24 sets `b` to the sum of the value in the variable `a` and the square root of 2.

Example 3–24. Arithmetic with the `expr` Command

```
set a 5
set b [expr { $a + sqrt(2) }]
```

Tcl also supports boolean operators such as `&&` (AND), `||` (OR), `!` (NOT), and comparison operators such as `<` (less than), `>` (greater than), and `==` (equal to).

Lists

A Tcl list is a series of values. Supported list operations include creating lists, appending lists, extracting list elements, computing the length of a list, sorting a list, and more. Example 3–25 sets `a` to a list with three numbers in it.

Example 3–25. Creating Simple Lists

```
set a { 1 2 3 }
```

You can use the `lindex` command to extract information at a specific index in a list. Indexes are zero-based. You can use the index `end` to specify the last element in the list, or the index `end-<<n>>` to count from the end of the list. Example 3–26 prints the second element (at index 1) in the list stored in `a`.

Example 3–26. Accessing List Elements

```
puts [lindex $a 1]
```

The `llength` command returns the length of a list. Example 3–27 prints the length of the list stored in `a`.

Example 3–27. List Length

```
puts [llength $a]
```
The `lappend` command appends elements to a list. If a list does not already exist, the list you specify is created. The list variable name is not specified with a dollar sign (“$”). Example 3–28 appends some elements to the list stored in a.

**Example 3–28. Appending to a List**

```
lappend a 4 5 6
```

**Arrays**

Arrays are similar to lists except that they use a string-based index. Tcl arrays are implemented as hash tables. You can create arrays by setting each element individually or with the `array set` command. To set an element with an index of `Mon` to a value of `Monday` in an array called `days`, use the following command:

```
set days(Mon) Monday
```

The `array set` command requires a list of index/value pairs. This example sets the array called `days`:

```
array set days { Sun Sunday Mon Monday Tue Tuesday 
               Wed Wednesday Thu Thursday Fri Friday Sat Saturday }
```

Example 3–29 shows how to access the value for a particular index.

**Example 3–29. Accessing Array Elements**

```
set day_abbreviation Mon
puts $days($day_abbreviation)
```

Use the `array names` command to get a list of all the indexes in a particular array. The index values are not returned in any specified order. Example 3–30 shows one way to iterate over all the values in an array.

**Example 3–30. Iterating Over Arrays**

```
foreach day [array names days] {
    puts "The abbreviation $day corresponds to the day \ 
         name $days($day)"
}
```

Arrays are a very flexible way of storing information in a Tcl script and are a good way to build complex data structures.
Control Structures

Tcl supports common control structures, including if-then-else conditions and for, foreach, and while loops. The position of the curly braces as shown in the following examples ensures the control structure commands are executed efficiently and correctly. Example 3–31 prints whether the value of variable a positive, negative, or zero.

**Example 3–31. If-Then-Else Structure**

```tcl
if { $a > 0 } {
    puts "The value is positive"
} elseif { $a < 0 } {
    puts "The value is negative"
} else {
    puts "The value is zero"
}
```

Example 3–32 uses a for loop to print each element in a list.

**Example 3–32. For Loop**

```tcl
set a { 1 2 3 }
for { set i 0 } { $i < [llength $a] } { incr i } {
    puts "The list element at index $i is [lindex $a $i]"
}
```

Example 3–33 uses a foreach loop to print each element in a list.

**Example 3–33. foreach Loop**

```tcl
set a { 1 2 3 }
foreach element $a {
    puts "The list element is $element"
}
```

Example 3–34 uses a while loop to print each element in a list.

**Example 3–34. while Loop**

```tcl
set a { 1 2 3 }
set i 0
while { $i < [llength $a] } {
    puts "The list element at index $i is [lindex $a $i]"
    incr i
}
```

You do not have to use the expr command in boolean expressions in control structure commands because they invoke the expr command automatically.
Procedures

Use the proc command to define a Tcl procedure (known as a subroutine or function in other scripting and programming languages). The scope of variables in a procedure is local to the procedure. If the procedure returns a value, use the return command to return the value from the procedure. Example 3–35 defines a procedure that multiplies two numbers and returns the result.

Example 3–35. Simple Procedure

```tcl
proc multiply { x y } {
    set product [expr { $x * $y }]
    return $product
}
```

Example 3–36 shows how to use the multiply procedure in your code. You must define a procedure before your script calls it.

Example 3–36. Using a Procedure

```tcl
proc multiply { x y } {
    set product [expr { $x * $y }]
    return $product
}
set a 1
set b 2
puts [multiply $a $b]
```

Define procedures near the beginning of a script. If you want to access global variables in a procedure, use the global command in each procedure that uses a global variable. Example 3–37 defines a procedure that prints an element in a global list of numbers, then calls the procedure.

Example 3–37. Accessing Global Variables

```tcl
proc print_global_list_element { i } {
    global my_data
    puts "The list element at index $i is [lindex $my_data $i]"
}
set my_data { 1 2 3}
print_global_list_element 0
```

File I/O

Tcl includes commands to read from and write to files. You must open a file before you can read from or write to it, and close it when the read and write operations are done. To open a file, use the open command; to close a file, use the close command. When you open a file, specify its name and the mode in which to open it. If you do not specify a mode, Tcl defaults to read mode. To write to a file, specify w for write mode as shown in Example 3–38.

Example 3–38. Open a File for Writing

```tcl
set output [open myfile.txt w]
```
Tcl supports other modes, including appending to existing files and reading from and writing to the same file.

The `open` command returns a file handle to use for read or write access. You can use the `puts` command to write to a file by specifying a filehandle, as shown in Example 3–39.

**Example 3–39. Write to a File**

```tcl
set output [open myfile.txt w]
puts $output "This text is written to the file."
close $output
```

You can read a file one line at a time with the `gets` command. Example 3–40 uses the `gets` command to read each line of the file and then prints it out with its line number.

**Example 3–40. Read from a File**

```tcl
set input [open myfile.txt]
set line_num 1
while { [gets $input line] >= 0 } {
    # Process the line of text here
    puts "$line_num: $line"
    incr line_num
}
close $input
```

**Syntax and Comments**

Arguments to Tcl commands are separated by white space, and Tcl commands are terminated by a newline character or a semicolon. As shown in “Substitutions” on page 3–19, you must use backslashes when a Tcl command extends more than one line.

Tcl uses the hash or pound character (`#`) to begin comments. The `#` character must begin a comment. If you prefer to include comments on the same line as a command, be sure to terminate the command with a semicolon before the `#` character. Example 3–41 is a valid line of code that includes a `set` command and a comment.

**Example 3–41. Comments**

```tcl
set a 1;# Initializes a
```

Without the semicolon, it would be an invalid command because the `set` command would not terminate until the new line after the comment.

The Tcl interpreter counts curly braces inside comments, which can lead to errors that are difficult to track down. Example 3–42 causes an error because of unbalanced curly braces.

**Example 3–42. Unbalanced Braces in Comments**

```tcl
# if { $x > 0 } {
if { $y > 0 } {
    # code here
}
```
External References

For more information about Tcl, refer to the following sources:

- Practical Programming in Tcl and Tk, Brent B. Welch
- Tcl and the TK Toolkit, John Ousterhout
- Effective Tcl/Tk Programming, Michael McLennan and Mark Harrison
- Quartus II Tcl example scripts at www.altera.com/support/examples/tcl/tcl.html
- Tcl Developer Xchange at tcl.activestate.com

Document Revision History

Table 3–4 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update, Updated supported version of Tcl in the section “Tool Command Language” on page 3–2, minor editorial changes</td>
</tr>
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<td>Removed LogicLock example, Added the incremental_compilation, insystem_source_probe, and rtl packages to Table 3-1 and Table 3-2, Added quartus_map to table 3-2.</td>
</tr>
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<td>9.0.0</td>
<td>Removed the “EDA Tool Assignments” section, Added the section “Compile All Revisions” on page 3–9, Added the section “Using the tclsh Shell” on page 3–20</td>
</tr>
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<td>Changed to 8½” × 11” page size. No change to content.</td>
</tr>
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<td>Updated references.</td>
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</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This section provides an overview of the I/O planning process, Altera FPGA pin terminology, as well as the various methods for importing, exporting, creating, and validating pin-related assignments using the Quartus® II software. This section also describes ways to use the Quartus II software to analyze signal integrity, including simultaneous switching noise, as well as interfaces with third-party PCB design tools.

This section includes the following chapters:

- **Chapter 4, I/O Management**
  This chapter provides an overview of the I/O planning process, Altera FPGA pin terminology, and the various methods for importing, exporting, creating, and validating pin-related assignments.

- **Chapter 5, Simultaneous Switching Noise (SSN) Analysis and Optimizations**
  This chapter describes the tools in the Quartus II software that allow you to estimate the SSN performance of your design both early in the design cycle and when your PCB is complete.

- **Chapter 6, Signal Integrity Analysis with Third-Party Tools**
  This chapter is intended for logic designers and board designers, and describes simulation and how to adjust designs to improve board-level timing and signal integrity. Also included is information about how to create accurate models of your design with the Quartus II software for use in simulation software.

- **Chapter 7, Mentor Graphics PCB Design Tools Support**
  This chapter discusses how the Quartus II software interacts with the Mentor Graphics I/O Designer software and the DxDesigner software to provide a completely cyclical FPGA-to-board integration design workflow.

- **Chapter 8, Cadence PCB Design Tools Support**
  This chapter addresses how the Quartus II software interacts with the Cadence Allegro Design Entry HDL software and the Allegro Design Entry CIS (Component Information System) software (also known as OrCAD Capture CIS) to provide a complete FPGA-to-board integration design workflow.
4. I/O Management

This chapter describes efficient management of the I/O pins in your target device. You must consider I/O standards, pin placement guidelines, and hardware capabilities in making pin-related assignments. You must begin I/O planning and PCB development early in the development cycle.

The Quartus II Pin Planner, Pin Advisor, and I/O assignment analysis features support early I/O planning and assignments. You can use the Pin Planner shown in Figure 4–1 to plan I/O assignments early, to create and validate pin assignments, and to generate a Pin-Out File (.pin) for use with third-party PCB tools. The Pin Advisor suggests pin planning steps. Use I/O assignment analysis to validate your I/O pin assignments.

Figure 4–1. Pin Planner
Understanding Altera Pin Terminology

Altera offers devices in a variety of package types. To describe pin terminology, this chapter uses a wire bond ball-grid array (BGA) package in its examples. On the top surface of the silicon die, there is a ring of bond pads that connect to the silicon to the I/O pins. In a wire bond BGA package, the device is placed in the package and copper wires connect the bond pads to the solder balls of the package. Figure 4–2 shows a cross section of a wire bond BGA package.

For more information about the BGA packages available for each Altera device, refer to the *Altera Device Package Information Data Sheet*.

**Figure 4–2. Wire Bond BGA**

![Wire Bond BGA Diagram](image)

**Package Pins**

BGA package pins are small solder balls arranged in a grid pattern on the bottom of the package. The Quartus II software identifies each pin with alphanumeric pin numbers using a coordinate system of letters and numbers corresponding with the pin row and column location.
Figure 4–3 shows the coordinate system to identify pin locations. The top row of pins is labeled A and continues alphabetically as you move down. The left-most column of pins is labeled 1 and increments by one as you move right. For example, pin number B4 represents the pin located in row B and column 4. The letters I, O, Q, S, X, and Z are never used in pin numbers. If the device contains more rows than letters of the alphabet, the alphabet is repeated, prefixed with the letter A.

Figure 4–3. Row and Column Labeling

For more information about the pin numbers for your Altera device, refer to the Pin-Out Files for Altera Devices page of the Altera website.

Pads

Package pins are connected to pads located on the perimeter of the top metal layer of the silicon die (refer to Figure 4–2). Figure 4–4 shows the numbering scheme for pads on the device. Each pad is identified by a pad ID, which is numbered starting at 0, and increments by one in a counterclockwise direction around the device.

Figure 4–4. Pad Number Ordering

To prevent signal integrity issues, the Quartus II software uses pin placement rules to validate your pin placements and pin-related assignments. You must understand the pad locations to which your pins were assigned, because some pin placement rules describe pad placement restrictions. For example, to ensure signal integrity, in certain devices there is a restriction on the number of I/O pins supported by a voltage reference (VREF) pad. There are also restrictions on the number of pads between single-ended input or output pins and a differential pin. The Quartus II software performs pin placement analysis. Design compilation fails and the Quartus II software reports an error if your design violates pin placement rules.

For more information about pin placement guidelines, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.
I/O Banks

I/O pins are organized into I/O banks designed to facilitate various supported I/O standards. Each I/O bank is numbered and has its own voltage source pins, called VCCIO, to offer the highest I/O performance. Depending on the device and I/O standards for the pins in the I/O bank, the specified voltage of the VCCIO pin is between 1.5 V and 3.3 V. Each I/O bank can support multiple pins with different I/O standards, however the pins must use the same VCCIO signal.

Figure 4–5 shows the I/O banks of a Stratix® II device. The pins in the I/O banks on the left and right side support high-speed I/O standards such as LVDS, whereas the pins on the top and bottom I/O banks support all single-ended I/O standards, including data strobe signaling (DQS). Pins belonging to the same I/O bank must use the same VCCIO signal.

For more information about the capabilities of each I/O bank, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

Figure 4–5. Stratix II I/O Banks  (1), (2), (3), (4)
VREF Groups

A VREF group is a group of pins that includes one dedicated VREF pin as required by voltage-referenced I/O standards. A VREF group is made up of a small number of pins, as compared to the I/O bank, to maintain the signal integrity of the VREF pin. One or more VREF groups exist in an I/O bank. The pins in a VREF group share the same VCCIO and VREF voltages.

For more information about I/O banks, VREF groups, and supported I/O standards, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

I/O Planning Overview

I/O planning involves anticipating and allocating I/O pins to match your device and PCB. The Quartus II Pin Planner, Pin Advisor, and I/O assignment analysis features assist you in I/O planning. Figure 4–6 shows the recommended design flow for creating I/O pin assignments for a design. I/O planning includes the following tasks:

1. Selecting a device that meets your logic and I/O requirements, based on the supported I/O standards for the device, I/O bank structure, supply voltage requirements such as VREF and VCCIO requirements in I/O banks, available pins for user I/O, power supply requirements, and other design requirements.

2. Preparing your design files. Design files contain the top-level ports or top-level interface information. You can click Early Pin Planning in the Task window (refer to Figure 4–6) to generate a top-level HDL wrapper file before defining design logic.

3. Importing any existing assignments from a Tcl Script File (.tcl), Comma-Separated Value File (.csv), or Quartus II Settings File (.qsf).

4. Creating I/O pin assignments to match your device and PCB, including location assignments, I/O standard, output loading, slew rate, and current strength assignments.

5. Validating your pin-related assignments. You can perform preliminary I/O assignment validation while creating the assignments with the live I/O check feature; you can perform a more a thorough validation with the I/O assignment analysis feature; and finally you can perform complete I/O assignment validation by running the Fitter with timing constraints.

6. Generating a validated .pin file for third-party PCB tools.
4–6 Chapter 4: I/O Management
I/O Planning Overview

Figure 4–6. Quartus II Software I/O Planning Flow

Notes to Figure 4–6:
(1) Use the live I/O check feature in the Pin Planner to dynamically validate pin assignments.
(2) To create the FPGA Xchange file (.fx), on the Processing menu, point to Start and click EDA Netlist Writer. The .pin is created at the <project_dir> level. The .fx is created at the <project_dir/board/> level.
(3) You must complete your design files and constraints before you begin full compilation. To learn how to create I/O timing constraints, refer to The Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.
(4) For more information, refer to the Area and Timing Optimization chapter in volume 2 of the Quartus II Handbook.
Selecting a Device

Before you begin I/O planning you must select a device family that supports the I/O resources, standards, clocking schemes required for your design.

For more information, refer to the various device handbooks available on the Literature and Technical Documentation page of the Altera website.

Working with Third-Party PCB Tools

If you have not designed your PCB, you must first create and validate your I/O assignments in the Quartus II software and then export the assignments to the PCB tool. If your PCB is partially designed, you must first create your device assignments in your PCB tool, and then import the assignments into the Quartus II software for validation. Figure 4–7 shows the recommended I/O planning flow to ensure that your pin assignments are valid.

Currently, only the Mentor Graphics® I/O Designer PCB tool and the Cadence Allegro PCB tool are supported in this reverse I/O planning flow.

Figure 4–7. I/O Planning Flow Using an FPGA Xchange File from a PCB Tool

Early I/O Planning with the Pin Planner

The Pin Planner allows you to plan early and assign nodes not yet defined in the design HDL or schematic, including interface IP core signals. The top-level file of your design instantiates the next level of hierarchy and includes port names and type. Top-level design files often specify interfaces for memory, high-speed I/O, device configuration, and debugging tools. Example 4–1 shows an example top-level Verilog HDL file that lists design input and output ports.
Click **Early Pin Planning** in the Task window to define or import interface IP core nodes or other user nodes not yet defined in the design. You can then generate a new top-level design that preserves the I/O pin assignments.

**Example 4–1. Top-Level Design File**

```verilog
module top (    clk_in,
                rst,
                a,
                z,
                b,
                c_in,
                d,
                e);

  input clk_in;
  input rst_;
  input a;
  input z;
  input [7:0] b;
  input [7:0] c_in;
  input [7:0] d;
  output reg [7:0] e;

  /* Instantiations of sub blocks */

endmodule
```

Early pin planning helps save time by accurately setting up critical I/O before implementing the rest of the design.

The following sections describe the typical steps of the early I/O planning flow:

- “Assigning Interface I/O in the Pin Planner”
- “Adding and Connecting Nodes” on page 4–9
- “Setting Up and Creating the Top-Level Design File” on page 4–10

For more information, refer to **Assigning Device I/O Pins in Pin Planner** in Quartus II Help.

**Assigning Interface I/O in the Pin Planner**

You can use the Pin Planner to assign signals to interface IP core ports not yet defined in your design. For example, if you know that your design requires a PHY to interface with high speed transceivers, you can define the parameters for a transceiver PHY IP core, and then assign signals to the interface IP nodes in the Pin Planner. Adding interface information directly in the Pin Planner allows you to assign required pins before manually defining the logic in a design file. Click **Early Pin Planning** in the Task window to create or import IP core or user nodes.
You can create complex interfaces with early pin planning. Figure 4–8 shows a schematic representation of the DDIO_IN and DDIO_OUT megafunctions defined in the Pin Planner. The ports in gray are internal ports, which are connected later in the design.

Figure 4–8. Connections Between Input and Output Megafunctions and User Nodes

Adding and Connecting Nodes

After you instantiate an IP core in your design, you can set up the user nodes in your design and customize information such as port direction and type. Before you create a top-level design file, you must connect the user nodes and ports to each other and to the rest of the design.

The user nodes you specify as user nodes are declared as virtual pins when you generate the top-level design file. The Pin Planner makes virtual pin assignments to user nodes so that they are not assigned to device pins. These virtual pins are not shown in the All Pins list or Groups list in the Pin Planner because they are not actual external ports of the design. Any new nodes you add to the design are declared as ports in the top-level design file and are shown in the All Pins list and Groups list.
Figure 4–9 shows the early pin planning interface.

**Figure 4–9. Connecting a User Node to an Interface Port**

---

**Setting Up and Creating the Top-Level Design File**

You can create a top-level design file after you create pin connections and add or modify user nodes and IP core nodes with the Pin Planner. If the internal logic is incomplete, generating the top-level design file allows you to validate your I/O assignments and provides a basis on which to build the rest of your design.

You must update the top-level design file whenever you change the top-level ports of the design, including any node changes made in the **Set Up Top-Level File** tab.
Example 4–2 shows a sample of a top-level HDL wrapper file representing the design in Figure 4–8.

Example 4–2. HDL Wrapper File Generated with the Early I/O Planning Flow

```verilog
module top
(
    reset,
    input_data,
    clk,
    output_data,
    ddio_in_dataout_h, // Internal
    ddio_in_dataout_l, // Internal
    ddio_out_datain_h, // Internal
    ddio_out_datain_l // Internal
);

input reset;
input[7:0] input_data;
input clk;
output[7:0] output_data;

output[7:0] ddio_in_dataout_h /* synthesis altera_attribute="-name VIRTUAL_PIN ON" */;
output[7:0] ddio_in_dataout_l /* synthesis altera_attribute="-name VIRTUAL_PIN ON" */;
input[7:0] ddio_out_datain_h /* synthesis altera_attribute="-name VIRTUAL_PIN ON" */;
input[7:0] ddio_out_datain_l /* synthesis altera_attribute="-name VIRTUAL_PIN ON" */;

ddio_in ddio_in_inst
(
    .aclr(reset),
    .datain(input_data),
    .inclock(clk),
    .dataout_h(ddio_in_dataout_h),
    .dataout_l(ddio_in_dataout_l)
);

ddio_out ddio_out_inst
(
    .aclr(reset),
    .datain_h(ddio_out_datain_h),
    .datain_l(ddio_out_datain_l),
    .outclock(clk),
    .dataout(output_data)
);
endmodule
```

After you generate the top-level design file and compile the design, use I/O assignment analysis as described in “Validating Pin Assignments with I/O Assignment Analysis” on page 4–25 and continue with your design flow by modifying or creating pin assignments with the Pin Planner.

Importing and Exporting Pin Assignments

If you have existing pin assignments in .tcl, .csv, .qsf, .fx, or .pin format from a different Quartus II project or from third-party PCB tools, you can transfer these assignments between the Quartus II software and other tools.
Importing and Exporting Assignments with the Quartus II Software

You can import and export pin-related assignments contained in .tcl, .csv, and .qsf files with the Quartus II software. You can import assignments from .tcl, .csv, and .qsf files and then view and change the assignments in the Pin Planner.

When you create pin assignments with the Pin Planner all your pin-related assignments are written to the .qsf as Tcl commands. You can import and export .qsf files between projects to transfer all assignments, including pin assignments.

Importing and Exporting Assignments with Third-Party PCB Tools

After creating and validating your pin assignments, you can export device and pin-related information from the Quartus II software to third-party PCB tools for board development. You must generate a .pin with the Quartus II software to export the correct pin locations and other important pin information.

The device .pin files display the pin name and pin number as well as detailed properties about each pin. Table 4–1 describes the columns in a .pin.

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pin Name/Usage</td>
<td>The name of the design pin, or whether the pin is GND or VCC pin</td>
</tr>
<tr>
<td>Location</td>
<td>The pin number of the location on the device package</td>
</tr>
<tr>
<td>Dir</td>
<td>The direction of the pin</td>
</tr>
<tr>
<td>I/O Standard</td>
<td>The name of the I/O standard to which the pin is configured</td>
</tr>
<tr>
<td>Voltage</td>
<td>The voltage level that is required to be connected to the pin</td>
</tr>
<tr>
<td>I/O Bank</td>
<td>The I/O bank to which the pin belongs</td>
</tr>
<tr>
<td>User Assignment</td>
<td>Y or N indicating if the location assignment for the design pin was user assigned (Y) or assigned by the Fitter (N)</td>
</tr>
</tbody>
</table>

For more information, refer to the Pin-Out Files for Altera Devices page of the Altera website.
To transfer device and pin-related information between the Quartus II software and the Mentor Graphics I/O Designer software you must generate an .fx. Importing assignments into the I/O Designer software requires both the .fx and a .pin generated by the Quartus II software. However, the Quartus II software requires only the .fx to import pin assignments back from the I/O Designer software.

For more information about the I/O Designer software and the DxDesigner interface, refer to the Mentor Graphics PCB Tools Support chapter in volume 2 of the Quartus II Handbook.

Creating Pin-Related Assignments

A pin-related assignment is any assignment applied to a top-level pin. For example, a pin location assignment assigns a top-level port or node to a pin number (location) on the targeted device. Other examples of pin-related assignments include assigning an I/O standard, assigning drive strength, or assigning a slew rate.

When you do not have complete information for all the top-level pins, you can reserve certain device pins temporarily until the I/O pins are defined in your design files. You can use reserved pins in the future, but they do not perform a function in your design. Reserved pins require a unique pin name and pin location. Using reserved pins as place holders for future design pins increases the accuracy of I/O assignment analysis.

Table 4–2 describes the tools and features in the Quartus II software for creating reserved pins and other pin-related assignments. Each tool and feature is described in more detail in the following sections. Altera recommends that you use the Pin Planner to create and edit pin-related assignments; however, depending on your design flow, you may find some of the other tools useful for working with pin-related assignments.

Table 4–2. Overview of Quartus II Tools and Features to Create Pin-Related Assignments (Part 1 of 2)

<table>
<thead>
<tr>
<th>Feature</th>
<th>Overview</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Pin Planner</strong></td>
<td>■ Create pin location assignments to one or more node names by dragging and dropping unassigned pins into the package view</td>
</tr>
<tr>
<td></td>
<td>■ Edit pin location assignments for one or more node names by dragging and dropping groups of pins in the package view</td>
</tr>
<tr>
<td></td>
<td>■ Visually analyze pin resources in the package view</td>
</tr>
<tr>
<td></td>
<td>■ Display various I/O pin types quickly</td>
</tr>
<tr>
<td></td>
<td>■ View the function of package pins with the Pin Legend window</td>
</tr>
<tr>
<td></td>
<td>■ Highlight various I/O resources</td>
</tr>
<tr>
<td></td>
<td>■ Make correct pin location decisions by referring to the Pad View window</td>
</tr>
<tr>
<td></td>
<td>■ Create, import, and edit IP cores for early I/O planning</td>
</tr>
<tr>
<td></td>
<td>■ Generate a top-level wrapper file without design files based on early I/O assignments</td>
</tr>
<tr>
<td></td>
<td>■ Configure board trace model assignments, instead of capacitive loading assignments, to generate Advanced I/O Timing results</td>
</tr>
<tr>
<td><strong>Tcl Scripts</strong></td>
<td>■ Create any pin-related assignments for multiple pins</td>
</tr>
<tr>
<td></td>
<td>■ Store and reapply all pin-related assignments with Tcl scripts</td>
</tr>
<tr>
<td></td>
<td>■ Create assignments from the command line</td>
</tr>
</tbody>
</table>
Creating Pin Assignments With the Pin Planner

During I/O planning, it can be cumbersome to try to correlate pin numbers with their relative location on the package and their pin properties. The Quartus II Pin Planner helps you visualize, plan, and assign device I/O pins in a graphical view of the target device package, as shown in Figure 4–1. You can quickly locate various I/O pins and assign them design elements or other properties to ensure compatibility with your PCB layout. The Quartus II software uses the assignments to place and route your design during device programming. The Pin Planner can also help with early pin planning by allowing you to plan for and assign nodes not yet defined in the design.

The Task window provides one-click access to common pin planning tasks. After clicking a task, you can view and highlight the results in the Report window. Use the Pin Planner to quickly identify VREF groups, edges, DQ/DQS pins, hard memory interface pins, PCIe hard IP interface pins, hard processor system pins, I/O modules, I/O banks, clock regions, and differential pin pairings. When deciding on a pin location, use the Pin Planner to gather information about available resources, as well as the functionality of each individual pin, I/O bank, and VREF group.

For more information, refer to Assigning Device I/O Pins in Pin Planner in Quartus II Help.

Finding Compatible Pin Locations with the Pin Finder

As device pin-counts and I/O capabilities continue to increase, it becomes more difficult to understand the capabilities of each I/O pin and to correctly assign your design I/Os. As you edit pin assignments to help your design fit or to achieve timing goals, you can use the tools in the Pin Planner to help you find appropriate pin locations for your design I/O nodes.

Verifying Pin Migration Compatibility

You can use the Pin Migration View window in Pin Planner to assist you in verifying whether your pin assignments migrate to a different device successfully. You can vertically migrate to a device with a different density while using the same device package, or migrate between packages with different densities and ball counts.

For more information, refer to Viewing Pin Migration Compatibility in Quartus II Help.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Overview</th>
</tr>
</thead>
<tbody>
<tr>
<td>Chip Planner</td>
<td>■ Create and change pin locations by dragging and dropping pins into the floorplan</td>
</tr>
<tr>
<td></td>
<td>■ Make correct pin location decisions by referring to the pad ID number and spacing</td>
</tr>
<tr>
<td></td>
<td>■ Display I/O banks, VREF groups, and differential pin pairing information</td>
</tr>
<tr>
<td>Synthesis Attributes</td>
<td>■ Embed pin-related assignments with attributes in the design files to pass assignments to the Quartus II software</td>
</tr>
</tbody>
</table>
When you select migration devices early in the design process, the Pin Planner displays only the pins that are available in the current device and in all migration devices. If you select migration devices later in your design cycle, there may be assignments for I/O nodes in your original design that do not have corresponding pins in a migration device. If no corresponding pin exists, the Compiler cannot honor the assignment and an error occurs when you try to recompile the design.

The Pin Migration View window helps you identify the difference in pins that can exist between migration devices. For example, Figure 4–10 shows the highlighted pin AC24 existed in the target EP2S30 device, but does not exist in one of the migration devices, resulting in a No Connect (NC).

![Figure 4–10. Pin Migration View](image)

The migration result for the pin function of highlighted PIN_AC23 is not an NC but a voltage reference VREFB1N2 even though the pin is an NC in one of the migration devices. VREF standards have a higher priority than an NC, thus the migration result display the voltage reference. Even if you do not use that pin for a port connection in your design, you must use the VREF standard for I/O standards that require it on the actual board for the migration device.

If one of the migration devices has pins intended for connection to VCC or GND and these same pins are I/O pins on a different device in the migration path, the Quartus II software ensures these pins are not used for I/O. Ensure that these pins are connected to the correct PCB plane.
Creating Pin Assignments With the Pin Planner

When migrating between two devices in the same package, pins that are not connected to the smaller die may be intended to connect to $V_{CC}$ or GND on the larger die. To facilitate migration, you can connect these pins to $V_{CC}$ or GND in your original design because the pins are not physically connected to the smaller die.

For more information about migration, refer to *AN90: SameFrame Pin-Out Design for FineLine BGA Packages*. For more information about designing for HardCopy series devices, refer to the *Designing HardCopy Series Devices* chapter in volume 1 of the *Quartus II Handbook*.

**Viewing Simultaneous Switching Noise (SSN) Results**

You can perform and SSN analysis of your design to estimate the voltage noise for each pin in the design. You can view the SSN results in the Pin Planner and adjust your I/O assignments to avoid potential signal integrity issues.

For more information about running the SSN Analyzer and viewing the results in the Pin Planner, refer to *Running the SSN Analyzer* in Quartus II Help.

For more information about the SSN Analyzer, refer to the *Simultaneous Switching Noise (SSN) Analysis and Optimization* chapter in volume 2 of the *Quartus II Handbook*.

**Creating Location Assignments**

You can create the following types of location assignments for your design and its reserved pins:

- Pin number
- I/O bank
- VREF group
- Edge

For more information about device support for I/O bank, VREF group, and edge location assignments, refer to *I/O bank, VREF group, and edge* in Quartus II Help.

You can assign your pins to a location with the Pin Planner. It is common to place a group of pins (or bus) with compatible I/O standards in the same I/O bank or VREF group. For example, two buses with two compatible I/O standards, such as 2.5-V and SSTL-II Class I, can be placed in the same I/O bank.
If your design contains a large bus that exceeds the pins available in a particular I/O bank, you can use edge location assignments to place the bus. Edge location assignments improve the circuit board routing ability of large buses, because they are close together near an edge. Figure 4–11 shows Altera device package edges.

**Figure 4–11. Die View and Package View of the Four Edges on an Altera Device**

---

**Creating Exclusive I/O Group Assignments**

You can create exclusive groups comprised of pins by creating and modifying the Exclusive I/O Group logic option with the Pin Planner. When you create exclusive I/O groups in your design and use the Quartus II software to map the signals onto device pins, the Fitter does not place the I/O pins belonging to one exclusive group in an I/O bank if the pins belong to another exclusive I/O group. To understand this, consider an example in which you have a set of signals assigned exclusively to a group called `group_a`, and another set of signals assigned to `group_b`. In both exclusive groups you have pins with different I/O standards. When you create these groups, the Quartus II software maps the pins of both groups in such a way that they are placed in different I/O banks.

**Changing the Slew Rate and Drive Strength**

As part of I/O planning you can set both the slew rate and drive strength of pins. Both slew rate and drive strength affect the outgoing signal integrity of the pin. Depending on your target device family, you can adjust the slew rate of a pin with either the Slew Rate or Slow Slew Rate logic option; you can adjust the drive strength of a pin with the Current Strength logic option. The settings you create for slew rate and drive strength apply during live I/O check, I/O assignment analysis, and full compilation.

For more information about slew rate support, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

If you want to make changes to slew rate or drive strength after you compile your design, you can use the Resource Property Editor to perform engineering change orders (ECOs). ECOs allow you to compile the changes to your design, without changing the synthesis or fitting results.
For more information about making post-compilation changes, refer to *About Post-Compilation Changes* in Quartus II Help.

For more information about ECOs, refer to the *Engineering Change Management with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*. For more information about the effect of I/O settings on signal integrity on the board, refer to *AN 476: Impact of I/O Settings on Signal Integrity in Stratix III Devices*.

### Assigning Locations for Differential Pins

When you use the Pin Planner to assign a differential I/O standard to a single-ended top-level pin in your design, it automatically recognizes the negative pin as part of the differential pin pair assignment and creates the negative pin for you. The Quartus II software writes the location assignment for the negative pin to the `.qsf` file; however, the I/O standard assignment is not added to the `.qsf` for the negative pin of the differential pair.

For example, Figure 4–12 shows a design in which you have a top-level pin defined as `lvds_in` to which you assign a differential I/O standard. The Pin Planner automatically creates the differential pin, `lvds_in(n)`, to complete the differential pin pair.

If you have a single-ended clock that feeds a PLL, assign the pin only to the positive clock pin of a differential pair in the target device. Single-ended pins that feed a PLL and are assigned to the negative clock pin device cause the design to not fit.

For more information about assigning locations for differential pins in HDL code with low-level I/O primitives, refer to “Creating Pin Assignments with Low-Level I/O Primitives” on page 4–22.

### Overriding I/O Placement Rules on Differential Pins

Each device family is constrained by I/O placement rules, as described in Table 4–3 on page 4–22. The pin placement rules ensure that noisy signals do not corrupt neighboring signals. For example, I/O placement rules define the allowed placement of single ended I/O with respect to differential pins, or how many output and bidirectional pins can be placed within a VREF group when using voltage referenced input standards. You can use the IO_MAXIMUM_TOGGLE_RATE assignment to
override I/O placement rules on pins, such as for system reset pins that do not switch during normal design activity. Setting a value of 0 MHz for this assignment causes the Fitter to recognize the pin at a DC state throughout device operation. The Fitter excludes the assigned pin from placement rule analysis. Do not assign an IO_MAXIMUM_TOGGLE_RATE of 0 MHz to any actively switching pin or your design may not function as intended.

Creating Pin Assignments with the Chip Planner

The floorplan of the device shows the pins in the same order as the pads of the device. Understanding the relative distance between a pad and related logic can help you meet your timing requirements. You can view the floorplan of the device in the Chip Planner and determine the distances between user I/O pads and \text{V}_{\text{CC}}, \text{GND}, \text{and } \text{V}_{\text{REF}} pads to avoid signal integrity issues.

For more information about creating pin location assignments in the Chip Planner, refer to \textit{Working with Assignments in the Chip Planner} in Quartus II Help.

For more information about pin placement guidelines, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

Creating Pin Assignments with Tcl Scripts

You can use Tcl scripts to create pin-related assignments as part of a script-based compilation flow. You can enter individual Tcl commands in the Tcl Console window of the Quartus II software. To run a Tcl script with the Quartus II software, type the following command at a system prompt:

\texttt{quartus\_sh \_t \_my\_tcl\_script.tcl}

Example 4–3 shows the \texttt{set\_location\_assignment} and \texttt{set\_instance\_assignment} Tcl commands used to create pin-related assignments to the input pin \texttt{address[10]}.

\textbf{Example 4–3. Tcl Commands to Create Pin-Related Assignments}

\begin{verbatim}
set\_location\_assignment PIN M20 -to address[10] -comment"Address pin to Second FPGA"
set\_instance\_assignment -name IO\_STANDARD "2.5 V" -to address[10]
set\_instance\_assignment -name CURRENT\_STRENGTH\_NEW "MAXIMUM CURRENT" -to address[10]
\end{verbatim}

For more information about creating and running Tcl scripts with the Quartus II software, refer to \textit{Creating and Running Tcl Scripts} in Quartus II Help.

For more information about using Tcl scripts to create pin-related assignments, refer to the \textit{Tcl Scripting} chapter in volume 2 of the \textit{Quartus II Handbook} and \textit{API Functions for Tcl} in Quartus II Help.
Creating Pin Assignments in HDL Code

You can use synthesis attributes or low-level I/O primitives to embed pin-related assignments directly in your HDL code. When you analyze and synthesize your HDL code, the information in the HDL code is converted into the appropriate assignments. There are two ways to specify pin-related assignments with HDL code:

- Using synthesis attributes for signal names that are top-level pins
- Using low-level I/O primitives, such as ALT_BUF_IN, to specify input, output, and differential buffers, and for setting parameters or attributes

Synthesis Attributes

Synthesis attributes allow you to embed pin-related assignments in your HDL code. During Analysis and Synthesis, the Quartus II software reads these synthesis attributes and translates them into assignments. The assignments are then populated in the Pin Planner. If you modify or delete these pin assignments in the Pin Planner and then recompile your design, any changes made in the Pin Planner take precedence over the assignments you made with synthesis attributes in your HDL code. Quartus II integrated synthesis supports the chip_pin, useioff, and altera_attribute synthesis attributes.

For more information about integrated synthesis, synthesis attributes, and syntax, refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook. For more information about synthesis attributes supported by third-party synthesis tools, contact your vendor.

chip_pin and useioff

Use the chip_pin and useioff synthesis attributes to create pin location assignments and to create Fast Input Register, Fast Output Register, and Fast Output Enable Register logic option assignments, respectively, in your HDL code. For all other assignments, including pin-related assignments, use the altera_attribute synthesis attribute as discussed in the “altera_attribute” section.

For more information, refer to useioff VHDL Synthesis Attribute, useioff Verilog HDL Synthesis Attribute, chip_pin VHDL Synthesis Attribute, and chip_pin Verilog HDL Synthesis Attribute in Quartus II Help.
Example 4–4 and Example 4–5 use the chip_pin and useioff attributes to embed location and Fast Input Register logic option assignments in both a Verilog HDL and VHDL design file using the synthesis attributes.

**Example 4–4. Verilog HDL Example**

```verilog
input my_pin1 /* synthesis chip_pin = "C1" useioff = 1 */;
```

**Example 4–5. VHDL Example**

```vhdl
entity my_entity is
    port(
        my_pin1: in std_logic
    );
end my_entity;

architecture rtl of my_entity is
    attribute useioff : boolean;
    attribute useioff of my_pin1 : signal is true;
    attribute chip_pin : string;
    attribute chip_pin of my_pin1 : signal is "C1";
begin -- The architecture body
end rtl;
```

**altera_attribute**

Use the altera_attribute synthesis attribute to create other pin-related assignments in your HDL code. The altera_attribute attribute is understood only by Quartus II integrated synthesis and supports all types of instance assignments.

For more information, refer to altera_attribute VHDL Synthesis Attribute, and altera_attribute Verilog HDL Synthesis Attribute in Quartus II Help.

Example 4–6 and Example 4–7 use the altera_attribute attribute to embed Fast Input Register logic option assignments and I/O standard assignments in both a Verilog HDL and a VHDL design file.

**Example 4–6. Verilog HDL Example**

```verilog
input my_pin1 /* synthesis altera_attribute = "-name FAST_INPUT_REGISTER ON; -name IO_STANDARD \"2.5 V\" = */ */;
```

**Example 4–7. VHDL Example**

```vhdl
entity my_entity is
    port(
        my_pin1: in std_logic
    );
end my_entity;

architecture rtl of my_entity is
begin
attribute altera_attribute : string;
attribute altera_attribute of my_pin1: signal is "-name FAST_INPUT_REGISTER ON;
-- The architecture body
end rtl;
```
Creating Pin Assignments with Low-Level I/O Primitives

You can create pin-related assignments for your top-level nodes using low-level I/O primitives, which allow you to create pin location assignments and set I/O standards, drive strengths, slew rates, and on-chip termination (OCT) value assignments. You can also use low-level differential I/O primitives to define both positive and negative pins of a differential pair in the HDL code for your design.

The pin-related assignments made with primitives do not automatically appear in the Pin Planner. When you use low-level I/O primitives to define various pin-related I/O assignments, the assignments are honored only after you perform a full compilation. After performing a full compilation, you can populate these assignments in the Pin Planner by back-annotating pin assignments.

For more information about back-annotating assignments, refer to Back-Annotating Assignments for A Project in Quartus II Help. For more information about differential I/O primitives, refer to Primitives in Quartus II Help.

For more information about using low-level I/O primitives in your design, refer to the Designing with Low-Level Primitives User Guide.

Validating Pin Assignments

The Quartus II software includes predefined I/O rules to guide you in pin placement and checks your pin-related assignments against these rules during pin planning. You must validate all pin-related assignments in your design. During a full compilation, the Quartus II software does not report illegal pin assignments until the Fitter stage. To preliminarily validate pin-related assignments against the predefined I/O rules, you can use the live I/O check feature or run I/O assignment analysis after performing analysis and synthesis. Typically, the I/O assignment analysis completes quickly. By preliminarily validating your pin-related assignments before fully compiling your design, you can avoid recompiling your design to fix pin-related assignment errors, thus reducing your overall design time. To fully validate pin-related assignments against all I/O timing checks, you must perform a full compilation.

Table 4–3 and Table 4–4 list a subset of the I/O rule checks performed when you run I/O assignment analysis.

For more information about each I/O rule, including which devices support which rules, refer to the device handbooks available on the Literature and Technical Documentation page of the Altera website.

<table>
<thead>
<tr>
<th>Rule</th>
<th>Description</th>
<th>HDL Required?</th>
</tr>
</thead>
<tbody>
<tr>
<td>I/O bank capacity</td>
<td>Checks the number of pins assigned to an I/O bank against the number of pins allowed in the I/O bank.</td>
<td>No</td>
</tr>
<tr>
<td>I/O bank V&lt;sub&gt;CCIO&lt;/sub&gt; voltage compatibility</td>
<td>Checks that no more than one V&lt;sub&gt;CCIO&lt;/sub&gt; is required for the pins assigned to the I/O bank.</td>
<td>No</td>
</tr>
</tbody>
</table>
### Table 4–3. Examples of I/O Rule Checks (Part 2 of 2)

<table>
<thead>
<tr>
<th>Rule</th>
<th>Description</th>
<th>HDL Required?</th>
</tr>
</thead>
<tbody>
<tr>
<td>I/O bank VREF voltage compatibility</td>
<td>Checks that no more than one VREF is required for the pins assigned to the I/O bank.</td>
<td>No</td>
</tr>
<tr>
<td>I/O standard and location conflicts</td>
<td>Checks whether the pin location supports the assigned I/O standard.</td>
<td>No</td>
</tr>
<tr>
<td>I/O standard and signal direction conflicts</td>
<td>Checks whether the pin location supports the assigned I/O standard and direction. For example, certain I/O standards on a particular pin location can only support output pins.</td>
<td>No</td>
</tr>
<tr>
<td>Differential I/O standards cannot have open drain turned on</td>
<td>Checks that open drain is turned off for all pins with a differential I/O standard.</td>
<td>No</td>
</tr>
<tr>
<td>I/O standard and drive strength conflicts</td>
<td>Checks whether the drive strength assignments are within the specifications of the I/O standard.</td>
<td>No</td>
</tr>
<tr>
<td>Drive strength and location conflicts</td>
<td>Checks whether the pin location supports the assigned drive strength.</td>
<td>No</td>
</tr>
<tr>
<td>SUSHOLD and location conflicts</td>
<td>Checks whether the pin location supports SUSHOLD. For example, dedicated clock pins do not support SUSHOLD.</td>
<td>No</td>
</tr>
<tr>
<td>WEAK_PULLUP and location conflicts</td>
<td>Checks whether the pin location supports WEAK_PULLUP (for example, dedicated clock pins do not support WEAK_PULLUP)</td>
<td>No</td>
</tr>
<tr>
<td>Electromigration check</td>
<td>Checks whether combined drive strength of consecutive pads exceeds a certain limit. For example, the total current drive for 10 consecutive pads on a Stratix II device cannot exceed 200 mA.</td>
<td>No</td>
</tr>
<tr>
<td>PCI_IO clamp diode, location, and I/O standard conflicts</td>
<td>Checks whether the pin location along with the I/O standard assigned supports PCI_IO clamp diode.</td>
<td>No</td>
</tr>
<tr>
<td>SERDES and I/O pin location compatibility check</td>
<td>Checks that all pins connected to a SERDES in your design are assigned to dedicated SERDES pin locations.</td>
<td>Yes</td>
</tr>
<tr>
<td>PLL and I/O pin location compatibility check</td>
<td>Checks whether pins connected to a PLL are assigned to the dedicated PLL pin locations.</td>
<td>Yes</td>
</tr>
</tbody>
</table>

### Table 4–4. SSN-Related Rules (Part 1 of 2)

<table>
<thead>
<tr>
<th>Rule</th>
<th>Description</th>
<th>HDL Required?</th>
</tr>
</thead>
<tbody>
<tr>
<td>I/O bank can not have single-ended I/O when DPA exists</td>
<td>Checks that no single-ended I/O pin exists in the same I/O bank as a DPA.</td>
<td>No</td>
</tr>
<tr>
<td>A PLL I/O bank does not support both a single-ended I/O and a differential signal simultaneously</td>
<td>Checks that there are no single-ended I/O pins present in the PLL I/O Bank when a differential signal exists.</td>
<td>No</td>
</tr>
<tr>
<td>Single-ended output is required to be a certain distance away from a differential I/O pin</td>
<td>Checks whether single-ended output pins are a certain distance away from a differential I/O pin.</td>
<td>No</td>
</tr>
<tr>
<td>Single-ended output has to be a certain distance away from a VREF pad</td>
<td>Checks whether single-ended output pins are a certain distance away from a VREF pad.</td>
<td>No</td>
</tr>
<tr>
<td>Single-ended input is required to be a certain distance away from a differential I/O pin</td>
<td>Checks whether single-ended input pins are a certain distance away from a differential I/O pin.</td>
<td>No</td>
</tr>
</tbody>
</table>
Validating Pin Assignments with the Live I/O Check Feature

The live I/O check feature provides live I/O rule checking capability to prevent you from creating pin placements that violate I/O fitting rules. When the live I/O check feature is turned on, pin-related assignment error and warning messages appear immediately in the Quartus II Messages window as you create pin-related assignments in the Pin Planner. This feature enhances your productivity by showing you warnings and errors as you create pin-related assignments so you can immediately correct basic errors before you proceed to the next step in your design flow.

The most basic I/O rules are the I/O buffer rules. The I/O buffer rules checked by the live I/O check feature include:

- $V_{CCIO}$ and $V_{REF}$ voltage compatibility rules
- Electromigration (current density) rules
- Simultaneous Switching Output (SSO) rules
- I/O property compatibility rules, such as drive strength compatibility, I/O standard compatibility, PCI.IO clamp diode compatibility, and I/O direction compatibility

When the live I/O check feature is turned on, the Quartus II software prevents you from assigning pins to unavailable locations. The following are examples of unassignable locations:

- An I/O bank or VREF group with no available pins
- The negative pin of a differential pair if the positive pin of the differential pair is assigned with a node name with a differential I/O standard
- Pin locations that do not support the I/O standard assigned to the selected node name
- For HSTL- and SSTL-type I/O standards, VREF groups of a different $V_{REF}$ voltage than the selected node name

You can turn on or turn off the live I/O check feature at any time. By default, the live I/O check feature is turned off. When the live I/O check feature is turned on, the Quartus II software immediately checks whether your new pin-related assignments pass the basic I/O buffer rules. The Live I/O Check Status window displays the total numbers of errors and warnings while you create and edit pin-related assignments. The Messages window shows detailed messages about any errors or warnings.

Although the live I/O check feature checks all the basic I/O buffer rules, you must run I/O assignment analysis to validate your pin-related assignments against the complete set of I/O system rules.
Validating Pin Assignments with I/O Assignment Analysis

Performing I/O assignment analysis allows you to validate your I/O assignments and surrounding logic for illegal assignments and violations of board layout rules early in the design process. The I/O assignment analysis feature also checks blocks that directly feed or are fed by resources such as a PLLs, LVDS, or gigabit transceiver blocks. You can check the legality of pin assignments before you compile your design, or allow the process to run automatically during compilation. If design files are available, you can perform more thorough legality checks on the I/O pins and surrounding logic in your design. Legality checks include proper \texttt{VREF} pin use, valid pin location assignments, and acceptable mixed I/O standards. Altera recommends that you run I/O assignment analysis each time you add or modify a pin-related assignment.

If you have partial or complete design files, you must perform Analysis and Synthesis to generate a synthesized (mapped) netlist before you can perform I/O assignment analysis.

Performing I/O assignment analysis directs the Fitter to read assignments from your mapped netlist and the \texttt{.qsf} to determine the legality of your pin-related assignments. These pin-related assignments include pin settings such as I/O standards, drive strength, and location assignments.

Incomplete I/O assignments trigger warnings during I/O assignment analysis. You can view the I/O Assignment Warnings report to find and resolve warnings generated during I/O assignment analysis. For example, you may receive a warning that some of the pins in the design are missing a drive strength or slew rate. Single-ended output and bidirectional pins default to the non-calibrated OCT setting if you do not assign drive strength and slew rate options to the pins, or if other OCT options are assigned to the pins. To resolve this issue, you can either assign drive strength or slew rate settings to the pins with the \texttt{Current Strength} or \texttt{Slew Rate} or \texttt{Slow Slew Rate} logic options, or assign the \texttt{Termination} logic option to the pins with a series setting. You cannot use drive strength and slew rate settings when a pin is assigned an OCT setting.

During I/O assignment analysis the Fitter automatically assigns suggested pin locations to unassigned pins in your design, based on your design constraints, so it can perform pin legality checks. For example, if you assign an edge location to a group of LVDS pins, the Fitter assigns pin locations for each LVDS pin in the specified edge location and then performs legality checks. To display the Fitter-placed pins use the Show Fitter Placements feature in the Pin Planner. To accept these suggested pin locations, you must back-annotate your pin assignments.

The following design flows show two different circumstances in which you can use I/O assignment analysis:
Validating Pin Assignments

Use the flow shown in Figure 4–13 if you must complete board layout before starting the design. This flow does not require design files and checks the legality of your pin assignments.

Use the flow shown in Figure 4–14 if your design is complete. This flow thoroughly checks the legality of your pin assignments against any design files provided.

Each flow involves creating pin assignments, running analysis, and reviewing the report file.

Running I/O Assignment Analysis without Design Files

During the early stages of developing a device, board layout engineers may request preliminary or final pin-outs. It is time consuming to manually check whether the pin-outs violate any design rules. To quickly perform basic checks on the legality of your pin assignments, you can use the Quartus II software to perform I/O assignment analysis.

If you create pin-related assignments in Mentor Graphics I/O Designer software, you can import a .fx into the Quartus II software.

Running I/O assignment analysis performs limited checks on pin assignments made in a design in which you specified a device, but does not yet include any HDL design files. For example, you can create a Quartus II project with only a target device specified and create pin-related assignments based on circuit board layout considerations that are already determined. Even though the Quartus II project does not yet contain any design files, you can reserve input and output pins and create pin-related assignments for each pin with the Pin Planner. After you assign an I/O standard to each reserved pin, run I/O assignment analysis to ensure that there are no I/O standard conflicts in each I/O bank. Figure 4–13 shows the work flow for assigning and analyzing pin-outs without design files.

Figure 4–13. Assigning and Analyzing Pin-Outs without Design Files
When you make and analyze pin-related assignments without design files, make sure you reserve the pins you intend to use as I/O pins, so the Fitter can determine each pin type. After performing I/O assignment analysis, correct any errors reported by the Fitter and rerun I/O assignment analysis until all errors are corrected.

Without a complete design, running I/O assignment analysis performs limited checks and cannot guarantee that your assignments do not violate design rules.

**Running I/O Assignment Analysis with Design Files**

If you have preliminary or complete design files, you can run I/O assignment analysis to help you determine if your design will fit. The rules checked during I/O assignment analysis depend on the completeness of the design. If you have a complete design, the legality of all pin-related assignments are thoroughly checked during I/O assignment analysis. If you have a partial design, which can be just the top-level wrapper file, the legality of those pin-related assignments for which there is enough information are checked during I/O assignment analysis. Figure 4–14 shows the work flow for assigning and analyzing pin-outs without design files.

**Figure 4–14. Assigning and Analyzing Pin-Outs with Design Files**
If you run I/O assignment analysis on incomplete design files, you may still encounter errors during full compilation. For example, you may assign a clock to a user I/O pin instead of assigning it to a dedicated clock pin, or design the clock to drive a PLL that you have not yet instantiated in the design. The checks run during I/O assignment analysis do not account for the logic that the pin drives, and do not check that only a dedicated clock input pin can drive the clock port of a PLL. To obtain better coverage, analyze as much of the design as possible, especially logic that connects to pins. For example, if your design includes PLLs or LVDS blocks, include these MegaWizard Plug-In Manager-generated files. To assign and analyze pin-related assignments successfully, after performing I/O assignment analysis, correct any errors reported by the Fitter and rerun I/O assignment analysis until all errors are corrected.

Figure 4–15 shows the compilation time benefit of performing I/O assignment analysis before running a full compilation.

Figure 4–15. Saving Compilation Time with the I/O Assignment Analysis

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Optimizing I/O Assignment Analysis with Output Enable Group Logic Option Assignments

Each device has a certain number of VREF pins, and each VREF pin supports a certain number of I/O pins. A VREF pin and its supported I/O pins are called a VREF bank. The VREF pins are used only for inputs with VREF I/O standards, such as HSTL- and SSTL-type I/O standards. VREF outputs do not require the VREF pin. When a voltage-referenced input is present in a VREF bank, only a certain number of outputs can be present in that VREF bank. For example, for devices in the Stratix II flip chip package, only 20 outputs can be present in a VREF bank when a VREF I/O standard input is present in that bank.

For more information about device VREF pins and their associated I/O pins, refer to the Pin-Out Files for Altera Devices page of the Altera website.

For interfaces that use bidirectional VREF I/O pins, your design must meet the output restriction for each I/O bank when the pins are driving in either direction. If a set of bidirectional signals are controlled by different output enables, they are treated as independent output enables during I/O assignment analysis, thus creating a situation in which the Fitter may determine that your design violates VREF restrictions. To treat
the set of bidirectional signals as a single output enable group so that the Fitter does not determine that the design violates the requirements for the maximum number of pins driving out of a VREF group, assign the **Output Enable Group** logic option assignment to the bidirectional signals. Assign an integer value for the **Output Enable Group** logic option assignment and assign the same integer value to all sets of signals that are driving in the same direction. Assigning this logic option to groups of signals is important in the case of external memory interfaces.

For example, in the case of a DDR2 interface in a Stratix II device, the device can have 30 pins in a VREF group. Each byte lane for a ×8 DDR2 interface has one DQS pin and eight DQ pins, for a total of nine pins per byte lane. The DDR2 interface uses **SSTL-18 Class I** as its I/O standard, which is a VREF I/O standard. In typical interfaces, each byte lane has its own output enable. In this example, the DDR2 interface has four byte lanes. Using 30 I/O pins in a VREF group, there are three byte lanes and an extra byte lane that supports the three remaining pins. If you do not use the **Output Enable Group** logic option assignment, the Fitter analyzes each byte lane as an independent group driven by a unique output enable during I/O assignment analysis. With this arrangement, the worst-case scenario is when the three pins are inputs, and the other 27 pins are outputs. In this case, the 27 output pins violate the 20 output pin limit.

In a DDR2 interface, all DQS and DQ pins are always driven in the same direction. Therefore, the Fitter reports an error that is not applicable to your design. Assigning the **Output Enable Group** logic option assignment to the DQS and DQ pins forces the Fitter to check these pins as a group driven by a common output enable during I/O assignment analysis. When you use the **Output Enable Group** logic option assignment, the DQS and DQ pins are checked as all input pins or all output pins and are not in violation of the I/O rules.

You can also use the **Output Enable Group** logic option assignment with pins that are driven only at certain times. For example, the data mask signal in DDR2 interfaces is an output signal, but it is driven only when the DDR2 is writing (bidirectional signals are outputs). To avoid errors during I/O assignment analysis, use the **Output Enable Group** logic option assignment to assign the data mask to the same value as the DQ and DQS signals.

You can also assign the **Output Enable Group** logic option to VREF input pins. If the VREF input pins are not active during the time the outputs are driving, add the VREF input pins to the output enable group, thus removing the VREF input pins from the VREF analysis. For example, the QVLD signal for an RLDRAM II interface is active only during a read. During a write, the QVLD pin is not active and does not count as an active VREF input pin in the VREF group. You can place the QVLD pins in the same output enable group as the RLDRAM II data pins.

**Understanding the I/O Assignment Analysis Report**

When I/O assignment analysis is complete, you can view detailed analysis reports and a `.pin`. The detailed messages in the reports help you quickly understand and resolve pin assignment errors. Each message includes a related node name and a description of the problem.

The Fitter section of the Compilation report contains information generated during I/O assignment analysis, including the following reports:

- I/O Assignment Warnings
- Resource Section
I/O Rules Section

The I/O Assignment Warnings report provides a list of pins and the Fitter warnings generated for the pins during I/O assignment analysis.

The Resource Section contains reports that categorize the pins as input pins, output pins, and bidirectional pins. You can view the utilization of each I/O bank in your device in the I/O Bank Usage report.

The I/O Rules Section includes detailed information about the I/O rules tested during I/O assignment analysis and contains the following reports:

- I/O Rules Summary report
- I/O Rules Details report
- I/O Rules Matrix report

The I/O Rules Summary report provides a quick summary of the number of I/O rules tested and how many applicable rules passed, failed, or were not checked due to other failing rules.

The I/O Rules Details report provides detailed information on all I/O rules. The Status column indicates whether applicable rules passed, failed, or could not be checked. All rules are given a level of severity rating to indicate their level of importance for an effective analysis.

The I/O Rules Matrix report (refer to Figure 4–16) provides a list of the I/O rules checked by the Fitter for each pin in the design. Rules that apply to the target device family either pass or fail for each pin. Rules marked Inapplicable are rules that do not apply to the target device family. You can ignore any rule marked Inapplicable.

Figure 4–16. I/O Rules Matrix
Validating Pin Assignments with Full Compilation

After performing preliminary pin assignment validation with the live I/O check feature and running I/O assignment analysis, you must perform a final I/O timing check of your design by performing a full compilation. To avoid costly board respins, you must include complete design files and constraints. With timing information, the Fitter makes intelligent placement and routing to achieve optimal timing performance in your design. Use the TimeQuest Timing Analyzer to create timing constraints for input, output, and bidirectional pins.

For more information about the TimeQuest analyzer, refer to the The Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.

Performing I/O Timing Analysis

Timing analysis is usually run during a full compilation of your design or when you perform an early timing estimate. You can also run timing analysis independently after you fully compile the design. For example, if you change the slew rates or drive strengths of some I/O pins with ECOs, you do not have to recompile the entire design, but only run timing analysis to verify the timing of your design.

As part of I/O planning, especially with high-speed designs, take board-level signal integrity and timing into account. When adding a device with high-speed interfaces to a board design, the quality of the signal at the far end of the board route, as well as the propagation delay, is vital for proper system operation.

As part of I/O planning, you must understand the I/O timing results reported by the Quartus II software after performing timing analysis on your design. If all your design files are complete and you have fully compiled your design, all the timing checks related to I/O timing are performed during timing analysis. Static timing analysis is performed when you compile your design in the Quartus II software. You must understand I/O timing and what factors affect I/O timing paths in your design. One important factor in I/O timing results is how accurately you specify the output loads of the output and bidirectional pins in your design. Incomplete I/O constraints can affect your I/O timing results.

The Quartus II software supports three different methods of I/O timing analysis:

- Advanced I/O timing using a user-defined board trace model to produce enhanced timing reports from accurate, “board-aware”, simulation models

  Advanced I/O timing allows you to configure a complete board trace model for each I/O standard or pin used in your design. With advanced I/O timing, the TimeQuest analyzer uses the results of simulations of the I/O buffer, package, and board trace model to generate more accurate I/O delays and extra reports to give insight into signal behavior at the system level. You can use these advanced timing reports as a guide to make changes to your I/O assignments and board design to improve timing and signal integrity.

  For more information about advanced I/O timing, including device support, refer to About Advanced I/O Timing in Quartus II Help.
I/O timing using a default or user-specified capacitive load without signal integrity analysis

The TimeQuest analyzer creates timing reports that measure \( t_{CO} \) to an I/O pin using a default or user-specified value for a capacitive load.

Full board routing simulation in third-party tools using Altera-provided or Quartus II software-generated IBIS or HSPICE I/O models

The IBIS and HSPICE Writers the simulation model files for use by third-party board simulation tools. The IBIS and HSPICE Writers in the Quartus II software can export accurate simulation models for use in applications such as Mentor Graphics HyperLynx and Synopsys HSPICE.

For devices that support advanced I/O timing, it is the default method of I/O timing analysis. For all other devices, you must use a default or user-specified capacitive load assignment to determine \( t_{CO} \) and power measurements.

For more information about advanced I/O timing support, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website. For more information about board-level signal integrity and tips on how to improve signal integrity in your high-speed designs, refer to the Altera Signal Integrity Center page of the Altera website.

For information about creating IBIS and HSPICE models with the Quartus II software and integrating those models into HyperLynx and HSPICE simulations, refer to the Signal Integrity Analysis with Third-Party Tools chapter in volume 2 of the Quartus II Handbook.

Enabling and Configuring Advanced I/O Timing

With the advanced I/O timing feature, you can expand upon the basic timing and power measurements made with the capacitive loading settings. The advanced I/O timing feature gives you the ability to fully define not only the capacitive load, but also any termination components and trace impedances in the board routing for any output pin or bidirectional pin in output mode. You can configure an overall board trace model for each I/O standard as well as customize the model for specific pins.

When you use the advanced I/O timing feature, the board trace model replaces any capacitive load setting you made because the load is included in the model. For timing measurements, the entire board trace model is taken into account when calculating I/O delays. For power measurements, an effective capacitive load is used based on the sum of the capacitive elements in the model, including the Near capacitance, Far capacitance, and Transmission line distributed capacitance elements of the model.

Defining Overall Board Trace Models

You can define an overall board trace model for each I/O standard in your design that is the default model for all pins that use a particular I/O standard. After configuring the overall board trace model, you can customize the model for specific pins using the Board Trace Model window in the Pin Planner.
Custom component value changes you make to selected pins in the Pin Planner take priority and are not affected by subsequent changes to a specific components for the entire design. Similarly, any changes you make to specific pins do not affect the component settings for the entire design.

When you define an overall board trace model you can specify the board trace, termination, and capacitive load parameters for each I/O standard. The default settings for components in the model for each I/O standard are device-specific and match the default test model used for calculating delay without advanced I/O timing. For differential I/O standards, the component values you set are used for both the positive and negative signals of a differential pin pair.

All the assignments for board trace models you specify are saved to the .qsf. You can also use Tcl commands to create board trace model assignments. Example 4–8 shows Tcl commands for specifying board trace model assignments.

### Example 4–8. Specifying Board Trace Models

```tcl
## setting the near end series resistance model of sel_p output pin to 25 ohms
set_instance_assignment -name BOARD_MODEL_NEAR_SERIES_R 25 -to sel_p
## Setting the far end capacitance model for sel_p output signal to 6 picofarads
set_instance_assignment -name BOARD_MODEL_FAR_C 6P -to sel_p
```

For more information about configuring component values for a board trace model, including a complete list of the supported unit prefixes and setting the values with Tcl scripts, refer to Board Trace Model Quartus II Help.

For more information about the default models used for measuring I/O delay, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

### Customizing the Board Trace Model in the Pin Planner

In the Pin Planner, you can view a graphical representation of the board trace model you configured with the Board Trace Model window. Initially, the settings you create for the overall board trace model match the settings in the Pin Planner. For differential signals, the Board Trace Model window displays the routing and components for both the positive and negative signals of the differential pair. Any changes you make for a differential signal pair must be performed on the positive signal of the pair. The settings must match between the positive and negative signals of a differential pair, so the changes are automatically reflected in the settings for the negative signal.

When editing board trace model assignments, for numerical values, use standard unit prefixes such as p, n, and k to represent pico, nano, and kilo, respectively. To short a series component or have an open circuit for a parallel component, select short or open, respectively, for the component value.
Configuring Board Trace Models

The Quartus II software provides board trace model templates for various I/O standards in which you can fill in various parameters. Figure 4–17 shows the template for a 2.5-V I/O standard. This model consists of near-end and far-end board component parameters.

Modeling of the near-end of the board trace includes the elements which are close to the device and modeling of the far-end includes the elements which are at the receiver end of the link, closer to the receiving device. The topology represented in the Quartus II board trace model is conceptual and does not necessarily match the board trace component for component. For example near-end model parameters can represent device-end discrete termination and breakout traces. Far-end modeling can represent the bulk of the board trace to discrete external memory components, and the far end termination network. The same circuit can be analyzed with near-end modeling of the entire board, including memory component termination, and far-end modeling of the actual memory component.

Figure 4–17. Board Trace Model

Figure 4–18 shows the template for the LVDS I/O standard. The far-end capacitance (Cf) represents the external-device or multiple-device capacitive load. If you have multiple devices on the far-end, you must find the equivalent capacitance at the far-end, taking into account all receiver capacitances. The far-end capacitance can be the sum of all the receiver capacitances.

For more information about the specifications for external device capacitance values, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.
The Quartus II software models lossless transmission lines, and does not require a transmission-line resistance value. Only distributed inductance (L) and capacitance (C) values are needed. The distributed L and C values of transmission lines must be entered on a per-inch basis, and can be obtained from the PCB vendor or manufacturer, the CAD Design tool, or a signal integrity tool such as the Mentor Graphics Hyperlynx software.

**Figure 4–18. Differential Board Trace Model**

Specifying Near-End vs Far-End Timing Analysis

With advanced I/O timing analysis, you have the option of selecting a near-end or far-end point for your I/O timing. With near-end timing, the timing is analyzed to the device pin. Figure 4–17 shows near-end timing ending at the vertical dashed line separating the device I/O pin and off-chip components.

By default, advanced I/O timing analysis analyzes output I/O timing to the device pin. When you choose a near-end endpoint, you can use the `set_output_delay` SDC timing constraint to account for the delay across the board. However, when you choose a far-end I/O timing endpoint, then advanced I/O timing analysis analyzes timing to the external device input, at the far end of the board trace. Whether you choose a near-end or far-end timing endpoint, the board trace models are taken into account during timing analysis.

For more information about calculating I/O timing to the near-end or far-end of the board trace, refer to *Using Advanced I/O Timing* in Quartus II Help.

**Understanding Advanced I/O Timing Analysis Reports**

When you perform advanced I/O analysis, the TimeQuest analyzer creates reports that signal integrity reports that provide board delay estimates and signal integrity data.
The TimeQuest analyzer section of the Compilation report contains information generated during advanced I/O timing analysis, including the following reports:

- Board Trace Model Assignments
- Signal Integrity Metrics

The Board Trace Model Assignments report summarizes the board trace model component settings for each output and bidirectional signal.

The Signal Integrity Metrics report contains all the signal integrity metrics calculated during advanced I/O timing analysis based on the board trace model settings for each output or bidirectional pin. The reports contain many metrics, including measurements at both the FPGA pin and at the far-end load of board delay, steady state voltages, and rise and fall times.

By default, the TimeQuest analyzer generates the Slow-Corner Signal Integrity Metrics report. To generate a Fast-Corner Signal Integrity Metrics report you must change the delay model.

For more information about the reports generated during advanced I/O timing analysis, refer to About TimeQuest Timing Analysis in Quartus II Help. For more information about the metrics calculated during advanced I/O timing analysis, including diagrams illustrating the metrics on output waveforms, refer to Signal Integrity Metrics in Quartus II Help. For more information about changing the delay model, refer to the Create Timing Netlist Dialog Box in Quartus II Help.

Adjusting I/O Timing and Power with Capacitive Loading

When calculating tCO and power for output and bidirectional pins, the TimeQuest analyzer and the PowerPlay Power Analyzer use a bulk capacitive load. You can adjust the value of the capacitive load per I/O standard to obtain tCO and power measurements that more accurately reflect the behavior of the output or bidirectional net on your PCB. Input pins ignore any capacitive load settings. You can adjust the capacitive load settings per I/O standard, in picofarads (pF), for your entire design. During compilation, the Compiler measures power and tCO measurements based on your settings. You can also adjust the capacitive load on an individual pin with the Output Pin Load logic option.

For more information about adjusting the value of the capacitive load for your entire design, refer to Setting Up and Running a Compilation in Quartus II Help. For more information about adjusting the value of the capacitive load on an individual pin, refer to Assigning Device I/O Pins in Quartus II Help.
Incorporating PCB Design Tools

Signal and pin assignments are initially made by the chip designer and it is up to the board designer to correctly transfer these assignments to the symbols in their system circuit schematics and board layout. As the board design progresses, pin reassignments may be requested or required to optimize the layout. These reassignments must in turn be communicated to the chip designer, so that the new assignments can be validated during I/O assignment analysis and processed through an updated placement-and-routing of the device.

The Quartus II software interacts with board layout tools by importing and exporting pin information files, including the .qsf, .pin, and .fx files.

For more information about incorporating PCB design tools, refer to the Cadence PCB Design Tools Support and Mentor Graphics PCB Design Tools Support chapters in volume 2 of the Quartus II Handbook.

Scripting Support

A Tcl script allows you to run procedures and determine settings described in this chapter. You can also run some of these procedures at a command prompt.

For detailed information about specific scripting command options and Tcl API packages, type the following command at a system command prompt to run the Quartus II Command-Line and Tcl API Help browser:

```
quartus_sh --qhelp
```

For more information about Quartus II scripting support, including examples, refer to the Tcl Scripting and Command-Line Scripting chapters in volume 2 of the Quartus II Handbook.

Running I/O Assignment Analysis

You can run I/O assignment analysis with a Tcl command or with a command-line command.

Enter the following in the Tcl console or a Tcl script:

```
execute_flow -check_ios
```

Type the following at a system command prompt:

```
quartus_fit <project name> --check_ios
```

For more information about running I/O assignment analysis, refer to “Understanding the I/O Assignment Analysis Report” on page 4–29.

Generating a Mapped Netlist

You can generate a mapped netlist with a Tcl command or with a command-line command.

Enter the following in the Tcl console or in a Tcl script:

```
execute_module -tool map
```

The execute_module command is in the flow package.
Type the following at a system command prompt:

```
quartus_map <project name> r
```

### Reserving Pins

You can reserve pins with a Tcl command.

Use the following Tcl command to reserve a pin:

```
set_instance_assignment -name RESERVE_PIN <value> -to <signal name>
```

Use one of the following valid reserved pin values:

- "AS BIDIRECTIONAL"
- "AS INPUT TRI-STATED"
- "AS OUTPUT DRIVING AN UNSPECIFIED SIGNAL"
- "AS OUTPUT DRIVING GROUND"
- "AS SIGNALPROBE OUTPUT"

Ensure you include the quotation marks when specifying the reserved pin value.

### Creating Location Assignments

You can create location assignments with a Tcl command.

Use the following Tcl command to assign a signal to a pin or device location:

```
set_location_assignment <location> -to <signal name> r
```

Valid locations are pin locations, I/O bank locations, or edge locations. Pin locations include pin names, such as PIN_A3. I/O bank locations include IOBANK_1 up to IOBANK_n, in which n is the number of I/O banks in the device.

Use one of the following valid edge location values:

- EDGE_BOTTOM
- EDGE_LEFT
- EDGE_TOP
- EDGE_RIGHT

For more information about location assignments, refer to “Creating Location Assignments” on page 4–16.

For more information about I/O banks in your device, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

### Creating Exclusive I/O Group Assignments

You can create exclusive I/O group assignments with a Tcl command.

Use the following Tcl command to create an exclusive I/O group assignments:

```
set_instance_assignment -name "EXCLUSIVE_IO_GROUP" -to pin
```
For more information about exclusive I/O group assignments, refer to “Creating Exclusive I/O Group Assignments” on page 4–17.

Changing the Slew Rate and Drive Strength

You can create slew rate and drive strength assignments with a Tcl command.

Use the following Tcl commands to create an slew rate and drive strength assignments:

```tcl
set_instance_assignment -name CURRENT_STRENGTH_NEW 8MA -to e[0]
set_instance_assignment -name SLEW_RATE 2 -to e[0]
```

For more information about slew rate and drive strength assignments, refer to “Creating Pin Assignments with the Chip Planner” on page 4–19.

Conclusion

The Quartus II software provides many tools and features to help you with I/O planning, including the ability to validate pin assignments in all design stages, even before the development of your design. The ability to import and export assignments between the Quartus II software and other PCB tools allows you to make iterative changes efficiently. Finally, the ability to enter a board trace model and create advanced timing reports based on how I/O signals are routed on a board truly makes the Quartus II software board-aware.

For more information about the Altera resources available for I/O planning, refer to the I/O Management, Board Development Support, and Signal Integrity Analysis Resource Center of the Altera website.

For more information about PCB designs for Altera high-speed FPGAs, refer to AN 315: Guidelines for Designing High-Speed FPGA PCBs, and the Board Design Resource Center of the Altera website.

Document Revision History

Table 4–5 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>■ Added information about overriding I/O placement rules.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Updated Pin Planner description for new task and report windows.</td>
</tr>
<tr>
<td>June 2012</td>
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<td>■ Removed survey link.</td>
</tr>
</tbody>
</table>
| November 2011 | 11.1.0 | ■ Minor updates and corrections.  
            |         | ■ Updated the document template.                |
| December 2010 | 10.0.1 | Template update                                 |
Table 4–5. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
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<tbody>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>- Reorganized and edited the chapter</td>
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<tr>
<td></td>
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<td>- Added links to Quartus II Help for procedural information previously included in the chapter</td>
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<td></td>
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<td>- Added information on rules marked Inapplicable in the I/O Rules Matrix Report</td>
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<td></td>
<td></td>
<td>- Added information on assigning slew rate and drive strength settings to pins to fix I/O assignment warnings</td>
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<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>- Reorganized entire chapter to include links to Quartus II help for procedural information previously included in the chapter</td>
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<tr>
<td></td>
<td></td>
<td>- Added documentation on near-end and far-end advanced I/O timing</td>
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<td>March 2009</td>
<td>9.0.0</td>
<td>- Updated “Pad View Window” on page 5–20</td>
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<tr>
<td></td>
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<td>- Added new figures:</td>
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<tr>
<td></td>
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<td>- Figure 5–15</td>
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<td></td>
<td></td>
<td>- Figure 5–16</td>
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<td></td>
<td></td>
<td>- Added new section “Viewing Simultaneous Switching Noise (SSN) Results” on page 5–17</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Added new section “Creating Exclusive I/O Group Assignments” on page 5–18</td>
</tr>
</tbody>
</table>

*For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.*
FPGA design has evolved from small programmable circuits to designs that compete with multimillion-gate ASICs. At the same time, the I/O counts on FPGAs and logic density requirements of designs have increased exponentially. The higher-speed interfaces in FPGAs, including high-speed serial interfaces and memory interfaces, require careful interface design on the PCB. Designers must address the timing and signal integrity requirements of these interfaces early in the design cycle. Simultaneous switching noise (SSN) often leads to the degradation of signal integrity by causing signal distortion, thereby reducing the noise margin of a system.

Today’s complex FPGA system design is incomplete without addressing the integrity of signals coming in to and out of the FPGA. Altera recommends that you perform SSN analysis early in your FPGA design and prior to the layout of your PCB with complete SSN analysis of your FPGA in the Quartus® II software. This chapter describes the Quartus II SSN Analyzer tool and covers the following topics:

- “Definitions”
- “Understanding SSN” on page 5–2
- “SSN Estimation Tools” on page 5–5
- “SSN Analysis Overview” on page 5–5
- “Optimizing Your Design for SSN Analysis” on page 5–8
- “Performing SSN Analysis and Viewing Results” on page 5–15
- “Decreasing Processing Time for SSN Analysis” on page 5–17

Definitions

The terminology used in this chapter includes the following terms:

**Aggressor**: An output or bidirectional signal that contributes to the noise for a victim I/O pin

**PDN**: Power distribution network

**QH**: Quiet high signal level on a pin

**QHN**: Quiet high noise on a pin, measured in volts

**QL**: Quiet low signal level on a pin

**QLN**: Quiet low noise on a pin, measured in volts

**SI**: Signal integrity (a superset of SSN, covering all noise sources)

**SSN**: Simultaneous switching noise
SSO: Simultaneous switching output (which are either the output or bidirectional pins)

Victim: An input, output, or bidirectional pin that is analyzed during SSN analysis. During SSN analysis, each pin is analyzed as a victim. If a pin is an output or bidirectional pin, the same pin acts as an aggressor signal for other pins.

**Understanding SSN**

SSN is defined as a noise voltage induced onto a single victim I/O pin on a device due to the switching behavior of other aggressor I/O pins on the device. SSN can be divided into two types of noise: voltage noise and timing noise.

Figure 5–1 shows a system with three pins. Two of the pins (A and C) are switching, while one pin (B) is quiet. If the pins are driven in isolation, the voltage waveforms at the output of the buffers appear without noise interference, as shown by the solid curves at the left of the figure. However, when the pins are switched simultaneously, the noise generated by pins A and C switching is injected onto the other pins, manifesting itself as a voltage noise on pin B and timing noise on pins A and C, as shown by the dotted curves in the figure.

Voltage noise is measured as the worst-case change in voltage of a signal due to SSN. When a signal is QH, it is measured as the change in voltage toward 0 V. When a signal is QL, it is measured as the change in voltage toward VCC.

In the Quartus II software, only voltage noise is analyzed. Voltage noise can be caused by SSOs under two worst-case conditions:

- The victim pin is high and the aggressor pins (SSOs) are switching from low to high
- The victim pin is low and the aggressor pins (SSOs) are switching from high to low
For outputs, the noise is computed at the far-end receiver for pin B (refer to Figure 5–2).

**Figure 5–2. Quiet High Output Noise Estimation**

For inputs, the noise is computed at the FPGA bumps as shown in for pin D (refer to Figure 5–3).

**Figure 5–3. Quiet Low Input Noise Estimation**

SSN can occur in any system, but the induced noise does not always result in failures. Voltage functional errors are caused by SSN on quiet victim pins only when the voltage values on the quiet pins change by a large enough voltage that the logic listening to that signal reads a change in the logic value. For QH signals, a voltage functional error occurs when noise events cause the voltage to fall below $V_{IH}$. Similarly, for QL signals, a voltage functional error occurs when noise events cause the voltage to rise above $V_{IL}$ (refer to Figure 5–4). Because $V_{IH}$ and $V_{IL}$ are different for
different I/O standards, and because signals have different quiet voltage values, the absolute amount of SSN, measured in volts, cannot be used to determine if a voltage failure occurs. Instead, to quantify whether an SSN event will cause a voltage error, the Quartus II software uses the amount of noise as a percent of signal margin when reporting noise margins in SSN analysis (refer to Figure 5–4).

**Figure 5–4. Reporting Noise Margins**

![Figure 5–4. Reporting Noise Margins](image)

Figure 5–4 shows four noise events, two on QH signals and two on QL signals. The two noise events on the right-side of the figure consume 50 percent of the signal margin and do not cause voltage functional errors. However, the two noise events on the left side of the figure consume 100 percent of the signal margin and can cause a voltage functional error.

**Figure 5–5. Synchronous Voltage Noise**

![Figure 5–5. Synchronous Voltage Noise](image)

Figure 5–5 illustrates a synchronous voltage noise event that does not result in a voltage functional error. Noise or glitches caused by aggressor signals are synchronously related to the victim pin outside of the sampling window of a receiver. The noise or glitches affect the switching time of a victim pin, but are not considered an input threshold violation failure.

For more information about the design factors that affect the noise margins during SSN analysis in the Quartus II software, refer to “SSN Analysis Overview”.
SSN Estimation Tools

Addressing SSN early in your FPGA design and PCB layout can help you avoid costly board respins and lost time, both of which can impact your time-to-market. Altera provides many tools for SSN analysis and estimation, including the following tools:

- SSN characterization reports
- An early SSN estimation (ESE) tool
- The SSN Analyzer in the Quartus II software

For more information about the SSN characterization reports and the ESE tool, including device support information, refer to the Signal Integrity Center page of the Altera website.

For more information about the devices for which you can run the SSN Analyzer, refer to About the SSN Analyzer in Quartus II Help.

The ESE tool is useful for preliminary SSN analysis of your FPGA design; for more accurate results, however, you must use the SSN Analyzer in the Quartus II software. Table 5–1 compares some of the differences between the ESE tool and the SSN Analyzer.

<table>
<thead>
<tr>
<th>ESE Tool</th>
<th>SSN Analyzer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Is not integrated with the Quartus II software.</td>
<td>Integrated with the Quartus II software, allowing you to perform preliminary SSN analysis while making I/O assignment changes in the Quartus II software.</td>
</tr>
<tr>
<td>QL and QH levels are computed assuming a worst-case pattern of I/O placements.</td>
<td>QL and QH levels are computed based on the I/O placements in your design.</td>
</tr>
<tr>
<td>No support for entering board information.</td>
<td>Supports board trace models and board layer information, resulting in a more accurate SSN analysis.</td>
</tr>
<tr>
<td>No graphical representation.</td>
<td>Integrated with the Quartus II Pin Planner, in which an SSN map shows the QL and QH levels on victim pins.</td>
</tr>
<tr>
<td>Good for doing an early SSN estimate. Does not require you to use the Quartus II software.</td>
<td>Requires you to create a Quartus II software project and provide the top-level port information.</td>
</tr>
</tbody>
</table>

SSN Analysis Overview

You can run the SSN Analyzer at different stages in your design cycle to obtain SSN results. The accuracy of the results depends on the completeness of your design information. Altera recommends that you start SSN analysis early in the design cycle to obtain preliminary results and make adjustments to your I/O assignments, and iterate through the design cycle to finally perform a fully constrained SSN analysis with complete information about your board.

Figure 5–6 shows the flows for both early pin-out and final pin-out SSN analysis. The early pin-out flow assumes conservative design rules initially, and then lets you analyze the design and iteratively apply tighter design rules until SSN analysis indicates your design meets SSN constraints. You must define pass criteria for SSN analysis as a percentage of signal margin in both the early pin-out flow and the final
SSN Analysis Overview

The pass criteria you define is specific to your design requirements. For example, a pass criterion you might define is a condition that verifies you have sufficient SSN margins in your design. You may require that the acceptable voltage noise on a pin must be below 70% of the voltage level for that pin. The pass criteria for the early-pin out flow may be higher than the final pin-out flow criteria, so that you do not spend too much time optimizing the on-FPGA portions of your design when the SSN metrics for the design may improve after the design is fully specified.

**Figure 5–6. Pin-Out Analysis**

Performing Early Pin-Out SSN Analysis

In the early stages of your design cycle, before you create pin location for your design, use the early pin-out flow (refer to Figure 5–6) to obtain preliminary SSN analysis results. In order to obtain useful SSN results, you must define the top-level ports of your design, but your design files do not have to be complete.

**Performing Early Pin-Out SSN Analysis with the ESE Tool**

If you know the I/O standards and signaling standards for your design, you can use the ESE tool to perform an initial SSN evaluation.

For more information about the ESE tool, refer to the *Signal Integrity Center* page of the Altera website.
Performing Early Pin-Out SSN Analysis with the SSN Analyzer

If you have complete information for the top-level ports of your design, you can use the SSN Analyzer to perform an initial SSN evaluation. Use the following steps to perform early pin-out SSN analysis:

1. Create a project in the Quartus II software.
2. Specify your top-level design information either in schematic form or in HDL code.
4. Create I/O assignments, such as I/O standard assignments, for the top-level ports in your design.
   - Do not create pin location assignments. The Fitter automatically creates optimized pin location assignments.
5. If you do not have completed design files and timing constraints, run I/O assignment analysis.
   - During I/O assignment analysis, the Fitter places all the unplaced pins on the device, and checks all the I/O placement rules.
6. Run the SSN Analyzer.

For more information about creating and managing projects, refer to the Managing Quartus II Projects chapter in volume 2 of the Quartus II Handbook. For more about generating a top-level design file in the Quartus II software and I/O assignment analysis, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

In the early stages of your design cycle, you may not have complete board information, such as board trace parameters, layer information, and the signal breakout layers. If you run the SSN Analyzer without this specific information, it uses default board trace models and board layer information for SSN analysis, and as a result the SSN Analyzer confidence level is low. If the noise amounts are larger than the pass criteria for early pin-out SSN analysis, verify whether the SSN noise violations are true failures or false failures. For example, sometimes the SSN Analyzer can determine whether pins are switching synchronously and use that information to filter false positives; however, it may not be able to determine all the synchronous groups. You can improve the SSN analysis results by adjusting your I/O assignments and other design settings. After you optimize your design such that it meets the pass criteria for the early pin-out flow, you can then begin to design your PCB.

For more information, refer to “Optimizing Your Design for SSN Analysis”.

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Volume 2: Design Implementation and Optimization
Performing Final Pin-Out SSN Analysis

You perform final pin-out SSN analysis after you place all the pins in your design, or the Fitter places them for you, and you have complete information about the board trace models and PCB layers. Even if your design achieves sufficient SSN results during early pin-out SSN analysis, you should run SSN analysis with the complete PCB information to ensure that SSN does not cause failures in your final design. You must specify the board parameters in the Quartus II software, including the PCB layer thicknesses, the signal breakout layers, and the board trace models, before you can run SSN analysis on your final assignments.

For more information, refer to “Optimizing Your Design for SSN Analysis”.

If the SSN analysis results meet the pass criteria for final pin-out SSN analysis, SSN analysis is complete. If the SSN analysis results do not meet the pass criteria, you must further optimize your design by changing the board and design parameters and then rerun the SSN Analyzer. If the design still does not meet the pass criteria, reduce the pass criteria for early pin-out SSN analysis, and restart the process. By reducing the pass criteria for early pin-out SSN analysis, you place a greater emphasis on reducing SSN through I/O settings and I/O placement. Changing the drive strength and slew rate of output and bidirectional pins, as well as adjusting the placement of different SSOs, can affect SSN results. Adjusting I/O settings and placement allows the design to meet the pass criteria for final pin-out SSN analysis after you specify the actual PCB board parameters.

Design Factors Affecting SSN Results

There are many factors that affect the SSN levels in your design. The two main factors are the drive strength and slew rate settings of the output and bidirectional pins in your design.

For more information about the factors that contribute to SSN voltage noise in your FPGA design and managing SSN in your system, refer to AN 472: Stratix II GX SSN Design Guidelines, AN 508: Cyclone III Simultaneous Switching Noise (SSN) Design Guidelines, and the Signal Integrity Center page of the Altera website.

Optimizing Your Design for SSN Analysis

The SSN Analyzer gives you flexibility to precisely define your system to obtain accurate SSN results. The SSN Analyzer produces a voltage noise estimate for each input, output, and bidirectional pin in the design. It allows you to estimate the SSN levels, comprised of QLN and QHN levels, for your FPGA pins. Performing SSN analysis helps you optimize your design for SSN during compilation.

Because the SSN Analyzer is integrated into the Quartus II software, it can automatically set up a system topology that matches your design. The SSN Analyzer accounts for different I/O standards and slew rate settings for each buffer in the design and models different board traces for each signal. Also, it correctly models the state of the unused pins in the design. The SSN Analyzer leverages any custom board trace assignments you set up for use by the advanced I/O timing feature.
The SSN Analyzer also models the package and vias in the design. Models for the different packages that Altera devices support are integrated into the Quartus II software. In the Quartus II software, you can specify different layers on which signals break out, each with its own thickness, and then specify which signal breaks out on which layer.

Figure 5–7 shows the circuit topology the SSN Analyzer automatically constructs. After constructing the circuit topology, the SSN Analyzer uses a simulation-based methodology to determine the SSN for each victim pin in the design.

### Figure 5–7. Circuit Topology for SSN Analysis

**Optimizing Pin Placements for Signal Integrity**

You can take advantage of a built-in SSN optimization feature in the Quartus II software with the SSN Optimization logic option.

The I/O placements in your design may be affected when you use this option. Setting this option to Normal compilation does not affect the \( f_{\text{MAX}} \) of your design during compilation, however setting this option to Extra effort level may impact your design \( f_{\text{MAX}} \).

In order to use the SSN Optimization logic option, Altera recommends that you do not create location assignments for your pins; instead, let the Fitter place the pins during compilation so that it places the pins to meet the timing performance of your design. To display the Fitter-placed pins use the Show Fitter Placements feature in the Pin Planner. To accept these suggested pin locations, you must back-annotate your pin assignments.
Figure 5–8 shows the results of turning on the **SSN Optimization** logic option for a design. The image on the left shows the placement of the pins without the **SSN Optimization** logic option, and the image on the right shows the adjustments the Fitter made to pin placements to reduce the amount of SSN in the design when the **SSN Optimization** logic option is turned on.

**Figure 5–8. SSN Analysis Results Before and After Using the SSN Optimization Logic Option**

For more information about creating project-wide logic option assignments, refer to *Setting Up and Running the Fitter* in Quartus II Help. For more information about the Show Fitter Placements feature, refer to *Show Commands* in Quartus II Help. For more information about back-annotating assignments, refer to *Back-Annotating Assignments for A Project* in Quartus II Help.

For more information about design optimization features, refer to the *Area, Timing, and Compilation Time Optimization* section in volume 2 of the *Quartus II Handbook*.

**Specifying Board Trace Model Settings**

The SSN Analyzer uses circuit models to determine voltage noise during SSN analysis. The circuit topology (refer to Figure 5–7) is incomplete without board trace information and PCB layer information. You must describe the board trace and PCB layer parameters in your design to accurately compute the SSN in your FPGA device. However, if you do not specify some or all of the board trace parameters and PCB layer information, the SSN Analyzer uses default parameters during SSN analysis. When you use the default parameters, the SSN confidence level is low.

For more information about the default parameters used by the SSN Analyzer and SSN confidence levels, refer to “Confidence Metric Details Report” on page 5–16.

The board trace models required for the SSN Analyzer include the board trace termination resistors, pin loads (capacitance), and transmission line parameters. You can define the board circuit models, which are also known as board trace models, in the Quartus II software. The board trace model settings are shared with the models used during advanced I/O timing.

For more information about defining board trace models and advanced I/O timing, refer to the *I/O Management* chapter in volume 2 of the *Quartus II Handbook*.
You can define an overall board trace model for each I/O standard in your design; this overall board trace model is the default model for all pins that use a particular I/O standard. After configuring the overall board trace model, you can customize the model for specific pins. The parameters you specify for the board trace model are also used in advanced I/O timing analysis with the TimeQuest Timing Analyzer. If you already specified the board trace models as part of your advanced I/O timing assignments, the same parameters are used during SSN analysis.

For more information about defining a board trace model for your entire design, refer to Using Advanced I/O Timing in Quartus II Help. For more information about configuring component values for a board trace model, including a complete list of the supported unit prefixes and setting the values with Tcl scripts, refer to Board Trace Model in Quartus II Help.

All the assignments for board trace models you specify are saved to the .qsf. You can also use Tcl commands to create board trace model assignments. Example 5–1 shows Tcl commands for specifying transmission line parameters.

Example 5–1. Specifying Board Trace Models

```
set_instance_assignment -name BOARD_MODEL_TLINE_L_PER_LENGTH "3.041E-7" -to e[0]
set_instance_assignment -name BOARD_MODEL_TLINE_LENGTH 0.1391 -to e[0]
set_instance_assignment -name BOARD_MODEL_TLINE_C_PER_LENGTH "1.463E-10" -to e[0]
```

The best way to calculate transmission line parameters is to use a two-dimensional solver to estimate the inductance per inch and capacitance per inch for the transmission line. The termination resistor topology information can be obtained from the PCB schematics. The near-end and far-end pin load (capacitance) values can be obtained from the PCB schematic and other device data sheets. For example, if you know that an FPGA pin is driving a DIMM, you can obtain the far-end loading information in the data sheet for your target device.

For more information, refer to the Device Family Data Sheet in the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

### Defining PCB Layers and PCB Layer Thickness

Every PCB is fabricated using a number of layers. To remove some of the pessimism from your SSN results, Altera recommends that you create assignments describing your PCB layers in the Quartus II software. You can specify the number of layers on your PCB, and their thickness. The PCB layer information is used only during SSN analysis and is not used in other processes run by the Quartus II software. If a custom PCB breakout region is not described you can select the default thickness, which directs the SSN Analyzer to use a single-layer PCB breakout region during SSN analysis.

For more information about specifying PCB layer information, refer to Running the SSN Analyzer in Quartus II Help.
All the assignments you create for the PCB layers are saved to the .qsf. You can also use Tcl commands to create PCB layer assignments. You can create any number of PCB layers, however, the layers must be consecutive. Example 5–2 shows Tcl commands for specifying PCB layer assignments.

Example 5–2. Specifying PCB Layer Assignments

<table>
<thead>
<tr>
<th>Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>set_global_assignment -name PCB_LAYER_THICKNESS 0.00099822M -section_id 1</td>
<td>Assigns thickness to layer 1</td>
</tr>
<tr>
<td>set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 2</td>
<td>Assigns thickness to layer 2</td>
</tr>
<tr>
<td>set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 3</td>
<td>Assigns thickness to layer 3</td>
</tr>
</tbody>
</table>

Figure 5–9 shows the layout cross-section of a PCB in the Cadence Allegro PCB tool. The cross-section shows the stackup information of a PCB, which tells you the number of layers used in your PCB. The PCB shown in this example consists of various signal and circuit layers on which FPGA pins are routed, as well as the power and ground layers.

```text
Figure 5–9. Snapshot of Stackup of a PCB Shown in the Allegro Board Design Environment
```

<table>
<thead>
<tr>
<th>Layer Name</th>
<th>Type</th>
<th>Material</th>
<th>Thickness (MIL)</th>
</tr>
</thead>
<tbody>
<tr>
<td>L4-SIGNAL</td>
<td>PLATED COPPER FOIL</td>
<td>1.9</td>
<td></td>
</tr>
<tr>
<td>L5-SIGNAL</td>
<td>FR-4</td>
<td>3.6</td>
<td></td>
</tr>
<tr>
<td>L8-SIGNAL</td>
<td>COPPER</td>
<td>1.2</td>
<td></td>
</tr>
<tr>
<td>L9-SIGNAL</td>
<td>FR-4</td>
<td>6</td>
<td></td>
</tr>
<tr>
<td>L6-GND</td>
<td>COPPER</td>
<td>1.2</td>
<td></td>
</tr>
<tr>
<td>L7-GND</td>
<td>FR-4</td>
<td>3</td>
<td></td>
</tr>
<tr>
<td>L8-PWR</td>
<td>COPPER</td>
<td>1.2</td>
<td></td>
</tr>
<tr>
<td>L9-PWR</td>
<td>FR-4</td>
<td>4</td>
<td></td>
</tr>
</tbody>
</table>

In this example, each of the four signal layers are a different thickness, with the depths shown in the Thickness (MIL) column. The layer thickness for each signal layer is computed as follows:

- Signal Layer 1 is the L4-SIGNAL, at thickness (1.9+3.6+1.2+3+1.2+4=) 14.9 mils
- Signal Layer 2 is the L5-SIGNAL, at thickness (0.6+6=) 6.6 mils
- Signal Layer 3 is the L8-SIGNAL, at thickness (0.6+4+1.2+3+1.2+4=) 14 mils
- Signal Layer 4 is the L9-SIGNAL, at thickness (0.6+6=) 6.6 mils
Figure 5–10 shows the results in the Quartus II software after you enter these PCB signal layers and thickness assignments.

**Figure 5–10. PCB Layers Specified in the Quartus II Software**

---

**Specifying Signal Breakout Layers**

Each user I/O pin in your FPGA device can break out at different layers on your PCB. In the Pin Planner, you can specify on which layers the I/O pins in your design break out. The breakout layer information is used only during SSN analysis and is not used in other processes run by the Quartus II software. To assign a pin to PCB layer, follow these steps:

1. On the Assignments menu, click **Pin Planner**.
2. If necessary, perform Analysis & Elaboration, Analysis & Synthesis, or fully compile the design to populate the Pin Planner with the node names in the design.
3. Right-click anywhere in the **All Pins** or **Groups** list, and then click **Customize Columns**.
4. Select the **PCB layer** column and move it from the **Available columns** list to the **Show these columns in this order** list.
5. Click **OK**.
6. In the **PCB layer** column, specify the PCB layer to which you want to connect the signal.
7. On the File menu, click **Save Project** to save the changes.

When you create PCB breakout layer assignments in the Pin Planner, you can assign the pin to any layer, even if you did not yet define the PCB layer.
Creating I/O Assignments

I/O assignments are required in FPGA design and are also used during SSN analysis to estimate voltage noise. Each input, output, or bidirectional signal in your design is assigned a physical pin location on the device using pin location assignments. Each signal has a physical I/O buffer that has a specific I/O standard, pin location, drive strength, and slew rate. The SSN Analyzer supports most I/O standards in a device family, such as the LVTTL and LVCMOS I/O standards.

The SSN Analyzer does not support differential I/O standards, such as the LVDS I/O standard and its variations, because differential I/O standards contribute a small amount of SSN.

For more information about supported I/O standards, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

For more information about creating and managing I/O assignments, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

Decreasing Pessimism in SSN Analysis

In the absence of specific timing information, the SSN Analyzer analyzes your design under worst-case conditions. Worst-case conditions include all pins acting as aggressor signals on all possible victim pins and all aggressor pins switching with the worst possible timing relationship. The results of SSN analysis under worst-case conditions are very pessimistic. You can improve the results of SSN Analysis by creating group assignments for specific types of pins. Use the following group assignments to decrease the pessimism in SSN analysis results:

- Assign pins to an output enable group—All pins in an output enable group must be either all input pins or all output pins. If all the pins in a group are always either all inputs or all outputs, it is impossible for an output pin in the group to cause SSN noise on an input pin in the group. You can assign pins to an output enable group with the Output Enable Group logic option.

- Assign pins to a synchronous group—I/O pins that are part of a synchronous group (signals that switch at the same time) may cause SSN, but do not result in any failures because the noise glitch occurs during the switching period of the signal. The noise, therefore, does not occur in the sampling window of that signal. You can assign pins to an output enable group with the Synchronous Group logic option. For example, in your design you have a bus with 32 pins that all belong to the same group. In a real operation, the bus switches at the same time, so any voltage noise induced by a pin on its groupmates does not matter, because it does not fall in the sampling window. If you do not assign the bus to a synchronous group, the other 31 pins can act as aggressors for the first pin in that group, leading to higher QL and QH noise levels during SSN analysis.

In some cases, the SSN Analyzer can detect the grouping for bidirectional pins by looking at the output enable signal of the bidirectional pins. However, Altera recommends that you explicitly specify the bidirectional groups and output groups in your design.
Performing SSN Analysis and Viewing Results

For more information about creating logic option assignments, refer to Assigning Device I/O Pins in Quartus II Help.

Excluding Pins as Aggressor Signals

The SSN Analyzer uses the following conditions to exclude pins as aggressor signals for a specific victim pin:

- A pin that is a complement of the victim pin. For example, any pin that is assigned a differential I/O standard cannot be an aggressor pin.
- A programming pin or JTAG pin because these pins are not active in user mode.
- Pins that have the same output enable signal as a bidirectional victim pin that the SSN Analyzer analyzes as an input pin. Pins with the same output enable signal also act as input pins and therefore cannot be aggressor pins at the same time. For information about grouping bidirectional pins, refer to “Performing SSN Analysis and Viewing Results”.
- Pins in the same synchronous group as a victim output pin. For information about grouping output pins, refer to “Performing SSN Analysis and Viewing Results”.
- A pin assigned the I/O Maximum Toggle Rate logic option with a frequency setting of zero. The SSN Analyzer does not consider pins with this setting as aggressor pins.

For more information about creating pin assignments with the Pin Planner, refer to Assigning Device I/O Pins in Quartus II Help.

Performing SSN Analysis and Viewing Results

You can perform SSN analysis either on your entire design, or you can limit the analysis to specific I/O banks.

If you know the problem area for SSN is within one I/O bank and you are changing pin assignments only in that bank, you can run SSN analysis for just that one I/O bank to reduce analysis time.

For more information, refer to Running the SSN Analyzer in Quartus II Help.

For more information about I/O bank numbering, refer to the appropriate device handbook available on the Literature and Technical Documentation page of the Altera website.

Understanding the SSN Reports

When SSN analysis is complete, you can view detailed analysis reports. The detailed messages in the reports help you understand and resolve SSN problems.

The SSN Analyzer section of the Compilation report contains information generated during SSN analysis, including the following reports:

- Summary
- Output Pins
Input Pins

Unanalyzed Pins

Confidence Metric Details

**Summary Report**

The Summary report summarizes the SSN Analyzer status and rates the SSN Analyzer confidence level as low, medium, or high. The confidence level depends on the completeness of your board trace model assignments. The more assignments you complete, the higher the confidence level. However, the confidence level does not always contribute to the accuracy of the QL and QH noise levels on a victim pin. The accuracy of QH and QL noise levels depends the accuracy of your board trace model assignments.

**Output Pins and Input Pins Reports**

The Output Pins report lists all of the output pins and bidirectional pins that are treated as output pins during SSN analysis. The Input Pins report lists all of the input pins and bidirectional pins that are treated as inputs during SSN analysis. Both reports list the location assignments for the pins treated as SSN outputs or inputs during SSN analysis, the QL and QH noise in volts, and what percentage the QL and QH margins are for the I/O standard used for that signal. The QH and QL noise margins that fall in the critical range (> 90%) are shown in red. The QH and QL noise margins that fall in the range of 70% to 90% are shown in gray.

**Unanalyzed Pins Report**

Not all pins are analyzed for SSN analysis. The following pins are not analyzed and are reported in the Unanalyzed Pins report:

- Pins assigned the LVDS I/O standard or any LVDS variations, such as the mini-LVDS I/O standard
- Pins created in the migration flow that cover power and supply pins in other packages
- The negative terminals of pseudo-differential I/O standards; the noise on differential standards is reported as the differential noise and is reported on the positive terminal

**Confidence Metric Details Report**

The Confidence Metric Details Report lists the values used during SSN Analysis for unspecified I/O, board, and PCB assignments.

**Viewing SSN Analysis Results in the Pin Planner**

After SSN analysis completes, you can analyze the results in the Pin Planner. In the Pin Planner you can identify the SSN hotspots in your device, as well as the QL and QH noise levels. The QL and QH results for each pin are displayed with a different color that represents whether the pin is below the warning threshold, below the critical threshold, or above the critical threshold. This color representation is also referred to as the SSN map of your FPGA device.
When you view the SSN map, you can customize which details to display, including input pins, output pins, QH signals, QL signals, and noise levels. You can also adjust the threshold levels for QH and QL noise voltages. Adjusting the threshold levels in the Pin Planner does not change the threshold levels reported during SSN analysis and does not change the data in any of the SSN reports.

You can also change I/O assignments and board trace information and rerun the SSN Analyzer to view the SSN analysis results based on those modified settings.

For more information, refer to Show SSN Analyzer Results and Running the SSN Analyzer in Quartus II Help.

Decreasing Processing Time for SSN Analysis

FPGA designs are getting larger in density, logic, and I/O count. The time it takes to complete SSN analysis and other Quartus II software processes affects your development time. Faster processing times can reduce your design cycle time. Use the following guidelines to reduce processing time:

- Direct the Quartus II software to use more than one processor for parallel executables, including the SSN Analyzer
- Perform SSN analysis after I/O assignment analysis if your design files and constraints are complete, and you are interested in generating the SSN results early in the design process and want to adjust I/O placements to see if you can obtain better results
- Perform SSN analysis after fitting if you want to view preliminary SSN results that do not take into account complete I/O assignment and I/O timing results
- Perform engineering change orders (ECOs) on your design, rather than recompiling the entire design, if you want to rerun SSN analysis after changing I/O assignments

For more information about using parallel processors, refer to Setting Up and Running Analysis and Synthesis and Compilation Process Settings Page in Quartus II Help. For more information about performing I/O assignment analysis, refer to Assigning Device I/O Pins in Quartus II Help. For more information about running the Fitter, refer to Setting Up and Running the Fitter in Quartus II Help.

For more information about performing ECOs on your design, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

Scripting Support

A Tcl script allows you to run procedures and determine settings described in this chapter. You can also run some of these procedures at a command prompt. The Quartus II software provides several packages to compile your design and create I/O assignments for analysis and fitting. You can create a custom Tcl script that maps the design and runs SSN analysis on your design.
For detailed information about specific scripting command options and Tcl API packages, type the following command at a system command prompt to run the Quartus II Command-Line and Tcl API Help browser:

```bash
grub --qhelp
```

For more information about Quartus II scripting support, including examples, refer to the Tcl Scripting and Command-Line Scripting chapters in volume 2 of the Quartus II Handbook and API Functions for Tcl in Quartus II Help.

### Optimizing Pin Placements for Signal Integrity

You can create an assignment that directs the Fitter to optimize pin placements for signal integrity with a Tcl command.

The following Tcl command directs the Fitter to optimize pin placement for signal integrity without affecting design fMAX:

```tcl
set_global_assignment -name OPTIMIZE_SIGNAL_INTEGRITY "Normal Compilation"
```

For more information, refer to “Optimizing Pin Placements for Signal Integrity” on page 5–9.

### Defining PCB Layers and PCB Layer Thickness

You can create PCB layer and thickness assignments with a Tcl command. shows Tcl commands for specifying PCB layer assignments.

**Example 5–3. Specifying PCB Layer Assignments**

```tcl
set_global_assignment -name PCB_LAYER_THICKNESS 0.00099822M -section_id 1
set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 2
set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 3
set_global_assignment -name PCB_LAYER_THICKNESS 0.00055372M -section_id 4
set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 5
set_global_assignment -name PCB_LAYER_THICKNESS 0.00034036M -section_id 6
set_global_assignment -name PCB_LAYER_THICKNESS 0.00082042M -section_id 7
```

These Tcl commands specify that there are seven PCB layers in the design, each with a different thickness. In each assignment, the letter M indicates the unit of measurement is millimeters. When you specify PCB layer assignments with Tcl commands, you must list the layers in consecutive order. For example, you would receive an error during SSN Analysis if your Tcl commands created the following assignments:

```tcl
set_global_assignment -name PCB_LAYER_THICKNESS 0.00099822M -section_id 1
set_global_assignment -name PCB_LAYER_THICKNESS 0.00082042M -section_id 7
```

To create assignments with the unit of measurement in mils, refer to the syntax in the following Tcl commands. These Tcl commands specify the same settings as shown in Figure 5–10 on page 5–13.

```tcl
set_global_assignment -name PCB_LAYER_THICKNESS 14.9MIL -section_id 1
set_global_assignment -name PCB_LAYER_THICKNESS 6.6MIL -section_id 2
set_global_assignment -name PCB_LAYER_THICKNESS 14MIL -section_id 3
set_global_assignment -name PCB_LAYER_THICKNESS 6.6MIL -section_id 4
```
For more information, refer to “Defining PCB Layers and PCB Layer Thickness” on page 5–11.

**Specifying Signal Breakout Layers**

You can create signal breakout layer assignments with a Tcl command. Example 5–4 shows Tcl commands for specifying signal breakout layer assignments:

**Example 5–4. Specifying Signal Breakout Layer Assignments**

```tcl
set_instance_assignment -name PCB_LAYER 10 -to e[2]
set_instance_assignment -name PCB_LAYER 3 -to e[3]
```

When you create PCB breakout layer assignments with Tcl commands, if you do not specify a PCB layer, or if you specify a PCB layer that does not exist, the SSN Analyzer breaks out the signal at the bottommost PCB layer.

If you create a PCB layer breakout assignment to a layer that does not exist, the SSN Analyzer will generate a warning message.

For more information, refer to “Specifying Signal Breakout Layers” on page 5–13.

**Decreasing Pessimism in SSN Analysis**

You can create output enable group and synchronous group assignments to help decrease pessimism during SSN Analysis with a Tcl command.

The following Tcl command assigns the bidirectional bus `DATAINOUT` to an output enable group:

```tcl
set_instance_assignment -name OUTPUT_ENABLE_GROUP 1 -to DATAINOUT
```

The following Tcl command assigns the bus `PCI_ADD_io` to a synchronous group:

```tcl
set_instance_assignment -name SYNCHRONOUS_GROUP 1 -to PCI_AD_io
```

For more information, refer to “Decreasing Pessimism in SSN Analysis” on page 5–14.

**Performing SSN Analysis**

You can perform SSN analysis with a command-line command. Use the `quartus_si` package that is provided with the Quartus II software.

Type the following command at a system command prompt to start the SSN Analyzer:

```bash
quartus_si <project name> r
```

To analyze just one I/O bank, type the following command at a system command prompt:

```bash
quartus_si <project revision> --bank = bank id
```

For example, to run analyze the I/O bank 2A type the following command:

```bash
quartus_si counter --bank=2A
```
Conclusion

To assist you with SSN Analysis, you can use the fast and accurate SSN Analyzer to help you estimate the SSN performance of your FPGA both early in the design cycle and when your PCB is complete. The SSN methodology discussed in this chapter gives you the tools you need to ensure your FPGA design meets your SSN requirements.

Document Revision History

Table 5–2 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Template update</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Reorganized and edited the chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added links to Quartus II Help for procedural information previously included in the chapter</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Added “Figure 6–9 shows the layout cross-section of a PCB in the Cadence Allegro PCB tool. The cross-section shows the stackup information of a PCB, which tells you the number of layers used in your PCB. The PCB shown in this example consists of various signal and circuit layers on which FPGA pins are routed, as well as the power and ground layers.” on page 6–12</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated for the Quartus II software 9.1 release</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Initial release</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
6. Signal Integrity Analysis with Third-Party Tools

Introduction

With the ever-increasing operating speed of interfaces in traditional FPGA design, the timing and signal integrity margins between the FPGA and other devices on the board must be within specification and tolerance before a single PCB is built. If the board trace is designed poorly or the route is too heavily loaded, noise in the signal can cause data corruption, while overshoot and undershoot can potentially damage input buffers over time.

As FPGA devices are used in high-speed applications, signal integrity and timing margin between the FPGA and other devices on the printed circuit board (PCB) are important aspects to consider to ensure proper system operation. To avoid time-consuming redesigns and expensive board respins, the topology and routing of critical signals must be simulated. The high-speed interfaces available on current FPGA devices must be modeled accurately and integrated into timing models and board-level signal integrity simulations. The tools used in the design of an FPGA and its integration into a PCB must be “board-aware”—able to take into account properties of the board routing and the connected devices on the board.

This chapter contains the following topics:

- “I/O Model Selection: IBIS or HSPICE” on page 6–3
- “FPGA to Board Signal Integrity Analysis Flow” on page 6–4
- “Simulation with IBIS Models” on page 6–7
- “Simulation with HSPICE Models” on page 6–16

The Quartus® II software provides methodologies, resources, and tools to ensure good signal integrity and timing margin between Altera® FPGA devices and other components on the board. Three types of analysis are possible with the Quartus II software:

- I/O timing with a default or user-specified capacitive load and no signal integrity analysis (default)
- The Quartus II Enable Advanced I/O Timing option utilizing a user-defined board trace model to produce enhanced timing reports from accurate “board-aware” simulation models
- Full board routing simulation in third-party tools using Altera-provided or generated Input/Output Buffer Information Specification (IBIS) or HSPICE I/O models
I/O timing using a specified capacitive test load requires no special configuration other than setting the size of the load. I/O timing reports from the Quartus II TimeQuest or the Quartus II Classic Timing Analyzer are generated based only on point-to-point delays within the I/O buffer and assume the presence of the capacitive test load with no other details about the board specified. The default size of the load is based on the I/O standard selected for the pin. Timing is measured to the FPGA pin with no signal integrity analysis details.

The **Enable Advanced I/O Timing** option expands the details in I/O timing reports by taking board topology and termination components into account. A complete point-to-point board trace model is defined and accounted for in the timing analysis. This ability to define a board trace model is an example of how the Quartus II software is “board-aware.”

In this case, timing and signal integrity metrics between the I/O buffer and the defined far end load are analyzed and reported in enhanced reports generated by the Quartus II TimeQuest Timing Analyzer.

For more information about defining capacitive test loads or how to use the **Enable Advanced I/O Timing** option to configure a board trace model, refer to the **I/O Management** chapter in volume 2 of the **Quartus II Handbook**.

This chapter focuses on the third type of analysis. The Quartus II software can export accurate HSPICE models with the built-in HSPICE Writer. You can run signal integrity simulations with these complete HSPICE models in Synopsys HSPICE. IBIS models of the FPGA I/O buffers are also created easily with the Quartus II IBIS Writer. You can integrate IBIS models into any third-party simulation tool that supports them, such as the Mentor Graphics® Hyperlynx software. With the ability to create industry-standard model definition files quickly, you can build accurate simulations that can provide data to help improve board-level signal integrity.

The I/O’s IBIS and HSPICE model creation available in the Quartus II software can help prevent problems before a costly board respin is required. In general, creating and running accurate simulations is difficult and time consuming. The tools in the Quartus II software automate the I/O model setup and creation process by configuring the models specifically for your design. With these tools, you can set up and run accurate simulations quickly and acquire data that helps guide your FPGA and board design.

The information about signal integrity in this chapter refers to board-level signal integrity based on I/O buffer configuration and board parameters, not simultaneous switching noise (SSN), also known as ground bounce or $V_{CC}$ sag. SSN is a product of multiple output drivers switching at the same time, causing an overall drop in the voltage of the chip’s power supply. This can cause temporary glitches in the specified level of ground or $V_{CC}$ for the device.

For more information about SSN and ways to prevent it, refer to **AN 315: Guidelines for Designing High-Speed FPGA PCBs**.
This chapter is intended for FPGA and board designers and includes details about the concepts and steps involved in getting designs simulated and how to adjust designs to improve board-level timing and signal integrity. Also included is information about how to create accurate models from the Quartus II software and how to use those models in simulation software.

The information in this chapter is meant for those who are familiar with the Quartus II software and basic concepts of signal integrity and the design techniques and components in good PCB design. Finally, you should know how to set up simulations and use your selected third-party simulation tool.

For information about basic signal integrity concepts and signal integrity details pertaining to Altera FPGA devices, refer to the Altera Signal Integrity Center.

I/O Model Selection: IBIS or HSPICE

The Quartus II software can export two different types of I/O models that are useful for different simulation situations. IBIS models define the behavior of input or output buffers through the use of voltage-current (V-I) and voltage-time (V-t) data tables. HSPICE models, often referred to as HSPICE decks, include complete physical descriptions of the transistors and parasitic capacitances that make up an I/O buffer along with all the parameter settings required to run a simulation. The HSPICE decks generated by the Quartus II software are preconfigured with the I/O standard, voltage, and pin loading settings for each pin in your design.

The choice of I/O model type is based on many factors. Table 6–1 shows a detailed comparison of the two I/O model types and information and examples of situations in which they might be used.

<table>
<thead>
<tr>
<th>Feature</th>
<th>IBIS Model</th>
<th>HSPICE Model</th>
</tr>
</thead>
<tbody>
<tr>
<td>I/O Buffer Description</td>
<td>Behavioral—I/O buffers are described by voltage-current and voltage-time tables in typical, minimum, and maximum supply voltage cases.</td>
<td>Physical—I/O buffers and all components in a circuit are described by their physical properties, such as transistor characteristics and parasitic capacitances, as well as their connections to one another.</td>
</tr>
<tr>
<td>Model Customization</td>
<td>Simple and limited—The model completely describes the I/O buffer and does not usually have to be customized.</td>
<td>Fully customizable—Unless connected to an arbitrary board description, the description of the board trace model must be customized in the model file. All parameters of the simulation are also adjustable.</td>
</tr>
<tr>
<td>Simulation Setup and Run Time</td>
<td>Fast—Simulations run quickly after set up correctly.</td>
<td>Slow—Simulations take time to set up and take longer to run and complete.</td>
</tr>
<tr>
<td>Simulation Accuracy</td>
<td>Good—for most simulations, accuracy is sufficient to make useful adjustments to the FPGA and/or board design to improve signal integrity.</td>
<td>Excellent—Simulations are highly accurate, making HSPICE simulation almost a requirement for any high-speed design where signal integrity and timing margins are tight.</td>
</tr>
<tr>
<td>Third-Party Tool Support</td>
<td>Excellent—Almost all third-party board simulation tools support IBIS.</td>
<td>Good—Most third-party tools that support SPICE support HSPICE. However, Synopsys HSPICE is required for simulations of Altera’s encrypted HSPICE models.</td>
</tr>
</tbody>
</table>
For more information about IBIS files created by the Quartus II IBIS Writer and IBIS files in general, as well as links to websites with detailed information, refer to AN 283: Simulating Altera Devices with IBIS Models.

FPGA to Board Signal Integrity Analysis Flow

Board signal integrity analysis can take place at any point in the FPGA design process and is often performed before and after board layout. If it is performed early in the process as part of a pre-PCB layout analysis, the models used for simulations can be more generic and can be changed as much as required to see how adjustments improve timing or signal integrity and help with the design and routing of the PCB. Simulations and the resulting changes made at this stage allow you to analyze “what if” scenarios to plan and implement your design better. To assist with early board signal integrity analysis, you can download generic IBIS model files for each device family and obtain HSPICE buffer simulation kits from the “Board Level Tools” section of the EDA Tool Support Resource Center.

Typically, if board signal integrity analysis is performed late in the design, it is used for a post-layout verification. The inputs and outputs of the FPGA are defined, and required board routing topologies and constraints are known. Simulations can help you find problems that might still exist in the FPGA or board design before fabrication and assembly. In either case, a simple process flow illustrates how to create accurate IBIS and HSPICE models from a design in the Quartus II software and transfer them to third-party simulation tools. Figure 6–1 shows this flow.
This chapter is organized around the type of model, IBIS or HSPICE, that you use for your simulations. When you understand the steps in the analysis flow, refer to the section of this chapter that corresponds to the model type you are using.

**Figure 6–1. Third-Party Board Signal Integrity Analysis Flow**

Create I/O and Board Trace Model Assignments

If your design uses a Stratix® III, Stratix II, or Cyclone® III device, you can configure a board trace model for output signals or for bidirectional signals in output mode and automatically transfer its description to HSPICE decks generated by the HSPICE Writer. This helps improve simulation accuracy.
To configure a board trace model, in the **Settings** dialog box, in the **TimeQuest Timing Analyzer** page, turn on the **Enable Advanced I/O Timing** option and configure the board trace model assignment settings for each I/O standard used in your design. You can add series or parallel termination, specify the transmission line length, and set the value of the far-end capacitive load. You can configure these parameters either in the Board Trace Model view of the Pin Planner, or click **Device and Pin Options** in the **Device** page of the **Settings** dialog box.

For information about how to use the **Enable Advanced I/O Timing** option and configure board trace models for the I/O standards used in your design, refer to the **I/O Management** chapter in volume 2 of the **Quartus II Handbook**.

The Quartus II software can generate IBIS models and HSPICE decks without having to configure a board trace model with the **Enable Advanced I/O Timing** option. In fact, IBIS models ignore any board trace model settings other than the far-end capacitive load. If any load value is set other than the default, the delay given by IBIS models generated by the IBIS Writer cannot be used to account correctly for the double counting problem. The load value mismatch between the IBIS delay and the $t_{CO}$ measurement of the Quartus II software prevents the delays from being safely added together. Warning messages displayed when the EDA Netlist Writer runs indicate when this mismatch occurs.

### Output File Generation

IBIS and HSPICE model files are not generated by the Quartus II software by default. To generate or update the files automatically during each project compilation, select the type of file to generate and a location where to save the file in the project settings. These settings can also be specified with commands in a Tcl script.

The IBIS and HSPICE Writers in the Quartus II software are run as part of the EDA Netlist Writer during normal project compilation. If either writer is turned on in the project settings, IBIS or HSPICE files are created and stored in the specified location. For IBIS, a single file is generated containing information about all assigned pins. HSPICE file generation creates separate files for each assigned pin. You can run the EDA Netlist Writer separately from a full compilation in the Quartus II software or at the command line. However, you must fully compile the project or perform I/O Assignment Analysis at least once for the IBIS and HSPICE Writers to have information about the I/O assignments and settings in the design.

### Customize the Output Files

The files generated by either the IBIS or HSPICE Writer are text files that you can edit and customize easily for design or experimentation purposes. IBIS files downloaded from the Altera website must be customized with the correct RLC values for the specific device package you have selected for your design. IBIS files generated by the IBIS Writer do not require this customization because they are configured automatically with the RLC values for your selected device. HSPICE decks require modification to include a detailed description of your board. With **Enable Advanced I/O Timing** turned on and a board trace model defined in the Quartus II software, generated HSPICE decks automatically include that model’s parameters. However, Altera recommends that you replace that model with a more detailed model that
describes your board design more accurately. A default simulation included in the generated HSPICE decks measures delay between the FPGA and the far-end device. You can make additions or adjustments to the default simulation in the generated files to change the parameters of the default simulation or to perform additional measurements.

**Set Up and Run Simulations in Third-Party Tools**

When you have generated the files, you can use them to perform simulations in your selected simulation tool. With IBIS models, you can apply them to input, output, or bidirectional buffer entities and quickly set up and run simulations. For HSPICE decks, the simulation parameters are included in the files. Open the files in Synopsys HSPICE and run simulations for each pin as required.

With HSPICE decks generated from the HSPICE Writer, the double counting problem is accounted for, which ensures that your simulations are accurate. Simulations that involve IBIS models created with anything other than the default loading settings in the Quartus II software must take the change in the size of the load between the IBIS delay and the Quartus II $t_{CO}$ measurement into account. Warning messages during compilation alert you to this change.

**Interpret Simulation Results**

If you encounter timing or signal integrity issues with your high-speed signals after running simulations, you can make adjustments to I/O assignment settings in the Quartus II software. These could include such things as drive strength or I/O standard, or making changes to your board routing or topology. After regenerating models in the Quartus II software based on the changes you have made, rerun the simulations to check whether your changes corrected the problem.

**Simulation with IBIS Models**

IBIS models provide a way to run accurate signal integrity simulations quickly. IBIS models describe the behavior of I/O buffers with voltage-current and voltage-time data curves. Because of their behavioral nature, IBIS models do not have to include any information about the internal circuit design of the I/O buffer. Most component manufacturers, including Altera, provide IBIS models for free download and use in signal integrity analysis simulation tools. You can download generic device family IBIS models from the Altera website for early design simulation or use the IBIS Writer to create custom IBIS models for your existing design.
Elements of an IBIS Model

An IBIS model file (.ibs) is a text file that describes the behavior of an I/O buffer across minimum, typical, and maximum temperature and voltage ranges with a specified test load. The tables and values specified in the IBIS file describe five basic elements of the I/O buffer. Figure 6–2 highlights each of these elements in the I/O buffer model.

![Figure 6–2. Five Basic Elements in IBIS Models](image)

The following elements correspond to each numbered block in Figure 6–2.

1. **Pulldown**—A voltage-current table describes the current when the buffer is driven low based on a pull-down voltage range of $-V_{CC}$ to $2V_{CC}$.

2. **Pullup**—A voltage-current table describes the current when the buffer is driven high based on a pull-up voltage range of $-V_{CC}$ to $V_{CC}$.

3. **Ground and Power Clamps**—Voltage-current tables describe the current when clamping diodes for electrostatic discharge (ESD) are present. The ground clamp voltage range is $-V_{CC}$ to $V_{CC}$, and the power clamp voltage range is $-V_{CC}$ to ground.

4. **Ramp and Rising/Falling Waveform**—A voltage-time ($dv/dt$) ratio describes the rise and fall time of the buffer during a logic transition. Optional rising and falling waveform tables can be added to more accurately describe the characteristics of the rising and falling transitions.

5. **Total Output Capacitance and Package RLC**—The total output capacitance includes the parasitic capacitances of the output pad, clamp diodes (if present), and input transistors. The package RLC is device package-specific and defines the resistance, inductance, and capacitance of the bond wire and pin of the I/O.

For more information about IBIS models and Altera-specific features, including links to the official IBIS specification, refer to AN 283: Simulating Altera Devices with IBIS Models.

Creating Accurate IBIS Models

There are two methods to obtain Altera device IBIS files for your board-level signal integrity simulations. You can download generic IBIS models from the Altera website or you can use the IBIS writer in the Quartus II software to create design-specific models.
Download IBIS Models

Altera provides IBIS models for almost all FPGA and FPGA configuration devices. Check the Altera IBIS Models page for information about whether models for your selected device are available. You can use the IBIS models from the website to perform early simulations of the I/O buffers you expect to use in your design as part of a pre-layout analysis.

Downloaded IBIS models have the RLC package values set to one particular device in each device family. To simulate your design with the model accurately, you must adjust the RLC values in the IBIS model file to match the values for your particular device package by performing the following steps:

1. Download and expand the ZIP file (.zip) of the IBIS model for the device family you are using for your design. The .zip file contains the .ibs file along with an IBIS model user guide and a model data correlation report.
2. Download the Package RLC Values spreadsheet for the same device family.
3. Open the spreadsheet and locate the row that describes the device package used in your design.
4. From the package’s I/O row, copy the minimum, maximum, and typical values of resistance, inductance, and capacitance for your device package.
5. Open the .ibs file in a text editor and locate the [Package] section of the file.
6. Overwrite the listed values copied with the values from the spreadsheet and save the file.

The .ibs file is now customized for your device package and can be used for any simulation. IBIS models downloaded and used for simulations in this manner are generic. They describe only a certain set of models listed for each device on the Altera IBIS Models page of the Altera website. To create customized models for your design, use the IBIS Writer as described in the next section.

Generate Custom IBIS Models with the IBIS Writer

If you have started your FPGA design and have created custom I/O assignments, such as drive strength settings or the enabling of clamping diodes for ESD protection, you can use the Quartus II IBIS Writer to create custom IBIS models to accurately reflect your assignments. IBIS models created with the IBIS Writer take I/O assignment settings into account.

If the Enable Advanced I/O Timing option is turned off, the generated .ibs files are based on the load value setting for each I/O standard on the Capacitive Loading page of the Device and Pin Options dialog box in the Device dialog box. With the Enable Advanced I/O Timing option turned on, IBIS models use an effective capacitive load based on settings found in the board trace model on the Board Trace Model page in the Device and Pin Options dialog box or the Board Trace Model view in the Pin Planner. The effective capacitive load is based on the sum of the Near capacitance, Transmission line distributed capacitance, and the Far capacitance settings in the board trace model. Resistances and transmission line inductance values are ignored.
If you made any changes from the default load settings, the delay in the generated IBIS model cannot safely be added to the Quartus II \( t_{CO} \) measurement to account for the double counting problem. This is because the load values between the two delay measurements do not match. When this happens, the Quartus II software displays warning messages when the EDA Netlist Writer runs to remind you about the load value mismatch.

For step-by-step instructions on how to generate IBIS models with the Quartus II software, refer to *Generating IBIS Output Files with the Quartus II Software* in Quartus II Help.

For more information about IBIS model generation, refer to the *AN 283: Simulating Altera Devices with IBIS Models* or to the Quartus II Help.

### Design Simulation Using the Mentor Graphics HyperLynx Software

You must integrate IBIS models downloaded from the Altera website ([www.altera.com](http://www.altera.com)) or created with the Quartus II IBIS Writer into board design simulations to accurately model timing and signal integrity. The HyperLynx software from Mentor Graphics is one of the most popular tools for design simulation. The HyperLynx software makes it easy to integrate IBIS models into simulations.

The HyperLynx software is a PCB analysis and simulation tool for high-speed designs, consisting of two products, LineSim and BoardSim. LineSim is an early simulation tool. Before any board routing takes place, LineSim is used to simulate “what if” scenarios to assist in creating routing rules and defining board parameters. BoardSim is a post-layout tool used to analyze existing board routing. Specific nets are selected from a board layout file and simulated in a manner similar to LineSim. With board and routing parameters, and surrounding signal routing known, highly accurate simulations of the final fabricated PCB are possible. This section focuses on LineSim. Because the process of creating and running simulations is very similar for both LineSim and BoardSim, the details of IBIS model use in LineSim applies to simulations in BoardSim.

Simulations in LineSim are configured using a schematic GUI to create connections and topologies between I/O buffers, route trace segments, and termination components. LineSim provides two methods for creating routing schematics: cell-based and free-form. Cell-based schematics are based on fixed cells consisting of typical placements of buffers, trace impedances, and components. Parts of the grid-based cells are filled with the desired objects to create the topology. A topology in a cell-based schematic is limited by the available connections within and between the cells.
A more robust and expandable way to create a circuit schematic for simulation is to use the free-form schematic format in LineSim as shown in Figure 6–3. The free-form schematic format makes it easy to place parts into any configuration and edit them as required. This section describes the use of IBIS models with free-form schematics, but the process is nearly identical for cell-based schematics.

**Figure 6–3. HyperLynx LineSim Free-Form Schematic Editor**

When you use HyperLynx software to perform simulations, you typically perform the following steps:

1. Create a new LineSim free-form schematic document and set up the board stackup for your PCB using the Stackup Editor. In this editor, specify board layer properties including layer thickness, dielectric constant, and trace width.

2. Create a circuit schematic for the net you want to simulate. The schematic represents all the parts of the routed net including source and destination I/O buffers, termination components, transmission line segments, and representations of impedance discontinuities such as vias or connectors.

3. Assign IBIS models to the source and destination I/O buffers to represent their behavior during operation.

4. Attach probes from the digital oscilloscope that is built in to LineSim to points in the circuit that you want to monitor during simulation. Typically, at least one probe is attached to the pin of a destination I/O buffer. For differential signals, you can attach a differential probe to both the positive and negative pins at the destination.

5. Configure and run the simulation. You can simulate a rising or falling edge and test the circuit under different drive strength conditions.
6. Interpret the results and make adjustments. Based on the waveforms captured in the digital oscilloscope, you can adjust anything in the circuit schematic to correct any signal integrity issues, such as overshoot or ringing. If necessary, you can make I/O assignment changes in the Quartus II software, regenerate the IBIS file with the IBIS Writer, and apply the updated IBIS model to the buffers in your HyperLynx software schematic.

7. Repeat the simulations and circuit adjustments until you are satisfied with the results. When the operation of the net meets your design requirements, implement changes to your I/O assignments in the Quartus II software and/or adjust your board routing constraints, component values, and placement to match the simulation.

For more information about HyperLynx software, including schematic creation, simulation setup, model usage, product support, licensing, and training, refer to HyperLynx Help or the Mentor Graphics website at www.mentor.com.

Configuring LineSim to Use Altera IBIS Models

You must configure LineSim to find and use the downloaded or generated IBIS models for your design. To do this, add the location of your .ibs file or files to the LineSim Model Library search path. Then you apply a selected model to a buffer in your schematic.

To add the Quartus II software’s default IBIS model location, <project directory>/board/ibis, to the HyperLynx LineSim model library search path, perform the following steps in LineSim:

1. From the Options menu, click Directories. The Set Directories dialog box appears (Figure 6–4). The Model-library file path(s) list displays the order in which LineSim searches file directories for model files.

Figure 6–4. LineSim Set Directories Dialog Box
2. Click **Edit**. A dialog box appears where you can add directories and adjust the order in which LineSim searches them (Figure 6–5).

**Figure 6–5. LineSim Select Directories Dialog Box**

3. Click **Add**

4. Browse to the default IBIS model location, `<project directory>/board/ibis`. Click **OK**.

5. Click **Up** to move the IBIS model directory to the top of the list. Click **Generate Model Index** to update LineSim’s model database with the models found in the added directory.

6. Click **OK**. The IBIS model directory for your project is added to the top of the Model-library file path(s) list.

7. To close the **Set Directories** dialog box, click **OK**.

**Integrating Altera IBIS Models into LineSim Simulations**

When the location for IBIS files has been set, you can assign the downloaded or generated IBIS models to the buffers in your schematic. To do this, perform the following steps:
1. Double-click a buffer symbol in your schematic to open the **Assign Models** dialog box (Figure 6–6). You can also click **Assign Models** from the buffer symbol’s right-click menu.

**Figure 6–6. LineSim Assign Model Dialog Box**

2. The pin of the buffer symbol you selected should be highlighted in the **Pins** list. If you want to assign a model to a different symbol or pin, select it from the list.

3. Click **Select**. The **Select IC Model** dialog box appears (Figure 6–7).

**Figure 6–7. LineSim Select IC Model Dialog Box**
4. To filter the list of available libraries to display only IBIS models, select .IBS. Scroll through the Libraries list, and click the name of the library for your design. By default, this is <project name>.ibs.

5. The device for your design should be selected as the only item in the Devices list. If not, select your device from the list.

6. From the Signal list, select the name of the signal you want to simulate. You can also choose to select by device pin number.

7. Click OK. The Assign Models dialog box displays the selected .ibs file and signal.

8. If applicable to the signal you chose, adjust the buffer settings as required for the simulation.

9. Select and configure other buffer pins from the Pins list in the same manner.

10. Click OK when all I/O models are assigned.

Running and Interpreting LineSim Simulations

You can now run any desired simulations and make adjustments to the I/O assignments or simulation parameters as required. For example, if you see too much overshoot in the simulated signal at the destination buffer after running a simulation (as shown in Figure 6–8), you could adjust the drive strength I/O assignment setting to a lower value. Regenerate the .ibs file, and run the simulation again to verify whether the change fixed the problem.

Figure 6–8. Example of Overshoot in HyperLynx with IBIS Models
If you see a discontinuity or other anomalies at the destination, such as slow rise and fall times (as shown in Figure 6–9), adjust the termination scheme or termination component values. After making these changes, rerun the simulation to check whether your adjustments solved the problem. In this case, it is not necessary to regenerate the .ibs file.

**Figure 6–9. Example of Signal Integrity Anomaly in HyperLynx with IBIS Models**

For more information about board-level signal integrity and to learn about ways to improve it with simple changes to your design, visit the Altera Signal Integrity Center.

**Simulation with HSPICE Models**

HSPICE decks are used to perform highly accurate simulations by describing the physical properties of all aspects of a circuit precisely. HSPICE decks describe I/O buffers, board components, and all of the connections between them, as well as defining the parameters of the simulation to be run. By their nature, to be effective, HSPICE decks are highly customizable and require a detailed description of the circuit under simulation. For devices that support advanced I/O timing, when **Enable Advanced I/O Timing** is turned on, the HSPICE decks generated by the Quartus II HSPICE Writer automatically include board components and topology defined in the Board Trace Model. Configure the board components and topology in the Pin Planner or in the **Board Trace Model** tab of the **Device and Pin Options** dialog box. All HSPICE decks generated by the Quartus II software include compensation for the double count problem. For more information about the double count problem, refer to “**The Double Counting Problem in HSPICE Simulations**” on page 6–17. You can simulate with the default simulation parameters built in to the generated HSPICE decks or make adjustments to customize your simulation.
Supported Devices and Signaling

Beginning with Quartus II software version 6.1 and later, the HSPICE Writer supports the devices and signaling defined in Table 6–2. Only Stratix III, Stratix II, and Cyclone III devices support the creation of a board trace model in the Quartus II software for automatic inclusion in an HSPICE deck. Other devices require the board description to be manually added to the HSPICE file.

Table 6–2. HSPICE Writer Device and Signaling Support

<table>
<thead>
<tr>
<th>Device</th>
<th>Input</th>
<th>Output</th>
<th>Single-Ended</th>
<th>Differential</th>
<th>Automatic Board Trace Model Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Stratix III</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>Stratix II GX (non-HSSI pins)</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Stratix II</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>HardCopy® II</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Cyclone III</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
</tr>
</tbody>
</table>

If you are using a Stratix II device for your design, you can turn on Enable Advanced I/O Timing and configure the board trace model for each I/O standard used in your design. Newer families have this feature turned on by default and it cannot be turned off. The HSPICE files include the board trace description you create in the Board Trace Model view in the Pin Planner or the Board Trace Model tab in the Device and Pin Options dialog box.

For more information about the Enable Advanced I/O Timing option and configuring board trace models for the I/O standards in your design, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

Accessing HSPICE Simulation Kits

You can access the available HSPICE models at the SPICE Models for Altera Devices web page and also with the Quartus II software’s HSPICE Writer tool. The Quartus II software HSPICE Writer tool removes many common sources of user error from the I/O simulation process. The HSPICE Writer tool automatically creates preconfigured I/O simulation spice decks that only require the addition of a user board model. All the difficult tasks required to configure the I/O modes and interpret the timing results are handled automatically by the HSPICE Writer tool.

The Double Counting Problem in HSPICE Simulations

Simulating I/Os using accurate models is extremely helpful for finding and fixing FPGA I/O timing and board signal integrity issues before any boards are built. However, the usefulness of such simulations is directly related to the accuracy of the models used and whether the simulations are set up and performed correctly. To ensure accuracy in models and simulations created for FPGA output signals, the timing hand-off between tCO timing in the Quartus II software and simulation-based board delay must be taken into account. If this hand-off is not handled correctly, the calculated delay could either count some of the delay twice or even miss counting some of the delay entirely.
Defining the Double Counting Problem

The double counting problem is inherent to the method output timing is analyzed versus the method used for HSPICE models. The timing analyzer tools in the Quartus II software measure delay timing for an output signal from the core logic of the FPGA design through the output buffer ending at the FPGA pin with a default capacitive load or a specified value for the selected I/O standard. This measurement is the t_{CO} timing variable as shown in Figure 6–10.

HSPICE models for board simulation measure t_{PD} (propagation delay) from an arbitrary reference point in the output buffer, through the device pin, out along the board routing, and ending at the signal destination.

It is apparent immediately that if these two delays were simply added together, the delay between the output buffer and the device pin would be counted twice in the calculation. A model or simulation that does not account for this double count would create overly pessimistic simulation results, because the double-counted delay can limit I/O performance artificially. To fix the problem, it might seem that simply subtracting the overlap between t_{CO} and t_{PD} would account for the double count. However, this adjustment would not be accurate because each measurement is based on a different load.

Input signals do not exhibit this problem because the HSPICE models for inputs stop at the FPGA pin instead of at the input buffer. In this case, simply adding the delays together produces an accurate measurement of delay timing.
The Solution to Double Counting

To adjust the measurements to account for the double-counting, the delay between the arbitrary point in the output buffer selected by the HSPICE model and the FPGA pin must be subtracted from either \( t_{CO} \) or \( t_{PD} \) before adding the results together. The subtracted delay must also be based on a common load between the two measurements. This is done by repeating the HSPICE model measurement, but with the same load used by the Quartus II software for the \( t_{CO} \) measurement. This second measurement, called \( t_{TESTLOAD} \), is illustrated with the top circuit in Figure 6–11.

With \( t_{TESTLOAD} \) known, the total delay for the output signal from the FPGA logic to the signal destination on the board, accounting for the double count, is calculated as shown in Equation 6–1.

**Equation 6–1.**

\[
t_{delay} = t_{CO} + (t_{PD} - t_{TESTLOAD})
\]

The preconfigured simulation files generated by the HSPICE Writer in the Quartus II software are designed to account for the double-counting problem based on this calculation automatically. Performing accurate timing simulations is easy without having to make adjustments for double counting manually.
HSPICE Writer Tool Flow

This section includes information to help you get started using the Quartus II software HSPICE Writer tool. The information in this section assumes you have a basic knowledge of the standard Quartus II software design flow, such as project and assignment creation, compilation, and timing analysis.

For additional information about standard design flows, refer to the appropriate sections of the *Quartus II Handbook*.

Applying I/O Assignments

The first step in the HSPICE Writer tool flow is to configure the I/O standards and modes for each of the pins in your design properly. In the Quartus II software, these settings are represented by assignments that map I/O settings, such as pin selection, and I/O standard and drive strength, to corresponding signals in your design.

The Quartus II software provides multiple methods for creating these assignments:

- Using the Pin Planner
- Using the assignment editor
- Manually editing the `.qsf` file
- By making assignments in a scripted Quartus II flow using Tcl

Enabling HSPICE Writer

You must enable the HSPICE Writer in the *Settings* dialog box of the Quartus II software (Figure 6–12) to generate the HSPICE decks from the Quartus II software.

Figure 6–12. EDA Tool Settings: Board Level Options Dialog Box
Enabling HSPICE Writer Using Assignments

You can also use HSPICE Writer in conjunction with a scripted Tcl flow. To enable HSPICE Writer during a full compile, include the lines shown in Example 6–1 in your Tcl script.

Example 6–1. Enable HSPICE Writer

```
set_global_assignment -name EDA_BOARD_DESIGN_SIGNAL_INTEGRITY_TOOL "HSPICE (Signal Integrity)"
set_global_assignment -name EDA_OUTPUT_DATA_FORMAT HSPICE -section_id eda_board_design_signal_integrity
set_global_assignment -name EDA_NETLIST_WRITER_OUTPUT_DIR <output_directory> -section_id eda_board_design_signal_integrity
```

As with command-line invocation, specifying the output directory is optional. If not specified, the output directory defaults to `board/hspice`.

Naming Conventions for HSPICE Files

HSPICE Writer automatically generates simulation files and names them using the following naming convention:

```
<device>_<pin #>_<pin_name>_<in/out>.sp
```

For bidirectional pins, two spice decks are produced; one with the I/O buffer configured as an input, and the other with the I/O buffer configured as an output.

The Quartus II software supports alphanumeric pin names that contain the underscore (_) and dash (-) characters. Any illegal characters used in file names are converted automatically to underscores.

The contents of the HSPICE files are described in detail in “Sample Output for I/O HSPICE Simulation Deck” on page 6–33 and “Sample Input for I/O HSPICE Simulation Deck” on page 6–28.

Invoking HSPICE Writer

After HSPICE Writer is enabled, the HSPICE simulation files are generated automatically each time the project is completely compiled. The Quartus II software also provides an option to generate a new set of simulation files without having to recompile manually. In the Processing menu, click Start EDA Netlist Writer to generate new simulation files automatically.

You must perform both Analysis & Synthesis and Fitting on a design before invoking the HSPICE Writer tool.
Invoking HSPICE Writer from the Command Line

If you use a script-based flow to compile your project, you can create HSPICE model files by including the commands shown in Example 6–2 in your Tcl script (.tcl file).

Example 6–2. Create HSPICE Model Files

```tcl
set_global_assignment -name EDA_BOARD_DESIGN_SIGNAL_INTEGRITY_TOOL "HSPICE (Signal Integrity)"
set_global_assignment -name EDA_OUTPUT_DATA_FORMAT HSPICE
set_global_assignment -name EDA_NETLIST_WRITER_OUTPUT_DIR <output_directory>
```

The `<output_directory>` option specifies the location where HSPICE model files are saved. By default, the `<project directory>/board/hspice` directory is used.

To invoke the HSPICE Writer tool through the command line, type the syntax shown in Example 6–3.

Example 6–3. Invoke HSPICE Writer

```bash
quartus_eda.exe <project_name> --board_signal_integrity=on --format=HSPICE --output_directory=<output_directory>
```

`<output_directory>` specifies the location where the generated spice decks will be written (relative to the design directory). This is an optional parameter and defaults to `board/hspice`.

Customizing Automatically Generated HSPICE Decks

HSPICE models generated by the HSPICE Writer can be used for simulation as generated. A default board description is included, and a default simulation is set up to measure rise and fall delays for both input and output simulations, which compensates for the double counting problem. However, Altera recommends that you customize the board description to more accurately represent your routing and termination scheme.

The sample board trace loading in the generated HSPICE model files must be replaced by your actual trace model before you can run a correct simulation. To do this, open the generated HSPICE model files for all pins you want to simulate and locate the section shown in Example 6–4.

Example 6–4. Sample Board Trace Section

```plaintext
* I/O Board Trace and Termination Description
  * - Replace this with your board trace and termination description
```

You must replace the example load with a load that matches the design of your PCB board. This includes a trace model, termination resistors, and, for output simulations, a receiver model. The spice circuit node that represents the pin of the FPGA package is called `pin`. The node that represents the far pin of the external device is called `load-in` (for output SPICE decks) and `source-in` (for input SPICE decks).
For an input simulation, you must also modify the stimulus portion of the spice file. The section of the file that must be modified is indicated in the comment block shown in Example 6–5.

**Example 6–5. Sample Source Stimulus Section**

```
* Sample source stimulus placeholder
* - Replace this with your I/O driver model
```

Replace the sample stimulus model with a model for the device that will drive the FPGA.

**Running an HSPICE Simulation**

Because simulation parameters are configured directly in the HSPICE model files, running a simulation requires only that you open an HSPICE file in the HSPICE user interface and start the simulation. The HSPICE user interface window is shown in Figure 6–13.

**Figure 6–13. HSPICE User Interface Window**

[Image of HSPICE User Interface Window]

Click **Open** and browse to the location of the HSPICE model files generated by the Quartus II HSPICE Writer. The default location for HSPICE model files is `<project directory>/board/hspice`. Select the `.sp` file generated by the HSPICE Writer for the signal you want to simulate. Click **OK**.

To run the simulation, click **Simulate**. The status of the simulation is displayed in the window and saved in an `.lis` file with the same name as the `.sp` file when the simulation is complete. Check the `.lis` file if an error occurs during the simulation requiring a change in the `.sp` file to fix.

**Interpreting the Results of an Output Simulation**

By default, the automatically generated output simulation spice decks are set up to measure three delays for both rising and falling transitions. Two of the measurements, `tpd_rise` and `tpd_fall`, measure the double-counting corrected delay from the FPGA pin to the load pin. To determine the complete clock-edge to load-pin delay, add these numbers to the Quartus II software reported default loading `tCO` delay.
The remaining four measurements, \texttt{tpd\_uncomp\_rise}, \texttt{tpd\_uncomp\_fall}, \texttt{t\_dblcnt\_rise}, and \texttt{t\_dblcnt\_fall}, are required for the double-counting compensation process and are not required for further timing usage. Refer to “Simulation Analysis” on page 6–33 for a description of these measurements.

**Interpreting the Results of an Input Simulation**

By default, the automatically generated input simulation SPICE decks are set up to measure delays from the source’s driver pin to the FPGA’s input pin for both rising and falling transitions. The propagation delay is reported by HSPICE measure statements as \texttt{tpd\_rise} and \texttt{tpd\_fall}. To determine the complete source driver pin-to-FPGA register delay, add these numbers to the Quartus II software reported \texttt{T\textsubscript{H}} and \texttt{T\textsubscript{SU}} input timing numbers.

**Viewing and Interpreting Tabular Simulation Results**

The \texttt{.lis} file stores the collected simulation data in tabular form. The default simulation configured by the HSPICE Writer produces delay measurements for rising and falling transitions on both input and output simulations. These measurements are found in the \texttt{.lis} file and named \texttt{tpd\_rise} and \texttt{tpd\_fall}. For output simulations, these values are already adjusted for the double count. To determine the complete delay from the FPGA logic to the load pin, add either of these measurements to the Quartus II \texttt{t\textsubscript{CO}} delay. For input simulations, add either of these measurements to the Quartus II \texttt{t\textsubscript{SU}} and \texttt{t\textsubscript{H}} delay values to calculate the complete delay from the far end stimulus to the FPGA logic. Other values found in the \texttt{.lis} file, such as \texttt{tpd\_uncomp\_rise}, \texttt{tpd\_uncomp\_fall}, \texttt{t\_dblcnt\_rise}, and \texttt{t\_dblcnt\_fall}, are parts of the double count compensation calculation. These values are not necessary for further analysis.

**Viewing Graphical Simulation Results**

You can view the results of the simulation quickly as a graphical waveform display using the AvanWaves viewer included with HSPICE. With the default simulation configured by the HSPICE Writer, you can view the simulated waveforms at both the source and destination in input and output simulations.
To see the waveforms for the simulation, in the HSPICE user interface window, click **AvanWaves**. The AvanWaves viewer opens and displays the **Results Browser** as shown in **Figure 6–14**.

**Figure 6–14. HSPICE AvanWaves Results Browser**

The **Results Browser** lets you select which waveform to view quickly in the main viewing window. If multiple simulations are run on the same signal, the list at the top of the **Results Browser** displays the results of each simulation. Click the simulation description to select which simulation to view. By default, the descriptions are derived from the first line of the HSPICE file, so the description might appear as a line of asterisks.

Select the type of waveform to view, by performing the following steps:

1. To see the source and destination waveforms with the default simulation, from the **Types** list, select **Voltages**.

2. On the **Curves** list, double-click the waveform you want to view. The waveform appears in the main viewing window.
You can zoom in and out and adjust the view as desired (Figure 6–15).

**Figure 6–15. AvanWaves Waveform Viewer**
Making Design Adjustments Based on HSPICE Simulations

Based on the results of your simulations, you can make adjustments to the I/O assignments or simulation parameters if required. For example, after you run a simulation and see overshoot or ringing in the simulated signal at the destination buffer as shown in the example in Figure 6–16, you can adjust the drive strength I/O assignment setting to a lower value. Regenerate the HSPICE deck, and run the simulation again to verify that the change fixed the problem.

Figure 6–16. Example of Overshoot in the AvanWaves Waveform Viewer
If there is a discontinuity or any other anomalies at the destination as shown in the example in Figure 6–17, adjust the board description in the Quartus II Board Trace Model (for Stratix II, Stratix III, or Cyclone III devices) or in the generated HSPICE model files to change the termination scheme or adjust termination component values. After making these changes, regenerate the HSPICE files if necessary, and rerun the simulation to verify whether your adjustments solved the problem.

Figure 6–17. Example of Signal Integrity Anomaly in the AvanWaves Waveform Viewer

For more information about board-level signal integrity and to learn about ways to improve it with simple changes to your FPGA design, refer to the Altera Signal Integrity Center.

Sample Input for I/O HSPICE Simulation Deck

The following sections examine a typical HSPICE simulation spice deck for an I/O of type input. Each section presents the simulation file one block at a time.

Header Comment

The first block of an input simulation spice deck is the header comment. The purpose of this block is to provide an easily readable summary of how the simulation file has been automatically configured by the Quartus II software.

This block has two main components: The first component summarizes the I/O configuration relevant information such as device, speed grade, and so on. The second component specifies the exact test condition that the Quartus II software assumes for the given I/O standard. Example 6–6 shows a header comment block.
Example 6–6. Header Comment Block

* Quartus II HSPICE Writer I/O Simulation Deck*

* This spice simulation deck was automatically generated by
* Quartus for the following IO settings:
* *
* Device: EP2S60F1020C3
* Speed Grade: C3
* Pin: AA4 (out96)
* Bank: IO Bank 6 (Row I/O)
* I/O Standard: LVTTL, 12mA
* OCT: Off
*
* Quartus II’s default I/O timing delays assume the following slow
* corner simulation conditions.
* *
* Specified Test Conditions For Quartus II Tco
* Temperature: 85C (Slowest Temperature Corner)
* Transistor Model: TT (Typical Transistor Corner)
* Vccn: 3.135V (Vccn_min = Nominal - 5%)
* Vccpd: 2.97V (Vccpd_min = Nominal - 10%)
* Load: No Load
* Vtt: 1.5675V (Voltage reference is Vccn/2)
*
* Note: The I/O transistors are specified to operate at least as
* fast as the TT transistor corner, actual production
* devices can be as fast as the FF corner. Any simulations
* for hold times should be conducted using the fast process
* corner with the following simulation conditions.
* Temperature: 0C (Fastest Commercial Temperature Corner **)
* Transistor Model: FF (Fastest Transistor Corner)
* Vccn: 1.98V (Vccn_hold = Nominal + 10%)
* Vccpd: 3.63V (Vccpd_hold = Nominal + 10%)
* Vtt: 0.95V (Vtt_hold = Vccn/2 - 40mV)
* Vcc: 1.25V (Vcc_hold = Maximum Recommended)
* Package Model: Short-circuit from pad to pin (no parasitics)
*
* Warnings:

Simulation Conditions

The simulation conditions block loads the appropriate process corner models for the
transistors. This condition is automatically set up for the slow timing corner and is
modified only if other simulation corners are desired. Example 6–7 shows a
simulation conditions block.

Example 6–7. Simulation Conditions Block

* Process Settings

.options brief
.inc ‘siii_tt.inc’ * TT process corner

Simulation Options

The simulation options block configures the simulation temperature and configures
HSPICE with typical simulation options. Example 6–8 shows a simulation options
block.
For a detailed description of these options, consult your HSPICE manual.

**Example 6–8. Simulation Options Block**

```plaintext
* Simulation Options
.options brief=0
.options badchr co=132 scale=1e-6 acct ingold=2 nomod dv=1.0
+     dcstep=1 absv=1e-3 absi=1e-8 probe csdf=2 accurate=1
+     converge=1
.temp 85
```

**Constant Definition**

The constant definition block of the simulation file instantiates the voltage sources that controls the configuration modes of the I/O buffer. **Example 6–9** shows a constant definition block.

**Example 6–9. Constant Definition Block**

```plaintext
* Constant Definition
voeb oeb 0 vc * Set to 0 to enable buffer output
vopdrain opdrain 0 0 * Set to vc to enable open drain
vrambh rambh 0 0 * Set to vc to enable bus hold
vrpullup rpullup 0 0 * Set to vc to enable weak pullup
vpcdp5 rpcdp5 0 rp5 * Set the IO standard
vpcdp4 rpcdp4 0 rp4
vpcdp3 rpcdp3 0 rp3
vpcdp2 rpcdp2 0 rp2
vpcdp1 rpcdp1 0 rp1
vpcdp0 rpcdp0 0 rp0
vpcdn4 rpcdn4 0 rn4
vpcdn3 rpcdn3 0 rn3
vpcdn2 rpcdn2 0 rn2
vpcdn1 rpcdn1 0 rn1
vpcdn0 rpcdn0 0 rn0
vdin din 0 0
```

*Where:*

- **Voltage source** `voeb` controls the output enable of the buffer and is set to `disabled` for inputs.
- **`vopdrain`** controls the open drain mode for the I/O.
- **`vrambh`** controls the bus hold circuitry in the I/O.
- **`vrpullup`** controls the weak pullup.
- The next 11 voltages sources control the I/O standard of the buffer and are configured through a later library call.
- **`vdin`** is not used on input pins because it is the data pin for the output buffer.
Buffer Netlist

The buffer netlist block (Example 6–10) of the simulation spice deck loads all the load models required for the corresponding input pin.

Example 6–10. Buffer Netlist Block

* IO Buffer Netlist
.include ‘vio_buffer.inc’

Drive Strength

The drive strength block (Example 6–11) of the simulation SPICE deck loads the configuration bits necessary to configure the I/O into the proper I/O standard and drive strengths. Although these settings are not relevant to an input buffer, they are provided to allow the SPICE deck to be modifiable to support bidirectional simulations.

Example 6–11. Drive Strength Block

* Drive Strength Settings
.lib ‘drive_select_hio.lib’ 3p3ttl_12ma

I/O Buffer Instantiation

The I/O buffer instantiation block of the simulation SPICE deck instantiates the necessary power supplies and I/O model components that are necessary to simulate the given I/O.
Example 6–12 shows I/O buffer instantiation.

**Example 6–12. I/O Buffer Instantiation**

I/O Buffer Instantiation

* Supply Voltages Settings
  .param vcn=3.135
  .param vpd=2.97
  .param vc=1.15

* Instantiate Power Supplies |
  vvcc       vcc       0     vc     * FPGA core voltage
  vvss       vss       0     0      * FPGA core ground
  vvccn      vccn      0     vcn    * IO supply voltage
  vvssn      vssn      0     0      * IO ground
  vvccpd     vccpd     0     vpd    * Pre-drive supply voltage

* Instantiate I/O Buffer
  xvio_buf din oeb opdrain die rambh
  + rpcdn4 rpcdn3 rpcdn2 rpcdn1 rpcdn0
  + rpcdp5 rpcdp4 rpcdp3 rpcdp2 rpcdp1 rpcdp0
  + rpullup vccn vccpd vcpad0 vio_buf

* Internal Loading on Pad
* - No loading on this pad due to differential buffer/support
*    circuitry

* I/O Buffer Package Model
* - Single-ended I/O standard on a Row I/O
  .lib ‘lib/package.lib’ hio
  xpkg die pin hio_pkg

**Board Trace and Termination**

The board trace and termination block of the simulation SPICE deck is provided only as an example (shown in Example 6–13). Replace this block with your own board trace and termination models.

**Example 6–13. Board Trace and Termination Block**

* I/O Board Trace and Termination Description
* - Replace this with your board trace and termination description

wtline pin vssn load vssn N=1 L=1 RLGCMODEL=tlinemodel
.MODEL tlinemodel W MODELTYPE=RLGC N=1 Lo=7.13n Co=2.85p
Rterm2 load vssn 1x

**Stimulus Model**

The stimulus model block of the simulation spice deck is provided only as a placeholder example (shown in Example 6–14). Replace this block with your own stimulus model. Options for this include an IBIS or HSPICE model, among others.

**Example 6–14. Stimulus Model Block**

* Sample source stimulus placeholder
* - Replace this with your I/O driver model

Vsource source 0 pulse(0 vcn 0s 0.4ns 0.4ns 8.5ns 17.4ns)
Simulation Analysis

The simulation analysis block (Example 6–15) of the simulation file is configured to measure the propagation delay from the source to the FPGA pin. Both the source and end point of the delay are referenced against the 50% VCCN crossing point of the waveform.

Example 6–15. Simulation Analysis Block

```
* Simulation Analysis Setup
* Print out the voltage waveform at both the source and the pin
  .print tran v(source)  v(pin)
  .tran 0.020ns 17ns
* Measure the propagation delay from the source pin to the pin
  * referenced against the 50% voltage threshold crossing point
  .measure TRAN tpd_rise TRIG v(source) val='vcn*0.5' rise=1 
  + TARG v(pin) val = 'vcn*0.5' rise=1
  .measure TRAN tpd_fall TRIG v(source) val='vcn*0.5' fall=1
  + TARG v(pin) val = 'vcn*0.5' fall=1
```

Sample Output for I/O HSPICE Simulation Deck

The following sections examine a typical HSPICE simulation SPICE deck for an I/O-type output. Each section presents the simulation file one block at a time.

Header Comment

The first block of an output simulation SPICE deck is the header comment, as shown in Example 6–16. The purpose of this block is to provide a readable summary of how the simulation file has been automatically configured by the Quartus II software.

This block has two main components:

- The first component summarizes the I/O configuration relevant information such as device, speed grade, and so on.

- The second component specifies the exact test condition that the Quartus II software assumes when generating tCO delay numbers. This information is used as part of the double-counting correction circuitry contained in the simulation file.

The SPICE decks are preconfigured to calculate the slow process corner delay but can also be used to simulate the fast process corner as well. The fast corner conditions are listed in the header under the notes section.
The final section of the header comment lists any warning messages that you must consider when you use the SPICE decks.

**Example 6–16. Header Comment Block**

* Quartus II HSPICE Writer I/O Simulation Deck
  *
  * This spice simulation deck was automatically generated by
  * Quartus II for the following IO settings:
  *
  * Device: EP2S60F1020C3
  * Speed Grade: C3
  * Pin: AA4 (out96)
  * Bank: IO Bank 6 (Row I/O)
  * I/O Standard: LVTTL, 12mA
  * OCT: Off
  *
  * Quartus’ default I/O timing delays assume the following slow
  * corner simulation conditions.
  *
  * Specified Test Conditions For Quartus II Tco
  * Temperature: 85°C (Slowest Temperature Corner)
  * Transistor Model: TT (Typical Transistor Corner)
  * Vccn: 3.135V (Vccn_min = Nominal - 5%)
  * Vccpd: 2.97V (Vccpd_min = Nominal - 10%)
  * Load: No Load
  * Vtt: 1.5675V (Voltage reference is Vccn/2)
  *
  * For C3 devices, the TT transistor corner provides an
  * approximation for worst case timing. However, for functionality
  * simulations, it is recommended that the SS corner be simulated
  * as well.
  *
  * Note: The I/O transistors are specified to operate at least as
  * fast as the TT transistor corner, actual production
  * devices can be as fast as the FF corner. Any simulations
  * for hold times should be conducted using the fast process
  * corner with the following simulation conditions.
  *
  * Temperature: 0°C (Fastest Commercial Temperature Corner
  **)
  *
  * Transistor Model: FF (Fastest Transistor Corner)
  * Vccn: 1.98V (Vccn_hold = Nominal + 10%)
  * Vccpd: 3.63V (Vccpd_hold = Nominal + 10%)
  * Vtt: 0.95V (Vtt_hold = Vccn/2 - 40mV)
  * Vcc: 1.25V (Vcc_hold = Maximum Recommended)
  *
  * Package Model: Short-circuit from pad to pin
  *
  * Warnings:

**Simulation Conditions**

The simulation conditions block (Example 6–17) loads the appropriate process corner models for the transistors. This condition is automatically set up for the slow timing corner and must be modified only if other simulation corners are desired.
Two separate corners cannot be simulated at the same time. Instead, simulate the base case using the Quartus corner as one simulation and then perform a second simulation using the desired customer corner. The results of the two simulations can be manually added together.

Example 6–17. Simulation Conditions Block

```
* Process Settings
.options brief
.inc 'sii_tt.inc' * typical-typical process corner
```

Simulation Options

The simulation options block (Example 6–18) configures the simulation temperature and configures HSPICE with typical simulation options.

For a detailed description of these options, consult your HSPICE manual.

Example 6–18. Simulation Options Block

```
* Simulation Options
.options brief=0
.options badchr co=132 scale=1e-6 acct ingold=2 nomod dv=1.0 +
    dcstep=1 absv=1e-3 absi=1e-8 probe csdf=2 accurate=1 +
    converge=1
.temp 85
```
**Constraint Definition**

The constant definition block (Example 6–19) of the output simulation SPICE deck instantiates the voltage sources that control the configuration modes of the I/O buffer.

**Example 6–19. Constant Definition Block**

```
* Constant Definition

voeb       oeb       0     0 * Set to 0 to enable buffer output
vopdrain   opdrain   0     0 * Set to vc to enable open drain
vrambh     rambh     0     0 * Set to vc to enable bus hold
vrpullup   rpullup   0     0 * Set to vc to enable weak pullup
vpci       rpci      0     0 * Set to vc to enable pci mode
vpcdp4     rpcdp4    0     rp4  * These control bits set the I/O standard
vpcdp3     rpcdp3    0     rp3
vpcdp2     rpcdp2    0     rp2
vpcdp1     rpcdp1    0     rp1
vpcdp0     rpcdp0    0     rp0
vpcdn4     rpcdn4    0     rn4
vpcdn3     rpcdn3    0     rn3
vpcdn2     rpcdn2    0     rn2
vpcdn1     rpcdn1    0     rn1
vpcdn0     rpcdn0    0     rn0
vdin       din       0     pulse(0 vc 0s 0.2ns 0.2ns 8.5ns 17.4ns)
```

*Where:*

- Voltage source `voeb` controls the output enable of the buffer.
- `vopdrain` controls the open drain mode for the I/O.
- `vrambh` controls the bus hold circuitry in the I/O.
- `vrpullup` controls the weak pullup.
- `vpci` controls the PCI clamp.
- The next ten voltage sources control the I/O standard of the buffer and are configured through a later library call. Stratix III and Cyclone III devices have more bits and so might have more voltage sources listed in the constant definition block. They also have slew rate and delay chain settings.
- `vdin` is connected to the data input of the I/O buffer.
- The edge rate of the input stimulus is automatically set to the correct value by the Quartus II software.

**I/O Buffer Netlist**

The I/O buffer netlist block (Example 6–20) loads all of the models required for the corresponding pin. These include a model for the I/O output buffer, as well as any loads that might be present on the pin.

**Example 6–20. I/O Buffer Netlist Block**

```
* I/O Buffer Netlist

.include 'hio_buffer.inc'
.include 'lvds_input_load.inc'
.include 'lvds_oct_load.inc'
```
Drive Strength

The drive strength block (Example 6–21) of the simulation spice deck loads the configuration bits for configuring the I/O to the proper I/O standard and drive strength. These options are set by the HSPICE Writer tool and are not changed for expected use.

Example 6–21. Drive Strength Block

* Drive Strength Settings
.lib ‘drive_select_hio.lib’ 3p3ttl_12ma

Slew Rate and Delay Chain

Stratix III and Cyclone III devices have sections for configuring the slew rate and delay chain settings (Example 6–22).

Example 6–22. Slew Rate and Delay Chain Settings

* Programmable Output Delay Control Settings
.lib ‘lib/output_delay_control.lib’ no_delay

* Programmable Slew Rate Control Settings
.lib ‘lib/slew_rate_control.lib’ slow_slow
I/O Buffer Instantiation

The I/O buffer instantiation block (Example 6–23) of the output simulation spice deck instantiates the necessary power supplies and I/O model components that are necessary to simulate the given I/O.

**Example 6–23. I/O Buffer Instantiation Block**

* I/O Buffer Instantiation

* Supply Voltages Settings
  .param vcn=3.135
  .param vpd=2.97
  .param vc=1.15

* Instantiate Power Supplies
  vvcc       vcc       0     vc     * FPGA core voltage
  vvss       vss       0     0      * FPGA core ground
  vvccn      vccn      0     vcn    * IO supply voltage
  vvssn      vssn      0     0      * IO ground
  vvccpd     vccpd     0     vpd    * Pre-drive supply voltage

* Instantiate I/O Buffer
  xhio_buf  din  oeb  opdrain  die  rambh
  + rpcdn4  rpcdn3  rpcdn2  rpcdn1  rpcdn0
  + rpcdp4  rpcdp3  rpcdp2  rpcdp1  rpcdp0
  + rpullup  vccn  vccpd  vcpad0  hio_buf

* Internal Loading on Pad
  * - This pad has an LVDS input buffer connected to it, along
  *   with differential OCT circuitry. Both are disabled but
  *   introduce loading on the pad that is modeled below.
  xlvds_input_load  die  vss  vccn  lvds_input_load
  xlvds_oct_load  die  vss  vccpd  vccn  vcpa0  vccn  lvds_oct_load

* I/O Buffer Package Model
  * - Single-ended I/O standard on a Row I/O
  .lib 'lib/package.lib' hio
  xpkg  die  pin  hio_pkg

Board and Trace Termination

The board trace and termination block (Example 6–24) of the simulation SPICE deck is provided only as an example. Replace this block with your specific board loading models.

**Example 6–24. Board Trace and Termination Block**

* I/O Board Trace And Termination Description
  * - Replace this with your board trace and termination description

wtline  pin  vssn  load  vssn  N=1  L=1  RLGC=MODEL=linemodel
.MODEL  linemodel  W  MODELTYPE=RLGC  N=1  Lo=7.13n  Co=2.85p
Rterm2  load  vssn  1x
Double-Counting Compensation Circuitry

The double-counting compensation circuitry block (Example 6–25) of the simulation SPICE deck instantiates a second I/O buffer that is used to measure double-counting. The buffer is configured identically to the user I/O buffer but is connected to the Quartus II software test load. The simulated delay of this second buffer can be interpreted as the amount of double-counting between the Quartus II software and HSPICE Writer simulated results.

As the amount of double-counting is constant for a given I/O standard on a given pin, consider separating the double-counting circuitry from the simulation file. In doing so, you can perform any number of I/O simulations while referencing the delay only once. For more information about the double-counting problem, refer to “The Double Counting Problem in HSPICE Simulations” on page 6–17.

Example 6–25. (Part 1 of 2) Double-Counting Compensation Circuitry Block

* Double Counting Compensation Circuitry
  * The following circuit is designed to calculate the amount of double counting between Quartus II and the HSPICE models. If you have not changed the default simulation temperature or transistor corner the double counting will be automatically compensated by this spice deck. In the event you wish to simulate an IO at a different temperature or transistor corner you will need to remove this section of code and manually account for double counting. A description of Altera’s recommended procedure for this process can be found in the Quartus II HSPICE Writer AppNote.

* Supply Voltages Settings
  .param vcn_tl=3.135
  .param vpd_tl=2.97

* Test Load Constant Definition
  vopdrain_tl opdrain_tl 0 0
  vrambh_tl rambh_tl 0 0
  vrpullup_tl rpullup_tl 0 0

* Instantiate Power Supplies
  vvccn_tl vccn_tl 0 vcn_tl
  vvssn_tl vssn_tl 0 0
  vvccpd_tl vccpd_tl 0 vpd_tl

* Instantiate I/O Buffer
  xhio_testload din oeb opdrain_tl die_tl rambh_tl
  + rpcdn4 rpcdn3 rpcdn2 rpcdn1 rpcdn0
  + rpcdp4 rpcdp3 rpcdp2 rpcdp1 rpcdp0
  + rpullup_tl vccn_tl vccpd_tl vcpad0_tl hio_buf

* Internal Loading on Pad
  xlvsd_input_testload die_tl vss vccn_tl lvds_input_load
  xlvsd_oct_testload die_tl vss vccpd_tl vccn_tl vcpad0_tl vccn_tl lvds_oct_load
Simulation Analysis Block

The simulation analysis block (Example 6–26) is set up to measure double-counting corrected delays. This is accomplished by measuring the uncompensated delay of the I/O buffer when connected to the user load, and when subtracting the simulated amount of double-counting from the test load I/O buffer.

Example 6–26. Simulation Analysis Block

* Simulation Analysis Setup

* Print out the voltage waveform at both the pin and far end load
  .print tran v(pin) v(load)
  .tran 0.020ns 17ns

* Measure the propagation delay to the load pin. This value will include some double counting with Quartus II’s Tco
  .measure TRAN tpd_uncomp_rise TRIG v(din) val='vc*0.5’ rise=1 + TARG v(load) val='vcn*0.5’ rise=1
  .measure TRAN tpd_uncomp_fall TRIG v(din) val='vc*0.5’ fall=1 + TARG v(load) val='vcn*0.5’ fall=1

* The test load buffer can calculate the amount of double counting
  .measure TRAN t_dblcnt_rise TRIG v(din) val='vc*0.5’ rise=1 + TARG v(pin_tl) val='vcn_tl*0.5’ rise=1
  .measure TRAN t_dblcnt_fall TRIG v(din) val='vc*0.5’ fall=1 + TARG v(pin_tl) val='vcn_tl*0.5’ fall=1

* Calculate the true propagation delay by subtraction
  .measure TRAN tpd_rise PARAM='tpd_uncomp_rise-t_dblcnt_rise’
  .measure TRAN tpd_fall PARAM='tpd_uncomp_fall-t_dblcnt_fall’

Advanced Topics

The information in this section describes some of the more advanced topics and methods employed when setting up and running HSPICE simulation files.

PVT Simulations

The automatically generated HSPICE simulation files are set up to simulate the slow process corner using low voltage, high temperature, and slow transistors. To ensure a fully robust link, Altera recommends that you run simulations over all process corners.

To perform process, voltage, and temperature (PVT) simulations, manually modify the spice decks in a two step process:
1. Remove the double-counting compensation circuitry from the simulation file. This is required as the amount of double-counting is dependant upon how the Quartus II software calculates delays and is not based on which PVT corner is being simulated. By default, the Quartus II software provides timing numbers using the slow process corner.

2. Select the proper corner for the PVT simulation by setting the correct HSPICE temperature, changing the supply voltage sources, and loading the correct transistor models.

A more detailed description of HSPICE process corners can be found in the family-specific HSPICE model documentation. This document is available online with the HSPICE models as described in “Accessing HSPICE Simulation Kits” on page 6–17.

**Hold Time Analysis**

Alterna recommends performing worst-case hold time analysis using the fast corner models, which use fast transistors, high voltage, and low temperature. This involves modifying the SPICE decks to select the correct temperature option, change the supply voltage sources, and load the correct fast transistor models. The values of these parameters are located in the header comment section of the corresponding simulation deck files.

For a truly worst-case analysis, combine the HSPICE Writer hold time analysis results with the Quartus II software fast timing model. This requires that you change the double-counting compensation circuitry in the simulations files to also simulate the fast process corners, as this is what the Quartus II software uses for the fast timing model.

This method of hold time analysis is recommended only for globally synchronous buses. Do not apply this method of hold-time analysis to source synchronous buses. This is because the source synchronous clocking scheme is designed to cancel out some of the PVT timing effects. If this is not taken into account, the timing results will not be accurate. Proper source synchronous timing analysis is beyond the scope of this document.

**I/O Voltage Variations**

Use each of the FPGA family datasheets to verify the recommended operating conditions for supply voltages. For current FPGA families, the maximum recommended voltage corresponds to the fast corner, while the minimum recommended voltage corresponds to the slow corner. These voltage recommendations are specified at the power pins of the FPGA and are not necessarily the same voltage that are seen by the I/O buffers due to package IR drops.

The automatically generated HSPICE simulation files model this IR effect pessimistically by including a 50-mV IR drop on the VCCPD supply when a high drive strength standard is being used.

**Correlation Report**

Correlation reports for the HSPICE I/O models are located in the family-specific HSPICE I/O buffer simulation kits. Refer to “Accessing HSPICE Simulation Kits” on page 6–17 for additional information.
Conclusion

As FPGA devices are used in more high-speed applications, it becomes increasingly necessary to perform board-level signal integrity analysis simulations. You must run such simulations to ensure good signal integrity between the FPGA and any connected devices. The Quartus II software helps to simplify this process with the ability to automatically generate I/O buffer description models easily with the IBIS and HSPICE Writers. IBIS models can be integrated into a third-party signal integrity analysis workflow using a tool such as Mentor Graphics HyperLynx software, generating quick and accurate simulation results. HSPICE decks include preconfigured simulations and only require descriptions of board routing and stimulus models to create highly accurate simulation results using Synopsys HSPICE. Either type of simulation helps prevent unnecessary board spins, increasing your productivity and decreasing your costs.

Document Revision History

Table 6–3 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
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<td>8.1.0</td>
<td>■ Changed to 8-1/2 x 11 page size.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information for Stratix III devices.</td>
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<td></td>
<td></td>
<td>■ Input signals for Cyclone III devices are supported.</td>
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<td>May 2008</td>
<td>8.0.0</td>
<td>■ Updated “Introduction” on page 12–1.</td>
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<td>■ Updated Figure 12–1.</td>
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<td>■ Updated Figure 12–3.</td>
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<td>■ Updated Figure 12–13.</td>
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<td>■ Updated “Output File Generation” on page 12–6.</td>
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<td>■ Updated “Simulation with HSPICE Models” on page 12–17.</td>
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<td>■ Updated “Invoking HSPICE Writer from the Command Line” on page 12–22.</td>
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<td>■ Added “Sample Input for I/O HSPICE Simulation Deck” on page 12–29.</td>
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<td></td>
<td>■ Added “Sample Output for I/O HSPICE Simulation Deck” on page 12–33.</td>
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<td></td>
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<td>■ Updated “Correlation Report” on page 12–41.</td>
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<td></td>
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<td>■ Added hyperlinks to referenced documents and websites throughout the chapter.</td>
</tr>
<tr>
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<td>■ Made minor editorial updates.</td>
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</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter discusses how the Quartus® II software interacts with the Mentor Graphics® I/O Designer software and the DxDesigner software to provide a complete FPGA-to-board design workflow.

With today’s large, high-pin-count and high-speed FPGA devices, good and correct PCB design practices are essential to ensure correct system operation. The PCB design takes place concurrently with the design and programming of the FPGA. The FPGA or ASIC designer initially creates signal and pin assignments, and the board designer must correctly transfer these assignments to the symbols in their system circuit schematics and board layout. As the board design progresses, Altera recommends reassigning pins to optimize the PCB layout. Ensure that you inform the FPGA designer of the pin reassignments so that the new assignments are included in an updated placement and routing of the design.

The Mentor Graphics I/O Designer software allows you to take advantage of the full FPGA symbol design, creation, editing, and back-annotation flow supported by the Mentor Graphics tools.

This chapter covers the following topics:

- Performing design flow between the Quartus II software, the Mentor Graphics I/O Designer software, and the DxDesigner software
- Setting up the Quartus II software to create the design flow files
- Creating an I/O Designer database project to incorporate the Quartus II software signal and pin assignment data
- Updating signal and pin assignment changes between the I/O Designer software and the Quartus II software
- Generating symbols in the I/O Designer software
- Creating symbols in the DxDesigner software from the Quartus II software output files without the use of the I/O Designer software

This chapter is intended for board design and layout engineers who want to start the FPGA board integration while the FPGA is still in the design phase. Alternatively, the board designer can plan the FPGA pin-out and routing requirements in the Mentor Graphics tools and pass the information back to the Quartus II software for placement and routing. Part librarians can also benefit from this chapter by learning how to use output from the Quartus II software to create new library parts and symbols.

The procedures in this chapter require the following software:

- The Quartus II software version 5.1 or later
- DxDesigner software version 2004 or later
To obtain and license the Mentor Graphics tools and for product information, support, and training, refer to the Mentor Graphics website (www.mentor.com).

FPGA-to-PCB Design Flow

You can create a design flow integrating an Altera® FPGA design from the Quartus II software, and a circuit schematic in the DxDesigner software. Figure 7–1 shows the design flow with and without the I/O Designer software.

Figure 7–1. Design Flow with and Without the I/O Designer Software

Note to Figure 7–1:
1. The Quartus II software generates the .fx in the output directory you specify in the Board-Level page of the Settings dialog box. However, the Quartus II software and the I/O Designer software can import pin assignments from an .fx located in any directory. Altera recommends working with a backup .fx to prevent overwriting existing assignments or importing invalid assignments.

To perform the design flow shown in Figure 7–1, follow these steps:

1. In the Quartus II software, set up the board-level assignment settings to generate an .fx for symbol generation.

2. Compile your design to generate the .fx and Pin-Out File (.pin). You can locate the generated .fx and .pin files in the Quartus II project directory.
3. Create a board design with the DxDesigner software and the I/O Designer software by performing the following steps:
   a. Create a new I/O Designer database based on the .fx and the .pin files.
   b. In the I/O Designer software, make adjustments to signal and pin assignments.
   c. Regenerate the .fx in the I/O Designer software to export the I/O Designer software changes to the Quartus II software.
   d. Generate a single or fractured symbol for use in the DxDesigner software.
   e. Add the symbol to the sym directory of a DxDesigner project, or specify a new DxDesigner project with the new symbol.
   f. Instantiate the symbol in your DxDesigner schematic and export the design to the board layout tool.
   g. Back-annotate pin changes created in the board layout tool to the DxDesigner software and back to the I/O Designer software and the Quartus II software.

4. Create a board design with the DxDesigner software without the I/O Designer software by performing the following steps:
   a. Create a new DxBoardLink symbol with the Symbol wizard and reference the .pin from the Quartus II software in an existing DxDesigner project.
   b. Instantiate the symbol in your DxDesigner schematic and export the design to a board layout tool.

   You can update these symbols with design changes with or without the I/O Designer software. If you use the Mentor Graphics I/O Designer software and you change symbols with the DxDesigner software, you must reimport the symbols into I/O Designer to avoid overwriting your symbol changes.

**Performing Simultaneous Switching Noise (SSN) Analysis of Your FPGA**

With the Quartus II software, you can extract pin assignment data and perform SSN analysis of your design for designs targeting the Stratix III device family. You can perform SSN analysis early in the board layout stage as part of your overall pin planning process; however, you do not have to perform SSN analysis to generate pin assignment data from the Quartus II software. You can use the SSN Analyzer tool in the Quartus II software to optimize the pin assignments for better SSN performance.

For more information about the SSN Analyzer, refer to the *Simultaneous Switching Noise (SSN) Analysis and Optimizations* chapter in volume 2 of the *Quartus II Handbook* and *About the SSN Analyzer* in Quartus II Help.
Setting Up the Quartus II Software

You can transfer pin and signal assignments from the Quartus II software to the Mentor Graphics tools by generating `.pin` and `.fx` files (refer to Figure 7–2). The `.pin` is an output file generated by the Quartus II Fitter that contains pin assignment information. You can use the Quartus II Pin Planner to set and change the assignments contained in the `.pin` and then transfer the assignments to the Mentor Graphics tools. You cannot, however, import pin assignment changes from the Mentor Graphics tools into the Quartus II software with the `.pin`.

The `.pin` lists all used and unused pins on your selected Altera device. It also provides the following basic information fields for each assigned pin on the device:

- Pin signal name and usage
- Pin number
- Signal direction
- I/O standard
- Voltage
- I/O bank
- User or Fitter-assigned

The `.fx` is an input/output file generated by the Quartus II software and the I/O Designer software that can be imported and exported from both programs. The `.fx` generated by the Quartus II software lists only assigned pins and provides the following advanced information fields for each pin on a device:

- Pin number
- I/O bank
- Signal name
- Signal direction
- I/O standard
- Drive strength (mA)
- Termination enabling
- Slew rate
- IOB delay
- Swap group
- Differential pair type

The `.fx` generated by the I/O Designer software lists all pins, including unused pins. In addition to the advanced information fields listed above, the `.fx` generated by the Mentor Graphics I/O Designer software also includes the following information fields:

- Device pin name
- Pin set
- Pin set position
Pin set group
Super pin set group
Super pin set position

For more information about .fx files and the information fields added by the Mentor Graphics software, refer to *FPGA Xchange-Format File (.fx) Definition* in Quartus II Help and Mentor Graphics website (www.mentor.com) respectively.

The I/O Designer software can also read from or update a Quartus II Settings File (.qsf). The design flow uses the .qsf in a similar manner to the .fx, but does not transfer pin swap group information between the I/O Designer software and the Quartus II software.

Because the .qsf contains additional information about your project that the Mentor Graphics I/O Designer software does not use, Altera recommends using the .fx instead of the .qsf.

For more information about the .qsf, refer to *Quartus II Settings File (.qsf) Definition* in Quartus II Help.

**Generating a .pin File**

The Quartus II software automatically generates the .pin after compiling your FPGA design or during I/O assignment analysis.

To start I/O assignment analysis, on the Processing menu, point to Start and then click **Start I/O Assignment Analysis**. The Quartus II Fitter generates the .pin and places the file in your Quartus II design directory with the name `<project name>.pin`.

The Quartus II software cannot import assignments from an existing .pin.

Figure 7–2 shows how to generate .pin and .fx files.

**Figure 7–2. Generating .pin and .fx Files**

Note to Figure 7–2:

1. For more information about the full design flow, which includes the I/O Designer software, the DxDesigner software, and the board layout tool flowchart details, refer to Figure 7–1.
For more information about pin and signal assignment transfer and the files that the Quartus II software can import and export, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

**Generating an .fx File**

You can generate an .fx in the Quartus II software for symbol generation in the Mentor Graphics I/O Designer software.

For more information about generating an .fx, refer to Generating FPGA Xchange-Format Files for Use with Other EDA Tools in Quartus II Help.

**Creating a Backup .qsf**

To create a backup .qsf of your current pin assignments, follow these steps:

1. On the Assignments menu, click Import Assignments. The Import Assignments dialog box appears.
2. In the Import Assignments dialog box, browse to your project and turn on Copy existing assignments into <project name>.qsf.bak.
3. Click OK.

For more information about pin and signal assignment transfer, and files the Quartus II software can import and export, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

**FPGA-to-Board Integration with the I/O Designer Software**

The Mentor Graphics I/O Designer software allows you to integrate your FPGA and PCB designs. Pin and signal assignment changes can be made anywhere in the design flow with either the Quartus II Pin Planner or the I/O Designer software. The I/O Designer software facilitates moving these changes, as well as synthesis, placement, and routing changes, between the Quartus II software, an external synthesis tool (if used), and a schematic capture tool such as the DxDesigner software.

This section describes how to use the I/O Designer software to transfer pin and signal assignment information to and from the Quartus II software with an .fx, and how to create symbols for the DxDesigner software.
Figure 7–3 shows the design flow using the I/O Designer software.

**Figure 7–3. Design Flow Using the I/O Designer Software**

For more information about the I/O Designer software, and to obtain usage, support, and product updates, use the Help menu in the I/O Designer software or refer to the Mentor Graphics website (www.mentor.com).

**I/O Designer Database Wizard**

An `.fpc` file stores all I/O Designer project information. You can create a new database incorporating information for the `.fx` and `.pin` files generated by the Quartus II software using the I/O Designer **Database Wizard**. You can also create a new, empty database and manually add the assignment information. If there is no signal or pin assignment information currently available, you can create an empty database containing only a selection of the target device. This action is useful if you know the signals in your design and the pins you want to assign. You can transfer this information at a later time to the Quartus II software for placement and routing.

Notes to Figure 7–3:

1. For more information about the full design flow including the Quartus II software flowchart details, refer to Figure 7–1 on page 7–2.
2. These are DxDesigner software-specific steps in the design flow and are not part of the I/O Designer flow.

For more information about the I/O Designer software, and to obtain usage, support, and product updates, use the Help menu in the I/O Designer software or refer to the Mentor Graphics website (www.mentor.com).
You can create an I/O Designer database with only a .pin or an .fx. However, if you are only using a .pin, you cannot import any I/O assignment changes made in the I/O Designer software back into the Quartus II software without first generating an .fx. If an .fx creates the I/O Designer database, the database may not contain all the available I/O assignment information. The .fx generated by the Quartus II software only lists pins with assigned signals. Because the .pin lists all device pins—whether signals are assigned to them or not—its use, along with the .fx, produces the most complete set of information for creating the I/O Designer database.

If you skip a step in the following process, you can complete the skipped step later. To return to a skipped step, on the Properties menu, click File. To create a new I/O Designer database using the Database wizard, follow these steps:

1. Start the I/O Designer software. The Welcome to I/O Designer dialog box appears. Select Wizard to create new database and click OK.

   If the Welcome to I/O Designer dialog box does not appear, you can access the wizard through the menu. To access the wizard, on the File menu, click Database Wizard.

2. Click Next. The Define HDL source file page appears.

   If no HDL files are available, or if the .fx contains your signal and pin assignments, you can skip Step 3 and proceed to Step 4.

   For more information about creating and using HDL files in the Quartus II software, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook, or refer to the I/O Designer Help.

3. If you have created a Verilog HDL or VHDL file in your Quartus II software design, you can add a top-level Verilog HDL or VHDL file in the I/O Designer software. Adding a file allows you to create functional blocks or get signal names from your design. You must create all physical pin assignments in I/O Designer if you are not using an .fx or a .pin. Click Next. The Database Name page appears.

4. In the Database Name page, type your database file name. Click Next. The Database Location window appears.

5. Add a path to the new or an existing database in the Location field, or browse to a database location. Click Next. The FPGA flow page appears.

6. In the Vendor menu, click Altera.

7. In the Tool/Library menu, click Quartus II 5.0, or a later version of the Quartus II software.

8. Select the appropriate device family, device, package, and speed (if applicable), from the corresponding menus. Click Next. The Place and route page appears.

   The Quartus II software version selections in the Tool/Library menu may not reflect the version of the Quartus II software currently installed in your system even if you are using the latest version of the I/O Designer software. The I/O Designer software uses the version number selection in this window to identify available or obsolete devices in that particular version of the Quartus II software. If you are unsure of the version to select,
use the latest version listed in the menu. If the device you are targeting does not appear in the device menu after making this selection, the device may be new and not yet added to the I/O Designer software. For I/O Designer software updates, contact Mentor Graphics or refer to their website (www.mentor.com).

9. In the **FPGAX file name** field, type or browse to the backup copy of the .fx generated by the Quartus II software.

10. In the **Pin report file name** field, type or browse to the .pin generated by the Quartus II software. Click **Next**.

You can also select a .qsf for update. The I/O Designer software can update the pin assignment information in the .qsf without affecting any other information in the file.

You can select a .pin without selecting an .fx for import. The I/O Designer software does not generate a .pin. To transfer assignment information to the Quartus II software, select an additional file and file type. Altera recommends selecting an .fx in addition to a .pin for transferring all the assignment information in the .fx and .pin files.

In some versions of the I/O Designer software, the standard file picker may incorrectly look for a .pin instead of an .fx. In this case, select All Files (*.*) from the **Save as type** list and select the file from the list.

11. The **Synthesis** page appears. On the **Synthesis** page, you can specify an external synthesis tool and a synthesis constraints file for use with the tool. If you do not use an external synthesis tool, click **Next**.

For more information about third-party synthesis tools, refer to *Volume 3: Verification* of the Quartus II Handbook.

12. On the **PCB Flow** page, you can select an existing schematic project or create a new project as a symbol information destination.

- To select an existing project, select **Choose existing project** and click **Browse** after the **Project Path** field. The **Select project** dialog box appears. Select the project.

- To create a new project, in the **Select project** dialog box, select **Create new empty project**. Type the project file name in the **Name** field and browse to the location where you want to save the file. Click **OK**.

If you have not specified a design tool to which you can send symbol information in the I/O Designer software, click **Advanced** in the **PCB Flow** page and select your design tool. If you select the DxDesigner software, you have the option to specify a Hierarchical Occurrence Attributes (.oat) file to import into the I/O Designer software. Click **Next** and then click **Finish** to create the database.

In I/O Designer version 2005 or later, the **Update Wizard** dialog box (refer to Figure 7–7 on page 7–13) appears if you are creating the database with the **Database** wizard. Use the **Update Wizard** dialog box to confirm creation of the I/O Designer database using the selected .fx and .pin files.
Use the I/O Designer software and your newly created database to make pin assignment changes, create pin swap groups, or adjust signal and pin properties in the I/O Designer GUI (Figure 7–4).

**Figure 7–4. Mentor Graphics I/O Designer Main Window**

For more information about using the I/O Designer software and the DxDesigner software, refer to the Mentor Graphics website (www.mentor.com) or refer to the I/O Designer software or the DxDesigner Help.
### Updating Pin Assignments from the Quartus II Software

As the design process continues, the FPGA designer must make changes to the logic design in the Quartus II software that places signals on different pins after recompiling the design, or manually with the Quartus II Pin Planner. These types of changes must be carried forward to the circuit schematic and board layout tools to ensure that signals connect to the correct pins on the FPGA. Updating the .fx and the .pin files in the Quartus II software facilitates this flow (Figure 7–5).

**Figure 7–5. Updating the I/O Designer Pin Assignments in the Design Flow**

To update the .fx in your selected output directory and the .pin in your project directory after making changes to the design, perform one of the following tasks:

- compile, or
- start EDA Netlist Writer.

You must rerun the I/O Assignment Analyzer whenever you make I/O changes in the Quartus II software. To rerun the I/O Assignment Analyzer, perform one of the following tasks:

- on the Processing menu, click **Start Compilation**, or
- on the Processing menu, click **Start I/O Assignment Analysis**.

For more information about setting up the .fx and running the I/O Assignment Analyzer, refer to the *I/O Management* chapter in volume 2 of the *Quartus II Handbook*.

If your I/O Designer database points to the .fx generated by the Quartus II software instead of a backup copy of the file, updating the file in the Quartus II software overwrites any changes made to the file by the I/O Designer software. If there are I/O Designer assignments in the .fx that you want to preserve, create a backup copy of the file before updating it in the Quartus II software, and verify that your I/O Designer database points to the backup copy. To point to the backup copy, perform the steps in the following section.

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**Note to Figure 7–5:**
(1) For more information about the full design flow, which includes the Quartus II software, the DxDesigner software, and the board layout tool flowchart details, refer to Figure 7–1 on page 7–2.
Whenever you update the .fx or the .pin in the Quartus II software, the I/O Designer database imports the changes. You must set up the locations for the files in the I/O Designer software.

1. To set up the file locations, on the File menu, click **Properties**. The project **Properties** dialog box appears (Figure 7–6).

   **Figure 7–6. Project Properties Dialog Box**

   ![Figure 7–6. Project Properties Dialog Box](image)

2. Under **FPGA Xchange**, click **Browse** to select the .fx file name and location.

3. To specify a Pin report file, under **Place and Route**, click **Browse** to select the .pin file name and location.

After you have set up these file locations, the I/O Designer software monitors these files for changes. If the .fx or .pin changes during the design flow, three indicators flash red in the lower right corner of the I/O Designer GUI (refer to **Figure 7–4 on page 7–10**). You can continue working or click on the indicators to open the I/O Designer **Update Wizard** dialog box. If you have made changes to your design in the Quartus II software that result in an updated .fx or .pin and the update indicators do not flash or you have previously canceled an indicated update, manually open the **Update Wizard** dialog box. To open the **Update Wizard** dialog box, on the File menu, click **Update**.
The I/O Designer **Update Wizard** dialog box lists the updated files associated with the database (Figure 7–7).

**Figure 7–7. Update Wizard Dialog Box**

The paths to the updated files have yellow exclamation points and the **Status** column shows **Not updated**, indicating that the database has not yet been updated with the newer information contained in the files. A checkmark to the left of any updated file indicates that the file updates the database. Turn on any files you want to use to update the I/O Designer database, and click **Next**. If you are not satisfied with the database update, on the Edit menu, click **Undo**.

You can update the I/O Designer database using the **.fx** and the **.pin** files simultaneously. Turning on the **.fx** and the **.pin** files for update causes the **Update Wizard** dialog box to provide options for using assignments from one file or the other exclusively or merging the assignments contained in both files into the I/O Designer database. Versions of the I/O Designer software older than version 2005 merge assignments contained in multiple files.
**Sending Pin Assignment Changes to the Quartus II Software**

In the same way that the FPGA designer can make adjustments that affect the PCB design, the board designer can make changes to optimize signal routing and layout that must be applied to the FPGA. The FPGA designer can take these required changes back into the Quartus II software to refit the logic to match the adjustments to the pin-out. The I/O Designer software accommodates this reverse flow as shown in Figure 7–8.

**Figure 7–8. Updating the Quartus II Pin Assignments in the Reverse Design Flow**

![Flowchart showing the reverse design flow from FPGA to I/O Designer](image)

**Notes to Figure 7–8:**

1. These are software-specific steps in the design flow and are not necessary for the reverse flow steps of the design.
2. For more information about the full design flow, which includes the complete I/O Designer software, the DxDesigner software, and the board layout tool flowchart details, refer to Figure 7–1 on page 7–2.

You can make pin assignment changes directly in the I/O Designer software, or the software can automatically update changes made in a board layout tool that are back-annotated to a schematic entry program such as the DxDesigner software. You must update the .fx to reflect these updates in the Quartus II software. To perform this update in the I/O Designer software, on the Generate menu, click **FPGA Xchange File**.

If your I/O Designer database points to the .fx generated by the Quartus II software instead of a backup copy, updating the file from the I/O Designer software overwrites any changes made to the file by the Quartus II software. If there are assignments from the Quartus II software in the file that you want to preserve, create a backup copy of the file before updating it in the I/O Designer software, and verify that your I/O Designer database points to the backup copy. To point to the backup copy, perform the steps in “Updating Pin Assignments from the Quartus II Software” on page 7–11.
You must import the updated .fx into the Quartus II software. To import the file, follow these steps:

1. Start the Quartus II software and open your project.
2. On the Assignments menu, click **Import Assignments**.
3. In the File name box, click **Browse** and from the **Files of type** list, select **FPGA Xchange Files (*.fx)**.
4. Select the .fx and click **Open**.
5. Click **OK**.

### Protecting Assignments in the Quartus II Software

To protect assignments in the Quartus II software, follow these steps:

1. Start the Quartus II software.
2. On the Assignments menu, click **Import Assignments**. The **Import Assignments** dialog box appears.
3. Turn on **Copy existing assignments into <project name>.qsf.bak before importing** before importing the .fx. This action automatically creates a backup copy of the Quartus II constraints file that contains all your current pin assignments.

### Generating Symbols for the DxDesigner Software

Along with circuit simulation, circuit board schematic creation is one of the first tasks required in the design of a new PCB. Schematics must understand how the PCB works, and to generate a netlist for a board layout tool for board design and routing. The I/O Designer software allows you to create schematic symbols based on the FPGA design exported from the Quartus II software.

Most FPGA devices contain hundreds of pins, requiring large schematic symbols that may not fit on a single schematic page. Symbol designs in the I/O Designer software can be split or fractured into various functional blocks, allowing multiple part fractures on the same schematic page or across multiple pages. In the DxDesigner software, these part fractures join together with the use of the \texttt{HETERO} attribute.

The I/O Designer software can generate symbols for use in various Mentor Graphics schematic entry tools, and can import changes back-annotated by board layout tools to update the database and feed updates back to the Quartus II software with the .fx. This section discusses symbol creation specifically for the DxDesigner software.

You can create schematic symbols with the I/O Designer software in the following ways:

- Manually
- Using the I/O Designer **Symbol** wizard
- Importing previously created symbols from the DxDesigner software

The I/O Designer **Symbol** wizard can be used as a design base that allows you to quickly create a symbol for manual editing at a later time. If you have created symbols in a DxDesigner project and want to apply a different FPGA design to them, you can manually import these symbols from the DxDesigner project. To import the symbols, start the I/O Designer software, and on the File menu, click **Import Symbol**.
For more information about importing symbols from the DxDesigner software into an I/O Designer database, refer to the I/O Designer Help.

Symbols created in the I/O Designer software are either functional, physical (PCB), or both. Signals imported into the database, usually from Verilog HDL or VHDL files, are the basis of a functional symbol. No physical device pins must be associated with the signals to generate a functional symbol. This section focuses on board-level PCB symbols with signals directly mapped to physical device pins through assignments in either the Quartus II Pin Planner or in the I/O Designer database.

For more information about manually creating, importing, and editing symbols in the I/O Designer software, as well as the different types of symbols the software can generate, refer to the I/O Designer Help.

**Setting Up the I/O Designer Software to Work with the DxDesigner Software**

To verify if you are set up to export symbols to a DxDesigner project, or to manually set up the I/O Designer software to work with the DxDesigner software, you must set the path to the DxDesigner executable, set the export type to DxDesigner, and set the path to a DxDesigner project directory.

To set these options, follow these steps:

1. Start the I/O Designer software.
2. On the Tools menu, click **Preferences**. The **Preferences** dialog box appears.
3. Click **Paths**, double-click on the **DxDesigner executable file** path field, and click **Browse** to select the location of the DxDesigner application (Figure 7–9).
4. Click **Apply**.

**Figure 7–9. Path Preferences Dialog Box**

5. Click **Symbol Editor** and click **Export**. In the Export type menu, under **General**, select **DxDesigner/PADS-Designer** (Figure 7–10).
6. Click **Apply** and click **OK**.

**Figure 7-10. Symbol Editor Export Preferences**

7. On the File menu, click **Properties**. The **Properties** dialog box appears.

8. Click the **PCB Flow** tab and click **Path to a DxDesigner project directory**.

9. Click **OK**.

If you do not have a new DxDesigner project in the **Database** wizard and a DxDesigner project, you must create a new database with the DxDesigner software, and point the I/O Designer software to this new project.

For more information about creating and working with DxDesigner projects, refer to the DxDesigner Help.

**Creating Symbols with the Symbol Wizard**

You can create, fracture, and edit FPGA symbols based on Altera devices with the I/O Designer **Symbol** wizard. To create a symbol based on a selected Altera FPGA device, follow these steps:

1. Start the I/O Designer software.
2. Click **Symbol Wizard** in the toolbar, or on the Symbol menu, click **Symbol Wizard**. The Symbol Wizard (1 of 6) page appears (Figure 7–11).

**Figure 7–11. Symbol Wizard**

3. On page 1 of the Symbol Wizard page, in the **Symbol name** field, type the symbol name. The DEVICE and PKG_TYPE fields are automatically populated with the device and package information. Under **Symbol type**, click **PCB**. Under **Use signals**, click **All**.

4. Click **Next**. The Symbol Wizard (2 of 6) page appears.

   If the DEVICE and PKG_TYPE fields are blank or incorrect, cancel the Symbol wizard and select the correct device information. On the File menu, click **Properties**. In the Properties window, click the FPGA Flow tab and enter the correct device information.

5. On page 2 of the Symbol Wizard page, select fracturing options for your symbol. If you are using the Symbol wizard to edit a previously created fractured symbol, you must turn on **Reuse existing fractures** to preserve your current fractures. Select other options on this page as appropriate for your symbol.

6. Click **Next**. The Symbol Wizard (3 of 6) page appears.

7. Additional fracturing options are available on page 3 of the Symbol Wizard page. After selecting the necessary options, click **Next**. The Symbol Wizard (4 of 6) page appears.

8. On page 4 of the Symbol Wizard page, select the options for the appearance of the symbols. Select the necessary options and click **Next**. The Symbol Wizard (5 of 6) page appears.
9. On page 5 of the **Symbol Wizard** page, define what information you want to label for the entire symbol and for individual pins. Select the necessary options and click **Next**. The **Symbol Wizard (6 of 6)** page appears.

10. On the final page of the **Symbol Wizard** page, add additional signals and pins that have not been placed in the symbol. Click **Finish** when you complete your selections.

You can view your symbol and any fractures you created with the Symbol Editor (Figure 7–12). You can edit parts of the symbol, delete fractures, or rerun the **Symbol wizard**.

**Figure 7–12. The I/O Designer Symbol Editor**

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If assignments in the I/O Designer database are updated, the symbols created in the I/O Designer software automatically reflect these changes. Assignment changes can be made in the I/O Designer software, with an updated .fx from the Quartus II software, or from a back-annotated change in your board layout tool.

**Exporting Symbols to the DxDesigner Software**

After you have completed your symbols, export the symbols to your DxDesigner project. To generate all the fractures of a symbol, on the Generate menu, click **All Symbols**. To generate a symbol for the currently displayed symbol in Symbol Editor, click **Current Symbol Only**. The `/sym directory` in your DxDesigner project saves each symbol in the database as a separate file. The symbols can be instantiated in your DxDesigner schematics.

For more information about working with DxDesigner projects, refer to the DxDesigner Help.

**Scripting Support**

The I/O Designer software features a command line Tcl interpreter. All commands issued through the GUI in the I/O Designer software translate into Tcl commands run by the tool. You can view the generated Tcl commands and run scripts, or type individual commands in the I/O Designer Console window.
This scripting support section includes commands that perform some of the operations described in this chapter.

If you want to change the .fx from which the I/O Designer software updates assignments, type the following command at an I/O Designer Tcl prompt:

```
set_fpga_xchange_file <file name>
```

You can type the following command to update the I/O Designer database with assignment updates made in the Quartus II software after specifying the .fx:

```
update_from_fpga_xchange_file
```

You can type the following command to update the .fx with changes made to the assignments in the I/O Designer software for transfer back into the Quartus II software:

```
generate_fpga_xchange_file
```

You can type the following command if you want to import assignment data from a .pin created by the Quartus II software:

```
set_pin_report_file -quartus_pin <file name>
```

You can run the I/O Designer Symbol wizard with the following command:

```
symbolwizard
```

You can set the DxDesigner project directory path where symbols are saved with the following command:

```
set_dx_designer_project -path <path>
```

For more information about Tcl scripting and Tcl scripting with the Quartus II software, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook. For more information about the Tcl scripting capabilities of the I/O Designer software as well as a list of available commands, refer to the I/O Designer Help.

**FPGA-to-Board Integration with the DxDesigner Software**

The Mentor Graphics DxDesigner software is a design entry tool for schematic capture. You can use it to create flat circuit schematics for all the PCB design types. You can also use the DxDesigner software to create hierarchical schematics that facilitate design reuse and a team-based design. You can use the DxDesigner software in the design flow alone or in conjunction with the I/O Designer software. However, if you use the DxDesigner software without the I/O Designer software, the design flow is one-way, using only the .pin generated by the Quartus II software.
You can only make signal and pin assignment changes in the Quartus II software and these changes reflect as updated symbols in a DxDesigner schematic. You cannot back-annotate changes made in a board layout tool or in a DxDesigner symbol to the Quartus II software. Figure 7–13 shows the design flow without the I/O Designer software.

**Figure 7–13. Design Flow Without the I/O Designer Software**

![Design Flow Without the I/O Designer Software](image)

**Note to Figure 7–13:**

(1) For more information about the full design flow, which includes the Quartus II software, the I/O Designer software, and the board layout tool flowchart details, refer to Figure 7–1 on page 7–2.

For more information about the DxDesigner software, including usage, support, training, and product updates, refer to the Mentor Graphics website (www.mentor.com), or choose Schematic Design Help Topics in the DxDesigner Help.

**DxDesigner Project Settings**

New projects in the DxDesigner software are set up to create FPGA symbols by default. However, if you are using the I/O Designer software with the DxDesigner software, you must enable the DxBoardLink Flow options for complete support and compatibility with the I/O Designer software.

You can enable the DxBoardLink flow design configuration during or after creating a new DxDesigner project.

To enable the DxBoardLink flow design configuration when creating a new DxDesigner project, follow these steps:

1. Start the DxDesigner software.
2. On the File menu, click **New** and click the **Project** tab. The **New** dialog box appears (Figure 7–14).

**Figure 7–14. New Project Dialog Box**

![New Project Dialog Box](image)

3. Click **More**. Turn on **DxBoardLink** (Figure 7–14).

To enable the DxBoardLink Flow design configuration in an existing project, click **Design Configurations** in the Design Configuration toolbar and turn on **DxBoardLink** (Figure 7–15).

**Figure 7–15. DxBoardLink Design Configuration**

![DxBoardLink Design Configuration](image)
DxDesigner Symbol Wizard

You can create schematic symbols in the DxDesigner software manually or with the Symbol wizard. The DxDesigner Symbol wizard is similar to the I/O Designer Symbol wizard, but with fewer fracturing options.

FPGA symbols based on Altera devices can be created, fractured, and edited with the DxDesigner Symbol wizard. To start the Symbol wizard, follow these steps:

1. Start the DxDesigner software.
2. Click Symbol Wizard in the toolbar, or on the File menu, click New. The New window appears. Click the File tab and create a new file of type Symbol Wizard.
3. Type the new symbol name in the name field and click OK. The Symbol Wizard page appears (Figure 7–16).

Figure 7–16. Wizard Task Selection

4. On the Wizard Task Selection page, choose to create a new symbol or modify an existing symbol. If you are modifying an existing symbol, specify the library path or alias, and select the existing symbol. If you are creating a new symbol, select DxBoardLink for the symbol source. The DxDesigner block type defaults to Module because the FPGA design does not have an underlying DxDesigner schematic. Choose whether or not to fracture the symbol. After making your selections, click Next. The New Symbol and Library Name page appears.

5. On the New Symbol and Library Name page, type a name for the symbol, an overall part name for all the symbol fractures, and a library name for the new library created for this symbol. By default, the part and library names are the same as the symbol name. Click Next. The Symbol Parameters page appears.
6. On the **Symbol Parameters** page, specify the appearance of the generated symbol and how it matches up with the grid you have set in your DxDesigner project schematic. After making your selections, click **Next**. The **DxBoardLink Pin List Import** page appears (Figure 7–17).

**Figure 7–17. DxBoardLink Pin List Import**

7. On the **DxBoardLink Pin List Import** page, in the **FPGA vendor** list, select **Altera Quartus**. In the **Pin-Out file to import** field, browse to and select the `.pin` from your Quartus II design project directory. You can also select choices from the Fracturing Scheme, Bus pin, and Power pin options. After making your selections, click **Next**. The **Symbol Attributes** page appears.

8. On the **Symbol Attributes** page, select to create or modify symbol attributes for use in the DxDesigner software. After making your selections, click **Next**. The **Pin Settings** page appears.

9. On the **Pin Settings** page, make any final adjustments to pin and label location and information. Each tabbed spreadsheet represents a fracture of your symbol. After making your selections, click **Save Symbol**.

After creating the symbol, you can examine and place any fracture of the symbol in your schematic. You can locate separate files of all the fractures you created in the library you specified or created in the `/syn` directory in your DxDesigner project. You can add the symbols to your schematics or you can manually edit the symbols or with the **Symbol** wizard.
Symbols created in the DxDesigner software can be edited and updated with newer versions of the .pin generated by the Quartus II software. However, you cannot fracture a symbol again because symbol fracturing is permanent. To create new fractures for your design, create a new symbol in the Symbol wizard, and perform the steps in “DxDesigner Symbol Wizard” on page 7–23.

Conclusion

Transferring a complex, high-pin-count FPGA design to a PCB for prototyping or manufacturing is a daunting process that can lead to errors in the PCB netlist or design, especially when multiple engineers are working on different parts of the project. The design workflow available when using the Quartus II software with the Mentor Graphics toolset assists the FPGA designer and the board designer in preventing errors and focusing their attention on the design.

Document Revision History

Table 7–1 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Removed Reference Document section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ General style editing.</td>
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<tr>
<td></td>
<td></td>
<td>■ Added a link to Help in “Performing Simultaneous Switching Noise (SSN) Analysis of Your FPGA.”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed Figure 8–4 on page 8–9 and Figure 8–5 on page 8–11.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Generating an .fx File.”</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Added minor information about simultaneous switching noise (SSN) analysis on “Performing Simultaneous Switching Noise (SSN) Analysis of Your FPGA.”</td>
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<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Was chapter 6 in the 8.1.0 release.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed Figures that were numbered 6-4, 6-6, 6-7, and 6-8 in v8.1.0.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8½” × 11” page size. No change to content.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Updated references.</td>
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</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
8. Cadence PCB Design Tools Support

This chapter addresses how the Quartus® II software interacts with the Cadence Allegro Design Entry HDL software and the Cadence Allegro Design Entry CIS (Component Information System) software (also known as OrCAD Capture CIS) to provide a complete FPGA-to-board integration design workflow.

With today’s large, high-pin-count and high-speed FPGA devices, good PCB design practices are important to ensure the correct operation of your system. The PCB design takes place concurrently with the design and programming of the FPGA. An FPGA or ASIC designer initially creates the signal and pin assignments and the board designer must transfer these assignments to the symbols used in their system circuit schematics and board layout correctly. As the board design progresses, you must perform pin reassignments to optimize the layout. You must communicate pin reassignments to the FPGA designer to ensure the new assignments are processed through the FPGA with updated placement and routing.

This chapter is intended for board design and layout engineers who want to begin the FPGA board integration process while the FPGA is still in the design phase. Part librarians can also benefit from this chapter by learning the method to use output from the Quartus II software to create new library parts and symbols.

This chapter discusses the following topics:

- Cadence tool description, history, and comparison.
- The general design flow between the Quartus II software and the Cadence Allegro Design Entry HDL software and the Cadence Allegro Design Entry CIS software.
- Generating schematic symbols from your FPGA design for use in the Cadence Allegro Design Entry HDL software.
- Updating Design Entry HDL symbols when making signal and pin assignment changes in the Quartus II software.
- Creating schematic symbols in the Cadence Allegro Design Entry CIS software from your FPGA design.
- Updating symbols in the Cadence Allegro Design Entry CIS software when making signal and pin assignment changes in the Quartus II software.
- Using Altera-provided device libraries in the Cadence Allegro Design Entry CIS software.

The procedures in this chapter require the following software:

- The Quartus II software version 5.1 or later
- The Cadence Allegro Design Entry HDL software or the Cadence Allegro Design Entry CIS software version 15.2 or later
The OrCAD Capture software with the optional CIS option version 10.3 or later (optional)

These programs are very similar because the Cadence Allegro Design Entry CIS software is based on the OrCAD Capture software. This chapter refers to the Cadence Allegro Design Entry CIS software; however, any procedural information can also apply to the OrCAD Capture software unless otherwise noted.

For more information about obtaining and licensing the Cadence tools and for product information, support, and training, refer to the Cadence website (www.cadence.com). For more information about the OrCAD Capture software and the CIS option, refer to the Cadence website (www.cadence.com). For more information about Cadence and OrCAD support and training, refer to the EMA Design Automation website (www.ema-eda.com).

Product Comparison

The Cadence and OrCAD design tools are different in their function and location of product information. Table 8–1 lists the Cadence and OrCAD products described in this chapter and provides information about changes, product information, and support.

Table 8–1. Cadence and OrCAD Product Comparison

<table>
<thead>
<tr>
<th>Description</th>
<th>Cadence Allegro Design Entry HDL</th>
<th>Cadence Allegro Design Entry CIS</th>
<th>OrCAD Capture CIS</th>
</tr>
</thead>
<tbody>
<tr>
<td>Former Name</td>
<td>Concept HDL Expert</td>
<td>Capture CIS Studio</td>
<td>—</td>
</tr>
<tr>
<td>History</td>
<td>More commonly known by its former name, Cadence renamed all board design tools in 2004 under the Allegro name.</td>
<td>Based directly on OrCAD Capture CIS, the Cadence Allegro Design Entry CIS software is still developed by OrCAD but sold and marketed by Cadence. EMA provides support and training.</td>
<td>The basis for Design Entry CIS is still developed by OrCAD for continued use by existing OrCAD customers. EMA provides support and training for all OrCAD products.</td>
</tr>
<tr>
<td>Vendor Design Flow</td>
<td>Cadence Allegro 600 series, formerly known as the Expert Series, for high-end, high-speed design.</td>
<td>Cadence Allegro 200 series, formerly known as the Studio Series, for small- to medium-level design.</td>
<td>—</td>
</tr>
</tbody>
</table>
FPGA-to-PCB Design Flow

You can create a design flow integrating an Altera FPGA design from the Quartus II software through a circuit schematic in the Cadence Allegro Design Entry HDL software or the Cadence Allegro Design Entry CIS software. Figure 8–1 shows the design flow with the Cadence Allegro Design Entry HDL software. Figure 8–2 shows the design flow with the Cadence Allegro Design Entry CIS software.

Figure 8–1. Design Flow with the Cadence Allegro Design Entry HDL Software
Figure 8–2. Design Flow with the Cadence Allegro Design Entry CIS Software

Figure 8–1 and Figure 8–2 show the possible design flows, depending on your tool choice. To create FPGA symbols using the Cadence Allegro PCB Librarian Part Developer tool, you must obtain the Cadence PCB Librarian Expert license. You can update symbols with changes made to the FPGA design using any of these tools.

To integrate an Altera FPGA design starting in the Quartus II software through to a circuit schematic in the Cadence Allegro Design Entry HDL software or the Cadence Allegro Design Entry CIS software, follow these steps:

1. In the Quartus II software, compile your design to generate a Pin-Out File (.pin) to transfer the assignments to the Cadence software.

2. If you are using the Cadence Allegro Design Entry HDL software for your schematic design, follow these steps:
   a. Open an existing project or create a new project in the Cadence Allegro Project Manager tool.
   b. Construct a new symbol or update an existing symbol using the Cadence Allegro PCB Librarian Part Developer tool.
   c. With the Cadence Allegro PCB Librarian Part Developer tool, edit your symbol or fracture it into smaller parts (optional).
   d. Instantiate the symbol in your Cadence Allegro Design Entry HDL software schematic and transfer the design to your board layout tool.

   or

If you are using the Cadence Allegro Design Entry CIS software for your schematic design, follow these steps:

a. Generate a new part in a new or existing Cadence Allegro Design Entry CIS project, referencing the .pin output file from the Quartus II software. You can also update an existing symbol with a new .pin.
b. Split the symbol into smaller parts as necessary.

c. Instantiate the symbol in your Cadence Allegro Design Entry CIS schematic and transfer the design to your board layout tool.

**Performing Simultaneous Switching Noise (SSN) Analysis of Your FPGA**

With the Quartus II software, you can extract pin assignment data and perform SSN analysis of your FPGA design for designs targeting the Stratix III device family. You can analyze SSN in your device early in the board layout stage as part of your overall pin planning process; however, you do not have to perform SSN analysis to generate pin assignment data from the Quartus II software. You can use the SSN Analyzer tool to optimize the pin assignments for better SSN performance of your device.

For more information about the SSN Analyzer, refer to *About the SSN Analyzer* in Quartus II Help and the *Simultaneous Switching Noise (SSN) Analysis and Optimizations* chapter in volume 2 of the *Quartus II Handbook*.

**Setting Up the Quartus II Software**

You can transfer pin and signal assignments from the Quartus II software to the Cadence design tools by generating the Quartus II project `.pin`. The `.pin` is an output file generated by the Quartus II Fitter containing pin assignment information. You can use the Quartus II Pin Planner to set and change the assignments in the `.pin` and then transfer the assignments to the Cadence design tools. You cannot, however, import pin assignment changes from the Cadence design tools into the Quartus II software with the `.pin`.

The `.pin` lists all used and unused pins on your selected Altera device. The `.pin` also provides the following basic information fields for each assigned pin on the device:

- Pin signal name and usage
- Pin number
- Signal direction
- I/O standard
- Voltage
- I/O bank
- User or Fitter-assigned

For more information about using the Quartus II Pin Planner to create or change pin assignment details, refer to the *I/O Management* chapter in volume 2 of the *Quartus II Handbook*.

**Generating a .pin File**

The Quartus II Fitter generates a `.pin` during a full compilation of your FPGA design, or when performing I/O assignment analysis on your design. You can locate the `.pin` in your Quartus II project directory with the name `<project name>`.pin.
For more information about pin and signal assignment transfer and the files that the Quartus II software can import and export, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

FPGA-to-Board Integration with the Cadence Allegro Design Entry HDL Software

The Cadence Allegro Design Entry HDL software is a schematic capture tool and is part of the Cadence 600 series design flow. Use the Cadence Allegro Design Entry HDL software to create flat circuit schematics for all types of PCB design. The Cadence Allegro Design Entry HDL software can also create hierarchical schematics to facilitate design reuse and team-based design. With the Cadence Allegro Design Entry HDL software, the design flow from FPGA-to-board is one-way, using only the .pin generated by the Quartus II software. You can only make signal and pin assignment changes in the Quartus II software and these changes reflect as updated symbols in a Cadence Allegro Design Entry HDL project. For more information about the design flow with the Cadence Allegro Design Entry HDL software, refer to Figure 8–1 on page 8–3.

Routing or pin assignment changes made in a board layout tool or a Cadence Allegro Design Entry HDL software symbol cannot be back-annotated to the Quartus II software.

For more information about the Cadence Allegro Design Entry HDL software and the Cadence Allegro PCB Librarian Part Developer tool, including licensing, support, usage, training, and product updates, refer to the Help in the software or to the Cadence website (www.cadence.com).

Creating Symbols

In addition to circuit simulation, circuit board schematic creation is one of the first tasks required when designing a new PCB. Schematics must understand how the PCB works, and to generate a netlist for a board layout tool for board design and routing. The Cadence Allegro PCB Librarian Part Developer tool allows you to create schematic symbols based on FPGA designs exported from the Quartus II software.

You can create symbols for the Cadence Allegro Design Entry HDL project with the Cadence Allegro PCB Librarian Part Developer tool, which is available in the Cadence Allegro Project Manager tool. Altera recommends using the Cadence Allegro PCB Librarian Part Developer tool to import FPGA designs into the Cadence Allegro Design Entry HDL software.

You must obtain a PCB Librarian Expert license from Cadence to run the Cadence Allegro PCB Librarian Part Developer tool. The Cadence Allegro PCB Librarian Part Developer tool provides a GUI with many options for creating, editing, fracturing, and updating symbols. If you do not use the Cadence Allegro PCB Librarian Part Developer tool, you must create and edit symbols manually in the Symbol Schematic View in the Cadence Allegro Design Entry HDL software.
If you do not have a PCB Librarian Expert license, you can automatically create FPGA symbols using the programmable IC (PIC) design flow found in the Cadence Allegro Project Manager tool. For more information about using the PIC design flow, refer to the Help in the Cadence design tools, or go to the Cadence website (www.cadence.com).

Before creating a symbol from an FPGA design, you must open a Cadence Allegro Design Entry HDL project with the Cadence Allegro Project Manager tool. If you do not have an existing Cadence Allegro Design Entry HDL project, you can create one with the Cadence Allegro Design Entry HDL software. The Cadence Allegro Design Entry HDL project directory with the name `<project name>.cpm` contains your Cadence Allegro Design Entry HDL projects.

While the Cadence Allegro PCB Librarian Part Developer tool refers to symbol fractures as slots, the other tools described in this chapter use different names to refer to symbol fractures. Table 8–2 lists the symbol fracture naming conventions for each of the tools addressed in this chapter.

### Table 8–2. Symbol Fracture Naming

<table>
<thead>
<tr>
<th></th>
<th>Cadence Allegro PCB Librarian Part Developer Tool</th>
<th>Cadence Allegro Design Entry HDL Software</th>
<th>Cadence Allegro Design Entry CIS Software</th>
</tr>
</thead>
<tbody>
<tr>
<td>During symbol generation</td>
<td>Slots</td>
<td>—</td>
<td>Sections</td>
</tr>
<tr>
<td>During symbol schematic instantiation</td>
<td>—</td>
<td>Versions</td>
<td>Parts</td>
</tr>
</tbody>
</table>
Cadence Allegro PCB Librarian Part Developer Tool

You can create, fracture, and edit schematic symbols for your designs using the Cadence Allegro PCB Librarian Part Developer tool. Symbols designed in the Cadence Allegro PCB Librarian Part Developer tool can be split or fractured into several functional blocks called slots, allowing multiple smaller part fractures to exist on the same schematic page or across multiple pages. Figure 8–3 shows how the Cadence Allegro PCB Librarian Part Developer tool fits into the design flow.

Figure 8–3. Cadence Allegro PCB Librarian Part Developer Tool in the Design Flow

Notes to Figure 8–3:
(1) For more information about the full design flow flowchart, refer to Figure 8–1 on page 8–3.
(2) Grayed out steps are not part of the FPGA symbol creation or update process.

To run the Cadence Allegro PCB Librarian Part Developer tool, you must open a Cadence Allegro Design Entry HDL project in the Cadence Allegro Project Manager tool. To open the Cadence Allegro PCB Librarian Part Developer tool, on the Flows menu, click Library Management, and then click Part Developer.

Import and Export Wizard

After starting the Cadence Allegro PCB Librarian Part Developer tool, use the Import and Export wizard to import your pin assignments from the Quartus II software.

Altera recommends using your PCB Librarian Expert license file. To point to your PCB Librarian Expert license file, on the File menu, click Change Product and then select the correct product license.

To access the Import and Export wizard, follow these steps:
1. On the File menu, click Import and Export.
2. Select Import ECO-FPGA, and then click Next.
3. In the **Select Source** page of the **Import and Export** wizard, specify the following settings:
   a. In the **Vendor** list, select **Altera**.
   b. In the **PnR Tool** list, select **quartusII**.
   c. In the **PR File** box, browse to select the `.pin` in your Quartus II project directory.
   d. Click **Simulation Options** to select simulation input files.
   e. Click **Next**.

4. In the **Select Destination** dialog box, specify the following settings:
   a. Under **Select Component**, click **Generate Custom Component** to create a new component in a library,
      or
      Click **Use standard component** to base your symbol on an existing component.

   ![Altera recommends creating a new component if you previously created a generic component for an FPGA device. Generic components can cause some problems with your design. When you create a new component, you can place your pin and signal assignments from the Quartus II software on this component and reuse the component as a base when you have a new FPGA design.](image)

   b. In the **Library** list, select an existing library. You can select from the cells in the selected library. Each cell represents all the symbol versions and part fractures for a particular part. In the **Cell** list, select the existing cell to use as a base for your part.
   c. In the **Destination Library** list, select a destination library for the component. Click **Next**.
   d. Review and edit the assignments you import into the Cadence Allegro PCB Librarian Part Developer tool based on the data in the `.pin` and then click **Finish**. The location of each pin is not included in the **Preview of Import Data** page of the **Import and Export** wizard, but input pins are on the left side of the created symbol, output pins on the right, power pins on the top, and ground pins on the bottom.

**Editing and Fracturing Symbol**

After creating your new symbol in the Cadence Allegro PCB Librarian Part Developer tool, you can edit the symbol graphics, fracture the symbol into multiple slots, and add or change package or symbol properties.

The Part Developer Symbol Editor contains many graphical tools to edit the graphics of a particular symbol. To edit the symbol graphics, select the symbol in the cell hierarchy. The **Symbol Pins** tab appears. You can edit the preview graphic of the symbol in the **Symbol Pins** tab.
Fracturing a Cadence Allegro PCB Librarian Part Developer package into separate symbol slots is useful for FPGA designs. A single symbol for most FPGA packages might be too large for a single schematic page. Splitting the part into separate slots allows you to organize parts of the symbol by function, creating cleaner circuit schematics. For example, you can create one slot for an I/O symbol, a second slot for a JTAG symbol, and a third slot for a power/ground symbol.

Figure 8–4 shows a part fractured into separate slots.

**Figure 8–4. Splitting a Symbol into Multiple Slots** *(Notes 1), (2)*

*Notes to Figure 8–4:*

1. Figure 8–4 represents a Cyclone device with JTAG or passive serial (PS) mode configuration option settings. Symbols created for other devices or other configuration modes may have different sets of configuration pins, but can be fractured in a similar manner.
2. The power/ground slot shows only a representation of power and ground pins because the device contains a large number of power and ground pins.

To fracture a part into separate slots, or to modify the slot locations of pins on parts fractured in the Cadence Allegro PCB Librarian Part Developer tool, follow these steps:

1. Start the Cadence Allegro Design Project Manager.
3. Click Part Developer.
4. Click the name of the package you want to change in the cell hierarchy.
5. Click Functions/Slots. If you are not creating new slots but want to change the slot location of some pins, proceed to Step 6. If you are creating new slots, click Add. A dialog box appears, allowing you to add extra symbol slots. Set the number of extra slots you want to add to the existing symbol, not the total number of desired slots for the part. Click OK.
6. Click Distribute Pins. Specify the slot location for each pin. Use the checkboxes in each column to move pins from one slot to another. Click OK.
7. After distributing the pins, click the Package Pin tab and click Generate Symbol(s).
8. Select whether to create a new symbol or modify an existing symbol in each slot. Click OK.

The newly generated or modified slot symbols appear as separate symbols in the cell hierarchy. Each of these symbols can be edited individually.

The Cadence Allegro PCB Librarian Part Developer tool allows you to remap pin assignments in the Package Pin tab of the main Cadence Allegro PCB Librarian Part Developer window. If signals remap to different pins in the Cadence Allegro PCB Librarian Part Developer tool, the changes reflect only in regenerated symbols for use in your schematics. You cannot transfer pin assignment changes to the Quartus II software from the Cadence Allegro PCB Librarian Part Developer tool, which creates a potential mismatch of the schematic symbols and assignments in the FPGA design. If pin assignment changes are necessary, make the changes in the Quartus II Pin Planner instead of the Cadence Allegro PCB Librarian Part Developer tool, and update the symbol as described in the following sections.

For more information about creating, editing, and organizing component symbols with the Cadence Allegro PCB Librarian Part Developer tool, refer to the Part Developer Help.

**Updating FPGA Symbols**

As the design process continues, you must make logic changes in the Quartus II software, placing signals on different pins after recompiling the design, or use the Quartus II Pin Planner to make changes manually. The board designer can request such changes to improve the board routing and layout. To ensure signals connect to the correct pins on the FPGA, you must carry forward these types of changes to the circuit schematic and board layout tools. Updating the .pin in the Quartus II software facilitates this flow. Figure 8–5 shows this part of the design flow.

**Figure 8–5. Updating the FPGA Symbol in the Design Flow**

![Diagram](image-url)

**Notes to Figure 8–5:**

1. For more information about the full design flow flowchart, refer to Figure 8–1 on page 8–3.
2. Grayed out steps are not part of the FPGA symbol update process.
To update the symbol using the Cadence Allegro PCB Librarian Part Developer tool after updating the .pin, follow these steps:

1. On the File menu, click **Import and Export**. The Import and Export wizard appears.

2. In the list of actions to perform, select **Import ECO - FPGA**. Click Next. The **Select Source** dialog box appears.

3. Select the updated source of the FPGA assignment information. In the **Vendor** list, select Altera. In the **PnR Tool** list, select quartusII. In the **PR File** field, click browse to specify the updated .pin in your Quartus II project directory. Click Next. The Select Destination window appears.

4. Select the source component and a destination cell for the updated symbol. To create a new component based on the updated pin assignment data, select **Generate Custom Component**. Selecting **Generate Custom Component** replaces the cell listed under the **Specify Library and Cell** name header with a new, nonfractured cell. You can preserve these edits by selecting **Use standard component and select the existing library and cell**. Select the destination library for the component and click Next. The **Preview of Import Data** dialog box appears.

5. Make any additional changes to your symbol. Click Next. A list of ECO messages appears summarizing the changes made to the cell. To accept the changes and update the cell, click Finish.

6. The main Cadence Allegro PCB Librarian Part Developer window appears. You can edit, fracture, and generate the updated symbols as usual from the main Cadence Allegro PCB Librarian Part Developer window.

If the Cadence Allegro PCB Librarian Part Developer tool is not set up to point to your PCB Librarian Expert license file, an error message appears in red at the bottom of the message text window of the Part Developer when you select the **Import and Export** command. To point to your PCB Librarian Expert license, on the File menu, click **Change Product**, and select the correct product license.

**Instantiating the Symbol in the Cadence Allegro Design Entry HDL Software**

To instantiate the symbol in your Cadence Allegro Design Entry HDL schematic after saving the new symbol in the Cadence Allegro PCB Librarian Part Developer tool, follow these steps:

1. In the Cadence Allegro Project Manager tool, switch to the board design flow.

2. On the Flows menu, click **Board Design**.

3. To start the Cadence Allegro Design Entry HDL software, click **Design Entry**.

4. To add the newly created symbol to your schematic, on the Component menu, click Add. The **Add Component** dialog box appears.

5. Select the new symbol library location, and select the name of the cell you created from the list of cells.

The symbol attaches to your cursor for placement in the schematic. To fracture the symbol into slots, right-click the symbol and choose **Version** to select one of the slots for placement in the schematic.
For more information about the Cadence Allegro Design Entry HDL software, including licensing, support, usage, training, and product updates, refer to the Help in the software or go to the Cadence website (www.cadence.com).

FPGA-to-Board Integration with Cadence Allegro Design Entry CIS Software

The Cadence Allegro Design Entry CIS software is a schematic capture tool (part of the Cadence 200 series design flow based on OrCAD Capture CIS). Use the Cadence Allegro Design Entry CIS software to create flat circuit schematics for all types of PCB design. You can also create hierarchical schematics to facilitate design reuse and team-based design using the Cadence Allegro Design Entry CIS software. With the Cadence Allegro Design Entry CIS software, the design flow from FPGA-to-board is unidirectional using only the .pin generated by the Quartus II software. You can only make signal and pin assignment changes in the Quartus II software. These changes reflect as updated symbols in a Cadence Allegro Design Entry CIS schematic project. Figure 8–2 on page 8–4 shows the design flow with the Cadence Allegro Design Entry CIS software.

Routing or pin assignment changes made in a board layout tool or a Cadence Allegro Design Entry CIS symbol cannot be back-annotated to the Quartus II software.

For more information about the Cadence Allegro Design Entry CIS software, including licensing, support, usage, training, and product updates, refer to the Help in the software, go to the Cadence (www.cadence.com) or go to the EMA Design Automation website (www.ema-eda.com).

Creating a Cadence Allegro Design Entry CIS Project

The Cadence Allegro Design Entry CIS software has built-in support for creating schematic symbols using pin assignment information imported from the Quartus II software.

To create a new project in the Cadence Allegro Design Entry CIS software, follow these steps:

1. On the File menu, point to New and click Project. The New Project wizard starts.

   When you create a new project, you can select the PC Board wizard, the Programmable Logic wizard, or a blank schematic.

2. Select the PC Board wizard to create a project where you can select which part libraries to use, or select a blank schematic.

   The Programmable Logic wizard only builds an FPGA logic design in the Cadence Allegro Design Entry CIS software.

Your new project is in the specified location and consists of the following files:

- OrCAD Capture Project File (.opj)
- Schematic Design File (.dsn)
Generating a Part

After you create a new project or open an existing project in the Cadence Allegro Design Entry CIS software, you can generate a new schematic symbol based on your Quartus II FPGA design. You can also update an existing symbol. The Cadence Allegro Design Entry CIS software stores component symbols in OrCAD Library File (.olb). When you place a symbol in a library attached to a project, it is immediately available for instantiation in the project schematic.

You can add symbols to an existing library or you can create a new library specifically for the symbols generated from your FPGA designs. To create a new library, follow these steps:

1. On the File menu, point to New and click Library in the Cadence Allegro Design Entry CIS software to create a default library named library1.olb. This library appears in the Library folder in the Project Manager window of the Cadence Allegro Design Entry CIS software.

2. To specify a desired name and location for the library, right-click the new library and select Save As. Saving the new library creates the library file.

You can now create a new symbol to represent your FPGA design in your schematic. To generate a schematic symbol, follow these steps:

1. Start the Cadence Allegro Design Entry CIS software.

2. On the Tools menu, click Generate Part. The Generate Part dialog box appears (Figure 8–6).

3. To specify the .pin from your Quartus II design, in the Netlist/source file type field, click Browse.

4. In the Netlist/source file type list, select Altera Pin File.
5. Type the new part name.

6. Specify the **Destination part library** for the symbol. Failing to select an existing library for the part creates a new library with a default name that matches the name of your Cadence Allegro Design Entry CIS project.

7. To create a new symbol for this design, select **Create new part**. If you updated your `.pin` in the Quartus II software and want to transfer any assignment changes to an existing symbol, select **Update pins on existing part in library**.

8. Select any other desired options and set **Implementation type** to `<none>`. The symbol is for a primitive library part based only on the `.pin` and does not require special implementation. Click **OK**.

9. Review the Undo warning and click **Yes** to complete the symbol generation.

You can locate the generated symbol in the selected library or in a new library found in the **Outputs** folder of the design in the Project Manager window (Figure 8–7). Double-click the name of the new symbol to see its graphical representation and edit it manually using the tools available in the Cadence Allegro Design Entry CIS software.

**Figure 8–7. Project Manager Window**

For more information about creating and editing symbols in the Cadence Allegro Design Entry CIS software, refer to the Help in the software.

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For more information about creating and editing symbols in the Cadence Allegro Design Entry CIS software, refer to the Help in the software.
Splitting a Part

After saving a new symbol in a project library, you can fracture the symbol into multiple parts called sections. Fracturing a part into separate sections is useful for FPGA designs. A single symbol for most FPGA packages might be too large for a single schematic page. Splitting the part into separate sections allows you to organize parts of the symbol by function, creating cleaner circuit schematics. For example, you can create one slot for an I/O symbol, a second slot for a JTAG symbol, and a third slot for a power/ground symbol. Figure 8–8 shows a part fractured into separate sections.

Figure 8–8. Splitting a Symbol into Multiple Sections

Notes to Figure 8–8:

1. Figure 8–8 represents a Cyclone device with JTAG or passive serial (PS) mode configuration option settings. Symbols created for other devices or other configuration modes might have different sets of configuration pins, but can be fractured in a similar manner.

2. The power/ground section shows only a representation of power and ground pins because the device contains a high number of power and ground pins.

Although symbol generation in the Design Entry CIS software refers to symbol fractures as sections, the other tools described in this chapter use different names to refer to symbol fractures.
To split a part into sections, select the part in its library in the Project Manager window of the Cadence Allegro Design Entry CIS software. On the Tools menu, click **Split Part** or right-click the part and choose **Split Part**. The **Split Part Section Input Spreadsheet** appears (Figure 8–9).

**Figure 8–9. Split Part Section Input Spreadsheet**

Each row in the spreadsheet represents a pin in the symbol. The **Section** column indicates the section of the symbol to which each pin is assigned. You can locate all pins in a new symbol in section 1. You can change the values in the **Section** column to assign pins to various sections of the symbol. You can also specify the side of a section on the location of the pin by changing the values in the **Location** column. When you are ready, click **Split**. A new symbol appears in the same library as the original with the name `<original part name>_Split1`.

View and edit each section individually. To view the new sections of the part, double-click the part. The Part Symbol Editor window appears and the first section of the part displays for editing. On the View menu, click **Package** to view thumbnails of all the part sections. To edit the section of the symbol, double-click the thumbnail.

For more information about splitting parts into sections and editing symbol sections in the Cadence Allegro Design Entry CIS software, refer to the Help in the software.
Instantiating a Symbol in a Design Entry CIS Schematic

After saving a new symbol in a library in your Cadence Allegro Design Entry CIS project, you can instantiate the new symbol on a page in your schematic. Open a schematic page in the Project Manager window of the Cadence Allegro Design Entry CIS software. To add the newly created symbol to your schematic on the schematic page, on the Place menu, click Part. The Place Part dialog box appears (Figure 8–10).

Select the new symbol library location and the newly created part name. If you select a part that is split into sections, you can select the section to place from the Part pop-up menu. Click OK. The symbol attaches to your cursor for placement in the schematic. To place the symbol, click on the schematic page.

For more information about using the Cadence Allegro Design Entry CIS software, refer to the Help in the software.

Altera Libraries for the Cadence Allegro Design Entry CIS Software

Altera provides downloadable .olb for many of its device packages. You can add these libraries to your Cadence Allegro Design Entry CIS project and update the symbols with the pin assignments contained in the .pin generated by the Quartus II software. You can use the downloaded library symbols as a base for creating custom schematic symbols with your pin assignments that you can edit or fracture. This method increases productivity by reducing the amount of time it takes to create and edit a new symbol.

To use the Altera-provided libraries with your Cadence Allegro Design Entry CIS project, follow these steps:

1. Download the library of your target device from the Download Center page found through the Support page on the Altera website (www.altera.com).
2. Create a copy of the appropriate .olb to maintain the original symbols. Place the copy in a convenient location, such as your Cadence Allegro Design Entry CIS project directory.

3. In the Project Manager window of the Cadence Allegro Design Entry CIS software, click once on the Library folder to select it. On the Edit menu, click Project or right-click the Library folder and choose Add File to select the copy of the downloaded .olb and add it to your project. You can locate the new library in the list of part libraries for your project.

4. On the Tools menu, click Generate Part. The Generate Part dialog box appears (Figure 8–11).

**Figure 8–11. Generate Part Dialog Box**

5. In the Netlist/source file field, click Browse to specify the .pin in your Quartus II design.

6. From the Netlist/source file type list, select Altera Pin File.

7. For Part name, type the name of the target device the same as it appears in the downloaded library file. For example, if you are using a device from the CYCLONE06.OLB library, type the part name to match one of the devices in this library such as ep1c6f256. You can rename the symbol in the Project Manager window after updating the part.

8. Set the Destination part library to the copy of the downloaded library you added to the project.

9. Select Update pins on existing part in library. Click OK.

10. Click Yes.
The symbol is updated with your pin assignments. Double-click the symbol in the Project Manager window to view and edit the symbol. On the View menu, click Package if you want to view and edit other sections of the symbol. If the symbol in the downloaded library is fractured into sections, you can edit each section but you cannot further fracture the part. You can generate a new part without using the downloaded part library if you require additional sections.

For more information about creating, editing, and fracturing symbols in the Cadence Allegro Design Entry CIS software, refer to the Help in the software.

**Conclusion**

Transferring a complex, high-pin-count FPGA design to a PCB for prototyping or manufacturing is a daunting process and can lead to errors in the PCB netlist or design, especially when different engineers are working on different parts of the project. The design workflow available when the Quartus II software is used with tools from Cadence assists the FPGA designer and the board designer in preventing such errors and focusing all attention on the design.

**Document Revision History**

Table 8–3 shows the revision history for this chapter.

![Table 8–3. Document Revision History](image)

For previous versions of the Quartus II Handbook, refer to the [Quartus II Handbook Archive](#).
This chapter provides guidelines for reviewing printed circuit board (PCB) schematics with the Quartus® II software. Altera FPGAs and CPLDs offer a multitude of configurable options to allow you to implement a custom application-specific circuit on your PCB.

Your Quartus II project provides important information specific to your programmable logic design, which you can use in conjunction with the device literature available on Altera’s website to ensure that you implement the correct board-level connections in your schematic.

This chapter highlights the important options in the Quartus II software, including Settings dialog box options, the Fitter report, and Messages window to which you should refer when creating and reviewing your PCB schematic. The Quartus II software also provides useful tools, such as the Pin Planner and the SSN Analyzer, to assist you during your PCB schematic review process.

The “Reviewing Quartus II Software Settings” section provides information about the settings you can make in the Quartus II software to help you review your PCB schematic. After verifying options in the Quartus II software, you can compile your design and use the data generated in the Fitter report, which is described in “Reviewing Device Pin-Out Information in the Fitter Report” on page 9–4 to verify settings in your PCB schematic. You should also ensure that you carefully review error and warning messages, as described in “Reviewing Compilation Error and Warning Messages” on page 9–6.

In addition to verifying your settings in the Settings dialog box and Fitter report, and checking messages, you can turn on additional settings, as described in “Using Additional Quartus II Software Features” on page 9–6 and “Running the HardCopy Design Readiness Check” on page 9–6.

Finally, Quartus II software tools, such as the Pin Planner and the SSN Analyzer, described in “Using Additional Quartus II Software Tools” on page 9–7, help you to verify proper I/O placement.

You should use this chapter in conjunction with Altera’s device family-specific literature.

For more information, refer to the Schematic Review Worksheets and the Pin Connection Guidelines pages of the Altera.com website.
Reviewing Quartus II Software Settings

The **Device** dialog box in the Quartus II software allows you to specify device-specific assignments and settings. You can use the **Device** dialog box to specify general project-wide options, including specific device and pin options, which help you to implement correct board-level connections in your PCB schematic.

The **Device** dialog box provides project-specific device information, including the target device and any migration devices you specify. Using migration devices can impact the number of available user I/O pins and internal resources, as well as require connection of some user I/O pins to power/ground pins to support migration.

If you want to use vertical migration, which allows you to use different devices with the same package, you can specify your list of migration devices in the **Migration Devices** dialog box. The Fitter places the pins in your design based on your targeted migration devices, and allows you to use only I/O pins that are common to all of the migration devices.

For more information about the **Migration Devices** dialog box in the Quartus II software, refer to *Migration Devices Dialog Box* in Quartus II Help.

If a migration device has pins that are power or ground, but the pins are also user I/O pins on a different device in the migration path, the Fitter ensures that these pins are not used as user I/O pins. You must ensure that these pins are connected to the appropriate plane on the PCB.

If you are migrating from a smaller device with NC (no-connect) pins to a larger device with power or ground pins in the same package, you can safely connect the NC pins to power or ground pins to facilitate successful migration.

**Device and Pins Options Dialog Box Settings**

You can verify important design-specific data in the **Device and Pin Options** dialog box when reviewing your PCB schematic, including options found on the **Configuration**, **Unused Pin**, **Dual-Purpose Pins**, and **Voltage** pages.

**Configuration Page Settings**

The **Configuration** page of the **Device and Pin Options** dialog box specifies the configuration scheme and configuration device for the target device. Use the **Configuration** page settings to verify the configuration scheme with the MSEL pin settings used on your PCB schematic and the I/O voltage of the configuration scheme.

Your specific configuration settings may impact the availability of some dual-purpose I/O pins in user mode. Refer to “**Dual-Purpose Pins Page Settings**” on page 9–3 for more information.
Unused Pin Page Settings

The Unused Pin page specifies the behavior of all unused pins in your design. Use the Unused Pin page to ensure that unused pin settings are compatible with your PCB. For example, if you reserve all unused pins as outputs driving ground, you must ensure that you do not connect unused I/O pins to VCC pins on your PCB. Connecting unused I/O pins to VCC pins may result in contention that could lead to higher than expected current draw and possible device overstress.

The Reserve all unused pins list shows available unused pin state options for the target device. The default state for each pin is the recommended setting for each device family.

When you reserve a pin as output driving ground, the Fitter connects a ground signal to the output pin internally. You should connect the output pin to the ground plane on your PCB, although you are not required to do so. Connecting the output driving ground to the ground plane is known as creating a virtual ground pin, which helps to minimize simultaneous switching noise (SSN) and ground bounce effects.

Dual-Purpose Pins Page Settings

The Dual-Purpose Pins page specifies how configuration pins should be used after device configuration completes. You can set the function of the dual-purpose pins by selecting a value for a specific pin in the Dual-purpose pins list. Pin functions should match your PCB schematic. The available options on the Dual-Purpose Pins page may differ depending on the selected configuration mode.

Voltage Page Settings

The Voltage page specifies the default VCCIO I/O bank voltage and the default I/O bank voltage for the pins on the target device. VCCIO I/O bank voltage settings made in the Voltage page are overridden by I/O standard assignments made on I/O pins in their respective banks. Refer to the “Reviewing Device Pin-Out Information in the Fitter Report” on page 9–4 for more details about the I/O bank voltages for your design.

Error Detection CRC Page Settings

The Error Detection CRC page specifies error detection cyclic redundancy check (CRC) use for the target device. When Enable error detection CRC is turned on, the device checks the validity of the programming data in the devices. Any changes made in the data while the device is in operation generates an error.

Turning on the Enable open drain on CRC error pin option allows the CRC ERROR pin to be set as an open-drain pin in some devices, which decouples the voltage level of the CRC ERROR pin from VCCIO voltage. You must connect a pull-up resistor to the CRC ERROR pin on your PCB if you turn on this option.

In addition to settings in the Device dialog box, you should verify settings in the Voltage page of the Settings dialog box.

For more information about the Device and Pins Options dialog box in the Quartus II software, refer to Device and Pin Options Dialog Box in Quartus II Help.
Voltage Page Settings

The Voltage page, under Operating Settings and Conditions in the Settings dialog box, allows you to specify voltage operating conditions for timing and power analyses. Ensure that the settings in the Voltage page match the settings in your PCB schematic, especially if the target device includes transceivers.

The Voltage page settings requirements differ depending on the settings of the transceiver instances in the design. Refer to the Fitter report for the required settings, and verify that the voltage settings are correctly set up for your PCB schematic.

For more information about voltage settings, refer to the Pin Connection Guidelines page of the Altera.com website.

Once you verify your settings in the Device and Settings dialog boxes, you can verify your device pin-out with the Fitter report.

Reviewing Device Pin-Out Information in the Fitter Report

After you compile your design, you can use the reports in the Resource section of the Fitter report to check your device pin-out in detail.

The Input Pins, Output Pins, and Bidirectional Pins reports identify all the user I/O pins in your design and the features enabled for each I/O pin. For example, you can find use of weak internal pull-ups, PCI clamp diodes, and on-chip termination (OCT) pin assignments in these sections of the Fitter report. You can check the pin assignments reported in the Input Pins, Output Pins, and Bidirectional Pins reports against your PCB schematic to determine whether your PCB requires external components.

These reports also identify whether you made pin assignments or if the Fitter automatically placed the pins. If the Fitter changed your pin assignments, you should make these changes user assignments because the location of pin assignments made by the Fitter may change with subsequent compilations.
Figure 9–1 shows the pins the Fitter chose for the OCT external calibration resistor connections (RUP/RDN) and the name of the associated termination block in the Input Pins report. You should make these types of assignments user assignments.

The I/O Bank Usage report provides a high-level overview of the VCCIO and VREF requirements for your design, based on your I/O assignments. Verify that the requirements in this report match the settings in your PCB schematic. All unused I/O banks, and all banks with I/O pins with undefined I/O standards, default the VCCIO voltage to the voltage defined in the Voltage page of the Device and Pin Options dialog box.

The All Package Pins report lists all the pins on your device, including unused pins, dedicated pins and power/ground pins. You can use this report to verify pin characteristics, such as the location, name, usage, direction, I/O standard and voltage for each pin with the pin information in your PCB schematic. In particular, you should verify the recommended voltage levels at which you connect unused dedicated inputs and I/O and power pins, especially if you selected a migration device. Use the All Package Pins report to verify that you connected all the device voltage rails to the voltages reported.

Errors commonly reported include connecting the incorrect voltage to the predriver supply (VCCPD) pin in a specific bank, or leaving dedicated clock input pins floating. Unused input pins that should be connected to ground are designated as GND+ in the Pin Name/Usage column in the All Package Pins report.

You can also use the All Package Pins report to check transceiver-specific pin connections and verify that they match the PCB schematic. Unused transceiver pins have the following requirements, based on the pin designation in the Fitter report:

- **GXB_GND**—Unused GXB receiver or dedicated reference clock pin. This pin must be connected to GXB_GND through a 10k Ohm resistor.
- **GXB_NC**—Unused GXB transmitter or dedicated clock output pin. This pin must be disconnected.
Some transceiver power supply rails have dual voltage capabilities, such as VCCA_L/R and VCCH_L/R, that depend on the settings you created for the ALTGX MegaWizard Plug-In Manager. Because these user-defined settings overwrite the default settings, you should use the All Package Pins report to verify that these power pins on the device symbol in the PCB schematics are connected to the voltage required by the transceiver. An incorrect connection may cause the transceiver to function not as expected.

If your design includes a memory interface, the DQS Summary report provides an overview of each DQ pin group. You can use this report to quickly confirm that the correct DQ/DQS pins are grouped together. This section also provides information on DLL usage.

Finally, the Fitter Device Options report summarizes some of the settings made in the Device and Pin Options dialog box. Verify that these settings match your PCB schematics.

**Reviewing Compilation Error and Warning Messages**

If your project does not compile without error or warning messages, you should resolve the issues identified by the Compiler before signing off on your pin-out or PCB schematic. Error messages often indicate illegal or unsupported use of the device resources and IP.

Additionally, you should cross-reference fitting and timing analysis warnings with the design implementation. Timing may be constrained due to nonideal pin placement. You should investigate if you can reassign pins to different locations to prevent fitting and timing analysis warnings. Ensure that you review each warning and consider its potential impact on the design.

**Running the HardCopy Design Readiness Check**

You should run the HardCopy Design Readiness Check for designs including a HardCopy device in the migration path. This tool checks for issues that must be addressed prior to handing off the design to the Altera HardCopy Design Center for the HardCopy back-end process, including possible issues with device resource and I/O usage. Verify all warning messages generated during the HardCopy Design Readiness Check.

For more information about the HardCopy Design Readiness Check and designing for HardCopy devices in the Quartus II software, refer to the Quartus II Support for HardCopy Series Devices chapter in volume 1 of the Quartus II Handbook.

**Using Additional Quartus II Software Features**

You can generate IBIS files, which contain models specific to your design and selected I/O standards and options, with the Quartus II software.

Because board-level simulation is important to verify, you should check for potential signal integrity issues. You can turn on the Board-Level Signal Integrity feature in the EDA Tool Settings page of the Settings dialog box.
For more information about signal integrity analysis in the Quartus II software, refer to the *Signal Integrity Analysis with Third-Party Tools* chapter in volume 3 of the *Quartus II Handbook*.

Additionally, using advanced I/O timing allows you to enter physical PCB information to accurately model the load seen by an output pin. This feature facilitates accurate I/O timing analysis.

For more information about advanced I/O timing, refer to the *I/O Management* chapter in volume 2 of the *Quartus II Handbook*.

### Using Additional Quartus II Software Tools

This section describes additional tools found in the Quartus II software, specifically the Pin Planner and the SSN Analyzer, and how you can use these tools to assist you with reviewing your PCB schematics.

#### Pin Planner

The Quartus II Pin Planner helps you visualize, plan, and assign device I/O pins in a graphical view of the target device package. You can quickly locate various I/O pins and assign them design elements or other properties to ensure compatibility with your PCB layout.

You can use the Pin Planner to verify the location of clock inputs, and whether they have been placed on dedicated clock input pins, which is recommended when your design uses PLLs.

You can also use the Pin Planner to verify the placement of dedicated SERDES pins. SERDES receiver inputs can be placed only on DIFFIO_RX pins, while SERDES transmitter outputs can be placed only on DIFFIO_TX pins.

The Pin Planner gives a visual indication of signal-to-signal proximity in the Pad View window, and also provides information about differential pin pair placement, such as the placement of pseudo-differential signals.

For more information about the Pin Planner, refer to the *I/O Management* chapter in volume 2 of the *Quartus II Handbook*.

#### SSN Analyzer

The SSN Analyzer supports pin planning by estimating the voltage noise caused by the simultaneous switching of output pins on the device. Because of the importance of the potential SSN performance for a specific I/O placement, you can use the SSN Analyzer to analyze the effects of aggressor I/O signals on a victim I/O pin.

For more information about the SSN Analyzer, refer to the *Simultaneous Switching Noise (SSN) Analysis and Optimizations* chapter in volume 2 of the *Quartus II Handbook*. 
Conclusion

This chapter describes guidelines and descriptions of settings to verify when reviewing your PCB schematic with the Quartus II software. You can use settings in the Settings dialog box; information in the Fitter report and Messages window; and the Pin Planner and SSN Analyzer during the PCB schematic review process.

Document Revision History

Table 9–1 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
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<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Minor update of Pin Planner description for task and report windows.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Changed to new document template. No change to content.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This section introduces features in the Quartus® II software that you can use to optimize area, timing, power, and compilation time when you design for programmable logic devices (PLDs).

This section includes the following chapters:

- **Chapter 10, Design Optimization Overview**
  This chapter summarizes features in the Quartus II software that you can use to achieve the highest design performance when you design for PLDs, especially high density FPGAs.

- **Chapter 11, Reducing Compilation Time**
  This chapter describes techniques for reducing the amount of time it takes to compile and recompile your design, accelerating your design process.

- **Chapter 12, Timing Closure and Optimization**
  This chapter describes a broad spectrum of Quartus II software features and design techniques to improve timing performance when designing for Altera® devices. This chapter also explains how and when to use some of the features described in other chapters of the *Quartus II Handbook*.

- **Chapter 13, Power Optimization**
  This chapter describes the power-driven compilation feature and flow in detail, as well as low power design techniques that can further reduce power consumption in your design.

- **Chapter 14, Area Optimization**
  This chapter describes design techniques to reduce resource usage.

- **Chapter 15, Analyzing and Optimizing the Design Floorplan with the Chip Planner**
  You can use the Chip Planner to perform design analysis and create a design floorplan. This chapter discusses how to analyze and optimize the design floorplan with the Chip Planner.

- **Chapter 16, Netlist Optimizations and Physical Synthesis**
  This chapter explains how the physical synthesis optimizations in the Quartus II software can improve your quality of results. This chapter also provides information about preserving and writing out a new netlist, and provides guidelines for applying the various options.
This chapter introduces features in Altera’s Quartus® II software that you can use to achieve the highest design performance when you design for programmable logic devices (PLDs), especially high density FPGAs.

Introduction

Physical implementation can be an intimidating and challenging phase of the design process. The Quartus II software provides a comprehensive environment for FPGA designs, delivering unmatched performance, efficiency, and ease-of-use.

In a typical design flow, you must synthesize your design with Quartus II integrated synthesis or a third-party tool, place and route your design with the Fitter, and use the TimeQuest timing analyzer to ensure your design meets the timing requirements. With the PowerPlay Power Analyzer, you ensure the design’s power consumption is within limits.

Initial Compilation: Required Settings

This section describes the basic assignments and settings for your initial compilation. Check the following settings before compiling your design in the Quartus II software. Significantly varied compilation results can occur depending on the assignments that you set.

Device Settings

Device assignments determine the timing model that the Quartus II software uses during compilation. Choose the correct speed grade to obtain accurate results and the best optimization. The device size and the package determine the device pin-out and the available resources in the device.

Device Migration Settings

If you anticipate a change to the target device later in the design cycle, either because of changes in your design or other considerations, plan for the change at the beginning of your design cycle. Whenever you select a target device, you can also list any other compatible devices you can migrate by clicking on the Migration Devices button in the Device dialog box. If you plan to move your design to a HardCopy® device, make sure to select the device from the HardCopy list under Companion device in the Device dialog box.

Selecting the migration device and companion device early in the design cycle helps to minimize changes to your design at a later stage.
I/O Assignments

The I/O standards and drive strengths specified for a design affect I/O timing. Specify I/O assignments so that the Quartus II software uses accurate I/O timing delays in timing analysis and Fitter optimizations.

If there is no PCB layout requirement, then you do not need to specify pin locations. If your pin locations are not fixed due to PCB layout requirements, then leave the pin locations unconstrained. If your pin locations are already fixed, then make pin assignments to constrain the compilation appropriately.

For more information about recommendations for making pin assignments that can have a large effect on your results in smaller macrocell-based architectures, refer to Optimizing Resource Utilization (Macrocell-Based CPLDs) in the Timing Closure and Optimization chapter in volume 2 of the Quartus II Handbook.

Use the Assignment Editor and Pin Planner to assign I/O standards and pin locations.

For more information about I/O standards and pin constraints, refer to the appropriate device handbook. For more information about planning and checking I/O assignments, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

For information about using the Assignment Editor, refer to About the Assignment Editor in Quartus II Help.

Timing Requirement Settings

You must use comprehensive timing requirement settings to achieve the best results for the following reasons:

- Correct timing assignments enable the software to work hardest to optimize the performance of the timing-critical parts of your design and make trade-offs for performance. This optimization can also save area or power utilization in non-critical parts of your design.
- If enabled, the Quartus II software performs physical synthesis optimizations based on timing requirements.
- Depending on the Fitter Effort setting, the Fitter can reduce runtime if your design meets the timing requirements.

For more information about optimization with physical synthesis, refer to Physical Synthesis Optimization in the Timing Closure and Optimization chapter in volume 2 of the Quartus II Handbook.

For more information about reducing runtime by changing Fitter effort, refer to Fitter Settings Page in the Quartus II Help.

Use your real requirements to get the best results. If you apply more demanding timing requirements than you need, then increased resource usage, higher power utilization, increased compilation time, or all of these may result.
The Quartus II TimeQuest Timing Analyzer determines if the design implementation meets the timing requirement. The Compilation Report shows whether your design meets the timing requirements, while the timing analysis reporting commands provide detailed information about the timing paths.

To create timing constraints for the TimeQuest analyzer, create a Synopsys Design Constraints File (.sdc). You can also enter constraints in the TimeQuest GUI. Use the write_sdc command, or the Constraints menu in the TimeQuest analyzer. Click Write SDC File to write your constraints to an .sdc. You can add an .sdc to your project on the Quartus II Settings page under Timing Analysis Settings.

If you already have an .sdc in your project, using the write_sdc command from the command line or using the Write SDC File option from the TimeQuest GUI allows you to create a new .sdc that combines the constraints from your current .sdc and any new constraints added through the GUI or command window, or overwrites the existing .sdc with your newly applied constraints.

Ensure that every clock signal has an accurate clock setting constraint. If clocks arrive from a common oscillator, then they are related. Ensure that you set up all related or derived clocks in the constraints correctly. You must constrain all I/O pins that require I/O timing optimization. Specify both minimum and maximum timing constraints as applicable. If your design contains more than one clock or contains pins with different I/O requirements, make multiple clock settings and individual I/O assignments instead of using a global constraint. 1

Make any complex timing assignments required in your design, including false path and multicycle path assignments. Common situations for these types of assignments include reset or static control signals (when the time required for a signal to reach a destination is not important) or paths that have more than one clock cycle available for operation in a design. These assignments enable the Quartus II software to make appropriate trade-offs between timing paths and can enable the Compiler to improve timing performance in other parts of your design.

For more information about timing assignments and timing analysis, refer to The Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook and the Quartus II TimeQuest Timing Analyzer Cookbook.

To ensure that you apply constraints or assignments to all design nodes, you can report all unconstrained paths in your design with the Report Unconstrained Paths command in the Task pane of the Quartus II TimeQuest Timing Analyzer or the report_ucp Tcl command.

**Partitions and Floorplan Assignments for Incremental Compilation**

The Quartus II incremental compilation feature enables hierarchical and team-based design flows in which you can compile parts of your design while other parts of your design remain unchanged and import parts of your design from separate Quartus II projects.
Using incremental compilation for your design with good design partitioning methodology helps to achieve timing closure. Creating design partitions on some of the major blocks in your design and assigning them to LogicLock™ regions, reduces Fitter time and improves the quality and repeatability of the results. LogicLock regions are flexible, reusable floorplan location constraints that help you place logic on the target device. When you assign entity instances or nodes to a LogicLock region, you direct the Fitter to place those entity instances or nodes inside the region during fitting.

For more information about LogicLock regions, refer to About LogicLock Regions in Quartus II Help.

Using incremental compilation helps you achieve timing closure block by block and preserve the timing performance between iterations, which aid in achieving timing closure for the entire design. Incremental compilation may also help reduce compilation times.

For more information, refer to the “Incremental Compilation” section in the Reducing Compilation Time chapter in volume 2 of the Quartus II Handbook.

If you plan to use incremental compilation, you must create a floorplan for your design. If you are not using incremental compilation, creating a floorplan is optional.

For more information about guidelines to create partition and floorplan assignments for your design, refer to the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

Physical Implementation

Most optimization issues involve preserving previous results, reducing area, reducing critical path delay, reducing power consumption, and reducing runtime. The Quartus II software includes advisors to address each of these issues and helps you optimize your design. Run these advisors during physical implementation for advice about your specific design.

You can reduce the time spent on design iterations by following the recommended design practices for designing with Altera® devices. Design planning is critical for successful design timing implementation and closure.

For more information, refer to the Design Planning with the Quartus II Software chapter in volume 1 of the Quartus II Handbook.

Trade-Offs and Limitations

Many optimization goals can conflict with one another, so you might need to resolve conflicting goals. For example, one major trade-off during physical implementation is between resource usage and critical path timing, because certain techniques (such as logic duplication) can improve timing performance at the cost of increased area. Similarly, a change in power requirements can result in area and timing trade-offs, such as if you reduce the number of available high-speed tiles, or if you attempt to shorten high-power nets at the expense of critical path nets.
In addition, system cost and time-to-market considerations can affect the choice of device. For example, a device with a higher speed grade or more clock networks can facilitate timing closure at the expense of higher power consumption and system cost.

Finally, not all designs can be realized in a hardware circuit with limited resources and given constraints. If you encounter resource limitations, timing constraints, or power constraints that cannot be resolved by the Fitter, consider rewriting parts of the HDL code.

For more information, refer to the Timing Closure and Optimization and Area Optimization chapters in volume 2 of the Quartus II Handbook.

Preserving Results and Enabling Teamwork

For some Quartus II Fitter algorithms, small changes to the design can have a large impact on the final result. For example, a critical path delay can change by 10% or more because of seemingly insignificant changes. If you are close to meeting your timing objectives, you can use the Fitter algorithm to your advantage by changing the fitter seed, which changes the pseudo-random result of the Fitter.

Conversely, if you cannot meet timing on a portion of your design, you can partition that portion and prevent it from recompiling if an unrelated part of the design is changed. This feature, known as incremental compilation, can reduce the Fitter runtimes by up to 70% if the design is partitioned, such that only small portions require recompilation at any one time.

When you use incremental compilation, you can apply design optimization options to individual design partitions and preserve performance in other partitions by leaving them untouched. Many optimization techniques often result in longer compilation times, but by applying them only on specific partitions, you can reduce this impact and complete iterations more quickly.

In addition, by physically floorplanning your partitions with LogicLock regions, you can enable team-based flows and allow multiple people to work on different portions of the design.

For more information, refer to Quartus II Incremental Compilation for Hierarchical and Team-Based Designs in volume 1 of the Quartus II Handbook and About Incremental Compilation in Quartus II Help.

Reducing Area

By default, the Quartus II Fitter might physically spread a design over the entire device to meet the set timing constraints. If you prefer to optimize your design to use the smallest area, you can change this behavior. If you require reduced area, you can enable certain physical synthesis options to modify your netlist to create a more area-efficient implementation, but at the cost of increased runtime and decreased performance.

For more information, refer to the Netlist Optimizations and Physical Synthesis, Timing Closure and Optimization, and Area Optimization chapters in volume 2 and the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
Reducing Critical Path Delay

To meet complex timing requirements involving multiple clocks, routing resources, and area constraints, the Quartus II software offers a close interaction between synthesis, timing analysis, floorplan editing, and place-and-route processes.

By default, the Quartus II Fitter tries to meet the specified timing requirements and stops trying when the requirements are met. Therefore, using realistic constraints is important to successfully close timing. If you under-constrain your design, you may get sub-optimal results. By contrast, if you over-constrain your design, the Fitter might over-optimize non-critical paths at the expense of true critical paths. In addition, you might incur an increased area penalty. Compilation time may also increase because of excessively tight constraints.

If your resource usage is very high, the Quartus II Fitter might have trouble finding a legal placement. In such circumstances, the Fitter automatically modifies some of its settings to try to trade off performance for area.

The Quartus II Fitter offers a number of advanced options that can help you improve the performance of your design when you properly set constraints. Use the Timing Optimization Advisor to determine which options are best suited for your design.

If you use incremental compilation, you can help resolve inter-partition timing requirements by locking down results, one partition at a time, or by guiding the placement of the partitions with LogicLock regions. You might be able to improve the timing on such paths by placing the partitions optimally to reduce the length of critical paths. Once your inter-partition timing requirements are met, use incremental compilation to preserve the results and work on partitions that have not met timing requirements.

In high-density FPGAs, routing accounts for a major part of critical path timing. Because of this, duplicating or retiming logic can allow the Fitter to reduce delay on critical paths. The Quartus II software offers push-button netlist optimizations and physical synthesis options that can improve design performance at the expense of considerable increases of compilation time and area. Turn on only those options that help you keep reasonable compilation times and resource usage. Alternately, you can modify your HDL to manually duplicate or adjust the timing logic.

Reducing Power Consumption

The Quartus II software has features that help reduce design power consumption. The PowerPlay power optimization options control the power-driven compilation settings for Synthesis and the Fitter.

For more information, refer to the Power Optimization chapter in volume 2 of the Quartus II Handbook.

Reducing Runtime

Many Fitter settings influence compilation time. Most of the default settings in the Quartus II software are set for reduced compilation time. You can modify these settings based on your project requirements.
The Quartus II software supports parallel compilation in computers with multiple processors. This can reduce compilation times by up to 15% while giving the identical result as serial compilation.

You can also reduce compilation time with your iterations by using incremental compilation. Use incremental compilation when you want to change parts of your design, while keeping most of the remaining logic unchanged.

**Using Quartus II Tools**

The following sections describe several Quartus II tools that you can use to help optimize your design.

**Design Analysis**

The Quartus II software provides tools that help with a visual representation of your design. You can use the RTL Viewer to see a schematic representation of your design before synthesis and place-and-route. The Technology Map Viewer provides a schematic representation of the design implementation in the selected device architecture after synthesis and place-and-route. It can also include timing information.

With incremental compilation, the Design Partition Planner and the Chip Planner allow you to partition and layout your design at a higher level. In addition, you can perform many different tasks with the Chip Planner, including: making floorplan assignments, implementing engineering change orders (ECOs), and performing power analysis. Also, you can analyze your design and achieve a faster timing closure with the Chip Planner. The Chip Planner provides physical timing estimates, critical path display, and a routing congestion view to help guide placement for optimal performance.

For more information, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Designs* and *Best Practices for Incremental Compilation Partitions and Floorplan Assignments* chapters in volume 1 and the *Engineering Change Management with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*.

**Advisors**

The Quartus II software includes several advisors to help you optimize your design and reduce compilation time. You can complete your design faster by following the recommendations in the Compilation Time Advisor, Incremental Compilation Advisor, Timing Optimization Advisor, Area Optimization Advisor, Resource Optimization Advisor, and Power Optimization Advisor. These advisors give recommendations based on your project settings and your design constraints.

For more information about advisors, refer to *Running Advisors in the Quartus II Software* in Quartus II Help.
Design Space Explorer

Use the Design Space Explorer (DSE) to find optimal settings in the Quartus II software. DSE automatically tries different combinations of netlist optimizations and advanced Quartus II software compiler settings, and reports the best settings for your design, based on your chosen primary optimization goal. You can try different seeds with the DSE if you are fairly close to meeting your timing or area requirements and find one seed that meets timing or area requirements. Finally, the DSE can run the different compilations on multiple computers in parallel, which shortens the timing closure process.

For more information, refer to About Design Space Explorer in Quartus II Help.

Conclusion

The Quartus II software includes a number of features and tools that you can use to optimize area, timing, power, and compilation time when you design for programmable logic devices (PLDs).

Document Revision History

Table 10–1 shows the revision history for this chapter.

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<th>Version</th>
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<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>Added the “Initial Compilation: Required Settings” on page 10–1 section. This section was moved from the Area Optimization chapter of the Quartus II Handbook. Minor updates to delineate division of Timing and Area optimization chapters.</td>
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<tr>
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<td>Removed survey link.</td>
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</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
11. Reducing Compilation Time

The Quartus® II software offers several features and techniques to help reduce compilation time.

This chapter describes techniques to reduce compilation time when designing for Altera® devices, and includes the following topics:

■ “Compilation Time Optimization Techniques”
■ “Compilation Time Advisor” on page 11–2
■ “Strategies to Reduce the Overall Compilation Time” on page 11–2
■ “Reducing Synthesis Time and Synthesis Netlist Optimization Time” on page 11–5
■ “Reducing Placement Time” on page 11–5
■ “Reducing Routing Time” on page 11–7
■ “Reducing Static Timing Analysis Time” on page 11–8
■ “Setting Process Priority” on page 11–8

Compilation Time Optimization Techniques

The Analysis and Synthesis and Fitter modules consume the majority of time in a compilation. The Analysis and Synthesis module includes physical synthesis optimizations performed during synthesis, if you have turned on physical synthesis optimizations. The Fitter includes two steps, placement and routing, and also includes physical synthesis if you turned on the physical synthesis option with Normal or Extra effort levels. The Flow Elapsed Time section of the Compilation Report shows the duration of the Analysis and Synthesis and Fitter modules. The Fitter Messages report in the Fitter section of the Compilation Report displays the elapsed time for placement and routing processes.

Placement is the process of finding optimum locations for the logic in your design. Placement includes Quartus II pre-Fitter operations, which place dedicated logic such as clocks, PLLs, and transceiver blocks. Routing is the process of connecting the nets between the logic in your design. Finding better placement for the logic in your design requires more compilation time. Good logic placement allows you to more easily meet your timing requirements and makes your design easier to route.
Example 11–1 shows examples of messages with each time component in two-digit format, and days shown only if applicable:

Example 11–1.

Info: Fitter placement operations ending: elapsed time = <days:hours:minutes:seconds>
Info: Fitter routing operations ending: elapsed time = <days:hours:minutes:seconds>

Example 11–2 shows an info message displayed while the Fitter is running (including Placement and Routing). The Message window displays this message every hour to indicate Fitter operations are progressing normally.

Example 11–2.

Info: Placement optimizations have been running for 4 hour(s)

Compilation Time Advisor

A Compilation Time Advisor is available to help you to reduce compilation time. Run the Compilation Time Advisor on the Tools menu by pointing to Advisors and clicking Compilation Time Advisor. You can find all the compilation time optimizing techniques described in this section in the Compilation Time Advisor as well.

Strategies to Reduce the Overall Compilation Time

This section discusses strategies to reduce overall compilation time, including the following topics:

- “Using Parallel Compilation with Multiple Processors”
- “Using Incremental Compilation” on page 11–4
- “Using the Smart Compilation Setting” on page 11–4

Using Parallel Compilation with Multiple Processors

The Quartus II software can detect the number of processors available on a computer and use multiple processors to reduce compilation time. You can also control the number of processors used during a compilation on a per user basis. The Quartus II software can use up to 16 processors to run algorithms in parallel and reduce compilation time. The Quartus II software turns on parallel compilation by default to enable the software to detect available multiple processors. You can specify the maximum number of processors that the software can use if you want to reserve some of the available processors for other tasks.

Do not consider processors with Intel Hyper-Threading as more than one processor. If you have a single processor with Intel Hyper-Threading enabled, you should set the number of processors to one. Altera recommends that you do not use the Intel Hyper-Threading feature for Quartus II compilations, because it can increase runtimes.
The software does not necessarily use all the processors that you specify during a given compilation. Additionally, the software never uses more than the specified number of processors, enabling you to work on other tasks on your computer without it becoming slow or less responsive.

If you have partitioned your design and enabled parallel compilation, the Quartus II software can use different processors to compile those partitions simultaneously during Analysis and Synthesis. This can cause higher peak memory usage during Analysis and Synthesis.

You can reduce the compilation time by up to 10% on systems with two processing cores and by up to 20% on systems with four cores. With certain design flows in which timing analysis runs alone, multiple processors can reduce the time required for timing analysis by an average of 10% when using two processors. This reduction can reach an average of 15% when using four processors.

The actual reduction in compilation time when using incremental compilation partitions depends on your design and on the specific compilation settings. For example, compilations with multi-corner optimization turned on benefit more from using multiple processors than do compilations without multi-corner optimization. The runtime requirement is not reduced for some other compilation goals, such as Analysis and Synthesis. The Fitter (`quartus_fit`) and the Quartus II TimeQuest Timing Analyzer (`quartus_sta`) stages in the compilation can, in certain cases, benefit from the use of multiple processors. The Flow Elapsed Time panel of the Compilation Report shows the average number of processors for these stages. The Parallel Compilation panel of the appropriate report, such as the Fitter report, shows a more detailed breakdown of processor usage. This panel is displayed only if parallel compilation is enabled.

Parallel compilation is available for Arria® series, Cyclone®, HardCopy III, HardCopy IV, MAX® II, MAX V (limited support), and Stratix® series devices.

For more information, refer to Processing Page (Options Dialog Box) in Quartus II Help.

For more information about how to control the number of processors used during compilation for a specific project, refer to Compilation Process Settings Page (Settings Dialog Box) in Quartus II Help.

You can also set the number of processors available for Quartus II compilation using the following Tcl command in your script:

```
set_global_assignment -name NUM_PARALLEL_PROCESSORS <value>
```

In this case, `<value>` is an integer from 1 to 16.

If you want the Quartus II software to detect the number of processors and use all the processors for the compilation, include the following Tcl command in your script:

```
set_global_assignment -name NUM_PARALLEL_PROCESSORS ALL
```

The use of multiple processors does not affect the quality of the fit. For a given Fitter seed on a specific design, the fit is exactly the same, regardless of whether the Quartus II software uses one processor or multiple processors. The only difference between compilations using a different number of processors is the compilation time.
Using Incremental Compilation

The incremental compilation feature can accelerate design iteration time by up to 70% for small design changes, and helps you reach design timing closure more efficiently. You can speed up design iterations by recompiling only a particular design partition and merging results with previous compilation results from other partitions. You can also use physical synthesis optimization techniques for specific design partitions while leaving other parts of your design untouched to preserve performance.

If you are using a third-party synthesis tool, you can create separate atom netlist files for the parts of your design that you already have synthesized and optimized so that you update only the parts of your design that change.

In the standard incremental compilation design flow, you can divide the top-level design into partitions, which the software can compile and optimize in the top-level Quartus II project. You can preserve fitting results and performance for completed partitions while other parts of your design are changing. Incremental compilation reduces the compilation time for each design iteration because the software does not recompile the unchanged partitions in your design.

The incremental compilation feature facilitates team-based design flows by enabling designers to create and optimize design blocks independently, when necessary, and supports third-party IP integration.

For more information about the full incremental compilation flow in the Quartus II software, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook. For more information about creating multiple netlist files in third-party tools for use with incremental compilation, refer to the appropriate chapter in Section IV. Synthesis in volume 1 of the Quartus II Handbook.

For more information about incremental compilation, refer to About Incremental Compilation in Quartus II Help.

Using the Smart Compilation Setting

Not all compilation processes are required for when recompiling your design. Smart compilation skips unnecessary Compiler stages, such as Analysis and Synthesis. This feature is different from incremental compilation, which can compile parts of your design while preserving results for unchanged parts.

This setting is especially useful when you perform multiple compilations during the optimization phase of your design process. Smart compilation requires more disk space than regular compilation. To turn on smart compilation, on the Assignments menu, click Settings. In the Category list, select Compilation Process Settings and turn on Use smart compilation.

For more information on how to use smart compilation, refer to Compilation Process Settings Page (Settings Dialog Box) in Quartus II Help.
Reducing Synthesis Time and Synthesis Netlist Optimization Time

You can reduce synthesis time without affecting the Fitter time by reducing your use of netlist optimizations and by using incremental compilation (with Netlist Type set to Post-Synthesis). For tips on reducing synthesis time when using third-party EDA synthesis tools, refer to your synthesis software’s documentation.

Settings to Reduce Synthesis Time and Synthesis Netlist Optimization Time

You can use Quartus II integrated synthesis to synthesize and optimize HDL designs, and you can use synthesis netlist optimizations to optimize netlists that were synthesized by third-party EDA software. When using Quartus II Integrated Synthesis, you can also enable specific options in the Physical Synthesis Optimizations window before performing Analysis and Synthesis. Netlist optimizations can cause the Analysis and Synthesis module to take much longer to run. Read the Analysis and Synthesis messages to determine how much time these optimizations take. The compilation time spent in Analysis and Synthesis is usually short compared to the compilation time spent in the Fitter.

If your design meets your performance requirements without synthesis netlist optimizations, turn off the optimizations to save time. If you require synthesis netlist optimizations to meet performance, you can optimize parts of your design hierarchy separately to reduce the overall time spent in Analysis and Synthesis.

Turn off settings that are not useful. In general, if you carry over compilation settings from a previous project, evaluate all settings and keep only those that you need.

Use Appropriate Coding Style to Reduce Synthesis Time

Your HDL coding style can also affect the synthesis time. For example, if you want to infer RAM blocks from your code, you must follow the guidelines for inferring RAMs. If RAM blocks are not inferred properly, the software implements those blocks as registers.

If you are trying to infer a large memory block, the software consumes more resources in the FPGA. This can cause routing congestion and increasing compilation time significantly. If you see high routing utilizations in certain blocks, it is a good idea to review the code for such blocks.

For more information about coding guidelines, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

Reducing Placement Time

The time required to place a design depends on two factors: the number of ways the logic in your design can be placed in the device and the settings that control how hard the Placer works to find a good placement. You can reduce the placement time in two ways:

■ Change the settings for the placement algorithm.
■ Use incremental compilation to preserve the placement for the unchanged parts of your design.
Sometimes there is a trade-off between placement time and routing time. Routing time can increase if the placer does not run long enough to find a good placement. When you reduce placement time, ensure that it does not increase routing time and negate the overall time reduction.

**Fitter Effort Setting**

The highest Fitter effort setting, **Standard Fit**, requires the most runtime, but does not always yield a better result than using the default **Auto Fit**. For designs with very tight timing requirements, both **Auto Fit** and **Standard Fit** use the maximum effort during optimization. Altera recommends using **Auto Fit** for reducing compilation time. If you are certain that your design has only easy-to-meet timing constraints, you can select **Fast Fit** for an even greater runtime savings.

**Placement Effort Multiplier Settings**

You can control the amount of time the Fitter spends in placement by reducing with the **Placement Effort Multiplier** option. On the Assignments menu, click **Settings**. Select **Fitter Settings**, and click **More Settings**. Under **Existing Option Settings**, select **Placement Effort Multiplier**. The default is **1.0**. Legal values must be greater than **0** and can be non-integer values. Numbers between **0** and **1** can reduce fitting time, but also can reduce placement quality and design performance.

**Physical Synthesis Effort Settings**

Physical synthesis options enable you to optimize your post-synthesis netlist and improve your timing performance. These options, which affect placement, can significantly increase compilation time.

If your design meets your performance requirements without physical synthesis options, turn them off to reduce compilation time. For example, if some or all of the physical synthesis algorithm information messages display an improvement of **0 ps**, turning off physical synthesis can reduce compilation time.

You also can use the **Physical synthesis effort** setting on the **Physical Synthesis Optimizations** page to reduce the amount of extra compilation time used by these optimizations.

The **Fast** setting directs the Quartus II software to use a lower level of physical synthesis optimization. Compared to the **Normal** physical synthesis effort level, using the **Fast** setting can cause a smaller increase in compilation time. However, the lower level of optimization can result in a smaller increase in design performance.

**Preserving Placement with Incremental Compilation**

Preserving information about previous placements can make future placements faster. The incremental compilation feature provides an easy-to-use method for preserving placement results. For more information, refer to “Using Incremental Compilation” on page 11–4.
Reducing Routing Time

The time required to route a design depends on three factors: the device architecture, the placement of your design in the device, and the connectivity between different parts of your design. The routing time is usually not a significant amount of the compilation time. If your design requires a long time to route, perform one or more of the following actions:

- Check for routing congestion.
- Turn off Fitter Aggressive Routability Optimization.
- Use incremental compilation to preserve routing information for parts of your design.

Identifying Routing Congestion in the Chip Planner

To identify areas of routing congestion in your design, open the Chip Planner from the Tools menu. To view the routing congestion in the Chip Planner, double-click the Report Routing Utilization command in the Tasks list. Click Preview in the Report Routing Utilization dialog box to preview the default congestion display. Change the Routing utilization type to display congestion for specific resources. The default display uses dark blue for 0% congestion and red for 100%. Adjust the slider for Threshold percentage to change the congestion threshold level.

Even if average congestion is not very high, your design may have areas where congestion is very high in a specific type of routing. You can use the Chip Planner to identify areas of high congestion for specific interconnect types. You can change the connections in your design to reduce routing congestion. If the area with routing congestion is in a LogicLock region or between LogicLock regions, change or remove the LogicLock regions and recompile your design. If the routing time remains the same, the time is a characteristic of your design and the placement. If the routing time decreases, consider changing the size, location, or contents of LogicLock regions to reduce congestion and decrease routing time.

Sometimes, routing congestion may be a result of the HDL coding style used in your design. After you identify congested areas using the Chip Planner, review the HDL code for the blocks placed in those areas to determine whether you can reduce interconnect usage by code changes.

The Quartus II compilation messages contain information about average and peak interconnect usage. Peak interconnect usage over 75%, or average interconnect usage over 60%, could be an indication that it might be difficult to fit your design. Similarly, peak interconnect usage over 90%, or average interconnect usage over 75%, are likely to have increased chances of not getting a valid fit.

For more information about identifying areas of congested routing using the Chip Planner, refer to the “Viewing Routing Congestion” subsection in the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

Preserving Routing with Incremental Compilation

Preserving the previous routing results for part of your design can reduce future routing time. Incremental compilation provides an easy-to-use methodology that preserves placement and routing results. For more information, refer to “Using Incremental Compilation” on page 11–4 and the references listed in the section.
Reducing Static Timing Analysis Time

If you are performing timing-driven synthesis, the Quartus II software runs the TimeQuest analyzer during Analysis and Synthesis. The Quartus II Fitter also runs the TimeQuest analyzer during placement and routing. If there are incorrect constraints in the Synopsys Design Constraints File (.sdc), the Quartus II software may spend unnecessary time processing constraints several times.

- If you do not specify false paths and multicycle paths in your design, the TimeQuest analyzer may analyze paths that are not relevant to your design.
- If you redefine constraints in the .sdc files, the TimeQuest analyzer may spend additional time processing them. To avoid this situation, look for indications that Synopsis design constraints are being redefined in the compilation messages, and update the .sdc file.
- Ensure that you provide the correct timing constraints to your design, because the software cannot assume design intent, such as which paths to consider as false paths or multicycle paths. When you specify these assignments correctly, the TimeQuest analyzer skips analysis for those paths, and the Fitter does not spend additional time optimizing those paths.

Setting Process Priority

It might be necessary to reduce the computing resources allocated to the compilation at the expense of increased compilation time. It can be convenient to reduce the resource allocation to the compilation with single processor machines if you must run other tasks at the same time.

For more information about setting process priority, refer to Processing Page (Options Dialog Box) in Quartus II Help.

Document Revision History

Table 11–1 shows the revision history for this chapter.

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<td></td>
<td></td>
<td>• Updated “Placement Effort Multiplier Settings” section.</td>
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<td>• Updated “Identifying Routing Congestion in the Chip Planner” section.</td>
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<td>• Updated “Using Parallel Compilation with Multiple Processors”.</td>
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Table 11–1. Document Revision History (Part 2 of 2)

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<td>• Template update.</td>
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<td>• Added details about peak and average interconnect usage.</td>
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<td>• Added new section “Reducing Static Timing Analysis Time”.</td>
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<td></td>
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<td>• Minor changes throughout chapter.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter describes techniques to improve timing performance when designing for Altera® devices.

The application techniques vary between designs. Applying each technique does not always improve results. Settings and options in the Quartus® II software have default values that provide the best trade-off between compilation time, resource utilization, and timing performance. You can adjust these settings to determine whether other settings provide better results for your design.

Initial Compilation: Optional Fitter Settings

The Fitter offers many optional settings; however, this section discusses important optional timing-optimization related Fitter settings only, which are the Optimize Hold Timing, Optimize Multi-Corner Timing, and Fitter Aggressive Routability Optimization settings.

For scripting and device family support information of the Optimize Hold Timing and Optimize Multi-Corner Timing settings, refer to the Fitter Settings Page (Settings Dialog Box) in Quartus II Help.

The settings required to optimize different designs could be different. The group of settings that work best for one design may not produce the best result for another design.

Optimize Hold Timing

The Optimize Hold Timing option directs the Quartus II software to optimize minimum delay timing constraints. By default, the Quartus II software optimizes hold timing for all paths for designs using devices newer than Arria® GX, Stratix® III, and Cyclone® III. By default, the Quartus II software optimizes hold timing only for I/O paths and minimum tPD paths for older devices.

When you turn on Optimize Hold Timing, the Quartus II software adds delay to paths to ensure that your design meets the minimum delay requirements. In the Fitter Settings pane, if you select I/O Paths and Minimum TPD Paths (the default choice for older devices such as Cyclone II and Stratix II devices if you turn on Optimize Hold Timing), the Fitter works to meet the following criteria:

- Hold times (tH) from the device input pins to the registers
- Minimum delays from I/O pins to I/O registers or from I/O registers to I/O pins
- Minimum clock-to-out time (tCO) from registers to output pins
If you select **All Paths**, the Fitter also works to meet hold requirements from registers to registers, shown in Figure 12–1, in which a derived clock generated with logic causes a hold time problem on another register. However, if your design has internal hold time violations between registers, correct the problems by making changes to your design, such as using a clock enable signal instead of a derived or gated clock.

**Figure 12–1. Optimize Hold Timing Option Fixing an Internal Hold Time Violation**

![Diagram showing how to fix an internal hold time violation](image)

For design practices that help eliminate internal hold time violations, refer to the *Recommended Design Practices* chapter of the *Quartus II Handbook*.

**Optimize Multi-Corner Timing**

Due to process variation and changes in operating conditions, delays on some paths can be significantly smaller than those in the slow corner timing model. This can result in hold time violations on those paths, and in rare cases, additional setup time violations.

Also, because of the small process geometries of the Arria GX, Cyclone III, Stratix III, and newer device families, the slowest circuit performance of designs targeting these devices does not necessarily occur at the highest operating temperature. The temperature at which the circuit is slowest depends on the selected device, the design, and the compilation results. Therefore, the Quartus II software provides the Arria GX, Cyclone III series, Stratix III, and newer device families with three different timing corners—Slow 85°C corner, Slow 0°C corner, and Fast 0°C corner. For other device families, two timing corners are available—Fast 0°C and Slow 85°C corner.

By default, the Fitter optimizes constraints with only the slow corner timing model. You can turn on the **Optimize multi-corner timing** option to instruct the Fitter to also optimize constraints considering all available timing corners, at the cost of a slight increase in runtime. By optimizing for all timing corners, you can create a design implementation that is more robust across process, temperature, and voltage variations. While optimizing for multi-corner timing, the Fitter chooses one of the two slow corners with more critical timing (depending on the chosen device), along with the fast corner.

The **Optimize multi-corner timing** option helps to create a design implementation that is more robust across process, temperature, and voltage variations. Turning on this option increases compilation time by approximately 10%.
Fitter Aggressive Routability Optimization

If you set the Fitter Aggressive Routability Optimizations logic option to Always, reducing wire utilization may affect the performance of your design. If improving timing is more important than reducing wire usage, then set this option to Automatically or Never.

Design Analysis

The initial compilation establishes whether the design achieves a successful fit and meets the specified timing requirements. This section describes how to analyze your design results in the Quartus II software.

Ignored Timing Constraints

The Quartus II software ignores illegal, obsolete, and conflicting constraints. You can view a list of ignored constraints by clicking Report Ignored Constraints in the Reports menu in the TimeQuest GUI or by typing the following command to generate a list of ignored timing constraints:

```
report_sdc -ignored -panel_name "Ignored Constraints"
```

You must analyze any constraints that the Quartus II software ignores. If necessary, correct the constraints and recompile your design before proceeding with design optimization.

For more information about the report_sdc command and its options, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.

I/O Timing (Including tPD)

TimeQuest analyzer supports the Synopsys Design Constraints (SDC) format for constraining your design. When using the TimeQuest analyzer for timing analysis, use the set_input_delay constraint to specify the data arrival time at an input port with respect to a given clock. For output ports, use the set_output_delay command to specify the data arrival time at an output port’s receiver with respect to a given clock. You can use the report_timing Tcl command to generate the I/O timing reports.

The I/O paths that do not meet the required timing performance are reported as having negative slack and are highlighted in red in the TimeQuest analyzer Report pane. In cases where you do not apply an explicit I/O timing constraint to an I/O pin, the Quartus II timing analysis software still reports the Actual number, which is the timing number that must be met for that timing parameter when the device runs in your system.

For more information about how timing numbers are calculated, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.
Register-to-Register Timing

This section contains the following sections:

- “Timing Analysis with the TimeQuest Timing Analyzer”
- “Tips for Analyzing Failing Paths” on page 12–5
- “Tips for Analyzing Failing Clock Paths that Cross Clock Domains” on page 12–5
- “Tips for Analyzing Paths from/to the Source and Destination of Critical Path” on page 12–6
- “Tips for Locating Multiple Paths to the Chip Planner” on page 12–8
- “Tips for Creating a .tcl Script to Monitor Critical Paths Across Compiles” on page 12–8
- “Global Routing Resources” on page 12–8

Timing Analysis with the TimeQuest Timing Analyzer

Analyze all valid register-to-register paths by using the appropriate constraints in the TimeQuest analyzer. In the TimeQuest analyzer, run the macro Report All Summaries to view all timing summaries. If any clock domains have failing paths, right-click the domain and go to Report Timing to get more details. Your design meets timing requirements when you do not have negative slack on any register-to-register path on any of the clock domains.

When you select a path listed in the TimeQuest Report Timing pane, the tabs in the corresponding path detail pane show a path summary of source and destination registers and their timing, statistics about the path delay, detailed information about the complete data path with all nodes in the path and the waveforms of the relevant signals. To locate a selected path in the Chip Planner or the Technology Map Viewer by using the shortcut menu, right-click a path, point to Locate, and click Locate in Chip Planner. The Chip Planner appears with the path highlighted. Similarly, if you know that a path is not a valid path, you can set it to be a false path using the shortcut menu.

To view the path details of any selected path, click the Data Path tab in the path details pane. The Data Path tab displays the details of the Data Arrival Path, as well as the Data Required Path. For a graphical view of the information, click the Waveform tab.

You can locate critical paths in the Chip Planner from the TimeQuest timing analysis report panel.

For more information about how timing analysis results are calculated, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.

You also can see the logic in a particular path by locating the logic in the RTL Viewer or Technology Map Viewer. These viewers allow you to view a gate-level or technology-mapped representation of your design netlist. To locate a timing path in one of the viewers, right-click a path in the report, point to Locate, and click Locate in RTL Viewer or Locate in Technology Map Viewer. When you locate a timing path in the Technology Map Viewer, the annotated schematic displays the same delay information that is shown when you use the List Paths command.
For more information about netlist viewers, refer to the *Analyzing Designs with Quartus II Netlist Viewers* chapter of the *Quartus II Handbook*.

**Tips for Analyzing Failing Paths**

When you are analyzing clock path failures, examine the reports and waveforms to determine if the correct constraints are being applied, and add multicycle or false paths as appropriate.

Focus on improving the paths that show the worst slack. The Fitter works hardest on paths with the worst slack. If you fix these paths, the Fitter might be able to improve the other failing timing paths in the design.

Check for particular nodes that appear in many failing paths. Look for paths that have common source registers, destination registers, or common intermediate combinational nodes. In some cases, the registers might not be identical, but are part of the same bus.

In the timing analysis report panels, clicking on the *From* or *To* column headings can help to sort the paths by the source or destination registers. Clicking first on *From*, then on *To*, uses the registers in the *To* column as the primary sort and the registers in the *From* column as the secondary sort. If you see common nodes, these nodes indicate areas of your design that might be improved through source code changes or *Quartus II* optimization settings. Constraining the placement for just one of the paths might decrease the timing performance for other paths by moving the common node further away in the device.

**Tips for Analyzing Failing Clock Paths that Cross Clock Domains**

When analyzing clock path failures, check whether these paths cross two clock domains. This is the case if the *From Clock* and *To Clock* in the timing analysis report are different. There can also be paths that involve a different clock in the middle of the path, even if the source and destination register clock are the same.

When you run Report Timing on your design, the report shows the launch clock and latch clock for each failing path. Check whether these failing paths between these clock domains should be analyzed synchronously. If the failing paths are not to be analyzed synchronously, they must be set as false paths. Also check the relationship between the launch clock and latch clock to make sure it is realistic and what you expect from your knowledge of the design. For example, the path can start at a rising edge and end at a falling edge, which reduces the setup relationship by one half clock cycle.

Review the clock skew reported in the Timing Report. A large skew may indicate a problem in your design, such as a gated clock or a problem in the physical layout (for example, a clock using local routing instead of dedicated clock routing). When you have made sure the paths are analyzed synchronously and that there is no large skew on the path, and that the constraints are correct, you can analyze the data path. These steps help you fine tune your constraints for paths across clock domains to ensure you get an accurate timing report.

Check if the PLL phase shift is reducing the setup requirement. You might be able to adjust this using PLL parameters and settings.
Paths that cross clock domains are generally protected with synchronization logic (for example, FIFOs or double-data synchronization registers) to allow asynchronous interaction between the two clock domains. In such cases, you can ignore the timing paths between registers in the two clock domains while running timing analysis, even if the clocks are related.

The Fitter attempts to optimize all failing timing paths. If there are paths that can be ignored for optimization and timing analysis, but the paths do not have constraints that instruct the Fitter to ignore them, the Fitter tries to optimize those paths as well. In some cases, optimizing unnecessary paths can prevent the Fitter from meeting the timing requirements on timing paths that are critical to the design. It is beneficial to specify all paths that can be ignored, so that the Fitter can put more effort into the paths that must meet their timing requirements instead of optimizing paths that can be ignored.

For more details about how to ignore timing paths that cross clock domains, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.

Evaluate the clock skew between the source clock and the destination clock to determine if that is reducing the available setup time. You can check the shortest and longest clock path reports to view what is causing the clock skew. Avoid using combinational logic in clock paths because it contributes to clock skew. Differences in the logic or in its routing between the source and destination can cause clock skew problems and result in warnings during compilation.

**Tips for Analyzing Paths from/to the Source and Destination of Critical Path**

In the project directory, run the Tcl command shown in Example 12–1 in a .tcl file to analyze the nodes in a critical path.

Example 12–1. report_timing Command

```tcl
set wrst_src <insert_source_of_worst_path_here>
set wrst_dst <insert_destination_of_worst_path_here>
report_timing -setup -npaths 50 -detail path_only -from $wrst_src -panel_name "Worst Path||wrst_src -> *"
report_timing -setup -npaths 50 -detail path_only -to $wrst_dst -panel_name "Worst Path||* -> wrst_dst"
report_timing -setup -npaths 50 -detail path_only -to $wrst_src -panel_name "Worst Path||* -> wrst_src"
report_timing -setup -npaths 50 -detail path_only -from $wrst_dst -panel_name "Worst Path||wrst_dst -> **"
```

Copy the node name from the From Node and To Node columns of the worst path into the first two variables, and then in the TimeQuest timing analyzer, in the Script menu, source the .tcl script.
Figure 12–2 shows a simplified example of what these reports analyzed.

The critical path of the design is in red. The script analyzes the path between the worst source and destination registers. The first `report_timing` command analyzes other path that the source is driving, as shown in green. The second `report_timing` command analyzes the critical path and other path going to the destination, shown in yellow. These commands report everything inside these two endpoints that are pulling them in different directions. The last two `report_timing` commands show everything outside of the endpoints pulling them in other directions. If any of these reports have slacks near the critical path, then the Fitter is balancing these paths with the critical path, trying to achieve the best slack. Figure 12–2 is quite simple compared to the critical path in most designs, but it is easy to see how this can get very complicated quickly.

These timing reports are useful for analyzing what is competing with the critical path, but not always good for examining how they might pull in different directions.
Tips for Locating Multiple Paths to the Chip Planner

The Chip Planner shows multiple paths in a timing report.

1. Run `report_timing` to show multiple paths. (Example 12–1)
2. Select multiple rows of timing report.
   
   Although you can select multiple rows, you can only select a single column.

3. Right-click, select **Locate Path**, and then click **Chip Planner**.
4. The **Locate History** window in the Chip Planner displays the selected paths and the worst path.
5. Double-click **Locate Paths** to show all paths at once, or select individual paths to view the path in the Chip Planner.

You can now view what is pulling on the critical path.

Tips for Creating a .tcl Script to Monitor Critical Paths Across Compiles

Many designs have the same critical paths show up after each compile, but some suffer from having critical paths bounce around between different hierarchies, changing with each compile.

In designs like this, create a `TQ_critical_paths.tcl` script in the project directory. For a given compile, view the critical paths and then write a generic `report_timing` command to capture those paths. For example, if several paths fail in a low-level hierarchy, you can add the following command as shown in Example 12–2:

Example 12–2. `report_timing` Command

```
report_timing -setup -npaths 50 -detail path_only -to "main_system:
main_system_inst|app_cpu:cpu|*" 
-panel_name "Critical Paths||s: * -> app_cpu"
```

If there is a specific path, such as a bit of a state-machine going to other `*count_sync*` registers, you can add a command as shown in Example 12–3:

Example 12–3. `report_timing` Command

```
report_timing -setup -npaths 50 -detail path_only -from "main_system:
main_system_inst|egress_count_sm:egress_inst|update" -to "*count_sync*" 
-panel_name "Critical Paths||s: egress_sm|update -> count_sync"
```

This tip helps you monitor paths that consistently fail and paths that are only marginal, so you can prioritize effectively.

Global Routing Resources

Global routing resources are designed to distribute high fan-out, low-skew signals (such as clocks) without consuming regular routing resources. Depending on the device, these resources can span the entire chip, or some smaller portion, such as a quadrant. The Quartus II software attempts to assign signals to global routing resources automatically, but you might be able to make more suitable assignments manually.
For details about the number and types of global routing resources available, refer to the relevant device handbook.

Check the global signal utilization in your design to ensure that the appropriate signals have been placed on the global routing resources. In the Compilation Report, open the Fitter report and click Resource Section. Analyze the Global & Other Fast Signals and Non-Global High Fan-out Signals reports to determine whether any changes are required.

You might be able to reduce skew for high fan-out signals by placing them on global routing resources. Conversely, you can reduce the insertion delay of low fan-out signals by removing them from global routing resources. Doing so can improve clock enable timing and control signal recovery/removal timing, but increases clock skew. Use the Global Signal setting in the Assignment Editor to control global routing resources.

**Optimizing Timing (LUT-Based Devices)**

This section contains guidelines that might help you if your design does not meet its timing requirements.

**Debugging Timing Failures in the TimeQuest Analyzer**

A Report Timing Closure Recommendations task is available in the Custom Reports section of the Tasks pane of the TimeQuest analyzer. Use this report to get more information and help on the failing paths in your design. This feature is available for the Arria II, Cyclone III, Stratix III, and newer devices.

When you run the Report Timing Closure Recommendations task, you get specific recommendations about failing paths in your design and changes that you can make to potentially fix the failing paths.


From the Report Timing Closure Recommendations dialog box, you can select paths based on the clock domain, filter by nodes on path, and choose the number of paths to analyze.

After running the Report Timing Closure Recommendations task in the TimeQuest analyzer, examine the reports in the Report Timing Closure Recommendations folder in the Report pane of the TimeQuest analyzer GUI. Each recommendation has star symbols (*) associated with it. Recommendations with more stars are more likely to help you close timing on your design.

The reports give you the most probable causes of failure for each path being analyzed. The reports are organized into sections, depending on the type of issues found in the design, such as large clock skew, restricted optimizations, unbalanced logic, skipped optimizations, coding style that has too many levels of logic between registers, or region or partition constraints specific to your project.
You will see recommendations that may help you fix the failing paths. For detailed analysis of the critical paths, run the `report_timing` command on specified paths. In the Extra Fitter Information tab of the Path report panel, you will also see detailed Fitter-related information that may help you visualize the issue and take the appropriate action if your constraints cause a specific placement.


**Timing Optimization Advisor**

The Timing Optimization Advisor guides you in making settings that optimize your design to meet your timing requirements. To run the Timing Optimization Advisor, on the Tools menu, point to Advisors and click Timing Optimization Advisor. This advisor describes many of the suggestions made in this section.

When you open the Timing Optimization Advisor after compilation, you can find recommendations to improve the timing performance of your design. Some of the recommendations in these advisors can contradict each other. Altera recommends evaluating these options and choosing the settings that best suit the given requirements.

Figure 12–3 shows the Timing Optimization Advisor after compiling a design that meets its frequency requirements, but requires setting changes to improve the timing.

When you expand one of the categories in the Timing Optimization Advisor, such as Maximum Frequency (fmax) or I/O Timing (tsu, tco, tpd), the recommendations are divided into stages. The stages show the order in which to apply the recommended settings. The first stage contains the options that are easiest to change, make the least drastic changes to your design optimization, and have the least effect on compilation.
time. Icons indicate whether each recommended setting has been made in the current project. In Figure 12–3, the checkmark icons in the list of recommendations for Stage 1 indicate recommendations that are already implemented. The warning icons indicate recommendations that are not followed for this compilation. The information icons indicate general suggestions. For these entries, the advisor does not report whether these recommendations were followed, but instead explains how you can achieve better performance. For a legend that provides more information for each icon, refer to the “How to use” page in the Timing Optimization Advisor.

There is a link from each recommendation to the appropriate location in the Quartus II GUI where you can change the settings. For example, consider the Synthesis Netlist Optimizations page of the Settings dialog box or the Global Signals category in the Assignment Editor. This approach provides the most control over which settings are made and helps you learn about the settings in the software. In some cases, you can also use the Correct the Settings button to automatically make the suggested change to global settings.

For some entries in the Timing Optimization Advisor, a button appears that allows you to further analyze your design and gives you more information. The advisor provides a table with the clocks in the design and indicates whether they have been assigned a timing constraint.

I/O Timing Optimization

This stage of design optimization focuses on I/O timing. Ensure that you have made the appropriate assignments described in the “Initial Compilation: Required Settings” section in the Design Optimization Overview chapter of the Quartus II Handbook. You must also ensure that resource utilization is satisfactory before proceeding with I/O timing optimization. The suggestions provided in this section are applicable to all Altera FPGA families and to the MAX II family of CPLDs.

Because changes to the I/O paths affect the internal register-to-register timing, complete this stage before proceeding to the register-to-register timing optimization stage as described in “Register-to-Register Timing Optimization Techniques (LUT-Based Devices)” on page 12–16.

The options presented in this section address how to improve I/O timing, including the setup delay (tSU), hold time (tH), and clock-to-output (tCO) parameters.

Improving Setup and Clock-to-Output Times Summary

Table 12–1 lists the recommended order in which to use techniques to reduce tSU and tCO times. Checkmarks indicate which timing parameters are affected by each technique. Reducing tSU times increases hold (tH) times.

<table>
<thead>
<tr>
<th>Technique</th>
<th>Affects tSU</th>
<th>Affects tCO</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ensure that the appropriate constraints are set for the failing I/Os</td>
<td>✔</td>
<td>✔</td>
</tr>
<tr>
<td>(refer to the “Initial Compilation: Required Settings” section in the</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Quartus II Handbook.)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Use timing-driven compilation for I/O (page 12–12)</td>
<td>✔</td>
<td>✔</td>
</tr>
<tr>
<td>Use fast input register (page 12–13)</td>
<td>✔</td>
<td>—</td>
</tr>
<tr>
<td>Use fast output register, fast output enable register, and fast OCT</td>
<td>—</td>
<td>✔</td>
</tr>
<tr>
<td>register (page 12–13)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Timing-Driven Compilation

This option moves registers into I/O elements if required to meet $t_{SU}$ or $t_{CO}$ assignments, duplicating the register if necessary (as in the case in which a register fans out to multiple output locations). This option is turned on by default and is a global setting. The option does not apply to MAX II series devices because they do not contain I/O registers.

The Optimize IOC Register Placement for Timing option affects only pins that have a $t_{SU}$ or $t_{CO}$ requirement. Using the I/O register is possible only if the register directly feeds a pin or is fed directly by a pin. This setting does not affect registers with any of the following characteristics:

- Have combinational logic between the register and the pin
- Are part of a carry or cascade chain
- Have an overriding location assignment
- Use the asynchronous load port and the value is not 1 (in device families where the port is available)

Registers with the characteristics listed are optimized using the regular Quartus II Fitter optimizations.

For more information, refer to Optimize IOC Register Placement for Timing logic option in Quartus II Help.

Fast Input, Output, and Output Enable Registers

Normally, with correct timing assignments, the Fitter already places the I/O registers in the correct I/O cell or in the core, to meet the performance requirement. However, you can place individual registers in I/O cells manually by making fast I/O assignments with the Assignment Editor.
For more information about the **Fast Input Register** option, **Fast Output Register** option, **Fast Output Enable Register** option, and **Fast OCT (on-chip termination) Register** option, refer to Quartus II Help.

In MAX II series devices, which have no I/O registers, these assignments lock the register into the LAB adjacent to the I/O pin if there is a pin location assignment for that I/O pin.

If the fast I/O setting is on, the register is always placed in the I/O element. If the fast I/O setting is off, the register is never placed in the I/O element. This is true even if the **Optimize IOC Register Placement for Timing** option is turned on. If there is no fast I/O assignment, the Quartus II software determines whether to place registers in I/O elements if the **Optimize IOC Register Placement for Timing** option is turned on.

You can also use the four fast I/O options (**Fast Input Register**, **Fast Output Register**, **Fast Output Enable Register**, and **Fast OCT Register**) to override the location of a register that is in a LogicLock region and force it into an I/O cell. If you apply this assignment to a register that feeds multiple pins, the register is duplicated and placed in all relevant I/O elements. In MAX II series devices, the register is duplicated and placed in each distinct LAB location that is next to an I/O pin with a pin location assignment.

**Programmable Delays**

You can use various programmable delay options to minimize the $t_{SU}$ and $t_{CO}$ times. For Arria, Cyclone, MAX II, MAX V, and Stratix series devices, the Quartus II software automatically adjusts the applicable programmable delays to help meet timing requirements. Programmable delays are advanced options to use only after you compile a project, check the I/O timing, and determine that the timing is unsatisfactory. For detailed information about the effect of these options, refer to the device family handbook or data sheet.

After you have made a programmable delay assignment and compiled the design, you can view the implemented delay values for every delay chain for every I/O pin in the **Delay Chain Summary** section of the Compilation Report.

You can assign programmable delay options to supported nodes with the Assignment Editor. You can also view and modify the delay chain setting for the target device with the Chip Planner and Resource Property Editor. When you use the Resource Property Editor to make changes after performing a full compilation, recompiling the entire design is not necessary; you can save changes directly to the netlist. Because these changes are made directly to the netlist, the changes are not made again automatically when you recompile the design. The change management features allow you to reapply the changes on subsequent compilations.

Although the programmable delays in newer devices are user-controllable, Altera recommends their use for advanced users only. However, the Quartus II software might use the programmable delays internally during the Fitter phase.

For more information about Stratix III programmable delays, refer to the *Stratix III Device Handbook* and *AN 474: Implementing Stratix III Programmable I/O Delay Settings in the Quartus II Software*. For more information about using the Chip Planner and Resource Property Editor, refer to the *Engineering Change Management with the Chip Planner* chapter of the *Quartus II Handbook*. 
For details about the programmable delay logic options available for Altera devices, refer to the following Quartus II Help topics:

- Decrease Input Delay to Input Register logic option
- Input Delay from Pin to Input Register logic option
- Decrease Input Delay to Internal Cells logic option
- Input Delay from Pin to Internal Cells logic option
- Decrease Input Delay to Output Register logic option
- Increase Delay to Output Enable Pin logic option
- Output Enable Pin Delay logic option
- Increase Delay to Output Pin logic option
- Delay from Output Register to Output Pin logic option
- Increase Input Clock Enable Delay logic option
- Input Delay from Dual-Purpose Clock Pin to Fan-Out Destinations logic option
- Increase Output Clock Enable Delay logic option
- Increase Output Enable Clock Enable Delay logic option
- Increase \( t_{<x} \) Delay to Output Pin logic option

**Use PLLs to Shift Clock Edges**

Using a PLL typically improves I/O timing automatically. If the timing requirements are still not met, most devices allow the PLL output to be phase shifted to change the I/O timing. Shifting the clock backwards gives a better \( t_{\text{H}} \) at the expense of \( t_{\text{SU}} \), while shifting it forward gives a better \( t_{\text{SU}} \) at the expense of \( t_{\text{H}} \) (refer to Figure 12–4). You can use this technique only in devices that offer PLLs with the phase shift option.

**Figure 12–4. Shift Clock Edges Forward to Improve \( t_{\text{SU}} \) at the Expense of \( t_{\text{H}} \)**

You can achieve the same type of effect in certain devices by using the programmable delay called *Input Delay from Dual Purpose Clock Pin to Fan-Out Destinations*.

For more information, refer to *Input Delay from Dual-Purpose Clock Pin to Fan-Out Destinations logic option* in Quartus II Help.

**Use Fast Regional Clock Networks and Regional Clocks Networks**

Altera devices have a variety of hierarchical clock structures. These include dedicated global clock networks, regional clock networks, fast regional clock networks, and periphery clock networks. The available resources differ between the various Altera device families.
For the number of clocking resources available in your target device, refer to the appropriate device handbook.

In general, fast regional clocks have less delay to I/O elements than regional and global clocks, and are used for high fan-out control signals. Regional clocks provide the lowest clock delay and skew for logic contained in a single quadrant. Placing clocks on these low-skew and low-delay clock nets provides better t\textsubscript{CO} performance.

**Spine Clock Limitations**

Global clock networks, regional clock networks, and periphery clock networks have an additional level of clock hierarchy known as spine clocks. Spine clocks drive the final row and column clocks to their registers; thus, the clock to every register in the chip is reached through spine clocks. Spine clocks are not directly user controllable.

If your project has high clock routing demands, due to limitations in the Quartus II software, you may see spine clock errors. These errors are often seen with designs using multiple memory interfaces and high-speed serial interface (HSSI) channels (especially PMA Direct mode).

To reduce these spine clock errors, you can constrain your design to better use your regional clock resources using the following techniques:

- If your design does not use LogicLock regions, or if the LogicLock regions are not aligned to your clock region boundaries, create additional LogicLock regions and further constrain your logic.

  Register packing, a Fitter optimization option, may ignore LogicLock regions. If this occurs, disable register packing for specific instances through the Quartus II Assignment Editor.

- Some periphery features may ignore LogicLock region assignments. When this happens, the global promotion process may not function properly. To ensure that the global promotion process uses the correct locations, assign specific pins to the I/Os using these periphery features.

- By default, some IP MegaCore functions apply a global signal assignment with a value of dual-regional clock. If you constrain your logic to a regional clock region and set the global signal assignment to **Regional** instead of **Dual-Regional**, you can reduce clock resource contention.

**Change How Hold Times are Optimized for MAX II Devices**

For MAX II devices, you can use the **Guarantee I/O Paths Have Zero Hold Time at Fast Corner** option to control how hold time is optimized by the Quartus II software.

For details, refer to **Guarantee I/O Paths Have Zero Hold Time at Fast Corner logic option** in Quartus II Help.
Register-to-Register Timing Optimization Techniques (LUT-Based Devices)

The next stage of design optimization is to improve register-to-register (f_{MAX}) timing. The following sections provide available options if the performance requirements are not achieved after compilation.

Coding style affects the performance of your design to a greater extent than other changes in settings. Always evaluate your code and make sure to use synchronous design practices.

For more details about synchronous design practices and coding styles, refer to the Recommended Design Practices chapter of the Quartus II Handbook.

When using the TimeQuest analyzer, register-to-register timing optimization is the same as maximizing the slack on the clock domains in your design. You can use the techniques described in this section to improve the slack on different timing paths in your design.

Before optimizing your design, understand the structure of your design as well as the type of logic affected by each optimization. An optimization can decrease performance if the optimization does not benefit your logic structure.

Optimize Source Code

In many cases, optimizing the design’s source code can have a very significant effect on your design performance. In fact, optimizing your source code is typically the most effective technique for improving the quality of your results and is often a better choice than using LogicLock or location assignments.

Be aware of the number of logic levels needed to implement your logic while you are coding. Too many levels of logic between registers could result in critical paths failing timing. Try restructuring the design to use pipelining or more efficient coding techniques. Also, try limiting high fan-out signals in the source code. When possible, duplicate and pipeline control signals. Make sure the duplicate registers are protected by a preserve attribute, to avoid merging during synthesis.

If the critical path in your design involves memory or DSP functions, check whether you have code blocks in your design that describe memory or functions that are not being inferred and placed in dedicated logic. You might be able to modify your source code to cause these functions to be placed into high-performance dedicated memory or resources in the target device. When using RAM/DSP blocks, enable the optional input and output registers.

Ensure that your state machines are recognized as state machine logic and optimized appropriately in your synthesis tool. State machines that are recognized are generally optimized better than if the synthesis tool treats them as generic logic. In the Quartus II software, you can check the State Machine report under Analysis & Synthesis in the Compilation Report. This report provides details, including state encoding for each state machine that was recognized during compilation. If your state machine is not recognized, you might have to change your source code to enable it to be recognized.
For coding style guidelines including examples of HDL code for inferring memory, functions, guidelines, and sample HDL code for state machines, refer to the Recommended HDL Coding Styles chapter of the Quartus II Handbook.

For additional HDL coding examples, refer to AN 584: Timing Closure Methodology for Advanced FPGA Designs.

Improving Register-to-Register Timing Summary

The choice of options and settings to improve the timing margin (slack) or to improve register-to-register timing depends on the failing paths in the design. To achieve the results that best approximate your performance requirements, apply the following techniques and compile the design after each step:

1. Ensure that your timing assignments are complete and correct. For details, refer to the “Initial Compilation: Required Settings” section in the Design Optimization Overview chapter of the Quartus II Handbook.

2. Ensure that you have reviewed all warning messages from your initial compilation and check for ignored timing assignments.

   For details and to fix any of these problems before proceeding with optimization, refer to the Design Optimization Overview chapter of the Quartus II Handbook.

3. Apply netlist synthesis optimization options.

4. To optimize for speed, apply the following synthesis options:
   - “Optimize Synthesis for Speed, Not Area” on page 12–19
   - “Flatten the Hierarchy During Synthesis” on page 12–20
   - “Set the Synthesis Effort to High” on page 12–20
   - “Change State Machine Encoding” on page 12–21
   - “Prevent Shift Register Inference” on page 12–22
   - “Use Other Synthesis Options Available in Your Synthesis Tool” on page 12–22

5. To optimize physical synthesis, apply the following options:
   - Perform physical synthesis for combinational logic
   - Perform automatic asynchronous signal pipelining
   - Perform register duplication
   - Perform register retiming
   - Perform logic to memory mapping

6. Try different Fitter seeds (page 12–22). If there are very few paths that are failing by small negative slack, then you can try with a different seed to see if there is a fit that meets constraints in the Fitter seed noise.

   Omit this step if a large number of critical paths are failing or if the paths are failing badly.
7. To control placement, make LogicLock assignments (page 12–23).

8. Make design source code modifications to fix areas of the design that are still failing timing requirements by significant amounts (page 12–16).

9. Make location assignments, or as a last resort, perform manual placement by back-annotating the design (page 12–25).

You can use the Design Space Explorer (DSE) to automate the process of running several different compilations with different settings.

For more information, refer to About Design Space Explorer in Quartus II Help.

If these techniques do not achieve performance requirements, additional design source code modifications might be required (page 12–16).

Physical Synthesis Optimizations

The Quartus II software offers physical synthesis optimizations that can help improve the performance of many designs regardless of the synthesis tool used. Physical synthesis optimizations can be applied both during synthesis and during fitting.

Physical synthesis optimizations that occur during the synthesis stage of the Quartus II compilation operate either on the output from another EDA synthesis tool or as an intermediate step in Quartus II integrated synthesis. These optimizations make changes to the synthesis netlist to improve either area or speed, depending on your selected optimization technique and effort level.

To view and modify the synthesis netlist optimization options, on the Assignments menu, click Settings. In the Category list, expand Compilation Process Settings and select Physical Synthesis Optimizations.

If you use a third-party EDA synthesis tool and want to determine if the Quartus II software can remap the circuit to improve performance, you can use the Perform WYSIWYG Primitive Resynthesis option. This option directs the Quartus II software to unmap the LEs in an atom netlist to logic gates and then map the gates back to Altera-specific primitives. Using Altera-specific primitives enables the Fitter to remap the circuits using architecture-specific techniques.

For more information, refer to Perform WYSIWYG Primitive Resynthesis logic option in Quartus II Help.

The Quartus II technology mapper optimizes the design to achieve maximum speed performance, minimum area usage, or balances high performance and minimal logic usage, according to the setting of the Optimization Technique option. Set this option to Speed or Balanced.

For more information, refer to Optimization Technique logic option in Quartus II Help.

The physical synthesis optimizations occur during the Fitter stage of the Quartus II compilation. Physical synthesis optimizations make placement-specific changes to the netlist that improve speed performance results for a specific Altera device.
The following physical synthesis optimizations are available during the Fitter stage for improving performance:

- Physical synthesis for combinational logic
- Automatic asynchronous signal pipelining
- Physical synthesis for registers
  - Register duplication
  - Register retiming

If you want the performance gain from physical synthesis only on parts of your design, you can apply the physical synthesis options on specific instances.

For more information, refer to *Physical Synthesis Optimizations Page (Settings Dialog Box)* in Quartus II Help.

To apply physical synthesis assignments for fitting on a per-instance basis, use the Quartus II Assignment Editor. The following assignments are available as instance assignments:

- Perform physical synthesis for combinational logic
- Perform register duplication for performance
- Perform register retiming for performance
- Perform automatic asynchronous signal pipelining

For information about making assignments, refer to *Working With Assignments in the Assignment Editor* in Quartus II Help.

**Turn Off Extra-Effort Power Optimization Settings**

If PowerPlay power optimization settings are set to Extra Effort, your design performance can be affected. If improving timing performance is more important than reducing power use, set the PowerPlay power optimization setting to Normal.

For more information, refer to *PowerPlay Power Optimization logic option* in Quartus II Help.

For more information about reducing power use, refer to the Power Optimization chapter of the Quartus II Handbook.

**Optimize Synthesis for Speed, Not Area**

The manner in which the design is synthesized has a large impact on design performance. Design performance varies depending on the way the design is coded, the synthesis tool used, and the options specified when synthesizing. Change your synthesis options if a large number of paths are failing or if specific paths are failing badly and have many levels of logic.

Set your device and timing constraints in your synthesis tool. Synthesis tools are timing-driven and optimized to meet specified timing requirements. If you do not specify a target frequency, some synthesis tools optimize for area.
Some synthesis tools offer an easy way to instruct the tool to focus on speed instead of area.

For more information, refer to *Optimization Technique logic option* in Quartus II Help.

You can also specify this logic option for specific modules in your design with the Assignment Editor while leaving the default *Optimization Technique* setting at *Balanced* (for the best trade-off between area and speed for certain device families) or *Area* (if area is an important concern). You can also use the *Speed Optimization Technique for Clock Domains* option in the Assignment Editor to specify that all combinational logic in or between the specified clock domain(s) is optimized for speed.

To achieve best performance with push-button compilation, follow the recommendations in the following sections for other synthesis settings. You can use the DSE to experiment with different Quartus II synthesis options to optimize your design for the best performance.

For information about setting timing requirements and synthesis options in Quartus II integrated synthesis and third-party synthesis tools, refer to the appropriate chapter in *Synthesis* of the *Quartus II Handbook*, or refer to your synthesis software documentation.

For more information about the Design Space Explorer, refer to *About Design Space Explorer* in Quartus II Help.

**Flatten the Hierarchy During Synthesis**

Synthesis tools typically let you preserve hierarchical boundaries, which can be useful for verification or other purposes. However, the best optimization results generally occur when the synthesis tool optimizes across hierarchical boundaries, because doing so often allows the synthesis tool to perform the most logic minimization, which can improve performance. Whenever possible, flatten your design hierarchy to achieve the best results. If you are using Quartus II incremental compilation, you cannot flatten your design across design partitions. Incremental compilation always preserves the hierarchical boundaries between design partitions. Follow Altera’s recommendations for design partitioning, such as registering partition boundaries to reduce the effect of cross-boundary optimizations.

For more information about using incremental compilation and recommendations for design partitioning, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter of the *Quartus II Handbook*.

**Set the Synthesis Effort to High**

Some synthesis tools offer varying synthesis effort levels to trade off compilation time with synthesis results. Set the synthesis effort to *high* to achieve best results when applicable.
Change State Machine Encoding

State machines can be encoded using various techniques. One-hot encoding, which uses one register for every state bit, usually provides the best performance. If your design contains state machines, changing the state machine encoding to one-hot can improve performance at the cost of area.

For more information, refer to State Machine Processing logic option in Quartus II Help.

Duplicate Logic for Fan-Out Control

Duplicating logic or registers can help improve timing in cases where moving a register in a failing timing path to reduce routing delay creates other failing paths or where there are timing problems due to the fan-out of the registers. Most often, timing failures occur not because of the high fan-out registers, but because of the location of those registers. Duplicating registers, where source and destination registers are physically close, can help improve slack on critical paths.

Many synthesis tools support options or attributes that specify the maximum fan-out of a register. When using Quartus II integrated synthesis, you can set the Maximum Fan-Out logic option in the Assignment Editor to control the number of destinations for a node so that the fan-out count does not exceed a specified value. You can also use the \texttt{maxfan} attribute in your HDL code. The software duplicates the node as required to achieve the specified maximum fan-out.

Logic duplication using Maximum Fan-Out assignments normally increases resource utilization and can potentially increase compilation time, depending on the placement and the total resource usage within the selected device. The improvement in timing performance that results because of Maximum Fan-Out assignments is very design-specific. This is because when you use the Maximum Fan-Out assignment, although the Fitter duplicates the source logic to limit the fan-out, it may not be able to control the destinations that each of the duplicated sources drive. Since the Maximum Fan-Out destination does not specify which of the destinations the duplicated source should drive, it is possible that it might still be driving logic located all around the device. To avoid this situation, you could use the Manual Logic Duplication logic option.

If you are using Maximum Fan-Out assignments, Altera recommends benchmarking your design with and without these assignments to evaluate whether they give the expected improvement in timing performance. Use the assignments only when you get improved results.

You can manually duplicate registers in the Quartus II software regardless of the synthesis tool used. To duplicate a register, apply the Manual Logic Duplication logic option to the register with the Assignment Editor.

Various Fitter optimizations may cause a small violation to the Maximum Fan-Out assignments to improve timing.

For more information, refer to Manual Logic Duplication logic option in Quartus II Help.


**Prevent Shift Register Inference**

In some cases, turning off the inference of shift registers increases performance. Doing so forces the software to use logic cells to implement the shift register instead of implementing the registers in memory blocks using the ALTSHIFT_TAPS megafunction. If you implement shift registers in logic cells instead of memory, logic utilization is increased.

**Use Other Synthesis Options Available in Your Synthesis Tool**

With your synthesis tool, experiment with the following options if they are available:

- Turn on register balancing or retiming
- Turn on register pipelining
- Turn off resource sharing

These options can increase performance, but typically increase the resource utilization of your design.

**Fitter Seed**

The Fitter seed affects the initial placement configuration of the design. Changing the seed value changes the Fitter results because the fitting results change whenever there is a change in the initial conditions. Each seed value results in a somewhat different fit, and you can experiment with several different seeds to attempt to obtain better fitting results and timing performance.

When there are changes in your design, there is some random variation in performance between compilations. This variation is inherent in placement and routing algorithms—there are too many possibilities to try them all and get the absolute best result, so the initial conditions change the compilation result.

Any design change that directly or indirectly affects the Fitter has the same type of random effect as changing the seed value. This includes any change in source files, **Analysis & Synthesis Settings**, **Fitter Settings**, or **Timing Analyzer Settings**. The same effect can appear if you use a different computer processor type or different operating system, because different systems can change the way floating point numbers are calculated in the Fitter.

If a change in optimization settings slightly affects the register-to-register timing or number of failing paths, you cannot always be certain that your change caused the improvement or degradation, or whether it could be due to random effects in the Fitter. If your design is still changing, running a seed sweep (compiling your design with multiple seeds) determines whether the average result has improved after an optimization change and whether a setting that increases compilation time has benefits worth the increased time (such as setting the **Physical Synthesis Effort** to **Extra**). The sweep also shows the amount of random variation to expect for your design.

If your design is finalized, you can compile your design with different seeds to obtain one optimal result. However, if you subsequently make any changes to your design, you might need to perform seed sweep again.

On the Assignments menu, select **Fitter Settings** to control the initial placement with the seed. You can use the DSE to perform a seed sweep easily.
You can use the following Tcl command from a script to specify a Fitter seed:

```tcl
set_global_assignment -name SEED <value>
```

For more information about compiling your design with different seeds using the Design Space Explorer (DSE seed sweep), refer to About Design Space Explorer in Quartus II Help.

### Set Maximum Router Timing Optimization Level

To improve routability in designs where the router did not pick up the optimal routing lines, set the **Router Timing Optimization Level** to **Maximum**. This setting determines how aggressively the router tries to meet the timing requirements. Setting this option to **Maximum** can increase design speed slightly at the cost of increased compilation time. Setting this option to **Minimum** can reduce compilation time at the cost of slightly reduced design speed. The default value is **Normal**.

For more information, refer to **Router Timing Optimization Level logic option** in Quartus II Help.

### LogicLock Assignments

Using LogicLock assignments to improve timing performance is only recommended for older Altera devices, such as the MAX II family. For other device families, especially for larger devices such as Arria and Stratix series devices, Altera does not recommend using LogicLock assignments to improve timing performance. For these devices, use the LogicLock feature for performance preservation and to floorplan your design.

LogicLock assignments do not always improve the performance of the design. In many cases, you cannot improve upon results from the Fitter by making location assignments. If there are existing LogicLock assignments in your design, remove the assignments if your design methodology permits it. Recompile the design, and then check if the assignments are making the performance worse.

When making LogicLock assignments, it is important to consider how much flexibility to give the Fitter. LogicLock assignments provide more flexibility than hard location assignments. Assignments that are more flexible require higher Fitter effort, but reduce the chance of design overconstraint. The following types of LogicLock assignments are available, listed in the order of decreasing flexibility:

- Auto size, floating location regions
- Fixed size, floating location regions
- Fixed size, locked location regions

For more information about using LogicLock regions, refer to the Analyzing and Optimizing the Design Floorplan with the Chip Planner chapter of the Quartus II Handbook.
If you are unsure of how big or where a LogicLock region should go, the **Auto/Floating** options are useful for your first pass. After you determine where a LogicLock region must go, modify the Fixed/Locked regions, as Auto/Floating LogicLock regions can hurt your overall performance. To determine what to put into a LogicLock region, refer to the timing analysis results and analyze the critical paths in the Chip Planner. The register-to-register timing paths in the Timing Analyzer section of the Compilation Report help you recognize patterns.

The following sections describe cases in which LogicLock regions can help to optimize a design.

**Hierarchy Assignments**

For a design with the hierarchy shown in Figure 12–5, which has failing paths in the timing analysis results similar to those shown in Table 12–2, mod_A is probably a problem module. In this case, a good strategy to fix the failing paths is to place the mod_A hierarchy block in a LogicLock region so that all the nodes are closer together in the floorplan.

![Figure 12–5. Design Hierarchy](image)

<table>
<thead>
<tr>
<th>From</th>
<th>To</th>
</tr>
</thead>
<tbody>
<tr>
<td>mod_A</td>
<td>reg1</td>
</tr>
<tr>
<td>mod_A</td>
<td>reg3</td>
</tr>
<tr>
<td>mod_A</td>
<td>reg4</td>
</tr>
<tr>
<td>mod_A</td>
<td>reg7</td>
</tr>
<tr>
<td>mod_A</td>
<td>reg0</td>
</tr>
</tbody>
</table>

Hierarchical LogicLock regions are also important if you are using an incremental compilation flow. Place each design partition for incremental compilation in a separate LogicLock region to reduce conflicts and ensure good results as the design develops. You can use the auto size and floating location regions to find a good design floorplan, but fix the size and placement to achieve the best results in future compilations.
For more information about using incremental compilation and recommendations for creating a design floorplan using LogicLock regions, refer to the *Quartus II IncrementalCompilation for Hierarchical and Team-Based Design* and *Best Practices for Incremental Compilation and Floorplan Assignments* chapters of the *Quartus II Handbook*, and *Analyzing and Optimizing the Design Floorplan with the Chip Planner* chapter of the *Quartus II Handbook*.

**Location Assignments**

If a small number of paths are failing to meet their timing requirements, you can use hard location assignments to optimize placement. Location assignments are less flexible for the Quartus II Fitter than LogicLock assignments. In some cases, when you are familiar with your design, you can enter location constraints in a way that produces better results.

Improving fitting results, especially for larger devices, such as Arria and Stratix series devices, can be difficult. Location assignments do not always improve the performance of the design. In many cases, you cannot improve upon the results from the Fitter by making location assignments.

**Metastability Analysis and Optimization Techniques**

Metastability problems can occur when a signal is transferred between circuitry in unrelated or asynchronous clock domains, because the designer cannot guarantee that the signal will meet its setup and hold time requirements. The mean time between failures (MTBF) is an estimate of the average time between instances when metastability could cause a design failure.

For more information about metastability and MTBF, refer to the *Understanding Metastability in FPGAs* white paper.

You can use the Quartus II software to analyze the average MTBF due to metastability when a design synchronizes asynchronous signals and to optimize the design to improve the MTBF. These metastability features are supported only for designs constrained with the TimeQuest analyzer, and for select device families.

If the MTBF of your design is low, refer to the Metastability Optimization section in the Timing Optimization Advisor, which suggests various settings that can help optimize your design in terms of metastability.

For details about the metastability features in the Quartus II software, refer to the *Managing Metastability with the Quartus II Software* chapter of the *Quartus II Handbook*. This chapter describes how to enable metastability analysis and identify the register synchronization chains in your design, provides details about metastability reports, and provides additional guidelines for managing metastability.
Optimizing Timing (Macrocell-Based CPLDs)

After resource optimization, design optimization focuses on timing. Ensure that you have made the appropriate assignments as described in the “Initial Compilation: Required Settings” section in the Design Optimization Overview chapter of the Quartus II Handbook. You must also ensure that the resource utilization is satisfactory before proceeding with timing optimization.

The following five timing parameters are primarily responsible for a design’s performance:

- **Setup time** ($t_{SU}$)—the propagation time for input data signals
- **Hold time** ($t_H$)—the propagation time for input data signals
- **Clock-to-output time** ($t_{CO}$)—the propagation time for output signals
- **Pin-to-pin delays** ($t_{PD}$)—the time required for a signal from an input pin to propagate through combinational logic and appear at an external output pin
- **Maximum clock frequency** ($f_{MAX}$)—the internal register-to-register performance

This section provides guidelines to improve the timing if the timing requirements are not met. Figure 12–6 shows the parts of the design that determine the $t_{SU}$, $t_H$, $t_{CO}$, $t_{PD}$, and $f_{MAX}$ timing parameters.

![Figure 12–6. Main Timing Parameters that Determine the System's Performance](image)

When you are analyzing a design to improve performance, be sure to consider the two major contributors to long delay paths:

- Excessive levels of logic
- Excessive loading (high fan-out)

When a MAX 7000 or MAX 3000 device signal drives more than one LAB, the programmable interconnect array (PIA) delay increases by 0.1 ns per additional LAB fan-out. Therefore, to minimize the added delay, concentrate the destination macrocells into fewer LABs, minimizing the number of LABs that are driven. The main cause of long delays in circuit design is excessive levels of logic.
## Improving Setup Time

Sometimes the $t_{SU}$ timing reported by the Quartus II Fitter does not meet your timing requirements. To improve $t_{SU}$ timing, refer to the following guidelines:

- Turn on the **Fast Input Register** option using the Assignment Editor. The **Fast Input Register** option allows input pins to directly drive macrocell registers by way of the fast-input path, thus minimizing the pin-to-register delay. This option is useful when a pin drives a D-type flipflop and there is no combinational logic between the pin and the register.

- Reduce the amount of logic between the input and the register. Excessive logic between the input pin and the register causes more delays. To improve setup time, Altera recommends reducing the amount of logic between the input pin and the register whenever possible.

- Reduce fan-out. The delay from input pins to macrocell registers increases when the fan-out of the pins increases. To improve the setup time, minimize the fan-out.

## Improving Clock-to-Output Time

To improve a design’s clock-to-output time, minimize the register-to-output-pin delay. To improve the $t_{CO}$ timing, refer to the following guidelines:

- Use the global clock. In addition to minimizing the delay from a register to an output pin, minimizing the delay from the clock pin to the register can also improve $t_{CO}$ timing. Always use the global clock for low-skew and speed-critical signals.

- Reduce the amount of logic between the register and output pin. Excessive logic between the register and the output pin causes more delay. Always minimize the amount of logic between the register and the output pin for faster clock-to-output time.

Table 12–3 lists the timing results for an EPM7064AETC100-4 device when you use a combination of the **Fast Input Register** option, global clock, and minimal logic. When you turn on the **Fast Input Register** option, $t_{SU}$ is improved ($t_{SU}$ decreases from 1.6 ns to 1.3 ns and from 2.8 ns to 2.5 ns). The $t_{CO}$ timing is improved when you use the global clock for low-skew and speed-critical signals ($t_{CO}$ decreases from 4.3 ns to 3.1 ns). However, if there is additional logic used between the input pin and the register or the register and the output pin, the $t_{SU}$ and $t_{CO}$ delays increase.

<table>
<thead>
<tr>
<th>Number of Registers</th>
<th>$t_{SU}$ (ns)</th>
<th>$t_{H}$ (ns)</th>
<th>$t_{CO}$ (ns)</th>
<th>Global Clock Used</th>
<th>Fast Input Register Option</th>
<th>D Input Location</th>
<th>Q Output Location</th>
<th>Additional Logic Between:</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>D Input Location &amp; Register</td>
</tr>
<tr>
<td>1</td>
<td>1.3</td>
<td>1.2</td>
<td>4.3</td>
<td>—</td>
<td>On</td>
<td>LAB A</td>
<td>LAB A</td>
<td>—</td>
</tr>
<tr>
<td>1</td>
<td>1.6</td>
<td>0.3</td>
<td>4.3</td>
<td>—</td>
<td>Off</td>
<td>LAB A</td>
<td>LAB A</td>
<td>—</td>
</tr>
<tr>
<td>1</td>
<td>2.5</td>
<td>0</td>
<td>3.1</td>
<td>✓</td>
<td>On</td>
<td>LAB A</td>
<td>LAB A</td>
<td>—</td>
</tr>
<tr>
<td>1</td>
<td>2.8</td>
<td>0</td>
<td>3.1</td>
<td>✓</td>
<td>Off</td>
<td>LAB A</td>
<td>LAB A</td>
<td>—</td>
</tr>
<tr>
<td>1</td>
<td>3.6</td>
<td>0</td>
<td>3.1</td>
<td>✓</td>
<td>Off</td>
<td>LAB A</td>
<td>LAB A</td>
<td>✓</td>
</tr>
<tr>
<td>1</td>
<td>2.8</td>
<td>0</td>
<td>7.0</td>
<td>✓</td>
<td>Off</td>
<td>LAB D</td>
<td>LAB A</td>
<td>—</td>
</tr>
</tbody>
</table>
Improving Propagation Delay (tPD)

Achieving fast propagation delay (tPD) timing is required in many system designs. However, if there are long delay paths through complex logic, achieving fast propagation delays can be difficult. To improve your design’s tPD, refer to the following guidelines:

- On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, and turn on Auto Parallel Expanders. Turning on the parallel expanders for individual nodes or sub-designs can increase the performance of complex logic functions. However, if the project’s pin or logic cell assignments use parallel expanders placed physically together with macrocells (which can reduce routability), parallel expanders can cause the Quartus II Fitter to have difficulties finding and optimizing a fit. Additionally, the number of macrocells required to implement the design increases and results in a no-fit error during compilation if the device resources are limited. For more information about turning on the Auto Parallel Expanders option, refer to the Area Optimization chapter of the Quartus II Handbook.

- Set Optimization Technique to Speed. By default, the Quartus II software sets the Optimization Technique option to Speed for MAX 7000 and MAX 3000 devices. Reset the Optimization Technique option to Speed only if you previously set it to Area. On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, and turn on Speed under Optimization Technique.

### Improving Maximum Frequency (fMAX)

Maintaining the system clock at or above a certain frequency is a major goal in circuit design. For example, if you have a fully synchronous system that must run at 100 MHz, the longest delay path from the output of any register to the inputs of the registers it feeds must be less than 10 ns. Maintaining the system clock speed can be difficult if there are long delay paths through complex logic. Altera recommends performing the following guidelines to increase your design’s clock speed (fMAX):

<table>
<thead>
<tr>
<th>Number of Registers</th>
<th>tsu (ns)</th>
<th>th (ns)</th>
<th>tco (ns)</th>
<th>Global Clock Used</th>
<th>Fast Input Register Option</th>
<th>D Input Location</th>
<th>Q Output Location</th>
<th>Additional Logic Between: D Input Location &amp; Register</th>
<th>Register &amp; Q Output Location</th>
</tr>
</thead>
<tbody>
<tr>
<td>16 with the same D and clock inputs</td>
<td>2.8</td>
<td>0</td>
<td>All 6.2</td>
<td>✓</td>
<td>Off</td>
<td>LAB D</td>
<td>LAB A, B</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>32 with the same D and clock inputs</td>
<td>2.8</td>
<td>0</td>
<td>All 6.4</td>
<td>✓</td>
<td>Off</td>
<td>LAB C</td>
<td>LAB A, B, C</td>
<td>—</td>
<td>—</td>
</tr>
</tbody>
</table>
On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, click More Settings, and turn on Auto Parallel Expanders. Turning on the parallel expanders for individual nodes or subdesigns can increase the performance of complex logic functions. However, if the project’s pin or logic cell assignments use parallel expanders placed physically together with macrocells (which can reduce routability), parallel expanders can cause the Quartus II compiler to have difficulties finding and optimizing a fit. Additionally, the number of macrocells required to implement the design also increases and can result in a no-fit error during compilation if the device’s resources are limited.

For more information about using the Auto Parallel Expanders option, refer to the Area Optimization chapter of the Quartus II Handbook.

Use global signals or dedicated inputs. Altera MAX 7000 and MAX 3000 devices have dedicated inputs that provide low skew and high speed for high fan-out signals. Minimize the number of control signals in the design and use the dedicated inputs to implement them.

Set Optimization Technique to Speed. By default, the Quartus II software sets the Optimization Technique option to Speed for MAX 7000 and MAX 3000 devices. Reset the Optimization Technique option to Speed only if you have previously set it to Area. You can reset the Optimization Technique option. In the Category list, select Analysis & Synthesis Settings, and turn on Speed under Optimization Technique.

Pipeline the design. Pipelining, which increases clock frequency ($f_{\text{MAX}}$), refers to dividing large blocks of combinational logic by inserting registers. When using RAM or DSP blocks, always enable the optional input and output registers.

**Optimizing Source Code—Pipelining for Complex Register Logic**

If the methods described in the preceding sections do not sufficiently improve your results, modify the design at the source to achieve the desired results. Using additional register stages (pipeline registers) consumes more device resources, but it also lowers the propagation delay between registers, allowing you to maintain high system clock speed.

For more information about pipelining registers and other examples of optimizing source code, refer to AN 584: Timing Closure Methodology for Advanced FPGA Designs.

**Scripting Support**

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II command-line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```
For more information about Tcl scripting, refer to the *Tcl Scripting* chapter of the *Quartus II Handbook*. For more information about all settings and constraints in the Quartus II software, refer to the *Quartus II Settings File Manual*. For more information about command-line scripting, refer to the *Command-Line Scripting* chapter of the *Quartus II Handbook*.

You can specify many of the options described in this section either in an instance, or at a global level, or both.

Use the following Tcl command to make a global assignment:

```
set_global_assignment -name <qsf variable name> <value>
```

Use the following Tcl command to make an instance assignment:

```
set_instance_assignment -name <qsf variable name> <value> -to <instance name>
```

If the `<value>` field includes spaces (for example, ‘Standard Fit’), you must enclose the value in straight double quotation marks.

## Initial Compilation Settings

Use the Quartus II Settings File (.qsf) variable name in the Tcl assignment to make the setting along with the appropriate value. The *Type* column indicates whether the setting is supported as a global setting, an instance setting, or both.

Table 12–4 lists the .qsf variable name and applicable values for the settings described in the “Initial Compilation: Required Settings” section in the *Design Optimization Overview* chapter of the *Quartus II Handbook*. Table 12–5 lists the advanced compilation settings.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Optimize IOC Register Placement For Timing</td>
<td>OPTIMIZE_IOC_REGISTER_PLACEMENT_FOR_TIMING</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Optimize Hold Timing</td>
<td>OPTIMIZE_HOLD_TIMING</td>
<td>OFF, IO PATHS AND MINIMUM TPD PATHS, ALL PATHS</td>
<td>Global</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Router Timing Optimization level</td>
<td>ROUTER_TIMING_OPTIMIZATION_LEVEL</td>
<td>NORMAL, MINIMUM, MAXIMUM</td>
<td>Global</td>
</tr>
</tbody>
</table>
### Resource Utilization Optimization Techniques (LUT-Based Devices)

Table 12–6 lists the `.qsf` file variable name and applicable values for the settings described in “Optimizing Timing (LUT-Based Devices)” on page 12–9.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th><code>.qsf</code> File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Auto Packed Registers (1)</td>
<td>AUTO_PACKED_REGISTERS_&lt;device family name&gt;</td>
<td>OFF, NORMAL, MINIMIZE AREA, MINIMIZE AREA WITH CHAINS, AUTO</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform WYSIWYG Primitive Resynthesis</td>
<td>ADV_NETLIST_OPT_SYNTH_WYSIWYG_REMAP</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Physical Synthesis for Combinational Logic for Reducing Area</td>
<td>PHYSICAL_SYNTHESIS_COMBO_LOGIC_FOR_AREA</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Physical Synthesis for Mapping Logic to Memory</td>
<td>PHYSICAL_SYNTHESIS_MAP_LOGIC_TO_MEMORY_FOR_AREA</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Optimization Technique</td>
<td>&lt;device family name&gt;:_OPTIMIZATION_TECHNIQUE</td>
<td>AREA, SPEED, BALANCED</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Speed Optimization Technique for Clock Domains</td>
<td>SYNTH_CRITICAL_CLOCK</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
<tr>
<td>State Machine Encoding</td>
<td>STATE_MACHINE_PROCESSING</td>
<td>AUTO, ONE-HOT, GRAY, JOHNSON, MINIMAL BITS, ONE-HOT, SEQUENTIAL, USER-ENCODE</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto RAM Replacement</td>
<td>AUTO_RAM_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto ROM Replacement</td>
<td>AUTO_ROM_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto Shift Register Replacement</td>
<td>AUTO_SHIFT_REGISTER_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto Block Replacement</td>
<td>AUTO_DSP_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Number of Processors for Parallel Compilation</td>
<td>NUM_PARALLEL_PROCESSORS</td>
<td>Integer between 1 and 16 inclusive, or ALL</td>
<td>Global</td>
</tr>
</tbody>
</table>

**Note to Table 12–6:**

(1) Allowed values for this setting depend on the device family that you select.
### I/O Timing Optimization Techniques (LUT-Based Devices)

Table 12–7 lists the .qsf file variable name and applicable values for the I/O timing optimization settings.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Optimize IOC Register Placement For Timing</td>
<td>OPTIMIZE_IOC_REGISTER_PLACEMENT_FOR_TIMING</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Fast Input Register</td>
<td>FAST_INPUT_REGISTER</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
<tr>
<td>Fast Output Register</td>
<td>FAST_OUTPUT_REGISTER</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
<tr>
<td>Fast Output Enable Register</td>
<td>FAST_OUTPUT_ENABLE_REGISTER</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
<tr>
<td>Fast OCT Register</td>
<td>FAST_OCT_REGISTER</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
</tbody>
</table>

### Register-to-Register Timing Optimization Techniques (LUT-Based Devices)

Table 12–8 lists the .qsf file variable name and applicable values for the settings described in “Register-to-Register Timing Optimization Techniques (LUT-Based Devices)” on page 12–16.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Perform WYSIWYG Primitive Resynthesis</td>
<td>ADV_NETLIST_OPT_SYNTH_WYSIWYG_REMAP</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform Physical Synthesis for Combinational Logic</td>
<td>PHYSICAL_SYNTHESIS_COMBO_LOGIC</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform Register Duplication</td>
<td>PHYSICAL_SYNTHESIS_REGISTER_DUPLICATION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform Register Retiming</td>
<td>PHYSICAL_SYNTHESIS_REGISTER_RETIMING</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform Automatic Asynchronous Signal Pipelining</td>
<td>PHYSICAL_SYNTHESISASYNCHRONOUSH_SIGNAL_PIPELINING</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Physical Synthesis Effort</td>
<td>PHYSICAL_SYNTHESIS_EFFORT</td>
<td>NORMAL, EXTRA, FAST</td>
<td>Global</td>
</tr>
<tr>
<td>Fitter Seed</td>
<td>SEED</td>
<td>&lt;integer&gt;</td>
<td>Global</td>
</tr>
<tr>
<td>Maximum Fan-Out</td>
<td>MAX_FANOUT</td>
<td>&lt;integer&gt;</td>
<td>Instance</td>
</tr>
<tr>
<td>Manual Logic Duplication</td>
<td>DUPLICATE_ATOM</td>
<td>&lt;node name&gt;</td>
<td>Instance</td>
</tr>
<tr>
<td>Optimize Power during Synthesis</td>
<td>OPTIMIZE_POWER_DURING_SYNTHESIS</td>
<td>NORMAL, OFF EXTRA_EFFORT</td>
<td>Global</td>
</tr>
<tr>
<td>Optimize Power during Fitting</td>
<td>OPTIMIZE_POWER_DURING_FITTING</td>
<td>NORMAL, OFF EXTRA_EFFORT</td>
<td>Global</td>
</tr>
</tbody>
</table>
### Document Revision History

Table 12–9 lists the revision history for this chapter.

**Table 12–9. Document Revision History (Part 1 of 2)**

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
</table>
| May 2013   | 13.0.0  | - Renamed chapter title from Area and Timing Optimization to “Timing Closure and Optimization”  
- Removed design and area/resources optimization information.  
- Added the following sections:  
  - “Fitter Aggressive Routability Optimization” on page 12–3  
  - “Tips for Analyzing Paths from/to the Source and Destination of Critical Path” on page 12–6,  
  - “Tips for Locating Multiple Paths to the Chip Planner” on page 12–8,  
  - “Tips for Creating a .tcl Script to Monitor Critical Paths Across Compiles” on page 12–8 |
- Minor text edits throughout the chapter. |
- Updated Table 13–6  
- Added the “Spine Clock Limitations” section  
- Removed the Change State Machine Encoding section from page 19  
- Removed Figure 13-5  
- Minor text edits throughout the chapter |
| May 2011    | 11.0.0  | - Reorganized sections in “Initial Compilation: Optional Fitter Settings” section  
- Added new information to “Resource Utilization” section  
- Added new information to “Duplicate Logic for Fan-Out Control” section  
- Added links to Help  
- Additional edits and updates throughout chapter |
Table 12–9. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
</table>
| December 2010 | 10.1.0  | • Added links to Help  
• Updated device support  
• Added “Debugging Timing Failures in the TimeQuest Analyzer” section  
• Removed Classic Timing Analyzer references  
• Other updates throughout chapter |
| August 2010 | 10.0.1  | Corrected link                                                          |
| July 2010   | 10.0.0  | • Moved Compilation Time Optimization Techniques section to new *Reducing Compilation Time* chapter  
• Removed references to Timing Closure Floorplan  
• Moved Smart Compilation Setting and Early Timing Estimation sections to new *Reducing Compilation Time* chapter  
• Added Other Optimization Resources section  
• Removed outdated information  
• Changed references to DSE chapter to Help links  
• Linked to Help where appropriate  
• Removed Referenced Documents section |
13. Power Optimization

The Quartus® II software offers power-driven compilation to fully optimize device power consumption. Power-driven compilation focuses on reducing your design’s total power consumption using power-driven synthesis and power-driven place-and-route. This chapter describes the power-driven compilation feature and flow in detail, as well as low power design techniques that can further reduce power consumption in your design. The techniques primarily target Arria® GX, Stratix® and Cyclone® series of devices, and HardCopy® II devices. These devices utilize a low-k dielectric material that dramatically reduces dynamic power and improves performance. Arria series, Stratix II, Stratix III, Stratix IV, and Stratix V device families include efficient logic structures called adaptive logic modules (ALMs) that obtain maximum performance while minimizing power consumption. Cyclone device families offer the optimal blend of high performance and low power in a low-cost FPGA.

For more information about a device-specific architecture, refer to the device handbook, available from the Literature and Technical Documentation page on the Altera website.

Altera provides the Quartus II PowerPlay Power Analyzer to aid you during the design process by delivering fast and accurate estimations of power consumption. You can minimize power consumption, while taking advantage of the industry’s leading FPGA performance, by using the tools and techniques described in this chapter.

For more information about the PowerPlay Power Analyzer, refer to the PowerPlay Power Analysis chapter in volume 3 of the Quartus II Handbook.

Total FPGA power consumption is comprised of I/O power, core static power, and core dynamic power. This chapter focuses on design optimization options and techniques that help reduce core dynamic power and I/O power. In addition to these techniques, there are additional power optimization techniques available for Stratix III and Stratix IV devices. These techniques include:

- Selectable Core Voltage (available only for Stratix III devices)
- Programmable Power Technology
- Device Speed Grade Selection

For more information about power optimization techniques available for Stratix III devices, refer to AN 437: Power Optimization in Stratix III FPGAs. For more information about power optimization techniques available for Stratix IV devices, refer to AN 514: Power Optimization in Stratix IV FPGAs.
Power Dissipation

This section describes the sources of power dissipation in Stratix III and Cyclone III devices. You can refine techniques that reduce power consumption in your design by understanding the sources of power dissipation.

Figure 13–1 shows the power dissipation of Stratix III and Cyclone III devices in different designs. All designs were analyzed at a fixed clock rate of 100 MHz and exhibited varied logic resource utilization across available resources.

Figure 13–1. Average Core Dynamic Power Dissipation

Notes to Figure 13–1:
(1) 103 different designs were used to obtain these results.
(2) 96 different designs were used to obtain these results.
(3) In designs using DSP blocks, DSPs consumed 5% of core dynamic power.

As shown in Figure 13–1, a significant amount of the total power is dissipated in routing for both Stratix III and Cyclone III devices, with the remaining power dissipated in logic, clock, and RAM blocks.

In Stratix and Cyclone device families, a series of column and row interconnect wires of varying lengths provide signal interconnections between logic array blocks (LABs), memory block structures, and digital signal processing (DSP) blocks or multiplier blocks. These interconnects dissipate the largest component of device power.

FPGA combinational logic is another source of power consumption. The basic building block of logic in the latest Stratix series devices is the ALM, and in Cyclone II, Cyclone III and Cyclone IV GX devices, it is the logic element (LE).

For more information about ALMs and LEs in Cyclone II, Cyclone III, Cyclone IV GX, Stratix II, Stratix III, Stratix IV, and Stratix V, devices, refer to the respective device handbook.
Memory and clock resources are other major consumers of power in FPGAs. Stratix II devices feature the TriMatrix memory architecture. TriMatrix memory includes 512-bit M512 blocks, 4-Kbit M4K blocks, and 512-Kbit M-RAM blocks, which are configurable to support many features. Stratix IV and Stratix III TriMatrix on-chip memory is an enhancement based upon the Stratix II FPGA TriMatrix memory and includes three sizes of memory blocks: MLAB blocks, M9K blocks, and M144K blocks. Stratix III, Stratix IV, and Stratix V devices feature Programmable Power Technology, an advanced architecture that enables a smooth trade-off between speed and power. The core of each Stratix III, Stratix IV, and Stratix V device is divided into tiles, each of which may be put into a high-speed or low-power mode. The primary benefit of Programmable Power Technology is to reduce static power, with a secondary benefit being a small reduction in dynamic power. Cyclone II devices have 4-Kbit M4K memory blocks, and Cyclone III and Cyclone IV GX devices have 9-Kbit M9K memory blocks.

Design Space Explorer

Design Space Explorer (DSE) is a simple, easy-to-use, design optimization utility that is included in the Quartus II software. DSE explores and reports optimal Quartus II software options for your design, targeting either power optimization, design performance, or area utilization improvements. You can use DSE to implement the techniques described in this chapter.

Figure 13–2 shows the DSE user interface. The Settings tab is divided into Project Settings and Exploration Settings.

Figure 13–2. Design Space Explorer User Interface
The **Search for Lowest Power** option, under **Exploration Settings**, uses a predefined exploration space that targets overall design power improvements. This setting focuses on applying different options that specifically reduce total design thermal power.

By default, the Quartus II PowerPlay Power Analyzer is run for every exploration performed by the DSE when the **Search for Lowest Power** option is selected. This helps you debug your design and determine trade-offs between power requirements and performance optimization.

For more information about the DSE, refer to *About Design Space Explorer* in Quartus II Help.

### Power-Driven Compilation

The standard Quartus II compilation flow consists of Analysis and Synthesis, placement and routing, Assembly, and Timing Analysis. Power-driven compilation takes place at the Analysis and Synthesis and Place-and-Route stages. Quartus II software settings that control power-driven compilation are located in the **PowerPlay power optimization** list on the **Analysis & Synthesis Settings** page, and the **PowerPlay power optimization** list on the **Fitter Settings** page. The following sections describes these power optimization options at the Analysis and Synthesis and Fitter levels.

#### Power-Driven Synthesis

Synthesis netlist optimization occurs during the synthesis stage of the compilation flow. The optimization technique makes changes to the synthesis netlist to optimize your design according to the selection of area, speed, or power optimization. This section describes power optimization techniques at the synthesis level.
The Analysis & Synthesis Settings page allows you to specify logic synthesis options. The PowerPlay power optimization option is available for all devices supported by the Quartus II software except MAX® 3000 and MAX 7000 devices. (Figure 13–3).

**Figure 13–3. Analysis & Synthesis Settings Page**

Table 13–1 shows the settings in the PowerPlay power optimization list. You can apply these settings on a project or entity level.

<table>
<thead>
<tr>
<th>Settings</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Off</td>
<td>No netlist, placement, or routing optimizations are performed to minimize power.</td>
</tr>
<tr>
<td>Normal compilation (Default)</td>
<td>Low compute effort algorithms are applied to minimize power through netlist optimizations as long as they are not expected to reduce design performance.</td>
</tr>
<tr>
<td>Extra effort</td>
<td>High compute effort algorithms are applied to minimize power through netlist optimizations. Max performance might be impacted.</td>
</tr>
</tbody>
</table>

The Normal compilation setting is turned on by default. This setting performs memory optimization and power-aware logic mapping during synthesis.
Memory blocks can represent a large fraction of total design dynamic power as described in "Reducing Memory Power Consumption" on page 13–14. Minimizing the number of memory blocks accessed during each clock cycle can significantly reduce memory power. Memory optimization involves effective movement of user-defined read/write enable signals to associated read-and-write clock enable signals for all memory types (Figure 13–4).

Figure 13–4 shows a default implementation of a simple dual-port memory block in which write-clock enable signals and read-clock enable signals are connected to VCC, making both read and write memory ports active during each clock cycle. Memory transformation effectively moves the read-enable and write-enable signals to the respective read-clock enable and write-clock enable signals. By using this technique, memory ports are shut down when they are not accessed. This significantly reduces your design’s memory power consumption. For more information about clock enable signals, refer to “Reducing Memory Power Consumption” on page 13–14. For Stratix III, Stratix IV, and Stratix V devices, the memory transformation takes place at the Fitter level by selecting the Normal compilation settings for the power optimization option.

In Stratix III, Cyclone III, Cyclone IV GX, and Stratix III devices, the specified read-during-write behavior can significantly impact the power of single-port and bidirectional dual-port RAMs. It is best to set the read-during-write parameter to “Don’t care” (at the HDL level), as it allows an optimization whereby the read-enable signal can be set to the inversion of the existing write-enable signal (if one exists). This allows the core of the RAM to shut down (that is, not toggle), which saves a significant amount of power.

The other type of power optimization that takes place with the Normal compilation setting is power-aware logic mapping. The power-aware logic mapping reduces power by rearranging the logic during synthesis to eliminate nets with high toggle rates.

The Extra effort setting performs the functions of the Normal compilation setting and other memory optimizations to further reduce memory power by shutting down memory blocks that are not accessed. This level of memory optimization can require extra logic, which can reduce design performance.
The **Extra effort** setting also performs power-aware memory balancing. Power-aware memory balancing automatically chooses the best memory configuration for your memory implementation and provides optimal power saving by determining the number of memory blocks, decoder, and multiplexer circuits required. If you have not previously specified target-embedded memory blocks for your design’s memory functions, the power-aware balancer automatically selects them during memory implementation.

**Figure 13–5** shows an example of a 4k × 4 (4k deep and 4 bits wide) memory implementation in two different configurations using M4K memory blocks available in Stratix II devices. The minimum logic area implementation uses M4K blocks configured as 4k × 1. This implementation is the default in the Quartus II software because it has the minimum logic area (0 logic cells) and the highest speed. However, all four M4K blocks are active on each memory access in this implementation, which increases RAM power. The minimum RAM power implementation is created by selecting **Extra effort** in the **PowerPlay power optimization** list. This implementation automatically uses four M4K blocks configured as 1k × 4 for optimal power saving. An address decoder is implemented by the RAM megafunction to select which of the four M4K blocks should be activated on a given cycle, based on the state of the top two user address bits. The RAM megafunction automatically implements a multiplexer to feed the downstream logic by choosing the appropriate M4K output. This implementation reduces RAM power because only one M4K block is active on any cycle, but it requires extra logic cells, costing logic area and potentially impacting design performance.

There is a trade-off between power saved by accessing fewer memories and power consumed by the extra decoder and multiplexor logic. The Quartus II software automatically balances the power savings against the costs to choose the lowest power configuration for each logical RAM. The benchmark data shows that the power-driven synthesis can reduce memory power consumption by as much as 60% in Stratix devices.

**Figure 13–5. 4K × 4 Memory Implementation Using Multiple M4K Blocks**
Memory optimization options can also be controlled by the **Low_Power_Mode** parameter in the **Default Parameters** page of the **Settings** dialog box. The settings for this parameter are **None**, **Auto**, and **ALL**. **None** corresponds to the **Off** setting in the **PowerPlay power optimization** list. **Auto** corresponds to the **Normal compilation** setting and **ALL** corresponds to the **Extra effort** setting, respectively. You can apply PowerPlay power optimization either on a compiler basis or on individual entities. The **Low_Power_Mode** parameter always takes precedence over the **Optimize Power for Synthesis** option for power optimization on memory.

You can also set the **MAXIMUM_DEPTH** parameter manually to configure the memory for low power optimization. This technique is the same as the power-aware memory balancer, but it is manual rather than automatic like the **Extra effort** setting in the **PowerPlay power optimization** list. You can set the **MAXIMUM_DEPTH** parameter for memory modules manually in the megafunction instantiation or in the MegaWizard™ Plug-In Manager for power optimization as described in “Reducing Memory Power Consumption” on page 13–14. The **MAXIMUM_DEPTH** parameter always takes precedence over the **Optimize Power for Synthesis** options for power optimization on memory optimization.

For step-by-step instructions on how to perform power-driven synthesis, refer to *Running a Power-Optimized Compilation* in Quartus II Help.
Power-Driven Fitter

The Fitter Settings page enables you to specify options for fitting (Figure 13–6). The **PowerPlay power optimization** option is available for Arria GX, Arria II GX, Cyclone II, Cyclone III, Cyclone IV, HardCopy series, Stratix II, Stratix II GX, Stratix III, Stratix IV, and Stratix V devices.

**Table 13–2. Power-Driven Fitter Option**

<table>
<thead>
<tr>
<th>Settings</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Off</td>
<td>No netlist, placement, or routing optimizations are performed to minimize power.</td>
</tr>
<tr>
<td>Normal compilation (Default)</td>
<td>Low compute effort algorithms are applied to minimize power through placement and routing optimizations as long as they are not expected to reduce design performance.</td>
</tr>
<tr>
<td>Extra effort</td>
<td>High compute effort algorithms are applied to minimize power through placement and routing optimizations. Max performance might be impacted.</td>
</tr>
</tbody>
</table>
The Normal compilation setting is selected by default and performs DSP optimization by creating power-efficient DSP block configurations for your DSP functions. For Stratix III, Stratix IV, and Stratix V devices, this setting, which is based on timing constraints entered for the design, enables the Programmable Power Technology to configure tiles as high-speed mode or low-power mode. Programmable Power Technology is always turned ON even when the OFF setting is selected for the Fitter PowerPlay power optimization option. Tiles are the combination of LAB and MLAB pairs (including the adjacent routing associated with LAB and MLAB), which can be configured to operate in high-speed or low-power mode. This level of power optimization does not have any affect on the fitting, timing results, or compile time. Also, for Stratix III devices, this setting enables the memory transformation as described in “Power-Driven Synthesis” on page 13–4.

For more information about Stratix III power optimization, refer to AN 437: Power Optimization in Stratix III FPGAs. For more information about Stratix IV power optimization, refer to AN 514: Power Optimization in Stratix IV FPGAs.

The Extra effort setting performs the functions of the Normal compilation setting and other place-and-route optimizations during fitting to fully optimize the design for power. The Fitter applies an extra effort to minimize power even after timing requirements have been met by effectively moving the logic closer during placement to localize high-toggling nets, and using routes with low capacitance. However, this effort can increase the compilation time.

The Extra effort setting uses a Value Change Dump File (.vcd) that guides the Fitter to fully optimize the design for power, based on the signal activity of the design. The best power optimization during fitting results from using the most accurate signal activity information. Signal activities from full post-fit netlist (timing) simulation provide the highest accuracy because all node activities reflect the actual design behavior, provided that supplied input vectors are representative of typical design operation. If you do not have a .vcd file, the Quartus II software uses assignments, clock assignments, and vectorless estimation values (PowerPlay Power Analyzer Tool settings) to estimate the signal activities. This information is used to optimize your design for power during fitting. The benchmark data shows that the power-driven Fitter technique can reduce power consumption by as much as 19% in Stratix devices. On average, you can reduce core dynamic power by 16% with the Extra effort synthesis and Extra effort fitting settings, as compared to the Off settings in both synthesis and Fitter options for power-driven compilation.

Only the Extra effort setting in the PowerPlay power optimization list for the Fitter option uses the signal activities (from .vcd files) during fitting. The settings made in the PowerPlay Power Analyzer Settings page in the Settings dialog box are used to calculate the signal activity of your design.

For more information about .vcd files and how to create them, refer to the PowerPlay Power Analysis chapter in volume 3 of the Quartus II Handbook.

For step-by-step instructions on how to perform power-driven fitting, refer to Running a Power-Optimized Compilation in Quartus II Help.
Area-Driven Synthesis

Using area optimization rather than timing or delay optimization during synthesis saves power because you use fewer logic blocks. Using less logic usually means less switching activity. The Quartus II integrated synthesis tool provides Speed, Balanced, or Area for the Optimization Technique option. You can also specify this logic option for specific modules in your design with the Assignment Editor in cases where you want to reduce area using the Area setting (potentially at the expense of register-to-register timing performance) while leaving the default Optimization Technique setting at Balanced (for the best trade-off between area and speed for certain device families). The Speed Optimization Technique can increase the resource usage of your design if the constraints are too aggressive, and can also result in increased power consumption.

The benchmark data shows that the area-driven technique can reduce power consumption by as much as 31% in Stratix devices and as much as 15% in Cyclone devices.

Gate-Level Register Retiming

You can also use gate-level register retiming to reduce circuit switching activity. Retiming shuffles registers across combinational blocks without changing design functionality. The Perform gate-level register retiming option in the Quartus II software enables the movement of registers across combinational logic to balance timing, allowing the software to trade off the delay between timing critical and noncritical timing paths.

Retiming uses fewer registers than pipelining. Figure 13–7 shows an example of gate-level register retiming, where the 10 ns critical delay is reduced by moving the register relative to the combinational logic, resulting in the reduction of data depth and switching activity.

Figure 13–7. Gate-Level Register Retiming
Gate-level register retiming makes changes at the gate level. If you are using an atom netlist from a third-party synthesis tool, you must also select the **Perform WYSIWYG primitive resynthesis** option to undo the atom primitives to gates mapping (so that register retiming can be performed), and then to remap gates to Altera primitives. When using Quartus II integrated synthesis, retiming occurs during synthesis before the design is mapped to Altera primitives. The benchmark data shows that the combination of WYSIWYG remapping and gate-level register retiming techniques can reduce power consumption by as much as 6% in Stratix devices and as much as 21% in Cyclone devices.

For more information about register retiming, refer to the **Netlist Optimizations and Physical Synthesis** chapter in volume 2 of the *Quartus II Handbook*.

### Design Guidelines

Several low-power design techniques can reduce power consumption when applied during FPGA design implementation. This section provides detailed design techniques for Cyclone II, Cyclone III, Cyclone IV GX, Stratix II, and Stratix III devices that affect overall design power. The results of these techniques might be different from design to design.

#### Clock Power Management

Clocks represent a significant portion of dynamic power consumption due to their high switching activity and long paths. Figure 13–1 on page 13–2 shows a 14% average contribution to power consumption for global clock routing in Stratix III devices and 16% in Cyclone III devices. Actual clock-related power consumption is higher than this because the power consumed by local clock distribution within logic, memory, and DSP or multiplier blocks is included in the power consumption for the respective blocks.

Clock routing power is automatically optimized by the Quartus II software, which enables only those portions of the clock network that are required to feed downstream registers. Power can be further reduced by gating clocks when they are not required. It is possible to build clock-gating logic, but this approach is not recommended because it is difficult to generate a glitch free clock in FPGAs using ALMs or LEs.

Arria GX, Arria II GX, Cyclone III, Cyclone IV, Stratix II, Stratix III, Stratix IV, and Stratix V devices use clock control blocks that include an enable signal. A clock control block is a clock buffer that lets you dynamically enable or disable the clock network and dynamically switch between multiple sources to drive the clock network. You can use the Quartus II MegaWizard Plug-In Manager to create this clock control block with the ALTCLKCTRL megafunction. Arria GX, Arria II GX, Cyclone III, Cyclone IV, Stratix II, Stratix III, Stratix IV, and Stratix V devices provide clock control blocks for global clock networks. In addition, Stratix II, Stratix III,
Stratix IV and Stratix V devices have clock control blocks for regional clock networks. The dynamic clock enable feature lets internal logic control the clock network. When a clock network is powered down, all the logic fed by that clock network does not toggle, thereby reducing the overall power consumption of the device. Figure 13–8 shows a 4-input clock control block diagram.

**Figure 13–8. Clock Control Block Diagram**

The enable signal is applied to the clock signal before being distributed to global routing. Therefore, the enable signal can either have a significant timing slack (at least as large as the global routing delay) or it can reduce the $f_{MAX}$ of the clock signal.

For more information about using clock control blocks, refer to the *Clock Control Block Megafunction User Guide (ALTCLKCTRL)*.

Another contributor to clock power consumption is the LAB clock that distributes a clock to the registers within a LAB. LAB clock power can be the dominant contributor to overall clock power. For example, in Cyclone III devices, each LAB can use two clocks and two clock enable signals, as shown in Figure 13–9. Each LAB’s clock signal and clock enable signal are linked. For example, an LE in a particular LAB using the `labclk1` signal also uses the `labclkena1` signal.

**Figure 13–9. LAB-Wide Control Signals**
To reduce LAB-wide clock power consumption without disabling the entire clock tree, use the LAB-wide clock enable to gate the LAB-wide clock. The Quartus II software automatically promotes register-level clock enable signals to the LAB-level. All registers within an LAB that share a common clock and clock enable are controlled by a shared gated clock. To take advantage of these clock enables, use a clock enable construct in the relevant HDL code for the registered logic.

**LAB-Wide Clock Enable Example**

The VHDL code in Example 13–1 makes use of a LAB-wide clock enable. This clock-gating logic is automatically turned into an LAB-level clock enable signal.

**Example 13–1.**

```vhdl
IF clk'event AND clock = '1' THEN
  IF logic_is_enabled = '1' THEN
    reg <= value;
  ELSE
    reg <= reg;
  END IF;
END IF;
```

For more information about LAB-wide control signals, refer to the *Stratix II Architecture*, *Cyclone III Device Family Overview*, or *Cyclone II Architecture* chapters in the respective device handbook.

**Reducing Memory Power Consumption**

The memory blocks in FPGA devices can represent a large fraction of typical core dynamic power. Memory consumes approximately 20% of the core dynamic power in typical Cyclone III and Stratix III device designs. Memory blocks are unlike most other blocks in the device because most of their power is tied to the clock rate, and is insensitive to the toggle rate on the data and address lines.

When a memory block is clocked, there is a sequence of timed events that occur within the block to execute a read or write. The circuitry controlled by the clock consumes the same amount of power regardless of whether or not the address or data has changed from one cycle to the next. Thus, the toggle rate of input data and the address bus have no impact on memory power consumption.

The key to reducing memory power consumption is to reduce the number of memory clocking events. You can achieve this through clock network-wide gating described in "Clock Power Management" on page 13–12, or on a per-memory basis through use of the clock enable signals on the memory ports. Figure 13–10 shows the logical view of the internal clock of the memory block. Use the appropriate enable signals on the memory to make use of the clock enable signal instead of gating the clock.

**Figure 13–10. Memory Clock Enable Signal**
Using the clock enable signal enables the memory only when necessary and shuts it down for the rest of the time, reducing the overall memory power consumption. You can use the MegaWizard Plug-In Manager to create these enable signals by selecting the Clock enable signal option for the appropriate port when generating the memory block function (Figure 13–11).

**Figure 13–11. MegaWizard Plug-In Manager RAM 2-Port Clock Enable Signal Selectable Option**

For example, consider a design that contains a 32-bit-wide M4K memory block in ROM mode that is running at 200 MHz. Assuming that the output of this block is only required approximately every four cycles, this memory block will consume 8.45 mW of dynamic power according to the demands of the downstream logic. By adding a small amount of control logic to generate a read clock enable signal for the memory block only on the relevant cycles, the power can be cut 75% to 2.15 mW.

You can also use the MAXIMUM_DEPTH parameter in your memory megafunction to save power in Cyclone II, Cyclone III, Cyclone IV GX, Stratix II, Stratix III, Stratix IV, and Stratix V devices; however, this approach might increase the number of LEs required to implement the memory and affect design performance.

You can set the MAXIMUM_DEPTH parameter for memory modules manually in the megafunction instantiation or in the MegaWizard Plug-In Manager (Figure 13–12). The Quartus II software automatically chooses the best design memory configuration for optimal power, as described in “Power-Driven Compilation” on page 13–4.
Memory Power Reduction Example

Table 13–3 shows power usage measurements for a 4K × 36 simple dual-port memory implemented using multiple M4K blocks in a Stratix II EP2S15 device. For each implementation, the M4K blocks are configured with a different memory depth.

Table 13–3. 4K × 36 Simple Dual-Port Memory Implemented Using Multiple M4K Blocks

<table>
<thead>
<tr>
<th>M4K Configuration</th>
<th>Number of M4K Blocks</th>
<th>ALUTs</th>
</tr>
</thead>
<tbody>
<tr>
<td>4K × 1 (Default setting)</td>
<td>36</td>
<td>0</td>
</tr>
<tr>
<td>2K × 2</td>
<td>36</td>
<td>40</td>
</tr>
<tr>
<td>1K × 4</td>
<td>36</td>
<td>62</td>
</tr>
<tr>
<td>512 × 9</td>
<td>32</td>
<td>143</td>
</tr>
<tr>
<td>256 × 18</td>
<td>32</td>
<td>302</td>
</tr>
<tr>
<td>128 × 36</td>
<td>32</td>
<td>633</td>
</tr>
</tbody>
</table>
Figure 13–13 shows the amount of power saved using the MAXIMUM_DEPTH parameter. For all implementations, a user-provided read enable signal is present to indicate when read data is required. Using this power-saving technique can reduce power consumption by as much as 60%.

![Power Savings Using the MAXIMUM_DEPTH Parameter](image)

As the memory depth becomes more shallow, memory dynamic power decreases because unaddressed M4K blocks can be shut off using a decoded combination of address bits and the read enable signal. For a 128-deep memory block, power used by the extra LEs starts to outweigh the power gain achieved by using a more shallow memory block depth. The power consumption of the memory blocks and associated LEs depends on the memory configuration.

The SOPC Builder and Qsys system do not offer specific power savings control for on-chip memory block. There is no read enable, write enable, or clock enable that you can enable in the on-chip RAM megafunction to shut down the RAM block in the SOPC Builder and Qsys system.

**Pipelining and Retiming**

Designs with many glitches consume more power because of faster switching activity. Glitches cause unnecessary and unpredictable temporary logic switches at the output of combinational logic. A glitch usually occurs when there is a mismatch in input signal timing leading to unequal propagation delay.

For example, consider an input change on one input of a 2-input XOR gate from 1 to 0, followed a few moments later by an input change from 0 to 1 on the other input. For a moment, both inputs become 1 (high) during the state transition, resulting in 0 (low) at the output of the XOR gate. Subsequently, when the second input transition takes place, the XOR gate output becomes 1 (high). During signal transition, a glitch is
produced before the output becomes stable, as shown in Figure 13–14. This glitch can propagate to subsequent logic and create unnecessary switching activity, increasing power consumption. Circuits with many XOR functions, such as arithmetic circuits or cyclic redundancy check (CRC) circuits, tend to have many glitches if there are several levels of combinational logic between registers.

Figure 13–14. XOR Gate Showing Glitch at the Output

Pipelining can reduce design glitches by inserting flipflops into long combinational paths. Flipflops do not allow glitches to propagate through combinational paths. Therefore, a pipelined circuit tends to have less glitching. Pipelining has the additional benefit of generally allowing higher clock speed operations, although it does increase the latency of a circuit (in terms of the number of clock cycles to a first result). Figure 13–15 shows an example where pipelining is applied to break up a long combinational path.

Figure 13–15. Pipelining Example

Pipelining is very effective for glitch-prone arithmetic systems because it reduces switching activity, resulting in reduced power dissipation in combinational logic. Additionally, pipelining allows higher-speed operation by reducing logic-level numbers between registers. The disadvantage of this technique is that if there are not many glitches in your design, pipelining can increase power consumption by adding unnecessary registers. Pipelining can also increase resource utilization. The benchmark data shows that pipelining can reduce dynamic power consumption by as much as 30% in Cyclone and Stratix devices.
**Architectural Optimization**

You can use design-level architectural optimization by taking advantage of specific device architecture features. These features include dedicated memory and DSP or multiplier blocks available in FPGA devices to perform memory or arithmetic-related functions. You can use these blocks in place of LUTs to reduce power consumption. For example, you can build large shift registers from RAM-based FIFO buffers instead of building the shift registers from the LE registers.

The Stratix device family allows you to efficiently target small, medium, and large memories with the TriMatrix memory architecture. Each TriMatrix memory block is optimized for a specific function. The M512 memory blocks available in Stratix II devices are useful for implementing small FIFO buffers, DSP, and clock domain transfer applications. M512 memory blocks are more power-efficient than the distributed memory structures in some competing FPGAs. The M4K memory blocks are used to implement buffers for a wide variety of applications, including processor code storage, large look-up table implementation, and large memory applications. The M-RAM blocks are useful in applications where a large volume of data must be stored on-chip. Effective utilization of these memory blocks can have a significant impact on power reduction in your design.

The latest Stratix and Cyclone device families have configurable M9K memory blocks that provide various memory functions such as RAM, FIFO buffers, and ROM.

For more information about using DSP and memory blocks efficiently, refer to the *Area and Timing Optimization* chapter in volume 2 of the *Quartus II Handbook*.

**I/O Power Guidelines**

Nonterminated I/O standards such as LVTTL and LVCMOS have a rail-to-rail output swing. The voltage difference between logic-high and logic-low signals at the output pin is equal to the V<sub>CCIO</sub> supply voltage. If the capacitive loading at the output pin is known, the dynamic power consumed in the I/O buffer can be calculated as shown in Equation 13–1:

**Equation 13–1. Capacitive loading at the output pin**

\[
P = 0.5 \times F \times C \times V^2
\]

In this equation, \( F \) is the output transition frequency and \( C \) is the total load capacitance being switched. \( V \) is equal to \( V_{CCIO} \) supply voltage. Because of the quadratic dependence on \( V_{CCIO} \), lower voltage standards consume significantly less dynamic power.

Transistor-to-transistor logic (TTL) I/O buffers consume very little static power. As a result, the total power consumed by a LVTTL or LVCMOS output is highly dependent on load and switching frequency.

When using resistively terminated I/O standards like SSTL and HSTL, the output load voltage swings by a small amount around some bias point. The same dynamic power equation is used, where \( V \) is the actual load voltage swing. Because this is much smaller than \( V_{CCIO} \), dynamic power is lower than for nonterminated I/O under similar conditions. These resistively terminated I/O standards dissipate significant static (frequency-independent) power, because the I/O buffer is constantly driving
current into the resistive termination network. However, the lower dynamic power of these I/O standards means they often have lower total power than LVCMOS or LVTTL for high-frequency applications. Use the lowest drive strength I/O setting that meets your speed and waveform requirements to minimize I/O power when using resistively terminated standards.

You can save a small amount of static power by connecting unused I/O banks to the lowest possible $V_{CCIO}$ voltage of 1.2 V.

Table 13–4 shows the total supply and thermal power consumed by outputs using different I/O standards for Stratix II devices. The numbers are for an I/O pin transmitting random data clocked at 200 MHz with a 10 pF capacitive load.

For this configuration, nonterminated standards generally use less power, but this is not always the case. If the frequency or the capacitive load is increased, the power consumed by nonterminated outputs increases faster than the power of terminated outputs.

<table>
<thead>
<tr>
<th>Standard</th>
<th>Total Supply Current Drawn from $V_{CCIO}$ Supply (mA)</th>
<th>Total On-Chip Thermal Power Dissipation (mW)</th>
</tr>
</thead>
<tbody>
<tr>
<td>3.3-V LVTTL</td>
<td>2.42</td>
<td>9.87</td>
</tr>
<tr>
<td>2.5-V LVCMOS</td>
<td>1.9</td>
<td>6.69</td>
</tr>
<tr>
<td>1.8-V LVCMOS</td>
<td>1.34</td>
<td>4.18</td>
</tr>
<tr>
<td>1.5-V LVCMOS</td>
<td>1.18</td>
<td>3.58</td>
</tr>
<tr>
<td>3.3-V PCI</td>
<td>2.47</td>
<td>10.23</td>
</tr>
<tr>
<td>SSTL-2 class I</td>
<td>6.07</td>
<td>4.42</td>
</tr>
<tr>
<td>SSTL-2 class II</td>
<td>10.72</td>
<td>5.1</td>
</tr>
<tr>
<td>SSTL-18 class I</td>
<td>5.33</td>
<td>3.28</td>
</tr>
<tr>
<td>SSTL-18 class II</td>
<td>8.56</td>
<td>4.06</td>
</tr>
<tr>
<td>HSTL-15 class I</td>
<td>6.06</td>
<td>3.49</td>
</tr>
<tr>
<td>HSTL-15 class II</td>
<td>11.08</td>
<td>4.87</td>
</tr>
<tr>
<td>HSTL-18 class I</td>
<td>6.87</td>
<td>4.09</td>
</tr>
<tr>
<td>HSTL-18 class II</td>
<td>12.33</td>
<td>5.82</td>
</tr>
</tbody>
</table>

When calculating I/O power, the PowerPlay Power Analyzer uses the default capacitive load set for the I/O standard in the **Capacitive Loading** page of the **Device and Pin Options** dialog box. For Stratix II devices, if **Enable Advanced I/O Timing** is turned on, I/O power is measured using an equivalent load calculated as the sum of the near capacitance, the transmission line distributed capacitance, and the far-end capacitance as defined in the **Board Trace Model** page of the **Device and Pin Options** dialog box or the Board Trace Model view in the Pin Planner. Any other components defined in the board trace model are not taken into account for the power measurement.

For Cyclone III, Cyclone IV GX, Stratix III, Stratix IV, and Stratix V, devices, Advanced I/O Timing, which uses the full board trace model, is always used.

For information about using Advanced I/O Timing and configuring a board trace model, refer to the **I/O Management** chapter in volume 2 of the **Quartus II Handbook**.

**Dynamically Controlled On-Chip Terminations**

Stratix V, Stratix IV and Stratix III FPGAs offer dynamic on-chip termination (OCT). Dynamic OCT enables series termination (RS) and parallel termination (RT) to dynamically turn on/off during the data transfer. This feature is especially useful when Stratix V, Stratix IV and Stratix III FPGAs are used with external memory interfaces, such as interfacing with DDR memories.

Compared to conventional termination, dynamic OCT reduces power consumption significantly as it eliminates the constant DC power consumed by parallel termination when transmitting data. Parallel termination is extremely useful for applications that interface with external memories where I/O standards, such as HSTL and SSTL, are used. Parallel termination supports dynamic OCT, which is useful for bidirectional interfaces (see Figure 13–16).

**Figure 13–16. Stratix III On-Chip Parallel Termination**

![Stratix III OCT](image)

The following is an example of power saving for a DDR3 interface using on-chip parallel termination.
The static current consumed by parallel OCT is equal to the $V_{CCIO}$ voltage divided by 100 $\Omega$. For DDR3 interfaces that use SSTL-15, the static current is $1.5 \, V / 100 \, \Omega = 15 \, mA$ per pin. Therefore, the static power is $1.5 \, V \times 15 \, mA = 22.5 \, mW$. For an interface with 72 DQ and 18 DQS pins, the static power is $90 \, pins \times 22.5 \, mW = 2.025 \, W$. Dynamic parallel OCT disables parallel termination during write operations, so if writing occurs 50% of the time, the power saved by dynamic parallel OCT is $50\% \times 2.025 \, W = 1.0125 \, W$.

For more information about dynamic OCT in Stratix IV and Stratix III devices, refer to the *Stratix III Device I/O Features* chapter in the *Stratix III Device Handbook* and the *Stratix IV Device I/O Features* chapter in the *Stratix IV Device Handbook*, respectively.

**Power Optimization Advisor**

The Quartus II software includes the Power Optimization Advisor, which provides specific power optimization advice and recommendations based on the current design project settings and assignments. The advisor covers many of the suggestions listed in this chapter. The following example shows how to reduce your design power with the Power Optimization Advisor.

**Power Optimization Advisor Example**

After compiling your design, run the PowerPlay Power Analyzer to determine your design power and to see where power is dissipated in your design. Based on this information, you can run the Power Optimization Advisor to implement recommendations that can reduce design power. Figure 13–17 shows the Power Optimization Advisor after compiling a design that is not fully optimized for power.

The Power Optimization Advisor shows the recommendations that can reduce power in your design. The recommendations are split into stages to show the order in which you should apply the recommended settings. The first stage shows mostly CAD setting options that are easy to implement and highly effective in reducing design power. An icon indicates whether each recommended setting is made in the current
In Figure 13–17, the checkmark icons for Stage 1 shows the recommendations that are already implemented. The warning icons indicate recommendations that are not followed for this compilation. The information icon shows the general suggestions. Each recommendation includes the description, summary of the effect of the recommendation, and the action required to make the appropriate setting.

There is a link from each recommendation to the appropriate location in the Quartus II user interface where you can change the setting. You can change the Power-Driven Synthesis setting by clicking Open Settings dialog box - Analysis & Synthesis Settings page. The Settings dialog box is shown with the Analysis & Synthesis Settings page selected, where you can change the PowerPlay power optimization settings.

After making the recommended changes, recompile your design. The Power Optimization Advisor indicates with green check marks that the recommendations were implemented successfully (Figure 13–18). You can use the PowerPlay Power Analyzer to verify your design power results.

The recommendations listed in Stage 2 generally involve design changes, rather than CAD settings changes as in Stage 1. You can use these recommendations to further reduce your design power consumption. Altera recommends that you implement Stage 1 recommendations first, then the Stage 2 recommendations.

**Conclusion**

The combination of a smaller process technology, the use of low-k dielectric material, and reduced supply voltage significantly reduces dynamic power consumption in the latest FPGAs. To further reduce your dynamic power, use the design recommendations presented in this chapter to optimize resource utilization and minimize power consumption.
Document Revision History

Table 13–5 shows the revision history for this chapter.

Table 13–5. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>Added a note to “Memory Power Reduction Example” on page 13–16 on Qsys and SOPC Builder power savings limitation for on-chip memory block.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Template update.</td>
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</table>
| July 2010      | 10.0.0  | • Was chapter 11 in the 9.1.0 release  
• Updated Figures 14-2, 14-3, 14-6, 14-18, 14-19, and 14-20  
• Updated device support  
• Minor editorial updates |
| November 2009  | 9.1.0   | • Updated Figure 11-1 and associated references  
• Updated device support  
• Minor editorial update |
| March 2009     | 9.0.0   | • Was chapter 9 in the 8.1.0 release  
• Updated for the Quartus II software release  
• Added benchmark results  
• Removed several sections  
• Updated Figure 13–1, Figure 13–17, and Figure 13–18 |
| November 2008  | 8.1.0   | • Changed to 8½” × 11” page size  
• Changed references to altsyncram to RAM  
• Minor editorial updates |
| May 2008       | 8.0.0   | • Added support for Stratix IV devices  
• Updated Table 9–1 and 9–9  
• Updated “Architectural Optimization” on page 9–22  
• Added “Dynamically-Controlled On-Chip Terminations” on page 9–26  
• Updated “Referenced Documents” on page 9–29  
• Updated references |

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
14. Area Optimization

This chapter describes techniques to reduce resource usage when designing for Altera® devices.

This chapter includes the following topics:

- “Resource Utilization” on page 14–1
- “Optimizing Resource Utilization (LUT-Based Devices)” on page 14–2
- “Optimizing Resource Utilization (Macrocell-Based CPLDs)” on page 14–12
- “Scripting Support” on page 14–17

Resource Utilization

Determining device utilization is important regardless of whether your design achieved a successful fit. If your compilation results in a no-fit error, resource utilization information is important for analyzing the fitting problems in your design. If your fitting is successful, review the resource utilization information to determine whether the future addition of extra logic or other design changes might introduce fitting difficulties. Also, review the resource utilization information to determine if it is impacting timing performance.

To determine resource usage, refer to the Flow Summary section of the Compilation Report. This section reports resource utilization, including pins, memory bits, digital signal processing (DSP) blocks, and phase-locked loops (PLLs). Flow Summary indicates whether your design exceeds the available device resources. More detailed information is available by viewing the reports under Resource Section in the Fitter section of the Compilation Report.

Flow Summary shows the overall logic utilization. The Fitter can spread logic throughout the device, which may lead to higher overall utilization.

As the device fills up, the Fitter automatically searches for logic functions with common inputs to place in one ALM. The number of packed registers also increases. Therefore, a design that has high overall utilization might still have space for extra logic if the logic and registers can be packed together more tightly.

The reports under the Resource Section in the Fitter section of the Compilation Report provide more detailed resource information. The Fitter Resource Usage Summary report breaks down the logic utilization information and provides other resource information, including the number of bits in each type of memory block. This panel also contains a summary of the usage of global clocks, PLLs, DSP blocks, and other device-specific resources.
You can also view reports describing some of the optimizations that occurred during compilation. For example, if you use Quartus® II integrated synthesis, the reports in the Optimization Results folder in the **Analysis & Synthesis** section include information about registers that integrated synthesis removed during synthesis. Use this report to estimate device resource utilization for a partial design to ensure that registers were not removed due to missing connections with other parts of the design.

If a specific resource usage is reported as less than 100% and a successful fit cannot be achieved, either there are not enough routing resources or some assignments are illegal. In either case, a message appears in the **Processing** tab of the Messages window describing the problem.

If the Fitter finishes unsuccessfully and runs much faster than on similar designs, a resource might be over-utilized or there might be an illegal assignment. If the Quartus II software seems to run for an excessively long time compared to runs on similar designs, a legal placement or route probably cannot be found. In the Compilation Report, look for errors and warnings that indicate these types of problems.

You can use the Chip Planner to find areas of the device that have routing congestion on specific types of routing resources. If you find areas with very high congestion, analyze the cause of the congestion. Issues such as high fan-out nets not using global resources, an improperly chosen optimization goal (speed versus area), very restrictive floorplan assignments, or the coding style can cause routing congestion. After you identify the cause, modify the source or settings to reduce routing congestion.

For more information about Fitter Resources Report, refer to *Fitter Resources Report* in Quartus II Help. For information about how to view routing congestion, refer to *Displaying Resources and Information* in Quartus II Help. For information about using the Chip Planner tool, refer to *About the Chip Planner* in Quartus II Help. For details about using the Chip Planner tool, refer to the *Analyzing and Optimizing the Design Floorplan with the Chip Planner* chapter of the Quartus II Handbook.

### Optimizing Resource Utilization (LUT-Based Devices)

After design analysis, the next stage of design optimization is to improve resource utilization. Complete this stage before proceeding to I/O timing optimization or register-to-register timing optimization. Ensure that you have already set the basic constraints “Initial Compilation: Required Settings” section in the **Design Optimization Overview** chapter of the *Quartus II Handbook* before proceeding with the resource utilization optimizations described in this section. If a design does not fit into a specified device, use the techniques in this section to achieve a successful fit.

After you optimize resource utilization and your design fits in the desired target device, optimize I/O timing as described in the *I/O Timing Optimization Techniques (LUT-Based Devices)* section in the **Timing Closure and Optimization** chapter of the Quartus II Handbook. These tips are valid for all FPGA families and the MAX® II family of CPLDs.
Using the Resource Optimization Advisor

The Resource Optimization Advisor provides guidance in determining settings that optimize resource usage. To run the Resource Optimization Advisor, on the Tools menu, point to Advisors, and click Resource Optimization Advisor.

The Resource Optimization Advisor provides step-by-step advice about how to optimize resource usage (logic element, memory block, DSP block, I/O, and routing) of your design. Some of the recommendations in these categories might conflict with each other. Altera recommends evaluating the options and choosing the settings that best suit your requirements.

For more information about the Resource Optimization Advisor, refer to Resource Optimization Advisor Command (Tools Menu) in Quartus II Help.

Resolving Resource Utilization Issues Summary

Resource utilization issues can be divided into the following three categories:

- Issues relating to I/O pin utilization or placement, including dedicated I/O blocks such as PLLs or LVDS transceivers (refer to “I/O Pin Utilization or Placement”).
- Issues relating to logic utilization or placement, including logic cells containing registers and LUTs as well as dedicated logic, such as memory blocks and DSP blocks (refer to “Logic Utilization or Placement” on page 14–4).
- Issues relating to routing (refer to “Routing” on page 14–9).

I/O Pin Utilization or Placement

Use the suggestions in the following sections to help you resolve I/O resource problems.

Use I/O Assignment Analysis

To help with pin placement, on the Processing menu, point to Start and click Start I/O Assignment Analysis. The Start I/O Assignment Analysis command allows you to check your I/O assignments early in the design process. You can use this command to check the legality of pin assignments before, during, or after compilation of your design. If design files are available, you can use this command to accomplish more thorough legality checks on your design’s I/O pins and surrounding logic. These checks include proper reference voltage pin usage, valid pin location assignments, and acceptable mixed I/O standards.

Common issues with I/O placement relate to the fact that differential standards have specific pin pairings and certain I/O standards might be supported only on certain I/O banks.

If your compilation or I/O assignment analysis results in specific errors relating to I/O pins, follow the recommendations in the error message. Right-click the message in the Messages window and click Help to open the Quartus II Help topic for this message.
Modify Pin Assignments or Choose a Larger Package

If a design that has pin assignments fails to fit, compile the design without the pin assignments to determine whether a fit is possible for the design in the specified device and package. You can use this approach if a Quartus II error message indicates fitting problems due to pin assignments.

If the design fits when all pin assignments are ignored or when several pin assignments are ignored or moved, you might have to modify the pin assignments for the design or select a larger package.

If the design fails to fit because insufficient I/Os pins are available, a successful fit can often be obtained by using a larger device package (which can be the same device density) that has more available user I/O pins.

For more information about I/O assignment analysis, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

Logic Utilization or Placement

Use the suggestions in the following sections to help you resolve logic resource problems, including logic cells containing registers and LUTs, as well as dedicated logic such as memory blocks and DSP blocks.

Optimize Source Code

If your design does not fit because of logic utilization, evaluate and modify the design at the source to achieve the desired results. You can often improve logic significantly by making design-specific changes to your source code. This is typically the most effective technique for improving the quality of your results.

If your design does not fit into available logic elements (LEs) or ALMs, but you have unused memory or DSP blocks, check if you have code blocks in your design that describe memory or DSP functions that are not being inferred and placed in dedicated logic. You might be able to modify your source code to allow these functions to be placed into dedicated memory or DSP resources in the target device.

Ensure that your state machines are recognized as state machine logic and optimized appropriately in your synthesis tool. State machines that are recognized are generally optimized better than if the synthesis tool treats them as generic logic. In the Quartus II software, you can check for the State Machine report under Analysis & Synthesis in the Compilation Report. This report provides details, including the state encoding for each state machine that was recognized during compilation. If your state machine is not being recognized, you might have to change your source code to enable it to be recognized.

For coding style guidelines, including examples of HDL code for inferring memory and DSP functions, refer to the “Instantiating Altera Megafuntions” and the “Inferring Multiplier and DSP Functions from HDL Code” sections of the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook. For guidelines and sample HDL code for state machines, refer to the “General Coding Guidelines” section of the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
For additional HDL coding examples, refer to *AN 584: Timing Closure Methodology for Advanced FPGA Designs*.

**Optimize Synthesis for Area, Not Speed**

If your design fails to fit because it uses too much logic, resynthesize the design to improve the area utilization. First, ensure that you have set your device and timing constraints correctly in your synthesis tool. Particularly when area utilization of the design is a concern, ensure that you do not over-constrain the timing requirements for the design. Synthesis tools generally try to meet the specified requirements, which can result in higher device resource usage if the constraints are too aggressive.

If resource utilization is an important concern, some synthesis tools offer an easy way to optimize for area instead of speed. If you are using Quartus II integrated synthesis, select Balanced or Area for the Optimization Technique. You can also specify an Optimization Technique logic option for specific modules in your design with the Assignment Editor in cases where you want to reduce area using the Area setting (potentially at the expense of register-to-register timing performance) while leaving the default Optimization Technique setting at Balanced (for the best trade-off between area and speed for certain device families) or Speed. You can also use the Speed Optimization Technique for Clock Domains logic option to specify that all combinational logic in or between the specified clock domain(s) is optimized for speed.

In some synthesis tools, not specifying an f\(_{\text{MAX}}\) requirement can result in less resource utilization.

In the Quartus II software, the Balanced setting typically produces utilization results that are very similar to those produced by the Area setting, with better performance results. The Area setting can give better results in some cases.

For information about setting the timing requirements and synthesis options in Quartus II integrated synthesis and other synthesis tools, refer to the appropriate chapter in *Synthesis* in volume 1 of the *Quartus II Handbook*, or your synthesis software’s documentation.

The Quartus II software provides additional attributes and options that can help improve the quality of your synthesis results.

**Restructure Multiplexers**

Multiplexers form a large portion of the logic utilization in many FPGA designs. By optimizing your multiplexed logic, you can achieve a more efficient implementation in your Altera device.

For more information about this option, refer to *Restructure Multiplexers logic option* in Quartus II Help.

For design guidelines to achieve optimal resource utilization for multiplexer designs, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*. 
Perform WYSIWYG Primitive Resynthesis with Balanced or Area Setting

The Perform WYSIWYG Primitive Resynthesis logic option specifies whether to perform WYSIWYG primitive resynthesis during synthesis. This option uses the setting specified in the Optimization Technique logic option. The Perform WYSIWYG Primitive Resynthesis logic option is useful for resynthesizing some or all of the WYSIWYG primitives in your design for better area or performance. However, WYSIWYG primitive resynthesis can be done only when you use third-party synthesis tools.

The Balanced setting typically produces utilization results that are very similar to the Area setting with better performance results. The Area setting can give better results in some cases. Performing WYSIWYG resynthesis for area in this way typically reduces register-to-register timing performance.

For information about this logic option, refer to Perform WYSIWYG Primitive Resynthesis logic option in Quartus II Help.

Use Register Packing

The Auto Packed Registers option implements the functions of two cells into one logic cell by combining the register of one cell in which only the register is used with the LUT of another cell in which only the LUT is used.

For more information, refer to Auto Packed Registers logic option in Quartus Help.

Remove Fitter Constraints

A design with conflicting constraints or constraints that are difficult to meet may not fit in the targeted device. For example, a design might fail to fit if the location or LogicLock assignments are too strict and not enough routing resources are available on the device.

To resolve routing congestion caused by restrictive location constraints or LogicLock region assignments, use the Routing Congestion task in the Chip Planner to locate routing problems in the floorplan, then remove any internal location or LogicLock region assignments in that area. If your design still does not fit, the design is over-constrained. To correct the problem, remove all location and LogicLock assignments and run successive compilations, incrementally constraining the design before each compilation. You can delete specific location assignments in the Assignment Editor or the Chip Planner. To remove LogicLock assignments in the Chip Planner, in the LogicLock Regions Window, or on the Assignments menu, click Remove Assignments. Turn on the assignment categories you want to remove from the design in the Available assignment categories list.

For more information about the Routing Congestion task in the Chip Planner, refer to Analyzing and Optimizing the Design Floorplan with the Chip Planner of the Quartus II Handbook.
Flatten the Hierarchy During Synthesis

Synthesis tools typically provide the option of preserving hierarchical boundaries, which can be useful for verification or other purposes. However, the Quartus II software optimizes across hierarchical boundaries so as to perform the most logic minimization, which can reduce area in a design with no design partitions.

If you are using Quartus II incremental compilation, you cannot flatten your design across design partitions. Incremental compilation always preserves the hierarchical boundaries between design partitions, and the synthesis does not flatten it across partitions. Follow Altera’s recommendations for design partitioning, such as registering partition boundaries to reduce the effect of cross-boundary optimizations.

For more information about using incremental compilation and recommendations for design partitioning, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design and Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapters in volume 1 of the Quartus II Handbook.

Retarget Memory Blocks

If your design fails to fit because it runs out of device memory resources, your design may require a certain type of memory that the device does not have. For example, a design that requires two M-RAM blocks cannot be targeted to a Stratix® EP1S10 device, which has only one M-RAM block. You might be able to obtain a fit by building one of the memories with a different size memory block, such as an M4K memory block.

If the memory block was created with the MegaWizard™ Plug-In Manager, open the MegaWizard Plug-In Manager and edit the RAM block type so it targets a new memory block size.

ROM and RAM memory blocks can also be inferred from your HDL code, and your synthesis software can place large shift registers into memory blocks by inferring the ALTSHIFT_TAPS megafunction. This inference can be turned off in your synthesis tool to cause the memory or shift registers to be placed in logic instead of in memory blocks. Also, for improved timing performance, you can turn this inference off to prevent registers from being moved into RAM.

For more information, refer to Auto RAM Replacement logic option, Auto ROM Replacement logic option, and Auto Shift Register Replacement logic option in Quartus II Help.

Depending on your synthesis tool, you can also set the RAM block type for inferred memory blocks. In Quartus II integrated synthesis, set the ramstyle attribute to the desired memory type for the inferred RAM blocks, or set the option to logic, to implement the memory block in standard logic instead of a memory block.
Consider the Resource Utilization by Entity report in the report file and determine whether there is an unusually high register count in any of the modules. Some coding styles can prevent the Quartus II software from inferring RAM blocks from the source code because of the blocks’ architectural implementation, and force the software to implement the logic in flipflops. As an example, a function such as an asynchronous reset on a register bank might make the resistor bank incompatible with the RAM blocks in the device architecture, so that the register bank is implemented in flipflops. It is often possible to move a large register bank into RAM by slight modification of associated logic.

For more information about memory inference control in other synthesis tools, refer to the appropriate chapter in *Synthesis* in volume 1 of the *Quartus II Handbook*, or your synthesis software’s documentation. For more information about coding styles and HDL examples that ensure memory inference, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

**Use Physical Synthesis Options to Reduce Area**

The physical synthesis options for fitting help you decrease resource usage. When you enable these options, the Quartus II software makes placement-specific changes to the netlist that reduce resource utilization for a specific Altera device.

The compilation time might increase considerably when you use physical synthesis options.

With the Quartus II software, you can apply physical synthesis options to specific instances, which can reduce the impact on compilation time. Physical synthesis instance assignments allow you to enable physical synthesis algorithms for specific portions of your design.

The following physical synthesis optimizations for fitting are available:

- Physical synthesis for combinational logic
- Map logic into memory

For more information, refer to *Physical Synthesis Optimizations Page (Settings Dialog Box)* in Quartus II Help.

**Retarget or Balance DSP Blocks**

A design might not fit because it requires too many DSP blocks. You can implement all DSP block functions with logic cells, so you can retarget some of the DSP blocks to logic to obtain a fit.

If the DSP function was created with the MegaWizard Plug-In Manager, open the MegaWizard Plug-In Manager and edit the function so it targets logic cells instead of DSP blocks. The Quartus II software uses the `DEDICATED_MULTIPLIER_CIRCUITRY` megafUNCTION parameter to control the implementation.
DSP blocks also can be inferred from your HDL code for multipliers, multiply-adders, and multiply-accumulators. You can turn off this inference in your synthesis tool. When you are using Quartus II integrated synthesis, you can disable inference by turning off the Auto DSP Block Replacement logic option for your entire project. On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, click More Settings, and turn off Auto DSP Block Replacement. Alternatively, you can disable the option for a specific block with the Assignment Editor.

For more information about disabling DSP block inference in other synthesis tools, refer to the appropriate chapter in Synthesis in volume 1 of the Quartus II Handbook, or your synthesis software’s documentation.

The Quartus II software also offers the DSP Block Balancing logic option, which implements DSP block elements in logic cells or in different DSP block modes. The default Auto setting allows DSP block balancing to convert the DSP block slices automatically as appropriate to minimize the area and maximize the speed of the design. You can use other settings for a specific node or entity, or on a project-wide basis, to control how the Quartus II software converts DSP functions into logic cells and DSP blocks. Using any value other than Auto or Off overrides the DEDICATED_MULTIPLIER_CIRCUITRY parameter used in megafunction variations.

For more details about the Quartus II logic options described in this section, refer to Auto DSP Block Replacement logic option and DSP Block Balancing logic option in Quartus II Help.

**Use a Larger Device**

If a successful fit cannot be achieved because of a shortage of LEs, ALMs, memory, or DSP blocks, you might require a larger device.

**Routing**

Use the suggestions in the following sections to help you resolve routing resource problems.

**Set Auto Packed Registers to Sparse or Sparse Auto**

The Auto Packed Registers option reduces LE or ALM count in a design. You can set this option in the Assignment Editor or by clicking More Settings on the Fitter Settings page in the Settings dialog box.

For more information, refer to Auto Packed Registers logic option in Quartus II Help.
Set Fitter Aggressive Routability Optimizations to Always

The Fitter Aggressive Routability Optimization option is useful if your design does not fit due to excessive routing wire utilization.

If there is a significant imbalance between placement and routing time (during the first fitting attempt), it might be because of high wire utilization. Turning on the Fitter Aggressive Routability Optimizations option can reduce your compilation time.

On average, this option can save up to 6% wire utilization, but can also reduce performance by up to 4%, depending on the device.

For more information, refer to Fitter Aggressive Routability Optimizations logic option in Quartus II Help.

Increase Router Effort Multiplier

The Router Effort Multiplier controls how quickly the router tries to find a valid solution. The default value is 1.0 and legal values must be greater than 0. Numbers higher than 1 help designs that are difficult to route by increasing the routing effort. Numbers closer to 0 (for example, 0.1) can reduce router runtime, but usually reduce routing quality slightly. Experimental evidence shows that a multiplier of 3.0 reduces overall wire usage by approximately 2%. Using a Router Effort Multiplier higher than the default value could be beneficial for designs with complex datapaths with more than five levels of logic. However, congestion in a design is primarily due to placement, and increasing the Router Effort Multiplier does not necessarily reduce congestion.

Any Router Effort Multiplier value greater than 4 only increases by 10% for every additional 1. For example, a value of 10 is actually 4.6.

For more information, refer to Router Effort Multiplier logic option in Quartus II Help.

Remove Fitter Constraints

A design with conflicting constraints or constraints that are difficult to achieve may not fit the targeted device. Conflicting or difficult-to-achieve constraints can occur when location or LogicLock assignments are too strict and there are not enough routing resources.

In this case, use the Routing Congestion task in the Chip Planner to locate routing problems in the floorplan, then remove all location and LogicLock region assignments from that area. If the local constraints are removed, and the design still does not fit, the design is over-constrained. To correct the problem, remove all location and LogicLock assignments and run successive compilations, incrementally constraining the design before each compilation. You can delete specific location assignments in the Assignment Editor or the Chip Planner. To remove LogicLock assignments in the Chip Planner, in the LogicLock Regions Window, or on the Assignments menu, click Remove Assignments. Turn on the assignment categories you want to remove from the design in the Available assignment categories list.
For more information about the **Routing Congestion** task in the Chip Planner, refer to the *Analyzing and Optimizing the Design Floorplan with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*. You can also refer to *About the Chip Planner* in Quartus II Help.

**Optimize Synthesis for Area, Not Speed**

In some cases, resynthesizing the design to improve the area utilization can also improve the routability of the design. First, ensure that you have set your device and timing constraints correctly in your synthesis tool. Ensure that you do not overconstrain the timing requirements for the design, particularly when the area utilization of the design is a concern. Synthesis tools generally try to meet the specified requirements, which can result in higher device resource usage if the constraints are too aggressive.

If resource utilization is important to improve the routing results in your design, some synthesis tools offer an easy way to optimize for area instead of speed. If you are using Quartus II integrated synthesis, on the Assignments menu, click **Settings**. In the **Category** list, select **Analysis & Synthesis Settings** and select **Balanced** or **Area** under **Optimization Technique**.

You can also specify this logic option for specific modules in your design with the Assignment Editor in cases where you want to reduce area using the **Area** setting (potentially at the expense of register-to-register timing performance). You can apply the setting to specific modules while leaving the default **Optimization Technique** setting at **Balanced** (for the best trade-off between area and speed for certain device families) or **Speed**. You can also use the **Speed Optimization Technique for Clock Domains** logic option to specify that all combinational logic in or between the specified clock domain(s) is optimized for speed.

In the Quartus II software, the **Balanced** setting typically produces utilization results that are very similar to those obtained with the **Area** setting, with better performance results. The **Area** setting can yield better results in some unusual cases.

In some synthesis tools, not specifying an $f_{\text{MAX}}$ requirement can result in less resource utilization, which can improve routability.

For information about setting the timing requirements and synthesis options in Quartus II integrated synthesis and other synthesis tools, refer to the appropriate chapter in *Synthesis* in volume 1 of the *Quartus II Handbook*, or your synthesis software’s documentation.
Optimize Source Code

If your design does not fit because of routing problems and the methods described in the preceding sections do not sufficiently improve the routability of the design, modify the design at the source to achieve the desired results. You can often improve results significantly by making design-specific changes to your source code, such as duplicating logic or changing the connections between blocks that require significant routing resources.

Use a Larger Device

If a successful fit cannot be achieved because of a shortage of routing resources, you might require a larger device.

Optimizing Resource Utilization (Macrocell-Based CPLDs)

The following recommendations help you take advantage of the macrocell-based architecture in the MAX 7000 and MAX 3000 devices to yield maximum speed, reliability, and device resource utilization while minimizing fitting difficulties.

After design analysis, the first stage of design optimization is to improve resource utilization. Complete this stage before proceeding to timing optimization. First, ensure that you have set the basic constraints described in “Initial Compilation: Required Settings” section in the Design Optimization Overview chapter of the Quartus II Handbook. If your design is not fitting into a specified device, use the techniques in this section to achieve a successful fit.

Use Dedicated Inputs for Global Control Signals

MAX 7000 and MAX 3000 devices have four dedicated inputs that you can use for global register control. Because the global register control signals can bypass the logic cell array and directly feed registers, product terms can be preserved for primary logic. Also, because each signal has a dedicated path into the LAB, global signals also can bypass logic and data path interconnect resources.

Because the dedicated input pins are designed for high fan-out control signals and provide low skew, always assign global signals (such as clock, clear, and output enable) to the dedicated input pins.

You can use logic-generated control signals for global control signals instead of dedicated inputs. However, the following list shows the disadvantages of using logic-generated control signals:

- More resources are required (logic cells, interconnect).
- More data skew is introduced.
- If the logic-generated control signals have high fan-out, the design can be more difficult to fit.

By default, the Quartus II software uses dedicated inputs for global control signals automatically. You can assign control signals to dedicated input pins in one of the following ways:
In the Assignment Editor, select one of the two following methods:

- Assign pins to dedicated pin locations.
- Assign a Global Signal setting to the pins.

On the Assignments menu, click Settings. On the Analysis & Synthesis Settings page, click More Settings, and in the Existing Option settings section, select Auto Global Register Control Signals.

Insert a GLOBAL primitive after the pins.

Reserve Device Resources

Because pin and logic option assignments can be necessary for board layout and performance requirements, and because full utilization of the device resources can increase the difficulty of fitting the design, Altera recommends leaving 10% of the logic cells and 5% of the I/O pins unused to accommodate future design modifications. Following the Altera-recommended device resource reservation guidelines for macrocell-based CPLDs increases the chance that the Quartus II software can fit the design during recompilation after changes or assignments have been made.

Pin Assignment Guidelines and Procedures

Sometimes user-specified pin assignments are necessary for board layout. This section describes pin assignment guidelines and procedures.

To minimize fitting issues with pin assignments, follow these guidelines:

- Assign speed-critical control signals to dedicated inputs.
- Assign output enables (OE) pin to appropriate locations.
- Estimate fan-in to assign output pins to the appropriate LAB.
- Assign output pins that require parallel expanders to macrocells numbered 4 to 16.

Altera recommends allowing the Quartus II software to select pin assignments automatically when possible. You can use the Quartus II Pin Advisor feature (accessible from the Tools menu) for pin connection guidelines.

For more information about the Pin Advisor, refer to Pin Advisor Command (Tools Menu) in Quartus II Help.

Control Signal Pin Assignments

Assign speed-critical control signals to dedicated input pins. Every MAX 7000 and MAX 3000 device has four dedicated input pins (GCLK1, OE2/GCLK2, OE1, and GCLRN). You can assign clocks to global clock dedicated inputs (GCLK1 and OE2/GCLK2), assign clear signals to the global clear dedicated input (GCLRN), and speed-critical OE signals to global OE dedicated inputs (OE1 and OE2/GCLK2).

Output Enable Pin Assignments

Occasionally, because the total number of required output enable pins is more than the dedicated input pins, output enable signals must be assigned to I/O pins.
To minimize possible fitting errors when assigning the output enable pins for MAX 7000 and MAX 3000 devices, refer to Pin-Out Files for Altera Devices on the Altera website (www.altera.com).

**Estimate Fan-In When Assigning Output Pins**

Macrocells with high fan-in can cause more placement problems for the Quartus II Fitter than those with low fan-in. The maximum fan-in per LAB should not exceed 36 in MAX 7000 and MAX 3000 devices. Therefore, estimate the fan-in of logic (such as an \(x\)-input AND gate) that feeds each output pin. If the total fan-in of logic that feeds each output pin in the same LAB exceeds 36, compilation can fail. To save resources and prevent compilation errors, avoid assigning pins that have high fan-in.

**Outputs Using Parallel Expander Pin Assignments**

Figure 14–1 shows how parallel expanders are used within a LAB. MAX 7000 and MAX 3000 devices contain chains that can lend or borrow parallel expanders. The Quartus II Fitter places macrocells in a location that allows them to lend and borrow parallel expanders appropriately.

As shown in Figure 14–1, only macrocells 2 through 16 can borrow parallel expanders. Therefore, assign output pins that might require parallel expanders to pins adjacent to macrocells 4 through 16. Altera recommends using macrocells 4 through 16 because they can borrow the largest number of parallel expanders.

![Figure 14–1. LAB Macrocells and Parallel Expander Associations](image-url)
Resolving Resource Utilization Problems

Excessive macrocell usage and lack of routing resources can cause resource utilization problems. Macrocell usage errors occur when the total number of macrocells in the design exceed the available macrocells in the device. Routing errors occur when the available routing resources are insufficient to implement the design. Check the Message window for the compilation results.

Messages in the Messages window are also copied in the Report Files. For more information about the message, right-click a message and click Help.

Resolving Macrocell Usage Issues

Occasionally, a design requires more macrocell resources than are available in the selected device, which results in the design not fitting. The following list provides tips for resolving macrocell usage issues as well as tips to minimize the number of macrocells used:

- On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, click More Settings, and turn off Auto Parallel Expanders. If the design’s clock frequency (f_{\text{MAX}}) is not an important design requirement, turn off parallel expanders for all or part of the project. The design usually requires more macrocells if parallel expanders are turned on.

- Change Optimization Technique from Speed to Area. Selecting Area instructs the compiler to give preference to area utilization rather than speed (f_{\text{MAX}}). On the Assignments menu, click Settings. In the Category list, change the Optimization Technique option in the Analysis & Synthesis Settings page.

- Use D-type flipflops instead of latches. Altera recommends always using D-type flipflops instead of latches in your design because D-type flipflops can reduce the macrocell fan-in, and thus reduce macrocell usage. The Quartus II software uses extra logic to implement latches in MAX 7000 and MAX 3000 designs because MAX 7000 and MAX 3000 macrocells contain D-type flipflops instead of latches.

- Use asynchronous clear and preset instead of synchronous clear and preset. To reduce product term usage, use asynchronous clear and preset in your design whenever possible. Using other control signals such as synchronous clear produces macrocells and pins with higher fan-out.

After following the suggestions in this section, if your project still does not fit the targeted device, consider using a larger device. When upgrading to a different density, the vertical package-migration feature of the MAX 7000 and MAX 3000 device families allows pin assignments to be maintained.
Resolving Routing Issues

Routing is another resource that can cause design fitting issues. For example, if the total fan-in into a LAB exceeds the maximum allowed, a no-fit error can occur during compilation. If your design does not fit the targeted device because of routing issues, consider the following suggestions:

- Use dedicated inputs/global signals for high fan-out signals. The dedicated inputs in MAX 7000 and MAX 3000 devices are designed for speed-critical and high fan-out signals. Always assign high fan-out signals to dedicated inputs/global signals.

- Change the Optimization Technique option from Speed to Area. This option can resolve routing resource and macrocell usage issues. Refer to “Resolving Macrocell Usage Issues” on page 14–15.

- Reduce the fan-in per cell. If you are not limited by the number of macrocells used in the design, you can use the Fan-in per cell (%) option to reduce the fan-in per cell. The allowable values are 20–100%; the default value is 100%. Reducing fan-in can reduce localized routing congestion but increase the macrocell count. You can set this logic option in the Assignment Editor or under More Settings in the Analysis & Synthesis Settings page of the Settings dialog box.

- On the Assignments menu, click Settings. In the Category list, select Analysis & Synthesis Settings, click More Options, and turn off Auto Parallel Expanders. By turning off the parallel expanders, you give the Quartus II software more fitting flexibility for each macrocell, allowing macrocells to be relocated. For example, you can move each macrocell (previously grouped together in the same LAB) to a different LAB to reduce routing constraints.

- Insert logic cells. Inserting logic cells reduces fan-in and shared expanders used per macrocell, increasing routability. By default, the Quartus II software automatically inserts logic cells when necessary. Otherwise, you can disable Auto Logic Cell as follows:
  1. On the Assignments menu, click Settings.
  2. In the Category list, select Analysis & Synthesis Settings.
  3. Under More Settings, turn off Auto Logic Cell Insertion. For more information, refer to “Using LCELL Buffers to Reduce Required Resources”.

- Change pin assignments. If you want to discard your pin assignments, you can let the Quartus II Fitter ignore some or all of the assignments.

  If you prefer reassigning pins to increase routing efficiency, refer to “Pin Assignment Guidelines and Procedures” on page 14–13.

Using LCELL Buffers to Reduce Required Resources

Complex logic, such as multilevel XOR gates, are often implemented with more than one macrocell. When this occurs, the Quartus II software automatically allocates shareable expanders—or additional macrocells (called synthesized logic cells)—to supplement the logic resources that are available in a single macrocell. You can also break down complex logic by inserting logic cells in the project to reduce the average fan-in and the total number of shareable expanders required. Manually inserting logic cells can provide greater control over speed-critical paths.
Instead of using the **Auto Logic Cell Insertion** option, you can manually insert logic cells. However, Altera recommends using the **Auto Logic Cell Insertion** option unless you know which part of the design is causing the congestion.

A good location to manually insert LCELL buffers is where a single complex logic expression feeds multiple destinations in your design. You can insert an LCELL buffer just after the complex expression; the Fitter extracts this complex expression and places it in a separate logic cell. Rather than duplicate all the logic for each destination, the Quartus II software feeds the single output from the logic cell to all destinations.

To reduce fan-in and prevent no-fit compilations caused by routing resource issues, insert an LCELL buffer after a NOR gate (Figure 14–2). The design in Figure 14–2 was compiled for a MAX 7000AE device. Without the LCELL buffer, the design requires two macrocells and eight shareable expanders, and the average fan-in is 14.5 macrocells. However, with the LCELL buffer, the design requires three macrocells and eight shareable expanders, and the average fan-in is just 6.33 macrocells.

**Figure 14–2. Reducing the Average Fan-In by Inserting LCELL Buffers**

---

**Scripting Support**

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II command-line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```
For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook. For more information about all settings and constraints in the Quartus II software, refer to the Quartus II Settings File Manual. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

You can specify many of the options described in this section either in an instance, or at a global level, or both.

Use the following Tcl command to make a global assignment:

```tcl
set_global_assignment -name <qsf variable name> <value>
```

Use the following Tcl command to make an instance assignment:

```tcl
set_instance_assignment -name <qsf variable name> <value> -to <instance name>
```

If the <value> field includes spaces (for example, ‘Standard Fit’), you must enclose the value in straight double quotation marks.

### Initial Compilation Settings

Use the Quartus II Settings File (.qsf) variable name in the Tcl assignment to make the setting along with the appropriate value. The Type column indicates whether the setting is supported as a global setting, an instance setting, or both.

Table 14–1 lists the advanced compilation settings.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Placement Effort Multiplier</td>
<td>PLACEMENT_EFFORT_MULTIPLIER</td>
<td>Any positive, non-zero value</td>
<td>Global</td>
</tr>
<tr>
<td>Router Effort Multiplier</td>
<td>ROUTER_EFFORT_MULTIPLIER</td>
<td>Any positive, non-zero value</td>
<td>Global</td>
</tr>
<tr>
<td>Router Timing Optimization level</td>
<td>ROUTER_TIMING_OPTIMIZATION_LEVEL</td>
<td>NORMAL, MINIMUM, MAXIMUM</td>
<td>Global</td>
</tr>
<tr>
<td>Final Placement Optimization</td>
<td>FINAL_PLACEMENT_OPTIMIZATION</td>
<td>ALWAYS, AUTOMATICALLY, NEVER</td>
<td>Global</td>
</tr>
</tbody>
</table>

### Resource Utilization Optimization Techniques (LUT-Based Devices)

Table 14–2 lists the .qsf file variable name and applicable values for the settings described in “Optimizing Resource Utilization (LUT-Based Devices)” on page 14–2.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Auto Packed Registers</td>
<td>AUTO_PACKED_REGISTERS_&lt;device family name&gt;</td>
<td>OFF, NORMAL, MINIMIZE AREA, MINIMIZE AREA WITH CHAINS, AUTO</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Perform WYSIWYG Primitive Resynthesis</td>
<td>ADV_NETLIST_OPT_SYNTH_WYSIWYG_REMAP</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
</tbody>
</table>
Table 14–2. Resource Utilization Optimization Settings (Part 2 of 2)

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>.qsf File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Physical Synthesis for Combinational Logic for Reducing Area</td>
<td>PHYSICAL_SYNTHESIS_COMBO_LOGIC_FOR_AREA</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Physical Synthesis for Mapping Logic to Memory</td>
<td>PHYSICAL_SYNTHESIS_MAP_LOGIC_TO_MEMORY_FOR_AREA</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Optimization Technique</td>
<td>&lt;device family name&gt;_OPTIMIZATION_TECHNIQUE</td>
<td>AREA, SPEED, BALANCED</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Speed Optimization Technique for Clock Domains</td>
<td>SYNTH_CRITICAL_CLOCK</td>
<td>ON, OFF</td>
<td>Instance</td>
</tr>
<tr>
<td>State Machine Encoding</td>
<td>STATE_MACHINE_PROCESSING</td>
<td>AUTO, ONE-HOT, GRAY, JOHNSON, MINIMAL BITS, ONE-HOT, SEQUENTIAL, USER-ENCODING</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto RAM Replacement</td>
<td>AUTO_RAM_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto ROM Replacement</td>
<td>AUTO_ROM_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto Shift Register Replacement</td>
<td>AUTO_SHIFT_REGISTER_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Auto Block Replacement</td>
<td>AUTO_DSP_RECOGNITION</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Number of Processors for Parallel Compilation</td>
<td>NUM_PARALLEL_PROCESSORS</td>
<td>Integer between 1 and 16 inclusive, or ALL</td>
<td>Global</td>
</tr>
</tbody>
</table>

Note to Table 14–2:
(1) Allowed values for this setting depend on the device family that you select.

Document Revision History

Table 14–3 lists the revision history for this document.

Table 14–3. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>
15. Analyzing and Optimizing the Design Floorplan with the Chip Planner

As FPGA designs grow larger in density, the ability to analyze the design for performance, routing congestion, and logic placement to meet the design requirements becomes critical. This chapter discusses how to analyze the design floorplan with the Chip Planner.

Design floorplan analysis is a valuable method for achieving timing closure and optimal performance in highly complex designs. With analysis capability, the Quartus II Chip Planner helps you close timing quickly on your designs. Using the Chip Planner together with LogicLock and Incremental Compilation enables you to compile your designs hierarchically, preserving the timing results from individual compilation runs. You can use LogicLock regions as part of an incremental compilation methodology to improve your productivity.

You can perform design analysis, as well as creating and optimizing the design floorplan with the Chip Planner. To make I/O assignments, use the Pin Planner.

For information about the Pin Planner, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

You can use the Design Partition Planner with the Chip Planner to customize the floorplan of your design. For more information, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design and the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapters in volume 1 of the Quartus II Handbook.

This chapter includes the following topics:

■ “Chip Planner Overview”
■ “LogicLock Regions” on page 15–3
■ “Using LogicLock Regions in the Chip Planner” on page 15–11
■ “Design Floorplan Analysis Using the Chip Planner” on page 15–12
■ “Scripting Support” on page 15–21

For a list of devices supported by the Chip Planner, refer to About the Chip Planner in Quartus II Help.

For more information about the Chip Planner, refer to the Altera Training page of the Altera website.
Chip Planner Overview

The Chip Planner provides a visual display of chip resources. The Chip Planner can show logic placement, LogicLock regions, relative resource usage, detailed routing information, fan-in and fan-out connections between nodes, timing paths between registers, delay estimates for paths, and routing congestion information.

You can also make assignment changes with the Chip Planner, such as creating and deleting resource assignments, and you can perform post-compilation changes such as creating, moving, and deleting logic cells and I/O atoms. With the Chip Planner, you can view and create assignments for a design floorplan, perform power and design analyses, and implement ECOs. With the Chip Planner and Resource Property Editor, you can change connections between resources and make post-compilation changes to the properties of logic cells, I/O elements, PLLs, and RAM and digital signal processing (DSP) blocks.

For details about how to implement ECOs in your design using the Chip Planner in the Quartus II software, refer to the *Engineering Change Management with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*.

Starting the Chip Planner

To start the Chip Planner, on the Tools menu, click **Chip Planner (Floorplan & Chip Editor)**. You can also start the Chip Planner by the following methods:

- Click the Chip Planner icon on the Quartus II software toolbar
- On the Shortcut menu in the following tools, click **Locate** and then click **Locate in Chip Planner (Floorplan and Chip Editor)**:
  - Design Partition Planner
  - Compilation Report
  - LogicLock Regions window
  - Technology Map Viewer
  - Project Navigator window
  - RTL source code
  - Node Finder
  - Simulation Report
  - RTL Viewer
  - Report Timing panel of the TimeQuest Timing Analyzer

Chip Planner Toolbar

The Chip Planner provides powerful tools for design analysis with a GUI. You can access Chip Planner commands from the View menu and the Shortcut menu, or by clicking the icons on the toolbar.
Chip Planner Presets, Layers, and Editing Modes

The Chip Planner models types of resource objects as unique display layers, and uses presets— which are predefined sets of layer settings—to control the display of resources. The Chip Planner provides a set of default presets, and you can create custom presets to customize the display for your particular needs. The Basic, Detailed, and Floorplan Editing presets provided with the Chip Planner are useful for general ECO and assignment-related activities, while the Design Partition Planner preset is optimized for specific activities.

The Chip Planner has two editing modes, which determine the types of operations that you can perform. The Assignment editing mode allows you to make assignment changes that are applied by the Fitter during the next place and route operation. The ECO editing mode allows you to make post-compilation changes, commonly referred to as engineering change orders (ECOs).

You should choose the editing mode appropriate for the work that you want to perform, and a preset that displays the resources that you want to view, in a level of detail appropriate for your design.

Locate History Window

As you optimize your design floorplan, you might have to locate a path or node in the Chip Planner many times. The Locate History window lists all the nodes and paths you have displayed using a Locate in Chip Planner (Floorplan and Chip Editor) command, providing easy access to the nodes and paths of interest to you. If you locate a required path from the TimeQuest Timing Analyzer Report Timing pane, the Locate History window displays the required clock path. If you locate an arrival path from the TimeQuest Timing Analyzer Report Timing pane, the Locate History window displays the path from the arrival clock to the arrival data. Double-clicking a node or path in the Locate History window displays the selected node or path in the Chip Planner.

For more information about the Chip Planner, refer to About the Chip Planner and Layers Settings Dialog Box in Quartus II Help. For more information about the ECO editing mode, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

LogicLock Regions

LogicLock regions are floorplan location constraints that help you place logic on the target device. When you assign entity instances or nodes to a LogicLock region, you direct the Fitter to place those entity instances or nodes within the region during fitting. Your floorplan can contain several LogicLock regions.
A LogicLock region is defined by its height, width, and location; you can specify the size or location of a region, or both, or the Quartus II software can generate these properties automatically. The Quartus II software bases the size and location of a region on the contents of the region and the timing requirements of the module. Table 15–1 describes the options for creating LogicLock regions.

### Table 15–1. Types of LogicLock Regions

<table>
<thead>
<tr>
<th>Property</th>
<th>Value</th>
<th>Behavior</th>
</tr>
</thead>
<tbody>
<tr>
<td>State</td>
<td>Floating (1), Locked</td>
<td>Floating allows the Quartus II software to determine the location of the region on the device. Floating regions are shown with a dashed boundary in the floorplan. Locked allows you to specify the location of the region. Locked regions are shown with a solid boundary in the floorplan. A locked region must have a fixed size.</td>
</tr>
<tr>
<td>Size</td>
<td>Auto (1), Fixed</td>
<td>Auto allows the Quartus II software to determine the appropriate size of a region given its contents. Fixed regions have a shape and size that you define.</td>
</tr>
<tr>
<td>Reserved</td>
<td>Off (1), On</td>
<td>Allows you to define whether the Fitter can use the resources within a region for entities that are not assigned to the region. If the reserved property is turned on, only items assigned to the region can be placed within its boundaries.</td>
</tr>
<tr>
<td>Origin</td>
<td>Any Floorplan Location</td>
<td>Specifies the location of the LogicLock region on the floorplan. For Arria series, Stratix series, Cyclone series, MAX II, and MAX V devices, the origin is located at the lower left corner of the LogicLock region. For other Altera® device families, the origin is located at the upper left corner of the LogicLock region.</td>
</tr>
</tbody>
</table>

**Note to Table 15–1:**

(1) Default value.

The Quartus II software cannot automatically define the size of a region if the location is locked. Therefore, if you want to specify the exact location of the region, you must also specify the size.

You can use the Design Partition Planner in conjunction with LogicLock regions to create a floorplan for your design. For more information about using the Design Partition Planner, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Designs* and the *Best Practices for Incremental Compilation Partition and Floorplan Assignments* chapters in volume 1 of the *Quartus II Handbook*.

### Creating LogicLock Regions

You can create LogicLock Regions with the Project Navigator, the LogicLock Regions window, the Design Partition Planner, the Chip Planner, and with Tcl commands.

#### Creating LogicLock Regions with the Project Navigator

After you perform either a full compilation or analysis and elaboration on the design, the Quartus II software displays the hierarchy of the design. On the View menu, **Utility Windows**, then **Project Navigator**. With the hierarchy of the design fully expanded, right-click on any design entity in the design, and click **Create New LogicLock Region** to create a LogicLock region and assign the entity to the new region.
Creating LogicLock Regions with the LogicLock Regions window

To create a LogicLock region with the LogicLock Regions window, on the Assignments menu, click **LogicLock Regions Window**. In the LogicLock Regions window, click **<<new>>**.

Creating LogicLock Regions with the Design Partition Planner

To create a LogicLock region and assign a partition to it with the Design Partition Planner, right-click the partition and then click **Create LogicLock Region**.

Creating LogicLock Regions with the Chip Planner

To create a LogicLock region in the Chip Planner, click **LogicLock Regions** then **Create LogicLock Region** command on the View menu, then click and drag on the Chip Planner floorplan to create a region of your preferred location and size.

Creating Nonrectangular LogicLock Regions

When you create a floorplan for your design, you may want to create nonrectangular LogicLock regions to exclude certain resources from the LogicLock region. You might also create a nonrectangular LogicLock region to place certain parts of your design around specific device resources to improve performance.

To create a nonrectangular region with the **Merge LogicLock Region** command, follow these steps:

1. In the Chip Planner, create two or more contiguous or non-contiguous rectangular regions as described in “Creating LogicLock Regions” on page 15–4.
2. Arrange the regions that you have created into the locations where you want the nonrectangular region to be.
3. Select all the individual regions that you want to merge by clicking each of them while pressing the Shift key.
4. Right-click the title bar of any of the LogicLock regions that you want to merge, point to **LogicLock Regions**, and then click **Merge LogicLock Region**. The individual regions that you select merge to create a single new region.

By default, the new LogicLock region has the same name as the component region containing the greatest number of resources; however, you can rename the new region. In the **LogicLock Regions Window**, the new region is shown as having a **Custom Shape**.
Figure 15–1 illustrates using the **Merge LogicLock Region** command to form a nonrectangular LogicLock region by merging two rectangular LogicLock regions.

**Figure 15–1. Using the Merge LogicLock Region command to create a nonrectangular region**

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**Hierarchical (Parent and Child) LogicLock Regions**

To further constrain module locations, you can define a hierarchy for a group of regions by declaring parent and child regions. The Quartus II software places a child region completely within the boundaries of its parent region; a child region must be placed entirely within the boundary of its parent. Additionally, parent and child regions allow you to further improve the performance of a module by constraining nodes in the critical path of a module.

To make one LogicLock region a child of another LogicLock region, in the LogicLock Regions window, select the new child region and drag and drop the new child region into its new parent region.

The LogicLock region hierarchy does not have to be the same as the design hierarchy.

You can create both auto-sized and fixed-sized LogicLock regions within a parent LogicLock region; however, the parent of a fixed-sized child region must also be fixed-sized. The location of a locked parent region is locked relative to the device; the location of a locked child region is locked relative to its parent region. If you change the parent’s location, the locked child’s origin changes, but maintains the same placement relative to the origin of its parent. The location of a floating child region can float within its parent. Complex region hierarchies might result in some LABs not being used, effectively increasing the resource utilization in the device. Do not create more levels of hierarchy than you need.
Placing LogicLock Regions

A fixed region must contain all resources required by the design block assigned to the region. Although the Quartus II software can automatically place and size LogicLock regions to meet resource and timing requirements, you can manually place and size regions to meet your design requirements. You should consider the following if you manually place or size a LogicLock region:

- LogicLock regions with pin assignments must be placed on the periphery of the device, adjacent to the pins. For the Arria series, Cyclone series, MAX II, MAX V, and Stratix series of devices, you must also include the I/O block within the LogicLock Region.

- Floating LogicLock regions can overlap with their ancestors or descendants, but not with other floating LogicLock regions.

Placing Device Resources into LogicLock Regions

A LogicLock region includes all device resources within its boundaries, including memory and pins. The Quartus II software does not include pins automatically when you assign an entity to a region—you can manually assign pins to LogicLock regions; however, this placement puts location constraints on the region. The software only obeys pin assignments to locked regions that border the periphery of the device. For the Arria series, Cyclone series, MAX II, MAX V, and Stratix series of devices, the locked regions must include the I/O pins as resources.

Pin assignments to LogicLock regions are effective only in fixed and locked regions. Pin assignments to floating regions do not influence the placement of the region.

Only one LogicLock region can claim a device resource. If a LogicLock region boundary includes part of a device resource, the Quartus II software allocates the entire resource to that LogicLock region. When the Quartus II software places a floating auto-sized region, it places the region in an area that meets the requirements of the contents of the LogicLock region.

If you want to import multiple instances of a module into a top-level design, you must ensure that the device has two or more locations with exactly the same device resources. (You can determine this from the applicable device handbook.) If the device does not have another area with exactly the same resources, the Quartus II software generates a fitting error during compilation of the top-level design.
LogicLock Regions Window

You can use the LogicLock Regions window to create LogicLock regions, assign nodes and entities to them, and modify the properties of a LogicLock region such as size, state, width, height, origin, and whether the region is a reserved region. The LogicLock Regions window also has a recommendations toolbar; select a LogicLock region from the drop-down list in the recommendations toolbar to display the relevant suggestions to optimize that LogicLock region. You can customize the LogicLock Regions window by dragging and dropping the columns to change their order; you can also show and hide optional columns by right-clicking any column heading and then selecting the appropriate columns in the shortcut menu.

The LogicLock Region Properties dialog box provides a summary of all LogicLock regions in your design. Use the LogicLock Region Properties dialog box to obtain detailed information about your LogicLock region, such as which entities and nodes are assigned to your region and which resources are required. The LogicLock Region Properties dialog box shows the properties of the current selected regions and allows you to modify them. To open the LogicLock Region Properties dialog box, double-click any region in the LogicLock Regions window, or right-click the region and click Properties.

For designs that target Arria series, Cyclone series, Stratix series, MAX II, and MAX V devices, the Quartus II software automatically creates a LogicLock region that encompasses the entire device. This default region is labelled Root_Region, and is locked and fixed.

For Arria series, Cyclone series, Stratix series, MAX II, and MAX V devices, the origin of the LogicLock region is located at the lower-left corner of the region. For all other supported devices, the origin is located at the upper-left corner of the region.
Reserved LogicLock Region

The Quartus II software honors all entity and node assignments to LogicLock regions. Occasionally, entities and nodes do not occupy an entire region, which leaves some of the region’s resources unoccupied. To increase the region’s resource utilization and performance, the Quartus II software’s default behavior fills the unoccupied resources with other nodes and entities that have not been assigned to another region. You can prevent this behavior by turning on Reserved on the General tab of the LogicLock Region Properties dialog box. When you turn on this option, your LogicLock region contains only the entities and nodes that you specifically assigned to your LogicLock region.

Excluded Resources

The Excluded Resources feature allows you to easily exclude specific device resources such as DSP blocks or M4K memory blocks from a LogicLock region. For example, you can assign a specific entity to a LogicLock region but allow the DSP blocks of that entity to be placed anywhere on the device. Use the Excluded Resources feature on a per-LogicLock region member basis.

To exclude certain device resources from an entity, in the LogicLock Region Properties dialog box, highlight the entity in the Design Element column, and click Edit. In the Edit Node dialog box, under Excluded Element Types, click the Browse button. In the Excluded Resources Element Types dialog box, you can select the device resources you want to exclude from the entity. When you have selected the resources to exclude, the Excluded Resources column is updated in the LogicLock Region Properties dialog box to reflect the excluded resources.

The Excluded Resources feature prevents certain resource types from being included in a region, but it does not prevent the resources from being placed inside the region unless you set the region’s Reserved property to On. To indicate to the Fitter that certain resources are not required inside a LogicLock region, define a resource filter. For more information about resource filters, refer to “LogicLock Resource Exclusions” in the Best Practices for Incremental Compilation Partitions and Floorplan Assignments chapter in volume 1 of the Quartus II Handbook.

Additional Quartus II LogicLock Design Features

To complement the LogicLock Regions window, the Quartus II software has additional features to help you design with LogicLock regions.

Analysis and Synthesis Resource Utilization by Entity

The Compilation Report contains an Analysis and Synthesis Resource Utilization by Entity section, which reports resource usage statistics, including entity-level information. You can use this feature to verify that any LogicLock region you manually create contains enough resources to accommodate all the entities you assign to it.
Quartus II Revisions Feature

When you evaluate different LogicLock regions in your design, you might want to experiment with different configurations to achieve your desired results. The Quartus II Revisions feature allows you to organize the same project with different settings until you find an optimum configuration.

To use the Revisions feature, on the Project menu, click Revisions. In the Revisions dialog box, you can create and specify revisions. You can create a revision from the current design or any previously created revisions. Each revision can have an associated description. You can use revisions to organize the placement constraints created for your LogicLock regions.

LogicLock Assignment Precedence

You can encounter conflicts during the assignment of entities and nodes to LogicLock regions. For example, an entire top-level entity might be assigned to one region and a node within this top-level entity assigned to another region. To resolve conflicting assignments, the Quartus II software maintains an order of precedence for LogicLock assignments. The following order of precedence, from highest to lowest, applies:

1. Exact node-level assignments
2. Path-based and wildcard assignments
3. Hierarchical assignments

For more information about LogicLock assignment precedence, refer to Understanding Assignment Priority in Quartus II Help.

Open the Priority dialog box by selecting Priority on the General tab of the LogicLock Regions Properties dialog box. You can change the priority of path-based and wildcard assignments with the Up and Down buttons in the Priority dialog box. To prioritize assignments between regions, you must select multiple LogicLock regions and then open the Priority dialog box from the LogicLock Regions Properties dialog box.

Virtual Pins

A virtual pin is an I/O element that is temporarily mapped to a logic element and not to a pin during compilation, and is then implemented as a LUT. Virtual pins should be used only for I/O elements in lower-level design entities that become nodes when imported to the top-level design. You can create virtual pins by assigning the Virtual Pin logic option to an I/O element.

You might use virtual pin assignments when you compile a partial design, because not all the I/Os from a partial design drive chip pins at the top level.

The virtual pin assignment identifies the I/O ports of a design module that are internal nodes in the top-level design. These assignments prevent the number of I/O ports in the lower-level modules from exceeding the total number of available device pins. Every I/O port that you designate as a virtual pin becomes mapped to either a logic cell or an adaptive logic module (ALM), depending on the target device.
The Virtual Pin logic option must be assigned to an input or output pin. If you assign this option to a bidirectional pin, tri-state pin, or registered I/O element, Analysis and Synthesis ignores the assignment. If you assign this option to a tri-state pin, the Fitter inserts an I/O buffer to account for the tri-state logic; therefore, the pin cannot be a virtual pin. You can use multiplexer logic instead of a tri-state pin if you want to continue to use the assigned pin as a virtual pin. Do not use tri-state logic except for signals that connect directly to device I/O pins.

In the top-level design, you connect these virtual pins to an internal node of another module. By making assignments to virtual pins, you can place those pins in the same location or region on the device as that of the corresponding internal nodes in the top-level module. You can use the Virtual Pin option when compiling a LogicLock module with more pins than the target device allows. The Virtual Pin option can enable timing analysis of a design module that more closely matches the performance of the module after you integrate it into the top-level design.

In the Node Finder, you can set Filter Type to Pins: Virtual to display all assigned virtual pins in the design. Alternatively, to access the Node Finder from the Assignment Editor, double-click the To field; when the arrow appears on the right side of the field, click the arrow and select Node Finder.

Using LogicLock Regions in the Chip Planner

You can easily create LogicLock regions in the Chip Planner and assign resources to them.

Viewing Connections Between LogicLock Regions in the Chip Planner

You can view and edit LogicLock regions using the Chip Planner. To view and edit LogicLock regions, select the Floorplan Editing mode in Layers Settings, or any Layers setting mode that has the User-assigned LogicLock regions setting enabled.

The Chip Planner shows the connections between LogicLock regions. By default, you can view each connection as an individual line. You can choose to display connections between two LogicLock regions as a single bundled connection rather than as individual connection lines. To use this option, open the Chip Planner and on the View menu, click Inter-region Bundles.

For more information about the Inter-region Bundles dialog box, refer to Inter-region Bundles Dialog Box in Quartus II Help.

Using LogicLock Regions with the Design Partition Planner

You can optimize timing in a design by placing entities that share significant logical connectivity close to each other on the device. By default, the Fitter usually places closely connected entities in the same area of the device; however, you can use LogicLock regions, together with the Design Partition Planner and the Chip Planner, to help ensure that logically connected entities retain optimal placement from one compilation to the next.
You can view the logical connectivity between entities with the Design Partition Planner, and the physical placement of those entities with the Chip Planner. In the Design Partition Planner, you can identify entities that are highly interconnected, and place those entities in a partition. In the Chip Planner, you can create LogicLock regions and assign each partition to a LogicLock region, thereby preserving the placement of the entities.

For more information about using LogicLock regions with design partitions, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* and the *Best Practices for Incremental Compilation Partition and Floorplan Assignments* chapters in volume 1 of the *Quartus II Handbook*. For more information about using the Design Partition Planner with the Chip Planner, refer to *About the Design Partition Planner* and *Using the Design Partition Planner* in Quartus II Help.

**Design Floorplan Analysis Using the Chip Planner**

The Chip Planner helps you visually analyze the floorplan of your design at any stage of your design cycle. With the Chip Planner, you can view post-compilation placement, connections, and routing paths. You can also create LogicLock regions and location assignments. The Chip Planner allows you to create new logic cells and I/O atoms and to move existing logic cells and I/O atoms in your design. You can also see global and regional clock regions within the device, and the connections between I/O atoms, PLLs and the different clock regions.

From the Chip Planner, you can launch the Resource Property Editor, which you can use to change the properties and parameters of device resources, and modify connectivity between certain types of device resources. The Change Manager records any changes that you make to your design floorplan so that you can selectively undo changes if necessary.

For more information about the Resource Property Editor and the Change Manager, refer to the *Engineering Change Management with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*, and to *About the Resource Property Editor* and *About the Change Manager* in Quartus II Help.

The following sections present Chip Planner floorplan views and design analysis procedures which you can use with any Chip Planner preset, unless a procedure requires a specific preset or editing mode.

**Chip Planner Floorplan Views**

The Chip Planner uses a hierarchical zoom viewer that shows various abstraction levels of the targeted Altera device. As you zoom in, the level of abstraction decreases, revealing more details about your design.

For more information about Chip Planner floorplan views, refer to the *Engineering Change Management with the Chip Planner* chapter in volume 2 of the *Quartus II Handbook*. 
Bird’s Eye View
The Bird’s Eye View displays a high-level picture of resource usage for the entire chip and provides a fast and efficient way to navigate between areas of interest in the Chip Planner.

The Bird’s Eye View is particularly useful when the parts of your design that you want to view are at opposite ends of the chip and you want to quickly navigate between resource elements without losing your frame of reference.

For more information about the Bird’s Eye View, refer to Bird’s Eye View and Displaying Resources and Information in Quartus II Help.

Properties Window
The Properties Window displays detailed properties of the objects (such as atoms, paths, LogicLock regions, or routing elements) currently selected in the Chip Planner. To display the Properties Window, click Properties on the View menu in the Chip Planner.

Viewing Architecture-Specific Design Information
By adjusting the Layers Settings in the Chip Planner, you can view the following architecture-specific information related to your design:

- **Device routing resources used by your design**—View how blocks are connected, as well as the signal routing that connects the blocks.
- **LE configuration**—View logic element (LE) configuration in your design. For example, you can view which LE inputs are used; if the LE utilizes the register, the look-up table (LUT), or both; as well as the signal flow through the LE.
- **ALM configuration**—View ALM configuration in your design. For example, you can view which ALM inputs are used, if the ALM utilizes the registers, the upper LUT, the lower LUT, or all of them. You can also view the signal flow through the ALM.
- **I/O configuration**—View device I/O resource usage. For example, you can view which components of the I/O resources are used, if the delay chain settings are enabled, which I/O standards are set, and the signal flow through the I/O.
- **PLL configuration**—View phase-locked loop (PLL) configuration in your design. For example, you can view which control signals of the PLL are used with the settings for your PLL.
- **Timing**—View the delay between the inputs and outputs of FPGA elements. For example, you can analyze the timing of the DATAB input to the COMBOUT output.

In addition, you can modify the following device properties with the Chip Planner:

- LEs and ALMs
- I/O cells
- PLLs
- Registers in RAM and DSP blocks
- Connections between elements
Placement of elements

For more information about LEs, ALMs, and other resources of an FPGA device, refer to the relevant device handbook.

Viewing Available Clock Networks in the Device

When you select a task with clock region layer preset enabled, you can display the areas of the chip that are driven by global and regional clock networks. This global clock display feature is available for Arria GX, Arria II, Arria V, Cyclone II, Cyclone III, Cyclone V, HardCopy II, HardCopy III, Stratix II, Stratix II GX, Stratix III, Stratix IV, and Stratix V device families.

Depending on the clock layers activated in the selected preset, the Chip Planner displays regional and global clock regions in the device, and the connectivity between clock regions, pins, and PLLs. Clock regions appear as rectangular overlay boxes with labels indicating the clock type and index. You can select each clock network region by clicking on the clock region. The clock-shaped icon at the top-left corner indicates that the region represents a clock network region. You can change the color in which the Chip Planner displays clock regions on the Options dialog box of the Tools menu.

The Layers Settings dialog box lists layers for different clock region types; when the selected device does not contain a given clock region, the option for that category is unavailable in the dialog box. You can customize the Chip Planner’s display of clock regions by creating a custom preset with selected clock layers enabled in the Layers Settings dialog box.

For more information about displaying clock regions, refer to Displaying Resources and Information in Quartus II Help.

Viewing Critical Paths

Critical paths are timing paths in your design that have a negative slack. These timing paths can span from device I/Os to internal registers, registers to registers, or from registers to device I/Os. The slack of a path determines its criticality; slack appears in the timing analysis report. Design analysis for timing closure is a fundamental requirement for optimal performance in highly complex designs. The analytical capability of the Chip Planner helps you close timing on complex designs.

Viewing critical paths in the Chip Planner helps you understand why a specific path is failing. You can see if any modification in the placement can reduce the negative slack. You can display details of a path (to expand/collapse the path to/from the connections in the path) by clicking Expand Connections in the toolbar, or by clicking on the “+/-” on the label.

You can locate failing paths from the timing report in the TimeQuest Timing Analyzer. To locate the critical paths, run the Report Timing task from the Custom Reports group in the Tasks pane of the TimeQuest Timing Analyzer. From the View pane, which lists the failing paths, right-click on any failing path or node, and select Locate Path. From the Locate dialog box, select Chip Planner to see the failing path in the Chip Planner.
To display paths in the floorplan, you must first make timing settings and perform a timing analysis.

For more information about performing static timing analysis with the Quartus II TimeQuest Timing Analyzer, refer to the "Quartus II TimeQuest Timing Analyzer" chapter in volume 3 of the "Quartus II Handbook.

### Viewing Routing Congestion

The **Report Routing Utilization** task allows you to determine the percentage of routing resources in use following a compilation. This feature can identify where there is a lack of routing resources, helping you to make design changes to meet routing congestion design requirements.

Open the Chip Planner from the Tools menu. To view the routing congestion in the Chip Planner, double-click the **Report Routing Utilization** command in the Tasks list. Click **Preview** in the **Report Routing Utilization** dialog box to preview the default congestion display. Change the **Routing utilization type** to display congestion for specific resources. The default display uses dark blue for 0% congestion (blue indicates zero utilization) and red for 100%. You can adjust the slider for **Threshold percentage** to change the congestion threshold level.

The routing congestion map uses the color and shading of logic resources to indicate relative resource utilization; darker shading represents a greater utilization of routing resources. Areas where routing utilization exceeds the threshold value specified in the **Report Routing Utilization** dialog box appear in red. The congestion map can help you determine whether you can modify the floorplan, or make changes to the RTL to reduce routing congestion.

To identify a lack of routing resources, it is necessary to investigate each routing interconnect type separately by selecting each interconnect type in turn in the **Routing Utilization Settings** dialog box.

The Quartus II compilation messages contain information about average and peak interconnect usage. Peak interconnect usage over 75%, or average interconnect usage over 60%, could be an indication that it might be difficult to fit your design. Similarly, peak interconnect usage over 90%, or average interconnect usage over 75%, are likely to have increased chances of not getting a valid fit.

For more information about displaying routing congestion, refer to **Displaying Resources and Information** in Quartus II Help.

### Viewing I/O Banks

The Chip Planner can show all of the I/O banks of the device. To see the I/O bank map of the device, select **Report All I/O Banks** in the Tasks pane.
Viewing High-Speed Serial Interfaces (HSSI)

For the Stratix V device family, the Chip Planner displays a detailed block view of the receiver and transmitter channels of the high-speed serial interfaces. To display the HSSI block view, select Report HSSI Block Connectivity. Figure 15–3 shows the blocks of a Stratix V HSSI receiver channel.

Figure 15–3. Stratix V HSSI receiver channel

Generating Fan-In and Fan-Out Connections

The ability to display fan-in and fan-out connections enables you to view the atoms that fan-in to or fan-out from the selected atom. To remove the connections displayed, use the Clear Unselected Connections icon in the Chip Planner toolbar.

Generating Immediate Fan-In and Fan-Out Connections

The ability to display immediate fan-in and fan-out connections enables you to view the resource that is the immediate fan-in or fan-out connection for the selected atom. For example, if you select a logic resource and choose to view the immediate fan-in for that resource, you can see the routing resource that drives the logic resource. You can generate immediate fan-ins and fan-outs for all logic resources and routing resources. To remove the displayed connections from the screen, click the Clear Unselected Connections icon in the toolbar.
Highlight Routing

The Show Physical Routing command in the Locate History pane enables you to highlight the routing resources used by a selected path or connection. Figure 15–4 shows the routing resources in use between two logic elements.

Figure 15–4. Highlight Routing

You can view and edit resources in the FPGA using the Resource Property Editor. For more information, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.
Show Delays

With the Show Delays command, you can view timing delays for paths located from TimeQuest Timing Analyzer reports. For example, you can view the delay between two logic resources or between a logic resource and a routing resource. Figure 15–5 shows the delay associated with a path located from a TimeQuest Timing Analyzer report.

Figure 15–5. Show Delays

Exploring Paths in the Chip Planner

You can use the Chip Planner to explore paths between logic elements. The following example uses the Chip Planner to traverse paths from the Timing Analysis report.

Locate Path from the Timing Analysis Report to the Chip Planner

To locate a path from the Timing Analysis report to the Chip Planner, perform the following steps:

1. Select the path you want to locate.
2. Right-click the path in the Timing Analysis report, point to **Locate Path**, and click **Locate in Chip Planner**. The path is displayed with its timing data in the Chip Planner main window and is listed in the Locate History window.

3. To view the routing resources taken for a path you have located in the Chip Planner, select the path and then click the **Highlight Routing** icon in the Chip Planner toolbar, or from the View menu, click **Highlight Routing**.

**Analyzing Connections for a Path**

To determine the connections between items in the Chip Planner, click the **Expand Connections** icon on the toolbar. To add the timing delays for paths located from the TimeQuest Timing Analyzer, click the **Show Delays** icon on the toolbar. Figure 15–6 shows the connections for a path located from the TimeQuest Timing Analyzer that are displayed in the Chip Planner. To see the constituent delays on the selected path, click on the “+” sign next to the path delay displayed in the Chip Planner.

**Figure 15–6. Path Analysis**

---

**Viewing Assignments in the Chip Planner**

You can view location assignments by selecting the appropriate layer set in the Chip Planner. To view location assignments, select the **Floorplan Editing** preset or any custom preset that displays block utilization, and the Assignment editing mode. See Figure 15–7.
The Chip Planner shows location assignments graphically, by displaying assigned resources in a particular color (gray, by default). You can create or move an assignment by dragging the selected resource to a new location.

**Figure 15–7. Viewing Assignments in the Chip Planner**

You can make node and pin location assignments to LogicLock regions and custom regions using the drag-and-drop method in the Chip Planner. The Fitter applies the assignments that you create during the next place-and-route operation.

For more information about managing assignments in the Chip Planner, refer to *Working With Assignments in the Chip Planner* in Quartus II Help.

**Viewing High-Speed and Low-Power Tiles in the Chip Planner**

To view a power map of designs that specify Stratix III, Stratix IV, or Stratix V devices, select the **Report High-Speed/Low-Power Tiles** command in the **Tasks** menu after running the Fitter. Stratix III, Stratix IV, or Stratix V devices have ALMs that can operate in either high-speed mode or low-power mode. The power mode is set during the fitting process in the Quartus II software. These ALMs are grouped together to form larger blocks, called “tiles.”

To learn more about power analyses and optimizations in Stratix III devices, refer to *AN 437: Power Optimization in Stratix III FPGAs*. To learn more about power analyses and optimizations in Stratix IV devices, refer to *AN 514: Power Optimization in Stratix IV FPGAs*. 
When you select the **Report High-Speed/Low-Power Tiles** command for Stratix III, Stratix IV, or Stratix V devices, the Chip Planner displays low-power and high-speed tiles in contrasting colors; yellow tiles operate in a high-speed mode, while blue tiles operate in a low-power mode (see Figure 15–8). When you select the **Power** task, you can perform all floorplanner-related functions for this task; however, you cannot edit tiles to change the power mode.

**Figure 15–8. Viewing High-Speed and Low Power Tiles in a Stratix III Device**

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**Scripting Support**

You can run procedures and specify the settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II command-line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```bash
quartus_sh --qhelp
```

Information about scripting command options is also available in *API Functions for Tcl* in Quartus II Help.

For more information about Tcl scripting, refer to the *Tcl Scripting* chapter in volume 2 of the *Quartus II Handbook*. For more information about command-line scripting, refer to the *Command-Line Scripting* chapter in volume 2 of the *Quartus II Handbook*. For information about all settings and constraints in the Quartus II software, refer to the *Quartus II Settings File Manual*. 
Initializing and Uninitializing a LogicLock Region

You must initialize the LogicLock data structures before creating or modifying any LogicLock regions and before executing any of the Tcl commands listed below.

Use the following Tcl command to initialize the LogicLock data structures:

```
initialize_logiclock
```

Use the following Tcl command to uninitialize the LogicLock data structures before closing your project:

```
uninitialize_logiclock
```

Creating or Modifying LogicLock Regions

Use the following Tcl command to create or modify a LogicLock region:

```
set_logiclock -auto_size true -floating true -region <my_region-name>
```

The command in the above example sets the size of the region to auto and the state to floating.

If you specify a region name that does not exist in the design, the command creates the region with the specified properties. If you specify the name of an existing region, the command changes all properties you specify and leaves unspecified properties unchanged.

For more information about creating LogicLock regions, refer to “Creating LogicLock Regions” on page 15–4.

Obtaining LogicLock Region Properties

Use the following Tcl command to obtain LogicLock region properties. This example returns the height of the region named `my_region`:

```
get_logiclock -region my_region -height
```

Assigning LogicLock Region Content

Use the following Tcl commands to assign or change nodes and entities in a LogicLock region. This example assigns all nodes with names matching `fifo*` to the region named `my_region`.

```
set_logiclock_contents -region my_region -to fifo*
```

You can also make path-based assignments with the following Tcl command:

```
set_logiclock_contents -region my_region -from fifo -to ram*
```

Save a Node-Level Netlist for the Entire Design into a Persistent Source File

Make the following assignments to cause the Quartus II Fitter to save a node-level netlist for the entire design into a `.vqm` file:

```
set_global_assignment-name LOGICLOCK_INCREMENTAL_COMPILE_ASSIGNMENT ON
set_global_assignment-name LOGICLOCK_INCREMENTAL_COMPILE_FILE <file name>
```
Any path specified in the file name is relative to the project directory. For example, specifying `atom_netlists/top.vqm` places `top.vqm` in the `atom_netlists` subdirectory of your project directory.

A `.vqm` file is saved in the directory specified at the completion of a full compilation.

The saving of a node-level netlist to a persistent source file is not supported for designs targeting newer devices such as Arria GX, Arria II, Cyclone III, MAX V, Stratix III, Stratix IV, or Stratix V.

### Setting LogicLock Assignment Priority

Use the following Tcl code to set the priority for a LogicLock region’s members. This example reverses the priorities of the LogicLock region in your design.

```tcl
set reverse [list]
for each member [get_logiclock_member_priority] {
    set reverse [insert $reverse 0 $member]
}
set_logiclock_member_priority $reverse
```

### Assigning Virtual Pins

Use the following Tcl command to turn on the virtual pin setting for a pin called `my_pin`:

```tcl
set_instance_assignment -name VIRTUAL_PIN ON -to my_pin
```

For more information about assigning virtual pins, refer to “Virtual Pins” on page 15–10.

For more information about Tcl scripting, refer to the `Tcl Scripting` chapter in volume 2 of the *Quartus II Handbook*.

### Document Revision History

Table 15–2 shows the revision history for this chapter.

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<td>May 2013</td>
<td>13.0.0</td>
<td>Updated “Viewing Routing Congestion” section</td>
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<tr>
<td></td>
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<td>Updated references to Quartus UI controls for the Chip Planner</td>
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<td>Updated “Viewing Routing Congestion”</td>
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<td>Updated “Locate History”</td>
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<tr>
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<tr>
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Table 15–2. Document Revision History (Part 2 of 2)

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<td>■ Updated device support information</td>
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<td>■ Removed references to Timing Closure Floorplan; removed “Design Analysis Using the Timing Closure Floorplan” section</td>
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<td>■ Added links to online Help topics</td>
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<td>■ Added “Using LogicLock Regions with the Design Partition Planner” section</td>
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<td>■ Updated “Viewing Critical Paths” section</td>
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<td>■ Updated “Creating Nonrectangular LogicLock Regions” section</td>
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<td>■ Added “Selected Elements Window” section</td>
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<td>■ Removed the following sections:</td>
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<td></td>
<td>■ Reserve LogicLock Region Design Analysis Using the Timing Closure Floorplan</td>
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*For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.*
The Quartus® II software offers physical synthesis optimizations to improve your design beyond the optimization performed in the normal course of the Quartus II compilation flow.

Physical synthesis optimizations can help improve the performance of your design regardless of the synthesis tool used, although the effect of physical synthesis optimizations depends on the structure of your design.

Netlist optimization options work with the atom netlist of your design, which describes a design in terms of Altera®-specific primitives. An atom netlist file can be an Electronic Design Interchange Format (.edf) file or a Verilog Quartus Mapping (.vqm) file generated by a third-party synthesis tool, or a netlist used internally by the Quartus II software. Physical synthesis optimizations are applied at different stages of the Quartus II compilation flow, either during synthesis, fitting, or both.

This chapter explains how the physical synthesis optimizations in the Quartus II software can modify your design’s netlist to improve the quality of results. This chapter also provides information about preserving compilation results through back-annotation and writing out a new netlist, and provides guidelines for applying the various options.

Because the node names for primitives in the design can change when you use physical synthesis optimizations, you should evaluate whether your design flow requires fixed node names. If you use a verification flow that might require fixed node names, such as the SignalTap® II Logic Analyzer, formal verification, or the LogicLock based optimization flow (for legacy devices), you must turn off physical synthesis options.

### WYSIWYG Primitive Resynthesis

If you use a third-party tool to synthesize your design, use the **Perform WYSIWYG primitive resynthesis** option to apply optimizations to the synthesized netlist.

The **Perform WYSIWYG primitive resynthesis** option directs the Quartus II software to un-map the logic elements (LEs) in an atom netlist to logic gates, and then re-map the gates back to Altera-specific primitives. Third-party synthesis tools generate either an .edf or .vqm atom netlist file using Altera-specific primitives. When you turn on the **Perform WYSIWYG primitive resynthesis** option, the Quartus II software can work on different techniques specific to the device architecture during the re-mapping process. This feature re-maps the design using the **Optimization Technique** specified for your project (Speed, Area, or Balanced).

The **Perform WYSIWYG primitive resynthesis** option has no effect if you are using Quartus II integrated synthesis to synthesize your design.
To turn on the Perform WYSIWYG primitive resynthesis option, perform the following steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
2. In the Category list, select Analysis and Synthesis Settings. The Analysis & Synthesis Settings page appears.
3. Turn on Perform WYSIWYG Primitive Resynthesis, and click OK.

If you want to perform WYSIWYG resynthesis on only a portion of your design, you can use the Assignment Editor to assign the Perform WYSIWYG primitive resynthesis logic option to a lower-level entity in your design. This logic option is available for all Altera devices supported by the Quartus II software except MAX 3000 and MAX 7000 devices.

The results of the remapping depend on the Optimization Technique you choose. To select an Optimization Technique, perform the following steps:

1. In the Category list, select Analysis & Synthesis Settings. The Analysis & Synthesis Settings page appears.
2. Under Optimization Technique, select Speed, Area, or Balanced to specify how the Quartus II technology mapper optimizes the design. The Balanced setting is the default for many Altera device families; this setting optimizes the timing critical parts of the design for speed and the rest of the design for area.
3. Click OK.

Refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook for details on the Optimization Technique option.

Figure 16–1 shows the Quartus II software flow for the WYSIWYG primitive resynthesis feature.

**Figure 16–1. WYSIWYG Primitive Resynthesis**
The **Perform WYSIWYG primitive resynthesis** option unmaps and remaps only logic cells, also referred to as LCELL or LE primitives, and regular I/O primitives (which may contain registers). Double data rate (DDR) I/O primitives, memory primitives, digital signal processing (DSP) primitives, and logic cells in carry/cascade chains are not remapped. Logic specified in an encrypted .vqm file or an .edf file, such as third-party intellectual property (IP), is not touched.

The **Perform WYSIWYG primitive resynthesis** option can change node names in the .vqm file or .edf file from your third-party synthesis tool, because the primitives in the atom netlist are broken apart and then remapped by the Quartus II software. The remapping process removes duplicate registers, but registers that are not removed retain the same name after remapping.

Any nodes or entities that have the Netlist Optimizations logic option set to Never Allow are not affected during WYSIWYG primitive resynthesis. You can use the Assignment Editor to apply the Netlist Optimizations logic option. This option disables WYSIWYG resynthesis for parts of your design.

Primitive node names are specified during synthesis. When netlist optimizations are applied, node names might change because primitives are created and removed. HDL attributes applied to preserve logic in third-party synthesis tools cannot be maintained because those attributes are not written into the atom netlist read by the Quartus II software.

If you use the Quartus II software to synthesize, you can use the Preserve Register (preserve) and Keep Combinational Logic (keep) attributes to maintain certain nodes in the design.

For more information about using these attributes during synthesis in the Quartus II software, refer to the Quartus II Integrated Synthesis chapter in volume 1 of the Quartus II Handbook.

### Performing Physical Synthesis Optimizations

The Quartus II design flow involves separate steps of synthesis and fitting. The synthesis step optimizes the logical structure of a circuit for area, speed, or both. The Fitter then places and routes the logic cells to ensure critical portions of logic are close together and use the fastest possible routing resources. While you are using this push-button flow, the synthesis stage is unable to anticipate the routing delays seen in the Fitter. Because routing delays are a significant part of the typical critical path delay, the physical synthesis optimizations available in the Quartus II software take those routing delays into consideration and focus timing-driven optimizations at those parts of the design. This tight integration of the fitting and synthesis processes is known as physical synthesis.

The following sections describe the physical synthesis optimizations available in the Quartus II software, and how they can help improve your performance results. Physical synthesis optimization options can be used with Arria series, Cyclone, HardCopy, and Stratix series device families.
If you are migrating your design to a HardCopy II device, you can target physical synthesis optimizations to the FPGA architecture in the FPGA-first flow or to the HardCopy II architecture in the HardCopy-first flow. The optimizations are mapped to the other device architecture during the migration process.

You cannot target optimizations to both device architectures individually because doing so results in a different post-fitting netlist for each device.

For more information about physical synthesis optimizations, refer to Physical Synthesis Optimizations Page (Settings Dialog Box) in Quartus II Help. For more information about using physical synthesis with HardCopy devices, refer to the Quartus II Support for HardCopy Series Devices chapter in volume 1 of the Quartus II Handbook.

You can choose the physical synthesis optimization options you want for your design during synthesis and fitting in the Physical Synthesis Optimizations page under the Compilation Process Settings page in the Settings dialog box. The settings include optimizations for improving performance and fitting in the selected device.

You can also set the effort level for physical synthesis optimizations. Normally, physical synthesis optimizations increase the compilation time; however, you can select the Fast effort level if you want to limit the increase in compilation time. When you select the Fast effort level, the Quartus II software performs limited register retiming operations during fitting. The Extra effort level runs additional algorithms to get the best circuit performance, but results in increased compilation time.

To optimize performance, the following options are available:

- **Perform physical synthesis for combinational logic**
- **Perform register retiming**
- **Perform automatic asynchronous signal pipelining**
- **Perform register duplication**

To optimize for better fitting, you can choose from the following options:

- **Perform physical synthesis for combinational logic**
- **Perform logic to memory mapping**

To view and modify the physical synthesis optimization options, perform the following steps:

1. On the Assignments menu, click **Settings**. The **Settings** dialog box appears.
2. In the **Category** list, select **Physical Synthesis Optimizations** under Compilation Process Settings. The Physical Synthesis Optimizations page appears.
3. Specify the options for performing physical synthesis optimizations.

Some physical synthesis options affect only registered logic and some options affect only combinational logic. Select options based on whether you want to keep the registers intact or not. For example, if your verification flow involves formal verification, you might have to keep the registers intact.
Performing Physical Synthesis Optimizations

All Physical Synthesis optimizations write results to the Netlist Optimizations report, which provides a list of atom netlist files that were modified, created, and deleted during physical synthesis. To access the Netlist Optimizations report, perform the following steps:

2. In the Compilation Report list, select Netlist Optimizations under Fitter.

Similarly, physical synthesis optimizations performed during synthesis write results to the synthesis report. To access this report, perform the following steps:

2. In the Compilation Report list, select Analysis & Synthesis.
3. In the Optimization Results folder, select Netlist Optimizations. The Physical Synthesis Netlist Optimizations table appears, listing the physical synthesis netlist optimizations performed during synthesis.

Nodes or entities that have the Netlist Optimizations logic option set to Never Allow are not affected by the physical synthesis algorithms. You can use the Assignment Editor to apply the Netlist Optimizations logic option. Use this option to disable physical synthesis optimizations for parts of your design.

Automatic Asynchronous Signal Pipelining

The Perform automatic asynchronous signal pipelining option on the Physical Synthesis Optimizations page in the Compilation Process Settings section of the Settings dialog box allows the Quartus II Fitter to perform automatic insertion of pipeline stages for asynchronous clear and asynchronous load signals during fitting when these signals negatively affect performance. You can use this option if asynchronous control signal recovery and removal times are not achieving their requirements.

The Perform automatic asynchronous signal pipelining option improves performance for designs in which asynchronous signals in very fast clock domains cannot be distributed across the chip fast enough due to long global network delays. This optimization performs automatic pipelining of these signals, while attempting to minimize the total number of registers inserted.

The Perform automatic asynchronous signal pipelining option adds registers to nets driving the asynchronous clear or asynchronous load ports of registers. These additional registers add register delays (adds latency) to the reset, adding the same number of register delays for each destination using the reset. The additional register delays can change the behavior of the signal in the design; therefore, you should use this option only if additional latency on the reset signals does not violate any design requirements. This option also prevents the promotion of signals to global routing resources.

The Quartus II software performs automatic asynchronous signal pipelining only if Enable Recovery/Removal analysis is turned on. If you use the TimeQuest Timing Analyzer, Enable Recovery/Removal analysis is turned on by default. Pipelining is allowed only on asynchronous signals that have the following properties:

- The asynchronous signal is synchronized to a clock (a synchronization register drives the signal)
- The asynchronous signal fans-out only to asynchronous control ports of registers

The Quartus II software does not perform automatic asynchronous signal pipelining on asynchronous signals that have the Netlist Optimization logic option set to Never Allow.

**Physical Synthesis for Combinational Logic**

To optimize the design and reduce delay along critical paths, you can turn on the Perform physical synthesis for combinational logic option, which swaps the look-up table (LUT) ports within LEs so that the critical path has fewer layers through which to travel. The Perform physical synthesis for combinational logic option also allows the duplication of LUTs to enable further optimizations on the critical path.

For more information about using the Perform physical synthesis for combinational logic option, refer to Physical Synthesis Optimizations Page (Settings Dialog Box) and to Setting Up and Running the Fitter in Quartus II Help.

The Perform physical synthesis for combinational logic option affects only combinational logic in the form of LUTs. These transformations might occur during the synthesis stage or the Fitter stage during compilation. The registers contained in the affected logic cells are not modified. Inputs into memory blocks, DSP blocks, and I/O elements (IOEs) are not swapped.

The Quartus II software does not perform combinational optimization on logic cells that have the following properties:

- Are part of a chain
- Drive global signals
- Are constrained to a single logic array block (LAB) location
- Have the Netlist Optimizations option set to Never Allow

If you want to consider logic cells with any of these conditions for physical synthesis, you can override these rules by setting the Netlist Optimizations logic option to Always Allow on a given set of nodes.

**Physical Synthesis for Registers—Register Duplication**

The Perform register duplication option on the Physical Synthesis Optimizations page in the Compilation Process Settings section of the Settings dialog box allows the Quartus II Fitter to duplicate registers based on Fitter placement information. You can also duplicate combinational logic when this option is enabled. A logic cell that fans out to multiple locations can be duplicated to reduce the delay of one path without degrading the delay of another. The new logic cell can be placed closer to critical logic without affecting the other fan-out paths of the original logic cell.

For more information about the Perform register duplication option, refer to Physical Synthesis Optimizations Page (Settings Dialog Box) and to Setting Up and Running the Fitter in Quartus II Help.
The Quartus II software does not perform register duplication on logic cells that have the following properties:

- Are part of a chain
- Contain registers that drive asynchronous control signals on another register
- Contain registers that drive the clock of another register
- Contain registers that drive global signals
- Contain registers that are constrained to a single LAB location
- Contain registers that are driven by input pins without a tSU constraint
- Contain registers that are driven by a register in another clock domain
- Are considered virtual I/O pins
- Have the Netlist Optimizations option set to Never Allow

For more information about virtual I/O pins, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

If you want to consider logic cells that meet any of these conditions for physical synthesis, you can override these rules by setting the Netlist Optimizations logic option to Always Allow on a given set of nodes.

**Physical Synthesis for Registers—Register Retiming**

The Perform Register Retiming option enables the movement of registers across combinational logic, allowing the Quartus II software to trade off the delay between timing-critical paths and non-critical paths. Register retiming can be done during Quartus II integrated synthesis or during the Fitter stages of design compilation.

Figure 16–2 shows an example of register retiming in which the 10-ns critical delay is reduced by moving the register relative to the combinational logic.

**Figure 16–2. Register Retiming Diagram**

Retiming can create multiple registers at the input of a combinational block from a register at the output of a combinational block. In this case, the new registers have the same clock and clock enable. The asynchronous control signals and power-up level are derived from previous registers to provide equivalent functionality. Retiming can also combine multiple registers at the input of a combinational block to a single register (Figure 16–3).
Performing Physical Synthesis Optimizations

To move registers across combinational logic to balance timing, perform the following steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
3. Specify your preferred option under Optimize for performance (physical synthesis) and Effort level.
4. Click OK.

For more information about the Optimize for performance (physical synthesis) options and effort levels, refer to Physical Synthesis Optimizations Page (Settings Dialog Box) in Quartus II Help.

If you want to prevent register movement during register retiming, you can set the Netlist Optimizations logic option to Never Allow. You can apply this option to either individual registers or entities in the design using the Assignment Editor.

In digital circuits, synchronization registers are instantiated on cross clock domain paths to reduce the possibility of metastability. The Quartus II software detects such synchronization registers and does not move them, even if register retiming is turned on.

The following sets of registers are not moved during register retiming:

- Both registers in a direct connection from input pin-to-register-to-register if both registers have the same clock and the first register does not fan-out to anywhere else. These registers are considered synchronization registers.

- Both registers in a direct connection from register-to-register if both registers have the same clock, the first register does not fan out to anywhere else, and the first register is fed by another register in a different clock domain (directly or through combinational logic). These registers are considered synchronization registers.

The Quartus II software assumes that a synchronization register chain consists of two registers. If your design has synchronization register chains with more than two registers, you must indicate the number of registers in your synchronization chains so that they are not affected by register retiming. To do this, perform the following steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
2. In the Category list, select Analysis & Synthesis Settings. The Analysis & Synthesis Setting page appears.

4. In the Name list, select Synchronization Register Chain Length and modify the setting to match the synchronization register length used in your design. If you set a value of 1 for the Synchronization Register Chain Length, it means that any registers connected to the first register in a register-to-register connection can be moved during retiming. A value of \( n > 1 \) means that any registers in a sequence of length \( 1, 2, \ldots n \) are not moved during register retiming.

The Quartus II software does not perform register retiming on logic cells that have the following properties:

- Are part of a cascade chain
- Contain registers that drive asynchronous control signals on another register
- Contain registers that drive the clock of another register
- Contain registers that drive a register in another clock domain
- Contain registers that are driven by a register in another clock domain

The Quartus II software does not usually retime registers across different clock domains; however, if you use the Classic Timing Analyzer and specify a global \( f_{\text{MAX}} \) requirement, the Quartus II software interprets all clocks as related. Consequently, the Quartus II software might try to retime register-to-register paths associated with different clocks.

To avoid this circumstance, provide individual \( f_{\text{MAX}} \) requirements to each clock when using Classic Timing Analysis. When you constrain each clock individually, the Quartus II software assumes no relationship between different clock domains and considers each clock domain to be asynchronous to other clock domains; hence no register-to-register paths crossing clock domains are retimed.

When you use the TimeQuest Timing Analyzer, register-to-register paths across clock domains are never retimed, because the TimeQuest Timing Analyzer treats all clock domains as asynchronous to each other unless they are intentionally grouped.

- Contain registers that are constrained to a single LAB location
- Contain registers that are connected to SERDES
- Are considered virtual I/O pins
- Registers that have the Netlist Optimizations logic option set to Never Allow

For more information about virtual I/O pins, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

If you want to consider logic cells that meet any of these conditions for physical synthesis, you can override these rules by setting the Netlist Optimizations logic option to Always Allow on a given set of registers.
Preserving Your Physical Synthesis Results

The Quartus II software generates the same results on every compilation for the same source code and settings on a given system, hence you do not need to preserve your results from compilation to compilation. When you make changes to the source code or to the settings, you usually get the best results by allowing the software to compile without using previous compilation results or location assignments. In some cases, if you avoid performing analysis and synthesis or `quartus_map`, and run the Fitter or another desired Quartus II executable instead, you can skip the synthesis stage of the compilation.

When you use the Quartus II incremental compilation flow, you can preserve synthesis results for a particular partition of your design by choosing a netlist type of post-synthesis. If you want to preserve fitting results between compilation runs, choose a netlist type of post-fit during incremental compilation.

The rest of this section is relevant only for those designs using older devices that do not support incremental compilation.

For information about the incremental compilation design methodology, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*, and to *About Incremental Compilation* in Quartus II Help.

You can preserve the resulting nodes from physical synthesis in older devices that do not support incremental compilation. You might need to preserve nodes if you use the LogicLock flow to back-annotate placement, import one design into another, or both. For all device families that support incremental compilation, use that feature to preserve results.

To preserve the nodes from Quartus II physical synthesis optimization options for older devices that do not support incremental compilation (such as Max II devices), perform the following steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
3. Turn on Save a node-level netlist of the entire design into a persistent source file. This setting is not available for Cyclone III, Stratix III, and newer devices.
4. Click OK.

The *Save a node-level netlist of the entire design into a persistent source file* option saves your final results as an atom-based netlist in .vqm file format. By default, the Quartus II software places the .vqm file in the atom_netlists directory under the current project directory. To create a different .vqm file using different Quartus II settings, in the Compilation Process Settings page, change the File name setting.
If you use the physical synthesis optimizations and want to lock down the location of all LEs and other device resources in the design with the **Back-Annotate Assignments** command, a `.vqm` file netlist is required. The `.vqm` file preserves the changes that you made to your original netlist. Because the physical synthesis optimizations depend on the placement of the nodes in the design, back-annotating the placement changes the results from physical synthesis. Changing the results means that node names are different, and your back-annotated locations are no longer valid.

You should not use a Quartus II-generated `.vqm` file or back-annotated location assignments with physical synthesis optimizations unless you have finalized the design. Making any changes to the design invalidates your physical synthesis results and back-annotated location assignments. If you require changes later, use the new source HDL code as your input files, and remove the back-annotated assignments corresponding to the Quartus II-generated `.vqm` file.

To back-annotate logic locations for a design that was compiled with physical synthesis optimizations, first create a `.vqm` file. When recompiling the design with the hard logic location assignments, use the new `.vqm` file as the input source file and turn off the physical synthesis optimizations for the new compilation.

If you are importing a `.vqm` file and back-annotated locations into another project that has any **Netlist Optimizations** turned on, you must apply the **Never Allow** constraint to make sure node names don’t change; otherwise, the back-annotated location or LogicLock assignments are invalid.

For newer devices, such as the Arria, Cyclone, or Stratix series, use incremental compilation to preserve compilation results instead of using logic back-annotation.

### Physical Synthesis Options for Fitting

The Quartus II software provides physical synthesis optimization options for improving fitting results. To access these options, perform the following steps:

1. On the Assignments menu, click **Settings**. The **Settings** dialog box appears.

2. In the **Category** list, select **Physical Synthesis Optimizations under Compilation Process Settings**. The **Physical Synthesis Optimizations** page appears.

3. Under **Optimize for fitting** (physical synthesis for density), there are two physical synthesis options available to improve fitting your design in the target device: **Physical synthesis for combinational logic** and **Perform logic to memory mapping** (Table 16–1).

<table>
<thead>
<tr>
<th>Option</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Physical Synthesis for Combinational Logic</strong></td>
<td>When you select this option, the Fitter detects duplicate combinational logic and optimizes combinational logic to improve the fit.</td>
</tr>
<tr>
<td><strong>Perform Logic to Memory Mapping</strong></td>
<td>When you select this option, the Fitter can remap registers and combinational logic in your design into unused memory blocks and achieves a fit.</td>
</tr>
</tbody>
</table>

For more information about physical synthesis optimization options, refer to **Physical Synthesis Optimizations Page (Settings Dialog Box)** in Quartus II Help.
Applying Netlist Optimization Options

The improvement in performance when using netlist optimizations is design dependent. If you have restructured your design to balance critical path delays, netlist optimizations might yield minimal improvement in performance. You may have to experiment with available options to see which combination of settings works best for a particular design. Refer to the messages in the compilation report to see the magnitude of improvement with each option, and to help you decide whether you should turn on a given option or specific effort level.

Turning on more netlist optimization options can result in more changes to the node names in the design; bear this in mind if you are using a verification flow, such as the SignalTap II Logic Analyzer or formal verification that requires fixed or known node names.

Applying all of the physical synthesis options at the Extra effort level generally produces the best results for those options, but adds significantly to the compilation time. You can also use the Physical synthesis effort level options to decrease the compilation time. The WYSIWYG primitive resynthesis option does not add much compilation time relative to the overall design compilation time.

To find the best results, you can use the Quartus II Design Space Explorer (DSE) to apply various sets of netlist optimization options.

For more information about DSE, refer to About Design Space Explorer in Quartus II Help.

Scripting Support

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II Command-Line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook and API Functions for Tcl in Quartus II Help. Refer to the Quartus II Settings File Manual for information about all settings and constraints in the Quartus II software. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

You can specify many of the options described in this section on either an instance or global level, or both.

Use the following Tcl command to make a global assignment:

```
set_global_assignment -name <QSF variable name> <value>
```

Use the following Tcl command to make an instance assignment:

```
set_instance_assignment -name <QSF variable name> <value> -to <instance name>
```
Synthesis Netlist Optimizations

Table 16–2 lists the Quartus II Settings File (.qsf) variable names and applicable values for the settings discussed in “WYSIWYG Primitive Resynthesis” on page 16–1. The .qsf file variable name is used in the Tcl assignment to make the setting along with the appropriate value. The Type column indicates whether the setting is supported as a global setting, an instance setting, or both.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>Quartus II Settings File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Perform WYSIWYG Primitive Resynthesis</td>
<td>ADV_NETLIST_OPT_SYNTH_WYSIWYG_REMAP</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Optimization Technique</td>
<td>&lt;Device Family Name&gt;_OPTIMIZATION_TECHNIQUE</td>
<td>AREA, SPEED, BALANCED</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Power-Up Don’t Care</td>
<td>ALLOW_POWER_UP_DONT_CARE</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Save a node-level netlist into a persistent source file</td>
<td>LOGICLOCK_INCREMENTAL_COMPILE_ASSIGNMENT</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td></td>
<td>LOGICLOCK_INCREMENTAL_COMPILE_FILE</td>
<td>&lt;file name&gt;</td>
<td>Global</td>
</tr>
<tr>
<td>Allow Netlist Optimizations</td>
<td>ADV_NETLIST_OPT_ALLOWED</td>
<td>&quot;ALWAYS ALLOW&quot;, DEFAULT, &quot;NEVER ALLOW&quot;</td>
<td>Instance</td>
</tr>
</tbody>
</table>

Physical Synthesis Optimizations

Table 16–3 lists the .qsf file variable name and applicable values for the settings discussed in “Performing Physical Synthesis Optimizations” on page 16–3. The .qsf file variable name is used in the Tcl assignment to make the setting, along with the appropriate value. The Type column indicates whether the setting is supported as a global setting, an instance setting, or both.

<table>
<thead>
<tr>
<th>Setting Name</th>
<th>Quartus II Settings File Variable Name</th>
<th>Values</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Physical Synthesis for Combinational Logic</td>
<td>PHYSICAL_SYNTHESIS_COMBO_LOGIC</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Automatic Asynchronous Signal Pipelining</td>
<td>PHYSICAL_SYNTHESISASYNCHRONOUS_SIGNAL_PIPELINING</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Perform Register Duplication</td>
<td>PHYSICAL_SYNTHESIS_REGISTER_DUPLICATION</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Perform Register Retiming</td>
<td>PHYSICAL_SYNTHESIS_REGISTER_RETIMING</td>
<td>ON, OFF</td>
<td>Global</td>
</tr>
<tr>
<td>Power-Up Don’t Care</td>
<td>ALLOW_POWER_UP_DONT_CARE</td>
<td>ON, OFF</td>
<td>Global, Instance</td>
</tr>
<tr>
<td>Power-Up Level</td>
<td>POWER_UP_LEVEL</td>
<td>HIGH, LOW</td>
<td>Instance</td>
</tr>
</tbody>
</table>
Incremental Compilation

For information about scripting and command line usage for incremental compilation as mentioned in “Preserving Your Physical Synthesis Results” on page 16–10, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.

Back-Annotating Assignments

You can use the `logiclock_back_annotate` Tcl command to back-annotate resources in your design. This command can back-annotate resources in LogicLock regions, and resources in designs without LogicLock regions.

For more information about back-annotating assignments, refer to “Preserving Your Physical Synthesis Results” on page 16–10.

The following Tcl command back-annotates all registers in your design:

```
logiclock_back_annotate -resource_filter "REGISTER"
```

The `logiclock_back_annotate` command is in the backannotate package.

Conclusion

Physical synthesis optimizations restructure and optimize your design netlist. You can take advantage of these Quartus II netlist optimizations to help improve your quality of results.

Document Revision History

Table 16–4 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Template update.</td>
</tr>
</tbody>
</table>
Table 16–4. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Added links to Quartus II Help in several sections.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed Referenced Documents section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Reformatted Document Revision History</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Added information to “Physical Synthesis for Registers—Register Retiming”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added information to “Applying Netlist Optimization Options”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Made minor editorial updates</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Was chapter 11 in the 8.1.0 release.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated the “Physical Synthesis for Registers—Register Retiming” and “Physical Synthesis Options for Fitting”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated “Performing Physical Synthesis Optimizations”</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Deleted Gate-Level Register Retiming section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated the referenced documents</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8½” × 11” page size. No change to content.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Updated “Physical Synthesis Optimizations for Performance on page 11-9</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added Physical Synthesis Options for Fitting on page 11-16</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
Programmable logic can accommodate changes to a system specification late in the design cycle. Last-minute design changes, commonly referred to as engineering change orders (ECOs), are small changes to the functionality of a design after the design has been fully compiled. This section describes how the Chip Planner feature in the Quartus® II software supports ECOs by allowing quick and efficient changes to your logic late in the design cycle.

This section includes the following chapter:

- **Chapter 17, Engineering Change Management with the Chip Planner**

  This chapter addresses the impact that ECOs have on the design cycle, discusses the design flow for performing ECOs, and describes how you can use the Chip Planner to perform ECOs.
17. Engineering Change Management with the Chip Planner

Programmable logic can accommodate changes to a system specification late in the design cycle. In a typical engineering project development cycle, the specification of the programmable logic portion is likely to change after engineering begins or while integrating all system elements. Last-minute design changes, commonly referred to as engineering change orders (ECOs), are small targeted changes to the functionality of a design after the design has been fully compiled. This chapter discusses the design flow for making ECOs, addresses the impact that ECOs have on the design cycle, and describes how you can use the Chip Planner to make ECOs.

The Chip Planner supports ECOs by allowing quick and efficient changes to your logic late in the design cycle. The Chip Planner provides a visual display of your post-place-and-route design mapped to the device architecture of your chosen FPGA and allows you to create, move, and delete logic cells and I/O atoms.

For a list of supported devices, refer to About the Chip Planner in Quartus® II Help.

In addition to making ECOs, the Chip Planner allows you to perform detailed analysis on routing congestion, relative resource usage, logic placement, LogicLock™ regions, fan-ins and fan-outs, paths between registers, and delay estimates for paths.

For more information about using the Chip Planner for design analysis, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

ECOs directly apply to atoms in the target device. As such, performing an ECO relies on your understanding of the device architecture of the target device. For more information about the architecture of your device, refer to the appropriate device handbook on the Literature page of the Altera website.

This chapter includes the following topics:

- “Engineering Change Orders” on page 17–2
- “ECO Design Flow” on page 17–4
- “The Chip Planner Overview” on page 17–5
- “Performing ECOs with the Chip Planner (Floorplan View)” on page 17–6
- “Performing ECOs in the Resource Property Editor” on page 17–7
- “Change Manager” on page 17–21
- “Scripting Support” on page 17–22
- “Common ECO Applications” on page 17–22
Engineering Change Orders

In the context of an FPGA design, you can apply an ECO directly to a physical resource on the device to modify its behavior. ECOs are typically made during the verification stage of a design cycle. When a small change is required on a design (such as modifying a PLL for a different clock frequency or routing a signal out to a pin for analysis) recompilation of the entire design can be time consuming, especially for larger designs. Because several iterations of small design changes can occur during the verification cycle, recompilation times can quickly add up. Furthermore, a full recompilation due to a small design change can result in the loss of previous design optimizations. Making ECOs, instead of performing a full recompilation on your design, limits the change only to the affected portions of logic.

This section discusses the areas in which ECOs have an impact on a system design and how the Quartus II software can help you optimize the design in these areas. The following topics are discussed in this section:

- “Performance Preservation”
- “Compilation Time”
- “Verification”
- “Change Modification Record”

Performance Preservation

You can preserve the results of previous design optimizations when you make changes to an existing design with one of the following methods:

- Incremental compilation
- Rapid recompile
- ECOs

Choose the method to modify your design based on the scope of the change. The methods above are arranged from the larger scale change to the smallest targeted change to a compiled design.

The incremental compilation feature allows you to preserve compilation results at an RTL component or module level. After the initial compilation of your design, you can assign modules in your design hierarchy to partitions. Upon subsequent compilations, incremental compilation recompiles changed partitions based on the chosen preservation levels.

The rapid recompilation feature leverages results from the latest post-fit netlist to determine the changes required to honor modifications you have made to the source code. If you turn on the rapid recompilation feature, the Compiler attempts to refit only the portion of the netlist that is related to the code modification.

ECOs provide a finer granularity of control compared to the incremental compilation and the rapid recompilation feature. All modifications are performed directly on the architectural elements of the device. You should use ECOs for targeted changes to the post-fit netlist.
In the Quartus II software versions 10.0 and later, the software does not preserve ECO modifications to the netlist when you recompile a design with the incremental compilation feature turned on. You can reapply ECO changes made during a previous compilation with the Change Manager.

For more information about the incremental compilation feature, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*.

### Compilation Time

In the traditional programmable logic design flow, a small change in the design requires a complete recompilation of the design. A complete recompilation of the design consists of synthesis and place-and-route. Making small changes to the design to reach the final implementation on a board can be a long process. Because the Chip Planner works only on the post-place-and-route database, you can implement your design changes in minutes without performing a full compilation.

### Verification

After you make a design change, you can verify the impact on your design. To verify that your changes do not violate timing requirements, perform static timing analysis with the Quartus II TimeQuest Timing Analyzer after you check and save your netlist changes in the Chip Planner.

For more information about the TimeQuest analyzer, refer to the *Quartus II TimeQuest Timing Analyzer* chapter in volume 3 of the *Quartus II Handbook*.

Additionally, you can perform a gate-level or timing simulation of the ECO-modified design with the post-place-and-route netlist generated by the Quartus II software.

### Change Modification Record

All ECOs made with the Chip Planner are logged in the Change Manager to track all changes. With the Change Manager, you can easily revert to the original post-fit netlist or you can pick and choose which ECOs to apply.

Additionally, the Quartus II software provides support for multiple compilation revisions of the same project. You can use ECOs made with the Chip Planner in conjunction with revision support to compare several different ECO changes and revert back to previous project revisions when required.
**ECO Design Flow**

Figure 17–1 shows the design flow for making ECOs.

For iterative verification cycles, implementing small design changes at the netlist level can be faster than making an RTL code change. As such, making ECO changes are especially helpful when you debug the design on silicon and require a fast turnaround time to generate a programming file for debugging the design.
A typical ECO application occurs when you uncover a problem on the board and isolate the problem to the appropriate nodes or I/O cells on the device. You must be able to correct the functionality quickly and generate a new programming file. By making small changes with the Chip Planner, you can modify the post-place-and-route netlist directly without having to perform synthesis and logic mapping, thus decreasing the turnaround time for programming file generation during the verification cycle. If the change corrects the problem, no modification of the HDL source code is necessary. You can use the Chip Planner to perform the following ECO-related changes to your design:

- Document the changes made with the Change Manager
- Easily recreate the steps taken to produce design changes
- Generate EDA simulation netlists for design verification

For more complex changes that require HDL source code modifications, the incremental compilation feature can help reduce recompilation time.

The Chip Planner Overview

The Chip Planner provides a visual display of device resources. It shows the arrangement and usage of the resource atoms in the device architecture that you are targeting. Resource atoms are the building blocks for your device, such as ALMs, LEs, PLLs, DSP blocks, memory blocks, or I/O elements.

The Chip Planner also provides an integrated platform for design analysis and for making ECOs to your design after place-and-route. The toolset consists of the Chip Planner (providing a device floorplan view of your mapped design) and two integrated subtools—the Resource Property Editor and the Change Manager.

For analysis, the Chip Planner can show logic placement, LogicLock regions, relative resource usage, detailed routing information, routing congestion, fan-ins and fan-outs, paths between registers, and delay estimates for paths. Additionally, the Chip Planner allows you to create location constraints or resource assignment changes, such as moving or deleting logic cells or I/O atoms with the device floorplan. For ECO changes, the Chip Planner enables you to create, move, or delete logic cells in the post-place-and-route netlist for fast programming file generation. Additionally, you can open the Resource Property Editor from the Chip Planner to edit the properties of resource atoms or to edit the connections between resource atoms. All changes to resource atoms and connections are logged automatically with the Change Manager.

Opening the Chip Planner

To open the Chip Planner, on the Tools menu, click Chip Planner. Alternatively, click the Chip Planner icon on the Quartus II software toolbar.

Optionally, you can open the Chip Planner by cross-probing from the shortcut menu in the following tools:

- Design Partition Planner
- Compilation Report
- LogicLock Regions window
Technology Map Viewer
Project Navigator window
RTL source code
Node Finder
Simulation Report
RTL Viewer
Report Timing panel of the TimeQuest Timing Analyzer

The Chip Planner Tasks and Layers

The Chip Planner allows you to set up tasks to quickly implement ECO changes or manipulate assignments for the floorplan of the device. Each task consists of an editing mode and a set of customized layer settings.

For more information about tasks and layers in the Chip Planner, refer to About the Chip Planner in Quartus II Help.

For more information about creating assignments and performing analysis with the Chip Planner, as well as the Chip Planner floorplan views, refer to the Analyzing and Optimizing the Design Floorplan chapter in volume 2 of the Quartus II Handbook.

For more information about making ECOs with the ECO mode, refer to “Performing ECOs with the Chip Planner (Floorplan View)” on page 17–6.

Performing ECOs with the Chip Planner (Floorplan View)

You can manipulate resource atoms in the Chip Planner when you select the ECO editing mode. The following ECO changes can be made with the Chip Planner Floorplan view:

- Create atoms
- Delete atoms
- Move existing atoms

To configure the properties of atoms, such as managing the connections between different LEs/ALMs, use the Resource Property Editor.

For more information about editing atom resource properties, refer to “Performing ECOs in the Resource Property Editor” on page 17–7.

To select the ECO editing mode in the Chip Planner, in the Editing Mode list at the top of the Chip Planner, select the ECO editing mode.
Creating, Deleting, and Moving Atoms

You can use the Chip Planner to create, delete, and move atoms in the post-compilation design.

FAQ For more information about creating, deleting, and moving atoms, refer to Creating, Deleting, and Moving Atoms in Quartus II Help.

Check and Save Netlist Changes

After making all the ECOs, you can run the Fitter to incorporate the changes by clicking the Check and Save Netlist Changes icon in the Chip Planner toolbar. The Fitter compiles the ECO changes, performs design rule checks on the design, and generates a programming file.

Performing ECOs in the Resource Property Editor

You can view and edit the following resources with the Resource Property Editor:

- “Logic Elements” on page 17–7
- “Adaptive Logic Modules” on page 17–10
- “FPGA I/O Elements” on page 17–12
- “PLL Properties” on page 17–24
- “FPGA RAM Blocks” on page 17–19
- “FPGA DSP Blocks” on page 17–20

Logic Elements

An Altera® LE contains a four-input LUT, which is a function generator that can implement any function of four variables. In addition, each LE contains a register fed by the output of the LUT or by an independent function generated in another LE.

You can use the Resource Property Editor to view and edit any LE in the FPGA. To open the Resource Property Editor for an LE, on the Project menu, point to Locate, and then click Locate in Resource Property Editor in one of the following views:

- RTL Viewer
- Technology Map Viewer
- Node Finder
- Chip Planner

FAQ For more information about LE architecture for a particular device family, refer to the device family handbook or data sheet.

You can use the Resource Property Editor to change the following LE properties:

- Data input to the LUT
- LUT mask or LUT equation
Logic Element Schematic View

Figure 17–2 shows how the LE appears in the Resource Property Editor. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Figure 17–2. Stratix LE Architecture (1)

Notes to Figure 17–2:
(1) For more information about the Stratix device's LE architecture, refer to the Stratix Device Handbook.

Logic Element Properties

Figure 17–3 shows an example of the properties that can be viewed for a selected LE in the Resource Property Editor. To view LE properties, on the View menu, click View Properties.
Modes of Operation

LUTs in an LE can operate in either normal or arithmetic mode.

When an LE is configured in normal mode, the LUT in the LE can implement a function of four inputs.

When the LE is configured in arithmetic mode, the LUT in the LE is divided into two 3-input LUTs. The first LUT generates the signal that drives the output of the LUT, while the second LUT generates the carry-out signal. The carry-out signal can drive only a carry-in signal of another LE.

For more information about LE modes of operation, refer to volume 1 of the appropriate device handbook.

Sum and Carry Equations

You can change the logic function implemented by the LUT by changing the sum and carry equations. When the LE is configured in normal mode, you can change only the sum equation. When the LE is configured in arithmetic mode, you can change both the sum and the carry equations.

The LUT mask is the hexadecimal representation of the LUT equation output. When you change the LUT equation, the Quartus II software automatically changes the LUT mask. Conversely, when you change the LUT mask, the Quartus II software automatically computes the LUT equation.

sload and sclr Signals

Each LE register contains a synchronous load (sload) signal and a synchronous clear (sclr) signal. You can invert either the sload or sclr signal feeding into the LE. If the design uses the sload signal in an LE, the signal and its inversion state must be the same for all other LEs in the same LAB. For example, if two LEs in a LAB have the sload signal connected, both LEs must have the sload signal set to the same value. This is also true for the sclr signal.

Register Cascade Mode

When register cascade mode is enabled, the cascade-in port feeds the input to the register. The register cascade mode is used most often when the design implements shift registers. You can change the register cascade mode by connecting (or disconnecting) the cascade in the port. However, if you create this port, you must ensure that the source port LE is directly above the destination LE.

Cell Delay Table

The cell delay table describes the propagation delay from all inputs to all outputs for the selected LE.
Logic Element Connections

To view the connections that feed in and out of an LE, on the View menu, click View Port Connections. Figure 17–4 shows the LE connections in the Connectivity window.

Figure 17–4. View LE Connections

Delete a Logic Element

To delete an LE, follow these steps:

1. Right-click the desired LE in the Chip Planner, point to Locate, and click Locate in Resource Property Editor.

2. You must remove all fan-out connections from an LE prior to deletion. To delete fan-out connections, right-click each connected output signal, point to Remove, and click Fanouts. Select all of the fan-out signals in the Remove Fan-outs dialog box and click OK.

3. To delete an atom after all fan-out connections are removed, right-click the atom in the Chip Planner and click Delete Atom.

Adaptive Logic Modules

Each ALM contains LUT-based resources that can be divided between two adaptive LUTs (ALUTs). With up to eight inputs to the two ALUTs, each ALM can implement various combinations of two functions. This adaptability allows the ALM to be completely backward-compatible with four-input LUT architectures. One ALM can implement any function with up to six inputs and certain seven-input functions. In addition to the ALUT-based resources, each ALM contains two programmable registers, two dedicated full adders, a carry chain, a shared arithmetic chain, and a register chain. The ALM can efficiently implement various arithmetic functions and shift registers with these dedicated resources.

You can implement the following types of functions in a single ALM:

- Two independent 4-input functions
- An independent 5-input function and an independent 3-input function
- A 5-input function and a 4-input function, if they share one input
- Two 5-input functions, if they share two inputs
- An independent 6-input function
- Two 6-input functions, if they share four inputs and share the same functions
Certain 7-input functions

You can use the Resource Property Editor to change the following ALM properties:
- Data input to the LUT
- LUT mask or LUT equation

**Adaptive Logic Module Schematic**

You can view and edit any ALM atom with the Resource Property Editor by right-clicking the ALM in the RTL Viewer, the Node Finder, or the Chip Planner, and clicking **Locate in Resource Property Editor** (Figure 17–5).

For a detailed description of the ALM, refer to the device handbooks of devices based on an ALM architecture.

**Figure 17–5. ALM Schematic (1)**

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**Note to Figure 17–5:**
(1) By default, the Quartus II software displays the used resources in blue and the unused in gray. For Figure 17–5, the used resources are in blue and the unused resources are in gray.

---

**Adaptive Logic Module Properties**

The properties that you can display for the ALM include an equations table that shows the name and location of each of the two combinational nodes and two register nodes in the ALM, the individual LUT equations for each of the combinational nodes, and the `combout`, `sumout`, `carryout`, and `shareout` equations for each combinational node.
Adaptive Logic Module Connections

On the View menu, click View Connectivity to view the input and output connections for the ALM.

FPGA I/O Elements

Altera FPGAs that have high-performance I/O elements, including up to six registers, are equipped with support for a number of I/O standards that allow you to run your design at peak speeds. Use the Resource Property Editor to view, change connectivity, and edit the properties of the I/O elements. Use the Chip Planner (Floorplan view) to change placement, delete, and create new I/O elements.

For a detailed description of the device I/O elements, refer to the applicable device handbook.

You can change the following I/O properties:
- Delay chain
- Bus hold
- Weak pull up
- Slow slew rate
- I/O standard
- Current strength
- Extend OE disable
- PCI I/O
- Register reset mode
- Register synchronous reset mode
- Register power up
- Register mode

Stratix V I/O Elements

The I/O elements in Stratix® V devices contain a bidirectional I/O buffer and I/O registers to support a complete embedded bidirectional single data rate (SDR) or double data rate (DDR) transfer (shown in Figure 17–6).

I/O registers are composed of the input path for handling data from the pin to the core, the output path for handling data from the core to the pin, and the output enable path for handling the output enable signal to the output buffer. These registers allow faster source-synchronous register-to-register transfers and resynchronization. The input path consists of the DDR input registers, alignment and synchronization registers, and half data rate blocks; you can bypass each block in the input path. The input path uses the deskew delay to adjust the input register clock delay across process, voltage, and temperature (PVT) variations.
By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Figure 17–6. Stratix V Device I/O Element Structure

For more information about I/O elements in Stratix V devices, refer to the *Stratix V Device Handbook*. 
Arria GX, Stratix, Stratix II, and Stratix GX I/O Elements

The I/O elements in Arria® GX, Stratix, Stratix II, and Stratix GX devices contain a bidirectional I/O buffer, six registers, and a latch for a complete bidirectional SDR or DDR transfer.

Figure 17–7 shows the Stratix and Stratix GX I/O element structure. The I/O element structure contains two input registers (plus a latch), two output registers, and two output enable registers. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Arria II GX, Stratix III, and Stratix IV I/O Elements

The I/O elements in Arria II GX, Stratix III, and Stratix IV devices contain a bidirectional I/O buffer and I/O registers to support a complete embedded bidirectional SDR or DDR transfer (shown in Figure 17–8). The I/O registers are composed of the input path for handling data from the pin to the core, the output path for handling data from the core to the pin, and the output enable path for handling the output enable signal for the output buffer. Each path consists of a set of delay elements that allow you to fine-tune the timing characteristics of each path for skew management. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Figure 17–8. Stratix III Device I/O Element and Structure

For more information about I/O elements in Stratix III and Stratix IV devices, refer to the Literature page of the Altera website.

For more information about programmable I/O elements in Stratix III devices, refer to AN 474: Implementing Stratix III Programmable I/O Delay Settings in the Quartus II Software.
**Cyclone and Cyclone II I/O Elements**

The I/O elements in Cyclone® and Cyclone II devices contain a bidirectional I/O buffer and three registers for complete bidirectional single data-rate transfer. Figure 17–9 shows the Cyclone and Cyclone II I/O element structure. The I/O element contains one input register, one output register, and one output enable register. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

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For more information about I/O elements in Cyclone II and Cyclone devices, refer to the *Cyclone II Device Handbook* and *Cyclone Device Handbook*, respectively.
Cyclone III I/O Elements

The I/O elements in Cyclone III devices contain a bidirectional I/O buffer and five registers for complete embedded bidirectional single data rate transfer. Figure 17–10 shows the Cyclone III I/O element structure. The I/O element contains one input register, two output registers, and two output-enable registers. The two output registers and two output-enable registers are utilized for double-data rate (DDR) applications. You can use the input registers for fast setup times and the output registers for fast clock-to-output times. Additionally, you can use the output-enable (OE) registers for fast clock-to-output enable timing. You can use I/O elements for input, output, or bidirectional data paths. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Figure 17–10. Cyclone III Device I/O Elements and Structure

For more information about I/O elements in Cyclone III devices, refer to the Cyclone III Device Handbook.
MAX II I/O Elements
The I/O elements in MAX® II devices contain a bidirectional I/O buffer. Figure 17–11 shows the MAX II I/O element structure. You can drive registers from adjacent LABs to or from the bidirectional I/O buffer of the I/O element. By default, the Quartus II software displays the used resources in blue and the unused resources in gray.

Figure 17–11. MAX II Device I/O Elements and Structure

For more information about I/O elements in MAX II devices, refer to the MAX II Device Handbook.
FPGA RAM Blocks

With the Resource Property Editor, you can view the architecture of different RAM blocks in the device, modify the input and output registers to and from the RAM blocks, and modify the connectivity of the input and output ports. Figure 17–12 shows an M9K RAM view in a Stratix III device.

Figure 17–12. M9K RAM View in a Stratix III Device

Note to Figure 17–12:
(1) By default, the Quartus II software displays the used resources in blue and the unused resources in gray. In Figure 17–12, the used resources are in blue and the unused resources are in gray.
FPGA DSP Blocks

Dedicated hardware DSP circuit blocks in Altera devices provide performance benefits for the critical DSP functions in your design. The Resource Property Editor allows you to view the architecture of DSP blocks in the Resource Property Editor for the Cyclone and Stratix series of devices. The Resource Property Editor also allows you to modify the signal connections to and from the DSP blocks and modify the input and output registers to and from the DSP blocks.

Figure 17–13 shows the DSP architecture in a Stratix III device.

**Figure 17–13. DSP Block View in a Stratix III Device (1)**

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**Note to Figure 17–13:**

(1) By default, the Quartus II software displays the used resources in blue and the unused resources in gray. In Figure 17–13, the used resources are in blue and the unused resources are in red.
Change Manager

The Change Manager maintains a record of every change you perform with the Chip Planner, the Resource Property Editor, the SignalProbe feature, or a Tcl script. Each row of data in the Change Manager represents one ECO.

The Change Manager allows you to apply changes, roll back changes, delete changes, and export change records to a Text File (.txt), a Comma-Separated Value File (.csv), or a Tcl Script File (.tcl). The Change Manager tracks dependencies between changes, so that when you apply, roll back, or delete a change, any prerequisite or dependent changes are also applied, rolled back, or deleted.

For more information about the Change Manager, refer to About the Change Manager in Quartus II Help.

Complex Changes in the Change Manager

Certain changes (including creating or deleting atoms and changing connectivity) can appear to be self-contained, but are actually composed of multiple actions. The Change Manager marks such complex changes with a plus icon in the Index column.

You can click the plus icon to expand the change record and show all the component actions preformed as part of that complex change.

For more information about complex change records and about managing changes with the Change Manager, refer to Examples of Managing Changes With the Change Manager in Quartus II Help.

Managing SignalProbe Signals

The SignalProbe pins that you create from the SignalProbe Pins dialog box are recorded in the Change Manager. After you have made a SignalProbe assignment, you can use the Change Manager to quickly disable SignalProbe assignments by selecting Revert to Last Saved Netlist on the shortcut menu in the Change Manager.

For more information about SignalProbe pins, refer to the Quick Design Debugging Using SignalProbe chapter in volume 3 of the Quartus II Handbook.

Exporting Changes

You can export changes to a .txt, a .csv, or a .tcl. Tcl scripts allow you to reapply changes that were deleted during compilation.

For more information about exporting changes, refer to Managing Changes With the Change Manager in Quartus II Help.

For more information about netlist types and the Quartus II incremental compilation feature, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Scripting Support

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. The Tcl commands for controlling the Chip Planner are located in the chip_planner package of the quartus_cdb executable.

A comprehensive list of Tcl commands for the Chip Planner can be found in About Quartus II Scripting in Quartus II Help.

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook. For more information about all settings and constraints in the Quartus II software, refer to the Quartus II Settings File Manual. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

Common ECO Applications

This section provides examples of how you might use an ECO to make a post-compilation change to your design. To help build your system quickly, you can use Chip Planner functions to perform the following activities:

- Adjust the drive strength of an I/O with the Chip Planner
- Modify the PLL properties with the Resource Property Editor (see “Modify the PLL Properties With the Chip Planner” on page 17–23)
- Modify the connectivity between new resource atoms with the Chip Planner and Resource Property Editor

Adjust the Drive Strength of an I/O with the Chip Planner

To adjust the drive strength of an I/O, follow the steps in this section to incorporate the ECO changes into the netlist of the design.
1. In the **Editing Mode** list at the top of the Chip Planner, select the ECO editing mode.

2. Locate the I/O in the **Resource Property Editor**, as shown in Figure 17–14.

**Figure 17–14. I/O in the Resource Property Editor**

3. In the **Resource Property Editor**, point to the **Current Strength** option in the **Properties** pane and double-click the value to enable the drop-down list.

4. Change the value for the **Current Strength** option.

5. Right-click the ECO change in the Change Manager and click **Check & Save All Netlist Changes** to apply the ECO change.

You can change the pin locations of input or output ports with the ECO flow. You can drag and move the signal from an existing pin location to a new location while in the Post Compilation Editing (ECO) task in the Chip Planner. You can then click **Check & Save All Netlist Changes** to compile the ECO.

### Modify the PLL Properties With the Chip Planner

You use PLLs to modify and generate clock signals to meet design requirements. Additionally, you can use PLLs to distribute clock signals to different devices in a design, reducing clock skew between devices, improving I/O timing, and generating internal clock signals.
The Resource Property Editor allows you to view and modify PLL properties to meet your design requirements. Using the Stratix PLL as an example, the rest of this section describes the adjustable PLL properties and the equations as a function of the adjustable PLL properties that govern the PLL output parameters.

Figure 17–15 shows a Stratix PLL in the Resource Property Editor.

**Figure 17–15. PLL View in a Stratix Device**

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**PLL Properties**

The Resource Property Editor allows you to modify PLL options, such as phase shift, output clock frequency, and duty cycle. You can also change the following PLL properties with the Resource Property Editor:

- Input frequency
- M \( V_{CO} \) tap
- M initial
- M value
- N value
- M counter delay
- N counter delay
- M2 value
- N2 value
- SS counter
- Charge pump current
- Loop filter resistance
- Loop filter capacitance
- Counter delay
- Counter high
- Counter low
- Counter mode
- Counter initial
- VCO tap

You can also view post-compilation PLL properties in the Compilation Report. To do so, in the Compilation Report, select Fitter and then select Resource Section.

**Adjusting the Duty Cycle**

Use Equation 17–1 to adjust the duty cycle of individual output clocks.

**Equation 17–1.**

\[
\text{High } \% = \frac{\text{Counter High}}{(\text{Counter High} + \text{Counter Low})}
\]

**Adjusting the Phase Shift**

Use Equation 17–2 to adjust the phase shift of an output clock of a PLL.

**Equation 17–2.**

\[
\text{Phase Shift} = (\text{Period } V_{CO} \times 0.125 \times \text{Tap } V_{CO}) + (\text{Initial } V_{CO} \times \text{Period } V_{CO})
\]

For normal mode, Tap \( V_{CO} \), Initial \( V_{CO} \), and Period \( V_{CO} \) are governed by the following settings:

- Tap \( V_{CO} = \text{Counter Delay} - M \times \text{Tap } V_{CO} \)
- Initial \( V_{CO} = \text{Counter Initial} - M \times \text{Initial} \)
- Period \( V_{CO} = \text{In Clock Period} \times N + M \)

For external feedback mode, Tap \( V_{CO} \), Initial \( V_{CO} \), and Period \( V_{CO} \) are governed by the following settings:

- Tap \( V_{CO} = \text{Counter Delay} - M \times \text{Tap } V_{CO} \)
- Initial \( V_{CO} = \text{Counter Initial} - M \times \text{Initial} \)
- Period \( V_{CO} = \frac{\text{In Clock Period} \times N}{(M + \text{Counter High} + \text{Counter Low})} \)
For a detailed description of the settings, refer to the Quartus II Help. For more information about Stratix device PLLs, refer to the *Stratix Architecture* chapter in volume 1 of the *Stratix Device Handbook*. For more information about PLLs in Arria GX, Cyclone, Cyclone II, and Stratix II devices, refer to the appropriate device handbook.

**Adjusting the Output Clock Frequency**

Use Equation 17–3 to adjust the PLL output clock in normal mode.

**Equation 17–3.**

\[
\text{Output Clock Frequency} = \text{Input Frequency} \cdot \frac{M \text{ value}}{N \text{ value} + \text{Counter High} + \text{Counter Low}}
\]

Use Equation 17–4 to adjust the PLL output clock in external feedback mode.

**Equation 17–4.**

\[
\text{OUTCLK} = \frac{M \text{ value} + \text{External Feedback Counter High} + \text{External Feedback Counter Low}}{N \text{ value} + \text{Counter High} + \text{Counter Low}}
\]

**Adjusting the Spread Spectrum**

Use Equation 17–5 to adjust the spread spectrum for your PLL.

**Equation 17–5.**

\[
\%\text{ spread} = \frac{M_2N_1}{M_1N_2}
\]

**Modify the Connectivity between Resource Atoms**

The Chip Planner and Resource Property Editor allow you to create new resource atoms and manipulate the existing connection between resource atoms in the post-fit netlist. These features are useful for small changes when you are debugging a design, such as manually inserting pipeline registers into a combinational path that fails timing, or routing a signal to a spare I/O pin for analysis. Use the following procedure to create a new register in a Cyclone III device and route register output to a spare I/O pin. This example illustrates how to create a new resource atom and modify the connections between resource atoms.

To create new resource atoms and manipulate the existing connection between resource atoms in the post-fit netlist, follow these steps:

1. Create a new register in the Chip Planner.
2. Locate the atom in the Resource Property Editor.
3. To assign a clock signal to the register, right-click the clock input port for the register, point to **Edit connection**, and click **Other**. Use the Node Finder to assign a clock signal from your design.
4. To tie the SLOAD input port to $V_{CC}$, right-click the clock input port for the register, point to **Edit connection**, and click **VCC**.
5. Assign a data signal from your design to the SDATA port.

6. In the Connectivity window, under the output port names, copy the port name of the register.

7. In the Chip Planner, locate a free I/O resource and create an output buffer.

8. Locate the new I/O atom in the Resource Property Editor.

9. On the input port to the output buffer, right-click, point to Edit connection, and click Other.

10. In the Edit Connection dialog box, type the output port name of the register you have created.

11. Run the ECO Fitter to apply the changes by clicking Check and Save Netlist Changes.

A successful ECO connection is subject to the available routing resources. You can view the relative routing utilization by selecting Routing Utilization as the Background Color Map in the Layers Settings dialog box of the Chip Planner. Also, you can view individual routing channel utilization from local, row, and column interconnects with the tooltips created when you position your mouse pointer over the appropriate resource. Refer to the device data sheet for more information about the architecture of the routing interconnects of your device.

Post ECO Steps

After you make an ECO change with the Chip Planner, you must perform static timing analysis of your design with the TimeQuest analyzer to ensure that your changes did not adversely affect the timing performance of your design.

For example, when you turn on one of the delay chain settings for a specific pin, you change the I/O timing. Therefore, to ensure that the design still meets all timing requirements, you should perform static timing analysis.

For more information about performing a static timing analysis of your design, refer to The Quartus II TimeQuest Timing Analyzer chapter in volume 3 of the Quartus II Handbook.

Conclusion

The Chip Planner allows you to analyze and modify your design floorplan. Also, ECO changes made with the Chip Planner do not require a full recompilation, eliminating the lengthy process of RTL modification, resynthesis, and another place-and-route cycle. In summary, the Chip Planner speeds design verification and timing closure.
Document Revision History

Table 17–1 shows the revision history for this chapter.

Table 17–1. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Updated chapter to new template</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed “The Chip Planner FloorPlan Views” section</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added Stratix V I/O elements in “FPGA I/O Elements” on page 17–12.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Added information to page 17–1.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information to “Engineering Change Orders” on page 17–2.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed heading from “Performance” to “Performance Preservation” on page 17–2.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated information in “Performance Preservation” on page 17–2.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed heading from “Documentation” to “Change Modification Record” on page 17–3.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed heading from “Resource Property Editor” to “Performing ECOs in the Resource Property Editor” on page 17–15.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed “Using Incremental Compilation in the ECO Flow” section. Preservation support for ECOs with the incremental compilation flow has been removed in the Quartus II software version 10.0.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed “Referenced Documents” section.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Updated device support list</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Made minor editorial updates</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Updated Figure 17–17.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Made minor editorial updates.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Chapter 15 was previously Chapter 13 in the 8.1.0 release.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>■ Corrected preservation attributes for ECOs in the section “Using Incremental Compilation in the ECO Flow” on page 15–32.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Changed to 8½&quot; x 11” page size.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Updated device support list</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Modified description for ECO support for block RAMs and DSP blocks</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Corrected Stratix PLL ECO example</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added an application example to show modifying the connectivity between resource atoms</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides additional information about the document and Altera.

About this Handbook

This handbook provides comprehensive information about the current version of the Altera® Quartus® II design software.

How to Contact Altera

To locate the most up-to-date information about Altera products, refer to the following table.

<table>
<thead>
<tr>
<th>Contact (1)</th>
<th>Contact Method</th>
<th>Address</th>
</tr>
</thead>
<tbody>
<tr>
<td>Technical support</td>
<td>Website</td>
<td><a href="http://www.altera.com/support">www.altera.com/support</a></td>
</tr>
<tr>
<td>Technical training</td>
<td>Website</td>
<td><a href="http://www.altera.com/training">www.altera.com/training</a></td>
</tr>
<tr>
<td></td>
<td>Email</td>
<td><a href="mailto:custrain@altera.com">custrain@altera.com</a></td>
</tr>
<tr>
<td>Product literature</td>
<td>Website</td>
<td><a href="http://www.altera.com/literature">www.altera.com/literature</a></td>
</tr>
<tr>
<td>Nontechnical support (general)</td>
<td>Email</td>
<td><a href="mailto:nacomp@altera.com">nacomp@altera.com</a></td>
</tr>
<tr>
<td>(software licensing)</td>
<td>Email</td>
<td><a href="mailto:authorization@altera.com">authorization@altera.com</a></td>
</tr>
</tbody>
</table>

Note to Table:

(1) You can also contact your local Altera sales office or sales representative.

Third-Party Software Product Information

Third-party software products described in this handbook are not Altera products, are licensed by Altera from third parties, and are subject to change without notice. Updates to these third-party software products may not be concurrent with Quartus II software releases. Altera has assumed responsibility for the selection of such third-party software products and its use in the Quartus II 12.0 software release. To the extent that the software products described in this handbook are derived from third-party software, no third party warrants the software, assumes any liability regarding use of the software, or undertakes to furnish you any support or information relating to the software. EXCEPT AS EXPRESSLY SET FORTH IN THE APPLICABLE ALTERA PROGRAM LICENSE SUBSCRIPTION AGREEMENT UNDER WHICH THIS SOFTWARE WAS PROVIDED TO YOU, ALTERA AND THIRD-PARTY LICENSORS DISCLAIM ALL WARRANTIES WITH RESPECT TO THE USE OF SUCH THIRD-PARTY SOFTWARE CODE OR DOCUMENTATION IN THE SOFTWARE, INCLUDING, WITHOUT LIMITATION, ANY WARRANTIES OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE, TITLE, AND NONINFRINGEMENT. For more information, including the latest available version of specific third-party software products, refer to the documentation for the software in question.
# Typographic Conventions

The following table shows the typographic conventions this document uses.

<table>
<thead>
<tr>
<th>Visual Cue</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Bold Type with Initial Capital Letters</strong></td>
<td>Indicate command names, dialog box titles, dialog box options, and other GUI labels. For example, <strong>Save As</strong> dialog box. For GUI elements, capitalization matches the GUI.</td>
</tr>
<tr>
<td><strong>Bold type</strong></td>
<td>Indicates directory names, project names, disk drive names, file names, file name extensions, software utility names, and GUI labels. For example, \qdesigns directory, <strong>D:</strong> drive, and <strong>chiptrip.gdf</strong> file.</td>
</tr>
<tr>
<td><strong>Italic Type with Initial Capital Letters</strong></td>
<td>Indicate document titles. For example, <em>Stratix IV Design Guidelines</em>.</td>
</tr>
<tr>
<td><strong>Italic type</strong></td>
<td>Indicates variables. For example, <code>n + 1</code>. Variable names are enclosed in angle brackets (<code>&lt;&gt;</code>). For example, <code>&lt;file name&gt;</code> and <code>&lt;project name&gt;.pof</code> file.</td>
</tr>
<tr>
<td><strong>Initial Capital Letters</strong></td>
<td>Indicate keyboard keys and menu names. For example, the <strong>Delete key</strong> and the <strong>Options menu</strong>.</td>
</tr>
<tr>
<td><strong>“Subheading Title”</strong></td>
<td>Quotation marks indicate references to sections in a document and titles of Quartus II Help topics. For example, “Typographic Conventions.”</td>
</tr>
<tr>
<td><strong>Courier type</strong></td>
<td>Indicates signal, port, register, bit, block, and primitive names. For example, <code>data1</code>, <code>tdi</code>, and <code>input</code>. The suffix <code>n</code> denotes an active-low signal. For example, <code>resetn</code>. Indicates command line commands and anything that must be typed exactly as it appears. For example, <code>c:\qdesigns\tutorial\chiptrip.gdf</code>. Also indicates sections of an actual file, such as a Report File, references to parts of files (for example, the AHDL keyword <code>SUBDESIGN</code>), and logic function names (for example, <code>TRI</code>).</td>
</tr>
<tr>
<td>←</td>
<td>An angled arrow instructs you to press the Enter key.</td>
</tr>
<tr>
<td>1., 2., 3., and a., b., c., and so on</td>
<td>Numbered steps indicate a list of items when the sequence of the items is important, such as the steps listed in a procedure.</td>
</tr>
<tr>
<td>▪ ▪ ▪</td>
<td>Bullets indicate a list of items when the sequence of the items is not important.</td>
</tr>
<tr>
<td>🌟</td>
<td>The hand points to information that requires special attention.</td>
</tr>
<tr>
<td>🕵️‍♂️</td>
<td>The question mark directs you to a software help system with related information.</td>
</tr>
<tr>
<td>🌐</td>
<td>The feet direct you to another document or website with related information.</td>
</tr>
<tr>
<td>🔗</td>
<td>The multimedia icon directs you to a related multimedia presentation.</td>
</tr>
<tr>
<td>🚧</td>
<td>A caution calls attention to a condition or possible situation that can damage or destroy the product or your work.</td>
</tr>
<tr>
<td>⚠️</td>
<td>A warning calls attention to a condition or possible situation that can cause you injury.</td>
</tr>
<tr>
<td>💌</td>
<td>The envelope links to the Email Subscription Management Center page of the Altera website, where you can sign up to receive update notifications for Altera documents.</td>
</tr>
<tr>
<td>📞</td>
<td>The feedback icon allows you to submit feedback to Altera about the document. Methods for collecting feedback vary as appropriate for each document.</td>
</tr>
<tr>
<td>🌐</td>
<td>The social media icons allow you to inform others about Altera documents. Methods for submitting information vary as appropriate for each medium.</td>
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</table>
The Quartus II Handbook Volume 3: Verification was revised on the following dates.

Chapter 1. Simulating Altera Designs
Revised: May 2013
Part Number: QII53025-13.0.0

Chapter 2. Mentor Graphics ModelSim and QuestaSim Support
Revised: November 2012
Part Number: QII53001-12.1.0

Chapter 3. Synopsys VCS and VCS MX Support
Revised: November 2012
Part Number: QII53002-12.1.0

Chapter 4. Cadence Incisive Enterprise Simulator Support
Revised: May 2013
Part Number: QII53003-13.0.0

Chapter 5. Aldec Active-HDL and Riviera-PRO Support
Revised: November 2012
Part Number: QII53023-12.1.0

Chapter 6. Timing Analysis Overview
Revised: June 2012
Part Number: QII53030-12.0.0

Chapter 7. The Quartus II TimeQuest Timing Analyzer
Revised: June 2012
Part Number: QII53018-12.0.0

Chapter 8. PowerPlay Power Analysis
Revised: November 2012
Part Number: QII53013-12.1.0

Chapter 9. System Debugging Tools Overview
Revised: June 2012
Part Number: QII53027-12.0.0

Chapter 10. Analyzing and Debugging Designs with the System Console
Revised: May 2013
Part Number: QII53028-13.0.0

Chapter 11. Debugging Transceiver Links
Revised: May 2013
Part Number: QII53029-13.0.0

Chapter 12. Quick Design Debugging Using SignalProbe
Revised: May 2013
Part Number: QII53008-13.0.0
Chapter 13. Design Debugging Using the SignalTap II Logic Analyzer
Revised: May 2013
Part Number: QII53009-13.0.0

Revised: June 2012
Part Number: QII53016-12.0.0

Chapter 15. In-System Modification of Memory and Constants
Revised: June 2012
Part Number: QII53012-12.0.0

Chapter 16. Design Debugging Using In-System Sources and Probes
Revised: June 2012
Part Number: QII53021-12.0.0

Chapter 17. Cadence Encounter Conformal Support
Revised: June 2012
Part Number: QII53011-12.0.0

Chapter 18. Quartus II Programmer
Revised: November 2012
Part Number: QII53022-12.1.0
As the design complexity of FPGAs continues to rise, accurate simulation is critical to your overall design efficiency. The Quartus II software provides a wide range of features that support RTL and gate-level simulation in industry standard EDA simulators.

This section includes the following chapters:

- **Chapter 1, Simulating Altera Designs**
  This chapter provides general guidelines to help you simulate Altera® designs in EDA simulators.

- **Chapter 2, Mentor Graphics ModelSim and QuestaSim Support**
  This chapter provides specific guidelines for simulation of Quartus® II designs with Mentor Graphics® ModelSim-Altera®, ModelSim, or QuestaSim software.

- **Chapter 3, Synopsys VCS and VCS MX Support**
  This chapter provides specific guidelines for simulation of Quartus® II designs with the Synopsys VCS or VCS MX software.

- **Chapter 4, Cadence Incisive Enterprise Simulator Support**
  This chapter provides specific guidelines for simulation of Quartus® II designs with the Cadence Incisive Enterprise (IES) software.

- **Chapter 5, Aldec Active-HDL and Riviera-PRO Support**
  This chapter provides specific guidelines for simulation of Quartus® II designs with the Aldec Active-HDL or Riviera-PRO software.
This document describes simulating designs that target Altera® devices. Simulation verifies design behavior before device programming. The Quartus® II software supports RTL and gate level design simulation in third-party EDA simulators.

### Altera Simulation Overview

Simulation involves setting up your simulator working environment, compiling simulation model libraries, and running your simulation. Generate simulation files in an automated or custom flow. Refer to Figure 1–1 and Table 1–3.

**Figure 1–1. Simulation in Quartus II Design Flow**

You can use the Quartus II NativeLink feature to automatically generate simulation files and scripts. NativeLink can launch your simulator a from within the Quartus II software. Use a custom flow for more control over all aspects of simulation file generation.

---

(1) Timing simulation is not supported for Arria® V, Cyclone® V, Stratix® V, and newer families.
Simulator Support

The Quartus II software supports specific versions of the following EDA simulators for RTL and gate-level simulation.

### Table 1–1. Supported Simulators

<table>
<thead>
<tr>
<th>Vendor</th>
<th>Simulator</th>
<th>Platform</th>
</tr>
</thead>
<tbody>
<tr>
<td>Aldec</td>
<td>Active-HDL</td>
<td>Windows</td>
</tr>
<tr>
<td>Aldec</td>
<td>Riviera-PRO</td>
<td>Windows, Linux</td>
</tr>
<tr>
<td>Cadence®</td>
<td>Incisive Enterprise</td>
<td>Linux</td>
</tr>
<tr>
<td>Mentor Graphics</td>
<td>ModelSim-Altera (provided)</td>
<td>Windows, Linux</td>
</tr>
<tr>
<td>Mentor Graphics</td>
<td>ModelSim PE</td>
<td>Windows</td>
</tr>
<tr>
<td>Mentor Graphics</td>
<td>ModelSim® SE</td>
<td>Windows, Linux</td>
</tr>
<tr>
<td>Mentor Graphics</td>
<td>QuestaSim</td>
<td>Windows, Linux</td>
</tr>
<tr>
<td>Synopsys</td>
<td>VCS/VCS MX</td>
<td>Linux</td>
</tr>
</tbody>
</table>

### Simulation Levels

Table 1–2 describes the supported Quartus II simulation levels.

### Table 1–2. Supported Simulation Levels

<table>
<thead>
<tr>
<th>Simulation Level</th>
<th>Description</th>
<th>Simulation Input</th>
</tr>
</thead>
<tbody>
<tr>
<td>RTL</td>
<td>Cycle-accurate simulation using Verilog HDL, SystemVerilog, and VHDL design source code with simulation models provided by Altera and other IP providers.</td>
<td>Design source/testbench&lt;br&gt;Altera simulation libraries&lt;br&gt;Altera IP plain text or IEEE encrypted RTL models&lt;br&gt;IP simulation models&lt;br&gt;Altera IPFS models&lt;br&gt;Altera IP BFM models&lt;br&gt;Qsys-generated models&lt;br&gt;Verification IP</td>
</tr>
<tr>
<td>Gate-level functional</td>
<td>Simulation using a post-synthesis or post-fit functional netlist testing the post-synthesis functional netlist, or post-fit functional netlist.</td>
<td>Testbench&lt;br&gt;Altera simulation libraries&lt;br&gt;Post-synthesis or post-fit functional netlist&lt;br&gt;Altera IP Bus BFM models</td>
</tr>
<tr>
<td>Gate-level timing</td>
<td>Simulation using a post-fit timing netlist, testing design’s functional and timing correctness. Not supported for Arria V, Cyclone V, or Stratix V devices.</td>
<td>Testbench&lt;br&gt;Altera simulation libraries&lt;br&gt;Post-fit timing netlist&lt;br&gt;Post-fit Standard Delay Output File (.sdo)</td>
</tr>
</tbody>
</table>

Gate-level timing simulation of an entire design can be slow and should be avoided. Gate-level timing simulation is not supported for Arria V, Cyclone V, or Stratix V devices. Rely on TimeQuest static timing analysis rather than on gate-level timing simulation.
## Simulation Flows

Table 1–3 describes the supported Quartus II simulation flows.

<table>
<thead>
<tr>
<th>Simulation Flow</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>NativeLink flow</td>
<td>The NativeLink automated flow supports a variety of design flows. NativeLink is not recommended if you require direct control over every aspect of simulation.</td>
</tr>
<tr>
<td><strong>Use NativeLink to generate simulation scripts to compile your design and simulation libraries, and to automatically launch your simulator, as described in “Setting Up Simulation (NativeLink Flow)” on page 1–8.</strong></td>
<td></td>
</tr>
<tr>
<td><strong>Specify your own compilation, elaboration, and simulation scripts for testbench and simulation model files that have not been analyzed by the Quartus II software.</strong></td>
<td></td>
</tr>
<tr>
<td><strong>Use NativeLink to supplement your scripts by automatically compiling:</strong></td>
<td></td>
</tr>
<tr>
<td></td>
<td><em>Design files</em></td>
</tr>
<tr>
<td></td>
<td><em>IP simulation model files</em></td>
</tr>
<tr>
<td></td>
<td><em>Altera simulation library models</em></td>
</tr>
<tr>
<td>Custom flows</td>
<td>Custom flows support manual control of all aspects of simulation, including the following:</td>
</tr>
<tr>
<td><strong>Manually compile and simulate testbench, design, IP, and simulation model libraries, or write scripts to automate compilation and simulation in your simulator.</strong></td>
<td></td>
</tr>
<tr>
<td><strong>Use the Simulation Library Compiler to compile simulation libraries for all Altera devices and supported third-party simulators and languages, as described in “Using IP and Qsys Simulation Setup Scripts (Custom Flow)” on page 1–12.</strong></td>
<td></td>
</tr>
<tr>
<td>Use the custom flow if you require any of the following:</td>
<td></td>
</tr>
<tr>
<td></td>
<td><em>Custom compilation commands for design, IP, or simulation library model files (for example, macros, debugging or optimization options, or other simulator-specific options).</em></td>
</tr>
<tr>
<td></td>
<td><em>Multi-pass simulation flows.</em></td>
</tr>
<tr>
<td></td>
<td><em>Flow that use dynamically generated simulation scripts.</em></td>
</tr>
<tr>
<td>Specialized flows</td>
<td>Altera supports specialized flows for various design variations, including the following:</td>
</tr>
<tr>
<td><strong>For simulation of Altera example designs, refer to the documentation for the example design or to the IP core user guide on the IP and Megafunctions Documentation section of the Altera website.</strong></td>
<td></td>
</tr>
<tr>
<td><strong>For simulation of Qsys designs, refer to Creating a System with Qsys chapter of the Quartus II Handbook.</strong></td>
<td></td>
</tr>
<tr>
<td><strong>For simulation of designs that include the Nios II embedded processor, refer to AN 351: Simulating Nios II Embedded Processors Designs.</strong></td>
<td></td>
</tr>
</tbody>
</table>
HDL Support

Table 1–4 describes Quartus II simulation support for hardware description languages:

<table>
<thead>
<tr>
<th>Language</th>
<th>Description</th>
</tr>
</thead>
</table>
| VHDL         | ■ For VHDL RTL simulation, compile design files directly in your simulator. To use Nativelink automation, analyze and elaborate your design in the Quartus II software, and then use the Nativelink simulator scripts to compile the design files in your simulator. You must also compile simulation models from the Altera simulation libraries and simulation models for the IP cores in your design. Use the Simulation Library Compiler or Nativelink to compile simulation models.  
■ For gate-level simulation, the EDA Netlist Writer generates a synthesized design netlist VHDL Output File (.vho). Compile the .vho in your simulator. You may also need to compile models from the Altera simulation libraries.  
■ IEEE 1364-2005 encrypted Verilog HDL simulation models are encrypted separately for each Altera-supported simulation vendor. If you want to simulate the model in a VHDL design, you need either a simulator that is capable of VHDL/Verilog HDL co-simulation, or any Mentor Graphics single language VHDL simulator. |
| Verilog HDL  | ■ For RTL simulation in Verilog HDL or SystemVerilog, compile your design files in your simulator. To use Nativelink automation, analyze and elaborate your design in the Quartus II software, and then use the Nativelink simulator scripts to compile your design files in your simulator. You must also compile simulation models from the Altera simulation libraries and simulation models for the IP cores in your design. Use the Simulation Library Compiler or Nativelink to compile simulation models.  
■ For gate-level simulation, the EDA Netlist Writer generates a synthesized design netlist Verilog Output File (.vo). Compile the .vo in your simulator. |
| SystemVerilog| ■ If your design is a mix of VHDL and Verilog/SystemVerilog files, you must use a mixed language simulator. Since Altera supports both languages, choose the most convenient language for any Altera IP in your design.  
■ Altera provides Stratix V, Arria V, Cyclone V and newer simulation model libraries and IP simulation models in Verilog HDL and IEEE encrypted Verilog. Your simulator’s co-simulation capabilities support VHDL simulation of these models using VHDL “wrapper” files. Altera provides the wrapper for Verilog models to instantiate these models directly from your VHDL design. |
| Mixed HDL    | You must convert schematics to HDL format before simulation. You can use the converted VHDL or Verilog HDL files for RTL simulation.                                                                           |
System and IP File Locations

The Quartus II software generates the following files for Altera IP cores:

Figure 1–2. System and IP Files Generated by MegaWizard Plug-In Manager and Qsys

MegaWizard-Generated IP Files
- `<instance name>_<EDA_tool_name>.sv`, `.v`, or `.vhd` - simulation model
- `<instance name>_<EDA_tool_name>_sim` (IP simulation files)
- `<sub_module_name>`
- `<EDA_tool_name>`
- `<IEEE_encrypted_Verilog_simulation_models>`

Qsys-Generated System and IP Files
- `<Qsys system name>_<EDA_tool_name>` - EDA simulation files
- `<symbol name>_<EDA_board_symbol_tool_name>` - symbols
- `<timing name>_<EDA_board_timing_tool_name>` - timing analysis tool files
- `<board name>_<EDA_board_tool_name>` - boundary scan tool files

The Quartus II software optionally generates the following files for other EDA tools:

Figure 1–3. Quartus II Generated Files for Other EDA Tools

- `<Quartus II Project Directory>`
- `<simulation>_<EDA_simulator>`
- `<symbols>`
- `<board>`
- `<bsp>`
- `<bsdl>`
Preparing for Simulation

Preparing for RTL or gate-level simulation involves compiling the RTL or gate-level representation of your design and testbench. You must also compile IP simulation models, models from the Altera simulation libraries, and any other model libraries required for your design.

Compiling Simulation Models

The Quartus II software includes simulation models for Altera megafUNCTIONs, primitives, library of parameterized modules (LPMs), IPFS models, and device family specific models in the `<Quartus II installation path>/eda/sim_lib` directory. These models include IEEE encrypted Verilog HDL models for both Verilog HDL and VHDL simulation in the simulators listed in Table 1–1. Before running simulation you must compile the appropriate simulation models from the Altera simulation libraries.

Use any of the following methods to compile Altera simulation models:

- Use the NativeLink feature to automatically compile your design, Altera IP, simulation model libraries, and testbench, as described in “Running RTL Simulation (NativeLink Flow)” on page 1–9.
- Run the Simulation Library Compiler to compile all RTL and gate-level simulation model libraries for your device, simulator, and design language, as described in “Using Simulation Library Compiler (Custom Flow)” on page 1–10.
- Compile Altera simulation models manually with your simulator, as described in Preparing for EDA Simulation in Quartus II Help.

After you compile the simulation model libraries, you can reuse these libraries in subsequent simulations to avoid having to compile them again.

For a complete list of the Altera simulation models, refer to *Altera Simulation Models* in Quartus II Help.

Generating IP Simulation Files for RTL Simulation

The Quartus II software supports both Verilog HDL and VHDL simulation of encrypted and unencrypted Altera IP cores. If your design includes Altera IP cores, you must compile any corresponding IP simulation models in your simulator along with the rest of your design and testbench. The Quartus II software generates and copies the simulation models for IP cores to your project design directory. For information about the location of IP simulation models for the IP cores in your design, refer to “Document Revision History” on page 1–13.

The Quartus II software can generate one or more of the files in Table 1–5 to support the IP core simulation. If generated, use these files to simulate your Altera IP core.
Preparing for Simulation

Generating IP Functional Simulation Models for RTL Simulation

Altera provides IPFS models for some Altera IP cores. To generate IPFS models, follow these steps:

- Turn on the **Generate Simulation Model** option when parameterizing the IP core in the MegaWizard Plug-In Manager.
- When you simulate your design, only compile the `.vo` or `.vho` for these IP cores in your simulator, rather than the corresponding HDL file, which may be encrypted to support only synthesis by the Quartus II software.

---

**Table 1–5. Altera IP Simulation Files**

<table>
<thead>
<tr>
<th>File Type</th>
<th>Description</th>
<th>File Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Simulator setup script</td>
<td>Simulator-specific script to compile, elaborate, and simulate Altera IP models and simulation model library files. Copy the commands into your simulation script, or edit these files to compile, elaborate, and simulate your design and testbench. Refer to “Using IP and Qsys Simulation Setup Scripts (Custom Flow)” on page 1–12.</td>
<td>Cadence</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ cds.lib</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ ncsim_setup.sh</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ hdl.var</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Mentor Graphics</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ msim_setup.tcl</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Synopsys</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ synopsys_sim.setup</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ vcs_setup.sh</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ vcsmx_setup.sh</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Aldec</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ rivierapro_setup.tcl</td>
</tr>
<tr>
<td>Quartus II Simulation IP File (.sip)</td>
<td>Contains IP core simulation library mapping information. .sip files enable NativeLink simulation and the Quartus II Archiver for IP cores.</td>
<td>&lt;design name&gt;.sip</td>
</tr>
<tr>
<td>IPFS models</td>
<td>IP Functional Simulation (IPFS) models are cycle-accurate VHDL or Verilog HDL models generated by the Quartus II software for some Altera IP cores. IPFS models support fast functional simulation of IP using industry-standard VHDL and Verilog HDL simulators. Refer to “Generating IP Functional Simulation Models for RTL Simulation” on page 1–7.</td>
<td>&lt;design name&gt;.vho</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&lt;design name&gt;.vo</td>
</tr>
<tr>
<td>IEEE encrypted models</td>
<td>Stratix V, Arria V, Cyclone V and newer simulation model libraries and IP simulation models are provided in Verilog HDL and IEEE encrypted Verilog HDL. VHDL simulation of these models is supported using your simulator’s co-simulation capabilities. IEEE encrypted Verilog HDL models are significantly faster than IPFS models.</td>
<td>&lt;design name&gt;.v</td>
</tr>
</tbody>
</table>

---

---

- Altera IP cores that do not require IPFS models for simulation lack the **Generate Simulation Model** option in the IP core parameter editor.
Many recently released Altera IP cores support RTL simulation using IEEE Verilog HDL encryption. IEEE encrypted models are significantly faster than IPFS models and can be simulated in both Verilog HDL and VHDL designs.

Generating an IPFS model for some AMPP megafuctions may require a license, refer to _AN 343: OpenCore Evaluation of AMPP Megafunctions._

Running a Simulation (NativeLink Flow)

The NativeLink feature integrates your EDA simulator with the Quartus II software and automates the following simulation steps:

- Set and reuse simulation settings
- Generate simulator-specific files and simulation scripts
- Compile Altera simulation libraries
- Launch your simulator automatically following Quartus II Analysis & Elaboration, Analysis & Synthesis, or after a full compilation.

Setting Up Simulation (NativeLink Flow)

Before running simulation using the NativeLink flow, you must specify settings for your simulator in the Quartus II software. To specify simulation settings in the Quartus II software, follow these steps:

1. Open a Quartus II project.
2. Click **Tools > Options** and specify the location of your simulator executable file.
3. Click **Assignments > Settings** and specify options on the **Simulation** page and the **More NativeLink Settings** dialog box. Specify default options for simulation library compilation, netlist and tool command script generation, and for launching RTL or gate-level simulation automatically following Quartus II processing.
4. If your design includes a testbench, turn on **Compile testbench** and then click **Test Benches** to specify options for each testbench. Alternatively, turn on **Use script to compile testbench** and specify the script file.
5. If you want to use a script to setup simulation, turn on **Use script to setup simulation**.

<table>
<thead>
<tr>
<th>Simulator</th>
<th>Path</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mentor Graphics ModelSim-Altera</td>
<td>&lt;drive letter&gt;:&lt;simulator install path&gt;\win32aloem (Windows) /&lt;simulator install path&gt;\bin (Linux)</td>
</tr>
<tr>
<td>Mentor Graphics ModelSim</td>
<td>&lt;drive letter&gt;:&lt;simulator install path&gt;\win32 (Windows) /&lt;simulator install path&gt;\bin (Linux)</td>
</tr>
<tr>
<td>Mentor Graphics QuestaSim</td>
<td>&lt;drive letter&gt;:&lt;simulator install path&gt;\bin (Linux)</td>
</tr>
<tr>
<td>Synopsys VCS/VCS MX</td>
<td>&lt;simulator install path&gt;\bin (Linux)</td>
</tr>
<tr>
<td>Cadence Incisive Enterprise</td>
<td>&lt;simulator install path&gt;\tools\bin (Linux)</td>
</tr>
<tr>
<td>Aldec Active-HDL</td>
<td>&lt;drive letter&gt;:&lt;simulator install path&gt;\bin (Windows)</td>
</tr>
<tr>
<td>Aldec Riviera-PRO</td>
<td>&lt;simulator install path&gt;\bin (Linux)</td>
</tr>
</tbody>
</table>
Running RTL Simulation (NativeLink Flow)

To run RTL simulation using the NativeLink flow, follow these steps:

1. Set up the simulation environment, as described in “Setting Up Simulation (NativeLink Flow)” on page 1–8.
2. Click Processing > Start > Analysis and Elaboration.
3. Click Tools > Run Simulation Tool > RTL Simulation.

NativeLink compiles simulation libraries and launches and runs your RTL simulator automatically according to the NativeLink settings.
4. Review and analyze the simulation results in your simulator. Correct any functional errors in your design. If necessary, re-simulate the design to verify correct behavior.

Running Gate-Level Simulation (NativeLink Flow)

To run gate-level simulation with the NativeLink flow, follow these steps:

1. Prepare for simulation, as described in “Preparing for Simulation” on page 1–6.
2. Set up the simulation environment, as described in “Setting Up Simulation (NativeLink Flow)” on page 1–8. To generate only a functional (rather than timing) gate-level netlist, click More EDA Netlist Writer Settings, and turn on Generate netlist for functional simulation only.
3. To synthesize the design, follow one of these steps:
   ■ To generate a post-fit functional or post-fit timing netlist and then automatically simulate your design according to your NativeLink settings, Click Processing > Start Compilation. Skip to step 6.
   ■ To synthesize the design for post-synthesis functional simulation only, click Processing > Start > Start Analysis and Synthesis.
4. To generate the simulation netlist, click Start EDA Netlist Writer.
5. Click Tools > Run Simulation Tool > Gate Level Simulation.
6. Review and analyze the simulation results in your simulator. Correct any unexpected or incorrect conditions found in your design. Simulate the design again until you verify correct behavior.

Running a Simulation (Custom Flow)

Use a custom simulation flow to support any of the following more complex simulation scenarios:

■ Custom compilation, elaboration, or run commands for your design, IP, or simulation library model files (for example, macros, debugging/optimization options, simulator-specific elaboration or run-time options)
■ Multi-pass simulation flows
■ Flows that use dynamically generated simulation scripts
Use these to compile libraries and generate simulation scripts for custom simulation flows:

- NativeLink-generated scripts—use NativeLink only to generate simulation script templates to develop your own custom scripts.
- Simulation Library Compiler—compile Altera simulation libraries for your device, HDL, and simulator. Generate scripts to compile simulation libraries as part of your custom simulation flow. This tool does not compile your design, IP, or testbench files.
- IP and Qsys simulation scripts—use the scripts generated for Altera IP cores and Qsys systems as templates to create simulation scripts. If your design includes multiple IP cores or Qsys systems, you can combine the simulation scripts into a single script, manually or by using the `ip-make-simscript` utility, described in “Generating Custom Simulation Scripts with ip-make-simscript” on page 1–12.

Use the following steps in a custom simulation flow:

1. “Preparing for Simulation” on page 1–6.
2. “Using Simulation Library Compiler (Custom Flow)” on page 1–10
5. Compile the design and testbench files in your simulator.
6. Run the simulation in your simulator.

Post-synthesis and post-fit gate-level simulations run significantly slower than RTL simulation. Altera recommends that you verify your design using RTL simulation for functionality and use the TimeQuest timing analyzer for timing. Timing simulation is not supported for Arria V, Cyclone V, Stratix V, and newer families.

For more information about running EDA simulation, refer to Running EDA Simulators in Quartus II Help.

**Using Simulation Library Compiler (Custom Flow)**

Simulation Library Compiler compiles all required Quartus II simulation library files for your HDL, device, and simulator. If your design includes IP cores generated with the classic IP file directory structure in Figure 1–2, you may need to compile additional library files.

If your design includes IP cores generated with the IP file directory structure illustrated in Figure 1–2, refer to “Generating Custom Simulation Scripts with ip-make-simscript” on page 1–12 to use the scripts in combination with the Simulation Library Compiler’s generated simulation scripts.

For detailed steps on using Simulation Library Compiler, refer to Preparing for EDA Simulation in Quartus II Help. For a complete list of the Altera simulation models, refer to Altera Simulation Models in Quartus II Help.
Using NativeLink-Generated Scripts (Custom Flow)

Use the NativeLink feature to generate simulation scripts to automate simulation steps. You can reuse these generated files and simulation scripts in a custom simulation flow. NativeLink optionally generates scripts for your simulator in the project subdirectory described in Table 1–7. To generate simulation scripts using the NativeLink feature, perform the following steps:

1. Click Assignments > Settings.
2. Under EDA Tool Settings, click Simulation.
3. Select the Tool name of your simulator.
4. Click More NativeLink Settings.
5. Turn on Generate third-party EDA tool command scripts without running the EDA tool.

Table 1–7. NativeLink Generated Scripts for RTL Simulation

<table>
<thead>
<tr>
<th>Simulator(s)</th>
<th>Simulation File</th>
<th>Use</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mentor Graphics</td>
<td>/simulation/modelsim/&lt;design&gt;.do</td>
<td>Source directly with your simulator.</td>
</tr>
<tr>
<td>ModelSim</td>
<td></td>
<td></td>
</tr>
<tr>
<td>QuestaSim</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Aldec Riviera Pro</td>
<td>/simulation/modelsim/&lt;design&gt;.do</td>
<td>Source directly with your simulator.</td>
</tr>
<tr>
<td>Synopsys VCS</td>
<td>/simulation/modelsim/&lt;revision name&gt;_&lt;rtl or gate&gt;.vcs</td>
<td>Add your testbench file name to this options file to pass the file to VCS using the -file option. If you specify a testbench file to NativeLink, and direct not to simulate, NativeLink generates an .sh script that runs VCS.</td>
</tr>
<tr>
<td>Synopsys VCS MX:</td>
<td>/simulation/scsim/&lt;revision name&gt;<em>&lt;vcsmx</em>&lt;rtl or gate&gt;_&lt;verilog or vhdl&gt;.tcl</td>
<td>Run this script at the command line using quartus_sh -t &lt;script&gt;.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Any testbench you specify with NativeLink is included in this script.</td>
</tr>
<tr>
<td>Cadence Incisive</td>
<td>/simulation/ncsim/&lt;revision name&gt;<em>&lt;ncsim</em>&lt;rtl or gate&gt;_&lt;verilog or vhdl&gt;.tcl</td>
<td>Run this script at the command line using quartus_sh -t &lt;script&gt;.</td>
</tr>
<tr>
<td>(NC SIM)</td>
<td></td>
<td>Any testbench you specify with NativeLink is included in this script.</td>
</tr>
</tbody>
</table>
Using IP and Qsys Simulation Setup Scripts (Custom Flow)

Altera IP cores and Qsys systems generate simulation setup scripts. Modify these scripts to set up supported simulators. Use the scripts to compile the required device libraries and system design files in the correct order, and then elaborate or load the top-level design for simulation. Also use the script to modify the top-level simulation environment independent of the IP simulation files that are over-written during regeneration.

These simulation scripts variables set up your simulation environment:

- **TOP_LEVEL_NAME**—the top-level entity of your simulation is often a testbench that instantiates your design, and then your design instantiates IP cores and/or Qsys systems. Set the value of `TOP_LEVEL_NAME` to the simulation the top-level entity.
- **QSYS_SIMDIR**—specifies the top-level directory containing the simulation files.
- Other variables control the compilation, elaboration, and simulation process.

Generating Custom Simulation Scripts with `ip-make-simscript`

Use the `ip-make-simscript` utility to generate simulation command scripts for multiple IP cores or Qsys systems. Specify all Simulation Package Descriptor files (.spd), each of which lists the required simulation files for the corresponding IP core or Qsys system. The MegaWizard Plug-In Manager and Qsys generate the .spd files.

This utility compiles IP simulation models into various simulation libraries. Use the `compile-to-work` option to compile all simulation files into a single work library. Use this option only if you require a simplified library structure.

When you specify multiple .spd files, the `ip-make-simscript` utility generates a single simulation script containing all required simulation information. The default value of `TOP_LEVEL_NAME` is the `TOP_LEVEL_NAME` defined in the IP core or Qsys .spd file. If this is not the top-level instance in your design, specify the top-level instance of your testbench or design.

Setting appropriate variables in the script, or edit the variable assignment directly in the script. If the simulation script is a tcl file that can be sourced in the simulator, set the variables before sourcing the script. If the simulation script is a shell script, pass in the variables as command-line arguments to shell script.

- To run `ip-make-simscript`, type the following at the command prompt:

  `<ACDS installation path>\quartus\sopc_builder\bin\ip-make-simscript`
The following are examples of options you can use with the utility:

Table 1–8.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
<th>Status</th>
</tr>
</thead>
<tbody>
<tr>
<td>--spd=&lt;file&gt;</td>
<td>Describes the list of compiled files and memory model hierarchy. If your</td>
<td>Required</td>
</tr>
<tr>
<td></td>
<td>design includes multiple IP cores or Qsys systems that include .spd files,</td>
<td></td>
</tr>
<tr>
<td></td>
<td>use this option for each file. For example:</td>
<td></td>
</tr>
<tr>
<td></td>
<td>ip-make-simscript --spd=ip1.spd --spd=ip2.spd</td>
<td></td>
</tr>
<tr>
<td>--output-directory=&lt;directory&gt;</td>
<td>Directory path specifying the location of output files. If unspecified,</td>
<td>Optional</td>
</tr>
<tr>
<td></td>
<td>the default setting is the directory from which ip-make-simscript is run.</td>
<td></td>
</tr>
<tr>
<td>--compile-to-work</td>
<td>Compiles all design files to the default work library. Use this option</td>
<td>Optional</td>
</tr>
<tr>
<td></td>
<td>only if you encounter problems managing your simulation with multiple</td>
<td></td>
</tr>
<tr>
<td></td>
<td>libraries.</td>
<td></td>
</tr>
<tr>
<td>--use-relative-paths</td>
<td>Uses relative paths whenever possible</td>
<td>Optional</td>
</tr>
</tbody>
</table>

Refer to Aldec Active-HDL and Riviera-PRO Support, Synopsys VCS and VCS MX Support, Cadence Incisive Enterprise Simulator Support, and Mentor Graphics ModelSim and QuestaSim Support for simulation script examples.

Document Revision History

Table 1–9 shows the revision history for this chapter.

Table 1–9. Document Revision History (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>■ Updated introductory section and system and IP file locations.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Revised chapter to reflect latest changes to other simulation</td>
</tr>
<tr>
<td></td>
<td></td>
<td>documentation.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Reorganization of chapter to reflect various simulation flows.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added NativeLink support for newer IP cores.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>■ Added information about encrypted Altera simulation model files.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about IP simulation and NativeLink.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>■ Added note to Figure 1–1 on page 1–2</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added new section “Converting Block Design Files (.bdf) to HDL Format</td>
</tr>
<tr>
<td></td>
<td></td>
<td>(.v/.vhd)” on page 1–4</td>
</tr>
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<td></td>
<td></td>
<td>■ Updated information in “Simulation Netlist Files”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated information in “Generating Gate-Level Timing Simulation</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Netlist Files”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated information in “Generating Post-Synthesis Simulation</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Netlist Files”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed information from “Generating Timing Simulation Netlist Files</td>
</tr>
<tr>
<td></td>
<td></td>
<td>with Different Timing Models”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed information from “Running the Simulation Library Compiler</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Through the GUI”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Table 1–1.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Simulating Qsys and SOPC Builder System Designs”</td>
</tr>
</tbody>
</table>

Document Revision History
For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides specific guidelines for simulation of Quartus® II designs with Mentor Graphics® ModelSim-Altera®, ModelSim, or QuestaSim software. Altera provides the entry-level ModelSim-Altera software, along with precompiled Altera simulation libraries, to simplify simulation of Altera designs. You can also refer to the following for more information about EDA simulation:

- For overview information, Simulating Altera Designs in the Quartus II Handbook and About Using EDA Simulators in Quartus II Help.
- For detailed GUI steps, Preparing for EDA Simulation and Running EDA Simulators in Quartus II Help.
- For support information, ModelSim-Altera Software page of the Altera website, Mentor Graphics ModelSim Simulation Design Examples page.

Quick Start Example (ModelSim Verilog)

You can adapt the following RTL simulation example to get started quickly with ModelSim:

1. Specify your EDA simulator and executable path in the Quartus II software:
   ```
   set_user_option -name EDA_TOOL_PATH_MODELSIM <modelsim executable path>
   set_global_assignment -name EDA_SIMULATION_TOOL "MODELSIM (verilog)
   ```

2. Compile simulation model libraries using one of the following:
   - Run NativeLink RTL simulation to compile required design files, simulation models, and run your simulator. Verify results in your simulator. Skip steps 3 through 5.
   - Use Simulation Library Compiler to compile all required simulation models.
   - Create and map Altera libraries manually:
     ```
     vlib <lib1>_ver
     vmap <lib1>_ver <lib1>_ver
     ```
     Then, compile Altera simulation models manually:
     ```
     vlog -work <lib1> <lib1>
     ```

3. Compile your design and testbench files:
   ```
   vlog -work work <design or testbench name>.v
   ```

4. Load the design:
   ```
   vsim -L work -L <lib1>_ver -L <lib2>_ver work.<testbench name>
   ```

5. Run the simulation in the ModelSim simulator.
In this chapter, “ModelSim” refers to ModelSim SE, PE, and DE, which share the same commands as QuestaSim. “ModelSim-Altera” refers to ModelSim-Altera Starter Edition and ModelSim-Altera Subscription Edition.

**ModelSim, ModelSim-Altera, and QuestaSim Guidelines**

The following guidelines apply to simulation of Altera designs in the ModelSim, ModelSim-Altera, or QuestaSim software.

**Using ModelSim-Altera Precompiled Libraries**

Precompiled libraries for both functional and gate-level simulations are provided for the ModelSim-Altera software. You should not compile these library files before running a simulation. No precompiled libraries are provided for ModelSim or QuestaSim. You must compile the necessary libraries to perform functional or gate-level simulation with these tools.

The precompiled libraries provided in `<ModelSim-Altera path>/altera` must be compatible with the version of the Quartus II software that is used to create the simulation netlist. To check whether the precompiled libraries are compatible with your version of the Quartus II software, refer to the `<ModelSim-Altera path>/altera/version.txt` file. This file shows which version and build of the Quartus II software was used to create the precompiled libraries.

For a list of precompiled library names for all functional and gate-level simulation models, refer to *ModelSim-Altera Precompiled Libraries* in Quartus II Help. For a list of all simulation model files, refer to *Altera Simulation Models* in Quartus II Help.

Encrypted Altera simulation model files shipped with the Quartus II software version 10.1 and later can only be read by ModelSim-Altera Edition Software version 6.6c and later. These encrypted simulation model files are located at the `<Quartus II System directory>/quartus/eda/sim_lib/<mentor>` directory.

**Disabling Timing Violation on Registers**

In certain situations, you may want to ignore timing violations on registers and disable the “X” propagation that occurs (for example, timing violations in internal synchronization registers in asynchronous clock-domain crossing).

By default, the `x_on_violation_option logic` option applying to all design registers is **On**, resulting in an output of “X” at timing violation. To disable “X” propagation at timing violations on a specific register, set the `x_on_violation_option logic` option to **Off** for that register. The following command is an example from the Quartus II Settings File (.qsf):

```
set_instance_assignment -name X_ON_VIOLATION_OPTION OFF -to \<register_name>
```
Passing Parameter Information from Verilog HDL to VHDL

You must use in-line parameters to pass values from Verilog HDL to VHDL. Example 2–1 shows modification to use in-line parameters.

Example 2–1. In-Line Parameter Passing Example

```verilog
lpm_add_sub#(.lpm_width(12), .lpm_direction("Add"),
.lpm_type("LPM_ADD_SUB"),
.lpm_hint("ONE_INPUT_IS_CONSTANT=NO,CIN_USED=NO")
) lpm_add_sub_component (  
.dataa (dataa),
.datab (datab),
.result (sub_wire0)
);
```

The sequence of the parameters depends on the sequence of the GENERIC in the VHDL component declaration.

Increasing Simulation Speed

By default, the ModelSim and QuestaSim software runs in a debug-optimized mode. To run the ModelSim and QuestaSim software in speed-optimized mode, add the following two vlog command-line switches:

```
vlog -fast -05
```

In this mode, module boundaries are flattened and loops are optimized, which eliminates levels of debugging hierarchy and may result in faster simulation. This switch is not supported in the ModelSim-Altera simulator.

Simulating Transport Delays

By default, the ModelSim and QuestaSim software filter out all pulses that are shorter than the propagation delay between primitives. Turning on the transport delay options in the ModelSim and QuestaSim software prevents the simulator from filtering out these pulses.

Table 2–1 describes the transport delay options.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>+transport_path_delays</td>
<td>Use when simulation pulses are shorter than the delay in a gate-level primitive. You must include the +pulse_e/number and +pulse_r/number options.</td>
</tr>
<tr>
<td>+transport_int_delays</td>
<td>Use when simulation pulses are shorter than the interconnect delay between gate-level primitives. You must include the +pulse_int_e/number and +pulse_int_r/number options.</td>
</tr>
</tbody>
</table>

The +transport_path_delays and +transport_path_delays options apply by default during NativeLink gate-level timing simulation.
For more information about either of these options, refer to the ModelSim-Altera Command Reference installed with the ModelSim and QuestaSim software.

The following ModelSim and QuestaSim software command shows the command line syntax to perform a gate-level timing simulation with the device family library:

```
vsim -t 1ps -L stratixii -sdftyp /i=filtref.vhd.sdo work.filtref_vhd_vec_tst \\
+transport_int_delays +transport_path_delays
```

### Viewing Error Messages

ModelSim and QuestaSim error and warning messages are tagged with a `vsim` or `vcom` code. To determine the cause and resolution for a `vsim` or `vcom` error or warning, use the `verror` command.

For example, ModelSim and QuestaSim may display the following error message:

```
# ** Error:
C:/altera_trn/DUALPORT_TRY/simulation/modelsim/DUALPORT_TRY.vho(31):
(vcom-1136) Unknown identifier "stratixiii".
```

In this case, type the following command:

```
verror 1136
```

A description of the error message appears as follows:

```
# vcom Message # 1136:
# The specified name was referenced but was not found. This indicates
# that either the name specified does not exist or is not visible at
# this point in the code.
```

### Generating Power Analysis Files

To generate a timing Value Change Dump File (.vcd) for power analysis, you must first generate a `<filename>_dump_all_vcd_nodes.tcl` script file in the Quartus II software. You can then run the script from the ModelSim, QuestaSim, or ModelSim-Altera software to generate a timing `<filename>.vcd`. You can use this `.vcd` for power analysis in the Quartus II PowerPlay power analyzer.

To use a `.vcd` for power analysis, follow these steps:

1. In the Quartus II software, click **Settings** on the Assignments menu.
2. Click **Simulation** under **EDA Tool Settings**.
3. Turn on **Generate Value Change Dump file script**, specify the type of output signals to include, and specify the top-level design instance name in your testbench.
4. On the Processing menu, click **Start Compilation**.
5. On the Tools menu, point to **Run EDA Simulation**, and then click **EDA Gate Level Simulation**. The Compiler creates the `<filename>_dump_all_vcd_nodes.tcl` file, the ModelSim simulation `<filename>_run_msim_gate_vhdl/verilog.do` file (including the `.vcd` and `.tcl` execution lines), and all other files for simulation. ModelSim then automatically runs the generated `.do` to start the simulation.
6. Break the simulation if your testbench does not have a break point. End the simulation to have ModelSim generate the `.vcd`. You can only generate the `.vcd` after simulation ends with the **End Simulation** function.
For more information about using the timing <filename>.vcd for power estimation, refer to the PowerPlay Power Analysis chapter in volume 3 of the Quartus II Handbook.

**Viewing a Simulation Waveform**

ModelSim-Altera, ModelSim, and QuestaSim automatically generate a Wave Log Format File (.wlf) following simulation. You can use the .wlf to generate a waveform view.

To view a waveform from a .wlf through ModelSim-Altera, ModelSim, or QuestaSim, perform the following steps:

1. Type `vsim` at the command line. The ModelSim/QuestaSim or ModelSim-Altera dialog box appears.
2. On the File menu, click Datasets. The Datasets Browser dialog box appears.
3. Click Open and browse to the directory that contains your .wlf.
4. Select the .wlf file and click Open, then click OK.
5. Click Done.
6. In the Object browser, select the signals that you want to observe.
7. On the Add menu, click Wave and then click Selected Signals.

You cannot view a waveform from a .vcd in ModelSim-Altera, ModelSim, or QuestaSim directly. The .vcd must first be converted to a .wlf.

1. Use the `vcd2wlf` command to convert the file. For example, type the following at the command-line:
   ```
   vcd2wlf <example>.vcd <example>.wlf
   ```
2. After you convert the .vcd to a .wlf, follow the procedures for viewing a waveform from a .wlf through ModelSim and QuestaSim.

You can also convert your .wlf to a .vcd by using the `wlf2vcd` command.

**Simulating with ModelSim-Altera Waveform Editor**

You can use the ModelSim-Altera Waveform Editor as a simple method to create stimulus vectors for simulation. You can create this design stimulus via interactive manipulation of waveforms from the wave window in ModelSim-Altera. With the ModelSim-Altera waveform editor, you can create and edit waveforms, drive simulation directly from created waveforms, and save created waveforms into a stimulus file.

For more information, refer to the Generating Stimulus with Waveform Editor chapter in the ModelSim SE User’s Manual available on the ModelSim website (www.model.com).
Simulation Setup Script Example

The Quartus II software can generate a `msim_setup.tcl` simulation setup script for IP cores in your design. The script compiles the required device library models, compiles the design files, and elaborates the design with or without simulator optimization. To run the script, type `source msim_setup.tcl` in the simulator Transcript window. Alternatively, if you are using the simulator at the command line, you can type the following command:

```bash
vsim -c -do msim_setup.tcl.
```

Example 2–2 shows the `top-level-simulate.do` custom top-level simulation script that sets the hierarchy variable `TOP_LEVEL_NAME` to `top_testbench` for the design, and sets the variable `QSYS_SIMDIR` to the location of the generated simulation files.

**Example 2–2. Example Top Level Simulation Script (top-level-simulate.do)**

```bash
# Set hierarchy variables used in the IP-generated files
set TOP_LEVEL_NAME "top_testbench"
set QSYS_SIMDIR "./ip_top_sim"

# Source generated simulation script which defines aliases used below
source $QSYS_SIMDIR/mentor/msim_setup.tcl

# dev_com alias compiles simulation libraries for device library files
dev_com

# com alias compiles IP simulation or Qsys model files and/or Qsys model
# files in the correct order
com

# Compile top level testbench that instantiates your IP
vlog -sv ./top_testbench.sv

# elab alias elaborates the top-level design and testbench
elab

# Run the full simulation
run -all
```

In this example, the top-level simulation files are stored in the same directory as the original IP core, so this variable is set to the IP-generated directory structure. The `QSYS_SIMDIR` variable provides the relative hierarchy path for the generated IP simulation files. The script calls the generated `msim_setup.tcl` script and uses the alias commands from the script to compile and elaborate the IP files required for simulation along with the top-level simulation testbench. You can specify additional simulator elaboration command options when you run the `elab` command, for example, `elab +nowarnTFMPC`. The last command run in the example starts the simulation.

Unsupported Features

The Quartus II software does not support the following simulation features:

- Altera does not support companion licensing for ModelSim AE.
- The USB software guard is not supported by versions earlier than Mentor Graphics ModelSim software version 5.8d.
- For ModelSim-Altera software versions prior to 5.5b, use the PCLS utility included with the software to set up the license.
Some versions of ModelSim and QuestaSim support SystemVerilog, PSL assertions, SystemC, and more. For more information about specific feature support, refer to Mentor Graphics literature.

For more information about the ModelSim-Altera Subscription Edition software, including pricing, refer to the ModelSim-Altera Software page of the Altera website. For more information about obtaining and setting up the license for the ModelSim-Altera Subscription Edition software, refer to the “Licensing Altera Software” section in the Altera Software Installation and Licensing Manual.

Document Revision History

Table 2–2 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
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<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Relocated general simulation information to Simulating Altera Designs.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>Added information about encrypted Altera simulation model files.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated power analysis information.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Updated “Software Requirements” on page 2–2</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Updated “Design Flow with ModelSim-Altera, ModelSim, or QuestaSim Software” on page 2–2</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Restructured “Simulating with the ModelSim-Altera Software” on page 2–4</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Restructured “Simulating with the ModelSim and QuestaSim Software” on page 2–5</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Restructured “Simulating Designs that Include Transceivers” on page 2–12</td>
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<tr>
<td></td>
<td></td>
<td>Changed section name from “ModelSim and QuestaSim Error Message Verification” to “ModelSim and QuestaSim Error Message Information” on page 2–18</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Changed section name from “Simulating with ModelSim-Altera Waveform” to “Simulating with ModelSim-Altera Waveform Editor” on page 2–20</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Changed to new document template</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Referenced Simulating Altera Designs chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added new section, “Simulating with ModelSim-Altera Waveform Editor” on page 2–20</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed Stratix V compilation information and linked to Quartus II Help</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Removed simulation library tables and linked to Quartus II Help</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added other links to Quartus II Help and ModelSim-Altera Help where appropriate and removed redundant information</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added QuestaSim support</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added Stratix V simulation information</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Minor editorial changes throughout</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed Referenced Documents section</td>
</tr>
</tbody>
</table>
Table 2–2. Document Revision History  (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
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<th>Changes</th>
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</thead>
</table>
| November 2009 | 9.1.0   | • Removed NativeLink information and referenced new *Simulating Designs with EDA Tools* chapter  
• Added Stratix IV transceiver simulation section  
• Reformatted transceiver simulation sections  
• Text edits throughout chapter |
| March 2009    | 9.0.0   | • Added the following sections:  
  ▪ “Compile Libraries Using the EDA Simulation Library Compiler” on page 2–17  
  ▪ “Generate Simulation Script from EDA Nettlist Writer” on page 2–77  
  ▪ “Viewing a Waveform from a .wlf File” on page 2–78  
Updated the following:  
  ▪ Table 2–1, Table 2–2, Table 2–5, Table 2–6, Table 2–7, Table 2–8, Table 2–9, Table 2–10  
  ▪ Figure 2–4 on page 2–81  
  ▪ All sections titled “Loading the Design” |
| November 2008 | 8.1.0   | • Updated the following:  
  ▪ Table 2–2, Table 2–3, Table 2–4, Table 2–5, Table 2–6  
  ▪ Removed *--zero_ic_delays* from *quartus_sta* option in “Generate Post-Synthesis Simulation Nettlist Files” on page 2–11  
  ▪ Removed steps to include the library when the simulation is run in VHDL mode from all procedures; this is no longer necessary  
  ▪ Added information about the Altera Simulation Library Compiler throughout the chapter  
  ▪ Added “Compile Libraries Using the Altera Simulation Library Compiler” on page 2–15  
  ▪ Added “Disabling Simulation” on page 2–72  
  ▪ Minor editorial updates  
  ▪ Updated entire chapter using 8½” × 11” chapter template |
| May 2008      | 8.0.0   | • Updated the following:  
  ▪ “Altera Design Flow with ModelSim-Altera or ModelSim Software” on page 2–3  
  ▪ “Simulation Libraries” on page 2–4  
  ▪ “Simulation Nettlist Files” on page 2–11  
  ▪ “Perform Simulation Using ModelSim-Altera Software” on page 2–15  
  ▪ “Perform Simulation Using ModelSim Software” on page 2–33  
  ▪ “Simulating Designs that Include Transceivers” on page 2–57  
  ▪ “Using the NativeLink Feature with ModelSim-Altera or ModelSim Software” on page 2–63  
  ▪ “Generating a Timing VCD File for PowerPlay” on page 2–68 |
This chapter provides specific guidelines for simulation of Quartus® II designs with the Synopsys VCS or VCS MX software. You can also refer to the following for more information about EDA simulation:

- For overview and version support information, *Simulating Altera Designs* in the *Quartus II Handbook* and *About Using EDA Simulators* in Quartus II Help.
- For detailed GUI steps, *Preparing for EDA Simulation* and *Running EDA Simulators* in Quartus II Help.
- The *VCS User Guide* installed with the VCS software, and the *Synopsys VCS Simulation Design Example* page.

**Quick Start Example (VCS Verilog)**

You can adapt the following RTL simulation example to get started quickly with VCS:

1. Specify your EDA simulator and executable path in the Quartus II software:
   ```
   set_user_option -name EDA_TOOL_PATH_VCS <VCS executable path>
   set_global_assignment -name EDA_SIMULATION_TOOL "VCS"
   ```
2. Compile simulation model libraries using one of the following:
   - Run NativeLink RTL simulation to compile required design files, simulation models, and run your simulator. Verify results in your simulator. Skip steps 3 through 4.
   - Use Simulation Library Compiler to generate the `simlib_comp.vcs` options file that contains VCS command-line arguments specifying required simulation models.
3. Modify the `simlib_comp.vcs` file to specify your design and testbench files.
4. Run the VCS simulator:
   ```
   vcs -R -file simlib_comp.vcs
   ```

**VCS and VCS MX Guidelines**

The following guidelines apply to simulating Quartus II designs in the VCS or VCS MX software:

- Do not specify the `-v` option for `altera_lnsim.sv` because it defines a systemverilog package.
- Add `-verilog` and `+verilog2001ext+.v` options to make sure all `.v` files are compiled as verilog 2001 files, and all other files are compiled as systemverilog files.
Add the -lca option for Stratix V and later families because they include IEEE-encrypted simulation files for VCS and VCS MX.

Add -timescale=1ps/1ps to ensure picosecond resolution.

Disabling Timing Violation on Registers

In certain situations, you may want to ignore timing violations on registers and disable the “X” propagation that occurs (for example, timing violations in internal synchronization registers in asynchronous clock-domain crossing).

By default, the x_on_violation_option logic option applying to all design registers is On, resulting in an output of “X” at timing violation. To disable “X” propagation at timing violations on a specific register, set the x_on_violation_option logic option to Off for that register. The following command is an example from the Quartus II Settings File (.qsf):

```
set_instance_assignment -name X_ON_VIOLATION_OPTION OFF -to \ <register_name>
```

Simulating Transport Delays

By default, the VCS software filters out all pulses that are shorter than the propagation delay between primitives. Turning on the transport delay options in the VCS software prevents the simulation tool from filtering out these pulses. Use the following options to ensure that simulation results include all signal pulses. Table 3–1 describes the transport delay options.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>+transport_path_delays</td>
<td>Use when simulation pulses are shorter than the delay in a gate-level primitive. You must include the +pulse_e/number and +pulse_r/number options.</td>
</tr>
<tr>
<td>+transport_int_delays</td>
<td>Use when simulation pulses are shorter than the interconnect delay between gate-level primitives. You must include the +pulse_int_e/number and +pulse_int_r/number options.</td>
</tr>
</tbody>
</table>

The +transport_path_delays and +transport_path_delays options apply by default during NativeLink gate-level timing simulation.

The following VCS software command runs a post-synthesis simulation:

```
vcs -R <testbench>.v <gate-level netlist>.v -v <Altera device family \ library>.v +transport_int_delays +pulse_int_e/0 +pulse_int_r/0 \ +transport_path_delays +pulse_e/0 +pulse_r/0
```

Generating Power Analysis Files

To run power analysis in the Quartus II software, you must first generate a Verilog Value Change Dump File (.vcd) in the Quartus II software, and then run the .vcd from the VCS software. You can then use this .vcd for power analysis in the Quartus II PowerPlay power analyzer. To use a .vcd for power analysis, follow these steps:
1. In the Quartus II software, click **Settings** on the Assignments menu.

2. Click **Simulation** under **EDA Tool Settings**.

3. Turn on **Generate Value Change Dump file script**, specify the type of output signals to include, and specify the top-level design instance name in your testbench.

4. On the Processing menu, click **Start Compilation**.

5. Include the script in your testbench file where the design under test (DUT) is instantiated:

   ```
   include <revision_name>_dump_all_vcd_nodes.v
   ```

   Include the script within the testbench module block. If you include the script outside of the testbench module block, syntax errors occur during compilation.

6. Run the simulation with the VCS command. Exit the VCS software when the simulation is finished and the `<revision_name>.vcd` file is generated in the simulation directory.

   For detailed instructions about generating a `.vcd` file and running power analysis, refer to the **PowerPlay Power Analysis** chapter in volume 3 of the *Quartus II Handbook*.

### Simulation Setup Script Example

The Quartus II software can generate a simulation setup script for IP cores in your design. The scripts for VCS and VCS MX are **vcs_setup.sh** (for Verilog HDL or SystemVerilog) and **vcsmx_setup.sh** (combined Verilog HDL and SystemVerilog with VHDL). The scripts contain shell commands that compile the required simulation models in the correct order, elaborate the top-level design, and run the simulation for 100 time units by default. You can run these scripts from a Linux command shell.

Read the generated `.sh` script to see the variables that are available for override when sourcing the script or redefining directly if you edit the script. To set up the simulation for a design such as Example 3–1, use the command-line to pass variable values to the shell script, as illustrated in Example 3–2 on page 3–4 and Example 3–3 on page 3–4. You can alternatively create a shell script that contains these commands.

**Example 3–1. Using Command-line to Pass Simulation Variables**

```
sh vcsmx_setup.sh
USER_DEFINED_ELAB_OPTIONS=+rad\nUSER_DEFINED_SIM_OPTIONS=+vcs+lic+wait
```
You can edit the .sh script to add simulator commands that compile the top-level simulation HDL file.

**Example 3–2. Example Top-Level Simulation Shell Script for VCS-MX**

```bash
# Run generated script to compile libraries and IP simulation files
# Skip elaboration and simulation of the IP variation
sh ./ip_top_sim/synopsys/vcsmx/vcsmx_setup.sh SKIP_ELAB=1 SKIP_SIM=1
QSYS_SIMDIR="./ip_top_sim"
# Compile top-level testbench that instantiates IP
vlogan -sverilog ./top_testbench.sv
# Elaborate and simulate the top-level design
vcs -lca -t ps <elaboration control options> top_testbench
simv <simulation control options>
```

**Example 3–3. Example Top-Level Simulation Shell Script for VCS**

```bash
# Run script to compile libraries and IP simulation files
sh ./ip_top_sim/synopsys/vcs/vcs_setup.sh
TOP_LEVEL_NAME="top_testbench"
# Pass VCS elaboration options to compile files and elaborate top-level
# passed to the script as the TOPLEVEL_NAME
USER_DEFINED_ELAB_OPTIONS="top_testbench.sv"
# Pass in simulation options and run the simulation for specified amount
time.
USER_DEFINED_SIM_OPTIONS="<simulation control options>"
```

---

**Document Revision History**

Table 3–2 shows the revision history for this chapter.

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<thead>
<tr>
<th>Date</th>
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<td>Removed survey link.</td>
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<tr>
<td>November 2011</td>
<td>11.0.1</td>
<td>Template update. Minor editorial updates.</td>
</tr>
</tbody>
</table>
| May 2011       | 11.0.0  | - Linked to Help for Stratix V Libraries
|                |         | - Added SystemVerilog HDL information                      |
|                |         | - Editorial updates throughout                             |
| December 2010  | 10.0.1  | Changed to new document template. No change to content.     |
| July 2010      | 10.0.0  | - Linked to Quartus II Help where appropriate               |
|                |         | - Added Stratix V simulation information                    |
|                |         | - Minor text edits                                         |
|                |         | - Removed VirSim references                                |
|                |         | - Removed Referenced Documents section                     |
| November 2009  | 9.1.0   | - Removed NativeLink information and referenced new *Simulating Designs with EDA Tools* chapter in volume 3 of the *Quartus II Handbook*
|                |         | - Added “RTL Functional Simulation for Stratix IV Devices” and “Gate-Level Timing Simulation for Stratix IV Devices” sections |
|                |         | - Minor text edits                                         |
Table 3–2. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
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<th>Changes</th>
</tr>
</thead>
</table>
| March 2009   | 9.0.0   | ■ Added support for Synopsys VCS MX software  
■ Changed chapter title to “Synopsys VCS and VCS MX Support”  
■ Major revision to “Compiling Libraries Using the EDA Simulation Library Compiler” on page 4–2  
■ Major revision to “RTL Functional Simulations” on page 4–2  
■ Added Table 3–4 on page 3–10 and Table 3–5 on page 3–11  
■ Added new section “Using DVE” on page 4–7  
■ Added new section “Generating a Simulation Script from the EDA Netlist Writer” on page 3–16  
■ Added new section “Viewing a Waveform from a .vpd or .vcd File” on page 4–13 |
| November 2008| 8.1.0   | ■ Added “Compile Libraries Using the EDA Simulation Library Compiler” on page 3–3  
■ Added information about the --simlib_comp utility  
■ Updated entire chapter using 8½” × 11” chapter template  
■ Minor editorial updates |

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides specific guidelines for simulation of Quartus® II designs with the Cadence Incisive Enterprise (IES) software. You can also refer to the following for more information about EDA simulation:

- For overview and version support information, Simulating Altera Designs in the Quartus II Handbook and About Using EDA Simulators in Quartus II Help.
- For detailed GUI steps, Preparing for EDA Simulation and Running EDA Simulators in Quartus II Help.

Quick Start Example (NC-Verilog)

You can adapt the following RTL simulation example to get started quickly with IES:

1. Specify your EDA simulator and executable path in the Quartus II software:
   ```
   set_user_option -name EDA_TOOL_PATH_NCSIM <ncsim executable path>
   set_global_assignment -name EDA_SIMULATION_TOOL "NC-Verilog (Verilog)"
   ```

2. Compile simulation model libraries using one of the following:
   - Run NativeLink RTL simulation to compile required design files, simulation models, and run your simulator. Verify results in your simulator. Skip steps 3 through 4.
   - Use Simulation Library Compiler to compile all required simulation models.
   - Map Altera simulation libraries by adding the following commands to a `cds.lib` file:
     ```
     include ${CDS_INST_DIR}/tools/inca/files/cds.lib
     DEFINE <lib1>_ver <lib1_ver>
     ```
   Then, compile Altera simulation models manually:
   ```
   vlog -work <lib1_ver>
   ```

3. Elaborate your design and testbench with IES:
   ```
   ncelab <work library>..<top-level entity name>
   ```

4. Run the simulation:
   ```
   ncsim <work library>..<top-level entity name>
   ```

Cadence Incisive Enterprise Guidelines

The following guidelines apply to simulation of Altera designs in the IES software.
Simulation Tool Interfaces

Altera supports both the IES GUI and command-line simulator interfaces. To start the IES GUI, type the following command at a command prompt:

```
nclaunch
```

Table 4–1 lists the ISE command-line programs supported for IES simulation.

<table>
<thead>
<tr>
<th>Program</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>ncvlog</td>
<td>ncvlog compiles your Verilog HDL code and performs syntax and static semantics checks.</td>
</tr>
<tr>
<td>ncvhdl</td>
<td>ncvhdl compiles your VHDL code and performs syntax and static semantics checks.</td>
</tr>
<tr>
<td>ncelab</td>
<td>Elaborates the design hierarchy and determines signal connectivity.</td>
</tr>
<tr>
<td>ncsdfc</td>
<td>Performs back-annotation for simulation with VHDL simulators.</td>
</tr>
<tr>
<td>ncsim</td>
<td>Runs mixed-language simulation. This program is the simulation kernel that performs event scheduling and executes the simulation code.</td>
</tr>
</tbody>
</table>

Elaborating Your Design

The simulator automatically reads the .sdo file during elaboration of the Quartus II-generated Verilog HDL or SystemVerilog HDL netlist file. The ncelab executable recognizes the embedded system task `$sdf_annotate` and automatically compiles and annotates the .sdo file (runs ncsdfc automatically). VHDL netlist files do not contain system task calls to locate your .sdf file; therefore, you must compile the standard .sdo file manually. Locate the .sdo file in the same directory where you run elaboration or simulation. Otherwise, the `$sdf_annotate` task cannot reference the .sdo file correctly. If you are starting an elaboration or simulation from a different directory, you can either comment out the `$sdf_annotate` and annotate the .sdo file with the GUI, or add the full path of the .sdo file.

If you use NC-Sim for post-fit VHDL functional simulation of a Stratix V design that includes RAM, an elaboration error might occur if the component declaration parameters are not in the same order as the architecture parameters. Use the `-namemap_mixgen` option with the ncelab command to match the component declaration parameter and architecture parameter names.

Back-Annotating Simulation Timing Data (VHDL Only)

You can back annotate timing information in a Standard Delay Output File (.sdo) for VHDL simulators. To back annotate the .sdo timing data at the command line, follow these steps:

1. To compile the .sdo with the ncsdfc program, type the following command at the command prompt:

   ```
   ncsdfc <project name>_vhd.sdo -output <output name>
   ```

   The ncsdfc program generates an `<output name>.sdf.X` compiled .sdo file.

   If you do not specify an output name, ncsdfc uses `<project name>.sdo.X`. 
2. Specify the compiled .sdf file for the project by adding the following command to an ASCII SDF command file for the project:

VICE_ID_SDF_FILE = "<project name>.sdf" SCOPE = <instance path>

Example 4–1 shows an example of an SDF command file.

Example 4–1. SDF Command File

```plaintext
// SDF command file sdf_file
COMPiled_SDF_FILE = "lpm_ram_dp_test_vhd.sdo.X",
SCOPE = :tb,
MTM_CONTROL = "TYPICAL",
SCALE_FACTORS = "1.0:1.0:1.0",
SCALE_TYPE = "FROM_MTM";
```

After you compile the .sdf file, type the following command to elaborate the design:

```
ncelab worklib.<project name>:entity -SDF_CMD_FILE <SDF Command File>
```

Disabling Timing Violations on Registers

In certain situations, you may want to ignore timing violations on registers and disable the “X” propagation that occurs (for example, timing violations in internal synchronization registers in asynchronous clock-domain crossing).

By default, the **x_on_violation_option logic** option applying to all design registers is **On**, resulting in an output of “X” at timing violation. To disable “X” propagation at timing violations on a specific register, set the **x_on_violation_option logic** option to **Off** for that register. The following command is an example from the Quartus II Settings File (.qsf):

```
set_instance_assignment -name X_ON_VIOLATION_OPTION OFF -to <register_name>
```

Simulating Pulse Reject Delays

By default, the IES software filters out all pulses that are shorter than the propagation delay between primitives. Setting the pulse reject delays options in the IES software prevents the simulation tool from filtering out these pulses. Use the following options to ensure that all signal pulses are seen in the simulation results.

Table 4–2 describes the pulse reject delay options.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-PULSE_R</td>
<td>Use when simulation pulses are shorter than the delay in a gate-level primitive. The argument is the percentage of delay for pulse reject limit for the path.</td>
</tr>
<tr>
<td>-PULSE_INT_R</td>
<td>Use when simulation pulses are shorter than the interconnect delay between gate-level primitives. The argument is the percentage of delay for pulse reject limit for the path.</td>
</tr>
</tbody>
</table>

The -PULSE_R and -PULSE_INT_R options apply by default during NativeLink gate-level timing simulation.
To perform a gate-level timing simulation with the device family library, type the following IES software command:

```
ncelab worklib.<project name>:entity -SDF_CMD_FILE <SDF Command File> \ 
-TIMESCALE 1ps/1ps -PULSE_R 0 -PULSE_INT_R 0
```

### Viewing a Simulation Waveform

IES generates a .trn file automatically following simulation. You can use the .trn for generating the SimVision waveform view.

To view a waveform from a .trn file through SimVision, follow these steps:
1. Type `simvision` at the command line. The Design Browser dialog box appears.
2. In the File menu, click Open Database and click the .trn file.
3. In the Design Browser dialog box, select the signals that you want to observe from the Hierarchy.
4. Right-click the selected signals and click Send to Waveform Window.

You cannot view a waveform from a .vcd file in SimVision, and the .vcd file cannot be converted to a .trn file.

### Simulation Setup Script Example

The Quartus II software can generate a ncsim_setup.sh simulation setup script for IP cores in your design. The script contains shell commands that compile the required device libraries, IP, or Qsys simulation models in the correct order. The script then elaborates the top-level design and runs the simulation for 100 time units by default. You can run these scripts from a Linux command shell.

To set up the simulation script for a design, you can use the command-line to pass variable values to the shell script, as illustrated in Example 4–2.

Read the generated .sh script to see the variables that are available for you to override when you source the script or that you can redefine directly in the generated .sh script. For example, you can specify additional elaboration and simulation options with the variables `USER_DEFINED_ELAB_OPTIONS` and `USER_DEFINED_SIM_OPTIONS`.

#### Example 4–2. Example Top-Level Simulation Shell Script for Incisive (NCSIM)

```bash
# Run script to compile libraries and IP simulation files
# Skip elaboration and simulation of the IP variation
sh ./ip_top_sim/cadence/ncsim_setup.sh SKIP_ELAB=1 SKIP_SIM=1 QSYS_SIMDIR="./ip_top_sim"

#Compile the top-level testbench that instantiates your IP
ncvlog -sv ./top_testbench.sv

#Elaborate and simulate the top-level design
ncelab <elaboration control options> top_testbench
ncsim <simulation control options> top_testbench
```
Document Revision History

Table 4–3 shows the revision history for this chapter.

Table 4–3. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
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<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>Added note about parameter mismatch workaround.</td>
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<tr>
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<td>Relocated general simulation information to Simulating Altera Designs.</td>
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<td>Removed survey link.</td>
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<td>Template update. Minor editorial updates.</td>
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<td>■ Changed chapter title</td>
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<td></td>
<td></td>
<td>■ Linked to Help for Stratix V Libraries</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added SystemVerilog HDL information</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Other minor changes throughout</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Changed to new document template. No change to content.</td>
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<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Linked to Help where appropriate</td>
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<td></td>
<td>■ Minor text edits</td>
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<tr>
<td></td>
<td></td>
<td>■ Removed Referenced Documents section</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Removed NativeLink information and referenced new Simulating Designs with EDA Tools chapter in volume 3 of the Quartus II Handbook</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “RTL Functional Simulation for Stratix IV Devices” and “Gate-Level Timing Simulation for Stratix IV Devices” sections</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor text edits</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ Removed “Compile Libraries Using the Altera Simulation Library Compiler”</td>
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<td></td>
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<td>■ Added “Compile Libraries Using the EDA Simulation Library Compiler” on page 4–5</td>
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<td></td>
<td>■ Added “Generate Simulation Script from EDA Netlist Writer” on page 4–35</td>
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<td>■ Added “Viewing a Waveform from a .trn File” on page 4–36</td>
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<td></td>
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<td>■ Added information about the --simlib_comp utility.</td>
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<td>■ Minor editorial updates.</td>
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<td></td>
<td></td>
<td>■ Updated entire chapter using 8½” × 11” chapter template.</td>
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<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Updated Table 4–1.</td>
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<td></td>
<td>■ Updated Figure 4–1.</td>
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<td></td>
<td>■ Updated “Compilation in Command-Line Mode” on page 4–9.</td>
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<td>■ Updated “Generating a Timing Netlist with Different Timing Models” on page 4–18</td>
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<td></td>
<td>■ Added “Disable Timing Violation on Registers” on page 4–20.</td>
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<td></td>
<td></td>
<td>■ Updated “Simulating Designs that Include Transceivers” on page 4–23.</td>
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<td></td>
<td>■ Updated “Performing a Gate Level Simulation Using NativeLink” on page 4–30.</td>
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<td></td>
<td></td>
<td>■ Added “Generating a Timing VCD File for PowerPlay” on page 4–33.</td>
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<td></td>
<td></td>
<td>■ Added hyperlinks to referenced documents throughout the chapter.</td>
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<td></td>
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<td>■ Minor editorial updates.</td>
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</table>
For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
5. Aldec Active-HDL and Riviera-PRO Support

This chapter provides specific guidelines for simulation of Quartus® II designs with the Aldec Active-HDL or Riviera-PRO software. You can also refer to the following for more information about EDA simulation:

- For overview and version support information, Simulating Altera Designs in the Quartus II Handbook and About Using EDA Simulators in Quartus II Help.
- For detailed GUI steps, Preparing for EDA Simulation and Running EDA Simulators in Quartus II Help.

Quick Start Example (Active-HDL VHDL)

You can adapt the following RTL simulation example to get started quickly with Active-HDL:

1. Specify your EDA simulator and executable path in the Quartus II software:
   ```
   set_user_option -name EDA_TOOL_PATH_ACTIVEHDL <active-hdl executable path> 
   set_global_assignment -name EDA_SIMULATION_TOOL "Active-HDL (VHDL)"
   ```

2. Compile simulation model libraries using one of the following:
   - Use NativeLink to compile required design files, simulation models, and run your simulator. Verify results in your simulator. Skip steps 3 through 6.
   - Use Simulation Library Compiler to compile all required simulation models.
   - Compile Altera simulation models manually:
     ```
     vlib <library1> <altera_library1>
     vcom -strict93 -dbg -work <library1> <lib1_component/pack.vhd> <lib1.vhd> 
     ```

3. Create and open the workspace:
   ```
   createdesign <workspace name> <workspace path> 
   opendesign -a <workspace name>.adf
   ```

4. Create the work library and compile the netlist and testbench files:
   ```
   vlib work 
   vcom -strict93 -dbg -work <output netlist> <testbench file>
   ```

5. Load the design:
   ```
   vsim +access+r -t 1ps +transport_int_delays +transport_path_delays \ 
   -L work -L <lib1> -L <lib2> work.<testbench module name> 
   ```

6. Run the simulation in the Active-HDL simulator.
Active-HDL and Riviera-PRO Guidelines

The following guidelines apply to simulating Altera designs in the Active-HDL or Riviera-PRO software.

Compiling SystemVerilog Files

If your design includes multiple SystemVerilog files, you must compile the System Verilog files together with a single `alog` command.

If you have Verilog files and SystemVerilog files in your design, it is recommended that you compile the Verilog files, and then compile only the SystemVerilog files in the single `alog` command.

Simulating Transport Delays

By default, the Active-HDL or Riviera-PRO software filters out all pulses that are shorter than the propagation delay between primitives. Turning on the transport delay options in the Active-HDL or Riviera-PRO software prevents the simulation tool from filtering out these pulses.

Table 5–1 describes the transport delay options.

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>+transport_path_delays</code></td>
<td>Use when simulation pulses are shorter than the delay in a gate-level primitive. You must include the <code>+pulse_e/number</code> and <code>+pulse_r/number</code> options.</td>
</tr>
<tr>
<td><code>+transport_int_delays</code></td>
<td>Use when simulation pulses are shorter than the interconnect delay between gate-level primitives. You must include the <code>+pulse_int_e/number</code> and <code>+pulse_int_r/number</code> options.</td>
</tr>
</tbody>
</table>

The `+transport_path_delays` and `+transport_path_delays` options apply by default during NativeLink gate-level timing simulation.

For more information about either of these options, refer to the Active-HDL online documentation installed with the Active-HDL software.

To perform a gate-level timing simulation with the device family library, type the Active-HDL command shown in Example 5–1.

Example 5–1.

```
vsim -t 1ps -L stratixii -sdftyp /i1=filtref_vhd.sdo \
work.filtref_vhd_vec_tst +transport_int_delays +transport_path_delays
```

Disabling Timing Violation on Registers

In certain situations, you may want to ignore timing violations on registers and disable the “X” propagation that occurs (for example, timing violations in internal synchronization registers in asynchronous clock-domain crossing).
By default, the `x_on_violation_option logic` option applying to all design registers is **On**, resulting in an output of “X” at timing violation. To disable “X” propagation at timing violations on a specific register, set the `x_on_violation_option logic` option to **Off** for that register. The following command is an example from the Quartus II Settings File (.qsf):

```
set_instance_assignment -name X_ON_VIOLATION_OPTION OFF -to \<register_name>
```

For VHDL designs, the back-annotating process is done by adding the `-sdftyp` option.

**Example**

```
vsim +access +r -t lps +transport_int_delays +transport_path_delays
-sdftyp <instance path to design>=<path to SDO file> -L adder -L work
-L lpm -L altera_mf work.adder_vhd_vec_tst
```

### Using Simulation Setup Scripts

The Quartus II software can generate a `rivierapro_setup.tcl` simulation setup script for IP cores in your design. The use and content of the script file is similar to the `msim_setup.tcl` file described in the *Mentor Graphics ModelSim and QuestaSim Support* chapter of the *Quartus II Handbook*.

### Document Revision History

Table 5–2 shows the revision history for this chapter.

**Table 5–2. Document Revision History**

<table>
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<th>Date</th>
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<tr>
<td>November 2012</td>
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<td>Template update.</td>
</tr>
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</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Linked to Help for Stratix V Libraries.</td>
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<tr>
<td></td>
<td></td>
<td>Reorganized and reformatted chapter</td>
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<td>July 2010</td>
<td>10.0.0</td>
<td>Linked to Quartus II Help</td>
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<tr>
<td></td>
<td></td>
<td>Revised simulation procedures</td>
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<td></td>
<td></td>
<td>Added Stratix V simulation information</td>
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<td></td>
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<td>Added Riviera-PRO support</td>
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<td>Minor text edits</td>
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<td>Removed Referenced Documents section</td>
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### Table 5-2. Document Revision History

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<th>Date</th>
<th>Version</th>
<th>Changes</th>
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</thead>
</table>
| November 2000 | 9.1.0   | - Updated Table 6–1  
- Removed Simulation Library tables and EDA Simulation Library Compiler sections and referenced new *Simulating Designs with EDA Tools* chapter  
- Added “RTL Functional Simulation for Stratix IV Devices” and “Gate-Level Timing Simulation for Stratix IV Devices” sections  
- Minor text edits |
| March 2009   | 9.0.0   | - Removed “Compile Libraries Using the Altera Simulation Library Compiler”  
- Added “Compile Libraries Using the EDA Simulation Library Compiler” on page 5–10  
- Added “Generate Simulation Script from EDA Netlist Writer” on page 5–51  
- Minor editorial updates |
| November 2008 | 8.1.0   | Added the following sections:  
- “Compile Libraries Using the Altera Simulation Library Compiler” on page 5–10  
- Added steps to the procedure “Performing an RTL Simulation Using NativeLink” on page 5–45 for using the Altera Simulation Library Compilation  
- Added steps to the procedure “Performing a Gate-Level Timing Simulation Using NativeLink” on page 5–47 for using the Altera Simulation Library Compilation  
- Minor editorial updates  
- Updated entire chapter using 8½” × 11” chapter template |
| May 2008     | 8.0.0   | Initial release |

For previous versions of the *Quartus II Handbook*, refer to the [*Quartus II Handbook Archive*](#).
As designs become more complex, advanced timing analysis capability requirements grow. Static timing analysis is a method of analyzing, debugging, and validating the timing performance of a design. The Quartus® II software provides the features necessary to perform advanced timing analysis for today’s system-on-a-programmable-chip (SOPC) designs.

Synopsys PrimeTime is an industry standard sign-off tool, used to perform static timing analysis on most ASIC designs. The Quartus II software provides a path to enable you to run PrimeTime on your Quartus II software designs, and export a netlist, timing constraints, and libraries to the PrimeTime environment.

This section explains the basic principles of static timing analysis, the advanced features supported by the Quartus II TimeQuest Timing Analyzer, and how you can use PrimeTime to analyze your Quartus II projects.

This section includes the following chapters:

- **Chapter 6, Timing Analysis Overview**
  This chapter describes static timing analysis in the context of the TimeQuest Timing Analyzer. The chapter focuses on the relationships and equations that are central to timing analysis.

- **Chapter 7, The Quartus II TimeQuest Timing Analyzer**
  This chapter describes the Quartus II TimeQuest Timing Analyzer, which is a powerful ASIC-style timing analysis tool that validates the timing performance of all logic in your design using an industry-standard constraint, analysis, and reporting methodology.
6. Timing Analysis Overview

Comprehensive static timing analysis involves analysis of register-to-register, I/O, and asynchronous reset paths. Timing analysis with the TimeQuest Timing Analyzer uses data required times, data arrival times, and clock arrival times to verify circuit performance and detect possible timing violations. The TimeQuest analyzer determines the timing relationships that must be met for the design to correctly function, and checks arrival times against required times to verify timing. This chapter is an overview of the concepts you need to know to analyze your designs with the TimeQuest analyzer.

For more information about the TimeQuest analyzer flow and TimeQuest examples, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.

TimeQuest Terminology and Concepts

Table 6–1 describes TimeQuest analyzer terminology.

<table>
<thead>
<tr>
<th>Term</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>nodes</td>
<td>Most basic timing netlist unit. Used to represent ports, pins, and registers.</td>
</tr>
<tr>
<td>cells</td>
<td>Look-up tables (LUT), registers, digital signal processing (DSP) blocks, memory blocks, input/output elements, and so on. (1)</td>
</tr>
<tr>
<td>pins</td>
<td>Inputs or outputs of cells.</td>
</tr>
<tr>
<td>nets</td>
<td>Connections between pins.</td>
</tr>
<tr>
<td>ports</td>
<td>Top-level module inputs or outputs; for example, device pins.</td>
</tr>
<tr>
<td>clocks</td>
<td>Abstract objects representing clock domains inside or outside of your design.</td>
</tr>
</tbody>
</table>

Notes to Table 6–1:
(1) For Stratix® devices, the LUTs and registers are contained in logic elements (LE) and modeled as cells.

Timing Netlists and Timing Paths

The TimeQuest analyzer requires a timing netlist to perform timing analysis on any design. After you generate a timing netlist, the TimeQuest analyzer uses the data to help determine the different design elements in your design and how to analyze timing.
The Timing Netlist

Figure 6–1 shows a sample design for which the TimeQuest analyzer generates a timing netlist equivalent.

Figure 6–1. Sample Design

Figure 6–2 shows the timing netlist for the sample design in Figure 6–1, including how different design elements are divided into cells, pins, nets, and ports.

Figure 6–2. The TimeQuest Analyzer Timing Netlist

Timing Paths

Timing paths connect two design nodes, such as the output of a register to the input of another register. Understanding the types of timing paths is important to timing closure and optimization. The TimeQuest analyzer uses the following commonly analyzed paths:

- Edge paths—connections from ports-to-pins, from pins-to-pins, and from pins-to-ports.
Clock paths—connections from device ports or internally generated clock pins to the clock pin of a register.

Data paths—connections from a port or the data output pin of a sequential element to a port or the data input pin of another sequential element.

Asynchronous paths—connections from a port or asynchronous pins of another sequential element such as an asynchronous reset or asynchronous clear.

Figure 6–3 shows path types commonly analyzed by the TimeQuest analyzer.

Data and Clock Arrival Times

After the TimeQuest analyzer identifies the path type, it can report data and clock arrival times at register pins.

The TimeQuest analyzer calculates data arrival time by adding the launch edge time to the delay from the clock source to the clock pin of the source register, the micro clock-to-output delay ($\mu t_{CO}$) of the source register, and the delay from the source register’s data output (Q) to the destination register’s data input (D).

The TimeQuest analyzer calculates data required time by adding the latch edge time to the sum of all delays between the clock port and the clock pin of the destination register, including any clock port buffer delays, and subtracts the micro setup time ($\mu t_{SU}$) of the destination register, where the $\mu t_{SU}$ is the intrinsic setup time of an internal register in the FPGA. Figure 6–4 shows the flow calculated for data arrival time and data required time.
Equation 6–1 shows the basic calculations for data arrival and data required times including the launch and latch edges.

**Equation 6–1. Data Arrival and Data Required Time Equations**

Data Arrival Time = Launch Edge + Source Clock Delay + $\mu t_{CO}$ + Register-to-Register Delay  
Data Required Time = Latch Edge + Destination Clock Delay - $\mu t_{SU}$

**Launch and Latch Edges**

All timing relies on one or more clocks. In addition to analyzing paths, the TimeQuest analyzer determines clock relationships for all register-to-register transfers in your design. Figure 6–5 shows the launch edge, which is the clock edge that sends data out of a register or other sequential element, and acts as a source for the data transfer. A latch edge is the active clock edge that captures data at the data port of a register or other sequential element, acting as a destination for the data transfer. In this example, the launch edge sends the data from register `reg1` at 0 ns, and the register `reg2` captures the data when triggered by the latch edge at 10 ns. The data arrives at the destination register before the next latch edge.

In timing analysis, and with the TimeQuest analyzer specifically, you create clock constraints and assign those constraints to nodes in your design. These clock constraints provide the structure required for repeatable data relationships. The primary relationships between clocks, in the same or different domains, are the setup relationship and the hold relationship. Figure 6–5 also shows the setup and hold relationships between a launch edge and a latch edge which are 10 ns apart.

**Figure 6–5. Launch and Latch Edges**

If you do not constrain the clocks in your design, the Quartus II software analyzes in terms of a 1 GHz clock to maximize timing based Fitter effort. To ensure realistic slack values, you must constrain all clocks in your design with real values.
Clock Setup Check

To perform a clock setup check, the TimeQuest analyzer determines a setup relationship by analyzing each launch and latch edge for each register-to-register path. For each latch edge at the destination register, the TimeQuest analyzer uses the closest previous clock edge at the source register as the launch edge. Figure 6–6 shows two setup relationships, setup A and setup B. For the latch edge at 10 ns, the closest clock that acts as a launch edge is at 3 ns and is labeled setup A. For the latch edge at 20 ns, the closest clock that acts as a launch edge is 19 ns and is labeled setup B. TimQuest analyzes the most restrictive setup relationship, in this case setup B; if that relationship meets the design requirement, then setup A meets it by default.

![Figure 6–6. Setup Check](image)

The TimeQuest analyzer reports the result of clock setup checks as slack values. Slack is the margin by which a timing requirement is met or not met. Positive slack indicates the margin by which a requirement is met; negative slack indicates the margin by which a requirement is not met. Equation 6–2 shows the TimeQuest analyzer clock setup slack time calculation for internal register-to-register paths.

**Equation 6–2. Clock Setup Slack for Internal Register-to-Register paths**

\[
\text{Clock Setup Slack} = \text{Data Required Time} - \text{Data Arrival Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay to Source Register} + \mu t_{CO} + \text{Register-to-Register Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} - \mu t_{SU} - \text{Setup Uncertainty}
\]

The TimeQuest analyzer performs setup checks using the maximum delay when calculating data arrival time, and minimum delay when calculating data required time.

Equation 6–3 shows the TimeQuest analyzer clock setup slack time calculation if the data path is from an input port to an internal register.

**Equation 6–3. Clock Setup Slack from Input Port to Internal Register**

\[
\text{Clock Setup Slack} = \text{Data Required Time} - \text{Data Arrival Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay} + \text{Input Maximum Delay} + \text{Port-to-Register Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} - \mu t_{SU} - \text{Setup Uncertainty}
\]
Equation 6–4 shows the TimeQuest analyzer clock setup slack time calculation if the data path is an internal register to an output port.

**Equation 6–4. Clock Setup Slack from Internal Register to Output Port**

\[
\text{Clock Setup Slack} = \text{Data Required Time} - \text{Data Arrival Time} \\
\text{Data Required Time} = \text{Latch Edge + Clock Network Delay to Output Port} - \text{Output Maximum Delay} \\
\text{Data Arrival Time} = \text{Launch Edge + Clock Network Delay to Source Register} + \mu t_{\text{CO}} + \text{Register-to-Port Delay}
\]

**Clock Hold Check**

To perform a clock hold check, the TimeQuest analyzer determines a hold relationship for each possible setup relationship that exists for all source and destination register pairs. The TimeQuest analyzer checks all adjacent clock edges from all setup relationships to determine the hold relationships. The TimeQuest analyzer performs two hold checks for each setup relationship. The first hold check determines that the data launched by the current launch edge is not captured by the previous latch edge. The second hold check determines that the data launched by the next launch edge is not captured by the current latch edge. From the possible hold relationships, the TimeQuest analyzer selects the hold relationship that is the most restrictive. The most restrictive hold relationship is the hold relationship with the smallest difference between the latch and launch edges and determines the minimum allowable delay for the register-to-register path. Figure 6–7 shows two setup relationships, setup A and setup B, and their respective hold checks. In this example, the TimeQuest analyzer selects hold check A2 as the most restrictive hold relationship.

**Figure 6–7. Hold Checks**

Equation 6–5 shows the TimeQuest analyzer clock hold slack time calculation.

**Equation 6–5. Clock Hold Slack for Internal Register-to-Register Paths**

\[
\text{Clock Hold Slack} = \text{Data Arrival Time} - \text{Data Required Time} \\
\text{Data Arrival Time} = \text{Launch Edge + Clock Network Delay to Source Register} + \mu t_{\text{CO}} + \text{Register-to-Register Delay} \\
\text{Data Required Time} = \text{Latch Edge + Clock Network Delay to Destination Register} + \mu t_{\text{H}} + \text{Hold Uncertainty}
\]
The TimeQuest analyzer performs hold checks using the minimum delay when calculating data arrival time, and maximum delay when calculating data required time.

Equation 6–6 shows the TimeQuest analyzer hold slack time calculation if the data path is from an input port to an internal register.

**Equation 6–6. Clock Hold Slack from Input Port to Internal Register**

\[
\text{Clock Hold Slack} = \text{Data Arrival Time} - \text{Data Required Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay} + \text{Input Minimum Delay} + \text{Pin-to-Register Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} + \mu t_H
\]

Equation 6–7 shows the TimeQuest analyzer hold slack time calculation if the data path is from an internal register to an output port.

**Equation 6–7. Clock Hold Slack from Internal Register to Output Port**

\[
\text{Clock Hold Slack} = \text{Data Arrival Time} - \text{Data Required Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay to Source Register} + \mu t_{CO} + \text{Register-to-Pin Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay} - \text{Output Minimum Delay}
\]

**Recovery and Removal Time**

Recovery time is the minimum length of time for the deassertion of an asynchronous control signal relative to the next clock edge; for example, signals such as clear and preset must be stable before the next active clock edge. The recovery slack calculation is similar to the clock setup slack calculation, but it applies to asynchronous control signals. **Equation 6–8** shows the TimeQuest analyzer recovery slack time calculation if the asynchronous control signal is registered.

**Equation 6–8. Recovery Slack if Asynchronous Control Signal Registered**

\[
\text{Recovery Slack Time} = \text{Data Required Time} - \text{Data Arrival Time}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} - \mu t_{SU}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay to Source Register} + \mu t_{CO} + \text{Register-to-Register Delay}
\]

**Equation 6–9** shows the TimeQuest analyzer recovery slack time calculation if the asynchronous control signal is not registered.

**Equation 6–9. Recovery Slack if Asynchronous Control Signal not Registered**

\[
\text{Recovery Slack Time} = \text{Data Required Time} - \text{Data Arrival Time}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} - \mu t_{SU}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay} + \text{Input Maximum Delay} + \text{Port-to-Register Delay}
\]
If the asynchronous reset signal is from a device I/O port, you must create an input delay constraint for the asynchronous reset port for the TimeQuest analyzer to perform recovery analysis on the path.

Removal time is the minimum length of time the deassertion of an asynchronous control signal must be stable after the active clock edge. The TimeQuest analyzer removal slack calculation is similar to the clock hold slack calculation, but it applies asynchronous control signals. Equation 6–10 shows the TimeQuest analyzer removal slack time calculation if the asynchronous control signal is registered.

**Equation 6–10. Removal Slack if Asynchronous Control Signal Registered**

\[
\text{Removal Slack Time} = \text{Data Arrival Time} - \text{Data Required Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay to Source Register} + \mu t_{CO} + \text{Register-to-Register Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} + \mu t_{H}
\]

Equation 6–11 shows the TimeQuest analyzer removal slack time calculation if the asynchronous control signal is not registered.

**Equation 6–11. Removal Slack if Asynchronous Control Signal not Registered**

\[
\text{Removal Slack Time} = \text{Data Arrival Time} - \text{Data Required Time}
\]

\[
\text{Data Arrival Time} = \text{Launch Edge} + \text{Clock Network Delay} + \text{Input Minimum Delay of Pin} + \text{Minimum Pin-to-Register Delay}
\]

\[
\text{Data Required Time} = \text{Latch Edge} + \text{Clock Network Delay to Destination Register} + \mu t_{H}
\]

If the asynchronous reset signal is from a device pin, you must assign the **Input Minimum Delay** timing assignment to the asynchronous reset pin for the TimeQuest analyzer to perform removal analysis on the path.
Multicycle Paths

Multicycle paths are data paths that require a non-default setup and/or hold relationship for proper analysis. For example, a register may be required to capture data on every second or third rising clock edge. Figure 6–8 shows an example of a multicycle path between the input registers of a multiplier and an output register where the destination latches data on every other clock edge.

Figure 6–8. Multicycle Path

Figure 6–9 shows a register-to-register path used for the default setup and hold relationship, the respective timing diagrams for the source and destination clocks, and the default setup and hold relationships, when the source clock, src_clk, has a period of 10 ns and the destination clock, dst_clk, has a period of 5 ns. The default setup relationship is 5 ns; the default hold relationship is 0 ns.

Figure 6–9. Register-to-Register Path and Default Setup and Hold Timing Diagram

To accommodate the system requirements you can modify the default setup and hold relationships with a multicycle timing exception.
Figure 6–10 shows the actual setup relationship after you apply a multicycle timing exception. The exception has a multicycle setup assignment of two to use the second occurring latch edge; in this example, to 10 ns from the default value of 5 ns.

**Figure 6–10. Modified Setup Diagram**

For more information about creating exceptions with multicycle paths, refer to The Quartus II TimeQuest Timing Analyzer chapter of the Quartus II Handbook.

**Metastability**

Metastability problems can occur when a signal is transferred between circuitry in unrelated or asynchronous clock domains because the designer cannot guarantee that the signal will meet setup and hold time requirements. To minimize the failures due to metastability, circuit designers typically use a sequence of registers, also known as a synchronization register chain, or synchronizer, in the destination clock domain to resynchronize the data signals to the new clock domain.

The mean time between failures (MTBF) is an estimate of the average time between instances of failure due to metastability.

The TimeQuest analyzer analyzes the potential for metastability in your design and can calculate the MTBF for synchronization register chains. The MTBF of the entire design is then estimated based on the synchronization chains it contains.

In addition to reporting synchronization register chains found in the design, the Quartus II software also protects these registers from optimizations that might negatively impact MTBF, such as register duplication and logic retiming. The Quartus II software can also optimize the MTBF of your design if the MTBF is too low.

For more information about metastability, its effects in FPGAs, and how MTBF is calculated, refer to the Understanding Metastability in FPGAs white paper. For more information about metastability analysis, reporting, and optimization features in the Quartus II software, refer to the Managing Metastability with the Quartus II Software chapter in volume 1 of the Quartus II Handbook.

**Common Clock Path Pessimism Removal**

Common clock path pessimism removal accounts for the minimum and maximum delay variation associated with common clock paths during static timing analysis by adding the difference between the maximum and minimum delay value of the common clock path to the appropriate slack equation.
Minimum and maximum delay variation can occur when two different delay values are used for the same clock path. For example, in a simple setup analysis, the maximum clock path delay to the source register is used to determine the data arrival time. The minimum clock path delay to the destination register is used to determine the data required time. However, if the clock path to the source register and to the destination register share a common clock path, both the maximum delay and the minimum delay are used to model the common clock path during timing analysis. The use of both the minimum delay and maximum delay results in an overly pessimistic analysis since two different delay values, the maximum and minimum delays, cannot be used to model the same clock path.

Figure 6–11 shows a typical register-to-register path with the maximum and minimum delay values shown.

**Figure 6–11. Common Clock Path**

Segment A is the common clock path between \texttt{reg1} and \texttt{reg2}. The minimum delay is 5.0 ns; the maximum delay is 5.5 ns. The difference between the maximum and minimum delay value equals the common clock path pessimism removal value; in this case, the common clock path pessimism is 0.5 ns. The TimeQuest analyzer adds the common clock path pessimism removal value to the appropriate slack equation to determine overall slack. Therefore, if the setup slack for the register-to-register path in Figure 6–11 equals 0.7 ns without common clock path pessimism removal, the slack would be 1.2 ns with common clock path pessimism removal.

You can also use common clock path pessimism removal to determine the minimum pulse width of a register. A clock signal must meet a register’s minimum pulse width requirement to be recognized by the register. A minimum high time defines the minimum pulse width for a positive-edge triggered register. A minimum low time defines the minimum pulse width for a negative-edge triggered register.
Clock pulses that violate the minimum pulse width of a register prevent data from being latched at the data pin of the register. To calculate the slack of the minimum pulse width, the TimeQuest analyzer subtracts the required minimum pulse width time from the actual minimum pulse width time. The TimeQuest analyzer determines the actual minimum pulse width time by the clock requirement you specified for the clock that feeds the clock port of the register. The TimeQuest analyzer determines the required minimum pulse width time by the maximum rise, minimum rise, maximum fall, and minimum fall times. Figure 6–12 shows a diagram of the required minimum pulse width time for both the high pulse and low pulse.

**Figure 6–12. Required Minimum Pulse Width**

With common clock path pessimism, the minimum pulse width slack can be increased by the smallest value of either the maximum rise time minus the minimum rise time, or the maximum fall time minus the minimum fall time. For Figure 6–12, the slack value can be increased by 0.2 ns, which is the smallest value between 0.3 ns (0.8 ns – 0.5 ns) and 0.2 ns (0.9 ns – 0.7 ns).

For more information, refer to *TimeQuest Timing Analyzer Page (Settings Dialog Box)* in Quartus II Help.

**Clock-As-Data Analysis**

The majority of FPGA designs contain simple connections between any two nodes known as either a data path or a clock path. A data path is a connection between the output of a synchronous element to the input of another synchronous element. A clock is a connection to the clock pin of a synchronous element. However, for more complex FPGA designs, such as designs that use source-synchronous interfaces, this simplified view is no longer sufficient. Clock-as-data analysis is performed in circuits with elements such as clock dividers and DDR source-synchronous outputs.
The connection between the input clock port and output clock port can be treated either as a clock path or a data path. Figure 6–13 shows a design where the path from port `clk_in` to port `clk_out` is both a clock and a data path. The clock path is from the port `clk_in` to the register `reg_data` clock pin. The data path is from port `clk_in` to the port `clk_out`.

**Figure 6–13. Simplified Source Synchronous Output**

With clock-as-data analysis, the TimeQuest analyzer provides a more accurate analysis of the path based on user constraints. For the clock path analysis, any phase shift associated with the phase-locked loop (PLL) is taken into consideration. For the data path analysis, any phase shift associated with the PLL is taken into consideration rather than ignored.

The clock-as-data analysis also applies to internally generated clock dividers. Figure 6–14 shows an internally generated clock divider. In this figure, waveforms are for the inverter feedback path, analyzed during timing analysis. The output of the divider register is used to determine the launch time and the clock port of the register is used to determine the latch time.

**Figure 6–14. Clock Divider**
Multicorner Analysis

The TimeQuest analyzer performs multicorner timing analysis to verify your design under a variety of operating conditions—such as voltage, process, and temperature—while performing static timing analysis.

To change the operating conditions or speed grade of the device used for timing analysis, use the `set_operating_conditions` command.

For more information about the `set_operating_conditions` and `get_available_operating_conditions` commands—including full syntax information, options, and example usage—refer to `set_operating_conditions` and `get_available_operating_conditions` in Quartus II Help.

If you specify an operating condition Tcl object, the -model, speed, -temperature, and -voltage options are optional. If you do not specify an operating condition Tcl object, the -model option is required; the -speed, -temperature, and -voltage options are optional.

To obtain a list of available operating conditions for the target device, use the `get_available_operating_conditions -all` command.

To ensure that no violations occur under various conditions during the device operation, perform static timing analysis under all available operating conditions. Table 6–2 shows the operating conditions for the slow and fast timing models for device families that support only slow and fast operating conditions.

### Table 6–2. Operating Conditions for Slow and Fast Models

<table>
<thead>
<tr>
<th>Model</th>
<th>Speed Grade</th>
<th>Voltage</th>
<th>Temperature</th>
</tr>
</thead>
<tbody>
<tr>
<td>Slow</td>
<td>Slowest speed grade in device density</td>
<td>$V_{cc}$ minimum supply $^{(1)}$</td>
<td>Maximum $T_J$ $^{(1)}$</td>
</tr>
<tr>
<td>Fast</td>
<td>Fastest speed grade in device density</td>
<td>$V_{cc}$ maximum supply $^{(1)}$</td>
<td>Minimum $T_J$ $^{(1)}$</td>
</tr>
</tbody>
</table>

**Note to Table 6–2:**

(1) Refer to the DC & Switching Characteristics chapter of the applicable device Handbook for $V_{cc}$ and $T_J$ values.

Example 6–1 shows how to set the operating conditions in Example 6–2 with a Tcl object.

**Example 6–1. Setting Operating Conditions with a Tcl Object**

```
set_operating_conditions 3_slow_1100mv_85c
```

Example 6–2 shows how to set the operating conditions for a Stratix III design to the slow timing model, with a voltage of 1100 mV, and temperature of 85° C.

**Example 6–2. Setting Operating Conditions with Individual Values**

```
set_operating_conditions -model slow -temperature 85 -voltage 1100
```
Example 6–3 shows how to use the set_operating_conditions command to generate different reports for various operating conditions.

**Example 6–3. Script Excerpt for Analysis of Various Operating Conditions**

```plaintext
# Specify initial operating conditions
set_operating_conditions -model slow -speed 3 -grade c -temperature 85 -voltage 1100

# Update the timing netlist with the initial conditions
update_timing_netlist

# Perform reporting

# Change initial operating conditions. Use a temperature of 0°C
set_operating_conditions -model slow -speed 3 -grade c -temperature 0 -voltage 1100

# Update the timing netlist with the new operating condition
update_timing_netlist

# Perform reporting

# Change initial operating conditions. Use a temperature of 0°C and a model of fast
set_operating_conditions -model fast -speed 3 -grade c -temperature 0 -voltage 1100

# Update the timing netlist with the new operating condition
update_timing_netlist

# Perform reporting
```

**Document Revision History**

Table 6–3 shows the revision history for this document.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Added social networking icons, minor text updates</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>Initial release.</td>
</tr>
</tbody>
</table>

For previous versions of the *Quartus II Handbook*, refer to the [Quartus II Handbook Archive](#).
The Quartus® II TimeQuest Timing Analyzer is a powerful ASIC-style timing analysis tool that validates the timing performance of all logic in your design using an industry-standard constraint, analysis, and reporting methodology. Use the TimeQuest analyzer GUI or command-line interface to constrain, analyze, and report results for all timing paths in your design.

This chapter contains the following sections:

- “Getting Started with the TimeQuest Analyzer”
- “Constraining and Analyzing with Tcl Commands” on page 7–7
- “Creating Clocks and Clock Constraints” on page 7–14
- “Creating I/O Requirements” on page 7–24
- “Creating Delay and Skew Constraints” on page 7–26
- “Creating Timing Exceptions” on page 7–27
- “Examples of Basic Multicycle Exceptions” on page 7–35
- “Application of Multicycle Exceptions” on page 7–44
- “Timing Reports” on page 7–57

For more information about basic timing analysis concepts and how they pertain to the TimeQuest analyzer, refer to the Timing Analysis Overview chapter in volume 3 of the Quartus II Handbook.

For more information about Altera resources available for the TimeQuest analyzer, refer to the TimeQuest Timing Analyzer Resource Center of the Altera website.

For more information about the TimeQuest analyzer, refer to the Altera Training page of the Altera website.
Getting Started with the TimeQuest Analyzer

This section provides an overview of the design steps for setting up your project for timing and analysis and to constrain your design with the TimeQuest analyzer.

Running the TimeQuest Analyzer

To run the TimeQuest analyzer directly from the Quartus II software GUI, click **TimeQuest Timing Analyzer** on the Tools menu.

For more information about the TimeQuest analyzer GUI, refer to *About TimeQuest Timing Analysis* in Quartus II Help.

To run the TimeQuest analyzer as a stand-alone GUI application, type the following command at the command prompt:

```
quartus_staw
```

To run the TimeQuest analyzer in command-line mode for easy integration with scripted design flows, type the following command at a system command prompt:

```
quartus_sta -s
```

Table 7–1 describes the available command-line options.

<table>
<thead>
<tr>
<th>Command-Line Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-h</td>
<td>--help</td>
</tr>
<tr>
<td>-t &lt;script file&gt;</td>
<td>--script=&lt;script file&gt;</td>
</tr>
<tr>
<td>-s</td>
<td>--shell</td>
</tr>
<tr>
<td>--tcl_eval &lt;tcl command&gt;</td>
<td>Evaluates the Tcl command &lt;tcl command&gt;.</td>
</tr>
<tr>
<td>--do_report_timing</td>
<td>For all clocks in the design, run the following commands:</td>
</tr>
<tr>
<td></td>
<td>report_timing -npaths 1 -to_clock $clock</td>
</tr>
<tr>
<td></td>
<td>report_timing -setup -npaths 1 -to_clock $clock</td>
</tr>
<tr>
<td></td>
<td>report_timing -hold -npaths 1 -to_clock $clock</td>
</tr>
<tr>
<td></td>
<td>report_timing -recovery -npaths 1 -to_clock $clock</td>
</tr>
<tr>
<td></td>
<td>report_timing -removal -npaths 1 -to_clock $clock</td>
</tr>
<tr>
<td>--force_dat</td>
<td>Forces an update of the project database with new delay information.</td>
</tr>
<tr>
<td>--lower_priority</td>
<td>Lowers the computing priority of the quartus_sta process.</td>
</tr>
<tr>
<td>--post_map</td>
<td>Uses the post-map database results.</td>
</tr>
<tr>
<td>--sdc=&lt;SDC file&gt;</td>
<td>Specifies the .sdc to use.</td>
</tr>
<tr>
<td>--report_script=&lt;script&gt;</td>
<td>Specifies a custom report script to call.</td>
</tr>
<tr>
<td>--speed=&lt;value&gt;</td>
<td>Specifies the device speed grade used for timing analysis.</td>
</tr>
<tr>
<td>--tq2hc</td>
<td>Generate temporary files to convert the TimeQuest analyzer .sdc file(s) to a PrimeTime .sdc that can be used by the HardCopy® Design Center.</td>
</tr>
<tr>
<td>--tq2pt</td>
<td>Generates temporary files to convert the TimeQuest analyzer .sdc file(s) to a PrimeTime .sdc.</td>
</tr>
<tr>
<td>-f &lt;argument file&gt;</td>
<td>Specifies a file containing additional command-line arguments.</td>
</tr>
</tbody>
</table>
For more information about steps to perform before opening the TimeQuest analyzer, refer to “Recommended Flow” on page 7–3.

For more information about using Tcl commands to constrain and analyze your design, refer to “Constraining and Analyzing with Tcl Commands” on page 7–7.

**Recommended Flow**

The Quartus II TimeQuest analyzer performs constraint validation to timing verification as part of the compilation flow. Figure 7–1 shows the recommended design flow to maximize the benefits of the TimeQuest Analyzer.

**Table 7–1. Summary of Command-Line Options (Part 2 of 2)**

<table>
<thead>
<tr>
<th>Command-Line Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>-c &lt;revision name&gt;</td>
<td>Specifies which revision and its associated .qsf to use.</td>
</tr>
<tr>
<td>-rev=&lt;revision_name&gt;</td>
<td></td>
</tr>
<tr>
<td>-multicorner</td>
<td>Specifies that all slack summary reports be generated for both slow- and fast-corners.</td>
</tr>
<tr>
<td>-multicorner[=on</td>
<td>off]</td>
</tr>
<tr>
<td>-voltage=&lt;value_in_mV&gt;</td>
<td>Specifies the device voltage, in mV used for timing analysis.</td>
</tr>
<tr>
<td>-temperature=&lt;value_in_C&gt;</td>
<td>Specifies the device temperature in degrees Celsius, used for timing analysis.</td>
</tr>
<tr>
<td>-parallel[=&lt;num_processors&gt;]</td>
<td>Specifies the number of computer processors to use on a multiprocessor system.</td>
</tr>
<tr>
<td>-64bit</td>
<td>Enables 64-bit version of the executable.</td>
</tr>
</tbody>
</table>

Figure 7–1. Design Flow with the TimeQuest Timing Analyzer

Create Quartus II Project and Specify Design Files

Perform Initial Compilation

Specify Timing Requirements

Perform Compilation

Verify Timing
Creating and Setting Up your Design

You must first create your project in the Quartus II software. Include all the necessary design files, including any existing Synopsys Design Constraints (.sdc) files that contain timing constraints for your design.

For more information, refer to Managing Files in a Project in Quartus II Help.

Performing an Initial Compilation

If you have never compiled your design, or you don’t have an .sdc file, and you want to use the TimeQuest analyzer to create one interactively, you must compile your design to create an initial design database before you specify timing constraints. You can either perform Analysis and Synthesis to create a post-map database, or perform a full compilation to create a post-fit database. Creating a post-map database is faster than a post-fit database, and is sufficient for creating initial timing constraints. The type of database you create determines the type of timing netlist generated by the TimeQuest analyzer; a post-map netlist if you perform Analysis and Synthesis or a post-fit netlist if you perform a full compilation.

If you are using incremental compilation, you must merge your design partitions after performing Analysis and Synthesis to create a post-map database.

For more information, refer to Setting up and Running Analysis and Synthesis and Setting up and Running a Compilation in Quartus II Help.

Specifying Timing Requirements

Before running timing analysis with the TimeQuest analyzer, you must specify initial timing constraints that describe the clock characteristics, timing exceptions, and signal transition arrival and required times. You can use the TimeQuest Timing Analyzer Wizard to enter initial constraints for your design, and then refine timing constraints with the TimeQuest analyzer GUI or with a Tcl script.

The Quartus II software assigns a default frequency of 1 GHz for clocks that have not been constrained, either in the TimeQuest GUI or an .sdc file, unless any constraint exists in the design. In that case, all unconstrained clocks remain unconstrained.

For more information, refer to Specifying Timing Constraints and Exceptions in Quartus II Help.

The .sdc must contain only SDC commands. Tcl commands to manipulate the timing netlist or control the compilation flow should be run as part of a separate Tcl script. After you create timing constraints, update the timing netlist to apply the new constraints. The TimeQuest analyzer applies all constraints to the netlist for verification and removes any invalid or false paths in the design from verification.

The constraints in the .sdc are read in sequence. You must first make a constraint before making any references to that constraint. For example, if a generated clock references a base clock, the base clock constraint must be made before the generated clock constraint.
The Quartus II Text Editor provides templates for SDC constraints. For more information, refer to “Using the Quartus II Templates” on page 7–6.

**Performing a Full Compilation**

After creating initial timing constraints, you must fully compile your design. When compilation is complete, you can open the TimeQuest analyzer to verify timing results and to generate summary, clock setup and clock hold, recovery, and removal reports for all defined clocks in the design.

**Verifying Timing**

The TimeQuest analyzer examines the timing paths in the design, calculates the propagation delay along each path, checks for timing constraint violations, and reports timing results as positive slack or negative slack. Negative slack indicates a timing violation. If you encounter violations along timing paths, use the timing reports to analyze your design and determine how best to optimize your design. If you modify, remove, or add constraints, you should perform a full compilation again. This iterative process helps resolve timing violations in your design.

For more information, refer to *Viewing Timing Analysis Results* in Quartus II Help.

Figure 7–2 shows the recommended flow for constraining and analyzing your design within the TimeQuest analyzer. Included are the corresponding Tcl commands for each step.

*Figure 7–2. The TimeQuest Timing Analyzer Flow*
SDC File Precedence

The Fitter and the TimeQuest analyzer process .sdc files in the order you specify in the Quartus II Settings File (.qsf). You can specify the files to process and the order they are processed from the Assignments menu. Click Settings, then TimeQuest Timing Analyzer, and specify a processing order in the SDC files to include in the project box.

If no .sdc files are listed in the .qsf, the Quartus II software looks for an .sdc named <current revision>.sdc in the project directory. An .sdc can also be added from a Quartus II IP File (.qip) included in the .qsf.

Figure 7–3 shows the order in which the Quartus II software searches for an .sdc.

Figure 7–3. .sdc File Order of Precedence

If you type the read_sdc command at the command line without any arguments, the TimeQuest analyzer reads constraints embedded in HDL files, then follows the .sdc file precedence order shown in Figure 7–3.

Using the Quartus II Templates

You can create an .sdc from constraint templates in the Quartus II software with the Quartus II Text Editor, or with your preferred text editor.

Creating a Constraint File with the Quartus II Text Editor

To insert constraints with the Quartus II Text Editor, follow these steps:

1. On the File menu, click New.
2. In the New dialog box, select the Synopsys Design Constraints File type from the Other Files group. Click OK.
3. Click the Insert Template button on the text editor menu, or, right-click in the blank .sdc file in the Quartus II Text Editor, then click Insert Template.
4. In the Insert Template dialog box, expand the TimeQuest section, then expand the SDC Commands section.
5. Expand a command category, for example, Clocks.
6. Select a command. The SDC constraint appears in the **Preview** pane.

7. Click **Insert** to paste the SDC constraint into the blank .sdc you created in step 2.

8. Repeat as needed with other constraints, or click **Close** to close the **Insert Template** dialog box.

You can now use any of the standard features of the Quartus II Text Editor to modify the .sdc or save the .sdc to edit in a text editor.

For more information on inserting a template with the Quartus II Text Editor, refer to *About the Quartus II Text Editor* in Quartus II Help.

### Constraining and Analyzing with Tcl Commands

You can use Tcl commands from the Quartus II software Tcl Application Programming Interface (API) to constrain, analyze, and collect information for your design. This section focuses on executing timing analysis tasks with Tcl commands; however, you can perform many of the same functions in the TimeQuest analyzer GUI. SDC commands are Tcl commands for constraining a design. SDC extension commands provide additional constraint methods and are specific to the TimeQuest analyzer. Additional TimeQuest analyzer commands are available for controlling timing analysis and reporting. These commands are contained in the following Tcl packages available in the Quartus II software:

- ::quartus::sta
- ::quartus::sdc
- ::quartus::sdc_ext

For more information about TimeQuest analyzer Tcl commands and a complete list of commands, refer to ::quartus::sta in Quartus II Help. For more information about standard SDC commands and a complete list of commands, refer to ::quartus::sdc in Quartus II Help. For more information about Altera extensions of SDC commands and a complete list of commands, refer to ::quartus::sdc_ext in Quartus II Help.

### Collection Commands

The TimeQuest analyzer Tcl commands often return port, pin, cell, or node names in a data set called a collection. In your Tcl scripts you can iterate over the values in collections to analyze or modify constraints in your design.

The TimeQuest analyzer supports collection commands that provide easy access to ports, pins, cells, or nodes in the design. Use collection commands with any valid constraints or Tcl commands specified in the TimeQuest analyzer.

Table 7–2 describes the collection commands supported by the TimeQuest analyzer.

<table>
<thead>
<tr>
<th>Command</th>
<th>Description of the collection returned</th>
</tr>
</thead>
<tbody>
<tr>
<td>all_clocks</td>
<td>All clocks in the design.</td>
</tr>
<tr>
<td>all_inputs</td>
<td>All input ports in the design.</td>
</tr>
<tr>
<td>all_outputs</td>
<td>All output ports in the design.</td>
</tr>
</tbody>
</table>
Wildcard Characters

To apply constraints to many nodes in a design, use the “*” and “?” wildcard characters. The “*” wildcard character matches any string; the “?” wildcard character matches any single character.

If you make an assignment to node `reg*`, the TimeQuest analyzer searches for and applies the assignment to all design nodes that match the prefix `reg` with any number of following characters, such as `reg`, `reg1`, `reg[2]`, `regbank`, and `reg12bank`.

If you make an assignment to a node specified as `reg?`, the TimeQuest analyzer searches and applies the assignment to all design nodes that match the prefix `reg` and any single character following; for example, `reg1`, `rega`, and `reg4`.

Adding and Removing Collection Items

Wildcards used with collection commands define collection items identified by the command. For example, if a design contains registers named `src0`, `src1`, `src2`, and `dst0`, the collection command `[get_registers src*]` identifies registers `src0`, `src1`, and `src2`, but not register `dst0`. To identify register `dst0`, you must use an additional command, `[get_registers dst*]`. To include `dst0`, you could also specify a collection command `[get_registers {src* dst*}]`.

To modify collections, use the `add_to_collection` and `remove_from_collection` commands. The `add_to_collection` command allows you to add additional items to an existing collection. Example 7–1 shows the `add_to_collection` command and arguments.

Example 7–1. `add_to_collection` Command

```
add_to_collection <first collection> <second collection>
```

The `add_to_collection` command creates a new collection that is the union of the two specified collections.
The `remove_from_collection` command allows you to remove items from an existing collection. **Example 7–2** shows the `remove_from_collection` command and arguments.

**Example 7–2. remove_from_collection Command**

```text
remove_from_collection <first collection> <second collection>
```

**Example 7–3** shows examples of how to add elements to collections.

**Example 7–3. Adding Items to a Collection**

```bash
#Setting up initial collection of registers
set regs1 [get_registers a*]

#Setting up initial collection of keepers
set kprs1 [get_keepers b*]

#Creating a new set of registers of $regs1 and $kprs1
set regs_union [add_to_collection $kprs1 $regs1]

#OR

# Creating a new set of registers of $regs1 and b*
# Note that the new collection appends only registers with name b*
# not all keepers
set regs_union [add_to_collection $regs1 b*]
```

In the Quartus II software, keepers are I/O ports or registers. An `.sdc` that includes `get_keepers` can only be processed as part of the TimeQuest analyzer flow and is not compatible with third-party timing analysis flows.

For more information about the `add_to_collection` and `remove_from_collection` commands—including full syntax information, options, and example usage—refer to `add_to_collection` and `remove_from_collection` in Quartus II Help.

**Using the `query_collection` Command**

You can display the contents of a collection with the `query_collection` command. **Example 7–4** shows how to report the contents of a collection:

**Example 7–4. query_collection Command**

```text
query_collection -report -all $regs_union
```

You can also examine collections and experiment with collections using wildcards in the TimeQuest analyzer by clicking **Name Finder** from the **View** menu.

**Using the `get_pins` Command**

The collection commands `get_pins` allow you to refine searches that include wildcard characters.

**Table 7–3** shows examples of search strings that use options to refine the search and wildcards. The examples in **Table 7–3** filter the following node and pin names to illustrate function:

- foo
The default method separates hierarchy levels of instances from nodes and pins with the pipe character (|). A match occurs when the levels of hierarchy match, and the string values including wildcards match the instance and/or pin names. For example, the command `get_pins <instance_name>|*|datac` returns all the datac pins for registers in a given instance. However, the command `get_pins *|datac` returns an empty collection because the levels of hierarchy do not match.

Use the -hierarchical matching scheme to return a collection of cells or pins in all hierarchies of your design.

For example, the command `get_pins -hierarchical *|*|datac` returns all the datac pins for all registers in your design. However, the command `get_pins -hierarchical *|*|datac` returns an empty collection because more than one pipe character (|) is not supported.

The -compatibility_mode option returns collections matching wildcard strings through any number of hierarchy levels. For example, an asterisk can match a pipe character when using -compatibility_mode.
Identifying the Quartus II Software Executable from the SDC File

To identify which Quartus II software executable is currently running you can use the $::TimeQuestInfo(nameofexecutable) variable from within an .sdc. Example 7–5 shows how to specify different SDC constraints for a specific Quartus II software executable.

Example 7–5. Identifying the Quartus II Executable

```tcl
#Identify which executable is running:
set current_exe $::TimeQuestInfo(nameofexecutable)
if { [string equal $current_exe "quartus_fit"] } {
    #Apply .sdc assignments for Fitter executable here
} else {
    #Apply .sdc assignments for non-Fitter executables here
}
if { ![string equal "quartus_sta" $::TimeQuestInfo(nameofexecutable)] } {
    #Apply .sdc assignments for non-TimeQuest executables here
} else {
    #Apply .sdc assignments for TimeQuest executable here
}
```

Examples of different executable names are quartus_map for Analysis & Synthesis, quartus_fit for Fitter, and quartus_sta for the TimeQuest analyzer.

Locating Timing Paths in Other Tools

You can locate paths and elements from the TimeQuest analyzer to other tools in the Quartus II software. Use the Locate Path command in the TimeQuest analyzer GUI or the locate command.

For more information about locating paths from the TimeQuest analyzer, refer to Viewing Timing Analysis Results and locate in Quartus II Help.

Example 7–6 shows how to locate ten paths from TimeQuest analyzer to the Chip Planner and locate all data ports in the Technology Map Viewer.

Example 7–6. Locating from the TimeQuest Analyzer

```tcl
# Locate in the Chip Planner all of the nodes in ten paths with the longest delay
locate [get_path -npaths 10] -chip
# locate all ports that begin with data to the Technology Map Viewer
locate [get_ports data*] -tmv
```
Design Constraints: An Example

Figure 7–4 shows an example circuit including two clocks, a phase-locked loop (PLL), and other common synchronous design elements.
Example 7–7 shows an .sdc file containing basic constraints for the circuit in Figure 7–4.

Example 7–7. Example Basic SDC Constraints

```plaintext
# Create clock constraints
create_clock -name clockone -period 10.000 [get_ports {clk1}]
create_clock -name clocktwo -period 10.000 [get_ports {clk2}]

# Create virtual clocks for input and output delay constraints
create_clock -name clockone_ext -period 10.000
create_clock -name clocktwo_ext -period 10.000
derive_pll_clocks

derive_clock_uncertainty

# set input and output delays
set_input_delay -clock { clockone_ext } -max 4 [get_ports {data1}]
set_input_delay -clock { clockone_ext } -min -1 [get_ports {data1}]
set_input_delay -clock { clockone_ext } -max 4 [get_ports {data2}]
set_input_delay -clock { clockone_ext } -min -1 [get_ports {data2}]
set_output_delay -clock { clocktwo_ext } -max 6 [get_ports {dataout}]
set_output_delay -clock { clocktwo_ext } -min -3 [get_ports {dataout}]
```

The .sdc in Example 7–7 contains the following basic constraints you should include for most designs:

- Definitions of clockone and clocktwo as base clocks, and assignment of those settings to nodes in the design.
- Definitions of clockone_ext and clocktwo_ext as virtual clocks, which represent clocks driving external devices interfacing with the FPGA.
- Automated derivation of generated clocks on PLL outputs.
- Derivation of clock uncertainty.
- Specification of two clock groups, the first containing clockone and its related clocks, the second containing clocktwo and its related clocks, and the third group containing the output of the PLL. This specification overrides the default analysis of all clocks in the design as related to each other. For more information about asynchronous clock groups, refer to “Asynchronous Clock Groups” on page 7–22.
- Specification of input and output delays for the design.

The following sections describe each of these constraint types in detail.
Creating Clocks and Clock Constraints

To ensure accurate static timing analysis results, you must specify all clocks and any associated clock characteristics in your design. The TimeQuest analyzer supports SDC commands that accommodate various clocking schemes and clock characteristics.

The TimeQuest analyzer supports the following types of clocks:

- Base clocks
- Virtual clocks
- Multifrequency clocks
- Generated clocks

Clocks are used to specify requirements for synchronous transfers and guide the Fitter optimization algorithms to achieve the best possible placement for your design.

Specify clock constraints first in the .sdc because other constraints may reference previously defined clocks. The TimeQuest analyzer reads SDC constraints and exceptions from top to bottom in the file.

Creating Base Clocks

Base clocks are the primary input clocks to the device. Unlike clocks from PLLs that are generated in the device, base clocks are generated by off-chip oscillators or forwarded from an external device. Define base clocks first because generated clocks and other constraints often reference base clocks.

To create clock settings for the signal from any register, port, or pin, use the create_clock command. You can create each clock with unique characteristics.

Example 7–8 shows how to create a 10 ns clock with a 50% duty cycle that is phase shifted by 90 degrees applied to port clk_sys.

Example 7–8. 100 MHz Shifted by 90 Degrees Clock Creation

```
create_clock -period 10 -waveform { 2.5 7.5 } [get_ports clk_sys]
```

Use the create_clock command to constrain all primary input clocks. The target for the create_clock command is usually a pin. To specify the pin as the target, use the get_ports command. Example 7–9 shows how to specify a 100 MHz requirement on a clk_sys input clock port.

Example 7–9. create_clock Command

```
create_clock -period 10 -name clk_sys [get_ports clk_sys]
```

You can apply multiple clocks on the same clock node with the -add option. Example 7–10 shows how to specify that two oscillators drive the same clock port on the device.

Example 7–10. Two Oscillators Driving the Same Clock Port

```
create_clock -period 10 -name clk_100 [get_ports clk_sys]
create_clock -period 5 -name clk_200 [get_ports clk_sys] -add
```
Creating Virtual Clocks

A virtual clock is a clock that does not have a real source in the design or that does not interact directly with the design. Virtual clocks are used in most I/O constraints, they represent the clock at the external device connected to the FPGA.

To create virtual clocks, use the `create_clock` command with no value specified for the `<targets>` option. Use virtual clocks for the reference clocks of `set_input_delay` and `set_output_delay` constraints.

Example 7–11 shows how to create a 10 ns virtual clock. The example is a virtual clock because no target is specified.

Example 7–11. Create Virtual Clock

```plaintext
create_clock -period 10 -name my_virt_clk
```

Figure 7–5 shows a design where a virtual clock is required for the TimeQuest analyzer to properly analyze the relationship between the external register and the registers in the design. Because the oscillator, `virt_clk`, does not interact with the Altera device, but acts as the clock source for the external register, you must declare the clock as a virtual clock. After you create the virtual clock, you can perform a register-to-register analysis between the register in the Altera device and the register in the external device.

Figure 7–5. Virtual Clock Board Topology
Example 7–12 shows how to create a 10 ns virtual clock named virt_clk with a 50% duty cycle where the first rising edge occurs at 0 ns. The virtual clock is then used as the clock source for an output delay constraint.

**Example 7–12. Virtual Clock Example**

```vhdl
# create base clock for the design
create_clock -period 5 [get_ports system_clk]

# create the virtual clock for the external register
create_clock -period 10 -name virt_clk

# set the output delay referencing the virtual clock
set_output_delay -clock virt_clk -max 1.5 [get_ports dataout]
set_output_delay -clock virt_clk -min 0.0 [get_ports dataout]
```

**I/O Interface Uncertainty**

Virtual clocks are recommended for I/O constraints because the derive_clock_uncertainty command can add different uncertainty values on clocks that interface with an external I/O port than uncertainty values between register paths fed by a clock inside the FPGA.

To specify I/O interface uncertainty, you must create a virtual clock and constrain the input and output ports with the set_input_delay and set_output_delay commands that reference the virtual clock. When the set_input_delay or set_output_delay commands reference a clock port or PLL output, the virtual clock allows the derive_clock_uncertainty command to apply separate clock uncertainties for internal clock transfers and I/O interface clock transfers.

Create the virtual clock with the same properties as the original clock that is driving the I/O port. Figure 7–6 shows a typical input I/O interface with clock specifications.

**Figure 7–6. I/O Interface Clock Specifications**
Example 7–13 shows the SDC commands to constrain the I/O interface shown in Figure 7–6.

**Example 7–13. SDC Commands to Constrain the I/O Interface**

```plaintext
# Create the base clock for the clock port
create_clock -period 10 -name clk_in [get_ports clk_in]

# Create a virtual clock with the same properties of the base clock
# driving the source register
create_clock -period 10 -name virt_clk_in

# Create the input delay referencing the virtual clock and not the base clock
# DO NOT use set_input_delay -clock clk_in <delay value>
# [get_ports data_in]
set_input_delay -clock virt_clk_in <delay value> [get_ports data_in]
```

For more information about clock uncertainty and clock transfers, refer to “Clock Uncertainty” on page 7–23.

**Creating Multifrequency Clocks**

You must create a multifrequency clock if your design has more than one clock source feeding a single clock node in your design. The additional clock may act as a low-power clock, with a lower frequency than the primary clock. If your design uses multifrequency clocks, use the `set_clock_groups` command to define clocks that are exclusive. For more information about creating exclusive clock groups, refer to “Creating Clock Groups” on page 7–22.

To create multifrequency clocks, use the `create_clock` command with the `-add` option to create multiple clocks on a clock node. Example 7–14 shows how to create a 10 ns clock applied to clock port `clk`, and then add an additional 15 ns clock to the same clock port. The TimeQuest analyzer uses both clocks when it performs timing analysis.

**Example 7–14. Multifrequency Clock Example**

```plaintext
create_clock -period 10 -name clock_primary -waveform { 0 5 } [get_ports clk]
create_clock -period 15 -name clock_secondary -waveform { 0 7.5 } [get_ports clk] -add
```

**Creating Generated Clocks**

Generated clocks are applied in the design when you modify the properties of a source synchronous clock signal, including phase, frequency, offset, and duty cycle. In the `.sdc`, generated clocks, which can be the outputs of PLLs or register clock dividers, are constrained after all base clocks. Generated clocks capture all clock delays and clock latency where the generated clock target is defined, ensuring that all base clock properties are accounted for in the generated clock.

Use the `create_generated_clock` command to constrain generated clocks in your design. The source of the `create_generated_clock` command should be a node in your design and not a previously constrained clock.
A common form of generated clock is a clock divider. Example 7–15 creates a base clock, clk_sys, then defines a generated clock clk_div_2, which is the clock frequency of clk_sys divided by two.

**Example 7–15. Clock Divider**

```plaintext
create_clock -period 10 -name clk_sys [get_ports clk_sys]
create_generated_clock -name clk_div_2 -divide_by 2 -source [get_ports clk_sys] [get_pins reg|regout]
```

When you use the `create_generated_clock` command, the -source option specifies a node with a clock used as a reference for your generated clock. Best practice is to specify the input clock pin of the target node for your new generated clock. You can also specify the target node of the reference clock. In Example 7–15, the -source option specifies the clock port clk feeding the clock pin of register reg.

If you have multiple base clocks feeding a node that is the source for a generated clock, you must define multiple generated clocks. Each generated clock is associated to one base clock using the -master_clock option in each generated clock statement.

The TimeQuest analyzer provides the `derive_pll_clocks` command to automatically generate clocks for all PLL clock outputs. The properties of the generated clocks on the PLL outputs match the properties defined for the PLL. For more information about deriving PLL clock outputs, refer to “Deriving PLL Clocks” on page 7–19.

For more information about the `create_generated_clock` and `derive_pll_clocks` commands—including for full syntax information, options, and example usage—refer to `create_generate_clock` and `derive_pll_clocks` in Quartus II Help.
The inverse of a clock divider is a clock multiplier. Figure 7–7 shows the effect of applying a multiplication factor to the generated clock.

**Figure 7–7. Multiplying a Generated Clock**

```plaintext
create_clock -period 10 -waveform { 0 5 } [get_ports clk]

# Creates a multiply-by-two clock
create_generated_clock -source [get_ports clk] -multiply_by 2 [get_registers \ clkmult|clkreg]
```

Figure 7–8 shows how to modify the generated clock by defining and shifting the edges.

**Figure 7–8. Edge Shifting a Generated Clock**

```plaintext
create_clock -period 10 -waveform { 0 5 } [get_ports clk]

# Creates a divide-by-two clock
create_generated_clock -source [get_ports clk] -edges { 1 3 5 } [get_registers \ clkdivA|clkreg]

# Creates a divide-by-two clock independent of the master clock’s duty cycle (now 50%)
create_generated_clock -source [get_ports clk] -edges { 1 1 5 } \ -edge_shift { 0 2.5 0 } [get_registers clkdivB|clkreg]
```

For information about creating generated clocks, refer to `create_generated_clocks` and `Specifying Timing Constraints and Exceptions` in Quartus II Help.

**Deriving PLL Clocks**

Use the `derive_pll_clocks` command to direct the TimeQuest analyzer to automatically search the timing netlist for all unconstrained PLL output clocks. The `derive_pll_clocks` command automatically creates generated clocks on the outputs of every PLL by calling the `create_generated_clock` command. The source for the `create_generated_clock` command is the input clock pin of the PLL.
Example 7–16 shows the command to create a base clock for the PLL input clock port and call derive_pll_clocks to create PLL output clocks.

**Example 7–16. Create Base Clock for PLL input Clock Ports**

```plaintext
create_clock -period 10.0 -name fpga_sys_clk [get_ports fpga_sys_clk]
derive_pll_clocks
```

If your design contains LVDS transmitters, LVDS receivers, or transceivers, Altera recommends using the `derive_pll_clocks` command. The command automatically constrains this logic in your design by adding the appropriate multicycle constraints to account for any deserialization factors.

For more information about the `derive_pll_clocks` command—including full syntax information, options, and example usage—refer to `derive_pll_clocks` and *Derive PLL Clocks* in Quartus II Help.

You can include the `derive_pll_clocks` command in your `.sdc`, which automatically detects any changes to the PLL settings. Each time the TimeQuest analyzer reads your `.sdc`, the appropriate `create_generated_clocks` command is generated for the PLL output clock pin.

Figure 7–9 shows a simple PLL design with a register-to-register path.

**Figure 7–9. Simple PLL Design**

Example 7–17 shows the messages generated by the TimeQuest analyzer when you use the `derive_pll_clocks` command to automatically constrain the PLL for the design shown in Figure 7–9.

**Example 7–17. derive_pll_clocks Command Messages**

```
Info: Deriving PLL Clocks:
Info: create_generated_clock -source pll_inst|altpll_component|pll|inclk[0] -divide_by 2 -name pll_inst|altpll_component|pll|clk[0]
```

The input clock pin of the PLL is the node `pll_inst|altpll_component|pll|inclk[0]` used for the `-source` option. The name of the output clock of the PLL is the PLL output clock node, `pll_inst|altpll_component|pll|clk[0]`. 
If the PLL is in clock switchover mode, multiple clocks are created for the output clock of the PLL; one for the primary input clock (for example, inclk[0]), and one for the secondary input clock (for example, inclk[1]). You should create exclusive clock groups for the primary and secondary output clocks.

For more information about creating exclusive clock groups, refer to “Creating Clock Groups” on page 7–22.

Automatically Detecting Clocks and Creating Default Clock Constraints

To automatically create clocks for all clock nodes in your design, use the `derive_clocks` command. The `derive_clocks` command is equivalent to using the `create_clock` command for each register or port feeding the clock pin of a register. The `derive_clocks` command creates clock constraints on ports or registers to ensure every register in your design has a clock constraints, and it applies one period to all base clocks in your design.

If there are no clock constraints in your design, the TimeQuest analyzer automatically creates default clock constraints for all detected unconstrained clock nodes to provide a complete clock analysis. The TimeQuest analyzer automatically creates clocks only when all synchronous elements have no associated clocks. For example, the TimeQuest analyzer does not create a default clock constraint if your design contains two clocks and you assigned constraints to one of the clocks. However, if you do not assign constraints to either clock, then the TimeQuest analyzer creates a default clock constraint.

Example 7–18 shows how the TimeQuest analyzer creates a base clock with a 1 GHz requirement for unconstrained clock nodes.

Example 7–18. Create Base Clock for Unconstrained Clock Nodes

```
derive_clocks -period 1
```

Do not use the `derive_clocks` command for final timing sign-off; instead, you should create clocks for all clock sources with the `create_clock` and `create_generated_clock` commands. If your design has more than a single clock, the `derive_clocks` command constrains all the clocks to the same specified frequency. To achieve a thorough and realistic analysis of your design’s timing requirements, you should make individual clock constraints for all clocks in your design.

You can also use the command `derive_pll_clocks -create_base_clocks` to create the input clocks for all PLL inputs automatically.

For more information about the `derive_clocks` command—including full syntax information, options, and example usage—refer to `derive_clocks` in Quartus II Help.
Creating Clock Groups

The TimeQuest analyzer assumes all clocks are related unless constrained otherwise. To specify clocks in your design that are exclusive or asynchronous, use the `set_clock_groups` command. The `set_clock_groups` command cuts timing between clocks in different groups, and performs the same analysis regardless of whether you specify -exclusive or -asynchronous. This is important if your design is migrating to HardCopy®, since other ASIC tools perform different analyses based on whether you specify-exclusive or -asynchronous.

For more information about the `set_clock_groups` command—including full syntax information, options, and example usage—refer to `set_clock_groups` in Quartus II Help.

Exclusive Clock Groups

Use the -exclusive option to declare that two clocks are mutually exclusive. You may want to declare clocks as mutually exclusive when multiple clocks are created on the same node or for multiplexed clocks. For example, a port can be clocked by either a 25-MHz or a 50-MHz clock. To constrain this port, you should create two clocks on the port, and then create clock groups to declare that they cannot coexist in the design at the same time. Declaring the clocks as mutually exclusive eliminates any clock transfers that may be derived between the 25-MHz clock and the 50-MHz clock. Example 7–19 shows how to create mutually exclusive clock groups.

Example 7–19. Create Mutually Exclusive Clock Groups

```plaintext
create_clock -period 40 -name clk_A [get_ports {port_A}]
create_clock -add -period 20 -name clk_B [get_ports {port_A}]
set_clock_groups -exclusive -group {clk_A} -group {clk_B}
```

A group is defined with the -group option. The TimeQuest analyzer excludes the timing paths between clocks for each of the separate groups.

If you apply multiple clocks to the same port, use the `set_clock_groups` command with the -exclusive option to place the clocks into separate groups and declare that the clocks are mutually exclusive. The clocks cannot physically exist in your design at the same time.

Asynchronous Clock Groups

Use the -asynchronous option to create asynchronous clock groups. Clocks contained within an asynchronous clock group are considered asynchronous to clocks in other clock groups; however, any clocks within a clock group are considered synchronous to each other.

For example, if your design has three clocks, clk_A, clk_B, and clk_C, and you establish that clk_A and clk_B are related to each other, but clock clk_C operates completely asynchronously, you can set up clock groups to define the clock behavior. If set_clock_groups is used with only one group, the clocks in that group are asynchronous with all other clocks in the design. For example, you can create a clock group containing only clk_C to ensure that clk_A and clk_B are synchronous with each other and asynchronous with clk_C. Because clk_C is the only clock in the constraint, it is asynchronous with every other clock in the design.
Example 7–20 shows how to create a clock group containing clocks \( \text{clk}_A \) and \( \text{clk}_B \) and a second unrelated clock group containing \( \text{clk}_C \).

**Example 7–20. Create Asynchronous Clock Groups**

```plaintext
set_clock_groups -asynchronous -group {\text{clk}_A \text{clk}_B} -group {\text{clk}_C}
```

For more information about the `set_clock_groups` command—including full syntax information, options, and example usage—refer to `set_clock_groups` in Quartus II Help.

**Accounting for Clock Effect Characteristics**

The clocks you create with the TimeQuest analyzer are ideal clocks that do not account for any board effects. You can account for clock effect characteristics with clock latency and clock uncertainty.

**Clock Latency**

There are two forms of clock latency, clock source latency and clock network latency. Source latency is the propagation delay from the origin of the clock to the clock definition point (for example, a clock port). Network latency is the propagation delay from a clock definition point to a register’s clock pin. The total latency at a register’s clock pin is the sum of the source and network latencies in the clock path.

To specify source latency to any clock ports in your design, use the `set_clock_latency` command.

The TimeQuest analyzer automatically computes network latencies; therefore, you only can characterize source latency with the `set_clock_latency` command. You must use the `-source` option.

For more information about the `set_clock_latency` command—including full syntax information, options, and example usage—refer to `set_clock_latency` in Quartus II Help.

**Clock Uncertainty**

The TimeQuest analyzer accounts for uncertainty clock effects for three types of clock-to-clock transfers; intraclock transfers, interclock transfers, and I/O interface clock transfers.

- Intraclock transfers occur when the register-to-register transfer takes place in the core of the device and the source and destination clocks come from the same PLL output pin or clock port.
- Interclock transfers occur when a register-to-register transfer takes place in the core of the device and the source and destination clocks come from a different PLL output pin or clock port.
- I/O interface clock transfers occur when data transfers from an I/O port to the core of the device or from the core of the device to the I/O port.
To manually specify clock uncertainty, or skew, for clock-to-clock transfers, use the set_clock_uncertainty command. You can specify the uncertainty separately for setup and hold, and you can specify separate rising and falling clock transitions. The TimeQuest analyzer subtracts setup uncertainty from the data required time for each applicable path and adds the hold uncertainty to the data required time for each applicable path.

To automatically apply interclock, intraclock, and I/O interface uncertainties, use the derive_clock_uncertainty command. The TimeQuest analyzer automatically applies clock uncertainties to clock-to-clock transfers in the design, and calculates both setup and hold uncertainties for each clock-to-clock transfer.

Any clock uncertainty constraints applied to source and destination clock pairs with the set_clock_uncertainty command have a higher precedence than the clock uncertainties derived with the derive_clock_uncertainty command for the same source and destination clock pairs. For example, if you use the set_clock_uncertainty command to set clock uncertainty between clka and clkb, the TimeQuest analyzer ignores the values for the clock transfer calculated with the derive_clock_uncertainty command. The TimeQuest analyzer reports the values calculated with the derive_clock_uncertainty command even if they are not used.

Use set_clock_uncertainty or derive_clock_uncertainty with the -overwrite option to overwrite previously applied clock uncertainty assignments. Use set_clock_uncertainty or derive_clock_uncertainty with the -add option to apply additional clock uncertainty to previously applied clock uncertainty. Use the remove_clock_uncertainty command to remove previous clock uncertainty assignments.

For more information about the set_clock_uncertainty, derive_clock_uncertainty, and remove_clock_uncertainty commands—including full syntax information, options, and example usage—refer to set_clock_uncertainty, remove_clock_uncertainty and derive_clock_uncertainty, in Quartus II Help.

Creating I/O Requirements

The TimeQuest analyzer reviews setup and hold relationships for designs in which an external source interacts with a register internal to the design. The TimeQuest analyzer supports input and output external delay modeling with the set_input_delay and set_output_delay commands. You can specify the clock and minimum and maximum arrival times relative to the clock.

You must specify timing requirements, including internal and external timing requirements, before you fully analyze a design. With external timing requirements specified, the TimeQuest analyzer verifies the I/O interface, or periphery of the device, against any system specification.

Input Constraints

Input constraints allow you to specify all the external delays feeding into the device. Specify input requirements for all input ports in your design.
You can use the `set_input_delay` command to specify external input delay requirements. Use the `-clock` option to reference a virtual clock. Using a virtual clock allows the TimeQuest analyzer to correctly derive clock uncertainties for interclock and intraclock transfers. The virtual clock defines the launching clock for the input port. The TimeQuest analyzer automatically determines the latching clock inside the device that captures the input data, because all clocks in the device are defined. Figure 7–10 shows an example of an input delay referencing a virtual clock.

**Figure 7–10. Input Delay**

---

**Equation 7–1. Input Delay Calculation**

\[
\begin{align*}
\text{input delay}_{\text{MAX}} &= (\text{cd}_{\text{extMAX}} - \text{cd}_{\text{altrMIN}}) + \text{tco}_{\text{extMAX}} + \text{dd}_{\text{MAX}} \\
\text{input delay}_{\text{MIN}} &= (\text{cd}_{\text{extMIN}} - \text{cd}_{\text{altrMAX}}) + \text{tco}_{\text{extMIN}} + \text{dd}_{\text{MIN}}
\end{align*}
\]

---

**Output Constraints**

Output constraints allow you to specify all external delays from the device for all output ports in your design.

You can use the `set_output_delay` command to specify external output delay requirements. Use the `-clock` option to reference a virtual clock. The virtual clock defines the latching clock for the output port. The TimeQuest analyzer automatically determines the launching clock inside the device that launches the output data, because all clocks in the device are defined. Figure 7–11 shows an example of an output delay referencing a virtual clock.

**Figure 7–11. Output Delay**
Equation 7–2 shows a typical output delay calculation.

**Equation 7–2. output Delay Calculation**

\[
\begin{align*}
\text{output delay}_{\text{MAX}} &= \text{dd}_{\text{MAX}} + \text{tsu}_\text{ext} + (\text{cd}_{\text{alt}}{\text{MAX}} - \text{cd}_{\text{ext}}{\text{MIN}}) \\
\text{output delay}_{\text{MIN}} &= (\text{dd}_{\text{MIN}} + \text{th}_\text{ext} + (\text{cd}_{\text{alt}}{\text{MIN}} - \text{cd}_{\text{ext}}{\text{MAX}}))
\end{align*}
\]

For more information about the `set_input_delay` and `set_output_delay` commands—including full syntax information, options, and example usage—refer to `set_input_delay` and `set_output_delay` in Quartus II Help.

### Creating Delay and Skew Constraints

The TimeQuest analyzer supports the Synopsys Design Constraint format for constraining timing for the ports in your design. These constraints allow the TimeQuest analyzer to perform a system static timing analysis that includes not only the device internal timing, but also any external device timing and board timing parameters.

#### Advanced I/O Timing and Board Trace Model Delay

The TimeQuest analyzer can use advanced I/O timing and board trace model assignments to model I/O buffer delays in your design.

If you change any advanced I/O timing settings or board trace model assignments, recompile your design before you analyze timing, or use the `-force_dat` option to force delay annotation when you create a timing netlist. Example 7–21 shows how to force delay annotation when creating a timing netlist.

**Example 7–21. Forcing Delay Annotation**

```plaintext
create_timing_netlist -force_dat
```

For more information about using advanced I/O timing, refer to Using Advanced I/O Timing in Quartus II Help.

For more information about advanced I/O timing, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

#### Maximum Skew

To specify the maximum path-based skew requirements for registers and ports in the design and report the results of maximum skew analysis, use the `set_max_skew` command in conjunction with the `report_max_skew` command.

By default, the `set_max_skew` command excludes any input or output delay constraints.

For more information about the `set_max_skew` and `report_max_skew` commands—including full syntax information, options, and example usage—refer to `set_max_skew` and `report_max_skew` in Quartus II Help.
Creating Timing Exceptions

Timing exceptions in the TimeQuest analyzer provide a way to modify the default timing analysis behavior to match the analysis required by your design. Specify timing exceptions after clocks and input and output delay constraints because timing exceptions can modify the default analysis.

Precedence

If a conflict of node names occurs between timing exceptions, the following order of precedence applies:

1. False path
2. Minimum delays and maximum delays
3. Multicycle path

The false path timing exception has the highest precedence. Within each category, assignments to individual nodes have precedence over assignments to clocks. Finally, the remaining precedence for additional conflicts is order-dependent, such that the assignments most recently created overwrite, or partially overwrite, earlier assignments.

False Paths

Specifying a false path in your design removes the path from timing analysis. Use the `set_false_path` command to specify false paths in your design. You can specify either a point-to-point or clock-to-clock path as a false path. For example, a path you should specify as false path is a static configuration register that is written once during power-up initialization, but does not change state again. Although signals from static configuration registers often cross clock domains, you may not want to make false path exceptions to a clock-to-clock path, because some data may transfer across clock domains. However, you can selectively make false path exceptions from the static configuration register to all endpoints.

Example 7–22 shows how to make false path exceptions from all registers beginning with A to all registers beginning with B.

Example 7–22. False Path

```plaintext
set_false_path -from [get_pins A*] -to [get_pins B*]
```

The TimeQuest analyzer assumes all clocks are related unless you specify otherwise. The “Creating Clock Groups” on page 7–22 describes how you can use clock groups. Clock groups are a more efficient way to make false path exceptions between clocks, compared to writing multiple `set_false_path` exceptions between every clock transfer you want to eliminate.

For more information about the `set_false_path` command—including full syntax information, options, and example usage—refer to `set_false_path` in Quartus II Help.
Minimum and Maximum Delays

To specify an absolute minimum or maximum delay for a path, use the `set_min_delay` command or the `set_max_delay` commands, respectively. Specifying minimum and maximum delay directly overwrites existing setup and hold relationships with the minimum and maximum values.

Use the `set_max_delay` and `set_min_delay` commands to create constraints for asynchronous signals that do not have a specific clock relationship in your design, but require a minimum and maximum path delay. You can create minimum and maximum delay exceptions for port-to-port paths through the device without a register stage in the path. If you use minimum and maximum delay exceptions to constrain the path delay, specify both the minimum and maximum delay of the path; do not constrain only the minimum or maximum value.

If the source or destination node is clocked, the TimeQuest analyzer takes into account the clock paths, allowing more or less delay on the data path. If the source or destination node has an input or output delay, that delay is also included in the minimum or maximum delay check.

If you specify a minimum or maximum delay between timing nodes, the delay applies only to the path between the two nodes. If you specify a minimum or maximum delay for a clock, the delay applies to all paths where the source node or destination node is clocked by the clock.

You can create a minimum or maximum delay exception for an output port that does not have an output delay constraint. You cannot report timing for the paths associated with the output port; however, the TimeQuest analyzer reports any slack for the path in the setup summary and hold summary reports. Because there is no clock associated with the output port, no clock is reported for timing paths associated with the output port.

To report timing with clock filters for output paths with minimum and maximum delay constraints, you can set the output delay for the output port with a value of zero. You can use an existing clock from the design or a virtual clock as the clock reference.

For more information about the `set_min_delay`, and `set_max_delay` commands—including full syntax information, options, and example usage—refer to `set_min_delay`, and `set_max_delay`, in Quartus II Help.

Delay Annotation

To modify the default delay values used during timing analysis, use the `set_annotated_delay` and `set_timing_derate` commands. You must update the timing netlist to see the results of these commands.

To specify different operating conditions in a single .sdc, rather than having multiple .sdc files that specify different operating conditions, use the `set_annotated_delay` command with the `-operating_conditions` option.

For more information about the `set_annotated_delay` and `set_timing_derate` commands—including full syntax information, options, and example usage—refer to `set_annotated_delay` and `set_timing_derate` in Quartus II Help.
**Multicycle Paths**

By default, the TimeQuest analyzer performs a single-cycle analysis, which is the most restrictive type of analysis. When analyzing a path, the setup launch and latch edge times are determined by finding the closest two active edges in the respective waveforms. For a hold analysis, the timing analyzer analyzes the path against two timing conditions for every possible setup relationship, not just the worst-case setup relationship. Therefore, the hold launch and latch times may be completely unrelated to the setup launch and latch edges. The TimeQuest analyzer does not report negative setup or hold relationships. When either a negative setup or a negative hold relationship is calculated, the TimeQuest analyzer moves both the launch and latch edges such that the setup and hold relationship becomes positive.

A multicycle constraint adjusts setup or hold relationships by the specified number of clock cycles based on the source (-start) or destination (-end) clock. An end setup multicycle constraint of 2 extends the worst-case setup latch edge by one destination clock period. If -start and -end values are not specified, the default constraint is -end.

Hold multicycle constraints are based on the default hold position (the default value is 0). An end hold multicycle constraint of 1 effectively subtracts one destination clock period from the default hold latch edge.

When the objects are timing nodes, the multicycle constraint only applies to the path between the two nodes. When an object is a clock, the multicycle constraint applies to all paths where the source node (-from) or destination node (-to) is clocked by the clock. When you adjust a setup relationship with a multicycle constraint, the hold relationship is adjusted automatically.

Table 7–4 shows the commands you can use to modify either the launch or latch edge times that the TimeQuest analyzer uses to determine a setup relationship or hold relationship.

**Table 7–4. Commands to Modify Edge Times**

<table>
<thead>
<tr>
<th>Command</th>
<th>Description of Modification</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>set_multicycle_path -setup -end &lt;value&gt;</code></td>
<td>Latch edge time of the setup relationship</td>
</tr>
<tr>
<td><code>set_multicycle_path -setup -start &lt;value&gt;</code></td>
<td>Launch edge time of the setup relationship</td>
</tr>
<tr>
<td><code>set_multicycle_path -hold -end &lt;value&gt;</code></td>
<td>Latch edge time of the hold relationship</td>
</tr>
<tr>
<td><code>set_multicycle_path -hold -start &lt;value&gt;</code></td>
<td>Launch edge time of the hold relationship</td>
</tr>
</tbody>
</table>

**Common Multicycle Variations**

Multicycle exceptions adjust the timing requirements for a register-to-register path, allowing the Fitter to optimally place and route a design in a device. Multicycle exceptions also can reduce compilation time and improve the quality of results, and can be used to change timing requirements. Two common multicycle variations are relaxing setup to allow a slower data transfer rate, and altering the setup to account for a phase shift.

**Relaxing Setup with set_multicycle_path**

A common type of multicycle exception occurs when the data transfer rate is slower than the clock cycle.
In this example, the source clock has a period of 10 ns, but a group of registers are enabled by a toggling clock, so they only toggle every other cycle. Since they are fed by a 10 ns clock, the TimeQuest analyzer reports a set up of 10 ns and a hold of 0 ns. However, since the data is transferring every other cycle, the relationships should be analyzed as if the clock were operating at 20 ns, which would result in a setup of 20 ns, while the hold remains 0 ns, in essence, extending the window of time when the data can be recognized.

Example 7–23 shows a pair of multicycle assignments that relax the setup relationship by specifying the \texttt{-setup} value of \(N\) and the \texttt{-hold} value as \(N-1\). You must specify the hold relationship with a \texttt{-hold} assignment to prevent a positive hold requirement.

**Example 7–23. Relaxing Setup while Maintaining Hold**

\begin{verbatim}
set_multicycle_path -setup -from src_reg* -to dst_reg* 2
set_multicycle_path -hold -from src_reg* -to dst_reg* 1
\end{verbatim}

Figure 7–12 shows how the exception relaxes the setup by two or three cycles.

**Figure 7–12. Relaxing Setup by Multiple Cycles**

This pattern can be extended to create larger setup relationships in order to ease timing closure requirements. A common use for this exception is when writing to asynchronous RAM across an I/O interface. The delay between address, data, and a write enable may be several cycles. A multicycle exception to I/O ports can allow extra time for the address and data to resolve before the enable occurs.

Example 7–24 shows how a relaxing the setup by three cycles can be achieved.

**Example 7–24. Three Cycle I/O Interface Exception**

\begin{verbatim}
set_multicycle_path -setup -to [get_ports {SRAM_ADD[*] SRAM_DATA[*]}] 3
set_multicycle_path -hold -to [get_ports {SRAM_ADD[*] SRAM_DATA[*]}] 2
\end{verbatim}
Accounting for a Phase Shift

In this example, the design contains a PLL that performs a phase-shift on a clock whose domain exchanges data with domains that do not experience the phase shift. For example, when the destination clock is phase-shifted forward and the source clock is not, the default setup relationship becomes that phase-shift.

Example 7–25 shows a circumstance where a PLL phase-shifts one output forward by a small amount, for example 0.2 ns.

Example 7–25. Cross Domain Phase-Shift

| create_generated_clock -source pll|inclk[0] -name pll|clk[0] pll|clk[0] |

The default setup relationship for this phase-shift is 0.2 ns, shown in Figure A, creating a scenario where the hold relationship is negative, which makes achieving timing closure nearly impossible.

Figure 7–13. Phase-Shifted Setup and Hold

Adding the constraint shown in Example Y allows the data to transfer to the following edge.

Example 7–26. Adjusting the Phase-Shift with a set_multicycle_path Constraint

| set_multicycle_path -setup -from [get_clocks clk_a] -to [get_clocks clk_b] 2 |

The hold relationship is derived from the setup relationship, making a multicycle hold constraint unnecessary. For a more complete example refer to “Same Frequency Clocks with Destination Clock Offset” on page 7–44.

For more information about the set_multicycle_path command—including full syntax information, options, and example usage—refer to set_multicycle_path in Quartus II Help.
Multicycle Clock Setup Check and Hold Check Analysis

You can modify the setup and hold relationship when you apply a multicycle exception to a register-to-register path. Figure 7–14 shows a register-to-register path with various timing parameters labeled.

**Figure 7–14. Register-to-Register Path**

Multicycle Clock Setup

The setup relationship is defined as the number of clock periods between the latch edge and the launch edge. By default, the TimeQuest analyzer performs a single-cycle path analysis, which results in the setup relationship being equal to one clock period (latch edge – launch edge). Applying a multicycle setup assignment, adjusts the setup relationship by the multicycle setup value. The adjustment value may be negative.

An end multicycle setup assignment modifies the latch edge of the destination clock by moving the latch edge the specified number of clock periods to the right of the determined default latch edge. Figure 7–15 shows various values of the end multicycle setup assignment and the resulting latch edge.

**Figure 7–15. End Multicycle Setup Values**
A start multicycle setup assignment modifies the launch edge of the source clock by moving the launch edge the specified number of clock periods to the left of the determined default launch edge. Figure 7–16 shows various values of the start multicycle setup assignment and the resulting launch edge.

**Figure 7–16. Start Multicycle Setup Values**

![Diagram showing various values of SMS for start multicycle setup assignment](image)

**Figure 7–17 shows the setup relationship reported by the TimeQuest analyzer for the negative setup relationship shown in Figure 7–16.**

**Figure 7–17. Start Multicycle Setup Values Reported by the TimeQuest Analyzer**

![Diagram showing setup relationship reported by TimeQuest analyzer](image)

**Multicycle Clock Hold**

The setup relationship is defined as the number of clock periods between the launch edge and the latch edge. By default, the TimeQuest analyzer performs a single-cycle path analysis, which results in the hold relationship being equal to one clock period (launch edge – latch edge). When analyzing a path, the TimeQuest analyzer performs two hold checks. The first hold check determines that the data launched by the current launch edge is not captured by the previous latch edge. The second hold check determines that the data launched by the next launch edge is not captured by the current latch edge. The TimeQuest analyzer reports only the most restrictive hold check. **Equation 7–3 shows the calculation that the TimeQuest analyzer performs to determine the hold check.**

**Equation 7–3. Hold Check**

\[
\text{hold check 1} = \text{current launch edge} - \text{previous latch edge} \\
\text{hold check 2} = \text{next launch edge} - \text{current latch edge}
\]
If a hold check overlaps a setup check, the hold check is ignored.

A start multicycle hold assignment modifies the launch edge of the destination clock by moving the latch edge the specified number of clock periods to the right of the determined default launch edge. Figure 7–18 shows various values of the start multicycle hold assignment and the resulting launch edge.

**Figure 7–18. Start Multicycle Hold Values**

![Start Multicycle Hold Values Diagram](image)

An end multicycle hold assignment modifies the latch edge of the destination clock by moving the latch edge the specified number of clock periods to the left of the determined default latch edge. Figure 7–19 shows various values of the end multicycle hold assignment and the resulting latch edge.

**Figure 7–19. End Multicycle Hold Values**

![End Multicycle Hold Values Diagram](image)
Figure 7–20 shows the hold relationship reported by the TimeQuest analyzer for the negative hold relationship shown in Figure 7–19.

**Figure 7–20. End Multicycle Hold Values Reported by the TimeQuest Analyzer**

<table>
<thead>
<tr>
<th>Source Clock</th>
<th>Destination Clock</th>
</tr>
</thead>
<tbody>
<tr>
<td>-10</td>
<td></td>
</tr>
<tr>
<td>0</td>
<td></td>
</tr>
<tr>
<td>10</td>
<td></td>
</tr>
<tr>
<td>20</td>
<td></td>
</tr>
</tbody>
</table>

**Examples of Basic Multicycle Exceptions**

This section describes the following examples of combinations of multicycle exceptions:

- “Default Settings” on page 7–35
- “End Multicycle Setup = 2 and End Multicycle Hold = 0” on page 7–38
- “End Multicycle Setup = 2 and End Multicycle Hold = 1” on page 7–41

Each example explains how the multicycle exceptions affect the default setup and hold analysis in the TimeQuest analyzer. The multicycle exceptions are applied to a simple register-to-register circuit. Both the source and destination clocks are set to 10 ns.

**Default Settings**

By default, the TimeQuest analyzer performs a single-cycle analysis to determine the setup and hold checks. Also, by default, the TimeQuest analyzer sets the end multicycle setup assignment value to one and the end multicycle hold assignment value to zero.
Figure 7–21 shows the source and the destination timing waveform for the source register and destination register, respectively where HC1 and HC2 are hold checks one and two and SC is the setup check.

**Figure 7–21. Default Timing Diagram**

Equation 7–4 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–4. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge} \\
= 10 \text{ ns} - 0 \text{ ns} \\
= 10 \text{ ns}
\]

The most restrictive setup relationship with the default single-cycle analysis, that is, a setup relationship with an end multicycle setup assignment of one, is 10 ns.
Figure 7–22 shows the setup report for the default setup in the TimeQuest analyzer with the launch and latch edges highlighted.

**Equation 7–5** shows the calculation that the TimeQuest analyzer performs to determine the hold check. Both hold checks are equivalent.

**Equation 7–5. Hold Check**

hold check 1  =  current launch edge – previous latch edge
= 0 ns – 0 ns
= 0 ns

hold check 2  =  next launch edge – current latch edge
= 10 ns – 10 ns
= 0 ns

The most restrictive hold relationship with the default single-cycle analysis, that a hold relationship with an end multicycle hold assignment of zero, is 0 ns.
Figure 7–23 shows the hold report for the default setup in the TimeQuest analyzer with the launch and latch edges highlighted.

### End Multicycle Setup = 2 and End Multicycle Hold = 0

In this example, the end multicycle setup assignment value is two, and the end multicycle hold assignment value is zero. Example 7–27 shows the multicycle exceptions applied to the register-to-register design for this example.

#### Example 7–27. Multicycle Exceptions

```plaintext
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -end 2
```

An end multicycle hold value is not required because the default end multicycle hold value is zero.

In this example, the setup relationship is relaxed by a full clock period by moving the latch edge to the next latch edge. The hold analysis is unchanged from the default settings.
Figure 7–24 shows the setup timing diagram. The latch edge is a clock cycle later than in the default single-cycle analysis.

**Figure 7–24. Setup Timing Diagram**

![Setup Timing Diagram](image)

Equation 7–6 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–6. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge} = 20 \text{ ns} - 0 \text{ ns} = 20 \text{ ns}
\]

The most restrictive setup relationship with an end multicycle setup assignment of two is 20 ns.

Figure 7–25 shows the setup report in the TimeQuest analyzer with the launch and latch edges highlighted.

**Figure 7–25. Setup Report**

![Setup Report](image)
Because the multicycle hold latch and launch edges are the same as the results of hold analysis with the default settings, the multicycle hold analysis in this example is equivalent to the single-cycle hold analysis. Figure 7–26 shows the timing diagram for the hold checks for this example. The hold checks are relative to the setup check. Usually, the TimeQuest analyzer performs hold checks on every possible setup check, not only on the most restrictive setup check edges.

**Figure 7–26. Hold Timing Diagram**

![Hold Timing Diagram](image)

Equation 7–7 shows the calculation that the TimeQuest analyzer performs to determine the hold check. Both hold checks are equivalent.

**Equation 7–7. Hold Check**

- **hold check 1** = current launch edge – previous latch edge
  = 0 ns – 10 ns
  = −10 ns

- **hold check 2** = next launch edge – current latch edge
  = 10 ns – 20 ns
  = −10 ns

The most restrictive hold relationship with an end multicycle setup assignment value of two and an end multicycle hold assignment value of zero is 10 ns.
Figure 7–27 shows the hold report for this example in the TimeQuest analyzer with the launch and latch edges highlighted.

**End Multicycle Setup = 2 and End Multicycle Hold = 1**

In this example, the end multicycle setup assignment value is two, and the end multicycle hold assignment value is one. Example 7–28 shows the multicycle exceptions applied to the register-to-register design for this example.

**Example 7–28. Multicycle Exceptions**

```
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -end 2
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -hold -end 1
```

In this example, the setup relationship is relaxed by two clock periods by moving the latch edge to the left two clock periods. The hold relationship is relaxed by a full period by moving the latch edge to the previous latch edge.
Figure 7–28 shows the setup timing diagram.

**Figure 7–28. Setup Timing Diagram**

![Setup Timing Diagram](image)

Equation 7–8 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–8. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge} \\
= 20 \text{ ns} - 0 \text{ ns} \\
= 20 \text{ ns}
\]

The most restrictive hold relationship with an end multicycle setup assignment value of two is 20 ns.

Figure 7–29 shows the setup report for this example in the TimeQuest analyzer with the launch and latch edges highlighted.

**Figure 7–29. Setup Report**

![Setup Report](image)
Figure 7–30 shows the timing diagram for the hold checks for this example. The hold checks are relative to the setup check.

**Figure 7–30. Hold Timing Diagram**

![Hold Timing Diagram](image)

Equation 7–9 shows the calculation that the TimeQuest analyzer performs to determine the hold check. Both hold checks are equivalent.

**Equation 7–9. Hold Check**

\[
\text{hold check 1} = \text{current launch edge} - \text{previous latch edge} = 0 \text{ ns} - 0 \text{ ns} = 0 \text{ ns}
\]

\[
\text{hold check 2} = \text{next launch edge} - \text{current latch edge} = 10 \text{ ns} - 10 \text{ ns} = 0 \text{ ns}
\]

The most restrictive hold relationship with an end multicycle setup assignment value of two and an end multicycle hold assignment value of one is 0 ns.
Creating Timing Exceptions

Figure 7–31 shows the hold report for this example in the TimeQuest analyzer with the launch and latch edges highlighted.

Figure 7–31. Hold Report

Application of Multicycle Exceptions

This section shows the following examples of applications of multicycle exceptions:

- “Same Frequency Clocks with Destination Clock Offset” on page 7–44
- “The Destination Clock Frequency is a Multiple of the Source Clock Frequency” on page 7–47
- “The Destination Clock Frequency is a Multiple of the Source Clock Frequency with an Offset” on page 7–50
- “The Source Clock Frequency is a Multiple of the Destination Clock Frequency” on page 7–52
- “The Source Clock Frequency is a Multiple of the Destination Clock Frequency with an Offset” on page 7–55

Each example explains how the multicycle exceptions affect the default setup and hold analysis in the TimeQuest analyzer. All of the examples are between related clock domains. If your design contains related clocks, such as PLL clocks, and paths between related clock domains, you can apply multicycle constraints.

Same Frequency Clocks with Destination Clock Offset

In this example, the source and destination clocks have the same frequency, but the destination clock is offset with a positive phase shift. Both the source and destination clocks have a period of 10 ns. The destination clock has a positive phase shift of 2 ns with respect to the source clock. Figure 7–32 shows an example of a design with same frequency clocks and a destination clock offset.
Creating Timing Exceptions

Figure 7–32. Same Frequency Clocks with Destination Clock Offset

Figure 7–33 shows the timing diagram for default setup check analysis performed by the TimeQuest analyzer.

Figure 7–33. Setup Timing Diagram

Equation 7–10 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–10. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge}
\]

\[
= 2 \text{ ns} - 0 \text{ ns}
\]

\[
= 2 \text{ ns}
\]

The setup relationship shown in Figure 7–33 is too pessimistic and is not the setup relationship required for typical designs. To correct the default analysis, you must use an end multicycle setup exception of two. Example 7–29 shows the multicycle exception used to correct the default analysis in this example.

**Example 7–29. Multicycle Exceptions**

```
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst]
  -setup -end 2
```
Figure 7–34 shows the timing diagram for the preferred setup relationship for this example.

**Figure 7–34. Preferred Setup Relationship**

![Preferred Setup Relationship Diagram](image)

Figure 7–35 shows the timing diagram for default hold check analysis performed by the TimeQuest analyzer with an end multicycle setup value of two.

**Figure 7–35. Default Hold Check**

![Default Hold Check Diagram](image)

Equation 7–11 shows the calculation that the TimeQuest analyzer performs to determine the hold check.

**Equation 7–11. Hold Check**

\[
\begin{align*}
\text{hold check 1} & = \text{current launch edge} - \text{previous latch edge} \\
& = 0 \text{ ns} - 2 \text{ ns} \\
& = -2 \text{ ns}
\end{align*}
\]

\[
\begin{align*}
\text{hold check 2} & = \text{next launch edge} - \text{current latch edge} \\
& = 10 \text{ ns} - 12 \text{ ns} \\
& = -2 \text{ ns}
\end{align*}
\]

In this example, the default hold analysis returns the preferred hold requirements and no multicycle hold exceptions are required.
Figure 7–36 shows the associated setup and hold analysis if the phase shift is −2 ns. In this example, the default hold analysis is correct for the negative phase shift of 2 ns, and no multicycle exceptions are required.

**Figure 7–36. Negative Phase Shift**

The Destination Clock Frequency is a Multiple of the Source Clock Frequency

In this example, the destination clock frequency value of 5 ns is an integer multiple of the source clock frequency of 10 ns. The destination clock frequency can be an integer multiple of the source clock frequency when a PLL is used to generate both clocks with a phase shift applied to the destination clock. Figure 7–37 shows an example of a design where the destination clock frequency is a multiple of the source clock frequency.

**Figure 7–37. Destination Clock is Multiple of Source Clock**
Figure 7–38 shows the timing diagram for default setup check analysis performed by the TimeQuest analyzer.

**Figure 7–38. Setup Timing Diagram**

![Setup Timing Diagram](image)

**Equation 7–12** shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–12. Setup Check**

\[
\text{setup check} = \text{current latch edge - closest previous launch edge} \\
= 5 \text{ ns} - 0 \text{ ns} \\
= 5 \text{ ns}
\]

The setup relationship shown in Figure 7–38 demonstrates that the data does not need to be captured at edge one, but can be captured at edge two; therefore, you can relax the setup requirement. To correct the default analysis, you must shift the latch edge by one clock period with an end multicycle setup exception of two. Example 7–30 shows the multicycle exception used to correct the default analysis in this example.

**Example 7–30. Multicycle Exceptions**

```
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -end 2
```
Figure 7–39 shows the timing diagram for the preferred setup relationship for this example.

**Figure 7–39. Preferred Setup Analysis**

![Preferred Setup Analysis Diagram](image)

Figure 7–40 shows the timing diagram for default hold check analysis performed by the TimeQuest analyzer with an end multicycle setup value of two.

**Figure 7–40. Default Hold Check**

![Default Hold Check Diagram](image)

Equation 7–13 shows the calculation that the TimeQuest analyzer performs to determine the hold check.

**Equation 7–13. Hold Check**

\[
\text{hold check 1} = \text{current launch edge} - \text{previous latch edge} = 0 \text{ ns} - 5 \text{ ns} = -5 \text{ ns}
\]

\[
\text{hold check 2} = \text{next launch edge} - \text{current latch edge} = 10 \text{ ns} - 10 \text{ ns} = 0 \text{ ns}
\]

In this example, hold check one is too restrictive. The data is launched by the edge at 0 ns and should check against the data captured by the previous latch edge at 0 ns, which does not occur in hold check one. To correct the default analysis, you must use an end multicycle hold exception of one.
The Destination Clock Frequency is a Multiple of the Source Clock Frequency with an Offset

This example is a combination of the previous two examples. The destination clock frequency is an integer multiple of the source clock frequency and the destination clock has a positive phase shift. The destination clock frequency is 5 ns and the source clock frequency is 10 ns. The destination clock also has a positive offset of 2 ns with respect to the source clock. The destination clock frequency can be an integer multiple of the source clock frequency with an offset when a PLL is used to generate both clocks with a phase shift applied to the destination clock. Figure 7–41 shows an example of a design in which the destination clock frequency is a multiple of the source clock frequency with an offset.

Figure 7–41. Destination Clock is Multiple of Source Clock with Offset

Figure 7–42 shows the timing diagram for default setup check analysis performed by the TimeQuest analyzer.

Figure 7–42. Setup Timing Diagram

Equation 7–14 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–14. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge}
\]

\[
= 2 \text{ ns} - 0 \text{ ns}
\]

\[
= 2 \text{ ns}
\]

The setup relationship shown in Figure 7–42 demonstrates that the data does not need to be captured at edge one, but can be captured at edge two; therefore, you can relax the setup requirement. To correct the default analysis, you must shift the latch edge by one clock period with an end multicycle setup exception of three.
Example 7–31 shows the multicycle exception used to correct the default analysis in this example.

**Example 7–31. Multicycle Exceptions**

```
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -end 3
```

Figure 7–43 shows the timing diagram for the preferred setup relationship for this example.

**Figure 7–43. Preferred Setup Analysis**

Figure 7–44 shows the timing diagram for default hold check analysis performed by the TimeQuest analyzer with an end multicycle setup value of three.

**Figure 7–44. Default Hold Check**
Equation 7–15 shows the calculation that the TimeQuest analyzer performs to determine the hold check.

**Equation 7–15. Hold Check**

<table>
<thead>
<tr>
<th>Description</th>
<th>Equation</th>
</tr>
</thead>
<tbody>
<tr>
<td>hold check 1 = current launch edge – previous latch edge</td>
<td>hold check 1 = 0 ns – 5 ns = –5 ns</td>
</tr>
<tr>
<td>hold check 2 = next launch edge – current latch edge</td>
<td>hold check 2 = 10 ns – 10 ns = 0 ns</td>
</tr>
</tbody>
</table>

In this example, hold check one is too restrictive. The data is launched by the edge at 0 ns and should check against the data captured by the previous latch edge at 2 ns, which does not occur in hold check one. To correct the default analysis, you must use an end multicycle hold exception of one.

### The Source Clock Frequency is a Multiple of the Destination Clock Frequency

In this example, the source clock frequency value of 5 ns is an integer multiple of the destination clock frequency of 10 ns. The source clock frequency can be an integer multiple of the destination clock frequency when a PLL is used to generate both clocks and different multiplication and division factors are used. Figure 7–45 shows an example of a design where the source clock frequency is a multiple of the destination clock frequency.

**Figure 7–45. Source Clock Frequency is Multiple of Destination Clock Frequency**

![Diagram showing a design with two clocks (clk0 and clk1) and two registers (REG1 and REG2) with combinational logic in between.]
Figure 7–46 shows the timing diagram for default setup check analysis performed by the TimeQuest analyzer.

**Figure 7–46. Default Setup Check Analysis**

![Timing Diagram](image)

Equation 7–16 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–16. Setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge} \\
= 10 \text{ ns} - 5 \text{ ns} \\
= 5 \text{ ns}
\]

The setup relationship shown in Figure 7–46 demonstrates that the data launched at edge one does not need to be captured, and the data launched at edge two must be captured; therefore, you can relax the setup requirement. To correct the default analysis, you must shift the launch edge by one clock period with a start multicycle setup exception of two.

**Example 7–32 shows the multicycle exception used to correct the default analysis in this example.**

**Example 7–32. Multicycle Exceptions**

```bash
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -start 2
```
Figure 7–47 shows the timing diagram for the preferred setup relationship for this example.

**Figure 7–47. Preferred Setup Check Analysis**

![Preferred Setup Check Analysis Diagram](image)

Figure 7–48 shows the timing diagram for default hold check analysis performed by the TimeQuest analyzer with a start multicycle setup value of two.

**Figure 7–48. Default Hold Check**

![Default Hold Check Diagram](image)

Equation 7–17 shows the calculation that the TimeQuest analyzer performs to determine the hold check.

**Equation 7–17. Hold Check**

\[
\begin{align*}
\text{hold check 1} & = \text{current launch edge} - \text{previous latch edge} \\
& = 0 \text{ ns} - 0 \text{ ns} \\
& = 0 \text{ ns}
\end{align*}
\]

\[
\begin{align*}
\text{hold check 2} & = \text{next launch edge} - \text{current latch edge} \\
& = 5 \text{ ns} - 10 \text{ ns} \\
& = -5 \text{ ns}
\end{align*}
\]

In this example, hold check two is too restrictive. The data is launched next by the edge at 10 ns and should check against the data captured by the current latch edge at 10 ns, which does not occur in hold check two. To correct the default analysis, you must use a start multicycle hold exception of one.
The Source Clock Frequency is a Multiple of the Destination Clock Frequency with an Offset

In this example, the source clock frequency is an integer multiple of the destination clock frequency and the destination clock has a positive phase offset. The source clock frequency is 5 ns and destination clock frequency is 10 ns. The destination clock also has a positive offset of 2 ns with respect to the source clock. The source clock frequency can be an integer multiple of the destination clock frequency with an offset when a PLL is used to generate both clocks, different multiplication and division factors are used, and a phase shift applied to the destination clock. Figure 7–49 shows an example of a design where the source clock frequency is a multiple of the destination clock frequency with an offset.

Figure 7–49. Source Clock Frequency is Multiple of Destination Clock Frequency with Offset

![Source Clock Frequency is Multiple of Destination Clock Frequency with Offset](image)

Figure 7–50 shows the timing diagram for default setup check analysis performed by the TimeQuest analyzer.

Figure 7–50. Setup Timing Diagram

![Setup Timing Diagram](image)

Equation 7–18 shows the calculation that the TimeQuest analyzer performs to determine the setup check.

**Equation 7–18. setup Check**

\[
\text{setup check} = \text{current latch edge} - \text{closest previous launch edge}
\]

\[
= 12 \text{ ns} - 10 \text{ ns}
\]

\[
= 2 \text{ ns}
\]

The setup relationship shown in Figure 7–50 demonstrates that the data is not launched at edge one, and the data that is launched at edge three must be captured; therefore, you can relax the setup requirement. To correct the default analysis, you must shift the launch edge by two clock periods with a start multicycle setup exception of three.
Example 7–33 shows the multicycle exception used to correct the default analysis in this example.

**Example 7–33. Multicycle Exceptions**

```
set_multicycle_path -from [get_clocks clk_src] -to [get_clocks clk_dst] -setup -start 3
```

Figure 7–51 shows the timing diagram for the preferred setup relationship for this example.

**Figure 7–51. Preferred Setup Check Analysis**

Figure 7–52 shows the timing diagram for default hold check analysis performed by the TimeQuest analyzer with a start multicycle setup value of three.

**Figure 7–52. Default Hold Check Analysis**

Equation 7–19 shows the calculation that the TimeQuest analyzer performs to determine the hold check.

**Equation 7–19. Hold Check**

\[
\text{hold check 1} = \text{current launch edge} - \text{previous latch edge} = 0\ \text{ns} - 2\ \text{ns} = -2\ \text{ns}
\]

\[
\text{hold check 2} = \text{next launch edge} - \text{current latch edge}
\]
In this example, hold check two is too restrictive. The data is launched next by the edge at 10 ns and should check against the data captured by the current latch edge at 12 ns, which does not occur in hold check two. To correct the default analysis, you must use a start multicycle hold exception of one.

### Timing Reports

The TimeQuest analyzer provides real-time static timing analysis result reports. The TimeQuest analyzer does not automatically generate reports; you must create each report individually in the TimeQuest analyzer GUI or with command-line commands. You can customize in which report to display specific timing information, excluding fields that are not required.

Table 7–5 shows some of the different command-line commands you can use to generate reports in the TimeQuest analyzer and the equivalent reports shown in the TimeQuest analyzer GUI.

#### Table 7–5. TimeQuest Analyzer Reports

<table>
<thead>
<tr>
<th>Command-Line Command</th>
<th>Report</th>
</tr>
</thead>
<tbody>
<tr>
<td>report_timing</td>
<td>Timing report</td>
</tr>
<tr>
<td>report_exceptions</td>
<td>Exceptions report</td>
</tr>
<tr>
<td>report_clock_transfers</td>
<td>Clock Transfers report</td>
</tr>
<tr>
<td>report_min_pulse_width</td>
<td>Minimum Pulse Width report</td>
</tr>
<tr>
<td>report_ucp</td>
<td>Unconstrained Paths report</td>
</tr>
</tbody>
</table>

For more information—including a complete list of commands to generate timing reports and full syntax information, options, and example usage—refer to `:quartus:sta` in Quartus II Help.

During compilation, the Quartus II software generates timing reports on different timing areas in the design. You can configure various options for the TimeQuest analyzer reports generated during compilation.

For more information about the options you can set to customize the reports, refer to `TimeQuest Timing Analyzer Page` in Quartus II Help.

You can also use the `TIMEQUEST_REPORT_WORST_CASE_TIMING_PATHS` assignment to generate a report of the worst-case timing paths for each clock domain. This report contains worst-case timing data for setup, hold, recovery, removal, and minimum pulse width checks.

Use the `TIMEQUEST_REPORT_NUM_WORST_CASE_TIMING_PATHS` assignment to specify the number of paths to report for each clock domain.

---

**Equation 7–19. Hold Check**

\[
5 \text{ ns} - 12 \text{ ns} = -7 \text{ ns}
\]

In this example, hold check two is too restrictive. The data is launched next by the edge at 10 ns and should check against the data captured by the current latch edge at 12 ns, which does not occur in hold check two. To correct the default analysis, you must use a start multicycle hold exception of one.
Example 7–34 shows an example of how to use the TIMEQUEST_REPORT_WORST_CASE_TIMING_PATHS and TIMEQUEST_REPORT_NUM_WORST_CASE_TIMING_PATHS assignments in the .qsf to generate reports.

Example 7–34. Generating Worst-Case Timing Reports

```%^#Enable Worst-Case Timing Report
set_global_assignment -name TIMEQUEST_REPORT_WORST_CASE_TIMING_PATHS ON
#Report 10 paths per clock domain
set_global_assignment -name TIMEQUEST_REPORT_NUM_WORST_CASE_TIMING_PATHS 10```

For more information about timing closure recommendations, refer to the Area and Timing Optimization chapter in volume 2 of the Quartus II Handbook.

**Document Revision History**

Table 7–6 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
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</table>
| June 2012  | 12.0.0  | ■ Reorganized chapter.  
■ Added “Using the Quartus II Templates” section on creating an SDC constraints file with the Insert Template dialog box.  
■ Added “Identifying the Quartus II Software Executable from the SDC File” section.  
■ Revised multicycle exceptions section. |
| November 2011 | 11.1.0  | ■ Consolidated content from the Best Practices for the Quartus II TimeQuest Timing Analyzer chapter.  
■ Changed to new document template. |
| May 2011   | 11.0.0  | ■ Updated to improve flow. Minor editorial updates. |
| December 2010 | 10.1.0  | ■ Changed to new document template.  
■ Revised and reorganized entire chapter.  
■ Linked to Quartus II Help. |
| July 2010  | 10.0.0  | Updated to link to content on SDC commands and the TimeQuest analyzer GUI in Quartus II Help. |
Table 7–6. Document Revision History  (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
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<th>Changes</th>
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<td>November 2009</td>
<td>9.1.0</td>
<td>Updated for the Quartus II software version 9.1, including:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about commands for adding and removing items from collections</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about the set_timing_derate and report_skew commands</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added information about worst-case timing reporting</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Updated for the Quartus II software version 8.1, including:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added the following sections:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “set_net_delay” on page 7–42</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “Annotated Delay” on page 7–49</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ “report_net_delay” on page 7–66</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated the descriptions of the -append and -file &lt;name&gt; options in tables throughout the chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated entire chapter using 8½” × 11” chapter template</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates</td>
</tr>
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</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
As FPGA designs grow larger and processes continue to shrink, power is an ever-increasing concern. When designing a PCB, the power consumed by a device must be accurately estimated to develop an appropriate power budget, and to design the power supplies, voltage regulators, heat sink, and cooling system.

The Quartus® II software allows you to estimate the power consumed by your current design during timing simulation. The power consumption of your design can be calculated using the Microsoft Excel-based power calculator, or the Simulation-Based Power Estimation features in the Quartus II software. This section explains how to use both.

This section includes the following chapter:

- **Chapter 8, PowerPlay Power Analysis**

  This chapter describes the Altera® Quartus II PowerPlay power analysis tool and how to use the tools to accurately estimate device power consumption.
This chapter describes how to use the Altera® Quartus® II PowerPlay Power Analysis tools to accurately estimate device power consumption.

As designs grow larger and process technology continues to shrink, power becomes an increasingly important design consideration. When designing a PCB, you must estimate the power consumption of a device accurately to develop an appropriate power budget, and to design the power supplies, voltage regulators, heat sink, and cooling system.

Figure 8–1 shows the ability of the PowerPlay Power Analysis tools to estimate power consumption from early design concept through design implementation.

Figure 8–1. PowerPlay Power Analysis

For more information about the PowerPlay suite of power analysis and optimizations tools, refer to About Power Estimation and Analysis in Quartus II Help. For more information about acquiring the PowerPlay EPE spreadsheet, refer to PowerPlay Early Power Estimators (EPE) and Power Analyzer on the Altera website.
This chapter discusses the following topics:

- “Types of Power Analyses” on page 8–2
- “Factors Affecting Power Consumption” on page 8–3
- “Creating PowerPlay EPE Spreadsheets” on page 8–6
- “PowerPlay Power Analyzer Flow” on page 8–8
- “Using Simulation Files in Modular Design Flows” on page 8–10
- “Using the PowerPlay Power Analyzer” on page 8–17

Types of Power Analyses

Understanding the uses of power analysis and the factors affecting power consumption helps you to use the PowerPlay Power Analyzer effectively. Power analysis meets two significant planning requirements:

- **Thermal planning**—Thermal power is the power that dissipates as heat from the FPGA. You must use a heatsink or fan to act as a cooling solution for your device. The cooling solution must be sufficient to dissipate the heat that the device generates. The computed junction temperature must fall within normal device specifications.

- **Power supply planning**—Power supply is the power needed to run your device. Power supplies must provide adequate current to support device operation.

The two types of analyses are closely related because much of the power supplied to the device dissipates as heat from the device; however, in some situations, the two types of analyses are not identical. For example, if you use terminated I/O standards, some of the power drawn from the power supply of the device dissipates in termination resistors rather than in the device.

Power analysis also addresses the activity of your design over time as a factor that impacts the power consumption of the device. Static power is the power consumption of the device regardless of your design activity. Dynamic power is the additional power consumption of the device due to signal activity or toggling.

- For power supply planning, you can use the PowerPlay EPE at the early stages of your design cycle. Alternatively, you can also use the PowerPlay Power Analyzer reports when your design is complete to get an estimate of your design power requirement.

- For system-on-a-chip (SoC) power estimation, you can use the HPS Power Calculator at the design implementation stage of your design cycle to include the hard processor system (HPS) power.

For more information about the HPS Power Calculator, refer to “Using the HPS Power Calculator” on page 8–7.
Factors Affecting Power Consumption

Understanding factors that affect power consumption allows you to use the PowerPlay Power Analyzer and interpret its results effectively.

Device Selection

Device families have different power characteristics. Many parameters affect the device family power consumption, including choice of process technology, supply voltage, electrical design, and device architecture.

Power consumption also varies in a single device family. A larger device consumes more static power than a smaller device in the same family because of its larger transistor count. Dynamic power can also increase with device size in devices that employ global routing architectures.

The choice of device package also affects the ability of the device to dissipate heat. This choice can impact your required cooling solution choice to comply to junction temperature constraints.

Process variation can affect power consumption. Process variation primarily impacts static power because sub-threshold leakage current varies exponentially with changes in transistor threshold voltage. So, you must consult device specifications for static power and not rely on empirical observation. Process variation has a weak effect on dynamic power.

Environmental Conditions

Operating temperature primarily affects device static power consumption. Higher junction temperatures result in higher static power consumption. The device thermal power and cooling solution that you use must result in the device junction temperature remaining within the maximum operating range for the device. The main environmental parameters affecting junction temperature are the cooling solution and ambient temperature.

Airflow

Airflow is a measure of how quickly the device removes heated air from the vicinity of the device and replaces air at ambient temperature. You can either specify airflow as “still air” when you are not using a fan, or as the linear feet per minute rating of the fan in the system. Higher airflow decreases thermal resistance.

Heat Sink and Thermal Compound

A heat sink allows more efficient heat transfer from the device to the surrounding area because of its large surface area exposed to the air. The thermal compound that interfaces the heat sink to the device also influences the rate of heat dissipation. The case-to-ambient thermal resistance ($\theta_{CA}$) parameter describes the cooling capacity of the heat sink and thermal compound employed at a given airflow. Larger heat sinks and more effective thermal compounds reduce $\theta_{CA}$. 
Junction Temperature

The junction temperature of a device is equal to:

\[ T_{\text{junction}} = T_{\text{ambient}} + P_{\text{thermal}} \cdot \theta_{JA} \]

in which \( \theta_{JA} \) is the total thermal resistance from the device transistors to the environment, having units of degrees Celsius per watt. The value \( \theta_{JA} \) is equal to the sum of the junction-to-case (package) thermal resistance (\( \theta_{JC} \)), and the case-to-ambient thermal resistance (\( \theta_{CA} \)) of your cooling solution.

Board Thermal Model

The junction-to-board thermal resistance (\( \theta_{JB} \)) is the thermal resistance of the path through the board, having units of degrees Celsius per watt. To compute junction temperature, you can use this board thermal model along with the board temperature, the top-of-chip \( \theta_{JA} \) and ambient temperatures.

Device Resource Usage

The number and types of device resources used greatly affects power consumption.

Number, Type, and Loading of I/O Pins

Output pins drive off-chip components, resulting in high-load capacitance that leads to a high-dynamic power per transition. Terminated I/O standards require external resistors that draw constant (static) power from the output pin.

Number and Type of Logic Elements, Multiplier Elements, and RAM Blocks

A design with more logic elements (LEs), multiplier elements, and memory blocks tends to consume more power than a design with fewer circuit elements. The operating mode of each circuit element also affects its power consumption. For example, a DSP block performing 18 × 18 multiplications and a DSP block performing multiply-accumulate operations consume different amounts of dynamic power because of different amounts of charging internal capacitance on each transition. The operating mode of a circuit element also affects static power.

Number and Type of Global Signals

Global signal networks span large portions of the device and have high capacitance, resulting in significant dynamic power consumption. The type of global signal is important as well. For example, Stratix II devices support several kinds of global clock networks that span either the entire device or a specific portion of the device (a regional clock network covers a quarter of the device). Clock networks that span smaller regions have lower capacitance and tend to consume less power. The location of the logic array blocks (LABs) driven by the clock network can also have an impact because the Quartus II software automatically disables unused branches of a clock.
Signal Activities

The final important factor in estimating power consumption is the behavior of each signal in your design. The two vital statistics are the toggle rate and the static probability.

The toggle rate of a signal is the average number of times that the signal changes value per unit of time. The units for toggle rate are transitions per second and a transition is a change from \(1\) to \(0\), or \(0\) to \(1\).

The static probability of a signal is the fraction of time that the signal is logic \(1\) during the period of device operation that is being analyzed. Static probability ranges from \(0\) (always at ground) to \(1\) (always at logic-high).

Dynamic power increases linearly with the toggle rate as you charge the capacitive load more frequently for logic and routing. The Quartus II software models full rail-to-rail switching. For high toggle rates, especially on circuit output I/O pins, the circuit can transition before fully charging the downstream capacitance. The result is a slightly conservative prediction of power by the PowerPlay Power Analyzer.

Static probabilities of their input signals can sometimes affect the static power that routing and logic consume. This effect is due to state-dependent leakage and has a larger effect on smaller process geometries. The Quartus II software models this effect on devices at 90 nm (or smaller) if it is important to the power estimate. The static power also varies with the static probability of a logic \(1\) or \(0\) on the I/O pin when output I/O standards drive termination resistors.

To get accurate results from the power analysis, the signal activities for analysis must represent the actual operating behavior of your design. Inaccurate signal toggle rate data is the largest source of power estimation error.
Creating PowerPlay EPE Spreadsheets

You can use PowerPlay EPE spreadsheets to perform a preliminary thermal analysis and power consumption estimate for your design. You can either enter the data manually, or use the tools in the Quartus II software to assist you with generating the device resources usage information.

For more information about generating a PowerPlay EPE File in the Quartus II software, refer to *Performing an Early Power Estimate Using the PowerPlay Early Power Estimator* in Quartus II Help.

The PowerPlay EPE spreadsheet includes the Import Data macro that parses the information in the PowerPlay EPE File, and then transfers the data into the spreadsheet. If you do not want to use the macro, you can manually transfer the information into the PowerPlay EPE spreadsheet.

After importing the PowerPlay EPE File information into the PowerPlay EPE spreadsheet, you can add additional device resource information at any time. If the existing Quartus II project represents only a portion of your full design, you must enter the additional device resources used in the final design manually.

**PowerPlay EPE File Generator Compilation Report**

After successfully generating the PowerPlay EPE File, you can locate a PowerPlay EPE File Generator report under the Compilation Report. The Compilation Report contains different sections, such as Summary, Settings, Generated Files, Confidence Metric Details, and Signal Activities. For more information about the PowerPlay EPE File Generator report, refer to “PowerPlay Power Analyzer Compilation Report” on page 8–20.

Table 8–1 lists the differences between the PowerPlay EPE and the Quartus II PowerPlay Power Analyzer.

<table>
<thead>
<tr>
<th>Characteristic</th>
<th>PowerPlay EPE</th>
<th>Quartus II PowerPlay Power Analyzer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Phase in the design cycle</td>
<td>Any time</td>
<td>Post-fit</td>
</tr>
<tr>
<td>Tool requirements</td>
<td>Spreadsheet program or the Quartus II software</td>
<td>The Quartus II software</td>
</tr>
<tr>
<td>Accuracy</td>
<td>Medium</td>
<td>Medium to very high</td>
</tr>
<tr>
<td>Data inputs</td>
<td>Resource usage estimates</td>
<td>Post-fit design</td>
</tr>
<tr>
<td></td>
<td>Clock requirements</td>
<td>Clock requirements</td>
</tr>
<tr>
<td></td>
<td>Environmental conditions</td>
<td>Signal activity defaults</td>
</tr>
<tr>
<td></td>
<td>Toggle rate</td>
<td>Environmental conditions</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Register transfer level (RTL) simulation results (optional)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Post-fit simulation results (optional)</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Signal activities per node or entity (optional)</td>
</tr>
</tbody>
</table>
The result of the PowerPlay Power Analyzer is only an estimation of power. Altera does not recommend using the result as a specification. The purpose of the estimation is to help you to establish guidelines for the power budget of your design. Altera recommends that you measure the actual power on the board. You must measure the total dynamic current of your design during device operation because the estimate is design dependent and depends on many variable factors, including input vector quantity, quality, and exact loading conditions of a PCB design. You must not base static power consumption on empirical observation. You must use the reported values by the PowerPlay Power Analyzer or data sheet because the tested devices might not exhibit worst-case behavior.

### Using the HPS Power Calculator

The HPS Power Calculator allows you to enable HPS power in the PowerPlay Power Analyzer. You can also use the HPS Power Calculator to estimate the HPS power for a given frequency.

If you want to estimate the HPS power without running the PowerPlay Power Analyzer, you can use the HPS Power Calculator to vary the frequency and view the HPS power estimation.

If you want the PowerPlay Power Analyzer to report the SoC power estimation for your device, including the HPS, then turn on the **Enable HPS** option in the **HPS Power Calculator** dialog box and then set the HPS parameters.

For more information about using the HPS Power Calculator, refer to the **Performing Power Analysis with the PowerPlay Power Analyzer** topic in Quartus II Help.
PowerPlay Power Analyzer Flow

The PowerPlay Power Analyzer supports accurate power estimations by allowing you to specify the important design factors affecting power consumption. Figure 8–2 shows the high-level PowerPlay Power Analyzer flow.

To obtain accurate I/O power estimates, the PowerPlay Power Analyzer requires you to synthesize your design and then fit your design to the target device. You must specify the electrical standard on each I/O cell and the capacitive load on each I/O standard in your design.

Operating Settings and Conditions

You can specify device power characteristics, operating voltage conditions, and operating temperature conditions for power analysis in the Quartus II software.

On the Operating Settings and Conditions page of the Settings dialog box, you can specify whether the device has typical power consumption characteristics or maximum power consumption characteristics.

For more information, refer to Operating Setting and Conditions Page (Settings Dialog Box) in Quartus II Help.

On the Voltage page of the Settings dialog box, you can view the operating voltage conditions for each power rail in the device, and specify supply voltages for power rails with selectable supply voltages.

For more information, refer to Voltage Page (Settings Dialog Box) in Quartus II Help.

On the Temperature page of the Settings dialog box, you can specify the thermal operating conditions of the device.
Signal Activities Data Sources

The PowerPlay Power Analyzer provides a flexible framework for specifying signal activities. The framework reflects the importance of using representative signal-activity data during power analysis. You can use the following sources to provide information about signal activity:

- Simulation results
- User-entered node, entity, and clock assignments
- User-entered default toggle rate assignment
- Vectorless estimation

The PowerPlay Power Analyzer allows you to mix and match the signal-activity data sources on a signal-by-signal basis. Figure 8–3 shows the priority scheme. The following sections describe the data sources.

Figure 8–3. Signal-Activity Data Source Priority Scheme

![Signal-Activity Data Source Priority Scheme Diagram]

**Note to Figure 8–3:**

(1) Vectorless estimation is available only for some device families. For more information, refer to Performing Power Analysis with the PowerPlay Power Analyzer.

Simulation Results

The PowerPlay Power Analyzer directly reads the waveforms generated by a design simulation. You can calculate the static probability and toggle rate for each signal from the simulation waveform. Power analysis is most accurate when you use representative input stimuli to generate simulations.

The PowerPlay Power Analyzer reads results generated by the following simulators:

- ModelSim®
- ModelSim-Altera
- QuestaSim
Signal activity and static probability information derive from a Verilog Value Change Dump File (.vcd). For more information, refer to “Signal Activities” on page 8–5.

For third-party simulators, use the EDA Tool Settings to specify the Generate Value Change Dump (VCD) file script option in the Simulation page of the Settings dialog box. These scripts instruct the third-party simulators to generate a .vcd that encodes the simulated waveforms. The Quartus II PowerPlay Power Analyzer reads this file directly to derive the toggle rate and static probability data for each signal.

Third-party EDA simulators, other than those listed, can generate a .vcd that you can use with the PowerPlay Power Analyzer. For those simulators, you must manually create a simulation script to generate the appropriate .vcd.

You can use a .vcd created for power analysis to optimize your design for power during fitting by utilizing the appropriate settings in the PowerPlay power optimization list, available in the Fitter Settings page of the Settings dialog box.

For more information about power optimization, refer to the Power Optimization chapter in volume 2 of the Quartus II Handbook. For more information about how to create a .vcd in other third-party EDA simulation tools, refer to Section I. Simulation in volume 3 of the Quartus II Handbook.

Using Simulation Files in Modular Design Flows

A common design practice is to create modular or hierarchical designs in which you develop each design entity separately, and then instantiate your design in a higher-level entity to form a complete design. You can perform simulation on a complete design or on each modular design for verification. The PowerPlay Power Analyzer supports modular design flows when reading the signal activities from simulation files. Figure 8–4 shows an example of a modular design flow.

Figure 8–4. Modular Simulation Flow
When specifying a simulation file, the software provides an associated design entity name, such that the PowerPlay Power Analyzer imports the signal activities derived from the simulation file (.vcd) for that design entity. The PowerPlay Power Analyzer also supports the specification of multiple .vcd files for power analysis, with each having an associated design entity name to enable the integration of partial design simulations into a complete design power analysis. When specifying multiple .vcd files for your design, more than one simulation file should contain signal-activity information for the same signal. When you apply multiple .vcd files to the same design entity, the signal activity used in the power analysis is the equal-weight arithmetic average of each .vcd. When you apply multiple simulation files to design entities at different levels in your design hierarchy, the signal activity in the power analysis derives from the simulation file that applies to the most specific design entity.

Figure 8–5 shows an example of a hierarchical design. The top-level module of your design, called Top, consists of three 8b/10b decoders, followed by a multiplexer. The software then encodes the output of the multiplexer again before being the output from your design. An error-handling module handles any 8b/10b decoding errors. The Top module contains the top-level entity of your design and any logic not defined as part of another module. The design file for the top-level module might be a wrapper for the hierarchical entities below it, or it might contain its own logic. The following usage scenarios show common ways that you can simulate your design and import the .vcd into the PowerPlay Power Analyzer.

**Figure 8–5. Example Hierarchical Design**

---

**Complete Design Simulation**

You can simulate the entire design and generate a .vcd from a third-party simulator. The PowerPlay Power Analyzer can then import the .vcd (specifying the top-level design). The resulting power analysis uses the signal activities information from the generated .vcd, including those that apply to submodules, such as decode [1-3], err1, mux1, and encode1.
Modular Design Simulation

You can independently simulate submodules of the top-level design, and then import all the resulting .vcd files into the PowerPlay Power Analyzer. For example, you can simulate the $8b10b\_dec$ independent of the entire design and the multiplexer, $8b10b\_rxerr$, and $8b10b\_enc$. You can then import the .vcd files generated from each simulation by specifying the appropriate instance name. For example, if the files produced by the simulations are $8b10b\_dec\_vcd$, $8b10b\_enc\_vcd$, $8b10b\_rxerr\_vcd$, and $mux\_vcd$, you can use the import specifications in Table 8–2.

### Table 8–2. Import Specifications

<table>
<thead>
<tr>
<th>File Name</th>
<th>Entity</th>
</tr>
</thead>
<tbody>
<tr>
<td>$8b10b_dec_vcd$</td>
<td>Top</td>
</tr>
<tr>
<td>$8b10b_dec_vcd$</td>
<td>Top</td>
</tr>
<tr>
<td>$8b10b_dec_vcd$</td>
<td>Top</td>
</tr>
<tr>
<td>$8b10b_rxerr_vcd$</td>
<td>Top</td>
</tr>
<tr>
<td>$8b10b_enc_vcd$</td>
<td>Top</td>
</tr>
<tr>
<td>$mux_vcd$</td>
<td>Top</td>
</tr>
</tbody>
</table>

The resulting power analysis applies the simulation vectors in each file to the assigned entity. Simulation provides signal activities for the pins and for the outputs of functional blocks. If the inputs to an entity instance are input pins for the entire design, the simulation file associated with that instance does not provide signal activities for the inputs of that instance. For example, an input to an entity such as $mux1$ has its signal activity specified at the output of one of the decode entities.

Multiple Simulations on the Same Entity

You can perform multiple simulations of an entire design or specific modules of a design. For example, in the process of verifying the top-level design, you can have three different simulation testbenches: one for normal operation, and two for corner cases. Each of these simulations produces a separate .vcd. In this case, apply the different .vcd file names to the same top-level entity, as shown in Table 8–3.

### Table 8–3. Multiple Simulation File Names and Entities

<table>
<thead>
<tr>
<th>File Name</th>
<th>Entity</th>
</tr>
</thead>
<tbody>
<tr>
<td>normal.vcd</td>
<td>Top</td>
</tr>
<tr>
<td>corner1.vcd</td>
<td>Top</td>
</tr>
<tr>
<td>corner2.vcd</td>
<td>Top</td>
</tr>
</tbody>
</table>

The resulting power analysis uses an arithmetic average that the signal activities calculated from each simulation file to obtain the final signal activities used. If a signal $err\_out$ has a toggle rate of zero toggles per second in normal.vcd, 50 toggles per second in corner1.vcd, and 70 toggles per second in corner2.vcd, the final toggle rate in the power analysis is 40 toggles per second.
Overlapping Simulations

You can perform a simulation on the entire design, and more exhaustive simulations on a submodule, such as 8b10b_rxerr. Table 8–4 lists the import specification for overlapping simulations.

Table 8–4. Overlapping Simulation Import Specifications

<table>
<thead>
<tr>
<th>File Name</th>
<th>Entity</th>
</tr>
</thead>
<tbody>
<tr>
<td>full_design.vcd</td>
<td>Top</td>
</tr>
<tr>
<td>error_cases.vcd</td>
<td>Top</td>
</tr>
</tbody>
</table>

In this case, the software uses signal activities from error_cases.vcd for all the nodes in the generated .vcd and uses signal activities from full_design.vcd for only those nodes that do not overlap with nodes in error_cases.vcd. In general, the more specific hierarchy (the most bottom-level module) derives signal activities for overlapping nodes.

Partial Simulations

You can perform a simulation in which the entire simulation time is not applicable to signal-activity calculation. For example, if you run a simulation for 10,000 clock cycles and reset the chip for the first 2,000 clock cycles. If the PowerPlay Power Analyzer performs the signal-activity calculation over all 10,000 cycles, the toggle rates are only 80% of their steady state value (because the chip is in reset for the first 20% of the simulation). In this case, you must specify the useful parts of the .vcd for power analysis. The Limit VCD Period option enables you to specify a start and end time when performing signal-activity calculations.

Node Name Matching Considerations

Node name mismatches happen when you have .vcd applied to entities other than the top-level entity. In a modular design flow, the gate-level simulation files created in different Quartus II projects might not match their node names with the current Quartus II project.

For example, if you have a file named 8b10b_enc.vcd, which the Quartus II software generates in a separate project called 8b10b_enc and is simulating the 8b10b encoder. You can import the .vcd into another project called Top, you might encounter name mismatches when applying the .vcd to the 8b10b_enc module in the Top project. This mismatch happens because the Quartus II software might name all the combinational nodes in the 8b10b_enc.vcd differently than in the Top project.

You can avoid name mismatching with only RTL simulation data, in which register names do not change, or with an incremental compilation flow that preserves node names along with a gate-level simulation.

To ensure accuracy, Altera recommends that you use an incremental compilation flow to preserve the node names of your design.

For more information about the incremental compilation flow, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Glitch Filtering

The PowerPlay Power Analyzer defines a glitch as two signal transitions so closely spaced in time that the pulse, or glitch, occurs faster than the logic and routing circuitry can respond. The output of a transport delay model simulator contains glitches for some signals. The logic and routing structures of the device form a low-pass filter that filters out glitches that are tens to hundreds of picoseconds long, depending on the device family.

Some third-party simulators use different models than the transport delay model as the default model. Different models cause differences in signal activity and power estimation. The inertial delay model, which is the ModelSim default model, filters out more glitches than the transport delay model and usually yields a lower power estimate.

Altera recommends that you use the transport simulation model when using the Quartus II software glitch filtering support with third-party simulators. Simulation glitch filtering has little effect if you use the inertial simulation model.

For more information about how to set the simulation model type for your specific simulator, refer to Quartus II Help.

Glitch filtering in a simulator can also filter a glitch on one LE (or other circuit element) output from propagating to downstream circuit elements to ensure that the glitch does not affect simulated results. Glitch filtering prevents a glitch on one signal from producing non-physical glitches on all downstream logic, which can result in a signal toggle rate and a power estimate that are too high. Circuit elements in which every input transition produces an output transition, including multipliers and logic cells configured to implement XOR functions, are especially prone to glitches. Therefore, circuits with such functions can have power estimates that are too high when you do not use glitch filtering.

Altera recommends that you use the glitch filtering feature to obtain the most accurate power estimates. For .vcd files, the PowerPlay Power Analyzer flows support two levels of glitch filtering.

To enable the first level of glitch filtering in the Quartus II software for supported third-party simulators, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, select Simulation under EDA Tool Settings.
3. Select the Tool name to use for the simulation.
4. Turn on Enable glitch filtering.

The second level of glitch filtering occurs while the PowerPlay Power Analyzer is reading the .vcd generated by a third-party simulator. To enable the second level of glitch filtering, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, select PowerPlay Power Analyzer Settings.
3. Under Input File(s), turn on Perform glitch filtering on VCD files.
The .vcd file reader performs complementary filtering to the filtering performed during simulation and is often not as effective. While the .vcd file reader can remove glitches on logic blocks, the file reader cannot determine how a given glitch affects downstream logic and routing, and may eliminate the impact of the glitch completely. Filtering the glitches during simulation avoids switching downstream routing and logic automatically.

When running simulation for design verification (rather than to produce input to the PowerPlay Power Analyzer), Altera recommends that you turn off the glitch filtering option to produce the most rigorous and conservative simulation from a functionality viewpoint. When performing simulation to produce input for the PowerPlay Power Analyzer, Altera recommends that you turn on the glitch filtering to produce the most accurate power estimates.

Node and Entity Assignments

You can assign toggle rates and static probabilities to individual nodes and entities in the design. Toggle assignments have the highest priority, overriding data from all other signal-activity sources.

You must use the Assignment Editor or Tcl commands to create the Power Toggle Rate and Power Static Probability assignments. You can specify the power toggle rate as an absolute toggle rate in transitions using the Power Toggle Rate assignment, or you can use the Power Toggle Rate Percentage assignment to specify a toggle rate relative to the clock domain of the assigned node for a more specific assignment made in terms of hierarchy level.

If you use the Power Toggle Rate Percentage assignment, and the node does not have a clock domain, the Quartus II software issues a warning and ignores the assignment.

For more information about how to use the Assignment Editor in the Quartus II software, refer to the Constraining Designs chapter in volume 2 of the Quartus II Handbook.

Assigning toggle rates and static probabilities to individual nodes and entities is appropriate for signals in which you have knowledge of the signal or entity being analyzed. For example, if you know that a 100 MHz data bus or memory output produces data that is essentially random (uncorrelated in time), you can directly enter a 0.5 static probability and a toggle rate of 50 million transitions per second.

The PowerPlay Power Analyzer treats bidirectional I/O pins differently. The combinational input port and the output pad for a pin share the same name. However, those ports might not share the same signal activities. For reading signal-activity assignments, the PowerPlay Power Analyzer creates a distinct name <node_name~output> when configuring the bidirectional signal as an output and <node_name~result> when configuring the signal as an input. For example, if a design has a bidirectional pin named MYPIN, assignments for the combinational input use the name MYPIN-result, and the assignments for the output pad use the name MYPIN-output.
When you create the logic assignment in the Assignment Editor, you cannot find the *MYPIN~result* and *MYPIN~output* node names in the Node Finder. Therefore, to create the logic assignment, you must manually enter the two differentiating node names to create the assignment for the input and output port of the bidirectional pin.

**Timing Assignments to Clock Nodes**

For clock nodes, the PowerPlay Power Analyzer uses timing requirements to derive the toggle rate when neither simulation data nor user-entered signal-activity data is available. *f_{MAX}* requirements specify full cycles per second, but each cycle represents a rising transition and a falling transition. For example, a clock *f_{MAX}* requirement of 100 MHz corresponds to 200 million transitions per second.

**Default Toggle Rate Assignment**

You can specify a default toggle rate for primary inputs and other nodes in your design. The PowerPlay Power Analyzer uses the default toggle rate when no other method specifies the signal-activity data.

The PowerPlay Power Analyzer specifies the toggle rate in absolute terms (transitions per second), or as a fraction of the clock rate in effect for each node. The toggle rate for a clock derives from the timing settings for the clock. For example, if the PowerPlay Power Analyzer specifies a clock with an *f_{MAX}* constraint of 100 MHz and a default relative toggle rate of 20%, nodes in this clock domain transition in 20% of the clock periods, or 20 million transitions occur per second. In some cases, the PowerPlay Power Analyzer cannot determine the clock domain for a node because either the PowerPlay Power Analyzer does not specify a clock domain for the node, or the clock domain is ambiguous. In these cases, the PowerPlay Power Analyzer substitutes and reports a toggle rate of zero.

**Vectorless Estimation**

For some device families, the PowerPlay Power Analyzer automatically derives estimates for signal activity on nodes with no simulation or user-entered signal-activity data. Vectorless estimation statistically estimates the signal activity of a node based on the signal activities of nodes feeding that node, and on the actual logic function that the node implements. Vectorless estimation cannot derive signal activities for primary inputs. Vectorless estimation is accurate for combinational nodes, but not for registered nodes. Therefore, the PowerPlay Power Analyzer requires simulation data for at least the registered nodes and I/O nodes for accuracy.

For more information, refer to *Performing Power Analysis with the PowerPlay Power Analyzer* in Quartus II Help.

The **PowerPlay Power Analyzer Settings** dialog box allows you disable vectorless estimation. When turned on, vectorless estimation requires priority over default toggle rates. Vectorless estimation does not override clock assignments.
Using the PowerPlay Power Analyzer

For flows that use the PowerPlay Power Analyzer, you must first synthesize your design, and then fit it to the target device. You must either provide timing assignments for all the clocks in your design, or use a simulation-based flow to generate activity data. You must specify the I/O standard on each device input or output and the capacitive load on each output in your design.

For more information about using the PowerPlay Power Analyzer, refer to Performing Power Analysis with the PowerPlay Power Analyzer in Quartus II Help.

Common Analysis Flows

You can use the analysis flows in this section with the PowerPlay Power Analyzer. However, vectorless activity estimation is only available for some device families.

Signal Activities from Full Post-Fit Netlist (Timing) Simulation

Timing Simulation flow provides the most accuracy because all node activities reflect actual design behavior, and, if supplied input vectors are representative of typical design operation. Results are better if the simulation filters glitches. The disadvantage of this method is that the simulation time is long.

Signal Activities from Full Post-Fit Netlist (Zero Delay) Simulation

You can use the zero delay simulation flow with designs for which signal activities from a full post-fit netlist (timing) simulation are not available. Zero delay simulation is as accurate as timing simulation in 95% of designs with no glitches. If your design has glitches, the power estimation may not be accurate. Altera recommends that you use the signal activities from a full post-fit netlist (timing) simulation to achieve an accurate power estimation of your design.

The following designs might exhibit glitches:
- Designs with many XOR gates (for example, an encryption core)
- Designs with arithmetic blocks without input and output registers (DSPs and carry chains)

For more information about creating zero delay simulation signal activities, refer to “Generating a .vcd from Full Post-Fit Netlist (Zero Delay) Simulation” on page 8–20.

Signal Activities from RTL (Functional) Simulation, Supplemented by Vectorless Estimation

In the functional simulation flow, simulation provides toggle rates and static probabilities for all pins and registers in your design. Vectorless estimation fills in the values for all the combinational nodes between pins and registers, giving good results. This flow usually provides a compilation time benefit when you use the third-party RTL simulator.
RTL simulation may not provide signal activities for all registers in the post-fitting netlist because synthesis loses some register names. For example, synthesis might automatically transform state machines and counters, thus changing the names of registers in those structures.

**Signal Activities from Vectorless Estimation and User-Supplied Input Pin Activities**

The vectorless estimation flow provides a low level of accuracy, because vectorless estimation for registers is not entirely accurate.

**Signal Activities from User Defaults Only**

The user defaults only flow provides the lowest degree of accuracy.

**Generating a .vcd**

In previous versions of the Quartus II software, you could use either the Quartus II simulator or an EDA simulator to perform your simulation. The Quartus II software no longer supports a built-in simulator, and you must use EDA simulators to perform simulation. Use the .vcd as the input to the PowerPlay Power Analyzer to estimate power for your design.

For more information about the supported third-party simulators, refer to “Simulation Results” on page 8–9.

To create a .vcd for your design, follow these steps:

1. On the Assignments menu, click Settings.
2. In the Category list, under EDA Tool Settings, click Simulation.
3. In the Tool name list, select your preferred EDA simulator.
4. In the Format for output netlist list, select Verilog HDL, or SystemVerilog HDL, or VHDL.
5. Turn on Generate Value Change Dump (VCD) file script.

   This option turns on the Map illegal HDL characters and Enable glitch filtering options. The Map illegal HDL characters option ensures that all signals have legal names and that signal toggle rates are available later in the PowerPlay Power Analyzer.

6. By turning on Enable glitch filtering, glitch filtering logic is the output when you generate an EDA netlist for simulation. This option is available regardless of whether or not you want to generate .vcd scripts. For more information about glitch filtering, refer to “Glitch Filtering” on page 8–14.
When performing simulation using ModelSim, the `+nospecify` option for the `vsim` command disables the `specify path delays and timing checks` option in ModelSim. By enabling glitch filtering on the Simulation page, the simulation models include specified path delays. Thus, ModelSim might fail to simulate a design if you enabled glitch filtering and specified the `+nospecify` option. Altera recommends that you remove the `+nospecify` option from the ModelSim `vsim` command to ensure accurate simulation for power estimation.

7. Click Script Settings.
   
   Select the signals that you want to output to the `.vcd`. With All signals selected, the generated script instructs the third-party simulator to write all connected output signals to the `.vcd`. With All signals except combinational lcell outputs selected, the generated script tells the third-party simulator to write all connected output signals to the `.vcd`, except logic cell combinational outputs.

   The file can become extremely large if you write all output signals to the file because the file size depends on the number of output signals being monitored and the number of transitions that occur.

8. Click OK.

9. Type a name for your testbench in the Design instance name box.

10. Compile your design with the Quartus II software and generate the necessary EDA netlist and script that instructs the third-party simulator to generate a `.vcd`.

   For more information about the NativeLink feature, refer to Section I. Simulation in volume 3 of the Quartus II Handbook.

11. Perform a simulation with the third-party EDA simulation tool. Call the generated script in the simulation tool before running the simulation. The simulation tool generates the `.vcd` and places it in the project directory.

**Generating a `.vcd` from ModelSim Software**

To produce a `.vcd` with the ModelSim software, follow these steps:

1. In the Quartus II software, on the Assignments menu, click Settings.

2. In the Category list, under EDA Tool Settings, click Simulation.

3. In the Tool name list, select your preferred EDA simulator.

4. In the Format for output netlist list, select Verilog HDL, or SystemVerilog HDL, or VHDL.

5. Turn on Generate Value Change Dump (VCD) file script.

6. To generate the `.vcd`, perform a full compilation.

7. In the ModelSim software, compile the files necessary for simulation.

8. Load your design by clicking Start Simulation on the Tools menu, or use the `vsim` command.
9. Use the `.vcd` script created in step 6 using the following command:
   
   ```
   source <design>_dump_all_vcd_nodes.tcl
   ```

10. Run the simulation (for example, run 2000ns or run -all).

11. Quit the simulation using the `quit -sim` command, if required.

12. Exit the ModelSim software. If you do not exit the software, the ModelSim software might end the writing process of the `.vcd` improperly, resulting in a corrupt `.vcd`.

---

### Generating a `.vcd` from Full Post-Fit Netlist (Zero Delay) Simulation

To successfully generate a `.vcd` from the full post-fit Netlist (zero delay) simulation, follow these steps:

1. Compile your design in the Quartus II software to generate the Netlist `<project_name>.vo`.

2. In `<project_name>.vo`, search for the include statement for `<project_name>.sdo`, comment the statement out, and save the file.

   - Altera recommends that you use the Standard Delay Format Output File (`.sdo`) for gate-level timing simulation. The `.sdo` contains the delay information of each architecture primitive and routing element in your design; however, you must exclude the `.sdo` for zero delay simulation.

3. Generate a `.vcd` for power estimation by performing the steps in “Generating a `.vcd”” on page 8–18.

   - For more information about how to create a `.vcd` in other third-party EDA simulation tools, refer to Section I. Simulation in volume 3 of the Quartus II Handbook.

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### Importance of `.vcd`

Altera recommends using a `.vcd` or a `.saf` generated by gate-level timing simulation for an accurate power estimation because gate-level timing simulation takes all the routing resources and the exact logic array resource usage into account.

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### Running the PowerPlay Power Analyzer Using the Quartus II GUI

To run the PowerPlay Power Analyzer using the Quartus II GUI, refer to Performing Power Analysis with the PowerPlay Power Analyzer in Quartus II Help.

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### PowerPlay Power Analyzer Compilation Report

The PowerPlay Power Analyzer section of the Compilation Report consists of the following sections.
Summary
The Summary section of the report shows the estimated total thermal power consumption of your design. This includes dynamic, static, and I/O thermal power consumption. The I/O thermal power consumption is the total I/O power the \( V_{CCIO} \) power supplies and some portion of the \( V_{CCINT} \) contribute. The report also includes a confidence metric that reflects the overall quality of the data sources for the signal activities. For example, a Low power estimation confidence value reflects that you have provided insufficient toggle rate data, or most of the signal-activity information used for power estimation is from default or vectorless estimation settings. For more information about the input data, refer to the PowerPlay Power Analyzer Confidence Metric report.

Settings
The Settings section of the report shows the PowerPlay Power Analyzer settings information of your design, including the default input toggle rates, operating conditions, and other relevant setting information.

Simulation Files Read
The Simulation Files Read section of the report lists the simulation output file that the .vcd used for power estimation. This section also includes the file ID, file type, entity, VCD start time, VCD end time, the unknown percentage, and the toggle percentage. The unknown percentage indicates the portion of the design module unused by the simulation vectors.

Operating Conditions Used
The Operating Conditions Used section of the report shows device characteristics, voltages, temperature, and cooling solution, if any, during the power estimation. This section also shows the entered junction temperature or auto-computed junction temperature during the power analysis.

Thermal Power Dissipated by Block
The Thermal Power Dissipated by Block section of the report shows estimated thermal dynamic power and thermal static power consumption categorized by atoms. This information provides you with estimated power consumption for each atom in your design.

Thermal Power Dissipation by Block Type (Device Resource Type)
This Thermal Power Dissipation by Block Type (Device Resource Type) section of the report shows the estimated thermal dynamic power and thermal static power consumption categorized by block types. This information is further categorized by estimated dynamic and static power and provides an average toggle rate by block type. Thermal power is the power dissipated as heat from the FPGA device.
Thermal Power Dissipation by Hierarchy

This Thermal Power Dissipation by Hierarchy section of the report shows estimated thermal dynamic power and thermal static power consumption categorized by design hierarchy. This information is further categorized by the dynamic and static power that was used by the blocks and routing in that hierarchy. This information is useful when locating problem modules in your design.

Core Dynamic Thermal Power Dissipation by Clock Domain

The Core Dynamic Thermal Power Dissipation by Clock Domain section of the report shows the estimated total core dynamic power dissipation by each clock domain, which provides designs with estimated power consumption for each clock domain in the design. If the clock frequency for a domain is unspecified by a constraint, the clock frequency is listed as “unspecified.” For all the combinational logic, the clock domain is listed as no clock with zero MHz.

Current Drawn from Voltage Supplies

The Current Drawn from Voltage Supplies section of the report lists the current drawn from each voltage supply. The $V_{CCIO}$ voltage supply is further categorized by I/O bank and by voltage. This section also lists the minimum safe power supply size (current supply ability) for each supply voltage. Minimum current requirement can be higher than user mode current requirement in cases in which the supply has a specific power up current requirement that goes beyond user mode requirement, such as the $V_{CCPD}$ power rail in Stratix III and Stratix IV devices, and the $V_{CCIO}$ power rail in Stratix IV devices.

Transceiver-based devices have multiple voltage supplies. The report also shows the static and dynamic current (in mA) drawn from each voltage supply. The Thermal Power Dissipation by Block Type report section, which contains a row that starts with “GXB Transceiver”, lists the total static and dynamic power consumed by the transceivers on all voltage supplies.

The I/O thermal power dissipation on the summary page does not correlate directly to the power drawn from the $V_{CCIO}$ voltage supply listed in this report. This is because the I/O thermal power dissipation value also includes portions of the $V_{CCINT}$ power, such as the I/O element (IOE) registers, which are modeled as I/O power, but do not draw from the $V_{CCIO}$ supply.

The reported current drawn from the I/O Voltage Supplies (ICCIO) as reported in the PowerPlay Power Analyzer report includes any current drawn through the I/O into off-chip termination resistors. This can result in ICCIO values that are higher than the reported I/O thermal power, because this off-chip current dissipates as heat elsewhere and does not factor in the calculation of device temperature. Therefore, total I/O thermal power does not equal the sum of current drawn from each VCCIO supply multiplied by VCCIO voltage.
**Confidence Metric Details**

The Confidence Metric is defined in terms of the total weight of signal activity data sources for both combinational and registered signals. Each signal has two data sources allocated to it; a toggle rate source and a static probability source.

The Confidence Metric Details section also indicates the quality of the signal toggle rate data to compute a power estimate. The confidence metric is low if the signal toggle rate data comes from poor predictors of real signal toggle rates in the device during an operation. Toggle rate data that comes from simulation, user-entered assignments on specific signals, or entities are reliable. Toggle rate data from default toggle rates (for example, 12.5% of the clock period) or vectorless estimation are relatively inaccurate. This section gives an overall confidence rating in the toggle rate data, from low to high. This section also summarizes how many pins, registers, and combinational nodes obtained their toggle rates from each of simulation, user entry, vectorless estimation, or default toggle rate estimations. This detailed information helps you understand how to increase the confidence metric, letting you determine your own confidence in the toggle rate data.

**Signal Activities**

The Signal Activities section lists toggle rates and static probabilities assumed by power analysis for all signals with fan-out and pins. This section also lists the signal type (pin, registered, or combinational) and the data source for the toggle rate and static probability. By default, this section reports all signal activities, but you can turn off the report with the **Write signal activities to report file** option on the PowerPlay Power Analyzer Settings page.

Altera recommends that you turn off the **Write signal activities to report file** option for a large design because of the large number of signals present. You can use the Assignment Editor to specify that activities for individual nodes or entities are reported by assigning an **on** value to those nodes for the **Power Report Signal Activities** assignment.

**Messages**

The Messages section lists the messages that the Quartus II software generates during the analysis.

**Specific Rules for Reporting**

In a Stratix GX device, the PowerPlay Power Analyzer grouped the XGM II state machine block power into the power for the GXB transceivers. Therefore, the PowerPlay Power Analyzer reports the power for the XGM II state machine block as zero Watts.
Avoiding Power Estimation and Hardware Measurement Mismatch

You can avoid power estimation and hardware measurement mismatch by applying the recommended settings for the PowerPlay Power Analyzer described in “Recommended Setting for PowerPlay Power Analyzer” and “Recommended Setting for PowerPlay EPE”. These recommended settings can improve the accuracy of the power estimation on your completed design.

The default setting and vectorless setting cannot correctly estimate the power consumption of your design because these settings differ from the user input file setting, as well as for board measurement, by approximately 30%.

**Recommended Setting for PowerPlay Power Analyzer**

Altera recommends that you use the settings in Table 8–5 for an accurate power estimation of your design. The recommended settings yield power estimation results which closely match the board measurement.

<table>
<thead>
<tr>
<th>Parameters</th>
<th>User input File with SDC Constraint (mΩ)</th>
<th>Board Measurement (mΩ)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Total Thermal Power Dissipation</td>
<td>1774.23</td>
<td>1840</td>
</tr>
<tr>
<td>Core Dynamic Power Dissipation</td>
<td>1216.62</td>
<td>1120</td>
</tr>
<tr>
<td>Core Static Power Dissipation</td>
<td>541.92</td>
<td>720</td>
</tr>
<tr>
<td>I/O Power Dissipation</td>
<td>15.69</td>
<td>-</td>
</tr>
</tbody>
</table>

**Recommended Setting for PowerPlay EPE**

Altera recommends that you use the settings in Table 8–6 for an accurate power estimation of your design. The recommended settings yield power estimation results which closely match the board measurement.

<table>
<thead>
<tr>
<th>Parameters</th>
<th>EPE (mΩ)</th>
<th>Board Measurement (mΩ)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Total Thermal Power Dissipation</td>
<td>1748</td>
<td>1840</td>
</tr>
<tr>
<td>Core Dynamic Power Dissipation</td>
<td>1192</td>
<td>1120</td>
</tr>
<tr>
<td>Core Static Power Dissipation</td>
<td>556</td>
<td>720</td>
</tr>
<tr>
<td>I/O Power Dissipation</td>
<td>16</td>
<td>-</td>
</tr>
</tbody>
</table>
Scripting Support

You can run procedures and create settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For more information about scripting command options, refer to the Quartus II Command-Line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook and API Functions for Tcl in Quartus II Help. For more information about all settings and constraints in the Quartus II software, refer to the Quartus II Settings File Reference Manual. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

Running the PowerPlay Power Analyzer from the Command–Line

The executable to run the PowerPlay Power Analyzer is `quartus_pow`. For a complete listing of all command–line options supported by `quartus_pow`, type the following command at a system command prompt:

```
quartus_pow --help or quartus_sh --qhelp
```

The following is an example of using the `quartus_pow` executable:

- To instruct the PowerPlay Power Analyzer to generate a PowerPlay EPE File, type the following command at a system command prompt:
  
  ```
quartus_pow sample --output_epe=sample.csv
  ```

- To instruct the PowerPlay Power Analyzer to generate a PowerPlay EPE File without performing the power estimate, type the following command at a system command prompt:
  
  ```
quartus_pow sample --output_epe=sample.csv --estimate_power=off
  ```

- To instruct the PowerPlay Power Analyzer to use a `.vcd` as input (`sample.vcd`), type the following command at a system command prompt:
  
  ```
quartus_pow sample --input_vcd=sample.vcd
  ```

- To instruct the PowerPlay Power Analyzer to use two `.vcd` files as input files (`sample1.vcd` and `sample2.vcd`), perform glitch filtering on the `.vcd` and use a default input I/O toggle rate of 10,000 transitions per second, type the following command at a system command prompt:
  
  ```
quartus_pow sample --input_vcd=sample1.vcd --input_vcd=sample2.vcd
  --vcd_filter_glitches=on --
  default_input_io_toggle_rate=10000transitions/s
  ```
To instruct the PowerPlay Power Analyzer to not use an input file, a default input I/O toggle rate of 60%, no vectorless estimation, and a default toggle rate of 20% on all remaining signals, type the following command at a system command prompt:

```
quartus_pow sample --no_input_file --
default_input_io_toggle_rate=60\% \--use_vectorless_estimation=off --default_toggle_rate=20\%
```

No command-line options are available to specify the information found on the PowerPlay Power Analyzer Settings Operating Conditions page. Use the Quartus II GUI to specify these options.

The `quartus_pow` executable creates a report file, `<revision name>.pow.rpt`. You can locate the report file in the main project directory. The report file contains the same information in “PowerPlay Power Analyzer Compilation Report” on page 8–20.

## Document Revision History

Table 8–7 lists the revision history for this chapter.

### Table 8–7. Document Revision History (Part 1 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Updated “Types of Power Analyses” on page 8–2, and “Confidence Metric Details” on page 8–23.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “Importance of .vcd” on page 8–20, and “Avoiding Power Estimation and Hardware Measurement Mismatch” on page 8–24</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Updated “Current Drawn from Voltage Supplies” on page 8–22.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “Using the HPS Power Calculator” on page 8–7.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>■ Template update.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>■ Added links to Quartus II Help, removed redundant material.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Moved “Creating PowerPlay EPE Spreadsheets” to page 8–6.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor edits.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>■ Removed references to the Quartus II Simulator.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Table 8–1 on page 8–6, Table 8–2 on page 8–13, and Table 8–3 on page 8–14.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Figure 8–3 on page 8–9, Figure 8–4 on page 8–10, and Figure 8–5 on page 8–12.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Updated “Creating PowerPlay EPE Spreadsheets” on page 8–6 and “Simulation Results” on page 8–10.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “Signal Activities from Full Post-Fit Netlist (Zero Delay) Simulation” on page 8–19 and “Generating a .vcd from Full Post-Fit Netlist (Zero Delay) Simulation” on page 8–21.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor changes to “Generating a .vcd from ModelSim Software” on page 8–21.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Figure 11–8 on page 11–24.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>■ This chapter was chapter 11 in version 8.1.</td>
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</table>
### Table 8–7. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>- Updated for the Quartus II software version 8.1.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Replaced Figure 11-3.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Replaced Figure 11-14.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>- Updated Figure 11–5.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated “Types of Power Analyses” on page 11–5.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated “Operating Conditions” on page 11–9.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>- Updated “Current Drawn from Voltage Supplies” on page 11–32.</td>
</tr>
</tbody>
</table>

For previous versions of the *Quartus II Handbook*, refer to the *Quartus II Handbook Archive*. 
The Altera® Quartus® II design software provides a complete design debugging environment that easily adapts to your specific design requirements. This handbook is arranged in chapters, sections, and volumes that correspond to the major tools available for debugging your designs. For a general introduction to features and the standard design flow in the software, refer to the *Introduction to the Quartus II Software* manual.

This section is an introduction to System Debugging Tools and includes the following chapters:

- **Chapter 9, System Debugging Tools Overview**
  
  This chapter compares the various system debugging tools and explains when to use each of them.

- **Chapter 10, Analyzing and Debugging Designs with the System Console**
  
  This chapter describes the System Console Toolkit and compares the different capabilities within the toolkit.

- **Chapter 11, Debugging Transceiver Links**
  
  This chapter explains what functions are available within the Transceiver Toolkit and helps you decide which tool best meets your debugging needs.

- **Chapter 12, Quick Design Debugging Using SignalProbe**
  
  This chapter provides detailed instructions about how to use SignalProbe to quickly debug your design.

  Use this chapter to verify your design more efficiently by routing internal signals to I/O pins quickly without affecting the design.

- **Chapter 13, Design Debugging Using the SignalTap II Logic Analyzer**
  
  This chapter describes how to debug your FPGA design during normal device operation without the need for external lab equipment. Use this chapter to learn how to examine the behavior of internal signals, without using extra I/O pins, while the design is running at full speed on an FPGA device.

- **Chapter 14, In-System Debugging Using External Logic Analyzers**
  
  This chapter explains how to use external logic analyzers to debug designs on Altera devices.

- **Chapter 15, In-System Modification of Memory and Constants**
  
  This chapter explains how to use the Quartus II In-System Memory Content Editor as part of your FPGA design and verification flow to easily view and debug your design in the hardware lab.
Chapter 16, Design Debugging Using In-System Sources and Probes

This chapter provides detailed instructions about how to use the In-System Sources and Probes Editor and Tcl scripting in the Quartus® II software to debug your design.
The Altera® system debugging tools help you verify your FPGA designs. As your product requirements continue to increase in complexity, the time you spend on design verification continues to rise. This chapter provides a quick overview of the tools available in the system debugging suite and discusses the criteria for selecting the best tool for your design.

The Quartus® II software provides a portfolio of system design debugging tools for real-time verification of your design. Each tool in the system debugging portfolio uses a combination of available memory, logic, and routing resources to assist in the debugging process. The tools provide visibility by routing (or “tapping”) signals in your design to debugging logic. The debugging logic is then compiled with your design and downloaded into the FPGA or CPLD for analysis. Because different designs can have different constraints and requirements, such as the number of spare pins available or the amount of logic or memory resources remaining in the physical device, you can choose a tool from the available debugging tools that matches the specific requirements for your design.

### System Debugging Tools

Table 9–1 summarizes the tools in the system debugging tool suite that are covered in this section.

<table>
<thead>
<tr>
<th>Tool</th>
<th>Description</th>
<th>Typical Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td>System Console</td>
<td>This is a Tcl console that communicates to hardware modules instantiated into your design. You can use it with the Transceiver Toolkit to monitor or debug your design.</td>
<td>You need to perform system-level debugging. For example, if you have an Avalon-MM slave or Avalon-ST interfaces, you can debug your design at a transaction level. The tool supports JTAG connectivity, but also supports PLI connectivity to a simulation model, as well as TCP/IP connectivity to the target FPGA you wish to debug.</td>
</tr>
<tr>
<td>Transceiver Toolkit</td>
<td>The Transceiver Toolkit allows you to test and tune transceiver link signal quality. You can use a combination of bit error rate (BER), bathtub curve, and eye contour graphs as quality metrics. Auto Sweeping of physical medium attachment (PMA) settings allows you to quickly find an optimal solution.</td>
<td>You need to debug or optimize signal integrity of your board layout even before the actual design to be run on the FPGA is ready.</td>
</tr>
</tbody>
</table>
With the exception of SignalProbe, each of the on-chip debugging tools uses the JTAG port to control and read back data from debugging logic and signals under test. System Console uses JTAG and other interfaces as well. The JTAG resource is shared among all of the on-chip debugging tools.

For all system debugging tools except System Console, the Quartus II software compiles logic into your design automatically to distinguish between data and control information and each of the debugging logic blocks when the JTAG resource is required. This arbitration logic, also known as the System-Level Debugging (SLD) infrastructure, is shown in the design hierarchy of your compiled project as sld_hub:sld_hub_inst. The SLD logic allows you to instantiate multiple debugging blocks into your design and run them simultaneously. For System Console, you must explicitly insert IP cores into your design to enable debugging.

To maximize debugging closure, the Quartus II software allows you to use a combination of the debugging tools in tandem to fully exercise and analyze the logic under test. All of the tools described in Table 9–1 have basic analysis features built in; that is, all of the tools enable you to read back information collected from the design nodes that are connected to the debugging logic. Out of the set of debugging tools, the SignalTap II Logic Analyzer, the LAI, and the SignalProbe feature are general purpose

### Table 9–1. Available Tools in the In-System Verification Tools Suite (Part 2 of 2)

<table>
<thead>
<tr>
<th>Tool</th>
<th>Description</th>
<th>Typical Usage</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>SignalTap® II Logic Analyzer</strong></td>
<td>This logic analyzer uses FPGA resources to sample tests nodes and outputs the information to the Quartus II software for display and analysis.</td>
<td>You have spare on-chip memory and you want functional verification of your design running in hardware.</td>
</tr>
<tr>
<td><strong>SignalProbe</strong></td>
<td>This tool incrementally routes internal signals to I/O pins while preserving results from your last place-and-routed design.</td>
<td>You have spare I/O pins and you would like to check the operation of a small set of control pins using either an external logic analyzer or an oscilloscope.</td>
</tr>
<tr>
<td><strong>Logic Analyzer Interface (LAI)</strong></td>
<td>This tool multiplexes a larger set of signals to a smaller number of spare I/O pins. LAI allows you to select which signals are switched onto the I/O pins over a JTAG connection.</td>
<td>You have limited on-chip memory, and have a large set of internal data buses that you would like to verify using an external logic analyzer. Logic analyzer vendors, such as Tektronics and Agilent, provide integration with the tool to improve the usability of the tool.</td>
</tr>
<tr>
<td><strong>In-System Sources and Probes</strong></td>
<td>This tool provides an easy way to drive and sample logic values to and from internal nodes using the JTAG interface.</td>
<td>You want to prototype a front panel with virtual buttons for your FPGA design.</td>
</tr>
<tr>
<td><strong>In-System Memory Content Editor</strong></td>
<td>This tool displays and allows you to edit on-chip memory.</td>
<td>You would like to view and edit the contents of on-chip memory that is not connected to a Nios II processor. You can also use the tool when you do not want to have a Nios II debug core in your system.</td>
</tr>
<tr>
<td><strong>Virtual JTAG Interface</strong></td>
<td>This megafunction allows you to communicate with the JTAG interface so that you can develop your own custom applications.</td>
<td>You have custom signals in your design that you want to be able to communicate with.</td>
</tr>
</tbody>
</table>
debugging tools optimized for probing signals in your register transfer level (RTL) netlist. In-System Sources and Probes, the Virtual JTAG Interface, System Console, Transceiver Toolkit, and In-System Memory Content Editor, in addition to being able to read back data from the debugging breakpoints, allow you to input values into your design during runtime.

Taken together, the set of on-chip debugging tools form a debugging ecosystem. The set of tools can generate a stimulus to and solicit a response from the logic under test, providing a complete debugging solution (Figure 9–1).

**Figure 9–1. Quartus II Debugging Ecosystem (1)**

![Quartus II Debugging Ecosystem Diagram](image_url)

Note to Figure 9–1:

(1) The set of debugging tools offer end-to-end debugging coverage.

The tools in the toolchain offer different advantages and different trade-offs. To understand the selection criteria between the different tools, the following sections analyze the tools according to their typical applications.

The first section, “Analysis Tools for RTL Nodes”, compares the SignalTap II Logic Analyzer, SignalProbe, and the LAI. These three tools are logically grouped since they are intended for debugging nodes from your RTL netlist at system speed.

The second section, “Stimulus-Capable Tools” on page 9–8, compares the System Console and Transceiver Toolkit, In-System Memory Content Editor, Virtual JTAG Interface megafunction, and In-System Sources and Probes. These tools are logically grouped since they offer the ability to both read and write transactions through the JTAG port.
Analysis Tools for RTL Nodes

The SignalTap II Logic Analyzer, the SignalProbe feature, and the LAI are designed specifically for probing and debugging RTL signals at system speed. They are general-purpose analysis tools that enable you to tap and analyze any routable node from the FPGA or CPLD. If you have spare logic and memory resources, the SignalTap II Logic Analyzer is useful for providing fast functional verification of your design running on actual hardware.

Conversely, if logic and memory resources are tight and you require the large sample depths associated with external logic analyzers, both the LAI and the SignalProbe feature make it easy to view internal design signals using external equipment.

The most important selection criteria for these three tools are the available resources remaining on your device after implementing your design and the number of spare pins available. You should evaluate your preferred debugging option early on in the design planning process to ensure that your board, your Quartus II project, and your design are all set up to support the appropriate options. Planning early can reduce time spent during debugging and eliminate the necessary late changes to accommodate your preferred debugging methodologies. The following two sections provide information to assist you in choosing the appropriate tool by comparing the tools according to their resource usage and their pin consumption.

The SignalTap II Logic Analyzer is not supported on CPLDs, because there are no memory resources available on these devices.

Resource Usage

Any debugging tool that requires the use of a JTAG connection requires the SLD infrastructure logic mentioned earlier, for communication with the JTAG interface and arbitration between any instantiated debugging modules. This overhead logic uses around 200 logic elements (LEs), a small fraction of the resources available in any of the supported devices. The overhead logic is shared between all available debugging modules in your design. Both the SignalTap II Logic Analyzer and the LAI use a JTAG connection.

SignalProbe requires very few on-chip resources. Because it requires no JTAG connection, SignalProbe uses no logic or memory resources. SignalProbe uses only routing resources to route an internal signal to a debugging test point.

The LAI requires a small amount of logic to implement the multiplexing function between the signals under test, in addition to the SLD infrastructure logic. Because no data samples are stored on the chip, the LAI uses no memory resources.

The SignalTap II Logic Analyzer requires both logic and memory resources. The number of logic resources used depends on the number of signals tapped and the complexity of the trigger logic. However, the amount of logic resources that the SignalTap II Logic Analyzer uses is typically a small percentage of most designs. A baseline configuration consisting of the SLD arbitration logic and a single node with basic triggering logic contains approximately 300 to 400 Logic Elements (LEs). Each additional node you add to the baseline configuration adds about 11 LEs. Compared with logic resources, memory resources are a more important factor to consider for
your design. Memory usage can be significant and depends on how you configure your SignalTap II Logic Analyzer instance to capture data and the sample depth that your design requires for debugging. For the SignalTap II Logic Analyzer, there is the added benefit of requiring no external equipment, as all of the triggering logic and storage is on the chip.

Figure 9–2 shows a conceptual graph of the resource usage of the three analysis tools relative to each other.

Figure 9–2. Resource Usage per Debugging Tool (*)

Note to Figure 9–2:
(1) Though resource usage is highly dependent on the design, this graph provides a rough guideline for tool selection.

The resource estimation feature for the SignalTap II Logic Analyzer and the LAI allows you to quickly judge if enough on-chip resources are available before compiling the tool with your design. Figure 9–3 shows the resource estimation feature for the SignalTap II Logic Analyzer and the LAI.

Figure 9–3. Resource Estimator

Pin Usage

The ratio of the number of pins used to the number of signals tapped for the SignalProbe feature is one-to-one. Because this feature can consume free pins quickly, a typical application for this feature is routing control signals to spare pins for debugging.

The ratio of the number of pins used to the number of signals tapped for the LAI is many-to-one. It can map up to 256 signals to each debugging pin, depending on available routing resources. The control of the active signals that are mapped to the spare I/O pins is performed via the JTAG port. The LAI is ideal for routing data buses to a set of test pins for analysis.

Other than the JTAG test pins, the SignalTap II Logic Analyzer uses no additional pins. All data is buffered using on-chip memory and communicated to the SignalTap II Logic Analyzer GUI via the JTAG test port.
Usability Enhancements

The SignalTap II Logic Analyzer, the SignalProbe feature, and the LAI tools can be added to your existing design with minimal effects. With the node finder, you can find signals to route to a debugging module without making any changes to your HDL files. SignalProbe inserts signals directly from your post-fit database. The SignalTap II Logic Analyzer and LAI support inserting signals from both pre-synthesis and post-fit netlists. All three tools allow you to find and configure your debugging setup quickly. In addition, the Quartus II incremental compilation feature and the Quartus II incremental routing feature allow for a fast turnaround time for your programming file, increasing productivity and enabling fast debugging closure.

Both LAI and the SignalTap II Logic Analyzer support incremental compilation. With incremental compilation, you can add a SignalTap II Logic Analyzer instance or an LAI instance incrementally into your placed-and-routed design. This has the benefit of both preserving your timing and area optimizations from your existing design, and decreasing the overall compilation time when any changes are necessary during the debugging process. With incremental compilation, you can save up to 70% compile time of a full compilation.

SignalProbe uses the incremental routing feature. The incremental routing feature runs only the Fitter stage of the compilation. This also leaves your compiled design untouched, except for the newly routed node or nodes. With SignalProbe, you can save as much as 90% compile time of a full compilation.

As another productivity enhancement, all tools in the on-chip debugging tool set support scripting via the `quartus_stp` Tcl package. For the SignalTap II Logic Analyzer and the LAI, scripting enables user-defined automation for data collection while debugging in the lab.

In addition, the JTAG server allows you to debug a design that is running on a device attached to a PC in a remote location. This allows you to set up your hardware in the lab environment, download any new `.sof` files, and perform any analysis from your desktop.

Table 9–2 compares common debugging features between these tools and provides suggestions about which is the best tool to use for a given feature.

<table>
<thead>
<tr>
<th>Feature</th>
<th>SignalProbe</th>
<th>Logic Analyzer Interface (LAI)</th>
<th>SignalTap II Logic Analyzer</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Large Sample Depth</td>
<td>N/A</td>
<td>✓</td>
<td>-</td>
<td>An external logic analyzer used with the LAI has a bigger buffer to store more captured data than the SignalTap II Logic Analyzer. No data is captured or stored with SignalProbe.</td>
</tr>
<tr>
<td>Ease in Debugging Timing Issue</td>
<td>✓</td>
<td>✓</td>
<td>-</td>
<td>External equipment, such as oscilloscopes and Mixed Signal Oscilloscopes (MSOs), can be used with either LAI or SignalProbe. When used with the LAI to provide you with access to timing mode, you can debug combined streams of data.</td>
</tr>
</tbody>
</table>
## Table 9–2. Suggested On-Chip Debugging Tools for Common Debugging Features  (Part 2 of 2)  

<table>
<thead>
<tr>
<th>Feature</th>
<th>SignalProbe</th>
<th>Logic Analyzer Interface (LAI)</th>
<th>SignalTap II Logic Analyzer</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Minimal Effect on Logic Design</td>
<td>✓</td>
<td>✓ (2)</td>
<td>✓ (2)</td>
<td>The LAI adds minimal logic to a design, requiring fewer device resources. The SignalTap II Logic Analyzer has little effect on the design, because it is set as a separate design partition. SignalProbe incrementally routes nodes to pins, not affecting the design at all.</td>
</tr>
<tr>
<td>Short Compile and Recompile Time</td>
<td>✓</td>
<td>✓ (2)</td>
<td>✓ (2)</td>
<td>SignalProbe attaches incrementally routed signals to previously reserved pins, requiring very little recompilation time to make changes to source signal selections. The SignalTap II Logic Analyzer and the LAI can take advantage of incremental compilation to refit their own design partitions to decrease recompilation time.</td>
</tr>
<tr>
<td>Triggering Capability</td>
<td>N/A</td>
<td>N/A</td>
<td>✓</td>
<td>The SignalTap II Logic Analyzer offers triggering capabilities that are comparable to commercial logic analyzers.</td>
</tr>
<tr>
<td>I/O Usage</td>
<td>—</td>
<td>—</td>
<td>✓</td>
<td>No additional output pins are required with the SignalTap II Logic Analyzer. Both the LAI and SignalProbe require I/O pin assignments.</td>
</tr>
<tr>
<td>Acquisition Speed</td>
<td>N/A</td>
<td>—</td>
<td>✓</td>
<td>The SignalTap II Logic Analyzer can acquire data at speeds of over 200 MHz. The same acquisition speeds are obtainable with an external logic analyzer used with the LAI, but may be limited by signal integrity issues.</td>
</tr>
<tr>
<td>No JTAG Connection Required</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>An FPGA design with the SignalTap II Logic Analyzer or the LAI requires an active JTAG connection to a host running the Quartus II software. SignalProbe does not require a host for debugging purposes.</td>
</tr>
<tr>
<td>No External Equipment Required</td>
<td>—</td>
<td>—</td>
<td>✓</td>
<td>The SignalTap II Logic Analyzer logic is completely internal to the programmed FPGA device. No extra equipment is required other than a JTAG connection from a host running the Quartus II software or the stand-alone SignalTap II Logic Analyzer software. SignalProbe and the LAI require the use of external debugging equipment, such as multimeters, oscilloscopes, or logic analyzers.</td>
</tr>
</tbody>
</table>

### Notes to Table 9–2:

(1) ✓ indicates the recommended tools for the feature.  
— indicates that while the tool is available for that feature, that tool may not give the best results.  
N/A indicates that the feature is not applicable for the selected tool. 
(2) When used with incremental compilation.
**Stimulus-Capable Tools**

The In-System Memory Content Editor, the In-System Sources and Probes, and the Virtual JTAG interface each enable you to use the JTAG interface as a general-purpose communication port. Though all three tools can be used to achieve the same results, there are some considerations that make one tool easier to use in certain applications than others. In-System Sources and Probes is ideal for toggling control signals. The In-System Memory Content Editor is useful for inputting large sets of test data. Finally, the Virtual JTAG interface is well suited for more advanced users who want to develop their own customized JTAG solution.

System Console provides system-level debugging at a transaction level, such as with Avalon-MM slave or Avalon-ST interfaces. You can communicate to a chip through JTAG, PLI connectivity for simulation models, and TCP/IP protocols. System Console is a Tcl console that you use to communicate with hardware modules that you have instantiated into your design.

**In-System Sources and Probes**

In-System Sources and Probes is an easy way to access JTAG resources to both read and write to your design. You can start by instantiating a megafunction into your HDL code. The megafunction contains source ports and probe ports for driving values into and sampling values from the signals that are connected to the ports, respectively. Transaction details of the JTAG interface are abstracted away by the megafunction. During runtime, a GUI displays each source and probe port by instance and allows you to read from each probe port and drive to each source port. The GUI makes this tool ideal for toggling a set of control signals during the debugging process.

A good application of In-System Sources and Probes is to use the GUI as a replacement for the push buttons and LEDs used during the development phase of a project. Furthermore, In-System Sources and Probes supports a set of scripting commands for reading and writing using `quartus_stp`. When used with the Tk toolkit, you can build your own graphical interfaces. This feature is ideal for building a virtual front panel during the prototyping phase of the design.

**In-System Memory Content Editor**

The In-System Memory Content Editor allows you to quickly view and modify memory content either through a GUI interface or through Tcl scripting commands. The In-System Memory Content Editor works by turning single-port RAM blocks into dual-port RAM blocks. One port is connected to your clock domain and data signals, and the other port is connected to the JTAG clock and data signals for editing or viewing.

Because you can modify a large set of data easily, a useful application for the In-System Memory Content Editor is to generate test vectors for your design. For example, you can instantiate a free memory block, connect the output ports to the logic under test (using the same clock as your logic under test on the system side), and create the glue logic for the address generation and control of the memory. At runtime, you can modify the contents of the memory using either a script or the In-System Memory Content Editor GUI and perform a burst transaction of the data contents in the modified RAM block synchronous to the logic being tested.
Virtual JTAG Interface Megafunction

The Virtual JTAG Interface megafunction provides the finest level of granularity for manipulating the JTAG resource. This megafunction allows you to build your own JTAG scan chain by exposing all of the JTAG control signals and configuring your JTAG Instruction Registers (IRs) and JTAG Data Registers (DRs). During runtime, you control the IR/DR chain through a Tcl API, or with System Console. This feature is meant for users who have a thorough understanding of the JTAG interface and want precise control over the number and type of resources used.

System Console

System Console is a framework that you can launch from the Quartus II software to start services for performing various debugging tasks. System Console provides you with Tcl scripts and a GUI to access either the Qsys system integration tool or SOPC Builder modules to perform low-level hardware debugging of your design, as well as identify a module by its path, and open and close a connection to a Qsys or SOPC Builder module. You can access your design at a system level for purposes of loading, unloading, and transferring designs to multiple devices.

System Console also allows you to access commands that allow you to control how you generate test patterns, as well as verify the accuracy of data generated by test patterns. You can use JTAG debug commands in System Console to verify the functionality and signal integrity of your JTAG chain. You can test clock and reset signals.

You can use System Console to access programmable logic devices on your development board, as well as bring up a board and verify stages of setup. You can also access software running on a Nios II processor, as well as access modules that produce or consume a stream of bytes.

Transceiver Toolkit runs from the System Console framework, and allows you to run automatic tests of your transceiver links for debugging and optimizing your transceiver designs. You can use the Transceiver Toolkit GUI to set up channel links in your transceiver devices, and then automatically run EyeQ and Auto Sweep testing to view a graphical representation of your test data.

Conclusion

The Quartus II on-chip debugging tool suite allows you to reach debugging closure quickly by providing you a with set of powerful analysis tools and a set of tools that open up the JTAG port as a general purpose communication interface. The Quartus II software further broadens the scope of applications by giving you a comprehensive Tcl/Tk API. With the Tcl/Tk API, you can increase the level of automation for all of the analysis tools. You can also build virtual front panel applications quickly during the early prototyping phase.

In addition, all of the on-chip debugging tools have a tight integration with the rest of the productivity features within the Quartus II software. The incremental compilation and incremental routing features enable a fast turnaround time for programming file generation. The cross-probing feature allows you to find and identify nodes quickly. The SignalTap II Logic Analyzer, when used with the TimeQuest Timing Analyzer, is a best-in-class timing verification suite that allows fast functional and timing verification.
Document Revision History

Table 9–3 shows the revision history for this chapter.

Table 9–3. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Maintenance release.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Maintenance release. Changed to new document template.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Initial release</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
The System Console performs low-level hardware debugging of Qsys systems. You can use the System Console to debug systems that include IP cores instantiated in your Qsys system, as well as for initial bring-up of your printed circuit board, and for low-level testing. You access the System Console from Qsys, under the Tools menu.

This chapter contains the following sections:

- “System Console Overview” on page 10–1
- “Setting Up the System Console” on page 10–5
- “Using the System Console” on page 10–8
- “System Console Examples” on page 10–28
- “On-Board USB Blaster II Support” on page 10–39
- “Document Revision History” on page 10–40

**System Console Overview**

The System Console allows you to use many types of services. When you interact with the Tcl console in the System Console, you have general commands related to finding and accessing instances of those services. Each service type has functions that are unique to that service.

The System Console GUI (shown in Figure 10–1 on page 10–2) consists of a main window with four separate panes:

- The **System Explorer** pane allows you to view a hierarchy of the System Console virtual file system in your design, including connections, devices, designs, design instances, and scripts.
- The **Tools** pane allows you to launch tools such as the GDB Server Control Panel and Transceiver Toolkit.
- The **Tcl Console** allows you to run commands and Tcl scripts from within the System Console.
- The **Messages** pane displays messages generated by the System Console.
For more information about the System Console GUI, refer to About System Console in Quartus® II Help and the Altera Training page of the Altera website.

Using the System Explorer Pane

The System Explorer pane contains the following information about the system being debugged.

- The connections folder displays information about the debug cables which are visible to the System Console.
- The designs folder displays information about quartus project designs which have been loaded into the System Console.
- The devices folder contains information about each device connected to the System Console.
- The design_instances folder contains design instances.
- The scripts folder stores scripts for easy execution.

You will be doing most of your work within the devices folder. Within the devices folder is a folder for each device currently connected to the System Console. Each device folder contains a (link) folder and sometimes contains a (files) folder.
The (link) folder shows debug agents (and other hardware) which the System Console is able to access, arranged by connection type. The (files) folder contains information about the design files loaded from the Quartus II project for the device. Folders under the design_instances folder are linked to the appropriate (files) node.

**Figure 10–2. System Explorer Pane**

- Figure 10–2 shows that the EP4SGX230 folder contains a (link) folder. The link folder contains a JTAG folder. The JTAG folder contains folders that describe the debug pipes and agents that are connected to the EP4SGX230 device via a JTAG connection.

- The (files) folder contains information about the design files loaded from the Quartus II project for the device. Instances within the design_instances folder are linked to the corresponding files in the (files) folder.

- Folders that have a context menu available show a small context menu badge. Right-click these folders to view the context menu.

- Folders that have informational messages available display a small informational message badge. Hover over these folders to see the informational message.

- Folders corresponding to debug agents have a clock status badge. The badge displays a green clock signal if the clock is running or a red clock signal if it is not.

**Finding and Referring To Services**

The System Console uses a virtual file system to organize the available services, which is similar to the /dev location on Linux systems. Instances of services are referred to by their unique service path in the file system. You can retrieve service paths for a particular service with the command `get_service_paths <service-type>`.

System Console commands require service paths to identify the service instance you want to access. The paths for different components can change between runs of the tool and between versions. Use `get_service_paths` and similar commands to obtain service paths rather than hard coding them into your Tcl scripts.
Most System Console service instances are automatically discovered when you start the System Console. The System Console automatically scans for all JTAG and USB-based service instances and retrieves their service paths. Some other services, such as those connected by TCP/IP, are not automatically discovered. You can use the add_service Tcl command to inform the System Console about those services.

**Accessing Services**

After you have a service path to a particular service instance, you can access the service for use.

The `open_service` command tells the System Console to start using a particular service instance. The `open_service` command works on every service type. The `open_service` command claims a service instance for exclusive use.

The `open_service` command does not tell the System Console which part of a service you are interested in. As such, service instances that you open are not safe for shared use among multiple users.

The `claim_service` command tells the System Console to start accessing a particular portion of a service instance. For example, if you use the master service to access memory, then use `claim_service` to tell the System Console that you only want to access the address space between 0x0 and 0x1000. The System Console then allows other users to access other memory ranges and denies them access to your claimed memory range. The `claim_service` command returns a newly created service path that you can use to access your claimed resources.

Not all services support the `claim_service` command.

You can access a service after you open or claim it. When you finish accessing a service instance, use the `close_service` command to direct the System Console to make resources available.

**Applying Services**

The System Console provides extensive portfolios of services for various applications, such as real-time on-chip control and debugging, and system measurement. Examples of how to use these services are provided in this chapter. Table 10–1 lists example applications included with the System Console and associated services.

The System Console functions by running Tcl commands that are described in Table 10–3 through Table 10–17.

<table>
<thead>
<tr>
<th>Application</th>
<th>Services Used</th>
</tr>
</thead>
<tbody>
<tr>
<td>Board Bring-Up</td>
<td>device, jtag_debug, sld</td>
</tr>
<tr>
<td>Processor Debug</td>
<td>processor, elf, bytestream, master</td>
</tr>
<tr>
<td>Active retrieval of dynamic information</td>
<td>bytestream, master, issp</td>
</tr>
<tr>
<td>Query static design information</td>
<td>marker, design</td>
</tr>
<tr>
<td>System Monitoring</td>
<td>monitor, master, dashboard</td>
</tr>
</tbody>
</table>
Setting Up the System Console

Set up the System Console according to the hardware you have available in your system. You can access available debug IP on your system with the System Console. The debug IP allows you to access the running state of your system. The following sections discuss setting up and using debug IP in further detail.

“System Console Examples” on page 10–28 provides you with detailed examples of using the System Console with debug IP.

Download the design files for the example designs from the On-chip Debugging Design Examples page on the Altera website.

These design examples demonstrate how to add debug IP blocks to your design and how to connect them before you can use the host application.

Qsys

You can use the System Console to help you debug Qsys systems. The System Console communicates with debug IP in your system design. You can instantiate debug IP cores using Qsys or the MegaWizard Plug-In Manager.

For more information about the Qsys system integration tool, refer to System Design with Qsys in volume 1 of the Quartus II Handbook.

Table 10–2 describes some of the IP cores you can use with the System Console to debug your system. When connected to the System Console, these components enable you to send commands and receive data.

Table 10–1. System Console Example Applications

<table>
<thead>
<tr>
<th>Application</th>
<th>Services Used</th>
</tr>
</thead>
<tbody>
<tr>
<td>Transceiver Toolkit Direct PHY Control</td>
<td>transceiver_reconfig_analog, alt_xcvr_reconfig_dfe, alt_xcvr_reconfig_eye_viewer</td>
</tr>
<tr>
<td>Transceiver Toolkit System Level Control</td>
<td>transceiver_channel_rx, transceiver_channel_tx, transceiver_debug_link</td>
</tr>
</tbody>
</table>

Table 10–2. Qsys Components for Communication with the System Console (Part 1 of 2)

<table>
<thead>
<tr>
<th>Component Name</th>
<th>Component Interface Types for Debugging</th>
</tr>
</thead>
<tbody>
<tr>
<td>Nios® II processor with JTAG debug enabled</td>
<td>Components that include an Avalon® Memory-Mapped (Avalon-MM) slave interface. The JTAG debug module can also control the Nios II processor for debug functionality, including starting, stopping, and stepping the processor.</td>
</tr>
<tr>
<td>JTAG to Avalon master bridge</td>
<td>Components that include an Avalon-MM slave interface.</td>
</tr>
<tr>
<td>USB Debug Master</td>
<td>Provides the same functionality as JTAG to Avalon master bridge, but is faster.</td>
</tr>
<tr>
<td>Avalon Streaming (Avalon-ST) JTAG Interface</td>
<td>Components that include an Avalon-ST interface.</td>
</tr>
<tr>
<td>JTAG UART</td>
<td>The JTAG UART is an Avalon-MM slave device that can be used in conjunction with the System Console to send and receive bytestreams.</td>
</tr>
<tr>
<td>TCP/IP</td>
<td>For more information, refer to AN624: Debugging with System Console over TCP/IP.</td>
</tr>
</tbody>
</table>
For more information about Qsys components, refer to the following web pages and documents:

- Nios II Processor page of the Altera website
- SPI Slave/JTAG to Avalon Master Bride Cores chapter in the *Embedded Peripherals IP User Guide*
- Avalon Verification IP Suite User Guide
- Avalon-ST JTAG Interface Core chapter in the *Embedded Peripherals IP User Guide*
- Virtual JTAG (sld_virtual_jtag) Megafunction User Guide
- JTAG UART Core chapter in the *Embedded Peripherals IP User Guide*
- System Design with Qsys in volume 1 of the *Quartus II Handbook*
- About Qsys in Quartus II Help

---

**Table 10–2. Qsys Components for Communication with the System Console (Part 2 of 2)**

<table>
<thead>
<tr>
<th>Component Name</th>
<th>Component Interface Types for Debugging</th>
</tr>
</thead>
<tbody>
<tr>
<td>In-System Sources and Probes</td>
<td>Provides Tcl support for ISSP.</td>
</tr>
</tbody>
</table>

**Note to Table 10–2:**

(1) The System Console can also send and receive bytestreams from any system-level debugging (SLD) node in Qsys components provided by Altera, a custom component, or part of your Quartus II project; however, this approach requires detailed knowledge of the JTAG commands.
Figure 10–3 illustrates examples of interfaces of the components that the System Console can use.

The System Console provides many different types of services. Different modules can provide the same type of service. For example, both the Nios II processor and the JTAG to Avalon Bridge master provide the master service; consequently, you can use the master commands to access both of these modules.

If your system includes a Nios II/f core with a data cache, it may complicate the debugging process. If you suspect the Nios II/f core writes to memory from the data cache at nondeterministic intervals; thereby, overwriting data written by the System Console, you can disable the cache of the Nios II/f core while debugging.

Altera recommends that you include the following components in your system:

- On-chip memory
- JTAG UART
- System ID core

If your system includes a Nios II/f core with a data cache, it may complicate the debugging process. If you suspect the Nios II/f core writes to memory from the data cache at nondeterministic intervals; thereby, overwriting data written by the System Console, you can disable the cache of the Nios II/f core while debugging.
The following describes how to start the System Console from a Nios II command shell.

1. On the Windows Start menu, point to All Programs, then Altera, then Nios II EDS <version>, and then click Nios II <version> Command Shell.
2. To start the System Console, type the following command:

   `system-console`

You can customize your System Console environment by adding commands to the configuration file called `system_console_rc.tcl`. This file can be located in either of the following locations:

- `<quartus_install_dir>/sopc_builder/system_console_macros/system_console_rc.tcl`, known as the global configuration file, which affects all users of the system
- `<$HOME>/system_console/system_console_rc.tcl`, known as the user configuration file, which only affects the owner of that home directory

On startup, the System Console automatically runs any Tcl commands in these files. The commands in the global configuration file run first, followed by the commands in the user configuration file.

### Using the System Console

The Quartus II software expands the framework of the System Console by allowing you to start services for performing different types of tasks, as described in the following sections of this chapter. These sections provide Tcl scripting commands, arguments, and a brief description of the command functions.

To use the System Console commands, you must connect to a system with a programming cable and with the proper debugging IP.

### Interactive Help

Typing `help help` into the System Console lists all available commands. Typing `help <command name>` provides the syntax of individual commands. The System Console provides command completion if you type the beginning letters of a command and then press the Tab key.

The System Console interactive help commands only provide help for enabled services; consequently, typing `help help` does not display help for commands supplied by disabled plug-ins.

### Console Commands

The console commands enable testing. You can use console commands to identify a module by its path, and to open and close a connection to it. The path that identifies a module is the first argument to most of the System Console commands. To exercise a module, follow these steps:

1. Identify a module by specifying the path to it, using the `get_service_paths` command.
2. Open a connection to the module using the `open_service` or `claim_service` command. (`claim_service` returns a new path for use).

3. Run Tcl and System Console commands to use a debug module that you insert to test another module of interest.

4. Close a connection to a module using the `close_service` command.

Table 10–3 describes the syntax of the console commands.

**Table 10–3. Console Commands (Part 1 of 3)**

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>get_service_types</code></td>
<td>—</td>
<td>Returns a list of service types that the System Console manages. Examples of service types include master, bytestream, processor, sld, jtag_debug, device, and plugin.</td>
</tr>
<tr>
<td><code>get_service_paths</code></td>
<td><code>&lt;service_type&gt;</code></td>
<td>Returns a list of paths to nodes that implement the requested service type.</td>
</tr>
<tr>
<td><code>open_service</code></td>
<td><code>&lt;service_type&gt;</code> <code>&lt;service_path&gt;</code></td>
<td>Opens the specified service type at the specified path. An open_service command is equivalent to a claim_service command without claims specified.</td>
</tr>
<tr>
<td><code>claim_service</code></td>
<td><code>&lt;service-type&gt;</code> <code>&lt;service-path&gt;</code> <code>&lt;claim-group&gt;</code> <code>&lt;claims&gt;</code></td>
<td>Provides finer control of the portion of a service you want to use. Run <code>help claim_service</code> to get a <code>&lt;service-type&gt;</code> list. Then run <code>help claim_service &lt;service-type&gt;</code> to get specific help on that service.</td>
</tr>
<tr>
<td><code>close_service</code></td>
<td><code>&lt;service_type&gt;</code> <code>&lt;service_path&gt;</code></td>
<td>Closes the specified service type at the specified path.</td>
</tr>
<tr>
<td><code>is_service_open</code></td>
<td><code>&lt;service_type&gt;</code> <code>&lt;service_path&gt;</code></td>
<td>Returns 1 if the service type provided by the path is open, 0 if the service type is closed.</td>
</tr>
<tr>
<td><code>get_services_to_add</code></td>
<td>—</td>
<td>Returns a list of all services that are instantiable with the add_service command.</td>
</tr>
<tr>
<td><code>add_service</code></td>
<td><code>&lt;service-type&gt;</code> <code>&lt;instance-name&gt;</code> <code>&lt;optional-parameters&gt;</code></td>
<td>Adds a service of the specified service type with the given instance name. Run <code>get_services_to_add</code> to retrieve a list of instantiable services. This command returns the path where the service was added. Run <code>help add_service &lt;service-type&gt;</code> to get specific help about that service type, including any parameters that might be required for that service.</td>
</tr>
<tr>
<td>Command</td>
<td>Arguments</td>
<td>Function</td>
</tr>
<tr>
<td>-------------------------------</td>
<td>---------------------------------------------------------------------------</td>
<td>--------------------------------------------------------------------------</td>
</tr>
<tr>
<td>add_service dashboard</td>
<td>&lt;name&gt;</td>
<td>Creates a new GUI dashboard in System Console desktop.</td>
</tr>
<tr>
<td>add_service gdbserver</td>
<td>&lt;Processor Service&gt; &lt;port number&gt;</td>
<td>Instantiates a gdbserver.</td>
</tr>
<tr>
<td>add_service nios2dpx</td>
<td>&lt;path to debug channel&gt; &lt;isDualHeaded&gt; &lt;Base Av St Channel number&gt;</td>
<td>Instantiates a Nios II DPX debug driver.</td>
</tr>
<tr>
<td>add_service pli_bytestream</td>
<td>&lt;instance_name&gt; &lt;port_number&gt;</td>
<td>Instantiates a PLI Bytestream service.</td>
</tr>
<tr>
<td>add_service pli_master</td>
<td>&lt;instance_name&gt; &lt;port_number&gt;</td>
<td>Instantiates a PLI Master service.</td>
</tr>
<tr>
<td>add_service pli_packets</td>
<td>&lt;instance_name&gt; &lt;port_number&gt;</td>
<td>Instantiates a PLI packet stream service.</td>
</tr>
<tr>
<td>add_service tcp</td>
<td>&lt;instance_name&gt; &lt;ip_addr&gt; &lt;port number&gt;</td>
<td>Instantiates a tcp service.</td>
</tr>
<tr>
<td>add_service transceiver_channel_rx</td>
<td>&lt;data_pattern_checker path&gt; &lt;transceiver path&gt; &lt;transceiver channel address&gt; &lt;reconfig path&gt; &lt;reconfig channel address&gt;</td>
<td>Instantiates a Transceiver Toolkit receiver channel.</td>
</tr>
<tr>
<td>add_service transceiver_channel_tx</td>
<td>&lt;data_pattern_generator path&gt; &lt;transceiver path&gt; &lt;transceiver channel address&gt; &lt;reconfig path&gt; &lt;reconfig channel address&gt;</td>
<td>Instantiates a Transceiver Toolkit transceiver channel.</td>
</tr>
<tr>
<td>add_service transceiver_debug_link</td>
<td>&lt;transceiver_channel_tx path&gt; &lt;transceiver_channel_rx path&gt;</td>
<td>Instantiates a Transceiver Toolkit debug link.</td>
</tr>
<tr>
<td>get_version</td>
<td></td>
<td>Returns the current System Console version and build number.</td>
</tr>
<tr>
<td>add_help</td>
<td>&lt;command&gt; &lt;help-text&gt;</td>
<td>Adds help text for a given command. Use this when you write a Tcl script procedure (proc) and then want to provide help for others to use the script.</td>
</tr>
<tr>
<td>get_claimed_services</td>
<td>&lt;claim-group&gt;</td>
<td>For the given claim group, returns a list of services claimed. The returned list consists of pairs of paths and service types. Each pair is one claimed service.</td>
</tr>
</tbody>
</table>
Plugins

Plugins allow you to customize how you use the System Console services and are enabled by default. Table 10–4 lists Plugin commands.

Table 10–4. Plugin Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>plugin_enable</td>
<td>&lt;plugin-path&gt;</td>
<td>Enables the plugin specified by the path. After a plugin is enabled, you can retrieve the &lt;service-path&gt; and &lt;service_type_name&gt; for additional services using the get_service_paths command.</td>
</tr>
<tr>
<td>plugin_disable</td>
<td>&lt;plugin-path&gt;</td>
<td>Disables the plugin specified by the path.</td>
</tr>
<tr>
<td>is_plugin_enabled</td>
<td>&lt;plugin-path&gt;</td>
<td>Returns a non-zero value when the plugin at the specified path is enabled.</td>
</tr>
</tbody>
</table>

Design Service Commands

Design Service commands allow you to use the System Console services to load and work with your design at a system level. Table 10–5 lists Design Service commands.

Table 10–5. Design Service Commands (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>design_load (1)</td>
<td>&lt;quartus-project-path&gt;, &lt;sof-file-path&gt;, or &lt;qpf-file-path&gt;</td>
<td>Loads a model of a Quartus II design into the System Console. Returns the design path. For example, if your Quartus II Project File (.qpf) file is in c:/projects/loopback, type the following command: design_load {c:\projects\loopback}\</td>
</tr>
<tr>
<td>design_instantiate</td>
<td>&lt;design-path&gt; &lt;instance-name&gt;</td>
<td>Instantiates a Quartus II design, which creates an instance. The instance name is optional. Returns the instance path.</td>
</tr>
</tbody>
</table>

Table 10–3. Console Commands (Part 3 of 3)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>refresh_connections</td>
<td>—</td>
<td>Scans for available hardware and updates the available service paths if there have been any changes.</td>
</tr>
<tr>
<td>send_message</td>
<td>&lt;level&gt; &lt;message&gt;</td>
<td>Sends a message of the given level to the message window. Available levels are info, warning, error, and debug.</td>
</tr>
</tbody>
</table>
Using the System Console

Programmable Logic Device (PLD) Commands

The PLD commands provide access to programmable logic devices on your board. Before you use these commands, identify the path to the programmable logic device on your board using the `get_service_paths` command described in Table 10–3.

Note to Table 10–5:
(1) Turn on the **Auto Usercode** option to have the System Console automatically instantiate and link designs after they have been loaded.

For more information about **Auto Usercode**, refer to the **General Page (Device and Pin Options Dialog Box)** in Quartus II Help.

### Table 10–5. Design Service Commands  (Part 2 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>design_link</td>
<td><code>&lt;design-instance-path&gt;</code></td>
<td>Creates a design instance if necessary and then links a Quartus II logical design with a physical device. For example, you can link a Quartus II design called <code>2c35_quartus_design</code> to a 2c35 device. After you create this link, the System Console creates the appropriate correspondences between the logical and physical submodules of the Quartus II project. Example 10–9 on page 10–38 shows a transcript illustrating the <code>design_load</code> and <code>design_link</code> commands. Note that the System Console does not verify that the link is valid; if you create an incorrect link, the System Console does not report an error.</td>
</tr>
<tr>
<td>design_extract_debug_files</td>
<td><code>&lt;design-path&gt;</code> <code>&lt;zip-file-name&gt;</code></td>
<td>Extract debug files from a SRAM Object File (.sof) to a zip file which can be emailed to Altera Support for analysis.</td>
</tr>
<tr>
<td>design_extract_dotty</td>
<td><code>&lt;design-path&gt;</code> <code>&lt;dot-file-name&gt;</code></td>
<td>Convert the project into a dotty file which shows what the System Console has extracted from the project. Use the Graphviz tools to display the file.</td>
</tr>
<tr>
<td>design_get_warnings</td>
<td><code>&lt;design-path&gt;</code></td>
<td>Gets the list of warnings for this design. If the design loads correctly, then an empty list returns.</td>
</tr>
<tr>
<td>design_update_debug_files</td>
<td><code>&lt;design-path&gt;</code> <code>&lt;list-of-files-to-update&gt;</code></td>
<td>Add files, or update files within the debug files section of the .sof.</td>
</tr>
</tbody>
</table>
Table 10–6 describes the PLD commands.

### Table 10–6. PLD Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>device_download_sof</td>
<td>&lt;service_path&gt;, &lt;sof-file-path&gt;</td>
<td>Loads the specified .sof file to the device specified by the path.</td>
</tr>
<tr>
<td>device_get_connections</td>
<td>&lt;service_path&gt;</td>
<td>Returns all connections which go to the device at the specified path.</td>
</tr>
<tr>
<td>device_get_design</td>
<td>&lt;device_path&gt;</td>
<td>Returns the design this device is currently linked to.</td>
</tr>
</tbody>
</table>

### Monitor Commands

You can use the Monitor commands to read many Avalon-MM slave memory locations at a regular interval. For example, if you want to perform 100 reads per second, every second, you get much better performance using the monitor service than if you call 100 separate `master_read_memory` commands every second. This is the primary difference between the monitor service and the master service.

Table 10–7 lists the commands usually called from the main program when setting up the monitor. Table 10–8 lists the commands called from within the monitor callback.

To use Monitor commands, you must create a new monitor, set its callback and interval, add ranges, and then set it to enabled. From within the callback you must use appropriate `read_data` commands to extract the data. Note that under heavy load, one or more monitor callbacks might be skipped.

### Table 10–7. Main Monitoring Commands (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>monitor_add_range</td>
<td>&lt;service-path&gt;, &lt;target-path&gt;, &lt;address&gt;, &lt;size&gt;</td>
<td>Adds a contiguous memory address into the monitored memory list.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&lt;service path&gt; is the value returned when you opened the service.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>&lt;target-path&gt; argument is the name of a master service to read. The address is within the address space of this service. &lt;target-path&gt; is returned from <code>[lindex [get_service_paths master] n]</code> where <code>n</code> is the number of the master service. &lt;address&gt; and &lt;size&gt; are relative to the master service.</td>
</tr>
<tr>
<td>monitor_set_callback</td>
<td>&lt;service-path&gt;, &lt;Tcl-expression&gt;</td>
<td>Defines a Tcl expression in a single string that will be evaluated after all the memories monitored by this service are read. Typically, this expression should be specified as a Tcl procedure call with necessary argument passed in.</td>
</tr>
<tr>
<td>monitor_set_interval</td>
<td>&lt;service-path&gt;, &lt;interval&gt;</td>
<td>Specifies the frequency of the polling action by specifying the interval between two memory reads. The actual polling frequency varies depending on the system activity. The monitor service will try to keep it as close to this specification as possible.</td>
</tr>
</tbody>
</table>
Under normal load, the monitor service reads the data after each interval and then calls the callback. If the value you read is timing sensitive, the `monitor_get_read_interval` command can be used to read the exact time between the intervals at which the data was read.

Under heavy load, or with a callback that takes a long time to execute, the monitor service skips some callbacks. If the registers you read do not have side effects (for example, they read the total number of events since reset), skipping callbacks has no effect on your code. The `monitor_read_data` command and `monitor_get_read_interval` command are adequate for this scenario.
If the registers you read have side effects (for example, they return the number of events since the last read), you must have access to the data that was read, but for which the callback was skipped. The `monitor_read_all_data` and `monitor_get_all_read_intervals` commands provide access to this data.

**Trace Commands**

The System Console trace system allows you to identify events of interest in the hardware and send details of those events to the host system. Table 10–9 lists the commands available in the System Console trace system.

To use the trace commands listed in this section and capture the data these commands produce, instantiate an Altera Trace System and suitable monitors (for example, an Avalon-ST Video Monitor) in your system. Open the trace system with the `claim_service trace <service-path> <library-name>` command. (You can determine the path to your instantiated trace system by running `get_service_paths trace`.)

The `claim_service trace` command returns a new service path that represents the instantiated and opened trace system. Use the returned service path in the commands below to access your hardware and retrieve events of interest.

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>trace_get_monitors</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns a Tcl list of monitor paths to the monitors (for example, an Avalon-ST Video Monitor) that are currently connected to the trace system specified by <code>&lt;service-path&gt;</code>.</td>
</tr>
<tr>
<td><code>trace_get_monitor_info</code></td>
<td><code>&lt;service-path&gt;</code> <code>&lt;monitor-path&gt;</code></td>
<td>Returns all the information that is known about the specified monitor on this trace system. The values returned are a serialized array of key / value pairs—they can be converted to an array with the <code>array set</code> command.</td>
</tr>
<tr>
<td><code>trace_read_monitor</code></td>
<td><code>&lt;service-path&gt;</code> <code>&lt;monitor-path&gt;</code> <code>&lt;index&gt;</code></td>
<td>Reads a configuration register from the monitor specified. The register at index 0 typically identifies the type of monitor, registers at higher indexes (which are word addresses), and provides more information about the monitor.</td>
</tr>
<tr>
<td><code>trace_write_monitor</code></td>
<td><code>&lt;service-path&gt;</code> <code>&lt;monitor-path&gt;</code> <code>&lt;index&gt;</code> <code>&lt;value&gt;</code></td>
<td>Writes a configuration register in the specified monitor. The register at word address 4 typically holds enable bits for the monitor, the meanings of registers at higher addresses are dependent on the type of monitor being accessed.</td>
</tr>
<tr>
<td><code>trace_set_max_db_size</code></td>
<td><code>&lt;service-path&gt;</code> <code>&lt;size&gt;</code></td>
<td>Sets the maximum database size (in bytes) that the driver stores in memory. When the database becomes larger than the specified size, the oldest events are dropped to prevent the application from running out of memory.</td>
</tr>
<tr>
<td><code>trace_get_max_db_size</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the current maximum database size (in bytes) for the trace system.</td>
</tr>
<tr>
<td><code>trace_get_db_size</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the approximate current size (in bytes) for the database.</td>
</tr>
</tbody>
</table>
### Board Bring-Up Commands

The board bring-up commands allow you to debug your design while the design is running in an FPGA. These commands are presented in the order that you would use them during board bring-up.

“Board Bring-Up with the System Console” on page 10–28 provides you with detailed examples of using the System Console for board bring-up.

The System Console is intended for debugging the basic hardware functionality of your Nios II processor, including its memories and pinout. If you are writing device drivers, you may want to use the System Console and the Nios II software build tools together to debug your code.

For more information about the hardware functionality and software debugging, refer to *Nios II Software Build Tools Reference* in the *Nios II Software Developer’s Handbook*.
### JTAG Debug Commands

You can use JTAG debug commands to verify the functionality and signal integrity of your JTAG chain. Your JTAG chain must function correctly to debug the rest of your system. To verify signal integrity of your JTAG chain, Altera recommends that you provide an extensive list of byte values. Table 10–10 lists these commands.

**Table 10–10. JTAG Commands**

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>jtag_debug_loop</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Loops the specified list of bytes through a loopback of <code>tdi</code> and <code>tdo</code> of a system-level debug (SLD) node. Returns the list of byte values in the order that they were received. Blocks until all bytes are received. Byte values are given with the 0x (hexadecimal) prefix and delineated by spaces.</td>
</tr>
<tr>
<td>jtag_debug_reset_system</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Issues a reset request to the specified service. Connectivity within your device determines which part of the system is reset.</td>
</tr>
</tbody>
</table>

### Clock and Reset Signal Commands

The next stage of board bring-up tests the clock and reset signals. Table 10–11 lists the three commands to verify these signals. Use these commands to verify that your clock is toggling and that the reset signal has the expected value.

**Table 10–11. Clock and Reset Commands**

<table>
<thead>
<tr>
<th>Command</th>
<th>Argument</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>jtag_debug_sample_clock</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the value of the clock signal of the system clock that drives the module's system interface. The clock value is sampled asynchronously; consequently, you may need to sample the clock several times to guarantee that it is toggling.</td>
</tr>
<tr>
<td>jtag_debug_sample_reset</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the value of the <code>reset_n</code> signal of the Avalon-ST JTAG Interface core. If <code>reset_n</code> is low (asserted), the value is 0 and if <code>reset_n</code> is high (deasserted), the value is 1.</td>
</tr>
<tr>
<td>jtag_debug_sense_clock</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the result of a sticky bit that monitors for system clock activity. If the clock has toggled since the last execution of this command, the bit is 1. Returns true if the bit has ever toggled and otherwise returns false. The sticky bit is reset to 0 on read.</td>
</tr>
</tbody>
</table>

### Avalon-MM Commands

The master service provides commands that allow you to access memory-mapped slaves via a suitable Avalon-MM master, which can be controlled by the host. You can use the commands listed in Table 10–12 to read and write memory with a master service.

Master services are provided either by System Console master components such as the JTAG Avalon Master or the USB Debug Master, by PLI or TCP masters, and by some processors which typically must be paused before they can be used for general purpose memory access.
SLD Commands

You can use the SLD commands to shift values into the instruction and data registers of SLD nodes and read the previous value. Table 10–12 lists these commands.

Claim Values for Claim Services

Each master claim consists of three parts: a base address, a size, and an access mode. The base address and size can be specified in decimal or hexadecimal (with preceding 0x). Valid access modes include the following:

- RO or READONLY gives read access to the specified addresses.
- RW or READWRITE gives read and write access to the specified addresses.
- EXC or EXCLUSIVE gives read and write access to the specified addresses.

If multiple RO and/or RW addresses have overlapping address ranges they are allowed to open at the same time. EXC and EXCLUSIVE claims do not allow other claims for the same memory range.

Table 10–12. Module Commands (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>master_write_memory</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;list_of_byte_values&gt;</td>
<td>Writes the list of byte values, starting at the specified base address.</td>
</tr>
<tr>
<td>master_write_8</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;list_of_byte_values&gt;</td>
<td>Writes the list of byte values, starting at the specified base address, using 8-bit accesses.</td>
</tr>
<tr>
<td>master_write_16</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;list_of_16_bit_words&gt;</td>
<td>Writes the list of 16-bit values, starting at the specified base address, using 16-bit accesses.</td>
</tr>
<tr>
<td>master_write_from_file</td>
<td>&lt;service-path&gt; &lt;file-name&gt; &lt;address&gt;</td>
<td>Writes the entire contents of the file through the master, starting at the specified address. The file is treated as a binary file containing a stream of bytes.</td>
</tr>
<tr>
<td>master_write_32</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;list_of_32_bit_words&gt;</td>
<td>Writes the list of 32-bit values, starting at the specified base address, using 32-bit accesses.</td>
</tr>
<tr>
<td>master_read_memory</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;size_in_bytes&gt;</td>
<td>Returns a list of &lt;size&gt; bytes. Read from memory starts at the specified base address.</td>
</tr>
<tr>
<td>master_read_8</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;size_in_bytes&gt;</td>
<td>Returns a list of &lt;size&gt; bytes. Read from memory starts at the specified base address, using 8-bit accesses.</td>
</tr>
<tr>
<td>master_read_16</td>
<td>&lt;service-path&gt; &lt;base-address&gt; &lt;size_in_multiples_of_16_bits&gt;</td>
<td>Returns a list of &lt;size&gt; 16-bit values. Read from memory starts at the specified base address, using 16-bit accesses.</td>
</tr>
</tbody>
</table>
## Processor Commands

These commands allow you to start, stop, and step through software running on a Nios II processor. The commands also allow you to read and write the registers of the processor. Table 10–13 lists the commands.

### Table 10–13. Processor Commands (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>elf_download</td>
<td><code>&lt;processor-service-path&gt;</code></td>
<td>Downloads the given Executable and Linking Format</td>
</tr>
<tr>
<td></td>
<td><code>&lt;master-service-path&gt;</code></td>
<td>File (.elf) to memory using the specified master</td>
</tr>
<tr>
<td></td>
<td><code>&lt;elf-file-path&gt;</code></td>
<td>service. Sets the processor’s program counter to the</td>
</tr>
<tr>
<td></td>
<td></td>
<td>.elf entry point.</td>
</tr>
<tr>
<td>processor_in_debug_mode</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns a non-zero value if the processor is in debug mode.</td>
</tr>
<tr>
<td>processor_reset</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Resets the processor and places it in debug mode.</td>
</tr>
<tr>
<td>processor_run</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Puts the processor into run mode.</td>
</tr>
<tr>
<td>processor_stop</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Puts the processor into stop mode.</td>
</tr>
<tr>
<td>processor_step</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Executes one assembly instruction.</td>
</tr>
</tbody>
</table>
### Bytestream Commands

These commands provide access to modules that produce or consume a stream of bytes. You can use the bytestream service to communicate directly to IP that provides bytestream interfaces, such as the Altera JTAG UART. Table 10–14 lists the commands.

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytestream_send</td>
<td>&lt;service-path&gt;</td>
<td>Sends the list of bytes to the specified bytestream service. Values argument is the list of bytes to send.</td>
</tr>
<tr>
<td>bytestream_receive</td>
<td>&lt;service-path&gt;</td>
<td>Returns a list of bytes currently available in the specified services receive queue, up to the specified limit. Length argument is the maximum number of bytes to receive.</td>
</tr>
</tbody>
</table>

### Marker Commands

These commands provide debugging information. Table 10–15 lists the commands.

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>marker_get_assignments</td>
<td>—</td>
<td>Returns all debug assignments in a key value ordered list.</td>
</tr>
<tr>
<td>marker_get_info</td>
<td>—</td>
<td>Returns all debug information in a key value ordered list.</td>
</tr>
<tr>
<td>marker_get_type</td>
<td>—</td>
<td>Returns the debug type of the marker.</td>
</tr>
</tbody>
</table>
In-System Sources and Probes Commands

You can use the In-System Sources and Probes (ISSP) commands to read source and probe data. Table 10–16 lists the commands. You use these commands with the In-System Sources and Probes that you insert into your project from the Quartus II software main menu.

Before you use the ISSP service, you must open it with commands similar to the following:

```
set issp [lindex [ get_service_paths issp ] 0]
set chan [claim_service issp $issp <library-name>]
set v [issp_read_probe_data $chan]
```

The valid values for probe claims include read_only, normal, and exclusive.

### Table 10–16. In-System Sources and Probes Tcl Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>issp_get_instance_info</td>
<td>&lt;service-path&gt;</td>
<td>Returns a list of the configurations of the In-System Sources and Probes instance, including: instance_index instance_name source_width probe_width</td>
</tr>
<tr>
<td>issp_read_probe_data</td>
<td>&lt;service-path&gt;</td>
<td>Retrieves the current value of the probes. A hex string is returned representing the probe port value.</td>
</tr>
<tr>
<td>issp_read_source_data</td>
<td>&lt;service-path&gt;</td>
<td>Retrieves the current value of the sources. A hex string is returned representing the source port value.</td>
</tr>
<tr>
<td>issp_write_source_data</td>
<td>&lt;service-path&gt; &lt;source-value&gt;</td>
<td>Sets values for the sources. The value can be either a hex string or a decimal value supported by System Console Tcl interpreter.</td>
</tr>
</tbody>
</table>

Dashboard Commands

The System Console dashboard allows you to create graphical tools that seamlessly integrate into the System Console. This section describes how to build your own dashboard with Tcl commands and the properties that you can assign to the widgets on your dashboard. The dashboard allows you to create tools that interact with live instances of an IP core on your device. Table 10–17 lists the dashboard Tcl commands available from the System Console.

Example 10–1 shows a Tcl command to create a dashboard. Run the Tcl command to return a path. You can then use the path on the commands listed in Table 10–17.
Example 10–1. Example of Creating a Dashboard

```
add_service dashboard my_new_dashboard "This is a New Dashboard" "Tools/My New Dashboard"
```

Table 10–17. Dashboard Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>dashboard_add</td>
<td>&lt;service-path&gt; &lt;id&gt; &lt;type&gt; &lt;group id&gt;</td>
<td>Allows you to add a specified widget to your GUI dashboard.</td>
</tr>
<tr>
<td>dashboard_remove</td>
<td>&lt;service-path&gt; &lt;id&gt;</td>
<td>Allows you to remove a specified widget from your GUI dashboard.</td>
</tr>
<tr>
<td>dashboard_set_property</td>
<td>&lt;service-path&gt; &lt;id&gt; &lt;property&gt; &lt;value&gt;</td>
<td>Allows you to set the specified properties of the specified widget that has been added to your GUI dashboard.</td>
</tr>
<tr>
<td>dashboard_get_property</td>
<td>&lt;service-path&gt; &lt;id&gt; &lt;type&gt;</td>
<td>Allows you to determine the existing properties of a widget added to your GUI dashboard.</td>
</tr>
<tr>
<td>dashboard_get_types</td>
<td></td>
<td>Returns a list of all possible widgets that you can add to your GUI dashboard.</td>
</tr>
<tr>
<td>dashboard_get_properties</td>
<td>&lt;widget type&gt;</td>
<td>Returns a list of all possible properties of the specified widgets in your GUI dashboard.</td>
</tr>
</tbody>
</table>

Specifying Widgets
You can specify the widgets that you add to your dashboard. Table 10–18 lists the widgets.

Note that `dashboard_add` performs a case-sensitive match against the widget type name.

Table 10–18. Dashboard Widgets

<table>
<thead>
<tr>
<th>Widget</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>group</td>
<td>Allows you to add a collection of widgets and control the general layout of the widgets.</td>
</tr>
<tr>
<td>button</td>
<td>Allows you to add a button.</td>
</tr>
<tr>
<td>tabbedGroup</td>
<td>Allows you to group tabs together.</td>
</tr>
<tr>
<td>fileChooserButton</td>
<td>Allows you to define button actions.</td>
</tr>
<tr>
<td>label</td>
<td>Allows you to add a text string.</td>
</tr>
<tr>
<td>text</td>
<td>Allows you to display text.</td>
</tr>
<tr>
<td>textField</td>
<td>Allows you to add a text field.</td>
</tr>
</tbody>
</table>
Example 10–2 is a Tcl script to instantiate a widget. In this example, the Tcl command adds a label to the dashboard. The first argument is the path to the dashboard. This path is returned by the `add_service` command. The next argument is the ID you assign to the widget. The ID must be unique within the dashboard. You use this ID later on to refer to the widget.

Following that argument is the type of widget you are adding, which in this example is a label. The last argument to this command is the group where you want to put this widget. In this example, a special keyword `self` is used. `Self` refers to the dashboard itself, the primary group. You can then add a group to `self`, which allows you to add other widgets to this group by using the ID of the new group, rather than using the ID of the `self` group.

Example 10–2. Example of Instantiating a Widget

```tcl
dashboard_add $dash mylabel label self
```

**Customizing Widgets**

You can change widget properties at any time. The `dashboard_set_property` command allows you to interact with the widgets you instantiate. This functionality is most useful when you change part of the execution of a callback. Example 10–3 shows how to change the text in a label.

In Example 10–3, the first argument is the path to the dashboard. Next is the unique ID of the widget, which then allows you to target an arbitrary widget. Following that is the name of the property. Each type of widget has a defined set of properties, discussed later. You can change the properties. In this example, `mylabel` is of the type `label`, and the example shows how to set its `text` property. The last argument is the value that the property takes when the command is executed.

Example 10–3. Example of Customizing a Widget

```tcl
dashboard_set_property $dash mylabel text "Hello World!"
```
Assigning Dashboard Widget Properties

In Table 10–19 through Table 10–31, the various properties are listed that you can apply to the widgets on your dashboard.

Table 10–19. Properties Common to All Widgets

<table>
<thead>
<tr>
<th>Property</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>enabled</td>
<td>Enables or disables the widget.</td>
</tr>
<tr>
<td>expandable</td>
<td>Allows the widget to be expanded.</td>
</tr>
<tr>
<td>expandableX</td>
<td>Allows the widget to be resized horizontally if there’s space available in the cell where it resides.</td>
</tr>
<tr>
<td>expandableY</td>
<td>Allows the widget to be resized vertically if there’s space available in the cell where it resides.</td>
</tr>
<tr>
<td>maxHeight</td>
<td>If the widget’s expandableY is set, this is the maximum height in pixels that the widget can take.</td>
</tr>
<tr>
<td>minHeight</td>
<td>If the widget’s expandableY is set, this is the minimum height in pixels that the widget can take.</td>
</tr>
<tr>
<td>maxWidth</td>
<td>If the widget’s expandableX is set, this is the maximum width in pixels that the widget can take.</td>
</tr>
<tr>
<td>minWidth</td>
<td>If the widget’s expandableX is set, this is the minimum width in pixels that the widget can take.</td>
</tr>
<tr>
<td>preferredHeight</td>
<td>The height of the widget if expandableY is not set.</td>
</tr>
<tr>
<td>preferredWidth</td>
<td>The width of the widget if expandableX is not set.</td>
</tr>
<tr>
<td>toolTip</td>
<td>A tooltip string that appears once the mouse hovers above the widget.</td>
</tr>
<tr>
<td>selected</td>
<td>The value of the checkbox, whether it is selected or not.</td>
</tr>
<tr>
<td>visible</td>
<td>Allows the widget to be displayed.</td>
</tr>
<tr>
<td>onChange</td>
<td>Allows for registering a callback function to be called when the value of the box changes.</td>
</tr>
<tr>
<td>options</td>
<td>Allows you to list available options.</td>
</tr>
</tbody>
</table>

Table 10–20. button Properties

<table>
<thead>
<tr>
<th>Property</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>onClick</td>
<td>A Tcl command to run, usually a proc, every time the button is clicked.</td>
</tr>
<tr>
<td>text</td>
<td>The text on the button.</td>
</tr>
</tbody>
</table>

Table 10–21. file Chooser Button Properties

<table>
<thead>
<tr>
<th>Property</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>text</td>
<td>The text on the button.</td>
</tr>
<tr>
<td>onChoose</td>
<td>A Tcl command to run, usually a proc, every time the button is clicked.</td>
</tr>
<tr>
<td>title</td>
<td>The text of file chooser dialog box title.</td>
</tr>
<tr>
<td>chooserButtonText</td>
<td>The text of file chooser dialog box approval button. By default, it is “Open”.</td>
</tr>
<tr>
<td>filter</td>
<td>The file filter based on extension. Only one extension is supported. By default, all file names are allowed. The filter is specified as [list filter_description file_extension], for example [list “Text Document (.txt)” “txt”].</td>
</tr>
</tbody>
</table>
### Table 10–21. fileChooserButton Properties

<table>
<thead>
<tr>
<th>Property</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>mode</td>
<td>Specifies what kind of files or directories can be selected. “files_only”, by default. Possible options are “files_only” and “directories_only”.</td>
</tr>
<tr>
<td>multiSelectionEnabled</td>
<td>Controls whether multiple files can be selected. False, by default.</td>
</tr>
<tr>
<td>paths</td>
<td>Returns a list of file paths selected in the file chooser dialog box. This property is read-only. It is most useful when used within the onclick script or a procedure when the result is freshly updated after the dialog box closes.</td>
</tr>
</tbody>
</table>

### Table 10–22. dial Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>max</td>
<td>The maximum value that the dial can show.</td>
</tr>
<tr>
<td>min</td>
<td>The minimum value that the dial can show.</td>
</tr>
<tr>
<td>tickSize</td>
<td>The space between the different tick marks of the dial.</td>
</tr>
<tr>
<td>title</td>
<td>The title of the dial.</td>
</tr>
<tr>
<td>value</td>
<td>The value that the dial’s needle should mark. It must be between min and max.</td>
</tr>
</tbody>
</table>

### Table 10–23. group Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>itemsPerRow</td>
<td>The number of widgets the group can position in one row, from left to right, before moving to the next row.</td>
</tr>
<tr>
<td>title</td>
<td>The title of the group. Groups with a title can have a border around them, and setting an empty title removes the border.</td>
</tr>
</tbody>
</table>

### Table 10–24. label Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>text</td>
<td>The text to show in the label.</td>
</tr>
</tbody>
</table>

### Table 10–25. led Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>color</td>
<td>The color of the LED. The options are: red_off, red, yellow_off, yellow, green_off, green, blue_off, blue, and black.</td>
</tr>
<tr>
<td>text</td>
<td>The text to show next to the LED.</td>
</tr>
</tbody>
</table>

### Table 10–26. text Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>editable</td>
<td>Controls whether the text box is editable.</td>
</tr>
<tr>
<td>htmlCapable</td>
<td>Controls whether the text box can format HTML.</td>
</tr>
<tr>
<td>text</td>
<td>The text to show in the text box.</td>
</tr>
</tbody>
</table>
Table 10–27. timeChart Properties

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>labelX</td>
<td>The label for the X axis.</td>
</tr>
<tr>
<td>labelY</td>
<td>The label for the Y axis.</td>
</tr>
<tr>
<td>latest</td>
<td>The latest value in the series.</td>
</tr>
<tr>
<td>maximumItemCount</td>
<td>The number of sample points to display in the historic record.</td>
</tr>
<tr>
<td>title</td>
<td>The title of the chart.</td>
</tr>
</tbody>
</table>

Table 10–28. table Properties (Part 1 of 2)

<table>
<thead>
<tr>
<th>Properties</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Table-wide Properties</strong></td>
<td></td>
</tr>
<tr>
<td>columnCount</td>
<td>The number of columns (Mandatory) (0, by default).</td>
</tr>
<tr>
<td>rowCount</td>
<td>The number of rows (Mandatory) (0, by default).</td>
</tr>
<tr>
<td>headerReorderingAllowed</td>
<td>Controls whether you can drag the columns (false, by default).</td>
</tr>
<tr>
<td>headerResizingAllowed</td>
<td>Controls whether you can resize all column widths. (false, by default). Note, each column can be individually configured to be resized by using the columnWidthResizable property.</td>
</tr>
<tr>
<td>rowSorterEnabled</td>
<td>Controls whether you can sort the cell values in a column (false, by default).</td>
</tr>
<tr>
<td>showGrid</td>
<td>Controls whether to draw both horizontal and vertical lines (true, by default).</td>
</tr>
<tr>
<td>showHorizontalLines</td>
<td>Controls whether to draw horizontal line (true, by default).</td>
</tr>
<tr>
<td>showVerticalLines</td>
<td>Controls whether to draw vertical line (true, by default).</td>
</tr>
<tr>
<td>rowIndex</td>
<td>Current row index. Zero-based. This value affects some properties below (0, by default).</td>
</tr>
<tr>
<td>columnIndex</td>
<td>Current column index. Zero-based. This value affects all column specific properties below (0, by default).</td>
</tr>
<tr>
<td>cellText</td>
<td>Specifies the text to be filled in the cell specified the current rowIndex and columnIndex (Empty, by default).</td>
</tr>
<tr>
<td>selectedRows</td>
<td>Control or retrieve row selection.</td>
</tr>
</tbody>
</table>
To view all the properties for a widget in the System Console, type:

```
% dashboard_get_properties <widget_type>
```

For example, the System Console returns all the properties for the dial widget when you type:

```
% dashboard_get_properties dial
```
System Console Examples

This section provides design examples for performing board bring-up, creating a simple dashboard, loading and linking a design, and programming a Nios II processor. The Debugging_using_SystemConsole.zip file contains design files for the board bring-up example. The Nios II Ethernet Standard zip files contain the design files for the Nios II processor example.

The instructions for these examples assume that you are familiar with the Quartus II software, Tcl commands, and Qsys.

Download the design files for the example designs from the On-Chip Debugging Design Examples page on the Altera website.

For an online demonstration of how to perform board bring-up using the System Console, refer to System Console: Faster Board Bring-Up and On-chip Debug on the Altera website.

Board Bring-Up with the System Console

You can perform low-level hardware debugging of Qsys systems with the System Console. Debug systems that include IP cores instantiated in your Qsys system or perform initial bring-up of your printed circuit board (PCB). This board bring-up example assumes you have a NEEK board (Cyclone III) and USB-Blaster cable. You can use different hardware for this design example, but additional setup time is required to change device and pin assignments.

Board Bring-Up Flow

The flow for this design example is the following:

1. “Setting Up the Board Bring-Up Design Example” on page 10–29
2. “Verifying JTAG Connectivity” on page 10–30
3. “Verifying Clock and Reset Signals” on page 10–31
4. “Verifying Memory and Other Peripheral Interfaces” on page 10–32
The Qsys design for this design example includes the following modules:

- **JTAG Avalon-MM**—Provides a connection between your development board and the Qsys system via JTAG interface.
- **On-chip memory**—Simplest type of memory for use in an FPGA-based embedded system. The memory is implemented on the FPGA; consequently, external connections on the circuit board are not necessary.
- **Parallel I/O (PIO) module**—Provides a memory-mapped interface between an Avalon-MM slave port and general purpose I/O ports.
- **Checksum Accelerator**—Calculates the checksum of a data buffer in memory. The Checksum Accelerator consists of the following:
  - **Checksum Calculator (checksum_transform.v)**
  - **Read Master (slave.v)**
  - **Checksum Controller (latency_aware_read_master.v)**

### How the Qsys Modules Work

The base address of the memory buffer and data length passes to the Checksum Controller via the Avalon-MM master. The Read Master continuously reads data from memory and passes the data to the Checksum Calculator. When the checksum calculations are complete, the Checksum Calculator issues a valid signal along with the checksum result to the Checksum Controller. The Checksum Controller sets the DONE bit in the status register and also asserts the interrupt signal. You should only read the result from the Checksum Controller when the DONE bit and interrupt signal are asserted.

### Setting Up the Board Bring-Up Design Example

To load the design example into the Quartus II software and program your device, follow these steps:

1. Extract the **Debugging_using_SystemConsole.zip** file to your local hard drive.
2. Open **Systemconsole_design_example.qpf** from step 1.
3. Verify the device name in the **Project Navigator**, Cyclone III: EP3C25, matches your device. If it does not match, change the pin assignments (LED, clock, and reset pins) in the `Systemconsole_design_example.qsf` file.

4. Compile the design. Right-click **Compile Design** under **Task** and click **Start**.

5. Program your device.
   a. On the Tools menu, click **Programmer**.
   b. Click **Hardware Setup**.
   c. Click the **Hardware Settings** tab.
   d. Under **Currently selected hardware**, click **USB-Blaster** and click **Close**.
      
      If you do not see the **USB-Blaster** option, then your device was not detected. Verify that the USB-Blaster driver is installed, the NEEK board is powered off, and the USB-Blaster connecting wire is intact.
      
      This design example has been validated using a USB-Blaster cable. If you do not have a USB-Blaster cable and you are using a different cable type, then select your cable from the **Currently selected hardware** options.
   e. Click **Auto Detect**, select EP3C25 device.
   f. Double-click your device under **File**, browse to your project folder from step 1 and open `Systemconsole_design_example.sof`.
   g. Turn on the **Program/Configure** option.
   h. Click **Start**.
   i. Close the Programmer.

6. Open Qsys. On the Tools menu, click **Qsys**.

7. Open the `Systemconsole_design_example.qsys` file in your project folder from step 1.

   It takes about 2 minutes for the design to load.

---

**Verifying JTAG Connectivity**

To debug any design with the System Console, begin by verifying JTAG connectivity.

1. Open System Console. From Qsys, under the Tools menu, click **System Console**.

2. In the Tcl Console, type `pwd`

   System Console returns the current working directory.

3. Type `get_service_types`

   System Console returns the available service types.

4. Find the available paths to the `jtag_debug` service. Type `get_service_paths jtag_debug`

5. Select a service path. Type `set jtag_debug_path [ lindex [get_service_paths jtag_debug] 0]`

6. Open the service. Type `open_service jtag_debug $jtag_debug_path`
7. Verify the JTAG chain. Type `jtag_debug_loop $jtag_debug_path [list 1 2 3 4 5 6 7 8 9 10]`.
   Sends a specified list of bytes through the loopback debug node and returns a list of bytes in the order received.

8. Issue a reset request to the system through the master_reset output of the JTAG to Avalon Master Bridge. Type `jtag_debug_reset_system $jtag_debug_path`.

**Figure 10–5. Master Reset of the JTAG to Avalon Master Bridge**

<table>
<thead>
<tr>
<th>JTAG to Avalon Master Bridge</th>
<th>JTAG_Master</th>
</tr>
</thead>
<tbody>
<tr>
<td>Clock Input</td>
<td>clk</td>
</tr>
<tr>
<td>Reset Input</td>
<td>clk_react</td>
</tr>
<tr>
<td>Avalon Memory Mapped Master</td>
<td>master</td>
</tr>
<tr>
<td>Reset Output</td>
<td>master_reset</td>
</tr>
</tbody>
</table>

**Example 10–4. Verifying JTAG Connectivity**

```
# jtag
C:/there/seventy/Desktop/System_Consol Debugging using SystemConsole
# get_service_types
alt_xcvr_alternate alt_xcvr_reconfig_q만큼 alt_xcvr_reconfig_dffe alt_xcvr_reconfig_eyes enable alt_xcvr_reconfig_masterids bytestream data_pattern_checker
data_pattern_generator design_instance device qserver jtag jtag_debug_loop master_monitor packet qsu processor slave qsu_trace_db
# get_service_paths jtag_debug
# get_service_paths jtag_debug $jtag_debug_path
# open_service jtag_debug $jtag_debug_path
# jtag_debug_loop $jtag_debug_path [list 1 2 3 4 5 6 7 8 9 10]
```

**Verifying Clock and Reset Signals**

To verify clock and reset signals, follow these steps:

1. Verify that the clock is toggling.
   - Type `jtag_debug_sample_clock $jtag_debug_path` Repeat several times. System Console returns a zero and then a one when the clock toggles.
   - Or type `jtag_debug_sample_clock` System Console returns a one if the clock has ever toggled.

2. Test the reset signal. Type `jtag_debug_sample_reset $jtag_debug_path`.

3. Close `jtag_debug` service. Type `close_service jtag_debug $jtag_debug_path`.
Example 10–5. Verifying Clock and Reset

```bash
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# jtag_debug_sample_clock $jtag_debug_path
# close_service jtag_debug $jtag_debug_path
```

Verifying Memory and Other Peripheral Interfaces

The Avalon-MM service accesses memory-mapped slaves via a suitable Avalon-MM master, which can be controlled by the host. You can use Tcl commands to read and write to memory with a master service. Master services are provided by System Console master components such as the JTAG Avalon master.

**Locating and Opening the Master Service**

1. Select the master service type and check for available service paths. Type `get_service_paths master`.
2. Set the master service path. Type `set master_service_path [ lindex [get_service_paths master] 0 ]`.
3. Open the master service. Type `open_service master $master_service_path`.

**Avalon-MM Slaves**

The address range for every Qsys component is specified under the **Address Map** tab. These addresses are used by the Avalon-MM master to communicate with slaves. Register mapping for all Altera components is specified in their respective Data Sheets.

For more information, refer to the **Data Sheets** page of the Altera website.

---

**Figure 10–6. Address Map**

[Diagram of the Address Map]
Testing the PIO component

The PIO connects to the LEDs of the Cyclone III board. Test if this component is operating properly by forcing values to it with the Avalon-MM master.

1. Write to memory locations of the PIO component.
   a. Type `master_write_8 $master_service_path 0x0 0x7`.
      
      To read back the register value, type `master_read_8 $master_service_path 0x0 0x1`.
   b. Type `master_write_8 $master_service_path 0x0 0x2`.
   c. Type `master_write_8 $master_service_path 0x0 0xe`.
   d. Type `master_write_8 $master_service_path 0x0 0x7`.

   Observe the LEDs turn on and off as you execute these Tcl commands. The LED is on if the register value is zero and off if the register value is one. LED 0, LED 1, and LED 2 connect to the PIO. LED 3 connects to the interrupt signal of the Checksum Accelerator.

Example 10–6. Verifying Clock and Reset Signals

```bash
% get_service_paths master
(/devices/EP4CE22I4Q20-0/linxJTAG/110:132 v1 #0)/phy_0/master
% est master_service_path [ index [get_service_paths master] 0]
/devices/EP4CE22I4Q20-0/linxJTAG/110:132 v1 #0)/phy_0/master
% open_service master $master_service_path
% master_write_8 $master_service_path 0x0 0x7
% master_read_8 $master_service_path 0x0 0x1
% master_write_8 $master_service_path 0x0 0x2
% master_write_8 $master_service_path 0x0 0xe
% master_read_8 $master_service_path 0x0 0x1
% master_write_8 $master_service_path 0x0 0x7
```
Testing On-chip Memory
Test the memory by using a recursive function that writes to incrementing memory addresses.

1. Run the design example Tcl script. Type `source set_memory_values.tcl`

Example 10–7. Testing On-chip Memory

```
% source set_memory_values.tcl

Sending memory values out after filling with data.
0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e
0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e
0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e 0x5e5e5e5e
```

Testing the Checksum Accelerator
The Checksum Accelerator calculates the checksum of a data buffer in memory. It calculates the value for a specified memory buffer, sets the DONE bit in the status register, and asserts the interrupt signal. You should only read the result from the controller when both the DONE bit and the interrupt signal are asserted. The host should assert the interrupt enable control bit in order to check the interrupt signal.

1. Pass base address of the memory buffer and data length to the Checksum Accelerator.
   a. Type `master_write_32 $master_service_path 0x24 0x80`
   b. Type `master_write_32 $master_service_path 0x2C 0x20`

Figure 10–8. Register Map for Checksum Component

<table>
<thead>
<tr>
<th>Offset (Bytes)</th>
<th>Hexadecimal value (after adding offset)</th>
<th>Register</th>
<th>Access</th>
<th>Bits (32 bits)</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>0x20</td>
<td>Status</td>
<td>Read/Write to clear</td>
<td>31-9 8 7-5 4 3 2 1 0</td>
</tr>
<tr>
<td>4</td>
<td>0x24</td>
<td>Address</td>
<td>Read/Write</td>
<td>Read Address</td>
</tr>
<tr>
<td>12</td>
<td>0x2C</td>
<td>Length</td>
<td>Read/Write</td>
<td>Length in bytes</td>
</tr>
<tr>
<td>24</td>
<td>0x38</td>
<td>Control</td>
<td>Read/Write</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Fixed Read Address bit</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Interrupt Enable</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>GO</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>INV</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>Clear</td>
<td></td>
</tr>
<tr>
<td>28</td>
<td>0x3C</td>
<td>Result</td>
<td>Read</td>
<td>Checksum result (upper 16 bits are zero)</td>
</tr>
</tbody>
</table>
2. Write clear to status and control registers.
   a. Type `master_write_32 $master_service_path 0x20 0x0`
   b. Type `master_write_32 $master_service_path 0x38 0x1`

3. Write Go to the control register. Type `master_write_32 $master_service_path 0x38 0x8`

4. Cross check if the checksum DONE bit is set. Type `master_read_32 $master_service_path 0x20 0x1`

5. Is the DONE bit set?
   - If yes, check the result. Type `master_read_16 $master_service_path 0x3C 0x1`
     You are finished with the Board Bring-Up Design Example.
   - If the result is zero and the JTAG chain works properly, the clock and reset
     signals work properly, and the memory works properly, then the problem is
     the Checksum Accelerator component.

6. Confirm if the DONE bit in the status register (bit 0) and interrupt signal are
   asserted.
   a. Type `master_read_32 $master_service_path 0x20 0x1`
      The check DONE bit should return a one.
   b. Type `master_write_32 $master_service_path 0x38 0x18`
   c. Check the Control Enable to see the interrupt signal. LED 3 (MSB) should be
      off. This indicates the interrupt signal is asserted.

   You have narrowed down the problem to the data path. View the RTL to check the
data path.

7. Open the `Checksum_transform.v` file from your project folder.

   <unzip
   dir>/Debugging_using_SystemConsole/ip/checksum_accelerator/checksum_accelerator.v

8. Notice that the data_out signal is grounded, uncommented line 87 and comment
   line 88. Fix the problem.

9. Save the file and regenerate the Qsys system.

10. Re-compile the design and reprogram your device.

11. Redo the above steps, starting with “Verifying Memory and Other Peripheral
    Interfaces” on page 10–32 or run the Tcl script included with this design example.
    - Type `source set_memory_and_run_checksum.tcl`

---

**Figure 10–9. Checksum File**

```vhdl
// first folding
assign first_folded_num = (initial_num[32] + initial_num[24:16] + initial_num[8:0]); // this result is at most 17 bits wide (16 bits with carry)

// second folding and optional inversion, this result is at most 16 bits wide
assign data_out = (invert == '1') ? (first_folded_num[16] + first_folded_num[8:0]) : (first_folded_num[16] + first_folded_num[8:0]);
```

---
Creating a Simple Dashboard

The following shows how to create a simple dashboard:

1. Create a Tcl file inside $HOME/system_console/scripts and name it dashboard_example.tcl.
2. Open the System Console from the command line by typing system-console. You should now see the System Console GUI, including the System Explorer.
3. Add the following lines to your Tcl file:

```tcl
namespace eval dashboard_example {
    set dash [add_service dashboard dashboard_example "Dashboard Example" "Tools/Example"]
    dashboard_set_property $dash self developmentMode true
    dashboard_add $dash mylabel label self
    dashboard_set_property $dash mylabel text "Hello World!"
    dashboard_add $dash mybutton button self
    dashboard_set_property $dash mybutton text "Start Count"
    dashboard_set_property $dash mybutton onClick {::dashboard_example::start_count 0}
    dashboard_set_property $dash self itemsPerRow 1
    dashboard_set_property $dash self visible true
}
```

4. Return to the System Console GUI. Under the System Explorer tree, locate the scripts, and right-click the node dashboard_example.tcl.
5. On the popup menu, click Execute. This command executes the Tcl script that you added to $HOME/system_console/scripts.
6. You should now see an internal window titled “Dashboard Example”, “Hello World!” in the middle of the dashboard window and a button named “Start Count”.
7. To add some behavior to the example dashboard, you can create a seconds counter. First create a proc inside the namespace_eval as follows:

```tcl
proc start_count { c } {
    incr c
    variable dash
```
8. Then add a line in the main script as shown by the following:

```tcl
dashboard_set_property $dash mybutton onclick {
::dashboard_example::start_count 0
}
```

9. As shown in this example, the above lines define a proc inside the namespace. When you click Start Count, the script calls the proc start_count with an argument of 0. The body of the proc is fairly simple. The proc increments the argument, sets the value of the label to the argument, and then directs Tcl to call this proc again in another 1000 milliseconds, using the recently incremented value as an argument.

Example 10–8 shows the entire script.

---

**Example 10–8. Example of Creating a Simple Dashboard**

```tcl
namespace eval dashboard_example {

proc start_count { c } {
incr c
variable dash
dashboard_set_property $dash mylabel text $c
after 1000 ::dashboard_example::start_count $c
}

set dash [add_service dashboard dashboard_example "Dashboard Example" "Tools/Example"]
dashboard_set_property $dash self developmentmode true
dashboard_add $dash mylabel label self
dashboard_set_property $dash mylabel text "Hello World!"
dashboard_add $dash mybutton button self
dashboard_set_property $dash mybutton text "Start Count"
dashboard_set_property $dash mybutton onclick {
::dashboard_example::start_count 0
}dashboard_set_property $dash self itemsperrow 1
dashboard_set_property $dash self visible true
```
Loading and Linking a Design

Example 10–9 shows how to load and link a Quartus II design.

Example 10–9. Loading and Linking a Design

```
% get_service_paths device
/devices/EP2C35@1#USB-0

% set device_path [lindex [get_service_paths device] 0]
/devices/EP2C35@1#USB-0

% design_load /projects/9.1/standard
QuartusDesignFactory elaborating /projects/9.1/standard
QuartusDesignFactory found SOF File at NiosII_cycloneII_2c35_standard.sof
QuartusDesignFactory found JDI File at NiosII_cycloneII_2c35_standard.jdi
QuartusDesignFactory found SOPC Info File at
/projects/9.1/standard/NiosII_cycloneII_2c35_standard_sopc.sopcinfo

% set design_path [lindex [get_service_paths design] 0]
/designs/standard/NiosII_cycloneII_2c35_standard.qpf

% design_link $design_path $device_path
Created a link from /designs/standard to /connections/USB-Blaster [USB-0]/EP2C35.
Created a link from /designs/standard/NiosII_cycloneII_2c35_standard_sopc.sopcinfo/cpu.data_master to /connections/USB-Blaster [USB-0]/EP2C35/cpu.
Created a link from /designs/standard/NiosII_cycloneII_2c35_standard_sopc.sopcinfo/cpu.data_master/jtag_uart.avalon_jtag_slave to /connections/USB-Blaster [USB-0]/EP2C35/jtag_uart
```

Nios II Processor Example

In this example, you program the Nios II processor on your board to run the count binary software example that is included in the Nios II installation. This is a simple program that uses an 8-bit variable to repeatedly count from 0x00 to 0xFF. The output of this variable is displayed on the LEDs on your board. After programming the Nios II processor, you use the System Console processor commands to start and stop the processor.

To run this example, perform the following steps:

1. Download the Nios II Ethernet Standard Design Example for your board from the Altera website.
2. Create a folder to extract the design. For this example, use C:\Count_binary.
3. Unzip the Nios II Ethernet Standard Design Example into C:\Count_binary.
4. In a Nios II command shell, change to the directory of your new project.
5. To program your board, type the following command in a Nios II command shell:
   ```
   nios2-configure-sof niosii_ethernet_standard_<board_version>.sof
   ```
6. Using Nios II Software Build Tools for Eclipse, create a new Nios II Application and BSP from Template using the Count Binary template and targeting the Nios II Ethernet Standard Design Example.
7. To build the executable and linkable format (ELF) file (.elf) for this application, right-click the Count Binary project and select Build Project.
For more information about creating Nios II applications, refer to the *Nios II Software Build Tools* chapter in the *Nios II Software Developer’s Handbook*.

8. Download the .elf file to your board by right-clicking **Count Binary** project and selecting **Run As, Nios II Hardware**.

   The LEDs on your board provide a new light show.

9. Start the System Console by typing `system-console` in your Nios II command shell.

10. Set the processor service path to the Nios II processor by typing the following command:
    ```
    set niosii_proc [lindex [get_service_paths processor] 0]
    ```

11. Open both services by typing the following commands:
    ```
    open_service processor $niosii_proc
    ```

12. Stop the processor by typing the following command:
    ```
    processor_stop $niosii_proc
    ```

   The LEDs on your board freeze.

13. Start the processor by typing the following command:
    ```
    processor_run $niosii_proc
    ```

   The LEDs on your board resume their previous activity.

14. Stop the processor by typing the following command:
    ```
    processor_stop $niosii_proc
    ```

15. Close the services by typing the following command:
    ```
    close_service processor $niosii_proc
    ```

   The `processor_step`, `processor_set_register`, and `processor_get_register` commands provide additional control over the Nios II processor.

**On-Board USB Blaster II Support**

The System Console supports an On-Board USB-Blaster™ II circuit via the USB Debug master command.

For more information about using the On-Board USB-Blaster II development kit for Stratix V devices, refer to the *All Development Kits* page on the Altera website.

Not all Stratix V boards support the On-Board USB-Blaster II. For example, the transceiver signal integrity board does not support the On-Board USB-Blaster II.
Document Revision History

Table 10–32 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2013</td>
<td>13.0.0</td>
<td>Updated Tcl command tables. Added board bring-up design example. Removed SOPC Builder content.</td>
</tr>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Re-organization of content.</td>
</tr>
<tr>
<td>August 2012</td>
<td>12.0.1</td>
<td>Moved Transceiver Toolkit commands to Transceiver Toolkit chapter.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Maintenance release. This chapter adds new System Console features.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>Maintenance release. This chapter adds new System Console features.</td>
</tr>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>Maintenance release. This chapter adds new System Console features.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Maintenance release. This chapter adds new commands and references for Qsys.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Initial release. Previously released as the System Console User Guide, which is being obsoleted. This new chapter adds new commands.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter describes using the Transceiver Toolkit to optimize high-speed serial links in your board design. The Transceiver Toolkit provides real-time control, monitoring, and debugging of the transceiver links running on your board.

Quick Start

Get started quickly with the following steps by downloading Transceiver Toolkit design examples from the On-Chip Debugging Design Examples website.

1. “Configuring the System”
2. “Loading a Design in Transceiver Toolkit”
3. “Identifying Transceiver Channels”
4. “Running Link Tests”

Transceiver Toolkit Overview

The Transceiver Toolkit tests high-speed serial links in your design in real time. You can control the transmitter or receiver channels to optimize transceiver settings and hardware features. The toolkit tests bit-error rate (BER) while running multiple links at the target data rate. You can also run auto sweep tests to identify the best PMA settings for each link. Figure 11–1 shows the Transceiver Toolkit GUI.

Figure 11–1. Transceiver Toolkit GUI
An “EyeQ” graph displays the receiver horizontal and vertical eye margin during testing. The toolkit supports testing of multiple devices across one or more boards simultaneously.

For more information about the System Console platform that supports Transceiver Toolkit and PCB bring-up, refer to Analyzing and Debugging Designs with the System Console in the Quartus II Handbook.

Transceiver Link Debugging Flow

You can use the Transceiver Toolkit in various ways. The following steps highlight one simple transceiver link debugging flow.

<table>
<thead>
<tr>
<th>Design configuration steps</th>
<th>Link debugging steps</th>
</tr>
</thead>
<tbody>
<tr>
<td>1. Use Qsys to define a system that includes IP cores to enable transceivers and debugging. Modify Altera design examples for use in your design.</td>
<td>1. Open the Transceiver Toolkit and load your design.</td>
</tr>
<tr>
<td>2. Use the Pin Planner to assign device I/O pins to match your device and board.</td>
<td>2. Link your design to the hardware you want to test. The toolkit automatically discovers links between transmitter and receiver of the same channel.</td>
</tr>
<tr>
<td>3. Compile your design with the Quartus II software.</td>
<td>3. (Optional) Create any additional links between transceiver and receiver channels.</td>
</tr>
<tr>
<td>4. Configure your development board correctly.</td>
<td>4. Save actions in a Tcl script to save time. The save feature only saves newly created channels and links.</td>
</tr>
<tr>
<td>5. Use the Quartus II Programmer to program the device.</td>
<td>5. Run auto sweep tests to determine the best BER with various combinations of PMA settings. Check <strong>Best Case</strong> in the Auto Sweep BER report for the best BER value. Specify the best PMA settings in <strong>Transmitter Channels</strong> and <strong>Receiver Channels</strong> tabs.</td>
</tr>
<tr>
<td>6.</td>
<td>6. Use the control panels to select the channel, PMA settings, and test pattern for the EyeQ test. Rerun EyeQ sweeps with different settings for best performance.</td>
</tr>
</tbody>
</table>

For more information about design configuration steps in the Quartus II software, refer to About Qsys, About Assigning Device I/O Pins, About Compilation, and Programming Devices in Quartus II Help.

Using the Transceiver Toolkit GUI

The Transceiver Toolkit GUI helps you to easily visualize and debug transceiver links in your design. To launch the GUI, click **Transceiver Toolkit** on the Tools menu. The GUI comprises the elements in Figure 11–1.
Control Panels
You can directly control and monitor transmitters, receivers, and links running on the board in real time. You can transmit a data pattern across the transceiver link, and then report the signal quality of the received data in terms of bit error rate or eye margin with EyeQ. Click Control Transmitter, Control Receiver, or Control Link to adjust transmitter or receiver settings while the channels are running. The Transceiver Toolkit can automatically identify the transceiver links in your design or you can identify them manually.

For more information, refer to Controlling and Monitoring Transceiver Channels and Control Channel and Control Link Panels in Quartus II Help.

Transceiver Auto Sweep
You can sweep ranges for your transceiver PMA settings and run tests automatically with the auto sweep feature. You can store a history of the test runs and keep a record of the best PMA settings. You can then apply the optimized settings to your final design.

For more information, refer to Auto Sweep Panel (Receiver/Transceiver) in Quartus II Help.

Transceiver EyeQ
The EyeQ graph allows you to visualize the estimated horizontal eye margin at the receiver. For Stratix V devices, the EyeQ graph also provides information about the vertical eye margin. The EyeQ graph displays a bathtub curve or eye diagram representing eye margin. The run list displays the statistics of each EyeQ test. You can right-click any of the test runs in the list, and then click Apply Settings to Device to quickly apply that PMA setting to your device. You can also click Export, Import, or Create Report.

There is an Auto sweep & EyeQ mode that allows you to create EyesQ graphs for each PMA setup in a sweep.

For more information, refer to EyeQ Panel (Receiver/Transceiver) in Quartus II Help.

Using the Transceiver Toolkit at the Command Line
You can alternatively use Tcl commands to access Transceiver Toolkit functions, rather than using the GUI. You can script various tasks, such as loading a project, creating design instances, linking device resources, and identifying high-speed serial links. You can save your project setup in a Tcl script for use in subsequent testing sessions. You can also build a custom test routine script.

Saving a Tcl Script
After you set up and define links that describe the entire physical system, you can click Save Tcl Script to save the setup for future use. To run the scripts, double-click script names in the System Explorer scripts folder. Refer to “Transceiver Toolkit Tcl Commands” on page 11–17 for the Tcl commands for the Transceiver Toolkit.
Using Altera Design Examples

Altera provides design examples to help you quickly learn to use the Transceiver Toolkit. You can experiment with these designs and modify them for your own application. Refer to “Using Altera Design Examples” on page 11–13 and to the readme.txt of each design example for more information. Download the Transceiver Toolkit design examples from the On-Chip Debugging Design Examples page of the Altera website.

Following the Online Demonstration

For an online demonstration of how to use the Transceiver Toolkit to run a high-speed link test with one of the design examples, refer to the Transceiver Toolkit Online Demo on the Altera website.

Required Files for Running the Transceiver Toolkit

To use the Transceiver Toolkit, you must have a SRAM Object File (.sof). All the required files are bundled into the .sof file.

Configuring the System

To debug transceiver links in the toolkit, you must first define a system that includes the Altera IP cores required to enable transceiver debugging. You can modify Altera design example components for use in your design. You can quickly edit and connect these components in the Qsys GUI. The Transceiver Toolkit requires one or more of the Altera IP cores in Table 11–1 to support its functions.

Table 11–1. Transceiver Toolkit IP Core Configuration

<table>
<thead>
<tr>
<th>IP Core</th>
<th>Function</th>
<th>Implementation Notes</th>
</tr>
</thead>
</table>
| Custom PHY               | Enables testing of all possible parallel data widths of transceivers.              | ■ Set lanes, group size, serialization factor, data rate, and input clock frequency to match your application.  
|                          |                                    | ■ Turn on Avalon data interfaces.                                                      |
|                          |                                    | ■ Disable 8B/10B.                                                                     |
|                          |                                    | ■ Set Word alignment mode to manual.                                                   |
|                          |                                    | ■ Disable rate match FIFO.                                                            |
|                          |                                    | ■ Disable byte ordering block.                                                        |
| Low Latency PHY          | Supports testing at more than 8.5 Gbps in GT devices or use of PMA direct mode (such as when using six channels in one quad). | ■ Set Phase compensation FIFO mode to EMBEDDED above certain data rates and set to NONE for PMA direct mode (Stratix IV designs only).  
|                          |                                    | ■ Turn on Avalon data interfaces.                                                      |
|                          |                                    | ■ Set serial loopback mode to enable serial loopback controls in the toolkit.          |
| Avalon®-ST Data Pattern Generator | Generates standard data test patterns at Avalon-ST source ports. | ■ Select PRBS7, PRBS15, PRBS23, PRBS31, high frequency, or low frequency patterns.  |
Table 11–1. Transceiver Toolkit IP Core Configuration

<table>
<thead>
<tr>
<th>IP Core</th>
<th>Function</th>
<th>Implementation Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>Avalon-ST Data Pattern Checker</td>
<td>Validates incoming data stream against test patterns accepted on Avalon streaming sink ports.</td>
<td>■ Specify a value for ST_DATA_W that matches the FPGA-fabric interface width.</td>
</tr>
<tr>
<td>Reconfiguration Controller</td>
<td>Controls PMA and other transceiver settings for Stratix V devices.</td>
<td>■ Connect this IP core to all PHY IP cores that you want controlled by the toolkit.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Connect reconfig_from_xcvr to reconfig_to_xcvr.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Enable Analog controls.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Enable EyeQ block (Stratix V devices only).</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Enable AEQ block (Stratix V devices only).</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Enable DFE block (Stratix V devices only).</td>
</tr>
<tr>
<td>JTAG to Avalon Master Bridge</td>
<td>Accepts encoded streams of bytes with transaction data and initiates Avalon-MM transactions.</td>
<td>—</td>
</tr>
</tbody>
</table>

Figure 11–2 shows a single-channel Transceiver Toolkit test configuration in loopback mode. Alternatively, you can transmit and receive from different devices rather than using loopback mode.

Figure 11–2. One Channel Loopback Mode
Figure 11–3 shows four-channel test setup.

**Figure 11–3. Four Channel Loopback Mode**
Figure 11–4 shows an expanded single channel example design for the Stratix V GX device.

Figure 11–4. Expanded Single Channel Example Design For Stratix V GX Device

You can also communicate with other devices that have the capability to generate and verify test patterns that Altera supports.

28-Gbps Transceiver Support

For high data rate applications, you can enable the 28-Gbps Advanced Transceiver Technology (ATT) channels on Stratix V GT devices.

For more information about 28-Gbps transceiver support, refer to Enabling 28-Gbps Advanced Transceiver Technology (ATT) Channels in Quartus II Help.

Conduit Mode Support

You can use the Data Pattern Generator and Data Pattern Checker with the Avalon-ST interface or with conduit interfaces, such as the parallel conduit interface. You can use conduit mode by clearing Enable Avalon interface on the Data Pattern Generator and Data Pattern Checker. Conduit mode requires conduit role adapters.
Serial Bit Comparator

The Serial Bit Comparator allows you to run EyeQ software features with any type of data, from PRBS patterns to user-design data, and without disrupting the data path. You must set up the Transceiver Reconfiguration Controller to use the Serial Bit Comparator. The Serial Bit Comparator is available for Stratix V devices only. To run user-design data, when using the Data Pattern Generator, you must use Bypass mode.

For single bit error checking, the Serial Bit Comparator is less accurate than the Data Pattern Checker. Do not use the Serial Bit Comparator if you are looking for an exact error rate. Use the Serial Bit Comparator for checking a large window of error.

Operational Limitations

The following are the hardware based restrictions for the Serial Bit Comparator.

- One BER counter per reconfiguration controller. Only one channel can be monitored at a time if the reconfiguration controller is shared for multiple channels.
- Reconfiguration controller use is limited while BER is running. For example, you cannot switch logical channels, run DFE one-time, run AEQ one-time, or recalibrate because it corrupts the error count.
- Potential to double count the number of errors.
- The bit error counter is not read in real-time because it is read through the memory mapped interface.

For more information about the Serial Bit Comparator, refer to EyeQ Panel in Quartus II Help.

For more information about using Bypass mode, refer to Auto Sweep Panel and Control Channel and Control Link Panels in Quartus II Help.

Controlling and Monitoring Transceiver Channels

To control and monitor transceiver channels after configuring, constraining, and programming your design into a device, follow these steps:

1. “Loading a Design in Transceiver Toolkit”
2. “Identifying Transceiver Channels”
3. “Running Link Tests”
4. “Viewing Results in the EyeQ Graph”

Loading a Design in Transceiver Toolkit

When you open the Transceiver Toolkit, the toolkit automatically loads the current Quartus II project. To load another design, select Load Design from the File menu. After the project is loaded, you can view design information in the System Explorer.
Linking Hardware Resources

When you open the Transceiver Toolkit, the toolkit automatically discovers connected hardware and designs. You can also manually link a design to connected hardware resources in the System Explorer.

You can identify and test the transceiver link between two Altera devices, or you can transmit a test pattern with a third-party device and monitor the data on an Altera device receiver channel. If a third-party device includes self-test capability, you can send the test pattern from the Altera device and monitor the signal integrity on the third-party receiver channel. If the third-party device supports reverse serial loopback, you can run the test entirely within the Transceiver Toolkit.

If you are using more than one Altera board, you can set up a test with multiple devices linked to the same design. This setup is useful when you want to perform a link test between a transmitter and receiver on two separate devices.

Prior to the Transceiver Toolkit version 11.1, you must manually load and link your design to hardware. In version 11.1 and later, the Transceiver Toolkit automatically links any device programmed with a project.

Linking One Design to One Device Connected By One USB-Blaster Download Cable

To link one design to one device by one USB-Blaster™ download cable, follow these steps:
1. Load the design for your Quartus II project.
2. Link each device to an appropriate design if the design has not auto-linked.
3. Create the link between channels on the device to test.

Linking Two Designs to Two Devices on Same Board Connected By One USB-Blaster Download Cable

To link two designs to two separate devices on the same board, connected by one USB-Blaster download cable, follow these steps:
1. Load the design for all the Quartus II project files you might need.
2. Link each device to an appropriate design if the design has not auto-linked.
3. Open the project for the second device.
4. Link the second device on the JTAG chain to the second design (unless the design auto-links).
5. Create a link between the channels on the devices you want to test.

Linking Two Designs to Two Devices on Separate Boards Connected to Separate USB-Blaster Download Cables

To link two designs to two separate devices on separate boards, connected to separate USB-Blaster download cables, follow these steps:
1. Load the design for all the Quartus II project files you might need.
2. Link each device to an appropriate design if the design has not auto-linked.
3. Create the link between channels on the device to test.

4. Link the device you connected to the second USB-Blaster download cable to the second design.

5. Create a link between the channels on the devices you want to test.

**Linking the Same Design on Two Separate Devices**

To link the same design on two separate devices, follow these steps:

1. In the Transceiver Toolkit, open the `.qpf` you are using on both devices.
2. Link the first device to this design instance.
3. Link the second device to the design.
4. Create a link between the channels on the devices you want to test.

**Linking Unrelated Designs**

Use a combination of the above steps to load multiple Quartus II projects and make links between different systems. You can perform tests on completely separate systems that are not related to one another. All tests run through the same tool instance.

**Verifying Hardware Connections**

After you load your design and link your hardware, verify that the channels are connected correctly and looped back properly on the hardware. Use the toolkit to send data patterns and receive them correctly. Verifying your link and correct channel before you perform Auto Sweep or EyeQ tests can save time in the work flow.

After you have verified that the transmitter and receiver are communicating with each other, you can create a link between the two transceivers so that you can perform Auto Sweep and EyeQ tests with this pair.

To verify your link and correct channel, refer to *Controlling and Monitoring Transceiver Channels* in Quartus II Help. For more information about Control Channel and Control Link settings, refer to *Control Channel and Control Link Panels* in Quartus II Help.
Identifying Transceiver Channels

The Transceiver Toolkit automatically displays existing transmitter and receiver channels. The toolkit identifies a channel automatically whenever a receiver and transmitter share a transceiver channel. For example, Figure 11–5 shows a transceiver link identified in the Transceiver Links tab between a transmitter and receiver.

Figure 11–5. Transceiver Link Identification

When you run link tests, channel color highlights indicate the test status:

- **Green**—Test is running successfully on connected device or devices.
- **Yellow**—Test is running but no data is locked.
- **Red**—Test is not running. Cannot communicate with channel.

For more information about channel color highlights, refer to About the Transceiver Toolkit in Quartus II Help.

You can select the link and click different buttons to control link tests. You can use the transmitter and receiver channels of the same device and loop them back on the far side of the board trace to check the signal integrity of your high-speed interface on the board trace, as shown in Figure 11–3. You can perform a physical link test without loopback by connecting one device transmitter channel to another device receiver channel.

Running Link Tests

Use the Transceiver Links tab options to control how you want to test the link. For example, use the Auto Sweep feature to sweep various transceiver settings parameters through a range of values to find the results that give the best BER value. You can also open Transmitter, Receiver and Link Control panels to manually control the PMA settings and run individual tests. You can change the various controls in the panels to suit your requirements.
Using DFE
To use the decision feedback equalization (DFE) flow, follow these steps:

1. Use the Auto Sweep flow to find optimal PMA settings while leaving the DFE setting OFF.

2. If BER = 0, take the best PMA setting achieved.

3. If BER > 0, then use this PMA setting and set minimum and maximum values obtained from Transceiver Auto Sweep to match this setting. Set the maximum DFE range to limits for each of the three DFE settings.

4. Run Create Report in the Auto Sweep panel to determine which DFE setting results in the best BER. Use these settings in conjunction with the PMA settings for the best results.

For more information about DFE, refer to Control Channel and Control Link Panels in Quartus II Help.

Running Simultaneous Sweep Tests
To run multiple link tests simultaneously in the Transceiver Toolkit, follow these steps:

1. In the control panel for the link you are working on, run either Transceiver Auto Sweep or EyeQ.

2. After you start the test, return to the Transceiver Toolkit control panel.

3. To return to the control panel, select the control panel tab.


5. Repeat step 2 until all tests are run.

The control panel shows which links and channel resources are in use to help identify which channels have tests started and their current run status.

Viewing Results in the EyeQ Graph
Some Altera devices include EyeQ circuitry that allows you to estimate the horizontal, (and for Stratix V devices, the vertical) eye margin at the receiver. With this feature, you can tune the PMA settings of your transceiver, which results in the best eye margin and BER at high data rates.

When PMA settings are suitable, the bathtub curve is wide, with sharp slopes near the edges. The curve may be up to 30 units wide. If the bathtub is narrow, maybe as small as two units wide, then the signal quality may be poor. The wider the bathtub curve, the wider the eye you have. The smaller the bathtub curve, the smaller the eye.

For more information about the EyeQ feature, refer to AN 605: Using the On-Chip Signal Quality Monitoring Circuitry (EyeQ) Feature in Stratix IV Transceivers.
Using Altera Design Examples

You can use the design examples as a starting point to work with a particular signal integrity development board. The design examples provide the components to quickly test the functionality of the receiver and transmitter channels in your design. You can easily change the transceiver settings in the design examples to see how they affect transceiver link performance. You can isolate and verify the high-speed serial links without debugging other logic in your design. You can modify and customize the design examples to match your intended transceiver design.

Download the Transceiver Toolkit design examples from the On-Chip Debugging Design Examples page of the Altera website.

Once you have downloaded the design examples, open the Quartus II software and restore the design example project archive. If you have access to the same development board with the same device as mentioned in the readme.txt file of the example, you can directly program the device with the provided programming file in that example. If you want to recompile the design, you must make your modifications to the configuration in Qsys, regenerate in Qsys, and recompile the design in the Quartus II software to generate a new programming file.

If you have the same board as mentioned in the readme.txt file, but a different device on your board, you must choose the appropriate device and recompile the design. For example, some early development boards are shipped with engineering sample devices.

You can make changes to the design examples so that you can use a different development board or different device. You make pin assignment changes and then recompile the design. If you have a different board, you must edit the necessary pin assignments and recompile the design examples.

Modifying Design Examples

You can modify a design example to experiment with various configurations that match your own design. For example, you can change data rate, number of lanes, PCS-PMA width, FPGA-fabric interface width, or input reference clock frequency. To modify the design examples, you modify the IP core parameters and regenerate the system in Qsys. Next, you modify the top-level design file, and reassign device I/O pins as necessary.

To modify the Low Latency PHY block in the design example, follow these steps:

1. Open the <project name>.qpf file in the Quartus II software.
2. In the Tools menu, click Qsys.
3. On the System Contents tab, right-click the phy block and click Edit. On the General tab, specify options for the PHY block to match your design requirement for number of lanes, data rate, PCS-PMA width, FPGA-fabric interface width, and input reference clock frequency.
4. Right-click the phy block and click Edit. Click the Additional Options tab. Specify a multiple of the FPGA-fabric interface data width for Avalon Data Symbol Size. The available values for Avalon Data Symbol Size are 8 or 10. Click Finish.
5. Delete the timing adapter from the design.
6. Right-click **data pattern generator** and click **Edit**. Specify a value for ST_DATA_W that matches the FPGA-fabric interface width.

7. Right-click **data pattern checker** and click **Edit**. Specify a value for ST_DATA_W that matches the FPGA-fabric interface width.

8. Right-click **transceiver configuration controller** and click **Edit**. Specify 2* number of lanes for the number of reconfigurations interfaces. Click **finish**.

9. From the Component Library, instantiate the data pattern generator and the data pattern checker components. The components are under **Debug and Performance** under **Peripherals**. Add one data pattern generator and data pattern checker for each transmitter and receiver lane.

10. Create connections for the data pattern generator and data pattern checker components. Right-click the net name in the **System Contents** tab and specify the connections listed in **Table 11–2**.

<table>
<thead>
<tr>
<th>From Block Name</th>
<th>From Net Name</th>
<th>To Block Name</th>
<th>To Net Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>clk_100</td>
<td>clk</td>
<td>data_pattern_generator</td>
<td>csr_clk</td>
</tr>
<tr>
<td>clk_100</td>
<td>clk_reset</td>
<td>data_pattern_generator</td>
<td>csr_clk_reset</td>
</tr>
<tr>
<td>master_0</td>
<td>master</td>
<td>data_pattern_generator</td>
<td>csr_slave</td>
</tr>
<tr>
<td>xcvr_*_phy_0</td>
<td>tx_clk_out0</td>
<td>data_pattern_generator</td>
<td>pattern_out_clk</td>
</tr>
<tr>
<td>xcvr_*_phy_0</td>
<td>tx_parallel_data0</td>
<td>data_pattern_generator</td>
<td>pattern_out</td>
</tr>
<tr>
<td>clk_100</td>
<td>clk</td>
<td>data_pattern_checker</td>
<td>csr_clk</td>
</tr>
<tr>
<td>clk_100</td>
<td>clk_reset</td>
<td>data_pattern_checker</td>
<td>csr_clk_reset</td>
</tr>
<tr>
<td>master_0</td>
<td>master</td>
<td>data_pattern_checker</td>
<td>csr_slave</td>
</tr>
<tr>
<td>xcvr_*_phy_0</td>
<td>rx_clk_out0</td>
<td>data_pattern_checker</td>
<td>pattern_in_clk</td>
</tr>
<tr>
<td>xcvr_*_phy_0</td>
<td>rx_parallel_data0</td>
<td>data_pattern_checker</td>
<td>pattern_in</td>
</tr>
</tbody>
</table>

11. On the System menu, click **Assign Base Addresses**.

12. Connect the reset port of timing adapters to clk_reset of clk_100.

13. To implement the changes to the system, click **Generate**.
14. If you modify the number of lanes in the PHY, you must update the top-level file accordingly. **Example 11–1** shows Verilog HDL code for a 2 channel design where input and output ports are declared in the top-level design.

**Example 11–1. I/O Declared in Top-level Two-channel Verilog HDL design**

```verilog
module low_latency_10g_1ch DUT ( 
    input  wire GXB_RXL11, 
    input  wire GXB_RXL12, 
    output wire GXB_TXL11, 
    output wire GXB_TXL12 
); 
.... 
low_latency_10g_1ch DUT ( 
    .... 
    .xcvr_low_latency_phy_0_tx_serial_data_export({GXB_TXL11, 
    GXB_TXL12}), 
    .xcvr_low_latency_phy_0_rx_serial_data_export({GXB_RXL11, 
    GXB_TXL12}), 
    .... 
); 
```

The example design includes the low latency PHY IP core. If you modify the PHY parameters, you must modify the top-level design with the correct port names. Qsys displays an example of the PHY in the design on the **HDL example** tab.

15. Use the Pin Planner to update pin assignments to match your board, as described in “Updating I/O Pin Assignments”.

16. Edit the design’s Synopsys Design Constraints (.sdc) to reflect the reference clock change. Ignore the reset warning messages.

17. Recompile the design.

**Updating I/O Pin Assignments**

You can use the Pin Planner to change pin assignments for development kits. Similarly, you can change the pin assignment for your own board and recompile the design examples. **Table 11–3**, **Table 11–4**, and **Table 11–5** indicate the correct I/O pin assignments for specific Altera development boards.

**Table 11–3. Stratix V GT Top-Level Pin Assignments**

<table>
<thead>
<tr>
<th>Top-Level Signal Name</th>
<th>I/O Standard</th>
<th>Location</th>
</tr>
</thead>
<tbody>
<tr>
<td>GXB_RXL11 (input)</td>
<td>1.4-V PCML</td>
<td>PIN_AA36</td>
</tr>
<tr>
<td>GXB_TXL11 (output)</td>
<td>1.4-V PCML</td>
<td>PIN_Y34</td>
</tr>
<tr>
<td>REFCLK_C70625mhz</td>
<td>LVDS</td>
<td>PIN_U31</td>
</tr>
<tr>
<td>SVGX_C100mhz</td>
<td>LVDS</td>
<td>PIN_AV7</td>
</tr>
</tbody>
</table>
Table 11–4 shows the pin-assignment edits for the Stratix IV Transceiver Signal Integrity Development Kit (DK-SI-4SGX230N). You must make these assignments before you recompile your design.

Table 11–4. Stratix IV GX Top-Level Pin Assignments (DK-SI-4SGX230N)

<table>
<thead>
<tr>
<th>Top-Level Signal Name</th>
<th>I/O Standard</th>
<th>Pin Number on DK-SI-4SGX230N Board</th>
</tr>
</thead>
<tbody>
<tr>
<td>REFCLK_GXB2_156M25 (input)</td>
<td>2.5 V LVTTL/LVCMOS</td>
<td>PIN_G38</td>
</tr>
<tr>
<td>S4GX_50M_CLK4P (input)</td>
<td>2.5 V LVTTL/LVCMOS</td>
<td>PIN_AR22</td>
</tr>
<tr>
<td>GXB1_RX1 (input)</td>
<td>1.4-V PCML</td>
<td>PIN_R38</td>
</tr>
<tr>
<td>GXB1_TX1 (output)</td>
<td>1.4-V PCML</td>
<td>PIN_P36</td>
</tr>
</tbody>
</table>

Table 11–5 shows pin assignments for the Stratix IV GX development kit (DK-DEV-4SGX230N). You must make these assignments before you recompile your design.

Table 11–5. Stratix IV GX Top-Level Pin Assignments (DK-DEV-4SGX230N)

<table>
<thead>
<tr>
<th>Top-Level Signal Name</th>
<th>I/O Standard</th>
<th>Pin Number on DK-DEV-4SGX230N Board</th>
</tr>
</thead>
<tbody>
<tr>
<td>REFCLK_GXB2_156M25 (input)</td>
<td>LVDS</td>
<td>PIN_AA2</td>
</tr>
<tr>
<td>S4GX_50M_CLK4P (input)</td>
<td>2.5 V LVTTL/LVCMOS</td>
<td>PIN_AC34</td>
</tr>
<tr>
<td>GXB1_RX1 (input)</td>
<td>1.4-V PCML</td>
<td>PIN_AU2</td>
</tr>
<tr>
<td>GXB1_TX1 (output)</td>
<td>1.4-V PCML</td>
<td>PIN_AT4</td>
</tr>
</tbody>
</table>

For more information about other development kits, read the readme.txt file accompanying the design examples. For more information about the pinouts refer to the corresponding reference manual on the Development Kits page of the Altera website.
Transceiver Toolkit Tcl Commands

You can view a list of the available Tcl commands for the Transceiver Toolkit from the Tcl Console window. From the list, you can select individual Tcl commands and view descriptions of the commands. Some commands have additional information about example usage in the description.

To view Tcl command descriptions from the Tcl Console window:

1. Type `help help`. All the available Tcl commands for the Transceiver Toolkit are displayed.
2. Identify the command of interest from the list and type `help <command name>`. The Tcl Console window displays the description for the specified command.

In Example 11–2, several Tcl commands are used together to locate and open an instance of the master service.

- `set`—Tcl command that sets the value of a variable.
- `m_path`—name of a variable.
- `lindex`—Tcl command that retrieves an element from a list.
- `get_service_paths`—returns a list of service instances by type.
- `master`—service type that accesses memory-mapped slaves.
- `n`—number of the master service.
- `open_service`—initiates a service instance.

**Example 11–2. Locate and Open Service Path**

```tcl
set m_path [lindex [get_service_paths master] n]
open_service master $m_path
```

Table 11–6 through Table 11–15 lists the Transceiver Toolkit commands.

**Table 11–6. Transceiver Toolkit Channel_rx Commands (Part 1 of 3)**

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_channel_rx_get_data</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns a list of the current checker data. The results are in the order of number of bits, number of errors, and bit error rate.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_dcgain</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the DC gain value on the receiver channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_dfe_tap_value</td>
<td><code>&lt;service-path&gt;</code>&lt;br&gt;<code>&lt;tap position&gt;</code></td>
<td>Gets the current tap value of the specified channel at the specified tap position.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_eqctrl</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the equalization control value on the receiver channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_pattern</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns the current data checker pattern by name.</td>
</tr>
<tr>
<td>transceiver_channel_rx_has_dfe</td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets whether this channel has the DFE feature available.</td>
</tr>
<tr>
<td>Command</td>
<td>Arguments</td>
<td>Function</td>
</tr>
<tr>
<td>----------------------------------------------</td>
<td>------------------------------------</td>
<td>----------------------------------------------------------------</td>
</tr>
<tr>
<td>transceiver_channel_rx_has_eyeq</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the EyeQ feature is available for the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_checking</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the checker is running.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_dfe_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the DFE feature is enabled on the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_locked</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the checker is locked onto the incoming data.</td>
</tr>
<tr>
<td>transceiver_channel_rx_reset_counters</td>
<td>&lt;service-path&gt;</td>
<td>Resets the bit and error counters inside the checker.</td>
</tr>
<tr>
<td>transceiver_channel_rx_reset</td>
<td>&lt;service-path&gt;</td>
<td>Resets the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_dcgain</td>
<td>&lt;service-path&gt; &lt;value&gt;</td>
<td>Sets the DC gain value on the receiver channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_dfe_enabled</td>
<td>&lt;service-path&gt; &lt;disable(0)/enable(1)&gt;</td>
<td>Enables or disables the DFE feature on the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_dfe_tap_value</td>
<td>&lt;service-path&gt; &lt;tap position&gt; &lt;tap value&gt;</td>
<td>Sets the current tap value of the specified channel at the specified tap position to the specified value.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_dfe_adaptive</td>
<td>&lt;service-path&gt;</td>
<td>Sets the mode of DFE adaptation.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_eqctrl</td>
<td>&lt;service-path&gt; &lt;value&gt;</td>
<td>Sets the equalization control value on the receiver channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_start_checking</td>
<td>&lt;service-path&gt;</td>
<td>Starts the checker.</td>
</tr>
<tr>
<td>transceiver_channel_rx_stop_checking</td>
<td>&lt;service-path&gt;</td>
<td>Stops the checker.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_eyeq_phase_step</td>
<td>&lt;service-path&gt;</td>
<td>Gets the current phase step of the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_pattern</td>
<td>&lt;service-path&gt; &lt;pattern-name&gt;</td>
<td>Sets the expected pattern to the one specified by the pattern name.</td>
</tr>
<tr>
<td>transceiver_channel_rx_has_eyeq</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the specified channel has the EyeQ feature.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_eyeq_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the EyeQ feature is enabled on the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_eyeq_enabled</td>
<td>&lt;service-path&gt; &lt;disable(0)/enable(1)&gt;</td>
<td>Enables or disables the EyeQ feature on the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_eyeq_phase_step</td>
<td>&lt;service-path&gt; &lt;phase step&gt;</td>
<td>Sets the phase step of the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_word_aligner_enabled</td>
<td>&lt;service-path&gt; &lt;disable(0)/enable(1)&gt;</td>
<td>Enables or disables the word aligner of the specified channel.</td>
</tr>
</tbody>
</table>
### Table 11–6. Transceiver Toolkit Channel rx Commands  (Part 3 of 3)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_channel_rx_is_word_aligner_enabled</td>
<td>&lt;service-path&gt; &lt;disable(0)/enable(1)&gt;</td>
<td>Gets whether the word aligner feature is enabled on the specified channel.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_locked</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the checker is locked onto the incoming signal.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_rx_locked_to_data</td>
<td>&lt;service-path&gt;</td>
<td>Returns 1 if transceiver in lock to data (LTD) mode. Otherwise 0.</td>
</tr>
<tr>
<td>transceiver_channel_rx_is_rx_locked_to_ref</td>
<td>&lt;service-path&gt;</td>
<td>Returns 1 if transceiver in lock to reference (LTR) mode. Otherwise 0.</td>
</tr>
<tr>
<td>transceiver_channel_rx_has_eyeq_1d</td>
<td>&lt;service-path&gt;</td>
<td>Detects whether the eye viewer pointed to by &lt;service-path&gt; supports 1D-EyeQ mode.</td>
</tr>
<tr>
<td>transceiver_channel_rx_set_1deye_mode</td>
<td>&lt;service-path&gt; &lt;disable(0)/enable(1)&gt;</td>
<td>Enables or disables 1D-EyeQ mode.</td>
</tr>
<tr>
<td>transceiver_channel_rx_get_1deye_mode</td>
<td>&lt;service-path&gt;</td>
<td>Returns the current on or off status of 1D-EyeQ mode.</td>
</tr>
</tbody>
</table>

### Table 11–7. Transceiver Toolkit Channel tx Commands  (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_channel_tx_disable_preamble</td>
<td>&lt;service-path&gt;</td>
<td>Disables the preamble mode at the beginning of generation.</td>
</tr>
<tr>
<td>transceiver_channel_tx_enable_preamble</td>
<td>&lt;service-path&gt;</td>
<td>Enables the preamble mode at the beginning of generation.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_number_of_preamble_beats</td>
<td>&lt;service-path&gt;</td>
<td>Returns the currently set number of beats to send out the preamble word.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_pattern</td>
<td>&lt;service-path&gt;</td>
<td>Returns the currently set pattern.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_preamble_word</td>
<td>&lt;service-path&gt;</td>
<td>Returns the currently set preamble word.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_preemph0t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis pre-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_preemph1t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis first post-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_preemph2t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis second post-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_get_vodctrl</td>
<td>&lt;service-path&gt;</td>
<td>Gets the $V_{DD}$ control value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_inject_error</td>
<td>&lt;service-path&gt;</td>
<td>Injects a 1-bit error into the generator’s output.</td>
</tr>
</tbody>
</table>
### Table 11–7. Transceiver Toolkit Channel _tx Commands (Part 2 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_channel_tx_is_generating</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the generator is running.</td>
</tr>
<tr>
<td>transceiver_channel_tx_is_preamble_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if preamble mode is enabled.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_number_of_preamble_beats</td>
<td>&lt;service-path&gt; &lt;number-of-preamble-beats&gt;</td>
<td>Sets the number of beats to send out the preamble word.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_pattern</td>
<td>&lt;service-path&gt; &lt;pattern-name&gt;</td>
<td>Sets the output pattern to the one specified by the pattern name.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_preamble_word</td>
<td>&lt;service-path&gt; &lt;preamble-word&gt;</td>
<td>Sets the preamble word to be sent out.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_preemph0t</td>
<td>&lt;service-path&gt; &lt;preemph0t value&gt;</td>
<td>Sets the pre-emphasis pre-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_preemph1t</td>
<td>&lt;service-path&gt; &lt;preemph1t value&gt;</td>
<td>Sets the pre-emphasis first post-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_preemph2t</td>
<td>&lt;service-path&gt; &lt;preemph2t value&gt;</td>
<td>Sets the pre-emphasis second post-tap value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_set_vodctrl</td>
<td>&lt;service-path&gt; &lt;vodctrl value&gt;</td>
<td>Sets the VOD control value on the transmitter channel.</td>
</tr>
<tr>
<td>transceiver_channel_tx_start_generation</td>
<td>&lt;service-path&gt;</td>
<td>Starts the generator.</td>
</tr>
<tr>
<td>transceiver_channel_tx_stop_generation</td>
<td>&lt;service-path&gt;</td>
<td>Stops the generator.</td>
</tr>
</tbody>
</table>

### Table 11–8. Transceiver Toolkit Debug_Link Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_debug_link_get_pattern</td>
<td>&lt;service-path&gt;</td>
<td>Gets the currently set pattern the link uses to run the test.</td>
</tr>
<tr>
<td>transceiver_debug_link_is_running</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the test is running on the link.</td>
</tr>
<tr>
<td>transceiver_debug_link_set_pattern</td>
<td>&lt;service-path&gt; &lt;data pattern&gt;</td>
<td>Sets the pattern the link uses to run the test.</td>
</tr>
<tr>
<td>transceiver_debug_link_start_running</td>
<td>&lt;service-path&gt;</td>
<td>Starts running a test with the currently selected test pattern.</td>
</tr>
<tr>
<td>transceiver_debug_link_stop_running</td>
<td>&lt;service-path&gt;</td>
<td>Stops running the test.</td>
</tr>
<tr>
<td>Command</td>
<td>Arguments</td>
<td>Function</td>
</tr>
<tr>
<td>----------------------------------------------</td>
<td>---------------------</td>
<td>--------------------------------------------------------------------------</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_logical_channel_address</td>
<td>&lt;service-path&gt;</td>
<td>Gets the transceiver logical channel address currently set.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_rx_dcgain</td>
<td>&lt;service-path&gt;</td>
<td>Gets the DC gain value on the receiver channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_rx_eqctrl</td>
<td>&lt;service-path&gt;</td>
<td>Gets the equalization control value on the receiver channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_tx_preemph0t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis pre-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_tx_preemph1t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis first post-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_tx_preemph2t</td>
<td>&lt;service-path&gt;</td>
<td>Gets the pre-emphasis second post-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_get_tx_vodctrl</td>
<td>&lt;service-path&gt;</td>
<td>Gets the VOD control value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_logical_channel_address</td>
<td>&lt;service-path&gt; &lt;logical channel address&gt;</td>
<td>Sets the transceiver logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_rx_dcgain</td>
<td>&lt;service-path&gt; &lt;dc_gain value&gt;</td>
<td>Sets the DC gain value on the receiver channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_rx_eqctrl</td>
<td>&lt;service-path&gt; &lt;eqctrl value&gt;</td>
<td>Sets the equalization control value on the receiver channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_tx_preemph0t</td>
<td>&lt;service-path&gt; &lt;preemph0t value&gt;</td>
<td>Sets the pre-emphasis pre-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_tx_preemph1t</td>
<td>&lt;service-path&gt; &lt;preemph1t value&gt;</td>
<td>Sets the pre-emphasis first post-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
</tbody>
</table>
Table 11–9. Transceiver Toolkit Reconfig_Analog Commands (Part 2 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>transceiver_reconfig_analog_set_tx_preemph2t</td>
<td>&lt;service-path&gt;</td>
<td>Sets the pre-emphasis second post-tap value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td></td>
<td>&lt;preemph2t value&gt;</td>
<td></td>
</tr>
<tr>
<td>transceiver_reconfig_analog_set_tx_vodctrl</td>
<td>&lt;service-path&gt;</td>
<td>Sets the VOD control value on the transmitter channel specified by the current logical channel address.</td>
</tr>
<tr>
<td></td>
<td>&lt;vodctrl value&gt;</td>
<td></td>
</tr>
</tbody>
</table>

Table 11–10. Transceiver Toolkit Decision Feedback Equalization (DFE) Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>alt_xcvr_reconfig_dfe_get_logical_channel_address</td>
<td>&lt;service-path&gt;</td>
<td>Gets the logical channel address that other alt_xcvr_reconfig_dfe commands use to apply.</td>
</tr>
<tr>
<td>alt_xcvr_reconfig_dfe_is_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the DFE feature is enabled on the previously specified channel.</td>
</tr>
<tr>
<td>alt_xcvr_reconfig_dfe_set_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Enables or disables the DFE feature on the previously specified channel.</td>
</tr>
<tr>
<td></td>
<td>&lt;disable(0)/enable(1)&gt;</td>
<td></td>
</tr>
<tr>
<td>alt_xcvr_reconfig_dfe_set_logical_channel_address</td>
<td>&lt;service-path&gt;</td>
<td>Sets the logical channel address that other alt_xcvr_reconfig_eye_viewer commands use.</td>
</tr>
<tr>
<td></td>
<td>&lt;logical channel address&gt;</td>
<td></td>
</tr>
<tr>
<td>alt_xcvr_reconfig_dfe_set_tap_value</td>
<td>&lt;service-path&gt;</td>
<td>Sets the tap value at the previously specified channel at specified tap position and value.</td>
</tr>
<tr>
<td></td>
<td>&lt;tap position&gt;</td>
<td></td>
</tr>
<tr>
<td></td>
<td>&lt;tap value&gt;</td>
<td></td>
</tr>
</tbody>
</table>

Table 11–11. Transceiver Toolkit Eye Monitor Commands (Part 1 of 3)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>alt_xcvr_custom_is_word_aligner_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Gets whether the word aligner feature is enabled on the previously specified channel.</td>
</tr>
<tr>
<td></td>
<td>&lt;disable(0)/enable(1)&gt;</td>
<td></td>
</tr>
<tr>
<td>alt_xcvr_custom_set_word_aligner_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Enables or disables the word aligner of the previously specified channel.</td>
</tr>
<tr>
<td></td>
<td>&lt;disable(0)/enable(1)&gt;</td>
<td></td>
</tr>
<tr>
<td>alt_xcvr_custom_is_rx_locked_to_data</td>
<td>&lt;service-path&gt;</td>
<td>Returns whether the receiver CDR is locked to data.</td>
</tr>
<tr>
<td>alt_xcvr_custom_is_rx_locked_to_ref</td>
<td>&lt;service-path&gt;</td>
<td>Returns whether the receiver CDR PLL is locked to the reference clock.</td>
</tr>
</tbody>
</table>
### Table 11–11. Transceiver Toolkit Eye Monitor Commands (Part 2 of 3)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>alt_xcvr_custom_is_serial_loopback_enabled</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns whether the serial loopback mode of the previously specified channel is enabled.</td>
</tr>
<tr>
<td><code>alt_xcvr_custom_set_serial_loopback_enabled</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Enables or disables the serial loopback mode of the previously specified channel.</td>
</tr>
<tr>
<td><code>alt_xcvr_custom_is_tx_pll_locked</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Returns whether the transmitter PLL is locked to the reference clock.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_get_logical_channel_address</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the logical channel address on which other <code>alt_reconfig_eye_viewer</code> commands will use to apply.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_get_phase_step</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the current phase step of the previously specified channel.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_is_enabled</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets whether the EyeQ feature is enabled on the previously specified channel.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_set_enabled</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Enables or disables the EyeQ feature on the previously specified channel. Setting a value of 2 enables both EyeQ and the Serial Bit Comparator.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_set_logical_channel_address</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Sets the logical channel address on which other <code>alt_reconfig_eye_viewer</code> commands will use to apply.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_set_phase_step</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Sets the phase step of the previously specified channel.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_has_ber_checker</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Detects whether the eye viewer pointed to by <code>&lt;service-path&gt;</code> supports the Serial Bit Comparator.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_is_enabled</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Detects whether the Serial Bit Comparator is enabled.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_start</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Starts the Serial Bit Comparator counters.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_stop</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Stops the Serial Bit Comparator counters.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_reset_counters</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Resets the Serial Bit Comparator counters.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_is_running</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets whether the Serial Bit Comparator counters are currently running or not.</td>
</tr>
</tbody>
</table>
**Table 11–11. Transceiver Toolkit Eye Monitor Commands (Part 3 of 3)**

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_ber_checker_get_data</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the current total bit, error bit, and exception counts for the Serial Bit Comparator.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_has_1deye</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Detects whether the eye viewer pointed to by <code>&lt;service-path&gt;</code> supports 1D-EyeQ mode.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_set_1deye_mode</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Enables or disables 1D-EyeQ mode.</td>
</tr>
<tr>
<td><code>alt_xcvr_reconfig_eye_viewer_get_1deye_mode</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Gets the enable or disabled state of 1D-EyeQ mode.</td>
</tr>
</tbody>
</table>

**Table 11–12. Channel Type Commands**

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>get_channel_type</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Reports the detected type (GX/GT) of channel <code>&lt;logical-channel-num&gt;</code> for the reconfiguration block located at <code>&lt;service-path&gt;</code>.</td>
</tr>
<tr>
<td><code>set_channel_type</code></td>
<td><code>&lt;service-path&gt;</code></td>
<td>Overrides the detected channel type of channel <code>&lt;logical-channel-num&gt;</code> for the reconfiguration block located at <code>&lt;service-path&gt;</code> to the type specified (0:GX, 1:GT).</td>
</tr>
</tbody>
</table>
Table 11–13. Loopback Commands

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
</table>
| loopback_get      | <service-path> | Returns the value of a setting or result on the loopback channel. Available results include:  
|                   |           | ■ Status—running or stopped.  
|                   |           | ■ Bytes—number of bytes sent through the loopback channel.  
|                   |           | ■ Errors—number of errors reported by the loopback channel.  
|                   |           | ■ Seconds—number of seconds since the loopback channel was started. |
| loopback_set      | <service-path> | Sets the value of a setting controlling the loopback channel. Some settings are only supported by particular channel types. Available settings include:  
|                   |           | ■ Timer—number of seconds for the test run.  
|                   |           | ■ Size—size of the test data.  
|                   |           | ■ Mode—mode of the test. |
| loopback_start    | <service-path> | Starts sending data through the loopback channel. |
| loopback_stop     | <service-path> | Stops sending data through the loopback channel. |

### Data Pattern Generator Commands

The Data Pattern Generator commands allow you to control data patterns that you generate for testing, debugging, and optimizing your design. You must first insert debug IP to enable these commands. Table 11–14 lists the Data Pattern Generator commands.

Table 11–14. Data Pattern Generator Commands (Part 1 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_pattern_generator_start</td>
<td>&lt;service-path&gt;</td>
<td>Starts the data pattern generator.</td>
</tr>
<tr>
<td>data_pattern_generator_stop</td>
<td>&lt;service-path&gt;</td>
<td>Stops the data pattern generator.</td>
</tr>
<tr>
<td>data_pattern_generator_is_generating</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the generator is running.</td>
</tr>
<tr>
<td>data_pattern_generator_inject_error</td>
<td>&lt;service-path&gt;</td>
<td>Injects a 1-bit error into the generator output.</td>
</tr>
</tbody>
</table>
Table 11–14. Data Pattern Generator Commands (Part 2 of 2)

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_pattern_generator_set_pattern</td>
<td>&lt;service-path&gt; &lt;pattern-name&gt;</td>
<td>Sets the output pattern specified by the &lt;pattern-name&gt;. In all, 6 patterns are available, 4 are pseudo-random binary sequences (PRBS), 1 is high frequency and 1 is low frequency. The PRBS7, PRBS15, PRBS23, PRBS31, HF (outputs high frequency, constant pattern of alternating 0s and 1s), and LF (outputs low frequency, constant pattern of 10b'1111100000 for 10-bit symbols and 8b'11111000 for 8-bit symbols) pattern names are defined. PRBS files are clear text and you can modify the PRBS files.</td>
</tr>
<tr>
<td>data_pattern_generator_get_pattern</td>
<td>&lt;service-path&gt;</td>
<td>Returns currently selected output pattern.</td>
</tr>
<tr>
<td>data_pattern_generator_get_available_patterns</td>
<td>&lt;service-path&gt;</td>
<td>Returns a list of available data patterns by name.</td>
</tr>
<tr>
<td>data_pattern_generator_enable_preamble</td>
<td>&lt;service-path&gt;</td>
<td>Enables the preamble mode at the beginning of generation.</td>
</tr>
<tr>
<td>data_pattern_generator_disable_preamble</td>
<td>&lt;service-path&gt;</td>
<td>Disables the preamble mode at the beginning of generation.</td>
</tr>
<tr>
<td>data_pattern_generator_is_preamble_enabled</td>
<td>&lt;service-path&gt;</td>
<td>Returns a non-zero value if preamble mode is enabled.</td>
</tr>
<tr>
<td>data_pattern_generator_set_preamble_word</td>
<td>&lt;service-path&gt; &lt;preamble-word&gt;</td>
<td>Sets the preamble word (could be 32-bit or 40-bit).</td>
</tr>
<tr>
<td>data_pattern_generator_get_preamble_word</td>
<td>&lt;service-path&gt;</td>
<td>Gets the preamble word.</td>
</tr>
<tr>
<td>data_pattern_generator_set_preamble_beats</td>
<td>&lt;service-path&gt; &lt;number-of-preamble-beats&gt;</td>
<td>Sets the number of beats to send out in the preamble word.</td>
</tr>
<tr>
<td>data_pattern_generator_get_preamble_beats</td>
<td>&lt;service-path&gt;</td>
<td>Returns the currently set number of beats to send out in the preamble word.</td>
</tr>
<tr>
<td>data_pattern_generator_fcnter_start</td>
<td>&lt;service-path&gt; &lt;max-cycles&gt;</td>
<td>Sets the max cycle count and starts the frequency counter.</td>
</tr>
<tr>
<td>data_pattern_generator_fcnter_check_status</td>
<td>&lt;service-path&gt;</td>
<td>Queries the data pattern generator for current status. Returns a bitmap indicating the status, with bits defined as follows: [0]-enabled, [1]-bypass enabled, [2]-avalon, [3]-sink ready, [4]-source valid, and [5]-frequency counter enabled.</td>
</tr>
<tr>
<td>data_pattern_generator_fcnter_report</td>
<td>&lt;service-path&gt; &lt;force-stop&gt;</td>
<td>Reports the current measured clock ratio, stopping the counting first depending on &lt;force-stop&gt;.</td>
</tr>
</tbody>
</table>
### Data Pattern Checker Commands

Data Pattern Checker commands allow you to verify the data patterns that you generate. You must first insert debug IP to enable these commands. Table 11–15 lists Data Pattern Checker commands.

<table>
<thead>
<tr>
<th>Command</th>
<th>Arguments</th>
<th>Function</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_pattern_checker_start</td>
<td>&lt;service-path&gt;</td>
<td>Starts the data pattern checker.</td>
</tr>
<tr>
<td>data_pattern_checker_stop</td>
<td>&lt;service-path&gt;</td>
<td>Stops the data pattern checker.</td>
</tr>
<tr>
<td>data_pattern_checker_is_checking</td>
<td>&lt;service-path&gt;</td>
<td>Returns a non-zero value if the checker is running.</td>
</tr>
<tr>
<td>data_pattern_checker_is_locked</td>
<td>&lt;service-path&gt;</td>
<td>Returns non-zero if the checker is locked onto the incoming data.</td>
</tr>
<tr>
<td>data_pattern_checker_set_pattern</td>
<td>&lt;service-path&gt;</td>
<td>Sets the expected pattern to the one specified by the &lt;pattern-name&gt;.</td>
</tr>
<tr>
<td>data_pattern_checker_get_pattern</td>
<td>&lt;service-path&gt;</td>
<td>Returns the currently selected expected pattern by name.</td>
</tr>
<tr>
<td>data_pattern_checker_get_available_patterns</td>
<td>&lt;service-path&gt;</td>
<td>Returns a list of available data patterns by name.</td>
</tr>
<tr>
<td>data_pattern_checker_get_data</td>
<td>&lt;service-path&gt;</td>
<td>Returns a list of the current checker data. The results are in the following order: number of bits, number of errors, and bit error rate.</td>
</tr>
<tr>
<td>data_pattern_checker_reset_counters</td>
<td>&lt;service-path&gt;</td>
<td>Resets the bit and error counters inside the checker.</td>
</tr>
<tr>
<td>data_pattern_checker_fcnter_start</td>
<td>&lt;service-path&gt;</td>
<td>Sets the max cycle count and starts the frequency counter.</td>
</tr>
<tr>
<td>data_pattern_checker_fcnter_report</td>
<td>&lt;service-path&gt;</td>
<td>Queries the data pattern checker for current status. Returns a bitmap indicating status, with bits defined as follows: [0]-enabled, [1]-locked, [2]-bypass enabled, [3]-avalon, [4]-sink ready, [5]-source valid, and [6]-frequency counter enabled.</td>
</tr>
<tr>
<td></td>
<td></td>
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<td></td>
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</tbody>
</table>
Table 11–16 lists the revision history for this handbook chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>Minor editorial updates. Added Tcl help information and removed Tcl command tables. Added 28-Gbps Transceiver support section. Added Figure 11-4.</td>
</tr>
<tr>
<td>August, 2012</td>
<td>12.0.1</td>
<td>General reorganization and revised steps in modifying Altera example designs</td>
</tr>
<tr>
<td>June, 2012</td>
<td>12.0.0</td>
<td>Maintenance release for update of Transceiver Toolkit features.</td>
</tr>
<tr>
<td>May, 2011</td>
<td>11.0.0</td>
<td>Added new Tcl scenario.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Changed to new document template. Added new 10.1 release features.</td>
</tr>
<tr>
<td>August 2010</td>
<td>10.0.1</td>
<td>Corrected links</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Initial release</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
This chapter provides detailed instructions about how to use SignalProbe to quickly debug your design. The SignalProbe incremental routing feature helps reduce the hardware verification process and time-to-market for system-on-a-programmable-chip (SOPC) designs.

Easy access to internal device signals is important in the design or debugging process. The SignalProbe feature makes design verification more efficient by routing internal signals to I/O pins quickly without affecting the design. When you start with a fully routed design, you can select and route signals for debugging to either previously reserved or currently unused I/O pins.

The SignalProbe feature supports the Arria® series, Cyclone® series, MAX® II, and Stratix® series, device families.

The Quartus® II software provides a portfolio of on-chip debugging tools. For an overview and comparison of all the tools available in the Quartus II software, refer to Section IV. System Debugging Tools in volume 3 of the Quartus II Handbook.

Debugging Using the SignalProbe Feature

The SignalProbe feature allows you to reserve available pins and route internal signals to those reserved pins, while preserving the behavior of your design. SignalProbe is an effective debugging tool that provides visibility into your FPGA.

You can reserve pins for SignalProbe and assign I/O standards after a full compilation. Each SignalProbe-source to SignalProbe-pin connection is implemented as an engineering change order (ECO) that is applied to your netlist after a full compilation.

To route the internal signals to the device’s reserved pins for SignalProbe, perform the following tasks:

1. Performing a Full Compilation, described on page 12–2.
3. Assigning SignalProbe Sources, described on page 12–2.
5. Performing a SignalProbe Compilation, described on page 12–3.
6. Analyzing the Results of a SignalProbe Compilation, described on page 12–4.
Performing a Full Compilation

You must complete a full compilation to generate an internal netlist containing a list of internal nodes to probe.

To perform a full compilation, on the Processing menu, click Start Compilation.

Reserving SignalProbe Pins

SignalProbe pins can only be reserved after a full compilation. You can also probe any unused I/Os of the device. Assigning sources is a simple process after reserving SignalProbe pins. The sources for SignalProbe pins are the internal nodes and registers in the post-compilation netlist that you want to probe.

Although you can reserve SignalProbe pins using many features within the Quartus II software, including the Pin Planner and the Tcl interface, you should use the SignalProbe Pins dialog box to create and edit your SignalProbe pins.

For more information, refer to About SignalProbe in Quartus II Help.

Assigning SignalProbe Sources

A SignalProbe source can be any combinational node, register, or pin in your post-compilation netlist. To find a SignalProbe source, in the Node Finder, use the SignalProbe filter to remove all sources that cannot be probed. You might not be able to find a particular internal node because the node can be optimized away during synthesis, or the node cannot be routed to the SignalProbe pin. For example, you cannot probe nodes and registers within Gigabit transceivers in Stratix IV devices because there are no physical routes available to the pins.

To probe virtual I/O pins generated in low-level partitions in an incremental compilation flow, select the source of the logic that feeds the virtual pin as your SignalProbe source pin.

For more information, refer to SignalProbe Pins Dialog Box and Add SignalProbe Pins Dialog Box in Quartus II Help.

Because SignalProbe pins are implemented and routed as ECOs, turning the SignalProbe enable option on or off is the same as selecting Apply Selected Change or Restore Selected Change in the Change Manager window. (If the Change Manager window is not visible at the bottom of your screen, on the View menu, point to Utility Windows and click Change Manager.)

For more information about the Change Manager for the Chip Planner and Resource Property Editor, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.
Adding Registers Between Pipeline Paths and SignalProbe Pins

You can specify the number of registers placed between a SignalProbe source and a SignalProbe pin. The registers synchronize data to a clock and control the latency of the SignalProbe outputs. The SignalProbe feature automatically inserts the number of registers specified into the SignalProbe path.

Figure 12–1 shows a single register between the SignalProbe source Reg_b_1 and SignalProbe SignalProbe_Output_2 output pin added to synchronize the data between the two SignalProbe output pins.

When you add a register to a SignalProbe pin, the SignalProbe compilation attempts to place the register to best meet timing requirements. You can place SignalProbe registers either near the SignalProbe source to meet fMAX requirements, or near the I/O to meet tCO requirements.

Figure 12–1. Synchronizing SignalProbe Outputs with a SignalProbe Register

To pipeline an existing SignalProbe connection, refer to Add SignalProbe Pins Dialog Box in Quartus II Help.

In addition to clock input for pipeline registers, you can also specify a reset signal pin for pipeline registers. To specify a reset pin for pipeline registers, use the Tcl command make_sp, as described in “Scripting Support” on page 12–6.

Performing a SignalProbe Compilation

Perform a SignalProbe compilation to route your SignalProbe pins. A SignalProbe compilation saves and checks all netlist changes without recompiling the other parts of the design. A SignalProbe compilation takes a fraction of the time of a full compilation to finish. The design’s current placement and routing are preserved.

To perform a SignalProbe compilation, on the Processing menu, point to Start and click Start SignalProbe Compilation.
Analyzing the Results of a SignalProbe Compilation

After a SignalProbe compilation, the results are available in the compilation report file. Each SignalProbe pin is displayed in the SignalProbe Fitting Result page in the Fitter section of the Compilation Report. To view the status of each SignalProbe pin in the SignalProbe Pins dialog box, on the Tools menu, click SignalProbe Pins.

The status of each SignalProbe pin appears in the Change Manager window (Figure 12–2). (If the Change Manager window is not visible at the bottom of your GUI, from the View menu, point to Utility Windows and click Change Manager.)

For more information about how to use the Change Manager, refer to the Engineering Change Management with the Chip Planner chapter in volume 2 of the Quartus II Handbook.

To view the timing results of each successfully routed SignalProbe pin, on the Processing menu, point to Start and click Start Timing Analysis.

SignalProbe Compilation Functions

After a full compilation, you can start a SignalProbe compilation either manually or automatically. A SignalProbe compilation performs the following functions:

- Validates SignalProbe pins
- Validates your specified SignalProbe sources
- Adds registers into SignalProbe paths, if applicable
- Attempts to route from SignalProbe sources through registers to SignalProbe pins

To run the SignalProbe compilation immediately after a full compilation, on the Tools menu, click SignalProbe Pins. In the SignalProbe Pins dialog box, click Start Check & Save All Nelist Changes.

To run a SignalProbe compilation manually after a full compilation, on the Processing menu, point to Start and click Start SignalProbe Compilation.

You must run the Fitter before a SignalProbe compilation. The Fitter generates a list of all internal nodes that can serve as SignalProbe sources.

Turn the SignalProbe enable option on or off in the SignalProbe Pins dialog box to enable or disable each SignalProbe pin.
Understanding the Results of a SignalProbe Compilation

After a SignalProbe compilation, the results appear in two sections of the compilation report file. The fitting results and status (Table 12–1) of each SignalProbe pin appears in the SignalProbe Fitting Result screen in the Fitter section of the Compilation Report (Figure 12–3).

<table>
<thead>
<tr>
<th>Status</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Routed</td>
<td>Connected and routed successfully</td>
</tr>
<tr>
<td>Not Routed</td>
<td>Not enabled</td>
</tr>
<tr>
<td>Failed to Route</td>
<td>Failed routing during last SignalProbe compilation</td>
</tr>
<tr>
<td>Need to Compile</td>
<td>Assignment changed since last SignalProbe compilation</td>
</tr>
</tbody>
</table>

The SignalProbe source to output delays screen in the Timing Analysis section of the Compilation Report displays the timing results of each successfully routed SignalProbe pin (Figure 12–4).
After a SignalProbe compilation, the processing screen of the Messages window also provides the results for each SignalProbe pin and displays slack information for each successfully routed SignalProbe pin.

Analyzing SignalProbe Routing Failures

A SignalProbe compilation can fail for any of the following reasons:

- **Route unavailable**—the SignalProbe compilation failed to find a route from the SignalProbe source to the SignalProbe pin because of routing congestion.
- **Invalid or nonexistent SignalProbe source**—you entered a SignalProbe source that does not exist or is invalid.
- **Unusable output pin**—the output pin selected is found to be unusable.

Routing failures can occur if the SignalProbe pin’s I/O standard conflicts with other I/O standards in the same I/O bank.

If routing congestion prevents a successful SignalProbe compilation, you can allow the compiler to modify routing to the specified SignalProbe source. On the Tools menu, click **SignalProbe Pins** and turn on **Modify latest fitting results during SignalProbe compilation**. This setting allows the Fitter to modify existing routing channels used by your design.

Turning on **Modify latest fitting results during SignalProbe compilation** can change the performance of your design.

Scripting Support

You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II command-line and Tcl API Help browser. To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```

The Tcl commands in this section are part of the `::quartus::chip_planner` Quartus II Tcl API. Source or include the `::quartus::chip_planner` Tcl package in your scripts to make these commands available.

For more information about Tcl scripting, refer to the *Tcl Scripting* chapter in volume 2 of the *Quartus II Handbook*. For more information about all settings and constraints in the Quartus II software, refer to the *Quartus II Settings File Reference Manual*. For more information about command-line scripting, refer to the *Command-Line Scripting* chapter in volume 2 of the *Quartus II Handbook*.

Making a SignalProbe Pin

To make a SignalProbe pin, type the following command:

```
-loc <loc> -pin_name <pin name> [-regs <regs>] [-reset <reset>] \ 
-src_name <source name>
```
Deleting a SignalProbe Pin

To delete a SignalProbe pin, type the following Tcl command:
```
delete_sp [-h | -help] [-long_help] -pin_name <pin name>
```

Enabling a SignalProbe Pin

To enable a SignalProbe pin, type the following Tcl command:
```
enable_sp [-h | -help] [-long_help] -pin_name <pin name>
```

Disabling a SignalProbe Pin

To disable a SignalProbe pin, type the following Tcl command:
```
disable_sp [-h | -help] [-long_help] -pin_name <pin name>
```

Performing a SignalProbe Compilation

To perform a SignalProbe compilation, type the following command:
```
quartus_sh --flow signalprobe <project name>
```

Script Example

Example 12–1 shows a script that creates a SignalProbe pin called sp1 and connects the sp1 pin to source node reg1 in a project that was already compiled.

```
Example 12–1. Creating a SignalProbe Pin Called sp1

package require ::quartus::chip_planner
project_open project
read_netlist
make_sp -pin_name sp1 -src_name reg1
check_netlist_and_save
project_close
```

Reserving SignalProbe Pins

To reserve a SignalProbe pin, add the commands shown in Example 12–2 to the Quartus II Settings File (.qsf) for your project.

```
Example 12–2. Reserving a SignalProbe Pin

set_location_assignment <location> -to <SignalProbe pin name>
set_instance_assignment -name RESERVE_PIN "AS SIGNALPROBE OUTPUT" -to <SignalProbe pin name>
```

Valid locations are pin location names, such as Pin_A3.

For more information about reserving SignalProbe pins, refer to “Reserving SignalProbe Pins” on page 12–2.
Common Problems When Reserving a SignalProbe Pin

If you cannot reserve a SignalProbe pin in the Quartus II software, it is likely that one of the following is true:

- You have selected multiple pins.
- A compilation is running in the background. Wait until the compilation is complete before reserving the pin.
- You have the Quartus II Web Edition software, in which the SignalProbe feature is not enabled by default. You must turn on TalkBack to enable the SignalProbe feature in the Quartus II Web Edition software.
- You have not set the pin reserve type to As Signal Probe Output. To reserve a pin, on the Assignments menu, in the Assign Pins dialog box, select As SignalProbe Output.
- The pin is reserved from a previous compilation. During a compilation, the Quartus II software reserves each pin on the targeted device. If you end the Quartus II process during a compilation, for example, with the Windows Task Manager End Process command or the UNIX kill command, perform a full recompilation before reserving pins as SignalProbe outputs.
- The pin does not support the SignalProbe feature. Select another pin.
- The current device family does not support the SignalProbe feature.

Adding SignalProbe Sources

Use the following Tcl commands to add SignalProbe sources.

To assign the node name to a SignalProbe pin, type the following Tcl command:

```
set_instance_assignment -name SIGNALPROBE_SOURCE <node name> -to \
<SignalProbe pin name>
```

The next command turns on SignalProbe routing. To turn off individual SignalProbe pins, specify OFF instead of ON with the following command:

```
set_instance_assignment -name SIGNALPROBE_ENABLE ON -to \
<SignalProbe pin name>
```

For more information about adding SignalProbe sources, refer to SignalProbe Pins Dialog Box and Add SignalProbe Pins Dialog Box in Quartus II Help.

Assigning I/O Standards

To assign an I/O standard to a pin, type the following Tcl command:

```
set_instance_assignment -name IO_STANDARD <I/O standard> -to \
<Signal Probe pin name>
```

For a list of valid I/O standards, refer I/O Standards to the in the Quartus II Help.
Adding Registers for Pipelining

To add registers for pipelining, type the following Tcl command:

```tcl
set_instance_assignment -name SIGNALPROBE_CLOCK <clock name> -to <SignalProbe pin name>
set_instance_assignment -name SIGNALPROBE_NUM_REGISTERS <number of registers> -to <SignalProbe pin name>
```

Running SignalProbe Immediately After a Full Compilation

To run SignalProbe immediately after a full compilation, type the following Tcl command:

```tcl
set_global_assignment -name SIGNALPROBE_DURING_NORMAL_COMPILATION ON
```

For more information about running SignalProbe automatically, refer to “SignalProbe Compilation Functions” on page 12–4.

Running SignalProbe Manually

To run SignalProbe as part of a scripted flow using Tcl, use the following in your script:

```tcl
execute_flow -signalprobe
```

To perform a Signal Probe compilation interactively at a command prompt, type the following command:

```bash
quartus_sh_fit --flow signalprobe <project name>
```

For more information about running SignalProbe manually, refer to “SignalProbe Compilation Functions” on page 12–4.

Enabling or Disabling All SignalProbe Routing

Use the Tcl command in Example 12–3 to turn on or turn off SignalProbe routing. When using this command, to turn SignalProbe routing on, specify ON. To turn SignalProbe routing off, specify OFF.

**Example 12–3. Turning SignalProbe On or Off with Tcl Commands**

```tcl
set spe [get_all_assignments -name SIGNALPROBE_ENABLE] 
foreach_in_collection asgn $spe {
    set signalprobe_pin_name [lindex $asgn 2]
    set_instance_assignment -name SIGNALPROBE_ENABLE -to <SignalProbe pin name> <ON|OFF> }
```

For more information about enabling or disabling SignalProbe routing, refer to page 12–4.

Allowing SignalProbe to Modify Fitting Results

To turn on Modify latest fitting results, type the following Tcl command:

```tcl
set_global_assignment -name SIGNALPROBE_ALLOW_OVERUSE ON
```
For more information, refer to “Analyzing SignalProbe Routing Failures” on page 12–6.

**Document Revision History**

Table 12–2 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tr>
<td>May 2013</td>
<td>13.0.0</td>
<td>Changed sequence of flow to clarify that you need to perform a full compilation before reserving SignalProbe pins. Affected sections are “Debugging Using the SignalProbe Feature” on page 12–1 and “Reserving SignalProbe Pins” on page 12–2. Moved “Performing a Full Compilation” on page 12–2 before “Reserving SignalProbe Pins” on page 12–2.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.0.2</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.0.1</td>
<td>Changed to new document template.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Revised for new UI.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed section SignalProbe ECO flows</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed support for SignalProbe pin preservation when recompiling with incremental compilation turned on.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed outdated FAQ section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added links to Quartus II Help for procedural content.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Removed all references and procedures for APEX devices.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Style changes.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Removed the “Generate the Programming File” section</td>
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<tr>
<td></td>
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<td>Removed unnecessary screenshots</td>
</tr>
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<td></td>
<td></td>
<td>Minor editorial updates</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Modified description for preserving SignalProbe connections when using Incremental Compilation</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added plausible scenarios where SignalProbe connections are not reserved in the design</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>Added “Arria GX” to the list of supported devices</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Removed the “On-Chip Debugging Tool Comparison” and replaced with a reference to the Section V Overview on page 13–1</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Added hyperlinks to referenced documents throughout the chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Minor editorial updates</td>
</tr>
</tbody>
</table>

For previous versions of the *Quartus II Handbook*, refer to the *Quartus II Handbook Archive*. 
13. Design Debugging Using the SignalTap II Logic Analyzer

Altera provides the SignalTap® II Logic Analyzer to help with the process of design debugging. This logic analyzer is a solution that allows you to examine the behavior of internal signals, without using extra I/O pins, while the design is running at full speed on an FPGA device.

The SignalTap II Logic Analyzer is scalable, easy to use, and is available as a stand-alone package or included with the Quartus® II software subscription. This logic analyzer helps debug an FPGA design by probing the state of the internal signals in the design without the use of external equipment. Defining custom trigger-condition logic provides greater accuracy and improves the ability to isolate problems. The SignalTap II Logic Analyzer does not require external probes or changes to the design files to capture the state of the internal nodes or I/O pins in the design. All captured signal data is conveniently stored in device memory until you are ready to read and analyze the data.

The topics in this chapter include:

- “Design Flow Using the SignalTap II Logic Analyzer” on page 13–5
- “SignalTap II Logic Analyzer Task Flow” on page 13–6
- “Configure the SignalTap II Logic Analyzer” on page 13–9
- “Define Triggers” on page 13–26
- “Compile the Design” on page 13–44
- “Program the Target Device or Devices” on page 13–48
- “Run the SignalTap II Logic Analyzer” on page 13–49
- “View, Analyze, and Use Captured Data” on page 13–54
- “Other Features” on page 13–60
- “Design Example: Using SignalTap II Logic Analyzers in SOPC Builder Systems” on page 13–65
- “Custom Triggering Flow Application Examples” on page 13–66
- “SignalTap II Scripting Support” on page 13–68
The SignalTap II Logic Analyzer is a next-generation, system-level debugging tool that captures and displays real-time signal behavior in a system-on-a-programmable-chip (SOPC) or any FPGA design. The SignalTap II Logic Analyzer supports the highest number of channels, largest sample depth, and fastest clock speeds of any logic analyzer in the programmable logic market. Figure 13–1 shows a block diagram of the components that make up the SignalTap II Logic Analyzer.

Figure 13–1. SignalTap II Logic Analyzer Block Diagram

Note to Figure 13–1:

(1) This diagram assumes that you compiled the SignalTap II Logic Analyzer with the design as a separate design partition using the Quartus II incremental compilation feature. This is the default setting for new projects in the Quartus II software. If incremental compilation is disabled or not used, the SignalTap II logic is integrated with the design. For information about the use of incremental compilation with SignalTap II, refer to “Faster Compilations with Quartus II Incremental Compilation” on page 13–44.

This chapter is intended for any designer who wants to debug an FPGA design during normal device operation without the need for external lab equipment. Because the SignalTap II Logic Analyzer is similar to traditional external logic analyzers, familiarity with external logic analyzer operations is helpful, but not necessary. To take advantage of faster compile times when making changes to the SignalTap II Logic Analyzer, knowledge of the Quartus II incremental compilation feature is helpful.

For information about using the Quartus II incremental compilation feature, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Hardware and Software Requirements

You need the following components to perform logic analysis with the SignalTap II Logic Analyzer:

- Quartus II design software
  
or
  Quartus II Web Edition (with the TalkBack feature enabled)
  
or
  SignalTap II Logic Analyzer standalone software, included in and requiring the Quartus II standalone Programmer software available from the Downloads page of the Altera website (www.altera.com)

- Download/upload cable

- Altera® development kit or your design board with JTAG connection to device under test

The Quartus II software Web Edition does not support the SignalTap II Logic Analyzer with the incremental compilation feature.

The memory blocks of the device store captured data and transfers the data to the Quartus II software waveform display with a JTAG communication cable, such as EthernetBlaster or USB-Blaster™. Table 13–1 summarizes features and benefits of the SignalTap II Logic Analyzer.

Table 13–1. SignalTap II Logic Analyzer Features and Benefits (Part 1 of 2)

<table>
<thead>
<tr>
<th>Feature</th>
<th>Benefit</th>
</tr>
</thead>
<tbody>
<tr>
<td>Multiple logic analyzers in a single device</td>
<td>Captures data from multiple clock domains in a design at the same time.</td>
</tr>
<tr>
<td>Multiple logic analyzers in multiple devices in a single JTAG chain</td>
<td>Simultaneously captures data from multiple devices in a JTAG chain.</td>
</tr>
<tr>
<td>Plug-In Support</td>
<td>Easily specifies nodes, triggers, and signal mnemonics for IP, such as the Nios® II processor.</td>
</tr>
<tr>
<td>Up to 10 basic or advanced trigger conditions for each analyzer instance</td>
<td>Enables sending more complex data capture commands to the logic analyzer, providing greater accuracy and problem isolation.</td>
</tr>
<tr>
<td>Power-Up Trigger</td>
<td>Captures signal data for triggers that occur after device programming, but before manually starting the logic analyzer.</td>
</tr>
<tr>
<td>State-based Triggering Flow</td>
<td>Enables you to organize your triggering conditions to precisely define what your logic analyzer captures.</td>
</tr>
<tr>
<td>Incremental compilation</td>
<td>Modifies the SignalTap II Logic Analyzer monitored signals and triggers without performing a full compilation, saving time.</td>
</tr>
<tr>
<td>Flexible buffer acquisition modes</td>
<td>The buffer acquisition control allows you to precisely control the data that is written into the acquisition buffer. Both segmented buffers and non-segmented buffers with storage qualification allow you to discard data samples that are not relevant to the debugging of your design.</td>
</tr>
<tr>
<td>MATLAB integration with included MEX function</td>
<td>Collects the SignalTap II Logic Analyzer captured data into a MATLAB integer matrix.</td>
</tr>
<tr>
<td>Up to 2,048 channels per logic analyzer instance</td>
<td>Samples many signals and wide bus structures.</td>
</tr>
<tr>
<td>Up to 128K samples in each device</td>
<td>Captures a large sample set for each channel.</td>
</tr>
</tbody>
</table>
The Quartus II software offers a portfolio of on-chip debugging solutions. For an overview and comparison of all tools available in the In-System Verification Tool set, refer to Section IV. In-System Design Debugging.

<table>
<thead>
<tr>
<th>Feature</th>
<th>Benefit</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fast clock frequencies</td>
<td>Synchronous sampling of data nodes using the same clock tree driving the logic under test.</td>
</tr>
<tr>
<td>Resource usage estimator</td>
<td>Provides estimate of logic and memory device resources used by SignalTap II Logic Analyzer configurations.</td>
</tr>
<tr>
<td>No additional cost</td>
<td>The SignalTap II Logic Analyzer is included with a Quartus II subscription and with the Quartus II Web Edition (with TalkBack enabled).</td>
</tr>
<tr>
<td>Compatibility with other on-chip debugging utilities</td>
<td>You can use the SignalTap II Logic Analyzer in tandem with any JTAG-based on-chip debugging tool, such as an In-System Memory Content editor, allowing you to change signal values in real-time while you are running an analysis with the SignalTap II Logic Analyzer.</td>
</tr>
</tbody>
</table>
Design Flow Using the SignalTap II Logic Analyzer

Figure 13–2 shows a typical overall FPGA design flow for using the SignalTap II Logic Analyzer in your design. A SignalTap II file (.stp) is added to and enabled in your project, or a SignalTap II HDL function, created with the MegaWizard™ Plug-In Manager, is instantiated in your design. The figure shows the flow of operations from initially adding the SignalTap II Logic Analyzer to your design to final device configuration, testing, and debugging.

Figure 13–2. SignalTap II FPGA Design and Debugging Flow
SignalTap II Logic Analyzer Task Flow

To use the SignalTap II Logic Analyzer to debug your design, you perform a number of tasks to add, configure, and run the logic analyzer. Figure 13–3 shows a typical flow of the tasks you complete to debug your design. Refer to the appropriate section of this chapter for more information about each of these tasks.

![SignalTap II Logic Analyzer Task Flow Diagram]

Add the SignalTap II Logic Analyzer to Your Design

Create an .stp or create a parameterized HDL instance representation of the logic analyzer using the MegaWizard Plug-In Manager. If you want to monitor multiple clock domains simultaneously, add additional instances of the logic analyzer to your design, limited only by the available resources in your device.

For information about creating an .stp, refer to Setting Up the SignalTap II Logic Analyzer in Quartus II Help.
Configure the SignalTap II Logic Analyzer

After you add the SignalTap II Logic Analyzer to your design, configure the logic analyzer to monitor the signals you want. You can manually add signals or use a plug-in, such as the Nios II processor plug-in, to quickly add entire sets of associated signals for a particular intellectual property (IP). You can also specify settings for the data capture buffer, such as its size, the method in which data is captured and stored, and the device memory type to use for the buffer in devices that support memory type selection.

For information about configuring the SignalTap II Logic Analyzer, refer to Setting Up the SignalTap II Logic Analyzer in Quartus II Help.

Define Trigger Conditions

The SignalTap II Logic Analyzer captures data continuously while the logic analyzer is running. To capture and store specific signal data, set up triggers that tell the logic analyzer under what conditions to stop capturing data. The SignalTap II Logic Analyzer allows you to define trigger conditions that range from very simple, such as the rising edge of a single signal, to very complex, involving groups of signals, extra logic, and multiple conditions. Power-Up Triggers allow you to capture data from trigger events occurring immediately after the device enters user-mode after configuration.

For information about defining trigger conditions, refer to Setting Up the SignalTap II Logic Analyzer in Quartus II Help.

Compile the Design

With the .stp configured and trigger conditions defined, compile your project as usual to include the logic analyzer in your design. Because you may need to change monitored signal nodes or adjust trigger settings frequently during debugging, Altera recommends that you use the incremental compilation feature built into the SignalTap II Logic Analyzer, along with Quartus II incremental compilation, to reduce recompile times.

For information about compiling your design, refer to Compiling a Design that Contains a SignalTap II Logic Analyzer in Quartus II Help.

Program the Target Device or Devices

When you debug a design with the SignalTap II Logic Analyzer, you can program a target device directly from the .stp without using the Quartus II Programmer. You can also program multiple devices with different designs and simultaneously debug them.

The SignalTap II Logic Analyzer supports all current Altera FPGA device families including Arria®, Cyclone®, HardCopy®, and Stratix® devices.

For instructions on programming devices in the Quartus II software, refer to Running the SignalTap II Logic Analyzer in Quartus II Help.
Run the SignalTap II Logic Analyzer

In normal device operation, you control the logic analyzer through the JTAG connection, specifying when to start looking for trigger conditions to begin capturing data. With Runtime or Power-Up Triggers, read and transfer the captured data from the on-chip buffer to the .stp for analysis.

For information about analyzing results from the SignalTap II Logic Analyzer, refer to Analyzing Data in the SignalTap II Logic Analyzer in Quartus II Help.

View, Analyze, and Use Captured Data

After you have captured data and read it into the .stp, that data is available for analysis and debugging. Set up mnemonic tables, either manually or with a plug-in, to simplify reading and interpreting the captured signal data. To speed up debugging, use the Locate feature in the SignalTap II node list to find the locations of problem nodes in other tools in the Quartus II software. Save the captured data for later analysis, or convert the data to other formats for sharing and further study.

For information about analyzing results from the SignalTap II Logic Analyzer, refer to Analyzing Data in the SignalTap II Logic Analyzer in Quartus II Help.

Embedding Multiple Analyzers in One FPGA

The SignalTap II Logic Analyzer Editor includes support for adding multiple logic analyzers by creating instances in the .stp. You can create a unique logic analyzer for each clock domain in the design.

For information about creating instances, refer to Running the SignalTap II Logic Analyzer in Quartus II Help.

Monitoring FPGA Resources Used by the SignalTap II Logic Analyzer

The SignalTap II Logic Analyzer has a built-in resource estimator that calculates the logic resources and amount of memory that each logic analyzer instance uses. Furthermore, because the most demanding on-chip resource for the logic analyzer is memory usage, the resource estimator reports the ratio of total RAM usage in your design to the total amount of RAM available, given the results of the last compilation. The resource estimator provides a warning if a potential for a “no-fit” occurs.

You can see resource usage of each logic analyzer instance and total resources used in the columns of the Instance Manager pane of the SignalTap II Logic Analyzer Editor. Use this feature when you know that your design is running low on resources.

The logic element value reported in the resource usage estimator may vary by as much as 10% from the actual resource usage.
Table 13–2 shows the SignalTap II Logic Analyzer M4K memory block resource usage for the listed devices per signal width and sample depth.

<table>
<thead>
<tr>
<th>Signals (Width)</th>
<th>Samples (Depth)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>256</td>
</tr>
<tr>
<td>8</td>
<td>&lt; 1</td>
</tr>
<tr>
<td>16</td>
<td>1</td>
</tr>
<tr>
<td>32</td>
<td>2</td>
</tr>
<tr>
<td>64</td>
<td>4</td>
</tr>
<tr>
<td>256</td>
<td>16</td>
</tr>
</tbody>
</table>

**Note to Table 13–2:**

(1) When you configure a SignalTap II Logic Analyzer, the Instance Manager reports an estimate of the memory bits and logic elements required to implement the given configuration.

### Using the MegaWizard Plug-In Manager to Create Your Logic Analyzer

You can create a SignalTap II Logic Analyzer instance by using the MegaWizard Plug-In Manager. The MegaWizard Plug-In Manager generates an HDL file that you instantiate in your design.

The State-based trigger flow, the state machine debugging feature, and the storage qualification feature are not supported when using the MegaWizard Plug-In Manager to create the logic analyzer. These features are described in the following sections:

- “Adding Finite State Machine State Encoding Registers” on page 13–14
- “Using the Storage Qualifier Feature” on page 13–18
- “State-Based Triggering” on page 13–30

For information about creating a SignalTap II instance with the MegaWizard Plug-In Manager, refer to *Setting Up the SignalTap II Logic Analyzer* in Quartus II Help.

### Configure the SignalTap II Logic Analyzer

There are many ways to configure instances of the SignalTap II Logic Analyzer. Some of the settings are similar to those found on traditional external logic analyzers. Other settings are unique to the SignalTap II Logic Analyzer because of the requirements for configuring a logic analyzer. All settings allow you to configure the logic analyzer the way you want to help debug your design.

Some settings can only be adjusted when you are viewing Run-Time Trigger conditions instead of Power-Up Trigger conditions. To learn about Power-Up Triggers and viewing different trigger conditions, refer to “Creating a Power-Up Trigger” on page 13–41.
Assigning an Acquisition Clock

Assign a clock signal to control the acquisition of data by the SignalTap II Logic Analyzer. The logic analyzer samples data on every positive (rising) edge of the acquisition clock. The logic analyzer does not support sampling on the negative (falling) edge of the acquisition clock. You can use any signal in your design as the acquisition clock. However, for best results, Altera recommends that you use a global, non-gated clock synchronous to the signals under test for data acquisition. Using a gated clock as your acquisition clock can result in unexpected data that does not accurately reflect the behavior of your design. The Quartus II static timing analysis tools show the maximum acquisition clock frequency at which you can run your design. Refer to the Timing Analysis section of the Compilation Report to find the maximum frequency of the logic analyzer clock.

For information about assigning an acquisition clock, refer to Working with Nodes in the SignalTap II Logic Analyzer in Quartus II Help.

Altera recommends that you exercise caution when using a recovered clock from a transceiver as an acquisition clock for the SignalTap II Logic Analyzer. Incorrect or unexpected behavior has been noted, particularly when a recovered clock from a transceiver is used as an acquisition clock with the power-up trigger feature.

If you do not assign an acquisition clock in the SignalTap II Logic Analyzer Editor, the Quartus II software automatically creates a clock pin called auto_stp_external_clk. You must make a pin assignment to this pin independently from the design. Ensure that a clock signal in your design drives the acquisition clock.

For information about assigning signals to pins, refer to the I/O Management chapter in volume 2 of the Quartus II Handbook.

Adding Signals to the SignalTap II File

While configuring the logic analyzer, add signals to the node list in the .stp to select which signals in your design you want to monitor. You can also select signals to define triggers. You can assign the following two types of signals to your .stp file:

- **Pre-synthesis**—These signals exist after design elaboration, but before any synthesis optimizations are done. This set of signals should reflect your Register Transfer Level (RTL) signals.
- **Post-fitting**—This signal exists after physical synthesis optimizations and place-and-route.

If you are not using incremental compilation, add only pre-synthesis signals to the .stp. Using pre-synthesis helps when you want to add a new node after you change a design. Source file changes appear in the Node Finder after you perform an Analysis and Elaboration. On the Processing Menu, point to Start and click Start Analysis & Elaboration.

For more information about incremental compilation, refer to About Incremental Compilation in Quartus II Help.
The Quartus II software does not limit the number of signals available for monitoring in the SignalTap II window waveform display. However, the number of channels available is directly proportional to the number of logic elements (LEs) or adaptive logic modules (ALMs) in the device. Therefore, there is a physical restriction on the number of channels that are available for monitoring. Signals shown in blue text are post-fit node names. Signals shown in black text are pre-synthesis node names. After successful Analysis and Elaboration, invalid signals are displayed in red. Unless you are certain that these signals are valid, remove them from the .stp for correct operation. The SignalTap II Status Indicator also indicates if an invalid node name exists in the .stp.

You can tap signals if a routing resource (row or column interconnects) exists to route the connection to the SignalTap II instance. For example, signals that exist in the I/O element (IOE) cannot be directly tapped because there are no direct routing resources from the signal in an IOE to a core logic element. For input pins, you can tap the signal that is driving a logic array block (LAB) from an IOE, or, for output pins, you can tap the signal from the LAB that is driving an IOE.

When adding pre-synthesis signals, make all connections to the SignalTap II Logic Analyzer before synthesis. Logic and routing resources are allocated during recompilation to make the connection as if a change in your design files had been made. Pre-synthesis signal names for signals driving to and from IOEs coincide with the signal names assigned to the pin.

In the case of post-fit signals, connections that you make to the SignalTap II Logic Analyzer are the signal names from the actual atoms in your post-fit netlist. You can only make a connection if the signals are part of the existing post-fit netlist and existing routing resources are available from the signal of interest to the SignalTap II Logic Analyzer. In the case of post-fit output signals, tap the COMBOUT or REGOUT signal that drives the IOE block. For post-fit input signals, signals driving into the core logic coincide with the signal name assigned to the pin.

Because NOT-gate push back applies to any register that you tap, the signal from the atom may be inverted. You can check this by locating the signal in either the Resource Property Editor or the Technology Map Viewer. The Technology Map viewer and the Resource Property Editor can also be used to help you find post-fit node names.

For information about cross-probing to source design files and other Quartus II windows, refer to the Analyzing Designs with Quartus II Netlist Viewers chapter in volume 1 of the Quartus II Handbook.

For more information about the use of incremental compilation with the SignalTap II Logic Analyzer, refer to “Faster Compilations with Quartus II Incremental Compilation” on page 13–44.

Signal Preservation

Many of the RTL signals are optimized during the process of synthesis and place-and-route. RTL signal names frequently may not appear in the post-fit netlist after optimizations. For example, the compilation process can add tildes (“~”) to nets that fan-out from a node, making it difficult to decipher which signal nets they actually represent. These process results can cause problems when you use the incremental compilation flow with the SignalTap II Logic Analyzer. Because you can
only add post-fitting signals to the SignalTap II Logic Analyzer in partitions of type **post-fit**, RTL signals that you want to monitor may not be available, preventing their use. To avoid this issue, use synthesis attributes to preserve signals during synthesis and place-and-route. When the Quartus II software encounters these synthesis attributes, it does not perform any optimization on the specified signals, forcing them to continue to exist in the post-fit netlist. However, if you do this, you could see an increase in resource utilization or a decrease in timing performance. The two attributes you can use are:

- **keep**—Ensures that combinational signals are not removed
- **preserve**—Ensures that registers are not removed

For more information about using these attributes, refer to the *Quartus II Integrated Synthesis* chapter in volume 1 of the *Quartus II Handbook*.

If you are debugging an IP core, such as the Nios II CPU or other encrypted IP, you might need to preserve nodes from the core to make them available for debugging with the SignalTap II Logic Analyzer. Preserving nodes is often necessary when a plug-in is used to add a group of signals for a particular IP.

If you use incremental compilation flow with the SignalTap II Logic Analyzer, pre-synthesis nodes may not be connected to the SignalTap II Logic Analyzer if the affected partition is of the post-fit type. A critical warning is issued for all pre-synthesis node names that are not found in the post-fit netlist.

For more information about node preservation or how to avoiding these warnings, refer to the *Working with Nodes in the SignalTap II Logic Analyzer* in Quartus II Help.

### Assigning Data Signals Using the Technology Map Viewer

You can easily add post-fit signal names that you find in the Technology map viewer. To do so, launch the Technology map viewer (post-fitting) after compiling your design. When you find the desired node, copy the node to either the active .stp for your design or a new .stp.

### Node List Signal Use Options

When a signal is added to the node list, you can select options that specify how the signal is used with the logic analyzer. You can turn off the ability of a signal to trigger the analyzer by disabling the **Trigger Enable** option for that signal in the node list in the .stp. This option is useful when you want to see only the captured data for a signal and you are not using that signal as part of a trigger.

You can turn off the ability to view data for a signal by disabling the **Data Enable** column. This option is useful when you want to trigger on a signal, but have no interest in viewing data for that signal.

For information about using signals in the node list to create SignalTap II trigger conditions, refer to “Define Triggers” on page 13–26.
Untappable Signals

Not all of the post-fitting signals in your design are available in the SignalTap II: post-fitting filter in the Node Finder dialog box. The following signal types cannot be tapped:

- **Post-fit output pins**—You cannot tap a post-fit output pin directly. To make an output signal visible, tap the register or buffer that drives the output pin. This includes pins defined as bidirectional.

- **Signals that are part of a carry chain**—You cannot tap the carry out (cout0 or cout1) signal of a logic element. Due to architectural restrictions, the carry out signal can only feed the carry in of another LE.

- **JTAGSignals**—You cannot tap the JTAG control (TCK, TDI, TDO, and TMS) signals.

- **ALTGXBMegafuntion**—You cannot directly tap any ports of an ALTGXB instantiation.

- **LVDS**—You cannot tap the data output from a serializer/deserializer (SERDES) block.

- **DQ, DQS Signals**—You cannot directly tap the DQ or DQS signals in a DDR/DDRII design.

Adding Signals with a Plug-In

Instead of adding individual or grouped signals through the Node Finder, you can add groups of relevant signals of a particular type of IP with a plug-in. The SignalTap II Logic Analyzer comes with one plug-in already installed for the Nios II processor. Besides easy signal addition, plug-ins also provide features such as pre-designed mnemonic tables, useful for trigger creation and data viewing, as well as the ability to disassemble code in captured data.

The Nios II plug-in, for example, creates one mnemonic table in the Setup tab and two tables in the Data tab:

- **Nios II Instruction (Setup tab)**—Capture all the required signals for triggering on a selected instruction address.

- **Nios II Instance Address (Data tab)**—Display address of executed instructions in hexadecimal format or as a programming symbol name if defined in an optional Executable and Linking Format (.elf) file.

- **Nios II Disassembly (Data tab)**—Displays disassembled code from the corresponding address.

For information about the other features plug-ins provide, refer to “Define Triggers” on page 13–26 and “View, Analyze, and Use Captured Data” on page 13–54.

To add signals to the .stp using a plug-in, perform the following steps after running Analysis and Elaboration on your design:

1. Right-click in the node list. On the Add Nodes with Plug-In submenu, choose the plug-in you want to use, such as the included plug-in named Nios II.

   > If the IP for the selected plug-in does not exist in your design, a message informs you that you cannot use the selected plug-in.
2. The **Select Hierarchy Level** dialog box appears showing the IP hierarchy of your design. Select the IP that contains the signals you want to monitor with the plug-in and click **OK**.

3. If all the signals in the plug-in are available, a dialog box might appear, depending on the plug-in selected, where you can specify options for the plug-in. With the Nios II plug-in, you can optionally select an `.elf` containing program symbols from your Nios II Integrated Development Environment (IDE) software design. Specify options for the selected plug-in as desired and click **OK**.

   To make sure all the required signals are available, in the Quartus II **Analysis & Synthesis** settings, turn on **Create debugging nodes for IP cores**.

   All the signals included in the plug-in are added to the node list.

**Adding Finite State Machine State Encoding Registers**

Finding the signals to debug Finite State Machines (FSM) can be challenging. Finding nodes from the post-fit netlist may be impossible, as FSM encoding signals may be changed or optimized away during synthesis and place-and-route. If you can find all of the relevant nodes in the post-fit netlist or you used the nodes from the pre-synthesis netlist, an additional step is required to find and map FSM signal values to the state names that you specified in your HDL.

The SignalTap II Logic Analyzer GUI can detect FSMs in your compiled design. The SignalTap II Logic Analyzer configuration automatically tracks the FSM state signals as well as state encoding through the compilation process. Shortcut menu commands from the SignalTap II Logic Analyzer GUI allow you to add all of the FSM state signals to your logic analyzer with a single command. For each FSM added to your SignalTap II configuration, the FSM debugging feature adds a mnemonic table to map the signal values to the state enumeration that you provided in your source code. The mnemonic tables enable you to visualize state machine transitions in the waveform viewer. The FSM debugging feature supports adding FSM signals from both the pre-synthesis and post-fit netlists.

*Figure 13–4* shows the waveform viewer with decoded signal values from a state machine added with the FSM debugging feature.

*Figure 13–4. Decoded FSM Mnemonics*

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For coding guidelines for specifying FSM in Verilog and VHDL, refer to the *Recommended HDL Coding Styles* chapter in volume 1 of the *Quartus II Handbook*.

For information about adding FSM signals to the configuration file, refer to *Setting Up the SignalTap II Logic Analyzer* in Quartus II Help.
Modifying and Restoring Mnemonic Tables for State Machines

When you add FSM state signals via the FSM debugging feature, the SignalTap II Logic Analyzer GUI creates a mnemonic table using the format `<StateSignalName>_table`, where `StateSignalName` is the name of the state signals that you have declared in your RTL. You can edit any mnemonic table using the Mnemonic Table Setup dialog box.

If you want to restore a mnemonic table that was modified, right-click anywhere in the node list window and select Recreate State Machine Mnemonics. By default, restoring a mnemonic table overwrites the existing mnemonic table that you modified. To restore a FSM mnemonic table to a new record, turn off Overwrite existing mnemonic table in the Recreate State Machine Mnemonics dialog box.

If you have added or deleted a signal from the FSM state signal group from within the setup tab, delete the modified register group and add the FSM signals back again.

For more information about using Mnemonics, refer to “Creating Mnemonics for Bit Patterns” on page 13–58.

Additional Considerations

The SignalTap II configuration GUI recognizes state machines from your design only if you use Quartus II Integrated Synthesis (QIS). The state machine debugging feature is not able to track the FSM signals or state encoding if you use other EDA synthesis tools.

If you add post-fit FSM signals, the SignalTap II Logic Analyzer FSM debug feature may not track all optimization changes that are a part of the compilation process. If the following two specific optimizations are enabled, the SignalTap II FSM debug feature may not list mnemonic tables for state machines in the design:

- If you have physical synthesis turned on, state registers may be resource balanced (register retiming) to improve fMAX. The FSM debug feature does not list post-fit FSM state registers if register retiming occurs.
- The FSM debugging feature does not list state signals that have been packed into RAM and DSP blocks during QIS or Fitter optimizations.

You can still use the FSM debugging feature to add pre-synthesis state signals.

Specifying the Sample Depth

The sample depth specifies the number of samples that are captured and stored for each signal in the captured data buffer. To specify the sample depth, select the desired number of samples to store in the Sample Depth list. The sample depth ranges from 0 to 128K.

If device memory resources are limited, you may not be able to successfully compile your design with the sample buffer size you have selected. Try reducing the sample depth to reduce resource usage.
Capturing Data to a Specific RAM Type

When you use the SignalTap II Logic Analyzer with some devices, you have the option to select the RAM type where acquisition data is stored. Once SignalTap II Logic Analyzer is allocated to a particular RAM block, the entire RAM block becomes a dedicated resource for the logic analyzer. RAM selection allows you to preserve a specific memory block for your design and allocate another portion of memory for SignalTap II Logic Analyzer data acquisition. For example, if your design has an application that requires a large block of memory resources, such as a large instruction or data cache, you would choose to use MLAB, M512, or M4k blocks for data acquisition and leave the M9k blocks for the rest of your design.

To select the RAM type to use for the SignalTap II Logic Analyzer buffer, select it from the RAM type list. Use this feature when the acquired data (as reported by the SignalTap II resource estimator) is not larger than the available memory of the memory type that you have selected in the FPGA.

Choosing the Buffer Acquisition Mode

The Buffer Acquisition Type Selection feature in the SignalTap II Logic Analyzer lets you choose how the captured data buffer is organized and can potentially reduce the amount of memory that is required for SignalTap II data acquisition. There are two types of acquisition buffer within the SignalTap II Logic Analyzer—a non-segmented buffer and a segmented buffer. With a non-segmented buffer, the SignalTap II Logic Analyzer treats entire memory space as a single FIFO, continuously filling the buffer until the logic analyzer reaches a defined set of trigger conditions. With a segmented buffer, the memory space is split into a number of separate buffers. Each buffer acts as a separate FIFO with its own set of trigger conditions. Only a single buffer is active during an acquisition. The SignalTap II Logic Analyzer advances to the next segment after the trigger condition or conditions for the active segment has been reached.

When using a non-segmented buffer, you can use the storage qualification feature to determine which samples are written into the acquisition buffer. Both the segmented buffers and the non-segmented buffer with the storage qualification feature help you maximize the use of the available memory space. Figure 13-5 illustrates the differences between the two buffer types.

Figure 13-5. Buffer Type Comparison in the SignalTap II Logic Analyzer (1), (2)

(1) Circular Buffer

(a) Circular Buffer

<table>
<thead>
<tr>
<th></th>
<th>Post-Trigger</th>
<th>Center Trigger</th>
<th>Pre-Trigger</th>
<th>Oldest Data Removed</th>
</tr>
</thead>
<tbody>
<tr>
<td>Newly Captured Data</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td></td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td></td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td></td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td></td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
</tbody>
</table>

(b) Segmented Buffer

<table>
<thead>
<tr>
<th></th>
<th>Segment 1</th>
<th>Segment 2</th>
<th>Segment 3</th>
<th>Segment 4</th>
</tr>
</thead>
<tbody>
<tr>
<td>All Trigger Level</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>1</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>Segment Trigger Level</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>1</td>
</tr>
</tbody>
</table>
Configure the SignalTap II Logic Analyzer

For more information about the storage qualification feature, refer to “Using the Storage Qualifier Feature” on page 13–18.

Non-Segmented Buffer

The non-segmented buffer (also known as a circular buffer) shown in Figure 13–5 (a) is the default buffer type used by the SignalTap II Logic Analyzer. While the logic analyzer is running, data is stored in the buffer until it fills up, at which point new data replaces the oldest data. This continues until a specified trigger event, consisting of a set of trigger conditions, occurs. When the trigger event happens, the logic analyzer continues to capture data after the trigger event until the buffer is full, based on the trigger position setting in the Signal Configuration pane in the .stp. To capture the majority of the data before the trigger occurs, select Post trigger position from the list. To capture the majority of the data after the trigger, select Pre-trigger position. To center the trigger position in the data, select Center trigger position. Alternatively, use the custom State-based triggering flow to define a custom trigger position within the capture buffer.

For more information, refer to “Specifying the Trigger Position” on page 13–41.

Segmented Buffer

A segmented buffer allows you to debug systems that contain relatively infrequent recurring events. The acquisition memory is split into evenly sized segments, with a set of trigger conditions defined for each segment. Each segment acts as a non-segmented buffer. If you want to have separate trigger conditions for each of the segmented buffers, you must use the state-based trigger flow. Figure 13–6 shows an example of a segmented buffer system.

Figure 13–6. Example System that Generates Recurring Events

Notes to Figure 13–5:

1. Both non-segmented and segmented buffers can use a predefined trigger (Pre-Trigger, Center Trigger, Post-Trigger) position or define a custom trigger position using the State-Based Triggering tab. Refer to “Specifying the Trigger Position” on page 13–41 for more details.

2. Each segment is treated like a FIFO, and behaves as the non-segmented buffer shown in (a).
The SignalTap II Logic Analyzer verifies the functionality of the design shown in Figure 13–6 to ensure that the correct data is written to the SRAM controller. Buffer acquisition in the SignalTap II Logic Analyzer allows you to monitor the \texttt{RDATA} port when \texttt{H'0F0F0F0F} is sent into the \texttt{RADDR} port. You can monitor multiple read transactions from the SRAM device without running the SignalTap II Logic Analyzer again. The buffer acquisition feature allows you to segment the memory so you can capture the same event multiple times without wasting allocated memory. The number of cycles that are captured depends on the number of segments specified under the Data settings.

To enable and configure buffer acquisition, select Segmented in the SignalTap II Logic Analyzer Editor and select the number of segments to use. In the example in Figure 13–6, selecting sixty-four 64-sample segments allows you to capture 64 read cycles when the \texttt{RADDR} signal is \texttt{H'0F0F0F0F}.

For more information about buffer acquisition mode, refer to Configuring the Trigger Flow in the SignalTap II Logic Analyzer in the Quartus II Help.

### Using the Storage Qualifier Feature

Both non-segmented and segmented buffers described in the previous section offer a snapshot in time of the data stream being analyzed. The default behavior for writing into acquisition memory with the SignalTap II Logic Analyzer is to sample data on every clock cycle. With a non-segmented buffer, there is one data window that represents a comprehensive snapshot of the data stream. Similarly, segmented buffers use several smaller sampling windows spread out over more time, with each sampling window representing a contiguous data set.

With carefully chosen trigger conditions and a generous sample depth for the acquisition buffer, analysis using segmented and non-segmented buffers captures a majority of functional errors in a chosen signal set. However, each data window can have a considerable amount of redundancy associated with it; for example, a capture of a data stream containing long periods of idle signals between data bursts. With default behavior using the SignalTap II Logic Analyzer, you cannot discard the redundant sample bits.

The Storage Qualification feature allows you to filter out individual samples not relevant to debugging the design. With this feature, a condition acts as a write enable to the buffer during each clock cycle of data acquisition. Through fine tuning the data that is actually stored in acquisition memory, the Storage Qualification feature allows for a more efficient use of acquisition memory in the specified number of samples over a longer period of analysis.

Use of the Storage Qualification feature is similar to an acquisition using a segmented buffer, in that you can create a discontinuity in the capture buffer. Because you can create a discontinuity between any two samples in the buffer, the Storage Qualification feature is equivalent to being able to create a customized segmented buffer in which the number and size of segment boundaries are adjustable. Figure 13–7 illustrates three ways the SignalTap II Logic Analyzer writes into acquisition memory.
You can only use the Storage Qualification feature with a non-segmented buffer. The MegaWizard Plug-In Manager instantiated flow only supports the Input Port mode for the Storage Qualification feature.

Figure 13–7. Data Acquisition Using Different Modes of Controlling the Acquisition Buffer

Notes to Figure 13–7:
(1) Non-segmented Buffers capture a fixed sample window of contiguous data.
(2) Segmented buffers divide the buffer into fixed sized segments, with each segment having an equal sample depth.
(3) Storage Qualification allows you to define a custom sampling window for each segment you create with a qualifying condition. Storage qualification potentially allows for a larger time scale of coverage.

There are six storage qualifier types available under the Storage Qualification feature:
- Continuous
- Input port
- Transitional
- Conditional
- Start/Stop
- State-based

Continuous (the default mode selected) turns the Storage Qualification feature off.

Each selected storage qualifier type is active when an acquisition starts. Upon the start of an acquisition, the SignalTap II Logic Analyzer examines each clock cycle and writes the data into the acquisition buffer based upon storage qualifier type and condition. The acquisition stops when a defined set of trigger conditions occur.
Trigger conditions are evaluated independently of storage qualifier conditions. The SignalTap II Logic Analyzer evaluates the data stream for trigger conditions on every clock cycle after the acquisition begins.

Trigger conditions are defined in “Define Trigger Conditions” on page 13–7.

The storage qualifier operates independently of the trigger conditions.

The following subsections describe each storage qualification mode from the acquisition buffer.

**Input Port Mode**

When using the Input port mode, the SignalTap II Logic Analyzer takes any signal from your design as an input. When the design is running, if the signal is high on the clock edge, the SignalTap II Logic Analyzer stores the data in the buffer. If the signal is low on the clock edge, the data sample is ignored. A pin is created and connected to this input port by default if no internal node is specified.

If you are using an .stp to create a SignalTap II Logic Analyzer instance, specify the storage qualifier signal using the input port field located on the Setup tab. You must specify this port for your project to compile.

If you use the MegaWizard Plug-In Manager flow, the storage qualification input port, if specified, appears in the MegaWizard-generated instantiation template. You can then connect this port to a signal in your RTL.

Figure 13–8 shows a data pattern captured with a segmented buffer. Figure 13–9 shows a capture of the same data pattern with the storage qualification feature enabled.

---

**Figure 13–8. Data Acquisition of a Recurring Data Pattern in Continuous Capture Mode (to illustrate Input port mode)**

| Node          | Type | Name          | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
| storage qualifier | data_out |            |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |   |

**Figure 13–9. Data Acquisition of a Recurring Data Pattern Using an Input Signal as a Storage Qualifier**

(1) Markers display samples when the logic analyzer paused a write into acquisition memory. These markers are enabled with the option “Record data discontinuities.”
Transitional Mode

In Transitional mode, you choose a set of signals for inspection using the node list check boxes in the Storage Qualifier column. During acquisition, if any of the signals marked for inspection have changed since the previous clock cycle, new data is written to the acquisition buffer. If none of the signals marked have changed since the previous clock cycle, no data is stored. Figure 13–10 shows the transitional storage qualifier setup. Figure 13–11 and Figure 13–12 show captures of a data pattern in continuous capture mode and a data pattern using the Transitional mode for storage qualification.

Figure 13–10. Transitional Storage Qualifier Setup

![Node List, Storage Enable, Transitional Enable, Storage Qualifier Dialog Box]

Figure 13–11. Data Acquisition of a Recurring Data Pattern in Continuous Capture Mode (to illustrate Transitional mode)

![Node List, Node, Name, Clock, Storage Enable, Transitional Enable, Storage Qualifier, Dialog Box]

Figure 13–12. Data Acquisition of Recurring Data Pattern Using a Transitional Mode as a Storage Qualifier

Redundant idle samples discarded

Conditional Mode

In Conditional mode, the SignalTap II Logic Analyzer evaluates a combinational function of storage qualifier enabled signals within the node list to determine whether a sample is stored. The SignalTap II Logic Analyzer writes into the buffer during the clock cycles in which the condition you specify evaluates TRUE.
You can select either **Basic AND**, **Basic OR**, or **Advanced** storage qualifier conditions. A **Basic AND** or **Basic OR** storage qualifier condition matches each signal to one of the following:
- Don’t Care
- Low
- High
- Falling Edge
- Rising Edge
- Either Edge

If you specify a **Basic AND** storage qualifier condition for more than one signal, the SignalTap II Logic Analyzer evaluates the logical AND of the conditions.

Any other combinational or relational operators that you may want to specify with the enabled signal set for storage qualification can be done with an advanced storage condition. Figure 13–13 details the conditional storage qualifier setup in the .stp.

You can specify storage qualification conditions similar to the manner in which trigger conditions are specified. For details about basic and advanced trigger conditions, refer to the sections “Creating Basic Trigger Conditions” on page 13–26 and “Creating Advanced Trigger Conditions” on page 13–27. Figure 13–14 and Figure 13–15 show a data capture with continuous sampling, and the same data pattern using the conditional mode for analysis, respectively.

**Figure 13–13. Conditional Storage Qualifier Setup**

**Figure 13–14. Data Acquisition of a Recurring Data Pattern in Continuous Capture Mode (to illustrate Conditional capture)**
Start/Stop Mode

The Start/Stop mode is similar to the Conditional mode for storage qualification. However, in this mode there are two sets of conditions, one for start and one for stop. If the start condition evaluates to TRUE, data begins is stored in the buffer every clock cycle until the stop condition evaluates to TRUE, which then pauses the data capture. Additional start signals received after the data capture has started are ignored. If both start and stop evaluate to TRUE at the same time, a single cycle is captured.

You can force a trigger by pressing the Stop button if the buffer fails to fill to completion due to a stop condition.

Figure 13–16 shows the Start/Stop mode storage qualifier setup. Figure 13–17 and Figure 13–18 show captures data pattern in continuous capture mode and a data pattern in using the Start/Stop mode for storage qualification.
State-Based
The State-based storage qualification mode is used with the State-based triggering flow. The state based triggering flow evaluates an if-else based language to define how data is written into the buffer. With the State-based trigger flow, you have command over boolean and relational operators to guide the execution flow for the target acquisition buffer. When the storage qualifier feature is enabled for the State-based flow, two additional commands are available, the `start_store` and `stop_store` commands. These commands operate similarly to the Start/Stop capture conditions described in the previous section. Upon the start of acquisition, data is not written into the buffer until a `start_store` action is performed. The `stop_store` command pauses the acquisition. If both `start_store` and `stop_store` actions are performed within the same clock cycle, a single sample is stored into the acquisition buffer.

For more information about the State-based flow and storage qualification using the State-based trigger flow, refer to the section “State-Based Triggering” on page 13–30.

Showing Data Discontinuities
When you turn on Record data discontinuities, the SignalTap II Logic Analyzer marks the samples during which the acquisition paused from a storage qualifier. This marker is displayed in the waveform viewer after acquisition completes.

Disable Storage Qualifier
You can turn off the storage qualifier quickly with the Disable Storage Qualifier option, and perform a continuous capture. This option is run-time reconfigurable; that is, the setting can be changed without recompiling the project. Changing storage qualifier mode from the Type field requires a recompilation of the project.

For a detailed explanation of Runtime Reconfigurable options available with the SignalTap II Logic Analyzer, and storage qualifier application examples using runtime reconfigurable options, refer to “Runtime Reconfigurable Options” on page 13–51.
Managing Multiple SignalTap II Files and Configurations

You may have more than one .stp in one design. Each file potentially has a different group of monitored signals. These signal groups make it possible to debug different blocks in your design. In turn, each group of signals can also be used to define different sets of trigger conditions. Along with each .stp, there is also an associated programming file (SRAM Object File [.sof]). The settings in a selected SignalTap II file must match the SignalTap II logic design in the associated .sof for the logic analyzer to run properly when the device is programmed. Use the Data Log feature and the SOF Manager to manage all of the .stp files and their associated settings and programming files.

The Data Log allows you to store multiple SignalTap II configurations within a single .stp. Figure 13–19 shows two signal set configurations with multiple trigger conditions in one .stp. To toggle between the active configurations, double-click on an entry in the Data Log. As you toggle between the different configurations, the signal list and trigger conditions change in the Setup tab of the .stp. The active configuration displayed in the .stp is indicated by the blue square around the signal specified in the Data Log. To store a configuration in the Data Log, on the Edit menu, click Save to Data Log or click Save to Data Log at the top of the Data Log.

Figure 13–19. Data Log

The SOF Manager allows you to embed multiple SOFs into one .stp. Embedding an SOF in an .stp lets you move the .stp to a different location, either on the same computer or across a network, without the need to include the associated .sof as a separate. To embed a new SOF in the .stp, right-click in the SOF Manager, and click Attach SOF File (Figure 13–20).

Figure 13–20. SOF Manager
As you switch between configurations in the Data Log, you can extract the SOF that is compatible with that particular configuration. You can use the programmer in the SignalTap II Logic Analyzer to download the new SOF to the FPGA, ensuring that the configuration of your .stp always matches the design programmed into the target device.

**Define Triggers**

When you start the SignalTap II Logic Analyzer, it samples activity continuously from the monitored signals. The SignalTap II Logic Analyzer “triggers”—that is, the logic analyzer stops and displays the data—when a condition or set of conditions that you specified has been reached. This section describes the various types of trigger conditions that you can specify using the SignalTap II Logic Analyzer on the Signal Configuration pane.

**Creating Basic Trigger Conditions**

The simplest kind of trigger condition is a basic trigger. Select this from the list at the top of the Trigger Conditions column in the node list in the SignalTap II Logic Analyzer Editor. If you select the Basic AND or Basic OR trigger type, you must specify the trigger pattern for each signal you have added in the .stp. To specify the trigger pattern, right-click in the Trigger Conditions column and click the desired pattern. Set the trigger pattern to any of the following conditions:

- Don’t Care
- Low
- High
- Falling Edge
- Rising Edge
- Either Edge

For buses, type a pattern in binary, or right-click and select Insert Value to enter the pattern in other number formats. Note that you can enter X to specify a set of “don’t care” values in either your hexadecimal or your binary string. For signals added to the .stp that have an associated mnemonic table, you can right-click and select an entry from the table to specify pre-defined conditions for the trigger.

For more information about creating and using mnemonic tables, refer to “View, Analyze, and Use Captured Data” on page 13–54, and to the Quartus II Help.

For signals added with certain plug-ins, you can create basic triggers easily using predefined mnemonic table entries. For example, with the Nios II plug-in, if you have specified an .elf from your Nios II IDE design, you can type the name of a function from your Nios II code. The logic analyzer triggers when the Nios II instruction address matches the address of the specified code function name.

Data capture stops and the data is stored in the buffer when the logical AND of all the signals for a given trigger condition evaluates to TRUE.
Chapter 13: Design Debugging Using the SignalTap II Logic Analyzer

Define Triggers

Creating Advanced Trigger Conditions

With the basic triggering capabilities of the SignalTap II Logic Analyzer, you can build more complex triggers with extra logic that enables you to capture data when a combination of conditions exist. If you select the Advanced trigger type at the top of the Trigger Conditions column in the node list of the SignalTap II Logic Analyzer Editor, a new tab named Advanced Trigger appears where you can build a complex trigger expression using a simple GUI. Drag-and-drop operators into the Advanced Trigger Configuration Editor window to build the complex trigger condition in an expression tree. To configure the operators’ settings, double-click or right-click the operators that you have placed and select Properties. Table 13–3 lists the operators you can use.

Table 13–3. Advanced Triggering Operators

<table>
<thead>
<tr>
<th>Name of Operator</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>Less Than</td>
<td>Comparison</td>
</tr>
<tr>
<td>Less Than or Equal To</td>
<td>Comparison</td>
</tr>
<tr>
<td>Equality</td>
<td>Comparison</td>
</tr>
<tr>
<td>Inequality</td>
<td>Comparison</td>
</tr>
<tr>
<td>Greater Than</td>
<td>Comparison</td>
</tr>
<tr>
<td>Greater Than or Equal To</td>
<td>Comparison</td>
</tr>
<tr>
<td>Logical NOT</td>
<td>Logical</td>
</tr>
<tr>
<td>Logical AND</td>
<td>Logical</td>
</tr>
<tr>
<td>Logical OR</td>
<td>Logical</td>
</tr>
<tr>
<td>Logical XOR</td>
<td>Logical</td>
</tr>
<tr>
<td>Reduction AND</td>
<td>Reduction</td>
</tr>
<tr>
<td>Reduction OR</td>
<td>Reduction</td>
</tr>
<tr>
<td>Reduction XOR</td>
<td>Reduction</td>
</tr>
<tr>
<td>Left Shift</td>
<td>Shift</td>
</tr>
<tr>
<td>Right Shift</td>
<td>Shift</td>
</tr>
<tr>
<td>Bitwise Complement</td>
<td>Bitwise</td>
</tr>
<tr>
<td>Bitwise AND</td>
<td>Bitwise</td>
</tr>
<tr>
<td>Bitwise OR</td>
<td>Bitwise</td>
</tr>
<tr>
<td>Bitwise XOR</td>
<td>Bitwise</td>
</tr>
<tr>
<td>Edge and Level Detector</td>
<td>Signal Detection</td>
</tr>
</tbody>
</table>

Note to Table 13–3:

(1) For more information about each of these operators, refer to the Quartus II Help.

Adding many objects to the Advanced Trigger Condition Editor can make the work space cluttered and difficult to read. To keep objects organized while you build your advanced trigger condition, use the shortcut menu and select Arrange All Objects. You can also use the Zoom-Out command to fit more objects into the Advanced Trigger Condition Editor window.
Examples of Advanced Triggering Expressions

The following examples show how to use Advanced Triggering:

- Trigger when bus outa is greater than or equal to outb (Figure 13–21).

Figure 13–21. Bus outa is Greater Than or Equal to Bus outb

- Trigger when bus outa is greater than or equal to bus outb, and when the enable signal has a rising edge (Figure 13–22).

Figure 13–22. Enable Signal has a Rising Edge
Trigger when bus `outa` is greater than or equal to bus `outb`, or when the enable signal has a rising edge. Or, when a bitwise AND operation has been performed between bus `outc` and bus `outd`, and all bits of the result of that operation are equal to 1 (Figure 13–23).

**Figure 13–23. Bitwise AND Operation**

\[
\text{Result: } outa \geq outb \lor \text{RED}((\text{enable}) \land (outc \land outd))
\]

---

**Trigger Condition Flow Control**

The SignalTap II Logic Analyzer offers multiple triggering conditions to give you precise control of the method in which data is captured into the acquisition buffers. Trigger Condition Flow allows you to define the relationship between a set of triggering conditions. The SignalTap II Logic Analyzer **Signal Configuration** pane offers two flow control mechanisms for organizing trigger conditions:

- **Sequential Triggering**—The default triggering flow. Sequential triggering allows you to define up to 10 triggering levels that must be satisfied before the acquisition buffer finishes capturing.
- **State-Based Triggering**—Allows you the greatest control over your acquisition buffer. Custom-based triggering allows you to organize trigger conditions into states based on a conditional flow that you define.

You can use sequential or state based triggering with either a segmented or a non-segmented buffer.

**Sequential Triggering**

Sequential triggering flow allows you to cascade up to 10 levels of triggering conditions. The SignalTap II Logic Analyzer sequentially evaluates each of the triggering conditions. When the last triggering condition evaluates to TRUE, the SignalTap II Logic Analyzer triggers the acquisition buffer. For segmented buffers, every acquisition segment after the first segment triggers on the last triggering condition that you have specified. Use the Simple Sequential Triggering feature with basic triggers, advanced triggers, or a mix of both. **Figure 13–24** illustrates the simple sequential triggering flow for non-segmented and segmented buffers.
The external trigger is considered as trigger level 0. The external trigger must be evaluated before the main trigger levels are evaluated.

**Figure 13–24. Sequential Triggering Flow**

![Sequential Triggering Flow Diagram]

**Notes to Figure 13–24:**

1. The acquisition buffer stops capture when all $n$ triggering levels are satisfied, where $n \leq 10$.
2. An external trigger input, if defined, is evaluated before all other defined trigger conditions are evaluated. For more information about external triggers, refer to “Using External Triggers” on page 13–43.

To configure the SignalTap II Logic Analyzer for Sequential triggering, in the SignalTap II editor on the Trigger flow control list, select **Sequential**. Select the desired number of trigger conditions from the Trigger Conditions list. After you select the desired number of trigger conditions, configure each trigger condition in the node list. To disable any trigger condition, turn on the trigger condition at the top of the column in the node list.

**State-Based Triggering**

Custom State-based triggering provides the most control over triggering condition arrangement. The State-Based Triggering flow allows you to describe the relationship between triggering conditions precisely, using an intuitive GUI and the SignalTap II Trigger Flow Description Language, a simple description language based upon conditional expressions. Tooltips within the custom triggering flow GUI allow you to describe your desired flow quickly. The custom State-based triggering flow allows for more efficient use of the space available in the acquisition buffer because only specific samples of interest are captured.
Figure 13–25 illustrates the custom State-based triggering flow. Events that trigger the acquisition buffer are organized by a state diagram that you define. All actions performed by the acquisition buffer are captured by the states and all transition conditions between the states are defined by the conditional expressions that you specify within each state.

Each state allows you to define a set of conditional expressions. Each conditional expression is a Boolean expression dependent on a combination of triggering conditions (configured within the Setup tab), counters, and status flags. Counters and status flags are resources provided by the SignalTap II Logic Analyzer custom-based triggering flow.

Within each conditional expression you define a set of actions. Actions include triggering the acquisition buffer to stop capture, a modification to either a counter or status flag, or a state transition.

Trigger actions can apply to either a single segment of a segmented acquisition buffer or to the entire non-segmented acquisition buffer. Each trigger action provides you with an optional count that specifies the number of samples captured before stopping acquisition of the current segment. The count argument allows you to control the amount of data captured precisely before and after triggering event.

Resource manipulation actions allow you to increment and decrement counters or set and clear status flags. The counter and status flag resources are used as optional inputs in conditional expressions. Counters and status flags are useful for counting the number of occurrences of particular events and for aiding in triggering flow control.

This SignalTap II custom State-based triggering flow allows you to capture a sequence of events that may not necessarily be contiguous in time; for example, capturing a communication transaction between two devices that includes a handshaking protocol containing a sequence of acknowledgements.

Notes to Figure 13–25:

(1) You are allowed up to 20 different states.
(2) An external trigger input, if defined, is evaluated before any conditions in the custom State-based triggering flow are evaluated. For more information, refer to “Using External Triggers” on page 13–43.
The State-Based Trigger Flow tab is the control interface for the custom state-based triggering flow. To enable this tab, select **State-based** on the Trigger Flow Control list. (Note that when Trigger Flow Control is specified as **Sequential**, the State-Based Trigger Flow tab is hidden.)

The State-Based Trigger Flow tab is partitioned into the following three panes:
- State Diagram Pane
- Resources Pane
- State Machine Pane

**State Diagram Pane**

The State Diagram pane provides a graphical overview of the triggering flow that you define. It shows the number of states available and the state transitions between the states. You can adjust the number of available states by using the menu above the graphical overview.

**State Machine Pane**

The State Machine pane contains the text entry boxes where you can define the triggering flow and actions associated with each state. You can define the triggering flow using the SignalTap II Trigger Flow Description Language, a simple language based on “if-else” conditional statements. Tooltips appear when you move the mouse over the cursor, to guide command entry into the state boxes. The GUI provides a syntax check on your flow description in real-time and highlights any errors in the text flow.

For a full description of the SignalTap II Trigger Flow Description Language, refer to “SignalTap II Trigger Flow Description Language” on page 13–33.

You can also refer to **SignalTap II Trigger Flow Description Language** in Quartus II Help.

The State Machine description text boxes default to show one text box per state. You can also have the entire flow description shown in a single text field. This option can be useful when copying and pasting a flow description from a template or an external text editor. To toggle between one window per state, or all states in one window, select the appropriate option under **State Display mode**.

**Resources Pane**

The Resources pane allows you to declare Status Flags and Counters for use in the conditional expressions in the Custom Triggering Flow. Actions to decrement and increment counters or to set and clear status flags are performed within the triggering flow that you define.

You can specify up to 20 counters and 20 status flags. Counter and status flags values may be initialized by right-clicking the status flag or counter name after selecting a number of them from the respective pull-down list, and selecting **Set Initial Value**. To specify a counter width, right-click the counter name and select **Set Width**. Counters and flag values are updated dynamically after acquisition has started to assist in debugging your trigger flow specification.
The configurable at runtime options in the Resources pane allows you to configure the custom-flow control options that can be changed at runtime without requiring a recompilation. Table 13–4 contains a description of options for the State-based trigger flow that can be reconfigured at runtime.

For a broader discussion about all options that can be changed without incurring a recompile refer to “Runtime Reconfigurable Options” on page 13–51.

### Table 13–4. Runtime Reconfigurable Settings, State-Based Triggering Flow

<table>
<thead>
<tr>
<th>Setting</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Destination of goto action</td>
<td>Allows you to modify the destination of the state transition at runtime.</td>
</tr>
<tr>
<td>Comparison values</td>
<td>Allows you to modify comparison values in Boolean expressions at runtime. In addition, you can modify the segment_trigger and trigger action post-fill count argument at runtime.</td>
</tr>
<tr>
<td>Comparison operators</td>
<td>Allows you to modify the operators in Boolean expressions at runtime.</td>
</tr>
<tr>
<td>Logical operators</td>
<td>Allows you to modify the logical operators in Boolean expressions at runtime.</td>
</tr>
</tbody>
</table>

You can restrict changes to your SignalTap configuration to include only the options that do not require a recompilation by using the menu above the trigger list in the Setup tab. Allow trigger condition changes only restricts changes to only the configuration settings that have the configurable at runtime specified. With this option enabled, to modify Trigger Flow conditions in the Custom Trigger Flow tab, click the desired parameter in the text box and select a new parameter from the menu that appears.

The runtime configurable settings for the Custom Trigger Flow tab are on by default. You may get some performance advantages by disabling some of the runtime configurable options. For details about the effects of turning off the runtime modifiable options, refer to “Performance and Resource Considerations” on page 13–47.

### SignalTap II Trigger Flow Description Language

The Trigger Flow Description Language is based on a list of conditional expressions per state to define a set of actions. Each line in Example 13–1 shows a language format. Keywords are shown in bold. Non-terminals are delimited by “<>” and are further explained in the following sections. Optional arguments are delimitred by “[]” (Example 13–1).
Examples of Triggering Flow descriptions for common scenarios using the SignalTap II Custom Triggering Flow are provided in “Custom Triggering Flow Application Examples” on page 13–66.

**Example 13–1. Trigger Flow Description Language Format**

```
state <State_label>:
  <action_list>
  if( <Boolean_expression> )
  <action_list>
  [else_if ( <boolean_expression> )
   <action_list>] (1)
  [else
   <action_list>]
```

**Note to Example 13–1:**
(1) Multiple else if conditions are allowed.

The priority for evaluation of conditional statements is assigned from top to bottom. The `<boolean_expression>` in an if statement can contain a single event, or it can contain multiple event conditions. The action_list within an if or an else if clause must be delimited by the begin and end tokens when the action list contains multiple statements. When the boolean expression is evaluated TRUE, the logic analyzer analyzes all of the commands in the action list concurrently. The possible actions include:

- Triggering the acquisition buffer
- Manipulating a counter or status flag resource
- Defining a state transition

**State Labels**

State labels are identifiers that can be used in the action goto.

```
state <state_label>: begins the description of the actions evaluated when this state is reached.
```

The description of a state ends with the beginning of another state or the end of the whole trigger flow description.

**Boolean_expression**

Boolean_expression is a collection of logical operators, relational operators, and their operands that evaluate into a Boolean result. Depending on the operator, the operand can be a reference to a trigger condition, a counter and a register, or a numeric value. Within an expression, parentheses can be used to group a set of operands.
Logical operators accept any boolean expression as an operand. The supported logical operators are shown in Table 13–5.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Syntax</th>
</tr>
</thead>
<tbody>
<tr>
<td>!</td>
<td>NOT operator</td>
<td>! expr1</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>AND operator</td>
<td>expr1 &amp;&amp; expr2</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Relational operators are performed on counters or status flags. The comparison value, the right operator, must be a numerical value. The supported relational operators are shown in Table 13–6.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Syntax (1) (2)</th>
</tr>
</thead>
<tbody>
<tr>
<td>&gt;</td>
<td>Greater than</td>
<td>&lt;identifier&gt; &gt; &lt;numerical_value&gt;</td>
</tr>
<tr>
<td>&gt;=</td>
<td>Greater than or Equal to</td>
<td>&lt;identifier&gt; &gt;= &lt;numerical_value&gt;</td>
</tr>
<tr>
<td>==</td>
<td>Equals</td>
<td>&lt;identifier&gt; == &lt;numerical_value&gt;</td>
</tr>
<tr>
<td>!=</td>
<td>Does not equal</td>
<td>&lt;identifier&gt; != &lt;numerical_value&gt;</td>
</tr>
<tr>
<td>&lt;=</td>
<td>Less than or equal to</td>
<td>&lt;identifier&gt; &lt;= &lt;numerical_value&gt;</td>
</tr>
<tr>
<td>&lt;</td>
<td>Less than</td>
<td>&lt;identifier&gt; &lt; &lt;numerical_value&gt;</td>
</tr>
</tbody>
</table>

Notes to Table 13–6:
(1) <identifier> indicates a counter or status flag.
(2) <numerical_value> indicates an integer.

**Action_list**

Action_list is a list of actions that can be performed when a state is reached and a condition is also satisfied. If more than one action is specified, they must be enclosed by begin and end. The actions can be categorized as resource manipulation actions, buffer control actions, and state transition actions. Each action is terminated by a semicolon (;).

**Resource Manipulation Action**

The resources used in the trigger flow description can be either counters or status flags. Table 13–7 shows the description and syntax of each action.

<table>
<thead>
<tr>
<th>Action</th>
<th>Description</th>
<th>Syntax</th>
</tr>
</thead>
<tbody>
<tr>
<td>increment</td>
<td>Increments a counter resource by 1</td>
<td>increment &lt;counter_identifier&gt;;;</td>
</tr>
<tr>
<td>decrement</td>
<td>Decrements a counter resource by 1</td>
<td>decrement &lt;counter_identifier&gt;;;</td>
</tr>
<tr>
<td>reset</td>
<td>Resets counter resource to initial value</td>
<td>reset &lt;counter_identifier&gt;;;</td>
</tr>
<tr>
<td>set</td>
<td>Sets a status Flag to 1</td>
<td>set &lt;register_flag_identifier&gt;;;</td>
</tr>
<tr>
<td>clear</td>
<td>Sets a status Flag to 0</td>
<td>clear &lt;register_flag_identifier&gt;;;</td>
</tr>
</tbody>
</table>
Buffer Control Action

Buffer control actions specify an action to control the acquisition buffer. Table 13–8 shows the description and syntax of each action.

Table 13–8. Buffer Control Action

<table>
<thead>
<tr>
<th>Action</th>
<th>Description</th>
<th>Syntax</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger</td>
<td>Stops the acquisition for the current buffer and ends analysis. This command is required in every flow definition.</td>
<td>trigger &lt;post-fill_count&gt;;</td>
</tr>
<tr>
<td>segment_trigger</td>
<td>Ends the acquisition of the current segment. The SignalTap II Logic Analyzer starts acquiring from the next segment on evaluating this command. If all segments are filled, the oldest segment is overwritten with the latest sample. The acquisition stops when a trigger action is evaluated. This action cannot be used in non-segmented acquisition mode.</td>
<td>segment_trigger &lt;post-fill_count&gt;;</td>
</tr>
<tr>
<td>start_store</td>
<td>Asserts the write_enable to the SignalTap II acquisition buffer. This command is active only when the State-based storage qualifier mode is enabled.</td>
<td>start_store</td>
</tr>
<tr>
<td>stop_store</td>
<td>De-asserts the write_enable signal to the SignalTap II acquisition buffer. This command is active only when the State-based storage qualifier mode is enabled.</td>
<td>stop_store</td>
</tr>
</tbody>
</table>

Both trigger and segment_trigger actions accept an optional post-fill count argument. If provided, the current acquisition acquires the number of samples provided by post-fill count and then stops acquisition. If no post-count value is specified, the trigger position for the affected buffer defaults to the trigger position specified in the Setup tab.

In the case of segment_trigger, acquisition of the current buffer stops immediately if a subsequent triggering action is issued in the next state, regardless of whether or not the post-fill count has been satisfied for the current buffer. The remaining unfilled post-count acquisitions in the current buffer are discarded and displayed as grayed-out samples in the data window.

State Transition Action

The State Transition action specifies the next state in the custom state control flow. It is specified by the goto command. The syntax is as follows:

```
goto <state_label>;
```

Using the State-Based Storage Qualifier Feature

When you select State-based for the storage qualifier type, the start_store and stop_store actions are enabled in the State-based trigger flow. These commands, when used in conjunction with the expressions of the State-based trigger flow, give you maximum flexibility to control data written into the acquisition buffer.
The start_store and stop_store commands can only be applied to a non-segmented buffer.

The start_store and stop_store commands function similar to the start and stop conditions when using the start/stop storage qualifier mode conditions. If storage qualification is enabled, the start_store command must be issued for SignalTap II to write data into the acquisition buffer. No data is acquired until the start_store command is performed. Also, a trigger command must be included as part of the trigger flow description. The trigger command is necessary to complete the acquisition and display the results on the waveform display.

The following examples illustrate the behavior of the State-based trigger flow with the storage qualification commands.

Figure 13–26 shows a hypothetical scenario with three trigger conditions that happen at different times after you click Start Analysis. The trigger flow description in Example 13–2, when applied to the scenario shown in Figure 13–26, illustrates the functionality of the storage qualification feature for the state-based trigger flow.

**Example 13–2. Trigger Flow Description 1**

State 1: ST1:
```
if ( condition1 )
  start_store;
else if ( condition2 )
  trigger value;
else if ( condition3 )
  stop_store;
```

**Figure 13–26. Capture Scenario for Storage Qualification with the State-Based Trigger Flow**

In this example, the SignalTap II Logic Analyzer does not write into the acquisition buffer until sample a, when Condition 1 occurs. Once sample b is reached, the trigger value command is evaluated. The logic analyzer continues to write into the buffer to finish the acquisition. The trigger flow specifies a stop_store command at sample c, m samples after the trigger point occurs.

The logic analyzer finishes the acquisition and displays the contents of the waveform if it can successfully finish the post-fill acquisition samples before Condition 3 occurs. In this specific case, the capture ends if the post-fill count value is less than m.
If the post-fill count value specified in Trigger Flow description 1 is greater than \( m \) samples, the buffer pauses acquisition indefinitely, provided there is no recurrence of Condition 1 to trigger the logic analyzer to start capturing data again. The SignalTap II Logic Analyzer continues to evaluate the `stop_store` and `start_store` commands even after the `trigger` command is evaluated. If the acquisition has paused, you can click **Stop Analysis** to manually stop and force the acquisition to trigger. You can use counter values, flags, and the State diagram to help you perform the trigger flow. The counter values, flags, and the current state are updated in real-time during a data acquisition.

**Figure 13–27** and **Figure 13–28** show a real data acquisition of the scenario. **Figure 13–27** illustrates a scenario where the data capture finishes successfully. It uses a buffer with a sample depth of 64, \( m = n = 10 \), and the post-fill count value = 5. **Figure 13–28** illustrates a scenario where the logic analyzer pauses indefinitely even after a trigger condition occurs due to a `stop_store` condition. This scenario uses a sample depth of 64, with \( m = n = 10 \) and post-fill count = 15.

**Figure 13–27. Storage Qualification with Post-Fill Count Value Less than \( m \) (Acquisition Successfully Completes)**
Figure 13–28. Storage Qualification with Post-Fill Count Value Greater than $m$ (Acquisition Indefinitely Paused)

Figure 13–29. Waveform After Forcing the Analysis to Stop
The combination of using counters, Boolean and relational operators in conjunction with the `start_store` and `stop_store` commands can give a clock-cycle level of resolution to controlling the samples that are written into the acquisition buffer. Example 13–3 shows a trigger flow description that skips three clock cycles of samples after hitting condition 1. Figure 13–30 shows the data transaction on a continuous capture and Figure 13–32 shows the data capture with the Trigger flow description in Example 13–3 applied.

Example 13–3. Trigger Flow Description 2

```
State 1: ST1
start_store
if (condition1)
begin
  stop_store;
  goto ST2;
end

State 2: ST2
if (c1 < 3)
  increment c1; //skip three clock cycles; c1 initialized to 0
else if (c1 == 3)
  begin
    start_store; //start_store necessary to enable writing to finish
                 //acquisition
    trigger;
  end
```

Figure 13–30. Continuous Capture of Data Transaction for Example 2

![Figure 13–30](image)

Figure 13–31. Capture of Data Transaction with Trigger Flow Description Applied

![Figure 13–31](image)
Specifying the Trigger Position

The SignalTap II Logic Analyzer allows you to specify the amount of data that is acquired before and after a trigger event. You can specify the trigger position independently between a Runtime and Power-Up Trigger. Select the desired ratio of pre-trigger data to post-trigger data by choosing one of the following ratios:

- **Pre**—Saves signal activity that occurred after the trigger (12% pre-trigger, 88% post-trigger).
- **Center**—Saves 50% pre-trigger and 50% post-trigger data.
- **Post**—Saves signal activity that occurred before the trigger (88% pre-trigger, 12% post-trigger).

These pre-defined ratios apply to both non-segmented buffers and segmented buffers.

If you use the custom-state based triggering flow, you can specify a custom trigger position. The `segment_trigger` and trigger actions accept a post-fill count argument. The post-fill count specifies the number of samples to capture before stopping data acquisition for the non-segmented buffer or a data segment when using the `trigger` and `segment_trigger` commands, respectively. When the captured data is displayed in the SignalTap II data window, the trigger position appears as the number of post-count samples from the end of the acquisition segment or buffer. Refer to Equation 13–1:

**Equation 13–1.**

\[
\text{Sample Number of Trigger Position} = (N - \text{Post-Fill Count})
\]

In this case, \(N\) is the sample depth of either the acquisition segment or non-segmented buffer.

For segmented buffers, the acquisition segments that have a post-count argument define use of the post-count setting. Segments that do not have a post-count setting default to the trigger position ratios defined in the **Setup** tab.

For more details about the custom State-based triggering flow, refer to “State-Based Triggering” on page 13–30.

Creating a Power-Up Trigger

Typically, the SignalTap II Logic Analyzer is used to trigger on events that occur during normal device operation. You start an analysis manually once the target device is fully powered on and the JTAG connection for the device is available. However, there may be cases when you would like to capture trigger events that occur during device initialization, immediately after the FPGA is powered on or reset. With the SignalTap II Power-Up Trigger feature, you arm the SignalTap II Logic Analyzer and capture data immediately after device programming.
Enabling a Power-Up Trigger

You can add a different Power-Up Trigger to each logic analyzer instance in the **SignalTap II Instance Manager** pane. To enable the Power-Up Trigger for a logic analyzer instance, right-click the instance and click **Enable Power-Up Trigger**, or select the instance, and on the **Edit** menu, click **Enable Power-Up Trigger**. To disable a Power-Up Trigger, click **Disable Power-Up Trigger** in the same locations. Power-Up Trigger is shown as a child instance below the name of the selected instance with the default trigger conditions specified in the node list. Figure 13–32 shows the SignalTap II Logic Analyzer Editor when Power-Up Trigger is enabled.

![Figure 13–32. SignalTap II Logic Analyzer Editor with Power-Up Trigger Enabled](image)

Managing and Configuring Power-Up and Runtime Trigger Conditions

When the Power-Up Trigger is enabled for a logic analyzer instance, you can create basic and advanced trigger conditions for the trigger as you do with a Run-Time Trigger. Power-Up Trigger conditions that you can adjust are color coded light blue, while Run-Time Trigger conditions you cannot adjust remain white. Since each instance now has two sets of trigger conditions—the Power-Up Trigger and the Run-Time Trigger—you can differentiate between the two with color coding. To switch between the trigger conditions of the Power-Up Trigger and the Run-Time Trigger, double-click the instance name or the Power-Up Trigger name in the **Instance Manager**.

You cannot make changes to Power-Up Trigger conditions that would normally require a full recompile with Runtime Trigger conditions, such as adding signals, deleting signals, or changing between basic and advanced triggers. To apply these changes to the Power-Up Trigger conditions, first make the changes using the Runtime Trigger conditions.
Any change made to the Power-Up Trigger conditions requires that you recompile the SignalTap II Logic Analyzer instance, even if a similar change to the Runtime Trigger conditions does not require a recompilation.

While creating or making changes to the trigger conditions for the Run-Time Trigger or the Power-Up Trigger, you may want to copy these conditions to the other trigger. This enables you to look for the same trigger during both power-up and runtime. To do this, right-click the instance name or the Power-Up Trigger name in the Instance Manager and click Duplicate Trigger, or select the instance name or the Power-Up Trigger name and on the Edit menu, click Duplicate Trigger.

You can also use In-System Sources and Probes in conjunction with the SignalTap II Logic Analyzer to force trigger conditions. The In-System Sources and Probes feature allows you to drive and sample values on to selected nets over the JTAG chain. For more information, refer to the Design Debugging Using In-System Sources and Probes chapter in volume 3 of the Quartus II Handbook.

### Using External Triggers

You can create a trigger input that allows you to trigger the SignalTap II Logic Analyzer from an external source. The external trigger input behaves like trigger condition 1, is evaluated, and must be TRUE before any other configured trigger conditions are evaluated. The logic analyzer supplies a signal to trigger external devices or other SignalTap II Logic Analyzer instances. These features allow you to synchronize external logic analysis equipment with the internal logic analyzer. Power-Up Triggers can use the external triggers feature, but they must use the same source or target signal as their associated Run-Time Trigger.

You can use external triggers to perform cross-triggering on a hard processor system (HPS). Use your processor debugger to configure the HPS to obey or disregard cross-trigger request from the FPGA, and to issue or not issue cross-trigger requests to the FPGA. Use your processor debugger in combination with the SignalTap II external trigger feature to develop a dynamic combination of cross-trigger behaviors.

For more information about setting up external triggers, refer to Signal Configuration Pane in Quartus II Help.

### Using the Trigger Out of One Analyzer as the Trigger In of Another Analyzer

An advanced feature of the SignalTap II Logic Analyzer is the ability to use the Trigger out of one analyzer as the Trigger in to another analyzer. This feature allows you to synchronize and debug events that occur across multiple clock domains.

To perform this operation, first turn on Trigger out for the source logic analyzer instance. On the Instance list of the Trigger out trigger, select the targeted logic analyzer instance. For example, if the instance named auto_signaltap_0 should trigger auto_signaltap_1, select auto_signaltap_1|trigger_in.

Turning on Trigger out automatically enables the Trigger in of the targeted logic analyzer instance and fills in the Instance field of the Trigger in trigger with the Trigger out signal from the source logic analyzer instance. In this example, auto_signaltap_0 is targeting auto_signaltap_1. The Trigger In Instance field of auto_signaltap_1 is automatically filled in with auto_signaltap_0|trigger_out.
Compile the Design

When you add an .stp to your project, the SignalTap II Logic Analyzer becomes part of your design. You must compile your project to incorporate the SignalTap II logic and enable the JTAG connection you use to control the logic analyzer. When you are debugging with a traditional external logic analyzer, you must often make changes to the signals monitored as well as the trigger conditions. Because these adjustments require that you recompile your design when using the SignalTap II Logic Analyzer, use the SignalTap II Logic Analyzer feature along with incremental compilation in the Quartus II software to reduce recompilation time.

For more information on reducing your recompilation burden with incremental compilation, refer to Using the Incremental Compilation Design Flow in Quartus II Help.

Faster Compilations with Quartus II Incremental Compilation

When you compile your design with an .stp, the sld_signaltap and sld_hub entities are automatically added to the compilation hierarchy. These two entities are the main components of the SignalTap II Logic Analyzer, providing the trigger logic and JTAG interface required for operation.

Incremental compilation enables you to preserve the synthesis and fitting results of your original design and add the SignalTap II Logic Analyzer to your design without recompiling your original source code. Incremental compilation is also useful when you want to modify the configuration of the .stp. For example, you can modify the buffer sample depth or memory type without performing a full compilation after the change is made. Only the SignalTap II Logic Analyzer, configured as its own design partition, must be recompiled to reflect the changes.

To use incremental compilation, first enable Full Incremental Compilation for your design if it is not already enabled, assign design partitions if necessary, and set the design partitions to the correct preservation levels. Incremental compilation is the default setting for new projects in the Quartus II software, so you can establish design partitions immediately in a new project. However, it is not necessary to create any design partitions to use the SignalTap II incremental compilation feature. When your design is set up to use full incremental compilation, the SignalTap II Logic Analyzer acts as its own separate design partition. You can begin taking advantage of incremental compilation by using the SignalTap II: post-fitting filter in the Node Finder to add signals for logic analysis.

Enabling Incremental Compilation for Your Design

Your project is fully compiled the first time, establishing the design partitions you have created. When enabled for your design, the SignalTap II Logic Analyzer is always a separate partition. After the first compilation, you can use the SignalTap II Logic Analyzer to analyze signals from the post-fit netlist. If your partitions are designed correctly, subsequent compilations due to SignalTap II Logic Analyzer settings take less time.

The netlist type for the top-level partition defaults to source. To take advantage of incremental compilation, specify the Netlist types for the partitions you wish to tap as Post-fit.
For more information about configuring and performing incremental compilation, refer to the *Quartus II Incremental Compilation for Hierarchical and Team-Based Design* chapter in volume 1 of the *Quartus II Handbook*.

**Using Incremental Compilation with the SignalTap II Logic Analyzer**

The SignalTap II Logic Analyzer is automatically configured to work with the incremental compilation flow. For all signals that you want to connect to the SignalTap II Logic Analyzer from the post-fit netlist, set the netlist type of the partition containing the desired signals to *Post-Fit* or *Post-Fit (Strict)* with a Fitter Preservation Level of *Placement and Routing* using the Design Partitions window. Use the **SignalTap II: post-fitting filter** in the **Node Finder** to add the signals of interest to your SignalTap II configuration file. If you want to add signals from the pre-synthesis netlist, set the netlist type to **Source File** and use the **SignalTap II: pre-synthesis filter** in the **Node Finder**. Do not use the netlist type **Post-Synthesis** with the SignalTap II Logic Analyzer.

Be sure to conform to the following guidelines when using post-fit and pre-synthesis nodes:

- Read all incremental compilation guidelines to ensure the proper partition of a project.
- To speed compile time, use only post-fit nodes for partitions specified as to preservation-level post-fit.
- Do not mix pre-synthesis and post-fit nodes in any partition. If you must tap pre-synthesis nodes for a particular partition, make all tapped nodes in that partition pre-synthesis nodes and change the netlist type to **source** in the design partitions window.

Node names may be different between a pre-synthesis netlist and a post-fit netlist. In general, registers and user input signals share common names between the two netlists. During compilation, certain optimizations change the names of combinational signals in your RTL. If the type of node name chosen does not match the netlist type, the compiler may not be able to find the signal to connect to your SignalTap II Logic Analyzer instance for analysis. The compiler issues a critical warning to alert you of this scenario. The signal that is not connected is tied to ground in the **SignalTap II data** tab.

If you do use incremental compile flow with the SignalTap II Logic Analyzer and source file changes are necessary, be aware that you may have to remove compiler-generated post-fit net names. Source code changes force the affected partition to go through resynthesis. During synthesis, the compiler cannot find compiler-generated net names from a previous compilation.

Altera recommends using only registered and user-input signals as debugging taps in your *.stp* whenever possible.

Both registered and user-supplied input signals share common node names in the pre-synthesis and post-fit netlist. As a result, using only registered and user-supplied input signals in your *.stp* limits the changes you need to make to your SignalTap II Logic Analyzer configuration.
You can check the nodes that are connected to each SignalTap II instance using the In-System Debugging compilation reports. These reports list each node name you selected to connect to a SignalTap II instance, the netlist type used for the particular connection, and the actual node name used after compilation. If incremental compile is turned off, the In-System Debugging reports are located in the Analysis & Synthesis folder. If incremental compile is turned on, this report is located in the Partition Merge folder. Figure 13–33 shows an example of an In-System Debugging compilation report for a design using incremental compilation.

Figure 13–33. Compilation Report Showing Connectivity to SignalTap II Instance

<p>| Connections to In-System Debugging Instance “auto_signaltap_0” |
|-----------------|---------------|</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Status</th>
<th>Pinlist Name</th>
<th>Netlist Type</th>
<th>Actual Name</th>
<th>Details</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>v<a href="25">0</a></td>
<td>pre-synthesis</td>
<td>connected</td>
<td>v[0][15]</td>
<td>V[0][15]</td>
<td>N/A</td>
</tr>
<tr>
<td>2</td>
<td>v<a href="25">0</a></td>
<td>pre-synthesis</td>
<td>connected</td>
<td>v[0][15]</td>
<td>V[0][15]</td>
<td>N/A</td>
</tr>
<tr>
<td>3</td>
<td>x[0]_tap[0]</td>
<td>pre-synthesis</td>
<td>connected</td>
<td>x[0]_tap[0]</td>
<td>x[0]_tap[0]</td>
<td>N/A</td>
</tr>
<tr>
<td>4</td>
<td>x[0]_tap[0]</td>
<td>pre-synthesis</td>
<td>connected</td>
<td>x[0]_tap[0]</td>
<td>x[0]_tap[0]</td>
<td>N/A</td>
</tr>
<tr>
<td>5</td>
<td>clk</td>
<td>post-fit</td>
<td>connected</td>
<td>clk</td>
<td>clk</td>
<td>N/A</td>
</tr>
<tr>
<td>6</td>
<td>@<a href="10">0</a></td>
<td>post-fit</td>
<td>connected</td>
<td>@<a href="10">0</a></td>
<td>@<a href="10">0</a></td>
<td>N/A</td>
</tr>
<tr>
<td>7</td>
<td>@<a href="10">0</a></td>
<td>post-fit</td>
<td>connected</td>
<td>@<a href="10">0</a></td>
<td>@<a href="10">0</a></td>
<td>N/A</td>
</tr>
<tr>
<td>8</td>
<td>@<a href="10">1</a></td>
<td>post-fit</td>
<td>connected</td>
<td>@<a href="10">1</a></td>
<td>@<a href="10">1</a></td>
<td>N/A</td>
</tr>
<tr>
<td>9</td>
<td>@<a href="10">1</a></td>
<td>post-fit</td>
<td>connected</td>
<td>@<a href="10">1</a></td>
<td>@<a href="10">1</a></td>
<td>N/A</td>
</tr>
<tr>
<td>10</td>
<td>x[1]_tap[1]</td>
<td>post-fit</td>
<td>connected</td>
<td>x[1]_tap[1]</td>
<td>x[1]_tap[1]</td>
<td>N/A</td>
</tr>
<tr>
<td>12</td>
<td>x[1]_tap[1]</td>
<td>post-fit</td>
<td>connected</td>
<td>x[1]_tap[1]</td>
<td>x[1]_tap[1]</td>
<td>N/A</td>
</tr>
<tr>
<td>14</td>
<td>x[1]_tap[1]</td>
<td>post-fit</td>
<td>connected</td>
<td>x[1]_tap[1]</td>
<td>x[1]_tap[1]</td>
<td>N/A</td>
</tr>
<tr>
<td>15</td>
<td>x[1]_tap[1]</td>
<td>post-fit</td>
<td>connected</td>
<td>x[1]_tap[1]</td>
<td>x[1]_tap[1]</td>
<td>N/A</td>
</tr>
</tbody>
</table>

To verify that your original design was not modified, examine the messages in the Partition Merge section of the Compilation Report. Figure 13–34 shows an example of the messages displayed.

Figure 13–34. Compilation Report Messages
Unless you make changes to your design partitions that require recompilation, only the SignalTap II design partition is recompiled. If you make subsequent changes to only the .stp, only the SignalTap II design partition must be recompiled, reducing your recompilation time.

**Preventing Changes Requiring Recompilation**

You can configure the .stp to prevent changes that normally require recompilation. To do this, select a lock mode from above the node list in the Setup tab. To lock your configuration, choose to allow only trigger condition changes, regardless of whether you use incremental compilation.

For more information about the use of lock modes, refer to **Setup Tab (SignalTap II Logic Analyzer)** in Quartus II Help.

**Timing Preservation with the SignalTap II Logic Analyzer**

In addition to verifying functionality, timing closure is one of the most crucial processes in successfully completing a design. When you compile a project with a SignalTap II Logic Analyzer without the use of incremental compilation, you add IP to your existing design. Therefore, you can affect the existing placement, routing, and timing of your design. To minimize the effect that the SignalTap II Logic Analyzer has on your design, Altera recommends that you use incremental compilation for your project. Incremental compilation is the default setting in new designs and can be easily enabled and configured in existing designs. With the SignalTap II Logic Analyzer instance in its own design partition, it has little to no affect on your design.

In addition to using the incremental compilation flow for your design, you can use the following techniques to help maintain timing:

- Avoid adding critical path signals to your .stp.
- Minimize the number of combinational signals you add to your .stp and add registers whenever possible.
- Specify an fMAX constraint for each clock in your design.

For an example of timing preservation with the SignalTap II Logic Analyzer, refer to the **Area and Timing Optimization** chapter in volume 2 of the Quartus II Handbook.

**Performance and Resource Considerations**

There is a necessary trade-off between the runtime flexibility of the SignalTap II Logic Analyzer, the timing performance of the SignalTap II Logic Analyzer, and resource usage. The SignalTap II Logic Analyzer allows you to select the runtime configurable parameters to balance the need for runtime flexibility, speed, and area. The default values have been chosen to provide maximum flexibility so you can complete debugging as quickly as possible; however, you can adjust these settings to determine whether there is a more optimal configuration for your design.

The following tips provide extra timing slack if you have determined that the SignalTap II logic is in your critical path, or to alleviate the resource requirements that the SignalTap II Logic Analyzer consumes if your design is resource-constrained.
If SignalTap II logic is part of your critical path, follow these tips to speed up the performance of the SignalTap II Logic Analyzer:

- **Disable runtime configurable options**—Certain resources are allocated to accommodate for runtime flexibility. If you use either advanced triggers or State-based triggering flow, disable runtime configurable parameters for a boost in \( f_{\text{MAX}} \) of the SignalTap II logic. If you are using State-based triggering flow, try disabling the **Goto state destination** option and performing a recompilation before disabling the other runtime configurable options. The **Goto state destination** option has the greatest impact on \( f_{\text{MAX}} \) as compared to the other runtime configurable options.

- **Minimize the number of signals that have Trigger Enable selected**—All signals that you add to the .stp have **Trigger Enable** turned on. Turn off **Trigger Enable** for signals that you do not plan to use as triggers.

- **Turn on Physical Synthesis for register retiming**—If you have a large number of triggering signals enabled (greater than the number of inputs that would fit in a LAB) that fan-in logic to a gate-based triggering condition, such as a basic trigger condition or a logical reduction operator in the advanced trigger tab, turn on **Perform register retiming**. This can help balance combinational logic across LABs.

If your design is resource constrained, follow these tips to reduce the amount of logic or memory used by the SignalTap II Logic Analyzer:

- **Disable runtime configurable options**—Disabling runtime configurability for advanced trigger conditions or runtime configurable options in the State-based triggering flow results in using fewer LEs.

- **Minimize the number of segments in the acquisition buffer**—You can reduce the number of logic resources used for the SignalTap II Logic Analyzer by limiting the number of segments in your sampling buffer to only those required.

- **Disable the Data Enable for signals that are used for triggering only**—By default, both the **data enable** and **trigger enable** options are selected for all signals. Turning off the **data enable** option for signals used as trigger inputs only saves on memory resources used by the SignalTap II Logic Analyzer.

Because performance results are design-dependent, try these options in different combinations until you achieve the desired balance between functionality, performance, and utilization.

For more information about area and timing optimization, refer the *Area and Timing Optimization* chapter in volume 2 of the *Quartus II Handbook*.

**Program the Target Device or Devices**

After you compile your project, including the SignalTap II Logic Analyzer, configure the FPGA target device. When you are using the SignalTap II Logic Analyzer for debugging, configure the device from the .stp instead of the Quartus II Programmer. Because you configure from the .stp, you can open more than one .stp and program multiple devices to debug multiple designs simultaneously.
The settings in an .stp must be compatible with the programming .sof used to program the device. An .stp is considered compatible with an .sof when the settings for the logic analyzer, such as the size of the capture buffer and the signals selected for monitoring or triggering, match the way the target device is programmed. If the files are not compatible, you can still program the device, but you cannot run or control the logic analyzer from the SignalTap II Logic Analyzer Editor.

When the SignalTap II Logic Analyzer detects incompatibility after analysis is started, a system error message is generated containing two CRC values, the expected value and the value retrieved from the .stp instance on the device. The CRC values are calculated based on all SignalTap II settings that affect the compilation.

To ensure programming compatibility, make sure to program your device with the latest .sof created from the most recent compilation. Checking whether or not a particular SOF is compatible with the current SignalTap II configuration is achieved quickly by attaching the SOF to the SOF manager. For more details about using the SOF manager, refer to “Managing Multiple SignalTap II Files and Configurations” on page 13–25.

Before starting a debugging session, do not make any changes to the .stp settings that would require recompiling the project. You can check the SignalTap II status display at the top of the Instance Manager pane to verify whether a change you made requires recompiling the project, producing a new .sof. This gives you the opportunity to undo the change, so that you do not need to recompile your project. To prevent any such changes, select Allow trigger condition changes only to lock the .stp.

Although the Quartus II project is not required when using an .stp, it is recommended. The project database contains information about the integrity of the current SignalTap II Logic Analyzer session. Without the project database, there is no way to verify that the current .stp matches the .sof that is downloaded to the device. If you have an .stp that does not match the .sof, incorrect data is captured in the SignalTap II Logic Analyzer.

For instructions on programming devices in the Quartus II software, refer to Running the SignalTap II Logic Analyzer in Quartus II Help.

**Run the SignalTap II Logic Analyzer**

After the device is configured with your design that includes the SignalTap II Logic Analyzer, perform debugging operations in a manner similar to when you use an external logic analyzer. You initialize the logic analyzer by starting an analysis. When your trigger event occurs, the captured data is stored in the memory buffer on the device and then transferred to the .stp with the JTAG connection.

You can also perform the equivalent of a force trigger instruction that lets you view the captured data currently in the buffer without a trigger event occurring. Figure 13–35 illustrates a flow that shows how you operate the SignalTap II Logic Analyzer. The flowchart indicates where Power-Up and Runtime Trigger events occur and when captured data from these events is available for analysis.
Run the SignalTap II Logic Analyzer

For information on running the analyzer from the Instance Manager pane, refer to Running the SignalTap II Logic Analyzer in Quartus II Help.

You can also use In-System Sources and Probes in conjunction with the SignalTap II Logic Analyzer to force trigger conditions. The In-System Sources and Probes feature allows you to drive and sample values on to selected signals over the JTAG chain. For more information, refer to the Design Debugging Using In-System Sources and Probes chapter in volume 3 of the Quartus II Handbook.
Runtime Reconfigurable Options

Certain settings in the .stp are changeable without recompiling your design when you use Runtime Trigger mode. Runtime Reconfigurable features are described in Table 13–9.

Table 13–9. Runtime Reconfigurable Features

<table>
<thead>
<tr>
<th>Runtime Reconfigurable Setting</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Basic Trigger Conditions and Basic Storage Qualifier Conditions</td>
<td>All signals that have the Trigger condition turned on can be changed to any basic trigger condition value without recompiling.</td>
</tr>
<tr>
<td>Advanced Trigger Conditions and Advanced Storage Qualifier</td>
<td>Many operators include runtime configurable settings. For example, all comparison operators are runtime-configurable. Configurable settings are shown with a white background in the block representation. This runtime reconfigurable option is turned on in the Object Properties dialog box.</td>
</tr>
<tr>
<td>Switching between a storage-qualified and a continuous acquisition</td>
<td>Within any storage-qualified mode, you can switch to continuous capture mode without recompiling the design. To enable this feature, turn on disable storage qualifier.</td>
</tr>
<tr>
<td>State-based trigger flow parameters</td>
<td>Table 13–4 lists Reconfigurable State-based trigger flow options.</td>
</tr>
</tbody>
</table>

Runtime Reconfigurable options can potentially save time during the debugging cycle by allowing you to cover a wider possible scenario of events without the need to recompile the design. You may experience a slight impact to the performance and logic utilization of the SignalTap II IP core. You can turn off Runtime re-configurability for Advanced Trigger Conditions and the State-based trigger flow parameters, boosting performance and decreasing area utilization.

You can configure the .stp to prevent changes that normally require recompilation. To do this, in the Setup tab, select Allow Trigger Condition changes only above the node list.
Example 13–4 illustrates a potential use case for Runtime Reconfigurable features. This example provides a storage qualified enabled State-based trigger flow description and shows how you can modify the size of a capture window at runtime without a recompile. This example gives you equivalent functionality to a segmented buffer with a single trigger condition where the segment sizes are runtime reconfigurable.

**Example 13–4. Trigger Flow Description Providing Runtime Reconfigurable “Segments”**

```plaintext
state ST1:
if ( condition1 && (c1 <= m) ) // each "segment" triggers on condition
  begin // m = number of total "segments"
    start_store;
    increment c1;
    goto ST2:
  End
else (c1 > m ) //This else condition handles the last
  begin
    start_store
    Trigger (n-1)
  end

state ST2:
if ( c2 >= n) //n = number of samples to capture in each
  begin
    reset c2;
    stop_store;
    goto ST1;
  end
else (c2 < n)
  begin
    increment c2;
    goto ST2;
  end

Note to Example 13–4:
(1) \( m \times n \) must equal the sample depth to efficiently use the space in the sample buffer.
```

Figure 13–36 shows a segmented buffer described by the trigger flow in Example 13–4.

During runtime, the values \( m \) and \( n \) are runtime reconfigurable. By changing the \( m \) and \( n \) values in the preceding trigger flow description, you can dynamically adjust the segment boundaries without incurring a recompile.

**Figure 13–36. Segmented Buffer Created with Storage Qualifier and State-Based Trigger (1)**

<table>
<thead>
<tr>
<th>Segment 1</th>
<th>Segment 2</th>
<th>Segment m</th>
</tr>
</thead>
<tbody>
<tr>
<td><img src="#" alt="Segment 1" /></td>
<td><img src="#" alt="Segment 2" /></td>
<td><img src="#" alt="Segment m" /></td>
</tr>
</tbody>
</table>

Note to Figure 13–36:
(1) Total sample depth is fixed, where \( m \times n \) must equal sample depth.
You can add states into the trigger flow description and selectively mask out specific states and enable other ones at runtime with status flags.

Example 13–5 shows a modified description of Example 13–4 with an additional state inserted. You use this extra state to specify a different trigger condition that does not use the storage qualifier feature. You insert status flags into the conditional statements to control the execution of the trigger flow.

Example 13–5. Modified Trigger Flow Description of Example 16–4 with Status Flags to Selectively Enable States

```vhdl
state ST1:
if (condition2 && f1)     //additional state added for a non-segmented
begin
    start_store;
    trigger
end
else if (! f1)
goto ST2;

state ST2:
if ( (condition1 && (c1 <= m)) && f2)   // f2 status flag used to mask state. Set f2
begin
    start_store;
    increment c1;
    goto ST3:
end
else (c1 > m )
start_store
    Trigger (n-1)
end

state ST3:
if ( c2 >= n)
begin
    reset c2;
    stop_store;
    goto ST1;
end
else (c2 < n)
begin
    increment c2;
    goto ST2;
end
```

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**SignalTap II Status Messages**

Table 13–10 describes the text messages that might appear in the SignalTap II Status Indicator in the *Instance Manager* pane before, during, and after a data acquisition. Use these messages to monitor the state of the logic analyzer or what operation it is performing.

**Table 13–10. Text Messages in the SignalTap II Status Indicator**

<table>
<thead>
<tr>
<th>Message</th>
<th>Message Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Not running</td>
<td>The SignalTap II Logic Analyzer is not running. There is no connection to a device or the device is not configured.</td>
</tr>
<tr>
<td>(Power-Up Trigger) Waiting for clock (1)</td>
<td>The SignalTap II Logic Analyzer is performing a Runtime or Power-Up Trigger acquisition and is waiting for the clock signal to transition.</td>
</tr>
<tr>
<td>Acquiring (Power-Up) pre-trigger data (1)</td>
<td>The trigger condition has not been evaluated yet. A full buffer of data is collected if using the non-segmented buffer acquisition mode and storage qualifier type is continuous.</td>
</tr>
<tr>
<td>Trigger In conditions met</td>
<td>Trigger In condition has occurred. The SignalTap II Logic Analyzer is waiting for the condition of the first trigger condition to occur. This can appear if Trigger In is specified.</td>
</tr>
<tr>
<td>Waiting for (Power-up) trigger (1)</td>
<td>The SignalTap II Logic Analyzer is now waiting for the trigger event to occur.</td>
</tr>
<tr>
<td>Trigger level &lt;x&gt; met</td>
<td>The condition of trigger condition x has occurred. The SignalTap II Logic Analyzer is waiting for the condition specified in condition x + 1 to occur.</td>
</tr>
<tr>
<td>Acquiring (power-up) post-trigger data (1)</td>
<td>The entire trigger event has occurred. The SignalTap II Logic Analyzer is acquiring the post-trigger data. The amount of post-trigger data collected is you define between 12%, 50%, and 88% when the non-segmented buffer acquisition mode is selected.</td>
</tr>
<tr>
<td>Offload acquired (Power-Up) data (1)</td>
<td>Data is being transmitted to the Quartus II software through the JTAG chain.</td>
</tr>
<tr>
<td>Ready to acquire</td>
<td>The SignalTap II Logic Analyzer is waiting for you to initialize the analyzer.</td>
</tr>
</tbody>
</table>

**Note to Table 13–10:**

(1) This message can appear for both Runtime and Power-Up Trigger events. When referring to a Power-Up Trigger, the text in parentheses is added.

In segmented acquisition mode, pre-trigger and post-trigger do not apply.

**View, Analyze, and Use Captured Data**

Once a trigger event has occurred or you capture data manually, you can use the SignalTap II interface to examine the data, and use your findings to help debug your design.

For information about what you can do with captured data, refer to *Analyzing Data in the SignalTap II Logic Analyzer* in Quartus II Help.
Capturing Data Using Segmented Buffers

Segmented Acquisition buffers allow you to perform multiple captures with a separate trigger condition for each acquisition segment. This feature allows you to capture a recurring event or sequence of events that span over a long period time efficiently. Each acquisition segment acts as a non-segmented buffer, continuously capturing data when it is activated. When you run an analysis with the segmented buffer option enabled, the SignalTap II Logic Analyzer performs back-to-back data captures for each acquisition segment within your data buffer. The trigger flow, or the type and order in which the trigger conditions evaluate for each buffer, is defined by either the Sequential trigger flow control or the Custom State-based trigger flow control. Figure 13–37 shows a segmented acquisition buffer with four segments represented as four separate non-segmented buffers.

The SignalTap II Logic Analyzer finishes an acquisition with a segment, and advances to the next segment to start a new acquisition. Depending on when a trigger condition occurs, it may affect the way the data capture appears in the waveform viewer. Figure 13–37 illustrates the method in which data is captured. The Trigger markers in Figure 13–37—Trigger 1, Trigger 2, Trigger 3 and Trigger 4—refer to the evaluation of the segment_trigger and trigger commands in the Custom State-based trigger flow. If you use a sequential flow, the Trigger markers refer to trigger conditions specified within the Setup tab.

If the Segment 1 Buffer is the active segment and Trigger 1 occurs, the SignalTap II Logic Analyzer starts evaluating Trigger 2 immediately. Data Acquisition for Segment 2 buffer starts when either Segment Buffer 1 finishes its post-fill count, or when Trigger 2 evaluates as TRUE, whichever condition occurs first. Thus, trigger conditions associated with the next buffer in the data capture sequence can preempt the post-fill count of the current active buffer. This allows the SignalTap II Logic Analyzer to accurately capture all of the trigger conditions that have occurred. Samples that have not been used appear as a blank space in the waveform viewer.
Figure 13–38 shows an example of a capture using sequential flow control with the trigger condition for each segment specified as Don’t Care. Each segment before the last captures only one sample, because the next trigger condition immediately preempts capture of the current buffer. The trigger position for all segments is specified as pre-trigger (10% of the data is before the trigger condition and 90% of the data is after the trigger position). Because the last segment starts immediately with the trigger condition, the segment contains only post-trigger data. The three empty samples in the last segment are left over from the pre-trigger samples that the SignalTap II Logic Analyzer allocated to the buffer.

For the sequential trigger flow, the Trigger Position option applies to every segment in the buffer. For maximum flexibility on how the trigger position is defined, use the custom state-based trigger flow. By adjusting the trigger position specific to your debugging requirements, you can help maximize the use of the allocated buffer space.

### Differences in Pre-fill Write Behavior Between Different Acquisition Modes

The SignalTap II Logic Analyzer uses one of the following three modes when writing into the acquisition memory:

- Non-segmented buffer
- Non-segmented buffer with a storage qualifier
- Segmented buffer

There are subtle differences in the amount of data captured immediately after running the SignalTap II Logic Analyzer and before any trigger conditions occur. A non-segmented buffer, running in continuous mode, completely fills the buffer with sampled data before evaluating any trigger conditions. Thus, a non-segmented capture without any storage qualification enabled always shows a waveform with a full buffer’s worth of data captured.

Filling the buffer provides you with as much data as possible within the capture window. The buffer gets pre-filled with data samples prior to evaluating the trigger condition. As such, SignalTap requires that the buffer be filled at least once before any data can be retrieved through the JTAG connection and prevents the buffer from being dumped during the first acquisition prior to a trigger condition when you perform a Stop Analysis.
For segmented buffers and non-segmented buffers using any storage qualification mode, the SignalTap II Logic Analyzer immediately evaluates all trigger conditions while writing samples into the acquisition memory. The logic analyzer evaluates each trigger condition before acquiring a full buffer’s worth of samples. This evaluation is especially important when using any storage qualification on the data set. The logic analyzer may miss a trigger condition if it waits until a full buffer’s worth of data is captured before evaluating any trigger conditions.

If the trigger event occurs on any data sample before the specified amount of pre-trigger data has occurred, then the SignalTap II Logic Analyzer triggers and begins filling memory with post-trigger data, regardless of the amount of pre-trigger data you specify. For example, if you set the trigger position to 50% and set the logic analyzer to trigger on a processor reset, start the logic analyzer, and then power on your target system, the logic analyzer triggers. However, the logic analyzer memory is filled only with post-trigger data, and not any pre-trigger data, because the trigger event, which has higher precedence than the capture of pre-trigger data, occurred before the pre-trigger condition was satisfied.

Figure 13–39 and Figure 13–40 on page 13–58 show the difference between a non-segmented buffer in continuous mode and a non-segmented buffer using a storage qualifier. The logic analyzer for the waveforms below is configured with a sample depth of 64 bits, with a trigger position specified as Post trigger position.

Figure 13–39. SignalTap II Logic Analyzer Continuous Data Capture

Note to Figure 13–39:

(1) Continuous capture mode with post-trigger position.
(2) Capture of a recurring pattern using a non-segmented buffer in continuous mode. The SignalTap II Logic Analyzer is configured with a basic trigger condition (shown in the figure as “Trig1”) with a sample depth of 64 bits.
Notice in Figure 13–39 that Trig1 occurs several times in the data buffer before the SignalTap II Logic Analyzer actually triggers. A full buffer’s worth of data is captured before the logic analyzer evaluates any trigger conditions. After the trigger condition occurs, the logic analyzer continues acquisition until it captures eight additional samples (12% of the buffer, as defined by the "post-trigger" position).

**Figure 13–40. SignalTap II Logic Analyzer Conditional Data Capture (1)**

**Note to Figure 13–40:**

(1) Conditional capture, storage always enabled, post-fill count.

(2) SignalTap II Logic Analyzer capture of a recurring pattern using a non-segmented buffer in conditional mode. The logic analyzer is configured with a basic trigger condition (shown in the figure as 'Trig1'), with a sample depth of 64 bits. The “Trigger in” condition is specified as ‘Don’t care’, which means that every sample is captured.

Notice in Figure 13–40 that the logic analyzer triggers immediately. As in Figure 13–39, the logic analyzer completes the acquisition with eight samples, or 12% of 64, the sample capacity of the acquisition buffer.

**Creating Mnemonics for Bit Patterns**

The mnemonic table feature allows you to assign a meaningful name to a set of bit patterns, such as a bus. To create a mnemonic table, right-click in the Setup or Data tab of an .stp and click **Mnemonic Table Setup**. You create a mnemonic table by entering sets of bit patterns and specifying a label to represent each pattern. Once you have created a mnemonic table, assign the table to a group of signals. To assign a mnemonic table, right-click on the group, click **Bus Display Format** and select the desired mnemonic table.

You use the labels you create in a table in different ways on the Setup and Data tabs. On the Setup tab, you can create basic triggers with meaningful names by right-clicking an entry in the Trigger Conditions column and selecting a label from the table you assigned to the signal group. On the Data tab, if any captured data matches a bit pattern contained in an assigned mnemonic table, the signal group data is replaced with the appropriate label, making it easy to see when expected data patterns occur.

**Automatic Mnemonics with a Plug-In**

When you use a plug-in to add signals to an .stp, mnemonic tables for the added signals are automatically created and assigned to the signals defined in the plug-in. To enable these mnemonic tables manually, right-click on the name of the signal or signal group. On the **Bus Display Format** shortcut menu, then click the name of the mnemonic table that matches the plug-in.
As an example, the Nios II plug-in helps you to monitor signal activity for your design as the code is executed. If you set up the logic analyzer to trigger on a function name in your Nios II code based on data from an .elf, you can see the function name in the **Instance Address** signal group at the trigger sample, along with the corresponding disassembled code in the **Disassembly** signal group, as shown in **Figure 13–41**. Captured data samples around the trigger are referenced as offset addresses from the trigger function name.

**Figure 13–41. Data Tab when the Nios II Plug-In is Used**

---

**Locating a Node in the Design**

When you find the source of an error in your design using the SignalTap II Logic Analyzer, you can use the node locate feature to locate that signal in many of the tools found in the Quartus II software, as well as in your design files. This lets you find the source of the problem quickly so you can modify your design to correct the flaw. To locate a signal from the SignalTap II Logic Analyzer in one of the Quartus II software tools or your design files, right-click on the signal in the .stp, and click **Locate in <tool name>**.

You can locate a signal from the node list with the following tools:

- Assignment Editor
- Pin Planner
- Timing Closure Floorplan
- Chip Planner
- Resource Property Editor
- Technology Map Viewer
- RTL Viewer
- Design File

For more information about using these tools, refer to each of the corresponding chapters in the *Quartus II Handbook*. 
Saving Captured Data

The data log shows the history of captured data and the triggers used to capture the data. The SignalTap II Logic Analyzer acquires data, stores it in a log, and displays it as waveforms. When the logic analyzer is in auto-run mode and a trigger event occurs more than once, captured data for each time the trigger occurred is stored as a separate entry in the data log. This allows you to review the captured data for each trigger event. The default name for a log is based on the time when the data was acquired. Altera recommends that you rename the data log with a more meaningful name.

The logs are organized in a hierarchical manner; similar logs of captured data are grouped together in trigger sets. To open the Data Log pane, on the View menu, select Data Log. To turn on data logging, turn on Enable data log in the Data Log (Figure 13–19). To recall and activate a data log for a given trigger set, double-click the name of the data log in the list.

You can use the Data Log feature for organizing different sets of trigger conditions and different sets of signal configurations. For more information, refer to “Managing Multiple SignalTap II Files and Configurations” on page 13–25.

Exporting Captured Data to Other File Formats

You can export captured data to the following file formats, for use with other EDA simulation tools:

- Comma Separated Values File (.csv)
- Table File (.tbl)
- Value Change Dump File (.vcd)
- Vector Waveform File (.vwf)
- Graphics format files (.jpg, .bmp)

To export the captured data from SignalTap II Logic Analyzer, on the File menu, click Export and specify the File Name, Export Format, and Clock Period.

Creating a SignalTap II List File

Captured data can also be viewed in an .stp list file. An .stp list file is a text file that lists all the data captured by the logic analyzer for a trigger event. Each row of the list file corresponds to one captured sample in the buffer. Columns correspond to the value of each of the captured signals or signal groups for that sample. If a mnemonic table was created for the captured data, the numerical values in the list are replaced with a matching entry from the table. This is especially useful with the use of a plug-in that includes instruction code disassembly. You can immediately see the order in which the instruction code was executed during the same time period of the trigger event. To create an .stp list file in the Quartus II software, on the File menu, select Create/Update and click Create SignalTap II List File.

Other Features

The SignalTap II Logic Analyzer has other features that do not necessarily belong to a particular task in the task flow.
Using the SignalTap II MATLAB MEX Function to Capture Data

If you use MATLAB for DSP design, you can call the MATLAB MEX function `alt_signaltap_run`, built into the Quartus II software, to acquire data from the SignalTap II Logic Analyzer directly into a matrix in the MATLAB environment. If you use the MATLAB MEX function in a loop, you can perform as many acquisitions in the same amount of time as you can when using SignalTap II in the Quartus II software environment.

The SignalTap II MATLAB MEX function is available only in the Windows version of the Quartus II software. It is compatible with MATLAB Release 14 Original Release Version 7 and later.

To set up the Quartus II software and the MATLAB environment to perform SignalTap II acquisitions, perform the following steps:

1. In the Quartus II software, create an .stp file.
2. In the node list in the Data tab of the SignalTap II Logic Analyzer Editor, organize the signals and groups of signals into the order in which you want them to appear in the MATLAB matrix. Each column of the imported matrix represents a single SignalTap II acquisition sample, while each row represents a signal or group of signals in the order they are organized in the Data tab.
3. Save the .stp and compile your design. Program your device and run the SignalTap II Logic Analyzer to ensure your trigger conditions and signal acquisition work correctly.
4. In the MATLAB environment, add the Quartus II binary directory to your path with the following command:
   ```
   addpath <Quartus install directory>\win
   ```

   You can view the help file for the MEX function by entering the following command in MATLAB without any operators:
   ```
   alt_signaltap_run
   ```

   Use the MATLAB MEX function to open the JTAG connection to the device and run the SignalTap II Logic Analyzer to acquire data. When you finish acquiring data, close the JTAG connection.

To open the JTAG connection and begin acquiring captured data directly into a MATLAB matrix called `stp`, use the following command:

```matlab
stp = alt_signaltap_run \
    ('<stp filename>',
    ('signed'|'unsigned'),
    '<instance names>',
    '<signalset name>',
    '<trigger name>')
```
When capturing data you must assign a filename, for example, `<stp filename>` as a requirement of the MATLAB MEX function. Other MATLAB MEX function options are described in Table 13–11.

<table>
<thead>
<tr>
<th>Option</th>
<th>Usage</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>signed</td>
<td>'signed'</td>
<td>The <code>signed</code> option turns signal group data into 32-bit two’s-complement signed integers. The MSB of the group as defined in the SignalTap II Data tab is the sign bit. The <code>unsigned</code> option keeps the data as an unsigned integer. The default is <code>signed</code>.</td>
</tr>
<tr>
<td>unsigned</td>
<td>'unsigned'</td>
<td></td>
</tr>
<tr>
<td>&lt;instance name&gt;</td>
<td>'auto_signaltap_0'</td>
<td>Specify a SignalTap II instance if more than one instance is defined. The default is the first instance in the .stp, auto_signaltap_0.</td>
</tr>
<tr>
<td>&lt;signal set name&gt;</td>
<td>'my_signalset'</td>
<td>Specify the signal set and trigger from the SignalTap II data log if multiple configurations are present in the .stp. The default is the active signal set and trigger in the file.</td>
</tr>
<tr>
<td>&lt;trigger name&gt;</td>
<td>'my_trigger'</td>
<td></td>
</tr>
</tbody>
</table>

You can enable or disable verbose mode to see the status of the logic analyzer while it is acquiring data. To enable or disable verbose mode, use the following commands:

```matlab
alt_signaltap_run('VERBOSE_ON');
alt_signaltap_run('VERBOSE_OFF');
```

When you finish acquiring data, close the JTAG connection with the following command:

```matlab
alt_signaltap_run('END_CONNECTION');
```

For more information about the use of MATLAB MEX functions in MATLAB, refer to the MATLAB Help.

### Using SignalTap II in a Lab Environment

You can install a stand-alone version of the SignalTap II Logic Analyzer. This version is particularly useful in a lab environment in which you do not have a workstation that meets the requirements for a complete Quartus II installation, or if you do not have a license for a full installation of the Quartus II software. The standalone version of the SignalTap II Logic Analyzer is included with and requires the Quartus II stand-alone Programmer which is available from the Downloads page of the Altera website (www.altera.com).

### Remote Debugging Using the SignalTap II Logic Analyzer

You can use the SignalTap II Logic Analyzer to debug a design that is running on a device attached to a PC in a remote location.

To perform a remote debugging session, you must have the following setup:

- The Quartus II software installed on the local PC
- Stand-alone SignalTap II Logic Analyzer or the full version of the Quartus II software installed on the remote PC
- Programming hardware connected to the device on the PCB at the remote location
- TCP/IP protocol connection
Equipment Setup
On the PC in the remote location, install the standalone version of the SignalTap II Logic Analyzer, included in the Quartus II standalone Programmer, or the full version of the Quartus II software. This remote computer must have Altera programming hardware connected, such as the EthernetBlaster or USB-Blaster.

On the local PC, install the full version of the Quartus II software. This local PC must be connected to the remote PC across a LAN with the TCP/IP protocol.

For information about enabling remote access to a JTAG server, refer to Using the JTAG Server in Quartus II Help.

Using the SignalTap II Logic Analyzer in Devices with Configuration Bitstream Security
Certain device families support bitstream decryption during configuration using an on-device AES decryption engine. You can still use the SignalTap II Logic Analyzer to analyze functional data within the FPGA. However, note that JTAG configuration is not possible after the security key has been programmed into the device.

Altera recommends that you use an unencrypted bitstream during the prototype and debugging phases of the design. Using an unencrypted bitstream allows you to generate new programming files and reconfigure the device over the JTAG connection during the debugging cycle.

If you must use the SignalTap II Logic Analyzer with an encrypted bitstream, first configure the device with an encrypted configuration file using Passive Serial (PS), Fast Passive Parallel (FPP), or Active Serial (AS) configuration modes. The design must contain at least one instance of the SignalTap II Logic Analyzer. After the FPGA is configured with a SignalTap II Logic Analyzer instance in the design, when you open the SignalTap II Logic Analyzer in the Quartus II software, you then scan the chain and are ready to acquire data with the JTAG connection.

Backward Compatibility with Previous Versions of Quartus II Software
You can open an .stp created in a previous version in a current version of the Quartus II software. However, opening an .stp modifies it so that it cannot be opened in a previous version of the Quartus II software.

If you have a Quartus II project file from a previous version of the software, you may have to update the .stp configuration file to recompile the project. You can update the configuration file by opening the SignalTap II Logic Analyzer. If you need to update your configuration, a prompt appears asking if you would like to update the .stp to match the current version of the Quartus II software.
SignalTap II Command-Line Options

To compile your design with the SignalTap II Logic Analyzer using the command prompt, use the `quartus_stp` command. Table 13–12 shows the options that help you use the `quartus_stp` executable.

Table 13–12. SignalTap II Command-Line Options

<table>
<thead>
<tr>
<th>Option</th>
<th>Usage</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stp_file</td>
<td><code>quartus_stp --stp_file &lt;stp_filename&gt;</code></td>
<td>Assigns the specified .stp to the USE_SIGNALTAP_FILE in the .qsf.</td>
</tr>
<tr>
<td>enable</td>
<td><code>quartus_stp --enable</code></td>
<td>Creates assignments to the specified .stp in the .qsf and changes ENABLE_SIGNALTAP to ON. The SignalTap II Logic Analyzer is included in your design the next time the project is compiled. If no .stp is specified in the .qsf, the --stp_file option must be used. If the --enable option is omitted, the current value of ENABLE_SIGNALTAP in the .qsf is used.</td>
</tr>
<tr>
<td>disable</td>
<td><code>quartus_stp --disable</code></td>
<td>Removes the .stp reference from the .qsf and changes ENABLE_SIGNALTAP to OFF. The SignalTap II Logic Analyzer is removed from the design database the next time you compile your design. If the --disable option is omitted, the current value of ENABLE_SIGNALTAP in the .qsf is used.</td>
</tr>
<tr>
<td>create_signaltap_hdl_file</td>
<td><code>quartus_stp --create_signaltap_hdl_file</code></td>
<td>Creates an .stp representing the SignalTap II instance in the design generated by the SignalTap II Logic Analyzer megafunction created with the MegaWizard Plug-In Manager. The file is based on the last compilation. You must use the --stp_file option to create an .stp properly. Analogous to the Create SignalTap II File from Design Instance(s) command in the Quartus II software.</td>
</tr>
</tbody>
</table>

Example 13–6 illustrates how to compile a design with the SignalTap II Logic Analyzer at the command line.

**Example 13–6.**

```
quartus_stp filtref --stp_file stp1.stp --enable
quartus_map filtref --source=filtref.bdf --family=CYCLONE
quartus_fit filtref --part=EP1C12Q240C6 --fmax=80MHz --tsu=8ns
quartus_asm filtref
```

The `quartus_stp --stp_file stp1.stp --enable` command creates the QSF variable and instructs the Quartus II software to compile the `stp1.stp` file with your design. The --enable option must be applied for the SignalTap II Logic Analyzer to compile properly into your design.
Example 13–7 shows how to create a new .stp after building the SignalTap II Logic Analyzer instance with the MegaWizard Plug-In Manager.

Example 13–7.

```
quartus_stp filtref --create_signaltap_hdl_file --stp_file stp1.stp
```

For information about the other command line executables and options, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

SignalTap II Tcl Commands

The `quartus_stp` executable supports a Tcl interface that allows you to capture data without running the Quartus II GUI. You cannot execute SignalTap II Tcl commands from within the Tcl console in the Quartus II software. They must be executed from the command line with the `quartus_stp` executable. To execute a Tcl file that has SignalTap II Logic Analyzer Tcl commands, use the following command:

```
quartus_stp -t <Tcl file>
```

For information about Tcl commands that you can use with the SignalTap II Logic Analyzer Tcl package, refer to `::quartus::stp` in Quartus II Help.

Example 13–8 is an excerpt from a script you can use to continuously capture data. Once the trigger condition is met, the data is captured and stored in the data log.

Example 13–8.

```
#opens signaltap session
open_session -name stp1.stp
#start acquisition of instance auto_signaltap_0 and auto_signaltap_1 at the same time
#calling run_multiple_end will start all instances
#run after run_multiple_start call
run_multiple_start
run -instance auto_signaltap_0 -signal_set signal_set_1 -trigger /trigger_1 -data_log log_1 -timeout 5
run -instance auto_signaltap_1 -signal_set signal_set_1 -trigger /trigger_1 -data_log log_1 -timeout 5
run_multiple_end
#close signaltap session
close_session
```

When the script is completed, open the .stp that you used to capture data to examine the contents of the Data Log.

Design Example: Using SignalTap II Logic Analyzers in SOPC Builder Systems

The system in this example contains many components, including a Nios processor, a direct memory access (DMA) controller, on-chip memory, and an interface to external SDRAM memory. In this example, the Nios processor executes a simple C program from on-chip memory and waits for you to press a button. After you press a button, the processor initiates a DMA transfer, which you analyze using the SignalTap II Logic Analyzer.
Custom Triggering Flow Application Examples

The custom triggering flow in the SignalTap II Logic Analyzer is most useful for organizing a number of triggering conditions and for precise control over the acquisition buffer. This section provides two application examples for defining a custom triggering flow within the SignalTap II Logic Analyzer. Both examples can be easily copied and pasted directly into the state machine description box by using the state display mode All states in one window.

For additional triggering flow design examples for on-chip debugging, refer to the On-chip Debugging Design Examples page on the Altera website.

Design Example 1: Specifying a Custom Trigger Position

Actions to the acquisition buffer can accept an optional post-count argument. This post-count argument enables you to define a custom triggering position for each segment in the acquisition buffer. Example 13–9 shows an example that applies a trigger position to all segments in the acquisition buffer. The example describes a triggering flow for an acquisition buffer split into four segments. If each acquisition segment is 64 samples in depth, the trigger position for each buffer will be at sample #34. The acquisition stops after all four segments are filled once.

Example 13–9.

```vhdl
if (c1 == 3 && condition1) 
    trigger 30;
else if (condition1) 
begin
    segment_trigger 30;
    increment c1;
end
```
Each segment acts as a non-segmented buffer that continuously updates the memory contents with the signal values. The last acquisition before stopping the buffer is displayed on the **Data** tab as the last sample number in the affected segment. The trigger position in the affected segment is then defined by \( N - \text{post count fill} \), where \( N \) is the number of samples per segment. **Figure 13–42** illustrates the triggering position.

**Figure 13–42. Specifying a Custom Trigger Position**

---

**Design Example 2: Trigger When triggercond1 Occurs Ten Times between triggercond2 and triggercond3**

The custom trigger flow description is often useful to count a sequence of events before triggering the acquisition buffer. **Example 13–10** shows such a sample flow. This example uses three basic triggering conditions configured in the SignalTap II **Setup** tab.
This example triggers the acquisition buffer when condition1 occurs after condition3 and occurs ten times prior to condition3. If condition3 occurs prior to ten repetitions of condition1, the state machine transitions to a permanent wait state.

**Example 13–10.**

```plaintext
state ST1:
    if ( condition2 )
        begin
            reset c1;
            goto ST2;
        end

    State ST2 :
    if ( condition1 )
        increment c1;
    else if (condition3 && c1 < 10)
        goto ST3;
    else if ( condition3 && c1 >= 10)
        trigger;

ST3:
    goto ST3;
```

**SignalTap II Scripting Support**

You can run procedures and make settings described in this chapter in a Tcl script. You can also run some procedures at a command prompt. For detailed information about scripting command options, refer to the Quartus II Command-Line and Tcl API Help browser. To run the Help browser, type the following at the command prompt:

```plaintext
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the *Tcl Scripting* chapter in volume 2 of the *Quartus II Handbook*.

You can also refer to *About Quartus II Tcl Scripting* in Quartus II Help.

**Conclusion**

As the FPGA industry continues to make technological advancements, outdated methodologies are replaced with new technologies that maximize productivity. The SignalTap II Logic Analyzer gives you the same benefits as a traditional logic analyzer, without the many shortcomings of a piece of dedicated test equipment. The SignalTap II Logic Analyzer provides many new and innovative features that allow you to capture and analyze internal signals in your FPGA, allowing you to quickly debug your design.
Document Revision History

Table 13–13 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes Made</th>
</tr>
</thead>
</table>
| May 2013   | 13.0.0  | - Added recommendation to use the state-based flow for segmented buffers with separate trigger conditions, information about Basic OR trigger condition, and hard processor system (HPS) external triggers.  
| June 2012  | 12.0.0  | Updated Figure 13–5 on page 13–16 and “Adding Signals to the SignalTap II File” on page 13–10.                                             |
| November 2011 | 11.0.1 | Template update.  
Minor editorial updates.                                                                                           |
| May 2011   | 11.0.0  | Updated the requirement for the standalone SignalTap II software.                                                                            |
| December 2010 | 10.0.1 | Changed to new document template.                                                                                                           |
| July 2010  | 10.0.0  | - Add new acquisition buffer content to the “View, Analyze, and Use Captured Data” section.  
- Added script sample for generating hexadecimal CRC values in programmed devices.  
- Created cross references to Quartus II Help for duplicated procedural content.    |
| November 2009 | 9.1.0  | No change to content.                                                                                                                      |
| March 2009 | 9.0.0   | - Updated Table 13–1  
- Updated “Using Incremental Compilation with the SignalTap II Logic Analyzer” on page 13–45  
- Added new Figure 13–33  
- Made minor editorial updates                                                                 |
| November 2008 | 8.1.0  | Updated for the Quartus II software version 8.1 release:  
- Added new section “Using the Storage Qualifier Feature” on page 14–25  
- Added description of start_store and stop_store commands in section “Trigger Condition Flow Control” on page 14–36  
- Added new section “Runtime Reconfigurable Options” on page 14–63                        |
| May 2008   | 8.0.0   | Updated for the Quartus II software version 8.0:  
- Added “Debugging Finite State machines” on page 14-24  
- Documented various GUI usability enhancements, including improvements to the resource estimator, the bus find feature, and the dynamic display updates to the counter and flag resources in the State-based trigger flow control tab  
- Added “Capturing Data Using Segmented Buffers” on page 14–16  
- Added hyperlinks to referenced documents throughout the chapter  
- Minor editorial updates                                                                 |

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
The Quartus II Logic Analyzer Interface (LAI) allows you to use an external logic analyzer and a minimal number of Altera-supported device I/O pins to examine the behavior of internal signals while your design is running at full speed on your Altera®- supported device.

The LAI connects a large set of internal device signals to a small number of output pins. You can connect these output pins to an external logic analyzer for debugging purposes. In the Quartus II LAI, the internal signals are grouped together, distributed to a user-configurable multiplexer, and then output to available I/O pins on your Altera-supported device. Instead of having a one-to-one relationship between internal signals and output pins, the Quartus II LAI enables you to map many internal signals to a smaller number of output pins. The exact number of internal signals that you can map to an output pin varies based on the multiplexer settings in the Quartus II LAI.

This chapter details the following topics:

- “Choosing a Logic Analyzer”
- “Debugging Your Design Using the LAI” on page 14–4
- “Working with LAI Files” on page 14–4
- “Controlling the Active Bank During Runtime” on page 14–7
- “Using the LAI with Incremental Compilation” on page 14–7

The term “logic analyzer” when used in this chapter includes both logic analyzers and oscilloscopes equipped with digital channels, commonly referred to as mixed signal analyzers or MSOs.

Refer to Devices and Adapters in Quartus II Help for a list of Altera-supported devices.

Choosing a Logic Analyzer

The Quartus II software offers the following two general purpose on-chip debugging tools for debugging a large set of RTL signals from your design:

- The SignalTap® II Logic Analyzer
- An external logic analyzer, which connects to internal signals in your Altera-supported device by using the Quartus II LAI
Choosing a Logic Analyzer

The Quartus II software offers a portfolio of on-chip debugging tools. For an overview and comparison of all tools available in the Quartus II software on-chip debugging tool suite, refer to Section V. In-System Debugging in volume 3 of the Quartus II Handbook.

Required Components

You must have the following components to perform analysis using the Quartus II LAI:

- The Quartus II software starting with version 5.1 and later
- The device under test
- An external logic analyzer
- An Altera communications cable
- A cable to connect the Altera-supported device to the external logic analyzer

Table 14–1 describes the advantages of each debugging tool.

<table>
<thead>
<tr>
<th>Feature and Description</th>
<th>Logic Analyzer Interface</th>
<th>SignalTap II Logic Analyzer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sample Depth</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>You have access to a wider sample depth with an external logic analyzer. In the SignalTap II Logic Analyzer, the maximum sample depth is set to 128 Kb, which is a device constraint. However, with an external logic analyzer, there are no device constraints, providing you a wider sample depth.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Debugging Timing Issues</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Using an external logic analyzer provides you with access to a “timing” mode, which enables you to debug combined streams of data.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Performance</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>You frequently have limited routing resources available to place and route when you use the SignalTap II Logic Analyzer with your design. An external logic analyzer adds minimal logic, which removes resource limits on place-and-route.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Triggering Capability</td>
<td>✓ ✓</td>
<td>— —</td>
</tr>
<tr>
<td>The SignalTap II Logic Analyzer offers triggering capabilities that are comparable to external logic analyzers.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Use of Output Pins</td>
<td>— ✓</td>
<td>— ✓</td>
</tr>
<tr>
<td>Using the SignalTap II Logic Analyzer, no additional output pins are required. Using an external logic analyzer requires the use of additional output pins.</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Acquisition Speed</td>
<td>— ✓</td>
<td>— ✓</td>
</tr>
<tr>
<td>With the SignalTap II Logic Analyzer, you can acquire data at speeds of over 200 MHz. You can achieve the same acquisition speeds with an external logic analyzer; however, you must consider signal integrity issues.</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Figure 14–1 shows the LAI and the hardware setup.

**Figure 14–1. LAI and Hardware Setup**

![Diagram of LAI and hardware setup]

**Notes to Figure 14–1:**

1. Configuration and control of the LAI using a computer loaded with the Quartus II software via the JTAG port.
2. Configuration and control of the LAI using a third-party vendor logic analyzer via the JTAG port. Support varies by vendor.
Debugging Your Design Using the LAI

Figure 14–2 shows the steps you must follow to debug your design with the Quartus II LAI.

Figure 14–2. LAI and Hardware Setup

Notes to Figure 14–1:
(1) Configuration and control of the LAI using a computer loaded with the Quartus II software via the JTAG port.
(2) Configuration and control of the LAI using a third-party vendor logic analyzer via the JTAG port. Support varies by vendor.

Working with LAI Files

The .lai file stores the configuration of an LAI instance. The .lai file opens in the LAI editor. The editor allows you to group multiple internal signals to a set of external pins. The configuration parameters are described in the following sections.

To create a new .lai file or open an existing .lai file, refer to Setting Up the Logic Analyzer Interface in Quartus II Help.
Configuring the File Core Parameters

After you create the .lai file, you must configure the .lai file core parameters by clicking on the Setup View list, and then selecting Core Parameters. Table 14-2 lists the .lai file core parameters.

Table 14-2. LAI File Core Parameters

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Pin Count</td>
<td>The Pin Count parameter signifies the number of pins you want dedicated to your LAI. The pins must be connected to a debug header on your board. Within the Altera-supported device, each pin is mapped to a user-configurable number of internal signals. The Pin Count parameter can range from 1 to 255 pins.</td>
</tr>
<tr>
<td>Bank Count</td>
<td>The Bank Count parameter signifies the number of internal signals that you want to map to each pin. For example, a Bank Count of 8 implies that you will connect eight internal signals to each pin. The Bank Count parameter can range from 1 to 255 banks.</td>
</tr>
</tbody>
</table>
| Output/Capture Mode | The Output/Capture Mode parameter signifies the type of acquisition you perform. There are two options that you can select:  
Combinational/Timing—This acquisition uses your external logic analyzer's internal clock to determine when to sample data. Because Combinational/Timing acquisition samples data asynchronously to your Altera-supported device, you must determine the sample frequency you should use to debug and verify your system. This mode is effective if you want to measure timing information, such as channel-to-channel skew. For more information about the sampling frequency and the speeds at which it can run, refer to the data sheet for your external logic analyzer.  
Registered/State—This acquisition uses a signal from your system under test to determine when to sample. Because Registered/State acquisition samples data synchronously with your Altera-supported device, it provides you with a functional view of your Altera-supported device while it is running. This mode is effective when you verify the functionality of your design. |
| Clock            | The Clock parameter is available only when Output/Capture Mode is set to Registered State. You must specify the sample clock in the Core Parameters view. The sample clock can be any signal in your design. However, for best results, Altera recommends that you use a clock with an operating frequency fast enough to sample the data you would like to acquire. |
| Power-Up State   | The Power-Up State parameter specifies the power-up state of the pins you have designated for use with the LAI. You have the option of selecting tri-stated for all pins, or selecting a particular bank that you have enabled. |

Mapping the LAI File Pins to Available I/O Pins

To configure the .lai file I/O pin parameters, select Pins in the Setup View list. To assign pin locations for the LAI, double-click the Location column next to the reserved pins in the Name column, and the Pin Planner opens.

For information about how to use the Pin Planner, refer to the Pin Planner section in the I/O Management chapter in volume 2 of the Quartus II Handbook.

Mapping Internal Signals to the LAI Banks

After you have specified the number of banks to use in the Core Parameters settings page, you must assign internal signals for each bank in the LAI. Click the Setup View arrow and select Bank n or All Banks.

To view all of your bank connections, click Setup View and select All Banks.
Using the Node Finder

Before making bank assignments, on the View menu, point to Utility Windows and click Node Finder. Find the signals that you want to acquire, then drag and drop the signals from the Node Finder dialog box into the bank Setup View. When adding signals, use SignalTap II: pre-synthesis for non-incrementally routed instances and SignalTap II: post-fitting for incrementally routed instances.

As you continue to make assignments in the bank Setup View, the schematic of your LAI in the Logical View of your .lai file begins to reflect your assignments. Continue making assignments for each bank in the Setup View until you have added all of the internal signals for which you wish to acquire data.

Compiling Your Quartus II Project

When you save your .lai file, a dialog box prompts you to enable the LAI instance for the active project. Alternatively, you can specify the .lai file your project uses in the Global Project Settings dialog box.

After you specify the name of your .lai file, you must compile your project. To compile your project, on the Processing menu, click Start Compilation.

To ensure that the LAI is properly compiled with your project, expand the entity hierarchy in the Project Navigator. (To display the Project Navigator, on the View menu, point to Utility Windows and click Project Navigator.) If the LAI is compiled with your design, the sld_hub and sld_multitap entities are shown in the project navigator (Figure 14–3).

Figure 14–3. Project Navigator

<table>
<thead>
<tr>
<th>Entity</th>
<th>Logic Cells</th>
<th>LC Registers</th>
</tr>
</thead>
<tbody>
<tr>
<td>Stratix EP1S10B672C7</td>
<td>136 (1)</td>
<td>81</td>
</tr>
<tr>
<td>sld_multitap:auto_lai_0</td>
<td>35 (1)</td>
<td>15</td>
</tr>
<tr>
<td>sld_hub:sld_hub_inst</td>
<td>100 (25)</td>
<td>65</td>
</tr>
</tbody>
</table>

Programming Your Altera-Supported Device Using the LAI

After compilation completes, you must configure your Altera-supported device before using the LAI.

You can use the LAI with multiple devices in your JTAG chain. Your JTAG chain can also consist of devices that do not support the LAI or non-Altera, JTAG-compliant devices. To use the LAI in more than one Altera-supported device, create an .lai file and configure an .lai file for each Altera-supported device that you want to analyze.

To configure a device or a set of devices for use with LAI, refer to Enabling the Logic Analyzer Interface in Quartus II Help.
**Controlling the Active Bank During Runtime**

When you have programmed your Altera-supported device, you can control which bank you map to the reserved .lai file output pins. To control which bank you map, in the schematic in the logical view, right-click the bank and click **Connect Bank** (Figure 14–4).

![Figure 14–4. Configuring Banks](image)

**Acquiring Data on Your Logic Analyzer**

To acquire data on your logic analyzer, you must establish a connection between your device and the external logic analyzer.

For more information about this process and for guidelines about how to establish connections between debugging headers and logic analyzers, refer to the documentation for your logic analyzer.

**Using the LAI with Incremental Compilation**

The Incremental Compilation feature in the Quartus II software allows you to preserve the synthesis and fitting results of your design. This is an effective feature for reducing compilation times if you only modify a portion of a design or you wish to preserve the optimization results from a previous compilation.

The Incremental Compilation feature is well suited for use with LAI since LAI comprises a small portion of most designs. Because LAI consists of only a small portion of your design, incremental compilation helps to minimize your compilation time. Incremental compilation works best when you are only changing a small portion of your design. Incremental compilation yields an accurate representation of your design behavior when changing the .lai file through multiple compilations.

For further details on how to use Incremental Compilation with the LAI, refer to **Enabling the Logic Analyzer Interface** in Quartus II Help.
Conclusion

As the device industry continues to make technological advancements, outdated debugging methodologies must be replaced with new technologies that maximize productivity. The LAI feature enables you to connect many internal signals within your Altera-supported device to an external logic analyzer with the use of a small number of I/O pins. This new technology in the Quartus II software enables you to use feature-rich external logic analyzers to debug your Altera-supported device design, ultimately enabling you to deliver your product in the shortest amount of time.

Document Revision History

Table 14–3 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>• Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>• Changed to new document template</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>• Minor editorial updates</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Changed to new document template</td>
</tr>
<tr>
<td>August 2010</td>
<td>10.0.1</td>
<td>• Corrected links</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>• Created links to the Quartus II Help</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Editorial updates</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Removed Referenced Documents section</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>• Removed references to APEX devices</td>
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<td>• Removed Figures 15–4, 15–5, and 15–11 from 8.1 version</td>
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<td>• Updated device support list on page 15–3</td>
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<td>• Added links to referenced documents throughout the chapter</td>
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<td>• Minor editorial updates</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
15. In-System Modification of Memory and Constants

This chapter explains how to use the Quartus® II In-System Memory Content Editor as part of your FPGA design and verification flow.

The In-System Memory Content Editor allows you to view and update memories and constants with the JTAG port connection.

The In-System Memory Content Editor allows access to dense and complex FPGA designs. When you program devices, you have read and write access to the memories and constants through the JTAG interface. You can then identify, test, and resolve issues with your design by testing changes to memory contents in the FPGA while your design is running.

Overview

This chapter contains the following sections:

- “Updating Memory and Constants in Your Design” on page 15–2
- “Creating In-System Modifiable Memories and Constants” on page 15–2
- “Running the In-System Memory Content Editor” on page 15–2

When you use the In-System Memory Content Editor in conjunction with the SignalTap II Logic Analyzer, you can more easily view and debug your design in the hardware lab.

For more information about the SignalTap II Logic Analyzer, refer to the Design Debugging Using the SignalTap II Logic Analyzer chapter in volume 3 of the Quartus II Handbook.

The ability to read data from memories and constants allows you to quickly identify the source of problems. The write capability allows you to bypass functional issues by writing expected data. For example, if a parity bit in your memory is incorrect, you can use the In-System Memory Content Editor to write the correct parity bit values into your RAM, allowing your system to continue functioning. You can also intentionally write incorrect parity bit values into your RAM to check the error handling functionality of your design.

The Quartus II software offers a variety of on-chip debugging tools. For an overview and comparison of all tools available in the Quartus II software on-chip debugging tool suite, refer to Section IV. System Debugging Tools in volume 3 of the Quartus II Handbook.
For a list of the types of memories and constants currently supported by the Quartus II software, refer to *Megafunctions/LPM* in Quartus II Help.

### Updating Memory and Constants in Your Design

To use the In-System Updating of Memory and Constants feature, perform the following steps:

1. Identify the memories and constants that you want to access.
2. Edit the memories and constants to be run-time modifiable.
3. Perform a full compilation.
4. Program your device.
5. Launch the In-System Memory Content Editor.

### Creating In-System Modifiable Memories and Constants

When you specify that a memory or constant is run-time modifiable, the Quartus II software changes the default implementation. A single-port RAM is converted to a dual-port RAM, and a constant is implemented in registers instead of look-up tables (LUTs). These changes enable run-time modification without changing the functionality of your design.

To enable your memory or constant to be modifiable, refer to *Setting up the In-System Memory Content Editor* in Quartus II Help.

If you instantiate a memory or constant megafunction directly with ports and parameters in VHDL or Verilog HDL, add or modify the `lpm_hint` parameter as follows:

**In VHDL code,** add the following:

```vhdl
lpm_hint => "ENABLE_RUNTIME_MOD = YES,
            INSTANCE_NAME = <instantiation name>";
```

**In Verilog HDL code,** add the following:

```verilog
defparam <megafunction instance name>.lpm_hint =
   "ENABLE_RUNTIME_MOD = YES,
    INSTANCE_NAME = <instantiation name>";
```

### Running the In-System Memory Content Editor

The In-System Memory Content Editor has three separate panes: the **Instance Manager**, the **JTAG Chain Configuration**, and the **Hex Editor**.

The **Instance Manager** pane displays all available run-time modifiable memories and constants in your FPGA device. The **JTAG Chain Configuration** pane allows you to program your FPGA and select the Altera® device in the chain to update.

Using the In-System Memory Content Editor does not require that you open a project. The In-System Memory Content Editor retrieves all instances of run-time configurable memories and constants by scanning the JTAG chain and sending a query to the specific device selected in the **JTAG Chain Configuration** pane.
If you have more than one device with in-system configurable memories or constants in a JTAG chain, you can launch multiple In-System Memory Content Editors within the Quartus II software to access the memories and constants in each of the devices. Each In-System Memory Content Editor can access the in-system memories and constants in a single device.

**Instance Manager**

When you scan the JTAG chain to update the **Instance Manager** pane, you can view a list of all run-time modifiable memories and constants in the design. The **Instance Manager** pane displays the Index, Instance, Status, Width, Depth, Type, and Mode of each element in the list.

You can read and write to in-system memory with the **Instance Manager** pane. For more information refer to **Instance Manager Pane** in Quartus II Help.

In addition to the buttons available in the **Instance Manager** pane, you can read and write data by selecting commands from the Processing menu, or the right-click menu in the **Instance Manager** pane or **Hex Editor** pane.

The status of each instance is also displayed beside each entry in the **Instance Manager** pane. The status indicates if the instance is **Not running**, **Offloading data**, or **Updating data**. The health monitor provides information about the status of the editor.

The Quartus II software assigns a different index number to each in-system memory and constant to distinguish between multiple instances of the same memory or constant function. View the **In-System Memory Content Editor Settings** section of the Compilation Report to match an index number with the corresponding instance ID.

**Editing Data Displayed in the Hex Editor Pane**

You can edit data read from your in-system memories and constants displayed in the **Hex Editor** pane by typing values directly into the editor or by importing memory files.

For more information, refer to **Working with In-System Memory Content Editor Data** in Quartus II Help.

**Importing and Exporting Memory Files**

The In-System Memory Content Editor allows you to import and export data values for memories that have the In-System Updating feature enabled. Importing from a data file enables you to quickly load an entire memory image. Exporting to a data file enables you to save the contents of the memory for future use.

For more information, refer to **Working with In-System Memory Content Editor Data** in Quartus II Help.
Scripting Support

The In-System Memory Content Editor supports reading and writing of memory contents via a Tcl script or Tcl commands entered at a command prompt. For detailed information about scripting command options, refer to the Quartus II command-line and Tcl API Help browser.

To run the Help browser, type the following command at the command prompt:

```
quartus_sh --qhelp
```

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook and API Functions for Tcl in Quartus II Help. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

The commonly used commands for the In-System Memory Content Editor are as follows:

- **Reading from memory:**
  ```
  read_content_from_memory
  [-content_in_hex]
  [-instance_index <instance index>]
  [-start_address <starting address>]
  [-word_count <word count>]
  ```

- **Writing to memory:**
  ```
  write_content_to_memory
  ```

- **Save memory contents to file:**
  ```
  save_content_from_memory_to_file
  ```

- **Update memory contents from File:**
  ```
  update_content_to_memory_from_file
  ```

For descriptions of the command options and scripting examples, refer to the Tcl API Help Browser and the API Functions for Tcl in Quartus II Help.

Programming the Device with the In-System Memory Content Editor

If you make changes to your design, you can program the device from within the In-System Memory Content Editor.

To program the device, refer to Setting up the In-System Memory Content Editor in Quartus II Help.

Example: Using the In-System Memory Content Editor with the SignalTap II Logic Analyzer

The following scenario describes how you can use the In-System Updating of Memory and Constants feature with the SignalTap II Logic Analyzer to efficiently debug your design in-system. You can use the In-System Memory Content Editor and the SignalTap II Logic Analyzer simultaneously with the JTAG interface.
Scenario: After completing your FPGA design, you find that the characteristics of your FIR filter design are not as expected.

1. To locate the source of the problem, change all your FIR filter coefficients to be in-system modifiable and instantiate the SignalTap II Logic Analyzer.

2. Using the SignalTap II Logic Analyzer to tap and trigger on internal design nodes, you find the FIR filter to be functioning outside of the expected cutoff frequency.

3. Using the **In-System Memory Content Editor**, you check the correctness of the FIR filter coefficients. Upon reading each coefficient, you discover that one of the coefficients is incorrect.

4. Because your coefficients are in-system modifiable, you update the coefficients with the correct data with the **In-System Memory Content Editor**.

In this scenario, you can quickly locate the source of the problem using both the In-System Memory Content Editor and the SignalTap II Logic Analyzer. You can also verify the functionality of your device by changing the coefficient values before modifying the design source files.

You can also modify the coefficients with the In-System Memory Content Editor to vary the characteristics of the FIR filter, for example, filter attenuation, transition bandwidth, cut-off frequency, and windowing function.

### Conclusion

The In-System Updating of Memory and Constants feature provides access to a device for efficient debugging in a hardware lab. You can use the In-System Memory and Content Editor with the SignalTap II Logic Analyzer to maximize the visibility into an Altera FPGA. By maximizing visibility and access to internal logic of the device, you can identify and resolve problems with your design more easily.

### Document Revision History

Table 15–1 shows the revision history of this chapter.

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<td>12.0.0</td>
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<td>November 2009</td>
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<td>Delete references to APEX devices</td>
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Table 15–1. Document Revision History

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<td>Added reference to Section V. In-System Debugging in volume 3 of the Quartus II Handbook on page 16-1</td>
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<td>Removed references to the Mercury device, as it is now considered to be a “Mature” device</td>
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<td>Minor editorial updates</td>
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For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
16. Design Debugging Using In-System Sources and Probes

This chapter provides detailed instructions about how to use the In-System Sources and Probes Editor and Tcl scripting in the Quartus® II software to debug your design.

Traditional debugging techniques often involve using an external pattern generator to exercise the logic and a logic analyzer to study the output waveforms during run time. The SignalTap® II Logic Analyzer and SignalProbe allow you to read or “tap” internal logic signals during run time as a way to debug your logic design. You can make the debugging cycle more efficient when you can drive any internal signal manually within your design, which allows you to perform the following actions:

- Force the occurrence of trigger conditions set up in the SignalTap II Logic Analyzer
- Create simple test vectors to exercise your design without using external test equipment
- Dynamically control run time control signals with the JTAG chain

The In-System Sources and Probes Editor in the Quartus II software extends the portfolio of verification tools, and allows you to easily control any internal signal and provides you with a completely dynamic debugging environment. Coupled with either the SignalTap II Logic Analyzer or SignalProbe, the In-System Sources and Probes Editor gives you a powerful debugging environment in which to generate stimuli and solicit responses from your logic design.

The Virtual JTAG Megafunction and the In-System Memory Content Editor also give you the capability to drive virtual inputs into your design. The Quartus II software offers a variety of on-chip debugging tools. For an overview and comparison of all the tools available in the Quartus II software on-chip debugging tool suite, refer to Section IV. System Debugging Tools in volume 3 of the Quartus II Handbook.

Overview

This chapter includes the following topics:

- “Design Flow Using the In-System Sources and Probes Editor” on page 16–4
- “Running the In-System Sources and Probes Editor” on page 16–7
- “Tcl interface for the In-System Sources and Probes Editor” on page 16–9
- “Design Example: Dynamic PLL Reconfiguration” on page 16–13

The In-System Sources and Probes Editor consists of the ALTSOURCE_PROBE megafuction and an interface to control the ALTSOURCE_PROBE megafuction instances during run time. Each ALTSOURCE_PROBE megafuction instance provides you with source output ports and probe input ports, where source ports drive selected signals and probe ports sample selected signals. When you compile...
your design, the ALTSOURCE_PROBE megafunction sets up a register chain to either drive or sample the selected nodes in your logic design. During run time, the In-System Sources and Probes Editor uses a JTAG connection to shift data to and from the ALTSOURCE_PROBE megafunction instances. Figure 16–1 shows a block diagram of the components that make up the In-System Sources and Probes Editor.

**Figure 16–1. In-System Sources and Probes Editor Block Diagram**

The ALTSOURCE_PROBE megafunction hides the detailed transactions between the JTAG controller and the registers instrumented in your design to give you a basic building block for stimulating and probing your design. Additionally, the In-System Sources and Probes Editor provides single-cycle samples and single-cycle writes to selected logic nodes. You can use this feature to input simple virtual stimuli and to capture the current value on instrumented nodes. Because the In-System Sources and Probes Editor gives you access to logic nodes in your design, you can toggle the inputs of low-level components during the debugging process. If used in conjunction with the SignalTap II Logic Analyzer, you can force trigger conditions to help isolate your problem and shorten your debugging process.

The In-System Sources and Probes Editor allows you to easily implement control signals in your design as virtual stimuli. This feature can be especially helpful for prototyping your design, such as in the following operations:

- Creating virtual push buttons
- Creating a virtual front panel to interface with your design
- Emulating external sensor data
- Monitoring and changing run time constants on the fly
The In-System Sources and Probes Editor supports Tcl commands that interface with all your ALTSOURCE_PROBE megafuction instances to increase the level of automation.

**Hardware and Software Requirements**

The following components are required to use the In-System Sources and Probes Editor:

- Quartus II software

or

- Quartus II Web Edition (with the TalkBack feature turned on)
- Download Cable (USB-Blaster™ download cable or ByteBlaster™ cable)
- Altera® development kit or user design board with a JTAG connection to device under test

The In-System Sources and Probes Editor supports the following device families:

- Arria® GX
- Stratix® series
- HardCopy® II
- Cyclone® series
- MAX® series
Design Flow Using the In-System Sources and Probes Editor

The In-System Sources and Probes Editor supports an RTL flow. Signals that you want to view in the In-System Sources and Probes editor are connected to an instance of the ALTSOURCE_PROBE megafunction. After you compile the design, you can control each ALTSOURCE_PROBE instance via the In-System Sources and Probes Editor pane or via a Tcl interface. The complete design flow is shown in Figure 16–2.

Figure 16–2. FPGA Design Flow Using the In-System Sources and Probes Editor

Configuring the ALTSOURCE_PROBE Megafunction

To use the In-System Sources and Probes Editor in your design, you must first instantiate the ALTSOURCE_PROBE megafunction variation file. You can configure the ALTSOURCE_PROBE megafunction with the MegaWizard™ Plug-In Manager. Each source or probe port can be up to 256 bits. You can have up to 128 instances of the ALTSOURCE_PROBE megafunction in your design.
To configure the ALTSOURCE_PROBE megafunction, performing the following steps:

1. On the Tools menu, click **MegaWizard Plug-In Manager**.
2. Select **Create a new custom megafunction variation**.
3. Click **Next**.
4. On page 2a of the MegaWizard Plug-In Manager, make the following selections:
   a. In the **Installed Plug-Ins** list, expand the **JTAG-accessible Extensions** folder and select **In-System Sources and Probes**.

   ![Verify that the currently selected device family matches the device you are targeting.](image)

   b. Select an output file type and enter the name of the ALTSOURCE_PROBE megafunction. You can choose AHDL (.tdf), VHDL (.vhd), or Verilog HDL (.v) as the output file type.

5. Click **Next**.
6. On page 3 of the MegaWizard Plug-In Manager, make the following selections:
   a. Under **Do you want to specify an Instance Index?**, turn on **Yes**.
   b. Specify the ‘**Instance ID**’ of this instance.
   c. Specify the width of the probe port. The width can be from 0 bit to 256 bits.
   d. Specify the width of the source port. The width can be from 0 bit to 256 bits.
7. On page 3 of the MegaWizard Plug-In Manager, you can click **Advanced Options** and specify other options, including the following:
   - **What is the initial value of the source port, in hexadecimal?**—Allows you to specify the initial value driven on the source port at run time.
   - **Write data to the source port synchronously to the source clock**—Allows you to synchronize your source port write transactions with the clock domain of your choice.
   - **Create an enable signal for the registered source port**—When turned on, creates a clock enable input for the synchronization registers. You can turn on this option only when the **Write data to the source port synchronously to the source clock** option is turned on.

   ![The In-System Sources and Probes Editor does not support simulation. You must remove the ALTSOURCE_PROBE megafunction instantiation before you create a simulation netlist.](image)
Instantiating the ALTSOURCE_PROBE Megafuction

The MegaWizard Plug-In Manager produces the necessary variation file and the instantiation template based on your inputs to the MegaWizard. Use the template to instantiate the ALTSOURCE_PROBE megafuction variation file in your design. The port information is shown in Table 16–1.

Table 16–1. ALTSOURCE_PROBE Megafuction Port Information

<table>
<thead>
<tr>
<th>Port Name</th>
<th>Required?</th>
<th>Direction</th>
<th>Comments</th>
</tr>
</thead>
<tbody>
<tr>
<td>probe[]</td>
<td>No</td>
<td>Input</td>
<td>The outputs from your design.</td>
</tr>
<tr>
<td>source_clk</td>
<td>No</td>
<td>Input</td>
<td>Source Data is written synchronously to this clock. This input is required if you turn on Source Clock in the Advanced Options box in the MegaWizard Plug-In Manager.</td>
</tr>
<tr>
<td>source_ena</td>
<td>No</td>
<td>Input</td>
<td>Clock enable signal for source_clk. This input is required if specified in the Advanced Options box in the MegaWizard Plug-In Manager.</td>
</tr>
<tr>
<td>source[]</td>
<td>No</td>
<td>Output</td>
<td>Used to drive inputs to user design.</td>
</tr>
</tbody>
</table>

You can include up to 128 instances of the ALTSOURCE_PROBE megafuction in your design, if your device has available resources. Each instance of the ALTSOURCE_PROBE megafuction uses a pair of registers per signal for the width of the widest port in the megafuction. Additionally, there is some fixed overhead logic to accommodate communication between the ALTSOURCE_PROBE instances and the JTAG controller. You can also specify an additional pair of registers per source port for synchronization.

Compiling the Design

When you compile your design with the In-System Sources and Probes megafuction instantiated, an instance of the ALTSOURCE_PROBE and SLD_HUB instances are added to your compilation hierarchy automatically. These instances provide communication between the JTAG controller and your instrumented logic.

You can modify the number of connections to your design by editing the ALTSOURCE_PROBE megafuction. To open the design instance you want to modify in the MegaWizard Plug-In Manager, double-click the instance in the Project Navigator. You can then modify the connections in the HDL source file. You must recompile your design after you make changes.

You can use the Quartus II incremental compilation feature to reduce compilation time. Incremental compilation allows you to organize your design into logical partitions. During recompilation of a design, incremental compilation preserves the compilation results and performance of unchanged partitions and reduces design iteration time by compiling only modified design partitions.

For more information about the Quartus II incremental compilation feature, refer to the Quartus II Incremental Compilation for Hierarchical and Team-Based Design chapter in volume 1 of the Quartus II Handbook.
Running the In-System Sources and Probes Editor

The In-System Sources and Probes Editor gives you control over all ALTSOURCE_PROBE megafuction instances within your design. The editor allows you to view all available run time controllable instances of the ALTSOURCE_PROBE megafuction in your design, provides a push-button interface to drive all your source nodes, and provides a logging feature to store your probe and source data.

To run the In-System Sources and Probes Editor, on the Tools menu, click In-System Sources and Probes Editor.

The In-System Sources and Probes Editor contains three panes:

- **JTAG Chain Configuration**—Allows you to specify programming hardware, device, and file settings that the In-System Sources and Probes Editor uses to program and acquire data from a device.

- **Instance Manager**—Displays information about the instances generated when you compile a design, and allows you to control data that the In-System Sources and Probes Editor acquires.

- **In-System Sources and Probes Editor**—Logs all data read from the selected instance and allows you to modify source data that is written to your device.

When you use the In-System Sources and Probes Editor, you do not need to open a Quartus II software project. The In-System Sources and Probes Editor retrieves all instances of the ALTSOURCE_PROBE megafuction by scanning the JTAG chain and sending a query to the device selected in the JTAG Chain Configuration pane. You can also use a previously saved configuration to run the In-System Sources and Probes Editor.

Each In-System Sources and Probes Editor pane can access the ALTSOURCE_PROBE megafuction instances in a single device. If you have more than one device containing megafuction instances in a JTAG chain, you can launch multiple In-System Sources and Probes Editor panes to access the megafuction instances in each device.

Programming Your Device With JTAG Chain Configuration

After you compile your project, you must configure your FPGA before you use the In-System Sources and Probes Editor. To configure a device to use with the In-System Sources and Probes Editor, perform the following steps:

1. Open the In-System Sources and Probes Editor.

2. In the JTAG Chain Configuration pane, point to Hardware, and then select the hardware communications device. You may be prompted to configure your hardware; in this case, click Setup.

3. From the Device list, select the FPGA device to which you want to download the design (the device may be automatically detected). You may need to click Scan Chain to detect your target device.

4. In the JTAG Chain Configuration pane, click to browse for the SRAM Object File (.sof) that includes the In-System Sources and Probes instance or instances. (The .sof may be automatically detected).

5. Click Program Device to program the target device.
**Instance Manager**

The **Instance Manager** pane provides a list of all ALTSOURCE_PROBE instances in the design and allows you to configure how data is acquired from or written to those instances.

The following buttons and sub-panes are provided in the **Instance Manager** pane:

- **Read Probe Data**—Samples the probe data in the selected instance and displays the probe data in the **In-System Sources and Probes Editor** pane.

- **Continuously Read Probe Data**—Continuously samples the probe data of the selected instance and displays the probe data in the **In-System Sources and Probes Editor** pane; you can modify the sample rate via the **Probe read interval** setting.

- **Stop Continuously Reading Probe Data**—Cancels continuous sampling of the probe of the selected instance.

- **Write Source Data**—Writes data to all source nodes of the selected instance.

- **Probe Read Interval**—Displays the sample interval of all the In-System Sources and Probe instances in your design; you can modify the sample interval by clicking **Manual**.

- **Event Log**—Controls the event log in the **In-System Sources and Probes Editor** pane.

The status of each instance is also displayed beside each entry in the **Instance Manager** pane. The status indicates if the instance is **Not running**, **Offloading data**, **Updating data**, or if an **Unexpected JTAG communication error** occurs. This status indicator provides information about the sources and probes instances in your design.

**In-System Sources and Probes Editor Pane**

The **In-System Sources and Probes Editor** pane allows you to view data from all sources and probes in your design. The data is organized according to the index number of the instance. The editor provides an easy way to manage your signals, and allows you to rename signals or group them into buses. All data collected from in-system source and probe nodes is recorded in the event log and you can view the data as a timing diagram.

**Reading Probe Data**

You can read data by selecting the ALTSOURCE_PROBE instance in the **Instance Manager** pane and clicking **Read Probe Data**. This action produces a single sample of the probe data and updates the data column of the selected index in the **In-System Sources and Probes Editor** pane. You can save the data to an event log by turning on the **Save data to event log** option in the **Instance Manager** pane.

If you want to sample data from your probe instance continuously, in the **Instance Manager** pane, click the instance you want to read, and then click **Continuously read probe data**. While reading, the status of the active instance shows **Unloading**. You can read continuously from multiple instances.
You can access read data with the shortcut menus in the **Instance Manager** pane.

To adjust the probe read interval, in the **Instance Manager** pane, turn on the **Manual** option in the **Probe read interval** sub-pane, and specify the sample rate in the text field next to the **Manual** option. The maximum sample rate depends on your computer setup. The actual sample rate is shown in the **Current interval** box. You can adjust the event log window buffer size in the **Maximum Size** box.

**Writing Data**

To modify the source data you want to write into the ALTSOURCE_PROBE instance, click the name field of the signal you want to change. For buses of signals, you can double-click the data field and type the value you want to drive out to the ALTSOURCE_PROBE instance. The In-System Sources and Probes Editor stores the modified source data values in a temporary buffer. Modified values that are not written out to the ALTSOURCE_PROBE instances appear in red. To update the ALTSOURCE_PROBE instance, highlight the instance in the **Instance Manager** pane and click **Write source data**. The **Write source data** function is also available via the shortcut menus in the **Instance Manager** pane.

The In-System Sources and Probes Editor provides the option to continuously update each ALTSOURCE_PROBE instance. Continuous updating allows any modifications you make to the source data buffer to also write immediately to the ALTSOURCE_PROBE instances. To continuously update the ALTSOURCE_PROBE instances, change the **Write source data** field from **Manually** to **Continuously**.

**Organizing Data**

The **In-System Sources and Probes Editor** pane allows you to group signals into buses, and also allows you to modify the display options of the data buffer.

To create a group of signals, select the node names you want to group, right-click and select **Group**. You can modify the display format in the Bus Display Format and the Bus Bit order shortcut menus.

The **In-System Sources and Probes Editor** pane allows you to rename any signal. To rename a signal, double-click the name of the signal and type the new name.

The event log contains a record of the most recent samples. The buffer size is adjustable up to 128k samples. The time stamp for each sample is logged and is displayed above the event log of the active instance as you move your pointer over the data samples.

You can save the changes that you make and the recorded data to a Sources and Probes File (**.spf**). To save changes, on the File menu, click **Save**. The file contains all the modifications you made to the signal groups, as well as the current data event log.

**Tcl interface for the In-System Sources and Probes Editor**

To support automation, the In-System Sources and Probes Editor supports the procedures described in this chapter in the form of Tcl commands. The Tcl package for the In-System Sources and Probes Editor is included by default when you run **quartus_stp**.
The Tcl interface for the In-System Sources and Probes Editor provides a powerful platform to help you debug your design. The Tcl interface is especially helpful for debugging designs that require toggling multiple sets of control inputs. You can combine multiple commands with a Tcl script to define a custom command set.

For more information about Tcl scripting, refer to the Tcl Scripting chapter in volume 2 of the Quartus II Handbook. For more information about settings and constraints in the Quartus II software, refer to the Quartus II Settings File Manual. For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

Table 16–2 shows the Tcl commands you can use instead of the In-System Sources and Probes Editor.

**Table 16–2. In-System Sources and Probes Tcl Commands**

<table>
<thead>
<tr>
<th>Command</th>
<th>Argument</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>start_insystem_source_probe</td>
<td>-device_name &lt;device name&gt; -hardware_name &lt;hardware name&gt;</td>
<td>Opens a handle to a device with the specified hardware. Call this command before starting any transactions.</td>
</tr>
<tr>
<td>get_insystem_source_probe_instance_info</td>
<td>-device_name &lt;device name&gt; -hardware_name &lt;hardware name&gt;</td>
<td>Returns a list of all ALTSOURCE_PROBE instances in your design. Each record returned is in the following format: (index, source width, probe width, instance name)</td>
</tr>
<tr>
<td>read_probe_data</td>
<td>-instance_index &lt;instance_index&gt; -value_in_hex (optional)</td>
<td>Retrieves the current value of the probe. A string is returned that specifies the status of each probe, with the MSB as the left-most bit.</td>
</tr>
<tr>
<td>read_source_data</td>
<td>-instance_index &lt;instance_index&gt; -value_in_hex (optional)</td>
<td>Retrieves the current value of the sources. A string is returned that specifies the status of each source, with the MSB as the left-most bit.</td>
</tr>
<tr>
<td>write_source_data</td>
<td>-instance_index &lt;instance_index&gt; -value &lt;value&gt; -value_in_hex (optional)</td>
<td>Sets the value of the sources. A binary string is sent to the source ports, with the MSB as the left-most bit.</td>
</tr>
<tr>
<td>end_interactive_probe</td>
<td>None</td>
<td>Releases the JTAG chain. Issue this command when all transactions are finished.</td>
</tr>
</tbody>
</table>

Example 16–1 shows an excerpt from a Tcl script with procedures that control the ALTSOURCE_PROBE instances of the design as shown in Figure 16–3. The example design contains a DCFIFO with ALTSOURCE_PROBE instances to read from and write to the DCFIFO. A set of control muxes are added to the design to control the flow of data to the DCFIFO between the input pins and the ALTSOURCE_PROBE
instances. A pulse generator is added to the read request and write request control
lines to guarantee a single sample read or write. The ALTSOURCE_PROBE instances,
when used with the script in Example 16–1, provide visibility into the contents of the
FIFO by performing single sample write and read operations and reporting the state
of the full and empty status flags.

Use the Tcl script in debugging situations to either empty or preload the FIFO in your
design. For example, you can use this feature to preload the FIFO to match a trigger
condition you have set up within the SignalTap II Logic Analyzer.

Figure 16–3. A DCFIFO Example Design Controlled by the Tcl Script in Example 16–1
Example 16–1. Tcl Script Procedures for Reading and Writing to the DCFIFO in Figure 16–3 (Part 1 of 2)

## Setup USB hardware - assumes only USB Blaster is installed and
## an FPGA is the only device in the JTAG chain

set usb [lindex [get_hardware_names] 0]
set device_name [lindex [get_device_names -hardware_name $usb] 0]

## write procedure : argument value is integer

proc write {value} {
    global device_name usb
    variable full

    start_insystem_source_probe -device_name $device_name -hardware_name $usb

    #read full flag
    set full [read_probe_data -instance_index 0]

    if {$full == 1} {end_insystem_source_probe
    return "Write Buffer Full"
    }

}
The In-System Sources and Probes Editor can help you create a virtual front panel during the prototyping phase of your design. You can create relatively simple, high functioning designs in a short amount of time. The following PLL reconfiguration example demonstrates how to use the In-System Sources and Probes Editor to provide a GUI to dynamically reconfigure a Stratix PLL.

Example 16–1. Tcl Script Procedures for Reading and Writing to the DC FIFO in Figure 16–3 (Part 2 of 2)

```tcl
###toggle select line, drive value onto port, toggle enable
###bits 7:0 of instance 0 is S_data[7:0]; bit 8 = S_write_req;
###bit 9 = Source_write_sel

## int2bits is custom procedure that returns a bitstring from an integer argument
write_source_data -instance_index 0 -value /[int2bits [expr 0x200 | $value]]
write_source_data -instance_index 0 -value [int2bits [expr 0x300 | $value]]

## clear transaction
write_source_data -instance_index 0 -value 0
end_insystem_source_probe

proc read {} {
    global device_name usb
    variable empty
    start_insystem_source_probe -device_name $device_name -hardware_name $usb

    ##read empty flag : probe port[7:0] reads FIFO output; bit 8 reads empty_flag
    set empty [read_probe_data -instance_index 1]
    if {[regexp {1........} $empty]} { end_insystem_source_probe return "FIFO empty" }

    ## toggle select line for read transaction
    ## Source_read_sel = bit 0; s_read_reg = bit 1
    ## pulse read enable on DC FIFO
    write_source_data -instance_index 1 -value 0x1 -value_in_hex
    write_source_data -instance_index 1 -value 0x3 -value_in_hex

    set x [read_probe_data -instance_index 1 ]
    end_insystem_source_probe
    return $x
}
```

**Design Example: Dynamic PLL Reconfiguration**

The In-System Sources and Probes Editor can help you create a virtual front panel during the prototyping phase of your design. You can create relatively simple, high functioning designs in a short amount of time. The following PLL reconfiguration example demonstrates how to use the In-System Sources and Probes Editor to provide a GUI to dynamically reconfigure a Stratix PLL.
Stratix PLLs allow you to dynamically update PLL coefficients during run time. Each enhanced PLL within the Stratix device contains a register chain that allows you to modify the pre-scale counters ($m$ and $n$ values), output divide counters, and delay counters. In addition, the ALTPLL_RECONFIG megafunction provides an easy interface to access the register chain counters. The ALTPLL_RECONFIG megafunction provides a cache that contains all modifiable PLL parameters. After you update all the PLL parameters in the cache, the ALTPLL_RECONFIG megafunction drives the PLL register chain to update the PLL with the updated parameters. Figure 16–4 shows a Stratix-enhanced PLL with reconfigurable coefficients.

Stratix II and Stratix III devices also allow you to dynamically reconfigure PLL parameters. For more information about these families, refer to the appropriate data sheet. For more information about dynamic PLL reconfiguration, refer to AN 282: Implementing PLL Reconfiguration in Stratix & Stratix GX Devices or AN 367: Implementing PLL Reconfiguration in Stratix II Devices.

The following design example uses an ALTSOURCE_PROBE instance to update the PLL parameters in the ALTPLL_RECONFIG megafunction cache. The ALTPLL_RECONFIG megafunction connects to an enhanced PLL in a Stratix FPGA to drive the register chain containing the PLL reconfigurable coefficients. This design example uses a Tcl/Tk script to generate a GUI where you can enter in new $m$ and $n$ values for the enhanced PLL. The Tcl script extracts the $m$ and $n$ values from the GUI, shifts the values out to the ALTSOURCE_PROBE instances to update the values in the
ALTPLL_RECONFIG megafunction cache, and asserts the reconfiguration signal on the ALTPLL_RECONFIG megafunction. The reconfiguration signal on the ALTPLL_RECONFIG megafunction starts the register chain transaction to update all PLL reconfigurable coefficients. A block diagram of a design example is shown in Figure 16–5. The Tk GUI is shown in Figure 16–6.

**Figure 16–5. Block Diagram of Dynamic PLL Reconfiguration Design Example**

![Block Diagram of Dynamic PLL Reconfiguration Design Example]

This design example was created using a Nios® II Development Kit, Stratix Edition. The file `sourceprobe_DE_dynamic_pll.zip` contains all the necessary files for running this design example, including the following:

- **Readme.txt**—A text file that describes the files contained in the design example and provides instructions about running the Tk GUI shown in Figure 16–6.

- **Interactive_Reconfig.qar**—The archived Quartus II project for this design example.

Download the **In-System Sources and Probes Example** from the On-chip Debugging Design Examples page of the Altera website.
Conclusion

The In-System Sources and Probes Editor provides stimuli and receives responses from the target design during run time. With the simple and intuitive interface, you can add virtual inputs to your design during run time without using external equipment. When used in conjunction with the SignalTap II Logic Analyzer, you can use the In-System Sources and Probes Editor to obtain greater control of the signals in your design, and thus help shorten the verification cycle.

Document Revision History

Table 16–3 shows the revision history for this chapter.

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2010</td>
<td>12.0.0</td>
<td>Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>10.1.1</td>
<td>Template update.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Minor corrections. Changed to new document template.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Minor corrections.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>■ Removed references to obsolete devices.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Style changes.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>No change to content.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>Changed to 8-1/2 x 11 page size. No change to content.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Documented that this feature does not support simulation on page 17–5</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Figure 17–8 for Interactive PLL reconfiguration manager</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added hyperlinks to referenced documents throughout the chapter</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Minor editorial updates</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
Section V. Formal Verification

The Quartus® II software easily interfaces with EDA formal design verification tools such as the Cadence Encounter Conformal and Synopsys Synplify software. In addition, the Quartus II software has built-in support for verifying the logical equivalence between the synthesized netlist from Synopsys Synplify and the post-fit Verilog Quartus Mapped (.vqm) files using Cadence Encounter Conformal software.

This section discusses formal verification, how to set-up the Quartus II software to generate the .vqm file and Cadence Encounter Conformal script, and how to compare designs using Cadence Encounter Conformal software.

This section includes the following chapter:

■ Chapter 17, Cadence Encounter Conformal Support
This chapter describes equivalence checking with the Cadence Encounter Conformal Logic Equivalence Check (LEC) software. The Quartus® II software provides formal verification support for Altera® designs through interfaces with the Conformal LEC software.

Logic equivalence checking uses Boolean arithmetic techniques to compare the logical equivalence of two versions of the same design. You can use the Conformal LEC software to verify the functional equivalence of a post-synthesis Verilog Quartus Mapping (.vqm) netlist from the Synopsys Synplify Pro software, a post-fit Verilog Output File (.vo) from the Quartus II software, or both. You can also use the Conformal LEC software to verify the functional equivalence of the register transfer level (RTL) source code and post-fit .vo with the Quartus II software when using Quartus II integrated synthesis.

This chapter discusses the following topics:

- “Formal Verification Design Flow” on page 17–2
- “RTL Coding Guidelines for Quartus II Integrated Synthesis” on page 17–4
- “Black Boxes in the Conformal LEC Flow” on page 17–8
- “Generating the Post-Fit Netlist Output File and the Conformal LEC Setup Files” on page 17–9
- “Understanding the Formal Verification Scripts for the Conformal LEC Software” on page 17–12
- “Comparing Designs Using the Conformal LEC Software” on page 17–15
- “Known Issues and Limitations” on page 17–16
- “Black Box Models” on page 17–18
- “Conformal Dofile/Script Example” on page 17–19
Formal Verification Versus Simulation

Formal verification is not a replacement for vector-based simulation. Formal verification only complements the existing vector-based simulation techniques to speed up the verification cycle. Vector-based simulation techniques of gate-level designs can take a considerable amount of time.

You can use vector-based simulation techniques to perform the following functions:

- Verify design functionality
- Verify timing specifications
- Debug designs

Formal Verification: What You Must Know

There might be an impact on area and performance during recompilation of your design with the Quartus II software if you use the formal verification flow for the Conformal LEC software. The following factors might affect the area and performance of your design:

- Preserving hierarchy
- Implementing ROM by logic elements (LEs)
- Enabling retiming

Before you consider using the formal verification flow in your design methodology, refer to “Known Issues and Limitations” on page 17–16.

Formal Verification Design Flow

Altera supports formal verification with the Conformal LEC software for the following two synthesis tools:

- “Quartus II Integrated Synthesis” on page 17–3
- “Synplify Pro” on page 17–3

The following sections describe the supported design flows for these synthesis tools.
Quartus II Integrated Synthesis

Figure 17–1 shows the design flow for formal verification with Quartus II integrated synthesis. This flow performs equivalence checking of the RTL source code and the post-fit netlist generated by the Quartus II software. The RTL source code can be in Verilog HDL or VHDL format. The Quartus II-generated post-fit netlist is in Verilog HDL format.

Figure 17–1. Formal Verification Using Quartus II Integrated Synthesis and the Conformal LEC Software

EDA Tool Support for Quartus II Integrated Synthesis

The formal verification flow using the Quartus II software and Conformal LEC software supports the following software versions and operating systems:

- The Quartus II software beginning with version 4.2
- The Conformal LEC software beginning with version 4.3.5A
- Linux operating system

Synplify Pro

Figure 17–2 shows the design flow for formal verification with Synplify Pro Synthesis performing equivalency checking for the post-synthesis netlist from Synplify Pro and the post-fit netlist generated by Quartus II software.
For more information about performing equivalence checking between RTL source code and post-synthesis netlists generated from the Synplify Pro software, refer to the Synplify Pro documentation.

Figure 17–2. Formal Verification Flow Using Synplify Pro and the Conformal LEC Software

RTL Coding Guidelines for Quartus II Integrated Synthesis

The Conformal LEC software compares the RTL source code against the Quartus II-generated post-fit netlist. The Conformal LEC software and Quartus II integrated synthesis parse and compile the RTL description differently. Quartus II integrated synthesis supports some RTL features that the Conformal LEC software does not support and vice versa. The style of the RTL source code is of particular concern because neither tool supports every construct, leading to potential formal verification mismatches. For example, different encoding mechanisms for state machine extraction can result in different structures. Therefore, Quartus II integrated synthesis and the Conformal LEC software must interpret the RTL source code in the same manner for successful verification.

The following section describes how you can identify and prevent problems that may occur in the formal verification flow.

For more information about RTL coding styles for Quartus II integrated synthesis, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.

Some of the coding guidelines apply to both Quartus II integrated synthesis and Synplify Pro flow, as indicated in each of the guidelines in the following sections.
**Synthesis Directives and Attributes**

You can use synthesis directives, also known as pragmas, to compare and verify the RTL source codes against the post-fit .vo netlist from the Quartus II software.

Quartus II integrated synthesis and the Conformal LEC software support the “synthesis” and “synopsys” trigger keywords. When Quartus II integrated synthesis does not recognize a keyword (such as “verplex”), the Quartus II software disables the keyword in the formal verification scripts produced for use with the Conformal LEC software. Therefore, you must use caution with unsupported pragmas because the unsupported pragmas can lead to verification mismatches.

Example 17–1 and Example 17–2 show that you can use Quartus II integrated synthesis to synthesize an RTL source code with the `read_comments_as_HDL()` synthesis directive.

**Example 17–1. Verilog HDL Example of Read Comments as HDL**

```verilog
// synthesis read_comments_as_HDL on
// my_rom lpm_rom (.address (address),
// .data (data));
// synthesis read_comments_as_HDL off
```

**Example 17–2. VHDL Example of Read Comments as HDL**

```vhdl
-- synthesis read_comments_as_HDL on
-- my_rom : entity lpm_rom
-- port map (  
-- address => address,  
-- data => data, );
-- synthesis read_comments_as_HDL off
```

The Conformal LEC software does not support the `read_comments_as_HDL` synthesis directive, and the directive does not affect the Conformal LEC software.

Table 17–1 lists supported pragmas and trigger keywords for formal verification.

**Table 17–1. Supported Pragmas and Trigger Keywords for Formal Verification**

<table>
<thead>
<tr>
<th>Pragmas</th>
<th>Trigger Keywords</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>full_case</code></td>
<td>synthesis</td>
</tr>
<tr>
<td><code>parallel_case</code></td>
<td>synopsys</td>
</tr>
<tr>
<td><code>pragma</code></td>
<td></td>
</tr>
<tr>
<td><code>synthesis_off</code></td>
<td></td>
</tr>
<tr>
<td><code>synthesis_on</code></td>
<td></td>
</tr>
<tr>
<td><code>translate_off</code></td>
<td></td>
</tr>
<tr>
<td><code>translate_on</code></td>
<td></td>
</tr>
</tbody>
</table>

Do not use Verilog 2001-style pragma declarations. The Quartus II software and the Conformal LEC software support this style of pragma differently.
Fixed-Output Registers

Quartus II integrated synthesis and Synplify Pro eliminate registers that have fixed output. Quartus II integrated synthesis issues a warning message and adds an entry to the corresponding report panel in the formal verification folder of the Analysis & Synthesis section of the Compilation Report. If the Conformal LEC software does not find the same optimizations, the result can lead to unmapped points in the golden netlist. Example 17–3 shows logic causing register outputs to be fixed at a constant value.

Example 17–3. Verilog HDL Example Showing Fixed Register Outputs

```verilog
module stuck_at_example(clk, a, b, c, d, out);
    input a, b, c, d, clk;
    output out;
    reg e, f, g;
    always @(posedge clk) begin
        e <= a and g; // e is stuck at 0
        g <= c and e; // g is stuck at 0
        f <= e | b;
    end
    assign out = f and d;
endmodule
```

In this module description, registers e and g are tied to logic 0. In this example, the Quartus II software generates the following warning message:

Warning: Reduced register "g" with stuck data_in port to stuck value GND
Warning: Reduced register "e" with stuck data_in port to stuck value GND

Example 17–4 shows that Quartus II integrated synthesis adds a command to the formal verification scripts to inform the Conformal LEC software that a register is stuck at a constant value.

Example 17–4. Conformal LEC Script Showing Commands for Instance Equivalence

```verbatim
// report floating signals
// Instance-constraints commands for constant-value registers removed during compilation
// add instance constraints 0 e -golden
// add instance constraints 0 g -golden
```

Quartus II integrated synthesis comments the command in the formal verification script to force the Conformal LEC software to treat the register as stuck at a constant value and potentially hides a compilation error. You must verify that input to the e and g registers is constant in your design and uncomment the command to obtain accurate results.

Altera recommends recoding your design to eliminate registers that have fixed output.
ROM, LPM_DIVIDE, and Shift Register Inference

For formal verification, Quartus II integrated synthesis implements ROM and shift registers with LEs instead of with dedicated on-chip memory resources. Using LEs can be less area efficient than inferring a megafunction that you can implement in a RAM block. The Quartus II software generates a warning message to indicate that the software does not infer the megafunction. Quartus II integrated synthesis also reports a suggested ROM or shift register instantiation that enables you to either use the MegaWizard™ Plug-In Manager to create the appropriate megafunction explicitly, or to isolate the corresponding logic in a separate entity that you can set as a black box. By setting black box properties on a module or a particular entity, you are directing the formal verification tool not to look inside the module or entity for formal verification. If you set the black box properties on the corresponding megafunction before synthesis, you can verify the megafunction with the Conformal LEC software. For more information about setting black box properties on a particular module, refer to Table 17–2 on page 17–9.

If your design contains division functionality, the Quartus II software infers an LPM_DIVIDE megafunction. The Quartus II software treats the inferred LPM_DIVIDE megafunction as a black box for formal verification.

RAM Inference

When the Quartus II software infers the ALTSYNTCRAM megafunction from the RTL source code, the Quartus II software generates the following warning message:

Created node "<mem_block_name>" as a RAM by generating altsyncram megafunction to implement register logic with M512 or M4K memory block or M-RAM. Expect to get an error or a mismatch for this block in the formal verification tool.

The Quartus II software generates this warning message because the memory block (altsyncram) is a new instance in the post-fit netlist. The Quartus II software handles the ALTSYNTCRAM megafunction as a black box by the formal verification tool. However, no such instance exists in the original RTL design, resulting in mismatch or error reporting in the formal verification tool.

Latch Inference

A combinational feedback loop implements a latch in Quartus II integrated synthesis. The Conformal LEC software infers a latch primitive in the Conformal LEC software library to implement a latch. This results in having a library on the golden side and a combinational loop with a cut point on the revised side, leading to verification mismatches. The Quartus II software issues a warning message whenever the Conformal LEC software infers a latch. The Quartus II software then adds an entry to the report panel in the Formal Verification folder of the Analysis & Synthesis report.

Altera recommends that you avoid latches in your design; however, if latches are necessary, Altera recommends using the LPM_LATCH megafunction.

For more information about latches, refer to the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook.
Combinational Loops

If your design consists of an intended combinational loop, you must define an appropriate cut point for both the RTL and the post-fit .vo netlist. You can find a warning indicating that a combinational loop exists in your design in the Formal Verification subfolder of the Quartus II software Analysis & Synthesis report.

For more information about issues with combinational loops, refer to “Known Issues and Limitations” on page 17–16.

Finite State Machine Coding Styles

When the Conformal LEC software infers a state machine, the state machine uses sequential encoding as the default encoding in the absence of user encoding. The Quartus II software selects the encoding most suited for the inferred state machine if you set the State Machine Processing setting to the default value (Auto). To do this, follow these steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
2. In the Category list, select Analysis & Synthesis Settings. The Analysis & Synthesis Settings page appears.
4. Under Existing Option Settings, in the Name list, select State Machine Processing. In the Setting list, select Auto.
5. In the More Analysis & Synthesis Settings dialog box, click OK.
6. Click OK.

Use the coding style described in the Recommended HDL Coding Styles chapter in volume 1 of the Quartus II Handbook when writing finite state machines (FSMs). The coding style in the specified chapter allows Quartus II integrated synthesis and the Conformal LEC software to infer a similar state machine for the same RTL source code.

Black Boxes in the Conformal LEC Flow

The Quartus II software generates a flattened netlist; however, you must treat the following modules in your design as black boxes:

- LPMs and megafuntions without formal verification models
- Encrypted IP functions
- Entities not implemented in Verilog HDL or VHDL

To perform equivalence checking of a design between its version, which consists of the modules listed above and its implemented version, the Conformal LEC software must treat these modules as black boxes. To facilitate the formal verification flow, the Quartus II software reconstructs the hierarchy of the black boxes with a port interface that is identical to the module on the golden side of your design.
If your golden netlist (.vqm netlist from Synplify Pro or RTL) includes any design entity not having a corresponding formal verification model, the software treats that entity as a black box with its boundary interface preserved. Table 17–2 on page 17–9 lists three types of black boxes with corresponding required actions.

The Quartus II-generated .vo contains the black box hierarchy when you make an EDA Formal Verification Hierarchy assignment with the value BLACKBOX.

If you do not make this assignment for a module, the Quartus II software implements that module in logic cells. When this happens, the .vo netlist no longer contains the black box hierarchy and does not preserve the port interface, resulting in a mismatch in the Conformal LEC software.

<table>
<thead>
<tr>
<th>Type of Black Box</th>
<th>Required Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>Altera library of parameterized modules (LPMs) and megafunctions.</td>
<td>No action required. The Quartus II software automatically creates a black box list of components and preserves the hierarchy.</td>
</tr>
<tr>
<td>Any parameterized entity other than the parameterized entities listed in the Guidelines for Creating a Design for Use with the Encounter Conformal and Quartus II Software topic in Quartus II Help.</td>
<td>You must designate the wrapper that instantiates the parameterized entity as a black box.</td>
</tr>
<tr>
<td>Non-parameterized entities that you want to designate as a black box.</td>
<td>You can designate the entity itself as a black box.</td>
</tr>
</tbody>
</table>

You can also use the Quartus II GUI to set the black box property on the entities, which the formal verification tool does not compare.

To preserve the boundary interface of an entity using the GUI, make an EDA Formal Verification Hierarchy assignment to the entity with the value BLACKBOX.

**Generating the Post-Fit Netlist Output File and the Conformal LEC Setup Files**

The following steps describe how to set up the Quartus II software environment to generate the post-fit .vo netlist and the Conformal LEC script for use in formal verification. With the exception of step 2, the steps are identical for both of the synthesis tools:

To create a new Quartus II project or open an existing project, follow these steps:

1. On the Assignments menu, click Settings. The Settings dialog box appears.
2. In the Category list, click EDA Tool Settings.
   If you are using Quartus II integrated synthesis, follow these steps:
   a. In the Category list, under EDA Tool Settings, select Design Entry/Synthesis.
      Select <None> from the Tool name list.
   b. In the Category list, under EDA Tool Settings, select Formal Verification.
      Select Conformal LEC from the Tool name list.
   If you are using Synplify Pro, follow these steps:
   a. In the Category list, under EDA Tool Settings, select Design Entry/Synthesis.
      Select Synplify Pro from the Tool name list.
   b. In the Category list, under EDA Tool Settings, select Formal Verification.
      Select Conformal LEC from the Tool name list.
3. In the Category list, click Incremental Compilation under Compilation Process Settings. The Incremental Compilation page appears.
4. Type the following Tcl command in the Quartus II software Tcl console to turn on the incremental compilation feature:
   ```tcl
   set_global_assignment -name INCREMENTAL_COMPILATION FULL_INCREMENTAL_COMPILATION
   ```
   Altera recommends that you turn on the incremental compilation feature for formal verification, and that your design does not contain any partition that you created. The incremental compilation feature is on by default.
5. In the Category list, click Physical Synthesis Optimizations under Compilation Process Settings. The Physical Synthesis Optimizations page appears.
6. Turn off Perform register retiming.
   If you do not turn off Perform register retiming, an error occurs during compilation: “Physical Netlist Optimization Register retiming is not supported by Formal Verification tool Conformal LEC”.
7. Under Optimize for fitting (physical synthesis for density), turn off Perform physical synthesis for combinational logic and Perform logic to memory mapping to prevent the software from mapping logic to RAMs.
   Retiming a design, either during the synthesis step or during the fitting step, usually results in moving and merging registers along the critical path and is not supported by the equivalence checking tools. Because equivalence checkers compare the logic cone terminating at registers, do not use retiming to move the registers during optimization in the Quartus II software.
   For more information about physical synthesis, refer to the Netlist Optimizations and Physical Synthesis chapter in volume 2 of the Quartus II Handbook.
8. Perform a full compilation of your design. On the Processing menu, click Start Compilation, or click the Start Compilation icon on the toolbar.
Quartus II Software Generated Files, Formal Verification Scripts, and Directories

After successful compilation, the Quartus II software generates a list of files, formal verification scripts, and directories in the `<project_directory>/fv/conformal/` directory (Table 17–3).

Table 17–3. Quartus II Software Compiler-Generated Files and Directories

<table>
<thead>
<tr>
<th>File or Directory</th>
<th>Name</th>
<th>Details</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&lt;proj rev&gt;.ctc</code></td>
<td>The <code>&lt;proj rev&gt;.ctc</code> references <code>&lt;proj rev&gt;.clg</code> and <code>&lt;proj rev&gt;.clr</code> that read the library files and black box descriptions. The <code>&lt;proj rev&gt;.ctc</code> also references the <code>&lt;proj rev&gt;.cmc</code> containing information about the mapped points. Use the <code>&lt;proj rev&gt;.ctc</code> with the Conformal LEC software.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.cec</code></td>
<td>The <code>&lt;proj rev&gt;.cec</code> contains information for instance equivalences.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.cep</code></td>
<td>The <code>&lt;proj rev&gt;.cep</code> contains information for black box pin equivalences in your design.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.cmp</code></td>
<td>The <code>&lt;proj rev&gt;.cmp</code> contains information for the black box pin mapping between the golden and revised sides. The Quartus II software calls the <code>&lt;proj rev&gt;.cmp</code> from the <code>&lt;proj rev&gt;.ctc</code> script file. By default, the line in which this file is called is commented out. This file is useful only for HardCopy II device family.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.cmc</code></td>
<td>The <code>&lt;proj rev&gt;.cmc</code> contains information about the additional points that the Quartus II software maps in addition to the points that the tool selects.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;_trivial.cmc</code></td>
<td>This <code>&lt;proj rev&gt;_trivial.cmc</code> contains mapping information for all the key points in your design. Sometimes, the Conformal LEC software performs incorrect key point mapping, resulting in formal verification mismatches. To overcome the verification mismatches, the Quartus II software writes out the <code>&lt;proj rev&gt;_trivial.cmc</code> that contains mapping information for all the key points in your design. Reading this file during the formal verification setup can result in increased run time. Therefore, the Quartus II software writes out the top-level script file <code>&lt;proj rev&gt;.ctc</code> with the command to read the <code>&lt;proj rev&gt;_trivial.cmc</code> commented out. If the formal verification results are not acceptable, you can uncomment the command and read the <code>&lt;proj rev&gt;_trivial.cmc</code>. The command in the <code>&lt;proj rev&gt;.ctc</code> is:</td>
<td></td>
</tr>
<tr>
<td></td>
<td>//Trivial mappings with same name registers</td>
<td></td>
</tr>
<tr>
<td></td>
<td>//read mapped points $PROJECT/fv/conformal/&lt;proj rev&gt;_trivial.cmc</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.clr</code></td>
<td>The <code>&lt;proj rev&gt;.clr</code> contains information about the macros and libraries for the revised design.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;proj rev&gt;.clg</code></td>
<td>The <code>&lt;proj rev&gt;.clg</code> contains information about the macros and libraries for the golden design.</td>
<td></td>
</tr>
<tr>
<td><code>&lt;project directory&gt;/fv/conformal/&lt;project rev&gt;_blackboxes</code></td>
<td>This directory contains top-level module descriptions for all the user-defined black box entities and contains modules with definitions other than Verilog HDL or VHDL, for example, in your design directory <code>&lt;project directory&gt;/fv/conformal/&lt;project rev&gt;_blackboxes</code></td>
<td></td>
</tr>
<tr>
<td><code>.vo netlist file</code></td>
<td><code>&lt;proj rev&gt;.vo</code></td>
<td>The Quartus II software-generated netlist for formal verification.</td>
</tr>
</tbody>
</table>
The script file contains the setup and constraints information to use with the formal verification tool. The `<entity>.v` in the `blackboxes` directory contains the module description of entities that you do not define in the formal verification library. The file also contains entities that you treat as black boxes. For example, if a reference to a black box for an instance of the ALTDPRAM megafuction in your design is present, the `blackboxes` directory does not contain a module description for the ALTDPRAM megafuction because you define the module description in the `altdpram.v` of the formal verification library. When a module does not have an RTL description, or the description exists only in the formal verification library and you do not want to compare the module with formal verification, a file containing only the top-level module description with port declaration is written out to the `blackboxes` directory and read into the Conformal LEC software. To learn more about black boxes, refer to “Black Boxes in the Conformal LEC Flow” on page 17–8.

### Understanding the Formal Verification Scripts for the Conformal LEC Software

The Quartus II software generates scripts to use with the Conformal LEC software. This section describes the details of the Conformal LEC commands in the scripts to help you compare the revised netlist with the golden netlist. Usually, you do not have to add anymore Conformal LEC constraints to verify your netlists.

You can view a sample Quartus II software generated script in “Conformal Dofile/Script Example” on page 17–19.

### Conformal LEC Commands in the Quartus II Software Generated Scripts

The value for the variable `QUARTUS` is the path to the Quartus II software installation directory:

```bash
setenv QUARTUS <Quartus Installation Directory>
```

The Quartus II software assigns the current working directory of your project to the `PROJECT` variable. Use this variable to change your project directory to the directory in which you install your design files when moving from a UNIX to a Windows environment, or vice versa:

```bash
setenv PROJECT <Quartus Project Directory>
```

The following command reads both the golden and the revised netlists, along with the appropriate library models:

```bash
read design <design files>
```

You must update your project location when you move the files from the Windows environment to the UNIX environment.

The post placement and routing netlist from the Quartus II software might contain net and instance names that are slightly different from net and instance names of the golden netlist. With the following command, the Quartus II software defines temporary substitute string patterns enabling the Conformal LEC software to map key points automatically when the names are different:

```bash
add renaming rule <rule>
```
The Conformal LEC software employs three name-based methods to map key points to compare the revised netlist with the golden netlist. Scripts set the correct method to get the best results.

```
set mapping method <mapping_rule>
```

The Quartus II software performs several optimizations, including optimizing the registers whose input is driven by a constant. Under these circumstances, for the formal verification software to compare the netlists properly, use the command set flatten model with the option seq_constant.

```
set flatten model <flattening_rule>
```

When you use the report black box command, verify that the software lists the following modules as black boxes, along with any of the modules that you treat as black boxes in the golden and revised netlists:

- LPMs and megafunctions without the formal verification models
- Encrypted IP functions
- Entities not implemented in Verilog HDL or VHDL

Use the following command to set the same implementation on multipliers for both the golden and revised netlists:

```
set multiplier implementation <implementation_name>
```

If combinational loops or instances of LPM_LATCH are present, the Quartus II software cuts the loop at the same point using the following command on both the golden and revised netlists:

```
add cut point
```

The Conformal LEC software does not always automatically map all the key points, or can incorrectly map some key points. To help the Conformal LEC software successfully complete the mapping process, the Quartus II software records optimizations performed on the netlist as a series of add mapped points in the Conformal LEC <file_name>.cmc script.

```
add mapped points <key_points>
```

When the software moves the inverter before the register to after the register, use the following command:

```
add mapped points <key_points> -invert
```
The following command reads in the mapped point information from the specified file:

```
read mapped points <file_name>.cmc
```

**Figure 17–3. Instance Equivalence**

During optimization, the Quartus II software might merge two registers into one (Figure 17–3). The Quartus II software informs the formal verification tool that the U1 and U2 registers are equivalent to each other using the following command:

```
add instance equivalence <instance_pathname ..> [-golden]
```

When register duplication happens, use the following command:

```
add instance equivalence <instance_pathname ..> [-revised]
```

When the software moves the inverter beyond the register along with either register duplication or merging, use the following command:

```
add instance equivalences <instance_pathname>
[-invert <instance_pathname>]
```

Sometimes, the software drives the register output to a constant, either logic 0 or logic 1. The Quartus II software sets the value of the register to a constraint using the add instance constraint command. For more information about this command, refer to “Fixed-Output Registers” on page 17–6.

```
add instance constraint <constraint_value>
```
Comparing Designs Using the Conformal LEC Software

This section describes using the Conformal LEC software to compare designs, and to prove logical equivalence between two versions of your design.

Running the Conformal LEC Software from the GUI

To run the Conformal LEC software from the GUI, follow these steps:

1. Open the Conformal LEC software.
2. On the File menu, click Do Dofile.
3. Select the `<path to project directory>/fv/conformal/<proj rev>.ctc`.

The Conformal LEC software GUI displays the comparison results. The Golden window displays the original RTL description or the post synthesis .vqm netlist from Synplify Pro, and the Revised window displays the information from the post-fit netlist generated by the Quartus II software. The message section at the bottom of the window reports the verification results and the number of unmapped and non-equivalent points found in your design.

To investigate the verification results, click the Mapping Manager icon in the toolbar, or on the Tools menu, click Mapping Manager. The Conformal LEC software reports the mapped, unmapped, and compared points in the Mapped Points, Unmapped Points, and Compared Points windows, respectively.

For more information about how to diagnose non-equivalent points, refer to the Conformal LEC software user documentation.

Running the Conformal LEC Software From a System Command Prompt

To run the Conformal LEC software without using the GUI, type the command shown in Example 17–5 at a system command prompt.

Example 17–5. Conformal LEC Command to Run Formal Verification

```
lec -dofile /<path to project directory>/fv/conformal/<proj rev>.ctc -nogui
```

To get a downloadable design example showing the formal verification flow with Quartus II software, refer to the Formal Verification Design Example page of the Altera website.

For more information about the latest debugging tips and solutions for formal verification flow between the Conformal LEC software and the Quartus II software, go to www.altera.com and perform an advanced search with keywords “formal verification”.

Known Issues and Limitations

The following known issues and limitations can occur when using the formal verification flow described in this chapter:

- When a port on a black box entity drives two or more signals in the black box, the Quartus II software pushes the connections outside of the black box, and creates the same number of ports on the black box. This problem occurs only in Stratix II and HardCopy II designs.

The Quartus II software names the additional ports on the black box as _unassoc_inputs_[] and _unassoc_outputs_[] (Figure 17–4). This issue occurs with reset and enable signals. Figure 17–4 shows an example in which the reset pin splits into two ports outside of the black box and the clkctrl block drives the _unassoc_inputs_[] port. In such situations, the Quartus II-generated .vo netlist has signals driving these black box ports, but the golden RTL does not contain any signals to drive the _unassoc_inputs_[] port, which results in a formal verification mismatch of the black box. The black box module definition that the Quartus II software generates in the <Quartus_project>/fv/conformal*_blackboxes directory contains these additional _unassoc_inputs_[] and _unassoc_outputs_[] ports. The Quartus II software reads this black box module on the golden and revised sides of your design, which results in unconnected ports on the golden side and formal verification mismatches.

Figure 17–4 shows the creation of the _unassoc_inputs_[] and _unassoc_outputs_[] ports for the reset signal.

Figure 17–4. Creation of _unassoc_inputs_[] and _unassoc_outputs_[]
Another common occurrence of this issue is in HardCopy II designs. Whenever a port drives large fan-out in the black box, the Quartus II software inserts a buffer on the net and moves the logic outside of the black box (Figure 17–5).

To fix the problem of _unassoc_input_[] ports causing black box mismatches, use Conformal LEC commands to change the type of the black box _unassoc_input_[] keypoint to a primary output keypoint, and then mark the appropriate pin equivalences. Similarly, to fix the problem of register mismatches due to _unassoc_output_[] pins from black boxes, use Conformal LEC commands to change the type of the blackbox _unassoc_output_[] keypoint to a primary input, and then mark the equivalent pins as such. You can view the commands to perform these actions in the <proj rev>.cep.

Figure 17–5 shows the creation of _unassoc_inputs_[] for a signal with large fan-out.

Figure 17–5. Creation of _unassoc_inputs_[] for a Signal with Large Fan-out

- In designs with combinational feedback loops, the Conformal LEC software can insert extra cut points in the revised netlist, causing unmapped points and ultimately verification mismatches.

- For Cyclone II designs, the Conformal LEC software might report non-equivalent flipflops and extra cut points for the revised (post-fit) design under the following conditions:
  - When your HDL source code instantiates the lpm_ff primitive with an asynchronous load signal aload (with or without any other asynchronous control signals) and;
  - When you use the asynchronous clear signal aclr and asynchronous set signal aset together.

To avoid this problem, ensure that a wrapper module or entity is present around the lpm_ff instantiation, and black box the module or entity that instantiates the lpm_ff primitive.

- For Stratix III designs, the Conformal LEC software creates cut points for the combinational loops on the golden side and might fail equivalence checking due to improper mapping. The combinational loops are due to logic around the registers emulating multiple sets, resets, or both. The Quartus II software reports these cut points with warning messages during mapping. You can add Conformal LEC commands manually to add cut points, which can result in proper mapping and formal verification.
To perform formal verification, the Quartus II software turns off certain synthesis optimization options (such as register retiming, optimization through black box hierarchy boundaries, and disabling the ROM and shift register inference), which can have an impact on the area resource and performance.

In the Quartus II software version 9.0 and earlier, turning on gate-level register retiming as part of a formal verification flow might impact area and resource utilization.

When you do not verify RAM and ROM instantiations, inferences, or both using formal verification.

Incremental compilation for formal verification does not support user-created design partitions.

Formal verification does not support clear box netlists due to unconnected ports on its WYSIWYG instances.

Formal verification does not support VHDL megafuction variations due to undriven ports on the megafuctions.

When a black box contains bidirectional ports, the Quartus II software does not reconstruct the hierarchy. Therefore, a flat netlist represents the black box, which results in formal verification mismatches.

You must treat ROMs as black boxes in your design before compilation with Quartus II integrated synthesis, because the Quartus II software might perform some optimizations on the ROM, resulting in formal verification mismatches.

The Conformal LEC software might report mismatches or cancel comparisons of some key points when the Quartus II software implements a DSP megafuction in LEs, due to implicit optimizations in the DSP and the complexity of the multiplier logic in terms of LEs.

Unused logic optimized in and around a black box by the Quartus II software can result in a black-box interface different from the interface in the synthesized .vqm netlist.

Black Box Models

The black box models are interface definitions of entities, such as primitives, atoms, LPMs, and megafuctions. These models have a parameterized interface, and do not contain any definition of behavior. These models work with the Conformal LEC software, which uses these models along with your design to generate black boxes for instances of the entity with varying sets of parameters in your design.

For a complete list of supported black box models, refer to Guidelines for Creating a Design for Use with the Encounter Conformal and Quartus II Software in Quartus II Help.
Conformal Dofile/Script Example

Example 17–6 shows an example script, generated by the Quartus II software. The example script lists some of the setup commands in Conformal LEC software.

Example 17–6. Conformal LEC Script (Part 1 of 2)

```bash
// Copyright (C) 1991-2008 Altera Corporation
// Your use of Altera Corporation’s design tools, logic functions
// and other software and tools, and its AMPP partner logi
// functions, and any output files from any of the foregoing
// (including device programming or simulation files), and any
// associated documentation or information are expressly subject
// to the terms and conditions of the Altera Program License
// Subscription Agreement, Altera MegaCore Function License
// Agreement, or other applicable license agreement, including,
// without limitation, that your use is for the sole purpose of
// programming logic devices manufactured by Altera and sold by
// Altera or its authorized distributors. Please refer to the
// applicable agreement for further details.

// Script generated by the Quartus II software
reset
set system mode setup
set log file mfs_3prm_1a.fv.log -replace
set naming rule "%s" -register -golden
set naming rule "%s" -register -revised
// Naming rules for Verilog
set naming rule "%L.%s" "%L[%d].%s" "%s" -instance
set naming rule "%L.%s" "%L[%d].%s" "%s" -variable
// Naming rules for VHDL
// set naming rule "%L:%s" "%L:%d:%s" "%s" -instance
// set naming rule "%L:%s" "%L:%d:%s" "%s" -variable
// set undefined cell black_box -both
// These are the directives that are not supported by the QIS RTL to gates FV flow
set directive off verplex ambit
set directive off assertion_library black_box clock_hold compile_off compile_on
set directive off dc_script_begin dc_script_end divider enum infer_latch
set directive off mem_rowselect multi_port multiplier operand state_vector template
add notranslate module alt3pram -golden
add notranslate module alt3pram -revised
setenv QUARTUS /data/quark/build/ajaishan/quartus
setenv PROJECT /net/quark/build/ajaishan/quartus_regtest/eda/fv/conformal/synplify/
stratix/mfs_3prm_1a_v1_/mfs_3prm_1a/qu_allopt
```
Example 17–6. Conformal LEC Script (Part 2 of 2)

read design
   $QUARTUS/eda/fv_lib/vhdl/dummy.vhd
   -map lpm $QUARTUS/eda/fv_lib/vhdl/lpms
   -map altera_mf $QUARTUS/eda/fv_lib/vhdl/mfs
   -map stratix $QUARTUS/eda/fv_lib/vhdl/stratix
   -vhdl -noelaborate -golden
read design
   -file $PROJECT/fv/conformal/mfs_3prm_1a.clg
   $PROJECT/p3rm_block.v
   $PROJECT/mfs_3prm_1a.v
   -verilog2k -merge none -golden
read design
   $QUARTUS/eda/fv_lib/vhdl/dummy.vhd
   -map lpm $QUARTUS/eda/fv_lib/vhdl/lpms
   -map altera_mf $QUARTUS/eda/fv_lib/vhdl/mfs
   -map stratix $QUARTUS/eda/fv_lib/vhdl/stratix
   -vhdl -noelaborate -revised
read design
   -file $PROJECT/fv/conformal/mfs_3prm_1a.clg
   $PROJECT/fv/conformal/mfs_3prm_1a.cmc
   // add ignored inputs _unassoc_inputs_* -all -revised
add renaming rule r1 "/\" "/" -revised
add renaming rule r2 "\" "/" -revised
set multiplier implementation rca -golden
set multiplier implementation rca -revised
set mapping method -name first
set mapping method -nounreach
set mapping method -noreport_unreach
set mapping method -nounreach
set mapping method -nodoc_name_match
set flatten model -seq_constant
set flatten model -nodff_to_dlat_zero
set flatten model -nodff_to_dlat_feedback
set flatten model -nooutput_z
set root module mfs_3prm_1a -golden
set root module mfs_3prm_1a -revised
report messages
report black box
report design data
   // report floating signals
dofile $PROJECT/fv/conformal/mfs_3prm_1a.cec
   // dofile $PROJECT/fv/conformal/mfs_3prm_1a.cep
   // during compilation
set system mode lec -nomap
read mapped points $PROJECT/fv/conformal/mfs_3prm_1a.cmc

// Trivial mappings with same name registers
// read mapped points $PROJECT/fv/conformal/mfs_3prm_1a_trivial.cmc
// dofile $PROJECT/fv/conformal/mfs_3prm_1a.cmp
map key points
remodel -seq_constant -repeat
add compare points -all
compare
usage
   // exit -f
Conclusion

Formal verification software enables verification of your design during all stages, from RTL to placement and routing. Verifying designs requires more time as designs increase in size. Formal verification helps to reduce the time needed for your design verification cycle.

Document Revision History

Table 17–4 lists the revision history for this chapter.

Table 17–4. Document Revision History

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Removed survey link.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>■ Updated “Black Boxes in the Conformal LEC Flow” on page 17–8 and “Known Issues and Limitations” on page 17–16.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Removed Figures.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>Changed to new document template. Removed Table 21-1.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>Updates for new GUI changes, and added link to Help.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>Updated “Black Boxes in the Encounter Conformal Flow” section.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>Updated Table 21-1.</td>
</tr>
<tr>
<td>November 2008</td>
<td>8.1.0</td>
<td>■ Changed to 8-1/2 x 11 page size.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added support for Stratix IV devices.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added support for Cadence Conformal LEC version 7.2 and Synplify Pro version 9.6.2.</td>
</tr>
<tr>
<td>May 2008</td>
<td>8.0.0</td>
<td>■ Added support for Cyclone III devices.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Black Boxes in the Encounter Conformal Flow” section.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Table 18–1 and Table 18–5.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
The Quartus® II software offers a complete software solution for system designers who design with Altera® FPGA and CPLD devices, including device programming. The Quartus II Programmer is part of the Quartus II software package that allows you to program Altera CPLD and configuration devices, and configure Altera FPGA devices. This section describes how you can use the Quartus II Programmer to program or configure your device after you successfully compile your design.

This section includes the following chapter:

- Chapter 18, Quartus II Programmer
This chapter describes how to program and configure Altera® CPLD, FPGA, and configuration devices with the Quartus® II Programmer.

The Quartus II software offers a complete solution for system designers who design with Altera FPGA and CPLD devices. After you compile your design, you can use the Quartus II Programmer to program or configure your device, to test its functionality on a circuit board.

This chapter contains the following sections:
- “Programming Flow”
- “Quartus II Programmer GUI” on page 18–3
- “Programming and Configuration Modes” on page 18–5
- “Scripting Support” on page 18–15

For more information about how to use the Quartus II Programmer GUI to program and configure your device, refer to Programming Devices in Quartus II Help.

Programming Flow

The following steps describe the general overview of the programming flow:

1. Compile your design, such that the Quartus II Assembler generates the programming or configuration file.
2. Convert the programming or configuration file to target your configuration device and, optionally, create secondary programming files.
3. Program and configure the FPGA, CPLD, or configuration device using the programming or configuration file with the Quartus II Programmer.
Table 18–1 lists the programming and configuration file formats supported by Altera FPGAs, CPLDs, and configuration devices.

<table>
<thead>
<tr>
<th>File Format</th>
<th>FPGA</th>
<th>CPLD</th>
<th>Configuration Device</th>
<th>Serial Configuration Device</th>
</tr>
</thead>
<tbody>
<tr>
<td>SRAM Object File (.sof)</td>
<td>✔</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>Programmer Object File (.pof)</td>
<td>—</td>
<td>✔</td>
<td>✔</td>
<td>✔</td>
</tr>
<tr>
<td>JEDEC JESD71 STAPL Format File (.jam)</td>
<td>✔</td>
<td>✔</td>
<td>✔</td>
<td>—</td>
</tr>
<tr>
<td>Jam Byte Code File (.jbc)</td>
<td>✔</td>
<td>✔</td>
<td>✔</td>
<td>—</td>
</tr>
</tbody>
</table>

Figure 18–1 shows the programming file generation flow.

For more information about Chain Description Files (.cdf), refer to About Programming in Quartus II Help.
Figure 18–2 shows the Quartus II Programmer programming flow.

**Figure 18–2. Programming Flow**

Quartus II Programmer GUI

The Quartus II Programmer GUI is a window in which you can add your programming and configuration files, specify programming options and hardware, and start the programming or configuration of the device.

To open the Programmer window, on the Tools menu, click **Programmer**. As you proceed through the programming flow, the Quartus II Message window reports the status of each operation.

If the Quartus II Programmer automatically detects devices with shared JTAG IDs, the Programmer prompts you to specify the correct device in the JTAG chain.

You must add a user defined device in the Quartus II software for each unknown device in the JTAG chain and specify the instruction register length for each device.
To edit the device details of an unknown device, follow these steps:
1. Double-click on the unknown device listed under the device column.
2. Click **Edit**.
3. Change the device **Name**.
4. Enter the **Instruction register Length**.
5. Click **OK**.
6. Save the **.cdf**.

For a description of the Programmer window, refer to *Programmer Window* in Quartus II Help. For a description of options in the Tools menu, refer to *Programmer Page (Options Dialog Box)* in Quartus II Help.

**Hardware Setup**

The Quartus II Programmer provides the flexibility to choose a download cable or programming hardware. Before you can program or configure your device, you must have the correct hardware setup.

For hardware settings, refer to *Setting Up Programming Hardware* in Quartus II Help.

For more information about programming hardware driver installation, refer to the Setting up Programming Hardware in Quartus II Software page on the Altera website.

**JTAG Settings**

The JTAG server allows the Quartus II Programmer to access the JTAG hardware. You can also access the JTAG download cable or programming hardware connected to a remote computer through the JTAG server of that computer. With the JTAG server, you can control the programming or configuration of devices from a single computer through other computers at remote locations. The JTAG server uses the TCP/IP communications protocol.

For more information about JTAG settings, refer to *Using the JTAG Server* in Quartus II Help.

**JTAG Chain Debugger Tool**

The JTAG Chain Debugger tool allows you to test the JTAG chain integrity and detect intermittent failures of the JTAG chain. In addition, the tool allows you to shift in JTAG instructions and data through the JTAG interface and step through the test access port (TAP) controller state machine for debugging purposes. You access the tool from the Tools menu on the main menu of the Quartus II software.

For more information, refer to *Using the JTAG Chain Debugger* in Quartus II Help.
Other Programming Tools

The following section describes other programming tools in more detail.

Stand-Alone Quartus II Programmer

Altera offers the free stand-alone Quartus II Programmer, which has the same full functionality as the Quartus II Programmer in the Quartus II software. The stand-alone Quartus II Programmer is useful when programming your devices with another workstation, so you do not need two full licenses. You can download the stand-alone Quartus II Programmer from the Download Center on the Altera website.

Programming and Configuration Modes

The following section describes the Quartus II Programmer and the Programmer configuration modes.

Configuration Modes

The Quartus II Programmer supports five configuration modes, including JTAG, passive serial (PS), active serial (AS), Configuration via Protocol (CvP), and in-socket modes (ISM).

Table 18–2 lists the programming and configuration modes supported by Altera devices.

Table 18–2. Programming and Configuration Modes

<table>
<thead>
<tr>
<th>Mode</th>
<th>FPGA</th>
<th>CPLD</th>
<th>Configuration Device</th>
<th>Serial Configuration Device</th>
</tr>
</thead>
<tbody>
<tr>
<td>JTAG</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td></td>
</tr>
<tr>
<td>PS</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td></td>
</tr>
<tr>
<td>AS</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓</td>
</tr>
<tr>
<td>CvP</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>✓</td>
</tr>
<tr>
<td>In-Socket Programming</td>
<td>—</td>
<td>✓ (1)</td>
<td>—</td>
<td>✓</td>
</tr>
</tbody>
</table>

Note to Table 18–2:

(1) MAX II CPLDs do not support in-socket programming mode.

For more information about programming and configuration modes, refer to About Programming in Quartus II Help.

For more information about CvP configuration mode, refer to the Configuration via Protocol (CvP) Implementation in Altera FPGAs User Guide.

For more information about JTAG, PS, and AS configuration modes and in-socket programming mode, refer to the Configuration Handbook, or the device handbook or data sheet for the respective FPGA, CPLD, or configuration device.

For a list of programming adapters available for Altera devices, refer to www.altera.com.
Design Security Keys

The Quartus II Programmer supports the generation of encryption key programming files and encrypted configuration files for Altera FPGAs that support the design security feature. You can also use the Quartus II Programmer to program the encryption key into the FPGA.

For more information about using the design security feature with the Quartus II software, refer to AN 341: Using the Design Security Feature in Stratix II and Stratix II GX Devices and AN 512: Using the Design Security Feature in Stratix III Devices.

Optional Programming or Configuration Files

The Quartus II software can generate optional programming or configuration files in various formats that you can use with programming tools other than the Quartus II Programmer. When you compile a design in the Quartus II software, the Assembler automatically generates either a .sof or .pof. The Assembler also allows you to convert FPGA configuration files to programming files for configuration devices.

For more information, refer to About Optional Programming Files in Quartus II Help.

For more information about the programming and configuration file formats, refer to file format topics in the Quartus II Help or the Configuration File Formats chapter of the Configuration Handbook. For more information about using the .jam and .jbc programming files with the Jam STAPL Player, Jam STAPL Byte-Code Player, and the quartus_jli command-line executable, refer to AN 425: Using Command-Line Jam STAPL Solution for Device Programming.

Secondary Programming Files

The Quartus II software generates programming files in various formats for use with different programming tools.

Table 18–3 lists the file types generated by the Quartus II software and supported by the Quartus II Programmer.

<table>
<thead>
<tr>
<th>File Type</th>
<th>Generated by the Quartus II Software</th>
<th>Supported by the Quartus II Programmer</th>
</tr>
</thead>
<tbody>
<tr>
<td>.sof</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>.pof</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>.jam</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>.jbc</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>JTAG Indirect Configuration File (.jic)</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>Serial Vector Format File (.svf)</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>In System Configuration File (.isc)</td>
<td>✓</td>
<td>—</td>
</tr>
</tbody>
</table>
For more information, refer to *Generating Secondary Programming Files* in Quartus II Help.

### Convert Programming Files Dialog Box

The Convert Programming Files dialog box in the Programmer allows you to convert programming files from one file format to another. For example, to store the FPGA data in configuration devices, you can convert the .sof data to another format, such as .pof, .hexout, .rnf, .rpdp, or .jic, and then program the configuration device.

To access the Convert Programming Files dialog box, on the main menu of the Quartus II software, click **File**, and then click **Convert Programming Files**. You can then configure multiple devices, such as combining multiple .sof files into one .pof.

Configure multiple devices with an external host, such as a microprocessor or CPLD. For example, you can combine multiple .sof files into one configuration file. For more information about converting programming files with the Quartus II software, refer to the *Configuration File Formats* chapter of the *Configuration Handbook*.

You can use the Advanced option in the Convert Programming Files dialog box to debug your configuration. You must choose the advanced settings that apply to your Altera device. You can direct the Quartus II software to enable or disable an advanced option by turning the option on or off in the Advanced Options dialog box.

---

**Table 18–3. File Types Generated by the Quartus II Software and Supported by the Quartus II Programmer (Part 2 of 2)**

<table>
<thead>
<tr>
<th>File Type</th>
<th>Generated by the Quartus II Software</th>
<th>Supported by the Quartus II Programmer</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hexadecimal (Intel-Format) Output File (.hexout)</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Raw Binary File (.rbi)</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Raw Binary File for Partial Reconfiguration (.rbi)</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Tabular Text File (.ttf)</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Raw Programming Data File (.rpdp)</td>
<td>✓</td>
<td>—</td>
</tr>
</tbody>
</table>
When you change settings in the Advanced Options dialog box, the change affects .pof, .jic, .rpd, and .rbf files. Table 18–4 lists the Advanced Options settings in more detail.

### Table 18–4. Advanced Options Settings

<table>
<thead>
<tr>
<th>Option Setting</th>
<th>Description</th>
</tr>
</thead>
</table>
| Disable EPCS ID check                               | - FPGA skips the EPCS silicon ID verification.  
- Default setting is unavailable (EPCS ID check is enabled).  
- Applies to the single- and multi-device AS configuration modes on all FPGA devices. |
| Disable AS mode CONF_DONE error check               | - FPGA skips the CONF_DONE error check.  
- Default setting is unavailable (AS mode CONF_DONE error check is enabled).  
- Applies to single- and multi-device (AS) configuration modes on all FPGA devices.  
- The CONF_DONE error check is disabled by default for Stratix V, Arria V, and Cyclone V devices for AS-PS multi device configuration mode. |
| Program Length Count adjustment                     | - Specifies the offset you can apply to the computed PLC of the entire bitstream.  
- Default setting is 0. The value should be an integer.  
- Applies to single- and multi-device (AS) configuration modes on all FPGA devices. |
| Post-chain bitstream pad bytes                      | - Specifies the number of pad bytes appended to the end of an entire bitstream.  
- Default value is set to 0 if the bitstream of the last device is uncompressed. Set to 2 if the bitstream of the last device is compressed. |
| Post-device bitstream pad bytes                     | - Specifies the number of pad bytes appended to the end of the bitstream of a device.  
- Default value is 0. No negative integer.  
- Applies to all single-device configuration modes on all FPGA devices. |
| Bitslice padding value                              | - Specifies the padding value used to prepare bitslice configuration bitstreams, such that all bitslice configuration chains simultaneously receive their final configuration data bit.  
- Default value is 1. Valid setting is 0 or 1.  
- Use only in 2, 4, and 8-bit PS configuration mode, when you use an EPC device with the decompression feature enabled.  
- Applies to all FPGA devices that support enhanced configuration devices. |

Table 18–5 lists symptoms you may encounter if a configuration fails, and describes the advanced options you must use to debug your configuration.

### Table 18–5. Failure Symptoms and Options Settings  (Part 1 of 2)

<table>
<thead>
<tr>
<th>Failure Symptoms</th>
<th>Disable EPCS ID Check</th>
<th>Disable AS Mode CONF_DONE Error Check</th>
<th>PLC Settings</th>
<th>Post-Chain Bitstream Pad Bytes</th>
<th>Post-Device Bitstream Pad Bytes</th>
<th>Bitslice Padding Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Configuration failure occurs after a configuration cycle.</td>
<td>—</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
</tr>
<tr>
<td>Decompression feature is enabled.</td>
<td>—</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>✓</td>
<td>—</td>
</tr>
</tbody>
</table>
Table 18–5. Failure Symptoms and Options Settings (Part 2 of 2)

<table>
<thead>
<tr>
<th>Failure Symptoms</th>
<th>Disable EPCS ID Check</th>
<th>Disable AS Mode CONF_DONE Error Check</th>
<th>PLC Settings</th>
<th>Post-Chain Bitstream Pad Bytes</th>
<th>Post-Device Bitstream Pad Bytes</th>
<th>Bitslice Padding Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Encryption feature is enabled.</td>
<td>—</td>
<td>☑</td>
<td>☑</td>
<td>✓ (Use only for multi-device chain)</td>
<td>✓ (Use only for single-device chain)</td>
<td>—</td>
</tr>
<tr>
<td>CONF_DONE stays low after a configuration cycle.</td>
<td>—</td>
<td>☑</td>
<td>✓ (Start with positive offset to the PLC settings)</td>
<td>✓ (Use only for multi-device chain)</td>
<td>✓ (Use only for single-device chain)</td>
<td>—</td>
</tr>
<tr>
<td>CONF_DONE goes high momentarily after a configuration cycle.</td>
<td>—</td>
<td>☑</td>
<td>✓ (Start with negative offset to the PLC settings)</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>FPGA does not enter user mode even though CONF_DONE goes high.</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓ (Use only for multi-device chain)</td>
<td>✓ (Use only for single-device chain)</td>
<td>—</td>
</tr>
<tr>
<td>Configuration failure occurs at the beginning of a configuration cycle.</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>Newly introduced EPCS, such as EPCS128.</td>
<td>✓</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>Failure in .pof generation for EPC device using Quartus II Convert Programming File Utility when the decompression feature is enabled.</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>—</td>
<td>✓</td>
</tr>
</tbody>
</table>

For more information about the Convert Programming Files dialog box, refer to Convert Programming Files Dialog Box in Quartus II Help.
Converting Programming Files for Partial Reconfiguration

Beginning from the Quartus II software version 12.1, the Convert Programming File dialog box supports the following programming file generation and option for Partial Reconfiguration:

- Partial-Masked SRAM Object File (.pmsf) output file generation, with .msf and .sof as input files.
- .rbf for Partial Reconfiguration output file generation, with a .pmsf as the input file.

The .rbf for Partial Reconfiguration file is only for Partial Reconfiguration.

- Providing the Enable decompression during Partial Reconfiguration option to enable the option bit for bitstream decompression during Partial Reconfiguration, when converting a .sof (full design .sof) to any supported file type.

For more information about Partial Reconfiguration, refer to the Design Planning for Partial Reconfiguration chapter in volume 1 of the Quartus II Handbook.

Generating .pmsf using a .msf and a .sof

You can generate a .pmsf with a .msf and a .sof in the Convert Programming Files dialog box.

To generate the .pmsf in the Convert Programming Files dialog box, follow these steps:

1. In the Convert Programming Files dialog box, under the Programming file type field, select Partial-Masked SRAM Object File (.pmsf).
2. In the File name field, specify the necessary output file name.
3. In the Input files to convert field, add necessary input files to convert. You can add only a .msf and .sof.
4. Click Generate to generate the .pmsf.
Figure 18–3. Generating .pmsf in the Convert Programming Files Dialog Box
Generating .rbf for Partial Reconfiguration Using a .pmsf

After you have successfully generated the .pmsf, you can now convert the .pmsf to a .rbf for Partial Reconfiguration in the Convert Programming Files dialog box.

To generate the .rbf for Partial Reconfiguration, follow these steps:

1. In the Convert Programming Files dialog box, in the Programming file type field, select Raw Binary File for Partial Reconfiguration (.rbf).
2. In the File name field, specify the output file name.
3. In the Input files to convert field, add input files to convert. You can add only a .pmsf.
4. After adding the .pmsf, select the .pmsf and click Properties. The PMSF File Properties dialog box appears.
5. Make your selection either by turning on or turning off the Compression, Enable SCRUB mode, and Write memory contents options.
   - Compression option—This option enables compression on Partial Reconfiguration bitstream.
   - Enable SCRUB mode option—The default of this option is based on AND/OR mode. This option is valid only when Partial Reconfiguration masks in your design are not overlapped vertically. Otherwise, you cannot generate the .rbf for Partial Reconfiguration.
   - Write memory contents option—This option is a workaround for initialized RAM/ROM in a Partial Reconfiguration region.

   For more information about these option, refer to the Design Planning for Partial Reconfiguration chapter in volume 1 of the Quartus II Handbook.

6. Click OK.
7. Click Generate to generate the .rbf for Partial Reconfiguration.
Figure 18–4. Generating .rbf for Partial Reconfiguration in the Convert Programming Files Dialog Box
Enable Decompression during Partial Reconfiguration Option

Beginning from the Quartus II software version 12.1, you can turn on the Enable decompression during Partial Reconfiguration option in the SOF File Properties: Bitstream Encryption dialog box, which can be accessed from the Convert Programming File dialog box. This option is available when converting a .sof to any supported programming file types listed in Table 18–3 on page 18–6.

This option is hidden for other targeted devices that do not support Partial Reconfiguration. To view this option in the SOF File Properties: Bitstream Encryption dialog box, the .sof must be targeted on an Altera device that supports Partial Reconfiguration.

If you turn on the Compression option when generating the .rbf for Partial Reconfiguration (Figure 18–4), then you must turn the Enable decompression during Partial Reconfiguration option on.

Figure 18–5 shows the SOF File Properties: Bitstream Encryption dialog box.

Figure 18–5. SOF File Properties: Enable Decompression During Partial Reconfiguration
Flash Loaders

Parallel and serial configuration devices do not support the JTAG interface. However, you can use a flash loader to program configuration devices in-system via the JTAG interface. You can use an FPGA as a bridge between the JTAG interface and the configuration device. The Quartus II software supports parallel and serial flash loaders.

For more information, refer to About Flash Loaders in Quartus II Help.

Scripting Support

In addition to the Quartus II Programmer GUI, you can use the Quartus II command-line executable quartus_pgm.exe to access programmer functionality from the command line and from scripts. The programmer accepts .pof, .sof, and .jic programming or configuration files and Chain Description Files (.cdf).

Example 18–1 shows a command that programs a device:

Example 18–1. Programming a Device

quartus_pgm –c byteblasterII –m jtag –o bpv;design.pof

Where:
■ -c byteblasterII specifies the ByteBlaster™ II download cable
■ -m jtag specifies the JTAG programming mode
■ -o bpv represents the blank-check, program, and verify operations
■ design.pof represents the .pof used for the programming

The Programmer automatically executes the erase operation before programming the device.

For more information about scripting command options, refer to About Quartus II Scripting in Quartus II Help.

The jtagconfig Debugging Tool

You can use the jtagconfig command-line utility (which is similar to the auto detect operation in the Quartus II Programmer) to check the devices in a JTAG chain and the user-defined devices.

For more information about the jtagconfig utility, type one of the following commands at the command prompt:

Example 18–2.

jtagconfig -h
jtagconfig --help

The help switch does not reference the -n switch. The jtagconfig -n command shows each node for each jtag device.
For more information about command-line scripting, refer to the Command-Line Scripting chapter in volume 2 of the Quartus II Handbook.

**Generating .pmsf using a .msf and a .sof**

You can generate a .pmsf with the quartus_cpf command by typing the following commands:

**Example 18–3.**

```
quartus_cpf -p <pr_revision.msf> <pr_revision.sof> <new_filename.pmsf>
```

**Generating .rbf for Partial Reconfiguration using a .pmsf**

You can generate a .rbf for Partial Reconfiguration with the quartus_cpf command by typing the following commands:

**Example 18–4.**

```
quartus_cpf –o foo.txt –c <pr_revision.pmsf> <pr_revision.rbf>
```

You must run this command in the same directory where the files are located.

**Conclusion**

The Quartus II Programmer offers you a wide variety of options to program and configure your Altera devices. With the Quartus II Programmer, the Quartus II software provides you with a complete solution for your FPGA or CPLD design prototyping, which you can also use in the production environment.

**Document Revision History**

Table 18–6 lists the revision history for this chapter.

**Table 18–6. Document Revision History (Part 1 of 2)**

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
</tr>
</thead>
<tbody>
<tr>
<td>November 2012</td>
<td>12.1.0</td>
<td>■ Updated Table 18–3 on page 18–6, and Table 18–4 on page 18–8.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Scripting Support” on page 18–15.</td>
</tr>
<tr>
<td>June 2012</td>
<td>12.0.0</td>
<td>■ Updated Table 18–5 on page 18–8.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated “Quartus II Programmer GUI” on page 18–3.</td>
</tr>
<tr>
<td>November 2011</td>
<td>11.1.0</td>
<td>■ Updated “Configuration Modes” on page 18–5.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Added “Optional Programming or Configuration Files” on page 18–6.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>■ Updated Table 18–2 on page 18–5.</td>
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</table>
Table 18–6. Document Revision History (Part 2 of 2)

<table>
<thead>
<tr>
<th>Date</th>
<th>Version</th>
<th>Changes</th>
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<tbody>
<tr>
<td>May 2011</td>
<td>11.0.0</td>
<td>▪ Added links to Quartus II Help.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Updated “Hardware Setup” on page 21–4 and “JTAG Chain Debugger Tool” on page 21–4.</td>
</tr>
<tr>
<td>December 2010</td>
<td>10.1.0</td>
<td>▪ Changed to new document template.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Updated “JTAG Chain Debugger Example” on page 20–4.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Added links to Quartus II Help.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Reorganized chapter.</td>
</tr>
<tr>
<td>July 2010</td>
<td>10.0.0</td>
<td>▪ Added links to Quartus II Help.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Deleted screen shots.</td>
</tr>
<tr>
<td>November 2009</td>
<td>9.1.0</td>
<td>▪ No change to content.</td>
</tr>
<tr>
<td>March 2009</td>
<td>9.0.0</td>
<td>▪ Added a row to Table 21–4.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Changed references from “JTAG Chain Debug” to “JTAG Chain Debugger”.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>▪ Updated figures.</td>
</tr>
</tbody>
</table>

For previous versions of the Quartus II Handbook, refer to the Quartus II Handbook Archive.
Additional Information

This chapter provides additional information about the document and Altera.

About this Handbook

This handbook provides comprehensive information about the current version of the Altera® Quartus® II design software.

How to Contact Altera

To locate the most up-to-date information about Altera products, refer to the following table.

<table>
<thead>
<tr>
<th>Contact (1)</th>
<th>Contact Method</th>
<th>Address</th>
</tr>
</thead>
<tbody>
<tr>
<td>Technical support</td>
<td>Website</td>
<td><a href="http://www.altera.com/support">www.altera.com/support</a></td>
</tr>
<tr>
<td>Technical training</td>
<td>Website</td>
<td><a href="http://www.altera.com/training">www.altera.com/training</a></td>
</tr>
<tr>
<td>Product literature</td>
<td>Email</td>
<td><a href="mailto:custrain@altera.com">custrain@altera.com</a></td>
</tr>
<tr>
<td>Nontechnical support (general)</td>
<td>Website</td>
<td><a href="http://www.altera.com/literature">www.altera.com/literature</a></td>
</tr>
<tr>
<td>(software licensing)</td>
<td>Email</td>
<td><a href="mailto:nacomp@altera.com">nacomp@altera.com</a></td>
</tr>
<tr>
<td>(software licensing)</td>
<td>Email</td>
<td><a href="mailto:authorization@altera.com">authorization@altera.com</a></td>
</tr>
</tbody>
</table>

Note to Table:
(1) You can also contact your local Altera sales office or sales representative.

Third-Party Software Product Information

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## Typographic Conventions

The following table shows the typographic conventions this document uses.

<table>
<thead>
<tr>
<th>Visual Cue</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Bold Type with Initial Capital Letters</strong></td>
<td>Indicate command names, dialog box titles, dialog box options, and other GUI labels. For example, <strong>Save As</strong> dialog box. For GUI elements, capitalization matches the GUI.</td>
</tr>
<tr>
<td><strong>bold type</strong></td>
<td>Indicates directory names, project names, disk drive names, file names, file name extensions, software utility names, and GUI labels. For example, <code>\qdesigns</code> directory, <strong>D:</strong> drive, and <strong>chiptrip.gdf</strong> file.</td>
</tr>
<tr>
<td><strong>Italic Type with Initial Capital Letters</strong></td>
<td>Indicate document titles. For example, <em>Stratix IV Design Guidelines</em>.</td>
</tr>
<tr>
<td><strong>italic type</strong></td>
<td>Indicates variables. For example, (n + 1). Variable names are enclosed in angle brackets (&lt;&gt;). For example, <code>&lt;file name&gt;</code> and <code>&lt;project name&gt;.pof</code> file.</td>
</tr>
<tr>
<td>Initial Capital Letters</td>
<td>Indicate keyboard keys and menu names. For example, the Delete key and the Options menu.</td>
</tr>
<tr>
<td>“Subheading Title”</td>
<td>Quotation marks indicate references to sections in a document and titles of Quartus II Help topics. For example, “Typographic Conventions.”</td>
</tr>
<tr>
<td><strong>Courier type</strong></td>
<td>Indicates signal, port, register, bit, block, and primitive names. For example, <code>data1</code>, <code>tdi</code>, and <code>input</code>. The suffix <code>n</code> denotes an active-low signal. For example, <code>resetn</code>. Indicates command line commands and anything that must be typed exactly as it appears. For example, <code>c:\qdesigns\tutorial\chiptrip.gdf</code>. Also indicates sections of an actual file, such as a Report File, references to parts of files (for example, the AHDL keyword <code>SUBDESIGN</code>), and logic function names (for example, <code>TRI</code>).</td>
</tr>
<tr>
<td>↔</td>
<td>An angled arrow instructs you to press the Enter key.</td>
</tr>
<tr>
<td>1., 2., 3., and a., b., c., and so on</td>
<td>Numbered steps indicate a list of items when the sequence of the items is important, such as the steps listed in a procedure.</td>
</tr>
<tr>
<td>■ ■ ■</td>
<td>Bullets indicate a list of items when the sequence of the items is not important.</td>
</tr>
<tr>
<td>🥤</td>
<td>The hand points to information that requires special attention.</td>
</tr>
<tr>
<td>🤔</td>
<td>The question mark directs you to a software help system with related information.</td>
</tr>
<tr>
<td>🦵</td>
<td>The feet direct you to another document or website with related information.</td>
</tr>
<tr>
<td>🎥</td>
<td>The multimedia icon directs you to a related multimedia presentation.</td>
</tr>
<tr>
<td>⚠</td>
<td>A caution calls attention to a condition or possible situation that can damage or destroy the product or your work.</td>
</tr>
<tr>
<td>⚠</td>
<td>A warning calls attention to a condition or possible situation that can cause you injury.</td>
</tr>
<tr>
<td>💌</td>
<td>The envelope links to the Email Subscription Management Center page of the Altera website, where you can sign up to receive update notifications for Altera documents.</td>
</tr>
<tr>
<td>📝</td>
<td>The feedback icon allows you to submit feedback to Altera about the document. Methods for collecting feedback vary as appropriate for each document.</td>
</tr>
<tr>
<td>🚀</td>
<td>The social media icons allow you to inform others about Altera documents. Methods for submitting information vary as appropriate for each medium.</td>
</tr>
</tbody>
</table>